

In proof by cases, should the implications for <sup>✓</sup>all the cases be true or is it sufficient if even one ~~✓~~ of the implications is true?

$$\begin{aligned} & (p \vee q \vee r) \rightarrow s \quad \checkmark (p \rightarrow s) \wedge (q \rightarrow s) \wedge (r \rightarrow s) \\ \equiv & \neg (p \vee q \vee r) \vee s \quad \times (p \rightarrow s) \vee (q \rightarrow s) \vee (r \rightarrow s) \\ \equiv & (\neg p \wedge \neg q \wedge \neg r) \vee s \\ \equiv & (\neg p \vee s) \wedge (\neg q \vee s) \wedge (\neg r \vee s) \\ \equiv & (p \rightarrow s) \wedge (q \rightarrow s) \wedge (r \rightarrow s) \\ \equiv & \end{aligned}$$

How to prove the following statement by induction?

$$F_{m+n} = F_m F_{n+1} + F_{m-1} F_n \text{ for all } m \geq 1 \text{ and } n \geq 0.$$

- ① Induction on  $m$  (with  $n$  arbitrary)
- ② Induction on  $n$  (with  $m$  arbitrary)
- ③ Induction on  $m+n$

② Base cases:

$$n = 0$$

$$\text{LHS} = F_m$$

$$\text{RHS} = F_m F_1 + F_{m-1} F_0 = F_m$$

$$n = 1$$

$$\text{LHS} = F_{m+1}$$

$$\text{RHS} = F_m F_2 + F_{m-1} F_1 = F_m + F_{m-1}$$

Induction

$$n \geq 2 \quad \begin{cases} F_{m+n-1} = F_m F_n + F_{m-1} F_{n-1} \\ \text{by P}^6 \quad F_{m+n-2} = F_m F_{n-1} + F_{m-1} F_{n-2} \end{cases}$$

$$\begin{aligned} F_{m+n} &= F_m (F_n + F_{n-1}) \\ &\quad + F_{m-1} (F_{n-1} + \cancel{F_{n-2}}) \\ &= F_m F_{n+1} + F_{m-1} F_n \end{aligned}$$

③ Induction on  $m+n$

Base  $m+n=1 \rightarrow m=1, n=0$

$m+n=2 \rightarrow m=1, n=1$

or  $m=2, n=0$

Induction Take  $m+n \geq 3$

True for  $m+n-1, m+n-2$

Case 1:  $m \geq 3$

$$m+n-1 = (m-1) + n$$

$$m+n-2 = (m-2) + n$$

Case 2:  $n \geq 2$

$$m+n-1 = m + (n-1)$$

$$m+n-2 = m + (n-2)$$

Case 3:  $m < 3$  and  $n < 2$

$m \leq 2$  and  $n \leq 1$

$m + n \geq 3$

$m = 2$  and  $n = 1$

Prove that:  $\gcd(F_{n+1}, F_n) = 1$  for all  $n \geq 0$ .

①  $p \mid F_{n+1}$  and  $p \mid F_n$

$F_{n-1} = F_{n+1} - F_n$  is divisible by  $p$

$F_{n-2} = F_n - F_{n-1}$  is divisible by  $p$

$\dots$   
 $F_1$  is divisible by  $p$

$p \mid 1$  ✓

② By induction on  $n$ .

Prove that:

$$\gcd(F_m, F_n) = F_{\gcd(m, n)} \text{ for all } m, n \text{ (not both zero).}$$

$$m \geq n \quad m = qn + r$$

$$\gcd(F_m, F_n) = \gcd(F_n, F_r)$$

$$m = n + k$$

$$F_m = F_n F_{k+1} + F_{n-1} F_k$$

$$\gcd(F_m, F_n) = \gcd(F_{n-1} F_k, F_n)$$

$$= \gcd(F_k, F_n)$$

$$\gcd(c + kb, b) = \gcd(c, b)$$

$$= \gcd(F_{m-n}, F_n) = \gcd(F_{m-2n}, F_n) \\ = \dots = \gcd(F_r, F_n) \\ = \gcd(F_n, F_r)$$

How to prove the following statement?

