Coco: Runtime Reasoning About Conflicting Commitments

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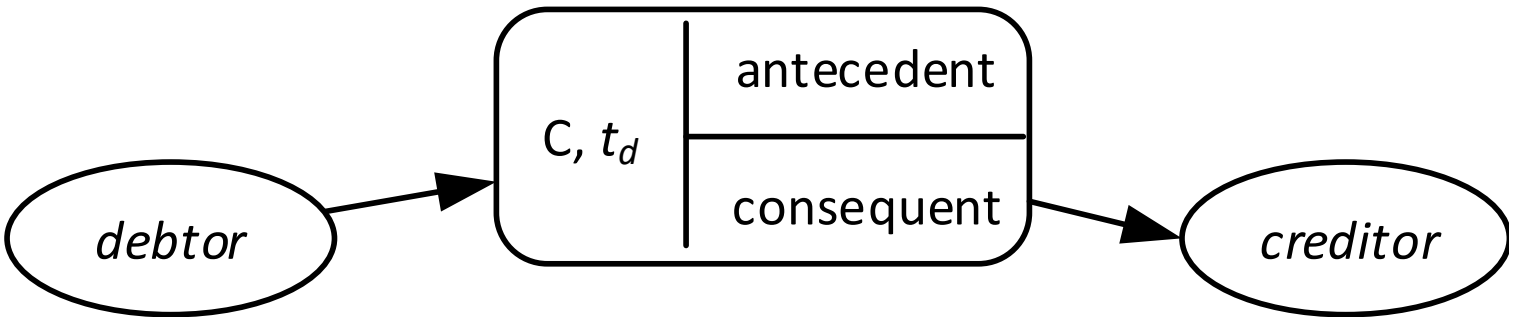
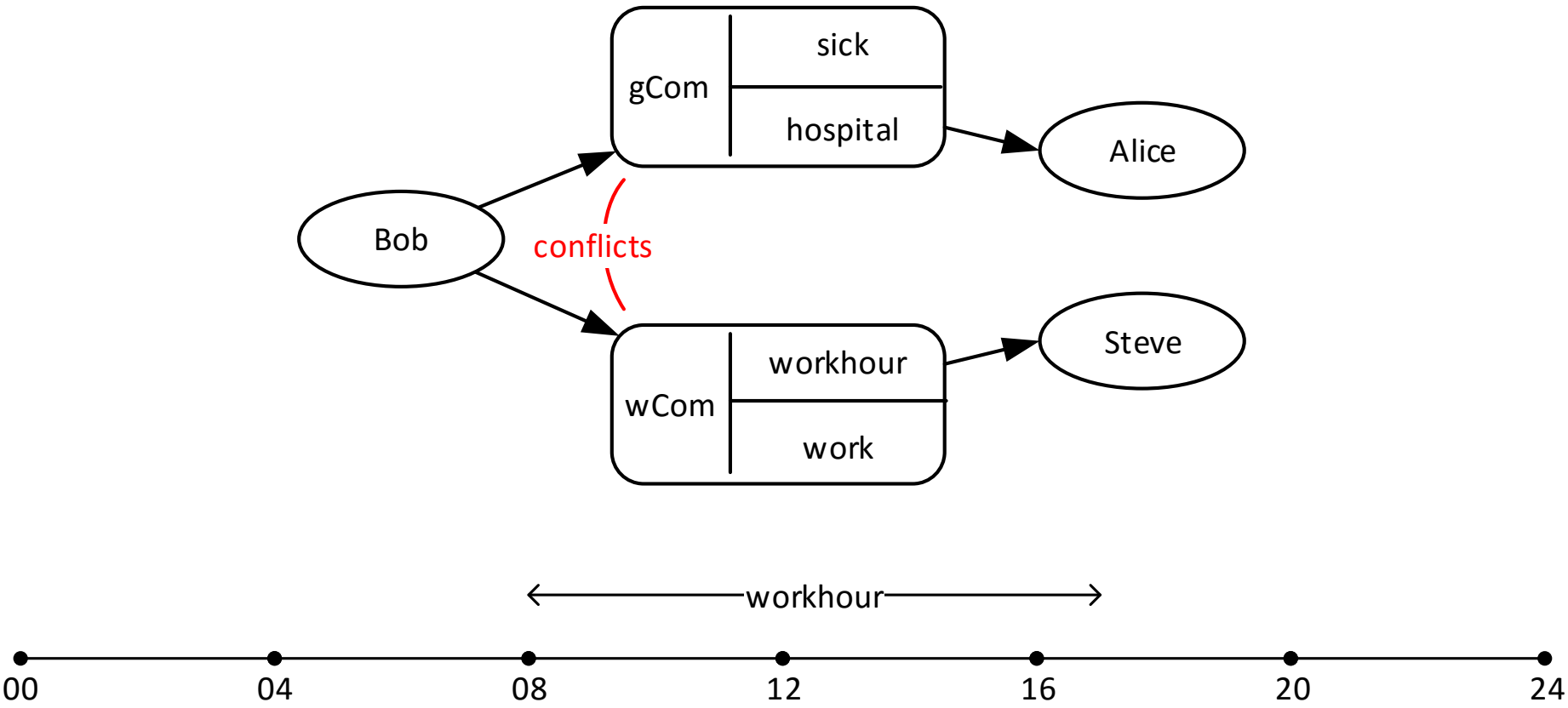
Outline

