

## Syntax

<b>Value</b>	$v ::= c \mid \lambda x.e$
<b>Expressions</b>	$e ::= v \mid x \mid e e$
<b>Basic Types</b>	$b ::= \text{int} \mid \text{bool}$
<b>Types</b>	$\tau ::= \{v:b \mid e\} \mid x:\tau \rightarrow \tau$
<b>Environment</b>	$\Gamma ::= \emptyset \mid x:\tau, \Gamma$

## Operational Semantic

$$\begin{aligned}
e_1 e_2 &\hookrightarrow e'_1 e_2 && \text{if } e_1 \hookrightarrow e'_1 \\
\lambda x.e e_x &\hookrightarrow e[e_x/x] \\
c e &\hookrightarrow c e' && \text{if } e \hookrightarrow e' \\
c v &\hookrightarrow [[c]](v)
\end{aligned}$$

## Erasing

$$\begin{aligned}
[[\{v:b \mid e\}]] &= b \\
[[x:\tau_x \rightarrow \tau]] &= [[\tau_x]] \rightarrow [[\tau]]
\end{aligned}$$

## Interpretations

$$\begin{aligned}
[[\{v:b \mid e_v\}]] &= \{e \mid e:b \wedge e \hookrightarrow^* v \Rightarrow e_v[e/v] \hookrightarrow^* \text{true}\} \\
[[x:\tau_x \rightarrow \tau]] &= \{e \mid e: [[\tau]] \rightarrow [[\tau_x]] \wedge \forall e_x \in [[\tau_x]]. e e_x \in [[\tau[e_x/x]]]\}
\end{aligned}$$

## Typing

$$\begin{array}{c}
\Gamma \vdash e:\tau \\
\\
\frac{(x, \{v:b \mid e\}) \in \Gamma}{\Gamma \vdash x:\{v:b \mid v=x\}} \quad \frac{(x, \tau) \in \Gamma \quad \tau \equiv x':\tau'_x \rightarrow \tau'}{\Gamma \vdash x:\tau} \\
\\
\frac{}{\Gamma \vdash c:ty(c)} \quad \frac{\Gamma \vdash e:\tau' \quad \Gamma \vdash \tau' \preceq \tau \quad \Gamma \vdash \tau}{\Gamma \vdash e:\tau} \\
\\
\frac{\Gamma, x:\tau_x \vdash e:\tau \quad \Gamma \vdash \tau_x}{\Gamma \vdash \lambda x.e:(x:\tau_x \rightarrow \tau)} \quad \frac{\Gamma \vdash e_1:(x:\tau_x \rightarrow \tau) \quad \Gamma \vdash e_2:\tau_x}{\Gamma \vdash e_1 e_2:\tau[e_2/x]} \\
\\
\Gamma \vdash \tau \\
\\
\frac{\Gamma, v:b \vdash e:\text{bool}}{\Gamma \vdash \{v:b \mid e\}} \quad \frac{\Gamma \vdash \tau_x \quad \Gamma, x:\tau_x \vdash \tau}{\Gamma \vdash x:\tau_x \rightarrow \tau} \\
\\
\Gamma \vdash \tau \preceq \tau \\
\\
\frac{\Gamma, v:b \vdash e_1 \Rightarrow e_2}{\Gamma \vdash \{v:b \mid e_1\} \preceq \{v:b \mid e_2\}} \quad \frac{\Gamma \vdash \tau'_x \preceq \tau_x \quad \Gamma, x:\tau'_x \vdash \tau \preceq \tau'}{\Gamma \vdash x:\tau_x \rightarrow \tau \preceq x:\tau'_x \rightarrow \tau'}
\end{array}$$

$$\Gamma \vdash e \Rightarrow e$$

$$\frac{\forall \theta. \Gamma \vdash \theta \wedge \theta \ e_1 \hookrightarrow^* \text{true} \Rightarrow \theta \ e_2 : s \hookrightarrow^* \text{true}}{\Gamma \vdash e_1 \Rightarrow e_2}$$

$$\Gamma \vdash \theta$$

$$\frac{\forall x \in \text{Dom}(\Gamma). \theta(x) \in \llbracket \theta \ \Gamma(x) \rrbracket}{\Gamma \vdash \theta}$$

