# Community Detection in Colorful Graphs

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### 1 Problem Set-up

Let  $P = (P_0, P_1, ..., P_L)$  and  $Q = (Q_0, Q_1, ..., Q_L)$  be discrete distributions over L colors and 0. Let  $\sigma_0 : [n] \to K$  be the true clustering. For a pair (i, j) in the same cluster  $(\sigma_0(i) = \sigma_0(j))$ , suppose  $A_{ij} \sim P$  and for a pair (i, j) in different clusters, we suppose  $A_{ij} \sim Q$ .

We observe the colorful matrix A but not P and Q and the goal is to recover the clusters. We refer to the setting where P, Q are known as the oracle setting.

Note that under the oracle setting, the log-likelihood of a clustering  $\sigma$  is

$$l(A \mid \sigma) = \sum_{i < j} \mathbf{1}_{\sigma(i) = \sigma(j)} \sum_{l=0}^{L} \log P_{l} \mathbf{1}_{A_{ij} = l} + \sum_{i < j} \mathbf{1}_{\sigma(i) \neq \sigma(j)} \sum_{l=0}^{L} \log Q_{l} \mathbf{1}_{A_{ij} = l}$$

$$= \sum_{i < j} \mathbf{1}_{\sigma(i) = \sigma(j)} \sum_{l=0}^{L} \log \frac{P_{l}}{Q_{l}} \mathbf{1}_{A_{ij} = l} + \sum_{i < j} \sum_{l=0}^{L} \log Q_{l} \mathbf{1}_{A_{ij} = l}$$

Let  $I_{tot} = -2 \log \sum_{l} \sqrt{P_{l}Q_{l}}$  be the 1/2 Renyi-divergence between the distributions P, Q.

## 2 Weak Recovery Under the Oracle Setting

[1] characterizes the minimax rate of weak recovery for Bernoulli P,Q. Their results and proofs can be extended in a straightforward manner to general discrete P,Q under the oracle setting.

**Proposition 2.1.** (Upper bound) Assume  $\frac{nI_{tot}}{K \log K} \to \infty$ . The maximum likelihood estimator  $\hat{\sigma}$  in the oracle setting achieves:

$$\sup_{\Theta(n,K,\beta,P,Q)} \mathbb{E}r(\widehat{\sigma},\sigma) \leq \left\{ \begin{array}{l} \exp\left(-(1+o(1))\frac{nI_{tot}}{2}\right), \quad K=2, \\ \exp\left(-(1+o(1))\frac{nI_{tot}}{\beta K}\right), \quad K \geq 3 \end{array} \right.$$

*Proof.* Proof of this proposition follows that of Theorem 3.2 in [1]. We describe only the parts that need to be modified.

Because we have a different likelihood function, our  $T(\sigma)$  takes on a different form from that of [1] at the bottom of page 8:

$$T(\sigma) = \sum_{i < j} \mathbf{1}_{\sigma(i) = \sigma(j)} \sum_{l=0}^{L} \log \frac{P_l}{Q_l} \mathbf{1}_{A_{ij} = l}$$

Let  $\sigma_0$  denote the true community assignment. We make a mistake if for some other community assignment  $\sigma$ , we get  $T(\sigma) > T(\sigma_0)$ . The key part of the proof is to bound

$$P_m \equiv P(\exists \sigma : T(\sigma) > T(\sigma_0), d_H(\sigma, \sigma_0) = m)$$

To that end, we bound the probability of error of a fixed  $\sigma$  *m*-distant from  $\sigma_0$  in Hamming distance. We will prove an analogue of Proposition 5.1 in [1]:

Let  $\sigma$  be an arbitrary assignment satisfying  $d(\sigma, \sigma_0) = m$ . Let  $X_i, Y_i$  be random variables such that

$$X_i = \log \frac{P_l}{Q_l}$$
 w.p.  $P_l$   $Y_i = \log \frac{P_l}{Q_l}$  w.p.  $Q_l$ 

and  $\alpha, \gamma$  be integers where

$$\alpha = |\{(i,j) : \sigma_0(i) = \sigma_0(j) \land \sigma(i) \neq \sigma(j)\}| \quad \gamma = |\{(i,j) : \sigma_0(i) \neq \sigma_0(j) \land \sigma(i) = \sigma(j)\}|$$

