

# **A hybrid VND method for the split delivery vehicle routing problem**

**Yury Kochetov**

Sobolev Institute of Mathematics,  
Novosibirsk State University,  
Novosibirsk, Russia

*jointly with Alexey Khmelev*

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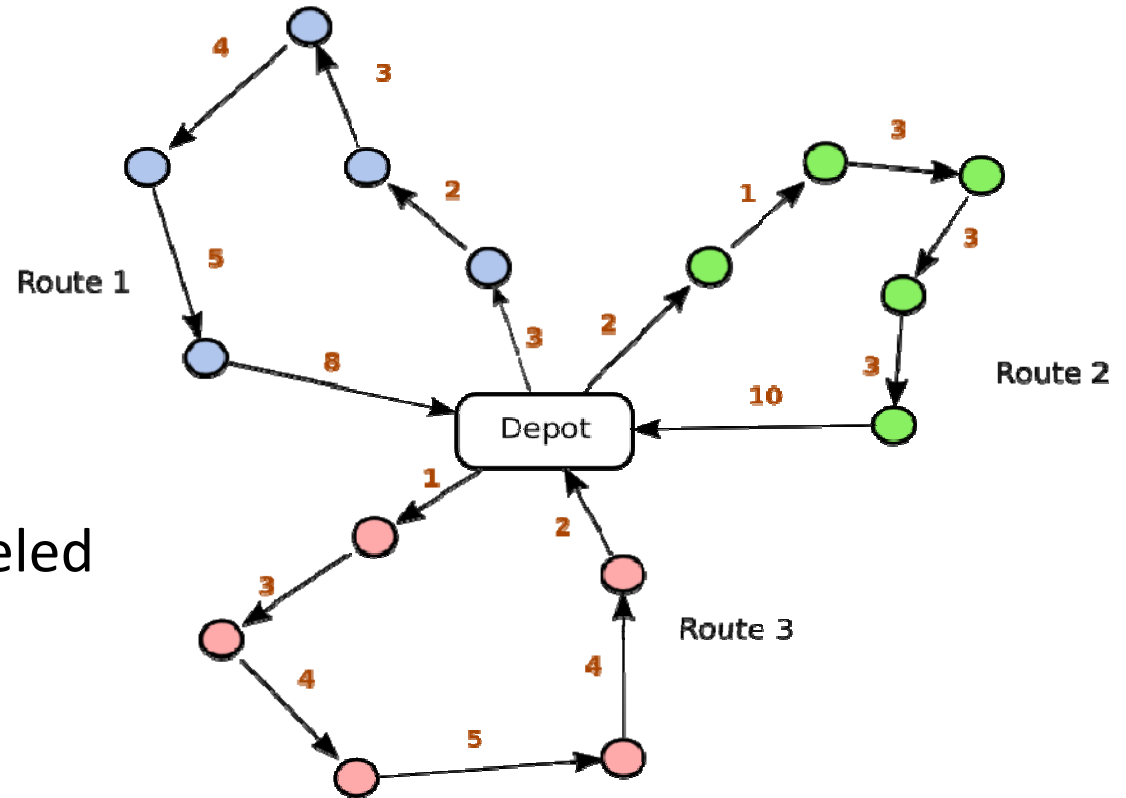
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# Overview

- SDVRP formulation
- Main idea of algorithm
- Coding solutions
- Variable neighborhood descent
- Tabu Search procedure
- Algorithm overview
- Computational results

# SDVRP Formulation

- Vehicles with capacity
- customers with demand
- **Objective:**
  - minimize the total distance traveled
- **Constraints:**
  - serve all customers
  - do not exceed the vehicle capacity
  - customers can be serviced by more than one vehicle



## Bounded Formulation (Frizzell & Giffin 1992)

$d_{ij}$  cost of traveling between customer  $i$  and customer  $j$

$w_i$  demand of customer  $i$

$Q$  capacity of vehicle

### **Variables:**

$x_{ijk}$  1 if the vehicle  $k$  travels directly from customer  $i$  to customer  $j$ ,  
0 otherwise

$f_{ik}$  the fraction of demand of customer  $i$  delivered by vehicle  $k$

$y_{ik}$  surplus variables for subtour elimination

$$\begin{aligned}
& \min \sum_{i=0}^n \sum_{j=0}^n \sum_{k=1}^v d_{ij} x_{ijk} \\
& \sum_{k=1}^v \sum_{i=0}^n x_{ijk} \geq 1; \\
& \sum_{i=0}^n x_{ijk} = \sum_{j=0}^n x_{ijk}; \\
& \sum_{k=1}^v f_{ik} = 1; \quad \sum_{i=1}^n w_i f_{ik} \leq Q; \quad f_{ik} \leq \sum_{j=0}^n x_{jik}; \\
& y_{ik} - y_{jk} + (n+1)x_{ijk} \leq n.
\end{aligned}$$

## ***Main idea***

Let  $\pi: \{1, \dots, n\} \rightarrow \{1, \dots, n\}$  is a permutation of customers.

$R^\pi$  is the set of all possible routes, obtained from the permutation  $\pi$ .

