OpenCL Buffers and Complete Examples

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Instructor Notes

- This is a brief lecture which goes into some more details on OpenCL memory objects
 - Describes various flags that can be used to change how data is handled between host and device, like page-locked I/O and so on
- The aim of this lecture is to cover required OpenCL host code for buffer management and provide simple examples
- Code for context and buffer management discussed in examples in this lecture serves as templates for more complicated kernels
 - This allows the next 3 lectures to be focused solely on kernel optimizations like blocking, thread grouping and so on
- Examples covered
 - Simple image rotation example
 - Simple non-blocking matrix-matrix multiplication

Topics

- Using OpenCL buffers
 - Declaring buffers
 - Enqueue reading and writing of buffers
- Simple but complete examples
 - Image Rotation
 - Non-blocking Matrix Multiplication

Creating OpenCL Buffers

- Data used by OpenCL devices is stored in a "buffer" on the device
- An OpenCL buffer object is created using the following function

- Data can implicitly be copied to the device using a host pointer parameter
 - In this case copy to device is invoked when kernel is enqueued

Memory Flags

 Memory flag field in clCreateBuffer() allows us to define characteristics of the buffer object

Memory Flag	Behavior
CL_MEM_READ_WRITE	
CL_MEM_WRITE_ONLY	Specifies memory read / write behavior
CL_MEM_READ_ONLY	
CL_MEM_USE_HOST_PTR	Implementations can cache the contents pointed to by host_ptr in device memory. This cached copy can be used when kernels are executed on a device.
CL_MEM_ALLOC_HOST_PTR	Specifies to the implementation to allocate memory from host accessible memory.
CL_MEM_COPY_HOST_PTR	Specifies to allocate memory for the object and copy the data from memory referenced by host_ptr.

Copying Buffers to Device

- clEnqueueWriteBuffer() is used to write a buffer object to device memory (from the host)
- Provides more control over copy process than using host pointer functionality of clCreateBuffer()
 - Allows waiting for events and blocking

```
cl_int clEnqueueWriteBuffer (
    cl_command_queue queue,
    cl_mem buffer,
    cl_bool blocking_read,
    size_t offset,
    void *ptr,
    cl_uint num_in_wait_list,
    const cl_event * event_wait_list,
    cl_event *event)
//Command queue to device
//Blocking/Non-Blocking Flag
//Blocking/Non-Blocking Flag
//Command queue to device
//Blocking/Non-Blocking Flag
//Blocking/Non-Blocking Flag
//Command queue to device
//Blocking/Non-Blocking Flag
//Blocking/Non-Blocking Flag
//Command queue to device
//Blocking/Non-Blocking Flag
//Command queue to device
//Blocking/Non-Blocking Flag
//Command queue to device
//Blocking/Non-Blocking Flag
//Blocking/Non-Blocking Flag
//Command queue
```

Copying Buffers to Host

- clEnqueueReadBuffer() is used to read from a buffer object from device to host memory
- Similar to clEnqueueWriteBuffer()

```
cl_int clEnqueueReadBuffer (
    cl_command_queue queue,
    cl_mem buffer,
    cl_bool blocking_read,
    size_t offset,
    void *ptr,
    cl_uint num_in_wait_list,
    const cl_event * event_wait_list,
    cl_event *event)
//Command queue to device
//OpenCL Buffer Object
//Blocking/Non-Blocking Flag
//Blocking/Non-Blocking Flag
//Size of data
//Size of data
//Number of events in wait list
//Number of events in wait list
//Array of events to wait for
//Event handler for this function
```

 The vector addition example discussed in Lecture 2 and 3 provide simple code snipped for moving data to and from devices

Example 1 - Image Rotation

- A common image processing routine
 - Applications in matching, alignment, etc.
- New coordinates of point (x₁,y₁) when rotated by an angle Θ around (x₀,y₀)

$$x_2 = \cos(\theta) * (x_1 - x_0) - \sin(\theta) * (y_1 - y_0) + x_0$$

$$y_2 = \sin(\theta) * (x_1 - x_0) + \cos(\theta) * (y_1 - y_0) + x_0$$

By rotating the image about the origin (0,0) we get

$$x_2 = \cos(\theta) * (x_1) - \sin(\theta) * (y_1)$$

 $y_2 = \sin(\theta) * (x_1) + \cos(\theta) * (y_1)$

 Each coordinate for every point in the image can be calculated independently

Original Image

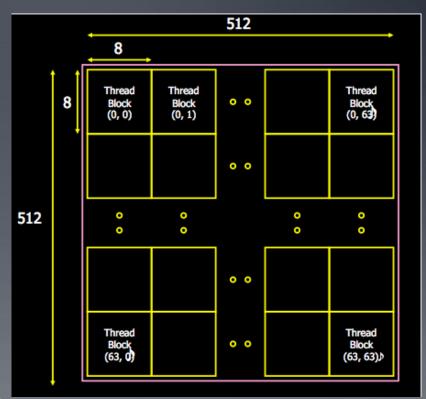


Rotated Image (90°)



