Threads

a new abstraction of program execution

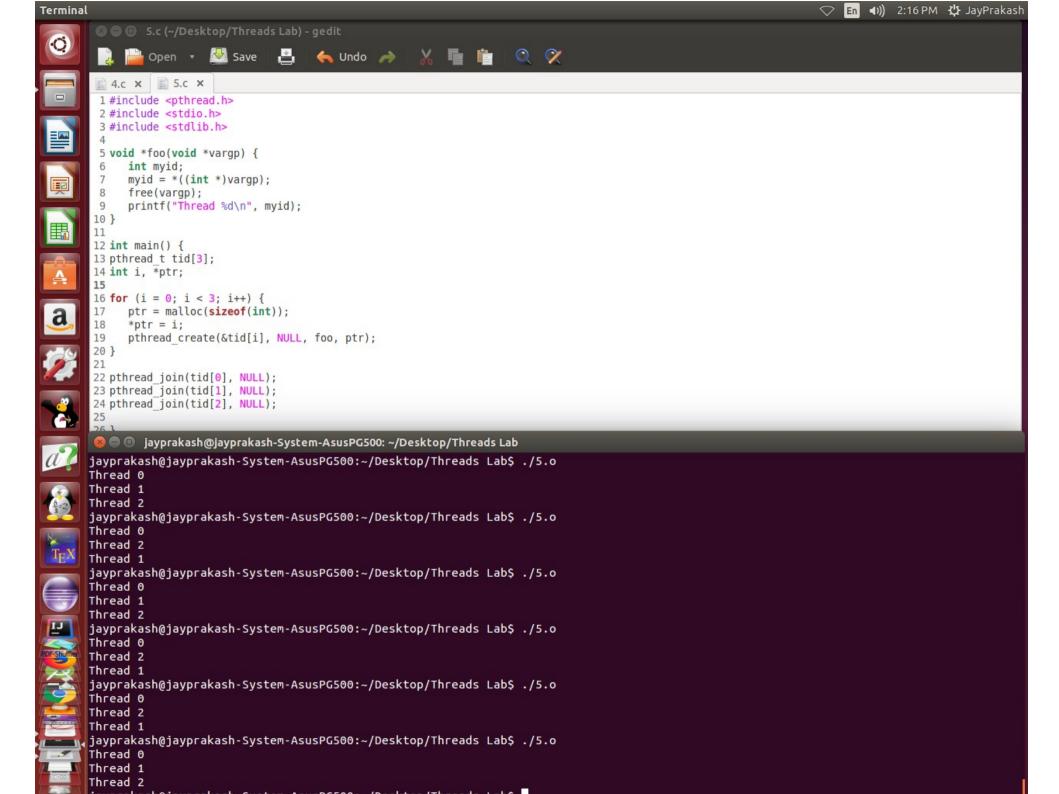
Disadvantage of MUTEX

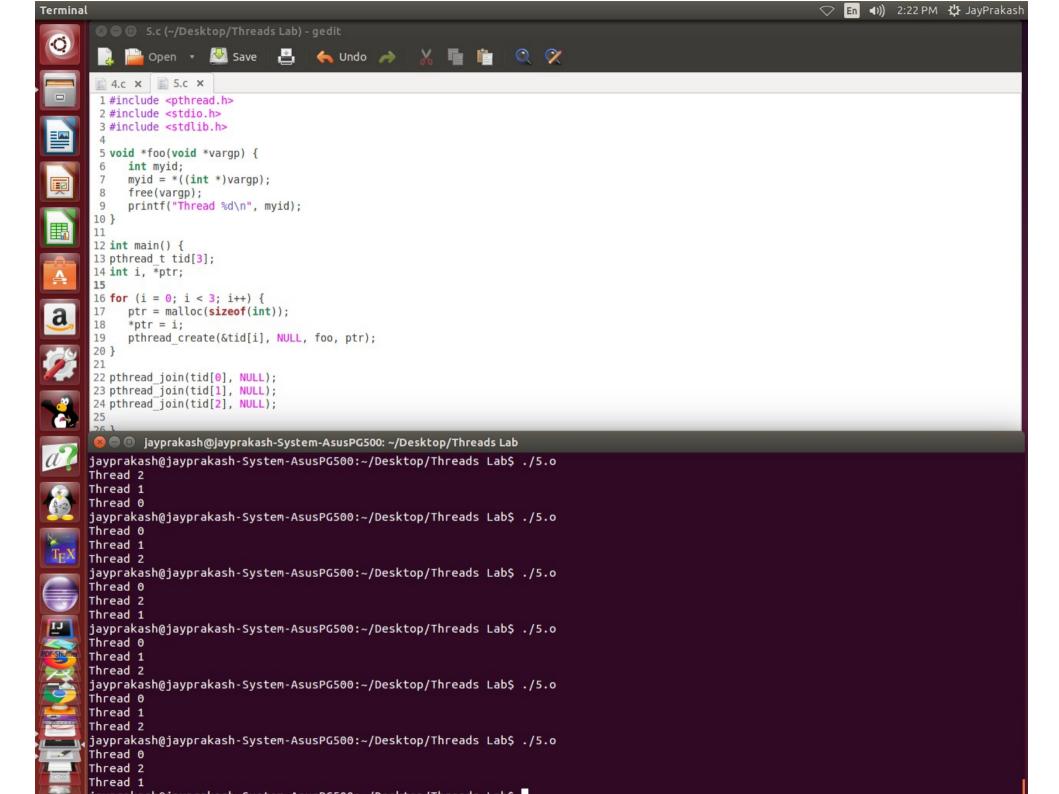
- Deadlock can occur if:
 - Thread locking the same MUTEX twice
 - Thread1 holding lock on mutex1 wants to get lock on mutex2 while thread2 holding lock on mutex2 wants to get lock on mutex1
 - Solution: As much as possible use pthread_mutex_trylock function
- Only one thread is allowed to lock mutex but in some situations such as Reader-Write problem or Dinning-Philosopher problem multiple locks should be possible
 - Solution: Reader-Writer Locks and Semaphores

Example Solution without race-condition

Q: Are there any race conditions in this code?

```
void *foo(void *vargp) {
    int myid;
   myid = *((int *)vargp);
    Free (varqp);
   printf("Thread %d\n", myid);
                                      No!
int main() {
   pthread t tid[2];
    int i, *ptr;
    for (i = 0; i < 2; i++) {
       ptr = Malloc(sizeof(int));
       *ptr = i;
       Pthread create(&tid[i], NULL, foo, ptr);
    Pthread join(tid[0], NULL);
    Pthread join(tid[1], NULL);
```





