OpenCL Programming & Optimization Case Study

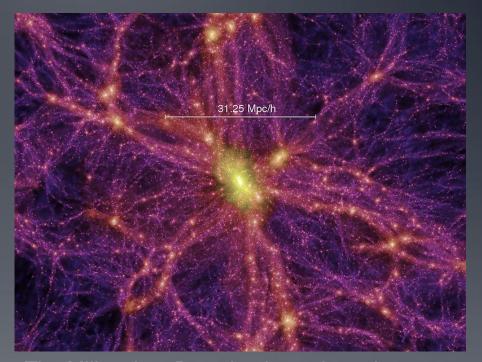
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Instructor Notes

- Lecture discusses parallel implementation of a simple embarrassingly parallel nbody algorithm
 - We aim to provide some correspondence between the architectural techniques discussed in the optimization lectures and a real scientific computation algorithm
 - We implement a N-body formulation of gravitational attraction between bodies
- Two possible optimizations on a naïve nbody implementation
 - Data reuse vs. non data reuse
 - Unrolling of inner loop
- Simple example based on the AMD Nbody SDK example
- Example provided allows user to:
 - Change data size to study scaling
 - Switch data reuse on or off
 - Change unrolling factor of n-body loop
 - Run example without parameters for help on how to run the example and define optimizations to be applied to OpenCL kernel

Topics

- N-body simulation algorithm on GPUs in OpenCL
- Algorithm, implementation
- Optimization Spaces
- Performance



The Millennium Run simulates the universe until the present state, where structures are abundant, manifesting themselves as stars, galaxies and clusters

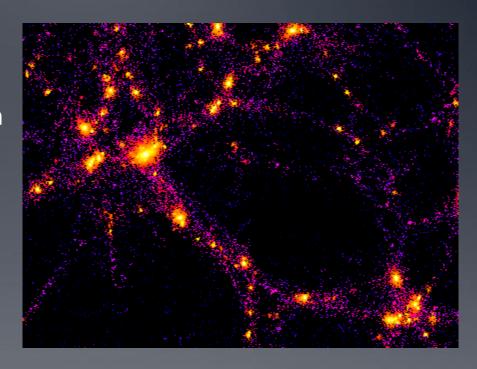
Source: http://en.wikipedia.org/wiki/N-body_simulation

Motivation

- We know how to optimize a program in OpenCL by taking advantage of the underlying architecture
 - We have seen how to utilize threads to hide latency
 - We have also seen how to take advantage of the different memory spaces available in today's GPUs.
- How do we apply this to an real world useful scientific computing algorithm?
- Our aim is to use the architecture specific optimizations discussed in previous lectures

N-body Simulation

- An n-body simulation is a simulation of a system of particles under the influence of physical forces like gravity
 - E.g.: An astrophysical system where a particle represents a galaxy or an individual star
- N² particle-particle interactions
 - Simple, highly data parallel algorithm
- Allows us to explore optimizations of both the algorithm and its implementation on a platform



Source: THE GALAXY-CLUSTER-SUPERCLUSTER CONNECTION http://www.casca.ca/ecass/issues/1997-DS/West/west-bil.html

Algorithm

- The gravitational attraction between two bodies in space is an example of an N-body problem
 - Each body represents a galaxy or an individual star, and bodies attract each other through gravitational force
- Any two bodies attract each other through gravitational forces (F)

$$F = G * \left(\frac{m_i * m_j}{\parallel r_{ij} \parallel^2}\right) * \frac{r_{ij}}{\parallel r_{ij} \parallel}$$

F = Resultant Force Vector between particles i and

G = Gravitational Constant

 $m_i = Mass of particle i$

 $m_i = \text{Mass of particle j}$

 $|\mathbf{r}_{ij}| = |\mathbf{D}|$ istance of particle i and j

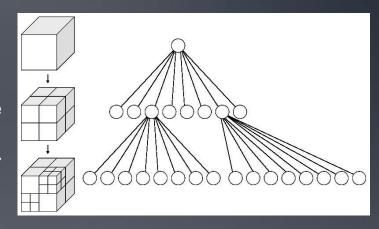
For each particle this becomes

$$F_i = (G * m_i) * \sum_{j=1 \to N} \left(\frac{m_j}{||r_{ij}||^2} * \left(\frac{r_{ij}}{||r_{ij}||} \right) \right)$$

