

Prim's Algorithm

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1  Algorithm Prim( $E, cost, n, t$ )
2  //  $E$  is the set of edges in  $G$ .  $cost[1 : n, 1 : n]$  is the cost
3  // adjacency matrix of an  $n$  vertex graph such that  $cost[i, j]$  is
4  // either a positive real number or  $\infty$  if no edge  $(i, j)$  exists.
5  // A minimum spanning tree is computed and stored as a set of
6  // edges in the array  $t[1 : n - 1, 1 : 2]$ .  $(t[i, 1], t[i, 2])$  is an edge in
7  // the minimum-cost spanning tree. The final cost is returned.
8  {
9      Let  $(k, l)$  be an edge of minimum cost in  $E$ ;
10      $mincost := cost[k, l]$ ;
11      $t[1, 1] := k$ ;  $t[1, 2] := l$ ;
12     for  $i := 1$  to  $n$  do // Initialize near.
13         if ( $cost[i, l] < cost[i, k]$ ) then  $near[i] := l$ ;
14         else  $near[i] := k$ ;
15      $near[k] := near[l] := 0$ ;
16     for  $i := 2$  to  $n - 1$  do
17     { // Find  $n - 2$  additional edges for  $t$ .
18         Let  $j$  be an index such that  $near[j] \neq 0$  and
19          $cost[j, near[j]]$  is minimum;
20          $t[i, 1] := j$ ;  $t[i, 2] := near[j]$ ;
21          $mincost := mincost + cost[j, near[j]]$ ;
22          $near[j] := 0$ ;
23         for  $k := 1$  to  $n$  do // Update near[ ].
24             if ( $near[k] \neq 0$ ) and ( $cost[k, near[k]] > cost[k, j]$ )
25                 then  $near[k] := j$ ;
26     }
27     return  $mincost$ ;
28 }

```