Nirbhay Sharma (B19CSE114)

Cryptography - Assignment-5

Oue-1

output considering first 16 bits (4 chars of hash) since $\frac{256}{64}=4$ so each char can be considered as 4 bits

- 1. x corresponding to emailed(sharma.59@iitj.ac.in) log(iterations) = 16
- 2. finding (x, y) log(iterations) = 9

```
iterations: 10922 -----
len: 1
m06awl3df@iitj.ac.in
iterations: 2758 -----
len: 1
l.hg0z97b@iitj.ac.in
                                        iterations: 523 -----
iterations: 50298 -----
                                        ('Omv14jsiq@iitj.ac.in', 'O13.v9ueo@iitj.ac.in')
len: 1
                                        iterations: 104 -----
di6y9j5v4@iitj.ac.in
                                        ('vqszyxgib@iitj.ac.in', '6yj5gnb7f@iitj.ac.in')
iterations: 30989 -----
                                        iterations: 161 -----
len: 1
                                        ('43motgvyn@iitj.ac.in', 'xg9pqi31t@iitj.ac.in')
qcmo09by3@iitj.ac.in
                                        iterations: 367 -----
iterations: 44514 ----
                                        ('3txy25m1d@iitj.ac.in', 'ehv427nlz@iitj.ac.in')
len: 1
                                        iterations: 345 -----
cxeoj8w45@iitj.ac.in
                                        len: 1
                                        ('52izdcspr@iitj.ac.in', '7qfgjz4d.@iitj.ac.in')
55417
                                        333
```

it is clearly visible that finding an x which has 16 bit common with (sharma.59@iitj.ac.in) is hard as compared to finding two values (x,y) such that their first 16 bits are matching as the first case is the case of second pre image matching and second case is the case for collision and it is always easy or less expensive (computationally) to find collisions than to find second pre image, and hence first case took many iterations to find out the answer.

output considering first 8 bits (2 chars of hash)

- 1. x corresponding to emailed(sharma.59@iitj.ac.in) log(iterations) = 8
- 2. finding (x,y) log(iterations) = 5

```
iterations: 55 --
len: 1
c1awylu7z@iitj.ac.in
iterations: 290 -
len: 1
                                         iterations: 39 -----
ve630l1ub@iitj.ac.in
                                         ('yg.8lv93w@iitj.ac.in', 'y3xwv6kco@iitj.ac.in')
                                         iterations: 43 -----
iterations: 46
                                         ('vji7k3g80@iitj.ac.in', '3caxep5lf@iitj.ac.in')
len: 1
zkgxuv0w.@iitj.ac.in
                                         iterations: 12 -----
                                         len: 1
                                         ('denq3hz8.@iitj.ac.in', '1a.Ofmjtc@iitj.ac.in')
iterations: 18 --
                                         iterations: 33 -----
len: 1
                                         len: 1
                                         ('yed4ub0xc@iitj.ac.in', 'enwyua94z@iitj.ac.in')
wi6m7t2ds@iitj.ac.in
                                         iterations: 16 -----
                                         len: 1
                                         ('uhcv196s4@iitj.ac.in', 'g0.7z681h@iitj.ac.in')
210
```

it is clearly visible that finding an x which has 8 bit common with (sharma.59@iitj.ac.in) is hard as compared to finding two values (x,y) such that their first 8 bits are matching as the first case is the case of second pre image matching and second case is the case for collision and it is always easy or less expensive (computationally) to find collisions than to find second pre image, and hence first case took many iterations to find out the answer.

