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INTERNATIONAL STANDARD

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Information technology — Programming languages, their environments and system software interfaces — Programming language Extended APL

Technologies de l'information — Langages de programmation, leurs environnements et interfaces logiciel système — Langage de programmation APL étendu



Reference number ISO/IEC 13751:2001(E)

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Case postale 56 • CH-1211 Geneva 20
Tel. + 41 22 749 01 11
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Foreword

ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and non-governmental, in liaison with ISO and IEC, also take part in the work.

International Standards are drafted in accordance with the rules given in the ISO/IEC Directives, Part 3.

In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.

Attention is drawn to the possibility that some of the elements of this International Standard may be the subject of patent rights. ISO and IEC shall not be held responsible for identifying any or all such patent rights.

International Standard ISO/IEC 13751 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Subcommittee SC 22, *Programming languages*, *their environments and system software interfaces*.

Annex A forms a normative part of this International Standard.

Introduction

APL stands for **A Programming Language**. It is a notation invented by K. E. Iverson in the late 1950s for the description of algorithms, and expanded on and made into the programming system $APL \setminus 360$ by Iverson and his colleagues Adin Falkoff, Larry Breed, Dick Lathwell, and Roger Moore in the mid-1960s.

This document, **Programming Language APL**, **Extended**, is a sequel to **Programming Language APL**, ISO 8485 (1989).

The principal differences that the reader will find here have to do with new features that have been added. These topics are:

without greatest common divisor least common multiple duplicate commute table join along first axis mixed arrays overbar in names underbar in names replicate character grades grades of arrays greater than rank one unique alpha as a name omega as a name ambivalent defined functions event handling n-wise reduction complex arithmetic left right function rank operator defined operators component file system

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enclose disclose enlist pick depth identical each first

An entry for each of these topics will be found in the index. Some new system commands have been added. Shared variable extensions have been added. Workspace Interchange Standard 2 is given, in which canonical representation vectors of type "E" are used to represent generalised arrays.

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INTERNATIONAL STANDARD

ISO/IEC 13751: 2001 (E)

Information technology — Programming languages, their environments and system software interfaces — Programming language Extended APL

1 Scope

This International Standard defines the programming language APL and the environment in which APL programs are executed. Its purpose is to facilitate interchange and promote portability of APL programs and programming skills. This International Standard specifies the syntax and semantics of APL programs and the characteristics of the environment in which APL programs are executed.

It also specifies requirements for conformance to this International Standard, including the publication of values and characteristics of implementation properties so that conforming implementations can be meaningfully compared.

This International Standard does not specify:

- implementation properties that are likely to vary with the particular equipment or operating system used;
- required values for implementation limits such as APL workspace size or numeric precision;
- the data structures used to represent APL objects;
- the facilities available through shared variables.

2 Normative References

The following normative documents contain provisions which, through reference in this text, constitute provisions of this International Standard. For dated references, subsequent amendments to, or revisions of, any of these publications do not apply. However, parties to agreements based on this International Standard are encouraged to investigate the possibility of applying the most recent editions of the normative documents indicated below. For undated references, the latest edition of the normative document referred to applies. Members of ISO and IEC maintain registers of currently valid International Standards.

ISO/IEC 2382-15:1999, Information technology – Vocabulary – Part 15: Programming languages.

ISO 8485:1989, Programming languages – APL.

International Register of Coded Character Sets To Be Used With Escape Sequences, Registered character set 68. (http://www.itscj.ipsj.or.jp/ISO-IR/068.pdf)

3 Form of this International Standard

This International Standard is a formal model of an APL machine, specified as a collection of finite sets, diagrams, and evaluation sequences, and objects constructed from finite sets, diagrams, and evaluation sequences.

The finite sets are the implementation-defined character-set, the implementation-defined set of numbers, and the enumerated sets array-type, class-names, keyboard-states, mode-names, required-character-set, and workspace-presence.

Diagrams are directed graphs used to designate syntactic forms.

Evaluation sequences are formal procedures that operate on finite sets, diagrams, other evaluation sequences and objects defined in the International Standard.

Objects are entities consisting of enumerated set members and other objects. The objects are list, array, defined-function, defined-operator, token, symbol, context, workspace, session, shared-variable, and system.

Each object has attributes describing its state. The attributes of an array, for example, are its typical element, its shape, and its ravel.

Objects often have defined properties derived from their attributes. The rank of an array, for example, is the shape of the shape of the array.

3.1 Form of Definitions

Defined terms in this International Standard are always set in **bold** and indexed. The index entry begins with the page number of the definition followed by the page numbers of all references to the term. If the definition and a use of a term occur on the same page, that page number will occur twice in the Index.

The following terms occur throughout the document and are not cross-indexed: **character**, **content**, **class**, **item**, and **number**.

Note: Terms in this document include both phrases such as implementation-parameter and words

such as nil.

Each definition in this International Standard takes one of four forms.

- 1. Regular definitions consist of the term being defined followed by a colon and the body of the definition. The term **Boolean** is defined in this way.
- 2. Members of an enumerated set are defined simply by being listed in the definition of the enumerated set. The term **nil** is defined in this way, as a member of the enumerated set **class-names**.
- 3. Diagrams are defined by directed graphs. The term **expression** is defined in this way.
- 4. Definitions of terms that designate evaluation sequences take the following form:
 - The term being defined, such as scan.
 - The forms that the evaluation mechanism used in the International Standard recognises as designating this term, such as $Z \leftarrow \mathbf{f} \setminus B$.
 - An informal introduction indicating the purpose of the procedure. The informal introduction is considered commentary on the International Standard.
 - An evaluation sequence expressed in a formal, though English-like, language defined in the subsection Evaluation Sequences. A conforming-implementation is required to emulate the behaviour described in the evaluation sequence as modified by the additional-requirements, if any.
 - Examples, which show effects of the procedure specified by the evaluation sequence.
 Examples are considered commentary on the International Standard.
 - Additional Requirements, giving aspects of the behaviour required of this operation that cannot conveniently be expressed in the evaluation sequence.

3.2 Named Arrays in Examples

In the examples in this International Standard, APL identifiers beginning with N, such as N234, represent numeric arrays whose shape and content are specified by the digits in the identifier. Each digit in the identifier specifies an element of the shape vector; each element of the array, when broken down into digits, gives the index of that element.

For example,

Example:

```
N4
1 2 3 4

N23
11 12 13
21 22 23
```

```
N234
111 112 113 114
121 122 123 124
131 132 133 134
211 212 213 214
221 222 223 224
231 232 233 234

N234[2;3;1]
231
```

3.3 Notes

This International Standard contains notes that comment on the text of the International Standard, pointing out the significance of definitions, noting relationships between definitions, and otherwise making the text approachable. These notes are set in a different type style than the text of the International Standard, and are prefixed with the word "Note:". The following is an example of a note.

Note: This is an example of a note.

Notes never set requirements for conformance. They may suggest desired properties, but such suggestions are not mandatory for conformance.

3.4 Cross-References

The heading levels in this International Standard are chapter, section, and subsection. When cross-references are given, they are always to a subsection title. In the Index, subsections are treated like definitions: the page on which a subsection begins is always the first entry in the Index; subsequent page numbers in the index show where references are made to the subsection.

3.5 General Definitions

For the purpose of this International Standard, the definitions given in ISO/DIS 2382/15 shall apply, together with the following:

Program: An application.

Note: The term is used in this International Standard to include everything from an APL expression to a collection of workspaces communicating via shared variables.

 Implementation: A combination of computer hardware and software that processes (APL) programs.

 $\textbf{Note:} \ \textit{An implementation is an instance of the object } \textbf{system} \ \textit{specified by this International Standard}.$

- Facility (of an implementation): A unit of behaviour. Every facility is one of:
 - Defined-Facility: A facility fully specified in this International Standard and not designated optional or implementation-defined.
 - Optional-Facility: A facility fully specified in this International Standard and designated optional.
 - Implementation-Defined-Facility: A facility not fully specified by this International Standard that is designated implementation-defined.
 - Consistent-Extension: A facility not specified in this International Standard that, for a construct this International Standard specifies as producing an error, gives some effect other than signalling the specified error.

4 Compliance

4.1 Conforming Implementations

An APL **implementation** conforms to this International Standard if it meets the following requirements for both behaviour and documentation.

4.1.1 Required Behaviour for Conforming Implementations

A conforming-implementation shall provide all defined-facilities and implementation-defined-facilities. Each such facility shall behave as specified by this International Standard.

A conforming-implementation may provide optional-facilities. If provided, an optional-facility shall behave as specified by this International Standard. Attempted use of an optional-facility that is not provided shall cause the conforming-implementation to signal an error. A conforming-implementation shall not replace the error signalled by a missing optional-facility with other behaviour.

A conforming-implementation may provide consistent-extensions. The presence of a consistent-extension shall not affect the behaviour of a conforming-program.

A **conforming-implementation** shall use algorithms that produce results that are the same as those produced by the evaluation sequences. Mathematical function algorithms shall have at least the accuracy that the algorithms given in the evaluation sequences would produce.

Note: The evaluation sequences used in this International Standard are intended to specify results, not implementation techniques.

The errors produced by the absence of an optional-facility cannot be replaced by consistent-extensions in a conforming-implementation, since this would affect the behaviour of conforming-programs that use the optional-facility.

4.1.2 Required Documentation for Conforming Implementations

A **conforming-implementation** shall provide a reference document that satisfies the following requirements for the documentation of **optional-facilities**, **implementation-defined-facilities**, and **consistent-extensions**:

4.1.2.1 Documentation of Optional-Facilities

A **conforming-implementation** shall document the presence or absence of each of the following **optional-facilities**:

- Shared-Variable-Protocol. A mechanism that permits one session to exchange data with other autonomous sessions.
- Statement-Separator-Facility. A mechanism for placing more than one statement on a line.
- Trace-and-Stop-Control. A mechanism that assists the testing and correction of defined functions.
- Complex-Arithmetic-Facility. A mechanism that permits the generation and manipulation of complex-arithmetic results.

4.1.2.2 Documentation of Implementation-Defined-Facilities

A **conforming-implementation** shall document the following aspects of **implementation-defined-facilities**:

- A description of the character-set. This shall include a table showing the correspondence between index positions in the atomic-vector and the members of the required-character-set. If the graphic symbols used in a conforming-implementation are dissimilar to those in Table 1, the correspondence between the graphic symbols used in the implementation and those in Table 1 shall be given.
- A description of the **numbers**. This shall include a characterisation of the internal representation used for **numbers**.
- Descriptions of the characteristics of each **implementation-algorithm**.
- The value of each **implementation-parameter**.
- A description of each internal-value-set.
- A description of additional event-types.

4.1.2.3 Consistent Extensions

A **conforming-implementation** shall document all **consistent-extensions** it provides. The documentation shall clearly indicate that the use of a **consistent-extension** prevents a **program** from conforming with this International Standard.

Note: Implementers of conforming-implementations should, in general, be wary of replacing limit-errors with consistent-extensions, since these errors are the only safeguards a conforming-program has when attempting to operate in a conforming-implementation whose implementation-parameters are inadequate to support it.

For example, if the **limit-error** on **identifier-length-limit** were not signalled, a **conforming-program** with identifiers longer than the local **identifier-length-limit** could malfunction without warning.

4.2 Conforming Programs

A **program** conforms to this International Standard if it meets the following requirements for both behaviour and documentation.

4.2.1 Required Behaviour for Conforming Programs

A **conforming-program** shall use only those **facilities** specified in this International Standard.

A conforming-program shall not use consistent-extensions.

A **conforming-program** shall not depend on the signalling of any error by a **conforming-implementation**.

Note: Conforming-programs cannot depend on error behaviour, since **consistent-extensions** that replace errors are permitted in **conforming-implementations**. A conforming program may make use of certain event handling features unreservedly (for example, $\square ES$), and others with some restrictions. Consider the following example:

$$Z \leftarrow 1$$
' $\square EA$ ' A $\square SVO$ B ' $Z \leftarrow (0 \ 1 \ 2 \ 1) [0 \ 1 \ 2 \ Z]$

This is conforming because it only depends on the values 0, 1, and 2, which have defined meanings when returned as the result of $\square SVO$. The first line can return anything, but the second line filters out all but 0, 1, and 2, hence eliminating any subtle dependencies on error behaviour or consistent extensions.

4.2.2 Required Documentation for Conforming Programs

A **conforming-program** shall document which **optional-facilities** it requires.

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A **conforming-program** shall document any specific requirements it has for **implementation-parameters**.

Note: A **conforming-program** may or may not work, and may or may not produce identical results on all **conforming-implementations**, because of inherent dependencies on **implementation-parameters** or **implementation-algorithms** such as the algorithm used for matrix inversion.

It is suggested that **conforming-programs** provide documentation detailing the values of **implementation-parameters** they require so that their suitability for a given **conforming-implementation** can be determined readily.

5 Definitions

5.1 Characters

- Character-Set: An implementation-defined finite set.
- Character: A member of the implementation-defined finite set character-set.
- Required-Character-Set: An enumerated set which is the union of a set of alphabetics with the sets of Numeric Characters and Special Characters designated in Table 1. The alphabetics may be any alphabet; this International Standard uses the Alphabetic Characters shown in Table 1 to represent the dual-case Latin alphabet. The character-set provided by a conforming-implementation shall contain all members of the required-character-set which results from the choice of alphabet.

Additional Requirement:

Members of the **required-character-set** are represented in this International Standard by graphic symbols in a particular typeface, as shown in **Table 1**. The graphic symbols associated with the **required-character-set** in a **conforming-implementation** are **implementation-defined**.

Note: There are conformance requirements associated with the **required-character-set**. A **conforming-implementation** must publish, as part of its required documentation, a table giving the correspondence between the graphic symbols in **Table 1** and the **atomic-vector**. Where the graphic symbols provided are not similar in appearance to those used in this International Standard, the correspondence between the graphic symbols in this International Standard and those provided by the implementation must also be provided.

Some of the graphic symbols given in Table 1 are not used to designate APL constructs in this International Standard. They are present to provide all keys on common terminal keyboards with corresponding symbols.

The names in **Table 1** are not to be considered part of this International Standard.

Alphabetic Characters

Α	В	C	D	E	F	G	H	Ι	J	Κ	L	Μ	N	0	P	Q	R	\mathcal{S}	T	U	V	W
Χ	Y	Z																				
<u>A</u>	<u>B</u>	<u>C</u>	\underline{D}	<u>E</u>	<u>F</u>	\underline{G}	<u>H</u>	<u>I</u>	\underline{J}	<u>K</u>	\underline{L}	\underline{M}	<u>N</u>	<u>0</u>	<u>P</u>	\underline{Q}	<u>R</u>	<u>S</u>	\underline{T}	<u>U</u>	\underline{V}	\underline{W}
<u>X</u>	<u>Y</u>	<u>Z</u>																				

Numeric Characters

0 1 2 3 4 5 6 7 8 9

Special Characters

α alpha	₹ del tilde	ľ	quote quad
↓ down arrow	∆ delta	ρ	rho
← left arrow		;	semicolon
→ right arrow	$\underline{\Delta}$ delta underbar	U	down shoe
↑ up arrow	· diaeresis	\subset	left shoe
– bar	♦ diamond	\supset	right shoe
blank	÷ divide	Λ	up shoe
{ left brace	\$ dollar sign	А	up shoe jot
} right brace	• dot	/	slash
[left bracket	€ epsilon	\	back slash
] right bracket	= equal	+	slash bar
∨ down caret	≥ greater-than or equal	+	back slash bar
→ down caret tilde	ı iota	*	star
< left caret	o jot		stile
> right caret	≤ less-than or equal	L	down stile
∧ up caret	× multiply	Γ	up stile
↑ up caret tilde	≠ not equal	\perp	up tack
o circle	overbar overbar	<u>Φ</u>	up tack jot
◊ circle backslash	(left parenthesis	\vdash	right tack
⊖ circle bar) right parenthesis	-	left tack
⊗ circle star	+ plus	Т	down tack
φ circle stile	☐ quad	φ	down tack jot
: colon	: quad divide	~	tilde
, comma	? query		underbar
, comma bar	o diaeresis jot	~	diaeresis tilde
∇ del	† quote	ω	omega
[†] del stile	! quote dot	=	equal underbar

Table 1: The Required Character Set

5.2 Numbers

- Number-Set: An implementation-defined finite set used to represent arithmetic quantities.
- **Number**: A member of the **number-set**.

Note: The **number-set** is an abstraction used in this document to represent the floating-point arithmetic quantities of an arbitrary computer.

5.2.1 Elementary Operations

Elementary-Operation: One of the following four implementation-algorithms.

- A **Plus** B.

Note: Plus implements addition.

- A Minus B.

Note: Minus implements subtraction.

A Times B.

Note: Times implements multiplication.

A Divided-by B.

Note: Divided-by implements division.

Note: Each elementary-operation is a mapping from the Cartesian product of the number-set with itself back onto the union of the number-set and the metaclass error, as described in the subsection Implementation-Algorithms. It is assumed that members of the metaclass error are returned by elementary-operations when results cannot be represented as numbers. In particular, exponent-overflow is returned when the result of an elementary-operation is too large in magnitude to be represented by a number, exponent-underflow is returned when the result of an elementary-operation is non-zero but too small in magnitude to be represented by a number other than zero, and domain-error is returned by an elementary-operation called with arguments for which its mathematical counterpart is undefined, such as one divided-by zero.

The treatment of exponent-overflow and exponent-underflow is not specified by this International Standard, except as suggestions to the implementer. Implementers should avoid having exponent-overflow and exponent-underflow occur in the intermediate results of implementation-algorithms wherever possible. In any case, exponent-underflow should not cause a limit-error.

The fact that these fundamental algorithms are effectively undefined in the International Standard is intentional. A definition general enough to cover all known and possible floating-point systems was felt to be less useful than the requirement that the internal representation for numbers in a conforming-implementation be described in the required-documentation for the implementation.

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5.2.2 Number Constants

Note: The following **numbers**, used in this International Standard, are defined here as terms to make clear the distinction between APL constant arrays represented as strings of APL digits, such as 1, and members of the the **implementation-defined** finite set containing **numbers**, such as **one**.

These terms are always set in **bold** when they are used in this formal sense. They are not cross-indexed.

- **Zero**: A **number** such that, for any **number** A, A **plus zero** is A.
- One: A number such that, for any number A, A times one is A.
- Negative-One: Zero minus one.
- Two: One plus one.
- Three: Two plus one.
- Four: Three plus one.
- Five: Four plus one.
- Six: Five plus one.
- One-half: One divided-by two.
- Imaginary-One: Negative-One to-the-power one-half.

5.2.3 Subsets of the Set of Numbers

Note: The following subsets of the numbers are used to define the domains of operations.

- Boolean: zero or one.

Note: False is represented by zero, true by one.

Note: A **counting-number** is a member of the subset of the **numbers** accessible under a successor function.

- Positive-Counting-Number: one or any number that can be generated by A plus one, where A is a positive-counting-number.
- Negative-Counting-Number: The number negative-one or any number that can be generated by N plus negative-one, where N is a negative-counting-number.
- Nonnegative-Counting-Number: A positive-counting-number or zero.

Counting-Number: A negative-counting-number or a nonnegative-counting-number.

Note: An **integer** is a member of the subset of the **numbers** accessible from **zero** and **one** through addition and subtraction.

- Positive-Integer: A positive-counting-number or a number that can be generated by A plus B, where A and B are positive-integers.
- Negative-Integer: A negative-counting-number or a number that can be generated by A plus B, where A and B are negative-integers.
- Nonnegative-Integer: A positive-integer or zero.
- Integer: A negative-integer or a nonnegative-integer.
- Positive-Number: A positive-integer, or a number other than zero that can be generated by A times B, by A divided-by B, or by A plus B, where A and B are positive-numbers.
- Negative-Number: A negative-integer, or a number other than zero that can be generated by A times B, by A divided-by B, or by C plus D, where A is a positive-number and B, C and D are negative-numbers.
- Nonnegative-Number: A positive-number or zero.
- Real-Number: A negative-number or a nonnegative-number.
- Complex-Number: A real-number, or a number that can be generated by A plus (B times imaginary-one), where A and B are real-numbers.
- Complex-Integer: A complex-number that can be generated by A plus (imaginary-one times B), where A and B are integers.
- Half-plane: Two numbers are in the same half-plane if the signs of either their real-parts or imaginary-parts are nonzero and are equal.
- **Unit-square**: A **unit-square** contains all numbers for which the integral part of the real parts are the same and the integral parts of the imaginary parts are the same.
- Arc: The arc of a number is zero if the number is zero, and otherwise is the imaginary-part of the natural-logarithm of the number, and has a value greater than minus pi and less than or equal to pi.
- Fraction: A number lying in the rectangle of the complex-plane bounded by the two parallel lines passing through zero and one at an angle of 135 degrees with the real axis, and the two parallel lines passing through one and negative-one at an angle of forty-five degrees with the real axis, including the numbers on the lines going through zero and negative-one, but not including the numbers on the two lines through one.

First-Quadrant-Associate: Either a nonnegative-number, or a number with both real-part and imaginary-part positive-numbers, and having the same magnitude as a given nonzero number, and separated from it by an angle of either 0, 90, 180, or 270 degrees.

Note: The numbers are assumed to be identical to the complex-numbers.

Example

Consider a normalised base-2 floating point number system with a two position exponent field and a two position mantissa field with the radix point between the digits. The decimal values representable with this system and the categories into which each falls are as follows:

-	Mantissa in Base 2	Value in Base 10	Boolean	Counting- Number	Integer
00	00	0	Y	Y	Y
00	10	1	Y	Y	Y
00	11	1.5			
01	10	2		Y	Y
01	11	3		Y	Y
10	10	4		Y	Y
10	11	6			Y
11	10	8			Y
11	11	12			Y

As can be seen in this example, **counting-numbers** form a dense subset of the **integers**; the largest **counting-number** is typically a power of the number system base.

A given number, such as a **counting-number**, may have several hardware representations. Except for their effect on system resources, these representations should be indistinguishable to a **conforming-program**.

5.2.4 Implementation Algorithms

An **Implementation-Algorithm** is an algorithm used in this International Standard whose behaviour is **implementation-defined**. The following implementation algorithms are used in this International Standard.

- Conjugate
- Cosine
- Current-Time
- Deal
- Display
- Divided-by

- Exponential
- Function-Display
- Gamma-Function
- Greatest-Common-Divisor
- Hyperbolic-Cosine
- Hyperbolic-Sine
- Hyperbolic-Tangent
- Inverse-Cosine
- Inverse-Hyperbolic-Cosine
- Inverse-Hyperbolic-Sine
- Inverse-Hyperbolic-Tangent
- Inverse-Sine
- Inverse-Tangent
- Magnitude
- Matrix-Divide
- Minus
- Modulo
- Natural-Logarithm
- Next-Definition-Line
- Numeric-Input-Conversion
- Numeric-Output-Conversion
- Pi-Times
- Plus
- Produce-Canonical-Representation-Vector
- Pseudorandom-Number-Generator
- Read-Keyboard
- Sine
- Tangent
- Times

- Time-Stamp
- To-the-Power
- Trace-Display
- Typical-Element-For-Mixed

The following **implementation-algorithms** take a **number** as an argument and return either a **number** or an **error**: **Cosine**, **Exponential**, **Gamma-Function**, **Hyperbolic-Cosine**, **Hyperbolic-Tangent**, **Inverse-Cosine**, **Inverse-Hyperbolic-Sine**, **Inverse-Hyperbolic-Tangent**, **Inverse-Sine**, **Inverse-Sine**, **Inverse-Tangent**, **Magnitude**, **Natural-Logarithm**, **Pi-Times**, **Sine**, **Tangent**.

The following **implementation-algorithms** take two **numbers** as arguments and return either a **number** or an **error**: **Divided-by**, **Minus**, **Modulo**, **Plus**, **Times**, **To-the-Power**.

The properties of the remaining **implementation-algorithms** are described in the chapters in which they are used.

5.2.5 Defined Operations

- A Equals B: An operation that, for A and B numbers, returns one if A and B are the same number, and zero otherwise.
- Direction of A: An operation that, for A a number, returns zero if A is zero, and otherwise returns the number determined by the radial projection of A onto the unit-circle.

Note: For A a real-number, the direction of A is either negative-one or zero or one.

- A is Greater-Than B: An operation that, for A and B real-numbers returns one if A is a positive-number and B is a negative-number, returns one if A minus B is a positive-number, and returns zero otherwise.
- A is Less-Than B: An operation that, for A and B real-numbers, returns B greater-than A.
- Negation of A: An operation that, for any number A, returns zero minus A.
- Magnitude of A: An operation that, for A a number, returns the non-negative real-number determined by rotating A onto the nonnegative real-axis.
- Open-Interval-Between A and B: An operation that, for any numbers A and B, returns a subset of the numbers as follows: if A is not greater-than B, the set of all numbers greater-than A and less-than B; otherwise, the open-interval-between B and A.
- Closed-Interval-Between A and B: An operation that, for any numbers A and B, returns
 a subset of the numbers consisting of A, B, and the open-interval-between A and B.
- Larger-Magnitude of A and B: An operation that, for any numbers A and B, returns the magnitude of A if it is greater-than that of B, and the magnitude of B otherwise.

- Distance-Between A and B: An operation that, for any two numbers A and B, returns the magnitude of A minus B.
- A is Tolerantly-Equal to B Within C: An operation that, given three numbers A, B, and C, with C greater than or equal to zero, returns a Boolean Z determined as follows:

If A equals B, then Z is one.

If A and B are not in the same **half-plane**, then Z is **zero**.

If the **distance-between** A and B is **less-than** or **equals** C **times** the **larger-magnitude** of A and B, then Z is **one**.

Otherwise, Z is **zero**.

Tolerant-Floor of A Within B: An operation that, for A a number and B a nonnegative-number, returns a complex-integer Z determined as follows:

Let A be a member of the set of **numbers** in the **unit-square** at the **complex-integer** C, and let D be A **minus** C.

If the sum of the real and imaginary parts of D is **tolerantly-less-than one within** B, then Z is C.

Otherwise, if the **imaginary-part** of D is **greater-than** the **real-part** of D, then Z is C **plus imaginary-one**.

Otherwise, Z is C plus one.

- A is Integral-Within B: An operation that, for a number A and a positive-number B, returns a Boolean Z determined as follows:

Let C stand for the **negation** of A.

Z is **one** if the **tolerant-floor** of C **within** B **equals** the **negation** of the **tolerant-floor** of A **within** B, and **zero** otherwise.

A is a Near-Integer: An operation that, for a number A, returns one if A is integral-within integer-tolerance, and zero otherwise.

Note: *The foregoing definition contains a forward reference to* **integer-tolerance**.

- Integer-Nearest-to A: An operation that, for a near-integer A, returns the tolerant-floor
 of A within integer-tolerance.
- A is Near-Boolean: An operation that, for a near-integer A, returns one if the integer-nearest-to A is a Boolean, and zero otherwise.

Note: Near-integers and near-Booleans include numbers whose magnitude is less-than integer-tolerance. The operation integer-nearest-to maps such numbers to zero.

- A is Real-Within B: An operation that, for a number A and a positive-number B, returns a Boolean Z determined as follows:
 - Let R stand for the **magnitude** of the **real-part** of A, and I for the **magnitude** of the **imaginary-part** of A.
 - Z is **one** if I is **less-than-or-equal** either to B, or to R **times** B, and is **zero** otherwise.
- A is Near-Real: An operation that, for a number A, returns one if A is real-within real-tolerance, and zero otherwise.

5.3 Objects

Note: The description of APL in this document is based on objects—abstract data structures that are described in terms of characters, numbers, and elementary set theory.

Each object has a small set of named attributes that take on values that are characters, numbers, or other objects.

Each object has additionally a small set of operations that can be performed upon it.

All the objects in the description are defined here, as are all the attributes of each object. The operations are introduced as they are needed. All objects, attributes, and operations are cross-referenced in the index.

5.3.1 Lists

Index: A nonnegative-counting-number less than or equal to index-limit.

List: An object with the following attributes:

- **Index-Set:** A finite set \mathcal{I} of **positive-counting-numbers** chosen so that for every subset of \mathcal{I} , except the empty set, the cardinality of that subset is in \mathcal{I} .
- Value-Set: A finite set in a specific correspondence with the index-set of the list.

Empty-List: A list, the cardinality of whose index-set is zero.

Nonempty-List: A list, the cardinality of whose index-set is greater-than zero.

Number-of-Items in L: The cardinality of the **index-set** of the **list** L.

Item X of L: An operation that, for X a member of the **index-set** of the **list** L, returns the member of the **value-set** of L associated with X by the correspondence between the **value-set** of L and the **index-set** of L.

Note: This form of indexing is always in origin one.

First-Item in L: An operation that for a **nonempty-list** L returns **item one** of L.

Last-Item in L: An operation that for a **nonempty-list** L, and for C the **number-of-items** in L, returns **item** C of L.

Rest-of L: An operation that, for a **nonempty-list** L whose **index-set** has cardinality C, returns a second **list** R whose **index-set** has cardinality C **minus one**, such that for each **item** J in the **index-set** of R, **item** J of R is **item** (J **plus one**) of L.

Product-of L: An operation defined on a **list** of **numbers** L as follows:

If the **number-of-items** in \mathcal{L} is **zero**, **one**.

If the **number-of-items** in L is **one**, the **first-item** in L.

If the number-of-items in L is greater-than one, first-item in L times the product-of the rest-of L.

Prefix: A (possibly empty) **list** P is a **prefix** of a **list** L if the **index-set** of P is a subset of the **index-set** of L, and each **item** of P is the same as the corresponding **item** of L.

5.3.2 Arrays

Array-Type: An enumerated set containing the members character, numeric, and mixed.

Array: An object with the following attributes:

- Shape-list: A list of nonnegative-counting-numbers.
- Ravel-list: A list. Each item of the list is either a character, a number, or an array.
 The number-of-items in the ravel-list of an array A is the same as the product-of the shape-list of A.
- Type: A member of the enumerated set array-type.
 If the type of A is character, the ravel-list of A is a list of characters; if the type of A is numeric, the ravel-list of A is a list of numbers.

Sufficient-Type of L under T: the following operation on a **type** T and a **list** L.

If T is **character** or **numeric**, return T

else if L is a **nonempty-list**

if all items of L are **numbers**, return the **type numeric**

if all items of L are **characters**, return the **type character**

otherwise, return the type mixed

else return the **type** of the **typical-element** of L.

Note: The foregoing definition is intended to provide for the cases where a mixed-array becomes a non-mixed array, and vice-versa.

Rank of A: An operation that, for an **array** or **array-of-vectors** A, returns the **number-of-items** in the **shape-list** of A.

Simple: An **array** is **simple** if each item of its **ravel-list** is either a **character** or a **number**.

Scalar: An array whose rank is zero.

Simple-scalar: A simple array whose rank is zero.

Max-shape-of B An operation that, for B, an **array**s defined as follows:

Set B1 to the **first-thingy** of B.

Set M1 to $\rho B1$.

If the **count-of** B is **one**, return M1.

Set B2 to the **remainder-of** B.

Set M2 to the **max-shape-of** B2.

If the **shape-list** of *M*1 differs from the **shape-list** of *M*2, signal **rank-error**.

Return *M*1 *√ M*2

First-thingy in A: An operation that for A, an **array**, returns an **array** B, defined as follows: If A is empty, set B1 to the **typical-element** of A.

Otherwise, set B1 to the **first-item** of the **ravel-list** of A.

If B1 is a **number** or a **character**, set B to an array, whose **ravel-list** contains the single item B1, and whose **shape-list** is empty.

Otherwise, set B to B1.

Numeric-Scalar with value I: For I a **number**, the **scalar** Z such that the **type** of Z is **numeric** and the **ravel-list** of Z is the **list** L such that the **number-of-items** in L is **one** and the **first-item** of L is L.

Count of A: For A an array, the **product-of** the **shape-list** of A.

Vector: An array whose rank is one.

Matrix: An array whose rank is two.

Length of A: For A a **vector**, the **count** of A.

Typical-Element of A: For A an **array**, if the **type** of A is **character**, the **character** blank; if the **type** of A is **numeric**, the number **zero**; if the **type** of A is **mixed**, the **typical-element** is determined by an **implementation-defined-algorithm**.

Empty: Of an array, the count of which is zero.

One-Element-Vector: A vector A is a one-element-vector if the length of A is one.

K is a **Valid-Axis** of A: An operation that, for A and K **arrays**, returns **one** if K is a **scalar** or **one-element-vector**, and the **first-item** in the **ravel-list** of K is a **near-integer**, the **integer-nearest-to** which is a member of the **index-set** of the **shape-list** of A; and returns **zero** otherwise.

Axis K of A: An operation that, for A an **array** and K a **valid-axis** of A, returns **item** K of the **shape-list** of A.

Array-of-vectors: An object with the following attributes:

- Shape-list: A list of nonnegative-counting-numbers.
- Ravel-list: Each vector V in the ravel-list of an array-of-vectors A has the type sufficient-type of the ravel-list of V under the type of A.

- **Type:** A member of the enumerated set **array-type**.

Note: The term array-of-vectors is used only as an expository device.

Along-Axis K of A: An operation that, for an **array** A with non-zero **rank** N, produces Z, an **array-of-vectors** of **rank** N-1 such that the **shape-list** of Z is the **shape-list** of A with **item** K omitted. Each **item** of the **ravel-list** of Z is a **vector** whose **length** is the same as **axis** K of A.

Note: Some operations on vectors extend to arrays of greater rank in a manner similar to **scalar-extension**. **Along-axis** is an expository device used in the definition of these operations.

Vector-Item I of A: An operation that, for an **array-of-vectors** A, returns **item** I of the **rayel-list** of A.

Ravel-Along-Axis K of A: An operation that, for an **array** A with non-zero **rank** N, produces an **array-of-vectors** Z such that the **shape-list** of Z is the **product-of** the **shape-list** of A with **item** K omitted, and the **ravel-list** of Z is the **ravel-list** of **along-axis** X of A.

First-Scalar in A: An operation that, for A a nonempty **array**, returns a **scalar** Z such that the **ravel-list** of Z is the **first-item** in the **ravel-list** of A and the **type** of Z is the **sufficient-type** of the **ravel-list** of Z under the **type** of A.

Remainder-of A: An operation that, for A a non-empty **vector**, returns a **vector** Z such that the **length** of Z is **negative-one plus** the **length** of A, the **ravel-list** of Z is the **rest-of** the **ravel-list** of A and the **type** of A is the **sufficient-type** of the **ravel-list** of A under the **type** of A.

Row I of A: An operation that, for A an **array** of **rank two** and I a **number**, returns **vector-item** I **along-axis two** of A.

Number-of-Rows in A: An operation that, for A an **array** of **rank two**, returns **item one** of the **shape-list** of A.

Integer-Array-Nearest-to A: An operation that, for A an **array** having the property that each **item** of the **ravel-list** of A is a **near-integer**, returns a **numeric array** Z such that the **shape-list** of Z is the same as the **shape-list** of A and each **item** of the **ravel-list** of Z is the **integer-nearest-to** the corresponding **item** of the **ravel-list** of A.

Boolean-Array-Nearest-to A: An operation that, for A an **array** having the property that each **item** of the **ravel-list** of A is a **near-Boolean**, returns the **integer-array-nearest-to** A.

5.3.3 Defined-Functions

Note: A defined-function represents a function defined by a user. The attributes of a defined-function are given here. Operations are described in the Defined Functions chapter. Defined-operators are a subclass of defined functions.

Defined-Function: An object with the following attributes:

- Canonical-Representation: A character array whose rank is two.
- Stop-Vector: A numeric vector.
- Trace-Vector: A numeric vector.

5.3.4 Tokens

Class-Names: A set containing the following members:

Assignment-Arrow Local-Name Axis-Error Monadic-Operator

Branch Nil

Branch-Arrow Niladic-Defined-Function
Character-Literal Niladic-Defined-Function-Name
Clear-State-Indicator Niladic-System-Function-Name

Colon Not-Copied
Command-Complete Not-Erased
Committed-Value Not-Found
Complete-Index-List Not-Saved
Constant Numeric-Literal
Defined-Function Partial-Index-List

Defined-Function-Name Primitive

Defined-Dyadic-Operator Primitive-Function
Defined-Dyadic-Operator-Name Rank-Error

Defined-Dyadic-Operator-Name Rank-Error Defined-Monadic-Operator Result-Name

Defined-Monadic-Operator-Name
Definition-Error
Distinguished-Identifier
Domain-Error
Dyadic-Operator
Elided-Index-Marker

Right-Argument-Name
Right-Axis-Bracket
Right-End-of-Statement
Right-Index-Bracket
Right-Operand-Name
Right-Parenthesis

Escape Semicolon
Implicit-Error Shared-Variable
Incorrect-Command Shared-Variable-Name
Index-Error Simple-Identifier
Index-Separator Small-Circle
Interrupt Syntax-Error

Label System-Function-Name Label-Name System-Variable-Name

Left-Argument-Name Unwind
Left-Axis-Bracket Value-Error
Left-End-of-Statement Variable
Left-Index-Bracket Variable-Name

Left-Operand-Name Left-Parenthesis Length-Error Limit-Error

Token: An object with the following attributes:

Class: A member of the set Class-Names.

Content: A number, a character, or an object, according to the class of the token, as indicated in Table 2.

Note: Tokens represent the internal objects manipulated by an implementation of APL.

The class of a token is a straightforward indication of the sort of object it represents. A token of class nil, for example, is used to return a value-error from a defined function that does not set its result name. Note that class-names is an enumerated set. This means that nil has no definition other than its literal appearance on this page. Its significance, like the significance of APL characters in this document, lies in the use of its name in evaluation sequences.

The content of a token varies with the class of the token. If specified, the content is a character, a list of characters, a list of numbers, an array, an index-list, a shared-variable, or a defined-function.

5.3.4.1 Metaclasses

Note: Metaclasses are subsets of the enumerated set **class-names**. They are used to shorten evaluation sequences by abstracting a sequence of tests on the **class** of a **token** into a single test for membership in a metaclass. "If B is an **error**," is equivalent to "If the **class** of B is in the metaclass **error**."

- Metaclass: A subset of the enumerated set class-names.
- Identifier: A metaclass containing simple-identifier and distinguished-identifier.
- Literal: A metaclass containing character-literal and numeric-literal.
- Lexical-Unit: A metaclass containing primitive, literal, and identifier.
- Value: A metaclass containing committed-value and constant.
- Delimiter: A metaclass containing primitive-function, branch-arrow, assignment-arrow, left-end-of-statement, right-end-of-statement, left-index-bracket, right-index-bracket, elided-index-marker, left-axis-bracket, right-axis-bracket, left-parenthesis, right-parenthesis, small-circle, and semicolon.
- Defined-Name: A metaclass containing shared-variable-name, variable-name, defined-function-name, defined-dyadic-operator-name, defined-monadic-operator-name, niladic-defined-function-name, result-name, left-argument-name, right-argument-name, label-name, and local-name.
- Defined-Operator: A metaclass containing defined-dyadic-operator and defined-monadic-operator.
- System-Name: A metaclass containing system-variable-name, system-function-name, and niladic-system-function-name.
- Classified-Name: A metaclass containing the members of system-name and defined-name.
- Syntactic-Unit: A metaclass containing the members of classified-name and delimiter.

Class-Name Content

Assignment-Arrow

Axis-Error

Branch An array

Branch-Arrow

Character-Literal A list of characters

Clear-State-Indicator

Colon

Command-Complete

Committed-ValueAn arrayComplete-Index-ListAn index-listConstantAn array

Defined-FunctionA defined-functionDefined-Function-NameA list of charactersDefined-Dyadic-OperatorA defined-functionDefined-Monadic-OperatorA list of charactersDefined-Monadic-Operator-NameA list of charactersDefined-Monadic-Operator-NameA list of characters

Definition-Error

Distinguished-Identifier A list of characters

Domain-Error

Dyadic-Operator A character

Elided-Index-Marker

Escape

Implicit-Error
Incorrect-Command

Index-Error Index-Separator Interrupt

Label An array

Label-NameA list of charactersLeft-Argument-NameA list of characters

Left-Axis-Bracket Left-End-of-Statement Left-Index-Bracket

Left-Operand-Name A list of characters

Left-Parenthesis Length-Error Limit-Error

Table 2: Relationship between Class-Name and Content

Class-Name	Content
Local-Name	A list of characters
Monadic-Operator	A character

Nil

Niladic-Defined-FunctionA defined-functionNiladic-Defined-Function-NameA list of charactersNiladic-System-Function-NameA list of characters

Not-Copied Not-Erased Not-Found Not-Saved Numeric-Literal

Numeric-LiteralA list of numbersPartial-Index-ListAn index-listPrimitiveA characterPrimitive-FunctionA character

Rank-Error

Result-NameA list of charactersRight-Argument-NameA list of characters

Right-Axis-Bracket Right-End-of-Statement Right-Index-Bracket

Right-Operand-Name A list of characters

Right-Parenthesis

Semicolon

Shared-VariableA shared-variableShared-Variable-NameA list of charactersSimple-IdentifierA list of characters

Small-Circle Syntax-Error

System-Function-Name A list of characters System-Variable-Name A list of characters

Unwind Value-Error

Variable An array

Variable-Name A list of characters

Table 2: (Continued)

- Error: A metaclass containing axis-error, domain-error, implicit-error, index-error, length-error, limit-error, rank-error, syntax-error, value-error, and interrupt.

Note: This metaclass includes only errors that occur in the evaluation of APL statements.

- Report: A metaclass containing incorrect-command, not-copied, not-erased, not-found, and not-saved.
- Exception: A metaclass containing branch, escape, clear-state-indicator, unwind, and the members of error and report.
- **Result**: A metaclass containing **nil** and the members of **exception** and **value**.

