# CS & IT





**Transaction & Concurrency** 

Control

**Discussion Notes** 



By- Vijay Agarwal sir

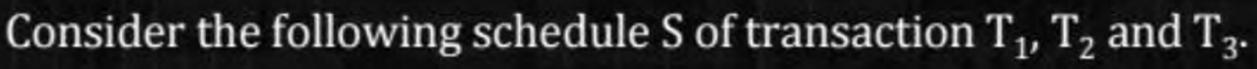


TOPICS TO BE COVERED

01 Question

02 Discussion







S:  $r_1(x)$ ;  $r_2(x)$ ;  $r_3(y)$ ;  $w_1(x)$ ;  $r_2(z)$ ;  $r_2(y)$ ;  $w_2(y)$ ;  $w_1(z)$ ;

Which one of the schedule below is the correct serialization of

the above.



$$T_2 \rightarrow T_1 \rightarrow T_3$$



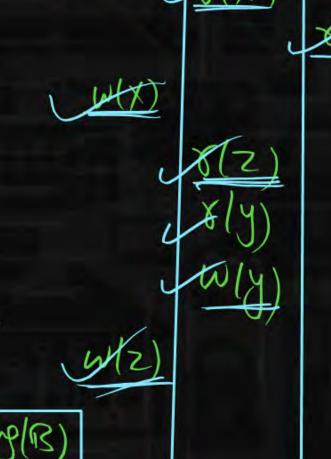
$$T_3 \rightarrow T_2 \rightarrow T_1$$



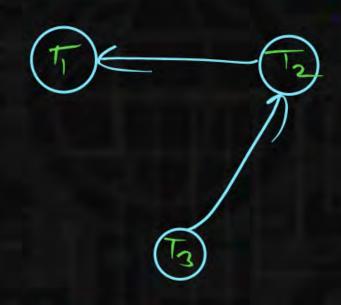
$$T_3 \rightarrow T_1 \rightarrow T_2$$



$$T_2 \rightarrow T_3 \rightarrow T_1$$

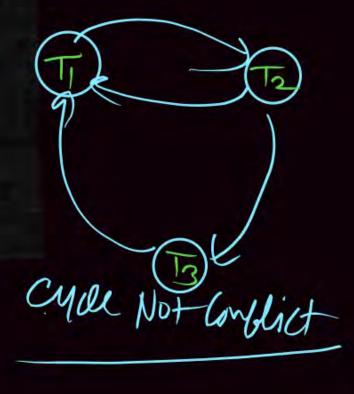


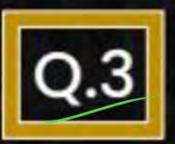
[MCQ]



Consider the transactions T<sub>1</sub>, T<sub>2</sub> and T<sub>3</sub> and the schedules S<sub>1</sub> and S<sub>2</sub> given below.  $T_1$ :  $r_1(A)$ ;  $r_1(C)$ ;  $w_1(A)$ ;  $w_1(C)$  $T_2$ :  $r_2(B)$ ;  $r_2(C)$ ;  $w_2(C)$ Grefit  $S_1$ :  $r_1(A)$ ;  $r_3(B)$ ;  $r_3(A)$ ;  $r_2(B)$ ;  $r_2(C)$ ;  $r_3(B)$ ;  $r_3(C)$ ;  $r_1(C)$ ;  $r_1(C)$ ;  $r_1(A)$ ; );  $r_3(B)$ ;  $r_2(B)$ ;  $r_3(A)$ ;  $r_1(C)$ ;  $r_2(C)$ ;  $w_3(B)$ ;  $w_1(A)$ ;  $w_2(C)$ ; Which one of the following statements about the schedule is TRUE? Only S<sub>1</sub> is conflict serializable. Only S2 is conflict serializable. Both S<sub>1</sub> and S<sub>2</sub> are conflict serializable. Neither  $S_1$  nor  $S_2$  is conflict serializable.