Oue-2

output of the code for RSA encryption and decryption

```
ma406@LAPTOP-N4BIN1J0:/mnt/d/coding assn sem6/cryptography-assn/assn5$ python3 q2_b19cse114.py
747901531573246085358591762583193586803063349028169408119730032917161126193217597903852619309992517475904996373826671808595213
3874349109354048353
      721829975462614001345278836477287822723444168173367114073620324640610182053001373316711411450157607842189131007438699349202268
5685478866994031023
private d: 10797154883679353051932038494943003358135277860405235893906167596348491411926963087912605238689426214925894347137837167
  19496082147717050132048318824655035195149
                      53985774418396765259660192474715016790676389302026179469530837981742457059634815439563026\underline{1934471310746294717356891858388}
0xac16c8e0 0x2f8b91c89 0x47ba2f320 0x21ac75e20 0x2d52e8000 0x1ffd869e1 0x609fdb0d9 0x2000000 0xeac8fd83 0x2d52e8000 0x1ffd869e1 0x
0x39517623d 0x1ffd869e1 0x2000000 0x2f8b91c89 0x4aedcc023 0x2000000 0x1ffd869e1 0x2000000 0x2b2fb2f87 0x3ec5f782f 0x3ec5f782f 0x2ec5f782f 0x2ec5f782f 0x2ec5f782f 0x2ec5f782f 0x3ec5f782f 0x2ec5f782f 0x2ec5f782f 0x3ec5f782f 0x3ec5f782f 0x3ec5f782f 0x3ec5f782f 0x3ec5f782f 0x3ec5f782f 0x3ec5f782f 0x2ec5f782f 0x3ec5f782f 0x3ec5f782f 0x2ec5f782f 0x2ec5f782f 0x2ec5f782f 0x2ec5f782f 0x2ec5f782f 0x3ec5f782f 0x3ec5f782f 0x3ec5f782f 0x2ec5f782f 0
 00 0x2d52e8000 0x272736815 0x2000000 0x4aedcc023 0x4e3e6b400 0x51acd0565 0x2540be400 0x2f8b91c89 0x272736815 0x4aedcc023 0x2000000
9e1 0x4e3e6b400 0x2000000 0x7b908fe9 0x7b908fe9 0xf9461400 0x2000000 0x84435aa0 0x3ec5f782f 0x2540be400 0x2d52e8000 0x41a700000 0x:
0x47ba2f320 0x2000000 0x1ffd869e1 0x3bff052e0 0x2540be400 0x2000000 0x2d52e8000 0x272736815 0x2000000 0x36bc9ac00 0x2f8b91c89 0x3
0x272736815 0x4aedcc023 0x2000000 0x4e3e6b400 0x3ec5f782f 0x2000000 0x2540be400 0x3ec5f782f 0x2000000 0x272736815 0x5cb278000 0x27
 x47ba2f320 0x236d590d3 0x2f8b91c89 0x4aedcc023 0x272736815 0x2000000 0x1ffd869e1 0x3bff052e0 0x2540be400 0x2000000 0x609fdb0d9 0x3
0x2b2fb2f87 0x1ffd869e1
```

Que-3

1. proving that a cyclic group of n elements has $\phi(n)$ generators

let C_n be a cyclic group of order n

```
so C_n can be written as = C_n = \langle a 
angle for some a \in C_n elemetrs of C_n can be written as \{g^k: g \in C_n, 0 \le k < n\}
```

so if we assume that some more generator exists which is g^k , from this we can follow that g^k can also generate C_n iff $k \perp n$ i.e gcd(k,n)=1 and from that we can follow that for all k that are coprime to n can be a generator and hence number of such k are $\phi(n)$

output of the code for sophie_germain_prime and diffie_hilmann_exchange

- according to definition sophie germain prime is a prime number p such that 2p+1 is also a prime (safe prime)

```
sharma406@LAPTOP-N4BIN1J0:/mnt/d/coding assn sem6/cryptography-assn/assn5$ python3 q3_b19cse114.py
sphie_germain_prime: 13344560613885683016956070482243112356334510034328720352835313621629930682301702354555
992166028532143590018205218612008422328107220067812906311886850798389
found generator: 929623431195408500820249953273817110352045311572590042129422185637614246640536714981862725
9669137018372438849005028247727386827350669215881710728639544328
key_at_alice end: 71094348791151363024389554286420996798449
key_at_bob end: 71094348791151363024389554286420996798449
```

• in the above image the generator is any random element from the group since Z_p^* is cyclic and p is prime (sophie germain prime) so every element of group Z_p^* is a generator (from the property) and hence any random element from $1 \to p-1$ is a generator

some more sophie germain primes are shown below -:

```
{'iter': 127, 'sgp':
2833841992718854119840701991506063663984854959536491, '2p+1':
\overline{2512372373502}78467291566823441946104168113405038325677726152289940110361041592446365834573\underline{2287629361614161}
5667683985437708239681403983012127327969709919072983}
{\daggeright 'iter': 126, 'sgp':
5260997875933215479312967870674178391549464815852293, '2p+1':
2517494186961143313348101855088267306182723164257657050373141657046853181654819108783378389997333889359
0521995751866430958625935741348356783098929631704587, 'generator':
2703369036523464116359060707183491969337103713835334
{'iter': 93, 'sgp':
0018205218612008422328107220067812906311886850798389, '2p+1':
2668912122777136603391214096448622471266902006865744070567062724325986136460340470911198433205706428718
0036410437224016844656214440135625812623773701596779, 'generator':
9296234311954085008202499532738171103520453115725900421294221856376142466405367149818627259669137018372
438849005028247727386827350669215881710728639544328}
```

