

## The AMS Estimator of $F_k$

So far, we have seen estimators for  $F_0, F_2$ . Today, we explore a general estimator for ~~any~~  $F_k$  where  $k \geq 2$ .

The motivation is we pick an element at random from the stream. We count the number of instances of the picked element from that point onwards. The basic estimator is  $m(r^k - (r-1)^k)$  where  $r$  is the number of instances and  $m$  the length of the stream.

However, as we do not know  $m$  beforehand, we have to pick one on the fly as we ~~see~~ saw in the week 1 problem set.

AMS Estimator on stream  $\sigma = (s_1, \dots, s_m)$  from  $\{1, \dots, n\}$ .

- Initialize: Set ~~random~~  $r = 0, a = 0$ .
- Process  $s_j$ :  
~~random~~
  - With prob.  $\frac{1}{j}$ , set  $a \leftarrow s_j, r \leftarrow 0$
  - If  $s_j = a$ ,  $r \leftarrow r + 1$
- Output:  $m(r^k - (r-1)^k)$ .

### Analysis

Let  $A, R$  be the r.v.'s for the values of  $a, r$  at the end of alg., respectively.

Let  $X$  denote the output.

We begin, by considering probabilities conditional on  $A = i$ .

If  $A=i$ , then  $R$  is equally likely to be any of the values  $\{1, \dots, f_i\}$  depending on which of the  $f_i$  occurrences of  $i$  was picked by the alg.

$$Pr(R = l | A = i) = \begin{cases} \frac{1}{f_i} & 1 \leq l \leq f_i \\ 0 & l > f_i \end{cases}.$$

$$\begin{aligned} \Rightarrow E(X | A=i) &= E(m(R^k - (R-1)^k) | A=i) \\ &= \sum_{l=1}^{f_i} \frac{1}{f_i} m(l^k - (l-1)^k) \\ &= \frac{m}{f_i} (f_i^k - 0^k) \quad (\text{by telescoping sums}) \\ &= m f_i^{k-1} \end{aligned}$$

Therefore, the general expectation is

$$\begin{aligned} E(X) &= \sum_{i=1}^n Pr(A=i) E(X | A=i) \\ &= \sum_{i=1}^n \frac{f_i}{m} m f_i^{k-1} = \sum_i f_i^k = F_k. \quad \checkmark \end{aligned}$$

So, the  $X$  produced is indeed an unbiased estimator for  $F_k$ .

To bound the variance,

$$\begin{aligned} Var(X) &= E(X^2) - E(X)^2 \leq E(X^2) = \sum_{i=1}^n \sum_{l=1}^{f_i} \frac{1}{f_i} m^2 (l^k - (l-1)^k)^2 \\ &= m \sum_{i=1}^n \sum_{l=1}^{f_i} (l^k - (l-1)^k)^2 \end{aligned}$$

Recall the mean value thm:  $\exists y \in [x-1, x]$  s.t.

$$f(x) - f(x-1) = \cancel{f'(y)} f'(y)$$

for continuous differentiable f's  $f$ . Apply this with  $f(x) = x^k$

$$\Rightarrow x^k - (x-1)^k = k y^{k-1} \leq k x^{k-1} \quad (\text{for } x \geq 1).$$

Applying this to the previous bnd.,

$$\begin{aligned} \text{Var}(X) &\leq m \sum_{i=1}^n \sum_{l=1}^{f_i} k l^{k-1} (l^k - (l-1)^k) \\ &= m \sum_{i=1}^n k f_i^{k-1} \sum_{l=1}^{f_i} (l^k - (l-1)^k) \\ &= mk \sum_{i=1}^n f_i^{k-1} \cdot f_i^k \quad \text{Again by telescoping sums} \\ &= k F_1 F_{2k-1}. \end{aligned}$$

If we want to apply the median-of-means technique from Lec 2, it would be useful if  $\text{Var}(X)$  can be expressed as a multiple of  $E(X^2)^2$  (or at least a bound of  $\text{Var}(X)$  can be expressed as such).

