

Compiler → The mother of all program.

It has to be 100% accurate & bug-free from the beginning. Otherwise all the programs will get crashed.

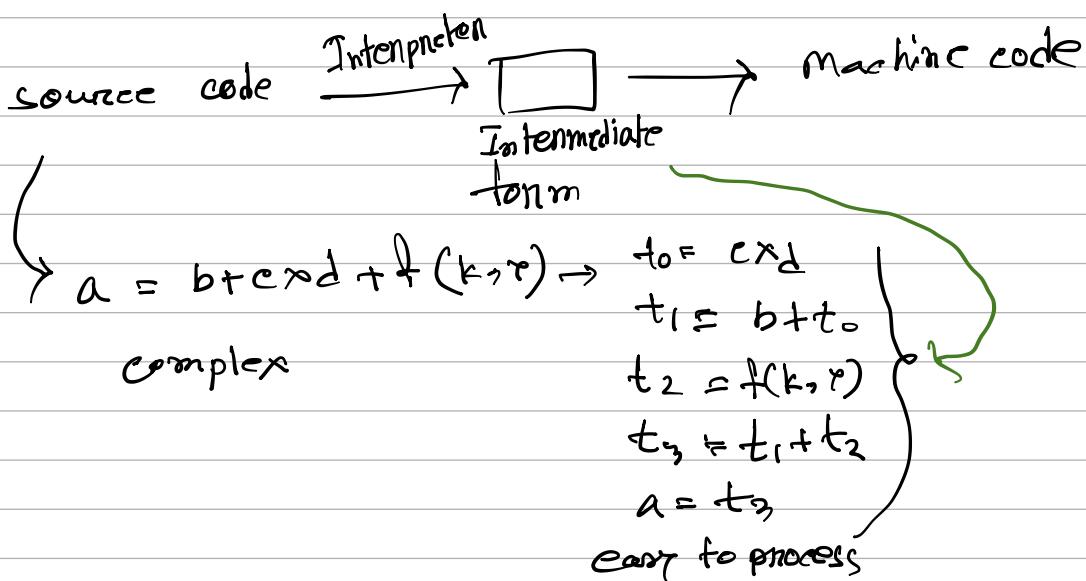
How did we write the computer itself?

What compiler compiled the first compiler?

↳ First compiler was written in machine code.

Only high-level languages need to be compiled.

What is the difference between compiler & interpreter?



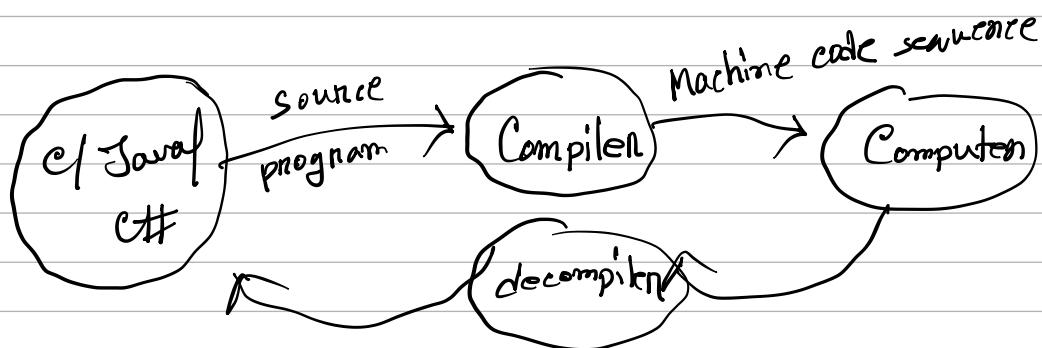
Interpreter is slower. It's more controlled environment.



Cuda → special version of C.

↳ Only executable by GPU not CPU.

Here compiler plays the role of translator.



✓ words → basic unit of meaning

✓ Grammar → How we arrange the words

✓ Semantics/meaning → Rice ate mice } words have meaning
but semantic is incorrect

Sequence of characters → tokens

↓
Syntax checking → Syntax " " " "

↖ Lexical analysis does the task

↓
Semantic Analysis

↓
Machine code generation

This should be the case. But sometime there is space for improvement.

$$\begin{aligned} a &= 3 + 10 \\ b &= a + 10 \end{aligned} \quad \left. \begin{array}{l} \hphantom{a = 3 + 10} \\ \hphantom{b = a + 10} \end{array} \right\} \rightarrow \begin{aligned} a &= 13 \\ b &= 23 \end{aligned}$$

This is done by a intermediate state called.
Machine independent optimization.
code

Advanced compilers add another step to the end which is "machine dependent code optimization."

