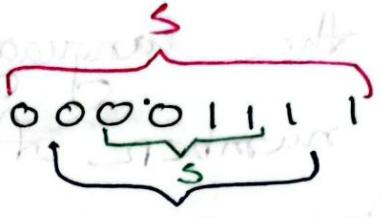


Context Free Grammar

CFG

→ We establish a recursive relationship

$O^m 1^m$. Here, @ number of O's and 1's same.



So, we can write it as,

$\begin{array}{l} S \\ \Rightarrow OS1 \\ \Rightarrow OOS11 \\ \Rightarrow OOOSS111 \\ \Rightarrow OOOO111 \end{array}$

At a moment that S will become ϵ ; $S \rightarrow \epsilon$

means, or

Thus the number of 0 and 1 remains same

CFG is a set of production rules,

- (i) $S \rightarrow OS1$
- (ii) $S \rightarrow \epsilon$

{Ex} # $O^n 1^{2n}$

Some strings can be $\underline{\underline{00}} \underline{\underline{1111}}$
 $\underline{\underline{000}} \underline{\underline{111111}}$

So, $S \rightarrow OS11$
 $S \rightarrow \epsilon$ or, $S \rightarrow OS11 | \epsilon$

$O^{3n} 1^{5n}$

$\begin{array}{c} 000 \quad 11111 \\ 000000 \quad 11111111 \end{array}$

$S \rightarrow OOOSS11111$
 $S \rightarrow \epsilon$

|

$S \rightarrow 000S11111$

Non-terminal Var Terminal Derivative Rule
 $\uparrow \quad \uparrow \quad \uparrow$

$CFG_2 \rightarrow 4 \text{ Tuple } (V, T, P, S)$ starting Var

~~# $S \rightarrow 0B1$~~

$$A \rightarrow 0A1 \mid B1$$

$$B \rightarrow 00$$

Derivation

$$\begin{aligned}
 & A \xrightarrow{\substack{A \\ \Rightarrow B1}} 0A1 \\
 & 0A1 \xrightarrow{\substack{A \\ \Rightarrow 001}} 00A11 \\
 & 00A11 \xrightarrow{\substack{A \\ \Rightarrow 000}} 000B111 \\
 & 000B111 \xrightarrow{\substack{B1 \\ \Rightarrow \#}} 0000111\#
 \end{aligned}$$

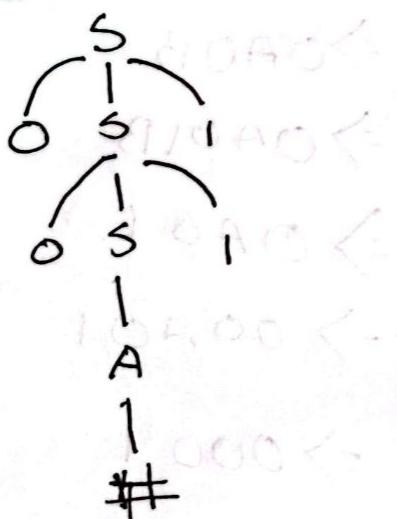
~~# $S \rightarrow OS1$~~

$$S \rightarrow A$$

$$A \rightarrow \#$$

Here, $V = \{A, A\}$
 $T = \{0, 1, \#\}$
 $\textcircled{S} \rightarrow \text{Starting}$

Parsed Tree



$$\begin{aligned}
 & S \xrightarrow{\substack{S \\ \Rightarrow OS1}} OS1 \\
 & OS1 \xrightarrow{\substack{S \\ \Rightarrow 00 S 11}} 00 S 11 \\
 & 00 S 11 \xrightarrow{\substack{S \\ \Rightarrow 00 A 11}} 00 A 11 ; \text{Recursion} \\
 & 00 A 11 \xrightarrow{\substack{A \\ \Rightarrow \#}} 00 \# 11
 \end{aligned}$$

