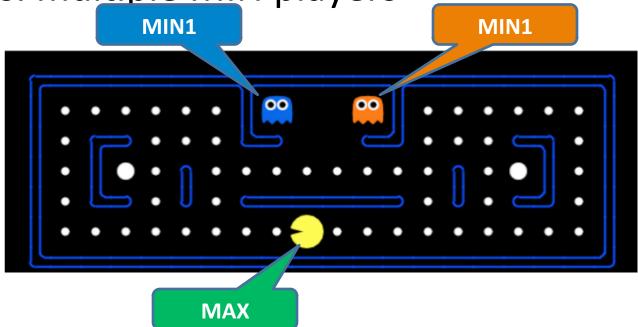
# Adversarial Search: More on Eval Functions, ExpectiMax

With slides from Dan Klein, Percy Liang, Luke Zettlemoyer

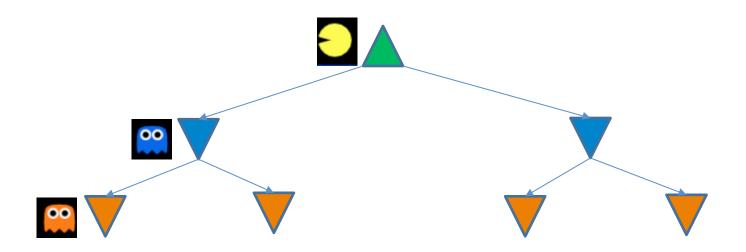
## Multiplayer Games: 1 vs. All

MAX vs. multiple MIN players

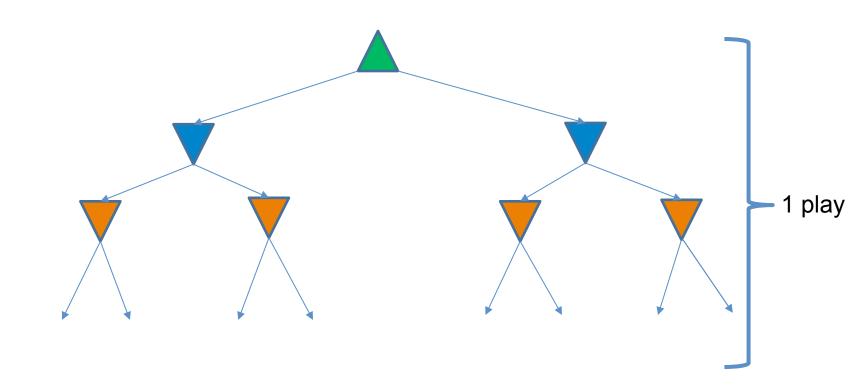


How does the algorithm change?