Tokening



Syntax checking



Semantic analysis



Machine independent
code optimization



Machine code



Machine dependent
code optimization

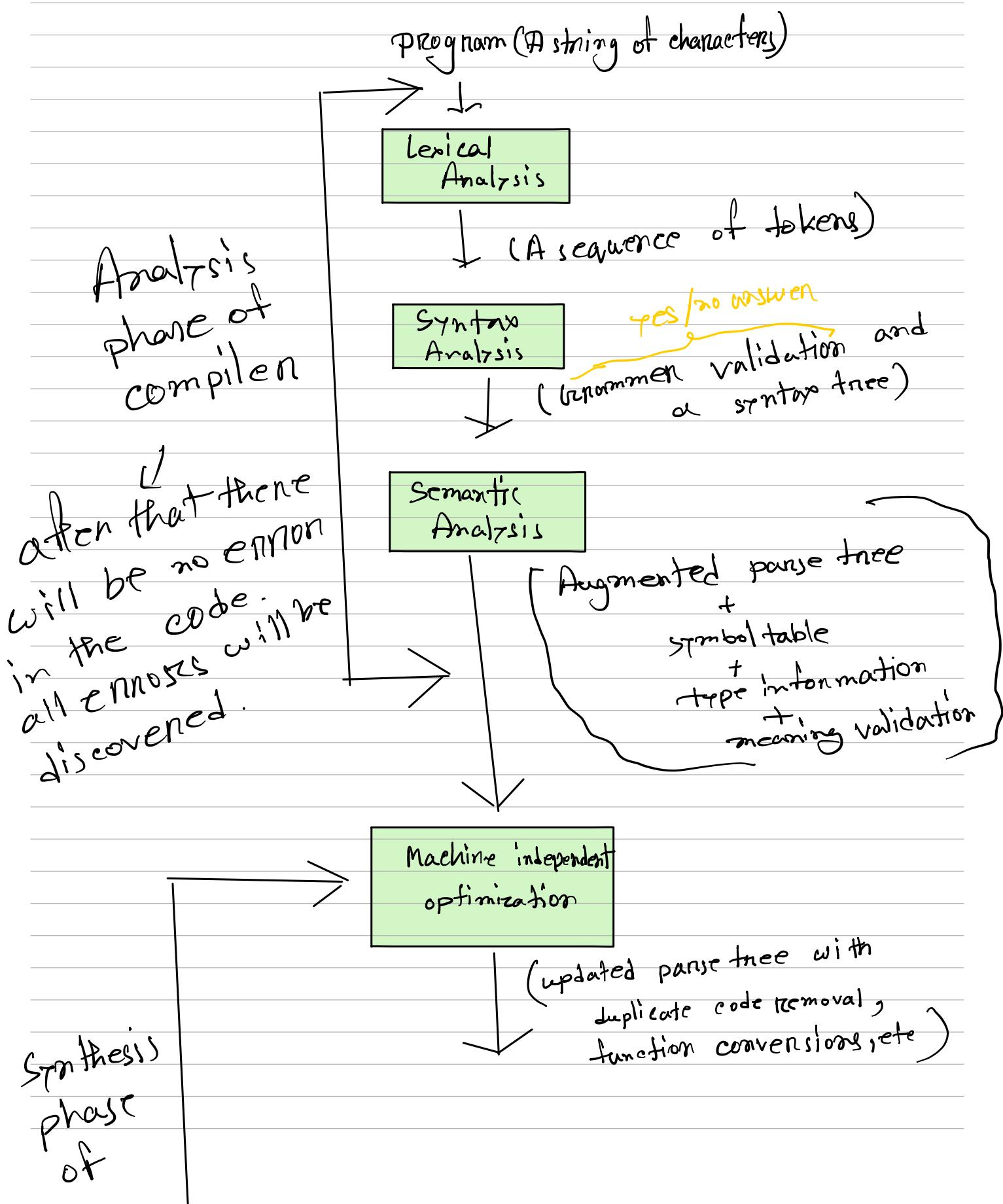


depends on the
architecture of
machine.

Lecture 2

Regular Expression \rightarrow DFA

In case of compiler, the translation process is one-way.



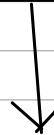
the
compiler

Machine code Generation



(sequence of machine instructions?)

Machine dependent code optimization



(updated sequence of instructions like
instruction conversions,
instruction swaps to avoid
pipeline stalls etc.)

Synthesis → breaking down complex steps in simpler
steps.

