

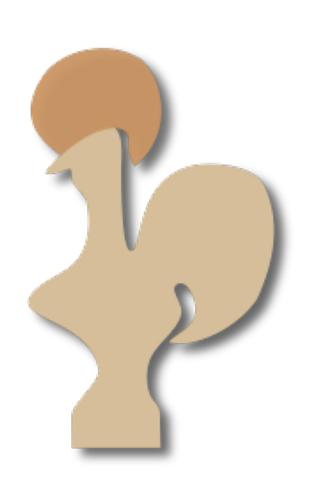
Towards a Translation

from

Liquid Haskell

0

Coq



Lykourgos Mastorou, **Niki Vazou**, and Michael Greenberg IMDEA Software Institute, Spain Stevens Institute of Technology, USA



Haskell Type: incr:: Int → Int

Code: incri=i+1



Refined Type: incr:: i:Int \rightarrow {v:Int | v = 1 + i }

Haskell Type: incr:: Int → Int

Code: incri = i + 1



Refined Type: incr:: i:Int \rightarrow {v:Int | v = 1 + i }

Haskell Type: incr:: Int → Int

Code: incri = i + 1



SMT says: i + 1 = 1 + i



Refined Type: $incr :: i:Int \rightarrow \{v:Int \mid v > i\}$

Haskell Type: incr:: Int → Int

Code: incri=i+1



SMT says: i+l>i



Liquid Haskell For Safe Indexing

```
get:: xs:[a] \rightarrow \{i:Int \mid 0 \le i < len xs\} \rightarrow [a]
```

```
toPair :: String → (String, String)
toPair input = (get input 42, drop input 42)
```



Liquid Haskell For Safe Indexing

```
get :: xs:[a] \rightarrow {i:Int | 0 \le i < len xs} \rightarrow [a]

toPair :: String \rightarrow (String, String)

toPair input = (get input' 42, drop input' 42)

where input' = input ++

replicate (42 - length input) '©'
```



Liquid Haskell For Theorem Proving

```
thm:: xs:[a] \rightarrow {i:Int | 0 \le i < len xs} \rightarrow {xs = get xs i ++ drop xs i}

thm xs 0 = ()

thm [] i = ()

thm (_:xs) i = thm xs (i-1)
```



Liquid Haskell For Theorem Proving

```
thm:: xs:[a] \rightarrow {i:Int | 0 \le i < len xs} \rightarrow {xs = get xs i ++ drop xs i}

thm xs 0 = ()

thm [] i = ()

thm (_:xs) i = thm xs (i-1)
```



Is that a proof?



Liquid Haskell For Theorem Proving

```
thm:: xs:[a] \rightarrow {i:Int | 0 \le i < len xs} \rightarrow {xs = get xs i ++ drop xs i}

thm xs 0 = ()

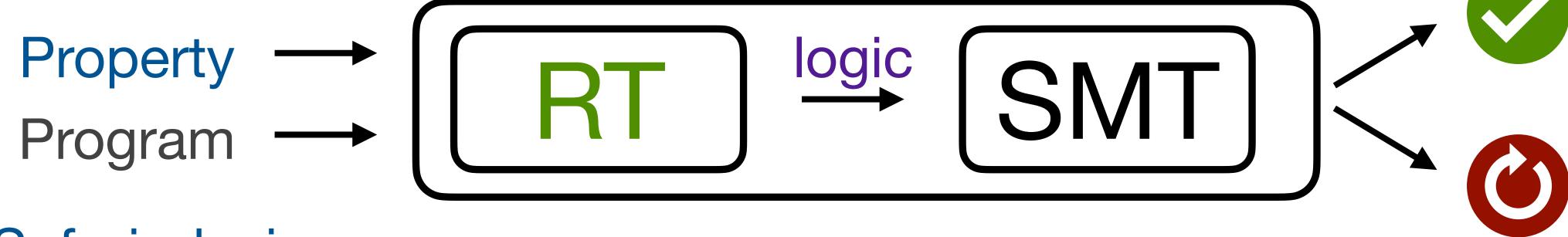
thm[] i = ()

thm (_:xs) i = thm xs i
```