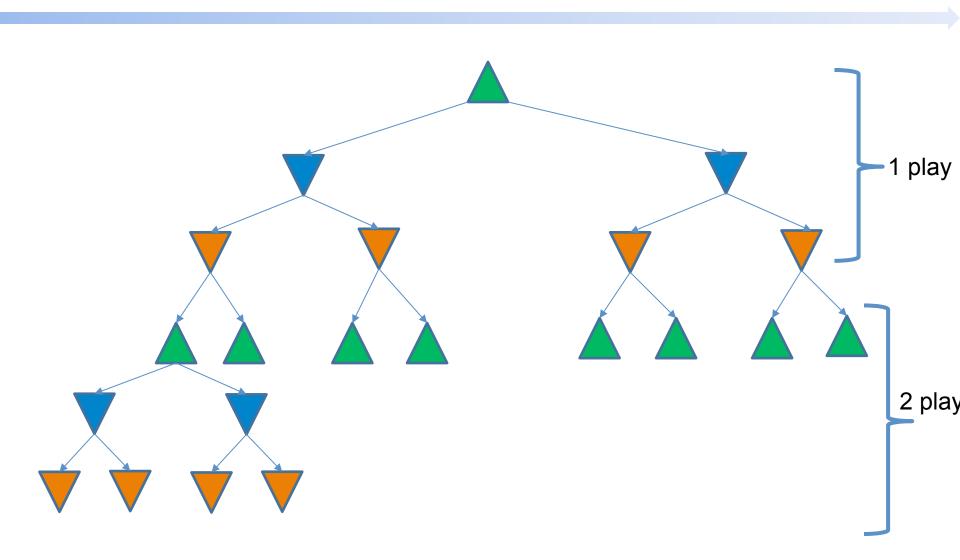
## Multiplayer Game Tree: 1 Max, 2 MINs

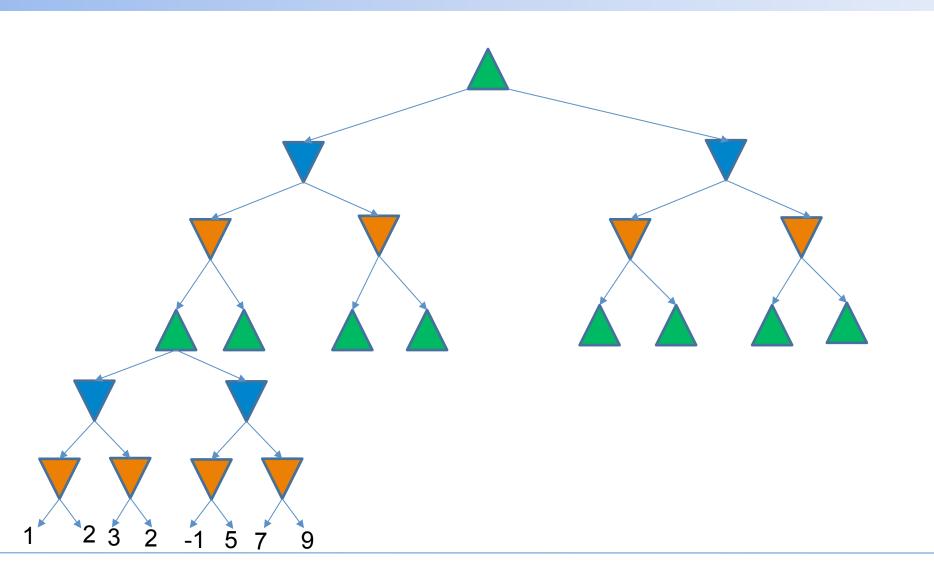


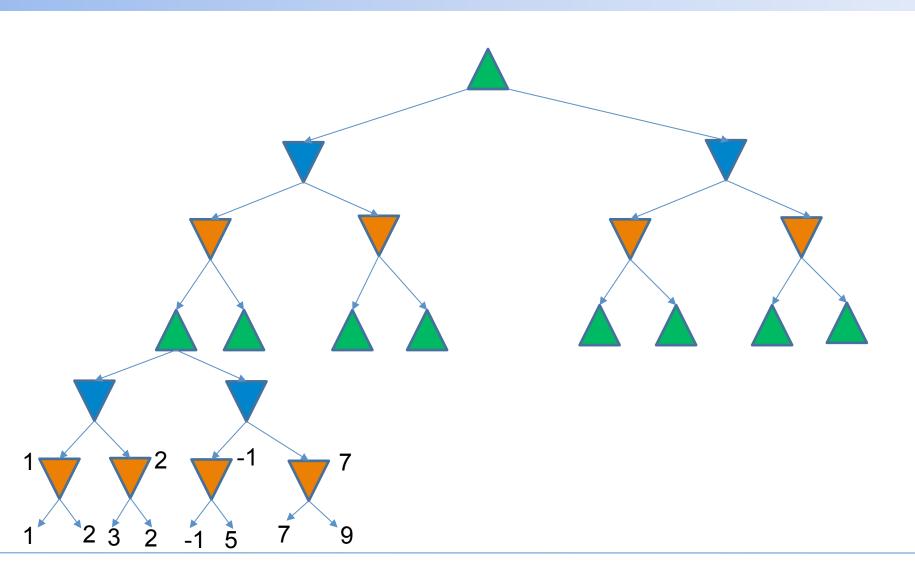
## Multiplayer Game Tree: depth 1

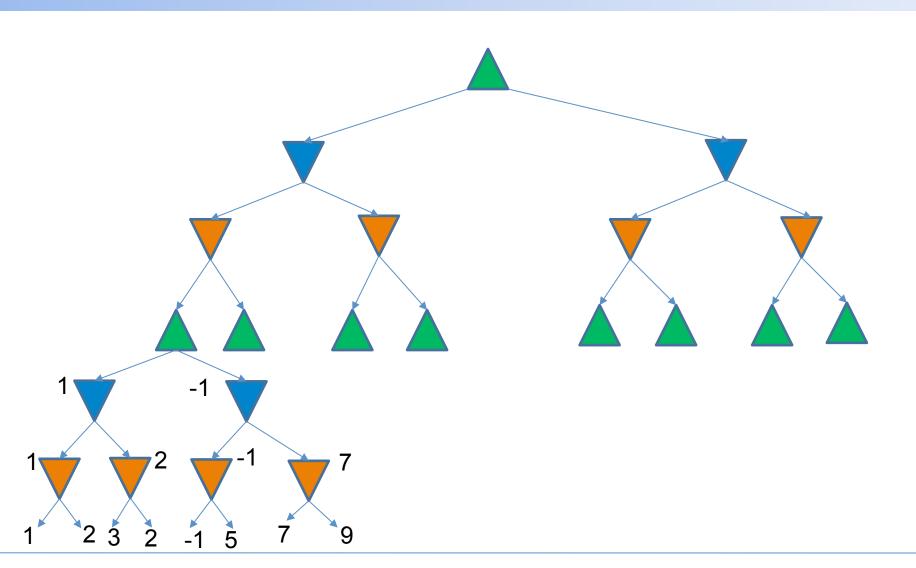


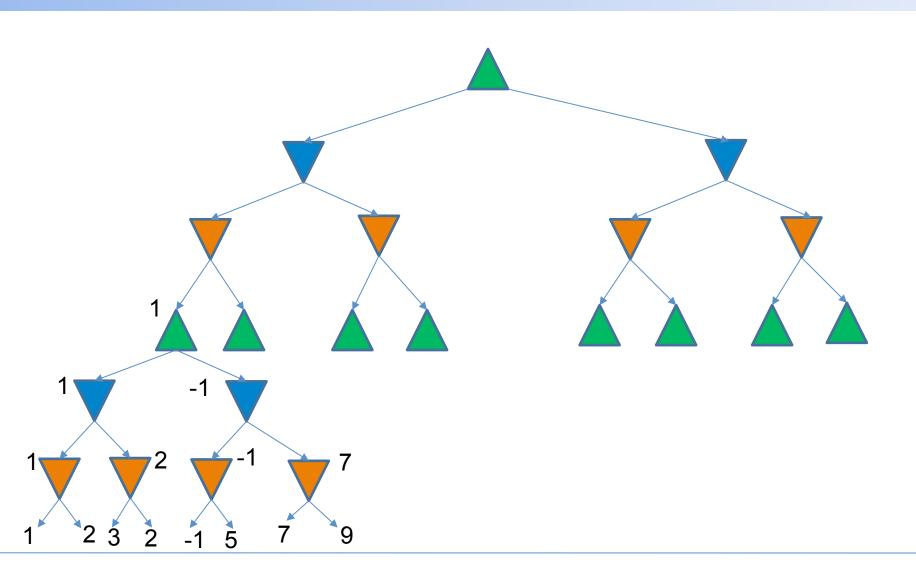
## Multiplayer Game Tree: depth 2

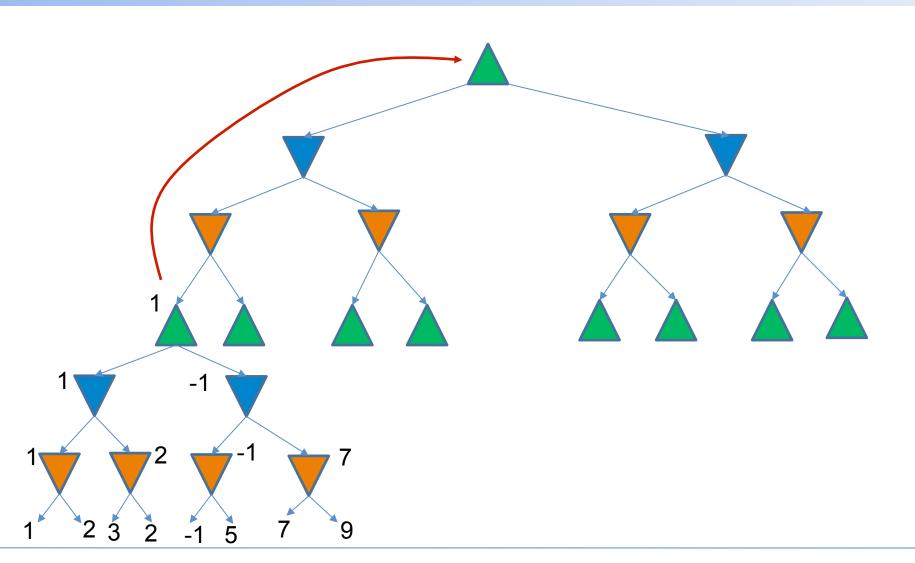












## MiniMax Algorithm: Multi-Min version

```
function MINIMAX-DECISION(state) returns an action
inputs: state, current state in game
v ← MAX-VALUE(state)
return the action in SUCCESSORS(state) with value v

function MAX-VALUE(state) returns a utility value
if TERMINAL-TEST(state) then return UTILITY(state)
v ← ∞
for a,s in SUCCESSORS(state) do
v ← MAX(v, MIN-VALUE(s, 1))
return v
```

```
Note:
MAX=Agent1
MIN1=Agent2
MIN2=Agent3