codes for question 1, 2, and 3

code-1

```
from Crypto.Hash import SHA256
import random,time,os

def find_hash(text):
    new_hash = SHA256.new()
    new_hash.update(bytes(text,'utf-8'))
```

```
return new hash.hexdigest()
def condition(text1,text2,first bits=4):
   return (text1[:first bits] == text2[:first bits])
def generate random(num):
   text = list('abcdefghijklmnopqrstuvwxyz0123456789.')
   random bits = ''.join(random.sample(text,num))
   return random bits
def write to file(filename,iterations,array):
   with open(filename, 'a') as f:
      string = f'iterations: {iterations} -----\nlen:
{len(array)}\n{" ".join(map(lambda x:str(x),array))}\n\n'
      f.write(string)
def generate pairs(size1,size2):
   p1 = generate random(size1) + '@iitj.ac.in'
   p2 = generate random(size2) + "@iitj.ac.in"
   return p1,p2
h1 = find hash('sharma.59@iitj.ac.in')
print(h1)
def n calls(h1,match value):
   t1 = time.perf counter();
   l = []
   iterations = 0
   max iterations = int(1e6)
   while (True):
      generated value = generate_random(9) + '@iitj.ac.in'
      h2 = find_hash(generated_value)
      iterations += 1
      if (condition(h1,h2,match_value)):
         l.append(generated_value)
         break;
      if (iterations > max_iterations):
         break;
   if (len(l) > 0):
      write_to_file('output1.txt',iterations,l)
   else:
      print('len = 0')
   t2 = time.perf_counter()
   print(f'completed in {(t2-t1)/60:.2f} Mins')
   return iterations;
def find pairs(values match):
   t1 = time.perf_counter();
```

```
l = []
   iterations = 0
   max iterations = int(1e6)
   values map = {}
   while (True):
      generated value1 = generate_random(9) + "@iitj.ac.in"
      h1 = find hash(generated value1)
      iterations += 1
      if (h1[:values match] in values map):
         l.append((values map[h1[:values match]],generated value1))
      else:
         values map[h1[:values match]] = generated value1;
      if (iterations > max_iterations):
         break:
   # print(len(l))
   # print(l)
   if (len(l) > 0):
      write_to_file('output2.txt',iterations,l)
   else:
      print('len = 0')
   t2 = time.perf counter()
   print(f'completed in {(t2-t1)/60:.2f} Mins')
   return iterations;
# check iterations(h1)
filenames = ['output1.txt','output2.txt']
for filename in filenames:
   if (os.path.exists(filename)):
      os.remove(filename)
output1 iterations = 0
output2_iterations = 0
n = 10;
for i in range(n):
   output1_iterations += n_calls(h1,4) / n
   output2 iterations += find pairs(4) / n
with open('output1.txt','a') as f:
   f.write(str(int(output1 iterations)))
with open('output2.txt','a') as f:
   f.write(str(int(output2_iterations)))
exit(0)
```