Ans: $0^n \# 1^n$

Ans: $00\#11$

$0^n \# 1^n$

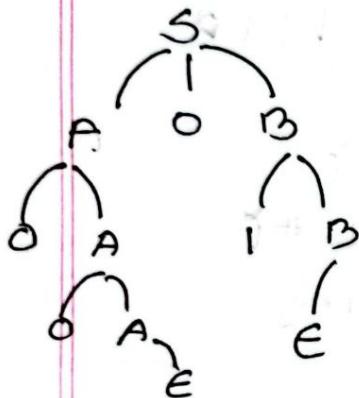
$S \rightarrow AOB$

$A \rightarrow OA|E$
 $B \rightarrow 1B|E$

Using Leftmost

$\begin{array}{l} S \\ \Rightarrow AOB \\ \Rightarrow OAOB \\ \Rightarrow OOA0B \\ \Rightarrow OOOB \\ \Rightarrow OOO1B \\ \Rightarrow OOO1 \end{array}$

Parsed Tree



$OOE01E$
 $\Rightarrow OOO1$

Without Using leftmost

$\begin{array}{l} S \\ \Rightarrow AOB \\ \Rightarrow OAOB \\ \Rightarrow OA01B \\ \Rightarrow OAO1 \\ \Rightarrow OOA01 \\ \Rightarrow OOO1 \end{array}$

Property of Grammar

Ambiguity

Ambiguous Grammar \rightarrow More than 1 parsed tree / left most / Right most

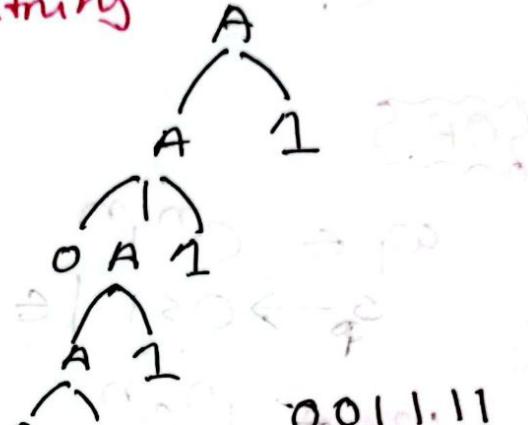
$A \rightarrow A1 | 0A1 | 01 ; "001111" \text{ derive this string}$

$\Rightarrow \overset{A}{A1} [A1]$

$\Rightarrow 0 \underset{\swarrow}{A} 1 1 [0A1]$

$\Rightarrow 0 A \underset{\swarrow}{1} 1 1 [A1]$

$\Rightarrow 001111 [01]$



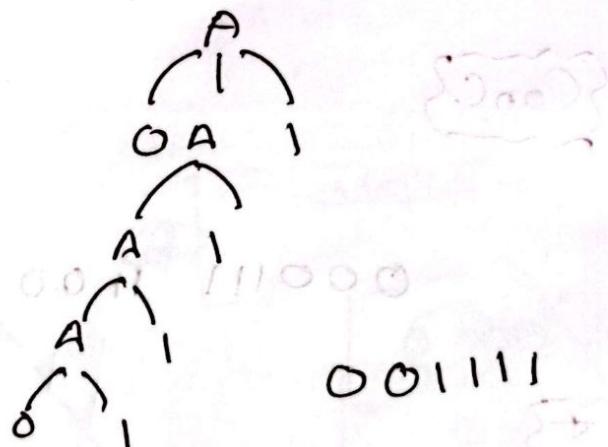
Another way

$\Rightarrow \overset{A}{0 A1} [0A1]$

$\Rightarrow 0 A \underset{\swarrow}{1} 1 [A1]$

$\Rightarrow 0 A 1 1 [A1]$

$\Rightarrow 001111 [01]$



Regular Operation

OR (+)

Concatenation (α)

Kleene Closure

{ORS}

$$\omega_1 \in 0^n 1^n$$

$$S_p \rightarrow 0 S_1 1 | \epsilon$$

$$\omega_2 \in 1^n 0^n$$

$$S_q \rightarrow 1 S_0 0 | \epsilon$$

{Conc.}

$$000111 \quad 1100$$

$$S_0, S_1 \rightarrow S_a \quad | \quad S_b$$

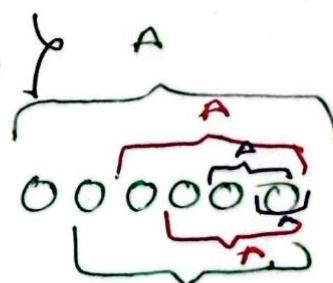
$$S_a \rightarrow 0 S_a 1 \quad | \quad \epsilon$$