....
MINK=AgentN
```

```
function MIN-VALUE(state, agentIndex) returns a utility value if TERMINAL-TEST(state) then return UTILITY(state)

v \leftarrow \infty

for a, s in SUCCESSORS(agentIndex, state) do

if (agentIndex >= numAgents) then

v_{temp} = \text{MAX-VALUE}(s)

else #another ghost plays

v_{temp} = \text{MIN-VALUE}(s, agentIndex+1))

v \leftarrow \text{MIN}(v, v_{temp})
```

## problem: collusion possible

- Previous slide (standard minimax analysis) assumes that each player operates to maximize only their own utility
- In practice, players could make alliances
  - Ex: C strong, A and B both weak
  - May be best for A and B to attack C rather than each other
- If game is not zero-sum (i.e., utility(A) = utility(B) then alliances can be useful even with 2 players
  - e.g., both cooperate to maximum the sum of the utilities
- Ignore this, assume non-cooperative games

## Minimax Algorithm: Analysis

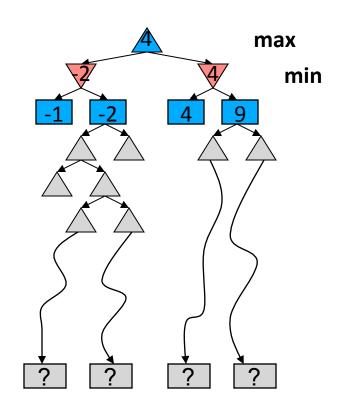
Complete depth-first exploration of the game tree

- Performance:
  - Max depth = d, b legal moves at each point
  - E.g., Chess: d ~ 100, b ~35

Criterion	Minimax
Time	O( <b>b</b> <sup>m</sup> )
Space	O(bm)

### **Resource Limits**

- Problem: In realistic games, cannot search to leaves!
- Solution: ? Hint: 2 complementary approaches?
   1. ?
   2. ?
- Example:
  - Suppose we have 100 seconds, can explore 10K nodes / sec
  - So can check 1M nodes per move
  - $\alpha$ - $\beta$  reaches about depth 8 decent chess program
- Guarantee of optimal play is gone
- More plies makes a BIG difference
- Use iterative deepening for an anytime algorithm