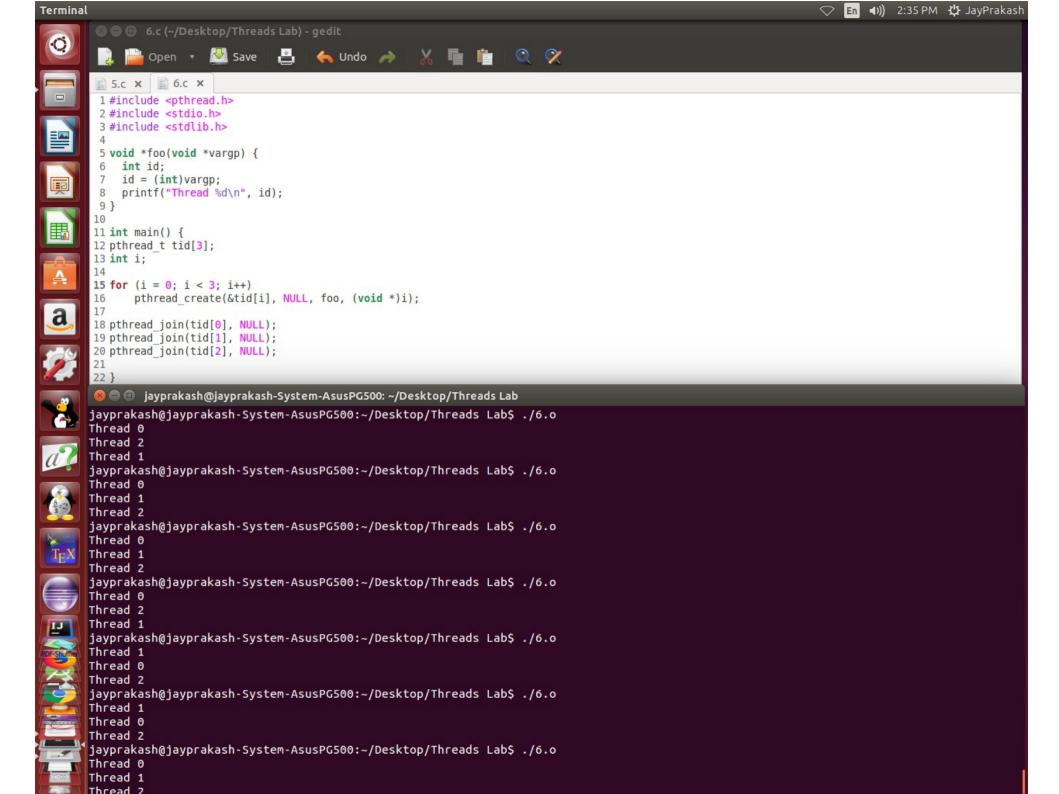
Question

 Outline an approach to avoid race conditions that does not use malloc/free

Example Solution without malloc()

Answer: Simply pass in the int directly!

```
void *foo(void *vargp) {
    int id;
    id = (int) vargp;
    printf("Thread %d\n", id);
int main() {
    pthread t tid[2];
    int i;
    for (i = 0; i < 2; i++)
        Pthread create(&tid[i], NULL, foo, (void *)i);
    Pthread join(tid[0], NULL);
    Pthread join(tid[1], NULL);
```



Previous example

- Pro: No added overhead due to malloc/free
- Con: Assumes that pointer datatype is at least bigger than size of int
 - May not be true on all systems

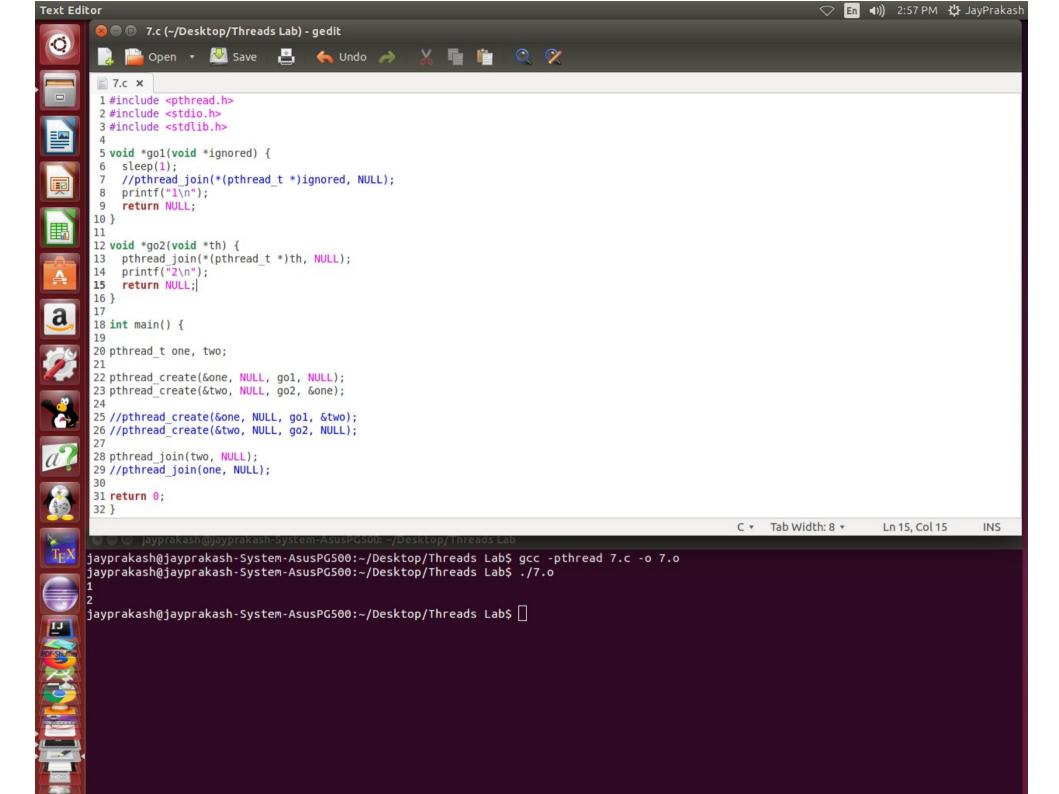
Threads are not hierarchical

- Unlike processes, a thread doesn't have a parent or children threads are peers
- Any thread can pthread_join any other thread