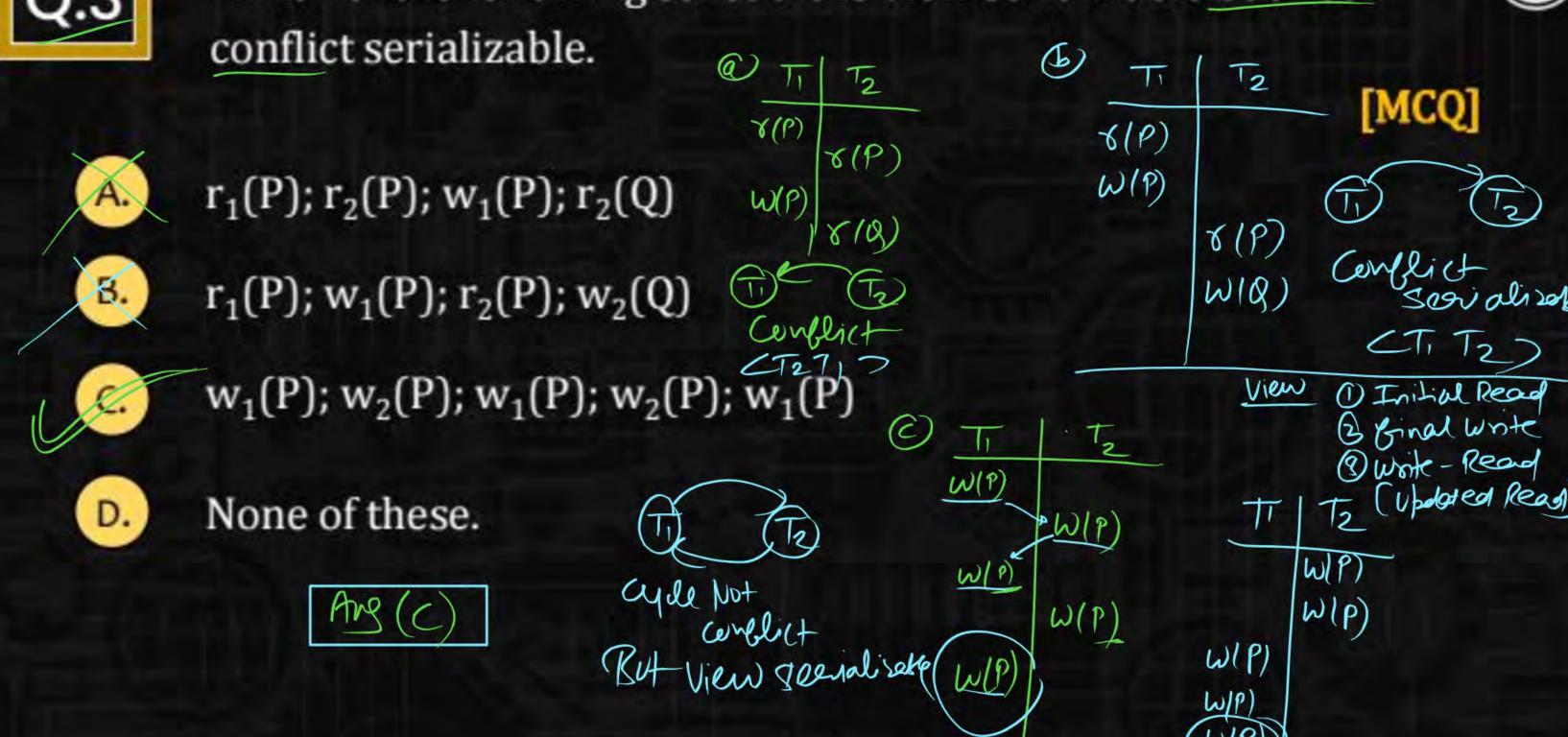
Consider the transactions  $T_1$ ,  $T_2$  and  $T_3$  and the schedules  $S_1$  and S<sub>2</sub> given below.  $\Gamma_1$ :  $r_1(A)$ ;  $r_1(C)$ ;  $w_1(A)$ ;  $w_1(C)$  $T_2$ :  $r_2(B)$ ;  $r_2(C)$ ;  $w_2(C)$  $T_3$ :  $r_3(B)$ ;  $r_3(A)$ ;  $w_3(B)$  $r_1(A)$ ;  $r_3(B)$ ;  $r_3(A)$ ;  $r_2(B)$ ;  $r_2(C)$ ;  $w_3(B)$ ;  $w_2(C)$ ;  $r_1(C)$ ;  $w_1(A)$ ; );  $r_3(B)$ ;  $r_2(B)$ ;  $r_3(A)$ ;  $r_1(C)$ ;  $r_2(C)$ ;  $w_3(B)$ ;  $w_1(A)$ ;  $w_2(C)$ ; Which one of the following statements about the schedule is TRUE? Only S<sub>1</sub> is conflict serializable. Only S2 is conflict serializable. Both S<sub>1</sub> and S<sub>2</sub> are conflict serializable. Neither S<sub>1</sub> nor S<sub>2</sub> is conflict serializable. M(B) W(C)