$$S_b \rightarrow 1 S_b 0 \quad | \quad \epsilon$$

*3

$$0^* = \{ \epsilon, 0, 00, 000, 0000, \dots \}$$

We can write 0^* as,



$$S_0, A \rightarrow 0 A \quad | \quad \epsilon$$

$(01001)^* = \{ \epsilon, 01001, 0100101001, \dots \}$

$0100010100101001010010100101001$

$\underbrace{\hspace{10em}}$

B

B $\rightarrow 01001 B | \epsilon$

ଏହାର କାର୍ଯ୍ୟର ନିଯମ ଲାଗି,

B $\rightarrow C B | \epsilon$

C $\rightarrow 01001$

{RegEx \rightarrow CFG}

$(0|10)^*$ | $\frac{10}{B}^*$

S $\rightarrow A B$

B $\rightarrow C$

C $\rightarrow 0C \epsilon$

A \rightarrow

{କେବୁଳ ବାବଦ କିମ୍ବା

(0/10) * -> A

B = 10^*
C = 0^*
D = 10

$$\left\{ \begin{array}{l} 1^* = 1S \\ 0^* = 0S \end{array} \right.$$

{ RegEx \rightarrow CFGs }

$$\frac{(0|10)^*}{A} \mid \frac{10^*}{B}$$

$$S \rightarrow A \mid B$$

$$B \rightarrow 1C$$

$$C \rightarrow 0C \mid E$$

$$A \rightarrow DA \mid E$$

$$D \rightarrow (0|10)$$

3PF

$$A = (0|10)^*$$

$$B = 10^*$$

$$C = 0^* \quad \{ B=1C \}$$

$$D = 0 \mid 10$$

$$E = D^*$$

{ Some Examples }

$$\# 0^n 1^n 2^n$$

$$S \rightarrow 0S2 \mid 1$$

$$\# 0^n 1^n 2^n$$

$$A \rightarrow S2$$

$$S \rightarrow 0S1 \mid E$$

$$\# 0^n 1^{n+1}$$



$$\Rightarrow$$

$$0^n 1^n \} 2^n 2$$

$$0^n 1^n 1 \} \text{connect}$$

$0^m 1^n 2^K$ where $K = m+n$

Hence $0^m 1^n 2^K$

$0^m 1^n 2^{m+n}$

$0^m \underbrace{1^n 2^m}_{\times} 2^m$

$s \Rightarrow 0s2$

$\Rightarrow 00s22$

$\Rightarrow 000s222$

$\Rightarrow 0000s2222$

$\Rightarrow 00001s2222$

$\Rightarrow 00\underbrace{012222}_s$

so, $s \rightarrow 0s2 | \times$

$\times \rightarrow 1x2 | \epsilon$

$0^m 1^n 2^K$ where, $n = m+k$

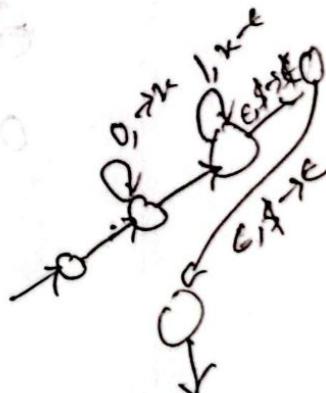
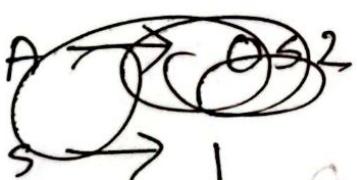
$0^m 1^{m+k} 2^K$

$0^m \underbrace{1^m}_{r} \underbrace{1^k 2^K}_{B}$

$s \rightarrow AB$

$A \rightarrow 0A1 | \epsilon$

$B \rightarrow 1B2 | \epsilon$



∴ $sA \leftarrow A$

∴ $sB \leftarrow B$

$0011 \rightarrow 0011$
 $000111 \rightarrow 000111$

$0^m 1^m \quad m > 2$

$s \rightarrow 00A11$

$A \rightarrow 0A1 | E$

$0^m 1^n \quad m > n$

$s \rightarrow 0A1 | A$
 $A \rightarrow 0A | 0$

$0^m 1^n 2^k 3^j \quad ; \quad k = m + n + j$

~~$0^m 1^n 2^m$~~

$0^m [1^n 2^m] 2^k 3^j$

~~$A \rightarrow 1A2$~~
 ~~$B \rightarrow 2B3$~~

$s \rightarrow 0A2^m | A$

$A \rightarrow 1A2 | G$

$B \rightarrow 2B3 | E$

$S[0^n1^n] ; S \rightarrow 0S1$
 $S[0^*] ; S \rightarrow 0S|\epsilon$
}
 }
 $S[(AB)^*] S \rightarrow \begin{array}{l} AB \\ A \rightarrow 011 \end{array}$
P(X)

