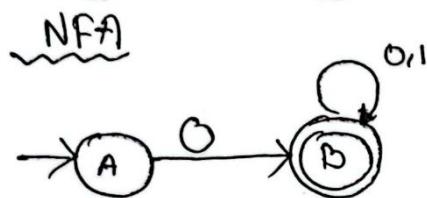


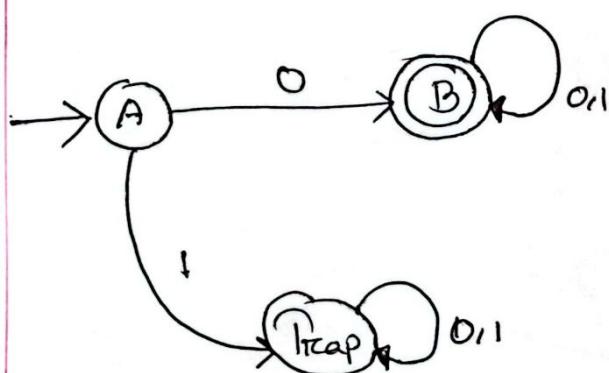
Conversion

NFA \rightarrow DFA



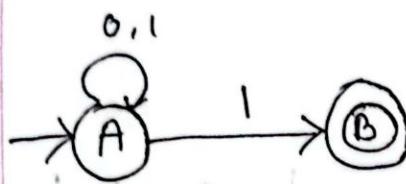
	0	1
A	B	\emptyset
B	B	B

DFA

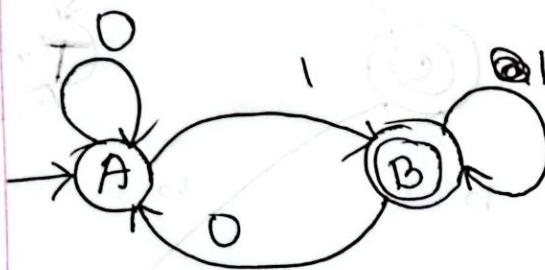


	0	1
A	B	Trap
B	B	B
Trap	Trap	Trap

NFA



DFA



	0	1
A	{A}	{A,B}
B	\emptyset	\emptyset

AB = single state

Hence AB: Relationship will be defined by $A \cup B$.

	0	1
A	{A}	{A,B}

AB	{A}	{AB}
----	-----	------

$$A \cup \emptyset = A$$

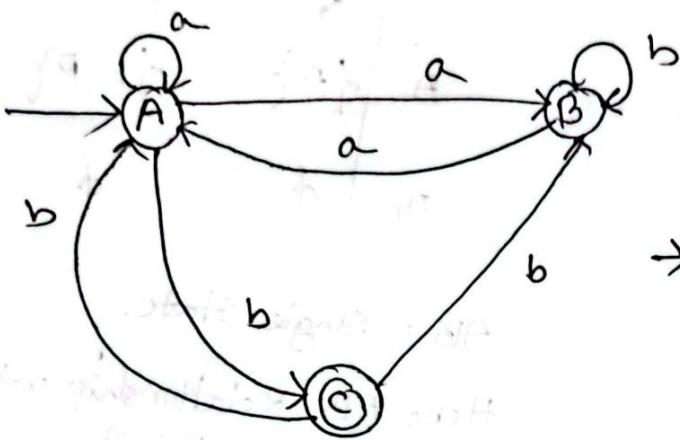
$$AB \cup \emptyset = AB$$

Here B \in states (एक जगह)

path $\overbrace{A}^0 \xrightarrow{1} B$ वाले B की

Path एकोरा \overbrace{A}^0 .

NFA

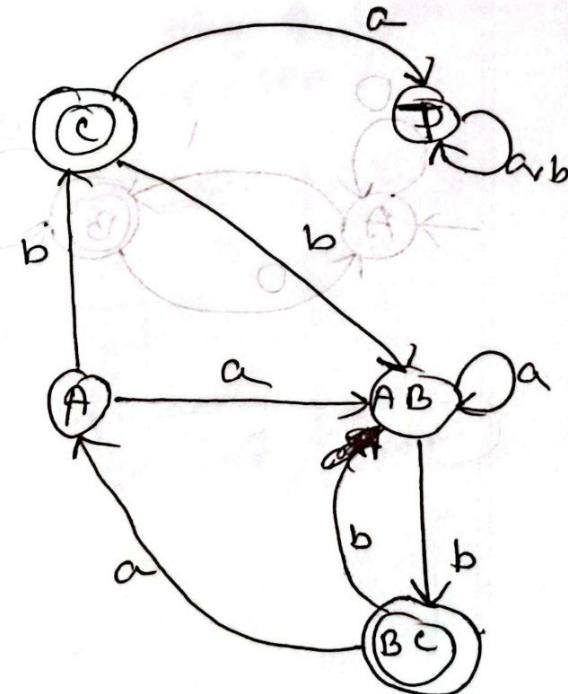


	a	b
A	{A,B}	{C}
B	{A}	{B}
C	∅	{A,B}

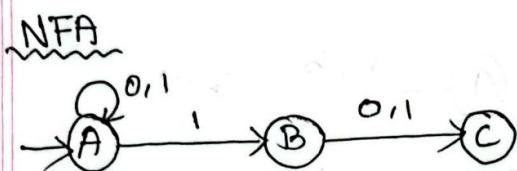
DFA

	a	b
A	AB	C
AB	AB	BC
BC	A	AB
C	T _{trap}	AB
T	T	T

$T = T_{trap}$



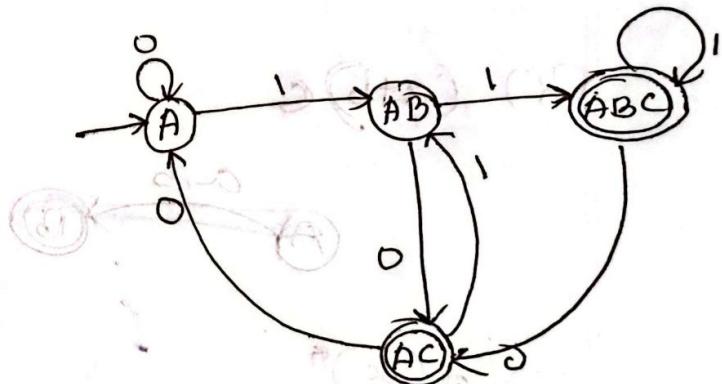
Design a NFA for a language that accepts all strings over $\{0,1\}$ in which the second last symbol is always "1". Then convert it to its equivalent DFA.



	0	1
A	A	$\{A, B\}$
B	C	C
C	\emptyset	\emptyset

DFA

	0	1
①	A	AB
②	AB	ABC
③	AC	AB
④	ABC	ABC
⑤	C	T
⑥	T	T
⑦	T	T



If we look closely there's no way to lead to C, so the last 2 will be unnecessary, means there'll be no states called C

