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1
      A Simple Sudoku Solver
       27th September, 2007
 3
      In Chapter 05
 5
      Explanations are added in Spring 2019 for CIS 623
 7
      A. The Sudoku game
 8
 9
         Given M, a 9 X 9 matrix of integers, where each cell
10
       is either empty (a blank) or a number k where 0 < k < 10,
11
       fill in the empty cells with the digits 1 to 9 so that
12
       each row, column and 3×3 box contains the numbers 1 to 9.
13
14
       B. Specification
15
16
           O. Introduce the basic data types
17
18
      > type Matrix a = [Row a]
19
      > type Row a
                             = [a]
20
21
      > type Grid
                             = Matrix Digit
22
      > type Digit
                             = Char
23
24
      > digits :: [Digit]
      > digits = ['1'..'9']
25
26
27
      > blank
                     :: Digit -> Bool
28
                     = (== '0')
      > blank
29
30
         Remarks:
31
32
         i. a sudoku `gameboard' is of type [[Char]].
33
              For example: egl is the sudoku grid given
34
              in Bird, Figure 5.1.
35
36
      > eg1 :: [[Char]]
37
      > eg1 = [
         ['0','0','4','0','0','5','7','0','0'],
['0','0','0','0','0','9','4','0','0'],
['3','6','0','0','0','0','0','0','8'],
38
39
40
          ['7','2','0','0','6','0','0','0','0','0'],
['0','0','0','4','0','2','0','0','0'],
['0','0','0','0','8','0','0','9','3'],
41
42
43
          ['4','0','0','0','0','0','0','5','6'],
44
          ['0','0','5','3','0','0','0','0','0','0'],
['0','0','6','1','0','0','9','0','0']]
45
46
47
48
      Another example, eg2 (solvable), is added
49
      below:
50
51
      > eg2 :: [[Char]]
52
      > eg2 = [
         ['5','3','0','0','7','0','0','0','0'],
['6','0','0','1','9','5','0','0','0'],
['0','9','8','0','0','0','0','0','6','0'],
['8','0','0','0','6','0','0','0','0','1'],
['4','0','0','8','0','3','0','0','1'],
['7','0','0','0','2','0','0','0','6'],
53
54
55
56
      >
57
58
      >
      > ['0','6','0','0','0','0','2','8','0'],
> ['0','0','0','4','1','9','0','0','5'],
> ['0','0','0','0','8','0','0','7','9']]
59
60
61
62
63
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64

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65
       1. Specify a solution (brute force)
 66
67
          We begin with a simple brute force solution
 68
           solve, using two subsidiary functions:
 69
 70
              solve :: Grid -> [Grid]
 71
              solve = filter valid . completions
 72
 73
          where the function completions will take the
74
          given grid as input, complete it by filling in
75
          every possible choice for the blank entries,
76
          and, the function valid will test, for each of
          the filled grids, if it contains duplicates
77
78
          in any row, box or column. By filtering out
79
          the invalid entries, we will obtain a solution
80
          to the given sudoku puzzle.
81
          We can define completions by a two way process:
82
83
84
     > solve1 :: Grid -> [Grid]
85
     > solve1 = filter valid . expand . choices
86
87
          where choices install the available digits for
88
          each cell. For example, choices eg1 will give
89
90
           *Main> choices eq1
91
           ["123456789","123456789","4","123456789","123456789",
92
           "5","7","123456789","123456789"],
93
           ["123456789", "123456789", "123456789", "123456789",
94
           "123456789", "9", "4", "123456789", "123456789"],
95
96
           ... 7 more lists (omitted)
97
           ]
98
99
          i. Step 1: implementation of the function choices
100
101
          The definition of the function choices is given
102
          below:
103
104
     > type Choices = [Digit]
105
106
     > choices :: Grid -> Matrix Choices
     > choices = map (map choice)
107
108
     > where choice d | blank d = digits
109
                        | otherwise = [d]
110
111
          For the function expand, we make use of cartesian
          product. Recall that, given two sets A and B, the
112
113
          cartesian product of A and B is:
114
115
          A \times B = \{(a,b): a \text{ is in set } A, b \text{ is in set } B\}
116
          For example, if we want to list all possible
117
          ways to fill the empty cells by 1,2 or 3
118
119
          in the matrix eg3 shown below:
120
121
122
           | 1 | | |
123
124
           | | 2 | |
                              <--- matrix eg3
125
126
           | 2 | | 3 |
127
128
```

```
129 > eg3 :: [[Char]]
130
     > eg3 = [
     > ['1','0','0'],
> ['0','2','0'],
> ['2','0','3']
131
132
133
134
135
     > digit123 :: [Digit]
> digit123 = ['1'..'3']
136
137
138
139
      > newchoices = map (map newchoice)
140
      > newchoice d = if blank d then digit123 else [d]
141
142
           we can compute (cp . map cp) (newchoices eg3)
143
           to list all possible ways (3^5 = 243 \text{ ways}) to
144
           fill in the matrix.
