## CS & IT



Database Management System

File Org & Indexing

DPP - 03 Discussion Notes



By- Vijay Agarwal sir



TOPICS TO BE COVERED

01 Question

02 Discussion



The order of a leaf node in a B+ tree is the maximum number of (value, data record pointer) pairs it can hold. Given that the block size is 1K bytes, data record pointer is 8 bytes long, the value field is 10 bytes long and a block pointer is 6 bytes, then what is the order of the leaf node?



A. 53

c. 55

P+lay + P+Rp+1Bp & Block Size

PX10+PX8+LX6 < 1024B 10P+8P+6 < 1024 18P < 1018

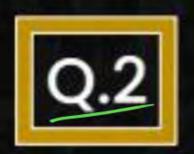
$$P = \left| \frac{1018}{18} \right| = \left( 56 \right) Amg$$

B. 54

56

Black Size: 1024B, Rp=8B, ky=10, Bp=6Byte

Ag (D).



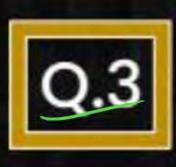
The order of a node in B tree is defined as the number of worders it can hold. What is the maximum number of keys that a B tree of order 4 and height 4 can have? \_\_\_\_\_\_ [NAT]

(Assume that the height of a root node is 1)



Or or or	(heignes/Level	max # Node	Max # Bp	max # 100yg		
De T	0/1	L	P	P-1	(4-1)	=3
	1/2	P	P2	P(P-1)	4×3	=  2
	2/3	p2	P3	bs (b-1)	42 x 3	- 48
Civento		P3	PY	p3(P-1)	43 X3	= 19
Criven						2

OPDER[P)=4



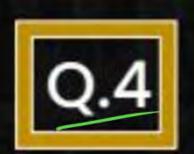
Given a block can hold either 3 records or 10 key pointers. A W database contains P records then how many blocks do we need to hold the data file and the dense index?

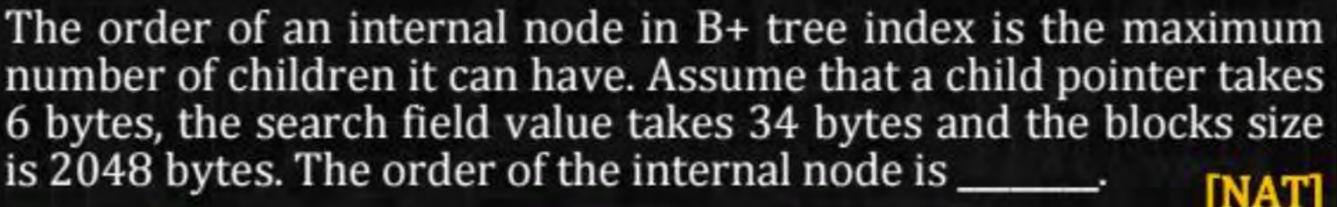
$$\frac{P}{30}$$

$$\frac{P}{3}$$

$$\frac{13P}{30}$$

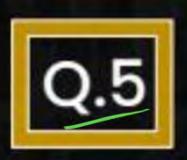
$$\frac{13P}{30}$$
To stone Record #Blocks =  $\frac{P}{10}$ 
Required =  $\frac{P}{3}$ 



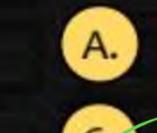




$$P \times RP + (P-1) \log C = Resct$$
 $\Rightarrow P \times G + (P-1) \approx C = 2048$ 
 $GP + 34P - 34 \leq 2048$ 
 $40P \leq 2082$ 
 $P = \left| \frac{2082}{40} \right| = \left( \frac{52}{52} \right) Ang$ 

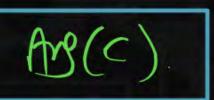


Assume a disk with block size B = 1024 Bytes, A block pointer is  $P_B = 12$  bytes long and a record pointer is  $P_R = 18$  bytes long. A file has 1,00,000 patients records of size 100 bytes. Suppose the file is ordered by the key field PID and we want to construct a secondary (dense) index on non-key field DeptID (14 bytes), then minimum of how many blocks are required to store index file assuming an unspanned organisation?



3000

3125



в. 3100



None of the above

Block Size = 1024 Byte R=18Byte key=14Byte, #Records=100.000 18+14=32 Block bacter of Index file = RUUGSize = 1024 = 25 732 Index entries Per Blocks Ly Total # Index entones = 100,000 (#DR Records) Total # of Index Blodg = [100,000] = 3125 Index Block.