For each  $r \in R^\pi$  vehicle can travel from customer  $i$  to customer  $j$  only if  $y_i < y_j$

## ***Theorem 1***

Let the distance matrix  $(d_{ij})$  satisfies the triangle inequality and let  $R^* = (r^1, \dots, r^F)$  is the set of routes in optimal solution. Then there is a permutation  $\pi$  such that  $r \in R^\pi$  for each  $r \in R^*$

## ***Theorem 2***

If  $(d_{ij})$  satisfies the triangle inequality, there always exists an optimal SDVRP solution such that  $y_{ik_1} = y_{ik_2}$  for all  $i, k_1, k_2$ .

## ***Theorem 3***

SDVRP is NP-hard for given permutation  $\pi$ .

## ***Algorithmic idea***

1. Apply local search for permutations
2. For each permutation we apply heuristics to obtain set of routes

## ***Encoding solutions***

### **Permutation U without split demands**

List of customer visits:	5 3 6 2 1 4
Amounts delivered:	9 5 7 4 8 6

### **Permutation V with 3 split demands**

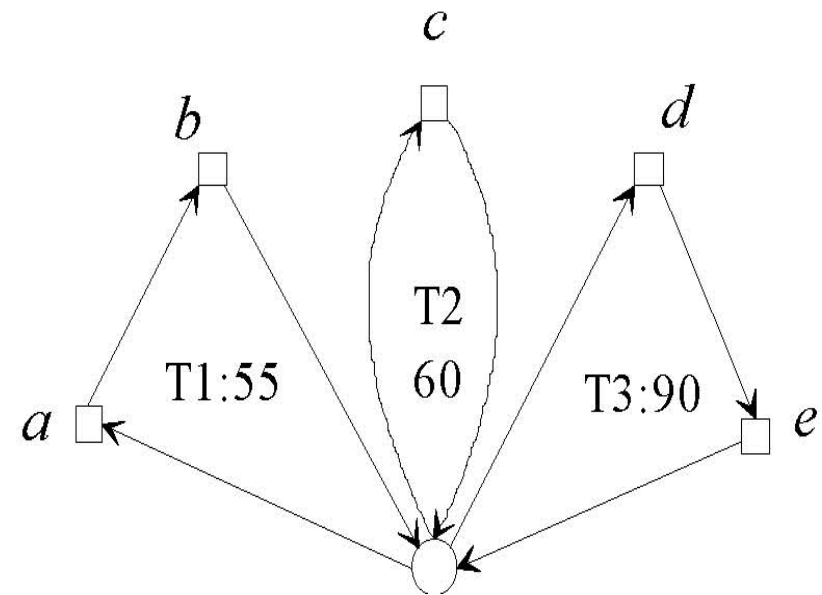
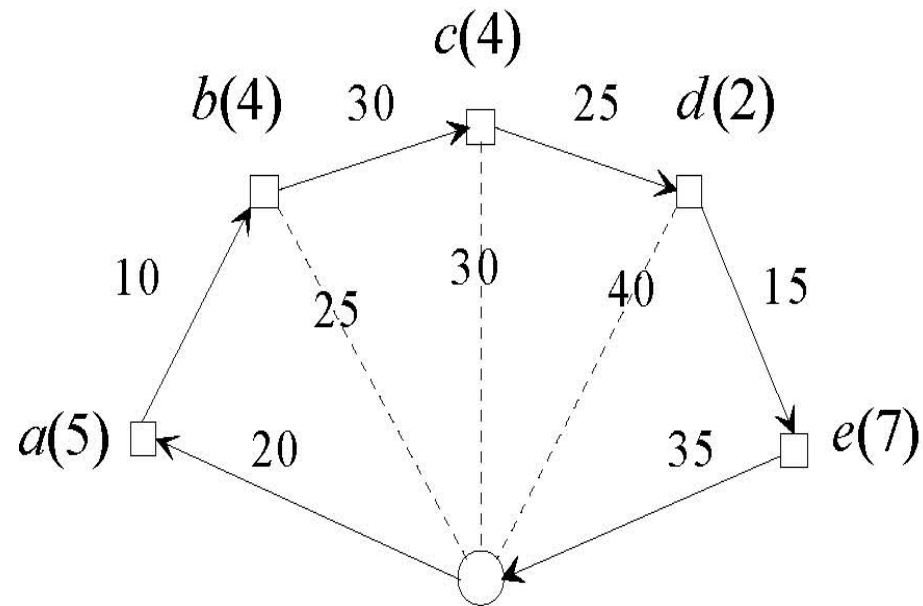
List of customer visits:	5 3 3 6 6 2 1 1 4
Amounts delivered:	9 1 4 6 1 4 5 3 6