Regular Expression \rightarrow FA

Some Important Rules

$$(a+b) \longrightarrow$$



$$(ab) \longrightarrow$$

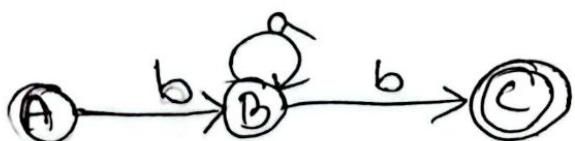


$$a^*$$

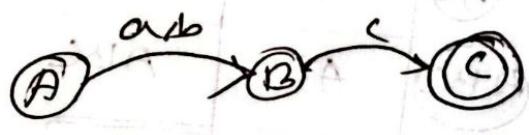
$$\longrightarrow$$



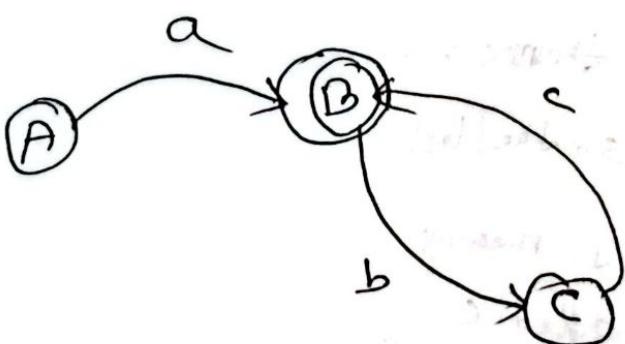
(i) ba^*b



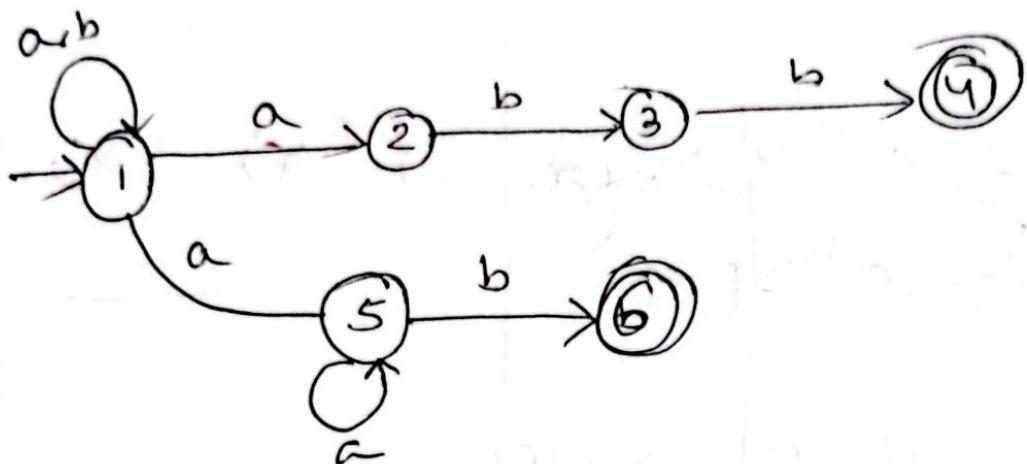
(ii) $(a+b)c$



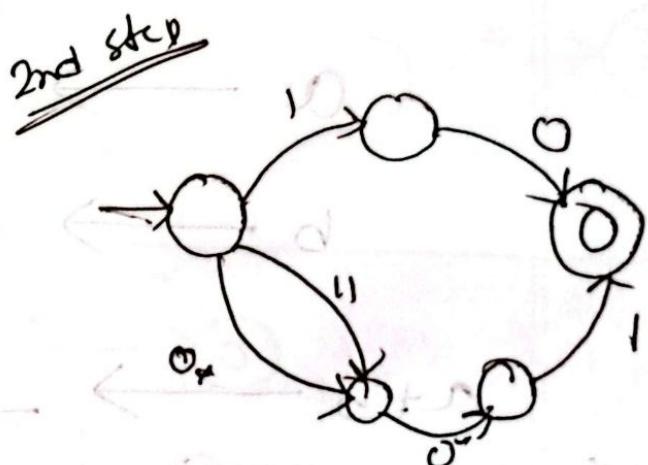
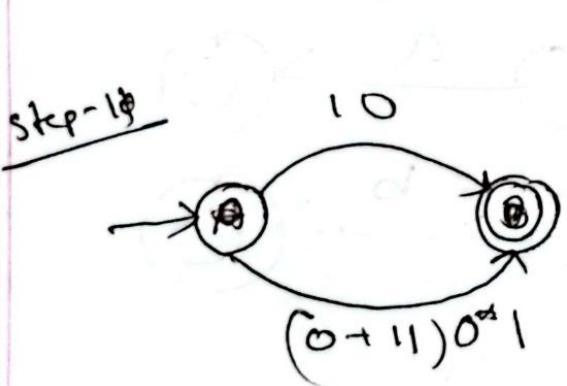
(iii) $a(bc)^*$



