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COURSE CODE:CSA0689.

SUB.NAME:DESIGN AND ANALYSIS OF

ALGORITHM.

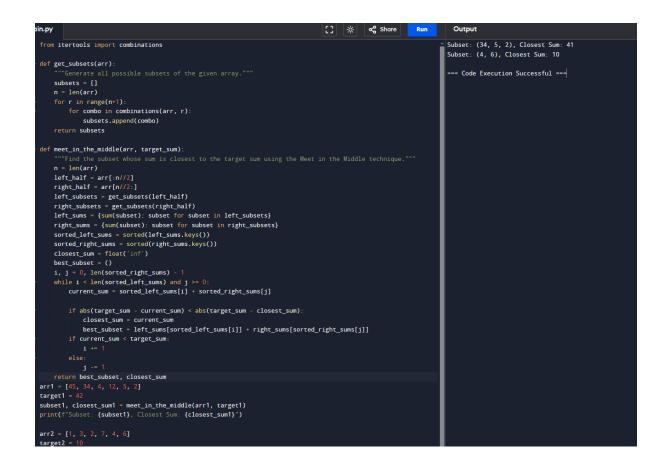
1.To Implement the Median of Medians algorithm ensures that you handle the worst-case time complexity efficiently while finding the k-th smallest element in an unsorted array. arr = [12, 3, 5, 7, 19] k = 2 Expected Output:5 arr = [12, 3, 5, 7, 4, 19, 26] k = 3 Expected Output:5 arr = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10] k = 6 Expected Output:6.

```
[] ☆ ∝ Share Run
                                                                                                                                                              Output
  main.py
   1 def partition(arr, low, high, pivot):
2 for i in range(low, high):
                                                                                                                                                            Expected Output: 5
                                                                                                                                                            Expected Output: 6
                      arr[i], arr[high] = arr[high], arr[i]
                                                                                                                                                             === Code Execution Successful ===
           pivot = arr[high]
            i = low -
           for j in range(low, high):
    if arr[j] <= pivot:</pre>
           arr[i], arr[j] = arr[j], arr[i]
arr[i+1], arr[high] = arr[high], arr[i+1]
  15 def find_median(arr):
 16 arr.sort()
17 return arr[len(arr) // 2]
18
19 def median_of_medians(arr, k):
for i in range(0, len(arr), 5):
    group = arr[i:i+5]
                 medians.append(find_median(group))
           \label{eq:median_of_medians} \begin{tabular}{ll} median_of_medians & len(medians)//2 + 1) \\ pivot_index & = partition(arr, 0, len(arr) - 1, median_of_medians_pivot) \\ \end{tabular}
           if pivot_index == k - 1
                return arr[pivot_index]
                 f pivot_index > k - 1:
return median_of_medians(arr[:pivot_index], k)
                 return median_of_medians(arr[pivot_index + 1:], k - pivot_index - 1)
      print("Expected Output:", median_of_medians(arr1, k1))
 42
43
```

2. To Implement a function median_of_medians(arr, k) that takes an unsorted array arr and an integer k, and returns the k-th smallest element in the array. arr = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10] k = 6 arr = [23, 17, 31, 44, 55, 21, 20, 18, 19, 27] k = 5 Output: An integer representing the k-th smallest element in the array.

```
1 def partition(arr, low, high, pivot_index):
       pivot_value = arr[pivot_index]
arr[pivot_index], arr[high] = arr[high], arr[pivot_index]
       store_index = low
                                                                                                                           === Code Execution Successful
       for i in range(low, high):
              arr[store_index], arr[i] = arr[i], arr[store_index]
                store_index -
       arr[store_index], arr[high] = arr[high], arr[store_index]
        return store_index
11 def median_of_medians(arr, k):
           return arr[0]
       sublists = [arr[i:i+5] for i in range(0, n, 5)]
medians = [sorted(sublist)[len(sublist)//2] for sublist in sublists]
       pivot = sorted(medians)[len(medians)//2]
else:
          pivot = median_of_medians(medians, len(medians)//2)
21
22
23
24
       pivot_index = arr.index(pivot)
       pivot_index = partition(arr, 0, len(arr) - 1, pivot_index)
       if k == pivot_index +
            return arr[pivot_index]
       elif k < pivot_index + '
26
27
28
            return median_of_medians(arr[:pivot_index], k)
            return median_of_medians(arr[pivot_index+1:], k - pivot_index - 1)
29 arr1 = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]
30 k1 = 6
   print(median_of_medians(arr1, k1))
35 print(median_of_medians(arr2, k2))
```