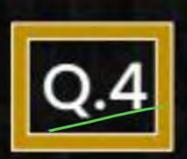




Which of the following schedule is view serializable but not







## Consider the following transactions T<sub>1</sub> and T<sub>2</sub>:



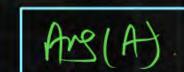
$T_1$	T <sub>2</sub>	
Read(A);	Read(A);	
Update $A = A + 100$ ;		
	Update A = A - 50;	
Write(A);		
	Write(A);	

It Two Dibbellent transaction Perform write Operation of between these 2 write operation there is No Read operation, then there 2" write operation,

overworten me tisst write operation

Lost of updation (ti)

Lest update Problem



Lost update problem

В.

Dirty read problem

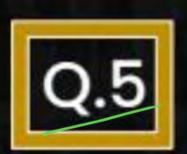


Unrepeatable read problem

The above transaction has \_\_\_\_



Incorrect summary problem



#### Consider the following schedule.



Time	T <sub>1</sub>	T <sub>2</sub>
t <sub>o</sub>	Read Item(A);	
$t_1$		Read Item(A);
t <sub>2</sub>	La I TE	$A = \underline{A + X};$
t <sub>3</sub>		Write Item(A);
$(t_4)$	Read Item(A);	

Which of the following concurrency problem exists in the above given schedule?

[MCQ]



Dirty read



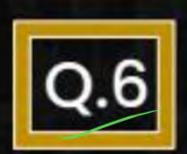
Unrepeatable read



Lost update

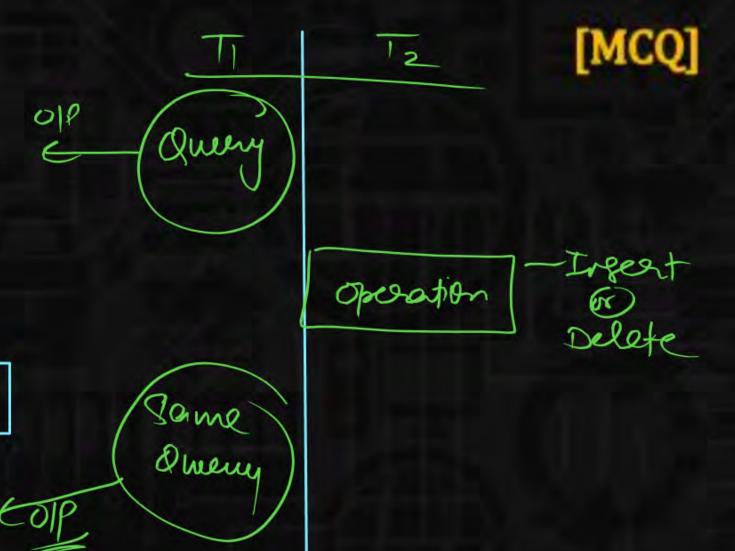


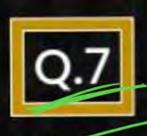
Both a and b



Problem occurs when a transaction reads data from a database, then another transaction reads the same database data, and this particular data is deleted by an operation of the first transaction.

- A. Dirty read
- B. Unrepeatable read problem
- Phantom read
- D. Lost update problem

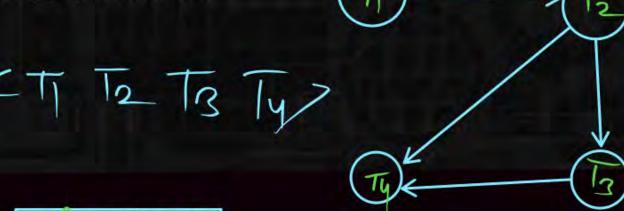




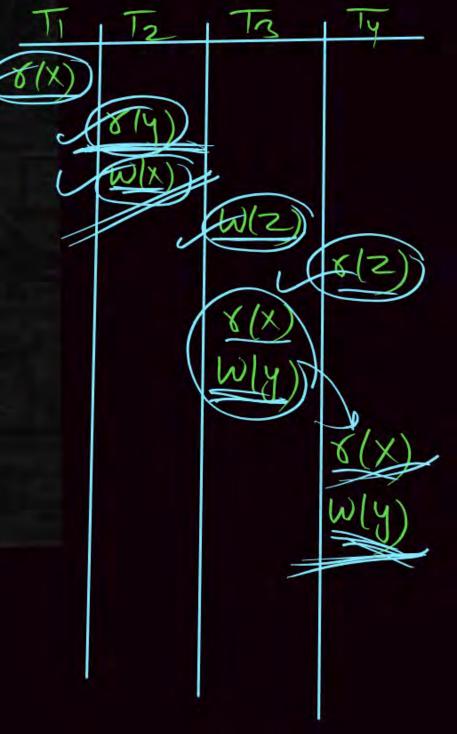
Consider a schedule S:

