# Multi-client LAN/WAN Performance Analysis of Ninf: a High-Performance Global Computing System



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http://ninf.etl.go.jp/

## Towards Global Computing Infrastructure

#### Rapid increase in speed and availability of network

→ Computational and Data Resources are collectively employed to solve large-scale problems.

Global Computing (Metacomputing, The "Grid")

**Ninf** (Network Infrastructure for Global Computing)

c.f., NetSolve, Legion, RCS, Javelin, Globus etc.

### Overview of Ninf

Other Global Computing Systems, e.g., NetSolve via Adapters Ninf DB Server makes available multiple remote Meta **Computing and DB servers** Server Client Library for Fortran, C and Internet Java. Meta Meta Multiple calls in a network are Server, Server, coordinated by **Metaservers**. WWW-based interfaces **Ninf** → NinfCalc+ **Ninf Client Library** Computational Java-based global computing Server infrastructure Ninf\_call("linpack", ..); → Ninflet Ninf RPC **Program** 

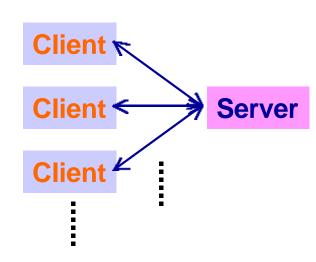
# Issues of Global Computing systems

- Communication Performance (throughput, latency)
  - − High latency → deterrent to performance?
- Computing Server Selection
  - Criteria to select servers?
  - Can supercomputers/MPPs be effectively shared by multiple remote clients? (Scheduling characteristics)
- Sharing by Multiple clients
  - Global computing servers on supercomputer/MPP OSes effectively handle multiple parallel tasks?
- Remote Library Design and Reuse
  - Allocate processors in a task-parallel / data-parallel manner ?
  - The right "remote library" design, decomposition, API, etc.

### **Outline of Benchmarks**

Single-Client Benchmarks:

- Client <---> Server
- Ninf Baseline Performance → LAN, Linpack
- Multi-Client Benchmarks :
  - Communication Performance
    - → LAN / WAN (Single-site / Multi-site)
  - Remote Library Design and Reuse
    - → Task Parallel / Data Parallel
  - Robustness of Computational Server (Cray J90)
  - Ratio of Computation and Communication of Library
    - → Linpack / EP