Then

$$P(T(\sigma) \ge T(\sigma_0)) \le P\left(\sum_{i=1}^{\alpha} X_i - \sum_{i=1}^{\gamma} Y_i < 0\right) \le \exp(-\frac{\gamma + \alpha}{2}I)$$
 (2.1)

Lemma 5.3 from [1] bounds  $\alpha, \gamma$  and Proposition 5.2 bounds the number of  $\sigma$ 's (up to equivalent classes) such that  $d_H(\sigma, \sigma_0) = m$ . These pieces together bounds  $P_m$ . The rest of the proof follows [1] exactly starting from Page 16.

We devote the rest of the proof toward proving equation 2.1.

$$T(\sigma_0) - T(\sigma') = \sum_{i < j} \mathbf{1}_{\sigma_0(i) = \sigma_0(j) \land \sigma'(i) \neq \sigma'(j)} \sum_{l=1}^L \mathbf{1}_{A_{ij} = l} \log \frac{P_l}{Q_l}$$
$$- \sum_{i < j} \mathbf{1}_{\sigma_0(i) \neq \sigma_0(j) \land \sigma'(i) = \sigma'(j)} \sum_{l=1}^L \mathbf{1}_{A_{ij} = l} \log \frac{P_l}{Q_l}$$
$$= \sum_{i=1}^{\alpha} X_i - \sum_{i=1}^{\gamma} Y_i$$

$$P(\sum_{i=1}^{\gamma} Y_i - \sum_{i=1}^{\alpha} X_i > 0) \le \mathbb{E}\left(e^{-t\sum_{i=1}^{\alpha} X_i} e^{t\sum_{i=1}^{\gamma} Y_i}\right)$$

$$\le \mathbb{E}\left(e^{-tX_1\alpha} e^{tY_1\gamma}\right)$$

$$\le \left(\mathbb{E}e^{-tX_i} \mathbb{E}e^{tY_1}\right)^{(1-w)\alpha + w\gamma} \frac{\left(\mathbb{E}e^{tY_1}\right)^{(1-w)(\gamma - \alpha)}}{\left(\mathbb{E}e^{-tX_1}\right)^{w(\gamma - \alpha)}}$$

We will show that when t = 1/2 and w = 1/2, the fraction term equals 1 and the first term equals  $\exp(-(1/2\alpha + 1/2\gamma)I)$ .

Note that

$$\mathbb{E}e^{-tX_1} = \sum_{l} P_l e^{-t\log\frac{P_l}{Q_l}}$$

$$= \sum_{l} P_l \left(\frac{Q_l}{P_l}\right)^t$$

$$= \sum_{l} \sqrt{P_l Q_l} \quad (\text{if } t = 1/2)$$

$$\mathbb{E}e^{tY_1} = \sum_{l} Q_l e^{t \log \frac{P_l}{Q_l}}$$

$$= \sum_{l} Q_l \left(\frac{P_l}{Q_l}\right)^t$$

$$= \sum_{l} \sqrt{P_l Q_l} \quad \text{if } t = 1/2$$

$$P(\sum_{i=1}^{\gamma} Y_i - \sum_{i=1}^{\alpha} X_i > 0) \le \left(\sum_{l} \sqrt{P_l Q_l}\right)^{\alpha + \gamma}$$
$$\le \exp\left(-\frac{1}{2}I\right)^{\alpha + \gamma}$$
$$\le \exp\left(-\frac{(\alpha + \gamma)}{2}I\right)$$

## 3 Fixed L Case

Suppose L is fixed. Suppose also that  $\lim_{n\to\infty} P_0 \wedge Q_0 \to 1$  so that  $P_l, Q_l \to 0$  for  $l \neq 0$ .

**Proposition 3.1.** Suppose  $\frac{nI}{K} \to \infty$ . If weak recovery (consistency) is possible under the oracle setting, then it is possible when P, Q are unknown.

