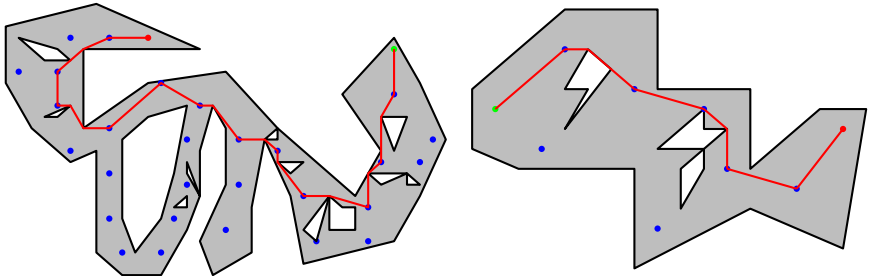


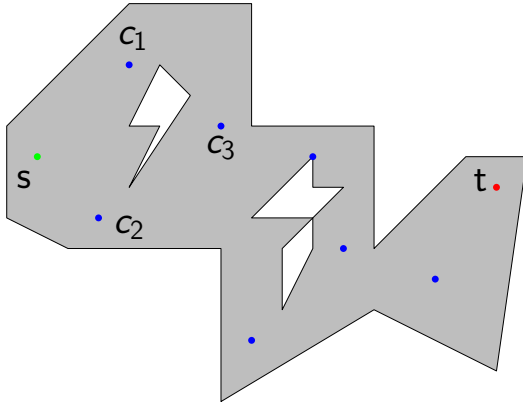
Save the Robot – Computational Geometry

February 3rd, 2016

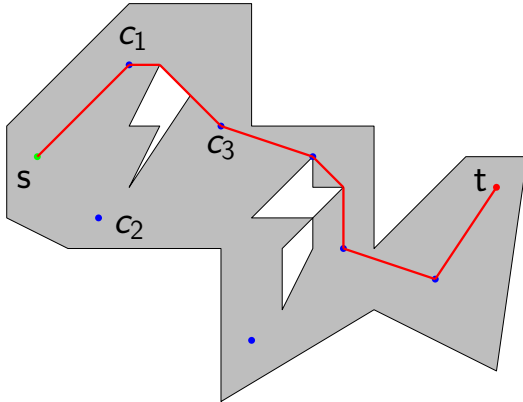
Niklas Baumstark, Samuel Groß

Institute of Theoretical Informatics





- Electric robot is in a warehouse polygon P , starting at green location s
- Range of r , recharge at blue charging stations $c \in C$
- What is the shortest path to red exit location t



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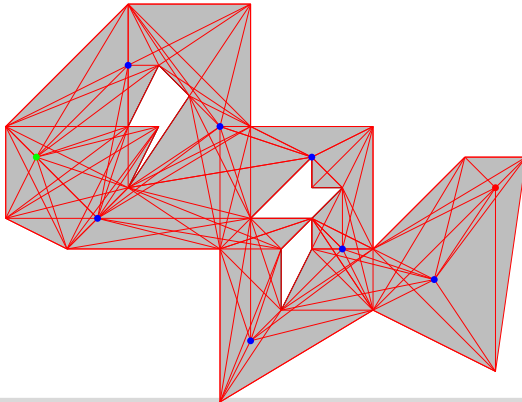
Computing visibility graphs

Computing shortest route

Possible optimizations

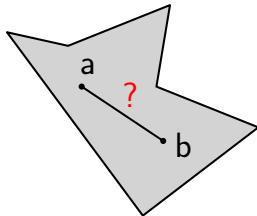
Given a set of points S inside P , the **visibility graph** of S is the undirected graph $G_S = (S, E)$ with edges $E = \{\{a, b\} \mid \overline{ab} \subseteq P\}$

We are interested in the point set $S = P \cup C \cup \{s, t\}$



Computing G_S – First Attempt

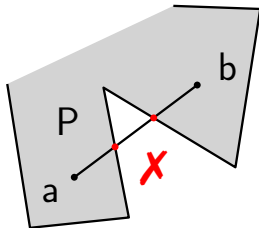
- For each pair $a, b \in S$, check if \overline{ab} is inside P
- $\Theta(|S|^2)$ segment-in-polygon tests
- How to check if segment ab is inside P ?