$$\forall a, b, c \in \mathbb{N} \left[ \left( \gcd(a, b) = 1 \right) \Rightarrow \exists x \in \mathbb{N} \left( \gcd(a + bx, c) = 1 \right) \right]$$

$$a = p_1^{e_1} p_2^{e_2} \dots p_r^{e_r}$$

$$e_i > 0$$

$$b = q_1^{f_1} q_2^{f_2} \dots q_s^{f_s}$$

$$f_i > 0$$

$$c = p_1^{u_1} p_2^{u_2} \dots p_r^{u_r} q_1^{v_1} q_2^{v_2} \dots q_s^{v_s}$$

$$\boxed{\pi_1^{w_1} \dots \pi_t^{w_t}}$$

$$u_i \geq 0, v_j \geq 0, w_k > 0$$



Consider the following function with  $m$  and  $n$  non-negative integers.

```
int g ( int m, int n)
{
    if ((m == 0) || (n == 0)) return 1;
    return g(m,n-1) + g(m-1,n);
}
```

Express the return value of  $g(2,n)$  as a function of  $n$ .

$$g(0, n) = 1 \quad \forall n \geq 0$$

$$g(1, n) = g(1, n-1) + g(0, n)$$

$$= g(1, n-1) + 1$$

$$= g(1, n-2) + 2$$

$$\dots$$

$$= g(1, 0) + n = n + 1$$

$$g(2, n) = g(2, n-1) + g(1, n) = g(2, n-1) + n + 1$$

$$= g(2, n-2) + (n) + (n+1) = \dots = g(2, 0) + 2 + \dots + n + 1$$

$$= 1 + 2 + \dots + n + 1 = \frac{(n+1)(n+2)}{2}$$

```
for (i=1; i<=n; ++i) L[i] = 0;
```

```
for (i=1; i<=n; ++i)  
    for (j=i; j<=n; j+=i)  
        L[j] = 1 - L[j];
```

After this,  $L[i] = 1$  for which values of  $i$ ?

$$i = p_1^{e_1} \cdots p_r^{e_r}$$

# of factors of  $i$

$$= (e_1 + 1)(e_2 + 1) \cdots (e_r + 1)$$

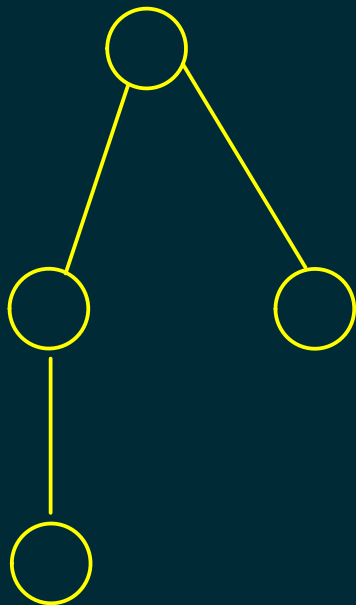
# Binomial Trees



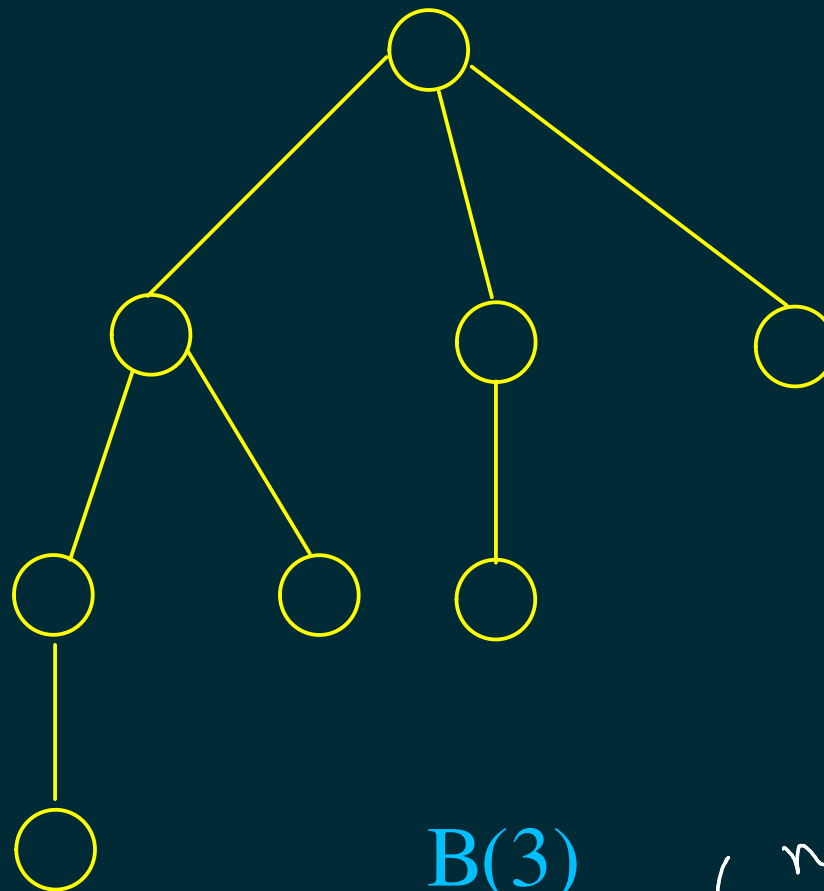
B(0)