5.3.4.2 Index-List

 Index-List: A (possibly empty) list consisting of tokens whose class is either constant or elided-index-marker.

5.3.5 Symbols

Symbol: An object with the following attributes:

- Name: A list of characters.
- Referent-List: A list of tokens.

5.3.6 Contexts

 Mode-Names: An enumerated set containing the members immediate-execution, execute, function-definition, quad-input, and defined-function.

Context: An object with the following attributes:

- Mode: An element of the enumerated set mode-names.
- Stack: A list of tokens.
- Current-Line: A list of characters.
- Current-Statement: A list of tokens.
- Current-Function: If mode is defined-function, a defined-function; otherwise undefined.
- Current-Line-Number: If mode is defined-function, an index; otherwise undefined.

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5.3.7 Workspaces

Note: The workspace is the basic organisational unit in an APL system. A workspace contains data, programs, execution status, and environmental information.

- Workspace-Presence: An enumerated set containing absent and present.

Workspace: An object with the following attributes:

- Owner: A user-identification.
- Workspace-Name: A list of characters.
- Symbol-Table: A list of all symbols whose names are distinct.
- State-Indicator: A list of contexts.
- Existential-Property: A member of the enumerated set workspace-presence.

Clear-Workspace: A workspace with the following values:

- Owner: this-owner.
- Workspace-Name: The clear-workspace-identifier.
- Symbol-Table: A list of all symbols whose names are distinct and whose referent-lists
 each consist of the list whose only member is nil.
- State-Indicator: An empty list of contexts.
- Existential-Property: absent.

5.3.8 Sessions

Note: A user interacts with an APL system through a session, an abstraction that represents a hypothetical machine capable of carrying out the evaluation sequences in the International Standard.

- Session-Identification: Either a number or a list of characters, depending upon the implementation-parameter session-identification-type.
- User-Identification: Either a number or a list of characters, depending upon the implementation-parameter user-identification-type.

 Keyboard-States: An enumerated set containing the members open-keyboard and locked-keyboard.

Session: An object with the following attributes:

- Active-Workspace: A workspace.
- This-Session: A session-identification.
- This-Owner: A user-identification.
- Attention-Flag: A member of the enumerated set Boolean.
- **Keyboard-State:** A member of the enumerated set **keyboard-states**.
- Current-Prompt: A character vector.
- Quote-Quad-Prompt: A character vector.
- Event-Time: A nonnegative-number.
- Event-Message: A character array giving, in an implementation-defined form, the
 event-report and any other information the implementation deems helpful in locating
 the error's source.
- Event-Type: A two-element integer vector corresponding to the most recent event.
 Defined values are:
 - 0 0 No error
 - 0 1 Undefined event

There may be additional **implementation-defined** values as well as values supplied by a program via $\Box ES$.

- Current-Context: The first-item in the state-indicator of the active-workspace.
- Current-Stack: The stack of the current-context.
- Symbol-Named-By T: The symbol in the symbol-table of the active-workspace whose
 name is the same as the content of the token T.
- Current-Referent of T: The first-item in the referent-list of the symbol-named-by T.
- Current-Class of T: The class of the current-referent of T.
- Current-Content of T: The content of the current-referent of T.
- Comparison-Tolerance: The current-content of $\square CT$.
- **Random-Link**: The **current-content** of $\square RL$.

- **Print-Precision**: The **current-content** of $\square PP$.
- Index-Origin: The current-content of $\square IO$.
- Latent-Expression: The current-content of $\square LX$.

System-Parameter: Any of comparison-tolerance, random-link, print-precision, index-origin, or latent-expression.

The initial values of **system-parameters** in a **clear-workspace** are **implementation-defined**.

5.3.9 Shared-Variables

Note: A shared-variable is a variable shared between two sessions.

Shared Variables are an **optional-facility**.

Shared-Variable: An object with the following attributes:

- Session-A: A session-identification.
- Session-A-Active: A Boolean.
- Session-A-ACV: A Boolean vector of length four.

Note: ACV stands for access control vector.

- Session-B: A session-identification.
- Session-B-Active: A Boolean.
- Session-B-ACV: A Boolean vector of length four.
- Shared-Name: An identifier.
- Shared-Value: A token, either a constant or nil.
- State: An integer, either zero, one, or two.
- Session-A-Event: A Boolean.
- Session-B-Event: A Boolean.

Note: The operations in the chapter Shared Variables use the shared-variable attributes as follows:

- Session-A: The session-identification of the first session to offer the shared-variable.

- Session-A-Active: A Boolean; one if session-A is currently sharing this shared-variable, zero
 if not.
- Session-A-ACV: The contribution of session-A to the ACV.
- Session-B: The session-identification of the session with which session-A offered to share the shared-variable; this may be the general-offer while state is one.
- Session-B-Active: A Boolean; one if session-B is currently sharing this shared-variable, zero if not
- **Session-B-ACV**: The contribution of **session-B** to the ACV.
- Shared-Name: An identifier, the name session-A designated for this shared-variable.
- Shared-Value: A token of class constant or nil; the current value of this shared-variable.
- State: An integer, either zero, one, or two, used in combination with the ACV to determine which
 operations must be delayed.

5.3.10 Systems

Note: A **system** represents a set of active APL users, a library, and, optionally, a shared variable facility.

System: An object with the following attributes:

- Library: A list of workspaces in which each combination of possible values for the attributes owner and workspace-name occurs exactly once.
- Active-Users: A list of sessions.
- Shared-Variable-List: A list of shared-variables.
- Implementation-Parameters: The following quantities, referred to in this International Standard by name, whose values are implementation-defined:
 - Atomic-Vector: An implementation-defined character vector containing every member of the required-character-set exactly once.
 - Initial-Comparison-Tolerance: A member of the internal-value-set for comparison-tolerance that is the value of comparison-tolerance in a clearworkspace.
 - Initial-Index-Origin: A member of the internal-value-set for index-origin that is the value of index-origin in a clear-workspace.
 - Initial-Latent-Expression: A member of the internal-value-set for latent-expression that is the value of latent-expression in a clear-workspace.
 - Initial-Print-Precision: A member of the internal-value-set for print-precision that
 is the value of print-precision in a clear-workspace.
 - Initial-Random-Link: A member of the internal-value-set for random-link that is the value of random-link in a clear-workspace.

- Initial-Event-Message: The empty-event-message.
- Reduction-Style: One of the two symbolics Enclose-Reduction-Style or Insert-Reduction-Style, indicating which of the two styles of reduction the system uses.
- **Initial-Event-Type**: The two-element integer vector (0,0).
- Clear-Workspace-Identifier: A list of characters.
- Positive-Number-Limit: The real-number greater-than all other real-numbers.
- Negative-Number-Limit: The real-number less-than all other real-numbers.
- Positive-Counting-Number-Limit: The counting-number greater than all other counting-numbers.
- Negative-Counting-Number-Limit: The counting-number less-than all other counting-numbers.
- Index-Limit: The index greater than all other indices. This value specifies the
 maximum value of any item of the shape-list of any array, ignoring storage
 limitations.
- Count-Limit: An index not greater than index-limit that specifies the maximum value for the count of an array, ignoring storage limitations.
- Rank-Limit: An index not greater than count-limit that specifies the maximum value for the rank of any array, ignoring storage limitations.
- Workspace-Name-Length-Limit: A positive-counting-number that specifies the maximum number of characters in a workspace-name.
- Identifier-Length-Limit: A positive-counting-number not greater than count-limit that specifies the maximum number of characters in an identifier.
- Quote-Quad-Output-Limit: A nonnegative-counting-number that specifies the maximum number of characters that can be used in a prompt set by Quote Quad Output.
- Comparison-Tolerance-Limit: The largest real-number permitted by the implementation for the system parameter comparison-tolerance.
- Integer-Tolerance: A real-number not greater than comparison-tolerance-limit used to determine whether a given number is to be considered integral.
- Real-Tolerance: A nonnegative-number not greater than comparison-tolerance-limit used to determine whether a given number is to be considered real.
- Full-Print-Precision: The smallest positive-counting-number which, when used as
 the value of print-precision, causes numeric-output-conversion to produce a unique
 numeric-scalar-literal for every number.
- Print-Precision-Limit: The largest positive-counting-number permitted by the implementation for the system parameter print-precision.

Note: Print-precision-limit must be at least **full-print-precision** if every unique **number** is to have a unique decimal representation.

 Exponent-Field-Width: A positive-counting-number giving the number of places, including sign and trailing blanks, used for the exponent field in the result of dyadic format.

- Session-Identification-Type: A member of the enumerated set array-type, either character or numeric.
- User-Identification-Type: A member of the enumerated set array-type, either character or numeric.
- Indent-Prompt: A list of characters used to indicate to the user that the session is in immediate-execution mode.
- Quad-Prompt: A list of characters used to indicate to the user that the session is in quad-input mode.
- Function-Definition-Prompt: A list of characters used to indicate to the user that the session is in function-definition mode.
- Line-Limit: A positive-counting-number that specifies the maximum number of lines permitted in a defined function, ignoring storage limitations.
- Definition-Line-Limit: A positive-number that specifies the maximum value of a line number in function-definition mode.
- General-Offer: A reserved session-identification used to indicate that the offerer of
 a shared-variable is willing to share the proffered variable with any other session.
 This implementation-parameter is required only if the optional-facility shared-variable-protocol is present.

Any action that would cause a limit specified by an implementation parameter to be exceeded shall signal a **limit-error**.

5.4 Evaluation Sequences

The evaluation sequences that define APL operations in the remainder of this International Standard are written in English in the imperative mood. The English phrases used in evaluation sequences are restricted to the set specified in this subsection.

Indention is used in evaluation sequences to indicate scope, typically of the consequent of an implication.

For example, in the evaluation sequence fragment below, the indented text is evaluated only if both A and B are **vectors**; the "otherwise" clause is evaluated only if at least one of A or B is not a **vector**.

```
...If A is a vector and B is a vector,If A is empty and ......Otherwise, return ...
```

Expressions in APL are used in evaluation sequences. A given evaluation sequence uses only APL operations that have been specified earlier in the International Standard. Where indices are generated or used by APL expressions in evaluation sequences, they are evaluated with origin **one**. The APL relational operations are never used in evaluation sequences unless they are qualified with the value to be used for **comparison-tolerance**.

5.4.1 Evaluation Sequence Phrases

Note: The following phrases are used in evaluation sequences. They are not set in **bold** type nor cross-referenced when employed in evaluation sequences.

- For all A, C: An evaluation sequence phrase used to specify that the action or condition specified by the consequent C is to be performed or checked for every value in the antecedent A.

```
Example: ... For all \mathcal I in the index-set of A, Set item \mathcal I of the ravel-list of A to zero. ...
```

- For form F, C:
- For pattern F, C:

An evaluation sequence phrase used to specify that the actions listed in the consequent C are to be performed only if the pattern or form F is the one that caused this evaluation sequence to be selected.

```
Example: For form A \phi B If B is a scalar ... ... For form A \Theta B If ...
```

- If A, C: An evaluation sequence phrase used to specify that the actions listed in the consequent C are to be performed only if the value of the antecedent A is one.
- If T is an mc, C: For T a token and mc a metaclass, an evaluation sequence phrase equivalent to "If the class of T is in the metaclass mc, C."
- Let A stand for B: An evaluation sequence phrase used to indicate that the name A is to be an abbreviation for the phrase B in subsequent evaluation sequence lines.
- Otherwise, C: An evaluation sequence phrase used to indicate that the actions listed in consequent C are to be performed only if the antecedent in the immediately preceding if phrase was zero. If a consequent is an indented paragraph, the immediately preceding if statement is the one at the same level of indention as the otherwise phrase.
- Repeat: An evaluation sequence phrase used to indicate that the block of text indented after the repeat is to be executed repeatedly until a return or signal phrase is encountered. The end of a repeated block is indicated by the parenthetic remark "(End of repeated block)."

- Return X: An evaluation sequence phrase used to specify that evaluation of this evaluation sequence is to stop and that a token is to be returned to the caller of the evaluation sequence. If X is a token, then X is returned; if X is an array, a token of class constant and content X is returned.
- Set A to B: An evaluation sequence phrase used to specify that the referent of A is to be assigned the value B.
- **Signal** *X*: An evaluation sequence phrase used to specify that evaluation of this evaluation sequence is to stop and that a **token** whose **class** is *X*, where *X* is an **error**, is to be returned.
- Using O, C: An evaluation sequence phrase used to indicate that the consequent C is to be evaluated against the specific object O. This construct is used, for example, in the description of shared variables to indicate which shared variable is to be changed.
- Wait until: An evaluation sequence phrase that indicates that the session is waiting for a condition to hold before continuing.

5.4.2 Diagrams

Note: Diagrams are used in this International Standard to indicate permissible sequences of characters or of tokens.

- Character-Diagram: A graph that designates a subset of the set of all lists of characters.
- Token-Diagram: A graph that designates a subset of the set of all lists of tokens.
- Thread D with A: An evaluation sequence phrase used to indicate that a search is to be made for a route through the diagram D that corresponds to the list or part of a list A. A character-diagram is threaded with a list of characters by setting a list-cursor to the first item in the list, and a diagram-cursor to the arrow-tail → , in the diagram, then progressing along paths in the diagram to the arrow-head → . At a junction, the diagram-cursor may enter the alternate path only if it can do so by making an acute-angle turn from its current direction. It may not back up into the alternate path.

For the diagram-cursor to advance along a path labelled with a graphic symbol such as \div , that graphic symbol must be pointed to by the list-cursor, which then also advances. For the diagram-cursor to advance along a path labelled with a diagram name, that diagram must itself be threaded, using the same list-cursor and a new diagram-cursor. If in either case the diagram-cursor cannot advance, the diagram-cursor is set back to the previous junction, the list-cursor is set back correspondingly, and a new path is tried; or if the diagram-cursor is now at the arrow-tail, the diagram cannot be threaded with the list.

If a route is found through a diagram for a given **list** of **characters**, it is unique. In this case, the diagram can be **rethreaded**, following the same route without entering blind alleys. For example, characters are collected into identifier tokens through actions performed while a **character-diagram** is being **rethreaded**.

Token-diagrams are **threaded** exactly like **character-diagrams**, except that the items pointed to by the list-cursor are **tokens** and are matched either by their **class** or by both their **class** and their **content**.

- Rethread D with A: An evaluation sequence phrase used to indicate that the route found
 by a previous threaded phrase is to be threaded again, in order to effect certain actions
 through when phrases.
- When A, C: An evaluation sequence phrase used during a **rethread** phrase to indicate that when the antecedent A is true, the consequent C is to be performed.
- A Matches D: An operation that, for diagram D and list A, is true if A is a member of the set of lists designated by D, and false otherwise.

A **list** matches a diagram if it can be **threaded** in such a way that there are no items remaining in the **list** when the final exit path is taken.

Note: The question of whether a **list matches** a diagram is different from the question of whether a **list** can **thread** a diagram, since there will normally be items left over in a **list** once a diagram has been **threaded**—in collecting the digits in a number, for example, the **digit** diagram removes only one digit at a time from a **list** of **characters**.

5.5 Other Terms

- **Side-effect**: Any effect an operation has other than returning a result.
- Atomic: The property of an operation with side-effects to produce its side-effects only
 if it completes without error.

Note: For example, the specification of items 3, 4, and 5 of A in the APL line

$$B \leftarrow A[3 + 5] \leftarrow 2$$

is a side effect, since the result placed in B is the scalar 2; further, since **indexed assignment** is specified to be **atomic**, no change to A will occur if the indexed assignment fails with, for example, an index error.

6 Syntax and Evaluation

6.1 Introduction

Note: This chapter specifies the rules for evaluating lines. These rules are used by the subsections execute, quad input, immediate-execution, and defined-function-control.

The rules are described in three subsections, evaluate-line, evaluate-statement, and reduce-statement.

The data structures and procedures used to describe syntax and evaluation in this International Standard are strictly expository devices; they are not intended to represent required or desirable implementation techniques. The order of reporting errors implied by the diagram-matching actions in the model is not required.

6.1.1 Evaluate-Line

Form: Evaluate-Line

Informal Description: Evaluate-line is the principal procedure in the evaluation of APL lines. It decomposes the **list** of **characters** named **current-line** into **lexical-units** according to the **character-diagram** named **line**.

Some diagrams referred to in **line** have their exit paths flagged with asterisks. Once the diagram **line** has been threaded, the rethreading pass is used to gather certain character sequences (diagrams ending in two asterisks), and either create tokens (diagrams ending in one asterisk) or discard the character sequences gathered (diagrams ending in three asterisks).

Evaluate-line calls **evaluate-statement** to convert the **lexical-units** into a **result**.

Note that the **result** can be a **constant** (from, for example, 2+2), **committed-value** (from $A\leftarrow 1$), **nil** (ϕ ' '), **branch** (ϕ ' $\rightarrow 3$ '), **escape** (\rightarrow), **unwind** (\rightarrow), **error** (2-), **clear-state-indicator** (from) SIC) or **interrupt** (from **signal-interrupt**).

Note: The handling of these result classes by the different callers of evaluate-line specifies the treatment of exceptions, and should be analysed carefully.

Evaluate-line is called from evaluation sequences in the subsections defined-functioncontrol, execute, execute-alternate, immediate-execution, and quad-input. If the optional-facility statement-separator-facility is implemented, evaluate-line evaluates successive statements in current-line beginning with the leftmost statement and continuing until evaluation produces an exception or the rightmost statement has been evaluated.

Evaluation Sequence:

Set C to the **empty-list** of **characters**.

Set current-statement to the empty-list of tokens.

Thread line with current-line; if current-line does not match line, signal syntaxerror.

Rethread **line** with **current-line**, taking the following actions:

When a **character-diagram** ending in '-***' is recognised, set C to the **empty-list** of characters.

When a **character-diagram** ending in '-**' is recognised, append to C as a new last item the character just passed.

When a **character-diagram** ending in '—*→' is recognised,

Append to current-statement as a new last item a token with class given by the name of the **character-diagram** and with **content** C.

Set C to the **empty-list** of **characters**.

When the optional-facility statement-separator-facility is implemented and the character-diagram statement-separator is recognised,

Set Z to evaluate-statement.

Ignore here, do in porser If Z is an exception, return Z.

If Z is a constant, display Z.

Set current-statement to the empty-list of tokens.

Set C to the **empty-list** of **characters**.

When the **character-diagram line** is matched,

Set Z to evaluate-statement.

Return Z.

6.1.2 **Character-Diagrams**

The diagrams in this subsection are **character-diagrams**: the APL Graphic Symbols such as ∇ in these character-diagrams match corresponding characters in the requiredcharacter-set.

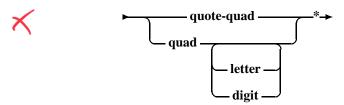
Line - identifier comment numeric-literal – primitive -- character-literal -- space --statement-separator -**Identifier** – simple-identifier – distinguished-identifier -Simple-identifier literal-identifier direct-identifier Literal-identifier 🗕 letter letter digit – underbar overbar -

Direct-identifier

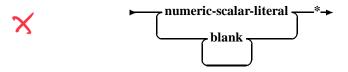
X

X

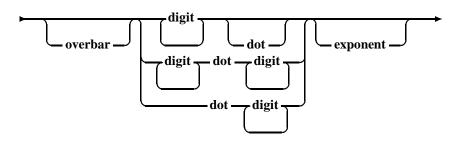
Distinguished-identifier



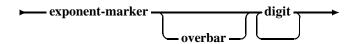
Numeric-Literal



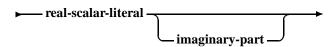
Real-scalar-literal



Exponent



Numeric-scalar-literal



Imaginary-part

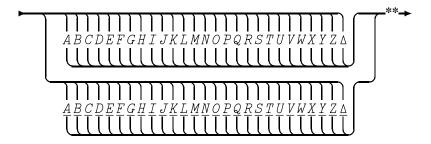
Character-literal × - quote quote nonquote Comment - lamp any Any quote nonquote -**Primitive** — ideogram —*→ **Space** — blank -***→ Nonquote ideogram digitblank -- letter -– lamp -– del – del-tilde -- quad quote-quad diamond

Statement-separator

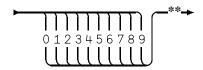
► diamond →

Letter

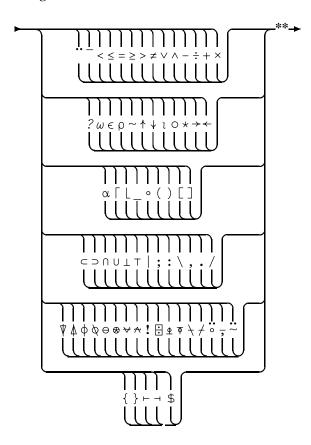
X



Digit



Ideogram



Quote

--- ' -**→

Exponent-marker

► E **-** →**

Complex-marker

__ *J* **_******→**

Dot

____ . _**-

Underbar
Overbar
─ ── [─] —** →
Blank
▶ —— —**→
Del
▶
Del-tilde
▶ ——₩—**- >
Lamp
▶─── A ──**─►
Quad
— □ —** →
Quote-quad
▶ —— □ —**→
Diamond
── ◇ ─ ** →
Example
The example introduced in this subsection is continued through the syntax analysis portion of thi International Standard.

After **evaluate-line** has processed the **current-line** $ABC \leftarrow FN \ \Box \phi \ [1+0] \ DEF \ [1;5 6] \times 3.45 E4, \rho' ABC' \cap COMMENT$

current-statement looks like this:

Identification	Content	Class
T01	' <i>ABC</i> '	simple-identifier
<i>T</i> 02	1 ←1	primitive
<i>T</i> 03	$^{\shortmid}FN$ $^{\backprime}$	simple-identifier
<i>T</i> 04	¹ 🗌 ¹	distinguished-identifier
<i>T</i> 05	'ф'	primitive
<i>T</i> 06	1 [1	primitive
<i>T</i> 07	111	numeric-literal
<i>T</i> 08	1 + 1	primitive
<i>T</i> 09	101	numeric-literal
T10	'] '	primitive
T11	$^{\dagger}DEF$ †	simple-identifier
T12	1 [1	primitive
T13	111	numeric-literal
T14	1 , 1	primitive
T15	15 61	numeric-literal
<i>T</i> 16	'] '	primitive
T17	1 × 1	primitive
<i>T</i> 18	'3.45 <i>E</i> 4'	numeric-literal
<i>T</i> 19	1,1	primitive
<i>T</i> 20	¹ p ¹	primitive
T21	'' <i>ABC</i> ''	character-literal

Each **token** has a **content**, which is a **list** of **characters**, and a **class**, which is in the metaclass **lexical-unit**. The **tokens** produced are numbered T01 to T21 for later reference. Note that comments and blanks between tokens are discarded during this tokenisation process.

Note: As the diagram **line** shows, identifiers and numeric literals are separated by one or more spaces, character-literals, or primitives. 1 ABC'A' is a line, 1ABC'A' is not. No such separation rule applies to primitive function symbols. $10^{\circ}AB'$ is a line.

The sequence 3+.4 is parsed numeric-literal, primitive, numeric-literal but the sequence 3+.× is parsed numeric-literal, primitive, primitive, primitive.

A comment may appear at the end of a line or alone on a line.

The statement-separator-facility is an optional-facility.

Parsing an identifier token signals a limit-error if the number-of-items in the list of characters is greater than identifier-length-limit.

6.1.3 Evaluate-Statement

Evaluate-Statement

Informal Description: Evaluate-statement is performed on current-statement, a list of

tokens found in the current-context.

It uses bind-token-class to convert identifiers to classified-names and constants.

It uses literal-conversion to convert literals to constants.

It uses the **token-diagram** named **statement** to verify that the statement is properly formed and to resolve certain ambiguous **tokens**. For example, the **token** containing] is resolved to either **right-axis-bracket** or **right-index-bracket**.

It calls the evaluation sequence in the subsection **reduce-statement** to convert the **current-statement** to a **result**, which it returns to its caller.

Evaluate-statement is called by evaluate-line.

Evaluation Sequence:

For every index I in the index-set of current-statement,

Let T stand for item I of current-statement.

If T is an **identifier**, set Q to **bind-token-class** of T

If T is a **literal**, set Q to the **literal-conversion** of T.

If T is a **primitive**, set Q to T.

If Q is an **exception**, return Q.

Otherwise, set T to Q.

Thread **statement** with **current-statement**; if **statement** cannot be matched, signal **syntax-error**.

Rethread **statement** with **current-statement**, taking the following action:

When any **token-diagram** ending in '—*→' is threaded,

Replace the **token** in **current-statement** that matched the diagram with a **token** having the same **content**, but having a **class** given by the name of the **token-diagram**.

Append to **current-statement** as a new first **item** a **left-end-of-statement token**.

Append to current-statement as a new last item a right-end-of-statement token.

Set Z to **reduce-statement**.

If mode is defined-function and current-line-number is in current-trace-vector, set Z to trace-display of Z.

Return Z.

Additional Requirement:

This International Standard permits two orders of evaluation for expressions in brackets, as follows:

The syntax evaluator used here distinguishes axis brackets from index brackets in order to classify a function immediately to the right as monadic or dyadic. For example, the evaluation of $\varphi[\rho X] + X$ is specified to proceed by first calling monadic plus, then evaluating monadic rho.

The permitted alternative is to evaluate any expression in brackets before determining whether a function immediately to the right is monadic or dyadic. In the example given, this alternative behaviour would evaluate monadic rho before discovering that plus is used

Threading statement ereate wrappe that porseses each statement delimited by P

monadically.

Note: The distinction can be observed only through side effects.

6.1.4 Bind-Token-Class

Bind-Token-Class of T

Informal Description: Bind-token-class is used to bind each **identifier** in **current-statement** to its current **syntactic-unit**; if the **class** of a **token** changes between this point and the time the **token** is used (as it would for example in $F \square EX^{!}F^{!}$, assuming F were initially a **defined-function**), the change will be detected and reported as a **syntax-error** in the appropriate **phrase-evaluator**.

This prebinding limits the International Standard to defining the meaning of statements only when the syntax class of their tokens does not change in mid-statement. **Conforming-implementations** may, of course, relax these rules. **Conforming-programs** must abide by them.

Evaluation Sequence:

Let **f** stand for the **content** of T.

If T is a simple-identifier,

If the **current-class** of T is

defined-monadic-operator, return a **token** of **class defined-monadic-operator-name** and **content f**.

defined-dyadic-operator, return a **token** of **class defined-dyadic-operator-name** and **content f**.

defined-function, return a **token** of **class defined-function-name** and **content f**. **niladic-defined-function**, return a **token** of **class niladic-defined-function-name** and **content f**.

nil or variable, return a token of class variable-name and content f.

shared-variable, return a token of class shared-variable-name and content f.

label, return a token of class constant and content the current-content of T.

Note: Other cases cannot occur.

If T is a **distinguished-identifier**.

If both forms $Z \leftarrow \mathbf{f}$ and $Z \leftarrow \mathbf{f} \leftarrow B$ occur in the **form-table**, return a **token** of **class system-variable-name** and **content** \mathbf{f} .

If the form $Z \leftarrow \mathbf{f}$ occurs in the **form-table**, but the form

 $Z \prec f \prec B$ does not, return a **token** of **class niladic-system-function-name** and **content f**.

If either form $Z \leftarrow f$ B or form $Z \leftarrow A f$ B occurs in the **form-table**, return a **token** of **class system-function-name** with **content f**.

Otherwise, signal syntax-error.

6.1.5 Literal-Conversion

Literal-Conversion of T.

Informal Description: Literal-conversion converts T, a token of class literal, to a token of class constant. The content of such a token is converted from a list of characters to an array.

Evaluation Sequence:

If T is a **character-literal**, generate Z, a **character vector** such that the **ravel-list** of Z is the **content** of T with the initial and final quotes removed, and each pair of adjacent quotes in T replaced by a single quote.

If T is a **numeric-literal**, generate Z, a **numeric vector** such that the **ravel-list** of Z is a **list** of **numbers**, each of which is obtained by calling the **implementation-algorithm numeric-input-conversion** for the corresponding **numeric-scalar-literal** in the **numeric-literal**.

If the **length** of Z is **greater-than one**, return Z.

Otherwise, return **first-scalar** in Z.

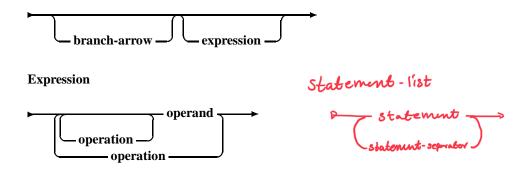
Note: A quote character is represented in a **character-literal** by two adjacent quote characters. A single character between quotes is a scalar. All other cases are **character** vector literals. For example, the **character** literal '' is the empty **character** vector literal and the character literal ''' is the **character** scalar "quote".

A numeric literal containing a single number is a scalar. A numeric literal containing two or more numbers is a vector.

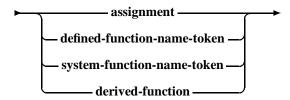
6.1.6 Statement-Analysis Token-Diagrams

Paths containing **ideograms** such as * in these **token-diagrams** match **tokens** whose **class** is **primitive** and whose **content** is the ideogram. Paths containing the word **token**, such as **shared-variable-name token** in **operand**, match **tokens** with the given **class**.

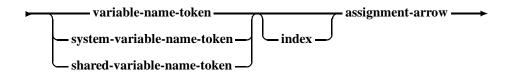
Statement



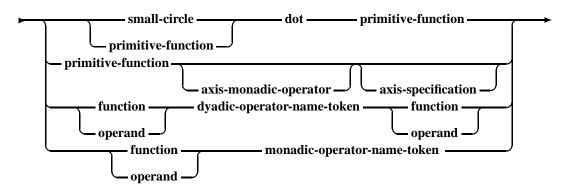
Operation



Assignment



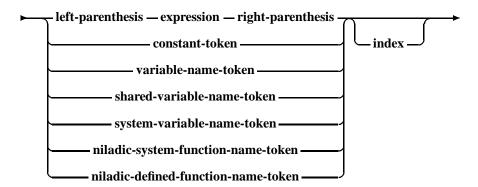
Derived-Function



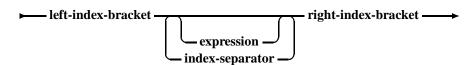
Axis-Specification

left-axis-bracket — expression — right-axis-bracket —

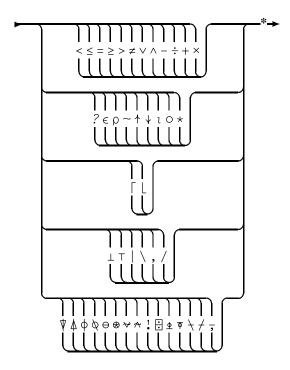
Operand



Index



Primitive-Function



Function primitive-function derived-function defined-function-name-token system-function-name-token-**Dyadic-Operator** — primitive-dyadic-operator — - defined-dyadic-operator-name-token – **Monadic-Operator** — primitive-monadic-operator — - defined-monadic-operator-name-token -Axis-Monadic-Operator **Primitive-Monadic-Operator** – axis-monadic-operator -- diaeresis-tilde -**Primitive-Dyadic-Operator**

─ diaeresis-jot —*→

Diaeresis-Jot

Diaeresis-Tilde	
<u>~</u> ~ —*→	
Left-Parenthesis — (—*→	
Right-Parenthesis	
▶) —* →	
Left-Axis-Bracket [*→	
Right-Axis-Bracket	
─ -*→	
Branch-Arrow	
─ → * →	
Assignment-Arrow	
Assignment-Arrow	
Assignment-Arrow ←—*→	
Assignment-Arrow	
Assignment-Arrow	
Assignment-Arrow	

Small-Circle



Example

From the **list** of **characters** in **current-line**,

 $ABC \leftarrow FN \ \Box \phi [1+0] \ DEF[1;5 6] \times 3.45E4, \rho'ABC' \cap COMMENT$

evaluate-line generated a **list** of **tokens** and stored it in **current-statement**. Here, **evaluate-statement** has replaced that **list** with a new **list** of **tokens**.

Identification	Content	Class
T00		left-end-of-statement
<i>T</i> 01	' <i>ABC</i> '	variable-name
<i>T</i> 02		assignment-arrow
<i>T</i> 03	$^{!}FN$ $^{!}$	defined-function-name
T04	' 🗌 '	system-variable-name
<i>T</i> 05	' ф '	primitive-function
<i>T</i> 06		left-axis-bracket
<i>T</i> 07	1	constant
<i>T</i> 08	' + '	primitive-function
<i>T</i> 09	0	constant
T10		right-axis-bracket
T11	'DEF'	variable-name
T12		left-index-bracket
T13	1	constant
T14		index-separator
T15	5 6	constant
T16		right-index-bracket
T17	1 × 1	primitive-function
T18	34500	constant
T19	1,1	primitive-function
<i>T</i> 20	'ρ'	primitive-function
T21	' <i>ABC</i> '	constant
T22		right-end-of-statement

The new list, shown above, begins with a left-end-of-statement token and ends with a right-end-of-statement token. Old identifier tokens have a new class and the same content; old literal tokens are now constants whose content is an appropriate array; old primitive tokens are now either primitive-functions whose content is the character identifying the primitive function, or are grouping signs such as right-axis-bracket with no specified content.

6.2 Reduce-Statement

Reduce-Statement

Informal Description: Reduce-statement converts **current-statement**, a **list** of **syntactic-units**, to a **result** by decomposing it into shorter lists of **syntactic-units** called **phrases**, then calling procedures termed **phrase-evaluators** to convert these phrases into **tokens**.

The letters Z, K, and J in this evaluation sequence refer to the graphic symbols found in the **resultant-prefix** column of **Table 3**.

Evaluation Sequence:

Set current-stack to the empty-list of tokens.

Repeat:

Find the first entry in the **phrase-table** whose **pattern** matches a **prefix** of **current-stack**.

If there is no matching entry,

If current-statement is empty, signal syntax-error.

Otherwise,

Remove the last token from current-statement.

Append it to **current-stack** as a new first **item**.

If there is a matching entry,

Let **p** stand for the **prefix** of the **current-stack** that matched the entry.

Let **r** stand for the **resultant-prefix** of the entry.

Let **s** stand for the **phrase-evaluator** of the entry.

Call s and set y to the token it returns.

If **s** is **process-end-of-statement**, return **y**.

If **y** is an **exception**, return **y**.

Otherwise, replace \mathbf{p} with \mathbf{r} in which \mathbf{y} has been substituted for Z, K or J according to whether \mathbf{y} is a **result**, a **complete-index-list** or a **partial-index-list**.

(End of repeated block)

The graphic symbols in the **Pattern** and **Resultant-Prefix** columns of **Table 3** designate lists of **syntactic-units**. Each graphic symbol matches **tokens** of the designated **classes** or **metaclasses**.

Note: The graphic symbols employed are chosen to be suggestive of the **list** of **characters** that give rise to such phrases.

The build-index-list entries are called with B bound to a value and I bound to a partial-index-list; they return either another partial-index-list (J) or a complete-index-list (K). Brackets are not passed as arguments or returned by build-index-list; they are preserved by reduce-statement to make the patterns more obvious.

Example

This example is continued from evaluate-statement.

The line being evaluated is

T00 is left-end-of-statement, and T22 is right-end-of-statement.

The **tokens** T00 through T22 initially form the columns of **Figure 1**. The rows of the figure show the actions taken because of the pattern or lack of pattern in **current-stack**.

Consider step 5: the line fragment , ρ ' ABC' matches the phrase X F B as follows: X matches token T19 because , is a **primitive-function**. F matches token T20 because ρ is a **primitive-function**. B matches token T21 because ABC is a **constant**. Token T22, **right-end-of-statement**, is not considered because the pattern X F B is concerned with only the first three tokens in **current-stack**.

The **phrase-evaluator** associated with X F B is **evaluate-monadic-function**. This **phrase-evaluator**, seeing that ρ is a primitive function, searches the **form-table** for $Z \leftarrow \rho B$, and calls the corresponding evaluation sequence, **shape**. Shape returns a **constant**, the one-element-vector containing **three**.

This becomes Z in the **resultant-prefix** column for X F B. Current-stack now holds three **tokens**: (**primitive-function**; ,), (**constant**; ,3) and (**right-end-of-statement**;). Since no entry in the **phrase-table** has an entry whose **pattern** matches a **prefix** of **current-stack**, the **token** T18 (**constant**; 34500) is added to **current-stack**. A **prefix** of **current-stack** now matches a pattern (A F B matches (**constant**; 34500), (**primitive-function**; ,), (**constant**; ,3)), so the corresponding **phrase-evaluator** (**evaluate-dyadic-function**) is called.

Pattern	Phrase Evaluator	Resultant-Prefix
(B)	Remove-Parentheses	Z
N	Evaluate-Niladic-Function	Z
X F B	Evaluate-Monadic-Function	X Z
X F [C] B	Evaluate-Monadic-Function	X Z
X F M B	Evaluate-Monadic-Operator	X Z
X F M [C] B	Evaluate-Monadic-Operator	X Z
A F M B	Evaluate-Monadic-Operator	X Z
A F M [C] B	Evaluate-Monadic-Operator	X Z
A F B	Evaluate-Dyadic-Function	Z
A F [C] B	Evaluate-Dyadic-Function	Z
X F D G B	Evaluate-Dyadic-Operator	Z
A F D G B	Evaluate-Dyadic-Operator	Z
$A \circ D G B$	Evaluate-Dyadic-Operator	Z
A [K]	Evaluate-Indexed-Reference	Z
V [K] ←B	Evaluate-Indexed-Assignment	Z
V ←B	Evaluate-Assignment	Z
V	Evaluate-Variable	Z
]	Build-Index-List	J]
; I	Build-Index-List	J
; B I	Build-Index-List	$\mathcal J$
Γ	Build-Index-List	[<i>K</i>
$[B\ I$	Build-Index-List	[<i>K</i>
L R	Process-End-of-Statement	_
L B R	Process-End-of-Statement	_
$L \rightarrow B R$	Process-End-of-Statement	_
L →R	Process-End-of-Statement	-

Legend:

A, B , Z match result .	X matches assignment-arrow, branch-arrow, defined-function-name, index-separator, left-axis-bracket, left-end-of-statement, left-index-bracket left-parenthesis, primitive-function,		
D matches dyadic-operator.			
F, G match defined-function-name, primitive-function, or system-function-name.			
${\cal I}$, ${\cal J}$ match partial-index-list .	system-function-name, or right-axis-bracket.		
C, K match complete-index-list.	(matches left-parenthesis.		
L matches left-end-of-statement. M matches monadic-operator. N matches niladic-defined-function-name or niladic-system-function-name. R matches right-end-of-statement. V matches variable-name, system-variable-name, or shared-variable-name.	 matches right-parenthesis. matches left-axis-bracket or left-index-bracket. matches right-axis-bracket or right-index-bracket. matches small-circle. matches index-separator. matches assignment-arrow. 		
			→ matches branch-arrow .

Table 3: The Phrase Table.

```
1 GET NEXT TOKEN T22
                                                                      ()
 2 GET NEXT TOKEN T21
                                                                     (R)
 3 GET NEXT TOKEN T20
                                                                   (B R)
 4 GET NEXT TOKEN T19
                                                                (FBR)
 5 EVALUATE MONADIC FUNCTION
                                                             (\langle X F B \rangle R)
 6 GET NEXT TOKEN T18
                                                             (\langle X Z \rangle R)
 7 EVALUATE DYADIC FUNCTION
                                                           ( < A F B >
                                                                      R)
 8 GET NEXT TOKEN T17
                                                            (Z
                                                                      R)
 9 GET NEXT TOKEN T16
                                                          (FB)
                                                                      R)
10 BUILD INDEX LIST
                                                      (<]>FB
                                                                      R)
11 GET NEXT TOKEN T15
                                                     (J ] F B
                                                                      R)
12 GET NEXT TOKEN T14
                                                   (B I ] F B
                                                                     R)
13 BUILD INDEX LIST
                                               (<; B I>] F B
                                                                     R)
14 GET NEXT TOKEN T13
                                                        \supset FB
                                                 (J
                                                                     R)
15 GET NEXT TOKEN T12
                                                           FB
                                              (B I
                                                                     R)
16 BUILD INDEX LIST
                                           (< [BI>
                                                       F B
                                                                     R)
17 GET NEXT TOKEN T11
                                           ( [ K
                                                       FB
                                                                     R)
                                                           FB
18 EVALUATE VARIABLE
                                                                     R)
                                        (\langle V \rangle [K])
                                                          FB
19 EVALUATE INDEXED REFERENCE
                                        ( < A [ K ]
                                                                      R)
20 EVALUATE DYADIC FUNCTION
                                                           F B>
                                                                     R)
                                         (< A
21 GET NEXT TOKEN T10
                                                                     R)
                                          (Z
                                      (<]>B
22 BUILD INDEX LIST
                                                                      R)
23 GET NEXT TOKEN TO9
                                     (I ] B
                                                                     R)
24 GET NEXT TOKEN TO8
                                   (B I ] B
                                                                     R)
25 GET NEXT TOKEN TO7
                                (F B I ] B
                                                                     R)
                             ( < A F B > I ] B
26 EVALUATE DYADIC
                                               FUNCTION
                                                                      R)
27 GET NEXT
                              (Z
                                     I \supset B
                                               TOKEN TO6
                                                                     R)
                          (< [B]
28 BUILD INDEX
                                     I > \exists B
                                               LIST
                                                                     R)
29 GET NEXT
                           ( [ C
                                       ] B
                                               TOKEN TO5
                                                                     R)
                         (F \ [ \ C
                                       ] B
30 GET NEXT
                                               TOKEN TO4
                                                                      R)
                      (< V > F \ [ \ C
31 EVALUATE
                                       ] B
                                               VARIABLE
                                                                     R)
                                       ] B>
32 EVALUATE
                      ( < A F [ C ]
                                               DYADIC FUNCTION
                                                                     R)
33 GET NEXT
                       (Z
                                               TOKEN TO3
                                                                     R)
                      (F B
                                 GET NEXT TOKEN TO2
34
                                                                     R)
35
                  (\langle X F B \rangle
                                 EVALUATE MONADIC FUNCTION
                                                                     R)
                                 GET NEXT TOKEN TO1
36
                    (X Z)
                                                                     R)
                (<V \leftarrow B>
                                 EVALUATE ASSIGNMENT
37
                                                                     R)
                  (Z
                                 GET NEXT TOKEN TOO
38
                                                                     R)
39 *
               (< L B
                                 END OF STATEMENT
                                                                     R>)
```

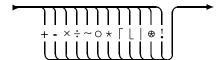
Figure 1: Statement Evaluation.

6.3 The Phrase Evaluators

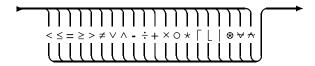
Informal Description: Each **phrase-evaluator** takes a **phrase** and reduces it to a single **token**. The **form-table** used by the **phrase-evaluators** is given as **Table 4**.

6.3.1 Diagrams

Primitive-Monadic-Scalar-Function



Primitive-Dyadic-Scalar-Function



6.3.2 Remove-Parentheses

Pattern (B)

Evaluation Sequence:

If *B* is **nil**, signal **value-error**. If *B* is a **branch**, signal **value-error**. Otherwise, return *B*.

Note: $({\scriptstyle \pm}{}^{\dagger}{}^{\dagger})$ and $({\scriptstyle \pm}{}^{\dagger}{}^{\dagger}{}^{\dagger})$ fail with a **value-error**. $(A \leftarrow 3)$ and $({\scriptstyle \pm}{}^{\dagger}{}^{\dagger}A \leftarrow 3{}^{\dagger})$ do not display, since the **token** returned by **remove-parentheses** is a **committed-value**, not a **constant**; for the same reason, $(\square \leftarrow 3)$ displays the value **three** only once. $(\rightarrow 3)$ fails with a **syntax-error** in **evaluate-statement**.

 (\Rightarrow) also fails in evaluate-statement, but $(\diamond ! \Rightarrow !)$ succeeds; evaluate-statement threads line successfully, and remove-parentheses receives and returns an escape token.

6.3.3 Evaluate-Niladic-Function

Pattern N

Evaluation Sequence:

Let \mathbf{n} stand for the **content** of N.

If N is a **niladic-defined-function-name**.

If the current-class of N is niladic-defined-function.

Search the **form-table** for $Z \leftarrow DFN$.

Call the corresponding evaluation sequence, passing **n** as the value of *DFN*.

Return the **token** it returns.

Otherwise, signal syntax-error.

If N is a **niladic-system-function-name**,

Search the **form-table** for $Z \leftarrow \mathbf{n}$.

If it is not found, signal syntax-error.

Otherwise, call the corresponding evaluation sequence.

Return the **token** it returns.

Note: This phrase evaluator checks for a change in syntax class.

6.3.4 Evaluate-Monadic-Function

Pattern X F B

Pattern X F[C] B

Evaluation Sequence:

If B is not a **value**, signal **value-error**.

Let \mathbf{f} stand for the **content** of F.

For pattern X F B

If F is a **defined-function-name**,

If the current-class of F is defined-function,

Search the **form-table** for $Z \leftarrow DFN B$.

Call the corresponding evaluation sequence, passing \mathbf{f} as the value of *DFN*.

Return the **token** it returns.

Otherwise, signal syntax-error.

If F is a **primitive-function** or a **system-function-name**,

If F matches **primitive-monadic-scalar-function** and B is not a **scalar**, perform **monadic-scalar-extension** as follows:

Return a **numeric array** Z such that the **shape-list** of Z is the **shape-list** of B and for all I in the **index-set** of the **ravel-list** of Z, **item** I of the **ravel-list** of Z is **f** (**item** I of the **ravel-list** of B).

Otherwise,

Search the **form-table** for $Z \leftarrow \mathbf{f} B$.

If it is not found, signal valence-error.

Otherwise, call the corresponding evaluation sequence.

Return the **token** it returns.

For pattern X F[C] B

If *F* is not a **primitive-function**, signal **syntax-error**.

If index-origin is nil, signal implicit-error.

Let C1 stand for the **first-item** in the **index-list** C.

If C1 is not a scalar or one-element-vector, signal axis-error.

If any **item** of the **ravel-list** of C1 is not a **number**, signal **axis-error**.

Set *K* to *C*1 plus (one minus index-origin).

Search the **form-table** for $Z \leftarrow \mathbf{f}[K]$ B.

If it is not found, signal syntax-error.

Otherwise, call the corresponding evaluation sequence and return the **token** it returns.

Additional Requirement:

The order in which the individual items of Z are produced during **monadic-scalar-extension** is not specified by this International Standard.

If any call of **f** signals an error during **monadic-scalar-extension**, that error is returned as the result of **evaluate-monadic-function**.

Note: This phrase evaluator checks for a change in syntax class.

6.3.5 Evaluate-Monadic-Operator

Pattern X F M B

Pattern A F M B

Pattern X F M[C] B

Pattern A F M[C] B

Evaluation Sequence:

If *B* is not a **value**, signal **value-error**.

Let **m** stand for the **content** of *M*.

Let f stand for the **content** of F.

For pattern $X \ F \ M \ B$

If M is a primitive-monadic-operator

Search the **form-table** for $Z \leftarrow \mathbf{f} \mathbf{m} B$

If it is not found, signal syntax-error.

Otherwise, call the corresponding evaluation sequence, passing token F as the value of \mathbf{f} .

Return the token it returns.

Otherwise,

Search the **form-table** for $Z \leftarrow \mathbf{f}$ *DFN B*.

If it is not found, signal syntax-error.

Otherwise, call the corresponding evaluation sequence, passing **token** F as the value of **f** and the **token** DFN as the value of **m**.

Return the **token** it returns.

For pattern A F M B

If A is not a **value**, signal **value-error**.

If m is a primitive-monadic-operator

Search the **form-table** for $Z \leftarrow A$ **f m** B.

If it is not found, signal syntax-error.

Call the corresponding evaluation sequence, passing **token** F as the value of \mathbf{f} .

Return the **token** it returns.

Otherwise.

Search the **form-table** for $Z \leftarrow A$ F DFN B.

If it is not found, signal **syntax-error**.

Call the corresponding evaluation sequence, passing **token** F as the value of **f** and the **token** DFN as the value of **m**.

Return the token it returns.

For pattern $X \in M[C]$ B or $A \in M[C]$ B

If M is not an axis-operator, signal syntax-error

If F is not a **primitive-function**, signal **syntax-error**.

If index-origin is nil, signal implicit-error.

Let C1 stand for the **first-item** in the **index-list** C.

If C1 is not a scalar or one-element-vector, signal axis-error.

If any **item** of the **ravel-list** of C1 is not a **number**, signal **axis-error**.

Set *K* to *C*1 **plus** (**one minus index-origin**).

For pattern $X \in M[C]$ B

Search the **form-table** for $Z \leftarrow \mathbf{f} \mathbf{m} [K] B$.

If it is not found, signal syntax-error.

Otherwise, call the corresponding evaluation sequence, passing **token** F as the value of \mathbf{f} .

Return the **token** it returns.

For pattern A F M[C] B

If A is not a value, signal value-error

Search the **form-table** for $Z \leftarrow A$ **f m** [K] B.

If it is not found, signal **syntax-error**.

Otherwise, call the corresponding evaluation sequence, passing **token** F as the value of \mathbf{f} .

Return the **token** it returns.

6.3.6 Evaluate-Dyadic-Function

Pattern A F B

Pattern A F[C] B

Evaluation Sequence:

If A is not a **value**, signal **value-error**.

If *B* is not a **value**, signal **value-error**.

Let **f** stand for the **content** of F.

For pattern A F B

If F is a **defined-function-name**,

If the **current-class** of F is **defined-function**,

Search the **form-table** for $Z \leftarrow A DFN B$.

Call the corresponding evaluation sequence, passing \mathbf{f} as the value of DFN.

Return the **token** it returns.

Otherwise, signal syntax-error.

If F is a **primitive-function** or a **system-function-name**,

If F matches **primitive-dyadic-scalar-function** and A and B are not both **scalars**, perform **dyadic-scalar-extension** as follows:

If the **rank** of A differs from the **rank** of B,

If A is a scalar or one-element-vector and B is not a scalar, return $((\rho B)\rho A)$ **f** B.

If B is a scalar or one-element-vector, return A f $(\rho A)\rho B$.

Otherwise, signal rank-error.

If the **shape-list** of *A* differs from the **shape-list** of *B*, signal **length-error**.

Return Z, an array such that the **shape-list** of Z is the same as the **shape-list** of A, the **type** of Z is **numeric**, and the **ravel-list** of Z is such that, for all Z in the **index-set** of the **ravel-list** of Z, **item** Z of the **ravel-list** of Z is (**item** Z of the **ravel-list** of Z).

Otherwise, search the **form-table** for $Z \leftarrow AfB$.

If it is not found, signal valence-error.

Call the corresponding evaluation sequence and return the **token** it returns.

For pattern A F[C] B

If *F* is not a **primitive-function**, signal **syntax-error**.

If **index-origin** is **nil**, signal **implicit-error**.

Let C1 stand for the **first-item** in the **index-list** C.

If C1 is not a **scalar** or **one-element-vector**, signal **axis-error**.

If any **item** of the **ravel-list** of C1 is not a **number**, signal **axis-error**.

Set *K* to *C*1 plus (one minus index-origin).

Search the **form-table** for $Z \leftarrow A \mathbf{f}[K] B$.

If it is not found, signal syntax-error.

Otherwise, call the corresponding evaluation sequence, passing **token** F as the value of \mathbf{f} .

Return the **token** it returns.

Additional Requirement:

There is an intentional forward reference to **shape** and **reshape** in the description of **dyadic-scalar-extension**.

The order in which the individual items of Z are produced during **dyadic-scalar-extension** is not specified by this International Standard.

If any call of **f** signals an error during **dyadic-scalar-extension**, that error is returned as the result of **evaluate-dyadic-function**.

Note: Dyadic-scalar-extension is intentionally stricter than it is in existing systems. For example, $(1 \ 101)+13$ signals a rank-error and $1 \ 2 + 10$, 1 signals a length-error.

This phrase evaluator checks for a change in syntax class.

An ambivalent function must be called either monadically or dyadically; hence, during syntax analysis and execution of a line, there is never a reference to an ambivalent function as any given instance is either monadic or dyadic.

6.3.7 Evaluate-Dyadic-Operator

Pattern X F D G B

Pattern A F D G B

Pattern $A \circ D G B$

Evaluation Sequence:

If B is not a **value**, signal **value-error**.

Let **d** stand for the **content** of *D*.

Let \mathbf{f} stand for the **content** of F.

Let **g** stand for the **content** of *G*.

For pattern X F D G B

If D is a **primitive-dyadic-operator**

Search the **form-table** for $Z \leftarrow \mathbf{f} \mathbf{d} \mathbf{g} B$.

If it is not found, signal syntax-error.

If it is found, call the corresponding evaluation sequence, passing **token f** as the value of F and **token g** as the value of G.

Return the token it returns.

Otherwise.

Search the **form-table** for $Z \leftarrow \mathbf{f} DFN \mathbf{g} B$.

If it is not found, signal **syntax-error**.

If it is found, call the corresponding evaluation sequence, passing **token f** as the value of F, **token g** as the value of G, and **token** DFN as the value of \mathbf{d} .

Return the **token** it returns.

For pattern A F D G B

If A is not a **value**, signal **value-error**.

If d is a primitive-dyadic-operator,

Search the **form-table** for $Z \leftarrow A \mathbf{f} \mathbf{d} \mathbf{g} B$.

If it is not found, signal syntax-error.

If it is found, call the corresponding evaluation sequence, passing **token f** as the value of F and **token g** as the value of G.

Return the **token** it returns.

Otherwise,

Search the **form-table** for $Z \leftarrow A \mathbf{f} DFN \mathbf{g} B$.

If it is not found, signal syntax-error.

If it is found, call the corresponding evaluation sequence, passing **token f** as the value of F and **token g** as the value of F, and **token** F as the value of F.

Return the **token** it returns.

For pattern $A \circ D G B$

If *D* is not **dot**, signal **syntax-error**.

Otherwise,

Search the **form-table** for $Z \leftarrow A \circ \mathbf{d} \mathbf{g} B$.

If it is not found, signal syntax-error.

If it is found, call the corresponding evaluation sequence, passing **token g** as the value of G.

Return the **token** it returns.

6.3.8 Evaluate-Indexed-Reference

Pattern $A \lceil K \rceil$

Evaluation Sequence:

If A is not a **value**, signal **value-error**.

If **index-origin** is **nil**, signal **implicit-error**.

If the **number-of-items** in the **index-list** K differs from the **rank** of A, signal **rank-error**.

If the rank of A is greater-than one,

Search the **form-table** for $Z \leftarrow A[I]$.

Call the corresponding evaluation sequence, passing K as the value of I.

Return the token it returns.

Otherwise,

If **first-item** in *K* is an **elided-index-marker**, return *A*.

Otherwise,

Set *X* to **first-item** in the **index-list** *K*.

If any **item** of the **ravel-list** of *X* is not a **near-integer**, signal **domain-error**.

Generate X1, a **numeric array** with the **shape-list** of X such that each **item** of the **ravel-list** of X1 is (**one minus index-origin**) **plus** the **integer-nearest-to** X. If any **item** of the **ravel-list** of X1 is not in the **index-set** of A, signal **index-error**. Return Z, an array with the **shape-list** of X1, such that for each **integer** I in the **index-set** of I, **item** I of the **ravel-list** of I is **item** I of the **ravel-list** of I. The **type** of I is the **sufficient-type** of the **ravel-list** of I under the **type** of I.

Note: Since an **index-list** will never have zero **items**, indexing will always signal a **rank-error** when argument A is a **scalar**.

6.3.9 Evaluate-Assignment

Pattern $V \leftarrow B$

Evaluation Sequence:

If *B* is not a **value**, signal **value-error**.

If V is a shared-variable-name,

If the **current-class** of *V* is **shared-variable**,

Search the **form-table** for $Z \leftarrow SHV \leftarrow B$.

Call the corresponding evaluation sequence, passing **token** *V* as the value of *SHV*.

Return the **token** it returns.

Otherwise, signal syntax-error.

If V is a system-variable-name,

Search the **form-table** for $Z \leftarrow \mathbf{q} \leftarrow B$, where \mathbf{q} is the **content** of V.

If it is not found, signal syntax-error.

Otherwise,

Call the corresponding evaluation sequence

Return the token it returns.

If *V* is a **variable-name**.

If the **current-class** of *V* is **nil** or **variable**,

Set the current-referent of V to a token whose class is variable and whose content is the content of B.

Return a **token** whose **class** is **committed-value** and whose **content** is *B*.

Otherwise, signal **syntax-error**.

Note: The phrase $ABC \leftarrow 2 \rightarrow 3$ yields value-error. The phrase $V \leftarrow \square SVR$ 'V' where V was a shared-variable yields syntax-error.

6.3.10 Evaluate-Indexed-Assignment

Pattern $V[K] \leftarrow B$

Evaluation Sequence:

If index-origin is nil, signal implicit-error.

If *B* is not a **value**, signal **value-error**.

If V is a shared-variable-name,

If the **current-class** of *V* is **shared-variable**,

Search the **form-table** for $Z \leftarrow SHV[I] \leftarrow B$.

Call the corresponding evaluation sequence, passing **token** V as the value of SHV, and K as the value of I.

Return the **token** it returns.

Otherwise, signal syntax-error.

If V is a **system-variable-name**,

Search the **form-table** for $Z \leftarrow \mathbf{q}[I] \leftarrow B$, where \mathbf{q} is the **content** of V.

If it is not found, signal syntax-error.

Otherwise,

Call the corresponding evaluation sequence, passing K as the value of I.

Return the **token** it returns.

If V is a variable-name,

If the **current-class** of *V* is **nil**, signal **value-error**.

If the **current-class** of *V* is a **variable**,

Search the **form-table** for $Z \leftarrow V[I] \leftarrow B$.

Call the corresponding evaluation sequence, passing K as the value of I.

Return the **token** it returns.

Otherwise, signal syntax-error.

6.3.11 Evaluate-Variable

Pattern V

Evaluation Sequence:

If V is a shared-variable-name,

If the **current-class** of *V* is **shared-variable**,

Search the **form-table** for $Z \leftarrow SHV$.

Call the corresponding evaluation sequence, passing **token** *V* as the value of *SHV*.

Return the token it returns.

Otherwise, signal syntax-error.

If V is a **system-variable-name**,

If the **current-class** of V is **nil**, signal **value-error**.

Search the **form-table** for $Z \leftarrow \mathbf{q}$, where \mathbf{q} is the **content** of V.

Call the corresponding evaluation sequence.

Return the token it returns.

If V is a variable-name,

If the **current-class** of *V* is **nil**, signal **value-error**.

If the **current-class** of *V* is **variable**, return the **current-content** of *V*.

Otherwise, signal syntax-error.

6.3.12 Build-Index-List

Pattern]

Pattern: I

Pattern; BI

Pattern [I

Pattern [B I

Evaluation Sequence:

For pattern]

Return J, a partial-index-list with content the index-list of length zero.

For pattern: I

Return J, a **partial-index-list** with **content** Z, an **index-list** such that the **first-item** in Z is an **elided-index-marker** and the **rest-of** Z is I.

For pattern; B I

Return J, a **partial-index-list** with **content** Z, an **index-list** such that the **first-item** in Z is B and the **rest-of** Z is I.

For pattern [I]

Return J, a **complete-index-list** with **content** Z, an **index-list** such that the **first-item** in Z is an **elided-index-marker** and the **rest-of** Z is J.

For pattern $[B \ I]$

Return J, a **complete-index-list** with **content** Z, an **index-list** such that the **first-item** in Z is B and the **rest-of** Z is I.

6.3.13 Process-End-of-Statement

Pattern L. R

Pattern L A R

Pattern $L \rightarrow R$

Pattern $L \rightarrow A R$

Evaluation Sequence:

For pattern L R

Return a token whose class is nil.

For pattern L A R

Return A.

For pattern $L \rightarrow R$ Process-End-of-Statement

Return a **token** whose **class** is **escape**.

Process-End-of-Statement

For pattern $L \rightarrow A R$

If the rank of A is greater-than one, signal rank-error.

If A is **empty**, return a **token** whose **class** is **nil**.

Otherwise, set A1 to the **first-scalar** in A.

If A1 is not a **near-integer**, signal **domain-error**.

Return a **token** whose **class** is **branch** and whose **content** is the **numeric-scalar** with value the **integer-nearest-to** A1.

6.4 The Form Table

The **form-table** is the **list** of all **lists** of **syntactic-units** for which evaluation sequences exist.

The following matching rules apply in the form-table.

- A, B, Z match constant.
- *I*, *K* match **complete-index-list**.
- \mathbf{f}, \mathbf{g} match $\mathbf{primitive}$ -function or $\mathbf{defined}$ -function, but not $\mathbf{defined}$ -operator.
- A given ideogram, such as *, matches a **primitive-function token** that contains it.
- A given **distinguished-identifier**, such as $\square IO$, matches any **system-variable-name token** or **system-function-name token** that contains it.

The behaviour of operations in the **form-table** that do not create new **contexts** is **atomic**.

This behaviour is observable only for those operations that have **side-effects**. For example, if any of the elements of an argument array is not in the domain of **roll**, the value of the system parameter **random-link** following execution will be as it was when **roll** was called.

Form	Operation Name	Page
$Z \leftarrow + B$	Conjugate	76
$Z \leftarrow -B$	Negative	76
$Z \leftarrow \times B$	Direction	77
$Z \leftarrow \div B$	Reciprocal	77
$Z \leftarrow L B$	Floor	78
$Z \leftarrow \Gamma B$	Ceiling	78
$Z \leftarrow * B$	Exponential	79
$Z \leftarrow \otimes B$	Natural Logarithm	79
$Z \leftarrow \mid B$	Magnitude	80
$Z \leftarrow ! B$	Factorial	81
$Z \leftarrow \circ B$	Pi times	82
$Z \leftarrow \sim B$	Not	83
$Z \leftarrow A + B$	Plus	84
$Z \leftarrow A - B$	Minus	84
$Z \leftarrow A \times B$	Times	85
$Z \leftarrow A \div B$	Divide	85
$Z \leftarrow A \mid B$	Maximum	86
$Z \leftarrow A \perp B$	Minimum	86
$Z \leftarrow A \times B$	Power	87
$Z \leftarrow A \otimes B$	Logarithm	88
$Z \leftarrow A \mid B$	Residue	89
$Z \leftarrow A : B$	Binomial	90
$Z \leftarrow A \circ B$	Circular Functions	91
$Z \leftarrow A \wedge B$	And/LCM	93
$Z \leftarrow A \vee B$	Or/GCD	94
$Z \leftarrow A \nsim B$	Nand	94
$Z \leftarrow A \ \forall \ B$	Nor	95
$Z \leftarrow A = B$	Equal	96
$Z \leftarrow A < B$	Less than	97
$Z \leftarrow A \leq B$	Less than or equal to	98
$Z \leftarrow A \neq B$	Not equal	99
$Z \leftarrow A \geq B$	Greater than or equal to	100
$Z \leftarrow A > B$	Greater than	101
$Z \leftarrow$, B	Ravel	102
$Z \leftarrow \rho B$	Shape	103
$Z \leftarrow \iota B$	Index Generator	104
$Z \leftarrow G$	Table	105
$Z \leftarrow A \rho B$	Reshape	107
$Z \leftarrow A$, B	Join	109
$Z \leftarrow A + B$	Join	109

Table 4: The Form Table

Form	Operation Name	Page
$Z \leftarrow \mathbf{f}/B$	Reduction	110
$Z \leftarrow \mathbf{f}/[K] B$	Reduction	110
$Z \leftarrow \mathbf{f} \neq B$	Reduction	110
$Z \leftarrow \mathbf{f} \neq [K] B$	Reduction	110
$Z \leftarrow \mathbf{f} \setminus B$	Scan	113
$Z \leftarrow \mathbf{f} \setminus [K] B$	Scan	113
$Z \leftarrow \mathbf{f} + B$	Scan	113
$Z \leftarrow \mathbf{f} \setminus [K] B$	Scan	113
$Z \leftarrow N \mathbf{f} / B$	N-wise Reduction	115
$Z \leftarrow N \mathbf{f} / [K] B$	N-wise Reduction	115
$Z \leftarrow N \mathbf{f} \neq B$	N-wise Reduction	115
$Z \leftarrow N \cdot \mathbf{f} \neq [K] B$	N-wise Reduction	115
$Z \leftarrow \mathbf{f} \stackrel{\sim}{\sim} B$	Duplicate	118
$Z \leftarrow A \mathbf{f} \stackrel{\dots}{\sim} B$	Commute	118
$Z \leftarrow A \circ \cdot \mathbf{f} B$	Outer Product	120
$Z \leftarrow A \mathbf{f} \cdot \mathbf{g} B$	Inner Product	121
$Z \leftarrow \mathbf{f} \stackrel{\circ}{\circ} \mathbf{y} B$ $Z \leftarrow A \mathbf{f} \stackrel{\circ}{\circ} \mathbf{y} B$	Rank	124
	Rank	125
$Z \leftarrow ? B$	Roll	127
$Z \leftarrow A B$	Grade Up	129
$Z \leftarrow V B$	Grade Down	131
$Z \leftarrow \phi B$	Reverse	132
$Z \leftarrow \Theta B$	Reverse	132
$Z \leftarrow \phi[K] B$	Reverse	132
$Z \leftarrow \Theta[K] B$	Reverse	132
$Z \leftarrow \Diamond B$	Monadic Transpose	133
$Z \leftarrow \boxminus B$	Matrix Inverse	134
$Z \leftarrow \Phi B$	Execute	135
$Z \leftarrow \cup B$	Unique	136
$Z \leftarrow A$, $[K]$ B	Join Along an Axis	137
$Z \leftarrow A i B$	Index of	140
$Z \leftarrow A \in B$	Member of	141
$Z \leftarrow A ? B$	Deal	142
$Z \leftarrow A / B$	Replicate	143
$Z \leftarrow A \neq B$	Replicate	143
$Z \leftarrow A / [K] B$	Replicate	143
$Z \leftarrow A \neq [K] B$	Replicate	143
$Z \leftarrow A \setminus B$	Expand	145
$Z \leftarrow A + B$	Expand	145
$Z \leftarrow A \setminus [K] B$	Expand	145
$Z \leftarrow A + [K] B$	Expand	145

Table 4: (Continued)

Form	Operation Name	Page
$Z \leftarrow A \phi B$	Rotate	147
$Z \leftarrow A \ominus B$	Rotate	147
$Z \leftarrow A \phi [K] B$	Rotate	147
$Z \leftarrow A \ominus [K] B$	Rotate	147
$Z \leftarrow A \perp B$	Base Value	149
$Z \leftarrow A + B$	Representation	151
$Z \leftarrow A \triangleleft B$	Dyadic Transpose	153
$Z \leftarrow A \uparrow B$	Take	155
$Z \leftarrow A \downarrow B$	Drop	156
$Z \leftarrow A \oplus B$	Matrix Divide	157
$Z \leftarrow A[I]$	Indexed Reference	158
$Z \leftarrow V[I] \leftarrow B$	Indexed Assignment	159
$Z \leftarrow A \sim B$	Without	161
$Z \leftarrow A \rightarrow B$	Left	161
$Z \leftarrow A \vdash B$	Right	162
$Z \leftarrow A \ \forall B$	Character Grade Down	163
$Z \leftarrow A \land B$	Character Grade Up	164
$Z \leftarrow \square TS$	Time Stamp	171
$Z \leftarrow \Box AV$	Atomic Vector	172
$Z \leftarrow \square LC$	Line Counter	172
$Z \leftarrow \square EM$	Event Message	173
$Z \leftarrow \Box ET$	Event Type	174
$Z \leftarrow \square DL B$	Delay	174
$Z \leftarrow \square NC B$	Name Class	175
$Z \leftarrow \square EX B$	Expunge	176
$Z \leftarrow \square NL B$	Name List	177
$Z \leftarrow \square STOP B$	Query Stop	178
$Z \leftarrow \Box TRACE B$	Query Trace	179
$\square ES$ B	Monadic Event Simulation	180
$Z \leftarrow A \square NL B$	Name List	180
$Z \leftarrow A \square STOP B$	Set Stop	181
$Z \leftarrow A \square TRACE B$	Set Trace	182
$Z \leftarrow A \square EA B$	Execute Alternate	183
$A \square ES B$	Dyadic Event Simulation	184
$Z \leftarrow A \square TF B$	Transfer Form	185
$Z \leftarrow \square CT \leftarrow B$	Comparison Tolerance	187
$Z \leftarrow \square CT$	Comparison Tolerance	187
$Z \leftarrow \square RL \leftarrow B$	Random Link	188
$Z \leftarrow \square RL$	Random Link	188
$Z \leftarrow \square PP \leftarrow B$	Print Precision	189
$Z \leftarrow \square PP$	Print Precision	189

Table 4: (Continued)

Form	Operation Name	Page
$Z \leftarrow \Box IO \leftarrow B$	Index Origin	190
$Z \leftarrow \square IO$	Index Origin	190
$Z \leftarrow \square LX \leftarrow B$	Latent Expression	191
$Z \leftarrow \square LX$	Latent Expression	191
$Z \leftarrow \Box LX[I] \leftarrow B$	Latent Expression	191
$Z \leftarrow DFN$	Call-Defined-Function	200
$Z \leftarrow DFN B$	Call-Defined-Function	200
$Z \leftarrow A DFN B$	Call-Defined-Function	200
$Z \leftarrow \mathbf{f} DFN B$	Call-Defined-Function	200
$Z \leftarrow A \mathbf{f} DFN B$	Call-Defined-Function	200
$Z \leftarrow \mathbf{f} DFN \mathbf{g} B$	Call-Defined-Function	200
$Z \leftarrow A \mathbf{f} DFN \mathbf{g} B$	Call-Defined-Function	200
$Z \leftarrow \Box FX B$	Function Fix	203
$Z \leftarrow \square CR B$	Character Representation	204
$Z \leftarrow SHV$	Shared Variable Reference	217
$Z \leftarrow SHV \leftarrow B$	Shared Variable Assignment	218
$Z \leftarrow SHV[I] \leftarrow B$	Shared Variable Indexed Assignment	219
$Z \leftarrow \square SVC B$	Shared Variable Access Control Inquiry	219
$Z \leftarrow \square SVQ B$	Shared Variable Query	221
$Z \leftarrow \square SVO B$	Shared Variable Degree of Coupling	222
$Z \leftarrow A \square SVO B$	Shared Variable Offer	223
$Z \leftarrow \square SVR B$	Shared Variable Retraction	224
$Z \leftarrow A \square SVC B$	Shared Variable Access Control Set	225
$Z \leftarrow \square SVS B$	Shared Variable State Inquiry	226
$\square SVE \leftarrow B$	Shared Variable Event	227
$Z \leftarrow \square SVE$	Shared Variable Event	227
$Z \leftarrow \Phi B$	Monadic Format	233
$Z \leftarrow A \Phi B$	Dyadic Format	237
$Z \leftarrow \square$	Quad Input	244
$Z \leftarrow \square$	Quote Quad Input	245
$Z \leftarrow \square \leftarrow B$	Quad Output	245
$Z \leftarrow \square \leftarrow B$	Quote Quad Output	246

Table 4: (Concluded)

7 Scalar Functions

Note: The primitive functions described in this chapter are called **scalar-functions**. All **scalar-functions** have uniform behaviour with respect to the structure of their argument arrays. The shape of the result of a **scalar-function** is determined solely by the shapes of its arguments.

This section defines scalar-functions individually for scalar arguments. Their common behaviour is described in this informal description by the expository device of an implicit operator, called the scalar-extension-operator.

If the argument to a monadic scalar function is not a scalar, a monadic scalar extension operator can be thought of as being invoked to produce a derived function which, in turn, applies the monadic scalar function to every element of the argument array, producing a result array of the same shape as the argument. The order in which the elements of the argument array are presented to the monadic scalar function is not specified by this International Standard. Monadic scalar functions never signal rank-error or length-error.

If either of the arguments of a dyadic scalar function is not a scalar, a dyadic scalar extension operator can be thought of as being invoked to produce a derived function, which provides pairs of scalars to the scalar function as follows:

The dyadic scalar extension operator first tests whether the two argument arrays have the same shape. Arguments to a dyadic scalar function must have the same shape. If they do not, and the argument of lesser rank is a scalar or one-element vector, the argument of lesser rank is reshaped to the shape of the argument of greater rank.

If the arguments cannot be made to have the same shape, the dyadic scalar extension operator signals a rank-error if the arguments are of different ranks and a length-error otherwise.

When the dyadic scalar extension operator succeeds, it produces a derived function which generates a scalar for each position in its result **array** by applying the subject **scalar-function** to pairs of scalars selected from corresponding positions in the argument arrays. The order in which the elements of the result **array** are produced is not specified by this International Standard.

Because the derived function produced by either scalar extension operator never calls its scalarfunction argument for empty arrays, domain-error can never be signalled for empty array arguments or for arrays reshaped by scalar extension to empty.

The type of all empty results produced by the functions derived from monadic and dyadic scalar extension is a property of the function to scalar extension. Since all scalar functions specified in this International Standard produce **numeric** results, the type of all empty results produced by scalar extension is specified as **numeric**.

For example, 11 \sqrt{5} and -11 return empty **numeric** results rather than signalling an error.

7.1 Monadic Scalar Functions

Note: The definitions in this section cover only scalar arguments. The phrase-evaluator evaluate-monadic-function handles non-scalar cases. All scalar-functions yield scalar results when applied to scalar arguments.

Note that, in this International Standard, roll is not a scalar-function.

7.1.1 Conjugate

$$Z \leftarrow + B$$

Informal Description: Z is the conjugate of B. Geometrically, it is the reflection of the number about the real axis.

Evaluation Sequence:

If B is not a **number**, signal **domain-error**. Return B with its **imaginary-part** negated.

Example:

7.1.2 Negative

$$Z \leftarrow -B$$

Informal Description: Z is the **negation** of B.

Evaluation Sequence:

If B is not a **number**, signal **domain-error**. Return **zero minus** B.

7.1.3 Direction

$$Z \leftarrow \times B$$

Informal Description: Z is **zero** if B is **zero**, and otherwise is the **number** with **magnitude one** that has the same **direction** as B. Geometrically, it is the number determined by the radial projection of B onto the **unit-circle**. For real numbers, this is identical with the **sign** or **signum** function, and takes on only the values **one**, **zero**, and **negative one**.

Evaluation Sequence:

If B is not a **number**, signal **domain-error**. If B is **zero**, return **zero**; otherwise, return B **divided-by** the **magnitude** of B.

Examples:

7.1.4 Reciprocal

$$Z \leftarrow \div B$$

Informal Description: Z is $1 \div B$.

Evaluation Sequence:

If B is not a **number**, signal **domain-error**. If B is **zero**, signal **domain-error**. Return **one divided-by** B.

Examples:

domain-error

7.1.5 Floor

$$Z \leftarrow \mid B$$

Informal Description: Z is the complex-integer U associated with the unit-square containing B, unless the sum of the fractional-parts of the real-part and imaginary-part of B is greater-than-or-equal-to one, in which case it is the nearer of U plus one or U plus imaginary-one, In case of a tie, it is U plus one. For real-numbers, Z is the greatest integer tolerantly less than or equal to B. Uses comparison-tolerance

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

If B is not a **number**, signal **domain-error**.

Return the **tolerant-floor** of *B* within comparison-tolerance.

Examples:

- In the following, **comparison-tolerance** is $1E^-10$.

Note: The following article describes the design of the floor function for complex arguments: McDonnell, E. E., "Complex Floor", APL Congress 73, North Holland Publishing Co., 1973

7.1.6 Ceiling

$$Z \leftarrow \Gamma B$$

Informal Description: Z is the **negation** of the **floor** of the **negation** of B. For B a **real-number**, Z is the least **integer** tolerantly greater than or equal to B. Uses **comparison-tolerance**.

Evaluation Sequence:

If **comparison-tolerance** is **nil**, signal **implicit-error**. If B is not a **number**, signal **domain-error**. Return - L - B.

Example:

- In the following, **comparison-tolerance** is $1E^{-}10$.

```
Γ<sup>-</sup>3.1416 3.1416 5.00000000001
```

7.1.7 Exponential

$$Z \leftarrow * B$$

Informal Description: Z is e raised to the power B, where e is the base of the natural logarithms.

Evaluation Sequence:

If B is not a **number**, signal **domain-error**. Return the **exponential** of B.

Examples:

```
* 1E50 2 1 0 1 2
0 0.135335 0.367879 1 2.71828 7.38906
* .693147
```

7.1.8 Natural Logarithm

$$Z \leftarrow \otimes B$$

Informal Description: Z is the natural logarithm of B.

Evaluation Sequence:

If B is not a **number**, signal **domain-error**. If B is **zero**, signal **domain-error**. Return the **natural-logarithm** of B.

```
$1828459045 2 1E<sup>-</sup>50 1E50
1 0.693147 -115.129 115.129

**1
```

7.1.9 Magnitude

$$Z \leftarrow \mid B$$

Informal Description: Z is the magnitude of B.

Evaluation Sequence:

If B is not a **number**, signal **domain-error**. Return the **magnitude** of B.

7.1.10 Factorial

 $Z \leftarrow ! B$

Informal Description: Z is the **gamma-function** of B+1. If B is a **nonnegative-integer**, this is factorial B.

Evaluation Sequence:

If B is not a **number**, signal **domain-error**. If B is a **negative-integer**, signal **domain-error**. Set B1 to B **plus one**. Return **gamma-function** of B1.

Examples:

Note: The gamma-function is defined in, for example, the National Bureau of Standards Handbook of Mathematical Functions, U.S. Government Printing Office, Washington D.C., 1964.

See also Hart, J. F., Computer Approximations, Robert C. Krieger Publishing Company, Huntington, NY, 1978.

7.1.11 Pi times

 $Z \leftarrow \circ B$

Informal Description: Z is π times B.

Evaluation Sequence:

If B is not a **number**, signal **domain-error**. Return **pi-times** B.

Example:

01 10 100 3.14159 31.4159 314.159

7.1.12 Not

$$Z \leftarrow \sim B$$

Informal Description: Z is the Boolean complement of B.

Evaluation Sequence:

If B is not **near-Boolean**, signal **domain-error**. If the **integer-nearest-to** B is **one**, return **zero**. Otherwise, return **one**.

Example:

- For the following, the **implementation-parameter integer-tolerance** is $1E^-10$.

$$\sim$$
0 1 1 E^- 11 .999999999999 1 0 1 0

7.2 Dyadic Scalar Functions

Note: The definitions in this section cover only scalar arguments. The phrase-evaluator evaluate-dyadic-function handles non-scalar cases. All scalar-functions yield scalar results when applied to scalar arguments.

The outer product operator, which has not yet been formally introduced at this point in the document, is used in the examples in this section as a convenient way of generating tables. The use of outer product in this section is limited to vector arguments. The same results could be obtained from each example, although not so compactly, by supplying the elements of the left argument one at a time, starting from the leftmost, as left arguments to the scalar function.

For example,

Example:

is equivalent to

7.2.1 Plus

$$Z \leftarrow A + B$$

Informal Description: Z is A plus B.

Evaluation Sequence:

If either of A or B is not a **number**, signal **domain-error**. Return A **plus** B.

Example:

7.2.2 Minus

$$Z \leftarrow A - B$$

Informal Description: Z is A minus B.

Evaluation Sequence:

If either of A or B is not a **number**, signal **domain-error**. Return A **minus** B.

7.2.3 Times

$$Z \leftarrow A \times B$$

Informal Description: Z is A times B.

Evaluation Sequence:

If either of A or B is not a **number**, signal **domain-error**. Return A **times** B.

Example:

7.2.4 Divide

$$Z \leftarrow A \div B$$

Informal Description: Z is A divided by B.

Evaluation Sequence:

If either of A or B is not a **number**, signal **domain-error**. If B is **zero** and A is not **zero**, signal **domain-error**. If B is **zero** and A is **zero**, return **one**. Otherwise, return A **divided-by** B.

7.2.5 Maximum

$$Z \leftarrow A \Gamma B$$

Informal Description: Z is the larger of A and B.

Evaluation Sequence:

If either of A or B is not a **near-real number**, signal **domain-error**. Set A1 to the **real-number** nearest to A. Set B1 to the **real-number** nearest to B. If A1 is **greater-than** B1, return A1. Otherwise, return B1.

Example:

7.2.6 Minimum

$$Z \leftarrow A \perp B$$

Informal Description: Z is the smaller of A and B.

Evaluation Sequence:

If either of A or B is not a **near-real number**, signal **domain-error**. Set A1 to the **real-number** nearest to A. Set B1 to the **real-number** nearest to B. If A1 is **greater-than** B1, return B1. Otherwise, return A1.

7.2.7 Power

$$Z \leftarrow A * B$$

Informal Description: *Z* is *A* raised to the *B*th power.

Evaluation Sequence:

If either of A or B is not a **number**, signal **domain-error**.

If *A* is **zero** and *B* is **zero**, return **one**.

If *A* is **zero** and the **real-part** of *B* is a **positive-number**, return **zero**, otherwise signal **domain-error**.

Return A to-the-power B.

Examples:

- In the following, **print-precision** is 12.

```
2*32
4294967296
4*0.5
2
-8*÷3
1J1.73205080757
```

Additional Requirement:

The foregoing example shows that, when the optional complex arithmetic facility is implemented, the **implementation-algorithm** should yield the principal value of the nth (odd n) root of a negative number, not the real negative root. If it is not implemented, a domain error should be signalled.

7.2.8 Logarithm

 $Z \leftarrow A \otimes B$

Informal Description: Z is the logarithm of B to the base A.

Evaluation Sequence:

If either of A or B are not **numbers** signal **domain-error**. If A and B are **equal**, return **one**. If A is **one**, signal **domain-error**. Set A1 to the **natural-logarithm** of A. Set B1 to the **natural-logarithm** of B. Return B1 **divided-by** A1.

Example:

10 2 10 0.1 2 65536 1E15 1E15 0.30103 16 15 15

7.2.9 Residue

$$Z \leftarrow A \mid B$$

Informal Description: *Z* is *B* modulo *A*. Uses **comparison-tolerance**.

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

If either of A or B is not a **number**, signal **domain-error**.

If *A* is **zero**, return *B*.

If comparison-tolerance is not zero, and B divided-by A is integral-within comparison-tolerance, return zero.

Otherwise set Z to B **modulo** A.

If Z is A, return **zero**.

Otherwise, return Z.

Examples:

– In the following, **print-precision** is 16 and **comparison-tolerance** is $1E^-10$.

- In the following, **comparison-tolerance** is **zero**.

Additional Requirement:

The range of residue is those numbers that are the product of a **fraction** and A, except when A is **zero**, in which case the range is the single **number** B.

Note: The implementation-algorithm P modulo Q provides an exact modulo operation for real-numbers P and Q. It evaluates $R \leftarrow P - (\times P) \times |Q \times L| P \div Q$ exactly, and returns R if $(\times R) = \times Q$, or R + Q otherwise.