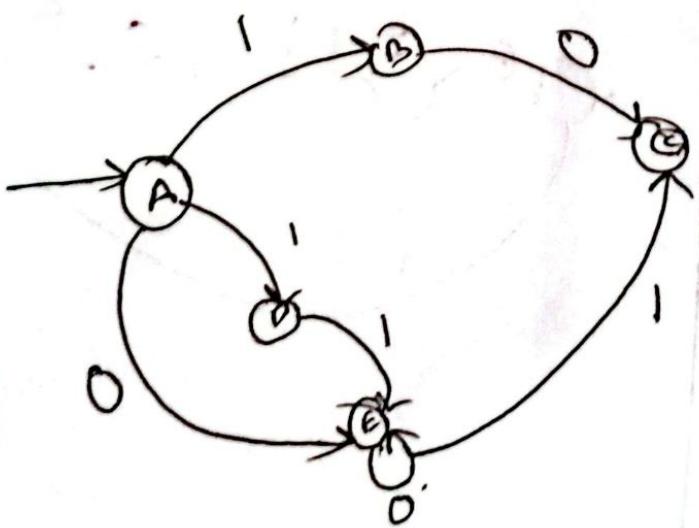
$(a|b)^*$ $(abb|a^*b)$



$10 + (0+11) 0^* 1$



Final step



RegEx to NFA

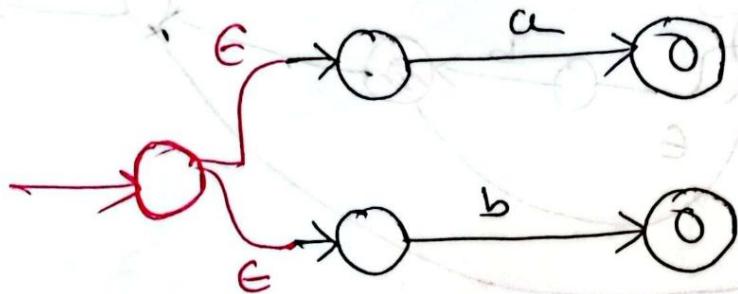
a



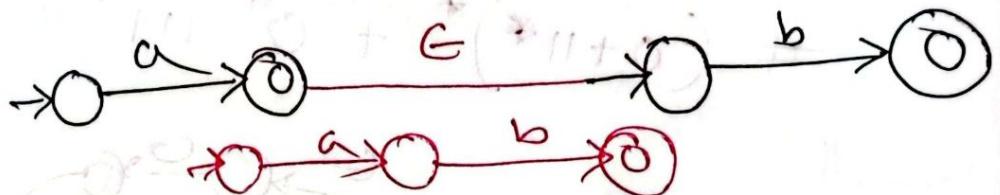
b



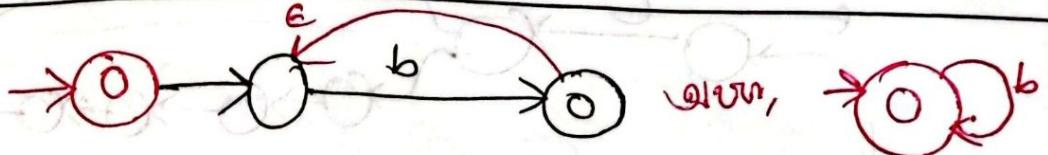
$a+b$



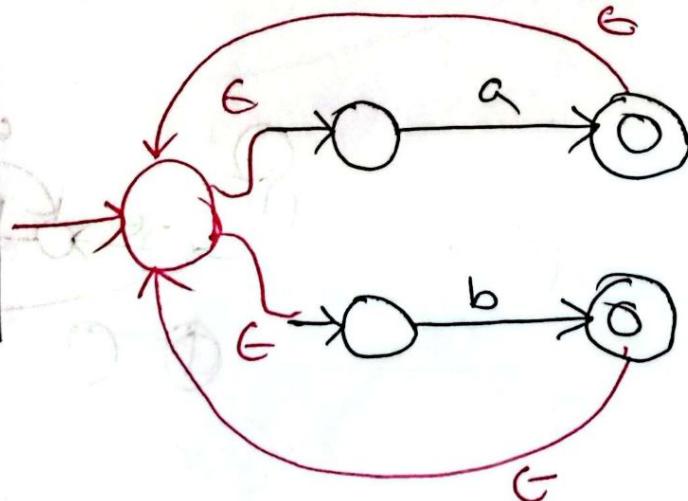
ab



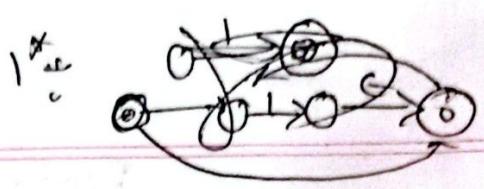
b^*



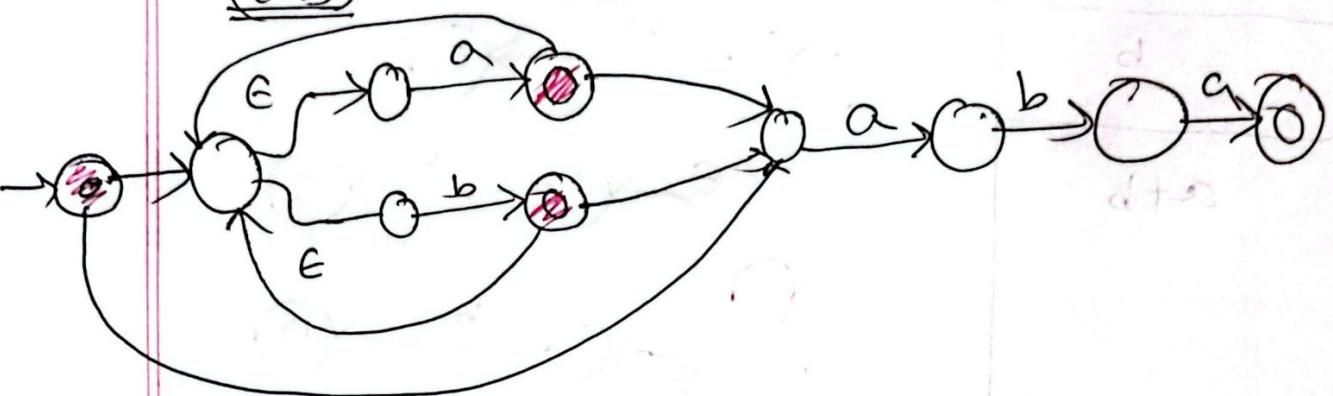
$(a+b)^4$



$\#(a+b)^*$ aba

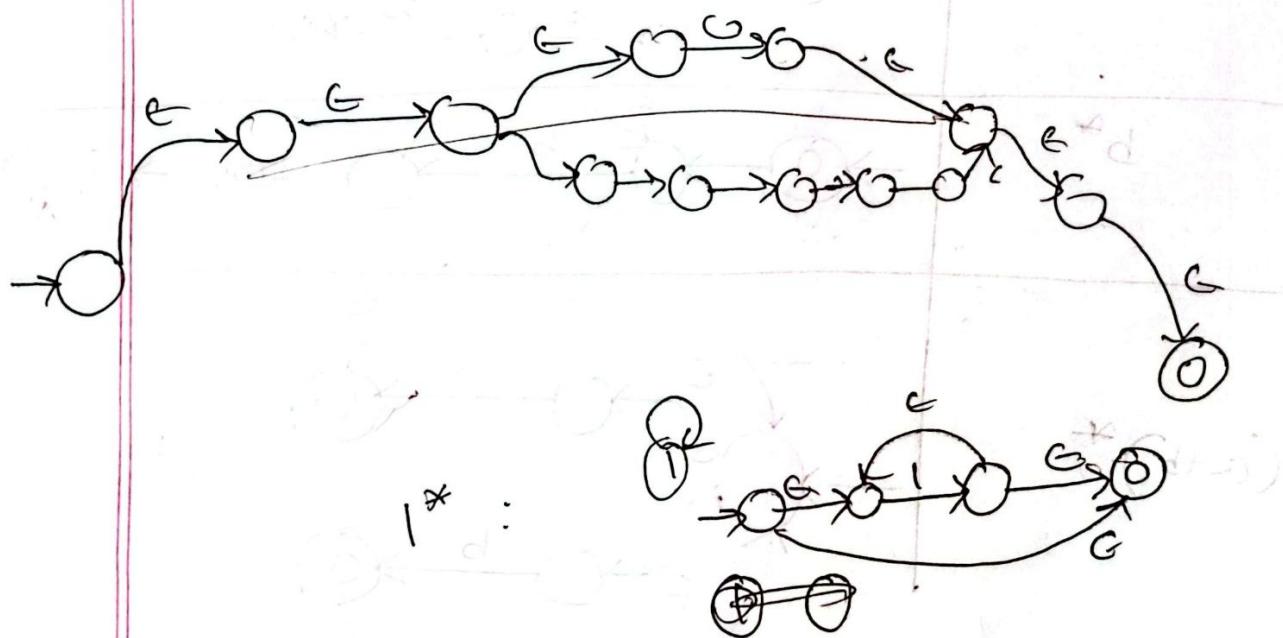


$(a+b)^*$



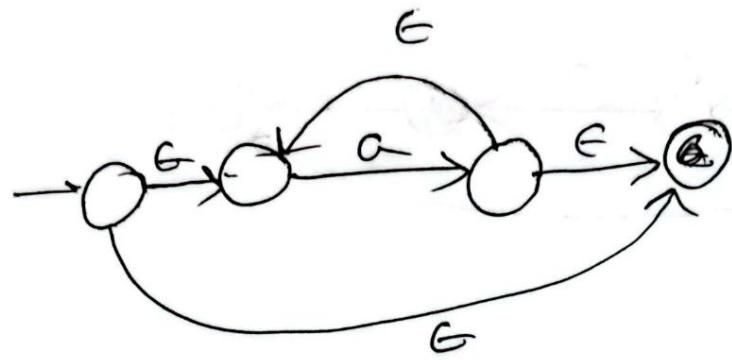
aba

$\#(0+11^*)^* + 0^* 111^*$

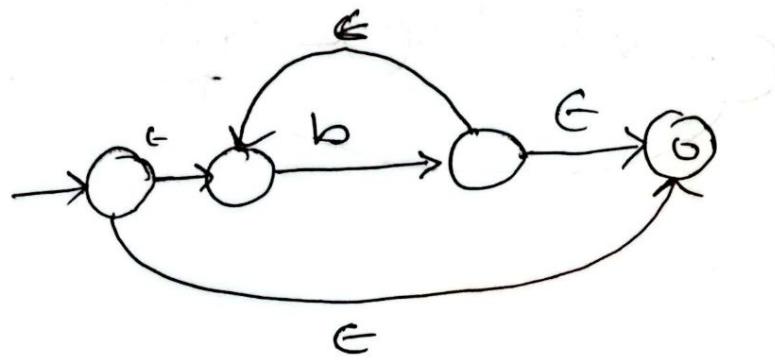


1^* :