No interaction

HUGE UNREADABLE ERROR MSG

Is Practical

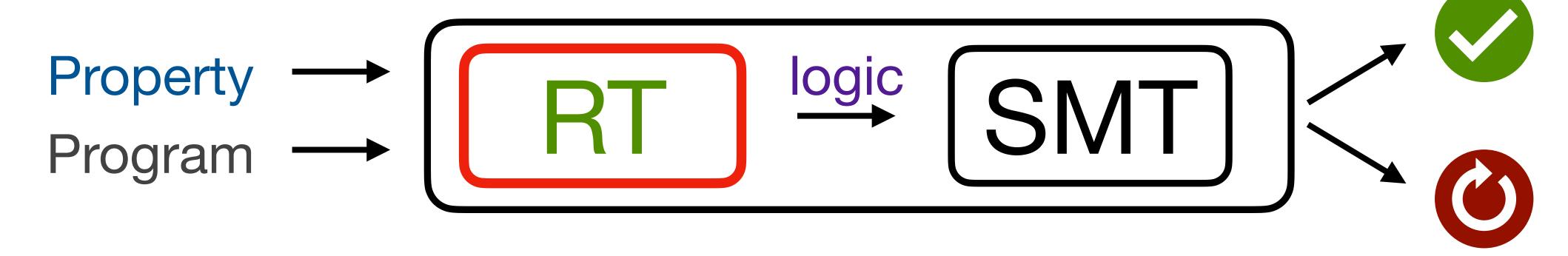


Safe-indexing

getxsi

INFO ⇒
i≤len xs

Is Practical, But Not Very Trustworthy

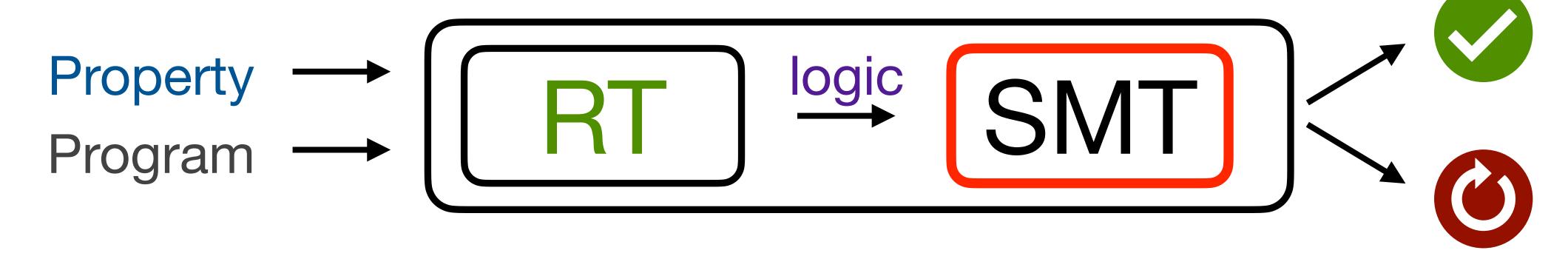




Function Extensionality is Tricky Negative Occurrences are Unsound

How to Safely Use Extensionality in Liquid Haskell by Niki Vazou and Michael Greenberg. Haskell 2022.