### **Permutation E with dummy customers**

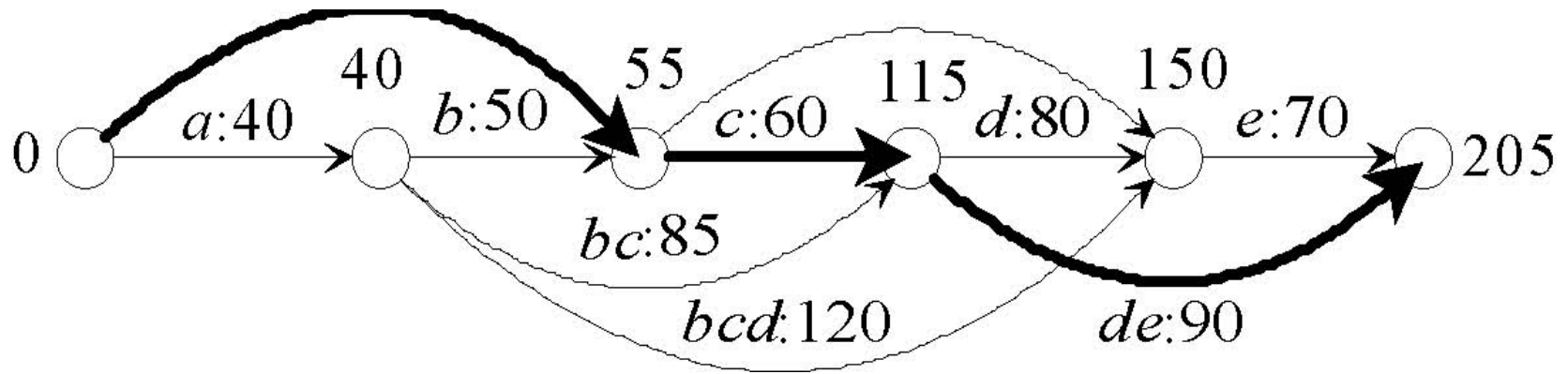
List of customer visits:	5 3 3 6 -- 6 2 1 -- 1 4
Amounts delivered:	9 1 4 6 5 1 4 5 7 3 6



## Decoding procedures: Exact (Prins 2004)



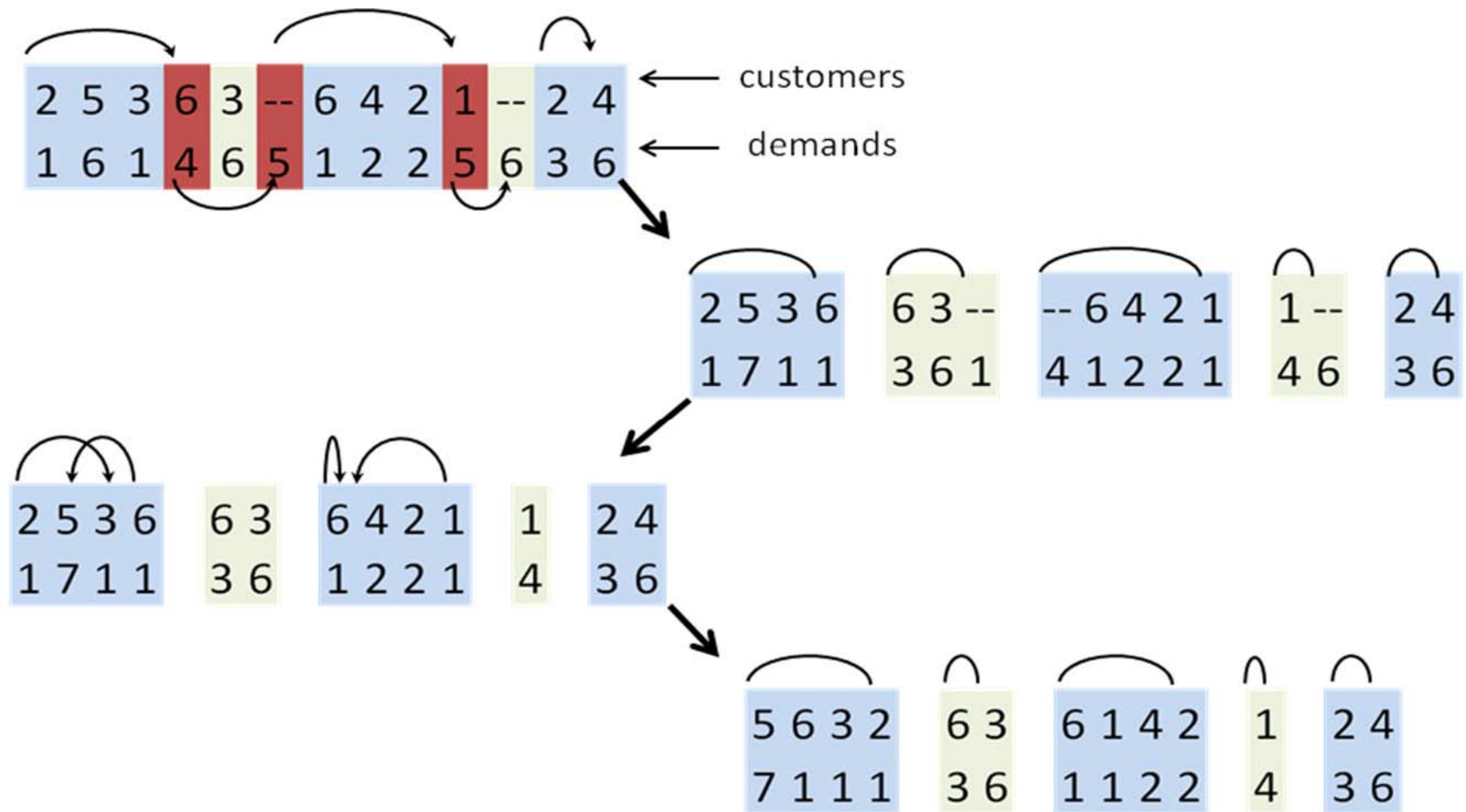
Permutation  $S = (a, b, c, d, e)$  and optimal splitting for VRP, cost 205