## Constants

For each constant  $c$ ,

1.  $\emptyset \vdash c : ty(c)$
2. If  $ty(c) = x : \tau_x \rightarrow \tau$ , then for each  $v$  such that  $\emptyset \vdash v : \tau_x$   $\llbracket c \rrbracket(v)$  is defined and  $\vdash \llbracket c \rrbracket(v) : \tau$   $[v/x]$
3. If  $ty(c) = \{v : b \mid e\}$ , then  $e [c/v] \hookrightarrow^* \text{true}$  and  $\forall c' \ c' \neq c. \neg(e [c'/v] \hookrightarrow^* \text{true})$

## Semantic Typing

$$\begin{aligned} \Gamma \vdash e \in \tau &\doteq \forall \theta. \Gamma \vdash \theta \Rightarrow \theta \ e \in \llbracket \theta \ \tau \rrbracket \\ \Gamma \vdash \tau_1 \subseteq \tau_2 &\doteq \forall \theta. \Gamma \vdash \theta \Rightarrow \llbracket \theta \ \tau_1 \rrbracket \subseteq \llbracket \theta \ \tau_2 \rrbracket \end{aligned}$$

**Lemma 1** .

1. If  $\Gamma \vdash \tau_1 \preceq \tau_2$  then  $\Gamma \vdash \tau_1 \subseteq \tau_2$
2. If  $\Gamma \vdash e : \tau$  then  $\Gamma \vdash e \in \tau$

**Lemma 2** If  $e \hookrightarrow e'$  then  $\Gamma \vdash \tau [e'/x] \preceq \tau [e/x]$

**Lemma 3 (Substitution)** If  $\Gamma \vdash e_x : \tau_x$ , then

1. If  $\Gamma, x : \tau_x, \Gamma' \vdash \tau_1 \preceq \tau_2$  then  $\Gamma, [e_x/x] \Gamma' \vdash [e_x/x] \tau_1 \preceq [e_x/x] \tau_2$
2. If  $\Gamma, x : \tau_x, \Gamma' \vdash e : \tau$  then  $\Gamma, [e_x/x] \Gamma' \vdash [e_x/x] e : [e_x/x] \tau$
3. If  $\Gamma, x : \tau_x, \Gamma' \vdash \tau$  then  $\Gamma, [e_x/x] \Gamma' \vdash [e_x/x] \tau$

**Lemma 4 (Preservation)** If  $\Gamma \vdash e : \tau$  and  $e \hookrightarrow e'$  then  $\Gamma \vdash e' : \tau$ .

**Lemma 5 (Progress)** If  $\emptyset \vdash e : \tau$  and  $e \neq v$  then there exists an  $e'$  such that  $e \hookrightarrow e'$ .

## Interpretations

$$\text{Fin } (e) \doteq \exists v. e \hookrightarrow^* v$$

$$\begin{aligned} \llbracket x \rrbracket &= x \\ \llbracket c \rrbracket &= c \end{aligned}$$

$$\begin{aligned} \llbracket \lambda x. e \rrbracket &= f \\ \llbracket e_1 \ e_2 \rrbracket &= \llbracket e_1 \rrbracket (\llbracket e_2 \rrbracket) \end{aligned}$$

**Claim 1**

$$\left\{ \bigwedge_{(x, \{b:v|e\}) \in \Gamma} (\text{Fin } (x) \Rightarrow \llbracket e[x/v] \rrbracket) \Rightarrow \llbracket e_1 \rrbracket \Rightarrow \llbracket e_2 \rrbracket \right\} \Rightarrow \{\Gamma \vdash e_1 \Rightarrow e_2\}$$