

CS & IT ENGINEERING

Operating System

Process Synchronization /
Coordination

DPP 07 (Discussion Notes)



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TOPICS TO BE COVERED



01 Question

02 Discussion

Q.1

What is a compulsory step before using semaphore?

[MCQ]



A.

Deciding final value of semaphore X

B.

Initialization of semaphore

C.

Defining number of operations to be performed

D.

All of the above. X

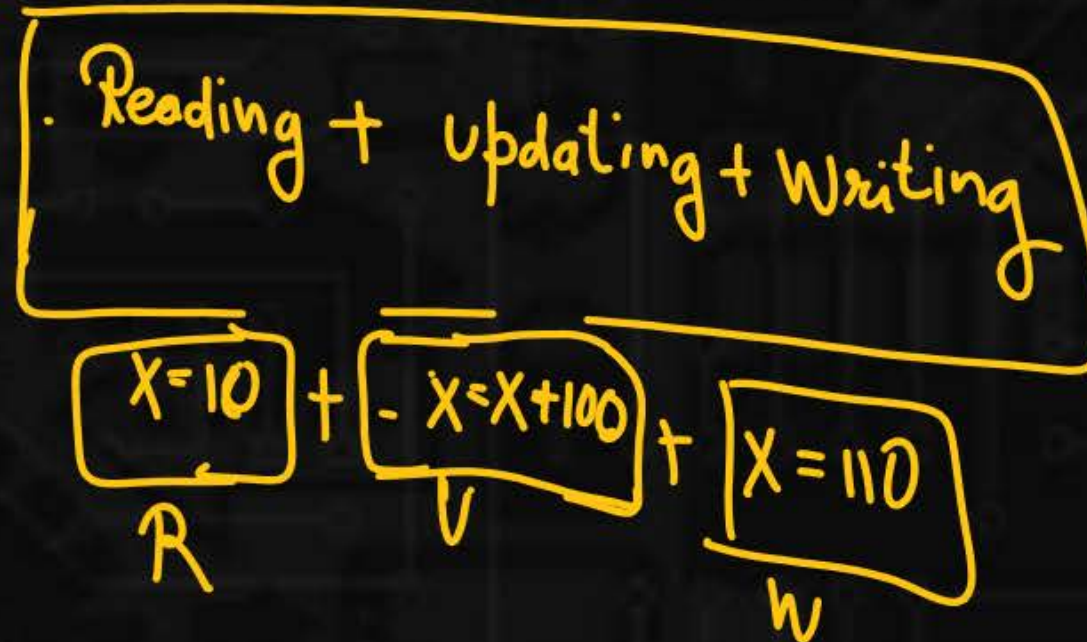
Q.2

Consider the following statements, which of the following is correct?



[MCQ]

- A. Semaphore are atomic in nature and implemented in user's mode.
- ☒ B. Semaphore are atomic in nature and implemented in kernel mode.
- ☒ C. Semaphore are non-atomic in nature and implemented in user's mode.
- ☒ D. Semaphore are non-atomic in nature and implemented in kernel mode.



Variable: `int x = 10`
110

Q.3

If a semaphore's value is "- 3", then what does magnitude of "-ve" value indicate ____?



[MCQ]

A.

Number of successful up operation.

B.

Number of successful down operation.

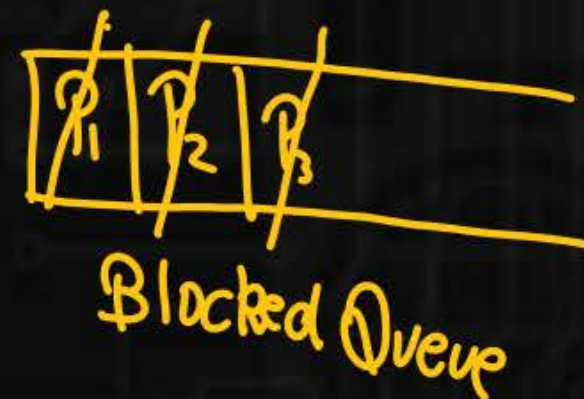
C.

Number of blocked processes.

D.

Number of unblocked processes. X

$$S = -3$$



Q.4

Processes x_1 and x_2 uses flag_critical in the following function to achieve mutual exclusion. Assuming flag_critical is initialized FALSE initially.



[MCQ]

get_access

```
{
  1.  $x_1$ 
  2.  $x_2$ 
  if (flag_critical == FALSE)
  {
    -false-
    flag_critical = TRUE;
    Critical_section();
    flag_critical = FALSE;
  }
}
```

x_1 is preempted

A.

(i) is true and (ii) is false. ✗

B.

(ii) is true (i) is false. ✓

C.

Both (i) and (ii) are true.

D.

Both (i) and (ii) are false.

Consider the following statement:

- (i) The above routine may lead to deadlock. ✗
- (ii) It is possible for processes x_1 and x_2 to access critical section concurrently. ✓

1. M.E

2. Hold & Wait

3. Circular Wait

4. No preemption.

Q.5

Consider the code given below, used by the processes x_1 and x_2 to access critical section. The initial value of shared Boolean variable P and Q are false



[MCQ]

x_1
while (P == Q);
<critical section>
P = Q;

x_2
while (P != Q);
<critical section>
P = !Q;

Initialized Value.

P = Q = False

P = True

False

Select the true statements from the following:

- (i) Process x_1 can go into critical section just after one entry by process x_2 into its critical section. ✓
- (ii) Mutual exclusion is not ensured. ✗
- (iii) Process x_1 can go into critical section many times without single entry of x_2 into its critical section. ✗
- (iv) None of the above

$x_2 x_1 x_2 x_1 x_2 x_1 x_2 \dots$

A.

(i) & (ii) ✗

C.

Only (i) ✓

B.

(ii) & (iii) ✗

D.

(i), (ii) & (iii)

Q.6



Consider the two function P_i and P_j that share a variable Q with an initial value '3' execute concurrently:

$P_i()$	$P_j()$
{ 3 1. $R = Q * 2 := 6$ 2. $Q = R;$ }	{ 3. $S = Q + 1 := 4$ 4. $Q = S;$ }
$Q=6$	$Q=4$

- $\checkmark \checkmark \checkmark \checkmark \Rightarrow Q=8$
• $\checkmark \checkmark \checkmark \checkmark \Rightarrow Q=7$
- $\checkmark \checkmark \checkmark \checkmark \Rightarrow Q=4$
• $\checkmark \checkmark \checkmark \checkmark \Rightarrow Q=6$

[NAT]

Initial $Q = 3$

What are the different possible value for variable Q at the end of execution of both process P_i and P_j ?

4

$Q = 4, 6, 7, 8$

Q.7



Match the following statements

List I	List II
A. Critical section	1. Ensuring that only one process can execute C.S.
B. Synchronization	2. atomic operation are used to ensure co-operation between processes.
C. Mutual exclusion	3. Section of code that only one process can access at once.

[MCQ]

Matches:

A-3, B-2, C-1

A.

A

B

C

1

2

3

B.

3

2

1

C.

2

3

1

D.

1

3

2

Q.8



Let S be a binary semaphore variable. Let $S = 1$ initially.

Assume that no blocked processes exist in the system. The following operations are performed on semaphore S .

6 P, 8 V, 12 P, 11 V, 19 P

The number of blocked processes after executing these operations are 19.

[NAT]

$S = 0$

$S = 1$

S	1 0	0 1	1 0	0	0
P/V	6P	8V	12P	11V	19P
Blocked Processes	5	0	11	0	19