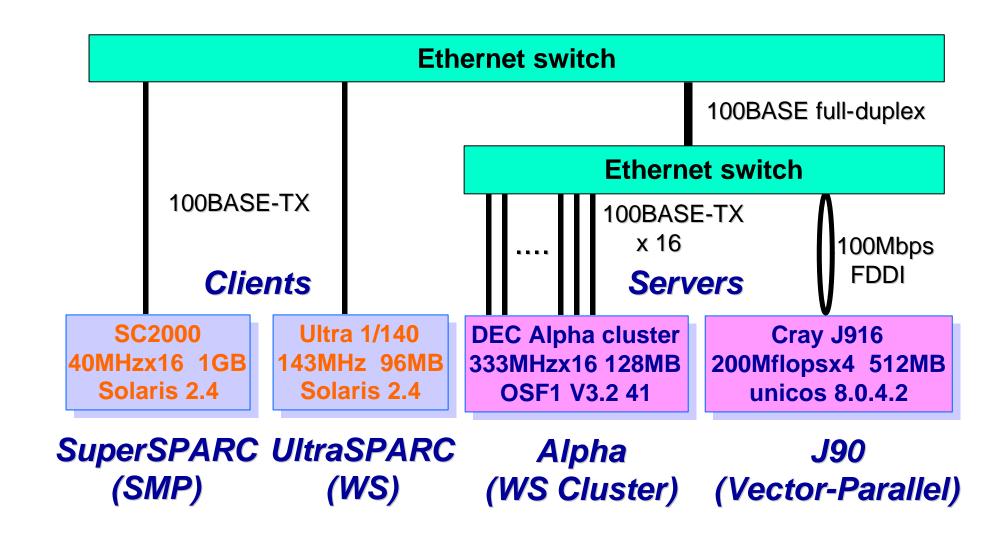


### **Benchmarks**

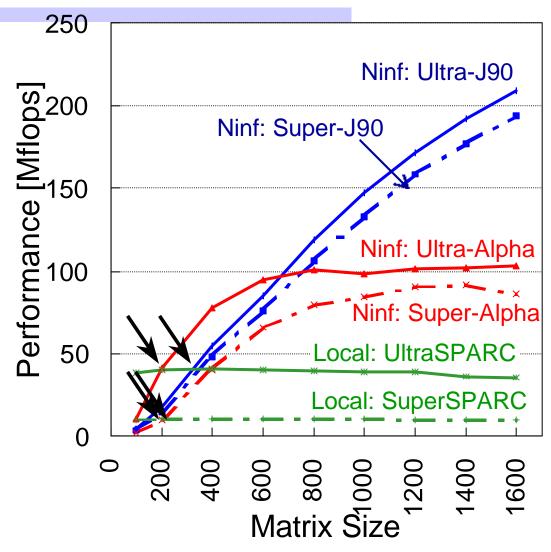
- Linpack : Gaussian Elimination
  - Computation:  $2/3n + ^32n$  [flops]
  - Communication:  $8n + ^220n + O(1)$  [bytes]
  - Performance:  $(2/3n + ^32n) / ^2$  lapsed Time [flops]
- **EP (NASPB Kernel)**: Random Number Generation
  - Computation: 2 n+1 [ops]
  - Communication: O(1) [bytes]
  - Performance:
     2 / Elapsed Time [ops]

**Elapsed Time** = Communication + Computation

# LAN Single-client Benchmarking Environment (at ETL)



# LAN Single Client Linpack Results



- Ninf is faster than Local at n= 150 ~ 300
- For Ninf\_call to J90, Ninf performance is not saturated.

(J90's Local achieves 600Mflops when n=1600)

Ninf performance quickly overtakes Local.

The effects of client machine's performance difference are small.

# Multi-client Benchmarks (LAN, WAN)

#### A Model Client Program

**Linpack** and **EP** are both repeatedly called:

- Each client performs a Ninf\_call on the interval of s seconds with probability p o s = 3, p = 1/2 chosen.
- Number of clients : c , problem size : n.

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\rightarrow c = 1, 2, 4, 8, 16, Linpack: n = 600, 1000, 1400 EP: n = 24(Sample)
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#### Parallel Processing on the server (J90 4PE)