*Proof.* It is clear that weak recovery is possible iff  $\frac{nI_{tot}}{K\log K} \to \infty$ .

$$I_{tot} = -2\log\left(\sum_{l=1}^{L} \sqrt{P_l Q_l} + \sqrt{(1 - \sum_{l=1}^{L} P_l)(1 - \sum_{l=1}^{L} Q_l)}\right)$$

$$= -2\log\left(1 - \frac{1}{2}\left(\sum_{l=0}^{L} (\sqrt{P_l} - \sqrt{Q_l})^2\right)\right)$$

$$= (1 + o(1))\sum_{l=0}^{L} (\sqrt{P_l} - \sqrt{Q_l})^2$$

Since  $I_{tot} = \omega(\frac{K \log K}{n})$  by hypothesis, it must be that, for some l,  $(\sqrt{P_l} - \sqrt{Q_l})^2 = \omega(\frac{K \log K}{n})$ .

We choose such an l and consider an estimator  $\hat{\sigma}_l$  that uses only the information  $\mathbf{1}_{A_{ij}=l}$ .

Since the Renyi-divergence  $I_l$  of  $Ber(P_l)$  and  $Ber(Q_l)$  is

$$I_l = (1 + o(1)) \left( (\sqrt{P_l} - \sqrt{Q_l})^2 + (\sqrt{1 - P_l} - \sqrt{1 - Q_l})^2 \right)$$

We have that  $\frac{nI_l}{K \log K} \to \infty$  and weak consistency is thus achievable with  $\hat{\sigma}_l$ .

Although weak consistency is achievable with the estimator  $\hat{\sigma}_l$  that considers only  $\mathbf{1}_{A_{ij}=l}$ , the estimator converges at  $\exp(-\frac{nI_l}{\beta K})$  and therefore does not converge at the same rate. It is easy to see that  $I_l \geq I_{tot} \geq \frac{1}{L}I_l(1-o(1))$ .

We propose a procedure to remove the  $\frac{1}{L}$  factor and get oracle rate of convergence:

- 1. Split data into two halves.
- 2. For each color l = 1, ..., L and for l = 0, cluster the first half of the data based on  $\mathbf{1}_{A_{ij}=l}$  and get  $\widehat{\sigma}_l$ . Use the second half to verify the accuracy, either by likelihood evaluation or by clustering the second half as well and testing consistency.
- 3. Cluster all data with the best l and use the rough clustering to estimate  $\widehat{P}_l$  and  $\widehat{Q}_l$ .
- 4. Plug in the estimated  $\hat{P}_l$  and  $\hat{Q}_l$  into the maximum likelihood estimator.

As a first step toward analyzing such a procedure, we ask the following question: suppose  $\sigma^*$  is an imperfect clustering of points (2,...,n) such that  $d_H(\sigma_0,\sigma^*)=\gamma$ . We want to assign a cluster to node 1 based on  $\sigma^*$ . We will

pretend we can see P, Q exactly and assign node 1 by maximum likelihood. We will for simplicity also assume two clusters each of size n/2 for now. Assume without loss of generality that  $\sigma_0(1) = 1$ .

The assignment is

$$\underset{k=1,2}{\operatorname{argmax}} \sum_{i:\sigma^*(i)=k} \sum_{l=0}^{L} \log \frac{P_l}{Q_l} \mathbf{1}_{A_{1j}=l}$$

We can think of this as assigning a weight of  $\log \frac{P_l}{Q_l}$  to edges with color l and then assigning node 1 to the cluster that maximizes the weights of the edges (1, j) for j in that cluster.

What is the probability that point 1 will be assigned the wrong cluster using  $\sigma^*$ ? There are four categories of points:

- 1. j such that  $\sigma_0(j) = 1$  and  $\sigma^*(j) = 1$ . There are  $n/2 \gamma/2$  of these.  $A_{1j} \sim P$ .
- 2. j such that  $\sigma_0(j) = 2$  and  $\sigma^*(j) = 1$ . There are  $\gamma/2$  of these.  $A_{1j} \sim Q$ .
- 3. j such that  $\sigma_0(j) = 1$  and  $\sigma^*(j) = 2$ . There are  $\gamma/2$  of these.  $A_{1j} \sim P$ .
- 4. j such that  $\sigma_0(j) = 2$  and  $\sigma^*(j) = 2$ . There are  $n/2 \gamma/2$  of these.  $A_{1j} \sim Q$ .