code-2

```
import random
import sys
sys.setrecursionlimit(1030)
def extended eucledian(a,b):
    if (a\%b == 0): return (b,0,1)
    val = [[a,1,1,0],[b,a//b,0,1]]
    while val[-1][0] != 1:
        new val = []
        new val.append(val[-2][0] % val[-1][0])
        if (\text{new val}[0] == 0): return (\text{val}[-1][0], \text{val}[-1][-2], \text{val}[-1][-1])
        new val.append(val[-1][0] // new val[0])
        new val.append(val[-2][-2] - val[-1][1] * val[-1][-2])
        new val.append(val[-2][-1] - val[-1][1] * val[-1][-1])
        val.append(new val)
    return (val[-2][1],val[-1][-2],val[-1][-1]);
class rsa algorithm(object):
    def init (self):
        # self.p = 10888869450418352160768000001
        # self.g = 265252859812191058636308479999999
        self.p = self.generate prime();
        self.q = self.generate prime();
        # self.p = 155;
        \# self.q = 21711;
        print('P: ',self.p)
        print('Q: ',self.q)
        self.phi n = (self.p-1)*(self.q-1)
        self.n = self.p * self.q
    def find e(self):
        for i in range(3, 100, 2):
            gcd,x,y = extended_eucledian(self.phi_n,i)
            if (gcd == 1):
                return (i,y)
    def rsa encryption(self,msg):
        self.e,self. d = self.find e();
        # print('inside d:',self. d)
        if (self. d < 0):
            print('d < 0, making it positive !!!')</pre>
            self. d = self.phi n + self. d
        encoded string = []
        encoded_msg = bytes(msg,'utf-8')
        for char in encoded_msg:
encoded_string.append(self.modulo_exponentiation(char,self.e,self.n))
        return ''.join(map(lambda x: bytearray(hex(x), 'utf-8').decode() +
   ,encoded string))[:-1]
    def rsa_decryption(self,cipher):
        decoded msg = []
        for i in cipher.split(' '):
            cipher msg = i.strip();
```

```
cipher msg = int(cipher msg.encode(),16)
            # print(f'cipher msg: {cipher msg}')
decoded msg.append(self.modulo exponentiation(cipher msg,self. d,self.n))
        # print(decoded msg)
        # return
        return ''.join(map(lambda x:chr(x),decoded msg))
    def square multiply(self,bas,exp):
        if (exp == 0): return 1;
        if (exp == 1): return bas;
        if (exp & 1):
            return bas * self.square_multiply(bas * bas, (exp-1)//2)
        else:
            return self.square multiply(bas * bas,exp//2)
    def modulo exponentiation(self,bas,exp,N):
        if (exp == 0):
            return 1;
        if (exp == 1):
            return bas % N;
        t = self.modulo exponentiation(bas, exp // 2,N);
        t = (t * t) % N;
        if (exp % 2 == 0):
            return t;
        else:
            return ((bas % N) * t) % N;
    def generate random 512(self):
        return int("1"+ ''.join([random.choice(["0","1"]) for in
range(510)]) + "1",2)
    def find_s_t(self,num):
        new num = num - 1;
        s = 0;
        while (new_num & 1 == 0):
            s += 1;
            new num >>= 1;
        return (s,new_num)
    def miller rabin test(self,num,s,t):
        new num = num - 1;
        aa = random.randint(2,new num)
        b i = self.modulo exponentiation(aa,t,num);
        if (b_i == 1 \text{ or } b_i == num - 1):
            return 1;
        pow_2 = 0;
        while (True):
            pow_2 += 1
            if (pow 2 == s+1):
                break;
            b_i = self.modulo_exponentiation(b_i,2,num);
            # if (b i == 1): return 0;
```

```
if (b i == num - 1): return 1;
        return 0:
    def check prime(self,num):
        s,t = self.find s t(num);
        for i in range (100):
            vote = self.miller rabin test(num,s,t)
            if (vote == 0):
                return False:
        return True;
    def generate prime(self):
        a = False;
        final prime = None;
        # iterations = 0;
        while (not a):
            # iterations+=1
            final prime = self.generate random 512()
            a = self. check prime(final prime)
        # print(iterations)
        return final prime;
    def fast modulo exponentiation(self,base,power,N):
        dp = [-1]
original_text = "Nirbhay Sharma is a good boy !!! and he studies at IIT
Jodhpur and he likes to do exercise and yoga";
rsa = rsa_algorithm();
# print(rsa.generate_prime())
# print(rsa.generate prime())
# exit(0)
enc = rsa.rsa encryption(original text);
print('private d:',rsa._rsa_algorithm__d)
print('public n: ',rsa.n)
print('public e:',rsa.e)
print('#'*25 + ' original_text ' + '#'*25)
print(original_text)
print('#'*25 + ' encrypted text ' + '#'*25)
print(enc)
dec = rsa.rsa_decryption(enc)
```

```
print('#'*25 + ' decrypted text ' + '#'*25)
print(dec)
```

