

Lecture - 4

Introduction to Database System

Basic Definitions

Data

Recorded facts with implicit meaning

Database

Coherent collection of related data built for a purpose.

Mini-world

A segment of the real world stored in the database. [Ex → student grades and transcripts at a university]

Database Management System (DBMS)

Software for creating/maintaining database

Database System

DBMS software plus data (may include applications).

Types of Databases and Databases App.

Traditional Apps → Numeric and textual databases

More Recent Apps

i) Multimedia db

ii) Geographic Information Systems (GIS)

iii) Biological and genome db

iv) Data Warehouses

v) Mobile db

vi) Real time and active db

Managing Data

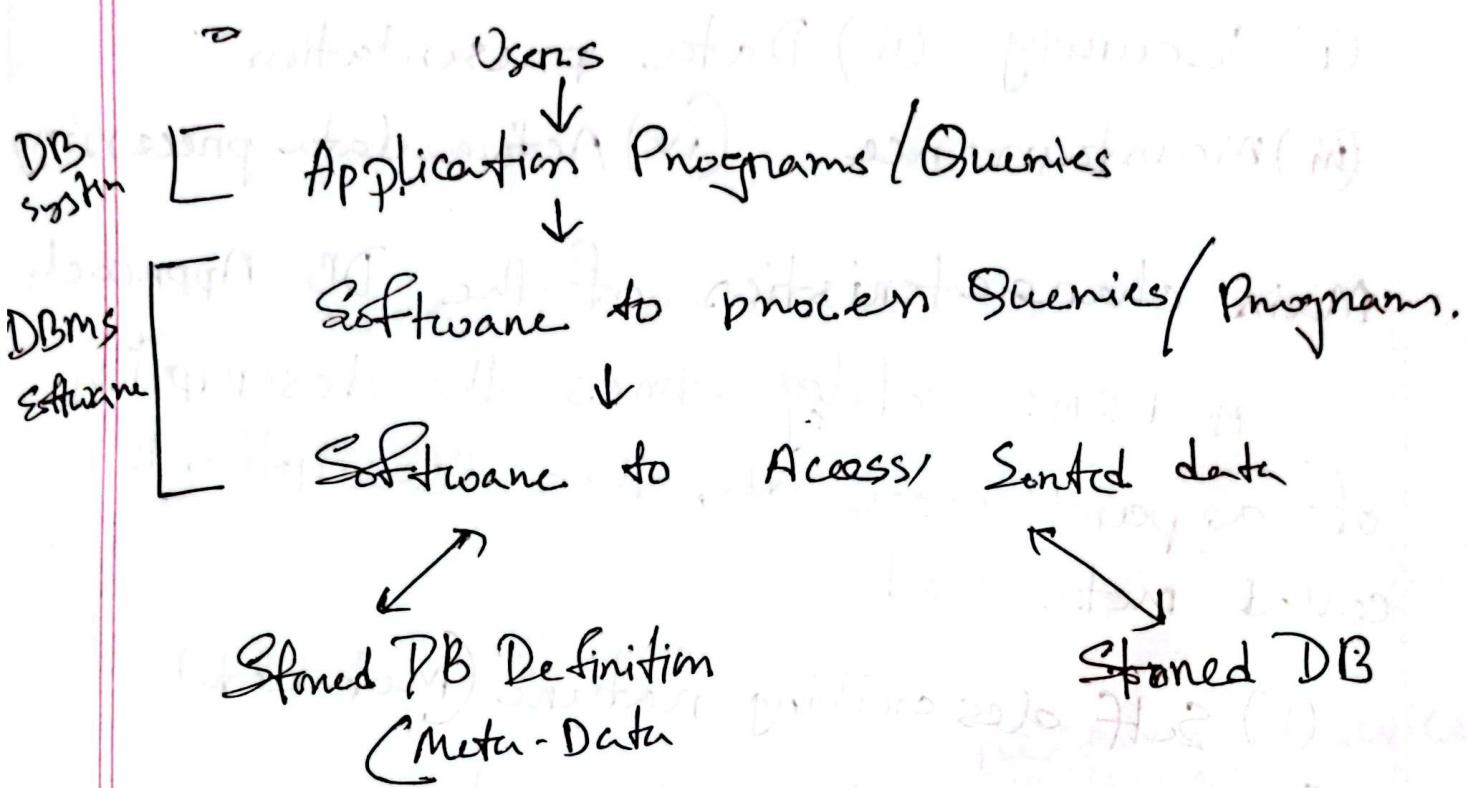
File Based Approach

Uses specific application programs, which performs services to end users, each managing its own data

Database Approach

Uses DBMS for data collection and sharing among multiple programs.

A simplified db system environment



DBMS Facilitates

- i Defining DB's
- ii Construct or load the initial db contents on a secondary storage medium
- iii Data retrieval, modification, web access.
- iv Multi-user processing

DBMS Functionalities

- (i) Security (ii) Data presentation
- (iii) Maintenance (iv) Active data processing

Main characteristics of the DB Approach

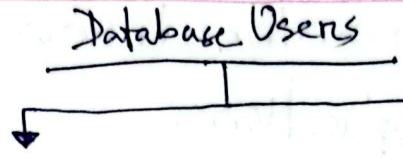
A DBMS catalog stores the description of a particular db, this description is called meta-data

- (i) Self describing nature (meta-data)
- (ii) Program data independence
- (iii) Data abstraction
- (iv) Support for multiple data views
- (v) Multi users transactions
- (vi) Concurrency control
- (vii) Restricting unauthorized access to data
- (viii) Providing storage structures for efficient query processing.

OLTP (Online Transaction Processing) is a major part of db applications.

Advantages

- (i) Data redundancy control
- (ii) Data sharing
- (iii) Restricted Access
- (iv) Query optimization
- (v) Providing backup and recovery services.
- (vi) Providing multiple interfaces to different classes of users.
- (vii) Representing complex relationships among data.
- (viii) Enforcing integrity constraints on the db.



Actors on the Scene

where db use হয়ে

workers behind the scene

DBMS দ্বাৰা and এবং
কম্পার্যুলেন্স দ্বাৰা

Actors on the scene

i) Database Administrators: DB কে access কৰিব

and নির্দলীয় কৰিব কৰতে, DB এর নির্মাণ,
পরিবর্তন কৰে Data গুলোকে উন্মোচন কৰতে।
They update and control db.

ii) Database Designers: They design db and
the structure. Db কে Logical and physical
design কৰিব - according to end-users

iii) Software Engineers: DB ফর্ম অপারেট কৰিব

এবং apps কে DB এর connection কৰিব

যোগ

#DB End Usages

They use the older form queries, responses and same of them update the db content. End-user can be categorized into:

(i) General Usages: They are for primary uses. It's active. It's primary uses. It's active.

(ii) Active: They are for secondary uses. It's active.

(iii) Interactive: They know what they want.

(iv) Sophisticated: They don't care about the details.

(v) Standalone: Mostly manifolds processed by user.

(vi) Ready-to-use packages app.

When not to use DBMS

(i) Too costly

- (i) High setup costs and Hardware needs.
- (ii) Overhead for security, recovery and concurrency.

(ii) Unnecessary

- (i) Simple, unchanging db.
- (ii) Q/R access to data by multiple user is not required

(iii) Infeasible (cost)

- (i) Limited storage (embedded systems)

(iv) Won't work

- (i) If real time speed needed (Phonesystems)
- (ii) Complex data (Genome db)
- (iii) If special tasks needed to be handled (GIS tools)

Data Models

A set of concepts to describe the structure of a db, the operations for manipulating these structures, and certain constraints that the db should obey.

#Categories

(i) Conceptual (High-level, semantic)

বাৰ প্ৰযোগ কৰিবলৈ সুস্থিৰ (বেলুন দেখা কৰা হৈলাম এবং [Entity-based or obj based মডেল]

(ii) Physical (low-level, Internal)

কৰিবলৈ নোবেলুন দেখা কৰা হৈলাম - new data is being scanned.

প্ৰযোগ কৰিবলৈ DBMS design & administration

প্ৰযোগ কৰিবলৈ কৰা হৈলাম - transaction (DBMS)

(iii) Implementation (Representational)

In between Conceptual and Physical Model.

Business Related DBMS বৈ উচ্চ (৩২)

12	3	03	foreign
11	2	02	Abnormal
10	1	01	Normal
			Class

- (iv) is called extension updated
- (v) db structure changes saying this db is collection of all data in the db
- (vi) current date is in the db of any time
- ① The current date is in the db of any time

Database State

Name	Count	DB#
Courses		

Name	Number	Class
Students		

- ② Schema is called extension
- ③ Changes very frequently
- ④ DB structure, data types, and the constraints on the db.
- ⑤ The description of a db. This includes description of the db structure, data types, and the constraints on the db.