Commitment Schema

- For effective interactions, agents enter into commitments



Commitment Instances



Problem

How to detect and resolve runtime conflicts?

Contributions

- Representations for commitments and dominance
- Decision procedures for determining compliance of actions with commitment instances
  - apply ASP to identify nondominated commitments
  - apply Alechina et al.'s [2013] techniques to determine compliance of agent actions with nondominated commitment instances

Outline

An Example Commitment Schema

Definition of the gCom schema

1 $schema(gCom)$
2 $guardianR(C, G, T) \rightarrow dbt(gCom, G, T)$
3 $guardianR(C, G, T) \rightarrow crd(gCom, C, T)$
4 $[sick(C, T) \wedge crd(gCom, C, T)] \rightarrow ant(gCom, T)$
5 $[bring(G, C, Ped, T) \wedge pedR(C, Ped, T) \wedge dbt(gCom, G, T) \wedge crd(gCom, C, T)]$
$\rightarrow con(gCom, T)$
6 $dDuration(gCom, 3)$

Schema Predicates

Schema predicates	
$schema(S)$	$ant(S, T)$
$dbt(S, P, T)$	$con(S, T)$
$crd(S, P, T)$	$dDuration(S, T)$
Instance predicates	
$sInst(Si)$	$violated(Si, T)$
$isInstOf(Si, S)$	$satisfied(Si, T)$
$dbtI(Si, P)$	$becomesDetached(Si, T)$
$crdI(Si, P)$	$becomesViolated(Si, T)$
$antI(Si, T)$	$becomesSat(Si, T)$
$conI(Si, T)$	$conflicting(Si_1, Si_2, T)$
$dDurationI(Si, T)$	$dominates(Si_1, Si_2, T)$
$detached(Si, T)$	$dominated(Si, T)$
$sameDbtI$	$conINotHold(Si, T_1, T)$
$sameCrdI$	