Linpack: <u>1PE ver.</u> --- Task Parallel

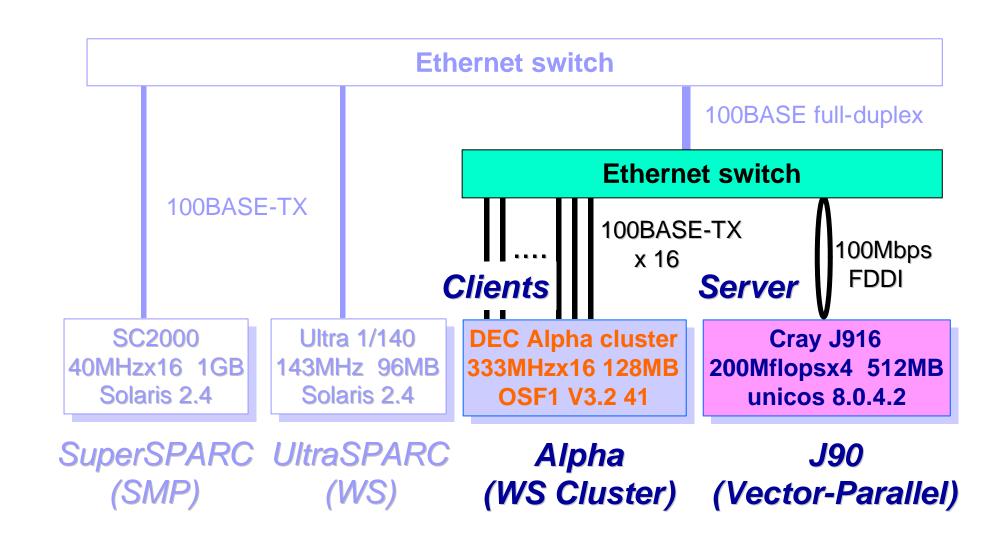
1PE Execution and Parallel Processing

4PE ver. --- Data Parallel

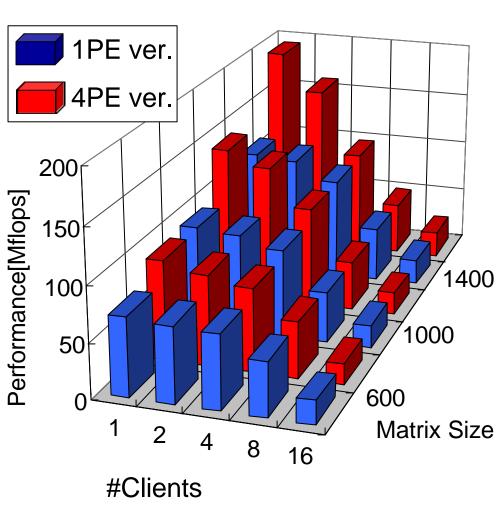
4PE Execution and Single Processing

– EP: <u>1PE ver.</u> --- Task Parallel

# LAN Multi-client Benchmarking Environment (at ETL)

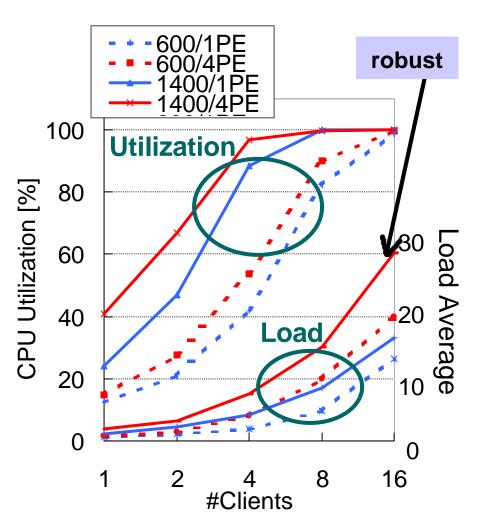


# LAN Linpack Benchmark Results Average Performance



- 4PE ver. exhibits higher performance for a small c,
- while there is little
   performance edge for the
   1PE ver. for a larger c.
  - Numerical core speed
  - Small overhead of switching parallel tasks on J90
  - Overlapped communication
- Average response and waiting time didn't differ depending on n, c, or 1PE / 4PE ver.

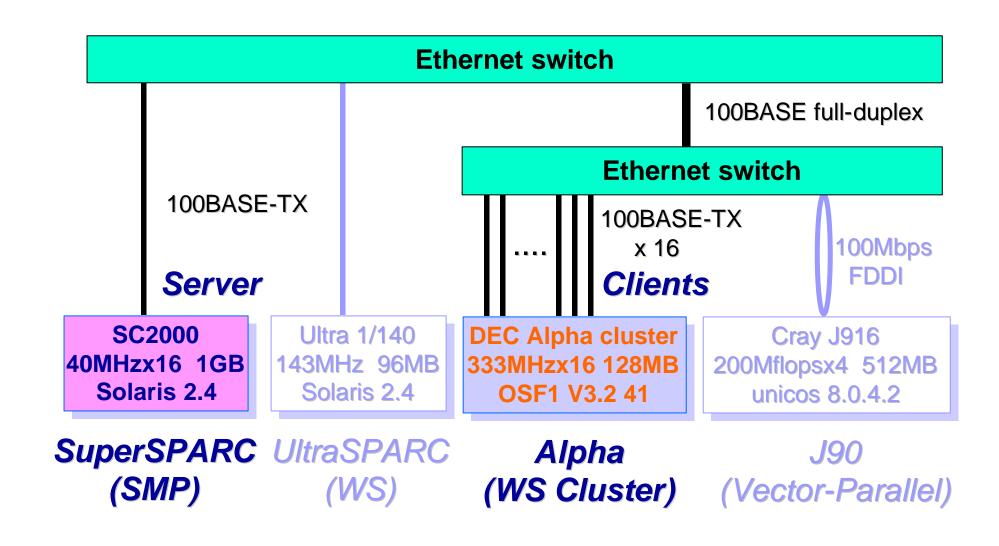
# LAN Linpack Benchmark Results CPU Utilization and Load Average