Database Schema

- combine the description of data with the data values. [XML, key-value storage]
- (i) Self describing data Models

Three Schema Architecture

- Proposed to support DBMS characteristics of.
 - (i) Program-data independence
 - (ii) Support of multiple views of the data
- Has been useful in explaining subsystem organization.

(i) Internal Schema (Physical Level)

It is the physical storage of the data, detailing how data is stored in the computer.

(ii) Conceptual Schema (Logical Level)

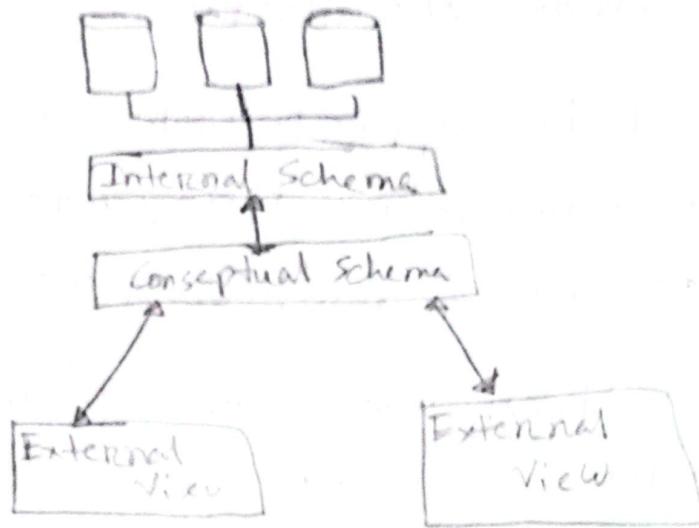
This level represents the overall structure of the entire db for the organization. It's like a blueprint showing all entities, relationships, and data types without concern for how it is physically stored.

(iii) External Schema (View Level)

The external schema is the user's view of the db. It allows different users to see only the data relevant to them.

Stonned DB

End-users



Data Independence

The ability to change the schema at one level without impacting the schema at the next higher level. 2 types.

(i) Logical Data Independence :

The capacity to change the conceptual schema without having to change the external schemas and their associated application programs.

[db to new column (vsy program, on datatype change in app v (various))]

(ii) Physical Data Independence

Internal schema නේ පරිගණක වලද යුතු conceptual schema නේ තොරතු ප්‍රතිඵල නී මාරුව. [db හෝ file නිස් නොවා ඇත්ති පරිගණක නොවා වා, එක් නොවා index create කර වා, conceptual schema unchanged වෙයි].

{DBMS Language}

(i) Data Definition Language (DDL)

To define db schema. DDL යුතු නො-db අනු පරිගණක ප්‍රතිඵල ප්‍රතිඵල Table → creation, dropping table, adding col and changing datatype. පරිගණක නොවැයි.

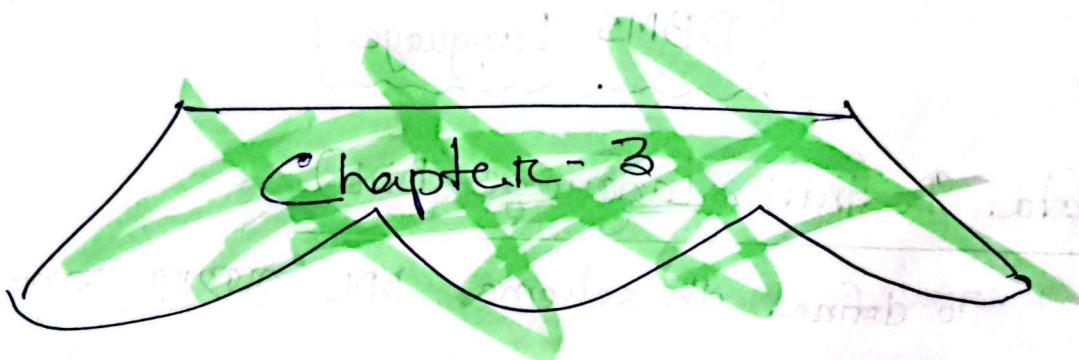
(ii) Data Manipulation Language (DML)

Used to manipulate data by inserting, deleting, updating and retrieving data. DML යුතුවන් general-purpose programming language (Java) හෝ DBMS embedded වාට්, 2nd මෘදු ප්‍රограм්ම (DBMS db පරිගණක නොවා).

Ex → Insert, Delete, Update, Select → Data

{ DBMS Interfaces }

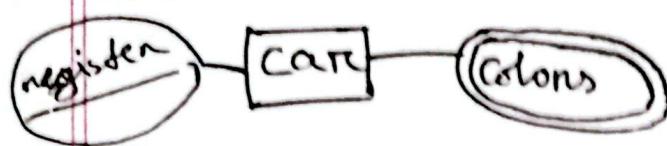
- (i) Stand-Alone Interfaces
- (ii) Programmatic n
- (iii) User-Friendly n
- (iv) Mobile n
- (v) Parametric n
- (vi) DBA-Specific n



ER Diagram - Entity Relationship Diagram

Main important things:

1. Entity → current object or db store
• Regular/Strong, Weak
2. Attributes
2. Relationships → Between 2 entity
Binary, ternary, n-ary / strong, weak
3. Properties of Entity.
• simple, composite, multivalued



