CS325: Homework #2

1. Solve recurrence relation using three methods

$$T(n) \begin{cases} T(0) = c_1 \\ T(1) = c_2 \\ 2(T(\frac{n}{2})) \text{ if } n \text{ is even} \\ 2T(\frac{n}{2}) + c_3 \text{ if } n \text{ is odd} \end{cases}$$

1a. Substitution

odd: $2T(\frac{n}{2}) + cn$

$T(n) = 2T(\frac{n}{2}) + cn$

$2[2T(\frac{n}{4}) + c(\frac{n}{2}] + cn$

$T(n) = 4T(\frac{n}{4}) + cn + cn$

$= 4[2T(\frac{n}{8}) + c(\frac{n}{4})] + 2cn$

$T(n) = 8T(\frac{n}{8}) + 8n + cn + cn$

$= 2^3 T(\frac{n}{2^3}) + cn + cn + cn$

k^{th} $= 2^k T(\frac{n}{2^k}) + k \cdot cn$

$\frac{n}{2^k} = 1$ $T(n) = 2^{\log_2 n} T(\frac{n}{2^{\log_2 n}}) + (2^{\log_2 n} - 1)\, c_3$

$= nT(1) + (n-1) \cdot c_3$

$= nc_2 + (n-1)c_3$

$T(n) = \Theta(n)$

$T(\frac{n}{2}) = 2T(\frac{\frac{n}{2}}{2}) + c \cdot (\frac{n}{2})$

$= 2T(\frac{n}{4}) + c \cdot (\frac{n}{2})$

$T(\frac{n}{4}) = 2T(\frac{\frac{n}{4}}{2}) + c \cdot (\frac{n}{4})$

$= 2T(\frac{n}{8}) + c(\frac{n}{4})$

1B. master method

$T(n) = 2T(\frac{n}{2}) + c_3$

$a = 2 \quad b = 2 \quad k = 0 \quad p = 0$

$\log_2 2 = 1 > k$, $\Theta(n^{\log_b a}) = \Theta(n^{\log_2 2})$

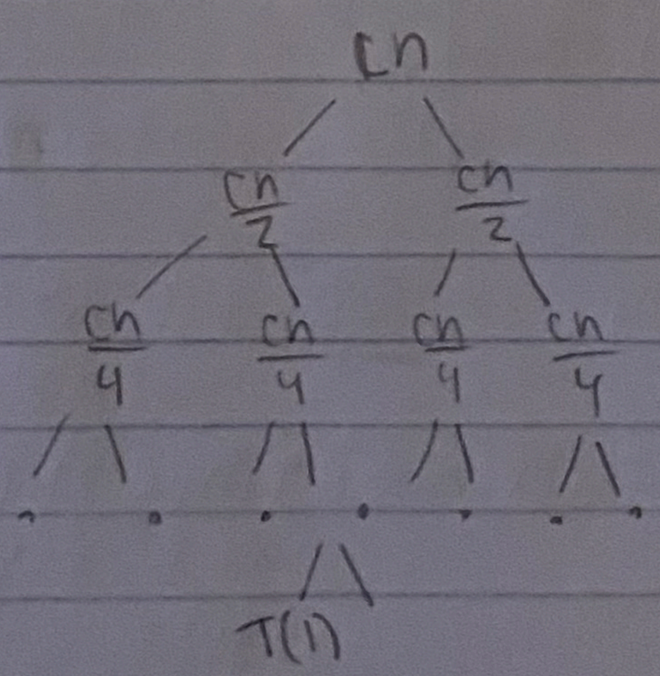
case 1: $\Theta(n)$

1c. tree

height $2^{\log_2 n}$

$n^{\log_2 2} = n^1$

$\Theta(n)$



level 0, 2^0

level 1, 2^1

level 2, 2^2

level i, 2^i

$n \cdot T(1)$

2. Solve recurrence, using master theorem

2A:

$T(n) = 4T(n/2) + n$

$a=4 \quad b=2 \quad k=1 \quad p=0$

$\log_2 4 = 2 > k$

Case 1: $\Theta(n^2)$

2B:

$T(n) = 2T(n/4) + n^2$

$a=2 \quad b=4 \quad k=2 \quad p=0$

$\log_4 2 = .5 < k$

Case 3: $\Theta(n^2)$

3. Writing Pseudocode

```
function kth element (Arr 1, Arr 2, k)
    # initialize values + 3rd arr = size of Arr1 + Arr2
loop through Arr 1 and Arr2
    if Arr1[i] < Arr2[j]
        Arr3 = Arr 1
            iterate through Arr
    else    Arr 3[j] <
        Arr 3 = Arr 2
            iterate through Arr
while more elements in Arr 1
    add to Arr 3
    iterate
while mor elements in Arr 2
    add to Arr 3
    iterate
return merged Arr(Arr 3[kth element])
```