

Type Theory in 15min

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The 3 Steps of Programming

Step 1: Coding

```
printf("Hello!"); // hello.c
```

Step 2: Compilation

```
> gcc -o hello.exec hello.c
```

Step 3: Execution

```
> ./hello.exec  
> Hello!
```

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Step 3: Execution

Many times & takes long.

The 3 Steps of Programming

Step 1: Coding

```
printf("Hello!"); // hello.c
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Step 2: Compilation

Few times & fast.

Step 3: Execution

Many times & takes long.

The 3 Steps of Programming

Step 1: Coding

Full of **errors**!

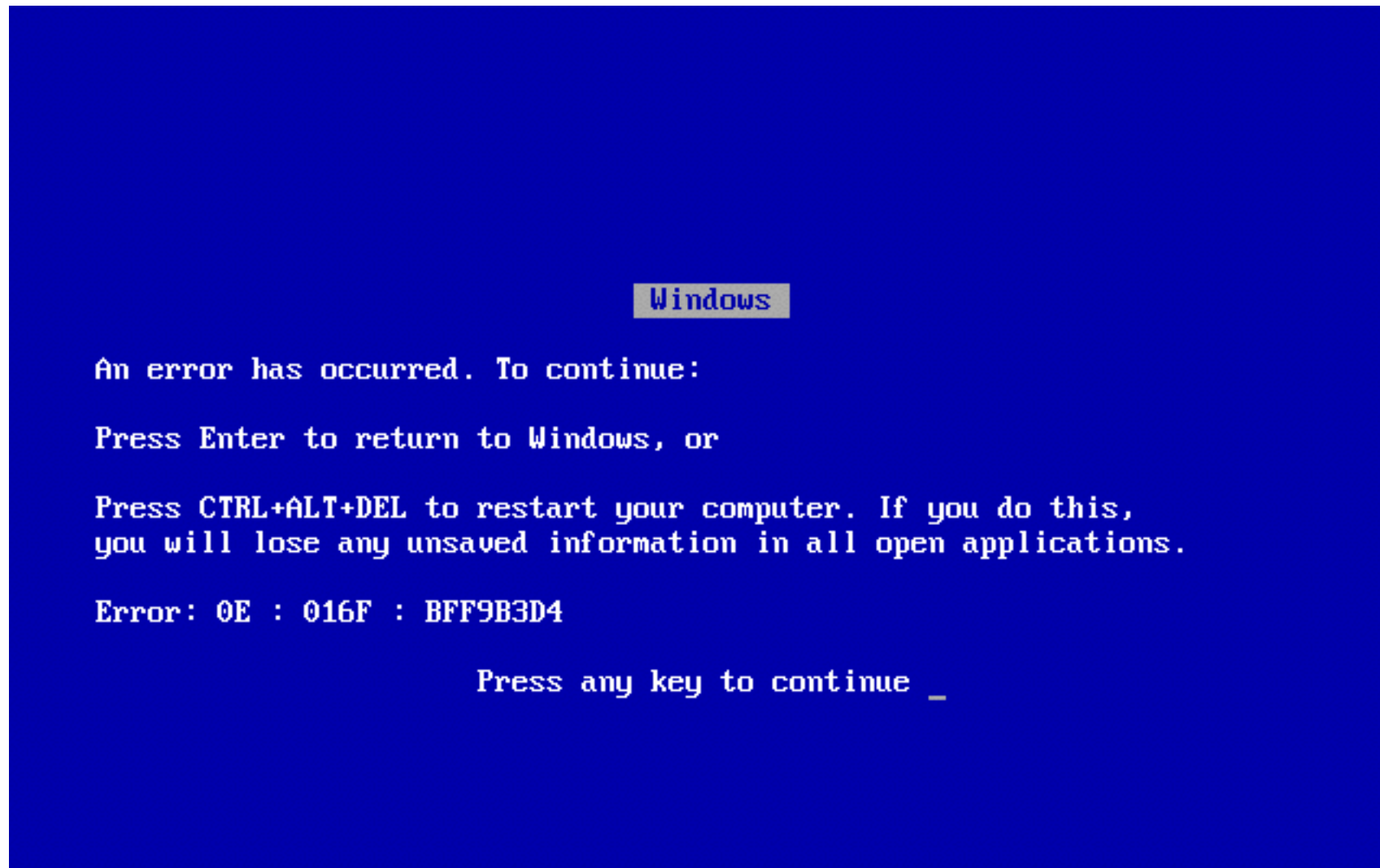
Step 2: Compilation

Few times & fast.

Step 3: Execution

Many times & takes long.

Code Errors crash at execution!



Types to detect code errors!

Code Errors

```
printf(1+"Hey!");
```

What will the above print?

A.

> ey!

B.

> Hey!

C.

> 1+Hey!

D.

> 1Hey!

Code Errors

```
printf(1+"Hey!");
```

What will the above print?

A.