Entity →

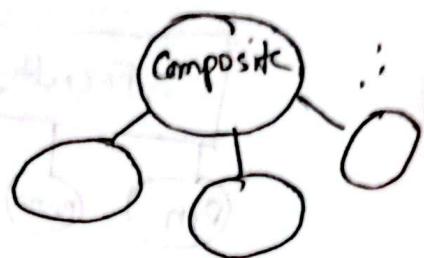
Strong

Weak

Attributes →

Single

Multivalued



Relationships →

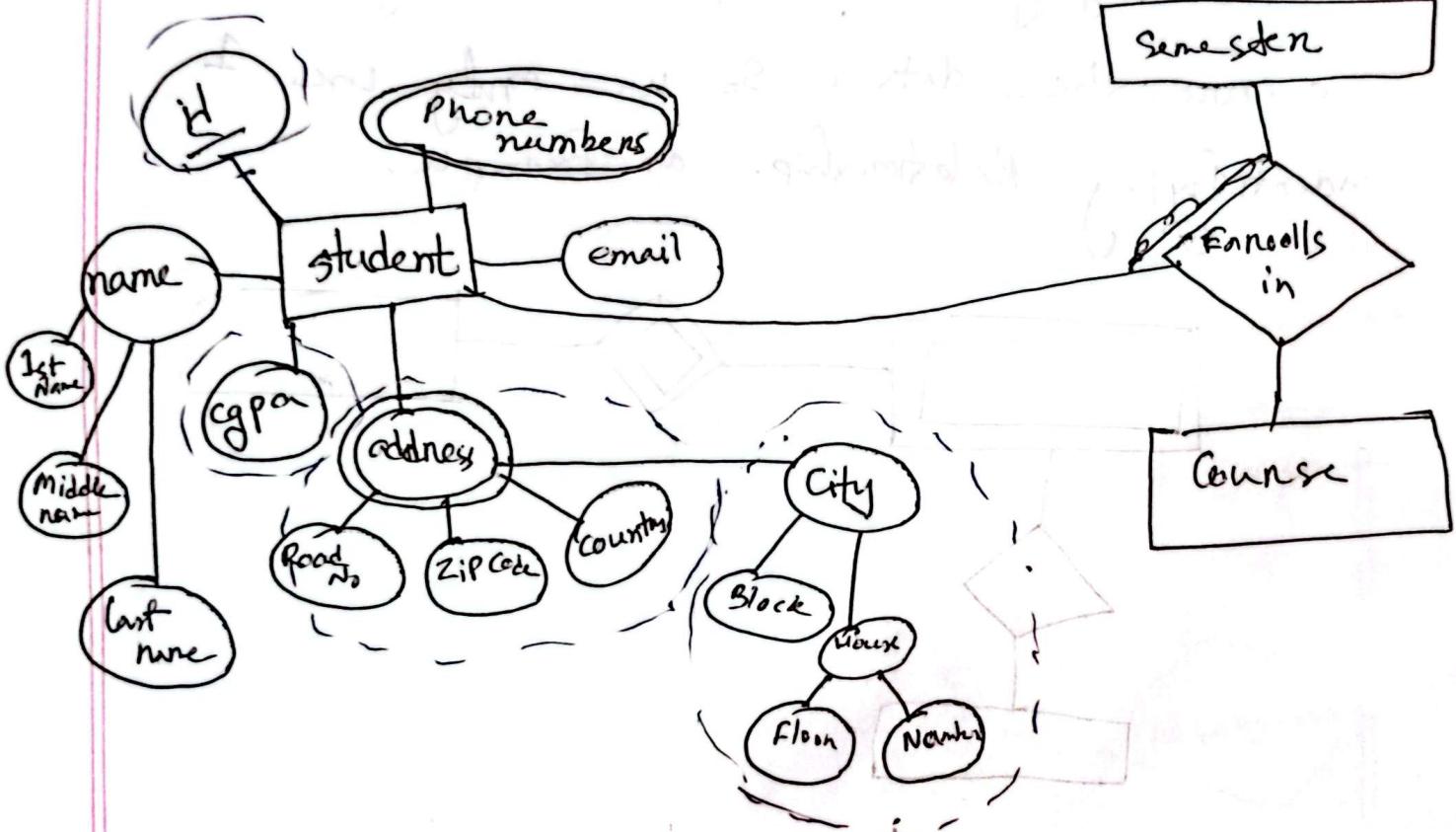
Strong →

Weak → MV →

Entity 1

Relationship name

Entity 2

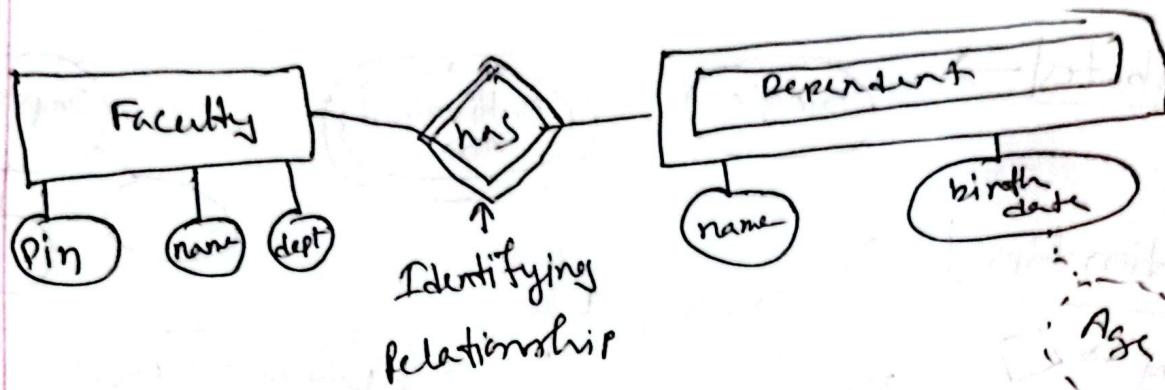


Value of address for an add
→ address → add

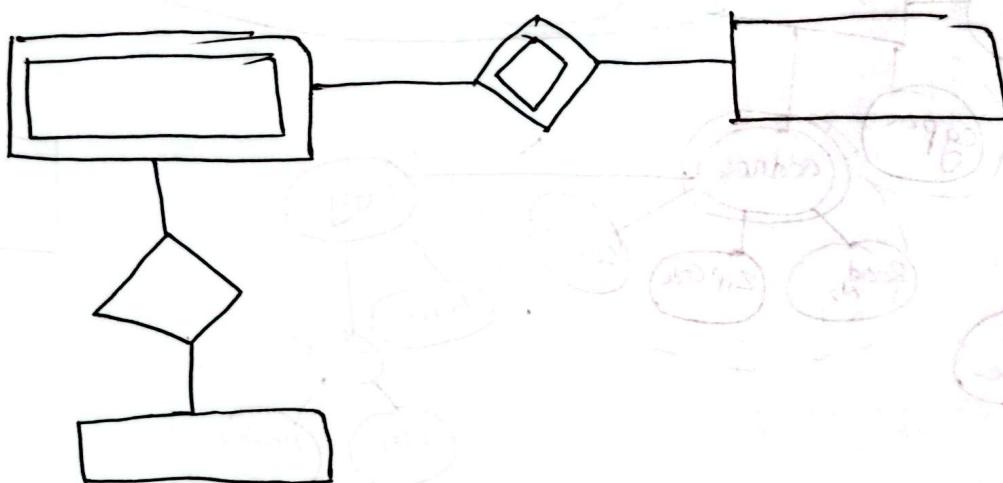
Key attribute
→ unique.
Underline
id.
Underline

④

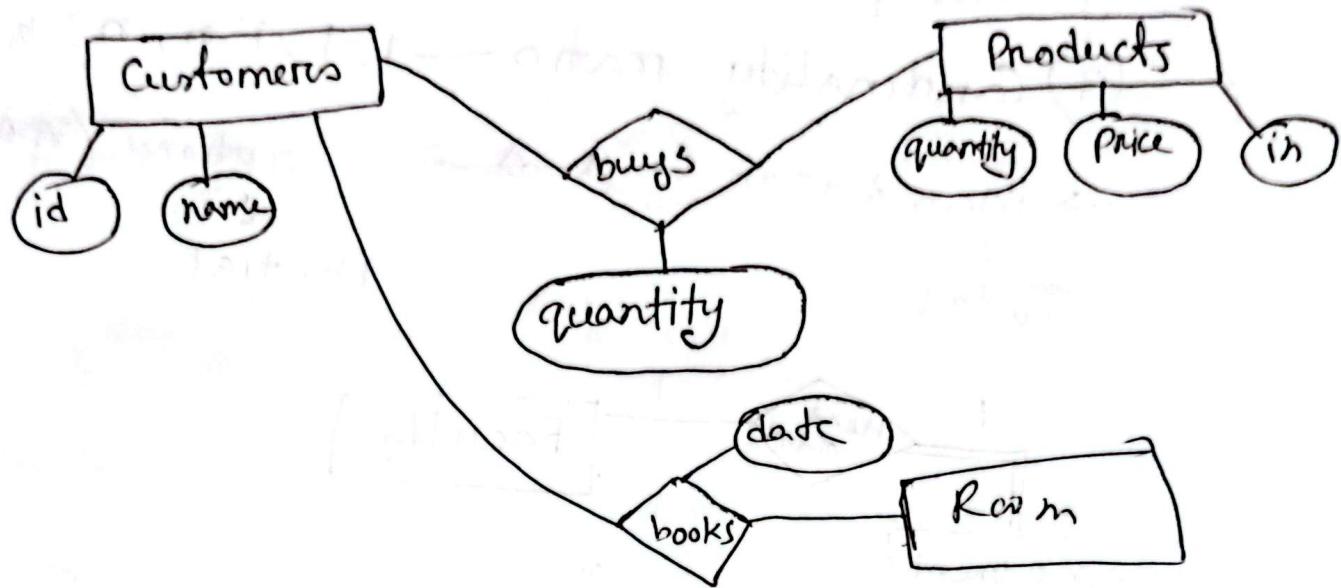
Weak Entity → user can entity via partial unique attributes



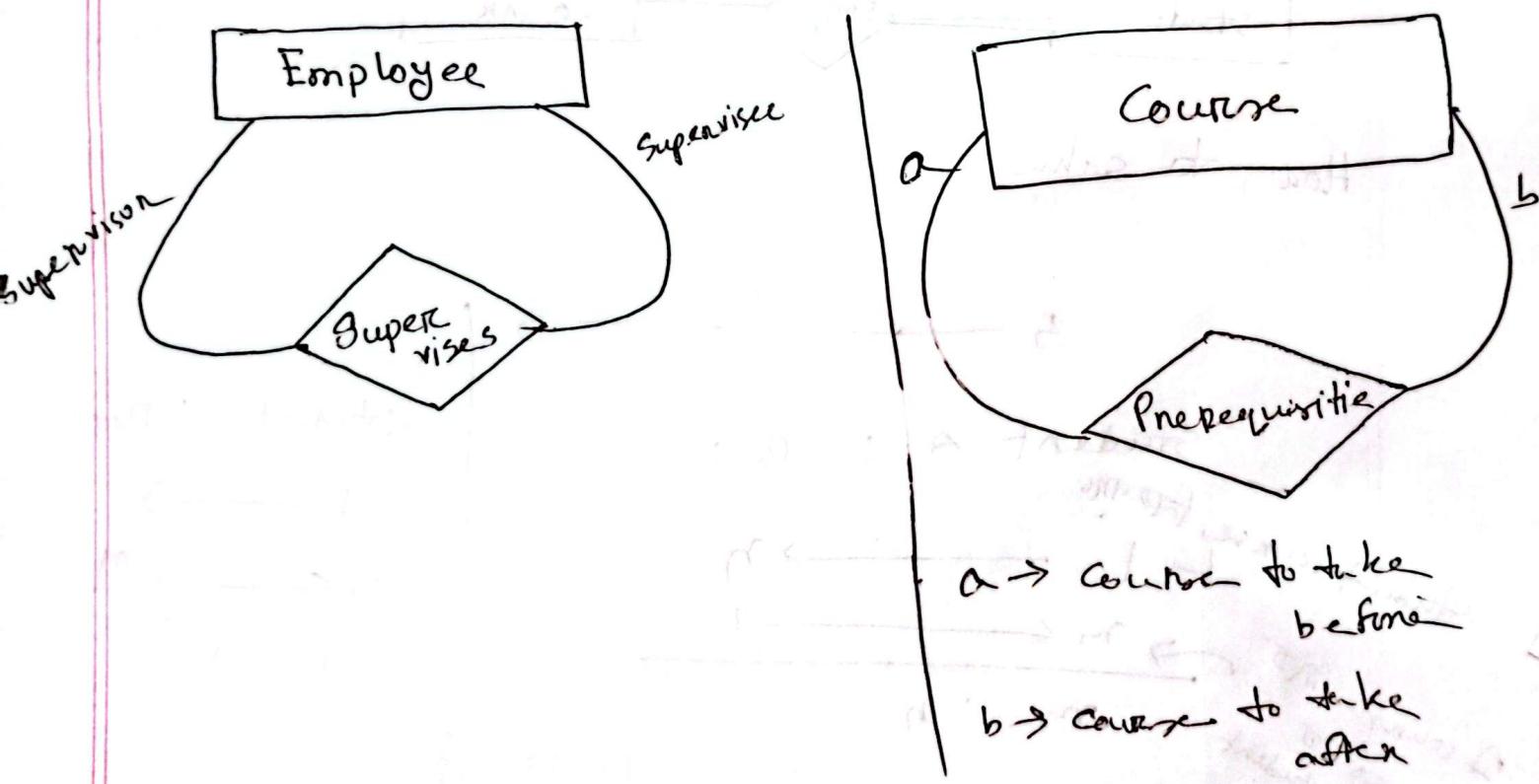
Identifying Relationships kinda assigns a id to secure the data. So we only use 1 identifying Relationship, as Example