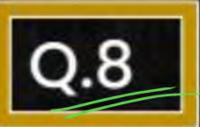
 $r_1(x)$ ,  $r_2(y)$ ,  $w_2(x)$ ,  $w_3(z)$ ,  $r_4(z)$ ,  $r_3(x)$ ,  $w_3(y)$ ,  $r_4(x)$ ,  $w_4(y)$ Choose the correct statements for the above schedule S.

- A. The schedule S is not serializable.
- The schedule S is conflict serializable with schedule S as  $T_1 \rightarrow T_2 \rightarrow T_3 \rightarrow T_4$ .
  - C. The schedule S is not view serializable.
  - D. None of the above.



[MCQ]





Consider the below schedule.



**S:**  $r_1(A)$ ,  $r_2(B)$ ,  $w_2(A)$ ,  $r_3(A)$ ,  $w_1(B)$   $w_3(A)$  Choose the correct statement from the following.





S is conflict serializable schedule.



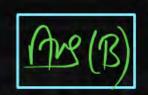
S is not conflict serializable schedule.

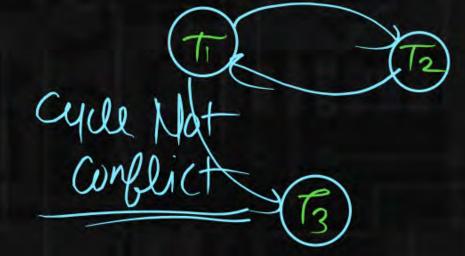


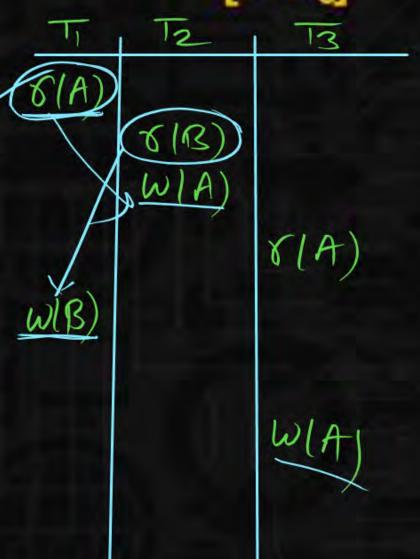
S may or may not be view serializable schedule.



None of these







### View Serializable

3 Toonsaction => 31 => 6 Serial Schedule

Dunny

- (1) Fritial Read
- 2 final write
- (3) write-Read (Updated-Read)

CT, T2 T3)