Image Rotation

- Input: To copy to device
 - Image (2D Matrix of floats)
 - Rotation parameters
 - Image dimensions
- Output: From device
 - Rotated Image
- Main Steps
 - Copy image to device by enqueueing a write to a buffer on the device from the host
 - Run the Image rotation kernel on input image
 - Copy output image to host by enqueueing a read from a buffer on the device



The OpenCL Kernel

- Parallel portion of the algorithm off-loaded to device
 - Most thought provoking part of coding process
- Steps to be done in Image Rotation kernel
 - Obtain coordinates of work item in work group
 - Read rotation parameters
 - Calculate destination coordinates
 - Read input and write rotated output at calculated coordinates
- Parallel kernel is not always this obvious.
 - Profiling of an application is often necessary to find the bottlenecks and locate the data parallelism
- In this example grid of output image decomposed into work items
 - Not all parts of the input image copied to the output image after rotation, corners of I/P image could be lost after rotation

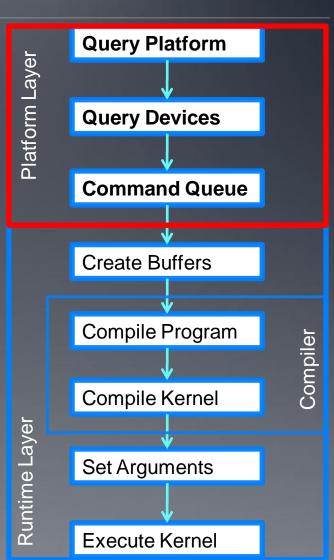
OpenCL Kernel

```
kernel void image rotate(
       _global float * src_data, __global float * dest_data, __//Data in global memory
     int W, int H,
                                                            //Image Dimensions
     float sinTheta, float cosTheta)
                                                            //Rotation Parameters
     //Thread gets its index within index space
     const int ix = get_global_id(0);
     const int iy = get_global_id(1);
     //Calculate location of data to move into ix and iy- Output decomposition as
mentioned
     float xpos = ( ((float) ix)*cosTheta + ((float)iy)*sinTheta);
     float ypos = ( ((float) iy)*cosTheta - ((float)ix)*sinTheta);
     if (( ((int)xpos>=0) && ((int)xpos< W)))
           && (((int)ypos>=0) && ((int)ypos< H)))
          dest data[iy*W+ix]=
                src_data[(int)(floor(ypos*W+xpos))];
```

Step0: Initialize Device

- Declare context
- Choose a device from context
- Using device and context create a command queue

```
ciErrNum = clGetDevicelDs (0,
CL_DEVICE_TYPE_GPU,
1, &device, cl_uint *num_devices)
```

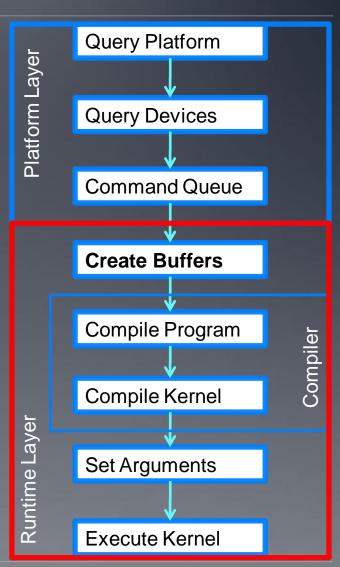


Step1: Create Buffers

- Create buffers on device
 - Input data is read-only
 - Output data is write-only

Transfer input data to the device

```
ciErrNum = clEnqueueWriteBuffer (
myqueue , d_ip, CL_TRUE,
0, mem_size, (void *)src_image,
0, NULL, NULL)
```



