

# ~~কম্পিউটার~~ practice problem solve করতে হবে,

## Syntax Analyzer Error Detection Mechanisms

Syntax Analyser have many error recovery methods

### 1. Panic Mode Recovery

- input string থেকে একটা একটা করে বাক্য দিয়ে বিশ্লেষণের try

করতে,

id \* id

- When an error is detected, the parser discards the input symbols until it finds a synchronizing token
- This method is simple and ensures quick recovery, preventing infinite loops.
- Drawback: it may skip a large portion of the code, losing multiple valid statements

### 2. Phrase Level Recovery:

- When an error is detected, the parser attempts to replace a small portion of the input with something that allows parsing to continue.
- example: ~~miss~~ inserting a missing semicolon or bracket.
- Advantage: Only minor modifications are needed.
- Drawback: Might introduce incorrect assumptions, leading to further errors

### 3. Error Productions

- The grammar is augmented with productions that describes common errors. When these patterns are recognized, specific error messages can be shown.

- example: if ~~grammar~~ programmers often forget a closing bracket, a rule like  $E \rightarrow (E|E)$  could help catch and handle such errors.

- Advantage: Allows detection of specific errors at the parsing stage.

- Drawback: Adding too much error rules makes the grammar ~~rule~~ complex.

- last  $\Delta$ ; or ~~यदि~~ if  $a = b$  or ~~निरा~~ ~~होता~~ If  $a = b$  actual rule  $\Delta$  match ~~न~~ ~~करता~~  $\Delta$  rule  $\Delta$  match ~~करता~~ ~~होता~~

### 4. Global Corrections

- The parser analyzes the entire code and makes the minimum number of changes to make it syntactically correct.

- example: if an opening brace  $\{$  is missing, then the parser might suggest adding one where it seems most appropriate.

- Advantage: Produces a more meaningful error message.

~~- Drawback: Requires complex~~

- Drawback: Computationally expensive and not practical for all compilers.

- ફાઇન instructions સારા.

- Structure જે સારા match થાય તો code modify કરવા ત્ર્ય  
કરતા, મર્યાદા સારા authentic થાય

method

↳ under a class

function

↳ not under a class

## Symbol Table

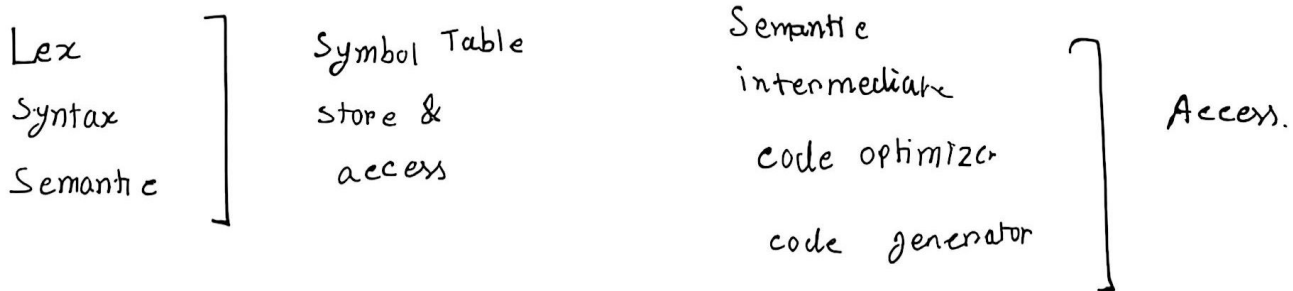
- Symbol table compiler এর সব phase এর সাথে connected থাকে,

- Symbol table o identifier এর info রাখে,  
While → while  
INT → int

☐ Why ID এর info রে symbol table এ রাখছি ; কারণ lexeme এর token info চক্কর রাখছি না,

⇒ - as memory space waste হবে সবকিছু info store করতে পারে.

- Searching time বেশি লাগবে



Semantic Action

↪ Grammar rule match করলে তা action perform হবে;

# array চক্কর মাঝে না [because চক্কর index এ থাকে তা জন্য মাঝে না.] link list, stack, queue we চক্কর মাঝে না symbol table এ.

# hashtable এর forward chaining we হয়.

↪ Searching এর জন্য best option.

■ Why we store only ID related info in the symbol table

■ Why these are called scope table.

■ Why symbol table important.

