

EE621: Markov Chains and Queueing Systems

Quiz #2 (20 Marks)

Time: 4:30 - 6 pm

Date: 18/02/2023

_____ x _____ x _____ x _____ x _____

[Q.1] For an irreducible DTMC $\{X_n\}_{n \geq 0}$ on $\mathcal{S} = \{0, 1, \dots\}$ with TPM P , prove or disprove:

$$\text{g.c.d. } \{n: f_{jj}^{(n)} > 0\} = \text{g.c.d. } \{n: p_{jj}^{(n)} > 0\}.$$

5 Marks

_____ x _____ x _____ x _____

[Q.2] Consider an irreducible $\{X_n\}_{n \geq 0}$ on $\mathcal{S} = \{0, 1, \dots\}$ with TPM P . Prove or disprove: If \exists a prob. mass function $\alpha > 0$ on \mathcal{S} satisfying $\alpha_i p_{ij} = \alpha_j p_{ji} \quad \forall i, j \in \mathcal{S}$, then the DTMC is positive recurrent.

5 Marks

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[Q.3] Consider the discrete time queueing system with two (instead of one) servers. In each slot, arrival occurs w.p. p in each slot. If a server is working, then prob. of it finishes service in a slot is q . The two servers are independent. The servers never idle as long as customers are waiting.

(a) Draw a state transition diagram. 5 Marks

(b) Find conditions under which the DTMC is positive recurrent. 5 Marks

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Solution

Q.1

Note that $p_{jj}^{(n)} > f_{jj}^{(n)}$.

$$\Rightarrow \{n: f_{jj}^{(n)} > 0\} \subseteq \{n: p_{jj}^{(n)} > 0\}$$

$$\Rightarrow \gcd \{n: f_{jj}^{(n)} > 0\} \geq \gcd \{n: p_{jj}^{(n)} > 0\}. \quad (1)$$

consider any n such that $p_{jj}^{(n)} > 0$ & $f_{jj}^{(n)} = 0$.

But note that

$$p_{jj}^{(n)} = \sum_{k=1}^n f_{jj}^{(k)} p_{jj}^{(n-k)}$$

$\therefore p_{jj}^{(n)} > 0$, then $\exists k \in \{1, 2, \dots, n-1\}$ s.t.

$$f_{jj}^{(k)} p_{jj}^{(n-k)} > 0.$$

$$\Rightarrow f_{jj}^{(k)} > 0 \text{ and } p_{jj}^{(n-k)} > 0.$$

Applying same steps to $p_{jj}^{(n-k)}$, we see that

any $n \in \{n: p_{jj}^{(n)} > 0\}$ is sum of elements of $\{n: f_{jj}^{(n)} > 0\}$.

Thus, g.c.d. $\{n: f_{jj}^{(n)} > 0\}$ also divides every element of $\{n: p_{jj}^{(n)} > 0\}$.

But this implies

$$\gcd \{n: p_{jj}^{(n)} > 0\} \geq \gcd \{n: f_{jj}^{(n)} > 0\} \quad (2)$$

The result follows from (1) & (2).

— x — x — x —

Q.2 We need to show that $\bar{\alpha}$ satisfies

$$\bar{\alpha} = \bar{\alpha} P.$$

Note that $\forall i, j \in S$,

$$\alpha_i p_{ij} = \alpha_j p_{ji}$$

$$\Rightarrow \sum_{j \in S} \alpha_i p_{ij} = \sum_{j \in S} \alpha_j p_{ji}$$

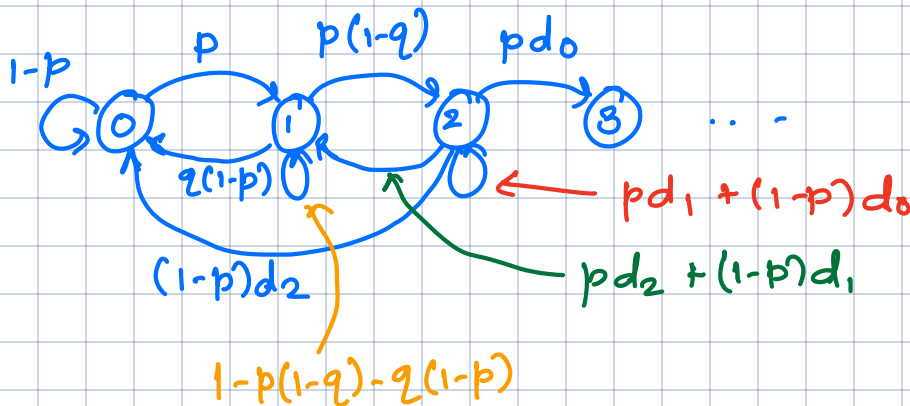
$$\Rightarrow \alpha_i = \sum_{j \in S} \alpha_j p_{ji}.$$

This proves the required.

— x — x — x —

Q.3 Let $d_i, i=0,1,2$ denote the probability of i departures when both servers are working.

$$d_0 = (1-q)^2; \quad d_1 = q(1-q) \quad \& \quad d_2 = q^2$$



Note that $\forall k \geq 2$,

$$P_{k,k+1} = pd_0$$

$$P_{k,k} = pd_1 + (1-p)d_0$$

$$P_{k,k-1} = pd_2 + (1-p)d_1$$

$$P_{k,k-2} = (1-p)d_2$$

Now consider $\mathcal{X} = \{0, 1\}$ and $Q(k) = k$.

Now find the drift.

$$\begin{aligned} & E[Q(X_{k+1}) - Q(X_k) | X_k = k] \\ &= (k+1)pd_0 + k(pd_1 + (1-p)d_0) + (k-1)(pd_2 + (1-p)d_1) \\ &\quad + (k-2)(1-p)d_2 - k \\ &= k[\cancel{pd_0} + \cancel{pd_1} + (1-p)d_0 + \cancel{pd_2} + (1-p)d_1 + (1-p)d_2 - 1] \\ &\quad + pd_0 - pd_2 - (1-p)d_1 - 2(1-p)d_2 \\ &= k[\underbrace{d_0 + d_1 + d_2 - 1}_{=0}] + \cancel{pd_0} - \cancel{pd_2} - d_1 + \cancel{pd_1} - 2d_2 + 2\cancel{pd_2} \\ &= p[d_0 + d_1 + d_2] - d_1 - 2d_2 \\ &= p - d_1 - 2d_2. \end{aligned}$$

Note that whenever $p < d_1 + 2d_2$, we can choose

$$\varepsilon = \frac{d_1 + 2d_2 - p}{2} \text{ and then note that } \text{Drift}(k) < -\varepsilon$$

$$\forall k \geq 2.$$

Also show that $E[Q(X_{k+1}) | X_k = 0]$ & $E[Q(X_{k+1}) | X_k = 1]$ are $< \infty$.

Thus, Foster's criteria implies the DTMC is positive recurrent for $p < d_1 + 2d_2$ ($p < q(1-q)$).

Note the $d_1 + 2d_2$ is the expected number of departures in a slot when both servers are working.