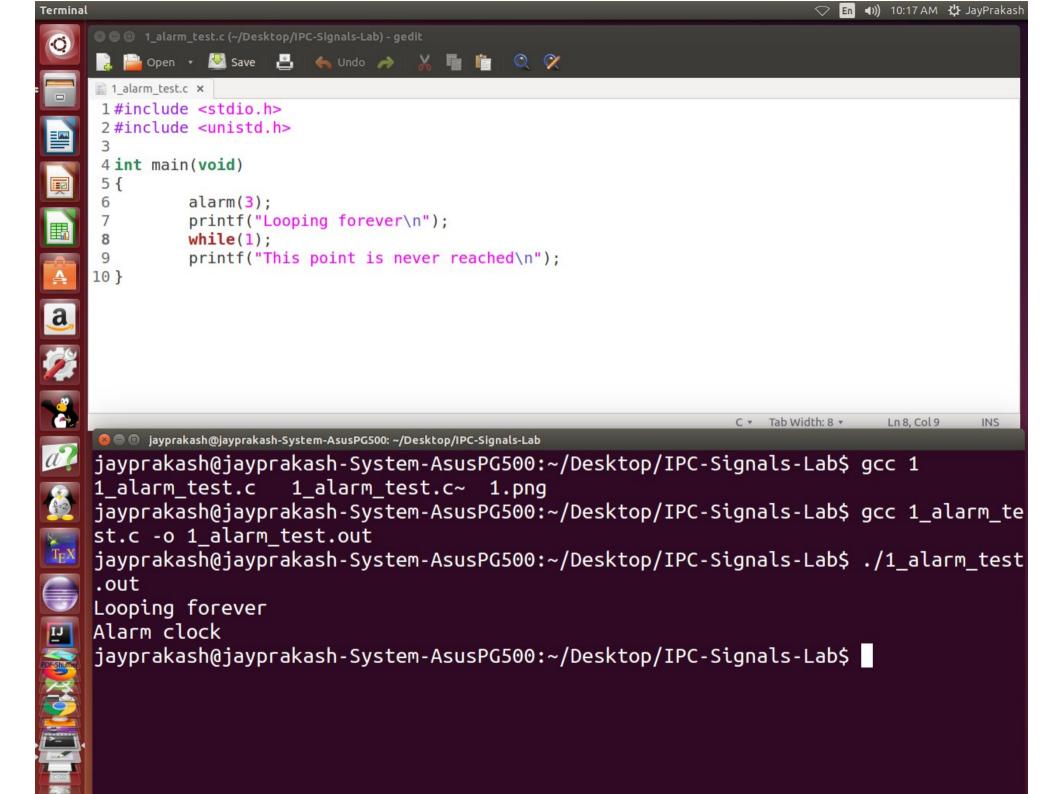
Requesting an Alarm Signal SIGALRM using alarm() System Call