■ Why it is preferred to use hashtable in symbol table instead of other data structures

■ କେଉଁ scenario ରେ କେଉଁ ସିଦ୍ଧାନ୍ତ symbol table implement କରା ଉଚିତ, pros and cons of both.

■ କେଉଁ Variable କେଉଁ scope ରେ ଥାଏ ତାହା କିପରି ସ୍ମରଣ କରାଯାଏ?

■ Difference between Static scoping and Dynamic scoping.

Symbol Table ଉପରେ 2 ଡାକ୍ତର implement କରାଯାଇ ଯାଏ ,

- ⇒ - list of hashtables (ଅନେକଗୁଡ଼ା hashtable list ଉପରେ ଥାଏ)
- Hashtable of lists. (କେଉଁ hashtable multiple list)

### List of Hashtables:

```
1. int ab
2. void main (int b) {
3.     int c
4.     if (b == c) {
5.         float d
6.         (type mismatch) d = "0"
7.     }
8.     int bd;
9. }
10. int a
```

function scope & local variable

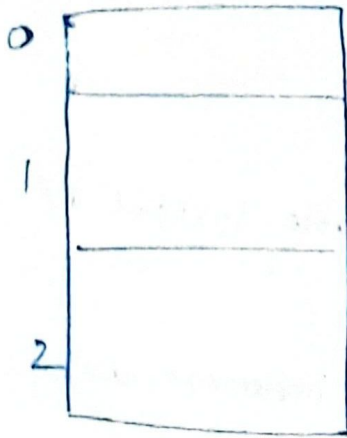
if scope

ascii value sum କିମ୍ବା bucket count  
କିମ୍ବା mod କିମ୍ବା hash function.  
exam ଏ custom hashfunction ଉପରେ  
ଧ୍ୟାନ କର.

function  
scope

Global Scope.





# Searching  $\rightarrow$  child  $\rightarrow$  parent  
এ হয়

# linked list এ key:  
value pair এ  
সহায়

hashfunction ( )

→ তামার index generate করে  
হয়, হয়ে number index এ  
ভিত্তি করে.

index = hashfunction.

ab = 0

ab এর জন্য index = 0  
generate হলো.

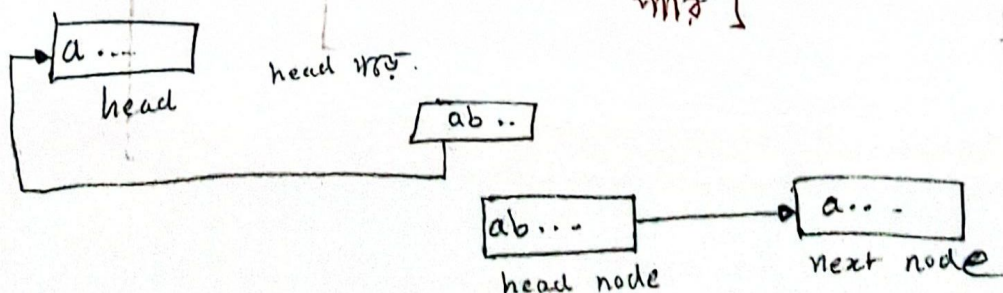
ab already exist করে কিনা  
search করে.

hashtable এ collision হলে, solve করে  
way.

→ linear probing (next available index)

→ double hashing (hash করে index generate করা)  
(length 3 হা কাজ করতে ন)

→ forward chaining (Efficient way) [ Link list দিয়ে add করে  
যাতি ]



- প্রত্যেকটি hashtable এর  
এক scope এর  
উপর,
- Per scope hashtable হবে.
- প্রত্যেকটি hashtable এর  
তার scope এর উপর  
base করে রাখা হয়  
বিভিন্ন তার scope table  
বলি.

defined variables.

- bucket\_count = 3  
→ length of the hashtable.
- level = 0 initialize করে.  
→ scope track এর জন্য  
use হয়  
→ প্রত্যেক scope,  
denote করে

## List of Hashtables:

⇒ ଭଜନକାରୀ hashtable list ଏବଂ ସମସ୍ତ ସଂକଳନ.

⇒ ଏହା ଏକ action ସଂକଳନ.

i) Scope entry

ii) Process a Declaration

iii) Process a Use

iv) Scope exit.