*Auxiliary graph of possible trips for  $Q = 10$  and shortest path in boldface*

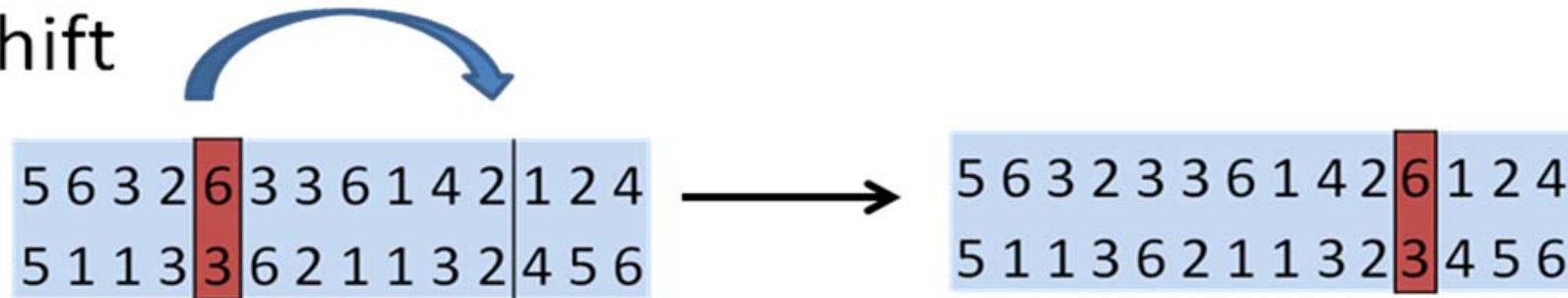
*(Bellman's algorithm for directed acyclic graphs)*