Thus,

$$\sum_{j:\sigma^*(j)=1} \sum_{l=0}^{L} \log \frac{P_l}{Q_l} \mathbf{1}_{A_{1j}=l} \leq \sum_{j:\sigma^*(j)=2} \sum_{l=0}^{L} \log \frac{P_l}{Q_l} \mathbf{1}_{A_{1j}=l}$$

if and only if

$$\begin{split} \sum_{j=1}^{n/2-\gamma/2} X_i + \sum_{j=1}^{\gamma/2} Y_i &\leq \sum_{j=1}^{n/2-\gamma/2} Y_j + \sum_{j=1}^{\gamma/2} X_j \\ \sum_{j=1}^{n/2-\gamma} X_i &\leq \sum_{j=1}^{n/2-\gamma} Y_j \end{split}$$

where  $X_j = \log \frac{P_l}{Q_l}$  with probability  $P_l$  and  $Y_j = \log \frac{P_l}{Q_l}$  in the probability  $Q_l$ . Note that  $\mathbb{E}[X_j] > \mathbb{E}[Y_j]$ .

Standard Chernoff argument yield that the probability of error is at most  $\exp(-(1-2\frac{\gamma}{n})\frac{n}{2}I)$ . This seems to indicate that refinement based on  $\sigma^*$  will achieve that oracle rate so long as  $\sigma^*$  is consistent, i.e.,  $\gamma = o(n)$ .

#### 4 $L \to \infty$ Case.

In this case, suppose that

$$\frac{n}{K \log K} \sum_{l=0}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2 \to \infty$$

Again, we want to know when we can reduce the multi-color problem to a Bernoulli problem. More precisely, we would like to say that there exists a subset of colors  $C \subset \{0,...,l\}$  such that  $\frac{n}{K \log K} I_{tot} \to \infty$  implies that  $\frac{n}{K \log K} I_C \to \infty$ . This would mean that clustering based on the labels

$$\widetilde{A}_{ij} = I_{A_{ij} \in C}$$

is consistent whenever consistency is possible under the oracle setting.

The current results in this section are:

- 1. Choosing C as a singleton is inconsistent.
- 2. Choosing  $C = \{l : P_l > Q_l\}$  is inconsistent in general. It is consistent under further assumptions.

Let us prove the first bullet point.

**Proposition 4.1.** There exist distributions P,Q such that, for any l, we have that  $\frac{n}{K \log K} I_{tot} \to \infty$  but  $\frac{n}{K \log K} I_{C} \to c < \infty$  for  $C = \{l\}$ .

*Proof.* We define P, Q as such

$$P_l = \left\{ \begin{array}{ll} \frac{1}{n} & \text{for } l = 1, ..., \log n \\ 0 & \text{for } l > \log n \end{array} \right. \qquad Q_l = \left\{ \begin{array}{ll} \frac{1}{2n} & \text{for } l = 1, ..., 2\log n \\ 0 & \text{for } l > 2\log n \end{array} \right.$$

with  $P_0 = 1 - \frac{\log n}{n}$  and  $Q_0 = 1 - \frac{\log n}{n}$ .

$$I_{tot} = \left(\sum_{l=1}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2\right) (1 + o(1)) \quad \text{(by lemma 5.1)}$$

$$\geq \left(\log n \left(\sqrt{\frac{1}{n}} - \sqrt{\frac{1}{2n}}\right)^2\right) (1 + o(1))$$

$$= \frac{\log n}{n} (1 - \sqrt{1/2})^2 (1 + o(1))$$

Therefore, for fixed K,  $\frac{n}{K \log K} I_{tot} \to \infty$ .

Let us now upper bound  $I_C$ . Clearly, if  $C = \{0\}$ , then  $I_C = 0$ . Therefore, let  $C = \{l\}$  for some l > 0. First, suppose  $l \in \{1, ..., \log n\}$ .

$$I_C = (\sqrt{P_l} - \sqrt{Q_l})^2 (1 + o(1))$$

$$= (\sqrt{1/n} - \sqrt{1/2n})^2 (1 + o(1))$$

$$\leq \frac{1}{n} (1 + o(1))$$

If  $l \in \{\log n + 1, ..., 2 \log n\}$ , then it is clear that  $I_C = \frac{1}{2n}(1 + o(1))$ . Therefore,  $\frac{nI_C}{K \log K} \to c < \infty$ 

Now we move on to the second point.