Lemma For  $x_1, \dots, x_n \geq 0$  reals and  $k \geq 1$  real. Then

$$\left( \sum_i x_i \right) \left( \sum_i x_i^{2k-1} \right) \leq n^{1-1/k} \left( \sum_i x_i^k \right)^2.$$

Cor  $F_1 \cdot F_{2k-1} \leq n^{1-1/k} F_k^2.$

### PP of Lemma

Let  ~~$M$~~   $M = \max_i x_i$ . Then

$$M^{k-1} = (M^k)^{k-1/k} \leq \left( \sum_i x_i^k \right)^{(k-1)/k}$$

As the fn  $x \mapsto x^k$  is convex,

$$\left( \frac{1}{n} \sum_i x_i \right)^k \leq \frac{1}{n} \sum_i x_i^k$$

$$\frac{1}{n} \sum_i x_i^k \leq \left( \frac{1}{n} \sum_i x_i^k \right)^{1/k}$$

$$\begin{aligned} \Rightarrow \left( \sum_i x_i \right) \left( \sum_i x_i^{2k-1} \right) &\leq \left( \sum_i x_i \right) \left( M^{k-1} \sum_i x_i^k \right) \\ &\leq \left( \sum_i x_i \right) \left( \sum_i x_i^k \right)^{(k-1)/k} \left( \sum_i x_i^k \right) \\ &\leq n \left( \frac{1}{n} \sum_i x_i^k \right)^{1/k} \left( \sum_i x_i \right)^{2k-1/k} \\ &= n^{1-1/k} \left( \sum_i x_i^k \right)^2. \end{aligned}$$

$$\Rightarrow E(X) = F_k \quad \text{and} \quad \text{Var}(X) \leq k n^{1-1/k} F_k^2.$$

We can apply median-of-means to get a  $(\epsilon, \delta)$ -approximation using

$$O\left( \frac{1}{\epsilon^2} \log\left(\frac{1}{\delta}\right) \frac{\text{Var}(X)}{E(X)^2} (\log m + \log n) \right)$$

$$= O\left( \frac{1}{\epsilon^2} \log\left(\frac{1}{\delta}\right) k n^{1-1/k} (\log m + \log n) \right)$$

Note, as  $n \rightarrow \infty$ , this isn't that great cause each  $f_i$  can be calculated precisely with  $O(n \log m)$  space. The alg. is also not efficient for calculating  $F_2$  & rather its just a good exercise in streaming algorithm analysis.

Using the Johnson-Lindenstrauss Lemma to calculate  $F_2$ .

The goal of calculating  $F_2$  can be thought of as calculating  $\|f\|^2$  where  $f = (f_1, \dots, f_n)$ . The premise of the predicament is that storing all of  $f$  is tedious ( $O(n)$ ) and we want to store something smaller.

The idea is to store instead  $y = Lf$  where  $L$  is a  $k \times n$  matrix so  $y \in \mathbb{R}^k$  and then ~~use~~ by the JL Lemma for  $k = \Omega\left(\frac{\ln(1/\delta)}{\epsilon^2 - \epsilon^3}\right)$

$$\Pr\left((1-\epsilon)\|f\|^2 \leq \|y\|^2 \leq (1+\epsilon)\|f\|^2\right) > 1-\delta$$

$$\downarrow$$

$$= \Pr\left(\left|\|y\|^2 - F_2\right| \leq \epsilon F_2\right) > 1-\delta.$$

To generate a streaming algorithm out of this, notice

$$f = \sum_{j=1}^m x_j \quad \text{where} \quad x_j = (0, \dots, 0, \underbrace{1}_{s_j \text{th position}}, 0, \dots, 0)$$

By linearity  $y = Lf = \sum_{j=1}^m Lx_j = \sum_{j=1}^m L_{s_j}$  where  $L_k$  is the  $k$ th row of  $L$ .

So the algorithm here is generate  $L = \left[ \text{Gaussian}\left(0, \frac{1}{n}\right) \right]^{k \times n}$  matrix and  $y = 0$ .

On stream element  $s_j$ , ~~add~~  $y \leftarrow y + L s_j$ .

Return  $\|y\|$ .

\* The downside to this method which has the same asymptotic complexity to the Tug-of-War Alg. seen in the first lectures is that we store  $y \in \mathbb{R}^k$ , so we would need to also pay attention to the decimal precision needed.

### Hoeffding's Inequality

Thm. If  $X_1, \dots, X_n$  are independent r.v. such that  $X_i \in [a_i, b_i]$  a.s.