## Decoding procedures: Greedy

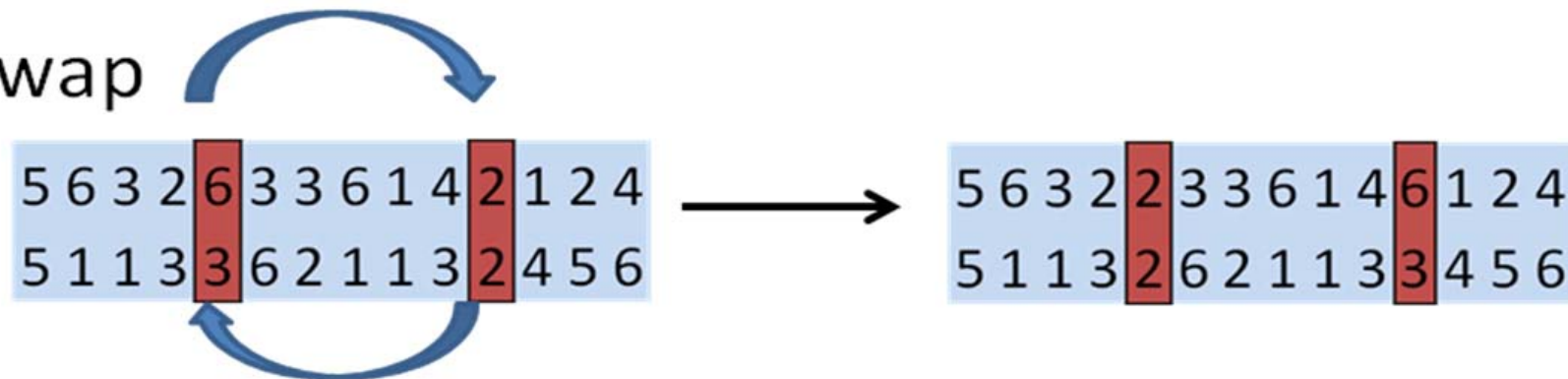


## Local search moves

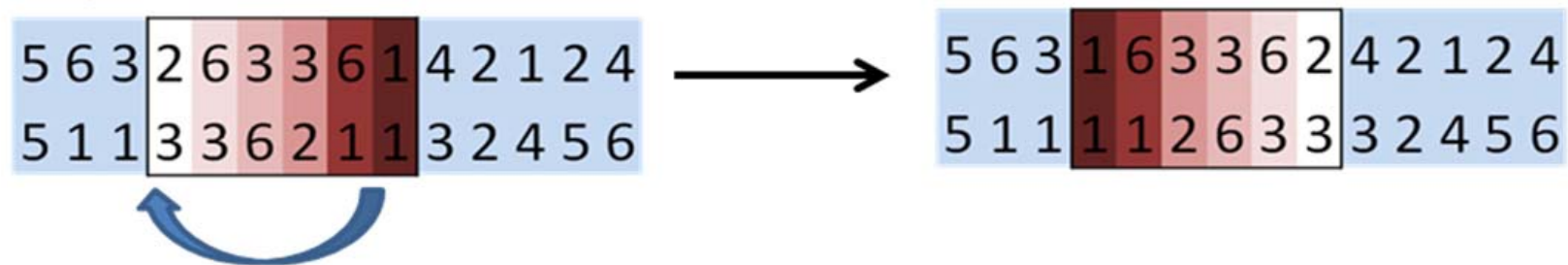
- Shift



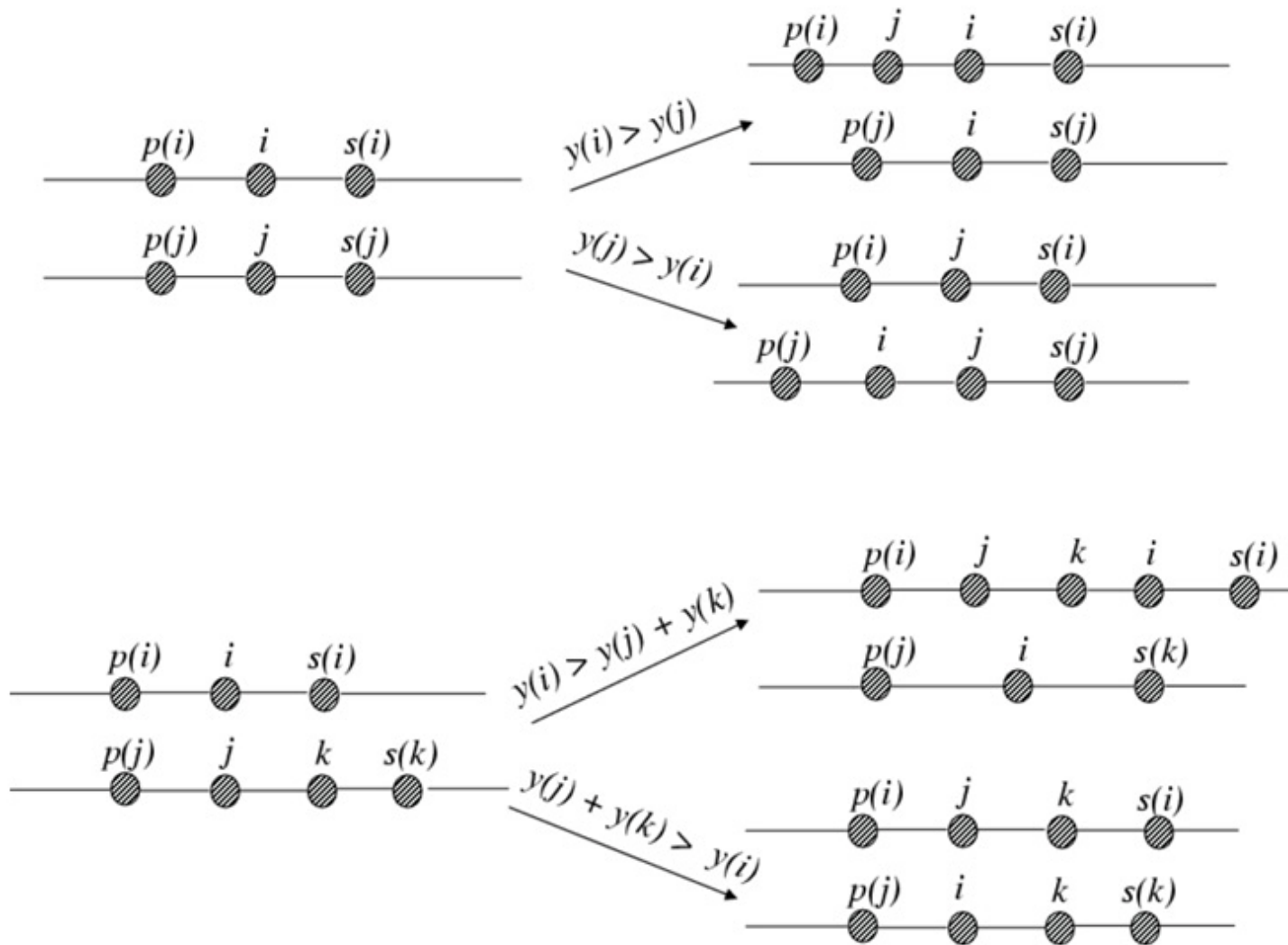
- Swap



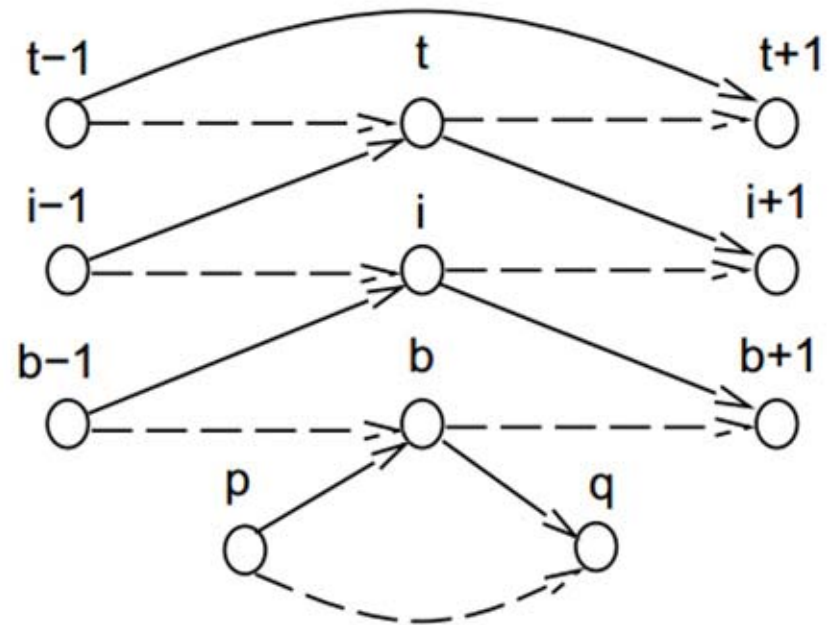
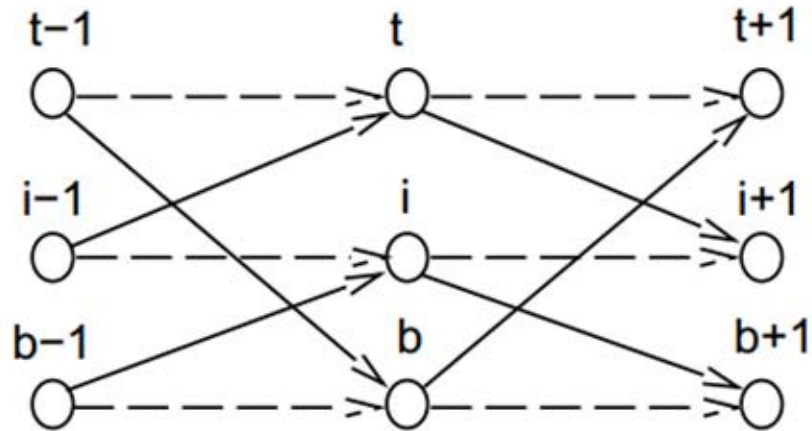
- 2-Opt