of $w|w$ contains at least three 1's }

We can ans in 2 way.

$\epsilon^* | \epsilon^* | \epsilon^* | \epsilon^*$
 $S \rightarrow A | A | A | A$
 $A \rightarrow 0A | 1A | \epsilon$
}
 }
 $0^* | 0^* | 0^* | \epsilon^*$
 $S \rightarrow B | B | B | A$
 $A \rightarrow 0A | 1A | G$
 $B \rightarrow 0B | \epsilon$

of $w|w$ contains starts and ends with same symbol }

$0\epsilon^* | \epsilon^* |$

$S \rightarrow 0AO | 1AI$

$A \rightarrow 0A | 1A | \epsilon$

of $w|w$ the length of w is odd }

$\overset{\epsilon}{\underset{A}{\epsilon}} (\underset{B}{\epsilon}\epsilon)^*$

$S \rightarrow AB$

$A \rightarrow 0|1$

$B \rightarrow AAB | \epsilon$

$\{w | w \text{ The length of } w \text{ is odd and its middle symbol is } 0\}$

$s \rightarrow 0s0 | 0s1 | 1s0 | 1s1 | 0$

Derivation, let, $w = 00101$

Both sides cover w & w^R .
else add s .

1010010 can also be derived.

10101

101100101

101100101

$\{w | w = w^R \text{ As, } w \text{ is a palindrome}\}$

$s \rightarrow 0s0 | 1s1 | 1| 0 | \epsilon$

#

S-

A-

B-

O

For C

A

A \rightarrow P

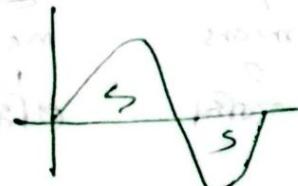
Q \rightarrow Z

P \rightarrow C

1 0 1

Num of 0 = Num of 1 100101

~~s → SOSI SIS0 | E~~
~~s → SOSI | SIS0 | E~~
~~s → OSIS | ISOS | E~~



$s \rightarrow OSIS | ISOS | E$

$s \rightarrow OSIS | ISOS | SS | E$

We may use

$s \rightarrow SOSI | SIS0 | E$

$s \rightarrow OSIS | ISOS | SS | E$

$s \rightarrow OAS | IBS | E$

$A \rightarrow 1 | OAA$

$B \rightarrow 0 | IBB$

$O^i | j^{2^k}$; where $i=j$

For Case-A

$A \rightarrow O^i | j^{2^k}$

$O^{i+j} | j^{2^k}$

$A \rightarrow PQ$

$Q \rightarrow 2Q | E$

$P \rightarrow OPI | E$

$PQ \leftarrow A$

or $j \neq k$

For case-B

$O^{i+j} | j^{2^k}$

$O^{ij}2^k$ where $\frac{i=j}{A}$ or $\frac{j \neq k}{B}$

For Case-A

$$\frac{O^{ij}2^k}{P^j Q^k}$$

$$A \rightarrow PQ$$

$$B \rightarrow 2Q|E$$

$$P \rightarrow OPI|E$$

For Case-B

$$\frac{O^i}{w} \frac{1^j 2^k}{x}$$

$$B \rightarrow wx$$

$$w \rightarrow ow|E$$

$$x \rightarrow y|z$$

$$y \rightarrow 1y2|y1$$

$$y1 \rightarrow 2y1|E$$

Here, $j \neq k$ means
 $\frac{j \neq k}{Y}$ or $\frac{j > k}{Z}$
means $\frac{2}{\text{तरीका}}$ means $\frac{1}{\text{एकत्र}}$

