

Queue object

buf



lock's *wait*
queue



lock's *blocked*
queue



There are two kinds of threads -- consumer threads $C1$, $C2$, ..., that remove characters from *buf*, and producer threads $P1$, $P2$, ..., that add characters to *buf*. For our purposes, *buf.size()* returns the number of characters in the buffer.

The buffer *buf* is initially empty.

1. Consumer $C1$ enters the synchronized block for the *get* method
2. *buf.size() == 0* is true
3. *wait()* is executed, placing $C1$ on the lock's *wait queue*

Show the status of *buf*, lock's *wait queue* and lock's *blocked queue*.

Queue object

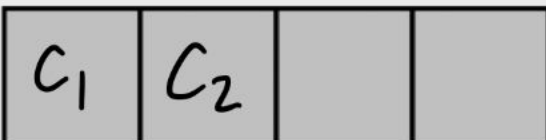
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1. Consumer 2 (*C2*) is just about to enter the synchronized block for the *get* method, but has not acquired the lock.
2. Producer *P1* enters the synchronized method *put*, acquires the lock, places the character “c” into *buf*, and calls *notify()*.
3. *C1* is woken up by the notify and must reacquire the lock before proceeding. Thus both *C1* and *C2* are competing for the lock.

Show the status of *buf*, lock's *wait queue* and lock's *blocked queue*.

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1. One of *C1* and *C2* is non-deterministically chosen to get the lock. Let's say *C2* gets the lock. It gets to enter the method since *C1* is awake it is put on the *blocked queue*, not back on the *wait queue*.
2. *C2* gets the character and releases the lock which is then acquired by *C1*.

Is there a character in *buf* for *C1* to get?

No

What will happen in the program as written?

Buffer is empty so *C1* will still be kept on hold

What would have happened if the *while* loop was not in the *get()* method?

If there was no *while* loop then *C1* will try to get a value out of the buffer which will result in getting an index out of bounds error

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Let's look at a scenario that shows the need for `notifyAll` instead of `notify` in the code.

To make this easy, assume a buffer size of 1. Producer and consumer threads are named as before. *buf* is initially empty.

1. *P1* puts a "c" into the buffer.
2. *P2* attempts a put, checks the while loop and performs a *wait()*
3. *P3* attempts a put, checks the *while* loop and performs a *wait()*

Show the status of *buf*, lock's *wait queue* and lock's *blocked queue*.

Queue object

buf

<i>c</i>			
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lock's *wait*
queue

P_2	P_3		
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lock's *blocked*
queue

C_2	C_3		
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4. The following happen at time step 4:

- $C1$ attempt to get 1 character and enters the *get* method;
- $C2$ attempts to get 1 character but blocks on entry to the *get* method;
- $C3$ attempts to get 1 character but blocks on entry to the *get* method;

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5. The following happen at time step 5.

- a. *C1* is executing the *get* method, gets the character, calls *notify* and exits the method (releasing the lock and giving *C2* and *C3* a chance to acquire it);
- b. The *notify* wakes up *P2*
- c. BUT, *C2* enters the method before *P2* can (*P2* must reacquire the lock), so *P2* blocks on entry to the put method;
- d. *C2* checks the wait loop, sees there are no more characters in the buffer and so it waits (releasing the lock in the process)
- e. *C3* enters the method after *C2*, but before *P2*, checks the wait loop, sees there are no more characters in the buffer, and so it waits

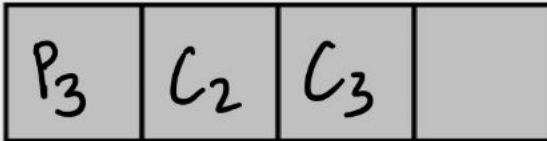
Show the status of *buf*, lock's *wait queue* and lock's *blocked queue*.

Queue object

buf



lock's *wait queue*



lock's *blocked queue*



Queue object

buf

<i>d</i>			
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lock's *wait*
queue

P_3	C_2	C_3	
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lock's *blocked*
queue

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6. The following happen at time step 6.

- Now P_3 , C_2 and C_3 are all waiting!
- P_2 acquires the lock, puts a “ d ” in the buffer, calls *notify* and exits the method

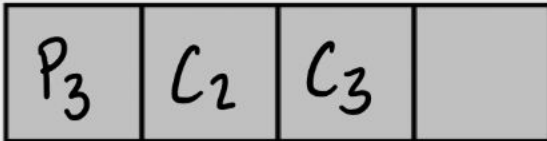
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7. The following happens at time step 7.

- P*2's notification wakes up *P*3 (any thread can be woken up)
- P*3 checks the wait loop condition. There is already a character (“*d*”) in the buffer and so it waits.

Show the status of *buf*, lock's *wait queue* and lock's *blocked queue*.

Is it possible for any thread to be woken up by another notify? No because all three threads are in the lock's wait queue

What would have happened if in 6b a notifyAll() was called?

Would've released the other 3 waiting threads