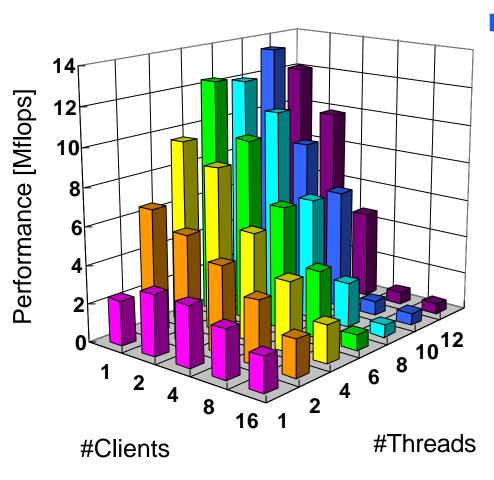
- 4PE ver. exhibits higher Utilization and Load.
- → The optimized parallel library would be appropriate for J90.
- Utilization and Load are higher as n and c increase.
- The J90 Ninf server continued to work flawlessly.

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(even for n=1400, 4PE ver., max. load average : 30)
```

# LAN SMP Multi-client Benchmarking Environment (at ETL)



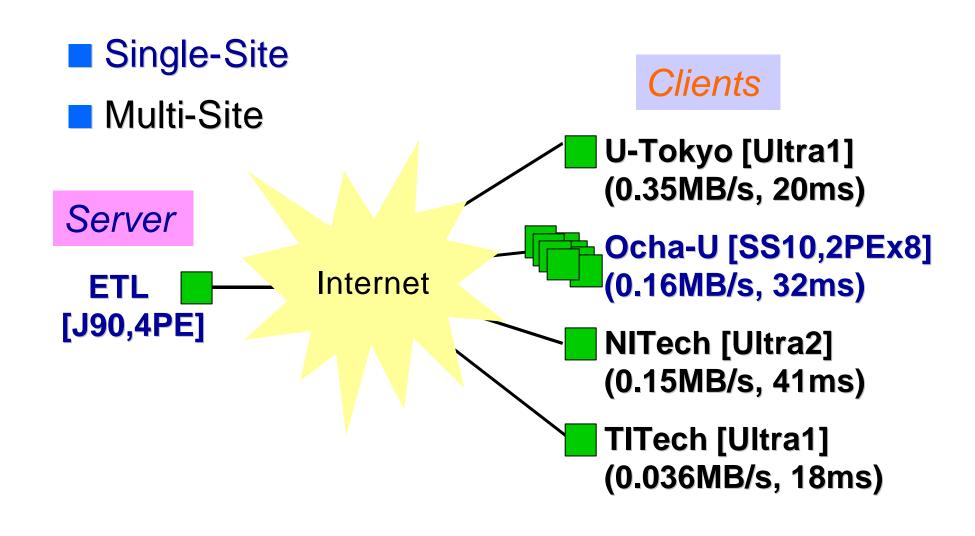
# Average Performance of Multi-thread LAN Linpack on SPARC SMP



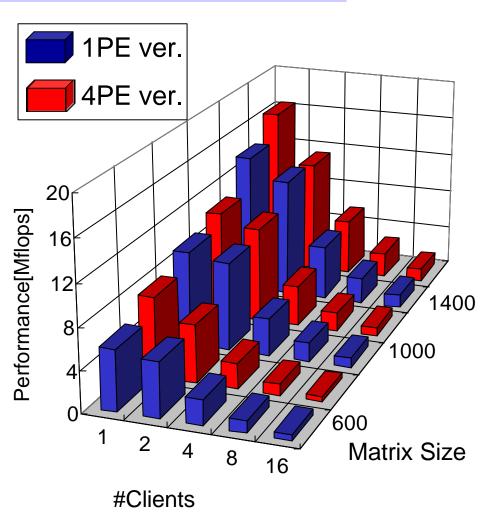
- Highly multi-threaded ver. exhibit notable slowdown.
  - Solaris on SPARC SMP is optimized for handling the requests.
  - do not co-schedule multiple threads.

thread-switching overhead. (cache / TLB misses etc.)

# WAN Multi-client Benchmarking Environment

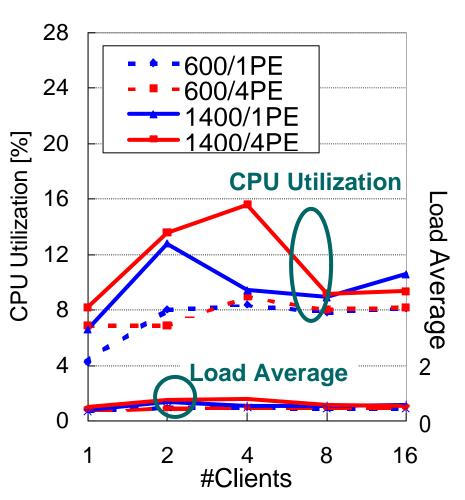


# Single-Site WAN Linpack Benchmark Results Average Performance



- Ave. performance deteriorates as c increases.
- $\rightarrow$  The network throughput saturates even for c=1 or 2.
- Exhibited almost the same characteristics as LAN.
- → Better to use the parallel library versions for WAN clients as well for J90.

# Single-Site WAN Linpack Benchmark Results CPU Utilization and Load Average

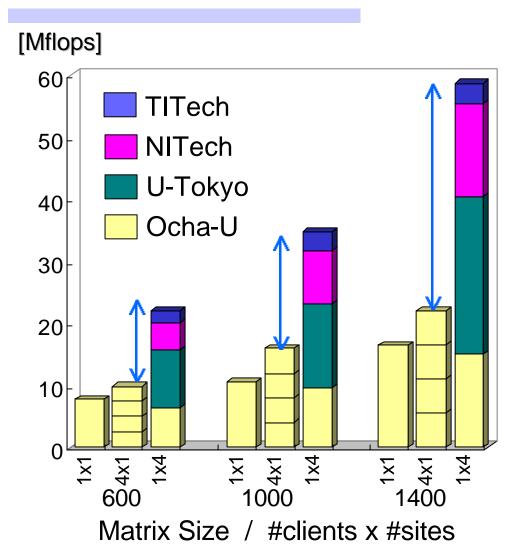