$$z \rightarrow 1z2|z1$$

$$z1 \rightarrow 1z1|E$$

so, Main Ans:

$$S \rightarrow A|B$$

$$A \rightarrow PQ$$

$$B \rightarrow 2Q|E$$

$$P \rightarrow OPI|E$$

$$B \rightarrow wx$$

$$w \rightarrow ow|E$$

$$x \rightarrow y|z$$

$$y \rightarrow 1y2|y1|1z2|z1$$

$$y1 \rightarrow 2y1|E$$

$$z \rightarrow 1z2|z1$$

$$z1 \rightarrow 1z1|E$$

$PQ \leftarrow A$

$S|QE \leftarrow B$

$\therefore PQE \leftarrow S$

#

$$L = \omega_1 \# \omega_2 \text{ where,}$$

num. of 0's in ω_1 = no. of 1's in ω_2

$$S \rightarrow 1S \mid S0 \mid 0S1 \mid \#$$

Q) Give a context-free grammar for the following language. For, $\{w \in \{0,1\}^*\}$

(a) ~~if~~ =

(a) w contains odd number of 1's

(b) w starts and ends with same characters.

(c) $w = u0^{2i}v1^{3i}$ where $u \in B, v \in A$ and $i \geq 0$

(a) RE $\rightarrow (\overbrace{0^* 1 0^*}^B 1)^* 0^* \overbrace{0^*}^C$

$$A \rightarrow B \mid C1C$$

$$C \rightarrow 0C \mid \epsilon$$

$$B \rightarrow DB \mid E$$

$$D \rightarrow C1C1$$

(b) $0 \leq 0 \mid 1 \leq 1 \mid 0 \mid 1$

$w \rightarrow 0p0 \mid 1p1 \mid 0 \mid 1$

$p \rightarrow qp \mid \epsilon$

$q \rightarrow 0 \mid 1$

(c) $w = \overbrace{u}^w 0^{2i} v 1^{3i}$

$s \rightarrow w x$

$x \rightarrow 00 \times 111 \mid A$

~~$w \rightarrow opo$ and A, B (funct)~~

(d) $\overbrace{w}^w \underbrace{0^{2i}}_x \underbrace{1^{3i}}_v$

$s \rightarrow w x A$

$x \rightarrow 00 \times 111 \mid \epsilon$

$s \mid s/a \leftarrow A$

$\epsilon \mid s0 \leftarrow C$

$\epsilon \mid sa \leftarrow D$

$\epsilon \mid sA \leftarrow E$

$\emptyset \emptyset$
 $0, \epsilon -$
 $1, n -$

$a =$

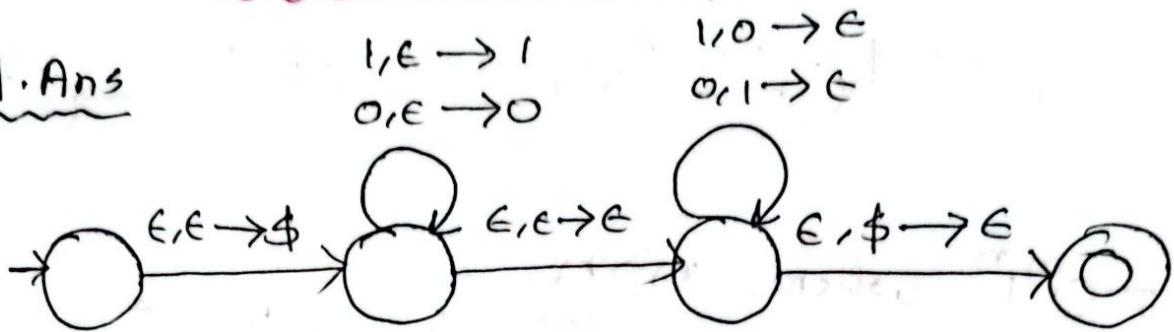
$\{ \text{Ex} \}$

for

x
 x
 x
 x
 $\$$

Spring - 2022 } Final

1. Ans



2. Ans

(a) $A = \{w \in \{0, 1\}^*: \text{at least two } 0's\}$

$1^* 0 1^* 0 \epsilon^*$

$$\begin{aligned}s &\rightarrow P O P O Y \\ P &\rightarrow 1 P | \epsilon \\ Y &\rightarrow O Y | 1 Y | \epsilon\end{aligned}$$

