```
What is a Monad ?
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 3
     1. Background
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 5
     (see Hutton, https://www.youtube.com/watch?v=t1e8qqXLbsU)
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7
     on Maybe monads
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9
10
11
     2. IO Monad
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13
     a. Recall that, for any type a, we have a corresponding type
14
15
         IO a
16
17
         where their members are actions (or programs) that yield a
18
         result of type a . Note that the notation () denotes the
19
         type with the single value () (the 0-ary tuple).
20
21
         Hence, many IO actions are represented as functions of
22
         the type:
23
24
         t -> IO () (t is a type such as String, Char etc.)
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26
        Examples include: putChar, putStr, putStrLn, print, return
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28
         IO actions are members in IO a
29
30
         Examples include: getChar, getLine etc.
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33
    b. One way to think of the type IO a is
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35
         type IO a = World -> (a, World)
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37
         Hence, when we compose multiple actions, it is not obvious
38
         that associative laws hold. We will revisit this point later.
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    c. IO monad
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         IO and Maybe examples of monads. In essence, Monad is a type
43
         class where a list of operations must be provided and the
44
         operations must satisfy a number of algebraic laws (called the
45
        Monad laws, see bird 243).
46
47
         Here is the definition of Monad
48
         [from: https://wiki.haskell.org/Monad]
49
50
    class Monad m where
51
52
       (>>=) :: m a -> ( a -> m b) -> m b
       (>>) :: m a -> m b
53
                                     -> m b
54
                                     -> m a
       return :: a
55
       fail :: String -> m a
56
57
     and all instances of Monad should obey the Monad Laws:
58
59
    return a >>= k
                                     = ka
60
    m >>= return
                                     = m
61
             >>= (\x -> k x >>= h) = (m >>= k) >>= h
62
63
64
     It is equivalent to say that the following operation (>=>)
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66
     (>=>) :: Monad m => (a -> m b) -> (b -> m c) -> a -> m c
```

```
(m >=> n) x = do
                         y <- m x
68
69
                         n y
70
71
     satisfy
72
73
     1. (>=>) has a left identity (which is, the return function)
74
     2. (>=>) has a right identity (which is, the return function)
75
     3. (>=>) is an associative operations
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77
     Exercise:
78
79
     Verify 1, 2 and 3.
80
81
     (see <a href="https://wiki.haskell.org/Monad laws">https://wiki.haskell.org/Monad laws</a>)
82
83
    Note:
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85
     1. Monad laws are also discussed in Bird Chapter 10,
86
     section 2. Note that the term Monoid is not the same
87
     as Monads. Also, Parser is a Monad defined in Chapter 11
88
     and the source code Parsing.lhs (given by Bird).
89
90
91
     We will not go through some {\tt of} the materials stated
92
    therein (e.g. leapfrog rule) this semester.
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94
95
```