

2D Matrices & Basic Sorting Algorithms

- Introduction - A 2D matrix is a specific type of 2D array that has a rectangular grid of numbers, where each number is called an element.
- It is a mathematical structure that consists of a set of numbers arranged in rows & columns.
- 2D matrix can be declared as:

mat [N][M]
 ↕ ↑ ↑
 name no. of no. of
 rows columns

mat[i][j] → i^{th} row, j^{th} column

			0 1 2 ← col. index
0	(0,0)	(0,1)	(0,2)
1	(1,0)	(1,1)	(1,2)
2	(2,0)	(2,1)	(2,2) → (N-1, M-1) index
	↑		
Row index		NXM	
		↑ ↑	rows columns

Total no. of elements = $N * M$

		0 - - - - - M-1
0	(0,0)	(0,M-1)
N-1	(N-1,0)	(N-1,M-1)
		NXM

- Print row-wise sum:

Given a 2D matrix mat [N][M], print the row-wise sum.

$$\text{mat}[3][4] = \{ \{1, 2, 3, 4\},$$

$$\{5, 6, 7, 8\},$$

$$\{9, 10, 11, 12\} \}$$

O/P: 10 26 42

26

42 (N*M) : 12

(12) : 32

Approach :- Traverse each row

- While traversing take sum of all elements present in that row.

Approach 2nd: printing sum of all elements

def sumRow (mat, N, M):

for i in range(0, N): # iterate over all rows
sum = 0

iterate over that row
and get sum

for j in range(0, M):

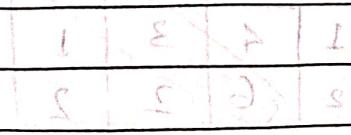
sum += mat[i][j]

After method of +9 and mat & sum with loop variable 9

+9 is also added at the end of loop iteration

TC: O(N*M)

SC: O(1)



10

0

0

0

0

0

0

0

0

1 2 3 4

5 6 7 8

9 10 11 12

• Print column wise sum:

```
def mat1_colRow(mat, N, M):
    for i in range(0, M):
        sum = 0
        for j in range(0, N):
            sum += mat[i][j]
        print(sum)
```

TC: $O(N * M)$

SC: $O(1)$

• Print diagonals of a square matrix:

Given a 2D square matrix, print diagonals
(cols = rows)

Q) How many main diagonals are there in a square matrix?

→ 2 ($M \times M$) number of main diagonals

1. Principal diagonal: From top left to bottom right
2. Anti-diagonal: From top right to bottom left.

	0	1	2
0	1	5	8
1	4	3	1
2	6	5	2

Principal
diagonal

O/p: 1 3 2

	0	1	5	8
0	1	4	3	1
1	4	3	1	
2	6	2	2	

Anti
diagonal

O/p: 8 3 6

Principal Diagonal

```
def printDiagonal (mat, N):
    i = 0
```

```
    while i < N:
```

```
        print (mat[i][i])
```

```
        i += 1
```

TC : O(N)

SC : O(1)

Anti-diagonal

```
def printDiagonal (mat, N):
    i = 0
```

```
    for i in range (N):
```

```
        while i < N:
```

```
            print (mat[i][N-i-1])
```

```
i += 1  
j -= 1
```

TC : O(N)

SC : O(1)



Introduction to Sorting:

Sorting is arranging data in a particular manner.

e.g.: 1, 7, 2, 9, 24

ele. → factors

1 → 1

7 → 2

2 → 2

9 → 3

24 → 8

∴ 1, 7, 2, 9, 24 → Yes, increasing order of factors

So, arranging data in increasing or decreasing order on the basis of any parameter is known as sorting.

Why is sorting important?

Helps in:

- Organized data
- Analysing data
- Searching
- Presenting

li.sort()

default → Inc. order
of value.

* Minimize the cost to empty array:

Given an array of n integers, minimize the cost of to empty given array where cost of removing an element is equal to sum of all elements left in an array.

Note: Add cost of removed element & remove it.