(b) $L = \{w \in \{0, 1\}^*: w = 0^{3i} \vee 1^{2i} \text{ where } i \geq 0 \text{ and } \vee \in A\}$

$s' \rightarrow 0 0 0 s' 1 1$

$s' \rightarrow s$

$s \rightarrow P O P O Y$

$P \rightarrow 1 P | \epsilon$

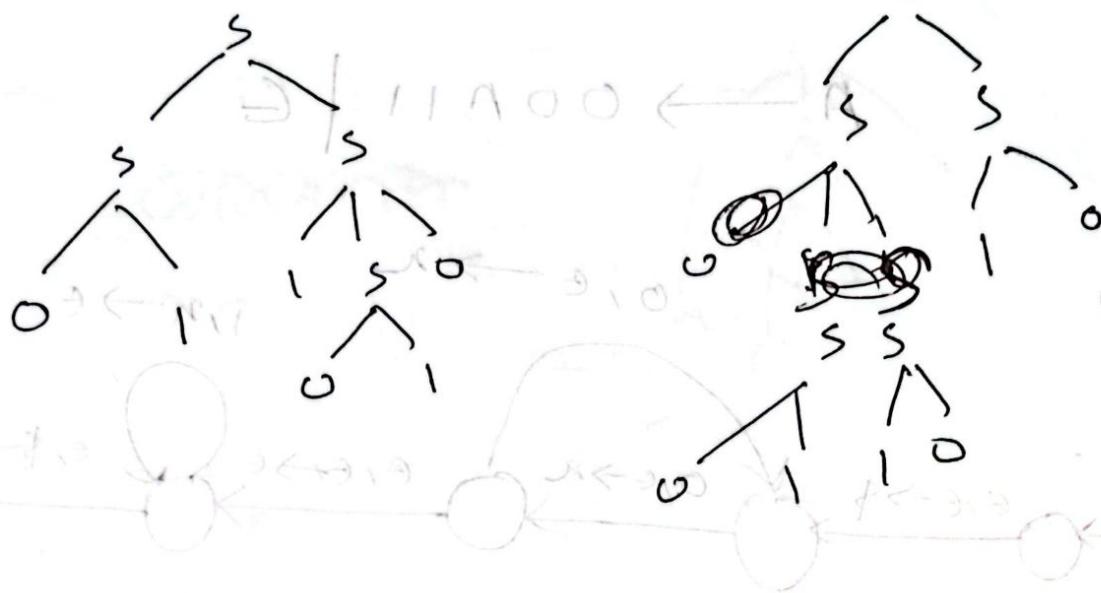
$Y \rightarrow O Y | 1 Y | \epsilon$

3.AM

$S \rightarrow OS1$	$ $	150	$ $	SS	$ $	01	$ $	10
---------------------	-----	-------	-----	------	-----	------	-----	------

② $\begin{array}{c} \text{S} \\ \text{---+---} \\ \text{O} \quad \text{I} \end{array}$ $\begin{array}{c} \text{S} \\ \text{---+---} \\ \text{O} \quad \text{I} \end{array}$ $\begin{array}{c} \text{S} \\ \text{---+---} \\ \text{O} \quad \text{I} \end{array}$

$1A0 \leftarrow 2$ ③

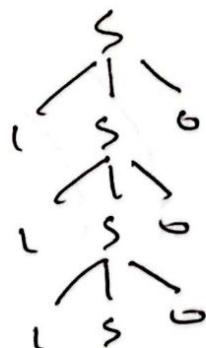


④

$111000 ; S \Rightarrow 150$

$\Rightarrow 115000$

$= 111000$



Summen 22

1. Am

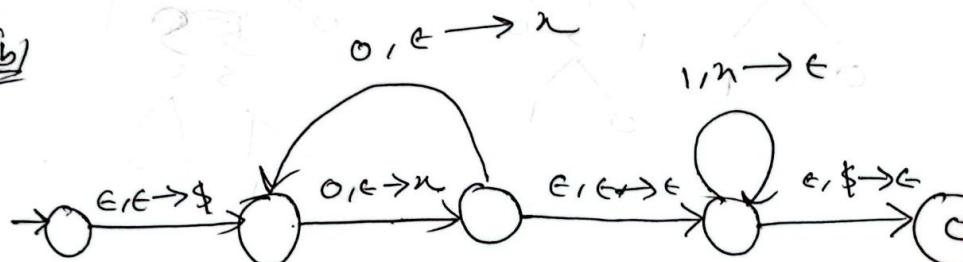
$L = \{ w \in \Sigma^* ; w = 0^n 1^m ; n \text{ is odd} \}$