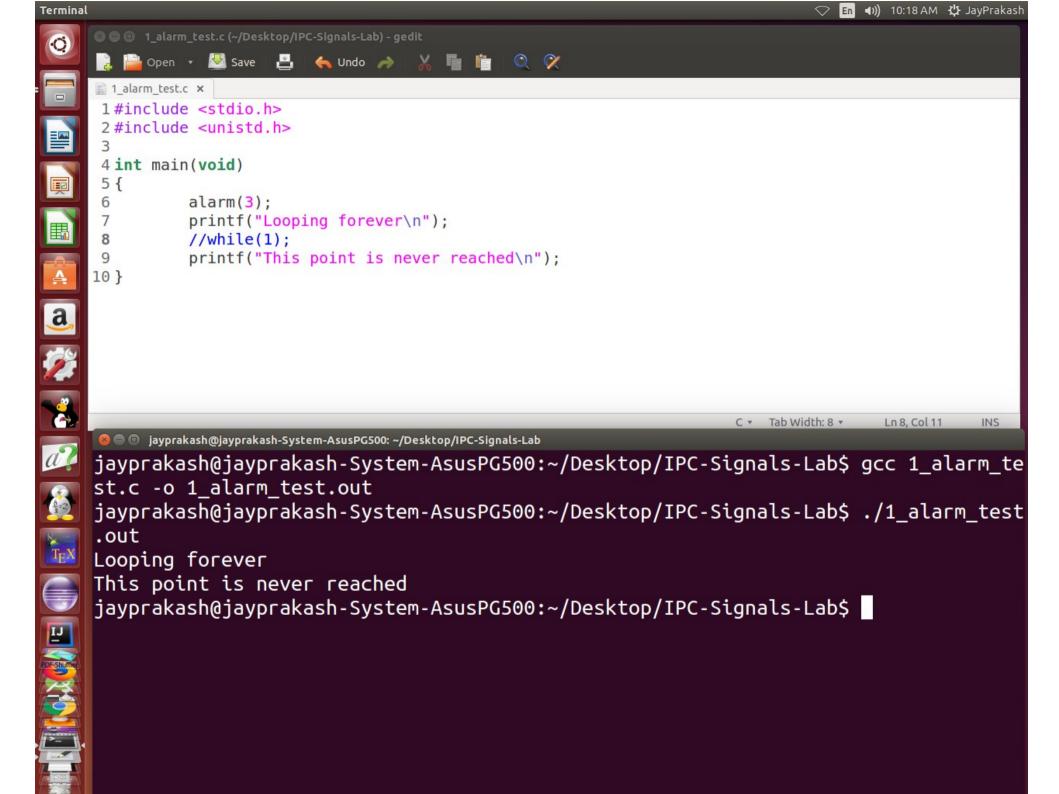
System Call: unsigned int alarm(unsigned int count)

- alarm() instructs the kernel to send the SIGALRM signal to the calling process after counting seconds.
 - If an alarm had already been scheduled, that alarm is overwritten.
 - If count is 0, any pending alarm requests are cancelled.
- The default handler for this signal displays the message "Alarm clock" and terminates the process.
- alarm() returns the number of seconds that remain until the alarm signal is sent

System Call alarm() example

 Set Alarm for 3 second with default action handler which will display a default message "Alarm clock" and exit the program <u>alarm test.c</u>





Handling Signals (Overriding Default Action)

System Call: void(*signal(int sigCode, void (*func)(int))) (int)

- signal() allows a process to specify the action that it will take when a particular signal is received.
 - sigCode specifies the signal number (as per the table shown in earlier slides) that is to be reprogrammed
 - func may be one of several values:
 - SIG_IGN, which indicates that the specified signal should be ignored and discarded
 - SIG_DFL, which indicates that the kernel's default handler should be used.
 - an address of a user-defined function, which indicates that the function should be executed when the specified signal arrives.

Handling Signals

- The valid signal numbers are stored in /usr/include/signal.h or /usr/include/bits/signum.h
 - The signals SIGKILL and SIGSTP may not be reprogrammed.
- signal() returns the previous func value associated with sigCode if successful; otherwise, it returns a value of -1
- signal() system call can be used to override the default action
- A child process inherits the signal settings from its parent during a fork()