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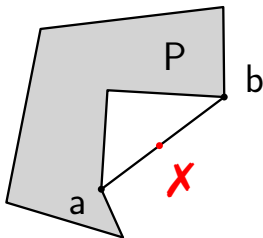
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→ output **NO**



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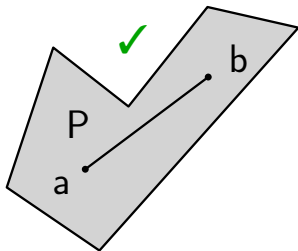
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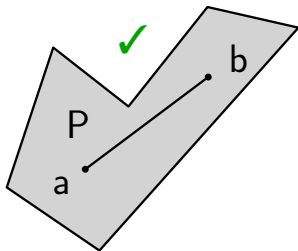
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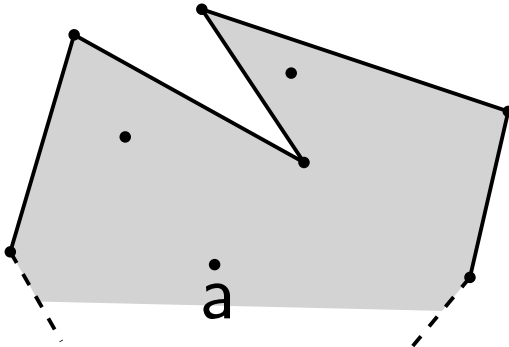
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Runtime: $\Theta(|S|^2 \cdot |P|)$, dominates runtime of the algorithm

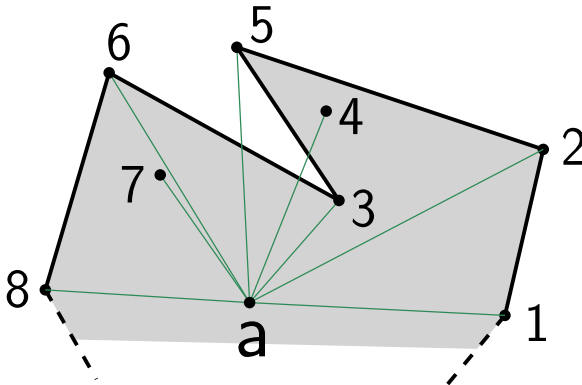
Computing G_S – Second Attempt

- For each $a \in S$, compute points b with \overline{ab} inside P in one pass
- Use **Circle-Sweep**: Sort all points/vertices by their angle around a
- Maintain edges of P intersecting the sweep line, sorted by distance