Is Practical, But Not Very Trustworthy

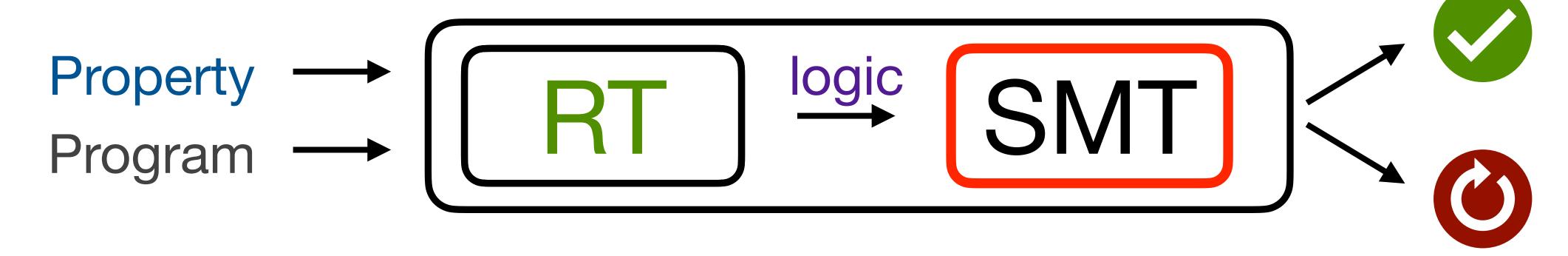


X Errors in RT

Function Extensionality is Tricky Negative Occurrences are Unsound

X Errors in SMT

Is Practical, But Not Very Trustworthy

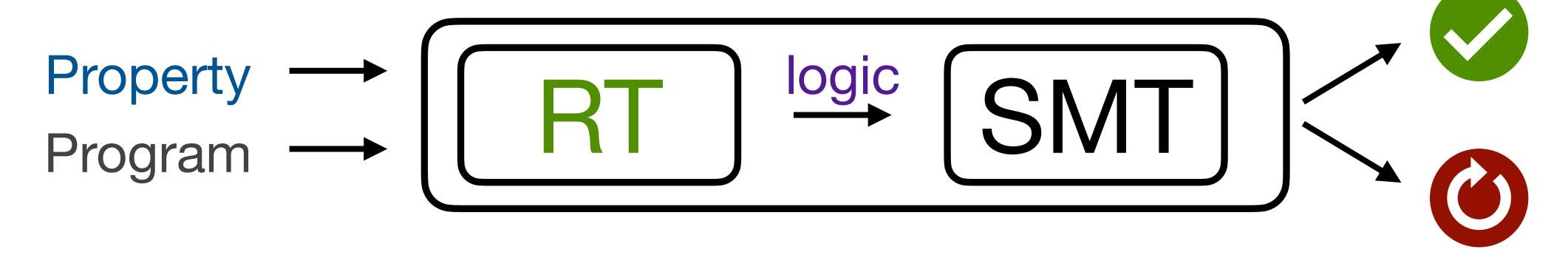


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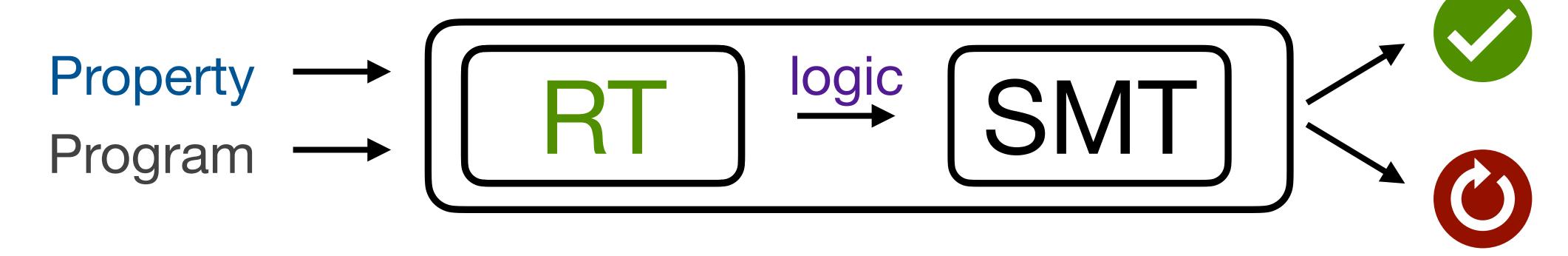


X Errors in RT

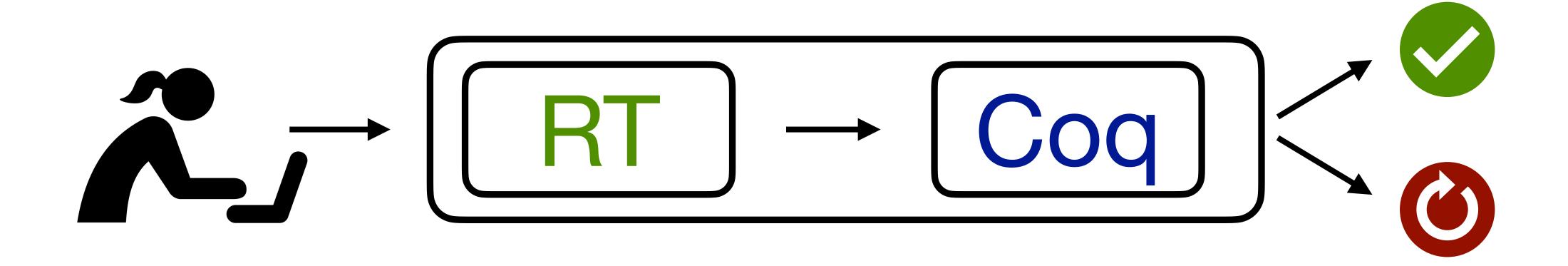
Function Extensionality is Tricky Negative Occurrences are Unsound