- Utilization and Load remain low even for c=16.
- → It is difficult for the global computing servers to process numerous requests.

(for communication intensive tasks)

- → A good load balancing algorithm is needed
  - server load
  - network traffic and topology
  - client location etc.

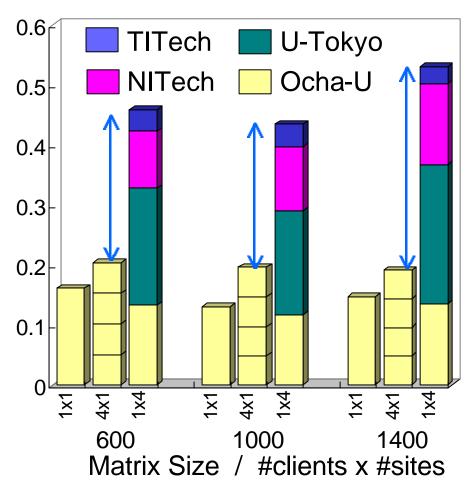
# Single/Multi-site WAN Linpack Benchmark Results Average Performance (c = 4, 4PE ver.)



- Total performance is higher as *n* increases.
- Aggregate performance from multi-site is substantially higher than from single-site for the same *c*.

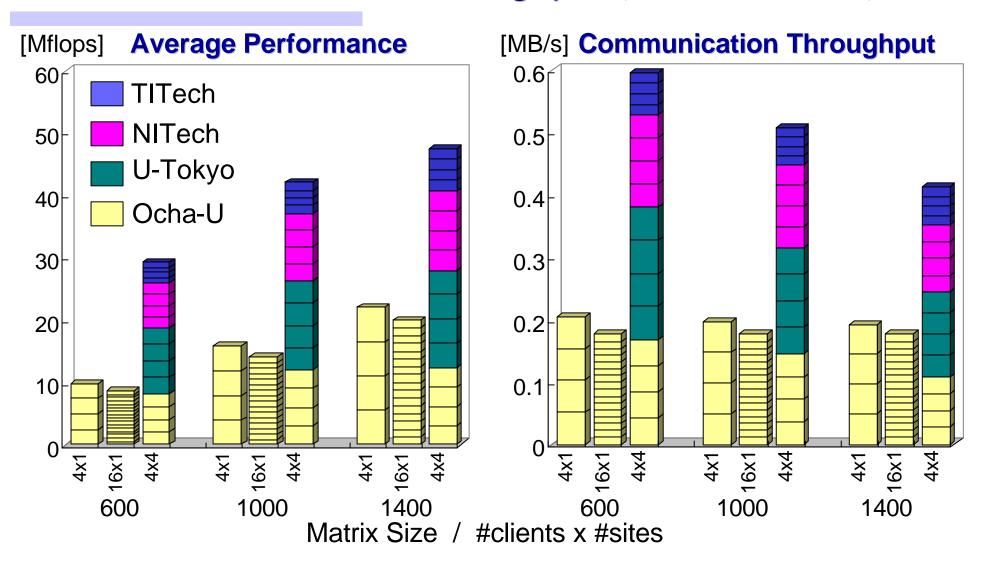
# Single/Multi-site WAN Linpack Benchmark Results Communication Throughput (c = 4, 4PE ver.)

#### [MB/s]



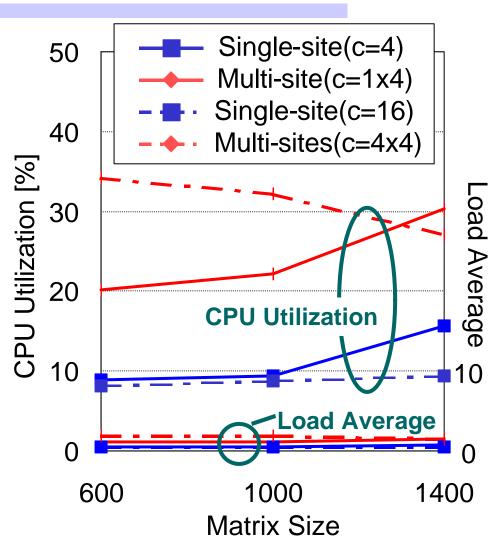
- Total communication throughput is nearly constant for each *n*.
- Aggregate communication throughput from multi-site is substantially higher than from single-site.
- → For communication intensive tasks, point-to-point bandwidth is the dominant factor in performance. (not latency)

# Single/Multi-site WAN Linpack Benchmark Results Performance and Throughput (c = 16, 4PE ver.)



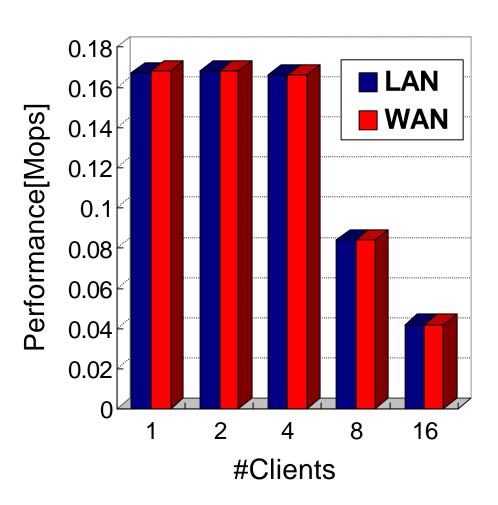
#### Single/Multi-site WAN Linpack Benchmark Results

# **CPU Utilization and Load Average**



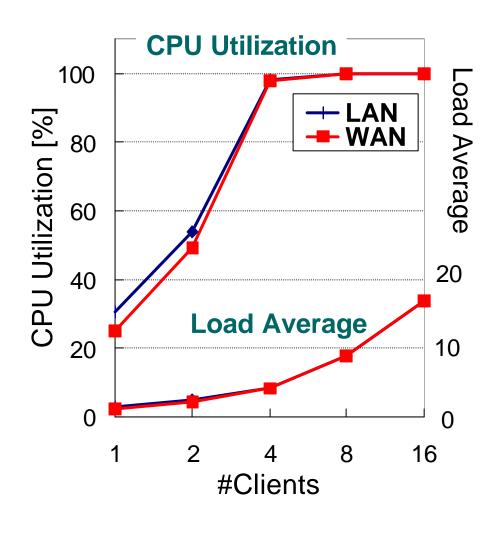
