

CSC2400 / Chapter 8 Homework

Due: Monday, April 20, 2020

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1. (6 points) What does dynamic programming have in common with divide-and-conquer? What is a principal difference between them?

Each problem is split into sub problems and solved individually. The principal difference is that Dynamic stores the sub problems generally in a table so the sub problems don't need to be solved again, it then uses the table to solve the original problem.

2. (10 points) Solve the instance 10, 25, 1, 1, 5, 1, 25, 10, 10 of the coin-row problem using the dynamic programming algorithm technique. Show the solution array and also the final output.

Given: $n=9$

	0	1	2	3	4	5	6	7	8	9
C	10	25	1	1	5	1	25	10	10	

Step 1: A

	0	1	2	3	4	5	6	7	8	9
A	0	10								

Step 2: $A[0] = 0$, $A[1] = C[1]$

Step 3: $A[I] = \text{Max} \begin{cases} C[I] + A[I-2] \\ \text{or} \\ A[I-1] \end{cases}$
i starts at 2 and goes through $n=9$

	0	1	2	3	4	5	6	7	8	9
A	0	25	25	26	31	33	58	58	68	

Step 4: return $A[N]$

	0	1	2	3	4	5	6	7	8	9
A	0	25	25	26	31	33	58	58	68	

→ return 68

3. **(10 points)** Using the change making DP algorithm, give change for amount **8** using the minimum number of coins of denominations **(1, 2, 3, 4)**. You may assume that there is an unlimited quantity of coins for each of the 4 denominations. Determine how many coins and also which coins will be used to make change.

Amount = 8 coins = {1, 2, 3, 4}

	0	1	2	3	4	5	6	7	8
i=0	Step 1								
f[0]=0									
	$f[i] = \min(f[i-1], f[i-2], f[i-3], f[i-4]) + 1$								
i=1									
$f[1] = f[1-1] + 1 = 1$									
i=2									
$f[2] = f[2-2] + 1 = f[0] + 1 = 1$									
i=3									
$f[3] = f[3-3] + 1 = f[0] + 1 = 1$									
i=4									
$f[4] = f[4-4] + 1 = f[0] + 1 = 1$									
i=5									
$f[5] = f[5-1] + 1 = f[1] + 1 = 2$									
i=6									
$f[6] = f[6-2] + 1 = f[2] + 1 = 2$									
i=7									
$f[7] = f[7-3] + 1 = f[3] + 1 = 2$									
i=8									
$f[8] = f[8-4] + 1 = f[4] + 1 = 2$									

	0	1	2	3	4	5	6	7	8
	0	1	1	1	1	2	2	2	2

Coin 4 produced $f[8] = 2$
 So we need 2 of coin '4' to get 8

4. Knapsack Problem.

- a. **(15 points)** Apply the bottom-up dynamic programming algorithm to the following instance of the knapsack problem:

- item 1, weight is 3 lbs, value is \$25
- item 2, weight is 2 lbs, value is \$20
- item 3, weight is 1 lb, value is \$15
- item 4, weight is 4 lbs, value is \$40
- item 5, weight is 5 lbs, value is \$50
- **Capacity of knapsack is 6 lbs total**

- b. (5 points) What is the maximum value of a feasible subset of the knapsack in part (a)?

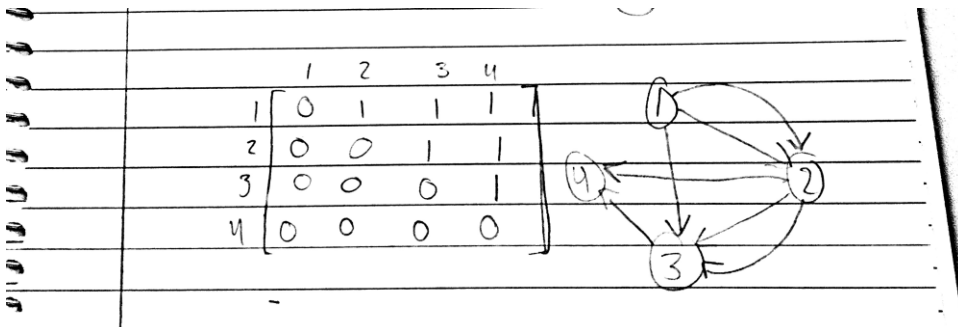
		Weight						
		0	1	2	3	4	5	6
0	0	0	0	0	0	0	0	0
1	0	0	0	0	0	0	0	0
2	0	0	0	0	0	0	0	0
3	0	0	0	0	0	0	0	0
4	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0

		0	1	2	3	4	5	6
0	0	0	0	0	0	0	0	0
1	0	0	0	25	25	25	25	25
2	0	0	20	25	25	45	45	45
3	0	15	20	35	40	45	60	60
4	0	15	20	35	40	55	60	60
5	0	15	20	35	40	55	65	65

\$65

5. (15 points) Apply Warshall's algorithm to find the transitive closure of the digraph defined by the following adjacency matrix:

$$\begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$



SCORE: _____ / 61