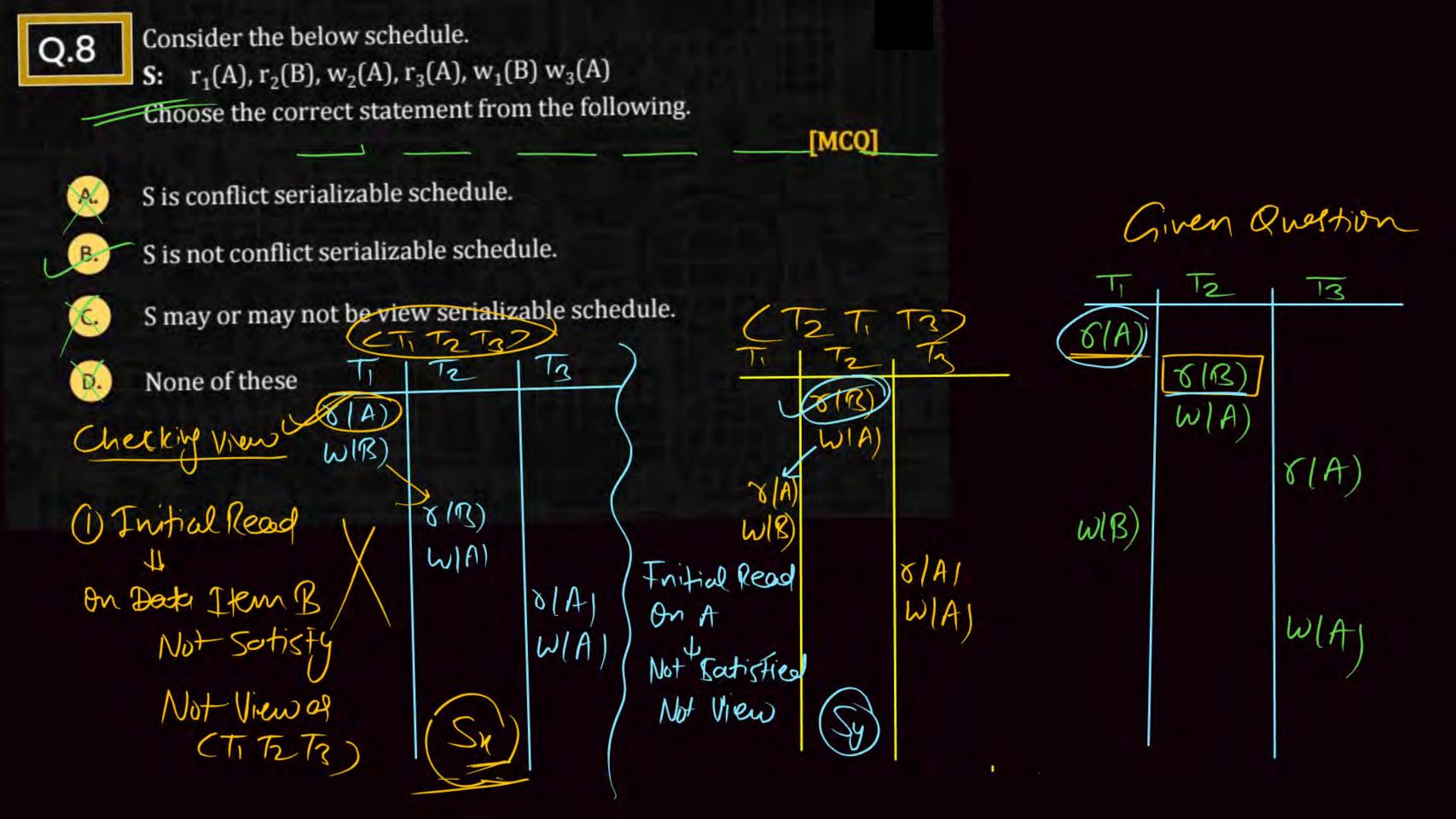
CTITS T2)

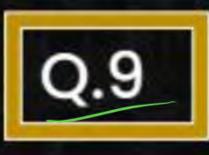
CT2 T1 T3)

CT2 T3 TID

(T3 T2 T1)

CT3 TI TZ)





#### Choose the correct statements from the following.





To test view serializability we make use of precedence graph



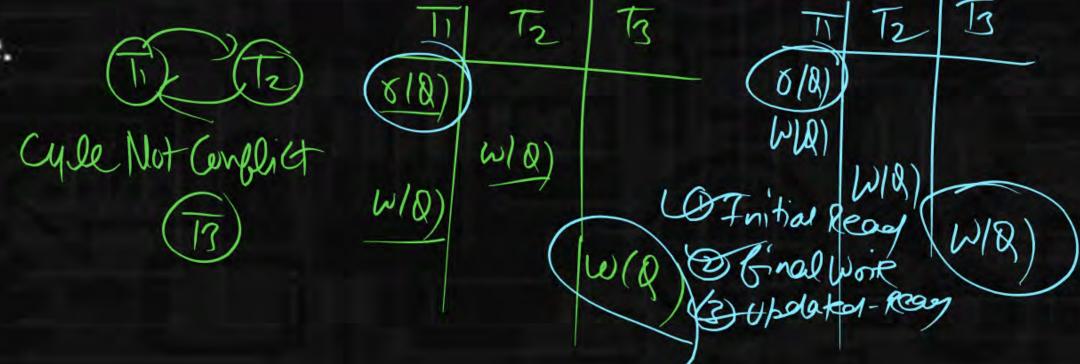
To test conflict serializability we make use of precedence graph.



If there exists no blind write and the schedule is not conflict serializable then we can conclude that it is not view serializable.



All of the above.



Precdence graph To Test Conflict Secriptizablity CNC => Cycle Not Conflict

5 May 60 May Not be View Schiolizable