(a)

$S \rightarrow 0 A 1$

$A \rightarrow 0 0 A 1 1 \mid \epsilon$

(b)



3. Am

0 0 1 1 1

(a)

$A \Rightarrow A 1$

$\Rightarrow A 1$

$\Rightarrow 0 A 1$

$\Rightarrow 0 0$

$\Rightarrow 0 0$

(c)

A

A

(d)

0 1 1 1 1

3. Ans $A \rightarrow A1 | OA1 | O1$

001111

(a) $A \Rightarrow A1$

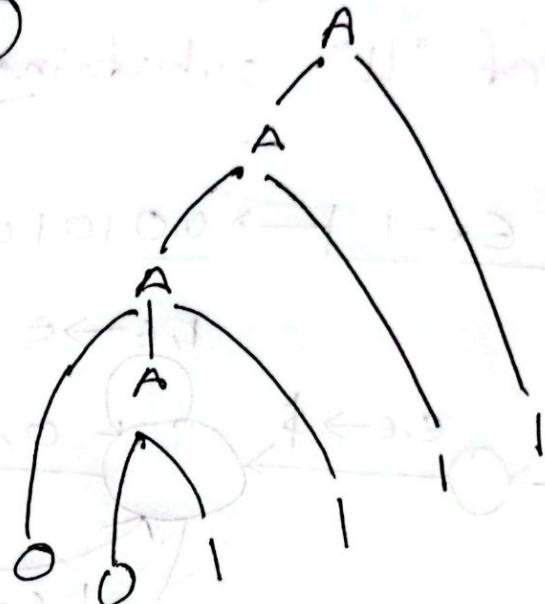
$\Rightarrow A11$

$\Rightarrow OA111$

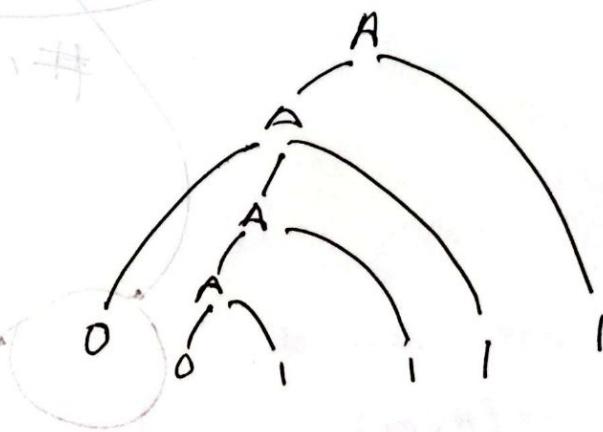
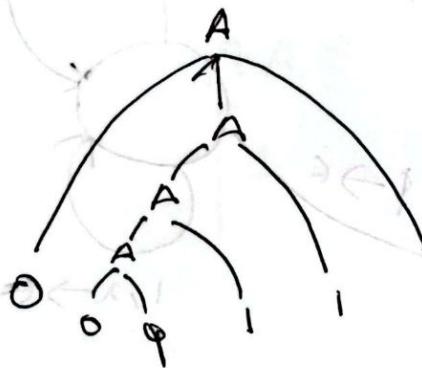
~~$\Rightarrow OOA111$~~

$\Rightarrow 001111$

(b)



(c)



(d)