X Errors in SMT

Is Practical, But Not Very Trustworthy



- X Errors in RT
 - Function Extensionality is Tricky Negative Occurrences are Unsound
- X Errors in SMT
- X No Constructive Proofs



- Keep Practical User Interface of LH
- Coq Soundness & Constructive Proofs
- Coq's Interactive Proof Development

Refinement Types
Functions
Semantic Subtyping
SMT Automation

Refinement Types

Functions

Semantic Subtyping SMT Automation

Refinement Types ⇒ Inductive Types

✓ Easy/Natural to Use $\{i:Int \mid 0 \le i\} \implies Nat$

X Do not break invariants

Refinement Types ⇒ **Inductive Types**

Easy/Natural to Use

$$\{i:Int \mid 0 \le i\} \implies Nat$$

m > 1

X Do not break invariants

(n-1) + m :: Nat

Refinement Types ⇒ **Inductive Types**

Easy/Natural to Use

$$\{i:Int \mid 0 \le i\} \implies Nat$$

m > 1

X Do not break invariants

(n-1) + m :: Nat

z. Nat

Refinement Types ⇒ Subset Types

```
\{i:Int \mid 0 \le i\} \implies \{i:Int \mid 0 \le i\}
42:: \{i:Int \mid 0 \le i\} \implies (42, evidence): \{i:Int \mid 0 \le i\}
```

Refinement Types ⇒ Subset Types

```
\{i:Int \mid 0 \le i\} \implies \{i:Int \mid 0 \le i\}
42:: \{i:Int \mid 0 \le i\} \implies (42, evidence): \{i:Int \mid 0 \le i\}
```

Challenge: Construct evidence!

Refinement Types ⇒ Subset Types Functions

Semantic Subtyping SMT Automation

Functions \(\Rightarrow \) Fixpoint

Only structural induction

Functions ⇒ Fixpoint

Only structural induction

```
 \begin{array}{l} ack :: m:Nat \rightarrow n:Nat \rightarrow Nat \,/\, [m,n] \\ ack \ 0 \ n = n + 1 \\ ack \ m \ n = if \ n == 0 \ then \ ack \ (m-1) \ 1 \ else \ ack \ (m-1) \ (ack \ m \ (n-1) \end{array}
```

Functions ⇒ Fixpoint

Only structural induction

```
ack :: m:Nat \rightarrow n:Nat \rightarrow Nat / [m,n] ack 0 n = n + 1 ack m n = if n == 0 then ack (m-1) 1 else ack (m-1) (ack m (n-1))
```

Functions ⇒ Program Fixpoint

Only opaque functions

Functions ⇒ Fixpoint

Only structural induction

```
ack :: m:Nat \rightarrow n:Nat \rightarrow Nat / [m,n]
ack 0 n = n + 1
ack m n = if n == 0 then ack (m-1) 1 else ack (m-1) (ack m (n-1)
```

Functions ⇒ Program Fixpoint

Only opaque functions

Functions ⇒ Equations

Function Definition & Induction Principle

Refinement Types ⇒ Subset Types
Functions ⇒ Equations
Semantic Subtyping

SMT Automation

Semantic Subtyping

```
4::{i:Int \mid i = 4} 2::{i:Int \mid i = 2}
{i:Int \mid i = 4} {i:Int \mid i = 2} {i:Int \mid i = 2}
```

4 :: Nat

2 :: Nat

if p then 4 else 2:: Nat

Semantic Subtyping

4 :: Nat

2 :: Nat

if p then 4 else 2:: Nat



if p then (ref 4 ...) else (ref 2 ...)

Semantic Subtyping ⇒ **Ref Custom Tactic**

Introduce explicit casts via the **ref** tactic at the places of semantic subtyping.

Concretely, because of bidirectional typing, join points, function calls and annotations.

