CS & IT



ENGINERING

Database Management System

File Org & Indexing

DPP - 01

Discussion Notes



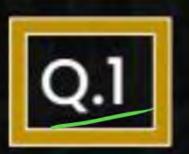
By- Vijay Agarwal sir

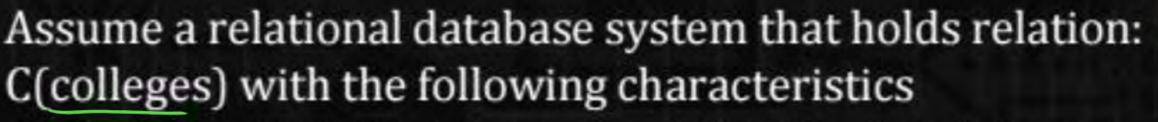


TOPICS TO BE COVERED

01 Question

02 Discussion







- Records are stored as <u>fixed length</u>, fixed format records, length is 256 bytes.
- There are 16384 records.
- Records contains key attribute CollegeNumber (C.N), length 22 bytes and other fields.
- Unspanned organization is used to store the information or record.

Let's suppose we want to build a sparse primary index on C.N then how many numbers of 4096-byte blocks are needed to store the primary index when block pointer size is 10 bytes ____?



7





8



9



Record Size = 256 Byte ley = 22 Byte & Bp=10Byte # Rewords = 16,384 Records Block Size = 4096 Byte

fixed Length unspanned 4 SPARSE Primary IndexU

Block foctor of DB File = | Block Size | = 1096 Byte = 212 - 24)

(BFDB) [Unsponned] = | Record Size | = 250 Byte = 28 - 24)

> 16 Record Per Blocks

Total # of DB Block = [16,384] = 214 = 210 = 1024 DB Block

One Index Record size: Size of + Size of = 22+10 = 32 Byte = 25 Byte

Block Size = 4096 Byte. Kay Block Binds = 22+10 = 32 Byte

Block factor of Index file = 4096 B = 212 = 27 = 128 Index Entries

[Bf.]

[Bf.]

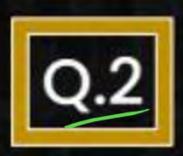
[Bf.]

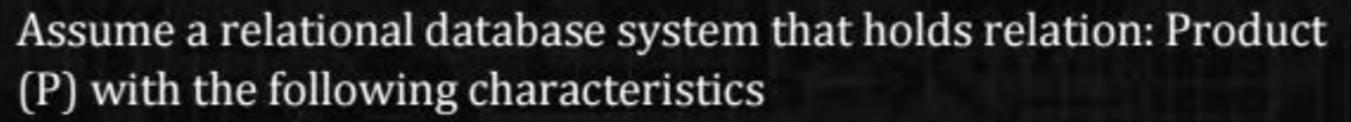
Primary Index > SPARSE

BG = 126 Index Entries Per Block

Sparse Index = Total # Index Entries = # of DB Blows = 1024

=) 8 Index Block. Ang







- Records are stored as fixed length, fixed format records, with the length of 256 bytes.
- There are 262144 records.
- Records contain attribute P.I (The identifier of the product involved), with the length <u>24 bytes</u>, and an attribute P.C (the cost of product), with the length <u>32 bytes</u> and other fields.
- Unspanned organization is used to store the record.

Assume that we want to build a dense secondary index on P.C, then how many numbers of 4096-byte blocks needed to store the dense secondary index. When record pointer size is 32 bytes?

Ang (4096

[NAT]

Rewords = 262144 [218], P.I = 24 Byte
Records Record Size = 256 Byte BOOUSIX=4096B BP=32 Byte Product Cost = 32 Byte (P.1) = key. Secondary > Denge One Index Record Size = Size oblicey + Bp = 32 + 32 = Gu Byte Rhode Size = 4096 Byte 4096 = 26 = 26 = 64 Index 64 = 26 = 26 = 26 = 64 Index entoing fer Blown Block factor of Index file Secondary Index Denge: No. of Index Entries = Number of DB Records = 262 144 [218]
Records

Total #Index Blocks = (262 144) = 218 = 212

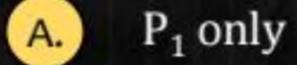
= 4096 Index Block

Are

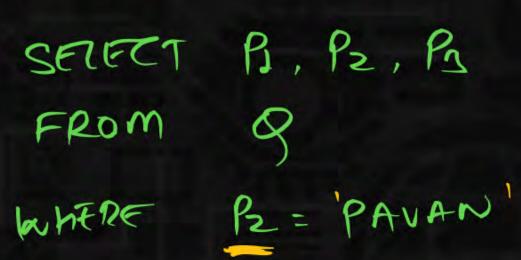


Consider a SQL statement SELECT P_1 , P_2 , P_3 from Q WHERE P_2 = 'Pavan' is frequently executed, which column(s) should be considered for indexing based only on the statement itself?