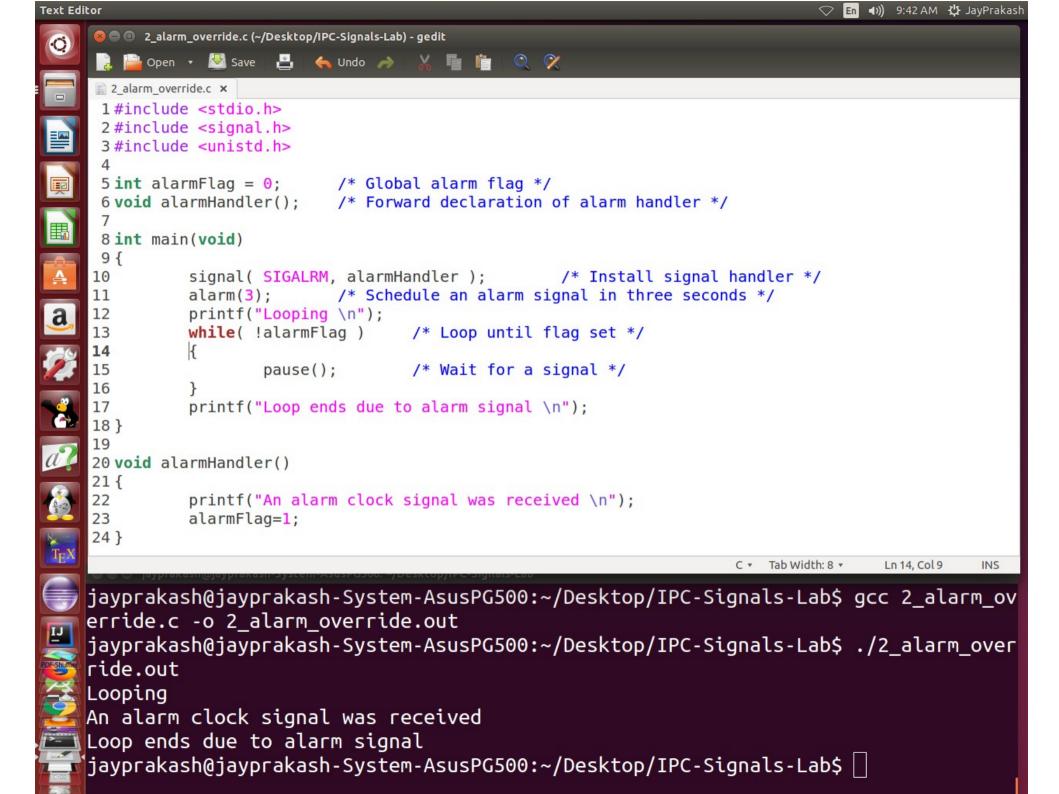
System Call pause()

System Call: int pause(void)

- pause() suspends the calling process and returns when a calling process receives a signal.
- It is most often used to wait efficiently for a signal.
- pause() doesn't return anything useful.

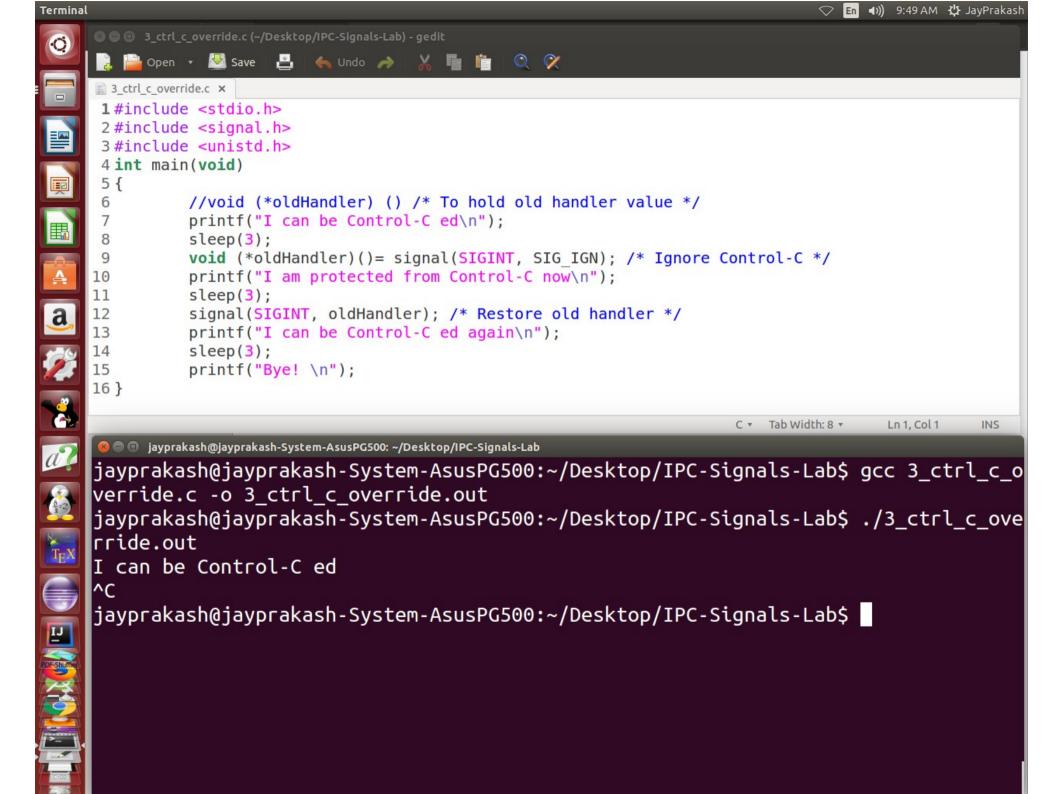
Catch SIGALRM of alarm() system call using signal() system call