145
146
           Note: One may like to verify that as follows:
147
148
           *main> length ((cp . map cp) (newchoices eg3))
149
           243
150
151
           The function cp (cartesian product) and the
152
           function expand (uses cp) are implemented as
153
           follows:
154
155
      > expand :: Matrix Choices -> [Grid]
156
     > expand = cp . map cp
157
158
     > cp :: [[a]] -> [[a]]
                  = [[]]
159
     > cp []
     > cp (xs:xss) = [x:ys | x \leftarrow xs, ys \leftarrow cp xss]
160
161
162
           i. Step 2: implementation of the function valid
163
164
              Let g be a sudoku grid where the entry at each
              cell is either 1,2, ...,9. Then g is a valid
165
166
              solution to a sudoku game if
167
168
              (1) there is no dups at each row of g, and
169
              (2) there is no dups at each col of g, and
170
              (3) there is no dups at each box of g
171
172
              Hence, we define the function valid as:
173
174
     > valid :: Grid -> Bool
175
     > valid g = all nodups (rows g) &&
176
                  all nodups (cols g) &&
177
                  all nodups (boxs g)
178
179
              We define the function nodups by recursion.
              The function x `notElem` xs checks if an
180
181
              element x is a member of xs.
182
183
                    :: Eq a => [a] -> Bool
     > nodups
                     = True
184
      > nodups []
185
      > nodups (x:xs) = x `notElem` xs && nodups xs
186
187
             The function rows and cols are defined as
188
             follows (Note the use of zipWith when
189
             defining cols):
190
191
     > rows :: Matrix a -> [Row a]
192
     > rows = id
```

```
193
194
     > cols
                     :: Matrix a -> [Row a]
195
                    = [[x] | x <- xs]
     > cols [xs]
196
     > cols (xs:xss) = zipWith (:) xs (cols xss)
197
198
            The definition of boxs is built
199
             from a sequence of operations.
200
201
     > boxs :: Matrix a -> [Row a]
202
     > boxs = map ungroup . ungroup . map cols .
203
               group . map group
204
205
            where
206
207
     > ungroup
                         = concat
208
     > group []
                         = []
209
     > group (x:y:z:xs) = [x,y,z]:group xs
210
211
             It is instructive to display the intermediate
212
             results for boxs eg1 to show how the sequence
213
            of functions works.
214
215
             *Main> map group egl
216
             *Main> (group . map group) eg1
217
             *Main> (map cols . group . map group) eg1
218
             *Main> (ungroup . map cols . group . map group) eg1
219
             *Main>
220
             (map ungroup . ungroup . map cols . group . map group) eg1
221
222
            the results will be shown in a separate document.
223
224
            To reinforce your understanding of the development
225
            of the solution, prove (see 5.2):
226
227
             rows . rows = id -- trivial
228
             cols . cols = id -- exercise
229
             boxs . boxs = id -- completed proof given in 5.2
230
231
232
233
       2. Improve the solution by Pruning
234
235
          Assuming about 20 of the 81 entries in a sudoku grid
236
          are fixed initially, then the number of choices generated
237
          by the function (expand . choices) is 9^{61}. As we need
238
          to avoid checking all of them, we use a pruning strategy:
239
240
          strategy:
241
242
           remove any choices from a cell c that already occur as
243
          singleton entries in the row, column and box containing c.
244
245
          We therefore seek a function
246
247
          prune :: Matrix [Digit] -> Matrix [Digit]
248
249
          so that
250
251
          filter valid . expand = filter valid . expand . prune
252
253
254
```

255 256

```
257
          i. The design of the prune function
258
259
              We begin by considering a special case where we
260
              prune by row. That is, we remove any choices from
261
              a cell c that already occur as singleton entries
262
              in the row containing c. This is acheived by the
263
              function pruneRow (with the help of the subsidiary
264
              function remove).