Refinement Types ⇒ Subset Types

Functions ⇒ Equations

Semantic Subtyping ⇒ Ref Custom Tactic

SMT Automation

SMT Automation

Or, how do we implement **ref**?

We want SMT logic + proof search.

SMT Automation

Or, how do we implement ref?

We want SMT logic + proof search.

Proof obligation for monotonicity of Ackerman:

ack m p < ack m n

SMT Automation

Or, how do we implement ref?

We want SMT logic + proof search.

SMT Coq + Coq Tactics = Sniper

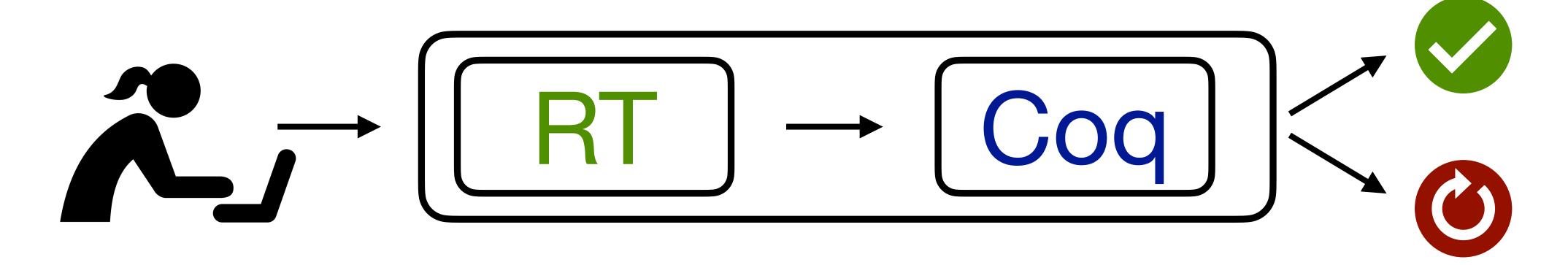
General automation in Coq through modular transformations, by Valentin Blot, Louise Prisque, Chantal Keller, and Pierre Vial. PxTP'21

Refinement Types ⇒ Subset Types

Functions ⇒ Equations

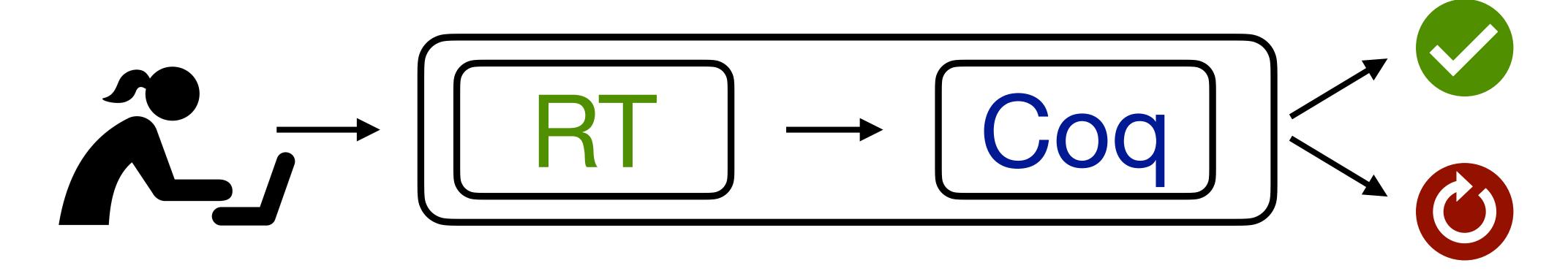
Semantic Subtyping ⇒ Ref Custom Tactic

SMT Automation ⇒ **Sniper**



- Practicality
- Soundness
- Interactivity

- **Refinement Types** ⇒ **Subset Types**
 - **Functions** ⇒ **Equations**
- **Semantic Subtyping** ⇒ **Ref Custom Tactic**
 - **SMT Automation** ⇒ **Sniper**



- Practicality
- Soundness
- Interactivity

- **Refinement Types** ⇒ **Subset Types**
 - **Functions** ⇒ **Equations**
- **Semantic Subtyping** ⇒ **Ref Custom Tactic**
 - **SMT Automation** ⇒ **Sniper**