# same level ଏବଂ  
same ସମୟ 2 variable  
ସଂକଳନ ସଂକଳନ ସଂକଳନ,

## Actions:

### 1. Scope entry:

⇒ level ++

⇒ create a hashtable

### 2. Process a Declaration:

⇒ ଏହା scope table ଏବଂ data  
insert କରାଯାଏ.

⇒ same scope ଏବଂ same ସମୟ 2  
variable ସଂକଳନ ସଂକଳନ ସଂକଳନ.

⇒ if identifier already exists in the  
current scope table then  
multiple declare variable.

⇒ if not put the information  
in the symbol table using  
hash function.

### 3. Process a Use:

H operation perform २३५७८ ३२५७८,

⇒ Searching

⇒ if it is found

⇒ if not found: Undeclared variable

⇒ look up the identifier in the current scope table

if it's not there: go to the parent-scope -- - finally

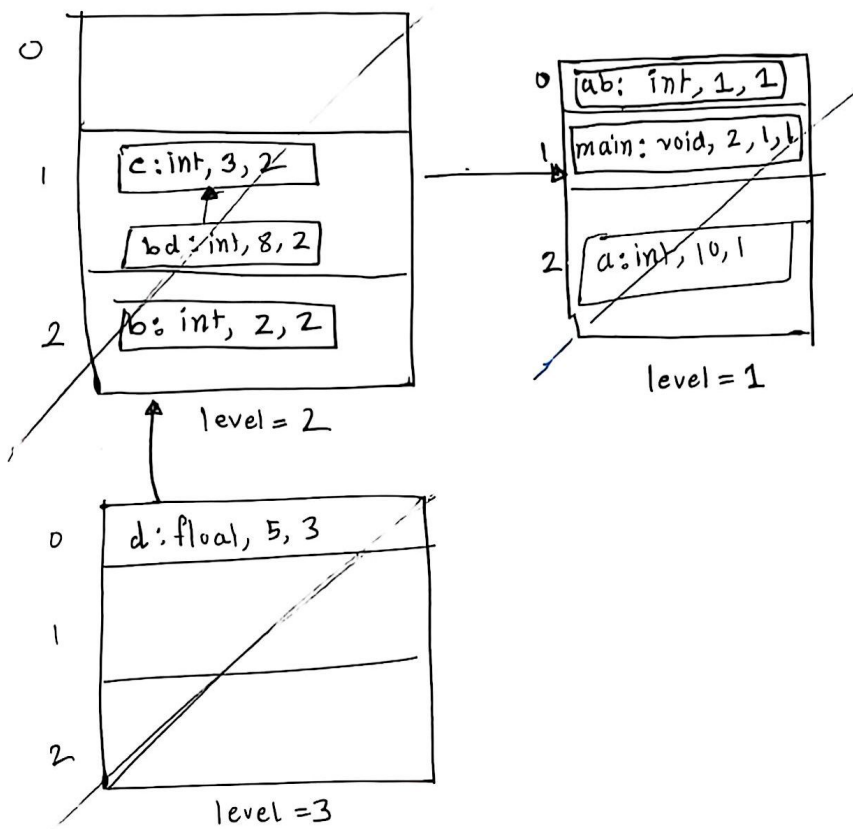
not found

### 4. Scope exits

⇒ Delete the hashtable

⇒ level --





level = 0 1 2 3 2 1 0

bucket\_count = 3

# memory head

insert karke,

`ab = 0`

`main = 1`

`b = 2`

`c = 1`

`d = 0`

`bd = 1`

`a = 2`

# arrow mark

child or

parent karte,

# key to actual lexeme

mark.; value to

ke ke info store

karke chahi.

# main function ke

global scope se

define karke so

ke info level 1 se

mark,

Hashtable: of lists.

Action  $\rightarrow$  4  $\rightarrow$  12-

1. Scope entry:

$\Rightarrow$  level++

2. Process a Declaration,

$\Rightarrow$  Search whether the identifier already exists in the symbol table with the same level number

if yes: multiple declare variable

if not: declare variable

3. Process a use:

$\Rightarrow$  check if the identifier, if found match the levels.