# Split moves (Boudia 2007)



# Ejection chains



## Variable neighborhood descent

1. Choose initial solution
2. Apply a local search method to find local optimum for current neighborhood
3. Switch neighborhood
4. Repeat steps 2 and 3 until it is local optimum for each neighborhood

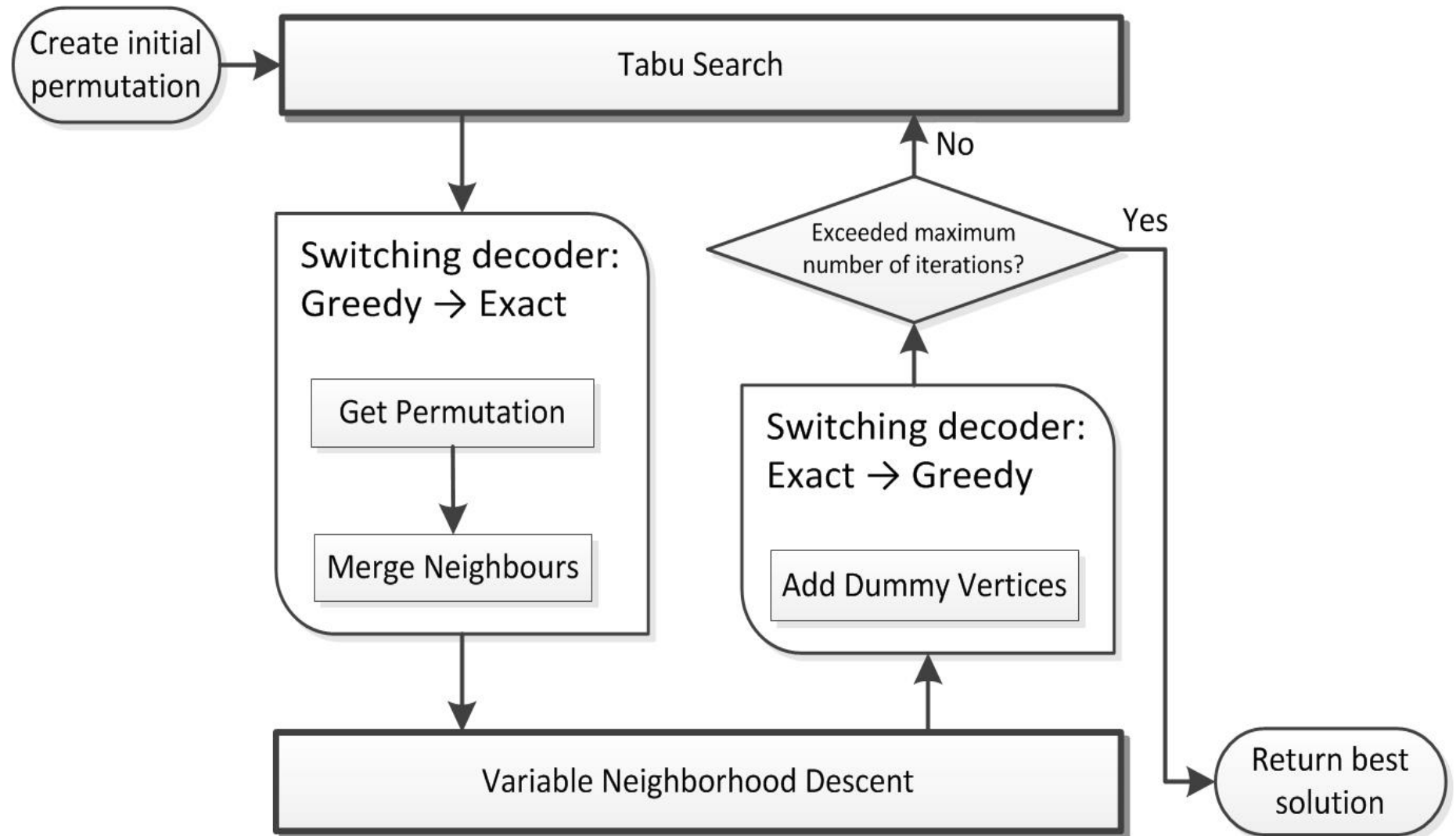
We use exact decoder.

