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ASSIGNMENT NO.:

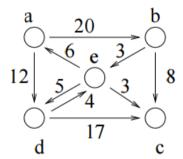
PROBLEM STATEMENT:-

Programming in C to find the path matrix using Warshall's algorithm.

THEORY:-

This is a classical algorithm by which we can determine whether there is a path from any vertex v_i to another vertex v_j either directly or through one or more intermediate vertices. In other words, we can test the reachability of all the pairs of vertices in a graph. The path matrix can be computed from the adjacency matrix A by $P = A + A^2 + A^3 + \dots + A^n$ where n = no. of vertices. This method is computationally not efficient at all. To compute the path matrix from a given graph, another elegant method is Warshall's algorithm. This algorithm treats the entries in the adjacency matrix as bit entries & performs AND (Δ) & OR (ν) Boolean operations on them. The heart of the algorithm is a trio of loops, which operates very much like the loops in the classic algorithms for matrix multiplication.

Example:



 $\begin{array}{c|c}
a & -20 & b \\
4 & 4 & 5
\end{array}$

Fig 1: Without negative cost cycle

fig 2: With negative cost cycle

ALGORITHM:-

Input:- A graph **G** whose pointer to its adjacency matrix is **GPTR** & vertices are labeled as 1,2, ...,N; N being the number of vertices in the graph.

Output:- The path matrix **a**.

Data structure:- Matrix representation of graph **G** with pointer as **GPTR**.

Algorithm for main() function:

Step 1: Input "Enter number of vertices"

Step 2: Read n

Step 3: Repeat through step 4 to step 8 for (i=0 to n) do

Step 4: Repeat through step 5 to step 7 for (j = 0 to n) do

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Step 5: Print the existence of path between vertices
Step 6: Read a[i][i]
Step 7: Next j
        [End of inner for loop]
Step 8: Next i
        [End of outer for loop]
Step 9: Call the method display (n,a)
Step 10: Repeat through step 11 to step 16 for (k=0 to n) do
Step 11: Repeat through step 12 to step 15 for (i=0 to n) do
Step 12: Repeat through step 13 to step 14 for (j=0 to n) do
Step 13: Set a[i][j]=a[i][j] OR (a[i][k] AND a[k][j])
Step 14: Next j
         [End of for loop]
Step 15: Next i
         [End of for loop]
Step 16: Next k
         [End of outer for loop]
Step 17: Call the method display(n,a)
Step 18: Stop
Algorithm for the method display():
```

```
Step 1: Repeat through step 2 to step for (i=0 to n) do
Step 2: Repeat through step 3 to step for (j=0 to n) do
Step 3: Print a[i][j]
Step 4: Next j
        [End of inner for loop]
Step 5: Next i
        [End of outer for loop]
Step 6: Stop
```

SOURCE CODE:-

```
#include<stdio.h>
void display(int n,int a[20][20]); //prototype declaration
void main()
int a[20][20],i,j,k,n; //variable declaration
printf("\nEnter the total number of vertices:\t");
scanf("%d",&n);
printf("\nThe existence of path between every pair of vertices:\n");
printf("1->There is a path between vertices\n0->There is no path between vertices\n");
//loop for taking inputs of the matrix from the user
for(i=0;i<=n;i++)
```

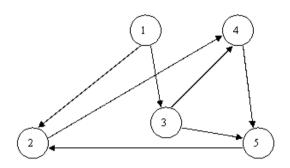
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for(j=0;j<=n;j++)
printf("\nEnter the existence of path between vertices %d & %d:\t",i,j);
scanf("%d",&a[i][j]);
printf("\n\nThe adjacency matrix is:\n");
display(n,a); //calling method display
//loop for finding the minimum distance
for(k=0;k<=n;k++)
{
for(i=0;i \le n;i++)
for(j=0;j<=n;j++)
a[i][j]=a[i][j] \parallel (a[i][k] \&\& a[k][j]);
printf("\n\nThe minimum distance between every pair of vertices:\n");
display(n,a);
}
void display(int n,int a[20][20])
   //method to display the matrix
int i,j;
for(i=0;i<=n;i++)
for(j=0;j<=n;j++)
printf(" %d ",a[i][j]);
printf("\n");
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INPUT-OUTPUT:-

The given graph is:



Enter the total no. of vertices: 5

The existence of path between every pair of vertices:

1->There is a path between vertices

0->There is no path between vertices

Enter the existence of path between vertices 1&1:	0
Enter the existence of path between vertices 1&2:	1
Enter the existence of path between vertices 1&3:	1
Enter the existence of path between vertices 1&4:	0
Enter the existence of path between vertices 1&5:	0
Enter the existence of path between vertices 2&1:	0
Enter the existence of path between vertices 2&2:	0
Enter the existence of path between vertices 2&3:	0
Enter the existence of path between vertices 2&4:	1
Enter the existence of path between vertices 2&5:	0
Enter the existence of path between vertices 3&1:	0
Enter the existence of path between vertices 3&2:	0
Enter the existence of path between vertices 3&3:	0
Enter the existence of path between vertices 3&4:	1
Enter the existence of path between vertices 3&5:	1
Enter the existence of path between vertices 4&1:	0
Enter the existence of path between vertices 4&2:	0
Enter the existence of path between vertices 4&3:	0
Enter the existence of path between vertices 4&4:	0
Enter the existence of path between vertices 4&5:	1

The adjacency matrix:

 $0\ 1\ 1\ 0\ 0$

 $0\,0\,0\,1\,0$

 $0\ 0\ 0\ 1\ 1$

 $0 \; 0 \; 0 \; 0 \; 1$

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The minimum distance between every pair of vertices:

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01011

01011

01011

01011

DISCUSSION:

- **A)** It is one of the most commonly used shortest path algorithm. A shortest path between two vertices is a path, which has the least no. of edges among several paths in between two vertices.
- **B)** It is an iterative process. The first iteration consists of finding the existence of path from one vertex to another vertex either directly or indirectly via any intermediate vertex or pivot vertex say v_i . The second iteration finds the existence of path from one vertex to another vertex with $v_1 \& v_2$ or both as pivot & so on.
- C) It is the most efficient method to compute the shortest path between every pair of vertices. It requires N_3 comparisons & has an order of complexity $O(N_3)$.

Floyd & Dijkstra are two other methods employed to determine the shortest path between vertices.