B(1)



B(2)

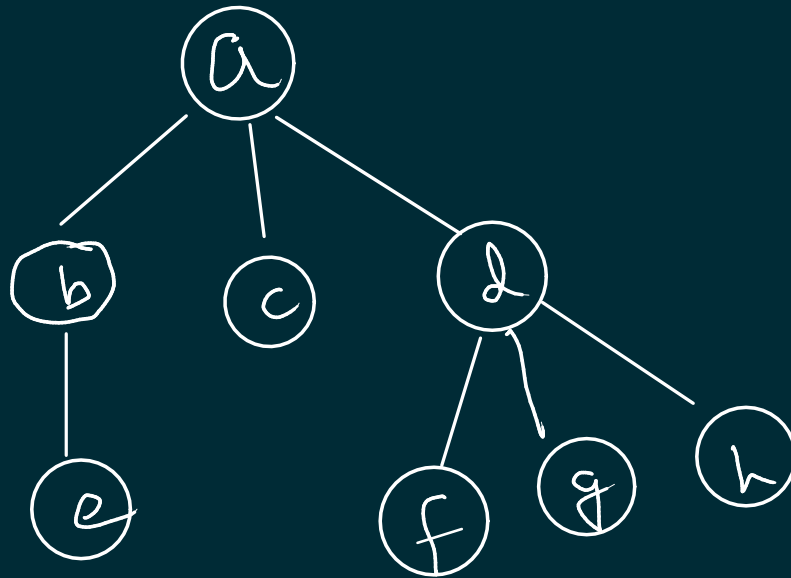


B(3)

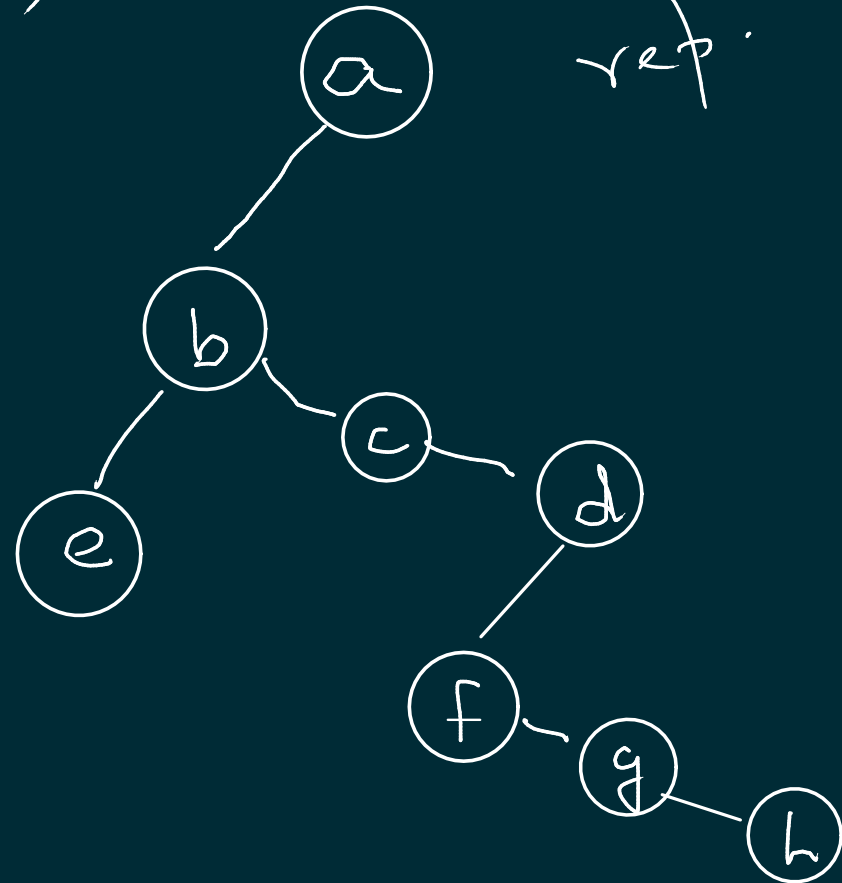
Prove that there are  $\binom{n}{i}$  nodes in B(n) at level i.