- An O(N²) algorithm since N*N interactions need to be calculated
- This method is known as an all-pairs N-body simulation

N-body Algorithms

- For large counts, the previous method calculates of force contribution of distant particles
 - Distant particles hardly affect resultant force
- Algorithms like Barnes Hut reduce number of particle interactions calculated
 - Volume divided into cubic cells in an octree
 - Only particles from nearby cells need to be treated individually
 - Particles in distant cells treated as a single large particle
- In this lecture we restrict ourselves to a simple all pair simulation of particles with gravitational forces



- A octree is a tree where a node has exactly 8 children
- Used to subdivide a 3D space

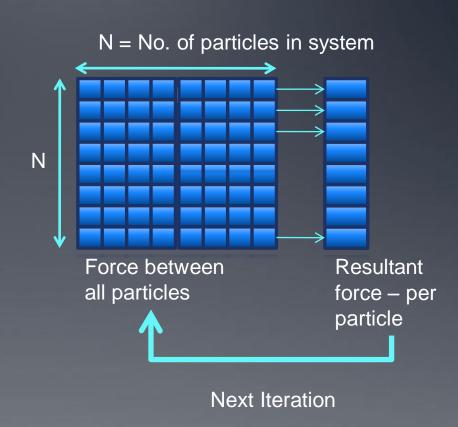
Basic Implementation – All pairs

- All-pairs technique is used to calculate close-field forces
- Why bother, if infeasible for large particle counts?
 - Algorithms like Barnes Hut calculate far field forces using near-field results
 - Near field still uses all pairs
 - So, implementing all pairs improves performance of both near and far field calculations
- Easy serial algorithm
 - Calculate force by each particle
 - Accumulate of force and displacement in result vector

```
for(i=0; i< n; i++)
   ax = ay = az = 0;
   for (j=0; j< n; j++)
     //Calculate Displacement
      dx=x[j]-x[i];
      dy=y[j]-y[i];
      dz=z[j]-z[i];
      invr = 1.0/sqrt(dx*dx+dy*dy+dz*dz
+eps);
      invr3 = invr*invr*invr;
      f=m[ | ]*invr3;
      ax += f*dx;
      ay += f*dy;
      az += f*dx:
```

Parallel Implementation

- Forces of each particle can be computed independently
 - Accumulate results in local memory
 - Add accumulated results to previous position of particles
- New position used as input to the next time step to calculate new forces acting between particles



Naïve Parallel Implementation

- Disadvantages of implementation where each work item reads data independently
 - No reuse since redundant reads of parameters for multiple work-items
 - Memory access= N reads*N threads= N²
- Similar to naïve non blocking matrix multiplication in Lecture 5

```
p items
/workgroup

N = No. of particles

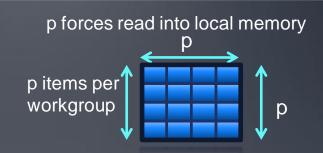
All N particles read in by each

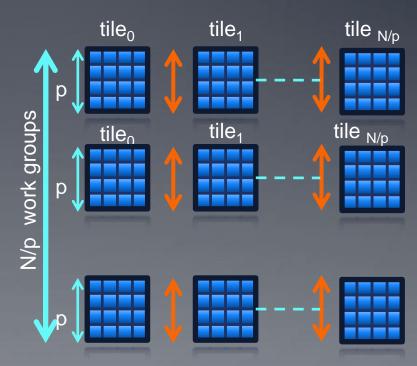
work item
```

```
kernel void nbody(
      global float4 * initial_pos,
      _global float4 * final_pos,
   Int N, __local float4 * result) {
int localid = get_local_id(0);
int globalid = get_global_id(0);
result [localid] = 0;
for( int i=0 ; i<N;i++) {
   //! Calculate interaction between
          GetForce(
                globalid, i,
                initial pos, final pos,
                &result [localid]);
   finalpos[globalid] = result[localid];
```

Local Memory Optimizations

- Data Reuse
 - Any particle read into compute unit can be used by all p bodies
- Computational tile:
 - Square region of the grid of forces consisting of size p
 - 2p descriptions required to evaluate all p² interactions in tile
 - p work items (in vertical direction) read in p forces
- Interactions on p bodies captured as an update to p acceleration vectors
- Intra-work group synchronization shown in orange required since all work items use data read by each work item