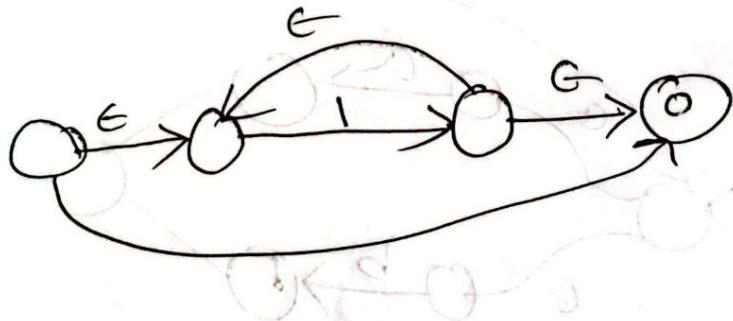
a^α



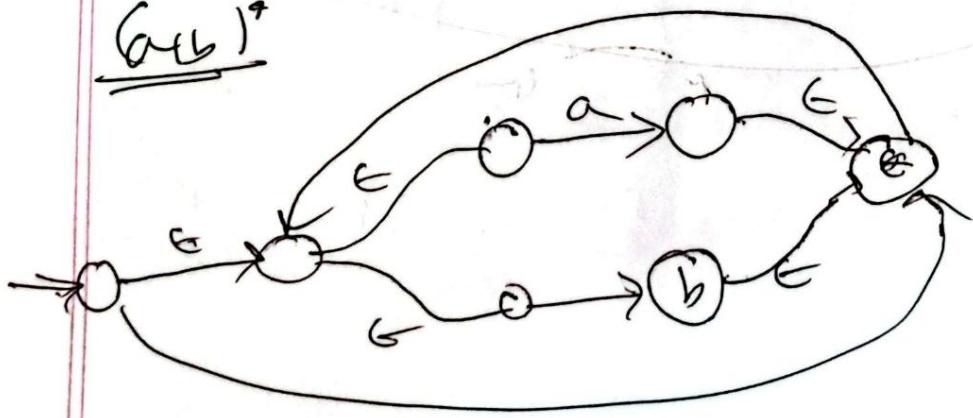
b^α



\perp^α



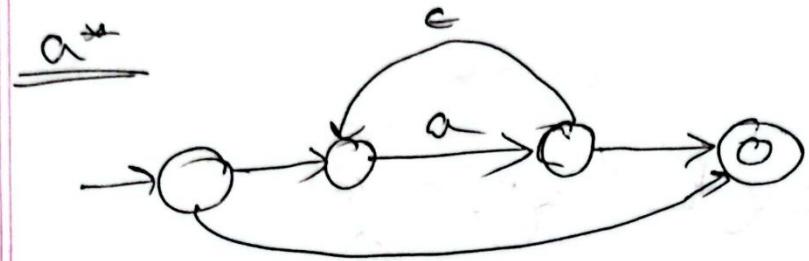
$(ab)^\alpha$



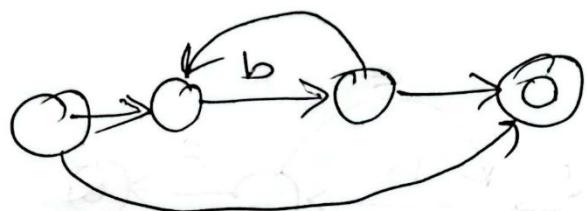
σ

$(\alpha + \beta)^n$

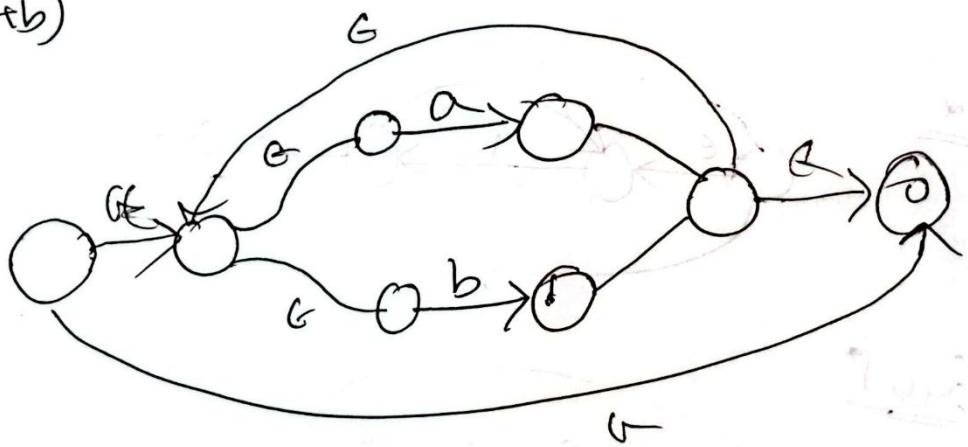
α^*



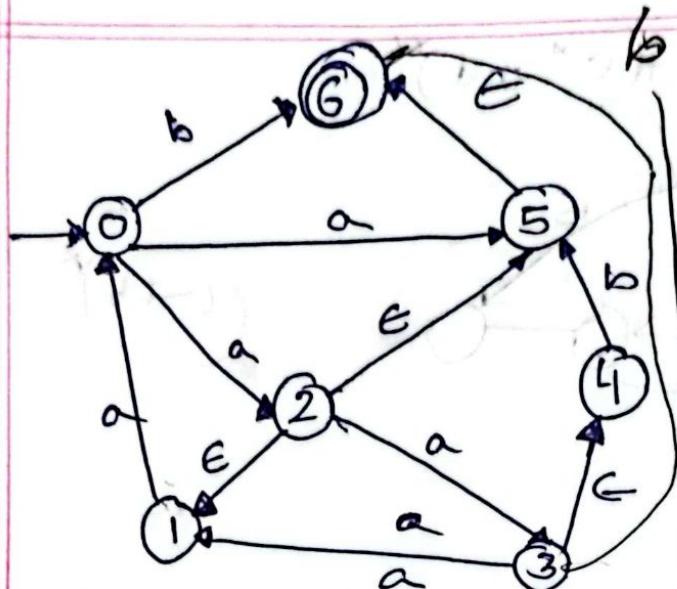
b^*



$(\alpha + \beta)^*$



subset Construction



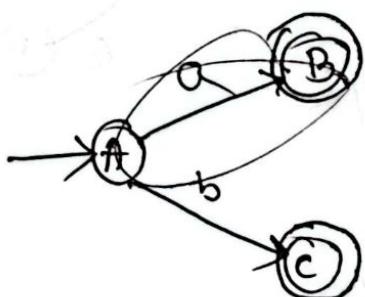
start state

$$\epsilon \text{ closure}(0) = \{0\} = A$$

$$\text{Move}(A, a) \xrightarrow{\epsilon} \{2, 5\} = B$$

$$\text{Move}(A, b) = \{6\} \xrightarrow{\epsilon} \{6\} = C$$

$$\text{Move}(B, a) = \{0, 3\}$$



start state

$$\epsilon \text{ closure}(0) = \{0\} = i$$

$$\text{Move}(A, a) = \{5, 2\} \xrightarrow{\epsilon} \{1, 2, 5, 6\} = B$$

$$\text{Move}(A, b) = \{6\} \xrightarrow{\epsilon} \{6\} = C$$

$$\text{Move}(B, a) = \{0, 3\} \xrightarrow{\epsilon} \{0, 3, 1\} = D$$

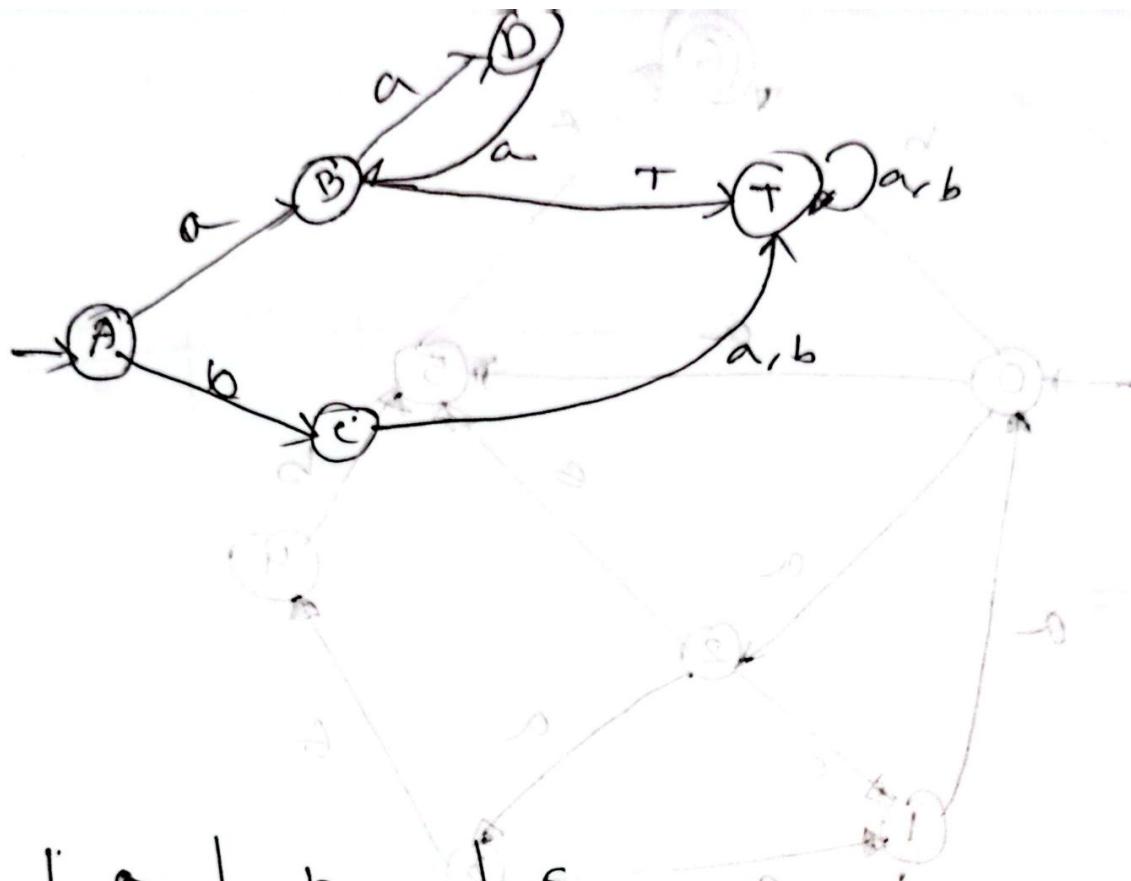
$$\text{Move}(B, b) = \{\} \rightarrow \emptyset = T$$

$$\text{Move}(C, a) = \{\} \rightarrow \emptyset$$

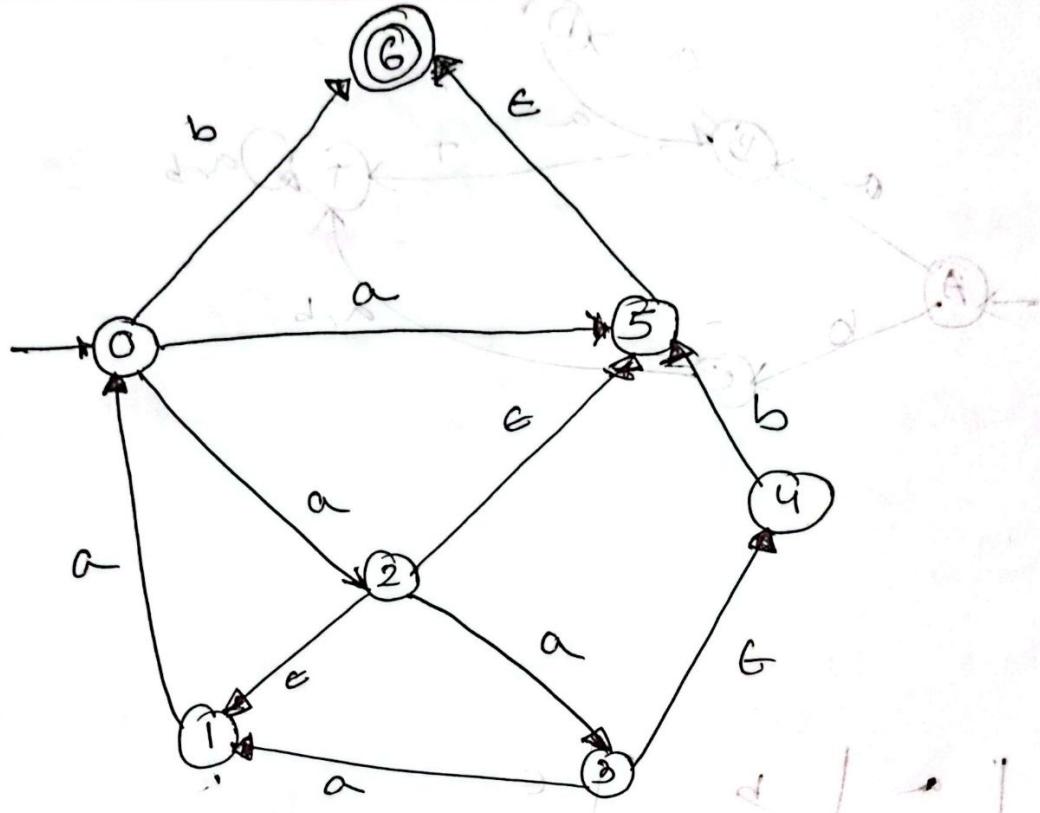
$$\text{Move}(C, b) = \{\} \rightarrow \emptyset$$

$$\text{Move}(D, a) = \{5, 2, 3\} \xrightarrow{\epsilon} \{2, 3, 5, 6\}$$

$$\text{Move}(D, b) =$$



	a	b	c	d	e	f	g	h	i
0	{5,2}	6			{5,2,6,1}				0
1	0	φ	φ	φ	φ	φ	φ	φ	0
2	1,3		φ		φ	φ	φ	φ	1
3			φ		φ	φ	φ	φ	2
4			φ		φ	φ	φ	φ	3
5			φ		φ	φ	φ	φ	4
6			φ		φ	φ	φ	φ	5



	a	b	e	d	f	g	h
0	5,2	6	φ	φ	0	1,4	0
1	0	φ	φ	φ	c,1	5	
2	3	φ	φ	1,5			
3	1	φ	φ	4			
4	φ	5	φ				
5	φ	φ	φ	6			
6	φ	φ	φ	φ			

	a	b
0	5,2	6
5,2	5,2	6
6		

A. Strategie: (0, 5, 2)

Wertfunktion: $\{8,8\} \rightarrow \{0,0\}$ mit

Strategie: $\{0, 5, 2\} \rightarrow \{0, 0\}$ mit

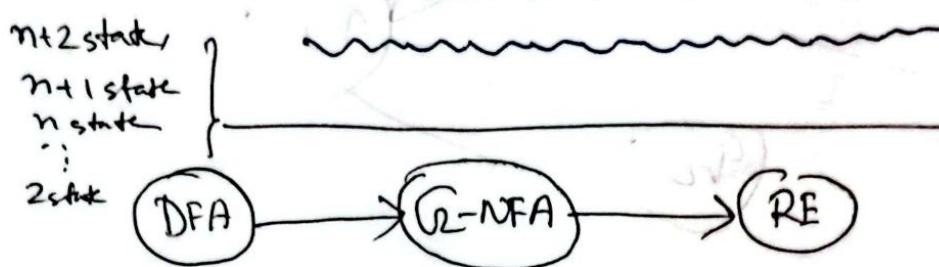
Wertfunktion:

$\{0, 0\} \rightarrow \{0, 0\}$

Finals

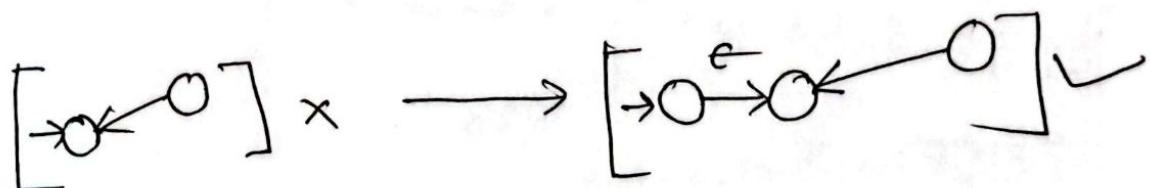
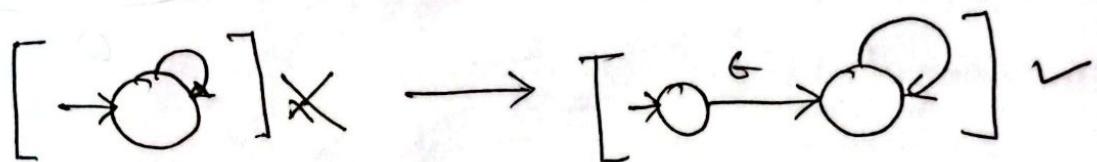
DFA To RE

State Elimination Method



N₂-NFA

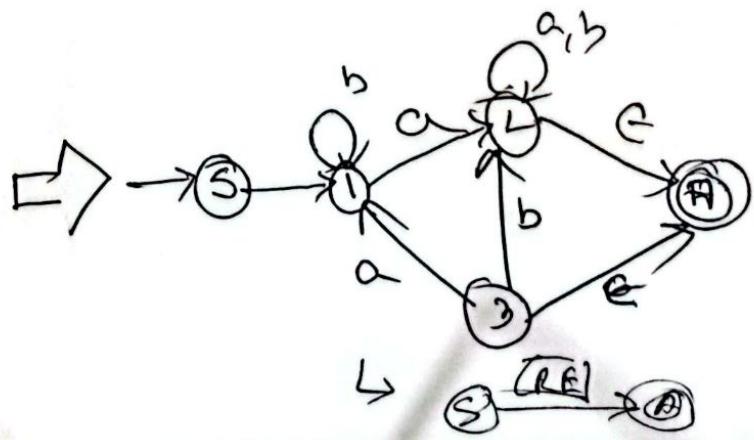
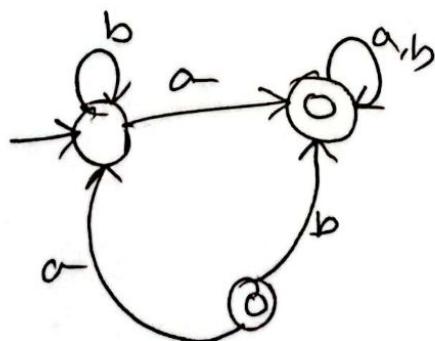
(i) No incoming arrow into starting state.



(ii) Only 1 accepting state

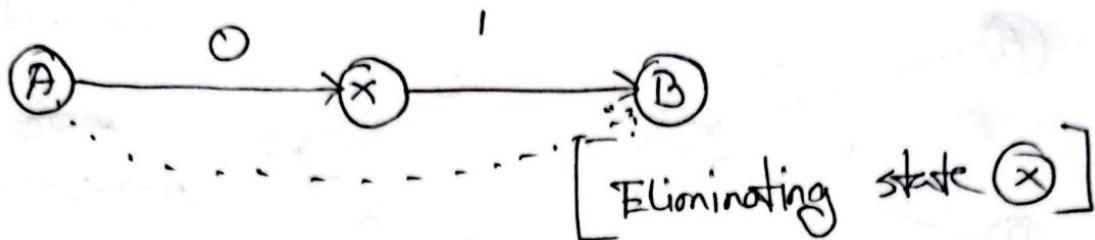
No outgoing arrow from accepting state

Ex

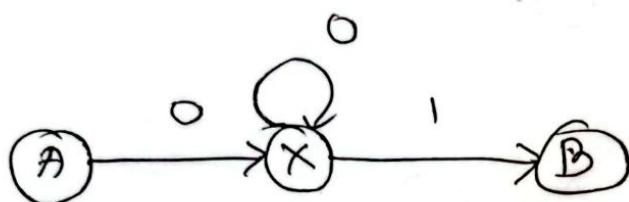


RE

#

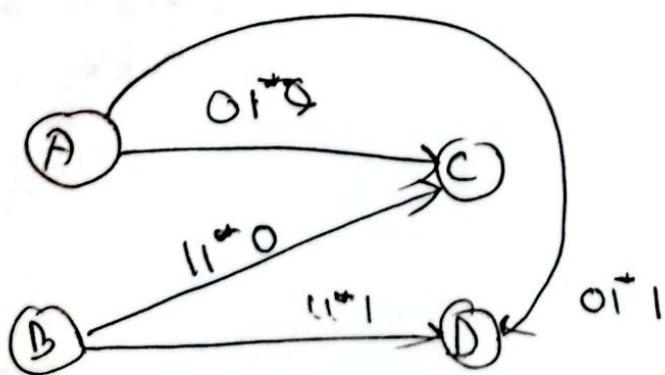
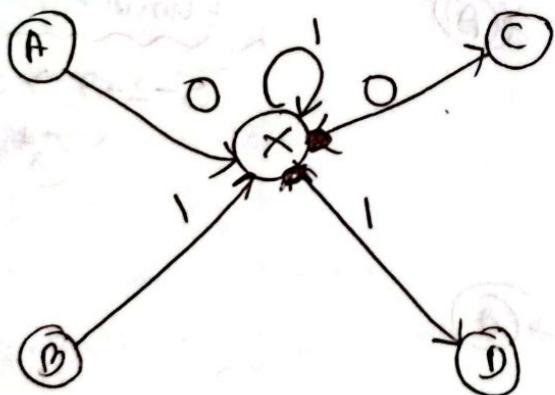


#

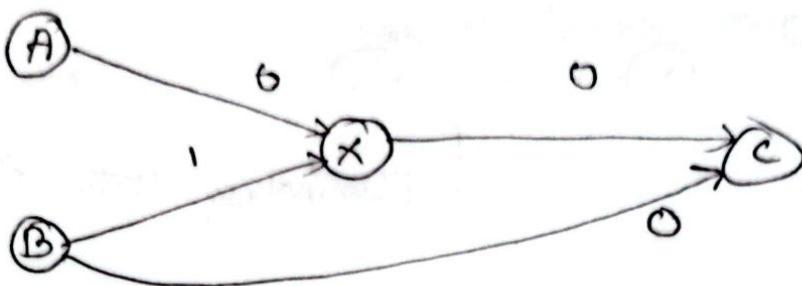


Total Transition

= # Incoming arrows (w/o self loop)
 x # Outgoing arrows (w/o self loop)

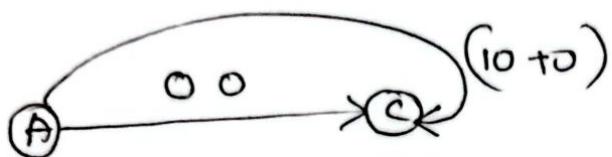


#

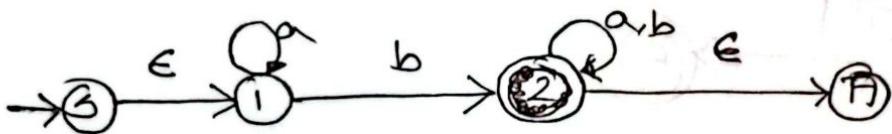
Eliminating \textcircled{X} ~~state X~~

$$A-X-C \rightarrow AC(00)$$

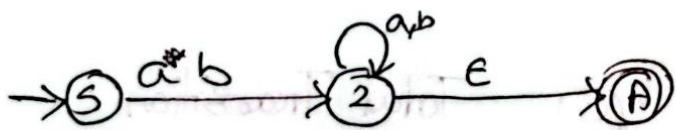
$$B-X-C \rightarrow BC(10)0$$



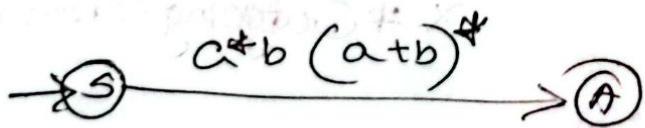
#

1st step

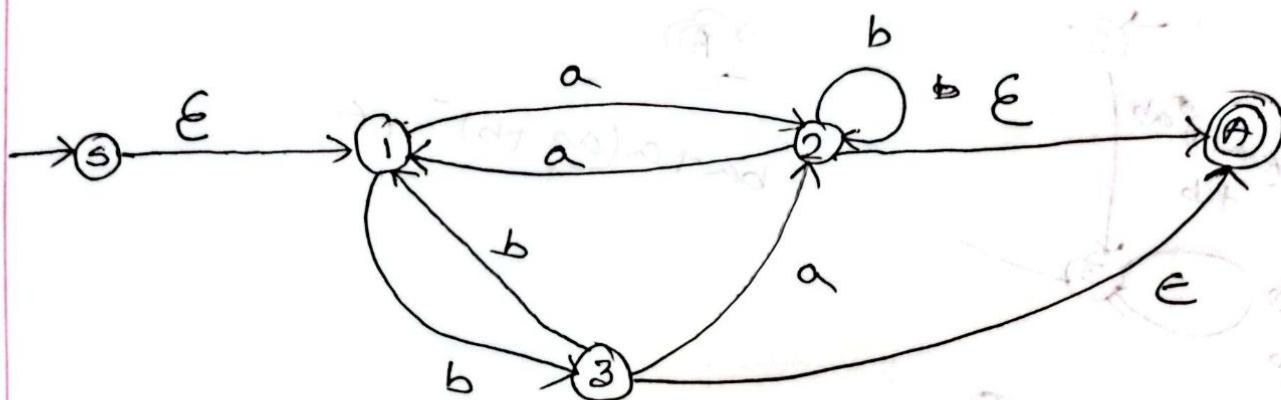
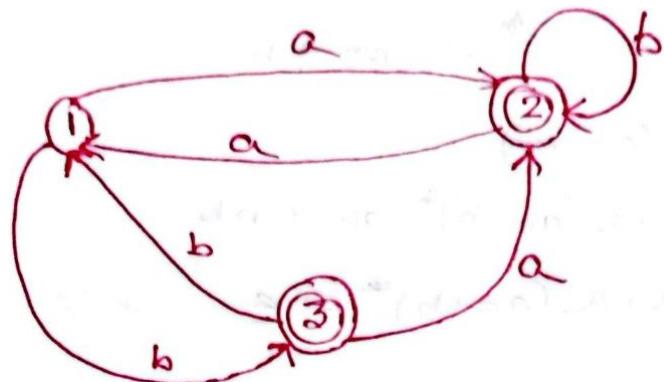
Eliminating 1
 $S-1-2 \rightarrow S-2(a^* b)$