Satisfaction and Violation

Satisfaction

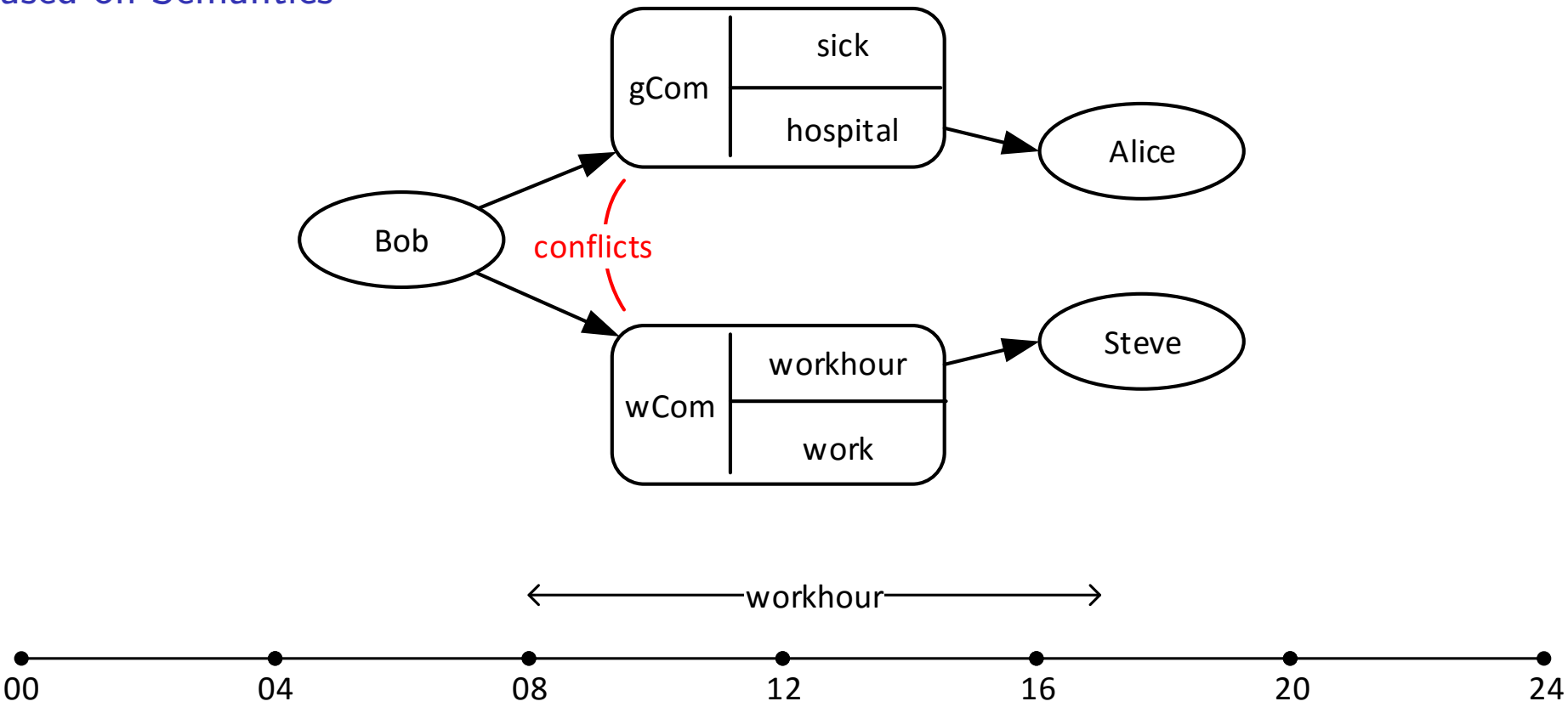
1 $[detached(Si, T') \wedge conI(Si, T) \wedge difference(T, T', 1)] \rightarrow becomesSat(Si, T)$

Violation

1 $[becomesDetached(Si, T_1) \wedge dDurationI(Si, T_2)$
$\wedge addition(T_1, T_2, T) \wedge conINotHold(Si, T_1, T)$
$\wedge \neg dominated(Si, T)] \rightarrow becomesViolated(Si, T)$

