

1 Glivenko-Cantelli Theorem

Theorem 1.1. Let $(X_i)_{i=1}^N$ be i.i.d. with distribution function F and $G_n(\omega, x) = \frac{1}{n} \sum_{i=1}^n 1_{X_i(\omega) \leq x}$ be its empirical distribution function. Then

$$\sup_x |G_n(\omega, x) - F(x)| \xrightarrow{a.s.} 0 \quad \text{as } n \rightarrow \infty \quad (1.1)$$

(i.e. $\|G_n(\omega, x) - F(x)\|_\infty \rightarrow 0$ or $G_n \xrightarrow{L_\infty} F$)

To prove theorem 1.1, we need a lemma proved in a later homework.

Definition 1.2. For a CDF $F : \mathbb{R} \rightarrow [0, 1]$, $x \in \mathbb{R}$ is an *atom* if

$$F(x) - F(x^-) = P(X = x) > 0 \quad (1.2)$$

Lemma 1.3. Let F_n, F be distribution functions. If

(a) For all $x \in \mathbb{Q}$: $F_n(x) \rightarrow F(x)$

(b) For each atom x of F : $F_n(x) \rightarrow F(x)$ and $F_n(x^-) \rightarrow F(x^-)$

then $\sup_x |F_n(x) - F(x)| \rightarrow 0$ almost surely.

Proof of theorem 1.1. Fix $x \in \mathbb{R}$. The events $\{X_i \leq x\}_{i=1}^n$ are i.i.d. with probability $= F(x)$. By SLLN, $G_n(\omega, x) \xrightarrow{a.s.} F(x)$ as $n \rightarrow \infty$. Consider $S = \mathbb{Q} \cup \{\text{atoms of } F\}$, which is countable (TODO: why?). Notice $P(G_n(\omega, x) \rightarrow F(x) \quad \forall x \in S) = 1$ so by lemma 1.3 $P(\sup_x |G_n(\omega, x) - F(x)| \rightarrow 0) = 1$. \square

2 Gambling on a favorable game

Suppose we are playing a game where we stake an amount $s \in \mathbb{R}$ and receive payoff

$$\begin{cases} +s, & \text{w.p. } \frac{1}{2} + \alpha \\ -s, & \text{w.p. } \frac{1}{2} - \alpha \end{cases} \quad (2.1)$$

Consider a strategy where at every time the stake s is equal to some proportion $q \in$

$[0, 1]$ of your current total. Let X_n denote your total fortune after n bets. Then

$$X_{n+1} = (1 - q)X_n + \begin{cases} 2qX_n, & \text{if win} \\ 0, & \text{if loose} \end{cases} \quad (2.2)$$

$$= (1 - q)X_n + 2qX_n \underbrace{1_{A_{n+1}}}_{\text{win } (n+1)\text{st bet}} \quad (2.3)$$

$$= (1 - q + 2q1_{A_{n+1}})X_n \quad (2.4)$$

$$\implies X_n = x_0 \prod_{i=1}^n (1 - q + 2q1_{A_i}) \quad (2.5)$$

$$\frac{\log X_n}{n} = \frac{\log x_0}{n} + \frac{1}{n} \sum_{i=1}^n \underbrace{\log(1 - q + 2q1_{A_i})}_{Y_i} \quad (2.6)$$

By SLLN, as $n \rightarrow \infty$

$$\frac{\log Y_n}{n} \xrightarrow{\text{a.s.}} \mathbb{E}Y \quad (2.7)$$

(Note: $\frac{1}{n} \log X_n \rightarrow c$ is slightly weaker than $x_n \approx e^{cn}$, $c = \text{"asymptotic growth rate"}$)
The optimal choice of q should maximize $\mathbb{E}Y$

$$\mathbb{E}Y = \left(\frac{1}{2} + \alpha\right) \log(1 + q) + \left(\frac{1}{2} - \alpha\right) \log(1 - q) \quad (2.8)$$

$$\approx 2\alpha q - \frac{1}{2}q^2 \quad \text{for } \alpha, q \text{ small} \quad (2.9)$$