We first state a useful lemma:

**Lemma 4.1.** Let  $C \subset \{0,...,l\}$  be a subset of the colors. Let  $P_C = \sum_{l \in C} P_l$  and  $Q_C$  be similarly defined.

Let  $d_C = P_C - Q_C$  and assume w.l.o.g.  $d_C > 0$ . Suppose

$$\frac{d_C}{P_C} \to 0 \qquad P_C \to 0$$

Then, we have that

$$I_C = \Theta\left(\frac{d_C^2}{P_C}\right)$$

Proof.

$$\begin{split} I_C &= (\sqrt{P_C} - \sqrt{Q_C})^2 (1 + o(1)) \\ &= \left( P_C (1 - \sqrt{\frac{Q_C}{P_C}})^2 \right) (1 + o(1)) \\ &= \left( P_C \left( 1 - \sqrt{1 - \frac{P_C - Q_C}{P_C}} \right)^2 \right) (1 + o(1)) \\ &= \left( P_C \left( 1 - (1 - \frac{1}{2} \frac{d_C}{P_C} (1 + o(1))) \right)^2 \right) (1 + o(1)) \\ &= P_C \left( \frac{1}{2} \frac{d_C}{P_C} \right)^2 (1 + o(1)) \\ &= \frac{1}{4} \frac{d_C^2}{P_C} (1 + o(1)) \end{split}$$

Now, we have the following unachievability proposition.

**Proposition 4.2.** There exist distributions P, Q such that  $\frac{n}{K \log K} I_{tot} \to \infty$  but  $\frac{n}{K \log K} I_C \to c < \infty$  for  $C = \{l : P_l \ge Q_l\}$ .

*Proof.* Let x, d > 0 such that  $x, d \to 0$ . Let P, Q be defined as the following:

$$P_{l} = \begin{cases} d & \text{if } l = 1\\ (x - d)2^{-l+1} & \text{if } l > 1 \end{cases} \qquad Q_{l} = \begin{cases} 0 & \text{if } l = 1\\ (x - d)2^{-l+1} & \text{if } l > 1 \end{cases}$$

We therefore have  $P_0 = 1 - x$  and  $Q_0 = 1 - (x - d)$ .

We have that

$$I_{tot} = \sum_{l=1}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2 (1 + o(1))$$
  
=  $d(1 + o(1))$ 

We then set  $d = \frac{\log n}{n}$  and  $x = \frac{\log^2 n}{n}$  so that  $I_C = \Theta(\frac{d^2}{x})$  by Lemma 4.1.

$$I_{tot} = \Theta\left(\frac{\log n}{n}\right) \qquad I_C = \Theta\left(\frac{1}{n}\right)$$

Finally, we have a condition under which we have achievability (ignoring computational difficulties). The intuition is that we need  $P_l$  and  $Q_l$  to be on the same scale.

**Proposition 4.3.** Let f(n) be some sequence such that  $f(n) \to 0$ . Let  $a_1, a_2, ...$  be constants such that  $\sum_{l=1}^{\infty} a_l < \infty$  and  $b_1, b_2, ...$  be constants such that  $\sum_{l=1}^{\infty} b_l < \infty$ 

Let  $P_l = a_l f(n)$  and  $Q_l = b_l f(n)$  for l > 0. Then, we have, so long as  $\frac{n}{K \log K} I_{tot} \to \infty$ , that

$$\frac{n}{K \log K} I_C \to \infty$$

for  $C = \{l : P_l \ge Q_l\}.$ 

*Proof.* First, it is clear that the total variation distance

$$||P - Q||_{TV} = \frac{1}{2} \sum_{l=0}^{\infty} |P_l - Q_l|$$

$$= \frac{1}{2} \left( |P_0 - Q_0| + \sum_{l=1}^{\infty} |a_l f(n) - b_l f(n)| \right)$$

$$= \frac{1}{2} \left( |1 - \sum_{l=1}^{\infty} a_l f(n) - (1 - \sum_{l=1}^{\infty} b_l f(n))| + ||a - b||_1 f(n) \right)$$

$$= \frac{1}{2} f(n) \left( |\sum_{l=1}^{\infty} (a_l - b_l)| + ||a - b||_1 \right)$$

Thus, the total variation, which we denote by d, is of the order  $d = \Theta(f(n))$ . Since  $I_{tot} \leq 2d(1+o(1))$ , we have that  $I_{tot} = O(f(n))$  as well.