# Alpha-Beta Pseudocode

```
inputs: state, current game state
α, value of best alternative for MAX on path to state
β, value of best alternative for MIN on path to state
returns: a utility value
```

```
function MAX-VALUE(state, \alpha, \beta)

if TERMINAL-TEST(state) then

return UTILITY(state)

v \leftarrow -\infty

for a, s in SUCCESSORS(state) do

v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(s, \alpha, \beta))

if v \ge \beta then return v

\alpha \leftarrow \text{MAX}(\alpha, v)

return v
```

```
function MIN-VALUE(state, \alpha, \beta)

if TERMINAL-TEST(state) then

return UTILITY(state)

v \leftarrow +\infty

for a, s in Successors(state) do

v \leftarrow \text{MIN}(v, \text{Max-Value}(s, \alpha, \beta))

if v \leq \alpha then return v

\beta \leftarrow \text{MIN}(\beta, v)

return v
```

# Alpha-Beta Pruning Properties

- This pruning has no effect on final result at the root
- Values of intermediate nodes might be wrong!
  - but, they are bounds
- Good child ordering improves effectiveness of pruning
- With "perfect ordering":
  - Time complexity drops to O(b<sup>m/2</sup>)
  - Doubles solvable depth!
  - Full search of, e.g. chess, is still hopeless...

## AlphaBeta Experiment

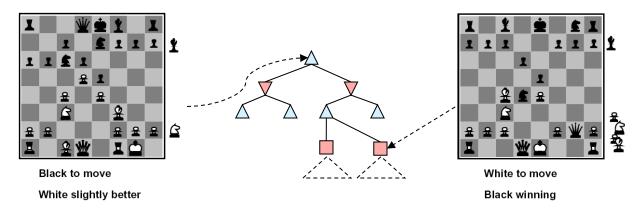
- time python pacman.py -p MinimaxAgent -l minimaxClassic -a depth=4 -q
- time python pacman.py -p AlphaBetaAgent -l minimaxClassic -a depth=4 -q

## More Alpha-Beta Pruning Examples

- Generic game tree visualization: http://homepage.ufp.pt/jtorres/ensino/ia/alfabeta.html
- Example Play: Connect Four: https://gimu.org/connect-four-js/

## (Better) Evaluation Functions

Evaluation functions score non-terminals in depth-limited search



- Ideal function: returns the actual minimax value of the position
- In practice: typically weighted linear sum of features:

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \ldots + w_n f_n(s)$$

• e.g.  $f_1(s)$  = (num white queens – num black queens), etc.

## Construction of an Evaluation Function

Usually a weighted sum of "features":

$$e(s) = \sum_{i=1}^{n} w_i f_i(s)$$

- Features may include
  - Number of pieces of each type
  - Number of possible moves
  - Number of squares controlled

## More Examples of Evaluation Functions

#### Reversi

- –Number squares held?
- —Better: number of squares held that cannot be flipped
- –Prefer valuable squares
  - NxN array w[i,j] of position values
  - Highest value: corners, edges
  - Lowest value: next to corner or edge

•s[i,j] = +1 player, 0 empty, -1 opponent 
$$score = \sum_{i,j} w[i,j]s[i,j]$$

## **Eval Function Approximation**



# Key idea: parameterized evaluation functions $\mathsf{Eval}(s;\mathbf{w}) \text{ depends on weights } \mathbf{w} \in \mathbb{R}^d$

Feature vector:  $\phi(s) \in \mathbb{R}^d$ 

$$\phi_1(s) = K - K'$$

$$\phi_2(s) = Q - Q'$$

Linear evaluation function:

$$\mathsf{Eval}(s; \mathbf{w}) = \mathbf{w} \cdot \phi(s)$$

## Chess Example: Revisited





#### Example: chess-

 $\begin{aligned} \mathsf{Eval}(s) &= \mathsf{material} + \mathsf{mobility} + \mathsf{king\text{-}safety} + \mathsf{center\text{-}control} \\ \mathsf{material} &= 10^{100}(K - K') + 9(Q - Q') + 5(R - R') + \\ &\quad 3(B - B' + N - N') + 1(P - P') \\ \mathsf{mobility} &= 0.1 (\mathsf{num\text{-}legal\text{-}moves} - \mathsf{num\text{-}legal\text{-}moves}') \\ \ldots \end{aligned}$ 