A relationship can have attributes too.



Recursive Relationships



Constraints

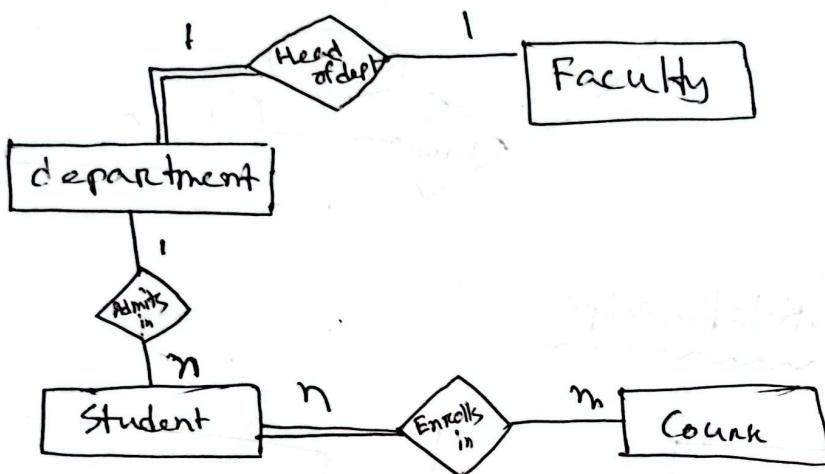
(i) Participation constraint → Relationship 1 for mandatory or optional

(ii) Cardinality ratio — 1:1, 1:n, n:m

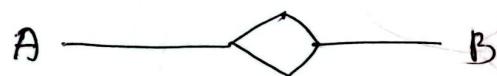
→ mandatory ($\rightarrow \square$) ; optional ($A \leftarrow B$)

Total

Partial



How to solve,



Student A : B course.

Student : Dept

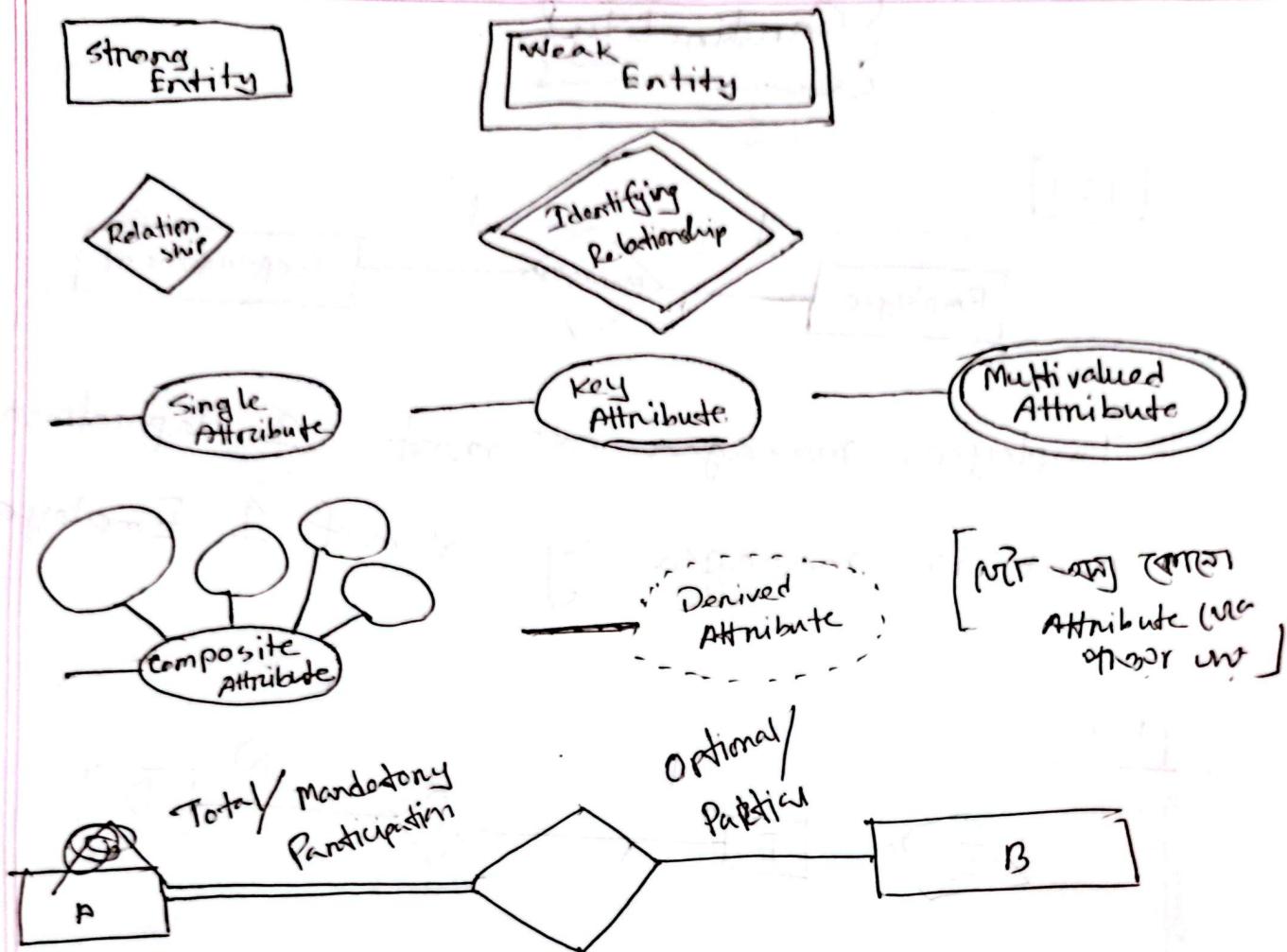
1 → 1

n → n

n : 1

1 student course form
n course form
m : n student

Summary of notation for ER diagram



ER Diagram Solved

Cardinality

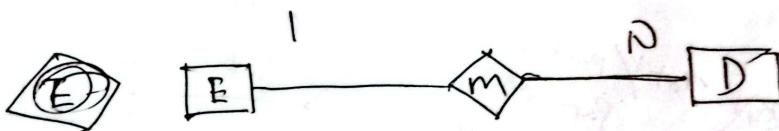
[1:1]



Employee manages at most 1 Department

Dept is managed by at most 1 Employee

[1:n]



E manage many ~~procedures~~ dept.

dept is managed by only 1 E.

[n:1]



From E manages only 1 Dept

Dept managed by many E

[M:N]

[E]

E manag.

Cours

* Student

* Student

* Course

M:N



E manages many ; Dept managed by many

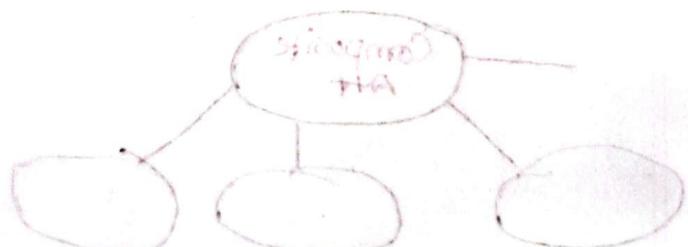


* Student dependent on course

* Student can must select most 2nd !