The order of a node in B tree is the maximum number of block

pointers it can hold. Given that the block size is 2K bytes, data a block pointer is 5 bytes long. The best possible order of B tree

node is 93 Ars

[NAT]

OPDERP: 
$$P + (P-1) | ceys + (P-1) | P \le R | cove size$$

$$\Rightarrow P + 5 + (P-1) | q + (P-1) | s \le 2048$$

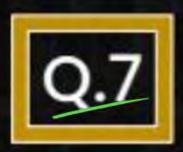
$$5P + 9P - 9 + 8P - 8 \le 2048$$

$$22P - 17 \le 2048$$

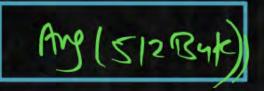
$$22P \le 2065$$

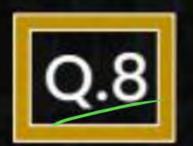
$$P = \frac{|2065|}{22} = 93 | Ang$$





The order of a leaf node (P) in a B+ tree is the maximum number of (value, data record pointer) pairs it can hold. Given that P=36, data record pointer is 8 bytes long, the search field is 6 bytes long and a block pointer is 8 bytes long. The permissible block size is 5/2 Byte [NAT]





## (Assume that the level of root node is1)



The order of different nodes in B<sup>+</sup> tree/B tree are given as-2 to P block pointers in root node.

$$\left\lceil \frac{P}{2} \right\rceil$$
 to P block pointers in internal node.

$$\left\lceil \frac{P}{2} \right\rceil$$
 -1 to (P-1) keys in leaf node.

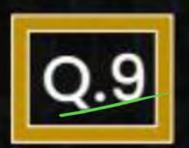
Let a and b be

The minimum number of keys in

B tree and B+ tree node of order

$$P = 5$$
 and level = 5. The value of  $(a + b)$  is

	Height/Lover	Min #Node	min & Bp	min #1 ceys.		
7	0/1	1	2	1	1	= 4
(P=5) [P2]=[5/2]	1/2	2	2 [2]	2(127-1)	2 x 2	(ca)
43	2/3	2[2]	2[3]	2[2] ([2]-1)		= 12 /2 (2)
P2-1	3/4	2527	2[1/2]	2[2]([3]-1)	2X(3X3)X	2 = 36
=3-	4/	2/2/3	2(12)	2 [3] ( [2]-1)	$2 \times (3 \times 3 \times $	161)
					b= 108	In Bt Tree all keys cue available in (lost (eve) leas Norde



(Assume that the level of root node is1)



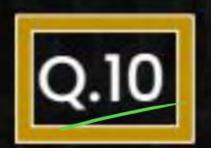
The order of different nodes in B<sup>+</sup> tree/B tree are given as-2 to P block pointers in root node.

$$\left\lceil \frac{P}{2} \right\rceil$$
 to P block pointer is internal node.

$$\left[\frac{P}{2}\right]$$
 -1 to (P-1) keys in leaf node.

Let a and b be the maximum number of keys in B tree and  $B^+$  tree node of order P = 5 and level = 5. The value of (a + b) is

Heigns (evel)	max thode	Max #130	(max # lays			
0/1	L	P	(P-1)	4	= 4	ORDER = (5=P)
	Pei	P2	P(P-1)	5×4	= 20	(p-1)=4
4/2	02	p3	p2(p-1)	5×5×4	- 100	
2/3	P					
3/4	b3	PY	P3(P-1)	5x5x5x4	= 500	
4(5)	py					
15)		PS	P'(P-1)	2x2x2x2x0		
			BTO	e a=312	3124	
			R <sup>†</sup> T			lable attest
				3:500	oll keys avois	().



Consider the keys (1–5000) are going to be interested into a B<sup>+</sup> tree. Assume, all the order are available before insertion. The orders P for B+ tree node is defined as-

2 to P pointer for root

 $\frac{P}{2}$  to P pointer for another node.



Maximum possible level, it we have non # key in Abode= (500)=(1250 Node) Minimum # 1ceys = (B) -1 = (B) -1 = 5-1-(4) (3)=1860c Min #BP= (3)=(3)=(5) (10) = Shage (SD) - 10 NOUR (2507 = 50 Node 1250 / 250 Node

250 Node

## Q.11 Consider the following statements:

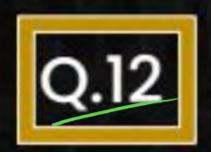


S<sub>1</sub>: In a B+ tree, data pointers are stored only at the leaf nodes of the tree.

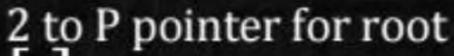
S<sub>2</sub>: In a B<sup>+</sup> tree, the leaf node has an entry for every value of the search field, along with the data pointer to the record.

Choose the correct statements.

D. Neither S<sub>1</sub> nor S<sub>2</sub> is true



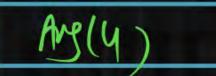
Consider the keys (1– 5000) are going to be interested into a B<sup>+</sup> tree. Assume, all the order are available before insertion. The orders P for B<sup>+</sup> tree node is defined as-



 $\frac{P}{2}$  to P pointer for another node.

The minimum possible levels in a B+ tree index for P = 9 is 4.

(Assume that level of the root node is 1)



Minimum # levels => Maximum Number of keys = (8)

4 Level Required.