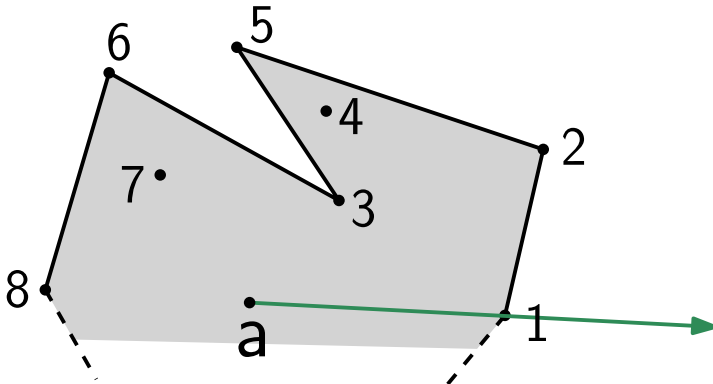
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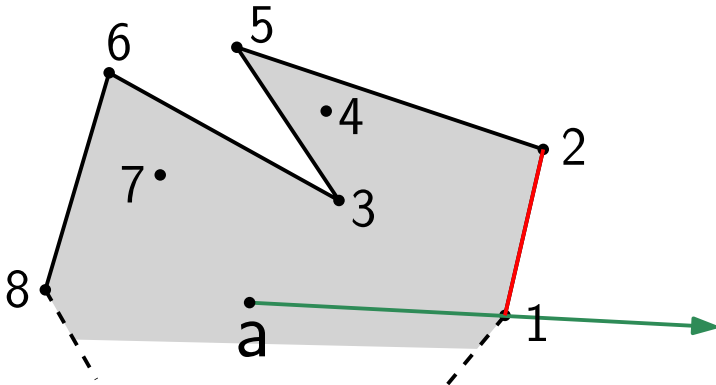
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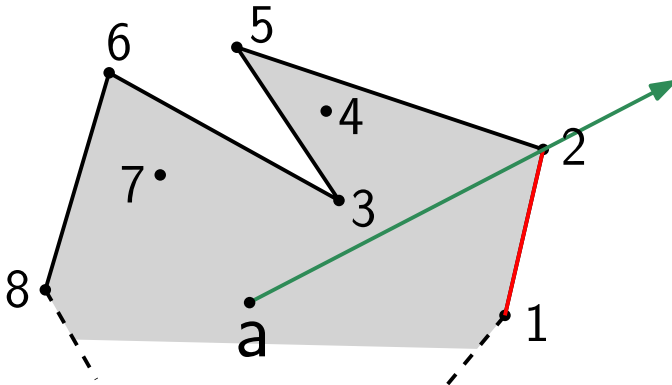
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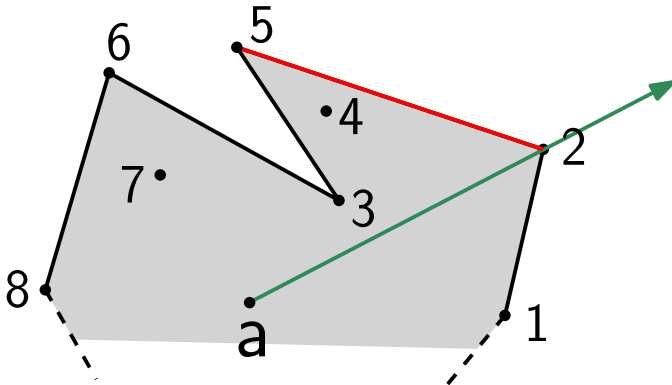
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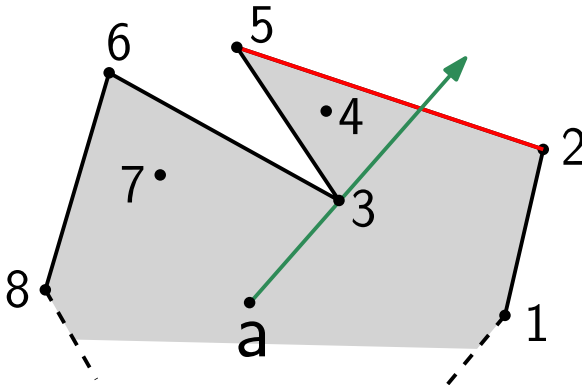
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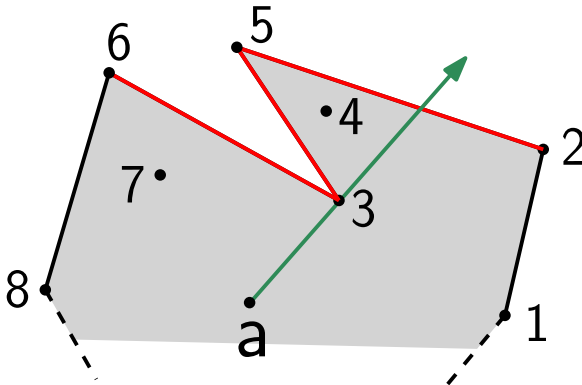