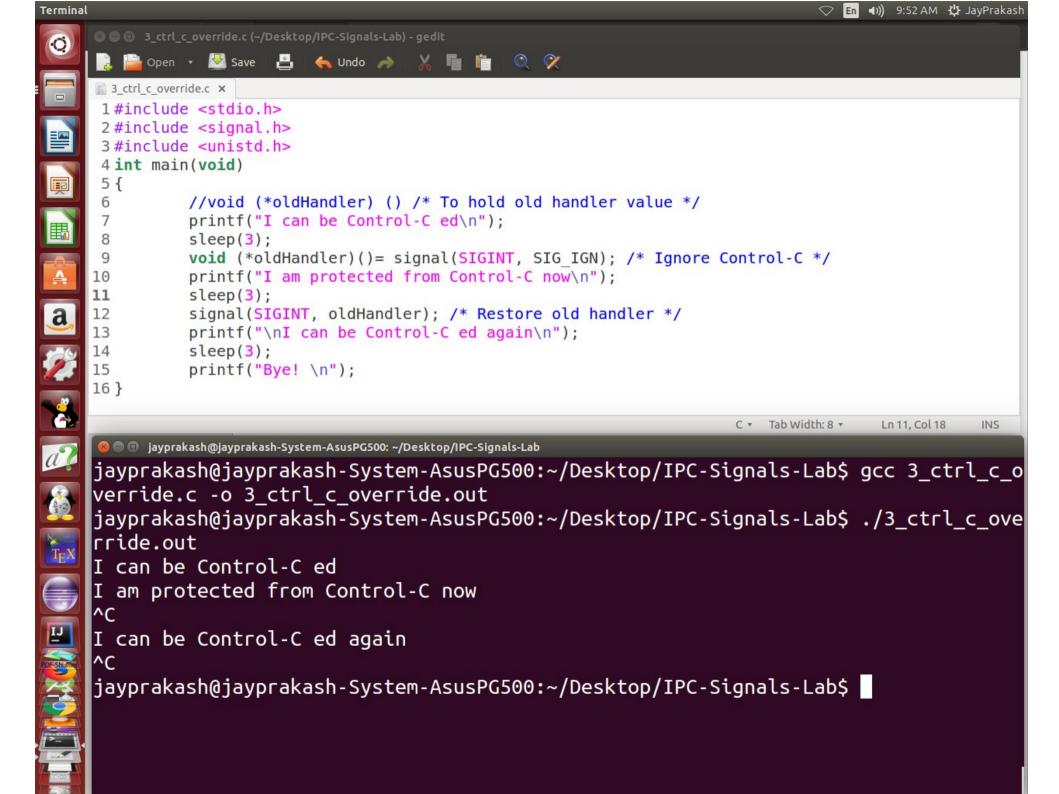
- Write your own handler to handle alarm signal SIGALRM
- Override default action using signal() system call
- alarm override.c

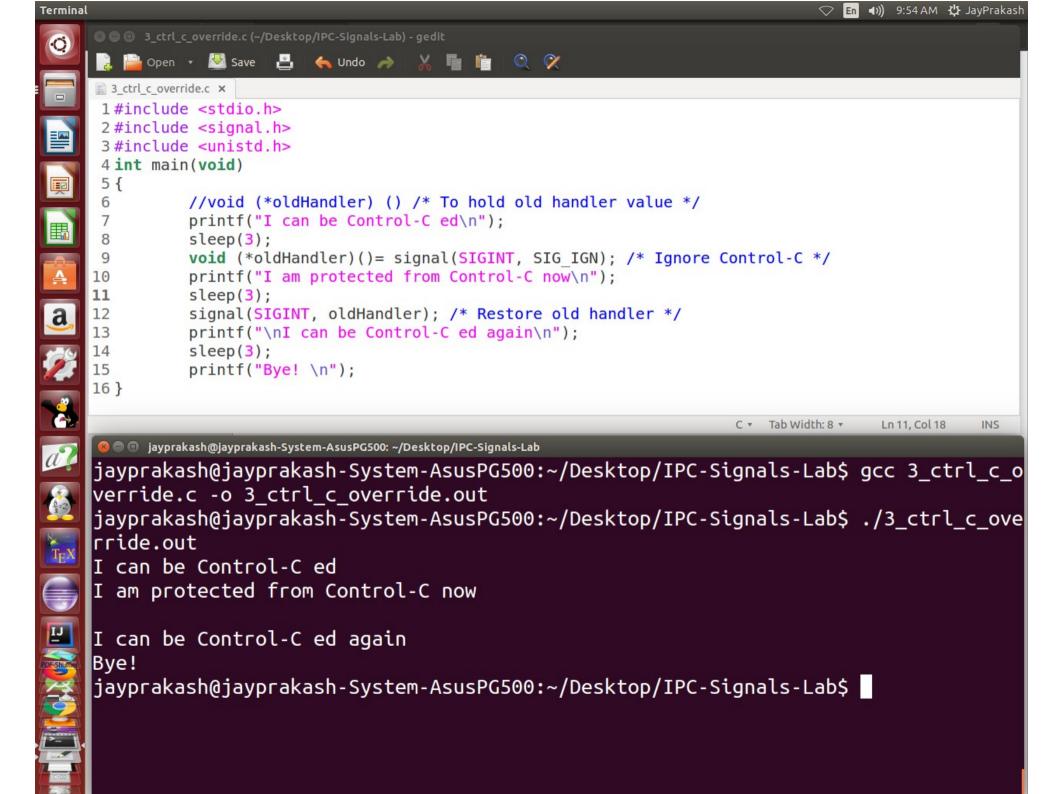


Protecting Critical Code from Ctrl-C attack

- Overriding may be used to protect critical pieces of code against Control-C attacks and other such signals.
- it can be restored after the critical code has executed.
- Here's the source code of a program that protects itself against SIGINT signals <u>ctrl c override.c</u>







- When you're in a shell and you execute a program that creates several children, a single Control-C from the keyboard will normally terminate the program and its children and then return you to the shell.
- In order to support this kind of behavior, UNIX introduced a few new concepts.
 - In addition to having a unique process ID number, every process is also a member of a process group.
 - Several processes can be members of the same process group.
 - When a process forks, the child inherits its process group from its parent.
 - A process may change its process group to a new value by using setpgid()
 - When a process execs, its process group remains the same.

- Every process can have an associated control terminal, which is typically the terminal where the process was started.
 - When a process forks, the child inherits its control terminal from its parent.
 - When a process execs, its control terminal stays the same.
- Every terminal can be associated with a single control process.
 - When a metacharacter such as a Control-C is detected, the terminal sends the appropriate signal to all of the processes in the process group of its control process.
- If a process attempts to read from its control terminal and is not a member of the same process group as the terminal's control process,
 - the process is sent a SIGTTIN signal, which normally suspends the process.

- When an interactive shell begins, it is the control process of a terminal and has that terminal as its control terminal.
- When a shell executes a foreground process:
 - the child shell places itself in a different process group before exec'ing the command and takes control of the terminal.
 - Any signals generated from the terminal thus go to the foreground command rather than to the original parent shell.
 - When the foreground command terminates, the original parent shell takes back control of the terminal.

- When a shell executes a background process:
 - the child shell places itself in a different process group before executing, but does not take control of the terminal.
 - Any signals generated from the terminal continue to go to the shell.
 - If the background process tries to read from its control terminal, it is suspended by a SIGTTIN signal.

Every process belongs to exactly one process group.