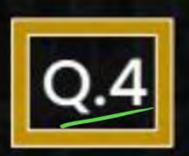




- B P_2 only
 - P_3 only
 - D. P_1 , P_2 and P_3



Ang (B)



Consider the following specification of system-

Disk block size = 2048 bytes

Block pointer size = 16 bytes

Record pointer size = 20 bytes long

file contains 30,000 records.

Each record of the file has the following fields: An extra/additional byte is used per record to represent end of the record.

What is the block factor of the database file assuming unspanned file organization?

1	P	
1		[MCQ
	AI CD	MCQ

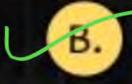
Fields	Size (in
	Bytes)
EmpName	5
EmpNum	10
DeptNum	9
Addr	20
PhNum	9
DOB	1
Sex	1
Job	3
Sal	5

63+1Byk

	16
A.)	10







D.

64

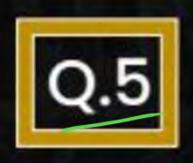
Record Size = 64 Byte

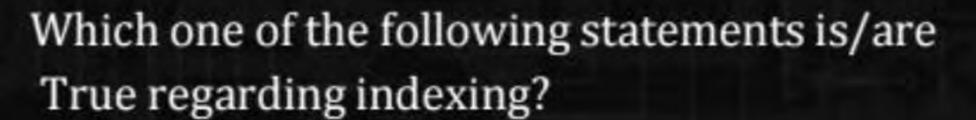
Block Size = 2048 Byte

Black factor of DB file = Black Size

(unspanned) Record Size

= 2048 Ryte = 2" = 25 64 Ryte = 25 = 25 32 Record to Block









A database file can contain multiple clustered indexes.



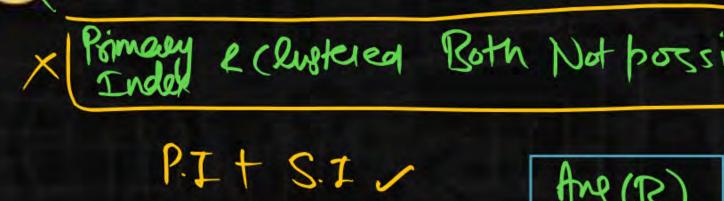
A database file can consist of only one clustered index with multiple secondary indexes.



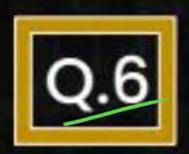
A database file can consist of multiple primary indexes.



A database file can consist of both primary and clustered index.



one Primary Index bossible





Consider a database of fixed-length records stored as an ordered file. The database has 25,000 records with each records being 100 bytes, of which the non-key attribute on which clustering index is formed occupies 10 bytes. The data file is completely block aligned.

Suppose, block size, of the file system is 512 bytes and a pointer to the block occupy 5 bytes. You may assume that a binary search on an index file of b block may take $\lceil \log_2 b \rceil$ accesses in worst case.

Given that a cluster consumes 2 blocks, the number of block accesses required to identify the desired data in the worst case is

[NA

#Records= 25,000 Record Size = 100 Byte key=10B Bp=5 Byte Block Size = 512 Byte Each Clusta Congume 2 Blocks Block forter of DR file = Block Size - 512B = 5 Record

Record Size - 512B Pel Block

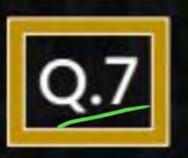
Pel Block Total # DB Block - [25000] = 5000 DATA BLOCK Each Cluster Congune = 2 Blooks Total Number of Index Entries = 5000 = 2500 Entries
One Index Record Size = 10+5 = 15 Byte

Block factor of Index Gile:
$$\left[\frac{512B}{15B}\right] = 34$$
 Index Entries Pal Block

Total # Index Block = $\left[\frac{2500}{34}\right] = 74$ Index Block

Blocks Access Requised = $\left[\frac{9000274}{11}\right] + \frac{11}{11}$

= $\left[\frac{9}{11}\right]$ App



Consider the following statements-



S₁: If the records of a relation X are physically ordered over a non-key field P and an index is build over the key-field of relation X, then the index is necessarily a secondary index over key attribute.

S₂: More than one secondary indexes are possible.

Which of the given statement(s) is/are CORRECT?

A. S_1 only

 S_2 only

Both S₁ and S₂ only

D. Neither S₁ nor S₂