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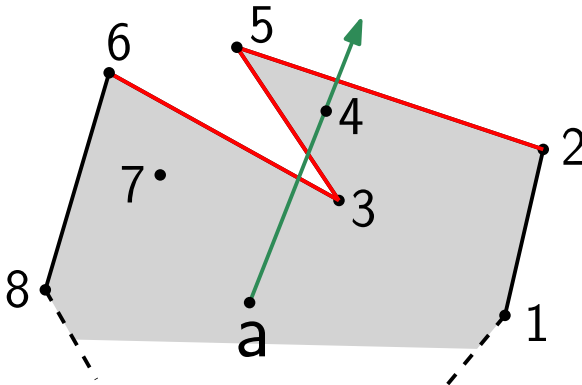
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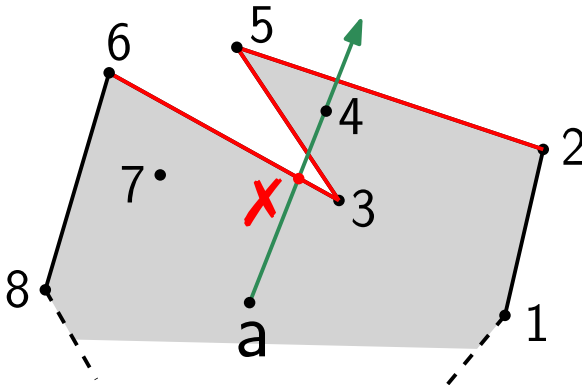
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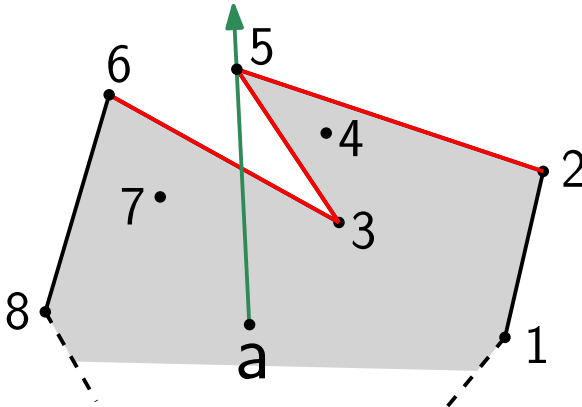
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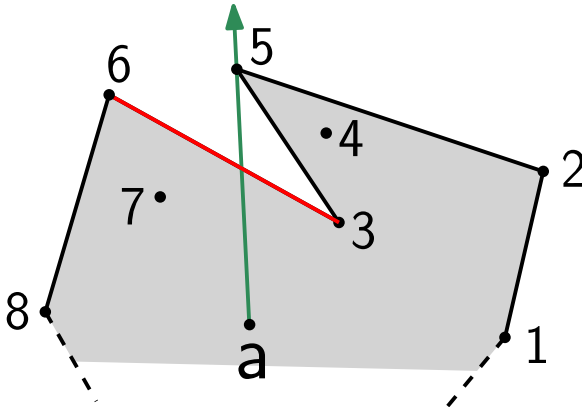
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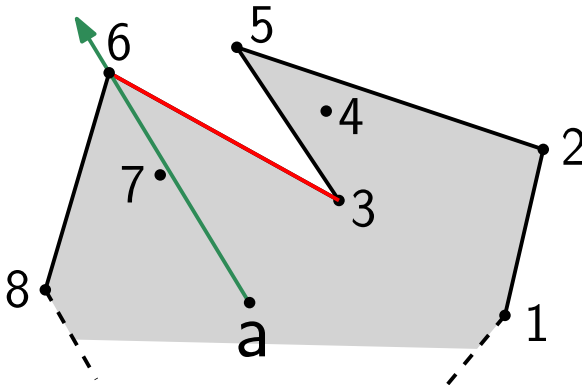
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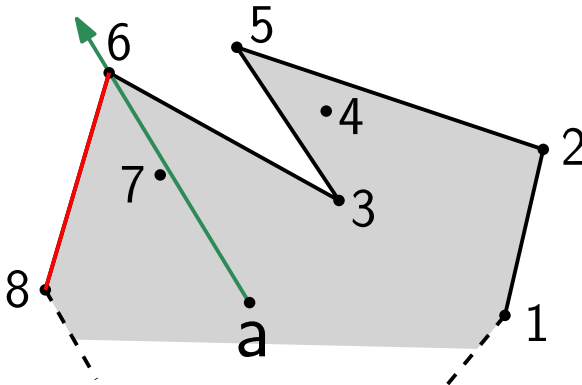