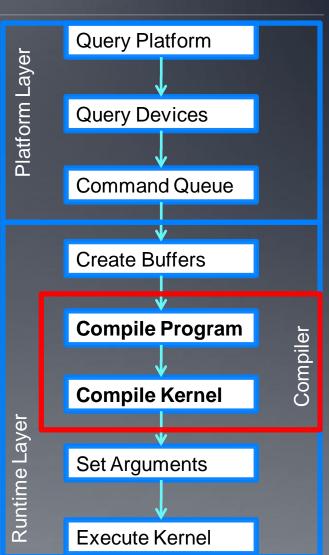
Step2: Build Program, Select Kernel

// create the program

// build the program

ciErrNum = clBuildProgram(myprog, 0, NULL, NULL, NULL, NULL);

//Use the "image rotate" function as the kernel



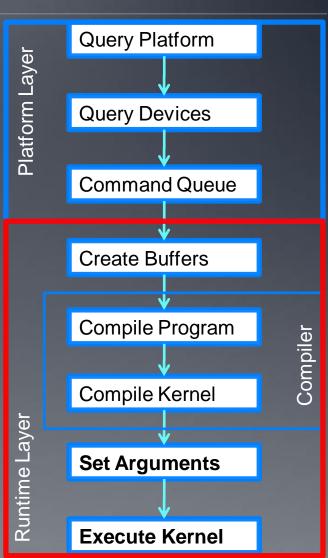
Step3: Set Arguments, Enqueue Kernel

//Set local and global workgroup sizes

```
size_t localws[2] = {16,16};
size_t globalws[2] = {W, H};//Assume divisible by 16
```

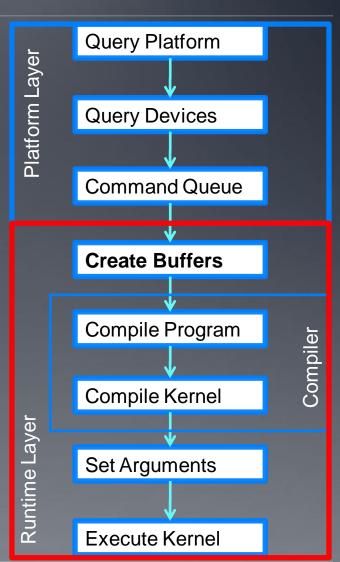
// execute kernel

```
clEnqueueNDRangeKernel(
myqueue, myKernel,
2, 0, globalws, localws,
0, NULL, NULL);
```



Step4: Read Back Result

- Only necessary for data required on the host
- Data output from one kernel can be reused for another kernel
 - Avoid redundant host-device IO



OpenCL Timing

- OpenCL provides "events" which can be used for timing kernels
 - Events will be discussed in detail in Lecture 11
- We pass an event to the OpenCL enqueue kernel function to capture timestamps
- Code snippet provided can be used to time a kernel
 - Add profiling enable flag to create command queue
 - By taking differences of the start and end timestamps we discount overheads like time spent in the command queue

```
cl_event event_timer;
clEnqueueNDRangeKernel(
myqueue, myKernel,
2, 0, globalws, localws,
0, NULL, &event_timer);
```

unsigned long starttime, endtime;

```
unsigned long elapsed =
(unsigned long)(endtime - starttime);
```

Example 2 Matrix Multiplication

Basic Matrix Multiplication

- Non-blocking matrix multiplication
 - Doesn't use local memory
 - Each element of matrix reads its own data independently
- Serial matrix multiplication

```
for(int i = 0; i < Ha; i++)
  for(int j = 0; j < Wb; j++) {
      c[i][j] = 0;
      for(int k = 0; k < Wa; k++)
           c[i][j] += a[i][k] + b[k][j]
}</pre>
```

- Reuse code from image rotation
 - Create context, command queues and compile program
 - Only need one more input memory object for 2nd matrix

Simple Matrix Multiplication

```
Wb
kernel void simpleMultiply(
          global float* c, int Wa, int Wb,
          global float* a, __global float* b) {
                                                                       col
  //Get global position in Y direction
  int row = get_global_id(1);
  //Get global position in X direction
  int col = get_global_id(0);
                                                        row
  float sum = 0.0f;
  //Calculate result of one element
  for (int i = 0; i < Wa; i++) {
        sum +=
           a[row*Wa+i] * b[i*Wb+col];
  c[row*Wb+col] = sum;
                                                        Wa
                                                                                Wh
```

Summary

- We have studied the use of OpenCL buffer objects
- A complete program in OpenCL has been written
- We have understood how an OpenCL work-item can be used to work on a single output element (seen with rotation and matrix multiplication)
 - While the previously discussed examples are correct data parallel programs their performance can be drastically improved
- Next Lecture
 - Study the GPU memory subsystem to understand how data must be managed to obtain performance for data parallel programs
 - Understand possible optimizations for programs running on data parallel hardware like GPUs