1. TAD REGISTRO

TAD GENEROTIPO

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igualdad observacional
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\begin{pmatrix}
(\forall r_1, r_2: \text{reg}) \ (r_1 =_{\text{obs}} r_2) \Leftrightarrow \left( \text{campos}(r_1) =_{\text{obs}} \text{campos}(r_2) \land_{\text{L}} \left( (\forall \text{c:Campos}) ((\forall \text{c:Campos}) \land_{\text{L}} ((\forall \text{c:Campos}) ((\forall \text{c:Campos})
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géneros reg

exporta reg, generadores, observadores, otras operaciones

usa NAT, STRING, CAMPO, TIPO

observadores básicos

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\begin{array}{lll} {\rm Campos: reg} & \longrightarrow {\rm conj}({\rm Campo}) \\ {\rm ValorEn: reg} & {\rm r} \times {\rm campo} & {\rm c} & \longrightarrow {\rm Nat} \\ {\rm PalabraEn: reg} & {\rm r} \times {\rm campo} & {\rm c} & \longrightarrow {\rm String} \end{array} \\ & & \left\{ \neg {\rm Nat?(tipo(c))} \wedge {\rm c} \in {\rm Campos(r)} \right\} \\ \end{array}
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generadores

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\begin{array}{llll} \text{NuevoRegistro} : \bullet & \longrightarrow \text{reg} \\ \text{agValor} : \text{reg} & r \times \text{campo} & c \times \text{nat} & n & \longrightarrow \text{reg} \\ \text{agPalabra} : \text{reg} & r \times \text{campo} & c \times \text{string} & s & \longrightarrow \text{reg} \\ \end{array} \\ \left\{ c \notin \text{campos(r)} \wedge \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text{Nat?(tipo(c))} \right\} \\ \left\{ c \notin \text{campos(r)} \wedge \neg \text
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otras operaciones

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noRepiten : reg r_1 \times \text{reg} r_2 \times \text{campo} c \longrightarrow bool { c \in \text{campos}(r_1) \land c \in \text{campos}(r_2)} Combinar : reg r_1 \times \text{reg} r_2 \longrightarrow \text{reg}
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CombinarDeAcuerdoA : reg $r_1 \times \text{reg}$ $r_2 \times \text{conj}(\text{campo})$ cs \longrightarrow reg **axiomas** $\forall r_1, r_2 : \text{reg}, \forall c : \text{campo}, \forall cs : \text{conj}(\text{campo}), \forall n : \text{nat}, \forall s : \text{string campos}(\text{nuevoReg}) \equiv \emptyset$

campos(AgValor(r_1,c,n) \equiv Ag($c,campos(r_1)$)

 $Nat?(String) \equiv False$

 $ElOtro(t) \equiv if Nat?(t)$ then String else Nat fi

Fin TAD