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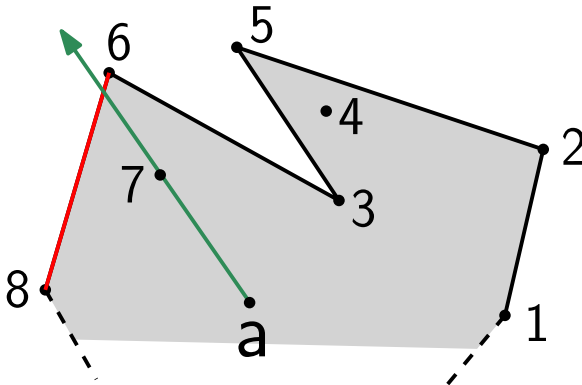
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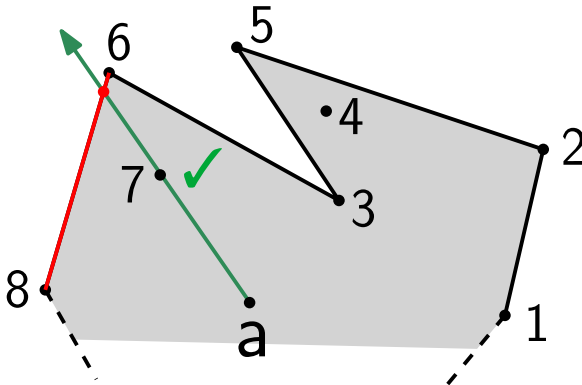
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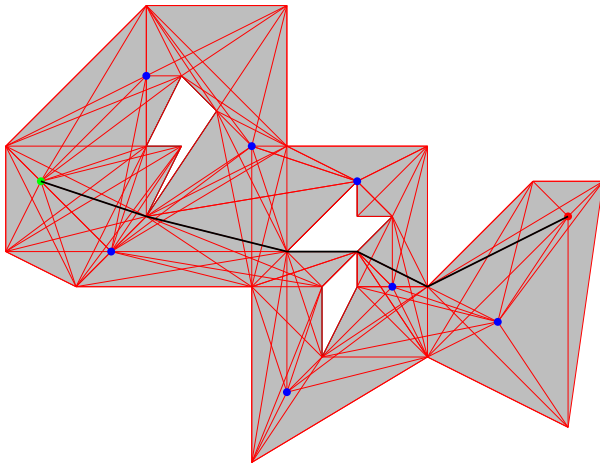
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- Runtime: $\Theta(|S| \cdot n \cdot \log n)$ where $n = |S| + |P|$
- Now that we have G_S , how to solve the problem on a general graph?

