

# Slides

# Introduction

# Organization

Professors:

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# Grades

$$\text{Final grade} = 0.75 \text{ Exam} + 0.25 \text{ Lab}$$

# Time schedule

- ▶ 14 weeks of lectures (3h each)
- ▶ 14 weeks of laboratories (2h each)
  - ▶ 5 laboratories
  - ▶ 7 seminars
  - ▶ 1 recuperari
  - ▶ 1 test

# Course structure

1. Introduction to probabilities

2. Pam

# Bibliography

1. Pam Pam
2. HamHam
3. Yoyo

# Introduction to probabilities



# Basic notions of probability

Random variable = the outcome of an experiment

Distribution (probability mass function)

Discrete distribution

Alphabet

## Basic properties

Two independent events:

$$P(A \cap B) = P(A) \cdot P(B)$$

## Chapter I: Discrete information sources

# Block diagram of a communication system

de pus poza

# What is information?

Example:

I tell you the following sentence: “your favorite football team lost the last match”.

Does this message carry information? How, why, how much?

Consider the following facts:

- ▶ the message carries information only because you didn't already know the result.
- ▶ if you already known the result, the message is useless (brings no information)
- ▶ since you didn't know the result, there were multiple results possible (win, equal or lose)
- ▶ the actual information in the message is that *lost* happened, and not *win* or *equal*
- ▶ if the result was to be expected, there is little information. If the result is highly unusual, there is more information in this message

## Information source

We will always consider information in a context similar to the above example.

We will use terminology from probability theory to define information:

- ▶ there is a *probabilistic source* that can produce a number of different *events*.
- ▶ each event has a certain probability. We know all the probabilities beforehand.
- ▶ at one time, an event is randomly selected according to its probability.
- ▶ afterwards, a new message can be selected, and so on  $\Rightarrow$  a stream of messages is produced.

The source is called an *information source* and the selected event is a *message*.

A message carries the information that **it** happened, and not the other possible message events that could have been selected.

The quantity of information is dependent in its probability.

## Discrete memoryless source

= is an information source where the messages are **independent** ,  
i.e. the choice of a message at one time does not depend on what  
were the previous message

Each message has a fixed probability. The set of probabilities is the  
*distribution* of the source.

$$S : \begin{pmatrix} s_1 & s_2 & s_3 \\ \frac{1}{2} & \frac{1}{4} & \frac{1}{4} \end{pmatrix}$$

Properties:

- ▶ Discrete: it can take a value from a discrete set (alphabet)
- ▶ Complete:  $\sum p(s_i) = 1$
- ▶ Memoryless: successive values are independent of previous values (e.g. successive throws of a coin)

A message from a DMS is also called a *random variable* in  
probabilistics.

## Examples

A coin is a discrete memoryless source (DMS) with two messages (head, tail):

$$S : \begin{pmatrix} s_1 & s_2 \\ \frac{1}{2} & \frac{1}{2} \end{pmatrix}$$

A dice is a discrete memoryless source (DMS) with six messages:

$$S : \begin{pmatrix} s_1 & s_2 & s_3 & s_4 & s_5 & s_6 \\ \frac{1}{6} & \frac{1}{6} & \frac{1}{6} & \frac{1}{6} & \frac{1}{6} & \frac{1}{6} \end{pmatrix}$$

Playing the lottery can be modeled as DMS:

$$S : \begin{pmatrix} s_1 & s_2 \\ 0.9999 & 0.0001 \end{pmatrix}$$

An extreme type of DMS containing the certain event:

$$S : \begin{pmatrix} s_1 & s_2 \\ 1 & 0 \end{pmatrix}$$



# Information

When a DMS provides a new message, it gives out some new information, i.e. the information that a particular message took place.

The information attached to a particular event (message) is rigorously defined as:

$$i(s_i) = -\log_2(p(s_i))$$

Properties:

- ▶  $i(s_i) \leq 0$
- ▶ lower probability means higher information
- ▶ higher probability means lower information
- ▶ a certain event brings no information:  $-\log(1) = 0$
- ▶ an event with probability 0 brings infinite information (but it never happens..)

# Entropy of a DMS

We usually don't care about a single message. We are interested in a large number of them (think millions of bits of data).

We are interested in the *average* information of a message from a DMS.

Definition: the entropy of a DMS source  $S$  is **the average information of a message**:

$$H(S) = \sum_i p_i i(s_i) = - \sum_i p_i \log(p_i)$$

where  $p_i = p(s_i)$  is the probability of message  $i$ .

...

## Example - Game

Game: I think of a number between 1 and 8. You have to guess it by asking yes/no questions.

- ▶ How much indetermination does the problem have?
- ▶ How is the best way to ask questions? Why?
- ▶ What if the questions are not asked in the best way?
- ▶ On average, what is the number of questions required to find the number?

## Example - Game v2

Suppose I choose a number according to the following distribution:

...

- ▶ On average, what is the number of questions required to find the number?
- ▶ What distribution makes guessing the number the most difficult?
- ▶ What distribution makes guessing the number the easiest?

## Entropy of a discrete memoryless source

# Properties of entropy

1. It
2. is
3. cool

## Sources with memory

A text can be considered as a sequence of symbols drawn from a memoryless source.

The distribution (frequencies) of letters in the Romanian language is close to: