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DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

EXAMINATION: END TERM**MONTH and YEAR: MAY 2025****Course: B.Tech.****Semester: 4th****Subject Code: CSE 223****Subject Name: Theory of Computation****Maximum Marks: 50****Duration: 03:00 Hour****Date: 14/05/2025****Time: 8:00 AM to 11:00 AM****Note:** Attempt all the questions. All questions carries equal marks.

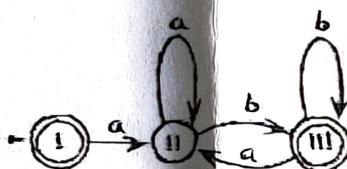
- 1 (a) Construct DFA for the following Language:

04 CO1

$L = \{w \mid w \in \{a, b\}^*, w \text{ starts with 'a' and has even length or } w \text{ starts with 'b' and has odd length}\}$

- (b) Consider following NFA that accepts language L.

04



Let L_1 be the language accepted by the same NFA by interchanging accepting and non-accepting states. Find L (Informally English language statement) and Clearly describe language $L \cap L_1$.

- (c) Show that $L = \{a^n b^m \mid n \neq m\}$ is non regular using pumping lemma.

02

- 2 (a) Consider the following context free grammar over alphabet {a, b, c}

 $G_1: S \rightarrow aSc \mid A$
 $G_2: S \rightarrow Sc \mid A$
 $A \rightarrow bA \mid \epsilon$
 $A \rightarrow Ab \mid B$
 $B \rightarrow aB \mid \epsilon$

The language generated by Grammar G_1 and G_2 are $L(G_1)$ and $L(G_2)$ respectively. Find the languages $L_1 = L(G_1) \cup L(G_2)$, $L_2 = L(G_1) \cap L(G_2)$ and $L_3 = L(G_1) - L(G_2)$. Also find and justify the type of language L_1 , L_2 and L_3 from the set = {Regular, CFL but not Regular, CSL but not CFL}.

- (b) Consider the following language L:

04

 $L = \{uvw^r \mid u, v, w \in \{0, 1\}^+, w^r \text{ is reverse of } w\}$

Find whether L is regular or not. If yes then construct finite automata or regular expression for it. Otherwise find the category of L and construct suitable automata/machine to recognize it.

- 3 (a) Find the Context Free Grammar for the following languages:

04 CO2

(i) $L = \{w \mid w \in \{a, b\}^* \text{ and each suffix of } w \text{ has at least as many } a's \text{ as } b's\}$

(ii) $L = \{w \mid w \in \{a, b\}^* \text{ and } \frac{n_a(w)}{2} \leq n_b(w) \leq \frac{3n_a(w)}{2}\}$

- (b) Show that if L is a DCFL, then there is a DPDA accepting the language

04

$L_1 = \{w_1 \# w_2 \mid w_1 \in L \text{ and } w_1 w_2 \in L\}$ where # is a symbol that does not occur in any string of L.

- 4 (a) Construct PDA for the following Language L
 $L = \{w_1 \# w_2 \mid w_1, w_2 \in \{a, b\}^* \text{ and either } w_2 \text{ contains suffix as } w_1^r \text{ or } w_1 \text{ contains prefix as } w_2^r\}$

04 CO2

- (b) Construct PDA equivalent to the following grammar $G(\{S, A, B, C\}, \{a, b\}, S, P)$
- $S \rightarrow bS / Aa$
 $A \rightarrow Ca / B$
 $B \rightarrow aA / Cb$
 $C \rightarrow Sb / \epsilon$

04 CO3

- 5 (a) Consider the following PDA $M = (\{q_0, q_1, q_2, q_3, q_f\}, \{0, 1, 2\}, \{0, 1, z\}, \delta, q_0, \{q_f\})$ with the following transitions:

$$\begin{aligned}\delta(q_0, a, z) &= \{(q_0, AZ)\}, \delta(q_0, a, A) = \{(q_0, AA)\}, \\ \delta(q_0, b, A) &= \{(q_1, \epsilon)\}, \delta(q_1, b, A) = \{(q_1, \epsilon)\}, \\ \delta(q_1, b, Z) &= \{(q_1, BZ)\}, \delta(q_1, b, B) = \{(q_1, BB)\}, \\ \delta(q_1, c, B) &= \{(q_2, \epsilon)\}, \delta(q_1, c, A) = \{(q_2, AA)\}, \\ \delta(q_2, c, B) &= \{(q_2, \epsilon)\}, \delta(q_2, c, Z) = \{(q_2, AZ)\}, \\ \delta(q_2, c, A) &= \{(q_2, AA)\}, \delta(q_2, d, A) = \{(q_3, \epsilon)\}, \\ \delta(q_3, d, A) &= \{(q_3, \epsilon)\}, \delta(q_3, \epsilon, Z) = \{(q_f, Z)\}\end{aligned}$$

Where Z indicate bottom marker of the stack. Find and justify the language accepted by the above PDA.

Show

CO4

- (b) Let L_1 and L_2 are two recursive languages. So that $L_1 \cup L_2$ and $L_1 \cdot L_2$ are also recursive.

04

- 6 (a) Construct Turing Machine for the following Language L
 $L = \{a^n x^m b^n y^m c^n \mid m, n \geq 1\}$

04 CO4

- (b) Consider the following Turing Machine for Language $L = \{a^n b^m a^n b^m \mid m, n \geq 1\}$. There are some transitions missing in TM on the places 1 to 5. Find all the missing transitions. (Note: Any location may have multiple missing transitions; B represents Blank symbol and head of TM initially points to B)

04 CO4