Conflict between Commitment Instances

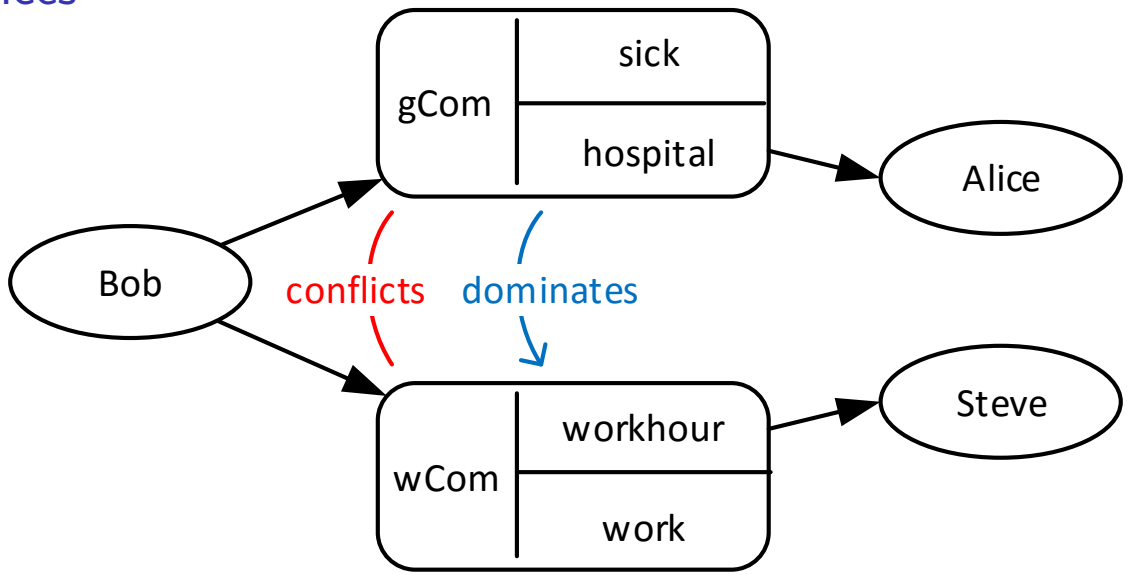
Based on Semantics

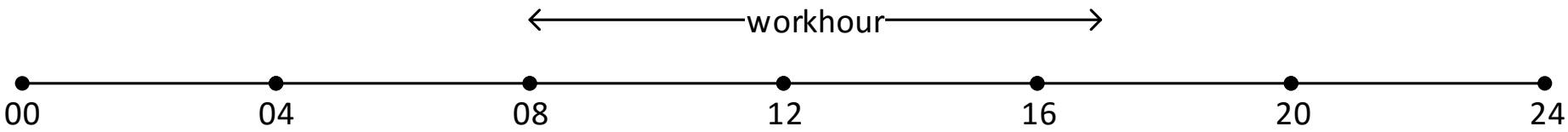


1 $[detached(Si_1, T) \wedge detached(Si_2, T) \wedge \neg(conI(Si_1, T) \leftrightarrow conI(Si_2, T))]$
$\rightarrow conflicting(Si_1, Si_2, T)$