```
> ey!
```


Code Errors vs. Correct Code

```
printf("%d", (1+1)+2);
```

What will the above print?

A.

> 4

B.

> No other alternative!

Types

Detect Errors Early (compile time)!

Types: Classify Data

e.g., “1 has type `Int`”

`1`:`Int`

Types: Classify Data

1: Int

2: Int

...

"Hey": String

Types: Combine Data

If $e_1 : \text{Int}$ and $e_2 : \text{Int}$, then $e_1 + e_2 : \text{Int}$.
where e_1 and e_2 are variables.

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If $e_1 : \text{Int}$ and $e_2 : \text{Int}$, then $e_1 + e_2 : \text{Int}$.
where e_1 and e_2 are variables.

$$e_1 : \text{Int} \quad e_2 : \text{Int}$$

$$e_1 + e_2 : \text{Int}$$

Types: Example I

$1:\text{Int}$ $1:\text{Int}$

$1+1:\text{Int}$

$e_1 := 1$

$e_2 := 1$

Types: Example II

$1+1:\text{Int}$ $2:\text{Int}$

$(1+1)+2:\text{Int}$

$e_1 := 1+1$

$e_2 := 2$

Types: Example II

1: Int 1: Int

1+1: Int 2: Int

(1+1)+2: Int

$e_1 := 1+1$

$e_2 := 2$

Types: Example III

~~1: Int "Hey!". String~~

~~1 + "Hey!": ??~~

$e_1 := 1$

$e_2 := \text{"Hey!"}$

Types: Example III

~~1: Int "Hey!".String~~

~~1+"Hey!":??~~

What rules are wrong?

Soundness

If e has a type, then e cannot crash!

Definitions

Expressions e : written by the programmer.

e.g., $(1+1)+2$

Values v : result of expressions.

e.g., 4

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Evaluation \hookrightarrow : from expressions to values.

$$(1+1)+2 \hookrightarrow 2+2 \hookrightarrow 4$$

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Values v : result of expressions.

Evaluation \hookrightarrow : from expressions to values.

Goes to \hookrightarrow^* : many steps of evaluation.

Crash crash : untyped crashing expression.

If e has a type, then e cannot crash!

If $e : t$, then $e \not\hookrightarrow^* \text{crash}$.

The two Steps of Soundness

Progress:

If $e_1 : t$, then $e_1 \hookrightarrow e_2$ or is a value.

$(1+1)+2 : \text{Int}$

$(1+1)+2 \hookrightarrow 2+2$

The two Steps of Soundness

Progress:

If $e_1 : t$, then $e_1 \hookrightarrow e_2$ or is a value.

Preservation:

If $e_1 : t$ and $e_1 \hookrightarrow e_2$, then $e_2 : t$.

$(1+1)+2 : \text{Int}$

$(1+1)+2 \hookrightarrow 2+2$

$2+2 : \text{Int}$

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If $e_1 : t$ and $e_1 \hookrightarrow e_2$, then $e_2 : t$.

$e : t \hookrightarrow e_1 : t \hookrightarrow \dots \hookrightarrow e_i : t$

If e_i is a value, we are done!

Since e_i has a type, $e_i \neq \text{crash}$!

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If $e_1 : t$, then $e_1 \hookrightarrow e_2$ or is a value.

Preservation:

If $e_1 : t$ and $e_1 \hookrightarrow e_2$, then $e_2 : t$.

Soundness:

If $e : t$, then $e \not\hookrightarrow^* \text{crash}$.

The 3 Steps of Programming & Types!

Step 1: Coding

Full of **errors**!

Step 2: Compilation

Few times & fast.

If $e:t$,

Step 3: Execution

Many times & takes long.

then $e \not\rightarrow^* \text{crash}$.

Thanks!