Threads are not hierarchical

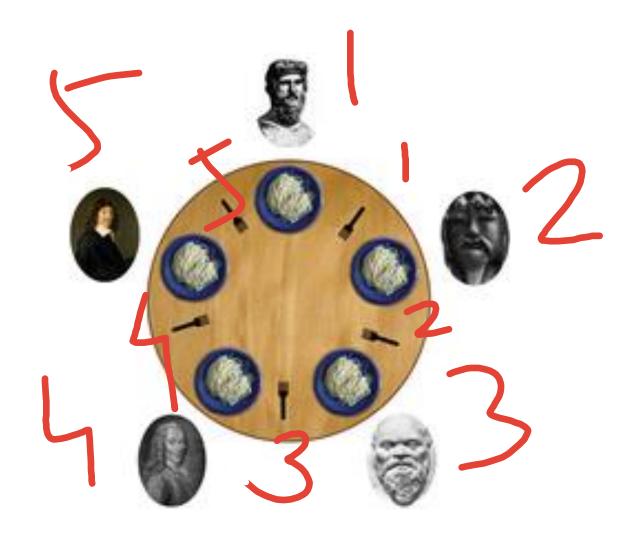
```
void *go1(void *ignored) {
    sleep(1);
   printf("1\n");
    return NULL;
void *go2(void *th) {
   pthread join(*(pthread t *)th, NULL);
   printf("2\n");
    return NULL;
int main() {
   pthread t one, two;
   pthread create(&one, NULL, go1, NULL);
   pthread create(&two, NULL, go2, &one);
   pthread join(two, NULL);
    return 0;
```

Prints 1 then 2



Dining-Philosopher Problem

- Five philosophers spend their time eating and thinking.
- They are sitting in front of a round table with spaghetti served.
- There are five plates at the table and five chopsticks set between the plates.
- Eating the spaghetti requires the use of two chopsticks which the philosophers pick up one at a time.
- Philosophers do not talk to each other.
- Semaphore chopstick [5] initialized to



Semaphores

- Allow multiple locks
- Semaphore S → integer variable
- Modified by two operations → wait() and signal()
 - wait() originally called P() for Dutch word "proberen" which means try
 - signal() originally called V() for Dutch word "verhogen" which means increase wait(S) {
 while S <= 0: // no-op

```
while S <= 0; // no-op
    S--;
}
signal(S) {
    S++;
}</pre>
```

Note wait() and signal() are atomic operation

Thread Synchronization using Semaphores

```
//Thread1:
                                       //Thread2:
                                       int t;
int t;
wait(sem)
                                       wait(sem)
sum = sum + x;
                                       sum = sum + y;
                                       t = sum;
t = sum;
• • •
                                        ...
                                       signal(sem);
signal(sem);
```

Dining-Philosopher Problem: Solution

```
The structure of Philosopher i:
do {
 wait ( chopstick[i] ); //lock
 wait (chopStick[ (i + 1) % 5] ); //lock
 // eat
 signal ( chopstick[i] ); //unlock
 signal (chopstick[ (i + 1) % 5] );
//unlock
 // think
} while (true);
```

- Problem is deadlock: when all philosopher decides to eat at the same time each will pick up one of the 2 chopsticks they need and wait for the other
- Solution: Allow to lock if both chopsticks are available at the same time
 - Lock the 1st chopstick
 - For 2nd chopstick, check if it can be locked otherwise release the 1st one

POSIX: Semaphores

• creating a semaphore:

int sem_init(sem_t *sem, int pshared, unsigned int value);

- initializes a semaphore object pointed to by sem to integer value of "value"
- pshared is a sharing option; a value of 0 means the semaphore is local to the calling process i.e. shared by all threads of the same process; a +ve number indicates it can be shared across multiple processes using shared memory instructions
- terminating a semaphore:

```
int sem_destroy(sem_t *sem);
```

- frees the resources allocated to the semaphore sem
- usually called after pthread_join()
- an error will occur if a semaphore is destroyed for which a thread is waiting

POSIX: Semaphores

semaphore control:

```
int sem_post(sem_t *sem); → same as signal()
sem_post atomically increases the value of a semaphore by 1, i.e.,
when 2 threads call sem_post simultaneously, the semaphore's value
will also be increased by 2 (there are 2 atoms calling)
int sem_wait(sem_t *sem);
sem_wait atomically decreases the value of a semaphore by 1; but
```

always waits until the semaphore has a non-zero value first

Semaphore Example

semaphore1.c

\$./semaphore1.out Starting thread, semaphore is unlocked. Hello from thread!