Dominance between Commitment Instances

Based on Preferences

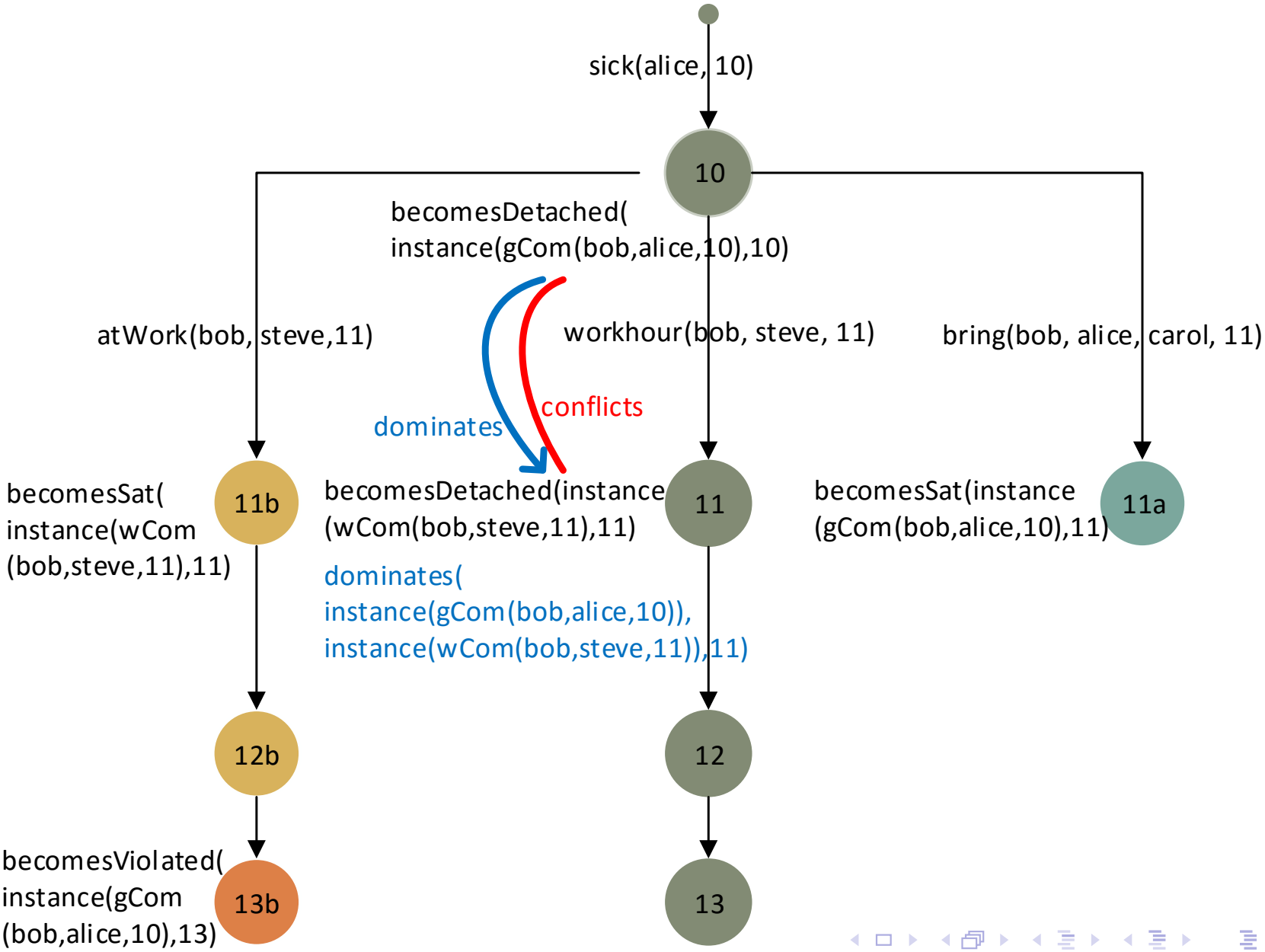




1 $[isInstOf(Si_1, gCom) \wedge isInstOf(Si_2, wCom) \wedge detached(Si_1, T) \wedge detached(Si_2, T) \wedge sameDbtI(Si_1, Si_2)] \rightarrow dominates(Si_1, Si_2, T)$

Outline

Possible Enactments



Conflict-detection, Formally

1 $[atLocation(P, L_1, T) \wedge (L_1 \neq L_2)] \rightarrow \neg atLocation(P, L_2, T)$

2 $[isInstOf(Si, gCom) \wedge conI(Si, T) \wedge dbtI(Si, G)] \rightarrow atLocation(G, hospitalLoc, T)$

3 $[employerR(Ee, Er, T) \wedge \neg atLocation(Ee, companyLoc, T)] \rightarrow \neg atWork(Ee, Er, T)$