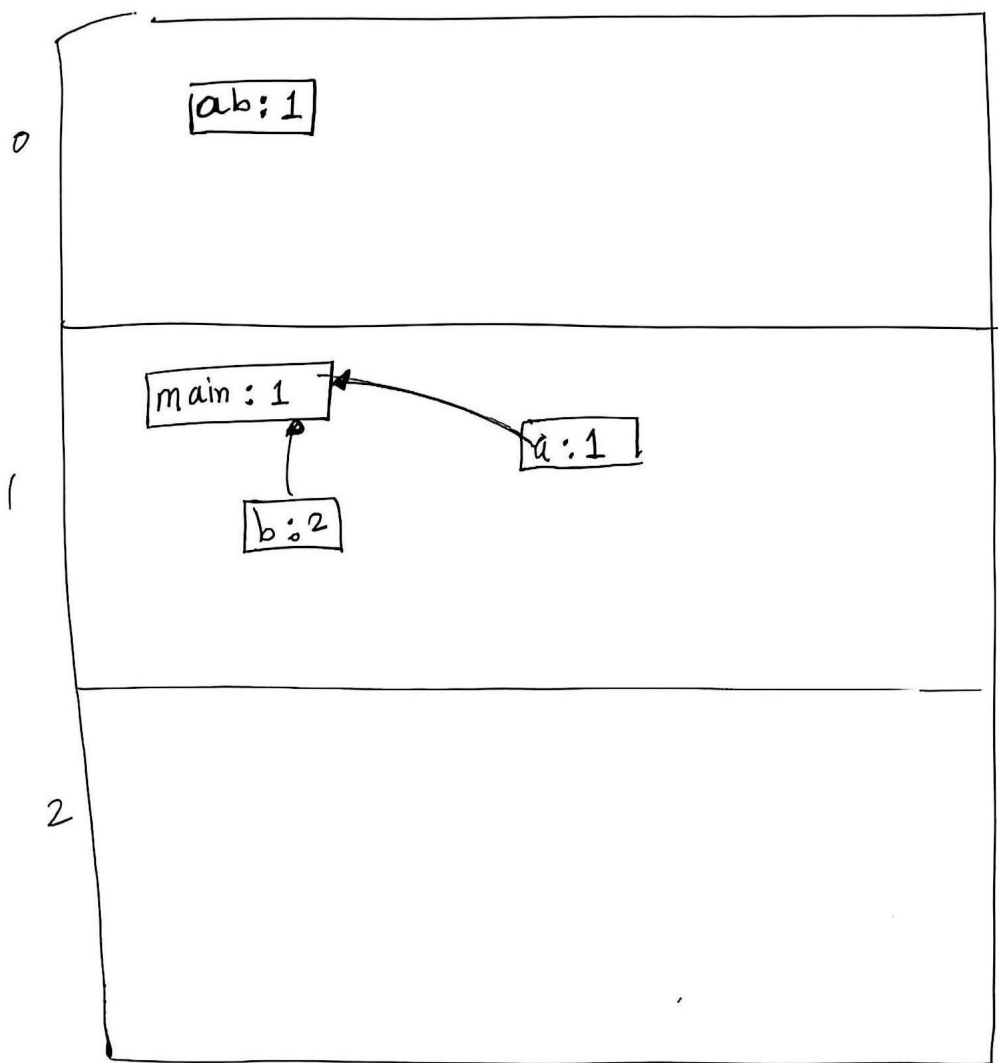
$\Rightarrow$  if not declared  $\Rightarrow$  undeclared variable

-1 $\rightarrow$ 12 $\rightarrow$ 12
search

4. Scope exit.

$\Rightarrow$  level--

$\Rightarrow$  node delete (  $\rightarrow$  level  $\rightarrow$  12  $\rightarrow$  12 node  
(  $\rightarrow$  12  $\rightarrow$  12 destroy )



level = 0 1 2 3 4 10

bucket-count = 3

ab = 0

main = 1

b = 1

c = 2

d = 2

bd = 2

# hashtable of list & linktable  
 ଦିଆଯାଇଥିବା data ଥିବାପରି ହେବ

# list of hashtable & ଅଧିକ ଯୋଗ

ହେଉଥିବା hash function ଥିବାପରି ଯୋଗାଯୋଗ

collision ହେବାର chances 0

ଓହ୍ଲା linked list ଦିଆଯାଇ

data ଥିବା not mandatory

level ଓ base ଯାହା different  
 different tables ଦିଆଯାଇ