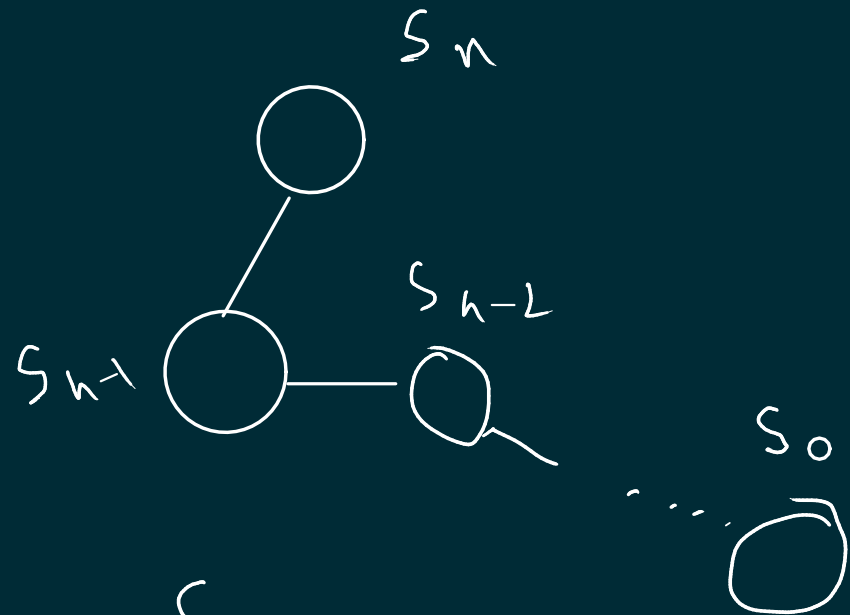
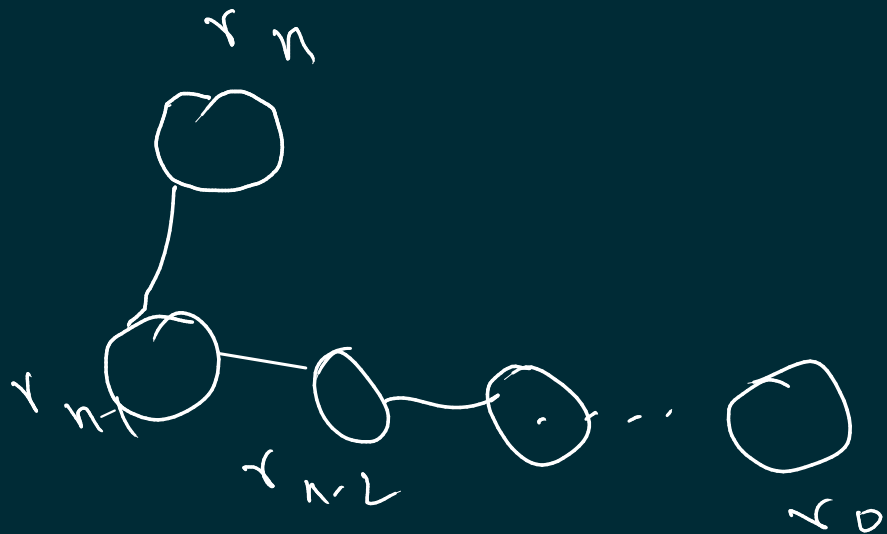
$\binom{n}{i} = 0$   
if  $n < i$

$B(n)$  contains  $2^n$  nodes (easy to prove by strong induction).  
You are given two disjoint copies of  $B(n)$ . How can you efficiently make a single copy of  $B(n+1)$ ?



First-child-next-sibling  
rep.





$O(1)$

time

Binomial  
heaps