* Course student at first

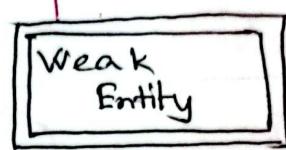
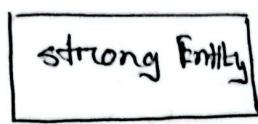
{student 1st
course 2nd
(course on unit)}



10 Ans

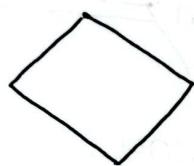
ER Model

Entity



→ No long Unique/
key attribute

Relationships



Weak Entity
→ Relationship.
It kinda assigns
a id to secure
the data. But
we only use 1
identifying relationship

Attributes

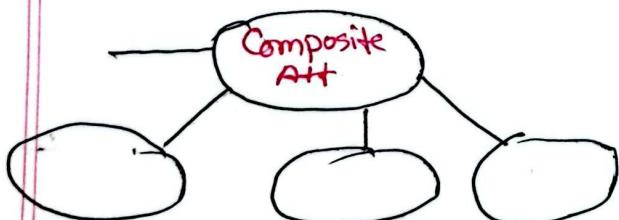
single attribute

Key Attribute

Multivalued
Attr

Unstructured attr
using string

→ Only 1 for (fixed) person



Derived Attr

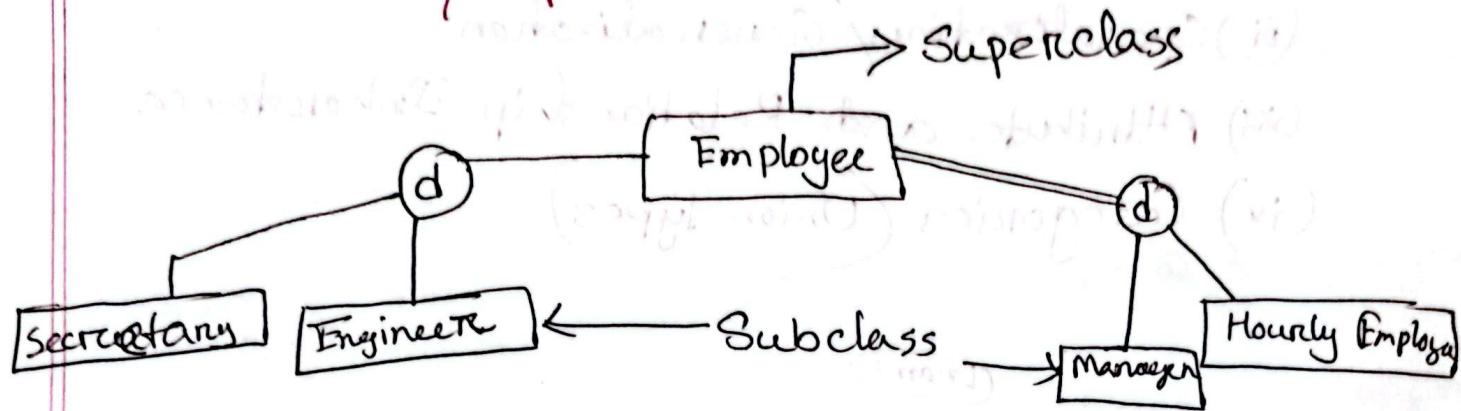
Att can be derived from

Chapter - 4

(ER → Extended ER)

ER → Extended ER

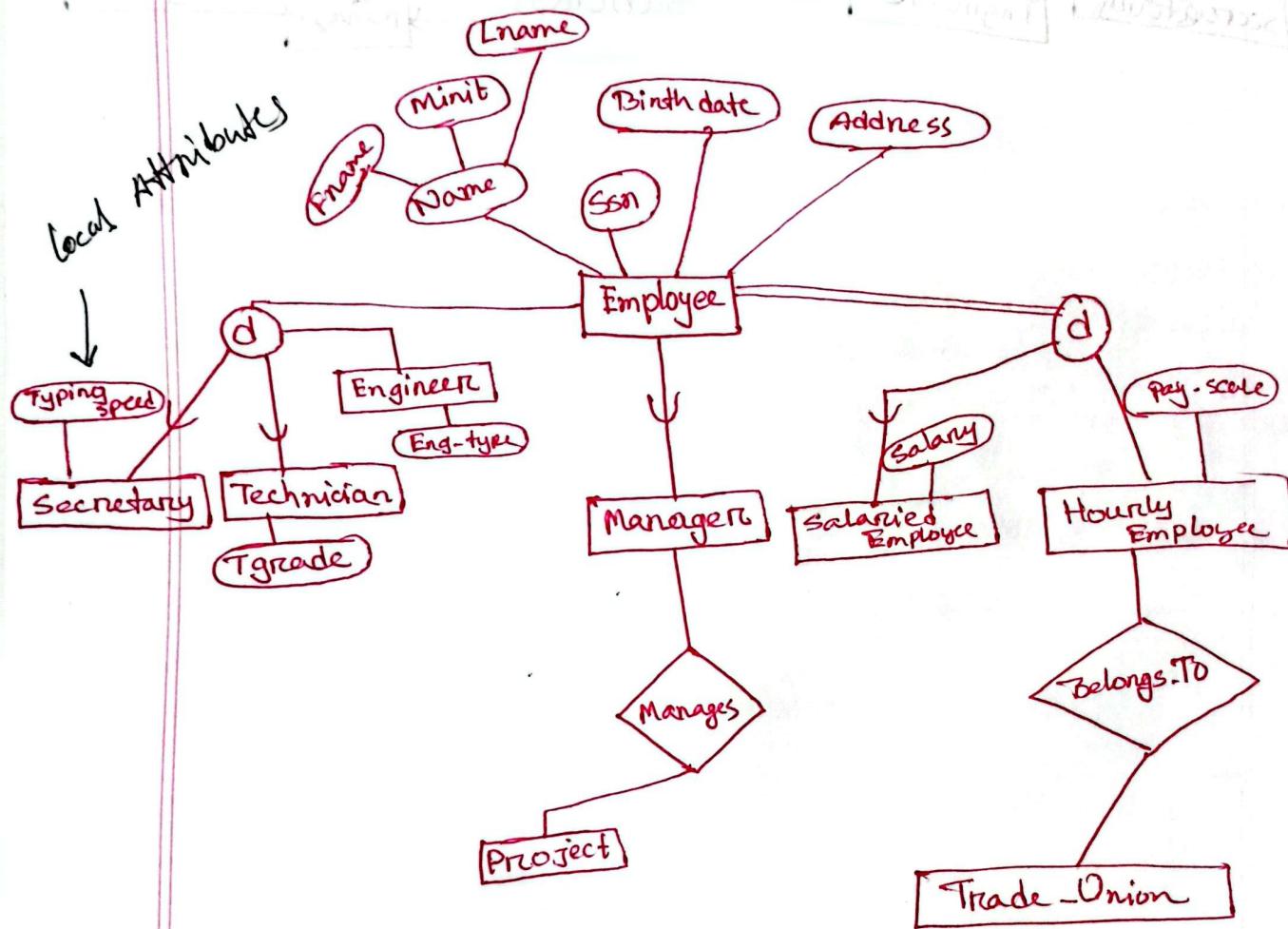
(i) Subclass / Superclass



Chapter - 4

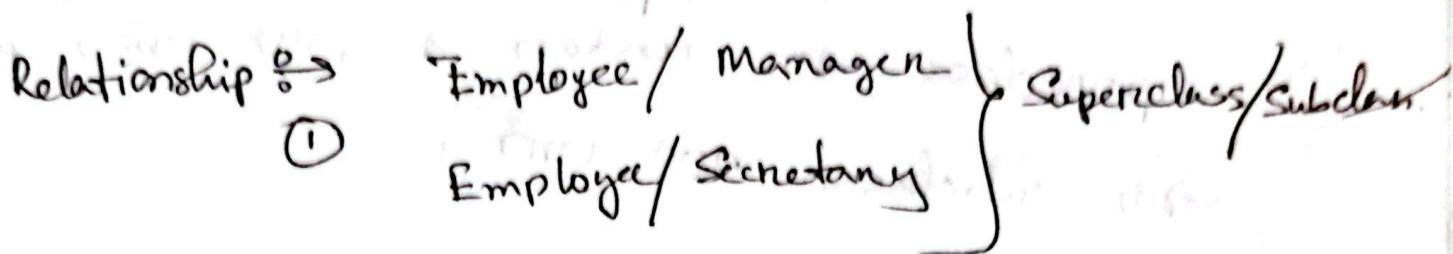
EER → Extended EER

- (i) Subclasses / Superclasses
- (ii) Specialization / Generalization
- (iii) Attribute and Relationship Inheritance
- (iv) Categories (Onion Types)



Here, Employee is the Superclass

Manager, Secretary etc are Subclass.



ii IS-A Relationship. Secretary IS-A Employee.