We claim that  $I_C = \Theta(f(n))$ . The proof would be complete if we can validate this claim. Without loss of generality, let's suppose that  $P_0 < Q_0$ .

$$I_C = P_C \left( 1 - \sqrt{\frac{Q_C}{P_C}} \right)^2 (1 + o(1))$$

Since  $\frac{Q_C}{P_C} = \frac{\sum_{l \in C} b_l}{\sum_{l \in C} a_l} = c < 1$ , we have that  $I_C = \Theta(P_C) = \Theta(f(n))$ .

#### 5 Technical Lemmas

**Lemma 5.1.** Let  $P = \{P_l\}_{l=0,...,\infty}$  and  $Q = \{Q_l\}_{l=0,...,\infty}$  be two discrete distributions and suppose  $P_0, Q_0 \to 1$ .

Then, we have that  $I \to 0$  and

$$I = (1 + o(1)) \sum_{l=1}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2$$

*Proof.* First, it is clear that if  $P_0, Q_0 \to 1$ , then

$$\sum_{l=0}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2 = (\sqrt{P_0} - \sqrt{Q_0})^2 + \sum_{l=1}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2$$

$$= (\sqrt{P_0} - \sqrt{Q_0})^2 + \sum_{l=1}^{\infty} P_l + \sum_{l=1}^{\infty} Q_l - 2\sum_{l=1}^{\infty} \sqrt{P_l Q_l}$$

$$\leq (\sqrt{P_0} - \sqrt{Q_0})^2 + \sum_{l=1}^{\infty} P_l + \sum_{l=1}^{\infty} Q_l$$

Therefore,  $\lim_{n\to\infty} \sum_{l=0}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2 = 0$ .

$$\begin{split} I &= -2\log\sum_{l=0}^{\infty}\sqrt{P_lQ_l}\\ &= -2\log\left(1-\frac{1}{2}\sum_{l=0}^{\infty}(\sqrt{P_l}-\sqrt{Q_l})^2\right)\\ &= (1+o(1))\sum_{l=0}^{\infty}(\sqrt{P_l}-\sqrt{Q_l})^2 \quad \text{(since the sum tends to 0)} \end{split}$$

We will show that  $(\sqrt{P_0} - \sqrt{Q_0})^2 = o(\sum_{l=1}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2)$  and the result follows immediately.

Let 
$$P' = 1 - P_0$$
 and  $Q' = 1 - Q_0$ .  

$$(\sqrt{P_0} - \sqrt{Q_0})^2 = (\sqrt{1 - P'} + \sqrt{1 - Q'})^2$$

$$= (1 - P') \left(1 - \sqrt{\frac{1 - Q'}{1 - P'}}\right)^2$$

$$= (1 - P') \left(1 - (1 - \frac{1}{2} \left(\frac{Q' - P'}{1 - P'}\right)(1 + o(1)))\right)^2$$

$$\leq (1 - P') \left(\frac{1}{2} \left(\frac{Q' - P'}{1 - P'}\right)(1 + o(1))\right)^2$$

$$\leq \frac{1}{4} \left(\frac{Q' - P'}{1 - P'}\right)^2 (1 + o(1)) \leq \frac{1}{4} (Q' - P')^2 (1 + o(1))$$

$$\sum_{l=1}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2 = \sum_{l=1}^{\infty} P_l + Q_l - 2\sqrt{P_lQ_l}$$

$$\geq P' + Q' - 2\sqrt{\left(\sum_{l=1}^{\infty} P_l\right)\left(\sum_{l=1}^{\infty} Q_l\right)}$$

$$= P' + Q' - 2\sqrt{P'Q'}$$

$$= (\sqrt{P'} - \sqrt{Q'})^2$$

$$= P' \left(1 - \sqrt{\frac{Q'}{P'}}\right)^2$$

$$= P' \left(1 - \sqrt{1 - \frac{P' - Q'}{P'}}(1 + o(1))\right)^2$$

$$\geq P' \left(\frac{1}{2} \frac{P' - Q'}{P'}(1 + o(1))\right)^2$$

$$\geq \frac{1}{4} \left(\frac{(P' - Q')^2}{P'}\right)(1 + o(1))$$

Thus, we have shown that

$$(\sqrt{P_0} - \sqrt{Q_0})^2 \le \frac{1}{4} (Q' - P')^2 (1 + o(1))$$
$$\sum_{l=1}^{\infty} (\sqrt{P_l} - \sqrt{Q_l})^2 \ge \frac{1}{4} \frac{(P' - Q')^2}{P'} (1 + o(1))$$