~~Computing visibility graphs~~

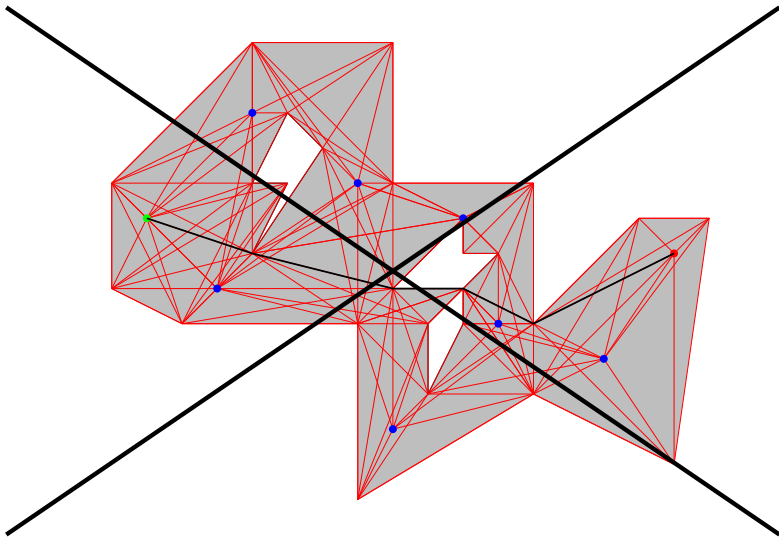
Computing shortest route

Possible optimizations

Computing shortest route



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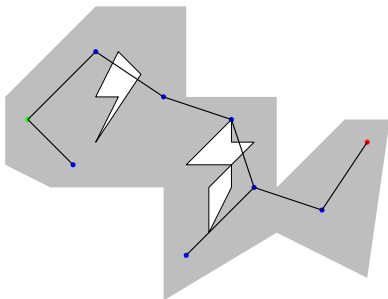
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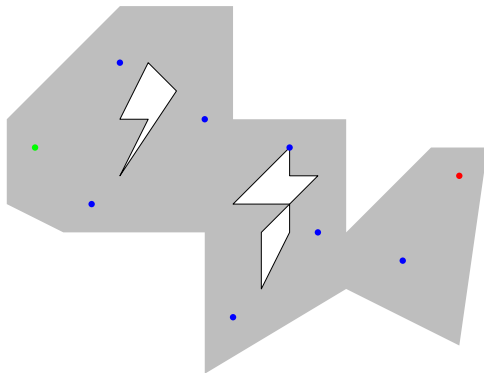
Computing shortest route