[Inherits] \rightarrow

iii Superclass \Rightarrow Entity \Rightarrow subclass \Rightarrow (optional)
[But not mandatory]

iv An entity in the superclass.

v For any entity to exist in the db, superclass

member \Rightarrow (optional)

vii your entity \Rightarrow superclass \Rightarrow optional

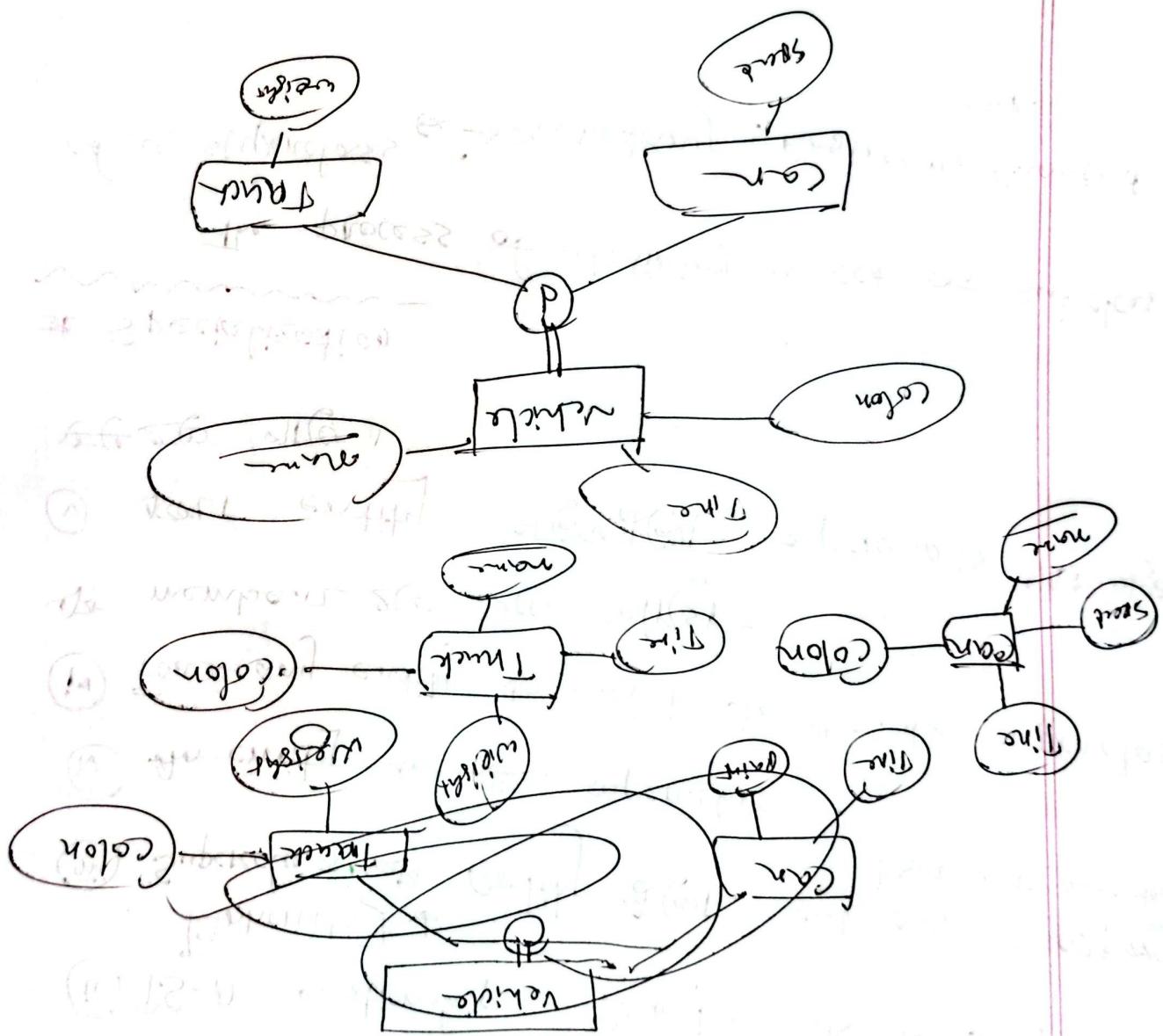
\Rightarrow (optional)

* Specialization

The process of defining a set of subclasses of a superclass

Ex \rightarrow Secretary, Engineer, Salaried

emp - -



superclasses are one
 subclasses are many →
 2nd tier superclasses
 3rd tier superclasses
 4th tier superclasses
 ↓
 Create object

Types of Specialization

Predicate Defined

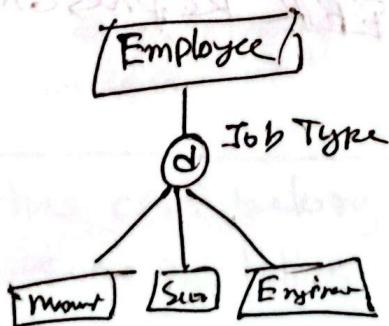
Membership is determined by a condition on rule. $\exists x \rightarrow \text{EMPLOYEE} \text{ Belongs to } \text{ENGINEER}$ if Jobtype = "Engineer".

Attribute defined

Membership is based on a specific attribute in the superclass. $\exists x \rightarrow \text{EMPLOYEE}$ specialization based on JobType.

User-Defined

Membership is manually determined by the db user. $\exists x \rightarrow \text{PROJECT}$ assigned to CRITICAL-PROJECT based on user decision.



#2 basic Constraints can apply to a specialization/generalization

(i) Disjointness Constraint

(ii) Completeness Constraint

(i) Disjointness Constraint

1. Disjoint Specialization (d)

- An entity can belong to only one subclass.
- Ex → • A student can Either be GRAD or Under grad.
- EER Representation → (d)

2. Overlapping Specialization (o)

- An entity can belong to multiple subclasses.
- Ex → • An Employee can be both Manager and Tech.
- EER Representation → (o)

(i)

(ii)

(iii)

(iv)

(ii) Completeness Constraint

1. Total Specialization (=)

- Every entity must belong to at least one subclass.

Ex → • Vehicle must be a car or truck

- EER Representation : Double line (\equiv)

2. Partial Specialization (-)

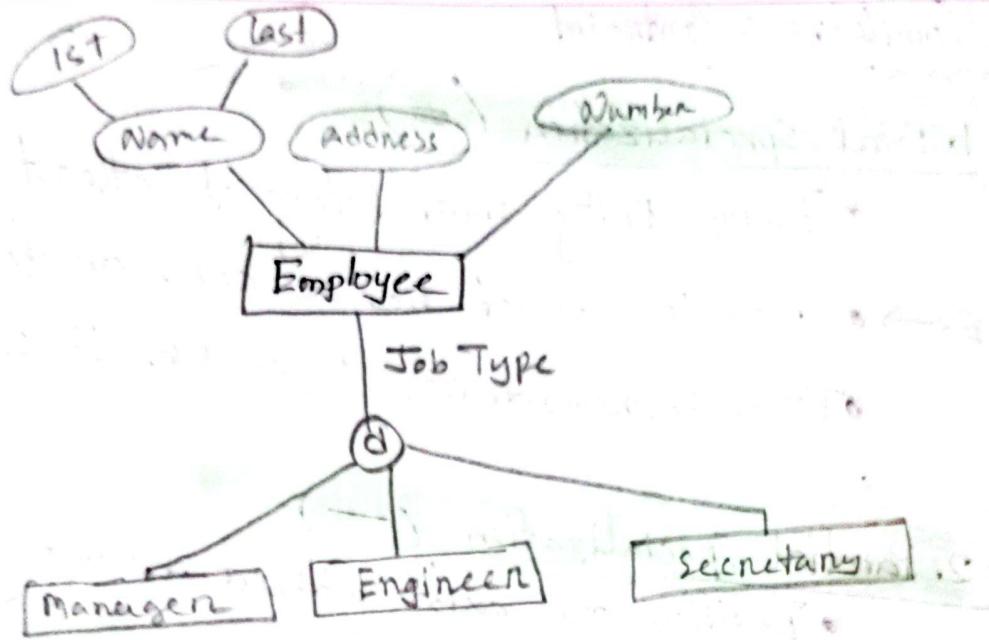
- Entities can exist without belonging to any subclass

Ex → • Some Employees may be neither Eng nor Tech.

