

Hashing Introduction

- * Dictionary: Stores key value pair.
- * Set: Stores only keys, don't store duplicates
- * Functions:
 - 1) `size()` : no. of keys
 - 2) `add(key)`: adds a key to hs, if 'key' is already present, it won't do anything
 - 3) `contains(key)`: check if a key is present or not.
 - 4) `remove(key)`: If key is present, it'll remove it.
- Notes:
 - keys can't be duplicate
 - order of insertion is not maintained
 - TCC): add, contains, remove $\rightarrow O(1)$
 - $\rightarrow O(1)$
 - (String len)

Q) Given an arr[N] elements & Q queries, for each query find the freq. of each element.

ex: arr[10]: 2 6 3 8 7 2 2 3 8 10

Queries	Freq	Idea 1:
2	3	For every Query [i]
8	2	Find Q[i] in arr, if
3	2	found freq + 1
5	0	

TC: $O(Q * N)$

* Idea 2: dict [key, value]
 ↓ ↓
 element freq.

- Steps:
 - 1) Create a dict : if $\text{arr}[i]$ is not present in dict \rightarrow add $(\text{arr}[i], 1)$

2) Otherwise \rightarrow update $(\text{arr}[i], \text{freq} + 1)$

TC : $O(N)$

dict	$[2, 1]^2, [3, 1]^3$
	$[6, 1]$
	$[10, 1]$
	$[3, 1]^2$
	$[8, 1]^2$

- 3) Search for $\text{query}[i]$ in dict. & return the value if found else 0

TC : $O(Q)$

Overall TC : $O(N+Q)$

SC : $O(N)$

```
def freq_query(queries, arr):
    freq_map = {}
    # Step 1: Create the dict
    for element in freq_map: arr:
        if element in freq_map:
            freq_map[element] += 1
        else:
            freq_map[element] = 1
```

Step 2: Handle queries

```
result = []
for q in queries:
    if q in freq_map:
        result.append(freq_map[q])
    else:
        result.append(0)
return result
```

Q) Given arr (N) elements, find the count of all distinct elements

arr [5]: 0 1 2 3 4
 arr [5]: 3 5 6 5 4 | 3 5 6 4 | set
 ans = 4 size(set) = 4

arr [5]: 0 1 2 3 4
 arr [5]: 1 1 1 2 2 | 1 2 | set
 size (set) = 2

Idea: 1. Create set (N) O : CF

2. Get length of set

def count_distinct (arr):

unique = set (arr)

return size(unique)

TC : O(N)

SC : O(N)

Q) Given a strings, find the length of the longest substring without repeating characters

abcabcbb — ③

bbbb — ①

pwwkew — ⑤

* Brute force approach:

1) Generate all substrings : $O(N^2)$

2) Check uniqueness : $O(N)$

Total : TC : $O(N^3)$

SC : O(1)

* Observations for optimization :

Get unique chars - set

Find substring - sliding window

0 1 2 3 4 5 6 7

a b c a b c b b

$s = 0, e = 0, \text{len} = 1$	a
$s = 0, e = 1, \text{len} = 2$	ab
$s = 0, e = 2, \text{len} = 3$	abc
$s = 1, e = 3, \text{len} = 3$	bca
$s = 2, e = 4, \text{len} = 3$	cab
$s = 3, e = 5, \text{len} = 3$	abc
$s = 5, e = 6, \text{len} = 2$	cb
$s = 7, e = 7, \text{len} = 1$	b

def longest_substring(s) :

 start = 0

 max_len = 0

 char_set = set()

 for e in range(len(s)):

 while s[e] in char_set:

 char_set.remove(s[start])

 start += 1

 char_set.add(s[e])

 max_len = max(max_len, e - start + 1)

 return max_len

TC: O(N) SC: N - len of str
 $m - 1 + n$ of sub-str $O(\min(N, m))$