#### More Evaluation Functions

#### Chess:

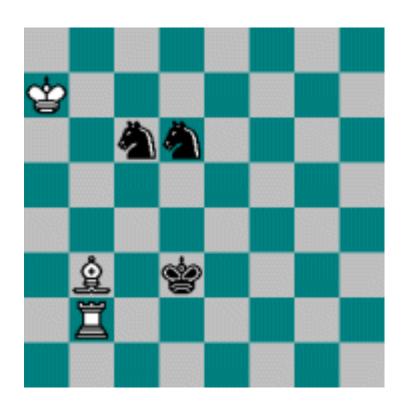
```
-eval(s) =
    w1 * material(s) +
    w2 * mobility(s) +
    w3 * king safety(s) +
    w4 * center control(s) + ...
```

- In practice MiniMax improves accuracy of heuristic eval function
- But one can construct pathological games where more search hurts performance!(Nau 1981)

## Idea: End-Game Databases

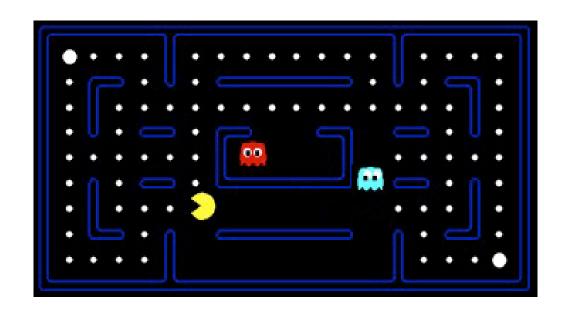
- Ken Thompson all 5 piece end-games
- Lewis Stiller all 6 piece end-games
  - Refuted common chess wisdom: many positions thought to be ties were really forced wins -- 90% for white
  - Is perfect chess a win for white?

## The MONSTER



White wins in 255 moves (Stiller, 1991)

## **Evaluation for Pacman**



What features would be good for Pacman?

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$$

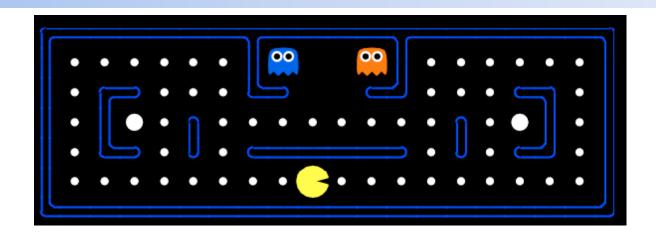
### Pacman Eval Functions

Use a linear combination of features.

```
Eval(state) = f1*w1 + f2*w2
```

- Example:
  - f1: distance to nearest food w1: ?
  - f2: distance to nearest ghost w2:?

#### Pacman Eval Functions Demo



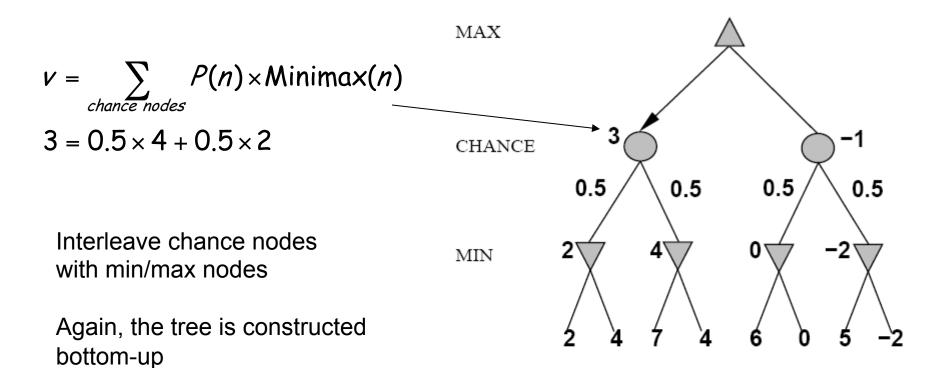
Alpha-beta pruning, simple eval function:

```
python pacman.py -p AlphaBetaAgent -a depth=4 -l minimaxClassic -q python pacman.py -p AlphaBetaAgent -a depth=4 -l smallClassic -q
```

Alpha-beta, better eval function:

```
python pacman.py -p AlphaBetaAgent -a depth=4,evalFn=better -l minimaxClassic -q python pacman.py -p AlphaBetaAgent -a depth=4,evalFn=better -l smallClassic -q
```

## Expected Minimax (ExpectiMax)



## Idea: Maximum Expected Utility

Why should we average utilities? Why not minimax?