- Computing reachability graph
- Computing final path on reachability graph

Possible optimizations

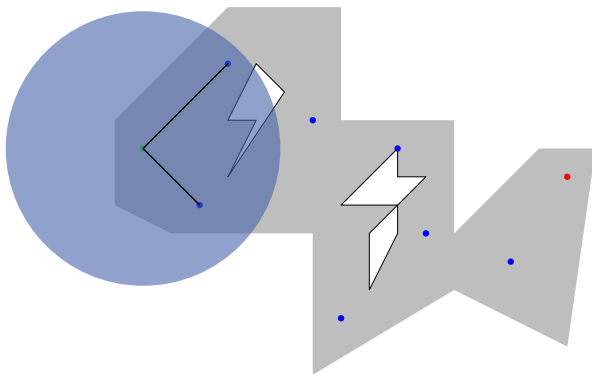
Computing reachability graph

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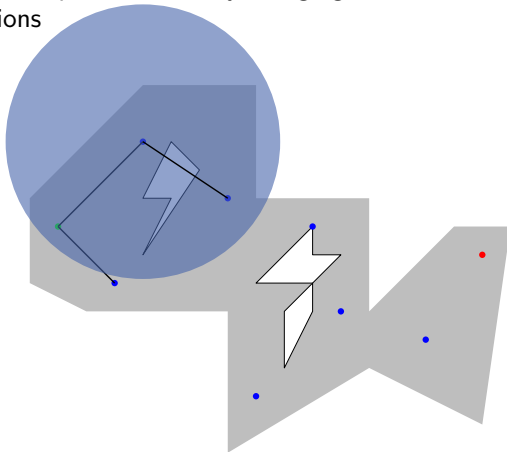
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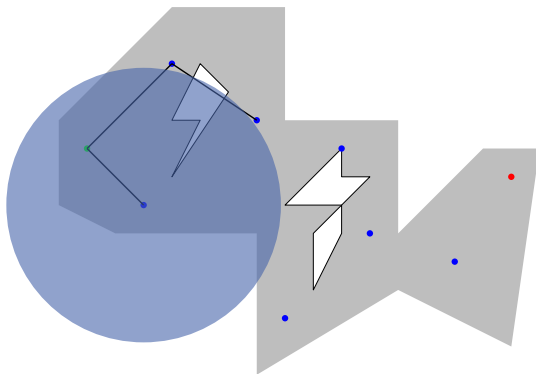
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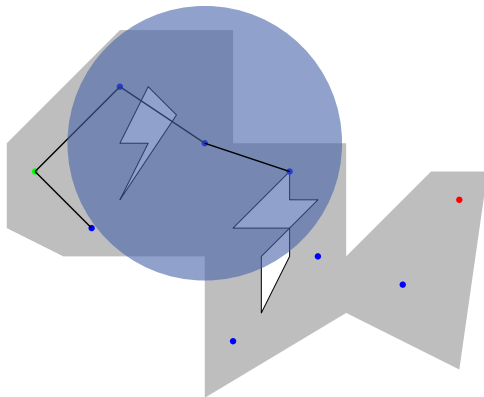
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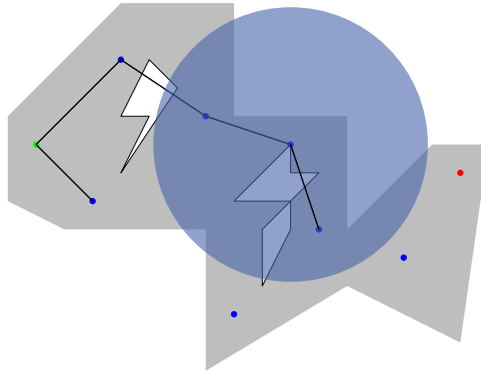
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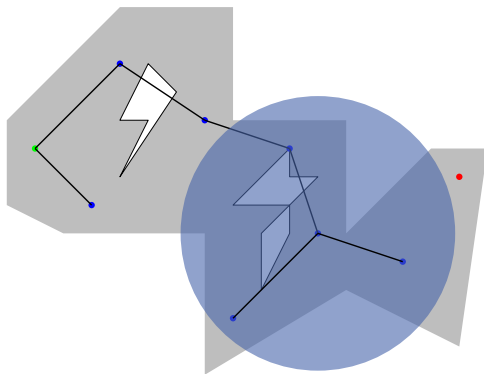
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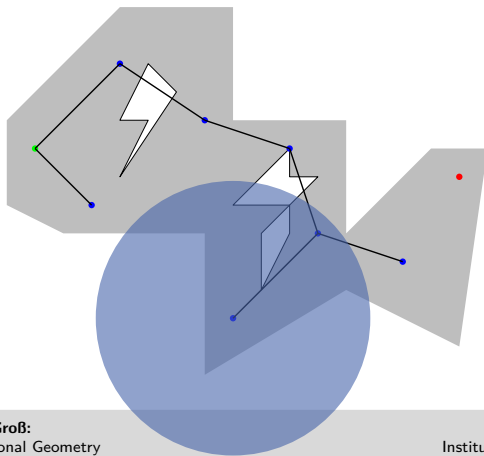
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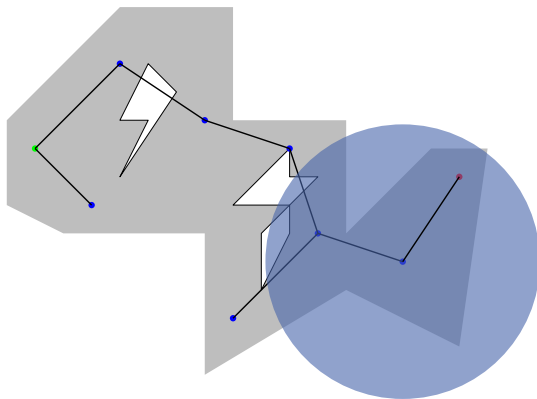
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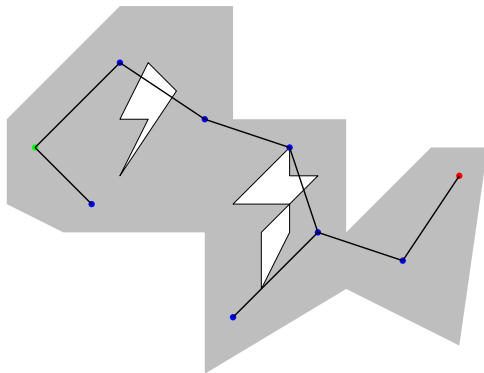
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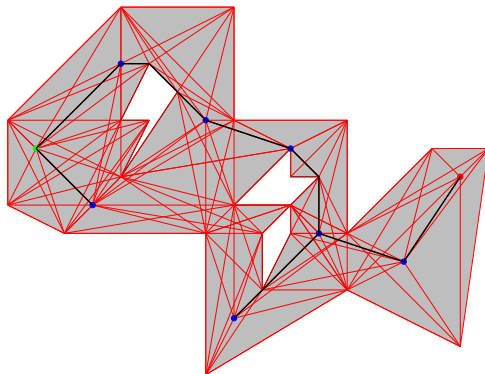
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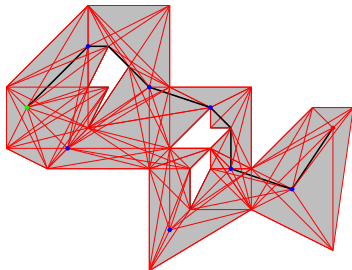
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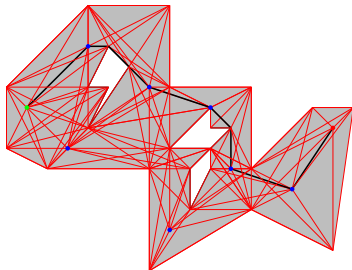
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- Need one final (shortest) path computation from start to end node