Since  $P' \to 0$ , the proof is complete.

#### 6 Continuous Distributions

In this section, we suppose that the weights of within-cluster edges are drawn from a density p and that of between-cluster edges are drawn from a density q.

In the following analysis, we will suppose that p, q are mixtures. The first components of both p, q are identical, and the second component of q is a mean-shift of that of p.

$$p(x) = \lambda f(x) + (1 - \lambda)g(x)$$
  
$$q(x) = \lambda f(x - \alpha) + (1 - \lambda)g(x)$$

for some constant  $\lambda \in (0,1)$ , some density g(x) supported on the interval [0,1], and some density f(x) supported on the interval  $[0,1-\alpha]$ .

For simplicity, we will suppose that g(x) = 1.

**Proposition 6.1.** Let  $I = -2\log \int \sqrt{p(x)q(x)}dx$  be the continuous Renyi divergence.

Suppose  $\alpha$  is sufficiently small, then, for some constants c, C, we have that

$$c\alpha^2 \le I \le C\alpha^2$$

*Proof.* First, denote  $H=\int (\sqrt{p(x)}-\sqrt{q(x)})^2 dx$  as the continuous Hellinger distance and we will first show that  $c\alpha^2 \leq H \leq C\alpha^2$ .

Note that  $f(x - \alpha) = f(x) + f'(c_x)\alpha$  for some  $c_x$  in between  $x - \alpha, x$ . Therefore,  $q(x) = p(x) - \lambda f'(c_x)\alpha$ .

$$\int (\sqrt{p(x)} - \sqrt{q(x)})^2 dx$$

$$= \int (\sqrt{p(x)} - \sqrt{p(x)} - \lambda f'(c_x)\alpha)^2 dx$$

$$= \int p(x) \left(1 - \sqrt{1 - \frac{\lambda f'(c_x)\alpha}{p(x)}}\right)^2 dx$$

Since f' is bounded and p(x) is bounded away from zero, we have that, for some  $\eta$  close to 0,

$$= \int p(x) \left( 1 - \left( 1 - \frac{1}{2} \frac{\lambda f'(c_x)\alpha}{p(x)} (1 - \eta) \right) \right)^2 dx$$

$$= \int p(x) \left( \frac{1}{2} \frac{\lambda f'(c_x)\alpha}{p(x)} \right)^2 (1 - \eta) dx$$

$$= \alpha^2 (1 - \eta) \int p(x) \left( \frac{1}{2} \frac{\lambda f'(c_x)}{p(x)} \right)^2 dx$$

Now, it remains to bound I in terms of H. Since

$$I = -2\log \int \sqrt{p(x)q(x)}dx$$

$$= -2\log \left(1 - \frac{1}{2}\int (\sqrt{p(x)} - \sqrt{q(x)})^2 dx\right)$$

$$= -2\log(1 - \frac{1}{2}H)$$

Thus, for  $\alpha$  small enough such that  $H \leq 1/2$ , we have that  $2H \geq I \geq H$ .  $\square$ 

Let the interval [0,1] be divided into L equally spaced sub-intervals. Let l index each of these sub-intervals and let B=1/L be the length of each sub-interval.

Let  $P_l = \int_{\text{Bin}_l} p(x) dx$  and  $Q_l = \int_{\text{Bin}_l} q(x) dx$ . We define the discretized Renyi divergence as  $\widetilde{I} = -2 \log \sum_{l=1}^L \sqrt{P_l Q_l}$  and  $\widetilde{H} = \sum_{l=1}^L (\sqrt{P_l} - \sqrt{Q_l})^2$ .