- Principle of maximum expected utility: an agent should chose the action which maximizes its expected utility, given its knowledge
  - General principle for decision making
  - Often taken as the definition of rationality
  - We'll see this idea over and over in this course!
- Let's decompress this definition...

## Reminder: Probabilities (More Later)

- A random variable represents an event whose outcome is unknown
- A probability distribution is an assignment of weights to outcomes
- Example: Traffic on freeway
  - Random variable: T = whether there's traffic
  - Outcomes: T in {none, light, heavy}
  - Distribution: P(T=none) = 0.25, P(T=light) = 0.50, P(T=heavy) = 0.25
- Some laws of probability (more later):
  - Probabilities are always non-negative
  - Probabilities over all possible outcomes sum to one
- As we get more evidence, probabilities may change:
  - P(T=heavy) = 0.25, P(T=heavy | Hour=8am) = 0.60
  - We'll talk about methods for reasoning and updating probabilities later



0.25



0.50



0.25

## Reminder: Expectations

 The expected value of a function of a random variable is th average, weighted by the probability distribution over outcomes



Example: How long to get to the airport?

Time: 20 min x

Probability: 0.25

30 min

x 0.50 60 min

0.25



35 min





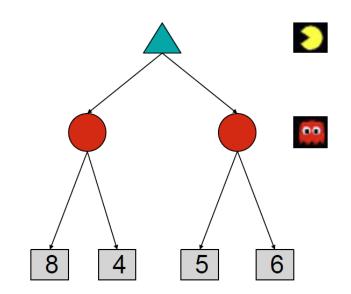


## ExpectiMax: Revised Pseudocode

```
def value(s)
  if s is a max node return maxValue(s)
  if s is an exp node return expValue(s)
  if s is a terminal node return evaluation(s)

def maxValue(s)
  values = [value(s') for s' in successors(s)]
```

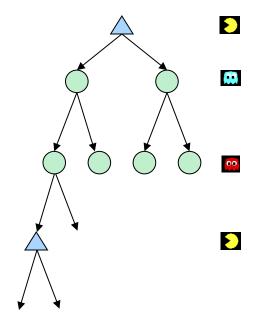
return max(values)



```
def expValue(s)
  values = [value(s') for s' in successors(s)]
  weights = [probability(s, s') for s' in successors(s)]
  return expectation(values, weights)
```

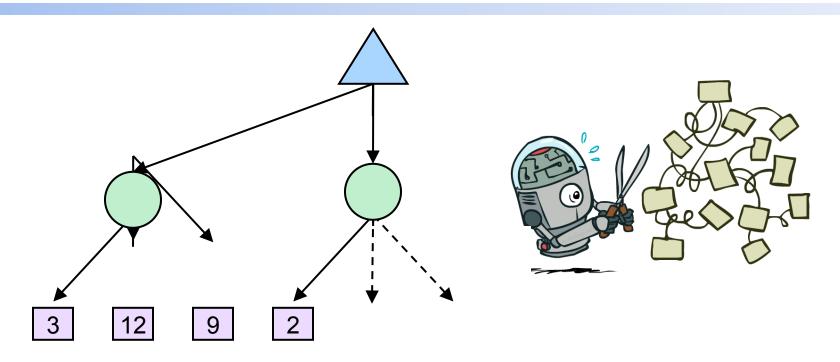
### What Probabilities to Use?

- In expectimax search, we have a probabilistic model of how the opponent (or environment) will behave in any state
  - Model could be a simple uniform distribution (roll a die)
  - Model could be sophisticated and require a great deal of computation
  - We have a chance node for any outcome out of our control: opponent or environment
  - The model might say that adversarial actions are likely!
- For now, assume each chance node magically comes along with probabilities that specify the distribution over its outcomes



Having a probabilistic belief about another agent's action does not mean that the agent is flipping any coins!