Way 1 : $\{2, 1, 4\}$ Initial condition

Remove 2 : $2 + 1 + 4 = 7$

$1, 8, 3 \leftarrow 1 + 4 = 5$ Condition $\rightarrow 7$

Remove 4 : $4 = 4$ Condition $\rightarrow 11$

Cost of removal = $7 + 5 + 4 = 16$

Way 2 : $\{2, 1, 4\}$

Remove 4 : $4 + 2 + 1 = 7$

2 : $2 + 1 = 3$

1 : $1 = 1$

Cost of removal = $7 + 3 + 1 = 11$

* Observation: $\text{arr}[4] = \{a, b, c, d\}$

Remove a : $a + b + c + d$

Remove b : $b + c + d$

Remove c : $c + d$

Removed : d

Cost = $a + 2b + 3c + 4d$

$\text{arr}[0] * 1 + \text{arr}[1] * 2 + \text{arr}[2] * 3$

+ $\text{arr}[3] * 4$

In general : $\text{arr}[i] * (i+1)$

★ Pseudocode:

```
def costmin(arr):
    arr.sort(reverse=True)
    total_cost = 0
    for i in range(0, len(arr)):
        total_cost += arr[i] * (i + 1)
    return total_cost
```

TC : #sorting + #operation $\rightarrow 6 \ 3 \ 1$
 $n \log n + n = O(n \ log n)$ total cost = 0
 $\Rightarrow O(n \ log n)$

SC : $O(n)$ i.e. #sort function
 $E = 1 + 2 + 3$ i = 1 3 * 2
 $i = 2 \ 1 * 3$

(15)

$$(1) = 1 + 2 + 3 = 6$$

Let's do a dry run [Arranged in降序] A

Sort the array in降序 is available

b + c + d = 6 By me

b + c = 5 By me

b = 3 By me

$$AP + DS + DS + SC = 1200$$

$$E = CS + AP + DS + SC + DS + SC$$

$$+ DS + SC$$

(1) * 6 = 600 : length of

★ Selection Sort →

arr = [5, 6, 4, 2] in降序排列

Pass 1 : 5 6 4 2

→ [2, 4, 5, 6] : FA1

Pass 2 : 2 6 4 5

→ [2, 4, 5, 6] : FA1

Pass 3 : 2 4 6 5

→ [2, 4, 5, 6] : FA1

2 4 5 6

→ [2, 4, 5, 6] : FA1

★ Pseudocode -

def selectionsort (arr, size):

i, j, mindex = 0 → (size-1)

no. of passes

for i in range (0, size-1):

mindex = i

compare & find smallest for j in range(i+1, size):

if arr[j] < arr[mindex]:

mindex = j

swap the smallest ele. with the first ele.

temp = arr[mindex]

arr[mindex] = arr[i]

arr[i] = temp

TC : O(N*N) | because it's a nested loop

SC : O(1)



Insertion Sorting :

Given an array $[N]$, and first $N-1$ elements are sorted sort the entire array.

arr [6]: [10, 3, 6, 8, 2, 5]

1. Assume that elements from 0 to i is sorted.
2. We have to insert $(i+1)$ element at its correct position in sorted array using insertion step.

0 1 2 3 4 5

arr [6]: [10, 3, 6, 8, 2, 5]

Step 1: 0 to 0 is sorted, place 1 correctly.

10 3 6 8 2 5

sorted

not sorted

Step 2: 0 to 1 is sorted, place 2 correctly.

3 10 6 8 2 5

10 3 6 8 2 5

Step 3: 0 to 2 is sorted, place 3 correctly.

3 6 10 8 2 5

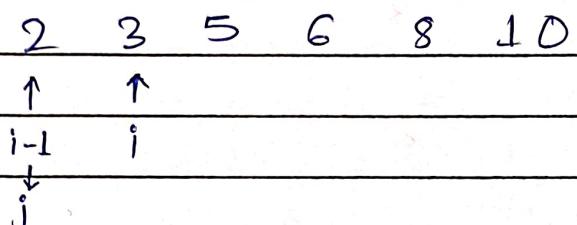
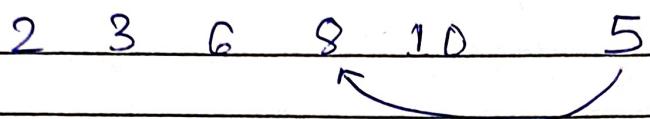
6 3 10 8 2 5

Step 4: 0 to 3 is sorted, place 4 correctly.

3 6 8 10 2 5

6 3 8 10 2 5

Step 5: 0 to 4 is sorted, place 5 correctly.



* Pseudocode -

```
def insertionSort(arr, n):
    for i in range(1, n):
        curele = arr[i]
        j = i - 1
```

while $j \geq 0$ and $arr[j] > curele$:

$$arr[j+1] = arr[j]$$

$$j = j - 1$$

$i=1$	$i=2$	$i=3$
$arr[0] = 3$	$arr[1] = 6$	$arr[2] = 8$
$curele = 3$	$curele = 6$	$curele = 8$
$j = 0 - 1$	$j = 1 \rightarrow 0$	$j = 2 \rightarrow 1$
$arr[0] = 3$	$arr[1] = 6$	$arr[2] = 8$

TC: $O(n^2)$

SC: $O(1)$

(In-place Algorithm)