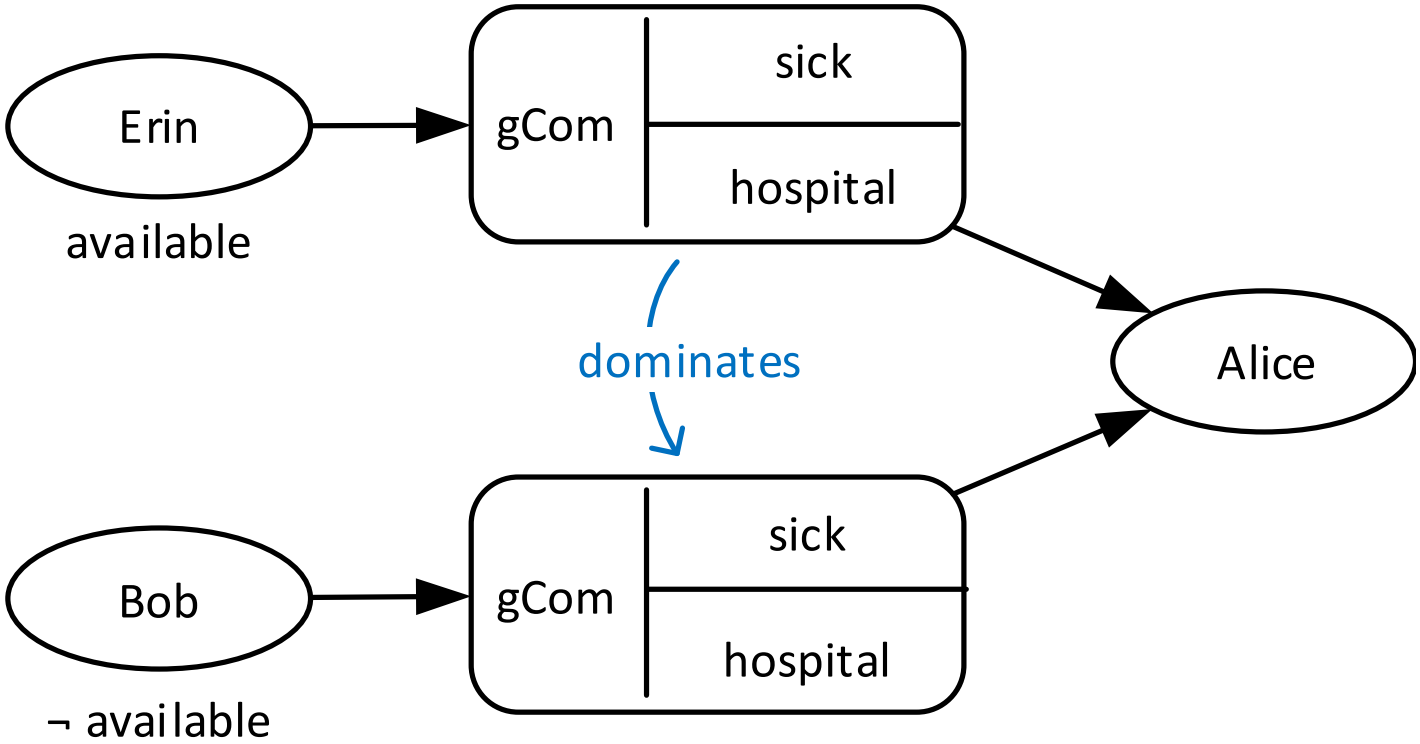
4 $[isInstOf(Si, wCom) \wedge dbtI(Si, Ee) \wedge crdI(Si, Er) \wedge \neg atWork(Ee, Er, T)] \rightarrow \neg conI(Si, T)$

5 $[isInstOf(Si, wCom) \wedge conI(Si, T) \wedge dbtI(Si, Ee) \wedge crdI(Si, Er)] \rightarrow atLocation(Ee, company, T)$

6 $[guardianR(C, G, T) \wedge pedR(C, Ped, T) \wedge \neg atLocation(G, hospitalLoc, T)] \rightarrow \neg bring(G, C, Ped, T)$