(i.e.  $\Pr(a_i \leq X_i \leq b_i) = 1$ ) then for  $S = \sum_i X_i$ ,

$$\Pr(S - E(S) \geq t) \leq \exp\left(-\frac{2t^2}{\sum_i (b_i - a_i)^2}\right)$$

~~Thm~~ Hoeffding's Lemma Let  $X$  be a r.v. with  $E(X) = 0$  and  $a \leq X \leq b$  a.s.

Then for all  $\lambda \in \mathbb{R}$ ,

$$E(e^{\lambda X}) \leq \exp\left(\frac{\lambda^2 (b-a)^2}{8}\right).$$

### Pf. of Lemma

As  $x \mapsto e^{\lambda x}$  is convex, ~~and~~ <sup>for</sup>  $a \leq x \leq b$ ,

$$e^{\lambda x} \leq \left( \frac{b-x}{b-a} \right) e^{\lambda a} + \left( \frac{x-a}{b-a} \right) e^{\lambda b}$$

$$\begin{aligned} \Rightarrow E(e^{\lambda X}) &\leq \left( \frac{b-E(X)}{b-a} \right) e^{\lambda a} + \left( \frac{E(X)-a}{b-a} \right) e^{\lambda b} \\ &= \left( \frac{b}{b-a} \right) e^{\lambda a} - \frac{a}{b-a} e^{\lambda b} \end{aligned}$$

Let  $h = \lambda(b-a)$ ,  $p = \frac{-a}{b-a}$ , and  $L(h) = -hp + \ln(1-p+pe^h)$

then we can see

$$E(e^{\lambda X}) \leq e^{L(h)}. \quad \text{Taylor expand } L(h) \text{ to get that } \begin{pmatrix} L(0) = 0 \\ L'(0) = 0 \\ L''(h) \leq \frac{1}{4} \forall h \end{pmatrix}$$

$$L(h) \leq \frac{1}{8} h^2 = \frac{1}{8} \lambda^2 (b-a)^2.$$

$$\Rightarrow E(e^{\lambda X}) \leq \exp\left(\frac{\lambda^2 (b-a)^2}{8}\right).$$

□

### Pf. of Thm

For any r.v.  $X$ ,  $X - E(X)$  has mean 0. ~~And~~ And if  $a \leq X \leq b$  a.s.

$a - E(X) \leq X - E(X) \leq b - E(X)$  a.s. So applying the Lemma,

$$E(e^{\lambda(X-E(X))}) \leq \exp\left(\frac{\lambda^2 ((b-E(X)) - (a-E(X)))^2}{8}\right) = \exp\left(\frac{\lambda^2 (b-a)^2}{8}\right).$$

Assume  $X_1, \dots, X_n$  indep. r.v. with  $a_i \leq X_i \leq b_i$  a.s. and let

$$S = \sum_i X_i. \quad \text{Then for } s, t \geq 0$$

$$\begin{aligned} \Rightarrow \Pr(S - E(S) \geq t) &= \Pr\left(e^{s(S - E(S))} \geq e^{st}\right) \\ &\leq \frac{1}{e^{st}} E\left(e^{s(S - E(S))}\right) \quad (\text{Markov's Ineq.}) \\ &= \frac{1}{e^{st}} \prod_i E\left(e^{s(X_i - E(X_i))}\right) \\ &\leq \frac{1}{e^{st}} \prod_i \exp\left(\frac{s^2 (b_i - a_i)^2}{8}\right) \\ &= \exp\left(-st + \frac{1}{8} s^2 \sum_{i=1}^n (b_i - a_i)^2\right) \end{aligned}$$

Define  $C := \sum_i (b_i - a_i)^2$ . Then the minimizing  $s$  satisfies

$$0 = \frac{d}{ds} \left(-st + \frac{C}{8} s^2\right) = -t + \frac{C}{4} s \quad \Rightarrow \quad s = \frac{4t}{C}.$$

$$\begin{aligned} \Rightarrow \Pr(S - E(S) \geq t) &\leq \exp\left(-\frac{4t^2}{C} + \frac{C}{8} \frac{16t^2}{C^2}\right) \\ &= \exp\left(-\frac{2t^2}{C}\right) = \exp\left(-\frac{2t^2}{\sum_i (b_i - a_i)^2}\right). \end{aligned}$$