code-3

```
import random, ison, os
def modulo exponentiation(bas,exp,N):
    if (exp == 0):
        return 1:
    if (exp == 1):
       return bas % N;
    t = modulo exponentiation(bas, exp // 2,N);
    t = (t * t) % N;
    if (exp % 2 == 0):
        return t;
    else:
        return ((bas % N) * t) % N;
def square multiply(bas,exp):
    if (exp == 0): return 1;
    if (exp == 1): return bas;
    if (exp & 1):
        return bas * square multiply(bas * bas, (exp-1)//2)
    else:
        return square multiply(bas * bas,exp//2)
def write to file(filename, dictionary):
    with open(filename, 'a') as f:
        f.write(str(dictionary) + "\n")
def find generator(prime p):
    while True:
        random a = random.randint(2,prime_p-1)
        rem = modulo exponentiation(random a,prime p-1,prime p)
        if (rem == 1): return random a;
class Prime_generator_machine(object):
    def __init__(self,bits_in_prime):
        self.bip = bits_in_prime;
    def modulo exponentiation(self,bas,exp,N):
        if (exp == 0):
            return 1;
        if (exp == 1):
            return bas % N;
        t = self.modulo exponentiation(bas, exp // 2,N);
        t = (t * t) % N;
        if (exp % 2 == 0):
            return t;
        else:
```

```
return ((bas % N) * t) % N;
    def generate random bits(self):
        return int("1"+ ''.join([random.choice(["0","1"]) for in
range(self.bip-2)]) + "1",2)
    def find s t(self,num):
        new num = num - 1;
        s = 0;
        while (new num & 1 == 0):
            s += 1;
            new num >>= 1;
        return (s, new num)
    def miller_rabin_test(self,num,s,t):
        new num = num - 1;
        aa = random.randint(2,new num)
        b i = self.modulo exponentiation(aa,t,num);
        if (b i == 1 \text{ or } b i == num - 1):
            return 1;
        pow 2 = 0;
        while (True):
            pow 2 += 1
            if (pow_2 == s+1):
                break:
            b i = self.modulo exponentiation(b i,2,num);
            # if (b i == 1): return 0;
            if (b i == num - 1): return 1;
        return 0;
    def _check_prime(self, num):
        s,t = self.find s t(num);
        for i in range(100):
            vote = self.miller_rabin_test(num,s,t)
            if (vote == 0):
                return False;
        return True;
    def generate_prime(self):
        a = False;
        final prime = None;
        # iterations = 0;
        while (not a):
            # iterations+=1
            final_prime = self.generate_random_bits()
            a = self._check_prime(final_prime)
        # print(iterations)
        return final_prime;
pgm = Prime_generator_machine(512)
def generate sophie germain prime(pgm,thresh):
```

```
sgp = None;
    iterations = 0;
    p dash = 0
    while True:
        p dash = pgm.generate prime()
        # print(f'p dash prime test: {pgm. check prime(p dash)}')
        iterations += 1
        print(f'iterations: {iterations}',end='\r')
        sophie germain prime = 2*p dash + 1
        if (pgm. check prime(sophie germain prime)):
            sgp = sophie germain prime
            break:
        if (iterations == thresh):
            break:
        # print(f'sophie germain prime test:
{pgm. check prime(sophie germain prime)}')
        # print('p dash:',p dash)
        # print('generated sophie germain prime:',sophie germain prime)
    # print('p dash:',p dash)
    return ((p dash,sqp,iterations) if sqp is not None else
("None", "None", iterations));
thresh = 1000
sophie germain prime, condition sophie germain, iterations =
generate sophie germain prime(pgm,thresh)
if condition sophie germain == "None":
    print(f'cannot find sophie germain prime till {thresh} iterations !!!
Try running again')
    exit(0)
print('sphie_germain_prime: ',sophie_germain_prime)
generator = find generator(sophie germain prime)
write_to_file('sophie-germain-prime.txt',
{"iter":iterations, "sgp":sophie_germain_prime, "2p+1":condition_sophie_germ
ain, "generator":generator})
# for a prime order group every element is a generator
print('found generator: ',generator)
# exit(0)
class Alice(object):
    def init (self,g):
        self.a = random.randint(3,10)
        self.g = g
    def _generate_key(self):
        return square_multiply(self.g,self.a)
    def generate final key(self,bob g random):
        return square_multiply(bob_g_random,self.a)
```

```
class Bob(object):
   def init (self,g):
       self.b = random.randint(3,10)
        self.g = g
   def generate key(self):
        return square multiply(self.g,self.b)
   def generate final key(self,alice g random):
       return square multiply(alice g random, self.b)
def diffie hilmann exchange(g):
   alice = Alice(g)
   bob = Bob(g)
   alice generated key = alice. generate key()
   bob_generated_key = bob._generate_key()
   final key at alice = alice. generate final key(bob generated key)
   final key at bob = bob. generate final key(alice generated key)
   print("key_at_alice end:",final_key_at_alice)
   print('key at bob end:',final key at bob)
   # print(final_key_at_alice==final_key_at_bob)
diffie hilmann exchange(23)
```