# **Expectimax Pruning?**

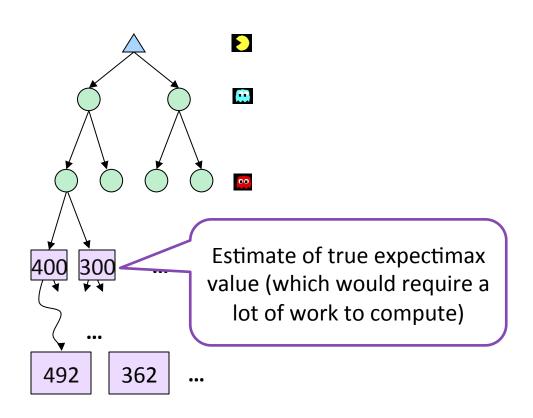


Theoretically possible, but:

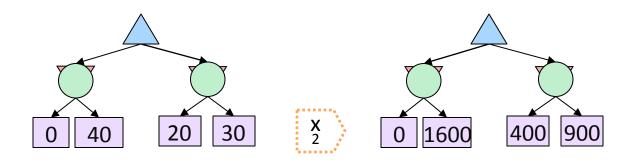
- Either accept that with some probability we will discard optimal solution
- OR: no guarantee that pruning will save much time

First algorithm (I know of) was published in 2009, makes a lot of assumptions: <a href="http://ieeexplore.ieee.org/document/5286476/">http://ieeexplore.ieee.org/document/5286476/</a>

## Depth-Limited Expectimax



## What Utilities to Use?



- For worst-case minimax reasoning, terminal function scale doesn't matter
  - We just want better states to have higher evaluations (get the ordering right)
  - We call this insensitivity to monotonic transformations
- For average-case expectimax reasoning, we need magnitudes to be meaningful

## **Expectimax for Pacman**

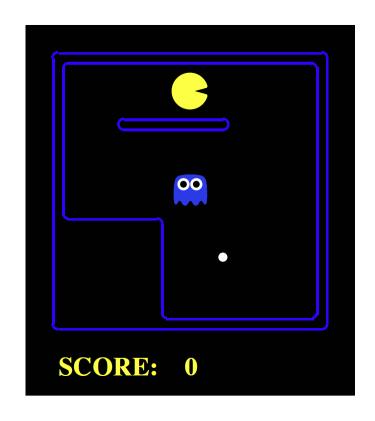
- Notice that we stepped away from thinking that the ghosts are trying to minimize pacman's score
- Instead, they are now a part of the environment
- Pacman has a belief (distribution) over how they will act

 Quiz: Can we model MiniMax as special case of ExpectiMax?

## **ExpectiMax for Pacman**

#### **Results from playing 5 games**

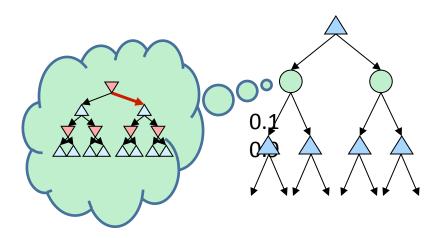
	Minimizing Ghost	Random Ghost
Minimax Pacman	Won 5/5 Avg. Score: 493	Won 5/5 Avg. Score: 483
Expectimax Pacman	Won 1/5 Avg. Score: -303	Won 5/5 Avg. Score: 503



Pacman does depth 4 search with an eval function that avoids trouble Minimizing ghost does depth 2 search with an eval function that seeks Pacman

## Quiz: Informed Probabilities

- Let's say you know that your opponent is actually running a depth 2 minimax, using the result 80% of the time, and moving randomly otherwise
- Question: What tree search should you use?



#### Answer: Expectimax!

- To figure out EACH chance node's probabilities, you have to run a simulation of your opponent
- This kind of thing gets very slow very quickly
- Even worse if you have to simulate your opponent simulating you...
- ... except for minimax, which has the nice property that it all collapses into one game tree

## Project 2: Assigned, Due Wed 2/15

- http://www.mathcs.emory.edu/~eugene/cs325/p2/
  - Use Piazza for discussions
  - Questions?