2nd step

Eliminating 2
 $S-2-A \rightarrow S-A$
 $(a^* b (a+b)^*)^*$

3rd step

#



Eliminating -① (3×2)

$$S-1-2 \rightarrow S-2; a$$

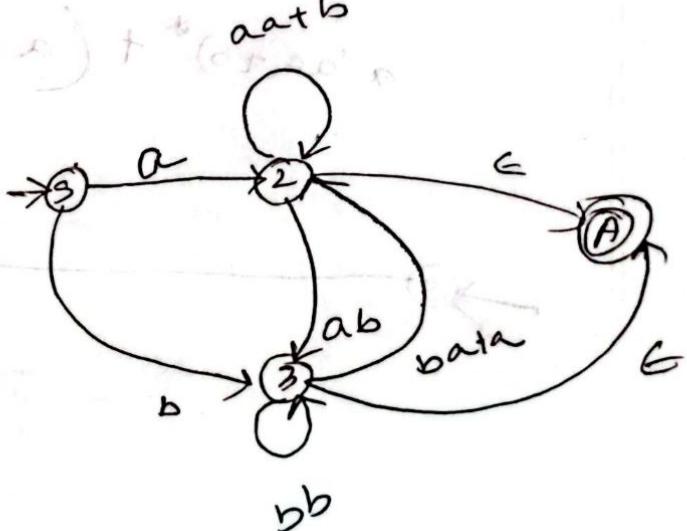
$$S-1-3 \rightarrow S-3; b$$

$$2-1-2 \rightarrow 2-2; aa+b$$

$$2-1-3 \rightarrow 2-3; ab$$

$$3-1-2 \rightarrow 3-2; ba+a$$

$$3-1-3 \rightarrow 3-3; bb$$



Eliminating -② 2x2

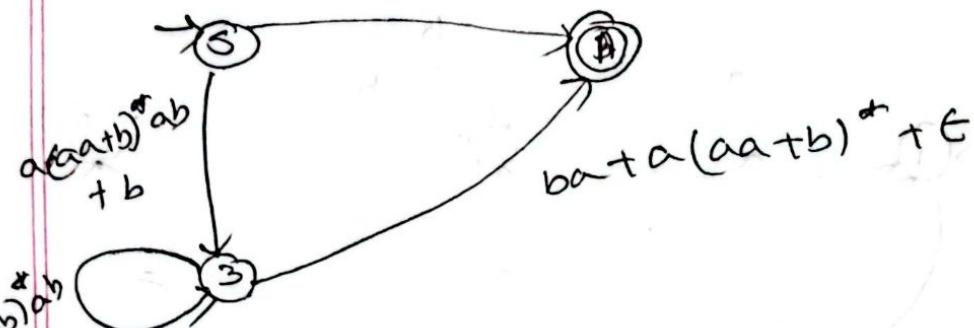
$$S-2-\beta \rightarrow S-\beta ; a(aatb)^* ab \cancel{+ b}$$

$$S-2-\alpha \rightarrow S-\alpha ; a(aa+b)$$

$$3-2-\beta \rightarrow 3-3 ; ba+a(aatb)^* ab + bb$$

$$3-2-\alpha \rightarrow 3-\alpha ; ba+a(aa+b)^* + \epsilon$$

$$a(aa+b)^* \cancel{ab}$$



Eliminating -⑤

$$a(aa+b)^* + (a(aatb)^* ab + b)$$

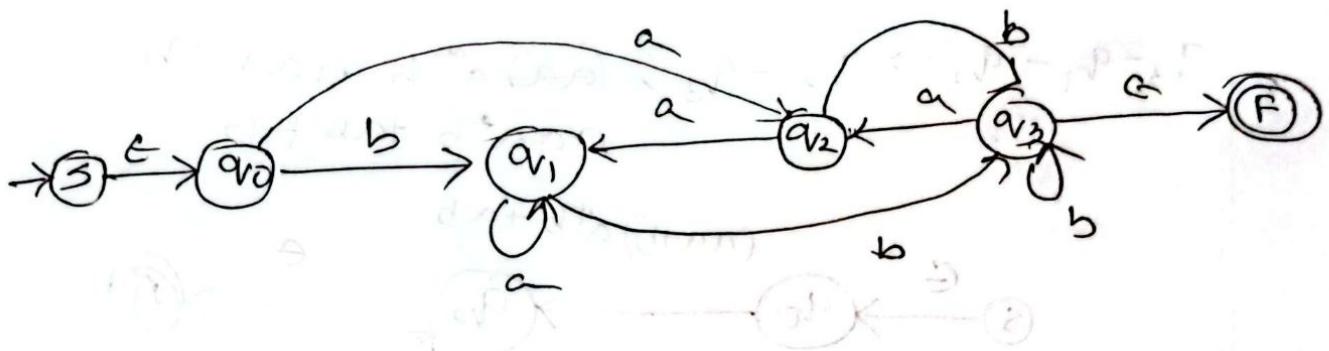
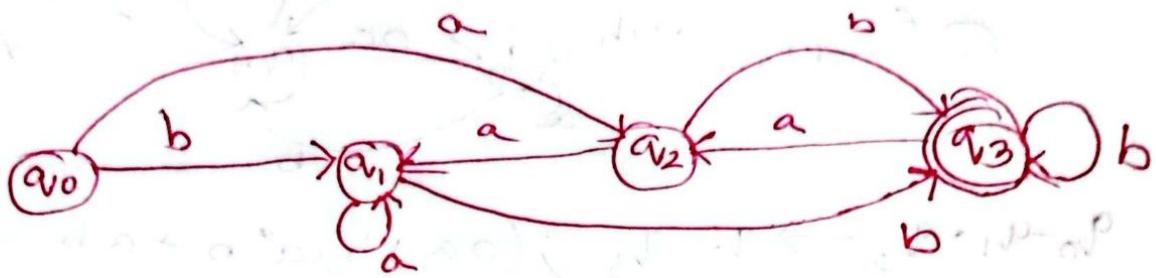
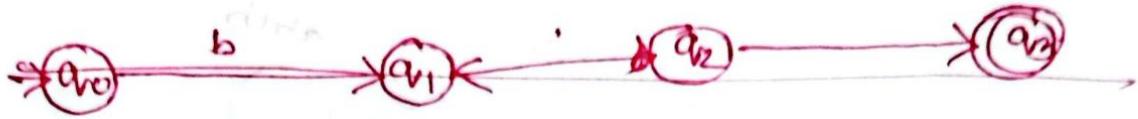
$$(a(aatb)^* ab + bb)^* (ba+a(aatb)^* + \epsilon)$$



$\begin{array}{l} (a(aatb)^* ab + bb)^* \\ (ba+a(aatb)^* + \epsilon) \end{array}$

$\begin{array}{l} d \leftarrow a-a \\ d+bb \leftarrow a-a \leftarrow a-a \\ d+a \leftarrow a-a \leftarrow a-a \\ d+d \leftarrow a-a \leftarrow a-a \end{array}$