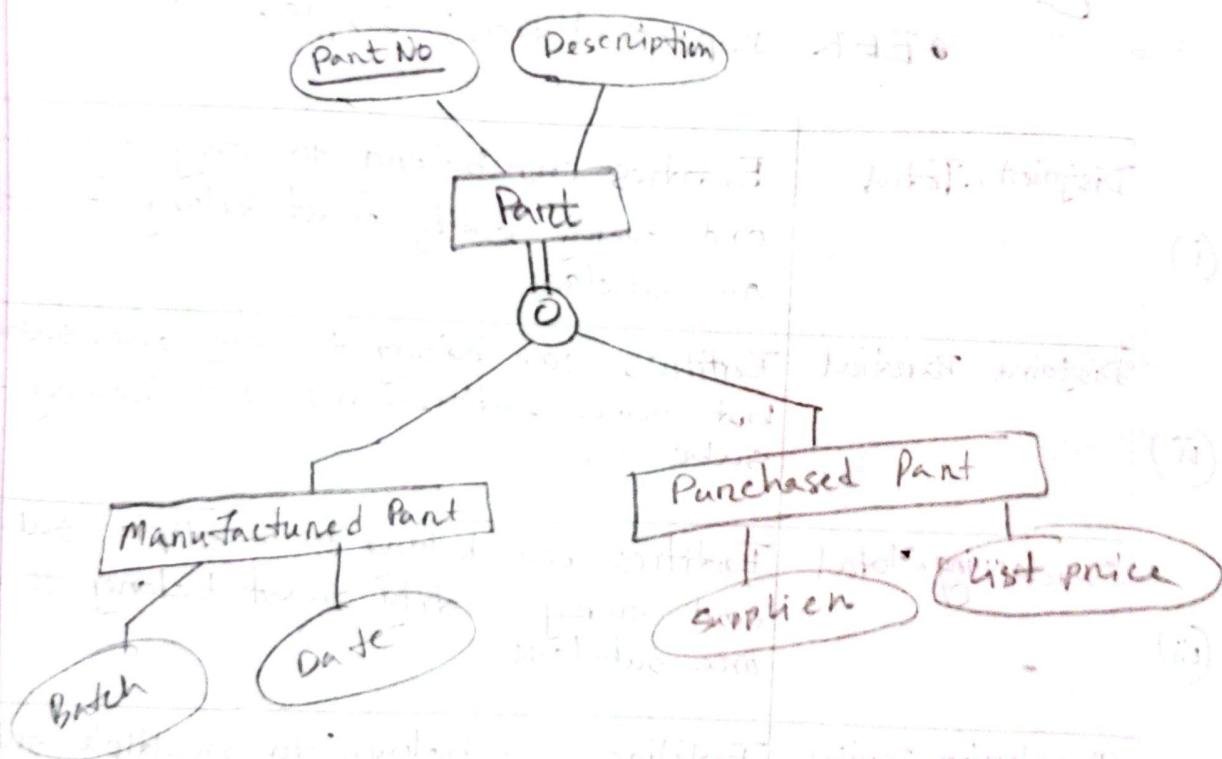
- EER Representation: Single line (-)

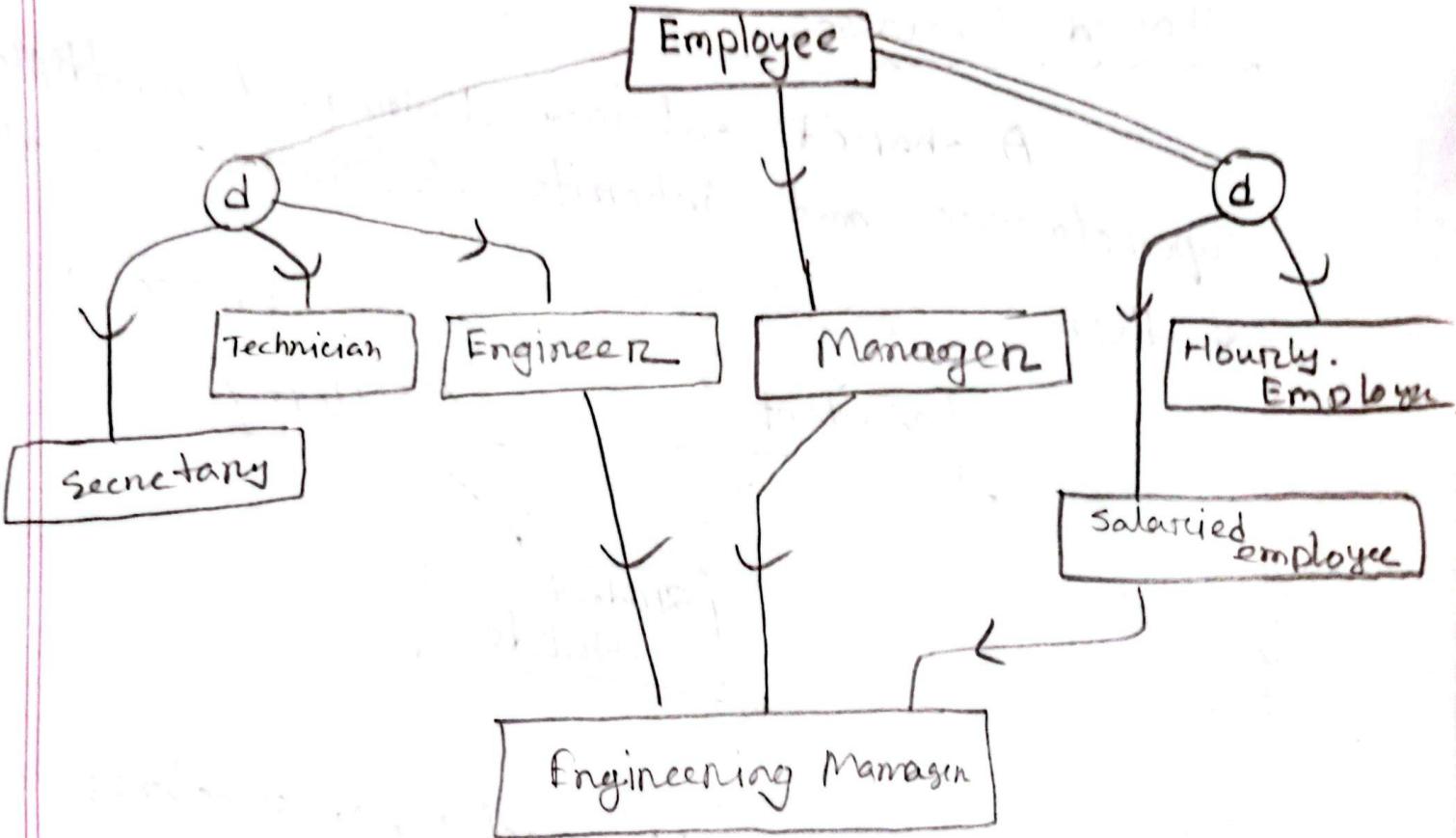
(i)	Disjoint, Total	Entities can belong to only one subclass, and every entity must belong to at least one subclass
(ii)	Disjoint, Partial	Entities can belong to only one subclass, but some entities may not belong to any subclass
(iii)	Overlapping, Total	Entities can belong to multiple subclasses, and every entity must belong to at least one subclass
(iv)	Overlapping, Partial	Entities can belong to multiple subclasses, and some entities may not belong to any subclass.

(i)



(ii)





Hierarchies, Lattice & Shared Subclasses

Hierarchy

Subclasses have only one superclass, forming a tree structure. (Single Inheritance)

EMPLOYEE → MANAGER → ENGINEERING MANAGER

Lattice

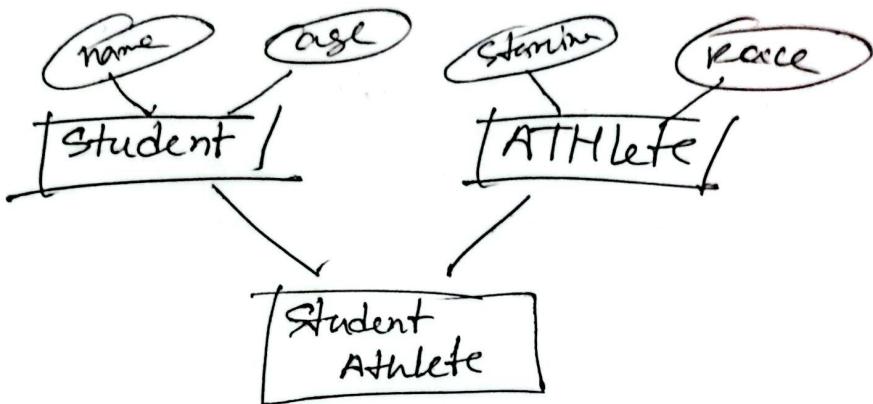
Subclasses can have multiple superclasses.