Hello from thread!

Hello from thread!

Hello from thread!

Semaphore locked.

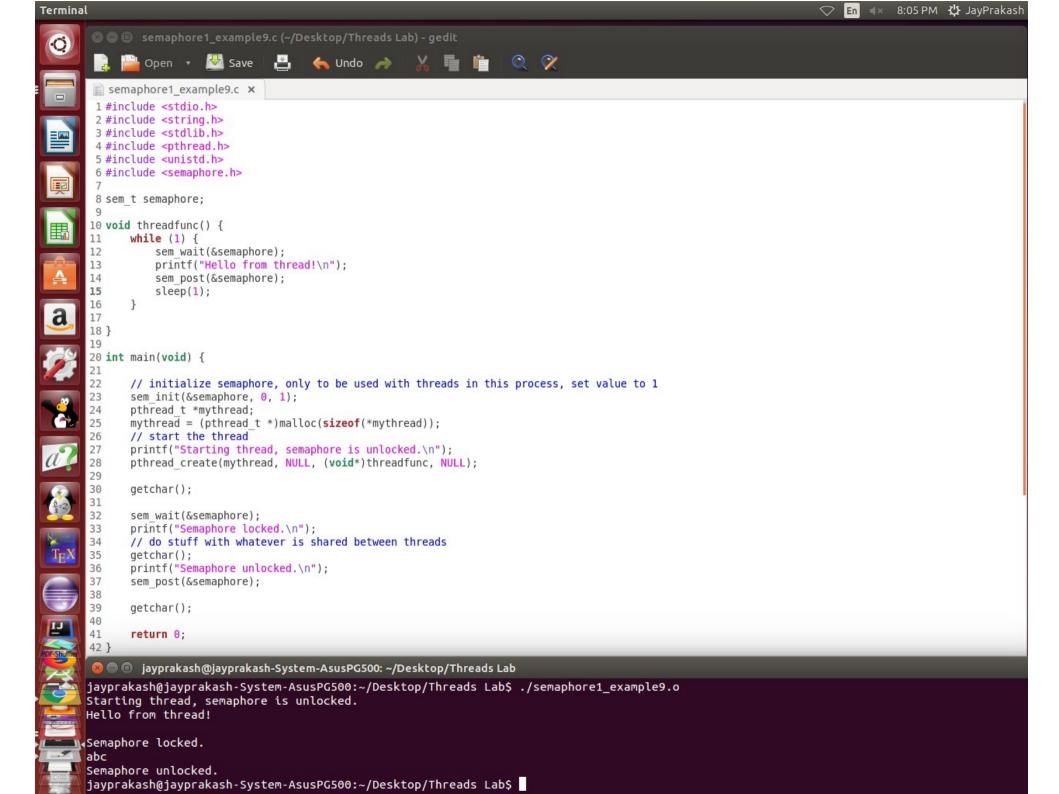
Semaphore unlocked.

Hello from thread!

Hello from thread!

Hello from thread!

\$



Dining-Philosopher Problem: Solution

dining-philosopher.c

Philosopher 1 is thinking

Philosopher 2 is thinking

Philosopher 3 is thinking

Philosopher 4 is thinking

Philosopher 5 is thinking

Philosopher 1 is Hungry

Philosopher 2 is Hungry

Philosopher 3 is Hungry

Philosopher 4 is Hungry

Philosopher 5 is Hungry

Philosopher 5 takes fork 4 and 5

Philosopher 5 is Eating

Philosopher 5 putting fork 4 and 5 down

Philosopher 5 is thinking

Philosopher 4 takes fork 3 and 4

Philosopher 4 is Eating

Philosopher 1 takes fork 5 and 1

Philosopher 1 is Eating

Philosopher 4 putting fork 3 and 4 down

Philosopher 4 is thinking

Philosopher 3 takes fork 2 and 3

Philosopher 3 is Eating

Philosopher 5 is Hungry

Philosopher 1 putting fork 5 and 1 down

Philosopher 1 is thinking

Philosopher 5 takes fork 4 and 5

Philosopher 5 is Eating

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