OpenCL Implementation

- Data reuse using local memory
 - Without reuse N*p items read per work group
 - With reuse p*(N/p) = N items read per work group
 - All work items use data read in by each work item
- SIGNIFICANT improvement: p is work group size (at least 128 in OpenCL, discussed in occupancy)
- Loop nest shows how a work item traverses all tiles
- Inner loop accumulates contribution of all particles within tile

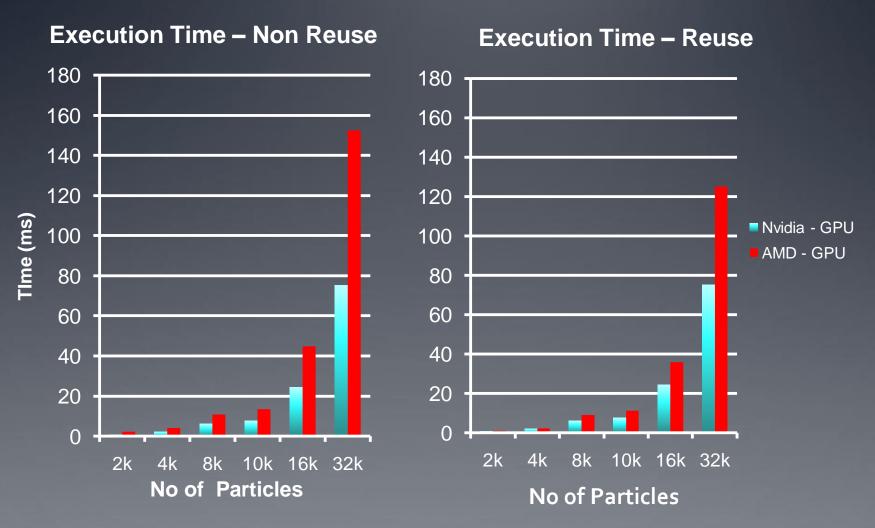
Kernel Code

```
for (int i = 0; i < numTiles; ++i)
     // load one tile into local memory
     int idx = i * localSize + tid;
     localPos[tid] = pos[idx];
    barrier(CLK_LOCAL_MEM_FENCE);
    // calculate acceleration effect due to each body
   for( int j = 0; j < localSize; ++j)
        float4 r = localPos[i] - myPos;
       float distSqr = r.x * r.x + r.y * r.y + r.z * r.z;
        float invDist = 1.0f / sqrt(distSqr + epsSqr);
       float s = localPos[j].w * invDistCube;
        acc += s * r;
     barrier(CLK_LOCAL_MEM_FENCE);
```

Performance

- Effect of optimizations compared for two GPU platforms
 - Exactly same code, only recompiled for platform
- Devices Compared
 - AMD GPU = 5870 Stream SDK 2.2
 - Nvidia GPU = GTX 480 with CUDA 3.1
- Time measured for OpenCL kernel using OpenCL event counters (covered in in Lecture 11)
 - Device IO and other overheads like compilation time are not relevant to our discussion of optimizing a compute kernel
 - Events are provided in the OpenCL spec to query obtain timestamps for different state of OpenCL commands

Effect of Reuse on Kernel Performance



Thread Mapping and Vectorization

- As discussed in Lecture 8
 - Vectorization allows a single thread to perform multiple operations in one instruction
 - Explicit vectorization is achieved by using vector data-types (such as float4) in the source
- Vectorization using float4 enables efficient usage of memory bus for AMD GPUs
 - Easy in this case due to the vector nature of the data which are simply spatial coordinates
- Vectorization of memory accesses implemented using float4 to store coordinates