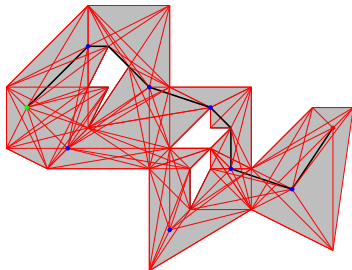
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$$\mathcal{O}((|C| + |P|)^2 + |C| \cdot |E| \cdot \log(|C| + |P|))$$



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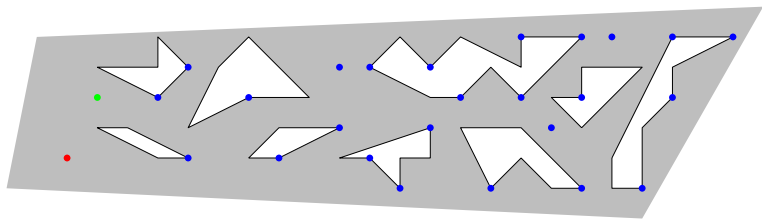
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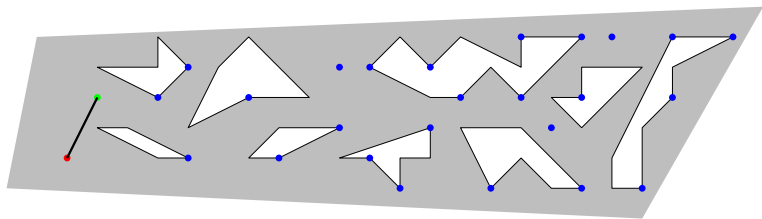
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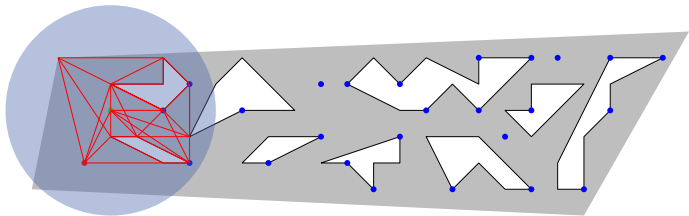


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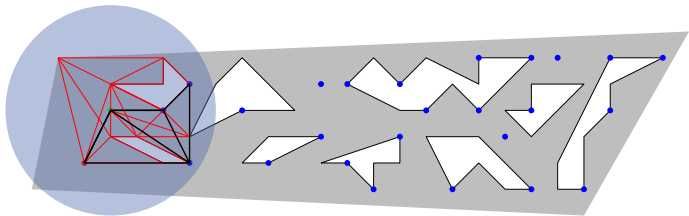
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Questions?