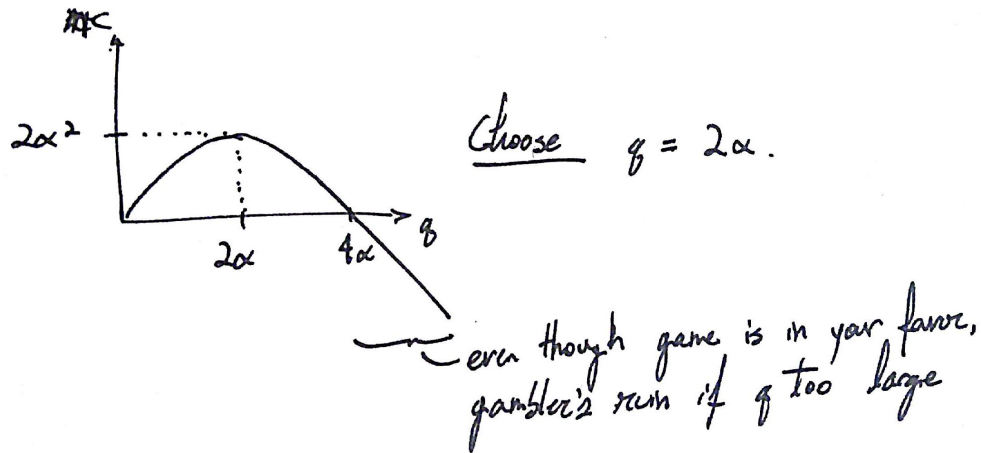


Figure 1: Asymptotic growth rate for different bet proportions q

\therefore the optimal choice is $q = 2\alpha$

However, notice that

$$\mathbb{E}X_n = x_0(1 + 2q\alpha)^n \rightarrow \infty \quad \text{but } X_n \xrightarrow{\text{a.s.}} 0 \text{ if } q \geq q_{\text{crit}} \approx 4\alpha \quad (2.10)$$

3 Almost-sure limits for maxima

Example 3.1. $(X_i)_1^n \stackrel{\text{iid}}{\sim} \text{Exp}(1)$, $P(X > x) = e^{-x}$.

Write $M_n = \sup_{1 \leq i \leq n} X_i$. Then

$$\limsup_n \frac{X_n}{\log n} \stackrel{\text{a.s.}}{=} 1 \quad (3.1)$$

and

$$\frac{M_n}{\log n} \xrightarrow{\text{a.s.}} 1 \quad (3.2)$$

Proof. Fix $\varepsilon > 0$. $P(X_n / \log n > 1 + \varepsilon) = \exp(-(1 + \varepsilon) \log n) = n^{-(1+\varepsilon)}$

$$\sum_{n \in \mathbb{N}} n^{-(1+\varepsilon)} < \infty \xRightarrow{(\text{BC } 1)} P\left(\frac{X_n}{\log n} \leq 1 + \varepsilon \text{ ult.}\right) = 1 \quad (3.3)$$

$$\iff \limsup_n \frac{X_n}{\log n} \stackrel{\text{a.s.}}{\leq} 1 + \varepsilon \quad (3.4)$$

$$\xRightarrow{\varepsilon \rightarrow 0} \limsup_n \frac{X_n}{\log n} \stackrel{\text{a.s.}}{\leq} 1 \quad (3.5)$$

To obtain a lower bound, $P(\underbrace{X_n / \log n \geq 1 - \varepsilon}_{\text{indep.}}) = n^{-(1-\varepsilon)}$ so

$$\sum_{n \in \mathbb{N}} n^{-(1-\varepsilon)} = \infty \xRightarrow{(\text{BC } 2)} P\left(\frac{X_n}{\log n} \geq 1 - \varepsilon \text{ i.o.}\right) = 1 \quad (3.6)$$

$$\xRightarrow{\varepsilon \rightarrow 0} \limsup_n \frac{X_n}{\log n} \stackrel{\text{a.s.}}{\geq} 1 \quad (3.7)$$

Together, we have $\limsup_n X_n / \log n \stackrel{\text{a.s.}}{=} 1$.

To prove the second part, we first prove a deterministic lemma:

Lemma 3.2 (Deterministic Lemma). *If $X_n \geq 0$ and $0 < b_n \uparrow \infty$, then $\limsup_n \frac{\max_{1 \leq i \leq n} X_i}{b_n} = \limsup_n \frac{X_n}{b_n}$.*

Proof. “ \geq ” is obvious. Fix j .

$$\limsup_n \frac{\max_{1 \leq i \leq n} X_i}{b_n} = \limsup_n \frac{\max_{j \leq i \leq n} X_i}{b_n} \quad b_n \uparrow \infty, x_j \text{ fixed} \quad (3.8)$$

$$\leq \lim_n \max_{j \leq i \leq n} \frac{x_i}{b_i} \quad \frac{x_i}{b_i} \geq \frac{x_i}{b_n} \quad (3.9)$$

$$= \sup_{i \geq j} \frac{x_i}{b_i} \quad (3.10)$$

Letting $j \rightarrow \infty$ shows “ \leq ”. □

Combining eq. (3.1) and lemma 3.2 imply $\limsup_n \frac{M_n}{\log n} \stackrel{\text{a.s.}}{=} 1$.

