

Two's Complement

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A trick to represent negative numbers without an extra symbol.


Consider the following exercise to

find $(-4) : 0 - 4$.

$$\begin{array}{r} \begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \end{array} \left[\begin{array}{cccc} \dots & 1 & 1 & 1 \\ \dots & \overset{10}{\cancel{0}} & \overset{10}{\cancel{0}} & \overset{10}{\cancel{0}} \\ \dots & 0 & 0 & 0 \\ \dots & 0 & 0 & 4 \end{array} \right] \leftarrow \begin{array}{l} \text{our row which} \\ \text{indicates borrowing} \end{array} \\ \hline \dots 9996 \end{array}$$

$$\text{So } -4 = (\dots 96)$$

$$\begin{array}{r} \text{Check: } \dots 9996 \\ \quad \quad \quad 4 \\ \hline \dots 0000 \quad \checkmark \end{array}$$

Now in binary: 

What is $4_{10} = 0100_2$

and $-4_{10} = ?$

$$\begin{array}{r} - \quad \dots 11 \\ \quad \quad \quad \cancel{10} \cancel{10} \cancel{10} 00_2 \\ - \quad \dots 00100_2 \\ \hline \quad \dots 11100_2 \end{array}$$

$$-4_{10} = (\dots 1100_2)$$

$$\left[\begin{array}{l} \text{check:} \\ \quad \quad \quad \dots 1100_2 \\ \quad \quad \quad + \quad \quad 0100 \\ \quad \quad \quad \hline \quad \quad \dots 00000 \end{array} \right] \checkmark$$

- Lets consider a truncated

Representation: for N bits,

The MSB denotes MSB \times infinite
repeating 1s, ex.

$$(1100)_2 = \dots 1100$$

So I can add things together.
with finite computation:

ex. $5 + (-4) :$

$$\begin{array}{r} 11 \\ 0101_2 \\ + \dots 1100_2 \\ \hline \dots 0001 \end{array}$$

$$= 0001_2 \rightarrow 5 + (-4) = 1 \checkmark$$

OK. So instead of subtracting from
zero every time, let's
consider that (in finite bits)

$$\begin{array}{r} 0000 \\ - 0101 \\ \hline \end{array} \Leftrightarrow \begin{array}{r} 10000 \\ - 0101 \\ \hline \end{array} \Leftrightarrow \begin{array}{r} 1111 + 1 \\ - 0101 \\ \hline \end{array}$$

and note that there's two cases:

(A) $\begin{array}{r} 1 \\ - 0 \\ \hline 1 \end{array}$

(B) $\begin{array}{r} 1 \\ - 1 \\ \hline 0 \end{array}$

No carries
involved!

↳

Going to add 1 at the end

see:

$$\begin{array}{r} 1111 \\ - 0101 \\ \hline 1010 + 1 \\ \hline \end{array}$$

$$-5 = (1011)_2$$

Trick: subtracting from zero

is equivalent to flip bits & add 1.

Ex. Find -6 in 4 bit 2's comp.

$$6_{10} = (0110)_2$$

$$\begin{aligned} -6_{10} &= +100\overset{1}{1} \\ &\quad \underline{1010} \\ &= (1010)_2 \end{aligned}$$

What about in reverse?

$$\begin{aligned} G_{10} &= -(-6_{10}) = +010\overset{1}{1} \\ &\quad \underline{0110} \quad \checkmark \end{aligned}$$

4 bit 2's Complement

S_3, S_2, S_1, S_0

0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
	1000
-7	1001
-6	1010
-5	1011
-4	1100
-3	1101
-2	1110
-1	1111

Same

$-8?$

$-(1001)_2 = +\frac{0110}{111}$

$-8 + 6$

$-8 + 7$

$+1$
 $+1$
 $+1$

⇒ weight factor representation.

$$S_3 \times (-2^3) + S_2 \times 2^2 + S_1 \times 2^1 + S_0 \times 2^0$$

In General: 2's complement

is a different numbering system
for different numbers of bits

MSB has a weight factor of

-2^{N-1} , all other bits in
position $n = 0, 1, \dots, N-2$ have wf 2^n

For N bits, we'll have a range
of $[-2^{N-1}, 2^{N-1}-1]$