A subclass can have more than one parent superclass,

Person / Employee
/ Engineering Manager

Shared Subclasses

A shared subclass belongs to multiple superclasses and inherits attributes from all of them.



- In a lattice or hierarchy, a subclass inherits attributes not only of its direct superclass, but also of all its predecessor superclasses.

Specialization

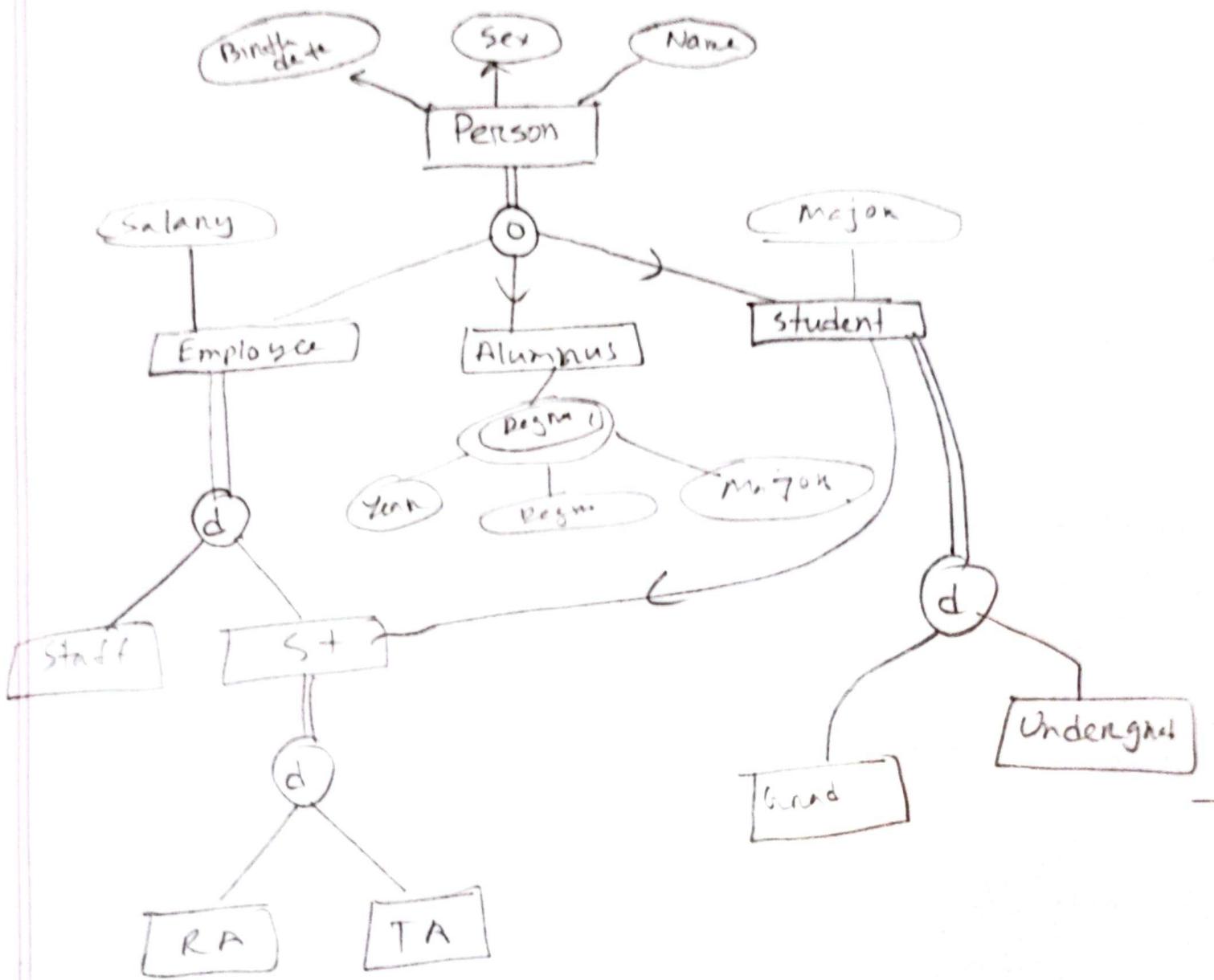
Start with one entity and break it into smaller, specific subclasses. This is a top-down conceptual refinement process.

Generalization

Start with many entities and combine them into a common superclass. This is a bottom-up conceptual refinement process.

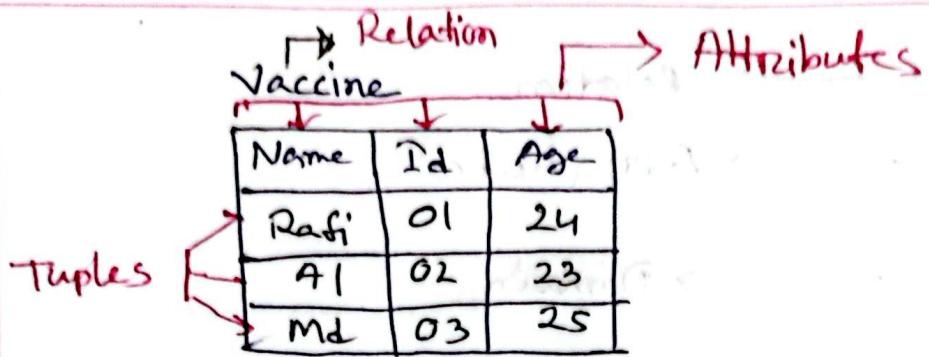
In practice, both processes are often used together.

Example → University Database



Chapter-5

DM and RDC



Schema of a Relation $\rightarrow R(A_1, A_2 \dots A_n)$

Vaccine (Name, id, age)

- Each att. has a domain on a set of valid strings (Numbers Fix or Var or [size])
 - The domain may have a data-type and/or format.
 - Tuples is an ordered set of values. ($\langle \dots \rangle$)
 - $\langle "Rafi", 01, 24 \rangle$
 - The relation state is a subset of the Cartesian product of the domains of its att.
 - $R(A_1, A_2)$ is a Relation schema.
- $$\text{dom}(A_1) = \{0, 1\} \quad \text{dom}(A_2) = \{a, b, c\}$$
- so, $\text{dom}(A_1) \times \text{dom}(A_2)$ is all possible combinations.

$\{ \langle 0, a \rangle, \langle 0, b \rangle, \langle 0, c \rangle, \langle 1, a \rangle, \langle 1, b \rangle, \langle 1, c \rangle \}$

Table \longrightarrow Relation

Col. Header \longrightarrow Attribute

All possible col. values \longrightarrow Domain

Row \longrightarrow Tuple

Table Definition \longrightarrow Schema of a relation

Populated Table \longrightarrow State of a Relation

• The tuples (rows) in a table have no specific order

• The attr. and their values are in a fixed order

Characteristics of Relations

Atomic Values Each value in a tuple is indivisible.

Valid Values Each value must come from the domain of its own att.

Tuple Format For a row $t = \langle v_1, v_2, \dots, v_n \rangle$ each v_i corresponds to a specific column A_i

Null Values Special values used for unknown or inapplicable data.

Notation

- $t[A_i]$ or $t.A_i$ Refers to the value of col. A_i in row t .
- $t[A_1, A_2, \dots, A_n]$ Refers to specific values of selected cols in row t .

Relational Integrity Constraints

→ 3 main types in the relational model

(i) Key Constraint

Each row in a table must be unique, identified by a key (like primary key)

(ii) Entity Integrity Constraints:

The primary key of a table can't be null, ensuring each record is uniquely identifiable.

(iii) Referential Integrity Constraints

($FK = PK$)
if (FK not in PK
null is allowed) Foreign keys must match a valid primary key in another table or be null, ensuring relationships between tables remain consistent.

Domain Constraint

(v2) col v have one datatype

Example 20:

student having one student table

Key Constraints

→ Attributes

Superkey : set of keys that are unique.

Key : KESK ; minimal superkey.

if any other superkey is set then consider

the attr which is not part of superkey

or min attr of key.

(Ex) $\{ \text{course_id}, \text{name} \}$; course_id is part of

course_name

course_id is unique or superkey minor.

one attr key

These can form database with unique or superkey minor.

and attributes present in base table projections

the above projection has been done

so we can say that it is a valid projection

unique filter is taken from first significant

projection and so that number which is not

present in previous will be inserted again

Key Type	Purpose	NULL _{allow.}	Example
Primary key	Uniquely identifies each row in the table.	NO	Student-ID in Student table
Foreign key	References the primary key in another table. (Not unique)	Yes, if no match	Student-ID in Course table
Candidate key (candidate for pk)	Potential key that could uniquely identify rows. <small>(Unique)</small>	Depends on use	{Serials & student}

Every primary key is a candidate key,
but not vice versa.