* Token

- keyword (if, while, def, for, etc)
- identifier (names you give to variables, functions, classes, etc)
- numbers (integers, fractions | numbers)
- Delimiter (;, :, , " ", etc)
- strings (whole thing inside quotation considered as a token)
- operators (+, *, -, /, etc)

* Lexical analysis does the job of taking sequence of words & characters, and returning the sequence of tokens.

* for each token we have a regular expression.

$$\text{alphabet } \{a, b\} = \{$$

three operations: (i) concatenation
(ii) kleen's closure

(iii) OR

* RE of identifier:

$$(\ldots [a \ldots z \ A \ldots Z])$$

* integer

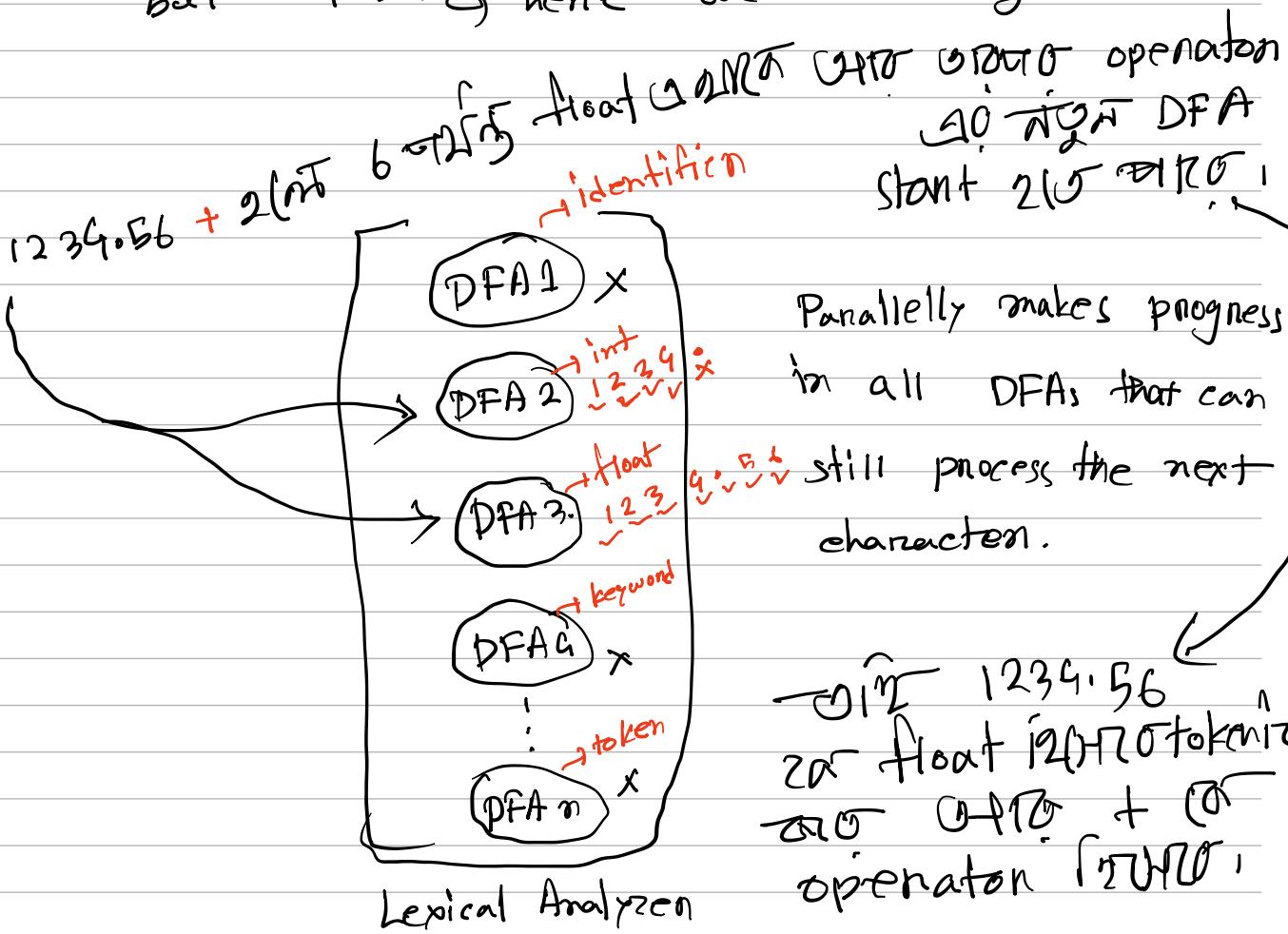
$$[0 \ldots 9] [0 \ldots 9]^*$$

0 | 1 | 2 | 3 | ...