3. Write a program to implement Meet in the Middle Technique. Given an array of integers and a target sum, find the subset whose sum is closest to the target. You will use the Meet in the Middle technique to efficiently find this subset. a) Set[] = {45, 34, 4, 12, 5, 2} Target Sum : 42 b) Set[] = {1, 3, 2, 7, 4, 6} Target sum = 10:



4. . Write a program to implement Meet in the Middle Technique. Given a large array of integers and an exact sum E, determine if there is any subset that sums exactly to E. Utilize the Meet in the Middle technique to handle the potentially large size of the array. Return true if there is a subset that sums exactly to E, otherwise return false. a) $E = \{1, 3, 9, 2, 7, 12\}$ exact Sum = 15 b) $E = \{3, 34, 4, 12, 5, 2\}$ exact Sum = 15.

```
[] ☆ & Share Run
                                                                                                                                  Output
 main.pv
    from itertools import combinations
    def get_all_subset_sums(arr):
        """Generate all possible subset sums of the given array."""
subset_sums = set()
                                                                                                                                False
                                                                                                                                === Code Execution Successful ==
        for r in range(n+1):
           for combo in combinations(arr, r):
    subset_sums.add(sum(combo))
11 def meet_in_the_middle(arr, target_sum):
        n = len(arr)
       left_half = arr[:n//2]
       right_half = arr[n//2:]
        left_sums = get_all_subset_sums(left_half)
       right_sums = get_all_subset_sums(right_half)
if target_sum in left_sums:
17
18
       if target_sum in right_sums:
       for left_sum in left_sums:
        if (target_sum - left_sum) in right_sums:
28 target_sum1 = 15
29 print(meet_in_the_middle(arr1, target_sum1))
32 target_sum2 =
33 print(meet_in_the_middle(arr2, target_sum2))
36 target_sum3 =
37 print(meet_in_the_middle(arr3, target_sum3))
```

5. Given two 2×2 Matrices A and B A=(1 7 B=(1 3 3 5) 7 5) Use Strassen's matrix multiplication algorithm to compute the product matrix C such that C=A×B. Test Cases: Consider the following matrices for testing your implementation: Test Case 1: A=(1 7 B=(6 8 3 5), 4 2) Expected Output: C=(18 14 62 66.

```
moin.py

1 def strassen_multiplication(A, B):
2 a, b, c, d = A[0][0], A[0][1], A[1][0], B[1][1]
3 e, f, g, h = B[0][0], B[0][1], B[1][0], B[1][1]
4 M1 = (a + d) * (e + h)
5 M2 = (c + d) * e
6 M3 = a * (f - h)
7 M4 = d * (g - e)
8 M5 = (a + b) * h
9 M6 = (c - a) * (e + f)
10 M7 = (b - d) * (g + h)
11 p = M1 + M4 - M5 + M7
12 q - M3 + M5
13 r = M2 + M4
14 s = M1 - M2 + M3 - M6
15 C = [[p, q], [r, s]]
16 return C
17
18 A = [[1, 7], [3, 5]]
19 B - [[6, 8], [4, 2]]
20 C = strassen_multiplication(A, B)
21 print("Resultant Matrix C.")
22 for row in C:
23 print("Resultant Matrix C.")
22 for row in C:
23 print("Resultant Matrix C.")
```

6. 6. Given two integers X=1234 and Y=5678: Use the Karatsuba algorithm to compute the product Z=X x Y Test Case 1: Input: x=1234,y=5678 Expected Output: z=1234×5678=7016652.