## Tabu search procedure

1. Choose an initial solution, tabu list is empty
  2. Move to the best available neighboring solution
  3. Update tabu list
  4. Repeat steps 2 and 3 until the stopping condition is met
- We use greedy decoder.
  - We use randomized 2-opt neighborhood.
  - Stopping condition is maximum number of iterations.
  - Diversification: we return the best solution.



# Algorithm overview



		TSVBA		MA MP		Hybrid			
File	Demand	Cost	Time	Cost	Time	Cost	Time	Saving	Best known
p1-50	[0.01-0.1]	466,74	20	460,79	12	461,87	82	-0,23	460,79
p2-75	[0.01-0.1]	614,09	136	600,06	19	600,32	170	-0,68	596,25
p3-100	[0.01-0.1]	741,60	1 944	726,81	37	745,62	317	-2,59	726,81
p4-150	[0.01-0.1]	891,10	2 641	875,61	100	885,93	951	-2,26	866,31
p5-199	[0.01-0.1]	1 069,24	11 216	1 018,71	356	1 048,63	1 337	-2,97	1 018,38
p6-120	[0.01-0.1]	990,59	2 736	976,57	73	1 003,75	450	-2,78	976,57
p7-100	[0.01-0.1]	658,99	462	649,73	35	637,46	243	-0,58	633,80
p1-50	[0.1-0.3]	753,98	23	751,41	10	766,94	36	-3,49	741,06
p2-75	[0.1-0.3]	1 085,70	97	1 074,46	34	1 084,41	69	-1,56	1 067,80
p3-100	[0.1-0.3]	1 416,35	161	1 392,85	78	1 403,99	100	-1,94	1 377,28
p4-150	[0.1-0.3]	1 929,91	755	1 878,71	148	1 904,92	226	-1,59	1 875,09
p5-199	[0.1-0.3]	2 408,16	1 544	2 340,14	347	2 366,39	399	-1,59	2 329,37
p6-120	[0.1-0.3]	2 755,74	464	2 720,38	144	<b>2 718,23</b>	146	<b>0,08</b>	2 720,38
p7-100	[0.1-0.3]	1 441,48	98	1 417,28	43	1 426,20	94	-0,74	1 415,78
p1-50	[0.1-0.5]	1 023,24	18	988,31	12	<b>988,05</b>	25	<b>0,03</b>	988,31
p2-75	[0.1-0.5]	1 458,59	68	1 413,80	37	1 436,98	63	-2,75	1 398,53
p3-100	[0.1-0.5]	1 886,70	145	1 845,30	28	1 859,68	122	-1,75	1 827,65
p4-150	[0.1-0.5]	2 647,17	470	2 561,65	225	2 581,26	259	-1,63	2 539,75
p5-199	[0.1-0.5]	3 296,69	1 217	3 191,25	475	3 216,43	535	-1,14	3 180,30
p6-120	[0.1-0.5]	4 010,80	341	3 934,39	163	<b>3 925,04</b>	147	<b>0,24</b>	3 934,39
p7-100	[0.1-0.5]	2 010,00	85	1 994,59	51	<b>1 989,40</b>	99	<b>0,26</b>	1 994,59