It remains to show $\liminf_n \frac{M_n}{\log n} \stackrel{\text{a.s.}}{=} 1$. But since $1 \stackrel{\text{a.s.}}{=} \limsup_n M_n / \log n \geq \liminf_n M_n / \log n$, it suffices to show $\liminf_n M_n / \log n \stackrel{\text{a.s.}}{\geq} 1$

Fix $\varepsilon > 0$.

$$P(M_n \leq (1 - \varepsilon) \log n) = P(X \leq (1 - \varepsilon) \log n)^n \quad (3.11)$$

$$= (1 - n^{-(1-\varepsilon)})^n \quad (3.12)$$

$$\leq \exp\left(-n^{-(1-\varepsilon)}\right)^n \quad 1 - x \leq e^{-x} \quad (3.13)$$

$$= \exp(-n^\varepsilon) \quad (3.14)$$

$$(3.15)$$

$$\sum_n P(M_n \leq (1 - \varepsilon) \log n) \leq 1 - \frac{1}{1 - e^\varepsilon} = \frac{1}{1 - e^\varepsilon} < \infty \quad (3.16)$$

$$\stackrel{\text{(BC 1)}}{\implies} M_n \geq (1 - \varepsilon) \log n \text{ ult. a.s.} \quad (3.17)$$

$$\implies \liminf_n \frac{M_n}{\log n} \stackrel{\text{a.s.}}{\geq} 1 - \varepsilon \quad (3.18)$$

$$\iff \liminf_n \frac{M_n}{\log n} \stackrel{\text{a.s.}}{\geq} 1 - \varepsilon \quad (3.19)$$

$$\stackrel{\varepsilon \rightarrow 0}{\implies} \liminf_n \frac{M_n}{\log n} \stackrel{\text{a.s.}}{\geq} 1 \quad (3.20)$$

□

Remark 3.3. Here $X_n / \log n \rightarrow 0$ in probability (i.e. $P(X_n / \log n \geq \varepsilon) = n^{-\varepsilon} \rightarrow 0$, but not a.s. (which requires showing $P(\lim_n X_n / \log n = 0) = 1$).

4 Second moment Strong Law of Large Numbers

Theorem 4.1 (Second moment SLLN). *Given $(X_i)_{i=1}^n$ with:*

- (a) $\mathbb{E}X_i = 0$ for all i
- (b) $\sup_i \mathbb{E}X_i^2 = B < \infty$
- (c) **Orthogonal:** $\mathbb{E}[X_i X_j] = 0$ for all $i \neq j$

Write $S_n = \sum_{i=1}^n X_i$. Then $S_n / n \xrightarrow{\text{a.s.}} 0$.

We first show a deterministic lemma used in the proof.

Lemma 4.2 (Deterministic Lemma). *Given $S_n \in \mathbb{R}$, to show $S_n / n \rightarrow 0$ it suffices to show \exists subsequence $n(j) \uparrow \infty$ such that:*

- (a) $S_{n(j)} / n(j) \rightarrow 0$ as $j \rightarrow \infty$

(b) $D_j/n(j) \rightarrow 0$ where $D_j := \max_{n(j) \leq n \leq n(j+1)} |S_n - S_{n(j)}|$

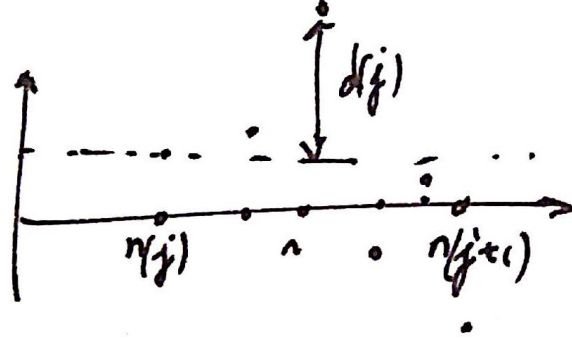


Figure 2: $d(j)$ and $n(j)$ as defined for lemma 4.2

Proof of lemma 4.2. Given n , for some j where $n(j) \leq n < n(j+1)$

$$\left| \frac{S_n}{n} \right| \leq \left| \frac{S_{n(j)}}{n(j)} \right| \leq \frac{|S_{n(j)}| + D_j}{n(j)} \xrightarrow{\text{a.s.}} 0 \quad (4.1)$$