The definition of "mod" in the IEEE standard for Binary Floating-Point Arithmetic (754) provides an example of this exact evaluation.

Implementations should avoid signalling **limit-error** in **residue**. If the operation B **divided-by** A causes **exponent-overflow**, return **zero**. If it causes **exponent-underflow**, and if A and B have the same signs, return B. If they have different signs, return **zero**.

7.2.10 Binomial

$$Z \leftarrow A : B$$

Informal Description: Z is $(gamma(1+B)) \div ((gamma(1+A)) \times gamma(1+B-A))$

If A and B are **nonnegative-integers**, Z is the number of combinations of B things taken A at a time.

Evaluation Sequence:

If either of A or B is not a **number**, signal **domain-error**.

Determine if each of A, B, and B-A is a **negative-integer**.

Select the appropriate case from the following table, where a **one** indicates that the corresponding value is a **negative-integer** and a **zero** indicates that it is not.

Ca	se		Rule
\boldsymbol{A}	В	B– A	
0	0	0	Return $(!B) \div (!A) \times !B - A$.
0	0	1	Return zero.
0	1	0	Signal domain-error.
0	1	1	Return $(-1*A)*A!A-B+1$.
1	0	0	Return zero.
1	0	1	(Case cannot arise.)
1	1	0	Return $(-1*B-A) \times (B+1)! (A+1)$.
1	1	1	Return zero.

Example:

	- ₄ - ₃	_2	⁻ 1 0	1 2	3 4	· .!	-4 -	3 -2	⁻ 1 0	1 2	3 4
1	-3	3	-1	0	0	0	0	0			
0	1	_2	1	0	0	0	0	0			
0	0	1	_1	0	0	0	0	0			
0	0	0	1	0	0	0	0	0			
1	1	1	1	1	1	1	1	1			
-4	-3	_2	_1	0	1	2	3	4			
10	6	3	1	0	0	1	3	6			
_20	_10	-4	_1	0	0	0	1	4			
35	15	5	1	0	0	0	0	1			

Note: The APL expressions in the **rule** column indicate the result required, not the algorithm to be used. For example, 64!65 should be 65 even if !65 signals **limit-error**.

7.2.11 Circular Functions

 $Z \leftarrow A \circ B$

Informal Description: Z is the result of applying a function designated by A to B.

Evaluation Sequence:

```
If A is not a near-integer, signal domain-error.
If B is not a number, signal domain-error.
Set A1 to the integer-nearest-to A.
If A1 is not in the closed-interval-between 12 and 12, signal domain-error.
If A1 is ^-12, return \star 0J1 \times B.
If A1 is -11, return 0J1 \times B.
If A1 is -10, return +B.
If A1 is -9, return B.
If A1 is ^{-}8, return -(^{-}1-B*2)*.5.
If A1 is -7,
  If B is negative-one or one, signal domain-error.
  Return the inverse-hyperbolic-tangent of B.
If A1 is ^{-}6,
  Return Z, the principal value of the inverse-hyperbolic-cosine of B, where Z is a
    nonnegative-number.
If A1 is 5, return the inverse-hyperbolic-sine of B.
If A1 is -4,
  If B is -1 return zero
  Otherwise return (B+1)\times((B-1)\div B+1)\star 0.5.
If A1 is \overline{\phantom{a}}3, return Z, the principal value in radians of the inverse-tangent of B, where
 Z is in the open-interval-between -\pi/2 and \pi/2.
If A1 is -2,
  Return Z, the principal value in radians of the inverse-cosine of B, where Z is either
    zero or a number in the open-interval-between zero and \pi.
A1 is ^{-}1,
  Return Z, the principal value in radians of the inverse-sine of B, where Z is either
    \pi/2, or a number in the open-interval-between -\pi/2 and \pi/2.
If A1 is 0.
  If B is not in the closed-interval-between negative-one and one, signal domain-
    error.
  Return (1-B*2)*0.5.
If A1 is 1, return the sine of B radians.
If A1 is 2, return the cosine of B radians.
  If B is an odd multiple of \pi/2, signal domain-error.
  Return the tangent of B radians.
If A1 is 4, return (1+B*2)*0.5.
If A1 is 5, return the hyperbolic-sine of B.
If A1 is 6, return the hyperbolic-cosine of B.
```

```
If A1 is 7, return the hyperbolic-tangent of B. If A1 is 8, return (-1-B*2)*.5. If A1 is 9, return the real-part of B. If A1 is 10, return |B|. If A1 is 11, return the imaginary-part of B. If A1 is 12, return the arc of B.
```

Examples:

Note: The APL expressions used for 00X, -40X, and 40X above indicate the result desired, not the algorithm to be used.

Note: The following article describes the reasons for the choices made in defining the circular functions for complex arguments,

Penfield, Paul, "Principal Values and Branch Cuts in Complex APL", APL81 Conference Proceedings, ACM, San Francisco, 1981

Note: Values of A greater than 7 in magnitude are required only for an implementation which includes the optional **complex-arithmetic-facility**.

7.2.12 **And/LCM**

```
Z \leftarrow A \wedge B
```

Informal Description: *Z* is the least common multiple of *A* and *B*. For Boolean arguments it is the Boolean product of *A* and *B*. Uses **comparison-tolerance**.

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

If either A or B is not a **number**, signal **domain-error**.

If both A and B are **near-boolean**,

Set A1 to the **integer-nearest-to** A.

Set B1 to the **integer-nearest-to** B.

If either A1 or B1 is **zero**, return **zero**.

Otherwise, return one.

Otherwise, set A1 to the **greatest-common-divisor** of A and B, using the **implementation-algorithm greatest-common-divisor**.

If A1 is **zero**, return **zero**.

Otherwise, return A times (B divided-by A1).

```
0 1 •.^0 1

0 0

0 1

30 ^ 36

180

3 ^ 3.6

18

-29J53 ^ -1J107

-853J-329
```

7.2.13 Or/GCD

$$Z \leftarrow A \vee B$$

Informal Description: Z is the greatest common divisor of A and B. For Boolean arguments, Z is the Boolean sum of A and B. Uses **comparison-tolerance**.

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

If either A or B is not a **number**, signal **domain-error**.

If both A and B are **near-boolean**,

Set A1 to the **nearest-integer** to A.

Set B1 to the **nearest-integer** to B.

If either A1 or B1 is **one**, return **one**.

Otherwise, return zero.

Otherwise, return the **greatest-common-divisor** of A and B, using the **implementation-defined-algorithm greatest-common-divisor**.

Examples:

7.2.14 Nand

$$Z \leftarrow A \nsim B$$

Informal Description: Z is the Boolean complement of the Boolean product of A and B.

Evaluation Sequence:

If either A or B is not **near-Boolean**, signal **domain-error**. Otherwise, return $\sim A \wedge B$, with **comparison-tolerance** set to **zero**.

7.2.15 Nor

$$Z \leftarrow A \nsim B$$

Informal Description: Z is the Boolean complement of the Boolean sum of A and B.

Evaluation Sequence:

If either A or B is not **near-Boolean**, signal **domain-error**. Otherwise, return $\sim A \vee B$, with **comparison-tolerance** set to **zero**.

7.2.16 Equal

$$Z \leftarrow A = B$$

Informal Description: Z is **one** if A and B are considered equal and **zero** otherwise. A and B are equal if they are the same **character**, or if they are both **numeric** and A is tolerantly equal to B within **comparison-tolerance**. Uses **comparison-tolerance**.

Evaluation Sequence:

If **comparison-tolerance** is **nil**, signal **implicit-error**.

If the **type** of *A* is not the same as the **type** of *B*, return **zero**.

If both A and B are **characters**,

If A is the same **character** as B, return **one**.

Otherwise, return zero.

If both A and B are **numbers**,

If A is tolerantly-equal to B within comparison-tolerance, return one.

Otherwise, return zero.

Example:

- In the following, **comparison-tolerance** is $1E^-13$.

```
4 = 4 + 5E^{-}13 \ 2E^{-}13 \ ^{-}2E^{-}13 \ ^{-}5E^{-}13
0 1 1 0
0 = ^{-}1E^{-}20 \ 1E^{-}20 \ 0
0 0 1
3 = ^{!}A3^{!}
0 0
```

Note: Comparisons of numbers whose signs differ are not affected by comparison-tolerance.

For any value of **comparison-tolerance** and any two real numbers A and B, exactly one of the expressions A < B, A = B, and A > B is **one**.

Equal should not signal a **limit-error**. For example, the result of **positive-number-limit** = **negative-number-limit** is **zero**. The following is a sample technique for handling exponent-overflow and exponent-underflow when scaling **comparison-tolerance**.

Set C to the larger of the magnitudes of A and B.

Set D to comparison-tolerance times C.

If exponent-underflow occurs,

Set A1 to A divided-by C.

Set B1 to B divided-by C.

Set C1 to the magnitude of A1 minus B1.

```
If C1 is greater-than comparison-tolerance, return zero. Otherwise, return one.

Set E to the magnitude of A minus B.

If exponent-overflow occurs, return zero.

If exponent-underflow occurs,

Set A1 to A divided-by C.

Set B1 to B divided-by C.

Set C1 to the magnitude of A1 minus B1.

If C1 is greater-than comparison-tolerance, return zero.

Otherwise, return one.

If E is not greater-than D, return one.

Otherwise, return zero.
```

7.2.17 Less than

$$Z \leftarrow A < B$$

Informal Description: Z is **one** if A is tolerantly less-than B, and **zero** otherwise. Uses **comparison-tolerance**.

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

If either of A or B is not a **near-real number**, signal **domain-error**.

Set A1 to the **real-number** nearest to A.

Set B1 to the **real-number** nearest to B.

If A1=B1, evaluated with the current value of **comparison-tolerance**, is **one**, return **zero**.

If A1 is **less-than** B1, return **one**.

Otherwise, return zero.

7.2.18 Less than or equal to

$$Z \leftarrow A \leq B$$

Informal Description: Z is **one** if A is less than or tolerantly equal to B, and **zero** otherwise. Uses **comparison-tolerance**.

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

If either of A or B is not a **near-real number**, signal **domain-error**.

Set A1 to the **real-number** nearest to A.

Set B1 to the **real-number** nearest to B.

If A1=B1, evaluated with the current value of **comparison-tolerance**, is **one**, return **one**.

If A1 is **less-than** B1, return **one**.

Otherwise, return zero.

7.2.19 Not equal

$$Z \leftarrow A \neq B$$

Informal Description: Z is **one** if A does not equal B, and **zero** otherwise. Uses **comparison-tolerance**.

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

Return $\sim A=B$, evaluated with the current value of **comparison-tolerance**.

7.2.20 Greater than or equal to

$$Z \leftarrow A \geq B$$

Informal Description: Z is **one** if A is greater than or tolerantly equal to B, and **zero** otherwise. Uses **comparison-tolerance**.

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

If either of A or B is not a **near-real number**, signal **domain-error**.

Set A1 to the **real-number** nearest to A.

Set B1 to the **real-number** nearest to B.

If A1=B1, evaluated with the current value of **comparison-tolerance**, is **one**, return **one**

If A1 is **greater-than** B1, return **one**.

Otherwise, return zero.

```
1 2 3 °.≥1 2 3
1 0 0
1 1 0
1 1 1
0 1 °.≥0 1
1 0
1 1
```

7.2.21 Greater than

$$Z \leftarrow A > B$$

Informal Description: Z is **one** if A is tolerantly greater than B, and **zero** otherwise. Uses **comparison-tolerance**.

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

If either of A or B is not a **near-real number**, signal **domain-error**.

Set A1 to the **real-number** nearest to A.

Set B1 to the **real-number** nearest to B.

If A1=B1, evaluated with the current value of **comparison-tolerance**, is **one**, return **zero**.

If A1 is **greater-than** B1, return **one**.

Otherwise, return zero.

8 Structural Primitive Functions

8.1 Introduction

Note: The functions in this chapter are used in the evaluation sequences of many non-scalar operations. They are defined here to avoid forward references.

8.2 Monadic Structural Primitive Functions

8.2.1 Ravel

$$Z \leftarrow , B$$

Informal Description: Z is a vector containing the elements of B in row-major order.

Evaluation Sequence:

Return Z, a **vector** such that the **ravel-list** of Z is the same as the **ravel-list** of B, the **type** of Z is the same as the **type** of B, and the **shape-list** of Z is a **list** of length **one** containing the **count** of B as its only **item**.

Examples:

```
,N22
11 12 21 22
,N222
111 112 121 122 211 212 221 222
,N2221
1111 1121 1211 1221 2111 2121 2211 2221
```

Note: Ravel always produces a vector result. The expressions **ravel-list** of \mathbb{Z} , **type** of \mathbb{Z} , and **shape-list** of \mathbb{Z} refer to attributes of an **array** object.

8.2.2 Shape

$$Z \leftarrow \rho B$$

Informal Description: Z is a **numeric** vector containing the shape of the array B.

Evaluation Sequence:

Return Z, an **array** such that the **type** of Z is **numeric**, the **ravel-list** of Z is the **shape-list** of B, and the **shape-list** of Z is a **list** whose only **item** contains the **number-of-items** in the **shape-list** of B.

Examples:

Note: Shape always produces a vector result. The expression **shape-list** of Z refers to an attribute of an **array** object.

8.2.3 Index Generator

$$Z \leftarrow \iota B$$

Informal Description: Z is a **numeric** vector of B consecutive ascending **integers**, the first of which is **index-origin**. Uses **index-origin**.

Evaluation Sequence:

If index-origin is nil, signal implicit-error.

If the rank of B is greater-than one, signal rank-error.

If the **count** of B is not **one**, signal **length-error**.

If *B* is not a **near-integer**, signal **domain-error**.

Set B1 to the **integer-nearest-to** B.

If *B*1 is not a **nonnegative-integer**, signal **domain-error**.

If B1 is **zero**, return an **empty numeric vector**.

Generate a **numeric vector** Z1 of length B1 such that the **ravel-list** of Z1 consists of the **integers** in the **closed-interval-between one** and B1 in ascending order. Generate Z, a **numeric array** with the **shape-list** of Z1 such that each **item** of the **ravel-list** of Z is (**index-origin minus one**) **plus** the corresponding element of Z1.

Return Z.

Examples:

- In the following example, **index-origin** is **zero**.

- In the following example, **index-origin** is **one**.

8.2.4 Table

$$Z \leftarrow \overline{,} B$$

Informal Description: *Z* is a **matrix** containing the elements of *B*. *Z* is formed by ravelling the subarrays of *B* along its first axis. For example, if *B* is a scalar, *Z* has shape 1 1; if *B* is a five-element vector, *Z* has shape 5 1; if *B* has shape 5 4 3 2, *Z* has shape 5 24.

Evaluation Sequence:

```
If B is a scalar set C to 1 1.
Else set C to a two-element list as follows:
Set the first-item of C to the first-item of the shape-list of B.
Set the last-item of C to the product-of the rest-of the shape-list of B.
Set Z to the C reshape of B.
Return Z.
```

```
,0
0
     ρ,0
1 1
     -N4
1
2
3
4
     ,N22
11 12
21 22
     -N222
111 112 121 122
211 212 221 222
     ,N2221
1111 1121 1211 1221
2111 2121 2211 2221
```

8.2.5 **Depth**

$$Z \leftarrow \equiv B$$

Informal Description: Z is a numeric scalar describing the level of nesting of B. For **simple-scalars**, this is 0. For arrays, this is one greater than the depth of the element with the greatest depth.

Note: An enclosed array is an array.

Evaluation Sequence:

If *B* is a **simple-scalar**, return **zero**. Return $1+\lceil / , \equiv B$.

```
=5

0

=1 2 3

1

=N234

1

='ABC',1 2 3

1

=-1 2 3

2

=,-1 2 3

2

='HERO',-2 3,-2 3ρ-5 8
```

8.2.6 Enlist

```
Z \leftarrow \epsilon B
```

Informal Description: If B is a **simple-array**, then enlist has the same effect as ravel. Otherwise, enlist has the effect of recursively raveling each element of B (in **ravel-list** order) and joining them together. Thus, Z will always be a **simple-array** of rank one, containing all items of B.

Note: Enlist is sometimes called "super ravel".

Evaluation Sequence:

```
If B is empty, return 0 \cap \in 1 \uparrow, B.
If B is a simple-array, return , B.
Set B1 to the first item of the ravel-list of B.
Return (\in B1), (\in 1 \downarrow, B).
```

Examples:

8.3 Dyadic Structural Primitive Functions

8.3.1 Reshape

$$Z \leftarrow A \rho B$$

Informal Description: Z is an array of shape , A whose elements are taken sequentially from , B repeated cyclically as required.

Evaluation Sequence:

If the rank of A is greater-than one, signal rank-error.

If any **item** of the **ravel-list** of A is not a **near-integer**, signal **domain-error**.

Set A1 to the **integer-array-nearest-to** , A.

If any **item** of the **ravel-list** of A1 is not a **nonnegative-counting-number**, signal **domain-error**.

Let RA stand for the **product-of** the **ravel-list** of A1.

Let CB stand for the **count** of B.

If *RA* is not **zero** and *CB* is **zero**, signal **length-error**.

Return an **array** Z such that the **type** of Z is the **sufficient-type** of the **ravel-list** of Z under the **type** of B, the **shape-list** of Z is the same as the **ravel-list** of A1, and the **ravel-list** of Z is a **list** with RA items such that for all I in the **index-set** of Z, item I of the **ravel-list** of Z is **item** $1+CB \mid I-1$ of the **ravel-list** of B.

Examples:

Note: For any X that is not **empty**, $" \circ X$ and $(\circ \circ) \circ X$ produce the same result: a **scalar** whose value is that of the **first-scalar** in X.

8.3.2 Join

$$Z \leftarrow A$$
 , B

$$Z \leftarrow A$$
 , B

Informal Description: If A and B are scalars or vectors, Z is the vector of length $(\rho, A) + \rho$, B whose first ρ , A elements are , A and whose last ρ , B elements are , B.

```
If either A or B has rank greater than one,
For form Z \leftarrow A, B
Z is A, [(\rho \rho A) \lceil \rho \rho B]B,
For form Z \leftarrow A, B
Z is A, [1]B,
```

as defined under Join Along an Axis.

Evaluation Sequence:

If A is a **scalar** and B is a **scalar**, return (,A), ,B. If A is a **scalar** and B is a **vector**, return (,A), B. If A is a **vector** and B is a **scalar**, return A, B. If A is a **vector** and B is a **vector**,

If A is **empty** and B is **empty** and the **type** of A differs from the **type** of B, return an **empty-list** whose **type** is determined by the **implementation-algorithm typical-element-for-mixed**.

If B is **empty**, return A.

If A is **empty**, return B.

Otherwise, return a **vector** Z, such that the **shape-list** of Z is $(\rho A) + \rho B$, the **ravel-list** of Z is a **list** whose first ρA **items** are the **ravel-list** of A and whose last ρB **items** are the **ravel-list** of B, and the **type** of Z is the **sufficient-type** of the **ravel-list** of Z under the **type mixed**.

Otherwise, when either A or B has rank greater than **one**, return A, [1]B for form A, B and A, $[(\rho \rho A) [\rho \rho B]B$ for form A, B.

Example:

Note: This subsection intentionally contains a forward reference to Join Along an Axis. Join and Join Along an Axis are defined separately because the description of Join Along an Axis requires APL operations that depend for their definitions in turn upon Join.

9 Operators

9.1 Introduction

Note: The forms in this chapter are referred to as **operators**.

9.2 Monadic Operators

9.2.1 Reduction

 $Z \leftarrow f/B$

 $Z \leftarrow f/[K] B$

 $Z \leftarrow f \neq B$

 $Z \leftarrow f \neq [K] B$

Informal Description: Z is the value produced by placing the dyadic function \mathbf{f} between adjacent items along a designated axis of B and evaluating the resulting expression. The axis designated determines how the subarrays are chosen. Uses **index-origin**.

There are two conforming definitions for Reduction; the **Implementation Parameter Reduction-Style** specifies which definition a particular implementation uses.

The **Enclose-Reduction-Style** style, informally known as the APL2 style, uses a definition based on enclose which, for all **f**, preserves the identity:

$$\rho \rho Z is 0 \Gamma^{-} 1 + \rho \rho B.$$

The **Insert-Reduction-Style** style, informally known as the Sharp/J style, does not universally preserve the above identity. Instead \mathbf{f} is inserted between each successive **cell** along the **frame** axis (the axis of reduction). The rank of \mathbf{f} controls the subsequent evaluation.

Note: Since reduction is used in the evaluation sequences of other operators, the choice of the Implementation Parameter Reduction-Style has a pervasive effect on the implementation, and its importance should not be underestimated.

Note: The functions listed in Table 5 are scalar functions, that is their left and right function ranks are zero.

Evaluation Sequence:

```
If Implementation Parameter Reduction-Style is Enclose-Reduction-Style:
  If f is not a dyadic-function, signal syntax-error.
  For form f/B
     If B is a scalar, return B.
     Otherwise, return \mathbf{f}/[-1 \uparrow \iota \rho \rho B] B.
  For form \mathbf{f} \neq B
     If B is a scalar, return B.
     Otherwise, return \mathbf{f} / [11] B.
  For forms \mathbf{f}/[K] B and \mathbf{f}/[K] B
     If K is not a valid-axis for B, signal axis-error.
     Otherwise, set K1 to the integer-nearest-to K.
     If B is a vector.
        If the length of B is zero, take the action designated in Table 5 for f.
        If the length of B is one, return a scalar Z such that the type of Z is the type
         of B and the ravel-list of Z is the ravel-list of B.
        If the length of B is greater-than one,
          Set B1 to the first-scalar in B.
          Set B2 to the remainder-of B.
          Return \subseteq B1 f \supset f /B2.
     If the rank of B is greater-than one, return an array Z such that the shape-list
      of Z is the shape-list of B with item K1 omitted, and the ravel-list of Z has
```

the property that if Z1 is an **item** of Z and B3 is the corresponding **vector-item along-axis** K1 of B, then Z1 is f/B3.

Otherwise Implementation Parameter Reduction-Style is Insert-Reduction-Style:

```
If f is not a dyadic-function, signal syntax-error.
```

```
For form f/B
   If B is a scalar, return B.
   Otherwise, return \mathbf{f}/[-1\uparrow \iota \rho B] B.
For form \mathbf{f} \neq B
   If B is a scalar, return B.
   Otherwise, return \mathbf{f} / \lceil \iota 1 \rceil B.
For forms \mathbf{f}/\lceil K \rceil B and \mathbf{f}/\lceil K \rceil B
   If K is not a valid-axis for B, signal axis-error.
   Otherwise, set K1 to the integer-nearest-to K.
   Set B1 to (AK1, (1ppB) \sim K1) \lozenge B.
   Set C1 to 1 \downarrow \rho B1.
```

If the number of C1-cells in B1 is **zero**, take the action designated in **Table 5** for **f** to choose a result item Z0. Return an **array** Z such that the **shape-list** of Z is C1, the **type** of Z is the **type** of Z0, and each item of the **ravel-list** of Z is the

ravel-list of Z0.

If the number of C1-cells in B1 is **one**, return an **array** Z such that the **shape-list** of Z is C1, the **type** of Z is the **type** of B1, and the **ravel-list** of Z is the **ravel-list** of B1.

If the number of C1-cells in B1 is **greater-than one**,

Set B2 to an array such that the **shape-list** of B2 is C1, the **type** of B2 is the **type** of B1, and the **ravel-list** of B2 is the **ravel-list** of the first C1-cell in B1. Set B3 to an array such that the **shape-list** of B3 is (($^-1+1+pB1$), C1), the **type** of B3 is the **type** of B1, and the **ravel-list** of B3 is the **ravel-list** of all the the C1-cells of B1 except the first.

Return *B*2 **f f** *F B*3.

Examples:

Additional Requirement:

If Z is **empty**, the **type** of Z is determined by the argument function \mathbf{f} .

When applied along an empty **axis** in an array, reduction produces an array whose shape is that of the argument array with the designated **axis** deleted. The elements of the array, if any, are determined by the argument function and **Table 5**. In some cases, an error is signalled.

```
+/2 0 ρ5.1
0 0
ρ+/2 0ρ5.1
2
```

Dyadic Function		Action
Plus Minus Times Divide	+ - × ÷	Return zero . Return zero . Return one . Return one .
Residue Minimum Maximum Power	 	Return zero . Return positive-number-limit . Return negative-number-limit . Return one .
Binomial And Or Less	! ^ V	Return one . Return zero . Return zero .
Not greater Equal Not less Greater	≤ = ≥ >	Return one . Return one . Return one . Return zero .
Not equal all others	≠	Return zero . Signal domain-error .

Table 5: Actions for the Reduction of an Empty Vector.

9.2.2 Scan

```
Z \leftarrow \mathbf{f} \setminus B
Z \leftarrow \mathbf{f} \setminus [K] B
Z \leftarrow \mathbf{f} + B
Z \leftarrow \mathbf{f} + [K] B
```

Informal Description: Z is an array having the same shape as B and containing the results produced by f reduction over all prefixes of a designated axis of B. Uses **index-origin**.

Evaluation Sequence:

```
If f is not a dyadic-function, signal syntax-error. For form \mathbf{f} \setminus B

If B is a scalar, return B.

Otherwise, return \mathbf{f} \setminus [\rho \rho B] B evaluated with index-origin set to one. For form \mathbf{f} \setminus B

If B is a scalar, return B.

Otherwise, return \mathbf{f} \setminus [1] B evaluated with index-origin set to one.
```

```
For forms \mathbf{f} \setminus [K]B and \mathbf{f} + [K]B
If K is not a valid-axis for B, signal axis-error.
Otherwise, set K1 to the integer-nearest-to K.
If B is a vector,
```

If the **count** of B is **less-than two**, return B.

Otherwise return Z, a **vector** such that the **shape-list** of Z is the **shape-list** of B, and the **ravel-list** of Z is such that **item** I of the **ravel-list** of Z is $\mathbf{f}/B[\iota I]$ for all I in the **index-set** of the **ravel-list** of B. The **type** of Z is the **sufficient-type** of the **ravel-list** of Z under the **type mixed**.

If the rank of B is greater-than one, each vector-item along-axis K1 of Z is $f \setminus B1$, where B1 is the corresponding vector-item along-axis K1 of B.

Examples:

Additional Requirement:

If the operator **reduction** signals an error when called by **scan**, **scan** returns the resultant error **token**.

The evaluation sequence above describes a quadratic algorithm for **scan**. If the function **f** is associative, **scan** may be implemented with a linear algorithm.

Note: Various error checks have been performed on K by the phrase evaluators, and **index-origin** is **one**, before this evaluation sequence is called.

9.2.3 N-wise Reduction

```
Z \leftarrow N \ f/B
Z \leftarrow N \ f/[K] B
Z \leftarrow N \ f \neq B
Z \leftarrow N \ f \neq [K] B
```

Informal Description: N-wise reduction is the dyadic invocation of a function derived by reduction. Z is the value produced by placing the primitive scalar dyadic function f between subarrays of B and evaluating the resulting expression. If N is negative the subarrays are reversed before application of f. Each subarray has length N along a common dimension. The axis designated determines how the subarrays are chosen. Uses **index-origin**.

Evaluation Sequence:

```
If the rank of N is greater-than one, signal rank-error.
```

If the length of the **ravel-list** of N is **greater-than one**, signal **length-error**.

If N is not a **near-integer**, signal **domain-error**.

Set N1 to the **integer-nearest-to** N.

Set M1 to the **magnitude** of N1.

If **f** is not a **dyadic-function**, signal **syntax-error**.

For form $N \mathbf{f}/B$

If B is a **scalar**.

If M1 is **greater-than two**, signal **domain-error**.

If M1 is **zero**, set R to **two** and take the action designated in **Table 6** for R and **f**.

Otherwise, return $(2-M1) \rho B$.

Otherwise, return N **f**/ $[\rho \rho B]$ B evaluated with **index-origin** set to **one**.

For form $N \mathbf{f} \neq B$

If B is a scalar,

If *M*1 is **greater-than two**, signal **domain-error**.

If M1 is **zero**, set R to **two** and take the action designated in **Table 6** for R and **f**. Otherwise, return $(2-M1) \rho B$.

Otherwise, return N **f**/[1] B evaluated with **index-origin** set to **one**.

For forms $N \mathbf{f} / [K]$ B and $N \mathbf{f} \neq [K]$ B

If *K* is not a **valid-axis** for *B*, signal **axis-error**.

Otherwise, set *K*1 to the **integer-nearest-to** *K*.

If *B* is a **vector**,

If M1 is greater-than one plus the length of B, signal domain-error.

If M1 is **zero**, set R to **one plus** the **length** of B and take the action designated in **Table 6** for R and \mathbf{f} .

Set B1 to M1 take B.

If N1 is less-than zero,

set B2 to $\phi B1$

```
otherwise set B2 to B1.

If M1 equals the length of B, return f/B2.

Otherwise,
Set B3 to f/B2.
Set B4 to one drop B.
Return B3, Nf / B4.
```

If the **rank** of B is **greater-than one**, return an **array** Z such that the **shape-list** of Z is the **shape-list** of B with **item** K1 replaced by **one plus** the value of **item** K1 **minus** M1, and the **ravel-list** of D has the property that if D is the **vector-item along-axis** D and D is the corresponding **vector-item along-axis** D of D, then D is D **if** D is D if D is D is D if D is D is D is D if D is D if D is D if D is D is

Dyadic Function		Action
Plus Minus Times Divide	+ - × ÷	Return R reshape zero . Return R reshape zero . Return R reshape one . Return R reshape one .
Residue Minimum Maximum Power	 	Return R reshape zero . Return R reshape positive-number-limit . Return R reshape negative-number-limit . Return R reshape one .
Logarithm Circular Binomial And	⊗ ∴ ∧	Signal domain-error . Signal domain-error . Return R reshape one . Return R reshape one .
Or Nand Nor Less	∨ * <	Return <i>R</i> reshape zero . Signal domain-error . Signal domain-error . Return <i>R</i> reshape zero .
Not greater Equal Not less Greater	<pre> </pre> <pre> <pre> <pre> </pre> <pre> <pre> </pre> <pre> </pre> > </pre></pre></pre>	Return R reshape one . Return R reshape one . Return R reshape one . Return R reshape zero .
Not equal all others	≠	Return R reshape zero . Signal domain-error .

Table 6: Actions for the N-wise Reduction of an Empty Vector.

9.2.4 Duplicate

$$Z \leftarrow \mathbf{f} \stackrel{::}{\sim} B$$

Informal Description: Z is B **f** B.

Evaluation Sequence:

Return B **f** B.

Example:

Note: If f is not an ambivalent-function, a valence-error will be signalled.

9.2.5 Commute

$$Z \leftarrow A \mathbf{f} \stackrel{\dots}{\sim} B$$

Informal Description: Z is B **f** A.

Evaluation Sequence:

Return B f A.

9.2.6 Each

$$Z \leftarrow \mathbf{f}$$
 B
 $Z \leftarrow A \mathbf{f}$ B

Informal Description: The operand function **f** is applied independently to corresponding **items** of the arguments, or (monadic case) independently to the **items** of the argument. The corresponding results are assembled in an array of the same shape as the argument(s).

Evaluation Sequence:

```
For form A\mathbf{f} B:
  If the rank of A differs from the rank of B,
     If A is a scalar or one-element-vector and B is not a scalar, set A to (\rho B)\rho A.
        If B is a scalar or one-element-vector, set B to (\rho A)\rho B.
        Otherwise, Signal rank-error.
  If the shape-list of A differs from the shape-list of B, Signal length-error.
For both forms:
  If B is not empty, Signal domain-error.
  Create Z, an array having the same shape-list as B.
  For each I in the index-set of the ravel-list of Z:
     For form \mathbf{f} \cdot B.
        Set X to item I of the ravel-list of B.
        Evaluate-Monadic-Function with fX giving token T.
     For form A\mathbf{f} B,
        Set X to item I of the ravel-list of A.
        Set Y to item I of the ravel-list of B.
        For each pair X and Y of corresponding items from A and B, Evaluate-Dyadic-
         Function with XfY giving token T.
     If T is an error, return T.
     Set Q to the content of T.
     Set item I of the ravel-list of Z as follows:
        If Q is a simple-scalar, the first-item of the ravel-list of Q.
        Otherwise, Q.
  If all items of the ravel-list of Z are nil, return nil.
```

If all items of the **ravel-list** of Z are **values**, return Z.

Otherwise, signal value-error.

9.3 Dyadic Operators

9.3.1 Outer Product

$$Z \leftarrow A \circ \mathbf{.f} B$$

Informal Description: Z is an array of shape (ρA) , ρB . The elements of Z are the result of applying the **dyadic-function f** to every possible combination of scalar arguments where the left argument is an element of A and the right an element of B.

Z is such that if I is an **index-list** that selects a single element of Z, the first ρA **items** of I are the **index-list** that would select from A the element used as the left argument to \mathbf{f} and the last ρB **items** of I are the **index-list** that would select from B the element used as the right argument to \mathbf{f} when producing the selected element of Z.

Evaluation Sequence:

If **f** is not a **dyadic-function**, signal **syntax-error**.

Return Z, an **array** such that the **shape-list** of Z is (ρA) , ρB , and the **ravel-list** of Z has the following property:

Let I stand for an **item** of the **index-set** of the **ravel-list** of A.

Let J stand for an **item** of the **index-set** of the **ravel-list** of B.

Let X stand for **item** I of the **ravel-list** of A.

Let Y stand for **item** J of the **ravel-list** of B.

Let *N* stand for the **count** of *B*.

Let P stand for $J+(N\times(I-1))$.

Set Q to X **f** Y.

If Q is a simple-scalar, set Q1 to the first-item of the ravel-list of Q.

Otherwise, set Q1 to Q.

Then, **item** P of the **ravel-list** of Z is Q1.

Set the **type** of Z to the **sufficient-type** of the **ravel-list** of Z under mixed.

Example:

```
10 20 30 °.+ 1 2 3
11 12 13
21 22 23
31 32 33
```

Additional Requirement:

If the **dyadic-function f** signals an error, outer product returns the resulting error **token**.

9.3.2 Inner Product

 $Z \leftarrow A \mathbf{f} \cdot \mathbf{g} B$

Informal Description: Z is an array of shape $(\rho A) [\iota 0 [-1 + \rho \rho A], (\rho B) [1 + \iota 0 [-1 + \rho \rho B]]$. The elements of Z are the results obtained from evaluating the expression $\mathbf{f}/X\mathbf{g}Y$ for all possible combinations of X and Y, where X is a **vector-item along-axis** $\rho \rho A$ of A and Y a **vector-item along-axis one** of B.

Evaluation Sequence:

If either **f** or **g** is not a **dyadic-function**, signal **syntax-error**.

If A is a scalar or one-element-vector and B is not, set A1 to $(100B) \rho A$.

If A and B are scalars or one-element-vectors, set A1 to A.

Otherwise, set *A*1 to *A*.

If *B* is a **scalar** or **one-element-vector** and *A* is not, set *B*1 to $(\rho A) [\rho \rho A] \rho B$.

If A and B are scalars or one-element-vectors, set B1 to B.

Otherwise, set B1 to B.

If the **last-item** in the **shape-list** of A1 is not the same as the **first-item** in the **shape-list** of B1, signal **length-error**.

If A1 and B1 are both vectors, return f/A1g B1.

Otherwise, set Z to an **array** such that the the **shape-list** of Z is $(\rho A1) [10 [-1+\rho\rho A1], (\rho B1) [1+10 [-1+\rho\rho B1]]$ and the **ravel-list** of Z has the following property:

Let I stand for an **item** of the **index-set** of the **ravel-along-axis** ($\rho \rho A1$) of A1.

Let X stand for vector-item I of the ravel-along-axis ($\rho \rho A1$) of A1.

Let J stand for an **item** of the **index-set** of the **ravel-along-axis one** of B1.

Let Y stand for vector-item J of the ravel-along-axis one of B1.

Let N stand for the **number-of-items** in the **ravel-along-axis one** of B1.

Let P stand for $(N \times (I-1)) + J$.

Set Q to $\mathbf{f}/X\mathbf{g}$ Y.

If Q is a **simple-scalar**, set Q1 to the **first-item** of the **ravel-list** of Q.

Otherwise, set Q1 to Q.

Then, **item** P of the **ravel-list** of Z is Q1.

Set the **type** of Z to the **sufficient-type** of Z under mixed.

Return Z.

Examples:

```
4 2 1+.×1 0 1

5

N22+.×0 1

12 22

N22+.×1 0

11 21

N22+.×2 2ρ0 1 1 0

12 11

22 21
```

Additional Requirement:

If scalar-function f or g signals an error, inner product returns the resulting error token.

Note: The evaluation sequence rule for when A1 and B1 are **vectors** holds if A1 or B1 is the empty vector. The result returned is $\mathbf{f}/\iota 0$.

The type of the result of inner product is numeric only because all permitted argument functions return numeric results.

9.3.3 Rank operator definitions

The following definitions are used in connection with the **rank** operator:

cell: The last **k items** of the **shape-list** of an **array** determine **rank-k cells** of the **array**. **frame**: For an array of **rank r**, its **frame** with respect to its **cells** of **rank k** is the **r-k** leading elements of its **shape-list**.

conform: The **shape** of a result is the **frame** of the argument (relative to the **cells** to which the function applies) catenated with the **shape** of the individual results produced by applying the function to the individual **cells**. If the results do not agree in **shape** they are brought to a common **shape** as follows:

If the **ranks** differ, they are brought to a common maximum **rank** by reshaping each individual result to introduce leading unit lengths.

If the individual **shapes** differ (after being brought to a common **rank**), each is brought to a common **shape** by using **take** on each individual result, using as the argument to **take** the **shape** which is the maximum over the **shapes**.

rank-vector: This is the right argument of the rank operator, and it is used to specify the rank of the cells to which the function left argument is to be applied. For the monadic case of a function this is a single value. For the dyadic case of a function it is in general a two-element vector, the first element specifying the rank of the cells of the left argument, and the second element specifying the rank of the cells of the right argument to which the function is to be applied. If only a single value is supplied for the dyadic case of a function, it is used to specify the cells of both the left and right arguments to which the function will be applied. In general, for an arbitrary function, the rank vector is a three-element vector, with the elements referring to the monadic rank, and the two dyadic ranks, in order. No matter whether a one-, two-, or three-element vector \mathbf{v} is supplied as the rank vector, the expression $\phi \circ \phi$ produces a canonical rank-vector of three elements, in which the first element defines the monadic rank, the second element defines the dyadic left rank, and the third element defines the dyadic right rank.

9.3.4 Rank operator deriving monadic function

$$Z \leftarrow \mathbf{f} \overset{\dots}{\circ} \mathbf{y} B$$

Informal Description: The result of $\mathbf{f} \circ \mathbf{y}$ is a function which, when applied to B, returns B, the result of applying the function \mathbf{f} to the **rank-y cells** of B.

Evaluation Sequence:

If y is a scalar, set y1 to y. Otherwise set y1 to y.

If y1 is not a vector, signal domain-error.

If y1 has more than three elements, signal length-error.

If any element of y1 is not a near-integer, signal domain-error.

Set y2 to $\phi 3\rho \phi y1$.

Set y3 to the first-item in y2.

Set y4 to the integer-nearest-to y3.

If y4 exceeds the rank of B, set y5 to the rank of B, otherwise set y5 to y4.

If y5 is **negative**, set y6 to $0 \lceil y5 \rceil$ plus the rank of B, otherwise set y6 to y5.

Apply **f** to the **rank-y6 cells** of *B*.

Conform the individual result **cells**. Let their common **shape** after conforming be \mathbf{q} , and let \mathbf{p} be the **frame** of B with respect to \mathbf{f} , that is, (**rank** of B) **minus** $\mathbf{y6}$, and return the overall result with **shape** \mathbf{p} , \mathbf{q} .

9.3.5 Rank operator deriving dyadic function

 $Z \leftarrow A \mathbf{f} \overset{\dots}{\circ} \mathbf{v} B$

Informal Description: The result of $\mathbf{f} \circ \mathbf{y}$ is a function which, when applied dyadically between A and B, returns Z, the result of applying the function \mathbf{f} between the **rank-l cells** of A and the **rank-r cells** of B, where \mathbf{l} and \mathbf{r} are given by \mathbf{y} .

Evaluation Sequence:

If y is a scalar, set y1 to ,y. Otherwise set y1 to y.

If **y1** is not a **vector**, signal **domain-error**.

If y1 has more than three elements, signal length-error.

If any element of y1 is not a near-integer, signal domain-error.

Set each element of y2 to the **integer-nearest-to** each corresponding element of y1.

Set y3 to $\phi 3\rho \phi y2$.

Set y4 to the second element of y3.

Set y5 to the third element of y3.

If y4 exceeds the rank of A, set y6 to the rank of A, otherwise set y6 to y4.

If y5 exceeds the rank of B, set y7 to the rank of B, otherwise set y7 to y5.

If **y6** is **negative**, set **y8** to the maximum of **zero** and **y6 plus** the **rank** of A, otherwise set **y8** to **y6**.

If **y7** is **negative**, set **y9** to the maximum of **zero** and **y7 plus** the **rank** of *B*, otherwise set **v9** to **v7**.

Set **y10** to the **shape** of *A* with the last **y8** items removed.

Set y11 to the **shape** of B with the last y9 items removed.

Execute the indicated action for whichever of the following cases is true:

Case y10 and y11 are empty:

Set A1 to A and B1 to B.

Case **y10** is empty and **y11** is nonempty:

Set A1 to A **conform** to B and B1 to B.

Case **y10** is nonempty and **y11** is empty:

Set A1 to A and B1 to B conform to A.