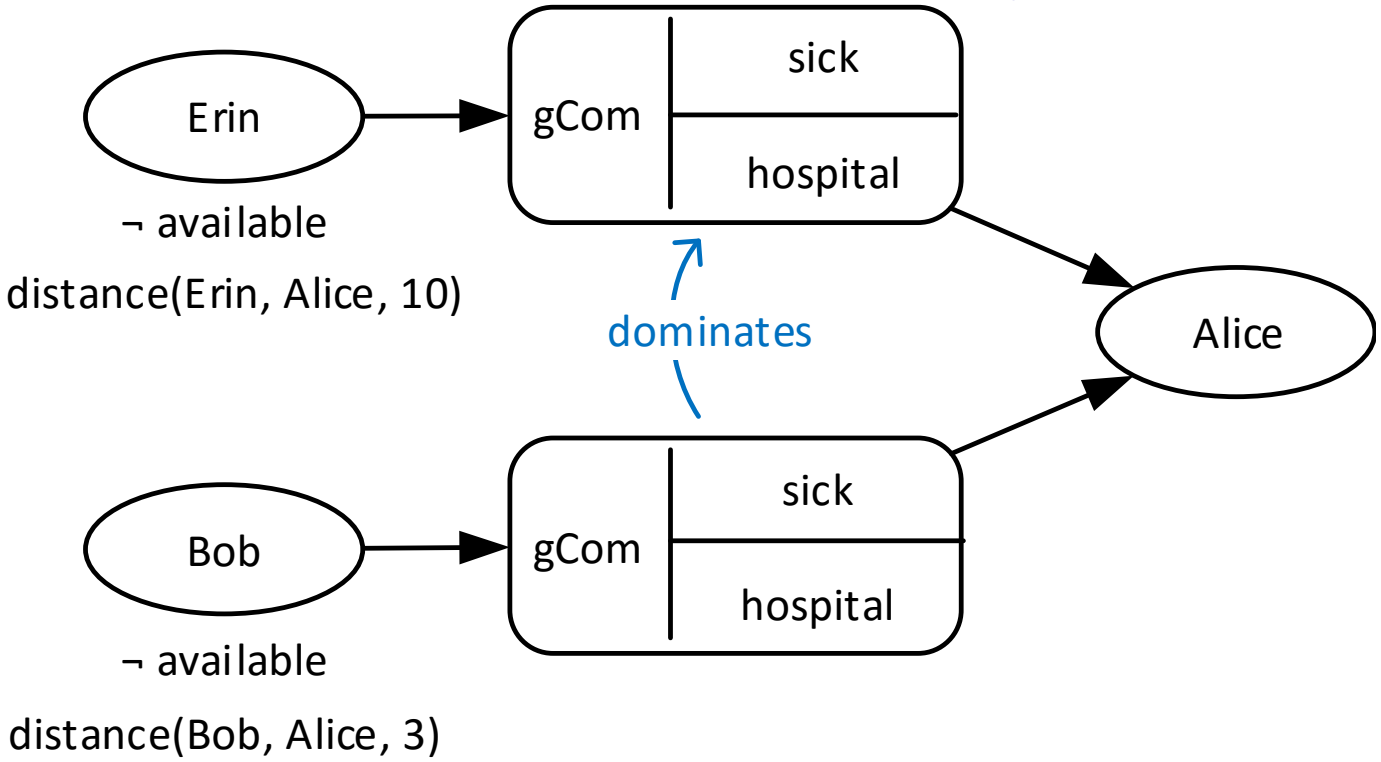
7 $[isInstOf(Si, gCom) \wedge dbtI(Si, G) \wedge crdI(Si, C) \wedge pedR(C, Ped, T) \wedge \neg bring(G, C, Ped, T)] \rightarrow \neg conI(Si, T)$

Example Basis for Dominance: Availability



1 $[isInstOf(Si_1, gCom) \wedge isInstOf(Si_2, gCom) \wedge detached(Si_1, T) \wedge detached(Si_2, T)$
$\wedge sameCrdl(Si_1, Si_2) \wedge dbtl(Si_1, G_1) \wedge dbtl(Si_2, G_2) \wedge (G_1 \neq G_2)$
$\wedge available(G_1, T) \wedge \neg available(G_2, T)] \rightarrow dominates(Si_1, Si_2, T)$

Example Basis for Dominance: Proximity



1 $[isInstOf(Si_1, gCom) \wedge isInstOf(Si_2, gCom) \wedge detached(Si_1, T) \wedge detached(Si_2, T)$
$\wedge sameCrdl(Si_1, Si_2) \wedge crdl(Si_1, C) \wedge dbtl(Si_1, G_1) \wedge dbtl(Si_2, G_2) \wedge (G_1 \neq G_2)$
$\wedge \neg available(G_1, T) \wedge \neg available(G_2, T) \wedge dist(G_1, C, D_1, T) \wedge dist(G_2, C, D_2, T)$
$\wedge (D_1 \leq D_2)] \rightarrow dominates(Si_1, Si_2, T)$

Outline

Conclusion and Future Directions

Coco provides

- Flexible formalization of instances
- Reasoning about conflicts between commitments and dominance at runtime

Subsequent work

- Human-subject study to elicit normative requirements
- Participants using Coco-based methodology
  - produce specifications with higher coverage and correctness;
  - expend equal time; and
  - feel Coco-based methodology is easy

Future directions

- Dealing with cycles in dominance
- Applying priority to a set of dominance relations

Thank you

Multiagent Systems and Service-Oriented Computing Laboratory, NCSU
`http://research.csc.ncsu.edu/mas/`