Step 3: Compute First & Follow's for all non terminal symbols.

↑ ↑
two functions

* First is the set of terminals that can arrive first in any string generated by the non-terminal.

* Follow's is the set of terminals that can appear immediately after a string generated by the non-terminal.

$E' \to E$
$E \to E+T$
$E \to T$
$T \to T*F$
$T \to F$
$F \to id$
$F \to (E)$

For $F \to id$, (must be in First set of F

For $T \to *$ will follow something

include ε

First	
F	{id, (}
T	{id, (}
E	{id, (}
E'	{id, (}

→ bottom up approach

→ we don't have any terminal starting T's production rule. But T can be F and F have id & (∴ Copy F's First for T.

* If we had : $T \to \epsilon$

then $T \to T*F$
$T \to \epsilon * F$
∴ $T \to *F$
∴ * will also be in T's First

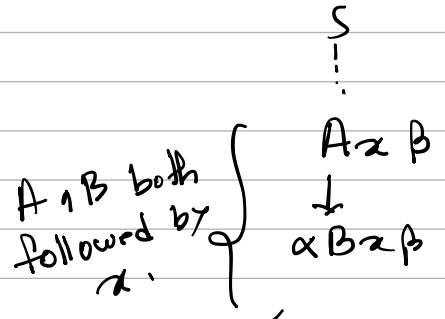
Follows → go top down

If we have:

Rule 1

$A \rightarrow \alpha B$

↳ what follows A follows B



If we have:

Rule 2

$A \rightarrow \alpha B\beta$

Here α, β means anything

↳ $F \rightarrow (E)$ (α = (, β =))

↳ $E \rightarrow E+T$ (α = E, β = T)

follows of B include first of β

↳ do not include ∈

Follows	
E'	{$} ← indicating end of file
E	{+,), $} (Rule 2, Rule 1)
T	{*, +,), $} (Rule 1)
F	{*, +,), $} (Rule 1)

For Rule 2: $A \rightarrow \alpha B\beta$ ↳ if β leads to ∈

it will lead to Rule 1

which means copying

A's follows