Case **y10** and **y11** are nonempty:

If **y10** is not the same length as **y11**, signal a **rank-error**.

If y10 does not match y11, signal a length-error.

Set A1 to A and B1 to B.

Apply **f** between each **rank-y8 cell** of *A*1 and each **rank-y9** cell of *B*1.

Conform the individual result **cells** to give the overall result.

```
0 1 2φ°0 1 'ABC'
ABC
BCA
CAB
     2 2 2T°1 0 N5
0 0 1
0 1 0
0 1 1
1 0 0
1 0 1
     N34+°2 N234
122 124 126 128
142 144 146 148
162 164 166 168
222 224 226 228
242 244 246 248
262 264 266 268
     N3, °1 N34
1 2 3 11 12 13 14
1 2 3 21 22 23 24
1 2 3 31 32 33 34
N2, 00 2 N34
1 11 12 13 14
1 21 22 23 24
1 31 32 33 34
2 11 12 13 14
2 21 22 23 24
2 31 32 33 34
```

10 Mixed Functions

10.1 Monadic Mixed Functions

10.1.1 Roll

 $Z \leftarrow ? B$

Informal Description: Z, for B a scalar integer, is **index-origin** plus R, where R is an integer selected pseudorandomly from the set of all nonnegative integers less than B. Each integer in the set has an equal chance of being selected. Uses **index-origin**. Uses and sets **random-link**.

Evaluation Sequence:

If index-origin is nil, signal implicit-error.

If random-link is nil, signal implicit-error.

Let J stand for **item** I of the **ravel-list** of B.

If B is not scalar, return a numeric array Z such that the shape-list of Z is the shape-list of B and, for all I in the index-set of the ravel-list of Z, item I of the ravel-list of Z is ?J.

If *B* is not a **near-integer**, signal **domain-error**.

Set B1 to the **integer-nearest-to** B.

If *B*1 is not a **positive-integer**, signal **domain-error**.

Using the **implementation-algorithm pseudorandom-number-generator**, whose only inputs are B1 and **random-link**, set Z0 to a **nonnegative-integer less-than** B1, then set **random-link** to a new value.

Return Z0 plus index-origin.

```
S\leftarrow \square RL
?4p1E30
1.69584E29 2.5949E29 5.38382E29 8.38274E29
\square RL \leftarrow S
?4p1E30
1.69584E29 2.5949E29 5.38382E29 8.38274E29
(?1E20 1E20)÷1E20
0.218959 0.678865
```

Additional Requirement:

The operation of **roll** is **atomic**: If **roll** signals an error, **random-link** shall be unchanged.

The result of ?B, where B is an **array**, shall be reproducible.

A **conforming-implementation** shall provide documentation describing the properties of its **pseudorandom-number-generator**.

Note: One class of appropriate algorithms is Lehmer's linear congruential method, described in Knuth, D. E., **Seminumerical Algorithms**, page 9.

Roll is often considered a scalar function. However, it does not have the property that the elements of its result array can be produced in parallel.

10.1.2 Grade Up

 $Z \leftarrow A B$

Informal Description: Z is, for B a vector, a permutation of $\mathfrak{1} \triangleright B$ for which $B \lceil Z \rceil$ is a monotone increasing sequence. The indices of identical elements of B occur in Z in ascending order. If B is a matrix, A B grades the rows, that is, it grades the base value of the rows, using a base larger than the magnitude of any of the elements. Higher rank arguments are graded as if their major cells were ravelled. Uses **index-origin**.

Evaluation Sequence:

If index-origin is nil, signal implicit-error.

If the **rank** of *B* is **zero**, signal **rank-error**.

If ρB is **zero**, return 10.

If any **item** of the **ravel-list** of B is not a **near-real number**, signal **domain-error**.

Otherwise, set B1 to an array having the same shape as B, and with each element having the value of the **real-number** nearest to that of the corresponding element of B.

If the **first-item** of ρB is **one**, return a **one-element-vector** Z such that the **type** of Z is **numeric** and the **ravel-list** of Z contains **index-origin**.

Set *K* to the length of the first axis of *B*1.

If *B*1 is a **vector**, set *B*2 to *B*1.

Otherwise,

Set *M* to $1+2\times\lceil/\mid B1$.

Set B2 to a list with K items, where item I is determined as follows:

Set B3 to the ravel-list of subarray I along the first axis of B1. Set B2[I] to $M \perp B3$.

Generate Z1, a permutation of ιK such that for I and J elements of ιK for which I is **less-than** J,

B2[Z1[I]] is not greater-than B2[Z1[J]] and

B2[Z1[I]] equals B2[Z1[J]] implies that Z1[I] is less-than Z1[J].

Generate Z, a **numeric array** with the **shape-list** of Z1 such that each **item** of the **ravel-list** of Z is (**index-origin minus one**) **plus** the corresponding element of Z1. Return Z.

Examples:

```
AV \leftarrow 1.1 \ 3.1 \ 1.1 \ 2.1 \ 5.1 \ 1 \ 3 \ 4 \ 2 \ 5 \ AAV \ 1 \ 4 \ 2 \ 3 \ 5 \ V[AV] \ 1.1 \ 1.1 \ 2.1 \ 3.1 \ 5.1 \ B \ 3 \ 1 \ 4 \ 2 \ 7 \ 9 \ 3 \ 2 \ 0 \ 3 \ 1 \ 4 \ AB \ 2 \ 1 \ 4 \ 3
```

- In the following example, index-origin is zero.

Additional Requirement:

The system-parameter comparison-tolerance is not an implicit argument of grade up.

10.1.3 Grade Down

$$Z \leftarrow V B$$

Informal Description: Grade down is like grade up except that the major cells are in descending order.

Evaluation Sequence:

Return $\triangle -B$.

Examples:

- In the following example, index-origin is zero.

Additional Requirement:

The system-parameter comparison-tolerance is not an implicit argument of grade down.

Note: Grade up and grade down are stable sort algorithms because they preserve the relative order of identical elements of B.

One appropriate algorithm for grade up and down appears in Woodrum, L. J., Internal Sorting with Minimal Comparing, IBM System Journal, Vol. 8, No. 3, p.189, 1969.

10.1.4 Reverse

```
Z \leftarrow \phi B
Z \leftarrow \Theta B
Z \leftarrow \phi [K] B
Z \leftarrow \Theta [K] B
```

Informal Description: Z is an array whose elements are those of B taken in reverse order along a specified axis. Uses **index-origin**

Evaluation Sequence:

```
For form φB

If B is scalar, return B.

Otherwise, return φ[ρρB] B, evaluated with index-origin set to one.

For form φB

If B is scalar, return B.

Otherwise, return φ[1] B, evaluated with index-origin set to one.

For forms φ[K] B and φ[K] B

If K is not a valid-axis for B, signal axis-error.

Otherwise, set K1 to the integer-nearest-to K.

If B is a vector, return B[(1+ρB)-1ρB], evaluated with index-origin set to one.

Otherwise, return an array Z, such that the type of Z is the type of B, the shape-list of Z is the shape-list of B, and the ravel-list of Z has the property that each vector-item along-axis K1 of Z is φ applied to the corresponding vector-item
```

Examples:

```
φN23
13 12 11
23 22 21

φ[2] N224
121 122 123 124
111 112 113 114

221 222 223 224
211 212 213 214
```

along-axis K1 of B.

Note: Various error checks have been performed on K by the phrase evaluators, and **index-origin** is **one**, before this evaluation sequence is called.

10.1.5 Monadic Transpose

$$Z \leftarrow \lozenge B$$

Informal Description: Z is B with the order of the axes reversed.

Evaluation Sequence:

Return $(\phi \iota \rho \rho B) \Diamond B$, evaluated with **index-origin** set to **one**.

Examples:

Note: This subsection contains a forward reference to dyadic transpose.

10.1.6 Matrix Inverse

 $Z \leftarrow \boxminus B$

Informal Description: Z is the result of applying a generalisation of the matrix inverse function to B. **Matrix inverse** is **matrix divide** with an appropriate identity matrix as a left argument.

Evaluation Sequence:

If the rank of B is greater-than two, signal rank-error. Return $((2\rho 1\rho \rho B)\rho 1, (1\rho \rho B)\rho 0) \exists B$

Note: This subsection contains a forward reference to matrix divide.

The following article describes the motivation for matrix inverse and an acceptable algorithm.

Jenkins, M. A., Domino—An APL Primitive Function for Matrix Inversion—Its Implementation and Applications. APL Quote-Quad Vol III No. 4, February 1972, pp. 4-15.

10.1.7 Execute

```
Z \leftarrow \Phi B
```

Informal Description: Z is the result of evaluating the **character** scalar or vector B as a line of APL.

Evaluation Sequence:

If the rank of B is greater-than one, signal rank-error.

If any **item** of the **ravel-list** of *B* is not a **character**, signal **domain-error**. Generate a new context in which

mode is execute.

current-line is the **ravel-list** of B,

current-function is 0 0p ' ',

current-line-number is one,

current-statement is the empty list of tokens, and

stack is the empty list of tokens.

Append the new context to the state-indicator of the active-workspace as a new first item.

Set Z to evaluate-line.

Remove the first **context** from the **state-indicator**.

Return Z.

Examples:

```
⊈¹T←3¹
3
    3
    A+•!!
value-error
```

Note: If an error is signalled during execute, the user should be able to determine from information provided by the system where the error occurred in the argument of execute as well as where the failing execute primitive occurred in the immediate-execution or defined-function line.

10.1.8 Unique

```
Z \leftarrow \cup B
```

Informal Description: B must be a **vector** or **scalar**. Z is a vector of the same **type** as B, containing one copy of each unique **item** in B. The order of the items in Z is the order they first occur in the **ravel** of B. Uses **comparison-tolerance**.

Evaluation Sequence:

```
If comparison-tolerance is nil signal implicit-error.
If the rank of B is greater-than one, signal rank-error.
Set B1 to an empty-list of the type of B.
Set B2 to the ravel-list of B.
Repeat:
If B2 is empty, return B1.
Set T to the vector-item one of B2.
Append T to B1.
Remove from B2 all items tolerantly-equal-to T, within comparison-tolerance. (End of Repeated Block).
```

```
U2 7 1 8 2 8 1 8 2 8 4 5 9 0 4 4 9 2 7 1 8 4 5 9 0 U'MISSISSIPPI' MISP
```

10.1.9 First

```
Z \leftarrow \uparrow B
```

Informal Description: Z is the first item of B in row-major order. If B is empty, the **typical-element** of B is returned.

Evaluation Sequence:

Return **first-thingy** of *B*.

Examples:

10.2 Dyadic Mixed Functions

10.2.1 Join Along an Axis

```
Z \leftarrow A, [K] B
```

Informal Description: Z is formed by joining A and B along a designated axis. There are two suboperations, catenate and laminate. The catenate suboperation joins the arrays along an existing axis, the laminate suboperation along a new axis. The choice of axis and the choice of operation is determined by K: if K is a **near-integer**, the operation is catenate and the axis is K; if K is not a **near-integer**, the operation is laminate, the new axis is K, and the axes greater than or equal to K are renumbered. Uses **index-origin**.

Evaluation Sequence:

```
If K is not a near-integer,
If A is a scalar, set A1 to (\rho B)\rho A.
Otherwise, set A1 to A.
If B is a scalar, set B1 to (\rho A)\rho B.
Otherwise, set B1 to B.
```

If the **rank** of A1 differs from the **rank** of B1, signal **rank-error**.

If K is not in the **open-interval-between zero** and (**one plus** the **rank** of A1), signal **axis-error**.

If the **shape-list** of *A*1 differs from the **shape-list** of *B*1, signal **length-error**.

Set T to $(1, \rho A1) [AK, \iota \rho \rho A1]$.

Set A2 to $T \circ A1$.

Set B2 to $T \circ B1$.

Return A2, $[\lceil K \rceil]$ B2, evaluated with **index-origin** set to **one** and with **comparison-tolerance** set to **integer-tolerance**.

If *K* is a **near-integer**, evaluate the following with **comparison-tolerance** set to **zero**:

If both A and B are scalars, signal axis-error.

Set K1 to the **integer-nearest-to** K.

If K1 is not in the **closed-interval-between one** and the larger of the **rank** of A and the **rank** of B, signal **axis-error**.

If A is a scalar, set A1 to $((\rho B) * K1 \neq \iota \rho \rho B) \rho A$

Otherwise, set A1 to A.

If B is a scalar, set B1 to $((\rho A) * K1 \neq \iota \rho \rho A) \rho B$.

Otherwise, set B1 to B.

If the rank of B1 minus the rank of A1 is one,

set A2 to $(1, \rho A1) [AK1, \iota \rho \rho A1] \rho A1$

Otherwise, set A2 to A1.

If the rank of A1 minus the rank of B1 is one,

set B2 to $(1, \rho B1) [AK1, \iota \rho \rho B1] \rho B1$

Otherwise, set B2 to B1.

If the **rank** of A2 differs from the **rank** of B2, signal **rank-error**.

If any **axis** of A2 except **axis** K1 differs from the corresponding **axis** of B2, signal **length-error**.

If B2 is empty, return $((\rho A2) + (\rho B2) \times K1 = \iota \rho \rho B2) \rho A2$

If A2 is empty, return $((\rho A2)+(\rho B2)\times K1=1\rho\rho B2)\rho B2$

If A2 and B2 are **vectors**, return A2, B2.

Otherwise, return Z, an **array** such that the

shape-list of Z is $(\rho A2) + (\rho B2) \times K1 = \iota \rho \rho B2$ and the **ravel-list** of Z has the property that each **vector-item along-axis** K1 of Z is A3, B3, where A3 is the corresponding **vector-item along-axis** K1 of A2 and B3 is the corresponding **vector-item along-axis** K1 of B2. The **type** of Z is the **sufficient-type** of the **ravel-list** of Z under the **type mixed**.

```
<u>-</u>M←2 3ρ'Δ'
\Delta\,\Delta\,\Delta
\Delta\,\Delta\,\Delta
          <u>-</u>H←3 3ρ'ο'
000
000
000
          M,[1]H
\Delta\,\Delta\,\Delta
\Delta\,\Delta\,\Delta
000
000
000
          □←L←2 4ρ'□'
M, L
\Delta \Delta \Delta \Box \Box \Box \Box
\Delta \Delta \Delta \Box \Box \Box \Box
         M, '+'
\Delta\Delta\Delta +
\Delta\Delta\Delta+
         M, '34'
ΔΔΔ3
\Delta\Delta\Delta4
          M, [1] '345'
\Delta\,\Delta\,\Delta
\Delta\,\Delta\,\Delta
345
          1 2 3,[.5] 4 5 6
1 2 3
4 5 6
          1 2 3,[1.5]4 5 6
1 4
2 5
3 6
          1 2 3, [1.5] 4
1 4
2 4
3 4
          (2 0 \rho 5), 'A'
Α
Α
          ρ3,[.5]''
2 0
```

Note: Various error checks have been performed on K by the phrase evaluators, and **index-origin** is **one**, before this evaluation sequence is called.

10.2.2 Index of

```
Z \leftarrow A \iota B
```

Informal Description: Z is a numeric array of shape ρB . Each element of Z is the least index in A of a value **tolerantly-equal** within **comparison-tolerance** to the corresponding item of B. Uses **index-origin**. Uses **comparison-tolerance**.

Evaluation Sequence:

If index-origin is nil, signal implicit-error.

If comparison-tolerance is nil, signal implicit-error.

If A is not a **vector**, signal **rank-error**.

Set Z to $+/\land \ B \circ . \ne A$, evaluated with the current value of **comparison-tolerance**.

Return Z plus index-origin.

Examples:

Note: In the following example, index-origin is zero.

```
□←A←2 2ρ1.1 3.1 5.1 4.1
1.1 3.1
5.1 4.1
3.1 4.1 5.11A
3 0
2 1
'ABC'13
```

Note: In the following example, index-origin is one.

10.2.3 Member of

$$Z \leftarrow A \in B$$

Informal Description: Z is a Boolean array with the shape of A. An element of Z is **one** if the corresponding element of A is **tolerantly-equal** to some element of B; otherwise, it is **zero**. Uses **comparison-tolerance**.

Evaluation Sequence:

If comparison-tolerance is nil, signal implicit-error.

Return $\vee/A \circ .=$, B, evaluated with the current value of **comparison-tolerance**.

Examples:

Note: *If* B *is empty, the result of* $A \in B$ *is* $(\rho A) \rho O$.

10.2.4 Deal

 $Z \leftarrow A ? B$

Informal Description: Z is **index-origin** plus R, where R is a vector of shape A obtained by making A pseudorandom selections without replacement from the set of nonnegative integers less than B. Uses **index-origin**. Uses and sets **random-link**.

Evaluation Sequence:

If **index-origin** is **nil**, signal **implicit-error**.

If random-link is nil, signal implicit-error.

If the rank of A or the rank of B is greater-than one, signal rank-error.

If either of A or B is neither a scalar nor a one-element-vector, signal length-error.

If either of A or B is not a **near-integer**, signal **domain-error**.

Set A1 to the **integer-nearest-to** A.

Set B1 to the **integer-nearest-to** B.

If either of A1 or B1 is not a **nonnegative-number**, signal **domain-error**.

If A1 is **greater-than** B1, signal **domain-error**.

If A1 is **zero**, return 10.

Using the **implementation-algorithm deal**, generate a **numeric vector** Z0 of **length** A1 whose elements are selected in a pseudorandom fashion without duplication from the **integers** in the **closed-interval-between zero** and B1 **minus one**, then set **random-link** to a new value.

Return Z0 plus index-origin.

Example:

```
12?300
2 3 42 94 70 105 215 9 110 298 201 5
```

Additional Requirement:

The selection of elements in Z is pseudorandom (see the operation **roll**). The elements of Z, their order, and the new value of **random-link** are determined by an **implementation-algorithm** whose only inputs are A1, B1 and **random-link**. The operation of **deal** is **atomic**: if **deal** signals an error, **random-link** shall be unchanged. The result of A?B, where A and B are **arrays**, shall be reproducible.

10.2.5 Replicate

 $Z \leftarrow A / B$ $Z \leftarrow A \neq B$ $Z \leftarrow A / [K] B$ $Z \leftarrow A \neq [K] B$

Informal Description: Z is an array formed by replicating the subarrays, along a specified axis of B, the number of times indicated by the corresponding element of the **near-integer vector** or **scalar** A. Uses **index-origin**.

Evaluation Sequence:

If A is a **scalar**, set A1 to , A. Otherwise, set A1 to A.

If B is a **scalar**, set B1 to $(\rho A1)\rho B$.

Otherwise, set B1 to B.

For form A/B

Return $A1/[\rho\rho B1]$ B1 evaluated with **index-origin** set to **one**.

For form $A \neq B$

Return A1/[1] B1 evaluated with **index-origin** set to **one**.

For forms A / [K] B and $A \neq [K]$ B

If *K* is not a **valid-axis** for *B*1, signal **axis-error**.

Otherwise, set *K*1 to the **integer-nearest-to** *K*.

If the rank of A1 is greater-than one, signal rank-error.

If the **count** of A1 is **one**, set A2 to $(\rho B1)[K1]\rho A1$.

Otherwise, set A2 to A1.

If $\rho A2$ is not the same as $(\rho B1)$ [K1], signal **length-error**.

If any **item** of the **ravel-list** of A2 is not a **near-integer**, signal **domain-error**.

Set A3 to the **integer-array-nearest-to** A2.

If any item of the ravel-list of A3 is not a nonnegative-integer, signal domain-error.

If *B*1 is a **vector**,

If every **item** of A3 is **near-Boolean** return B1 [$(+/A3) \rho \sqrt[4]{A3}$].

Otherwise, return $B1[(,A3\circ.\geq\iota[/0,A3)/,(\iota\rho B1)\circ.+0\times\iota[/0,A3]$

Otherwise, return an **array** Z such that

The **shape-list** of Z is $((K1 \neq \iota \rho \rho B1) \times \rho B1) + (K1 = \iota \rho \rho B1) \times +/A3$ evaluated with **comparison-tolerance** set to **zero**,

The **ravel-list** of Z has the property that if Z2 and B2 are corresponding **vector-items along-axis** K1 of Z and B1 respectively, then Z2 is A3/B2.

The **type** of Z is the **sufficient-type** of the **ravel-list** of Z under the **type** of B.

Examples:

```
101/123
1 3
     1 / 1 2 3
1 2 3
     3 2 1/1 2 3
1 1 1 2 2 3
     1 0 1/ 2
2 2
     0 0 1 0 0 1 0 /[2] N2714
1311 1312 1313 1314
1611 1612 1613 1614
2311 2312 2313 2314
2611 2612 2613 2614
     \rho 1/1
     \rho\rho(,1)/2
1
     3 4/1 2
1 1 1 2 2 2 2
```

Note: Various error checks have been performed on K by the phrase evaluators, and **index-origin** is **one**, before this evaluation sequence is called.

Note: *If the left argument is boolean, the function may be called* **compress**.

Note: 1/A for a scalar A, returns a vector.

10.2.6 Expand

```
Z \leftarrow A \setminus B
Z \leftarrow A \setminus B
Z \leftarrow A \setminus [K] B
Z \leftarrow A \setminus [K] B
```

Informal Description: Z is, for a **near-Boolean** scalar or vector A, an array of the same type as B containing subarrays of B along a specified axis.

For vector B, Z is such that $A/A \setminus B$ is B and, if P is the **typical-element** of B, $(\sim A)/A \setminus B$ is $(+/\sim A) \cap P$, similarly, $A/[K]A \setminus [K]B$ is B, $A \neq A \setminus B$ is B, and $(\sim A)/[K]A \setminus [K]B$ and $(\sim A) \neq A \setminus B$ are arrays of the **typical-element** of B. Uses **index-origin**.

Evaluation Sequence:

If the rank of A is greater-than one, signal rank-error. If any **item** of the **ravel-list** of *A* is not a **near-Boolean**, signal **domain-error**. Set A1 to the **Boolean-array-nearest-to**, A. If B is a scalar, set B1 to $(+/A1) \rho B$. Otherwise, set B1 to B. For form $A \setminus B$ Return $A1 \setminus [\rho \rho B1]$ B1 evaluated with **index-origin** set to **one**. For form A + BReturn $A1 \setminus [1]$ B1 evaluated with **index-origin** set to **one**. For forms $A \setminus [K]$ B and A + [K] B If *K* is not a **valid-axis** for *B*1, signal **axis-error**. Otherwise, set *K*1 to the **integer-nearest-to** *K*. If $(\rho B1)[K1]$ differs from +/A1, signal **length-error**. Return an **array** Z such that the **type** of Z is the **type** of B, the **shape-list** of Z is $((K1 \neq \iota \rho \rho B1) \times \rho B1) + (K1 = \iota \rho \rho B1) \times \rho A1$, evaluated with comparison-tolerance set to zero, and the **ravel-list** of Z has the property that A1/[K1]Z is B1 and $(\sim A1)/[K1]Z$ is an **array** consisting entirely of the **typical-element** of B.

Note: Various error checks have been performed on K by the phrase evaluators, and **index-origin** is **one**, before this evaluation sequence is called.

```
1 0 1 \1 3
1 0 3
     1 1 1 0 1 \'ABCD'
ABCD
     1 0 1\2
2 0 2
     1 0 1\[2] 2 2 ρ'ABCD'
AB
CD
     1 0 1 1\1 2 3
1 0 2 3
     1 0 1 1\3
3 0 3 3
     0 1 \3 1 ρ3.14 2E17 <sup>-</sup>47
0 3.14E0
0 2.00E17
0^{-4.70E1}
     0 0\5
0 0
    1 0 1 0\[2] N224
111 112 113 114
 0 0 0 0
121 122 123 124
 0 0 0 0
211 212 213 214
 0 0 0 0
221 222 223 224
 0 0 0
```

10.2.7 Rotate

 $Z \leftarrow A \Leftrightarrow B$ $Z \leftarrow A \Leftrightarrow B$ $Z \leftarrow A \Leftrightarrow [K] B$ $Z \leftarrow A \Leftrightarrow [K] B$

Informal Description: Z is an array of the same type and shape as B in which elements have been shifted cyclically along a specified axis. The amount and direction of shift is controlled by A. Uses **index-origin**.

Evaluation Sequence:

For form $A \phi B$

If B is a scalar and A is either a scalar or a vector of length one,

If any **item** of the **ravel-list** of A is not a **near-integer**, signal **domain-error**. Otherwise, return B.

Otherwise, return $A \phi [\rho \rho B]$ B evaluated with **index-origin** set to **one**.

For form $A \ominus B$

If B is a scalar and A is either a scalar or a vector of length one,

If any **item** of the **ravel-list** of A is not a **near-integer**, signal **domain-error**. Otherwise, return B.

Otherwise, return $A \phi [1]$ B evaluated with **index-origin** set to **one**.

For forms $A \Leftrightarrow [K] B$ and $A \Leftrightarrow [K] B$

If *K* is not a **valid-axis** for *B*, signal **axis-error**.

Otherwise, set *K*1 to the **integer-nearest-to** *K*.

If A is a scalar, set A1 to $((K1 \neq \iota \rho \rho B)/\rho B)\rho A$, evaluated with comparison-tolerance set to zero.

Otherwise, set A1 to A.

If A1 is a **one-element-vector** and B is a **vector**, set A2 to $(10) \rho A1$.

Otherwise, set A2 to A1.

If the rank of B minus the rank of A2 is not one, signal rank-error.

If the **shape-list** of A2 is not the same as the **shape-list** of B with **axis** K1 omitted, signal **length-error**.

If any **item** of the **ravel-list** of A2 is not a **near-integer**, signal **domain-error**.

Set A3 to the **integer-array-nearest-to** A2.

If A3 is a scalar and B is a vector, return $B[1+(\rho B)]^{-1}+A3+\iota\rho B$] evaluated with **comparison-tolerance** set to **zero**.

Otherwise return Z, an **array** such that the **shape-list** of Z is the same as the **shape-list** of B, the **type** of Z is the same as the **type** of B, and the **ravel-list** of Z has the property that if Z0 is a **vector-item along-axis** K1 of Z, A0 is the corresponding **item** of A3, and B0 is the corresponding **vector-item along-axis** K1 of B, then Z0 is A0 ΦB 0.

Examples:

```
3 4 5
4 5 1 2 3
     <sup>-</sup>1¢1 2 3 4 5
5 1 2 3 4
     _7φ'ABCDEF'
FABCDE
     10N33
21 22 23
31 32 33
11 12 13
     1¢[1]N33
21 22 23
31 32 33
11 12 13
     1 2 3φN34
12 13 14 11
23 24 21 22
34 31 32 33
     N23\phi[2]N243
141 112 123
111 122 133
121 132 143
131 142 113
221 232 243
231 242 213
241 212 223
211 222 233
```

Note: Various error checks have been performed on K by the phrase evaluators, and **index-origin** is **one**, before this evaluation sequence is called.

10.2.8 Base Value

```
Z \leftarrow A \perp B
```

Informal Description: Z is, for A and B numeric vectors, a number produced by regarding B as the representation of a number in the mixed radix number system specified by A. If A or B is an array of rank greater than **one** the value is like inner product: each **vector-item** along the last **axis** of A is applied to each **vector-item** along the first axis of B.

Evaluation Sequence:

```
Set A0 to 1\rho \varphi 1, \rho A.

Set B0 to 1\rho (\rho B), 1.

If A0 differs from B0,

If A is a scalar or one-element-vector, return (B0\rho A) \bot B.

If B is a scalar or one-element-vector, return A\bot A0\rho B.

Otherwise, signal length-error.

If A0 is the same as B0,

If any item of the ravel-list of A is not a number, signal domain-error.

If any item of the ravel-list of B is not a number, signal domain-error.

Return (\varphi (\varphi A) \rho \times \varphi , \varphi A, 1) + \cdot \times B.
```

```
1011 2 3
123
      24 60 60 11 2 3
3723
      □←A←2 3ρ10 10 10 12 60 60
10 10 10
12 60 60
      □<del><</del>B←3 2ρ1 4 2 5 3 6
1 4
2 5
3 6
      A \perp B
 123
        456
3723 14706
       -.001 10 10<sub>1</sub>1 2 3
123
      60 11 2 3
3723
      1113
0
      'A'⊥10
0
```

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ISO/IEC 13751:2001(E)

Note: The shape requirements of A and B are intentionally stricter in this definition than in several existing systems. Specifically, this International Standard does not require that if the last axis of A is one it be replicated to match the first axis of B, or that if the first axis of B is one it be replicated to match the last axis of A.

The first element of each left argument has no actual effect on the result of base value, but permits use of the same left argument for base value and representation. If A is a positive numeric vector and B a non-negative numeric scalar such that $B<\times/A$, then $A\perp A\top B$ is B.

10.2.9 Representation

```
Z \leftarrow A \top B
```

Informal Description: Z is the representation of B in mixed radix number system A.

Evaluation Sequence:

```
If A or B is empty, return ((\rho A), \rho B) \rho 0.
If any item of the ravel-list of A is a character, signal domain-error.
If any item of the ravel-list of B is a character, signal domain-error.
If A is scalar, return A \mid B, evaluated with comparison-tolerance set to zero.
If A is a vector and B is a scalar.
  Generate two numeric vectors Z and C that satisfy the following constraints:
     The length of Z is \rho A.
     The length of C is 1+\rho A.
     C[1+\rho A] is B.
     For all scalar indices I in \iota \rho A:
       Z[I] is A[I] + C[I+1].
       If A \lceil I \rceil is zero, C \lceil I \rceil is zero;
       Otherwise, C[I] is (C[I+1]-Z[I]) \div A[I].
  Return Z.
Otherwise, return Z1, an array such that the type of Z1 is numeric, the shape-list of
 Z1 is (\rho A), \rho B and the ravel-list of Z1 has the following property:
  Let I stand for an item of the index-set of the ravel-along-axis one of A.
  Let A1 stand for vector-item I of the ravel-along-axis one of A.
  Let J stand for an item of the index-set of the ravel-list of B.
  Let B1 stand for item J of the ravel-list of B.
  Let N stand for the count of B.
  Let P stand for J+(N\times(I-1)).
  Then, vector-item P of the ravel-list of Z1 along-axis one is A1 TB1.
```

Examples:

```
10 10 10T123
1 2 3
      10 10 10T123 456
1 4
2 5
3 6
      A←11 3ρ10 16 2
      '0123456789ABCDEF'[\dagged1+A\tau1000 1024]
0000001000
000000003E8
01111101000
0000001024
00000000400
10000000000
      2 2 2 T<sup>-</sup>1
1 1 1
      0 2 2 T<sup>-</sup>1
<sup>-</sup>1 1 1
     0 1 T3.75 - 3.75
     -4
3
0.75 0.25
```

Note: The shape of the result of **representation** is always (ρA) , ρB .

 ${\it The} \ {\bf system-parameter} \ {\bf comparison-tolerance} \ is \ not \ an \ implicit \ argument \ of \ representation.$

10.2.10 Dyadic Transpose

 $Z \leftarrow A \diamond B$

Informal Description: Z is an array formed by rearranging and possibly coalescing the axes of B according to the vector A. Each element of A corresponds to an axis of B by position and to an axis of B by value. The largest value in A determines the rank of B. All axes of B must be present in B. If B contains no repeated elements, the shape of B is B. Repeated elements in B select diagonals from B. The length of the corresponding axis of B is the shortest of the lengths of the designated axes of B. Uses **index-origin**.

Evaluation Sequence:

If index-origin is nil, signal implicit-error.

If A is a **scalar**, set A1 to , A.

Otherwise, set A1 to A.

If A1 is not a **vector**, signal **rank-error**.

If the **length** of A1 is not the same as the **rank** of B, signal **length-error**.

If any item of the ravel-list of A1 is not a near-integer, signal domain-error.

Set A2 to the integer-array-nearest-to A1. Generate A3, a numeric array with the shape-list of A2 such that each item of the ravel-list of A3 is (one minus index-origin) plus the corresponding element of A2.

If $\wedge/A3 \in 1 \lceil/0$, A3 evaluated with **comparison-tolerance** set to **zero** is not **one**, signal **domain-error**.

If $\land/(\iota \lceil /0, A3) \in A3$ evaluated with **comparison-tolerance** set to **zero** is not **one**, signal **domain-error**.

Return an **array** Z having the following properties:

The rank of Z is $\lceil / 0$, A3.

The **shape-list** of Z is such that for all I in the **index-set** of the **shape-list** of Z, **item** I of the **shape-list** of Z is $\lfloor / (A3=I)/\rho B$, evaluated with **comparison-tolerance** set to **zero**.

The **ravel-list** of Z is such that, for all J in the **index-set** of the **ravel-list** of Z, **item** J of the **ravel-list** of Z is **item** $1+(\rho B)\perp((\rho Z)\top J-1)[A3]$ of the **ravel-list** of B.

The **type** of Z is the **sufficient-type** of the **ravel-list** of Z under the **type** of B.

```
1 3 2\dag{N}234
111 121 131
112 122 132
113 123 133
114 124 134
211 221 231
212 222 232
213 223 233
214 224 234
     1 1¢N34
11 22 33
    3 1 2\dag{N}234
111 211
112 212
113 213
114 214
121 221
122 222
123 223
124 224
131 231
132 232
133 233
134 234
     ρρ(ι0)$5
0
```

10.2.11 Take

$$Z \leftarrow A \uparrow B$$

Informal Description: Z is an array having shape , A. Informally, Z is a corner of B. The choice of corner is based on the signs of the elements of A. If the absolute values of elements of A are greater than corresponding elements of A, then A is padded with the **typical-element** of A.

Evaluation Sequence:

If the rank of A is greater-than one, signal rank-error.

Set A1 to A.

If B is a scalar, set B1 to $((\rho A1)\rho 1)\rho B$.

Otherwise, set B1 to B.

If the **shape-list** of A1 does not match the **rank** of B1, signal **length-error**.

If any **item** of the **ravel-list** of A1 is not a **near-integer**, signal **domain-error**.

Set A2 to the **integer-array-nearest-to** A1.

Return an **array** Z such that:

The **shape-list** of Z is A2.

The **ravel-list** of Z has the property that for each **scalar** I in the **index-set** of the **ravel-list** of Z, $i t \in M$ of the **ravel-list** of Z is determined as follows:

Let J stand for $((|A2) \top I - 1) + (A2 < 0) \times A2 + \rho B1$ evaluated with **comparison-tolerance** set to **zero**.

If each element of J is in the **open-interval-between negative-one** and the corresponding element of $\rho B1$, **item** I of the **ravel-list** of Z is $(B1)[1+(\rho B1) \perp J]$.

Otherwise, item I of the ravel-list of Z is the typical-element of B.

The **type** of Z is the **sufficient-type** of the **ravel-of** Z under the **type** of B.

```
21N5

1 2

-21N5

4 5

-2 61N44

31 32 33 34 0 0

41 42 43 44 0 0

-4 4199

0 0 0 0

0 0 0 0

0 0 0 0

0 0 0 0

0 0 0 99

210

0 0
```

10.2.12 Drop

$$Z \leftarrow A \downarrow B$$

Informal Description: Z is a corner of B with shape $0 \lceil (\rho B) - | A$. The choice of corner is based on the signs of the elements of A.

Evaluation Sequence:

```
If the rank of A is greater-than one, signal rank-error.

Set A1 to A.

If B is a scalar, set B1 to ((\rho A1)\rho 1)\rho B.

Otherwise, set B1 to B.

If the shape-list of A1 does not match the rank of B1, signal length-error.

If any item of the ravel-list of A1 is not a near-integer, signal domain-error.

Set A2 to the integer-array-nearest-to A1.

Return (((A2<0)\times0 \Gamma A2+\rho B1)+(A2\geq0)\times0 \Gamma A2-\rho B1) \uparrow B1 evaluated with
```

Examples:

```
1+N5
2 3 4 5
2 -1+N44
31 32 33
41 42 43
01+5
0
000
11+5
5
0
000
11+5
5
0
000
11+5
5
000
11+5
5
000
11+5
5
```

comparison-tolerance set to zero.

10.2.13 Matrix Divide

 $Z \leftarrow A \ominus B$

Informal Description: Z is the least squares solution to a linear system specified by A and B.

Evaluation Sequence:

If the rank of A or the rank of B is greater-than two, signal rank-error.

Set A1 to $(2\uparrow(\rho A), 1 1)\rho A$.

Set B1 to $(2\uparrow(\rho B), 1 1)\rho B$.

If $1 \uparrow \rho A 1$ is not the same as $1 \uparrow \rho B 1$, signal **length-error**.

If $1 \uparrow \rho B 1$ is **less-than** $-1 \uparrow \rho B 1$, signal **length-error**.

If any **item** of the **ravel-list** of A is not a **number**, signal **domain-error**.

If any **item** of the **ravel-list** of B is not a **number**, signal **domain-error**.

Use the **implementation-algorithm matrix-divide** to generate an **array** Z, such that the **type** of Z is **numeric**, the **shape-list** of Z is $(-1 \uparrow \rho B1)$, $-1 \uparrow \rho A1$, and the **ravel-list** of Z has the property that for each **scalar** I in $\iota^-1 \uparrow \rho A1$, Z[;I] is such that it minimises $+/(A1[;I]-B1+.\times Z[;I]) \times 2$

If Z cannot be uniquely determined within a reasonable round-off criterion by the above constraints, signal **domain-error**.

Return $((1 \downarrow \rho B), 1 \downarrow \rho A) \rho Z$.

Additional Requirement:

The round-off criterion discussed above is not specified by this International Standard.

Note: The following article describes an acceptable algorithm for matrix division:

Jenkins, M. A., Domino—An APL Primitive Function for Matrix Inversion—Its Implementation and Applications. APL Quote-Quad Vol III No. 4, February 1972, pp 4-15.

10.2.14 Indexed Reference

```
Z \leftarrow A \lceil I \rceil
```

Informal Description: Z is an array consisting of elements of the array A selected and structured by the position and values of the arrays in the **index-list** I. Uses **index-origin**.

Evaluation Sequence:

```
If the number-of-items in the index-list I does not match the rank of A, signal rank-error.
```

In the following,

Let IX stand for (array) **item** X of I.

Let JX stand for (array) **item** X of J.

Let LX stand for (array) **item** X of L.

Let LY stand for (array) **item** X+1 of L.

Generate \mathcal{I} , an **index-list** having the same **count** as \mathcal{I} such that for every **item** X of the **index-set** of \mathcal{I} ,

If IX is an **elided-index-marker**, set JX to $1+\iota(\rho A)[X]$.

Otherwise,

If any **item** of the **ravel-list** of IX is not a **near-integer**, signal **domain-error**.

Set JX to the **integer-array-nearest-to** IX.

Set JX to JX minus index-origin.

If any element of JX is not an element of $1+\iota(\rho A)[X]$, signal **index-error**.

Set JX to $JX \times \times / X \downarrow \rho A$.

Generate L, an **index-list** with $1+\rho\rho A$ (array) **items** such that

Item $1+\rho\rho A$ of L is **one**.

For all X in $\iota \rho \rho A$, LX is JX $\circ \cdot + LY$.

Let K stand for **first-item** in the **list** of **arrays** L.

Return (A)[K].

Examples:

```
1 2 3[2]
2
N222[2 1;;2]
212 222
112 122
```

Note: Since an **index-list** has at least one **item**, indexing will always signal **rank-error** when argument A is a scalar.

The vector indexing on which this subsection is based is defined in the subsection **evaluate-indexed-reference**.

10.2.15 Indexed Assignment

```
Z \leftarrow V \lceil I \rceil \leftarrow B
```

Informal Description: Z is B. As a side effect, indexed assignment sets elements of the array V selected by the position and values of the arrays in the **index-list** I to corresponding elements of B. Note that, at this stage, V is known to be a **variable-name**, not a **value**. Uses **index-origin**.

Evaluation Sequence:

```
If index-origin is nil, signal implicit-error.
```

Set A to the **current-content** of V.

If the **number-of-items** in the **index-list** I does not match the **rank** of A, signal **rank-error**.

In the following,

Let IX stand for (array) **item** X of I.

Let JX stand for (array) **item** X of J.

Let LX stand for (array) **item** X of L.

Let LY stand for (array) **item** X+1 of L.

Generate J, an **index-list** having the same **count** as I such that for every **item** X of the **index-set** of I.

If IX is an **elided-index-marker**, set JX to $-1+\iota(\rho A)[X]$.

Otherwise,

If any **item** of the **ravel-list** of IX is not a **near-integer**, signal **domain-error**.

Set JX to the **integer-array-nearest-to** IX.

Set JX to JX minus index-origin.

If any element of JX is not an element of $1+\iota(\rho A)[X]$, signal **index-error**.

Set JX to $JX \times \times / X \downarrow \rho A$.

Generate L, an **index-list** with $1+\rho\rho A$ (array) **items** such that

Item $1+\rho\rho A$ of L is **one**.

For all X in $\iota \rho \rho A$, LX is $JX \circ \cdot + LY$.

Let *K* stand for **first-item** in the **list** of **arrays** *L*.

Set K1 to $((1 \neq \rho K)/\rho K)\rho K$ evaluated with **comparison-tolerance** set to **zero**.

Set B1 to $((1 \neq \rho B)/\rho B) \rho B$ evaluated with **comparison-tolerance** set to **zero**.

If B1 is a **scalar**, set B2 to $(\rho K1) \rho B1$.

Otherwise, set B2 to B1.

If the **rank** of *K*1 is not the same as the **rank** of *B*2, signal **rank-error**.

If any **item** of the **shape-list** of *K*1 is not the same as the corresponding **item** of the **shape-list** of *B*2, signal **length-error**.

Set *M* to **one**.

Repeat:

If M is not greater-than the count of K1,

Let *K*2 stand for **item** *M* of the **ravel-list** of *K*1.