Token $\xrightarrow{\text{we define}}$ RE $\xrightarrow{\text{converted into}}$ DFA

→ 1234
 * we start from first character and match with regular expressions by constructing DFA. If we end up in final state we can conclude that this can be identified as integer token.

but 1234_c here we can disagree

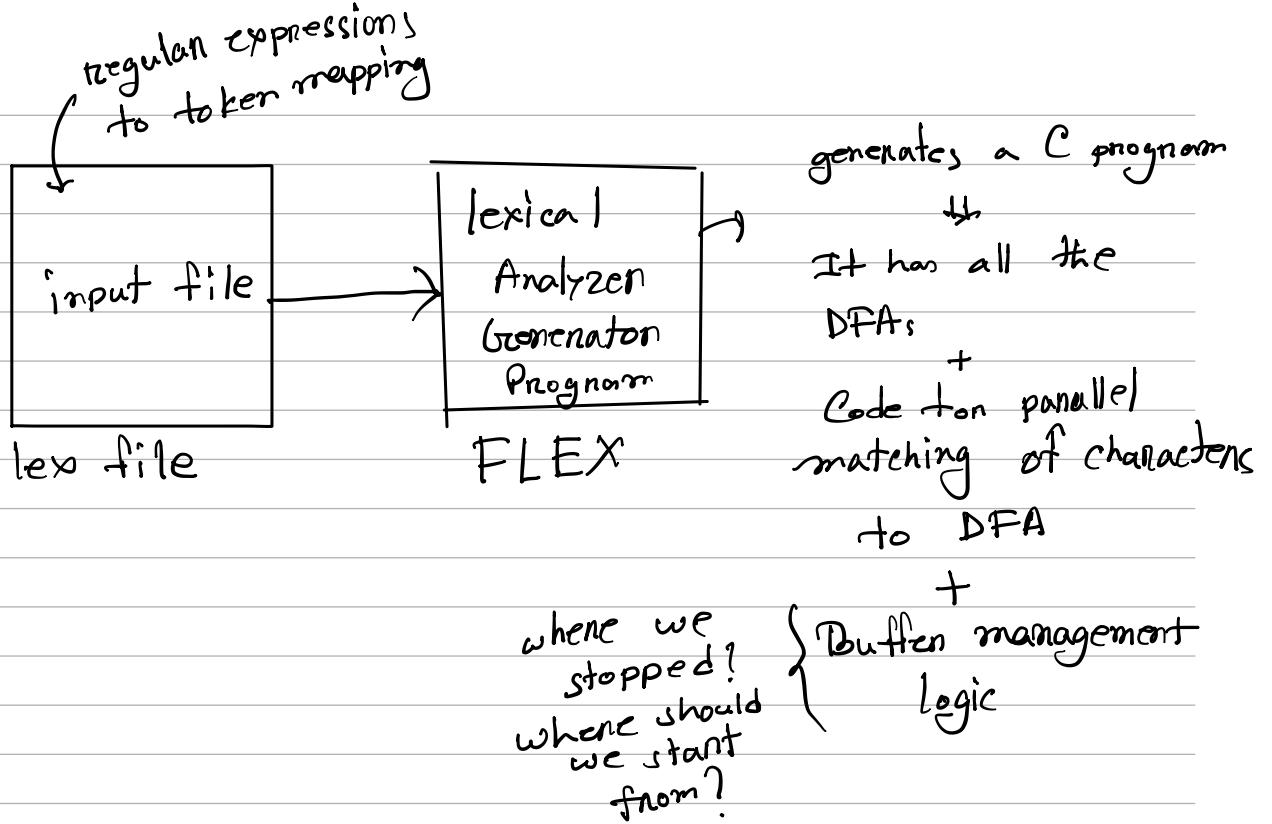


→ 1234.56
 → float 1234.56
 → operator +

→ tokenize

{ If → keyword
 Ifat → identifier?
 } hence till 'if', both the identifier DFA & keyword DFA is progressing parallelly but when we get 'a' we cannot start a new DFA cause identifier DFA is already progressed till 'if', this is why we don't consider the keyword & just consider the identifier token.

{ DFA for keyword must be placed before DFA for identifier → otherwise all if, while, for etc will be tokenized as identifier.



FLEX generates a `lex.yy.c` type file. It can be added to another. Output of lexical analysis is useful for syntax analysis.

→ In the lexical analysis stage, where is the error?

- ↳ when none of the DFAs can go to the end stage. [invalid token type]
- ↳ only identifiable error is this stage.