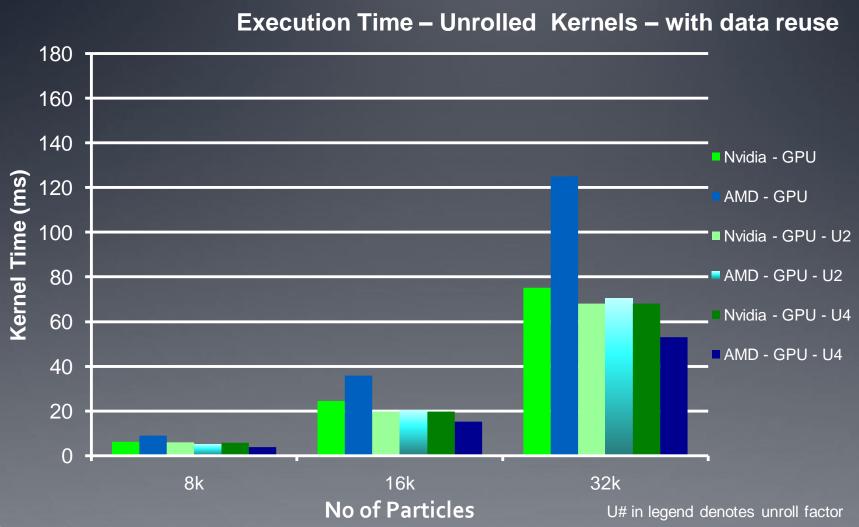
```
// Loop Nest Calculating per tile interaction
// Same loop nest as in previous slide
for (int i=0 ; i< numTiles; i++)
{
    for( int j=0; j<localsize; j ++)
}</pre>
```

}

Performance - Loop Unrolling

- We also attempt loop unrolling of the reuse local memory implementation
 - We unroll the innermost loop within the thread
- Loop unrolling can be used to improve performance by removing overhead of branching
 - However this is very beneficial only for tight loops where the branching overhead is comparable to the size of the loop body
 - Experiment on optimized local memory implementation
 - Executable size is not a concern for GPU kernels
- We implement unrolling by factors of 2 and 4 and we see substantial performance gains across platforms
 - Decreasing returns for larger unrolling factors seen

Effect of Unroll on Kernel Performance



Performance Conclusions

- From the performance data we see:
- Unrolling greatly benefits the AMD GPU for our algorithm
 - This can be attributed to better packing of the VLIW instructions within the loop
 - Better packing is enabled by the increased amount of computation uninterrupted by a conditional statement when we unroll a loop
- The Nvidia Fermi performance doesn't benefit as much from the reuse and tiling
 - This can be attributed to the caching which is possible in the Fermi GPUs. The caching could enable data reuse.
 - As seen the Fermi performance is very close for both with and without reuse
- Note: You can even run this example on the CPU to see performance differences

Provided Nbody Example

- A N-body example is provided for experimentation and explore GPU optimization spaces
- Stand-alone application based on simpler on AMD SDK formulation
 - Runs correctly on AMD and Nvidia hardware
- Three kernels provided
 - Simplistic formulation
 - Using local memory tiling
 - Using local memory tiling with unrolling
- Note: Code is not meant to be a high performance N-body implementation in OpenCL
 - The aim is to serve as an optimization base for a data parallel algorithm



Screenshot of provided N-body demo running 10k particles on a AMD 5870

References for Nbody Simulation

- AMD Stream Computing SDK and the Nvidia GPU Computing SDK provide example implementations of N-body simulations
- Our implementation extends the AMD SDK example. The Nvidia implementation implementation is a different version of the N-body algorithm
- Other references
 - Nvidia GPU Gems
 - http://http.developer.nvidia.com/GPUGems3/gpugems3_ch31.html
 - Brown Deer Technology
 - http://www.browndeertechnology.com/docs/BDT_OpenCL_Tutorial_NBody.html

Summary

- We have studied a application as an optimization case study for OpenCL programming
- We have understood how architecture aware optimizations to data and code improve performance
- Optimizations discussed do not use device specific features like Warp size, so our implementation yields performance improvement while maintaining correctness on both AMD and Nvidia GPUs

Other Resources

- Well known optimization case studies for GPU programming
- CUDPP and Thrust provide CUDA implementations of parallel prefix sum, reductions and parallel sorting
 - The parallelization strategies and GPU specific optimizations in CUDA implementations can usually be applied to OpenCL as well
 - CUDPP http://code.google.com/p/cudpp/
 - Thrust http://code.google.com/p/thrust/
- Histogram on GPUs
 - The Nvidia and AMD GPU computing SDK provide different implementations for 64 and 256 bin Histograms
- Diagonal Sparse Matrix Vector Multiplication OpenCL Optimization example on AMD GPUs and multi-core CPUs
 - http://developer.amd.com/documentation/articles/Pages/OpenCL-Optimization-Case-Study.aspx
- GPU Gems which provides a wide range of CUDA optimization studies