#1



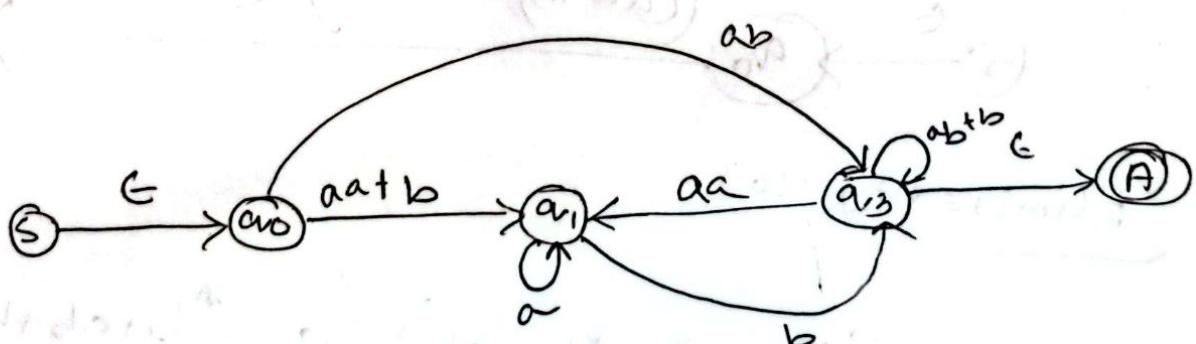
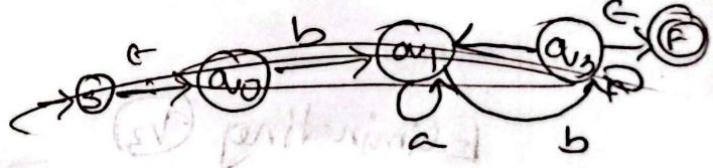
Eliminating q_2

$$q_3 - q_2 - q_1 \rightarrow q_3 - q_1 ; aa$$

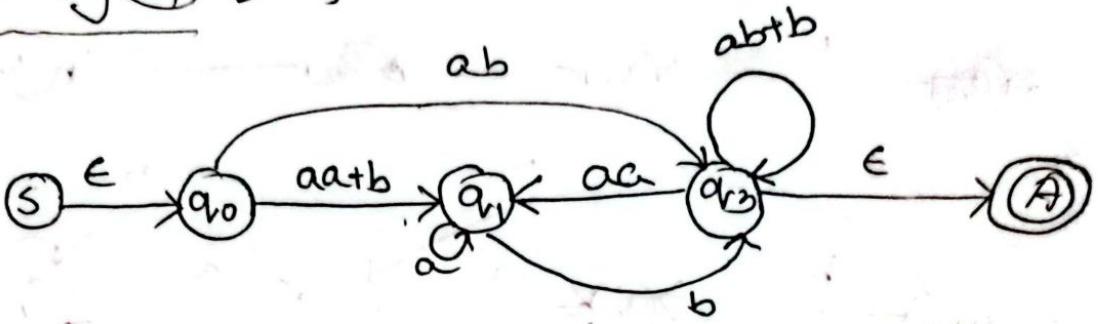
$$q_3 - q_2 - q_3 \rightarrow q_3 - q_3 ; ab + b$$

$$q_0 - q_2 - q_1 \rightarrow q_0 - q_1 ; aa + b$$

$$q_0 - q_2 - q_3 \rightarrow q_0 - q_3 ; ab$$

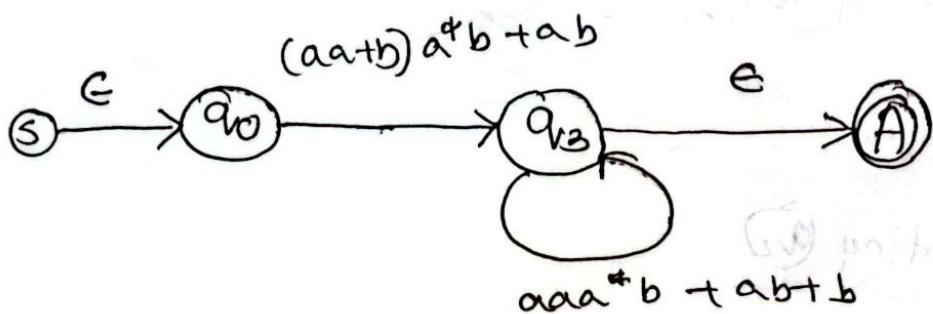


Eliminating q_1 [2x1]

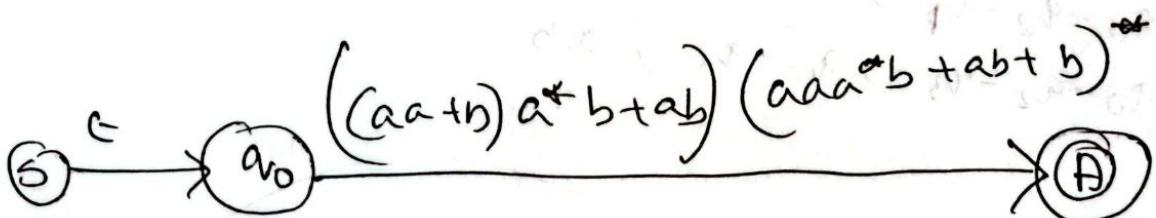


$$q_0 - q_1 + q_3 \rightarrow q_0 - q_3 ; (aa+b)a^*b + ab$$

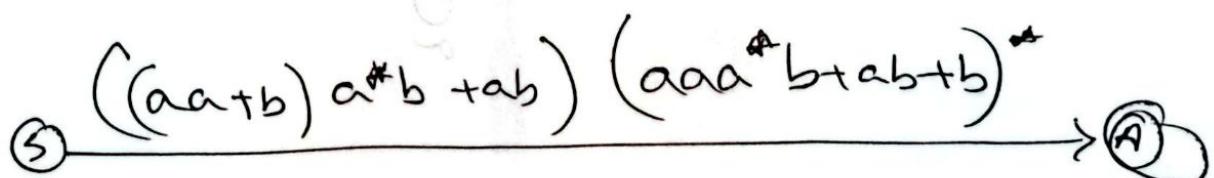
$$q_3 - q_1 - q_3 \rightarrow q_3 - q_3 ; \cancel{(aa)}a^*b + ab + b \\ aaa^*b + ab + b$$

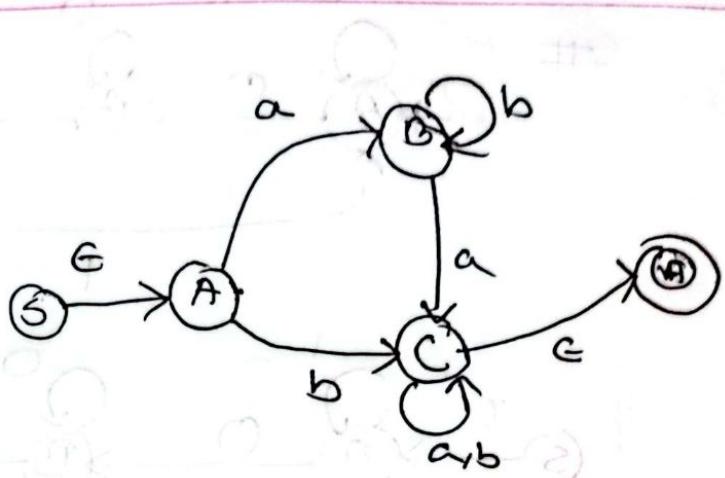
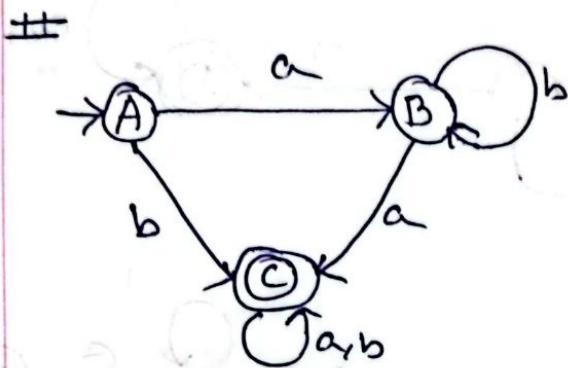