Set item *K*2 of the ravel-list of *A* to item *M* of the ravel-list of *B*.

Set *M* to *M* **plus one**.

Otherwise,

Set the **type** of *A* to the **sufficient-type** of the **ravel-list** of *A* under the **type mixed**. Set the **current-referent** of *V* to a **token** whose **class** is **variable** and whose **content** is *A*.

Return a **token** whose **class** is **committed-value** and whose **content** is *B*. (End of repeated block)

Examples:

```
X 2 3
      \square \leftarrow X[3\ 2] \leftarrow 4\ 5
4 5
      X
1 5 4
      111 112
121 122
211 212
221 222
      \square \leftarrow Y[2 1;;1] \leftarrow N12121
11111
11121
12111
12121
12111 112
12121 122
11111 212
11121 222
```

Additional Requirements:

The evaluation sequence requires that after the statement $A[1 \ 1 \ 1] \leftarrow 1 \ 2 \ 3$ has been evaluated the value of A[1] be 3.

Indexed assignment exhibits **atomic** behaviour. If **indexed assignment** signals an error, the value of *V* shall be unchanged.

Note: Since an **index-list** has at least one **item**, indexing will always signal **rank-error** when argument A is a scalar.

10.2.16 Without

$$Z \leftarrow A \sim B$$

Informal Description: Z is a **vector** whose elements are the elements of A without those that are also elements of B. Uses **comparison-tolerance**.

Evaluation Sequence:

If **comparison-tolerance** is **nil**, signal **implicit-error**. If the **rank** of A or B is **greater-than one**, signal **rank-error**. Set A1 to A; set B1 to B. Return $A1 \in B1 / A1$, evaluated with the current value of **comparison-tolerance**.

Examples:

10.2.17 Left

$$Z \leftarrow A \rightarrow B$$

Informal Description: Z is A.

Evaluation Sequence:

Return A.

10.2.18 Right

```
Z \leftarrow A \vdash B
```

Informal Description: Z is B.

Evaluation Sequence:

Return B.

Examples:

10.2.19 Character Grade Definitions

A precedes B in C order: An operation that, for character-vectors A and B and character array C of rank greater-than zero, returns a Boolean result determined by the following evaluation sequence:

Evaluation Sequence:

(end of repeated sequence).

```
Set AX to the rank of C.

Repeat

Set C1 to the array-of-vectors along-axis AX of C.

Generate AC, a numeric vector with the same length as A, such that for all I in the index-set of A, AC[I] is the minimum of V1A[I] over all items V in C1.

Generate BC, a numeric vector with the same length as B, such that for all I in the index-set of B, BC[I] is the minimum of V1B[I] over all items V in C1.

If ((AC<BC)11) < (AC>BC)11, return one.

Set AX to AX-1.

If AX equals zero, return zero.
```

A identifies-with B in C order: An operation that returns zero if A precedes B in C order or B precedes A in C order and otherwise returns one.

10.2.20 Character Grade Down

 $Z \leftarrow A \forall B$

Informal Description: Z is a permutation of the index set of the first axis of B which specifies a stable sort of the character subarrays crossing that axis into a non-increasing sequence according to an ordering relation given in the collating array A. Uses **indexorigin**.

Evaluation Sequence:

If index-origin is nil, signal implicit-error.

Set *IO* to the **numeric-scalar** with value **index-origin minus one**.

If any **item** of the **ravel-list** of *A* is not a **character**, signal **domain-error**.

If any **item** of the **ravel-list** of B is not a **character**, signal **domain-error**.

If A is a **scalar** signal **rank-error**.

If *A* is empty, return $IO + \iota 1 \uparrow \rho B$.

Set B1 to $((1 \uparrow \rho B), \times /1 \downarrow \rho B) \rho B$.

If $1 \uparrow \rho B$ is **zero**, return 10.

If $1 \uparrow \rho B$ is **one**, return a **one-element-vector** Z such that the **ravel-list** of Z contains **index-origin**.

Otherwise, generate Z, a permutation of $\iota 1 \uparrow \rho B$ such that for every pair (I and J) of elements of $\iota 1 \uparrow \rho B$ where I is **less-than** J either:

B1[Z[J];] precedes B1[Z[I];] in A-order or

B1[Z[I];] identifies-with B1[Z[J];] in A-order and Z[I] is less-than Z[J].

Return Z+IO.

10.2.21 Character Grade Up

 $Z \leftarrow A \land B$

Informal Description: Z is a permutation of the index set of the first axis of B which specifies a stable sort of the character subarrays crossing that axis into a non-decreasing sequence according to an ordering relation given in the collating array A. Uses **indexorigin**.

Evaluation Sequence:

If index-origin is nil, signal implicit-error.

Set *IO* to the **numeric-scalar** with value **index-origin minus one**.

If any **item** of the **ravel-list** of *A* is not a **character**, signal **domain-error**.

If any **item** of the **ravel-list** of *B* is not a **character**, signal **domain-error**.

If A is a **scalar** signal **rank-error**.

If *A* is empty, return $IO + \iota 1 \uparrow \rho B$.

Set B1 to $((1 \uparrow \rho B), \times /1 \downarrow \rho B) \rho B$.

If $1 \uparrow \rho B$ is **zero**, return 10.

If $1 \uparrow \rho B$ is **one**, return a **one-element-vector** Z such that the **ravel-list** of Z contains **index-origin**.

Otherwise, generate Z, a permutation of $\iota 1 \uparrow \rho B$ such that for every pair (I and J) of elements of $\iota 1 \uparrow \rho B$ where I is **less-than** J either:

B1[Z[I];] precedes B1[Z[J];] in A-order or

B1[Z[I];] identifies-with B1[Z[J];] in A-order and Z[I] is less-than Z[J].

Return Z+IO.

Note: In comparing rows of B1, differences at low index positions are more significant than those at higher index positions. Thus, if 'a' comes before 'b' in A-order, the rows 'aa', 'ab', 'ba', and 'bb' are in ascending A-order.

Every axis of the collating array A defines an ordering of the character set. A character's lowest index along an axis defines its position in the order described by that axis (low indices precede high indices). Characters not found in A are equal and occur after all characters in A.

The order given along A's last axis is the most significant. The result of character grade groups together rows of B which are equal under this (major) ordering. Within each such group, the result arranges B's rows according to the hierarchy of (minor) orderings along axes of A ending with the first (least significant) axis. Character grade is stable; its results are equal under all the A-orderings in the order that they occur in B.

Hierarchical sorting allows distinctions such as spelling, case, font, and length to carry their proper lexicographic weight. For example, the rows of W (below) are taken in order from Webster's New Collegiate Dictionary.

Examples:

```
(2 2p'ABBA') A 'AB'[?5 2p2] A A AND B ARE EQUIVALENT
1 2 3 4 5
    A+2 14p' abcdegiklmnrt ABCDEGIKLMNRT'
W
            W[(,A) AW;]
                         W[(, \Diamond A) \Delta W;] W[A \Delta W;]
Ab
                                      Ab
            aba
                          aba
AB
            abaca
                                      AB
                          abaca
aba
            abecedarian abecedarian aba
                        Ab
ABA
            black
                                      ABA
abaca
            black belt Abelian
                                     abaca
abecedarian blackball AB
                                      abecedarian
Abelian
          blacking
                        ABA
                                      Abelian
black
            AD
                         black
                                      black
                         black belt black belt
blackball
            Abelian
black belt
                          blackball
                                      Black Mass
            AB
blacking
             ABA
                         blacking
                                      blackball
Black Mass
            Black Mass
                         Black Mass
                                      blacking
```

Note: Character grade is equivalent to the function SRT4 defined in Smith, Howard J., "Sorting – A New/Old Problem", **APL Quote-Quad 9**, 1, June 1979, pp 123-127

10.2.22 Pick

$$Z \leftarrow A \supset B$$

Informal Description: *Z* is an item selected from some level of *B*. *A* is a scalar or vector. The items of *A* are selectors applying to one or more successive levels of *B*, such that the first (or only) item of *A* selects an item from *B*, the next (if any) selects an item from this item, and so on. The length of each item in *A* must agree with the rank of the object to which it is to be applied. Uses **index-origin**.

Evaluation Sequence:

If **index-origin** is **nil**, signal **implicit-error**.

If the rank of A is greater-than one, signal rank-error.

If A is empty, return B.

If the **count** of *a* is **one**:

Set A1 to the **first-thingy** of A.

If A1 is not a **simple array**, signal **domain-error**.

If the rank of A1 is greater-than one, signal rank-error

Set A2 to the **ravel-list** of A1.

If any item in A2 is not a **near-integer**, signal **domain-error**.

Set A3 to a list with each item the **integer-nearest-to** the corresponding element of A2.

If the **number-of-items** in A3 is not the same as the rank of B, signal **rank-error** Set A4 to a list, with each item equal to the corresponding item in A3, minus **index-origin**.

For every member X of the **index-set** of A4:

If item X of A4 is **less-than zero**, or **greater-than-or-equal-to** item X of the **shape-list** of B, signal **index-error**.

Return item $1+(\rho B) \perp A + \text{ of the } \mathbf{ravel-list}$ of B.

Otherwise, return $(1 \downarrow A) \supset (1 \uparrow A) \supset B$.

```
3>'OSCAR'
C
(,3)>10 11 12 13
12
(<3 2)>N44
32
1 5>(<'OSCAR'),(<'BEATA')
R
```

10.2.23 Identical

$$Z \leftarrow A \equiv B$$

Informal Description: Returns 1 if *A* and *B* match with respect to structure and content, and 0 otherwise.

Note: Identical is sometimes called "match".

Evaluation Sequence:

If A and B are both **simple-scalars**, return A=B.

If the **shape-list** of A is different from the **shape-list** of B, return **zero**.

If *A* is empty, return $(1 \uparrow, A) \equiv 1 \uparrow, B$.

Set *A*1 to the first item of the **ravel-list** of *A*, and *B*1 to the first item of the **ravel-list** of *B*.

Set Z1 to $A1 \equiv B1$.

If Z1 is **zero**, return **zero**.

Set A2 to the rest of the **ravel-list** of A, and B2 to the rest of the **ravel-list** of B.

If *A*2 is empty, return **one**.

Return $A2 \equiv B2$.

Examples:

10.2.24 Disclose

$$Z \leftarrow \supset B$$

Informal Description: For B, an array of shape A, all of whose items have the same **rank** K, Z is an array of shape A, M where M is the maximum over the shapes of the items in B, and each **rank**-K **cell** of Z is the result of M \uparrow of the corresponding item in B.

Evaluation Sequence:

```
Set M to \max-shape-of B.

Set M1 to (\rho B), M.

Set K to \rho M.

Set S to the index-set of the ravel-list of S.

Return an array S such that \rho S is S1, and such that for every member S2 of S3, the S3 the S4 the S5 to the ravel-list of S5.
```

Examples:

```
T \leftarrow (c'OSCAR'), (c'FRED')
\rho T
2
\supset T
OSCAR
FRED
\rho \supset T
2 = 5
```

10.2.25 Disclose with Axis

$$Z \leftarrow \supset [K] B$$

Informal Description: This is the same as Disclose, except that the function specifies which axes in the result Z, are new. Uses **index-origin**.

Evaluation Sequence:

```
Set Z1 to \supset B.
```

If index-origin is nil, signal implicit-error.

If A is not a **simple vector**, signal **axis-error**.

If any **thingy-in** A is not a **valid-axis-of** Z1, signal **axis-error**.

Set A1 to the **integer-array-nearest-to** A.

If any item of the **ravel-list** of A1 is not distinct, signal **axis-error**.

If the **count-of** A1 is not the **rank** of the **first-thingy-of** B, signal **axis-error**.

Return $(((\iota \rho \rho Z1) \sim A), A) \Diamond Z1$

10.2.26 Enclose

$$Z \leftarrow \subset B$$

Informal Description: Z is a scalar array whose only item is B.

Note: If B is a simple-scalar, Z is B.

Evaluation Sequence:

If B is a **simple-scalar**, return B.

Return an array whose **shape-list** is empty, and whose **ravel-list** has a single **item**, the array *B*, and whose **type** is **mixed**.

Examples:

10.2.27 Enclose with Axis

$$Z \leftarrow \subset [K] B$$

Informal Description: Z is B with the indicated axes eliminated and replace by a corresponding new level of nesting. K specifies the axis to be enclosed, and their order in the resulting items. Uses **index-origin**.

Evaluation Sequence:

```
If, for some N, K = (-N) \uparrow \iota \rho \rho B, return an array whose shape is (-N) \downarrow \rho B, and whose ravel-list contains the rank-N cells of B.

Otherwise, return \subset [(-\rho K) \uparrow \iota \rho \rho B] (((\iota \rho \rho B) \sim K), K) \lozenge B
```

11 System Functions

11.1 Introduction

Note: System functions are primitive functions whose names are **distinguished-identifiers** rather than **primitives**. System function names are not permitted in the **locals-list** of a **defined-function** header-line.

The system functions related to shared-variables and defined-functions are specified in those chapters.

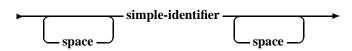
11.2 Definitions

Note: Some system functions take an argument which is treated as an identifier-array.

- Identifier-Array: A character array of rank two each of whose rows matches the character-diagram identifier-row.
- **Event-Report**: A representation of the most recent error.
- Event-Message: A character array giving, in an implementation-defined form, the
 event-report and any other information the implementation deems helpful in locating
 the error's source.
- Empty-Event-Message: An empty array, of an implementation-defined shape, indicating that the most recent event-type corresponds to no error.

11.3 Diagram

Identifier-Row



11.4 Niladic System Functions

Note: A **niladic system function** is an implementation-provided facility associated with a **distinguished-identifier**. The primary difference between a **niladic system function** and a **system variable** is that a **niladic system function** causes an error to be signalled if a value is assigned to it, or its name is included in the list of locals of a defined function header.

11.4.1 Time Stamp

 $Z \leftarrow \Box TS$

Informal Description: *Z* is a seven-element numeric vector representing the current date and time relative to an epoch and a time zone.

Evaluation Sequence:

Return the seven-element **numeric vector** result of the **implementation-algorithm time-stamp**.

Example:

Additional Requirements:

The first through sixth elements of the result returned by **time-stamp** are integral and represent respectively the current year, month, day, hour, minute, and second relative to the underlying epoch and time zone. The hour corresponds to the number of integral hours that have elapsed within the current day.

The seventh and last element is a quantity, not necessarily integral, representing the milliseconds that have elapsed within the current second.

The accuracy, resolution, epoch, and time zone for **time-stamp** are not specified by this International Standard.

11.4.2 Atomic Vector

 $Z \leftarrow \Box AV$

Informal Description: Z is an implementation-defined character vector containing every element of the required-character-set exactly once.

Evaluation Sequence:

Return atomic-vector.

Note: The elements of the **atomic-vector** are often ordered so that, in **index-origin zero**, $((\lceil 2 \otimes \rho \rfloor AV) \rho 2) \ T \rfloor AV \iota \ 'X'$ gives the bit representation of X. This behaviour is not required.

11.4.3 Line Counter

 $Z \leftarrow \Box LC$

Informal Description: Z is a vector of statement numbers of active defined functions ordered so that the most recently called function has the lowest index.

Evaluation Sequence:

Set Y to a **numeric vector** whose **length** is the same as the **number-of-items** in the **state-indicator**, such that for every I in $\iota \rho Y$, Y[I] is the **current-line-number** of **item** I of the **state-indicator**, if the **mode** of **item** I of the **state-indicator** is **defined-function**, and **zero** otherwise.

Set M to a **Boolean vector** whose **length** is the same as the **number-of-items** in the **state-indicator**, such that for every I in $1 \rho M$, M[I] is **one** if the **mode** of **item** I of the **state-indicator** is **defined-function**, and **zero** otherwise.

Return M/Y.

Note: The subsections that affect the value of \Box LC are discussed in the subsection **defined-function-control**.

11.4.4 Event Message

 $Z \leftarrow \square EM$

Informal Description: When an error occurs, three lines of information are usually displayed: the first line indicates what happened, the second line shows the line in which the error was found, and the third line contains one or more carets pointing to the place(s) in the second line where the error was detected. **Event-message** is commonly a character matrix corresponding to this three-line message. It is set according to the contents of **event-type**. If **event-type** is 0 0 then **event-message** is an empty three-row matrix. Otherwise, **event-message** is set to an appropriate message. In some cases its first row is blank (when **event-message** is set via $\square ES$) because there is no message for the given value of **event-type**. (See the description of **event-type** for proposed values and messages.)

event-message is local to a context.

Evaluation Sequence:

Return event-message.

Note: While the informal description speaks of a three line message of a particular format, this is an attempt to describe common usage, rather than prescribed behaviour. The format of **event-message** is implementation defined.

11.4.5 Event Type

 $Z \leftarrow \sqcap ET$

Informal Description: The default value for **event-type** is 0 0 (no error). When an event happens which halts execution **event-type** is set to a pair of integers specifying the class and sub-class of the event. **event-type** is local to a **context**.

Evaluation Sequence:

Return event-type.

Note: Classes and subclasses of event-type:

0 - no error

0 1 – unclassified event

There may be additional implementation-defined types.

11.5 Monadic System Functions

Note: *Refer to the* **Defined Functions** *chapter for definitions of* $\Box FX$ *and* $\Box CR$

11.5.1 Delay

 $Z \leftarrow \square DL \ B$

Informal Description: Z is the time, in seconds, for completion of this operation. As a side effect, $\square DL$ causes a delay in execution of at least B seconds.

Evaluation Sequence:

Set To to current-time seconds.

If B is not a **scalar**, signal **rank-error**.

If *B* is not a **near-real-number**, signal **domain-error**.

Set T1 to T0 plus B.

Wait until **current-time** is not **less-than** T1.

Return Z, a **numeric scalar** such that the **first-item** in the **ravel-list** of Z is **current-time minus** T0.

Note: *B* may be fractional or negative.

11.5.2 Name Class

```
Z \leftarrow \sqcap NC B
```

Informal Description: *Z* is a vector of name classes giving the usage of the identifier in the corresponding row of *B*. 0 means an identifier is available without current referent; 1, a label; 2, a variable or a shared variable; 3, a defined function; 4, a defined operator; 5, a system variable; 6, a system function.

Evaluation Sequence:

If the rank of B is greater-than two, signal rank-error.

If B is **empty**, return 10.

Set B1 to $(-2 \uparrow 1 1, \rho B) \rho B$.

If any row of *B*1 does not match **identifier-row**, signal **domain-error**.

Generate Z, a **numeric vector** such that the **length** of Z is the same as the **number-of-rows** in B1, and the **ravel-list** of Z has the property that for every I in $11 \uparrow \rho B$, the following holds:

Let ZI stand for **item** I of the **ravel-list** of Z.

Let N stand for the **token** that matched the **character-diagram simple-identifier** in the **identifier-row**.

If the **current-class** of N is **nil**, ZI is **zero**.

If the **current-class** of N is **label**, ZI is **one**.

If the current-class of N is variable or shared-variable, ZI is two.

If the current-class of N is defined-function or niladic-defined-function, ZI is three.

If the current-class of N is defined-operator, ZI is four.

If the current-class of N is system-variable, ZI is five.

If the current-class of N is system-function, ZI is six.

Otherwise, signal domain-error.

Return Z.

Example:

```
∇Z←TEST; R0; R21; R22; R31; R32

[1] R1: R21←1 □SVO 'R22'

[2] Z←□FX 1 3ρ'R31'

[3] Z←□FX 1 5ρ'R32 X'

[4] Z←□NC 6 3ρ'R0 R1 R21R22R31R32'

▼

TEST

0 1 2 2 3 3
```

11.5.3 Expunge

 $Z \leftarrow \sqcap EX B$

Informal Description: Z is a Boolean vector such that an element of Z is **one** if the identifier in the corresponding row of B is available for use when the operation completes. Expunge changes the **current-referent** of **symbols** whose **current-class** is **variable**, **shared-variable**, **defined-function**, **defined-operator**, or **niladic-defined-function** to **nil**, making them available for redefinition. The **symbols** to be changed are named by the rows of B, an **identifier-array**.

Evaluation Sequence:

If the rank of B is greater-than two, signal rank-error.

Set B1 to $(-2 \uparrow 1 1, \rho B) \rho B$.

For each **row** J of B1,

If row J of B1 does not match identifier-row, signal domain-error.

Let N stand for the **token** that matched the **character-diagram simple-identifier** in **identifier-row**.

If the current-class of N is shared-variable, retract N.

If the **current-class** of N is **variable**, set the **current-referent** of N to a **token** whose **class** is **nil**.

If the current-class of N is defined-function or niladic-defined-function, and if the current-referent of N is locally-erasable, set the current-referent of N to a token whose class is nil.

Return $\sim \times \square NC$ B.

Additional Requirement:

The operation of Expunge is **atomic**: if Expunge signals an error, **current-referents** shall be unchanged.

Note: Because of the definition of **locally-erasable**, a **defined-function** or **defined-operator** cannot be expunged by a **conforming-program** while it is **pendent** or **waiting**. This does not prevent a **conforming-implementation** from permitting such behaviour; it simply leaves the consequences to a **conforming-program** undefined.

11.5.4 Name List

 $Z \leftarrow \square NL \ B$

Informal Description: *Z* is an **identifier-array**. An identifier occurs in *Z* if its **name-class** is in *B*. Elements of *B* can be: 1, for labels; 2, for variables or shared variables; 3, for defined functions; 4, for defined operators; 5, for system variables; 6, for system functions.

Evaluation Sequence:

If the rank of B is greater-than one, signal rank-error.

If any **item** of the **ravel-list** of *B* is not a **near-integer**, signal **domain-error**.

Set B1 to the **integer-array-nearest-to** B.

If B1 contains any **numbers** other than **one**, **two**, **three**, **four**, **five**, or **six**, signal **domain-error**.

Generate Z, an **identifier-array** consisting of the **names** of **symbols** whose **current-class** is **label**, **variable**, **shared-variable**, **niladic-defined-function**, **defined-function**, **defined-operator**, **system-variable** or **system-function**. The **symbol names** are left-justified in the rows of Z.

Set Z to $((\square NC \ Z) \in B1) \neq Z$.

Remove any all-blank trailing columns from Z.

Return Z.

Note: The rank of Z is always two. If there are no symbols with a name-class designated in B, Z is O O O!.

11.5.5 Query Stop

 $Z \leftarrow \sqcap STOP B$

Informal Description: Z is an integer vector of line numbers in the function or operator whose name is B that have stop control set.

Evaluation Sequence:

If the rank of B is greater-than one, signal rank-error.

If B does not match **identifier-row**, signal **domain-error**.

Let N stand for the **token** that matched the **character-diagram simple-identifier** in **identifier-row**.

If the current-class of N is not defined-function, defined-operator or niladic-defined-function, signal domain-error.

Return the **stop-vector** of the **current-content** of N.

Additional Requirement:

Query Stop is part of the optional-facility Trace-and-Stop-Control.

Note: Query Stop is not origin sensitive.

If no lines have stop control set, the result is 10.

11.5.6 Query Trace

 $Z \leftarrow \Box TRACE B$

Informal Description: Z is an integer vector of line numbers in the function or operator whose name is B that are being traced.

Evaluation Sequence:

If the rank of B is greater-than one, signal rank-error.

If B does not match **identifier-row**, signal **domain-error**.

Let N stand for the **token** that matched the **character-diagram simple-identifier** in **identifier-row**.

If the current-class of N is not defined-function, defined-operator or niladic-defined-function, signal domain-error.

Return the **trace-vector** of the **current-content** of N.

Additional Requirement:

Query Trace is part of the **optional-facility Trace-and-Stop-Control**.

Note: Query Trace is not origin sensitive.

If no lines are being traced, the result is 10.

11.5.7 Monadic Event Simulation

 $\sqcap ES B$

Informal Description: An exception is created in the caller's **context**. *B* is either a pair of integers specifying the class and sub-class of the event, or a character vector.

If B is a pair of integers, **event-type** is set to B and **event-message** is set to an appropriate message. If the pair of integers is undefined, the message part of **event-message** is blank, but an exception is signalled and **event-type** is set. If B is O (no error), **event-type** is set to O and **event-message** is set to **empty-event-message**, but no exception is signalled.

If *B* is empty no action is taken. This allows for conditional signalling of an event.

Evaluation Sequence:

If the rank of B is greater-than one, signal a rank-error.

If B is **empty**, return **nil**.

If B is **numeric**,

If the rank of B is not one, signal rank-error.

If the **count** of **items** in *B* is not **two**, signal **length-error**.

If any **item** in B is not a **near-integer**, signal **domain-error**.

Set B1 to the **integer-nearest-to** B.

Set **event-type** to B1.

Set event-message to empty-event-message.

If B1 is not 0 0,

Set event-message as appropriate.

Signal event corresponding to event-type.

If *B* is **character**,

Set **event-type** to 0 1.

Set message portion of **event-message** to *B*.

Signal event corresponding to event-type.

11.6 Dyadic System Functions

11.6.1 Name List

 $Z \leftarrow A \sqcap NL B$

Informal Description: Z is an **identifier-array** of all names having a **name-class** in B that begin with a letter in A.

Evaluation Sequence:

If the rank of A is greater-than one, signal rank-error. If any item of the ravel-list of A does not match letter, signal domain-error. Set Z1 to $\square NL$ B. If $0 \in \rho Z1$, return Z1. Return $(Z1[;1] \in A) \neq Z1$.

11.6.2 **Set Stop**

 $Z \leftarrow A \sqcap STOP B$

Informal Description: Z is a numeric vector representing the stop controls in effect for the function or operator whose name is in B before $\square STOP$ began evaluation. As a side effect, $\square STOP$ sets stop controls for the function or operator whose name is in B at lines specified by A.

Evaluation Sequence:

If the rank of B is greater-than one, signal rank-error.

If B does not match identifier-row, signal domain-error.

Let N stand for the **token** that matched the **character-diagram simple-identifier** in **identifier-row**.

If the current-class of N is not defined-function, defined-operator or niladic-defined-function, signal domain-error.

If the rank of A is greater-than one, signal rank-error.

If any **item** of the **ravel-list** of A is not a **near-integer**, signal **domain-error**.

Set A1 to the **integer-array-nearest-to** A.

If any **item** of the **ravel-list** of A1 is not a **positive-integer**, signal **domain-error**.

Set \tilde{Z} to $\square STOPB$.

Let *M* stand for the **last-line-number** of the **current-content** of *N*.

Set the **stop-vector** of the **current-content** of N to $((\iota M) \in A1) / \iota M$, evaluated with **index-origin one**.

Return Z.

Additional Requirement:

Set Stop is part of the optional-facility Trace-and-Stop-Control.

Note: Stop controls are an attribute of **defined-function** objects. \square STOP affects the **current-referent** of B only.

11.6.3 Set Trace

 $Z \leftarrow A \sqcap TRACE B$

Informal Description: Z is a numeric vector representing the lines being traced in the function or operator whose name is in B before $\Box TRACE$ began evaluation. As a side effect, $\Box TRACE$ begins tracing the lines specified by A in the function or operator whose name is in B.

Evaluation Sequence:

If the rank of B is greater-than one, signal rank-error.

If B does not match **identifier-row**, signal **domain-error**.

Let N stand for the **token** that matched the **character-diagram simple-identifier** in **identifier-row**.

If the current-class of N is not defined-function, defined-operator, or niladic-defined-function, signal domain-error.

If the rank of A is greater-than one, signal rank-error.

If any **item** of the **ravel-list** of A is not a **near-integer**, signal **domain-error**.

Set A1 to the **integer-array-nearest-to** A.

If any **item** of the **ravel-list** of A1 is not a **positive-integer**, signal **domain-error**.

Set Z to $\square TRACE B$.

Let *M* stand for the **last-line-number** of the **current-content** of *N*.

Set the **trace-vector** of the **current-content** of N to $((\iota M) \in A1) / \iota M$, evaluated with **index-origin one**.

Return Z.

Additional Requirement:

Set Trace is part of the **optional-facility Trace-and-Stop-Control**.

Note: Trace controls are an attribute of **defined-function** objects. $\Box TRACE$ affects the **current-referent** of B only.

11.6.4 Execute Alternate

 $Z \leftarrow A \sqcap EA B$

Informal Description: A and B are both character vectors representing executable expressions. If execution of B is successful the result is B. If there is an error or an interrupt in the execution of B, execution is abandoned and A is executed. The execution of A is subject to normal error handling.

It is possible to get a **value-error** if the successful execution of B does not return a result and an explicit result is needed. A is not executed in this case.

Evaluation Sequence:

If the rank of either A or B is greater-than one signal rank-error.

If any **item** of the **ravel-list** of A is not a **character**, signal **domain-error**.

If any **item** of the **ravel-list** of *B* is not a **character**, signal **domain-error**.

Generate a new context in which

mode is execute.

current-line is the **ravel-list** of B,

current-function is 0 0 0 1 1,

current-line-number is one,

current-statement is the empty list of tokens,

stack is the empty list of tokens.

Append the new **context** to the **state-indicator** of the **active-workspace** as a new first **item**.

Set Z to **evaluate-line**.

Remove the first **context** from the **state-indicator**.

If Z is not a member of **error** return Z,

Note: See Conforming Programs in Compliance regarding hazards to conforming programs if this facility is used.

11.6.5 Dyadic Event Simulation

 $A \sqcap ES B$

Informal Description: An exception is created in the caller's **context**. *A* is a **character scalar** or **vector** and *B* is either a pair of integers specifying the class and sub-class of the event, or an empty **vector**.

Event-type is set to B and the message portion of **event-message** is set to A. If B is 0 0 (no error), **event-type** is set to 0 0 and **event-message** is set to **empty-event-message**, and no event is signalled.

If B is empty no action is taken. This allows for conditional signalling of an event (for example $A \square ES$ **cond** /B).

Evaluation Sequence:

If the rank of A is greater-than one, signal rank-error.

If any **item** of *A* is not a **character**, signal **domain-error**.

If the **rank** of *B* is not **one**, signal **rank-error**.

If *B* is empty, return **nil**.

If the **count** of *B* is not **two**, signal **length-error**.

If any **item** in B is not a **near-integer**, signal **domain-error**.

Set B1 to the **integer-nearest-to** B.

Set **event-type** to *B*1.

Set event-message to empty-event-message.

If B1 is not 0 0,

Set the message portion of **event-message** to *A*.

Set the rest of **event-message** as appropriate.

Signal event corresponding to **event-type**.

11.6.6 Transfer Form

 $Z \leftarrow A \sqcap TF B$

Informal Description: B is a character vector giving the name of the object for which the transfer form is desired. If there is no transfer form for the item named by B (for example, B names a system function, a label, or has no value) the result is an empty character vector.

A is either **one** or **two**. If it is **one** B may name only defined functions, defined operators, and simple homogeneous arrays. If it is **two**, B may also name **mixed** arrays.

For a given A, $\Box TF$ is its own inverse. For example, if ABC names a transferable object, $A \Box TF A \Box TF ABC'$ will reestablish the object named ABC and return the character list ABC'.

If *A* is **two**, the result is an executable character vector giving the name and contents of the object named by *B*. (Numeric conversions are carried out to **full-print-precision**.)

Evaluation Sequence:

If the rank of A is greater-than one signal rank-error.

If the **count** of *A* is not **one** signal **length-error**.

If any **item** of the **ravel-list** of A is not a **near-integer**, signal **domain-error**.

Set A1 to the **integer-nearest-to** A.

If A1 has the value **one**, set Z to **produce-canonical-representation-vector** of the object named by B.

If A1 has the value **two**,

If the rank of B is greater-than one signal rank-error.

If the **type** of B is not **character** signal **domain-error**.

Set *NC* to the **name-class** of *B*.

If *NC* is **two**:

Set B1 to ΦB

Set Z to an executable string that will recreate the object named by B

If *NC* is **three** or **four**:

Set B1 to the **character-representation** of the object named by B

If the **count** of the **shape** of B is not **zero**

Set Z to an executable string that will recreate the object named by B

Else set Z to an empty character vector.

If *NC* is **negative-one**:

If ΦB is successful, set Z to the name of the object created

Otherwise, set Z to an empty character vector

Otherwise signal domain-error

Otherwise, signal domain-error.

12 System Variables

12.1 Definitions

System-Variable-Symbol: A symbol, having as its name a distinguished-identifier, and as its referent-list a list of tokens representing the value assigned to an associated system-parameter in each context of the state-indicator.

Note: System parameters are values used implicitly and set as side effects by primitive operations.

 Internal-Value-Set: An implementation-defined set of values that the tokens in the referent-list of a system-variable-symbol can take on.

Note: When a system-variable-symbol is localised the initial class of its associated system-parameter is nil. If a primitive operation is invoked that requires a system parameter whose class is currently nil, the primitive operation signals implicit-error. Therefore, a conforming-program that localises system-variable-symbols should assign them values from their internal-value-set before calling primitive operations that require them.

The assignment operation for a system variable rejects attempts to set its associated system parameter to a value outside its **internal-value-set**.

12.2 Evaluation Sequences

12.2.1 Comparison Tolerance

$$Z \leftarrow \Box CT \leftarrow B$$
$$Z \leftarrow \Box CT$$

Informal Description: *Z* is the current value of **comparison-tolerance**.

The distance within which numbers are to be considered equal by certain primitives is a function of **comparison-tolerance**.

Evaluation Sequence:

For form $\square CT$

Return comparison-tolerance.

For form $\square CT \leftarrow B$

If the rank of B is greater-than one, signal rank-error.

If the **count** of *B* is not **one**, signal **length-error**.

Let B1 be the **first-scalar** in B.

If *B*1 is not a **nonnegative-number**, signal **domain-error**.

If B1 is not in the **internal-value-set** of **comparison-tolerance**, signal **limit-error**.

Set **comparison-tolerance** to *B*1.

Return a **token** whose **class** is **committed-value** and whose **content** is *B*.

Additional Requirement:

The internal-value-set for comparison-tolerance consists of nonnegative-numbers not greater than the implementation-parameter comparison-tolerance-limit.

The initial value of **comparison-tolerance** in a **clear-workspace** is that member of the **internal-value-set** for **comparison-tolerance** given by the **implementation-parameter initial-comparison-tolerance**.

Note: The following subsections reference comparison-tolerance: ceiling, equal, floor, greater than or equal to, greater than, index of, less than or equal to, less than, member of, not equal, residue, and, or and unique.

12.2.2 Random Link

```
Z \leftarrow \square RL \leftarrow BZ \leftarrow \square RL
```

Informal Description: Z is the current value of **random-link**; **random-link** is the current seed of the pseudorandom number generator.

Evaluation Sequence:

For form $\Box RL$

Return random-link.

For form $\square RL \leftarrow B$

If the rank of B is greater-than one, signal rank-error.

If the **count** of B is not **one**, signal **length-error**.

Set B1 to the **first-scalar** in B.

If *B*1 is not a **near-integer**, signal **domain-error**.

Set B2 to the **integer-nearest-to** B1.

If B2 is **less-than one**, signal **domain-error**.

If B2 is not in the **internal-value-set** for **random-link**, signal **limit-error**.

Set **random-link** to *B*2.

Return a **token** whose **class** is **committed-value** and whose **content** is *B*.

Additional Requirement:

The internal-value-set of random-link is implementation-defined.

The initial value of **random-link** in a **clear-workspace** is that member of the **internal-value-set** for **random-link** given by the **implementation-parameter initial-random-link**.

Note: The system parameter **random-link** is used and set by **roll** and **deal**. **Roll** gives a reference for a suitable algorithm.

12.2.3 Print Precision

$$Z \leftarrow \Box PP \leftarrow B$$
$$Z \leftarrow \Box PP$$

Informal Description: Z is the current value of **print-precision**, which controls the number of significant positions in the output form produced by **monadic format** and **numeric-output-conversion**.

Evaluation Sequence:

For form $\square PP$

Return **print-precision**.

For form $\square PP \leftarrow B$

If the rank of B is greater-than one, signal rank-error.

If the **count** of *B* is not **one**, signal **length-error**.

If *B* is not **near-integer**, signal **domain-error**.

Set *B*1 to the **integer-nearest-to** the **first-scalar** in *B*.

If *B*1 is **less-than one**, signal **domain-error**.

If B1 is not in the **internal-value-set** of **print-precision**, signal **limit-error**.

Set **print-precision** to *B*1.

Return a **token** whose **class** is **committed-value** and whose **content** is *B*.

Additional Requirements:

The internal-value-set of print-precision consists of the integer scalars in the closed-interval-between one and print-precision-limit.

The initial value of **print-precision** in a **clear-workspace** is that member of the **internal-value-set** for **print-precision** given by the **implementation-parameter initial-print-precision**.

Note: The system parameter print-precision is used by monadic format and numeric-output-conversion. Print-precision-limit should be greater than or equal to full-print-precision.

12.2.4 Index Origin

$$Z \leftarrow \square IO \leftarrow B$$
$$Z \leftarrow \square IO$$

Informal Description: Z is the current value of **index-origin**, which is the index associated with the first position of any axis of non-zero length.

Evaluation Sequence:

For form $\Box IO$ Return index-origin.
For form $\Box IO \leftarrow B$ If the rank of B is greater-than one, signal rank-error.
If the count of B is not one, signal length-error.
If B is not a near-integer, signal domain-error.
Set B1 to the integer-nearest-to the first-scalar in B.
If B1 is not in the internal-value-set of index-origin, signal limit-error.
Set index-origin to B1.
Return a token whose class is committed-value and whose content is B.

Additional Requirement:

The internal-value-set for index-origin is the scalar integer values zero and one.

The initial value of **index-origin** in a **clear-workspace** is that member of the **internal-value-set** for **index-origin** given by the **implementation-parameter initial-index-origin**.

Note: The following primitive operations refer to index-origin: deal, dyadic transpose, grade down, grade up, index generator, index of, indexed assignment, indexed reference, roll, and all functions that provide for axis specification when the form containing an axis is used.

12.2.5 Latent Expression

```
Z \leftarrow \Box LX \leftarrow B
Z \leftarrow \Box LX
Z \leftarrow \Box LX[I] \leftarrow B
```

Informal Description: Z is the current value of **latent-expression**, which is a **character** vector that is executed when a workspace is activated.

Evaluation Sequence:

Additional Requirement:

The **internal-value-set** for **latent-expression** is an **implementation-defined** subset of the set of all character vectors which includes the empty vector.

Return a **token** whose **class** is **committed-value** and whose **content** is B.

The initial value of **latent-expression** in a **clear-workspace** is that member of the **internal-value-set** for **latent-expression** given by the **implementation-parameter initial-latent-expression**.

13 Defined Functions

13.1 Introduction

Note: Algorithms written in APL are called defined functions. A defined function consists of a header line and zero or more body lines. The header line indicates the name and syntax class of the function and gives a list of names to be localised and, in the case of argument names, initialised. Each body line consists of an optional label followed by an APL line to be evaluated. Body lines are evaluated in the order in which they occur in the defined function unless a statement beginning with a right arrow is evaluated. In this chapter, the term defined function also includes defined operator.

A defined function is established in a workspace by the system function $\Box FX$, by actions in function definition mode, or by the system command) COPY.

In this International Standard, a **defined-function** is represented as an object with three attributes: **canonical-representation**, **trace-vector**, and **stop-vector**. A **defined-function** is called when its **name** occurs in a **prefix** of the **current-stack** matching one of the **patterns** in the **phrase-table**.

A defined-function is called by the operation call-defined-function. This operation prefixes the state-indicator with a new context, localises and initialises any local names in the header-line, and calls the evaluation sequence in the subsection defined-function-control.

A defined-function ends the evaluation if defined-function-control finds current-line-number set to a value not in the closed-interval-between one and last-line-number. At this point, the first item of the state-indicator is discarded and a token of class result is returned, via call-defined-function, to the phrase-evaluator that called it.

The returned **token** may be **constant** (for functions returning a value), **nil** (for functions that return no value), **unwind** (for functions that end through the evaluation of an escape arrow), or **clear-state-indicator** (as the result of an)SIC command).

If a line in a defined-function signals an error, immediate-execution is called to report the error and suspend the defined-function. Immediate-execution may return either a branch to indicate that evaluation of the defined-function should continue, or an escape or clear-state-indicator to indicate that the defined-function context should be removed from the state-indicator.

13.2 Definitions

Note: The following term defines the subset of arrays that can be used to create defined-functions.

- **Proper-Defined-Function**: An array A is a proper-defined-function if The **type** of A is **character**.

The rank of A is two.

A has at least one **row**.

The first **row** of *A* matches the **character-diagram header-line**.

No **identifier** in the first row of A is the **distinguished-identifier** \square or \square or is a **system-function-name** or **niladic-system-function-name**.

All rows after the first in A match the character-diagram body-line.

Note: The following definitions give properties of arrays that are **proper-defined-functions**.

- Niladic: A, an array that is a proper-defined-function, is niladic if the form of the first row of A is niladic-function-header.
- Monadic: A, an array that is a proper-defined-function, is monadic if the form of the first row of A is monadic-function-header.
- Ambivalent: A, an array that is a proper-defined-function, is ambivalent if the form
 of the first row of A is ambivalent-function-header.
- Monadic-Monadic-Operator: A, an array that is a proper-defined-function, is monadic-monadic-operator if the form of the first row of A is monadic-monadic-operator-header.
- Monadic-Dyadic-Operator: A, an array that is a proper-defined-function, is monadic-dyadic-operator if the form of the first row of A is monadic-dyadic-operator-header.
- Ambivalent-Monadic-Operator: A, an array that is a proper-defined-function, is ambivalent-monadic-operator if the form of the first row of A is ambivalent-monadic-operator-header.
- Ambivalent-Dyadic-Operator: A, an array that is a proper-defined-function, is ambivalent-dyadic-operator if the form of the first row of A is ambivalent-dyadic-operator-header.
- Value-Returning: A, an array that is a proper-defined-function, is value-returning
 if, when header-line is threaded with the first row of A, the diagram result is threaded.