		TSVBA		MA MP		Hybrid			
File	Demand	Cost	Time	Cost	Time	Cost	Time	Saving	Best known
p1-50	[0.1-0.9]	1 530,81	19	1 467,06	21	<b>1 461,64</b>	54	<b>0,37</b>	1 467,06
p2-75	[0.1-0.9]	2 164,74	62	2 102,58	46	2 112,65	123	-1,22	2 087,22
p3-100	[0.1-0.9]	2 874,86	125	2 780,95	84	2 789,66	132	-0,31	2 780,95
p4-150	[0.1-0.9]	4 151,90	452	4 045,87	245	4 050,19	534	-0,34	4 036,44
p5-199	[0.1-0.9]	5 066,24	109	4 941,22	726	<b>4 890,86</b>	608	<b>1,02</b>	4 941,22
p6-120	[0.1-0.9]	6 308,76	419	6 318,37	196	<b>6 227,87</b>	173	<b>1,43</b>	6 318,37
p7-100	[0.1-0.9]	3 157,48	98	3 113,72	52	<b>3 105,25</b>	143	<b>0,27</b>	3 113,72
p1-50	[0.3-0.7]	1 505,38	19	1 477,01	25	1 498,21	52	-1,44	1 477,01
p2-75	[0.3-0.7]	2 182,33	55	2 132,16	52	2 149,66	138	-0,82	2 132,16
p3-100	[0.3-0.7]	2 929,29	135	2 858,87	100	2 866,12	191	-0,25	2 858,87
p4-150	[0.3-0.7]	4 151,90	449	4 045,87	245	<b>4 037,89</b>	395	<b>0,20</b>	4 045,87
p5-199	[0.3-0.7]	5 281,50	119	5 155,36	750	<b>5 134,90</b>	745	<b>0,40</b>	5 155,36
p6-120	[0.3-0.7]	6 511,08	437	6 424,71	271	<b>6 349,54</b>	253	<b>0,78</b>	6 399,42
p7-100	[0.3-0.7]	3 200,62	65	3 155,69	91	<b>3 142,46</b>	153	<b>0,42</b>	3 155,69
p1-50	[0.7-0.9]	2 219,32	24	2 154,35	23	2 176,75	122	-1,04	2 154,35
p2-75	[0.7-0.9]	3 278,33	85	3 200,35	27	3 244,74	343	-1,39	3 200,35
p3-100	[0.7-0.9]	4 435,56	186	4 312,95	56	4 339,39	675	-0,61	4 312,95
p4-150	[0.7-0.9]	6 416,12	679	6 267,48	402	6 305,73	767	-0,61	6 267,48
p5-199	[0.7-0.9]	8 333,61	153	8 081,58	572	<b>8 075,79</b>	1 101	<b>0,07</b>	8 081,58
p6-120	[0.7-0.9]	10 186,06	30	10 063,47	298	10 031,42	523	-0,14	10 017,47
p7-100	[0.7-0.9]	4 996,88	153	4 919,48	180	4 934,64	296	-0,31	4 919,48

	TSVBA		EMIP + VRTP		Hybrid			
Instance	Cost	Time	Cost	Time	Cost	Time	Saving	Best known
c-SD01-008	228,28	0,00	228,28	0,70	228,28	0,37	0,00	228,28
c-SD02-016	708,28	0,02	714,40	54,40	714,40	2,58	-0,86	708,28
c-SD03-016	430,58	0,03	430,61	67,30	430,58	2,07	0,00	430,58
c-SD04-024	631,05	0,08	631,06	400,00	631,05	6,78	0,00	631,05
c-SD05-032	1 390,57	0,13	1 408,12	402,70	1 403,99	17,74	-0,97	1 390,57
c-SD06-032	831,24	0,14	831,21	408,30	831,24	17,58	0,00	831,21
c-SD07-040	3 640,00	0,09	3 714,40	403,20	3 640,00	30,76	0,00	3 640,00
c-SD08-048	5 068,28	0,14	5 200,00	404,10	5 108,28	54,87	-0,79	5 068,28
c-SD09-048	2 071,03	0,36	2 059,84	404,30	<b>2 048,67</b>	62,11	<b>0,54</b>	2 059,84
c-SD10-064	2 747,83	0,89	2 749,11	400,00	<b>2 716,96</b>	85,03	<b>1,12</b>	2 747,83
c-SD11-080	13 280,00	0,41	13 612,12	400,10	13 280,00	148,30	0,00	13 280,00
c-SD12-080	7 213,62	0,84	7 399,06	408,30	7 236,66	109,80	-0,32	7 213,62
c-SD13-096	10 110,58	1,20	10 367,06	404,50	10 110,58	191,14	0,00	10 110,58
c-SD14-120	10 802,87	2,31	11 023,00	5 021,70	<b>10 776,82</b>	237,35	<b>0,24</b>	10 802,87
c-SD15-144	15 153,45	3,20	15 271,77	5 042,30	<b>15 121,94</b>	289,28	<b>0,21</b>	15 153,45
c-SD16-144	3 446,43	7,59	3 449,05	5 014,70	<b>3 381,31</b>	258,50	<b>1,89</b>	3 446,43
c-SD17-160	26 493,56	7,27	26 665,76	5 023,60	26 584,90	421,31	-0,34	26 493,56
c-SD18-160	14 323,04	27,95	14 546,58	5 028,60	<b>14 288,57</b>	416,28	<b>0,24</b>	14 323,04
c-SD19-192	20 157,10	11,95	20 559,21	5 034,20	<b>20 109,05</b>	530,70	<b>0,24</b>	20 157,10
c-SD20-240	39 722,86	11,02	40 408,22	5 053,00	<b>39 697,18</b>	1 056,28	<b>0,06</b>	39 722,86
c-SD21-288	11 458,76	111,56	11 491,67	5 051,00	<b>11 292,96</b>	1 241,37	<b>1,45</b>	11 458,76

# Conclusion

- We found some new properties of the problem
- We developed new hybrid heuristic based on idea of different decoders
- It improves 21 best-known solutions in the 70 test instances with number of customers up to 280