265
266
     > pruneRow :: Row Choices -> Row Choices
267
     > pruneRow row = map (remove ones) row
     > where ones = [d | [d] <- row]</pre>
268
269
270
     > remove :: Choices -> Choices
271
     > remove xs [d] = [d]
272
     > remove xs ds = filter (`notElem` xs) ds
273
274
              Observe that the cols function (resp. the boxs function)
275
              rearranges the columns (resp. boxs) of the input g
276
              (sudoku grid) to rows of the output sudoku grid. Hence,
277
              the cases where we prune by columns (resp. by boxs) can
              be carried out in a similar manner. To be precise,
278
              the action `prune by rows', `prune by columns' and `prune
279
280
              by boxs' can be carried out by using the pruneBy function
281
              specified as
282
283
              pruneBy f = f . map pruneRow . f
284
285
              Hence, our prune function is:
286
     > prune :: Matrix Choices -> Matrix Choices
287
288
     > prune =
289
       pruneBy boxs . pruneBy cols . pruneBy rows
290
     > where pruneBy f = f . map pruneRow . f
291
292
              Remark:
293
294
              Now, Our solver can be expressed as:
295
296
              solve = ( search ) . choices
297
298
              We will incorporate our prune function
299
              as part of our search algorithm.
300
301
          ii. The verification of the prune function
302
303
              We can prove, by equation reasoning (section 5.3),
304
              that
305
306
           filter valid . expand = filter valid . expand . prune
307
308
          iii. An alternate approach
309
310
              one way to implement a search algorithm
              is to apply the prune function until
311
312
              one singleton choices are left. We can
              use the function many to implement this
313
314
              idea:
315
316
     > many :: (Eq a) => (a -> a) -> a -> a
317
     > many f x = if x==y then x else many f y
318
           where y = f x
319
320
     > solve3 = filter valid . expand . many prune . choices
```

322 Remark: 323 324 The result of many prune . choices is a matrix 325 of choices that can be put into one of three classes: 326 327 1. A complete matrix in which every entry is a 328 singleton choice. In this case expand will extract 329 a single grid that can be checked for validity. 330 2. A matrix that contains the empty choice somewhere. 331 In this case expand will produce the empty list. 332 3. A matrix that does not contain the empty choice but 333 does contain some entry with two or more choices. 334 335 In case 3, rather than carry out full expansion, a more 336 sensible idea is to make use of a partial expansion that installs the choices for just one of the entries, and to 337 338 start the pruning process again on each result. 339 340 3. Single-cell expansion 341 The following function expand1 implement a single-cell 342 343 expansion. 344 345 > expand1 :: Matrix Choices -> [Matrix Choices] 346 > expand1 rows = 347 > [rows1 ++ [row1 ++ [c]:row2] ++ rows2 | c <- cs]</pre> 348 349 (rows1,row:rows2) = break (any smallest) rows 350 (row1,cs:row2) = break smallest row 351 smallest cs = length cs == n 352 = minimum (counts rows) 353 354 > counts = filter (/=1) . map length . concat 355 356 To understand how it works, consider using the 357 examples eg1 and eg2 and examine: 358 359 expand1 [eg1] 360 expand2 [eg2] 361 362 363 We leave the details experiments to the reader. 364 365 discussions: how does the function resembles 366 `expanding a node in a search tree' ? 367 368 4. Final algorithm 369 370 The final algorithm uses the following search 371 algorithm: 372 373 > solve2 :: Grid -> [Grid] 374 > solve2 = search . choices 375 376 > search :: Matrix Choices -> [Grid] 377 > search cm -- cm: the list of possible choices 378 [not (safe pm) = [] -- return null if dups detected 379 |complete pm = [map (map head) pm] -- search's done 380 > |otherwise = (concat . map search . expand1) pm 381 > -- expand the search tree one step 382 > -- note the use of 'prune' at each 383 > -- step 384 > where pm = prune cm -- pm: results after pruning cm

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```
385
386
387
      > complete :: Matrix Choices -> Bool
      > complete = all (all single) -- test if `search' completes
388
389
390
      > single [_] = True
391
      > single _ = False
392
393
394
      > safe :: Matrix Choices -> Bool
      > -- test if each grid in the input list

> -- is `ok' (no duplicates)

> safe cm = all ok (rows cm) &&

> all ok (cols cm) &&
395
396
397
398
                    all ok (boxs cm)
399
      >
400
401
      > ok row = nodups [d | [d] <- row]</pre>
402
403
404
      5. Exercises
405
406
      Attempt Exercise A to H and study the given solution the text
407
      provides.
```