

## Homework 3

**Instructions:** Please submit your solutions via Gradescope by **Friday, 11 February 2022, 10:00am**. Make sure your name, your class group number, and the name of your class teacher is put on **every page** of your submission. Your submission should, ideally, be a **PDF** file.

### Exercise 3.1: Fibonacci numbers, recursively

4 pts

The Fibonacci numbers are a recursively-defined sequence of numbers, which arise in a surprising variety of real-world phenomena. The  $n$ th Fibonacci number is usually denoted by  $F_n$  and has the following recursive definition:

$$\begin{aligned} F_1 &= 1, \\ F_2 &= 1, \\ F_n &= F_{n-1} + F_{n-2}, \quad \text{for } n > 2. \end{aligned}$$

- (a) Write Python code that, given  $n \geq 1$  as an argument, implements the natural recursive algorithm for computing the  $n$ th Fibonacci number  $F_n$ .
- (b) Measure the time it takes for computing the  $n$ th Fibonacci number  $F_n$  for small values of  $n$  (up to say 40 or 45): for example, using the code snippet `stopwatch.py` provided on the course's Moodle page. Use this to explore the ratio between the running times for two consecutive values of  $n$ . What do you observe?
- (c) Let  $a, b > 0$  be positive constants. Then, the running time of the recursive algorithm is well captured by the following recurrence:

$$\begin{aligned} T(n) &= T(n-1) + T(n-2) + b, & \text{for } n \geq 3, \text{ and} \\ T(n) &= a & \text{for } n = 1, 2. \end{aligned}$$

Use induction to show that

$$T(n) = (a+b)F_n - b, \quad \text{for all } n \geq 1.$$

- (d) Assume  $a = b = 1$ . The number  $\phi = (1 + \sqrt{5})/2 \approx 1.618$  is known as the Golden ratio. Use Binet's formula for the  $n$ th Fibonacci number  $F_n$ , given as

$$F_n = \frac{1}{\sqrt{5}} (\phi^n - (-\phi)^{-n}),$$

to argue that  $T(n) = \Omega(\phi^n)$ .

**Exercise 3.2: Fibonacci numbers, iteratively****3 pts**

The natural recursive algorithm for computing the  $n$ th Fibonacci number  $F_n$  has exponential running time. From a time complexity perspective, that is really terrible. From a practical perspective, this means that you will not be able to compute the  $n$ th Fibonacci number  $F_n$  even for moderately sized values of  $n$  using the recursive algorithm.

Luckily, there is a more clever iterative algorithm for computing the  $n$ th Fibonacci number that runs in linear time.

- (a) Describe this algorithms in words.
- (b) Implement this algorithm in Python.
- (c) Argue, using big- $O$  notation, that the running time of your algorithm is  $O(n)$ .

**Exercise 3.3: Big  $O$ -notation and the sum rule****3 pts**

- (a) Show that, if  $f_1(n) = O(g_1(n))$  and  $f_2(n) = O(g_2(n))$ , then

$$f(n) = f_1(n) + f_2(n) = O(g_1(n) + g_2(n)).$$

- (b) For functions in part (a), do we also have  $f(n) = O(\max\{g_1(n), g_2(n)\})$ ?
- (c) Is it also true that if  $f_1(n) = \Omega(g_1(n))$  and  $f_2(n) = \Omega(g_2(n))$ , then

$$f(n) = f_1(n) + f_2(n) = \Omega(g_1(n) + g_2(n))?$$