We will show that  $\widetilde{H} \geq c\alpha^2$  for some constant c and for small enough  $\alpha$ . We will also show that, for any  $\alpha$ ,  $\lim_{L\to\infty} \widetilde{I} = I$ .

will also show that, for any  $\alpha$ ,  $\lim_{L\to\infty} \widetilde{I} = I$ . These facts combined show that  $\widetilde{I} \geq \widetilde{H} \geq c\alpha^2$  and therefore,  $\widetilde{I} - I = \eta\alpha^2$  where  $\lim_{L\to\infty} \eta = 0$ .

**Proposition 6.2.** Let  $P_l, Q_l, \widetilde{I}$  be defined as above. Suppose  $|p'(x)|, |q'(x)| \leq M$  for some absolute constant M.

Then,

$$\lim_{L \to \infty} \widetilde{I} = I$$

*Proof.* Let  $Bin_l = [a_l, b_l]$ .

$$P_{l} = \int_{\text{Bin}_{l}} p(x)dx$$

$$= \int_{a_{l}}^{b_{l}} p(x)dx$$

$$= \int_{a_{l}}^{b_{l}} p(a_{l}) + p'(c_{x})(x - a_{l})dx \quad \text{for } c_{x} \in [a_{l}, b_{l}]$$

$$= Bp(a_{l}) + B^{2}\xi_{l} \quad \text{where } |\xi_{l}| \leq M/2$$

Likewise, we have that  $Q_l = Bq(a_l) + B^2 \xi'_l$ .

$$\begin{split} \widetilde{I} &= -2\log \sum_{l=1}^{L} \sqrt{P_l Q_l} \\ &= -2\log \sum_{l=1}^{L} B \sqrt{(p(a_l) + B\xi_l)(q(a_l) + B\xi_l')} \\ &= -2\log \sum_{l=1}^{L} B \sqrt{p(a_l)q(a_l)} + \left(B^2 \frac{\xi_l}{p(a_l)} + B^2 \frac{\xi_l'}{q(a_l)}\right) (1+\eta) \end{split}$$

where  $\eta \to 0$  as  $B \to 0$ .  $p(a_l), q(a_l)$  are lower bounded and  $\xi_l, \xi'_l$  are bounded. Thus, we have that

$$\begin{split} \lim_{B \to 0} \widetilde{I} &= \lim_{B \to 0} -2 \log \left( \sum_{l=1}^{L} B \sqrt{p(a_l)q(a_l)} + \sum_{l=1}^{L} B^2 c_l (1+\eta) \right) \\ &= -2 \log \left( \int \sqrt{p(x)q(x)} dx \right) \end{split}$$

**Proposition 6.3.** Let  $\widetilde{I} = -2 \log \sum_{l=1}^{L} \sqrt{P_l Q_l}$ . Then, for sufficiently small  $\alpha$ ,  $\widetilde{I} \geq c\alpha^2$  for some absolute constant c.

Proof.

$$\widetilde{H} = \sum_{l=1}^{L} (\sqrt{P_l} - \sqrt{Q_l})^2$$

$$= \sum_{l=1}^{L} P_l \left( 1 - \sqrt{Q_l} P_l \right)^2$$

$$= \sum_{l=1}^{L} P_l \left( 1 - \sqrt{1 - \frac{P_l - Q_l}{P_l}} \right)^2$$

We bound the terms  $P_l - Q_l$  and  $P_l$  separately.

$$P_{l} - Q_{l} = \int_{\text{Bin}_{l}} p(x) - q(x)dx$$
$$= \int_{a_{l}}^{b_{l}} \lambda f'(c_{x})\alpha dx$$
$$= \lambda \xi_{l} \alpha B$$

where  $|\xi_l| \leq M$ .

$$P_l = \int_{\text{Bin}_l} p(x)dx \ge B(1-\lambda)$$

Thus, for small enough  $\alpha$ , we have that

$$\widetilde{H} = \sum_{l=1}^{L} P_l \left( c_l \alpha \right)^2$$

$$\geq c \alpha^2$$

for some constant c. The claim follows since  $\widetilde{I} \geq \widetilde{H}$ .

# 7 A More General Condition

## References

[1] Anderson Y Zhang and Harrison H Zhou. Minimax rates of community detection in stochastic block model.