Note: The following terms categorise **defined-functions** according to their presence on or absence from the **state-indicator**.

- Suspended: N, a simple-identifier, is suspended if it is the function-name of a context
 that immediately follows an immediate-execution context.
- Pendent: N, a simple-identifier, is pendent if it is the function-name of a context that
 does not immediately follow an immediate-execution context.

- Waiting: N, a simple-identifier, is waiting if a token whose class is defined-function-name and whose content is the content of N occurs in the stack of some context in the state-indicator of the active-workspace, and is not the first defined-function-name in that stack.

Note: For example, if F and G are defined function names, G is **waiting** during the evaluation of (F 2) in the line

 $(F \ 2) \ G \ 3.$

- Editable: N, a simple-identifier, is editable if the current-referent of N is nil, or if all the following are true:

The current-class of N is defined-function, defined-operator or niladic-defined-function.

The **number-of-items** in its **referent-list** is **one**.

N is neither **pendent** nor **waiting**.

N, if **suspended**, is **suspended** in only the **current-context**.

Note: This definition requires that a **conforming-implementation** be capable of editing the top function on the state indicator as long as it has not been **localised** and occurs nowhere else on the state indicator. This minimal definition is included to be certain that all **conforming-implementations** be capable of interactive program correction.

- Globally-Erasable: Globally-Erasable is defined for N, a simple-identifier, as follows: If no **context** of the **state-indicator** contains, in its **current-statement**, a **defined-name** with **content** that of N, then N is **globally-erasable**.

Otherwise, let L be the last (most global) **context** in whose **local-name-list** N appears. If there is no such **context** then N is not **globally-erasable**.

Otherwise, let C be the last (most global) **context** in whose **current-statement** a **defined-name** with **content** that of N appears.

If the index of L is greater-than or equal to that of C then N is **globally-erasable**.

Note: *Informally,* \mathbb{N} *is* **globally-erasable** *if it was localised before it was made waiting or pendent.*

- Locally-Erasable: Locally-Erasable is defined for N, a simple-identifier, as follows: If no **context** of the **state-indicator** contains, in its **current-statement**, a **defined-name** with **content** that of N, then N is **locally-erasable**.

Otherwise, let L be the first (most local) **context** in whose **local-name-list** N appears. If there is no such **context** then N is not **locally-erasable**.

Otherwise, let C be the first (most local) **context** in whose **current-statement** a **defined-name** with **content** that of N appears.

If the index of C is greater than that of L, then N is **locally-erasable**.

Note: Informally, N is **locally-erasable** if it has been localised since the last time it was made waiting or pendent. This definition is used in $\square EX$ to specify that the expression ($\square EX^{\dagger}F^{\dagger}$) F 1 be the same as 0 F 1; it is used in $\square FX$ to specify that the expression ($\square FX$ 1 $1 \rho^{\dagger}F^{\dagger}$) F 1 signal a **domain-error**.

Note: The following operations provide information about arrays that are proper-defined-functions.

- Function-Name of A: For A a proper-defined-function, the simple-identifier that matches function-name when the character-diagram header-line is threaded with the first row of A.
- Last-Line-Number of A: For A a proper-defined-function, negative-one plus the number-of-rows in A.
- Header-Name-List of A: For A a proper-defined-function, a list of identifiers developed as follows:

Thread the **character-diagram header-line** with the first **row** of A.

Return a **list** of **tokens** consisting of **identifiers** that matched any of the **character-diagrams local-name**, **result-name**, **left-argument-name**, **right-argument-name** right-operand-name, or left-operand-name.

 Label-Name-List of A: For A a proper-defined-function, a list of simple-identifiers developed as follows:

Set Z to the **empty-list** of **tokens**.

For each index I in the closed-interval-between one and last-line-number of A,

Thread the **character-diagram body-line** with **row** I **plus one** of A.

If **labelled-line** was matched, append the **simple-identifier token** that matched the **character-diagram label-name** to Z as a new last **item**.

Return Z.

- Local-Name-List of A: For A a proper-defined-function, a list of identifiers consisting
 of the label-name-list followed by the header-name-list.
- Identifier-Matching D in A: For A a proper-defined-function and D a character-diagram, a simple-identifier token developed as follows:

Thread the **character-diagram header-line** with the first **row** of A.

Return a **token** of **class simple-identifier** whose **content** is the **list** of **characters** that matched D.

Note: Identifier-Matching is used to refer to the arguments and result of an arbitrary defined function.

 Function-Line I of A: For A a proper-defined-function, a list of characters developed as follows:

Thread the **character-diagram body-line** with **row** I **plus one** of A.

Return the list of characters that matched the character-diagram line in body-line.

Note: The following operations affect symbols.

– Localise N:

Append a new first **item**, a **token** of **class nil**, to the **referent-list** of the **symbol-named-by** N.

– Delocalise N:

If the current-class of N is shared-variable, retract N.

Remove the first **item** from the **referent-list** of the **symbol-named-by** *N*.

- Current-Canonical-Representation: The canonical-representation of the current-function of the current-context in the active-workspace.
- Current-Function-Line I: Function-Line I of the current-canonical-representation.
- Current-Last-Line-Number: The last-line-number of the current-canonical-representation.
- Current-Stop-Vector: The stop-vector of the current-function of the current-context in the active-workspace.
- Current-Trace-Vector: The trace-vector of the current-function of the current-context in the active-workspace.
- Current-Local-Names: The local-name-list of the current-canonical-representation.
- Current-Right-Argument-Name: The identifier-matching right-argument-name in the current-canonical-representation.
- Current-Left-Argument-Name: The identifier-matching left-argument-name in the current-canonical-representation.
- Current-Result-Name: The identifier-matching result-name in the current-canonical-representation.
- Current-Left-Operand-Name: The identifier-matching left-operand-name in the current-canonical-representation.
- Current-Right-Operand-Name: The identifier-matching right-operand-name in the current-canonical-representation.

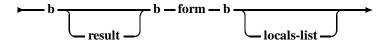
13.3 Diagrams

In these diagrams,

b stands for **permitted-blanks**.

r stands for required-blanks.

Header-Line

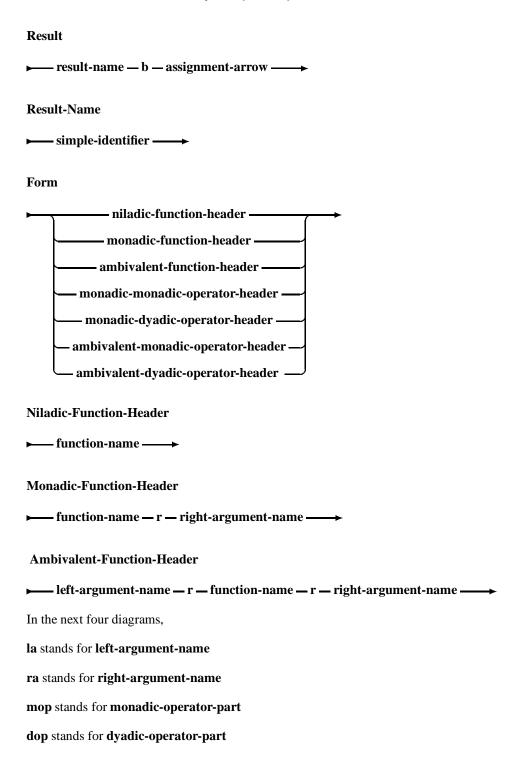


Example:

$$Z \leftarrow A F B; C; \square IO$$

Note: The header line indicates the name and syntactic class of the **defined-function**. It also gives a list of identifiers to be **localised** when the **defined-function** is called.

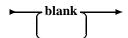
This International Standard does not provide for end-of-line comments in header lines.



Monadic-Monadic-Operator-Header
\longrightarrow b $-$ mop $-$ b $-$ ra \longrightarrow
Monadic-Dyadic-Operator-Header
\longrightarrow b $-dop - b - ra \longrightarrow$
Ambivalent-Monadic-Operator-Header
\longrightarrow la $-b$ $-mop$ $-b$ $-ra$ \longrightarrow
Ambivalent-Dyadic-Operator-Header
\longrightarrow la -b -dop -b -ra \longrightarrow
Monadic-Operator-Part
- ($-$ b $-$ left-operand-name $-$ r $-$ function-name $-$ b $-$)
Dyadic-Operator-Part
► (-b -left-operand-name -r -function-name -r -right-operand-name -b -)
Right-Argument-Name
► simple-identifier →
Left-Argument-Name
simple-identifier ——
Right-Operand-Name
► simple-identifier →
Left-Operand-Name
► simple-identifier →

Locals-List - local-marker — b — local-name — b -Colon **---**: ----Local-Marker **Function-Name** ► simple-identifier → Local-Name identifier ----Label-Name **──** simple-identifier **── Body-Line** - labelled-line -- line – Labelled-Line \longrightarrow label-name — b — colon — b — line \longrightarrow **Permitted-Blanks** blank -

Required-Blanks



13.4 Operations

13.4.1 Call-Defined-Function

 $Z \leftarrow DFN$

 $Z \leftarrow DFN B$

 $Z \leftarrow A DFN B$

 $Z \leftarrow \mathbf{f} DFN B$

 $Z \leftarrow A \mathbf{f} DFN B$

 $Z \leftarrow \mathbf{f} DFN \mathbf{g} B$

 $Z \leftarrow A \mathbf{f} DFN \mathbf{g} B$

Informal Description: Z is the value of the variable given as the result name in the header line of the defined function whose name is DFN. If there is no result name, Z is **nil**.

Evaluation Sequence:

Generate a new context in which

mode is defined-function,

current-line is the empty list of characters,

current-function is the **current-referent** of *DFN*,

current-line-number is one,

current-statement is the empty list of tokens, and

stack is the empty **list** of **tokens**.

Append the new **context** to the **state-indicator** of the **active-workspace** as a new first **item**.

Localise the current-local-names.

For form $Z \leftarrow A$ **f** DFN **g** B

If current-function is not ambivalent-defined-dyadic-operator, signal valence-error.

Set the **symbol-named-by current-left-argument-name** to A.

Set the **symbol-named-by current-left-operand-name** to **f**.

Set the symbol-named-by current-right-operand-name to g.

Set the **symbol-named-by current-right-argument-name** to *B*.

For form $Z \leftarrow \mathbf{f} DFN \mathbf{g} B$

If current-function is not monadic-defined-dyadic-operator, signal valenceerror.

Set the symbol-named-by current-left-operand-name to f.

Set the **symbol-named-by current-right-operand-name** to **g**.

Set the **symbol-named-by current-right-argument-name** to *B*.

For form $Z \leftarrow A$ **f** DFN B

If **current-function** is not **ambivalent-defined-monadic-operator**, signal **valence-error**.

Set the **symbol-named-by current-left-argument-name** to A.

Set the **symbol-named-by current-left-operand-name** to **f**.

Set the **symbol-named-by current-right-argument-name** to *B*.

For form $Z \leftarrow \mathbf{f} DFN B$

If current-function is not monadic-defined-monadic-operator, signal valence-

Set the symbol-named-by current-left-operand-name to f.

Set the **symbol-named-by current-right-argument-name** to *B*.

For form $Z \leftarrow A DFN B$

If current-function is not ambivalent, signal valence-error.

Set the **symbol-named-by current-left-argument-name** to A.

Set the **symbol-named-by current-right-argument-name** to *B*.

For form $Z \leftarrow DFN$

If current-function is not niladic, signal valence-error.

For all indices I in the closed-interval-between one and current-last-line-number, If row I plus one of canonical-representation matches the diagram labelled-line, set the current-referent of the simple-identifier that matches label-name to a token of class label and content the numeric-scalar with value I.

Set Z to **defined-function-control**.

Delocalise the current-local-names.

Remove the **first-item** in the **state-indicator** of the **active-workspace**.

If Z is **escape**, signal **unwind**.

Otherwise, return Z.

Note: Tokens of class escape are converted to tokens of class unwind so that immediate-execution can distinguish an escape issued in immediate-execution mode from one issued in a defined-function.

The context attribute current-function includes the function itself plus its trace and stop vectors.

13.4.2 Defined-Function-Control

Informal Description: This procedure controls the execution of a user defined function.

Evaluation Sequence:

Set Z to a **token** whose **class** is **nil**.

Let I stand for **current-line-number**.

Set *I* to **one**.

Repeat:

If I is in the current-stop-vector, set attention-flag to one.

If Z is an **error** or **attention-flag** is **one**,

Set N to **immediate-execution** with Z.

If N is **escape** or **clear-state-indicator**, return N.

If N is **branch**, set I to the **content** of N.

If I is not in the **closed-interval-between one** and **current-last-line-number**,

If the current-function is not value-returning, return nil.

If the **current-class** of the **current-result-name** is **variable** or **nil**, return the **current-referent** of the **current-result-name**.

Otherwise, signal value-error.

Set Z to evaluate-line of current-function-line I.

If Z is **escape**, **unwind**, or **clear-state-indicator**, return Z.

If Z is **branch**, set I to the **content** of Z.

If Z is **value** or **nil**, set I to I plus **one**.

(End of repeated block)

Note: The attention-flag is reset by the user facility enter.

The class of the token returned by evaluate-line for $\rightarrow 10$ is nil, not branch.

The value-error signalled from defined-function-control is signalled in the line that invoked the defined function; the error is introduced to prevent conforming-programs from exiting a defined-function with the current-class of the result-name set to such classes as defined-function, niladic-defined-function, and shared-variable.

13.4.3 Function Fix

 $Z \leftarrow \prod FX B$

Informal Description: Z is, for B the **canonical-representation** of a defined function or operator having **function-name** N, a **character** vector holding N. As a side effect, the defined function that B represents is established as the current referent of N. Rows of B that are all blanks give rise to blank lines in the established function.

Evaluation Sequence:

If the **rank** of *B* is not **two**, signal **rank-error**.

If *B* is an **empty array**, signal **length-error**.

If any **item** of the **ravel-list** of *B* is not a **character**, signal **domain-error**.

If B is not a **proper-defined-function**, signal **domain-error**.

If **identifiers** are duplicated in the **local-name-list** of *B*, signal **domain-error**.

Let N stand for the **simple-identifier** that matches **function-name** in the first **row** of B.

If the current-class of N is not defined-operator, defined-function, niladic-defined-function, or nil, signal domain-error.

If N is not **locally-erasable**, signal **domain-error**.

Set the current-referent of N to a defined-function or defined-operator for which canonical-representation is set from B,

stop-vector is 10,

trace-vector is 10.

Return a **character vector** Z such that the **length** of Z is the **number-of-items** in the **content** of N and the **ravel-list** of Z is the **content** of N.

13.4.4 Character Representation

 $Z \leftarrow \sqcap CR B$

Informal Description: Z is a **character** matrix representation of the defined function or operator named by B.

Evaluation Sequence:

If the rank of B is greater-than one, signal rank-error.

If the **length** of *B* is **zero**, signal **length-error**.

If B does not match the **character-diagram identifier-row**, signal **domain-error**.

If the current-class of B is not defined-operator, defined-function or niladic-defined-function, signal domain-error.

Otherwise return a **token** whose **class** is **constant** and whose **content** is the **canonical-representation** of the **current-content** of B.

Additional Requirement:

The formal model of APL used in this International Standard assumes that the internal representation of a defined function is precisely the **character** matrix satisfying the requirements for a **proper-defined-function** that was supplied as an argument to **function fix**. Representation of defined functions in this way in an implementation is neither required nor suggested.

The preservation of numeric literals and blanks as entered is desirable in the character representation, but it is not required by this International Standard. $\Box CR \Box FX \Box CR \Box FX$ X shall be the same as $\Box CR \Box FX X$, if X is a **proper-defined-function**.

If **numeric-literals** are stored as **numbers** in **defined-functions**, they are to be formatted with **print-precision** set to **full-print-precision**.

13.5 Function Editing

Note: The function editing facilities described here are included to provide the APL programmer with a standard method for entering and correcting defined functions.

13.5.1 Evaluate-Function-Definition-Request

Evaluate-Function-Definition-Request *Q.*

Note: This subsection provides minimal facilities for the establishment and alteration of **defined-functions**.

An identifier can be edited by the function editing facilities only if it is is **editable**. The header line of a **defined-function** can be altered only if the **function-name** is **locally-erasable**.

In function-definition mode, lines of the function being edited are temporarily associated with numbers in decimal-rational form by an unspecified mechanism. When a function is closed, the order of the lines given by the line number associated with each line is maintained.

The display of a function should be the same as the entries required to add the function to a **clear-workspace**.

This operation is called by immediate-execution when it recognises a function-definition-line.

Evaluation Sequence:

Thread **opening-request** with Q; if Q does not match **opening-request**, signal **definition-error**.

Rethread **opening-request** with *Q*, taking the following actions:

When **creation-request** is matched,

Set G to the **list** of **characters** that matched **creation-request**.

Set N to the **simple-identifier** that matched **function-name** in G.

If the **current-class** of N is not **nil** and G does not match **identifier-row**, signal **definition-error**.

When **change-request** is matched,

Set N to the **simple-identifier** that matched the diagram **subject-function** in **change-request**.

Set G to the **list** of **characters** that matched **initial-request**.

If the **current-class** of N is **nil**, set M to a new **defined-function** for which

Canonical-Representation is 0 0p 1 1

Stop-Vector is 10

Trace-Vector is 10.

Otherwise,

If N is not **editable**, signal **definition-error**.

Set M to the **current-referent** of N.

Generate a new context in which

mode is function-definition.

current-line is *G*,

current-function is *M*,

current-line-number is the **number-of-rows** in the **canonical-representation** of M

current-statement is the empty list of tokens, and

stack is the empty list of tokens.

Append the **context** as a new first **item** to the **state-indicator**.

Repeat:

Set Z to evaluate-editing-request.

If Z is **command-complete**,

Remove the **first-item** in the **state-indicator** of the **active-workspace**.

Return Z

If Z is an **error**, **display** Z.

Set current-prompt to function-definition-prompt of current-line-number.

Set current-line to read-keyboard.

(End of repeated block)

Additional Requirement:

The prompt in definition mode is a line number enclosed in brackets. Line numbers are always displayed in **decimal-rational** form. The line number of the header is always **zero**; it is not affected by **index-origin**.

13.5.2 Evaluate-Editing-Request

Evaluate-Editing-Request

Informal Description: This subsection acts on a single line of input to the editor.

Evaluation Sequence:

Let C stand for the **current-canonical-representation**.

Let N stand for the **simple-identifier** that matches **function-name** in **function-line**.

Let S stand for the **current-stop-vector**.

Let T stand for the **current-trace-vector**.

Thread **general-request** with **current-line**; if **general-request** does not match **current-line**, signal **definition-error**.

Otherwise, rethread **general-request** with **current-line**, taking the following actions: When **positioning-request** is threaded,

Set \bar{L} to the **numeric-input-conversion** of the **list** of **characters** that matched **line-number** in **positioning-request**.

If L is a **negative-number**, signal **definition-error**.

If L is greater-than definition-line-limit, signal limit-error.

Set **current-line-number** to L.

When **deletion-request** is threaded,

Set \mathcal{L} to the **numeric-input-conversion** of the **list** of **characters** that matched **line-number** in **deletion-request**.

If L is not a **positive-number**, signal **definition-error**.

If L is greater-than definition-line-limit, signal limit-error.

Set current-line-number to L.

Delete from C the row associated with L, if it exists.

Delete from S and T the item associated with L, if it exists.

When **display-request** is threaded, **display function-display** of C.

When **function-line** is threaded, if **function-line** does not match **permitted-blanks**, If **current-line-number** is **zero**,

If function-line does not match header-line, signal definition-error.

If N is not **locally-erasable**, signal **definition-error**.

Set the **row** of C associated with **current-line-number** to the **list** of **characters** that matched **function-line**.

If current-line-number is not an integer, or is greater-than current-last-line-

number, set the **items** of S and T associated with **current-line-number** to **zero**.

Set current-line-number to next-definition-line of current-line-number.

When **end-definition** is matched,

If C is not a **proper-defined-function**, signal **definition-error**.

If **identifiers** are duplicated in the **local-name-list** of *C*, signal **definition-error**.

Set the current-referent of N to a defined-function for which

Canonical-Representation is *C*.

Stop-Vector is S.

Trace-Vector is T.

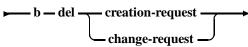
Return command-complete.

13.5.3 Diagrams

In these diagrams,

b stands for **permitted-blanks**.

Opening-Request



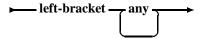
Creation-Request

header-line —

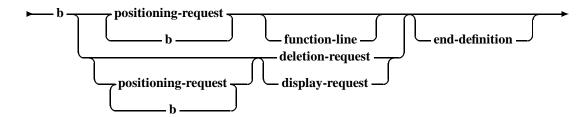
Change-Request

▶ — b — subject-function — b — initial-request — **→**

Initial-Request



General-Request



Positioning-Request

▶── left-bracket — b — line-number — right-bracket — →

Deletion-Request

► left-bracket — b — delta — b — line-number — right-bracket — →

Display-Request

 \longrightarrow left-bracket — b — quad — b — right-bracket \longrightarrow

Function-Line

─ body-line **─**

End-Definition

— b − del − b — →

Line-Number

▶ decimal-rational — b →

Subject-Function

── simple-identifier **──**

Delta

 \longrightarrow \triangle \longrightarrow

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ISO/IEC 13751: 2001 (E)

Left-Bracket
<u> </u>
Right-Bracket
├

14 Shared Variables

14.1 Informal Introduction

This section gives an informal introduction to the **Shared-Variable-Protocol**. The evaluation sequences which follow this introduction determine the required behaviour of **shared-variables** exactly.

The **Shared-Variable-Protocol** is an **optional-facility**. The **Shared-Variable-Protocol** is described in this document as an interface between cooperating **sessions**, but it may also be used to provide an interface between a **session** and an **auxiliary processor**. An **auxiliary processor** is a program, written in any programming language, that typically provides access to facilities in the underlying operating system, such as file systems.

A **shared-variable** can be shared by precisely two partners, each either a **session** or an **auxiliary processor**.

Conforming-programs that use the **Shared-Variable-Protocol** should document this fact, since the **Shared-Variable-Protocol** is not a **defined-facility** of this International Standard.

Operations

There are conceptually four operations involved in the **Shared-Variable-Protocol**: **offer, retract, set, and reference.**

Sharing

Offer is an operation provided to permit **sessions** to share **shared-variables**. A given **shared-variable** can be shared between at most two **sessions**. Once a **shared-variable** has been shared, assignments made to it in one **session** are visible in the other **session**.

In the evaluation sequences for the **Shared-Variable-Protocol**, sharing is accomplished by appending **tokens** of **class shared-variable** to the **shared-variable-list** which can be accessed by all **active-workspaces** in the **system**.

Any **Shared-Variable-Protocol** operation may cause a **session** to be delayed while another **session** changes a **shared-variable** or adds an item to the **shared-variable-list**.

Sharing begins when one **session** offers to share a **symbol** whose **current-class** is **nil** or **variable**. The **system** adds a new **shared-variable** to the **shared-variable** then replaces the **current-referent** of the offered **symbol**.

There are two forms of offers. A **specific-offer** identifies a particular **session** as the partner; the **general-offer** does not.

Two offers match if both **sessions** use the same **shared-name** and:

At least one of the offers is a specific-offer.

Either specific-offer identifies the other session.

Consequently, two **general-offers** do not match.

Once both **sessions** have shared the **shared-variable**, the **shared-variable** is the **current-referent** of a **symbol** in both **active-workspaces**, and a change made to the **shared-variable** by one **session** will be visible to the other **session**.

Note that two **sessions** can share more than one **shared-variable**, that a given **session** can share **shared-variables** with more that one other **session**, and that each **shared-variable** can be shared by only two **sessions**.

Retracting

Two sessions, coupled by a shared-variable, decouple when either retracts the shared-variable. This leaves the shared-variable in the same state it would have been had the remaining session made a specific-offer to the session that retracted. If the remaining session retracts, the shared-variable becomes inaccessible to the sessions.

Degree of Coupling

All **symbols** in the **symbol-table** of an **active-workspace** have a **degree of coupling**. This is **zero** if the **current-referent** is not a **shared-variable**, **one** if the **current-referent** is a **shared-variable** that is offered to another **session**, and **two** if **current-referent** is a **shared-variable** shared with another **session**.

Naming

For any session, there are two **identifiers** associated with a given **shared-variable**. The first is the **name** of the **symbol** whose **current-referent** is the **shared-variable**; the second is the **shared-name** of the **shared-variable**. The **shared-name** is the name used in matching offers to share.

Only the **name** of the **symbol** is required for invoking any of the **Shared-Variable-Protocol** system functions other than **Shared Variable Offer**.

Setting, Referencing, and Using

Each partner can set and reference the shared-variable.

A shared-variable is set when the operation shared-variable-assignment or shared-variable-indexed-assignment is called. A shared-variable is referenced when the operation shared-variable-reference is called. A shared-variable is used when it is either set or referenced. Both indexed assignment and indexed reference of shared-variables are possible.

Synchronisation

Synchronisation of **sessions** sharing a **shared-variable** is achieved by **access control**. With each **shared-variable** is associated an **access-control-vector**, which is designated here as the *ACV*. The *ACV* is a four element Boolean vector that is accessible to a **session** through the **Shared-Variable-Access-Control-Inquiry** operation.

The order of the elements in the ACV is relative to the viewing **session**. When **session** A, sharing a **shared-variable** with **session** B, views the ACV, its elements have the following meanings:

- If ACV[1] is one, once A has set the shared-variable B must use it before A can set it again.
- If ACV[2] is one, once B has set the shared-variable A must use it before B can set it again.
- If ACV[3] is one, once A has used the shared-variable B must set it before A can reference it again.
- If ACV[4] is one, once B has used the shared-variable A must set it before B can reference it again.

If a **session** attempts to use a **shared-variable** when not permitted to by the above rules, it is delayed until the partner has performed the required intermediate operation.

The rules can be summarised in a three-state state diagram. Let A be the **session** that made the initial offer. Then the states are as follows:

- State 0: The shared-variable has been referenced by the partner of the session that last set it.
- State 1: The shared-variable was last set by A and has not been referenced by B since then.
- State 2: The shared-variable was last set by B and has not been referenced by A since then.

State 0 is the state entered when the **shared-variable** is initially offered. Note that only an explicit assignment is considered a **set** to the **shared-variable**. The act of transferring the **current-content** of a **symbol** to the **shared-value** of a **shared-variable** when an offer

Action attempted:		SET by A			SET by B			REF by A			REF by B		
Current state:	0 1		+				+	3	3		4 +		
	2		+				2	+					4
Next state:		0	1	2	0	1	2	0	1	2	0	1	2

Figure 2: Shared Variable Access Rules.

is made is not a **set**. The state is not affected by retraction and re-sharing. **Figure 2** summarises the state transitions permitted. In the figure, the action is always permitted if there is a + in the row corresponding to the input state. The column in which the + appears determines the output state. A digit in the row indicates that the corresponding element in the ACV as seen by A must be zero.

For example, if the state of a **shared-variable** is 0 and the ACV is 0 0 1 0, **session** A can assign a new value to the variable (SET by A), but must wait until the variable is in state 2 to refer to it (REF by A).

Note that all digits lie along the diagonal in these matrices, indicating that only actions that do not change the state can be blocked by the ACV. That is, the ACV can only prevent a given session from referencing or setting a **shared-variable** twice in a row.

Other Rules

Retraction of all the **shared-variables** in a **session** takes place when a new **workspace** is loaded, the **active-workspace** is cleared, or the **session** ends. A **shared-variable** is automatically **retracted** when the **shared-variable** is erased, expunged, or replaced by a) *COPY* command, or, if it is local, when the function to which it is local terminates execution. The state of sharing is not saved in a library workspace.

The shared variable mechanism must appear to the user to be consistent with one in which there exists a single copy of each **shared-variable**. Specifically, this means:

- When a **shared-variable** is initially offered, its value, if any, becomes the shared-value of the **shared-variable**.
- When an offered **shared-variable** is matched, it retains its value if it has one. If it does not have a value, it takes on the value (if any) in the matcher's workspace.
- When a shared-variable is retracted, its value in the retractor's workspace is the value last set by either partner before the retraction.

Note that, since there is conceptually a single copy across a period of sharing, no use is implied in the above rules for making the value of a **shared-variable** available.

It must appear to the APL user that sessions communicating via shared-variables are

independently scheduled. Any synchronisation between the **sessions** must appear to be by means of the delays caused by accessing **shared-variables**. The procedures used to describe the **Shared-Variable-Protocol** are written as though synchronisation conflicts do not occur when two **sessions** are accessing the **shared-variable-list** or the same **shared-variable**.

A **limit-error** may occur during the retraction of an active **shared-variable** if there is not enough room in a **workspace** for the **shared-value** of the **shared-variable**. When the expected system action is to ignore the variable (as in the case of) *CLEAR* or delocalisation), a **limit-error** is not signalled.

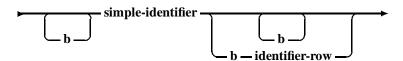
14.2 Definitions

- Identifier-Pair-Array: A character array of rank two in which each row matches the character-diagram identifier-pair-row.

14.3 Diagrams

In the following, **b** stands for **blank**.

Identifier-Pair-Row



14.4 Operations

14.4.1 Primary-Name

Primary-Name in *A*: For *A*, a **character vector** that matches **identifier-pair-row**, the first **simple-identifier** matched in **identifier-pair-row**.

14.4.2 Surrogate-Name

Surrogate-Name in A: For A, a **character vector** that matches **identifier-pair-row**, the last **simple-identifier** matched in **identifier-pair-row**.

14.4.3 Degree-of-Coupling

Degree-of-Coupling of *N*: For **symbol** *N*, a **number** defined as follows:

If the **current-class** of *N* is not **shared-variable**, **zero**.

Otherwise, using the **current-referent** of N, the sum of **session-A-active** and **session-B-active**.

14.4.4 Access-Control-Vector

Access-Control-Vector of N: An operation that, for N a **shared-variable**, is defined as follows:

Using N.

Let ACVA stand for **session-A-ACV** and ACVB stand for **session-B-ACV**.

If **this-session** is **session-A**, return *ACVA* \vee *ACVB*.

If **this-session** is **session-B**, return $(ACVA \lor ACVB)$ [2 1 4 3].

14.4.5 Offer

```
Offer S to P with value V: An operation defined as follows:
```

Let X stand for the **item** in the **shared-variable-list** with the smallest **index** such that

shared-name is S.

session-A is this-session or, if P is not general-offer, general-offer.

session-A-active is zero.

session-B is P.

session-B-active is one.

or

shared-name is S.

session-A is P.

session-A-active is one.

session-B is **this-session** or, if P is not **general-offer**, **general-offer**.

session-B-active is zero.

If such a shared-variable exists,

If **shared-value** is **nil**, set **shared-value** to *V*.

If session-A-active is zero.

Set session-B-event to one.

Set session-A-active to one.

Set session-A to this-session.

If session-B-active is zero,

Set session-A-event to one.

Set session-B-active to one.

Set session-B to this-session.

Return X.

Otherwise, create a new **shared-variable**, *Y*, as follows:

shared-name is S.

```
session-A is this-session.
session-A-active is one.
session-A-ACV is 0 0 0 0.
session-B is P.
session-B-active is zero.
session-B-ACV is 0 0 0 0.
shared-value is V.
state is zero.
session-A-event is zero.
session-B-event is one.
Append Y to the shared-variable-list as a new last item.
Return Y.
```

14.4.6 Retract

```
Retract N: An operation that, for N a shared-variable, is defined as follows: Using N,

If session-A is this-session, set session-A-active to zero.

If session-B is this-session, set session-B-active to zero.

Signal-event.

Clear-event.
```

14.4.7 Shared-Variable-Reset

```
For each item in the shared-variable-list, If session-A is this-session, set session-A-active to zero. If session-B is this-session, set session-B-active to zero. For each symbol \mathbb{N} of the symbol-table of the active-workspace,
```

For each **token** T in the **referent-list** of N,

If T is a shared-variable, set T to a token whose class is nil.

Note: In the model, a shared-variable is never reused. It is effectively discarded once session-A-active and session-B-active are both zero.

14.4.8 Report-State

Shared-Variable-Reset

```
Report-State for N: An operation that, for N a shared-variable, is defined as follows: Using N,

If state is zero, return 0 0 1 1.

Otherwise,

If state is one and session-A is this-session,

or if state is two and session-B is this-session,
```

Return 1 0 1 0. Otherwise return 0 1 0 1.

14.4.9 Signal-Event

Signal-Event for N: An operation that, for N a **shared-variable**, is defined as follows: Using N,

If session-A is this-session, set session-B-event to one. If session-B is this-session, set session-A-event to one.

14.4.10 Clear-Event

Clear-Event for N: An operation that, for N a **shared-variable**, is defined as follows: Using N,

If session-A is this-session, set session-A-event to zero. If session-B is this-session, set session-B-event to zero.

14.5 Shared Variable Forms

14.5.1 Shared Variable Reference

 $Z \leftarrow SHV$

Informal Description: Return the value of a shared variable.

Evaluation Sequence:

Using the **current-content** of SHV,

Let ACV stand for access-control-vector.

Let ASV stand for **report-state**.

Set A to $ACV \land ASV$.

If A[2] or A[3] is **one**, **signal-event**.

Wait until at least one of the following is true:

ACV[3] is zero.

State is two and session-A is this-session.

State is one and session-B is this-session.

If either of the following is true, set **State** to **zero**:

State is two and session-A is this-session.

State is one and session-B is this-session.

Clear-event.

Return Shared-Value.

Note: This operation is all that is required for **shared-variable-indexed-reference**, since the **phrase-evaluator** calls **indexed reference** with a value, not a name.

14.5.2 Shared Variable Assignment

```
Z \leftarrow SHV \leftarrow B
```

Informal Description: Z is B. As a side effect, the value B is assigned to the shared variable SHV.

Evaluation Sequence:

Using the current-content of SHV,

Let ACV stand for access-control-vector.

Let ASV stand for **report-state**.

Set A to $ACV \land ASV$.

If A[1] or A[4] is **one**, **signal-event**.

Wait until at least one of the following is true:

ACV[1] is zero.

State is not one and session-A is this-session.

State is not two and session-B is this-session.

Set shared-value to B.

If session-A is this-session, set state to one.

If session-B is this-session, set state to two.

Clear-event.

Return a **token** whose **class** is **committed-value** and whose **content** is *B*.

14.5.3 Shared Variable Indexed Assignment

```
Z \leftarrow SHV[I] \leftarrow B
```

Informal Description: Z is B. As a side effect, elements of the array SHV selected by the position and values of the arrays in the **index-list** I are replaced by elements of the array B.

Evaluation Sequence:

Using the **current-content** of SHV,

Let ACV stand for access-control-vector.

Let ASV stand for **report-state**.

Set A to $ACV \land ASV$.

If A[1] or A[4] is **one**, **signal-event**.

Wait until at least one of the following is true:

ACV[1] is zero.

session-A is this-session and state is not one.

session-B is this-session and state is not two.

If the current-class of shared-value is nil, signal value-error.

Set C to shared-value.

Search the **form-table** for $Z \leftarrow V[I] \leftarrow B$.

Call the corresponding evaluation sequence, passing C as the value of V.

Set Z to the **token** it returns.

If Z is an **exception**, return Z.

Otherwise,

Set **shared-value** to *C*.

If session-A is this-session, set state to one.

If session-B is this-session, set state to two.

Clear-event.

Return Z.

Note: Indexed assignment of a shared variable constitutes a set, not a reference, of the shared-variable.

14.6 Shared Variable System Functions

14.6.1 Shared Variable Access Control Inquiry

 $Z \leftarrow \square SVC B$

Informal Description: Each row of Z is the **access-control-vector** of the shared variable named by the corresponding row of B.

Evaluation Sequence:

14.6.2 Shared Variable Query

```
Z \leftarrow \sqcap SVQ B
```

Informal Description: *Z* is one of the following:

An array representing sessions offering to share variables with this session.

An array representing the names of shared variables offered to this session by another session.

Evaluation Sequence:

```
If session-identification-type is character, and B is the general-offer,
  Set Z to 0 00' '.
  Using each item of the shared-variable-list,
     If
        session-A is this-session and
        session-A-active is zero and
       session-B-active is one,
          Set S to session-B.
          Set Z to ((\rho Z) \lceil 0, \rho S), [1]((-1 \uparrow \rho Z) \lceil \rho S) \uparrow S evaluated with index-
            origin set to one.
     If
        session-A-active is one and
        session-B is this-session and
        session-B-active is zero.
          Set S to session-B.
          Set Z to ((\rho Z) \lceil 0, \rho S), [1] ((-1 \uparrow \rho Z) \lceil \rho S) \uparrow S evaluated with index-
            origin set to one.
If session-identification-type is numeric, and B is zero, signal domain-error.
If session-identification-type is numeric, and B is 10,
  Set Z to 10.
  Using each item of the shared-variable-list,
     If all of the following are true,
        session-A is this-session.
       session-A-active is zero.
       session-B-active is one.
     append session-B as a new last element to the vector \mathbb{Z}.
     If all of the following are true,
        session-A-active is one.
       session-B is this-session.
       session-B-active is zero.
     append session-A as a new last element to the vector Z.
  Return ((Z \iota Z) = \iota \rho Z)/Z.
Otherwise,
  If B is not a session-identification, signal domain-error.
  If B is this-session, signal domain-error.
```

```
Set Z to the empty character array of rank two.

Using each item of the shared-variable-list,

If all of the following are true,
session-A is this-session.
session-A-active is zero.
session-B is B.
session-B-active is one.
append shared-name as a new last row to the identifier-array Z.

If all of the following are true,
session-A is B.
session-A-active is one.
session-B is this-session.
session-B-active is zero.
append shared-name as a new last row to the identifier-array Z.

Return Z.
```

14.6.3 Shared Variable Degree of Coupling

```
Z \leftarrow \sqcap SVO B
```

Informal Description: Z is a numeric vector in which each element is the degree of coupling of the symbol named by the corresponding row of B.

Evaluation Sequence:

```
If the rank of B is greater-than two, signal rank-error.

Set B1 to (-2+1, 0B) 0B.

If any item of the ravel-list of B1 is not a character, signal domain-error.

Set Z to (1+0B1) 0.

For every I in 1+0B1,

Let B2 stand for row I of B1.

If B2 matches identifier-row,

Set I to the primary-name in B2.

Set item I of the ravel-list of I to degree-of-coupling of the symbol-named-by I.

Otherwise, signal domain-error.

Return I.
```

Note: Degree-of-coupling is a property of symbols, not shared-variables.

14.6.4 Shared Variable Offer

 $Z \leftarrow A \sqcap SVO B$

Informal Description: Z is a numeric vector representing the **degree-of-coupling** of the identifiers in the rows of the **identifier-pair-array** B after the operation. As a side effect, dyadic $\square SVO$ shares variables with another **session**. B is an **identifier-pair-array** representing the names of variables to be shared. A is a **session-identification**, representing the **session** with which the names in B are to be shared.

Evaluation Sequence:

If the rank of B is greater-than two, signal rank-error.

If any **item** of the **ravel-list** of *B* is not a **character**, signal **domain-error**.

If A is not a valid **session-identification**, signal **domain-error**.

If **this-session** is *A*, signal **domain-error**.

Set B1 to $(-2 \uparrow 1 1, \rho B) \rho B$.

For every I in $11 \uparrow \rho B1$,

Let B2 stand for **row** I of B1.

If B2 does not match identifier-pair-row, signal domain-error.

Otherwise,

Set N to the **primary-name** in B2.

Set S to the surrogate-name in B2.

If the current-class of N is neither shared-variable, variable nor nil, signal domain-error.

If the **current-class** of *N* is **variable** or **nil**,

Set V to the **current-referent** of N.

Set Y to offer S to A with value V.

Set the **current-referent** of N to Y.

Return $\square SVOB$.

Note: This operation does not allow for multiple **session-identifications**.

14.6.5 Shared Variable Retraction

 $Z \leftarrow \sqcap SVR B$

Informal Description: Elements of Z hold the degree of coupling that the identifier represented in the corresponding row of the **identifier-array** B had, before this retraction. As a side effect, $\square SVR$ retracts any shared-variables named in B.

Evaluation Sequence:

If the rank of B is greater-than two, signal rank-error. Set B1 to $(2+1 1, \rho B) \rho B$. Set B1 to $(2+1 1, \rho B) \rho B$. Set B1 to B1 to B1. If B1 is an error, return B1. For every B1 in B1 the B1, Let B2 stand for row B1 of B1. If B2 matches identifier-row, Set B1 to the simple-identifier in B1. If the current-class of B1 is shared-variable, Set B1 to the current-referent of B1. Set the current-referent of B1. Set the current-referent of B1. Set the current-referent of B1. Otherwise, signal domain-error. Return B1.

14.6.6 Shared Variable Access Control Set

 $Z \leftarrow A \sqcap SVC B$

Informal Description: Z is the **access-control-vector** of each of the shared variables represented by the **identifier-array** B when $\square SVC$ completes. As a side effect, $\square SVC$ sets to A the contribution this session makes to the combined access control vector of the shared variables.

Evaluation Sequence:

If the rank of B is greater-than two, signal rank-error.

Set B1 to $(-2 \uparrow 1 \ 1, \rho B) \rho B$.

If the rank of A is greater-than two, signal rank-error.