□

Proof of theorem 4.1. $\text{Var}(S_n) \leq nB$ so by Chebyshev's inequality

$$P\left(\frac{|S_n|}{n} \geq \epsilon\right) \leq \frac{nB}{n^2\epsilon^2} = \frac{B}{n\epsilon^2} \quad (4.2)$$

Take $a(j) = j^2$.

$$P\left(\frac{|S_{n(j)}|}{n(j)} \geq \epsilon\right) \leq \frac{B}{\epsilon^2 j^2} \quad (4.3)$$

$\sum_j P\left(\frac{|S_{n(j)}|}{n(j)} \geq \epsilon\right) < \infty$ so by Borel-Cantelli 1 $\limsup_n S_{n(j)}/n(j) \leq \epsilon$ a.s. Taking $\epsilon \rightarrow 0$ yields $\limsup_n S_{n(j)}/n(j) \stackrel{\text{a.s.}}{=} 0$, and since $S_{n(j)}/n(j) \geq 0$ for all j , $\liminf_n S_{n(j)}/n(j) \geq 0$ hence $S_{n(j)}/n(j) \stackrel{\text{a.s.}}{\rightarrow} 0$.

By lemma 4.2 it suffices to show $D_j/j^2 \stackrel{\text{a.s.}}{\rightarrow} 0$ for $D_j = \max_{j^2 \leq n \leq (j+1)^2} |S_n - S_{j^2}|$.

$$D_j^2 = \max_{j^2 \leq n \leq (j+1)^2} (S_n - S_{j^2})^2 \leq \sum_{n=j^2}^{(j+1)^2} (S_n - S_{j^2})^2 \quad (4.4)$$

$$\mathbb{E}D_j^2 \leq \sum_{n=j^2}^{(j+1)^2} \mathbb{E}(S_n - S_{j^2})^2 = \sum_{n=j^2}^{(j+1)^2} \text{Var}(S_n - S_{j^2}) = \sum_{n=j^2}^{(j+1)^2} \text{Var}\left(\sum_{i=j^2+1}^n X_i\right) \quad (4.5)$$

$$\leq \sum_{n=j^2}^{(j+1)^2} B(n - j^2) = B \sum_{i=1}^{2j+1} i = \frac{1}{2}(2j+1)(2j+2)B \quad (4.6)$$

By Chebyshev bound

$$P\left(\frac{D_j}{j} \geq \varepsilon\right) \leq \frac{\mathbb{E}D_j^2}{\varepsilon^2 j^4} \in O(j^{-2}) \quad (4.7)$$

By Borel-Cantelli 1, $D_j/j^2 \xrightarrow{\text{a.s.}} 0$. □

Remark 4.3. Theorem 4.1 does not rely on independence, only bounded variance and orthogonality.

5 Misc. MT

Definition 5.1. For a RV X and a non-negative integrable RV Y ($Y \geq 0$, $\mathbb{E}Y \leq \infty$), X is *dominated by* Y (written $X \ll Y$) means $|X_n| \stackrel{\text{a.s.}}{\leq} Y$.

Theorem 5.2 (Dominated Convergence Theorem). *If $X_n \xrightarrow{\text{a.s.}} X$ and $X_n \ll Y$, then:*

- (a) $\mathbb{E}X_n \rightarrow \mathbb{E}X$
- (b) $\mathbb{E}|X_n - X| \rightarrow 0$
- (c) $\mathbb{E}|X| < \infty$

Proof. Fix $\epsilon > 0$. Define $A_N = \{|X_n - X| \leq \epsilon \text{ for all } n \geq N\}$. Then $A_N \uparrow A$, $P(A) = 1$ implies $A_n^c \downarrow A^c$, $P(A^c) = 0$.

$$\mathbb{E}|X_N - X| = \mathbb{E}|X_N - X|1_{A_N} + \mathbb{E}|X_N - X|1_{A_N^c} \quad (5.1)$$

$$\leq \epsilon + \mathbb{E}|X_N - X|1_{A_N^c} \quad \text{TODO: Bernstein's theorem from previous lecture} \quad (5.2)$$

$$\limsup_N \mathbb{E}|X_N - X| \leq \epsilon + 0A_n^c \rightarrow A^c \text{ by monotone convergence} \quad (5.3)$$

True $\forall \epsilon$ so $\mathbb{E}|X_N - X| \rightarrow 0$. □

This proof may need Fatou's lemma.