- Utilization and Load are greater for multi-site. c.f., single site.
- The J90 server does not saturate for n and c.
  - Network bandwidth
     saturation again the cause.
- → Utilization and Load alone are NOT suitable criteria for load balancing of global computing.

# LAN/WAN EP Benchmark Results Average Performance (1PE ver.)



- Because of task-parallel execution, the performances decline when c=8, 16.
- In both LAN and WAN cases Ninf\_call performances are almost equivalent.

# LAN/WAN EP Benchmark Results CPU Utilization and Load Average



- Both LAN and WAN cases are almost equivalent for Utilization and Load.
- → Global computing can now be considered quite feasible for this class of applications:
  - parallel rendering / imaging
  - parameter sensitivity analysis

### Related Work

- RPC based systems → use existing programming languages
  - NetSolve [Casanova and Dongarra, Univ. Tennessee]
    - The same basic API as Ninf\_call (now interchangeable)
    - load-balancing with a daemon process called Agent.
  - RCS [Arbenz, ETH Zurich]
    - PVM-based
- Systems using parallel distributed language etc.
  - Legion [Grimshaw, Univ. Virginia]
    - An user distributes his programs written with the parallel object-oriented language Mentat.
  - Javelin [Schauser et al., UCSB]
    - High portability due to using Java and WWW.
- Global scheduling systems NWS, AppLeS, DQS

### Conclusion

- The current Global Computing systems are likely to work well in both LAN and WAN situations.
- The use of optimized parallel library would be sufficient for vector computing servers.
- In LAN, computing server performance dictates overall performance, while in WAN limitation of communication throughput is more significant. (Esp. communication-intensive applications)
- We expect multiple client requests will be issued from different sites, causing "false" lowering of load ave..

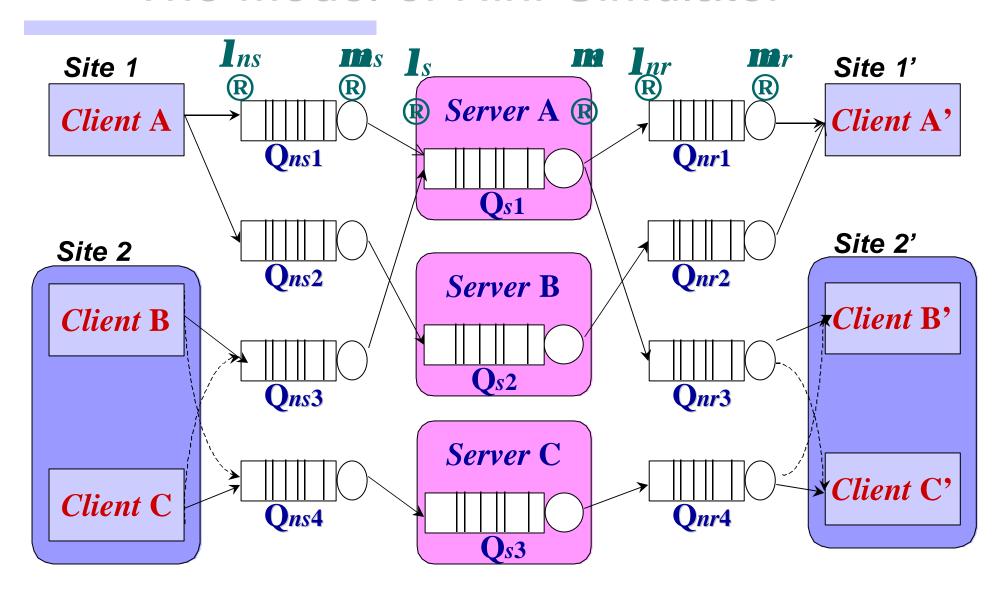




#### **Future Work**

- Load Balancing for the Global Computing
  - Meta Server (c.f. NWS)
  - Develop a theoretical model
    - Performance analysis for more practical situations
    - Global Computing Simulator
- Make the Ninf system more powerful
  - Guaranteeing Performance in multi-client global computing
  - Server Job-handling Methodology (FCFS → SJF)
  - Multi-Job Scheduling for MPP servers (FPFS, FPMPFS)
  - Extension of the API for MPP's network services

### The Model of Ninf Simulator



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### **Exhibition**

Electrotechnical Laboratory