If the first **item** of $(-1 \uparrow 4, \rho A)$ is not **four**, signal **length-error**.

If A is a scalar or a vector, set A1 to $((1 \uparrow \rho B1), 4) \rho A$.

Otherwise, set A1 to A.

If $1 \uparrow \rho A1$ is not the same as $1 \uparrow \rho B1$, signal **length-error**.

If any **item** of the **ravel-list** of B1 is not a **character**, signal **domain-error**.

If any **item** of the **ravel-list** of A1 is not a **near-Boolean**, signal **domain-error**.

Set A2 to the **Boolean-array-nearest-to** A1.

For every I in $11 \uparrow \rho B1$,

Let B2 stand for **row** I of B1.

If B2 matches **identifier-row**,

Set N to the **simple-identifier** in B2.

If the **current-class** of *N* is **shared-variable**,

Let ACV stand for **row** I of A2.

Using the **current-content** of *N*,

If **session-A** is **this-session**, set **session-A-ACV** to ACV.

If **session-B** is **this-session**, set **session-B-ACV** to $ACV[2 \ 1 \ 4 \ 3]$.

Signal-event.

Otherwise, signal domain-error.

Return $\square SVCB$.

14.6.7 Shared Variable State Inquiry

 $Z \leftarrow \sqcap SVS B$

Informal Description: Each row of Z is a boolean 4-element representation of the **state** of the shared variable named by the corresponding row of B, as seen by **this-session**. It may be combined with the **access-state-vector** to show which accesses are currently possible for the **shared-variable**.

Evaluation Sequence:

Return $((^1 \downarrow \rho B), 4) \rho Z$.

If the rank of B is greater-than two, signal rank-error. Set B1 to $(-2+1, \rho B)\rho B$. If any item of the ravel-list of B1 is not a character, signal domain-error. Set Z to $((1+\rho B1), 4)\rho 0$. For every I in $1+\rho B1$, evaluated with index-origin set to one, Let B2 stand for row I of B1. If B2 matches identifier-row, Set I to the simple-identifier in I in I to the current-class of I is shared-variable, Set row I of I to report-state of the current-referent of I. Otherwise, signal domain-error.

14.6.8 Shared Variable Event

```
\Box SVE \leftarrow B
Z \leftarrow \Box SVE
```

Informal Description: Assignment to the system variable starts a timer. The specified number of seconds is the maximum time the program will wait for a **shared-variable-event**.

Reference to the system variable suspends program execution until a **shared-variable-event** occurs, or until maximum time has elapsed. Then *Z* is the remaining time in the timer, and all outstanding **shared-variable-events** are cleared.

A **shared-variable-event** occurs whenever a new offer to share a variable occurs, or whenever certain operations are performed on a shared variable by the partner (see the various Evaluation Sequences in this chapter).

Evaluation Sequence:

For form $\square SVE \leftarrow B$

Set T0 to **current-time** in seconds.

If B is not a **scalar**, signal **rank-error**.

If any **item** of the **ravel-list** of B is not a **nonnegative-number**, signal **domain-error**.

Set **event-time** to T0 **plus** B.

For form $Z \leftarrow \square SVE$

Wait until at least one of the following is true:

current-time is not less-than event-time.

For any item N in the **shared-variable-list**,

Using N,

If session-A is this-session, session-A-event is one.

If session-B is this-session, session-B-event is one.

Return event-time minus current-time.

For each item N in the **shared-variable-list**, **clear-event** for N.

15 Formatting and Numeric Conversion

15.1 Introduction

The mapping of arrays of numbers into arrays of characters is referred to as formatting.

Note: The form, field width, and number of digits in a formatted number can be specified explicitly through the use of **dyadic format**, or left to be chosen by the system. **Conforming-programs** that depend upon a particular selection of formatting parameters should employ **dyadic format**.

15.2 Numeric Conversion

Note: Conversion involves the transformation of lists of characters to and from numbers. This conversion is performed by the implementation-algorithms numeric-input-conversion and numeric-output-conversion.

The accuracy and rounding properties of **numeric-input-conversion** and **numeric-output-conversion** are interdependent. The most significant property of these **implementation-algorithms** is the following:

For any **numeric scalar** X, when **print-precision** is set to **full-print-precision**, X shall be the same as $\Phi \Phi X$.

15.2.1 Numeric-Input-Conversion

Numeric-Input-Conversion of *C*.

Informal Description: Numeric input conversion converts numeric values represented in decimal notation as **lists** of **characters** into equivalent **numbers**—numeric quantities whose format is **implementation-defined**.

In the following, C stands for a **real-number**, or the **real-part** or **imaginary-part** of a **complex-number**.

Evaluation Sequence:

Let \mathcal{L} and \mathcal{G} be respectively the most negative and the most positive abstract numeric values represented by all **numeric-scalar-literals** producible by **numeric-output-conversion** for all **numbers** and all valid values of **print-precision**.

Let T be the set of all **lists** of **characters** that match the diagram **numeric-scalar-literal** and have abstract numeric values lying in the closed interval between L and G.

If C is not an element of T, signal **limit-error**.

Otherwise.

Let \mathbf{c} stand for the abstract numeric value of C.

If c is the same as some **number** N, return N.

If **c** is greater than **positive-number-limit**, return **positive-number-limit**.

If **c** is less-than **negative-number-limit**, return **negative-number-limit**.

Otherwise, **c** lies between two **numbers**. Return one that permits **numeric-input-conversion** to satisfy the following property:

For any two **numeric-scalar-literals** C1 and C2, if the abstract numeric value of C1 is less than the abstract numeric value of C2, then **numeric-input-conversion** of C1 is not **greater-than numeric-input-conversion** of C2.

Examples:

Note: In many programming languages, the syntactic form of a literal is used to determine the datatype of the corresponding number. This is not so in APL. Therefore, it is important that the input conversion routine not make distinctions between different forms of the same abstract numeric value. 2.5E4 should produce the same **number** as 25000 does and 002 should produce the same number as 2 does.

Every number can be obtained by converting some numeric-scalar-literal.

The set T in the evaluation sequence above may include numbers with real-part or imaginary-part having abstract numeric values less than negative-number-limit or greater than positive-number-limit, due to rounding by numeric-output-conversion.

There is no **number** \mathbb{N} for which the expression $\mathfrak{T} \mathbb{N}$ signals a **limit-error**.

15.2.2 Numeric-Output-Conversion

E Numeric-Output-Conversion of N.

E Numeric-Output-Conversion of N to P places.

Informal Description: Numeric output conversion converts numeric values represented as **numbers**—numeric quantities whose format is **implementation-defined**—into the same numeric quantities represented in decimal notation as **lists** of **characters**.

For E a **character-diagram**, and N and P **numbers**, numeric output conversion returns a **numeric-scalar-literal** that matches E and represents the abstract numeric value of the **real-part** of N either to P decimal places, if P is specified, or else to **print-precision** places.

Note: Any imaginary-part of the numeric input to Numeric-Output-Conversion is simply ignored; only the real-part of a complex-number is used.

Evaluation Sequence:

If E is minimal-decimal-exponential,

Return a **list** of **characters** whose abstract numeric value is a good approximation of N.

Note: The choice of approximation is **implementation-defined**, but can be characterised as follows:

Let DS be the set whose members D are lists of characters that match minimal-decimalexponential and have at most print-precision digits to the left of the exponent-marker. Let DT be the subset of DS for which $|N-\Phi D|$ is minimal.

Let DM be the subset of DT whose members have the fewest digits.

Choose the member of DM whose abstract numeric value has the largest magnitude.

If E is fixed-decimal,

Let DF be the set whose members are **lists** of **characters** that match **fixed-decimal** and have P digits to the right of the units digit.

Return a member D of DF for which $|N-\Phi D|$ is minimal.

If E is **decimal-exponential**,

Let DE be the set whose members are **lists** of **characters** that match diagram **decimal-exponential** and have P digits to the left of the **exponent-marker**.

Select a member D of DE for which $|N-\Phi D|$ is minimal.

Let *W* stand for **exponent-field-width**.

Return $(W+D \iota ' E') \uparrow D$.

Note: When P is not specified, E is always minimal-decimal-exponential.

Fixed-decimal and **decimal-exponential** cases are used only by **dyadic format**. They are permitted, but not required, to return results of arbitrarily high precision. They should, however, return accurate results if no more than **full-print-precision** significant digits are requested.

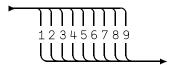
15.3 Diagrams

Note: The following **character-diagrams** characterise the outputs of both monadic and dyadic format as subsets of the set of **lists** of **characters** that match **numeric-scalar-literal**.

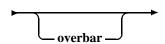
Zero-Digit



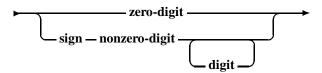
Nonzero-Digit



Sign



Decimal-Integer



Examples:

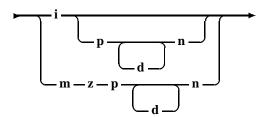
-44 0 622 1066 1215 1415 1685 1750 1776 1812 1867 1944

In the following diagrams,

- **b** stands for **blank**.
- d stands for digit.
- e stands for exponent-marker.
- i stands for decimal-integer.

- **m** stands for **overbar**.
- **n** stands for **nonzero-digit**.
- **p** stands for **dot**.
- s stands for sign.
- z stands for zero-digit.

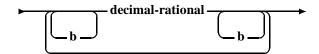
Decimal-Rational



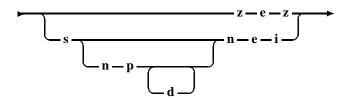
Examples:

98.6 212 -2.001 -0.00000000001

Decimal-Rational-Row



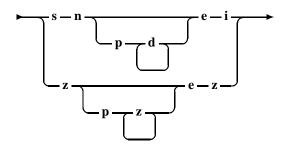
Minimal-Decimal-Exponential



Examples:

-1.234567E-890 1E1 -1.1111E1 1E-11111 0E0

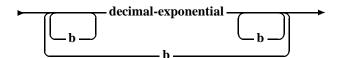
Decimal-Exponential



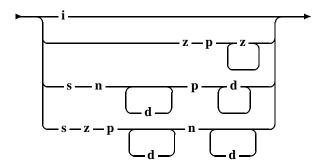
Examples:

-1.23456700000*E*-890 0.00000000*E*0 -1.0000*E*-10

Decimal-Exponential-Row



Fixed-Decimal



Examples:

 $0\ \ ^{-1}\ 12\ 0.00000000\ 12.30\ 0.1235\ \ ^{-0}.1235\ 0.1\ \ ^{-0}.00010000$

15.4 Operations

15.4.1 Monadic Format

 $Z \leftarrow \Phi B$

Informal Description: Z is a **character** array. If B is a **character** array, then Z is B. If B is a numeric array, then B is a **character** array arranged so that each row of B is a decimal representation of the corresponding row of B. The output form chosen for a column of B is determined by **print-precision** and the values of elements in that column. Uses **print-precision**

Evaluation Sequence:

If print-precision is nil, signal implicit-error.

If the **type** of B is **character**, return B.

If *B* is **empty**, return an **array** Z0 such that the **type** of Z0 is **character**, the **shape-list** of Z0 is the **shape-list** of *B*, and the **ravel-list** of Z0 is the **empty-list**.

If the rank of B is less-than two,

return, Φ ($^{-}2\uparrow1$ 1, ρB) ρB .

If the last **item** of the **shape-list** of *B* is not **one**,

return $(\overline{\tau} ((1 \downarrow \rho B), 1) \uparrow B), ' ', \overline{\tau} ((-\rho \rho B) \uparrow 1) \downarrow B$

Otherwise, choose U, E, and W, and return an **array** Z for which the **type** of Z is **character**, the **shape-list** of Z is $(-1 + \rho B)$, W, and one of the following statements is true:

Every **row** of Z is the **decimal-exponential-row** equivalent of the **minimal-decimal-exponential numeric-output-conversion** of the corresponding **row** of B, and has its **exponent-marker** at position E and its first digit at position U.

Every **row** of Z is the **decimal-rational-row** equivalent of the **minimal-decimal-exponential numeric-output-conversion** of the corresponding **row** of B, and has its units digit at position U.

Every **row** of Z is formatted independently for each element B1 of B. The results are padded with blanks to width W.

If B1 is a **character**, use B1.

If B1 is a **number**, and the **imaginary-part** of B1 is 0, use $(\overline{P}B1)$.

If B1 is a **number**, and the **imaginary-part** of B1 is not 0, then

Let B2 be the **real-part** of B1.

Let B3 be the **imaginary-part** of B1.

Let CI be the character representation of the Complex-Marker.

Use $((\Phi B2), CI, \Phi B3)$.

Note: This form of output must match the diagram numeric-scalar-literal.

Note: The quantities W, U and E and the choice of formatting form are not standardised here, since this is currently impractical. However, the following algorithm is recommended for all future implementations of APL:

If any of the following conditions holds for any element X of B

 $\it X~is~$ greater-than positive-counting-number-limit.

X is less-than negative-counting-number-limit.

The decimal-rational equivalent of the minimal-decimal-exponential numeric-output-conversion of X would have more than print-precision significant digits to the left of the decimal point.

The decimal-rational equivalent of the minimal-decimal-exponential numeric-output-conversion of X would have more than five zero digits to the right of the decimal point and to the left of the first nonzero digit.

then select decimal-exponential form.

Otherwise, select decimal-rational form.

If decimal-exponential form is selected,

Let S be one if any element of B is a negative-number.

Let M be the maximum number of characters to the left of the exponent-marker in the minimal-decimal-exponential numeric-output-conversion of any element of B.

Let N be the maximum number of characters to the right of the exponent-marker in the minimal-decimal-exponential numeric-output-conversion of any element of B.

Set U to S+1.

Set E to M+1.

Set W to M+1+*N*.

If decimal-rational form is selected,

Let M be the maximum number of characters to the left of the units position in the decimalrational equivalent of the minimal-decimal-exponential numeric-output-conversion of any number in B.

Let N be the maximum number of characters to the right of the units position in the decimalrational equivalent of the minimal-decimal-exponential numeric-output-conversion of any number in B.

Set U to M+1.

Set W to M+1+N.

Note: There is no uniformity in the display of an array containing mixed types, nor in the display of a column containing complex numbers. Suggested methods for displaying these are as follows: For an array containing mixed types, format each simple array, and provide enough space in each column to contain the widest element.

For a column containing at least one complex number, right-justify each number in the column, making no attempt to align other numeric elements.

Examples:

```
□PP←10
    '|',(φ5 <sup>6</sup> 7 8 9 10°.÷ι5),'|'
 5 2.5 1.666666667 1.25 1
   -3
 -<sub>6</sub>
                     1.5
                           1.2
| 7
   3.5 2.333333333 1.75 1.4
8
   4
         2.6666666667 2
                           1.6
9 4.5 3
                     2.25 1.8
10 5
         3.333333333 2.5
```

```
D \leftarrow 0.7 \ 0.8 \ 0.9, 7 \ 8 \ 9 \div 10
      \Phi D
0.7 0.8 0.9 0.7 0.8 0.9
      D-\Phi \Phi D
0 0 0 0 1.387778781E<sup>-</sup>17 0
      □PP←999
0.7 0.8 0.9 0.7 0.799999999999999 0.9
      D-\Phi \Phi D
0 0 0 0 0 0
      □PP←4
      M \leftarrow 0.78901 \times (10 \times 16) \circ . \times 1 1E^{-}10 1E^{-}6 0.01 0.1
      \Phi M
7.89E0 \ 7.89E^{-}10 \ 0.00000789
                                      0.0789 7.89E<sup>-</sup>1
7.89E1 7.89E 9 0.0000789
                                      0.789 7.89E0
7.89E2 \ 7.89E^{-}8 \ 0.000789
                                      7.89
                                               7.89E1
7.89E3 7.89E^{-7} 0.00789
                                     78.9
                                               7.89E2
                                   789
7.89E4 7.89E^{-6} 0.0789
                                               7.89E3
7.89E5 7.89E^{-}5 0.789
                                  7890
                                               7.89E4
      Φ5ρM
7.89 7.89E-10 0.00000789 0.0789 0.789
      Φ1 5ρM
7.89 7.89E-10 0.00000789 0.0789 0.789
```

Note: When print-precision has a value in its internal-value-set greater than or equal to full-print-precision, the result produced for N is unique and is unaffected by increases in the value of print-precision. At such settings of print-precision, N is the same as $\mathfrak{L} \mathbb{R} N$ for any numeric scalar N.

15.4.2 Dyadic Format

```
Z \leftarrow A \neq B
```

Informal Description: *Z* is a **character** array representing the numbers in *B* formatted according to the specifications in *A*. *A* consists of pairs of integers. One pair is associated with each column of *B*. The first integer in the pair specifies the field width—the number of columns in the character result—for the formatted numeric values.

If A has only two elements then all elements of B are formatted according to A.

Note: For complex elements of B, the precision specifier applies to each part of the representation.

A non-negative second integer indicates that the field is to have **fixed-decimal** form and gives the number of digits to the right of the decimal point, with zero indicating no decimal places and no decimal point. A negative second integer indicates that the field is to have decimal exponential form; the absolute value specifies the number of digits in the mantissa.

Evaluation Sequence:

```
If the rank of A is greater-than one, signal rank-error.
If \rho, A is not 2 \times 1 \uparrow 1, \rho B
  If \rho A is two, return ((2 \times 1 \uparrow \rho B) \rho A) \Phi B.
   Otherwise, signal length-error.
If any item of the ravel-list of A is not a near-integer, signal domain-error.
If any item of the ravel-list of B is not a real-number, signal domain-error.
Set A1 to the integer-array-nearest-to A.
If B is empty.
   Set W to +/((\rho A_1)\rho_1 0)/A_1.
  Return ((^1 \downarrow \rho B), W) \rho'.
If \rho A1 is not two, return ((2\uparrow A1) \bar{\sigma}((-1 \downarrow \rho B), 1) \uparrow B), (2 \downarrow A1) \bar{\sigma}((-1 \downarrow \rho B), 1) \uparrow B)
 \rho \rho B) \uparrow 1) \downarrow B.
If B is not a scalar,
   If A1[1] is not a positive-integer, signal domain-error.
   Return a character array Z such that the shape-list of Z is ( 1 \downarrow \rho B), A1[1] and
    the ravel-list of Z has the property that each vector-item along-axis \rho \rho Z of Z is
    A1 \Phi B0, where B0 is the corresponding (scalar) item of (-1 + \rho B) \rho B.
If B is a scalar.
   If A1[2] is not less-than zero, set R to the
   fixed-decimal numeric-output-conversion of B to A1[2] places.
   Otherwise, set R to the decimal-exponential numeric-output-conversion of B to
    A1[2] places.
   If A1[1] is less-than \rho R, signal domain-error.
   Otherwise, return (-A1[1]) \uparrow R.
```

Examples:

- In the following examples, **exponent-field-width** is 3.

```
D \leftarrow 1 \quad 0.1 \quad 0 \quad 0.1 \quad 1 \quad 0. \times 500.5
| \mid \mid , (9 \quad 37 \quad 16 \quad 35 \quad 13 \quad 0 \quad \Phi D), \mid \mid \mid
| \quad 5.00E \quad 1 \quad 5E \quad 1 \quad 0.500 \quad 0.5 \quad 1 \mid
| \quad 5.00E \quad 2 \quad 5E \quad 2 \quad 0.050 \quad 0.1 \quad 0 \mid
| \quad 0.00E0 \quad 0E0 \quad 0.000 \quad 0.0 \quad 0 \mid
| \quad 5.00E \quad 2 \quad 5E \quad 2 \quad 0.050 \quad 0.1 \quad 0 \mid
| \quad 5.00E \quad 1 \quad 5E \quad 1 \quad 0.500 \quad 0.5 \quad 1 \mid
| \quad \rho \quad 10 \quad 20 \quad 40 \quad 80 \quad \Phi 0 \quad 40 \quad 1
0 \quad 15
| \quad 31 \quad 0 \quad \Phi \quad 2*100
1267650600228229401496703205376
| \quad 31 \quad 20 \quad \Phi \quad .07
0.070000000000000000666
| \quad 41 \quad \Phi \quad .99 \quad .89 \quad 07.5 \quad 11.5
| \quad 1.0 \quad 0.9 \quad 0.0 \quad 7.511.5
```

Note: Like other axis lengths, the field width A[1] is limited only by **index-limit**.

16 Input and Output

16.1 Introduction

Note: A user interacts with an APL system through a **session**, an abstraction that represents a hypothetical machine capable of carrying out the evaluation sequences in this International Standard.

The protocol for the use of a session takes the form of a dialogue: the user makes an **entry** and passes control to the APL system. The system processes the entry, produces a response, and returns control to the user.

The user makes entries on a keyboard, and obtains responses by seeing them presented on a display-device. The combination of a keyboard and a display-device is intended to represent, abstractly, a terminal. This International Standard does not require the use of any particular terminal, keyboard, display or method of encoding **characters**. The ISO 2022.2 APL character encoding reproduced in Annex A is widely used.

The display-device is considered a window into an unbounded sequence of past entries and responses. The mapping between this unbounded sequence and the display-device is not specified by this International Standard.

The system obtains entries as **tokens** of **class constant** or **interrupt** by calling the session operation **read-keyboard**. It presents responses by calling the session operation **display** with a **token** of class **result**. The transformations from entry to token and from token to response are not specified by this International Standard, but certain desirable characteristics of these transformations are given as comments.

Because the dialogue protocol described here alternates between user and system, the session is always in one of two logical keyboard-states, open-keyboard or locked-keyboard, depending upon whether the user is making an entry or the system is producing a response. The user can request a change in keyboard-state from locked-keyboard to open-keyboard by using the signal-attention facility. In addition, the user can force a change in keyboard-state by signalling interrupt.

16.2 Definitions

16.2.1 User Facilities

A user should have the following facilities.

 Edit-Actions: A collection of implementation-defined facilities that permit the user to make entries.

Note: Typical **edit-actions** include inserting, deleting, replacing and superimposing graphic symbols within the current entry and combining display-lines from the presentation-space with the current entry.

- Enter: A facility, available to the user when the terminal is in open-keyboard state, for changing the keyboard-state to locked-keyboard and returning to the caller of read-keyboard either an interrupt or a character vector. Enter sets the value of the session attribute attention-flag to zero.
- Signal-Attention: A facility, available to the user when the terminal is in lockedkeyboard state, for setting the value of the session attribute attention-flag to one.
- Signal-Interrupt: A facility, available to the user at all times, for interrupting any
 evaluation sequence, causing it to return an interrupt token to the current caller of
 evaluate-line. Any atomic operation so interrupted has no effect on the state of the
 active-workspace.

Note: Signal-interrupt can be issued in **open-keyboard** state and in **locked-keyboard** state, although the particular means of invoking the operation may be different for the two states. Transmitting the superimposition of the graphics for O U and T is a typical means of signalling **interrupt** in **open-keyboard** state.

16.2.2 Implementation Algorithms

Read-Keyboard: An operation, available to the system when the session is in locked-keyboard state, that returns either an interrupt token or a token of class constant and content a character vector.

Note: Read-keyboard should display current-prompt, then change the keyboard-state to open-keyboard to permit the user to use edit-actions to enter a list of characters or to signal an interrupt.

A concept long honoured in APL is that of visual fidelity: the graphic symbols that appear on the display as the result of edit-actions should correspond exactly to those returned to the system as the result of read-keyboard.

- **Presentation Space**: an unbounded sequence of past entries and responses.
- **Display** T: An operation, available when the session is in **locked-keyboard** state, for presenting the **token** T on the **display-device**. T is either a **value** or an **error**.

```
Note: The following is a suggested, but not required, evaluation sequence for display T:
  If quote-quad-prompt is not the empty character vector,
    Append to the presentation-space as a new last item a display-line whose graphic symbols
      represent the characters in quote-quad-prompt.
     Set quote-quad-prompt to the empty character vector.
  If T is an error, append to the presentation-space as a new last item an indication of
     The class of the error.
     The point in current-line at which evaluation was stopped.
     If the mode is defined-function, the function-name and current-line-number of the current-
      context.
  If T is a value,
    Let A stand for the content of T.
    If A is numeric, display \Phi A.
     Otherwise,
     If A is a scalar or vector, append to the presentation-space as a new last item a display-line
      whose graphic symbols represent the characters in A.
     Otherwise,
       Set I to zero.
       Set M to ((\times/^-1 \downarrow \rho A), -1 \uparrow \rho A) \rho A
       Repeat the following 1 \uparrow \rho M times:
          Set I to I+1.
          Display row I of M.
          Append (-1+\rho\rho A) \mid +/\wedge \setminus 0=\phi (-1+\rho A) \top I blank display-lines to the end of the
           presentation-space.
       (End of repeated block)
```

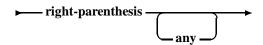
16.2.3 Prompts

Note: The following **implementation-defined arrays** are used to indicate to the user the type of entry required.

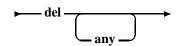
- Indent-Prompt: An implementation-defined character array, typically a vector of six blanks.
- Quad-Prompt: An implementation-defined character array, whose first row typically begins with
 : and whose second row is the indent-prompt.
- Function-Definition-Prompt: An implementation-defined character array, typically
 a vector beginning with [nn] where nn is a line-number.

16.3 Diagrams

System-Command-Line



Function-Definition-Line



16.4 Operations

16.4.1 Immediate-Execution

Immediate-Execution with *X*.

Informal Description: Immediate-execution mode is the primal mode, the first evaluation sequence processed by the system to enter into a dialogue with a user. *X* is either an **error token** or **nil**.

The system also calls **immediate-execution** when an **error** or the **attention-flag** causes a defined function to be suspended. In this case, **immediate-execution** displays a status indication before prompting for input.

Evaluation Sequence:

```
Display X.

Repeat:
Set current-prompt to indent-prompt.
Set E to read-keyboard.

If E is a constant,
Let Q stand for the content of E.

If Q matches system-command-line, set T to evaluate-system-command Q.

If Q matches function-definition-line, set T to evaluate-function-definition-request Q.

Otherwise,
Generate a new context in which
mode is immediate-execution,
current-line is Q,
current-function is Q or Q \cap Q \cap Q current-line-number is one,
```

current-statement is the empty **list** of **tokens**, and **stack** is the empty **list** of **tokens**.

Append the new **context** to the **state-indicator** of the **active-workspace** as a new first **item**.

Set T to **evaluate-line**.

Remove the **first-item** in the **state-indicator** of the **active-workspace**.

If T is a **branch**, a **clear-state-indicator**, or an **escape** and the **state-indicator** of the **active-workspace** is not an **empty-list**, return T.

If T is an **error** or a **constant**, **display** T.

(End of repeated block)

16.4.2 Quad Input

 $Z \leftarrow \square$

Informal Description: Z is an array provided by the user in response to a prompt.

Evaluation Sequence:

```
Repeat:
```

Set current-prompt to quad-prompt.

Set E to **read-keyboard**.

If E is an **interrupt**, return E.

Let Q stand for the **content** of E.

If some element of Q is not blank,

Generate a new context in which

mode is quad-input,

current-line is Q,

current-function is 0 0p 1 1,

current-line-number is one,

current-statement is the empty list of tokens, and

stack is the empty list of tokens.

Append the new **context** to the **state-indicator** of the **active-workspace** as a new first **item**.

Set Z to **evaluate-line**.

Remove the **first-item** in the **state-indicator**.

If Z is **escape**, signal **unwind**.

If Z is **unwind** or **clear-state-indicator**, return Z.

If Z is a value, return a token with class constant and content that of Z.

If Z is an **error**, **display** Z.

Otherwise, display value-error. (see note)

(End of repeated block)

Note: The **display value-error** line in the above evaluation sequence is introduced to permit certain **consistent-extensions**. In the evaluation sequence, errors are **displayed** rather than **signalled**, since signalling would terminate the Repeat loop and force Return.

Quad-input returns only tokens of class escape and unwind (from \rightarrow), clear-state-indicator (from SIC), constant (from $2 \times .5$) and interrupt. It reports an error and reprompts for all other results. Note that it reprompts without an error report if the input was empty or all blank.

16.4.3 Quote Quad Input

 $Z \leftarrow \Pi$

Informal Description: Z is a **character vector**. Input in response to \square is treated as a **character** value. Z is a vector whose length is the number of positions from the left margin up to the rightmost **character** of the input, including explicitly entered trailing blanks.

Evaluation Sequence:

Set **current-prompt** to **quote-quad-prompt**. Set **quote-quad-prompt** to the **empty character vector**. Set E to **read-keyboard**. Return E.

Note: The behaviour required here differs from that of some existing systems, in which a single character response to \square is returned as a scalar.

16.4.4 Quad Output

$$Z \leftarrow \square \leftarrow B$$

Informal Description: Z is B. As a side effect, the array B is displayed on the terminal.

Evaluation Sequence:

Display B.

Return a **committed-value** with **content** B.

16.4.5 Quote Quad Output

```
Z \leftarrow \Pi \leftarrow B
```

Informal Description: Z is B, and **quote-quad-prompt** is set to B.

Evaluation Sequence:

If any **item** of the **ravel-list** of *B* is not a **character**, signal **domain-error**.

If the rank of B is greater-than one, signal rank-error.

If the **count** of B is **greater-than quote-quad-output-limit**, signal **limit-error**.

If **quote-quad-prompt** is not empty, signal **limit-error**.

Set quote-quad-prompt to B.

Return a **committed-value** with **content** *B*.

Example:

```
∇Z←PROMPT X
[1] □←X←, ▼X
[2] Z←(ρX) ↓□
∇
PROMPT '1 CLEANSPACE DATE? '
1 CLEANSPACE DATE? 1966-11-27
1966-11-27
```

Note: Quote-quad output is difficult to standardise because it is implemented in several different ways in existing systems, and, in each system, is heavily used. There is, however, sufficient commonality to permit the inclusion of this restricted form of quote-quad output. Quote-quad output on a **conforming-implementation** should support the prompt function shown in the example.

17 System Commands

17.1 Introduction

Note: Any line of input in **immediate-execution** whose first non-blank **character** is a right parenthesis is considered to be a system command.

17.2 Definitions

- Global-Referent of T: For T a classified-name, the last-item in the referent-list of the symbol-named-by T.
- Global-Context The last context in the state-indicator of the active-workspace.
- Library-Workspace-Named A: For workspace-identifier A, the item of the library whose owner is this-owner, and whose workspace-name is A.
- Already-Exists: A workspace already-exists if its existential-property is present.
- Does-Not-Exist: A workspace does-not-exist if its existential-property is absent.

Note: The model of libraries used in this International Standard assumes that a workspace object exists in the library for all possible workspace-identifiers. Workspaces that have been dropped or that have never been saved are distinguished from those that have been saved by the setting of their existential-property: the existential-property of saved workspaces is present, while that of unsaved workspaces is absent.

Attempt-to-Erase A: For A, a simple-identifier, an operation defined as follows:
 If A is globally-erasable,

If the **current-class** of A is **shared-variable**, **retract** A.

Set A to nil.

Return **command-complete**.

Otherwise, signal **not-erased**.

Copy A from W: For simple-identifier A, and workspace-identifier W, an operation defined as follows:

Let *GA* stand for the **global-referent** of *A* in the **active-workspace**.

Let GB stand for the **global-referent** of A in the **library-workspace-named** W.

Set Z to attempt-to-erase GA.

If Z is an **exception**, signal **not-copied**.

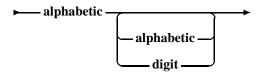
If GB is a **shared-variable**, set GA to the **shared-value** of GB.

Otherwise, set GA to GB.

Return **command-complete**.

17.3 Diagrams

Workspace-Identifier



Example:

CLEANSPACE

Alphabetic



17.4 Operations

17.4.1 Evaluate-System-Command

Evaluate-System-Command Q.

Note: The forms of system commands are not regular enough to permit much generalisation. Therefore, the forms themselves are used as the key to the evaluation sequences.

Evaluation Sequence:

Find an entry that matches Q in the **system-command evaluation sequences**.

If no such entry is found, signal **incorrect-command**. Otherwise, call the corresponding evaluation sequence. Return the **token** it returns.

17.5 Diagrams and Evaluation Sequences

In the following,

- p stands for permitted-blanks,
- r stands for required-blanks,
- w stands for workspace-identifier, and
- a stands for simple-identifier.

Clear Active Workspace

-p - p - CLEAR - p - CLEAR - p

Call Shared-Variable-Reset.

Set the active-workspace to the clear-workspace.

Set comparison-tolerance to initial-comparison-tolerance.

Set event-message to initial-event-message.

Set event-type to initial-event-type.

Set index-origin to initial-index-origin.

Set latent-expression to initial-latent-expression.

Set print-precision to initial-print-precision.

Set quote-quad-prompt to the empty character vector.

Set random-link to initial-random-link.

Return command-complete.

Copy Library Workspace Object

If the library-workspace-named w does-not-exist, signal not-found.

If the **global-referent** of **a** in **w** is not a **defined-function**, **defined-operator**, **niladic-defined-function**, **variable**, or **shared-variable**, signal **incorrect-command**.

Set Z to copy a from w.

Return Z.

Copy Library Workspace

$$-p - p - COPY - r - w - p$$

If the library-workspace-named w does-not-exist, signal not-found.

For each symbol S of the library-workspace-named w

Let A stand for the **name** of S.

If the global-referent of A in w is a defined-function, defined-operator, niladic-defined-function, variable, or shared-variable,

Set Z to copy A from w.

If Z is an **exception**, return Z.

Return command-complete.

Drop Library Workspace

$$-p - p - DROP - r - w - p \longrightarrow$$

If the library-workspace-named w does-not-exist, signal not-found.

Set the existential-property of the library-workspace-named w to absent.

Return command-complete.

Erase Global Referent

$$-p - p - ERASE - r - a - p$$

Let GA stand for the global-referent of a in the active-workspace.

If the **global-referent** of **a** is not a **defined-function**, **defined-operator**, **niladic-defined-function**, **variable**, or **shared-variable**, signal **incorrect-command**.

Set Z to attempt-to-erase GA.

Return Z.

List Global Function Names

$$\mathbf{p}$$
 \mathbf{p} \mathbf{FNS} \mathbf{p} $-$

For each symbol A whose global-referent is a defined-function or a niladic-defined-function, display the name of A.

Return command-complete.

List Library Directory

$$-p - 1 - p - LIB - p - p$$

For each workspace W whose existential-property is present and whose owner is this-owner, display the workspace-name of W.

Return command-complete.

Load Library Workspace

$$-p - 1 - p - LOAD - r - w - p \longrightarrow$$

If the library-workspace-named w does-not-exist, signal not-found.

Set the active-workspace to the library-workspace-named w.

Call Shared-Variable-Reset.

If **latent-expression** is not **nil**,

Generate a new context in which

mode is immediate-execution.

current-line is the ravel-list of latent-expression,

current-function is 0 0p'',

current-line-number is one,

current-statement is the empty list of tokens, and

stack is the empty **list** of **tokens**.

Append the new **context** to the **state-indicator** of the **active-workspace** as a new first **item**.

Set T to **evaluate-line**.

Remove the first item from the state-indicator of the active-workspace.

Return T.

Otherwise, return command-complete.

Note: Shared-variable-reset has the effect of retracting and expunging all variables that were shared in the library-workspace. This means that conforming-programs cannot depend upon the values of variables that were shared when a) SAVE command was issued.

List Global Names and their Name Class

For each symbol A whose global-referent is a variable, defined-function, niladic-defined-function, or defined-operator, display the name of A appended with a dot followed by the name-class (for example FN.3, OP.4).

Return command-complete.

List Global Operator Names

$$-p - p - p - OPS - p - PS - p$$

For each **symbol** A whose **global-referent** is a **defined-operator**, **display** the **name** of A.

Return command-complete.

Save Active Workspace

$$-p - p - SAVE - p - SAVE - p$$

Set *W* to the **workspace-name** of the **active-workspace**.

If W is clear-workspace-identifier, signal not-saved.

Replace the **library-workspace-named** W with the **active-workspace**.

Set the existential-property of the library-workspace-named *W* to present.

Return command-complete.

Note: See) LOAD for a discussion of shared variable values.

Save Active Workspace with Name

$$-p - p - SAVE - r - w - p \longrightarrow$$

If the workspace-name of the active-workspace is not w,

If the library-workspace-named w already-exists, signal not-saved.

Otherwise, set the workspace-name of the active-workspace to w.

Replace the **library-workspace-named w** with the **active-workspace**.

Set the **existential-property** of the **library-workspace-named w** to **present**.

Return command-complete.

Note: See) LOAD for a discussion of shared variable values.

List State Indicator

$$-p -) -p - SI - p -$$

Let N stand for the **number-of-items** in the **state-indicator** of the **active-workspace**. Set I to **zero**.

Repeat:

Set I to I plus one.

If I is greater-than N, return command-complete.

If the mode of item \mathcal{I} of the state-indicator is defined-function or defined-operator, display the function-name and current-line-number of item \mathcal{I} of the state-indicator, along with an indication of whether the function-name is suspended.

(End of repeated block)

Clear State Indicator

$$p - p - p - SIC - p$$
Return clear-state-indicator.

List State Indicator and Local Names

$$-p -) - p - SINL - p -$$

Let N stand for the **number-of-items** in the **state-indicator** of the **active-workspace**. Set I to **zero**.

Repeat:

Set I to I plus one.

If I is greater-than N, return command-complete.

If the mode of item \mathcal{I} of the state-indicator is defined-function or defined-operator, display the function-name, current-line-number, and local-name-list of item \mathcal{I} of the state-indicator, along with an indication of whether the function-name is suspended.

(End of repeated block)

List Global Variable Names

-p - p - VARS - p - VARS - p

For each **symbol** A whose **global-referent** is a **variable**, **display** the **name** of A. Return **command-complete**.

List Workspace Identification

-p - p - WSID - p - WSID

Display the **workspace-name** of the **active-workspace**. Return **command-complete**.

Change Workspace Identification

-p - p - WSID - r - w - p

Set the **workspace-name** of the **active-workspace** to **w**. Return **command-complete**.

Annex A (normative)

Component Files

Informal Description: This annex describes a minimal set of functions which a conforming implementation must provide. There are, therefore, many facilities (such as shared libraries, access control, etc.) which are intentionally excluded from this annex, although it is expected that they will be available on most systems.

In particular, file names are severely restricted and access is provided only to a default library, in order to ensure compatible behaviour, on all systems, of the functions defined here.

This annex does not addresss the migration of the component files themselves, although it is presumed that at the very least one will be able to retrieve the objects into a workspace, migrate the workspace, and recreate the file, using these functions.

The user interface (arguments, function name, and result) are described here. When an error is signalled it implies that suspension will occur in the caller (that is, by means of $\square ES$).

A.1 Definitions of arguments and results

- A—Any array
- CID—Component id, must be a positive integer
- FH—File handle, an integer
- FHL—A list of file handles
- FL—File names, a character matrix with rows giving names of files

- FNM—File name, a character vector composed of 1–8 characters, the first chosen from A-Z, and any others chosen from A-Z and 0–9
- NFNM— New file name, see FNM above

A.2 Definition of functions

- FL←CF_LIST A

If A is an empty vector, FL is a character matrix containing the names of files existing in the default library.

If A is not empty, an error will be signalled. An implementation may provide a consistent extension where A is a library identifier.

If there are no component files in the library, FL will be an empty character matrix.

- FH←CF_CREATE FNM

Creates a file with name *FNM* and returns a numeric handle for the file. If the operation is unsuccessful, an error will be signalled.

 $-Z \leftarrow CF_ERASE\ FNM(Z \text{ is } 0\ 0p0)$

Erase the file specified in *FNM*. If successful an empty numeric matrix will be returned. Success implies that the file no longer exists (either it has been erased or it never existed). If the operation is unsuccessful an error will be signalled.

- Z←NFNM CF_RENAME FNM(Z is 0 0p0)

The file named by *FNM* is renamed *NFNM*. If this is successful an empty numeric matrix is returned. If the operation is unsuccessful an error is signalled.

- FH←[FH] CF_OPEN FNM

The file named by *FNM* is opened and a file handle is returned. The left argument is optional. The presence of the left argument indicates that the file is to be opened with the specified file handle. If the operation is unsuccessful an error is signalled.

Note: If a left argument is present but the file cannot be opened with the specified FH, the open will be unsuccessful.

 $-Z \leftarrow CF_CLOSE\ FHL\ (Z \text{ is } 0\ 0\rho\ 0)$

Each file listed in FHL is closed. The presence of the result indicates that all files named in FHL are now closed. An error will be signalled if the close is unsuccessful (for example, an invalid FH or unable to close the file).

- FHL←CF_INUSE

The result FHL is a list of file handles (FH) that are currently valid. It will be empty if there are no valid file handles current.

 $-Z \leftarrow A$ CF_WRITE FH, CID(Z is 0 0p0)

Array A is written to component CID in file FH. The presence of the result indicates success. An error will be signalled if the operation is unsuccessful.

- CID← A CF_APPEND FH

Array A is appended to the file FH and CID is its component id in the file. An error is signalled if the operation is unsuccessful.

- A←CF_READ FH, CID

The result A is component CID of file FH. An error is signalled if the operation is unsuccessful.

- CID←CF_NEXT FH

The result is the component id that will be returned for the next successful use of *CF_APPEND* with file *FH*. An empty file will return a *CID* of one. An error is signalled if the operation is unsuccessful.

A.3 Errors

The functions described here must be written so that they do not suspend. Any errors encountered or detected must be passed to the context from which the function was invoked.

Bibliography

The following publications are referred to in this International Standard to provide general background and guidance for the benefit of implementers. The International Standard does not rely upon the content of these documents for its interpretation, and implementers are under no obligation to consult these documents or to use the information contained in them.

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