$011111 ; A \Rightarrow A1$

$\Rightarrow A11$

$\Rightarrow A111$

$\Rightarrow A1111$

$\Rightarrow 011111$

$\# L = \{0^i 1^j 2^k 3^m ; j=k+1, i=m-2\}$

$$\frac{0^i 1^j 2^k 3^m}{0^i \underline{1^{k+1}} 2^k 3^{i+2}}$$

$S \rightarrow 0S3$

$S \rightarrow P33$

$P \rightarrow 1P2$

$P \rightarrow 1$

$\# L = \{0^i 1^j 2^k 3^m 4^n ; j=i, m=n, k > 0\}$

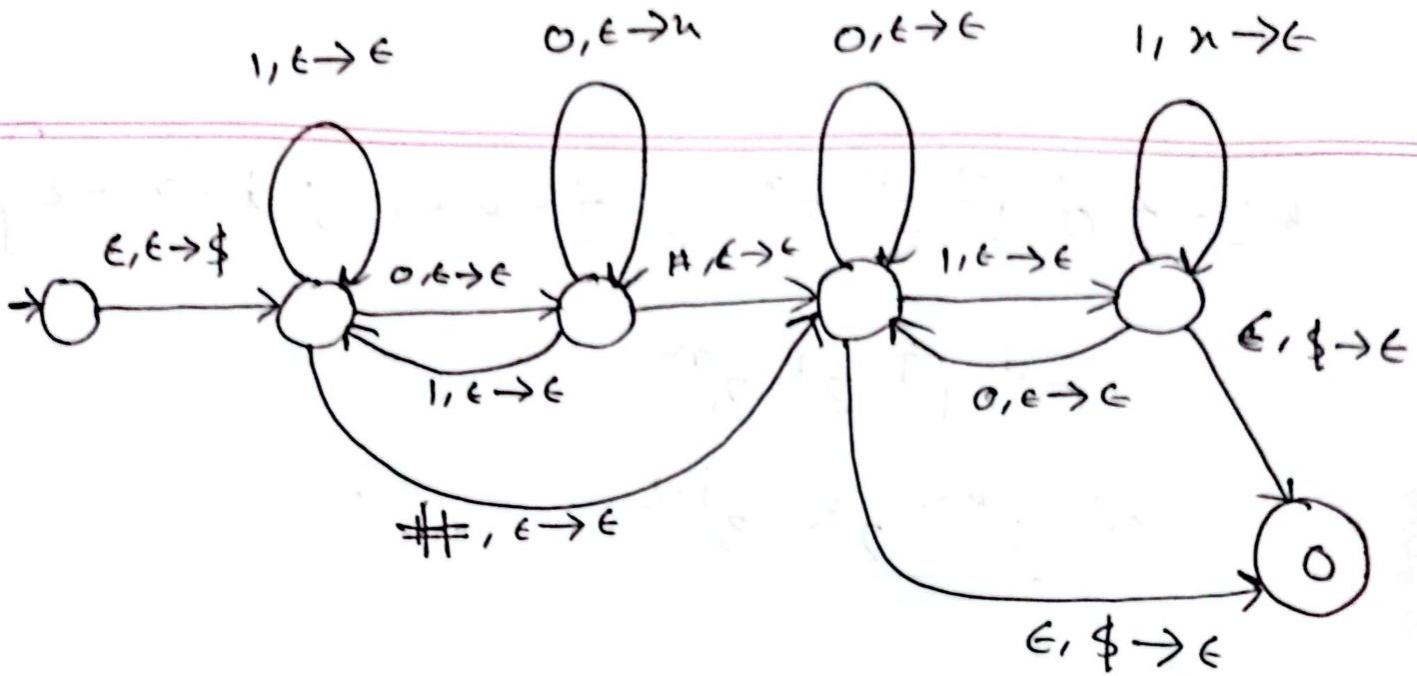
$$\frac{0^i 1^j 2^k 3^m 4^n}{0^i 1^j 2^k 3^m 4^n}$$

$S \rightarrow ABC$

$A \rightarrow 0A1 | \epsilon$

$C \rightarrow 3C4 | \epsilon$

$B \rightarrow 2B | \epsilon$



$\# L = \{ w_1 \# w_2 \mid |w_1| \text{ is double of } |w_2| \}$

$\boxed{w_1} \# \boxed{w_2}$

$S \Rightarrow A A S A$

$\Rightarrow A A \underline{A A} S \underline{A A}$

$S \rightarrow A A S A \#$

$A \rightarrow 0 \mid 1$

$\# |w_1| = m ; |w_2| = n ; 3m = 4n$
 $(m, n) = (4, 3)$

$\boxed{\dots \# \dots}$

$S \rightarrow A A A A S A A A \#$

$A \rightarrow 0 \mid 1$