Eliminating q_2



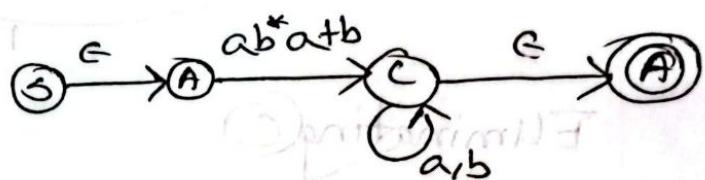
Eliminating q_0



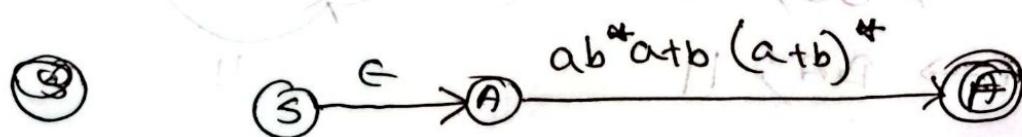


Eliminating ③

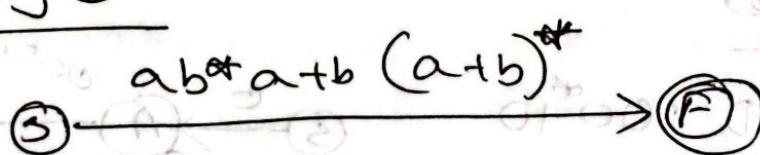
$$A-B-C \rightarrow AC; ab^*atb$$

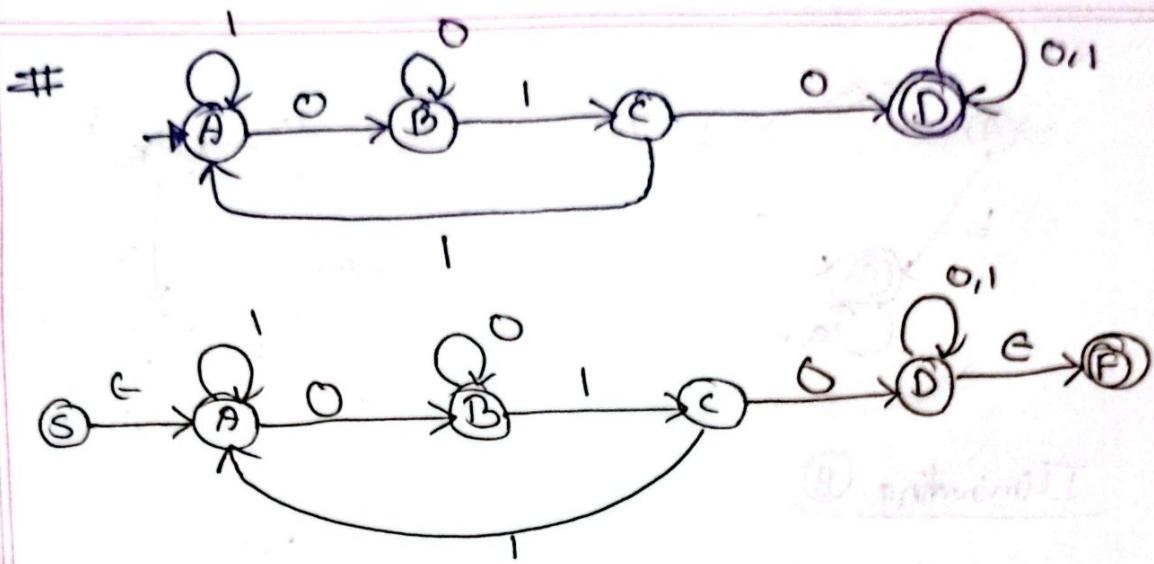


Eliminating ④



Eliminating ①

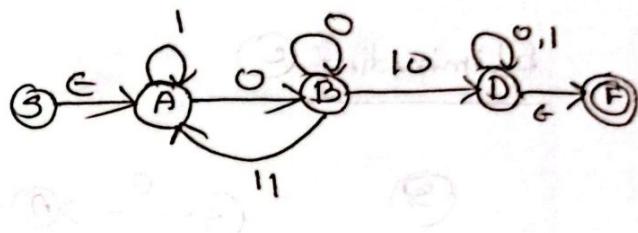




Eliminating C

$$B-C-D \rightarrow BD; 10$$

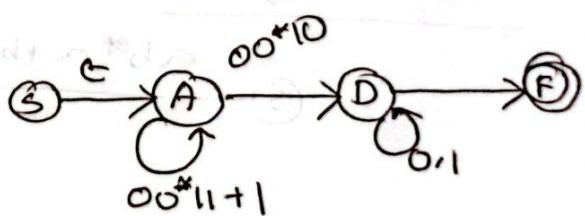
$$B-C-A \rightarrow BA; 11$$



Eliminating B

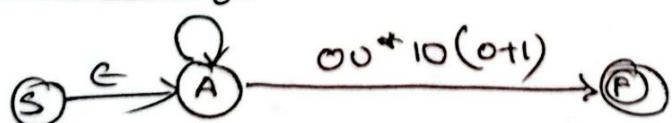
$$AB-D \rightarrow AD; 00^* 10$$

$$A-B-A \rightarrow AA; 00^* 11+1$$

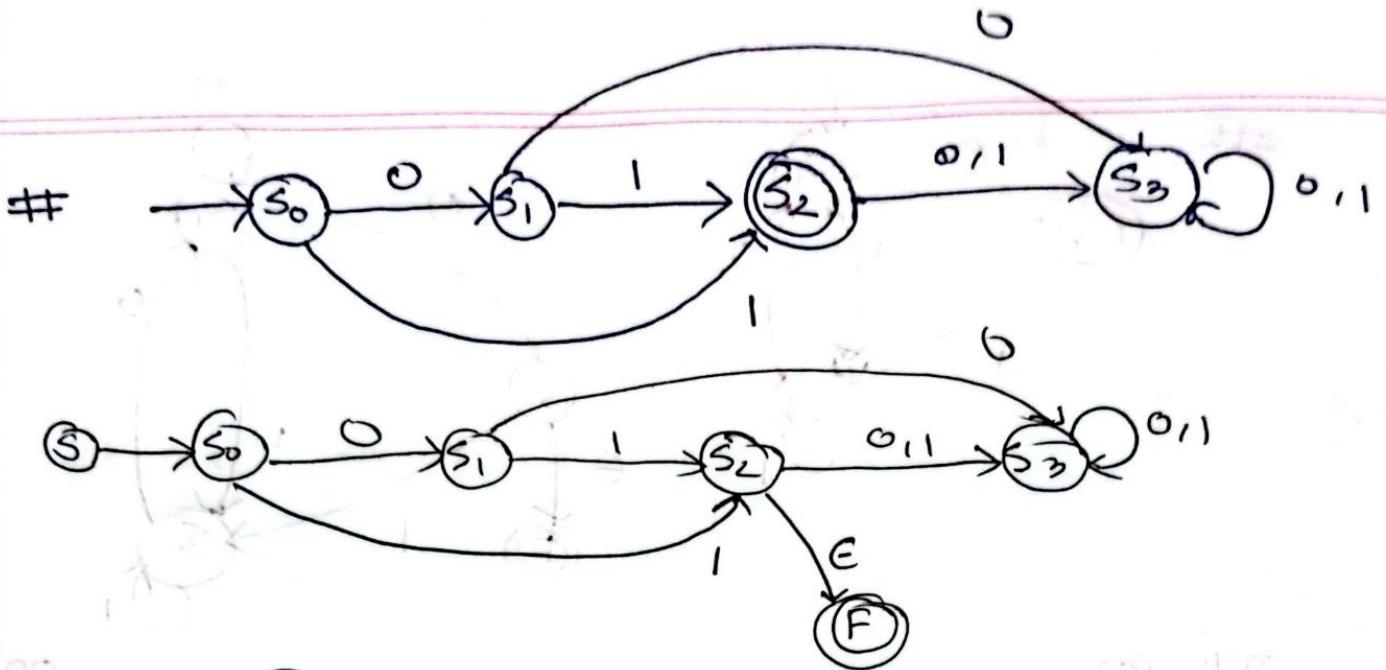


Eliminating D

$$00^* 11+1$$

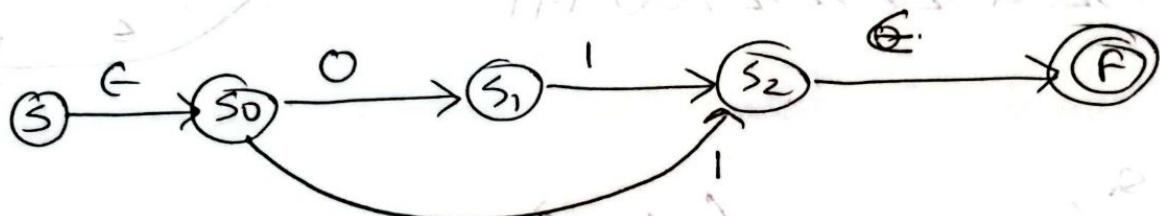


$$\therefore S \xrightarrow{(00^* 11+1)^* 00^* 10(0+1)} F$$



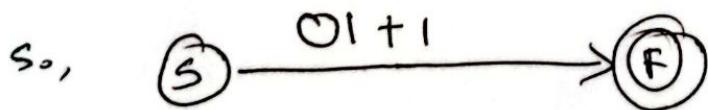
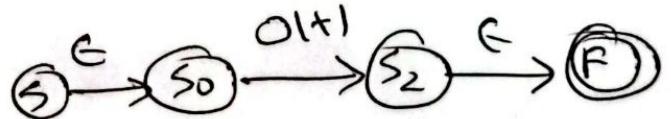
Eliminating S_3

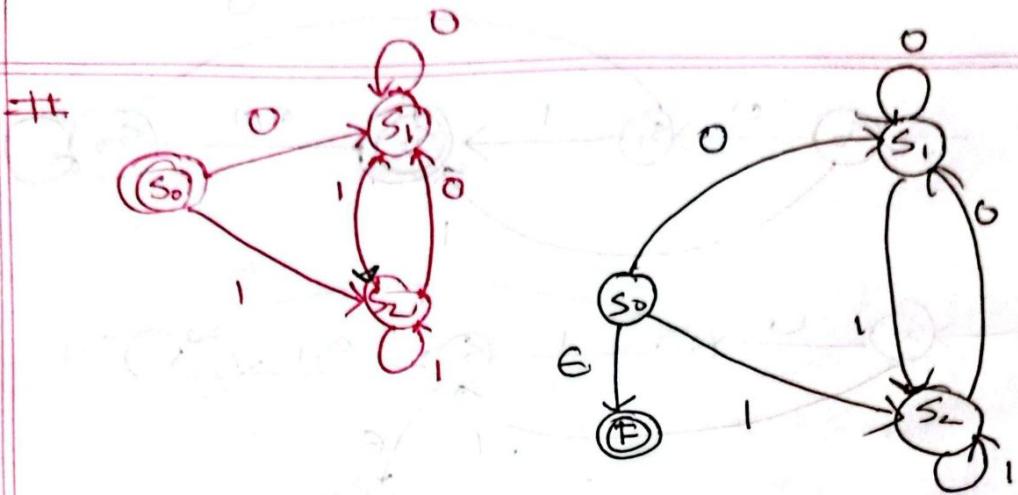
Trap state, S_0 , No work



Eliminating - S_1

$$S_0 - S_1 - S_2 \rightarrow S_0 - S_2; 01+1$$



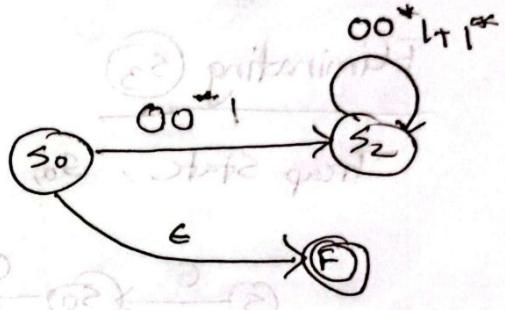


Elim- S_1

$$S_0 - S_1 - S_2 \rightarrow S_0 - S_2 ; 00^* 1$$

$$S_2 - S_1 - S_2 \rightarrow S_2 - S_2 ; 00^* 1 + 1^*$$

$$S_0, 00^* 1 (00^* 1 + 1) + E \rightarrow F$$



12 - posikování F

110 (zadní řada) - 10 = 00

1 + 10