24TH JULY COMPLEXITY ANALYSIS AND ARRAY

Q₁: Program to print the sum of all the elements present at even indices in the given array:

```
copy code
public class SumOfEvenIndicesElements {
  public static void main(String[] args) {
    int[] arr = {3, 20, 4, 6, 9};
    int sum = 0;

    for (int i = 0; i < arr.length; i += 2) {
        sum += arr[i];
    }

    System.out.println("Output: " + sum);
}
</pre>
```

Q2: Program to traverse the array using a for-each loop and print all even elements:

```
java
Copy code
public class PrintEvenElements {
   public static void main(String[] args) {
     int[] arr = {34, 21, 54, 65, 43};

     System.out.print("Output: ");
     for (int num : arr) {
        if (num % 2 == 0) {
            System.out.print(num + " ");
        }
}
```

```
}
```

Q3: Program to calculate the maximum element in the array:

```
java
Copy code
public class MaxElement {
    public static void main(String[] args) {
        int[] arr = {34, 21, 54, 65, 43};
        int max = arr[0];

        for (int i = 1; i < arr.length; i++) {
            if (arr[i] > max) {
                max = arr[i];
            }
        }
        System.out.println("Output: " + max);
      }
}
```

Q4: Program to find the second largest element in the given array:

```
copy code
public class SecondLargestElement {
  public static void main(String[] args) {
    int[] arr = {34, 21, 54, 65, 43};
    int largest = arr[0];
    int secondLargest = Integer.MIN_VALUE;

for (int i = 1; i < arr.length; i++) {
    if (arr[i] > largest) {
        secondLargest = largest;
        largest = arr[i];
    } else if (arr[i] > secondLargest && arr[i] != largest) {
```

```
secondLargest = arr[i];
}

System.out.println("Output: " + secondLargest);
}
```

Q5: Program to find the first peak element in the array:

```
java
Copy code
public class FirstPeakElement {
   public static void main(String[] args) {
     int[] arr = {1, 3, 2, 6, 5};

     for (int i = 1; i < arr.length - 1; i++) {
        if (arr[i] > arr[i - 1] && arr[i] > arr[i + 1]) {
            System.out.println("Output: " + arr[i]);
            break;
        }
     }
}
```

Q6: Java code to count positive, negative, odd, even, and zero numbers from user inputs:

```
import java.util.Scanner;

public class NumberCounter {
   public static void main(String[] args) {
      Scanner scanner = new Scanner(System.in);

      System.out.print("Enter the number of rows (m): ");
      int m = scanner.nextInt();
```

```
System.out.print("Enter the number of columns (n): ");
int n = scanner.nextInt();
int positiveCount = 0;
int negativeCount = 0;
int oddCount = 0;
int evenCount = 0:
int zeroCount = 0;
System.out.println("Enter " + (m * n) + " integer inputs:");
for (int i = 0; i < m * n; i++) {
  int num = scanner.nextInt();
  if (num > 0) {
    positiveCount++;
  } else if (num < 0) {
    negativeCount++;
  if (num % 2 == 0) {
    evenCount++:
  } else {
    oddCount++:
  if (num == 0) {
    zeroCount++;
System.out.println("Number of positive numbers = " + positiveCount);
System.out.println("Number of negative numbers = " + negativeCount);
System.out.println("Number of odd numbers = " + oddCount);
System.out.println("Number of even numbers = " + evenCount);
System.out.println("Number of 0 = " + zeroCount);
```

```
scanner.close();
}
```

Q7: Java code to print elements above the secondary diagonal in a user-inputted square matrix:

```
```java
import java.util.Scanner;
public class SecondaryDiagonalElements {
 public static void main(String[] args) {
 Scanner scanner = new Scanner(System.in);
 System.out.print("Enter the size of the square matrix: ");
 int size = scanner.nextInt();
 int[][] matrix = new int[size][size];
 System.out.println("Enter the elements of the square matrix:");
 for (int i = 0; i < size; i++) {
 for (int j = 0; j < size; j++) {
 matrix[i][j] = scanner.nextInt();
 System.out.println("Elements above the secondary diagonal:");
 for (int i = 0; i < size; i++) {
 for (int j = i + 1; j < size; j++) {
 System.out.print(matrix[i][j] + " ");
 scanner.close();
```

```
}
```

Q8: Java code to print elements of both the diagonals in a user-inputted square matrix:

```
```java
import java.util.Scanner;
public class DiagonalElements {
  public static void main(String[] args) {
     Scanner scanner = new Scanner(System.in);
     System.out.print("Enter the size of the square matrix: ");
     int size = scanner.nextInt();
     int[][] matrix = new int[size][size];
     System.out.println("Enter the elements of the square matrix:");
     for (int i = 0; i < size; i++) {
       for (int j = 0; j < size; j++) {
          matrix[i][j] = scanner.nextInt();
     System.out.println("Elements of the secondary diagonal:");
     for (int i = 0; i < size; i++) {
       System.out.print(matrix[i][size - i - 1] + " ");
     }
     System.out.println("\nElements of the primary diagonal:");
     for (int i = 0; i < size; i++) {
       System.out.print(matrix[i][i] + " ");
```

```
scanner.close();
}
```

Q9: Java code to find the largest element of a given 2D array of integers:

```
```java
import java.util.Scanner;
public class LargestElement {
 public static void main(String[] args) {
 Scanner scanner = new Scanner(System.in);
 System.out.print("Enter the number of rows: ");
 int rows = scanner.nextInt();
 System.out.print("Enter the number of columns: ");
 int columns = scanner.nextInt();
 int[][] arr = new int[rows][columns];
 System.out.println("Enter the elements of the array:");
 for (int i = 0; i < rows; i++) {
 for (int j = 0; j < columns; j++) {
 arr[i][j] = scanner.nextInt();
 }
 int largest = arr[0][0];
 for (int i = 0; i < rows; i++) {
 for (int j = 0; j < columns; j++) {
```

```
if (arr[i][j] > largest) {
 largest = arr[i][j];
}

System.out.println("Largest element in the array: " + largest);
scanner.close();
}
```

Q<sub>10</sub>: Java function to display the elements of the middle row and middle column of a square matrix:

```
printMiddleRowAndColumn(matrix);
 scanner.close();
public static void printMiddleRowAndColumn(int[][] matrix) {
 int size = matrix.length;
 int middleRow = size / 2:
 int middleColumn = size / 2;
 System.out.println("Elements of the middle row:");
 for (int j = 0; j < size; j++) {
 System.out.print(matrix[middleRow][j] + "");
 }
 System.out.println("\nElements of the middle column:");
 for (int i = 0; i < size; i++) {
 System.out.print(matrix[i][middleColumn] + " ");
```

11. Analyze the time complexity of the following Java code and suggest a way to improve it: int sum = 0; for (int i = 1; i < n; i+)  $\{$  for(int j = 1; j < i; jH)  $\{$  sUm+;  $\}$   $\}$ 

Ans: The time complexity of the given Java code is  $O(n^2)$ . It consists of two nested loops, one iterating from 1 to n and the other from 1 to i. As a result, the total number of iterations will be the sum of the first n natural numbers, which is  $O(n^2)$ .

To improve the time complexity, we can optimize the code by using a mathematical formula to directly calculate the sum of the first n natural numbers, which is n \* (n + 1) / 2. This way, we can achieve a linear time complexity of O(n) for calculating the sum, rather than O( $n^2$ ) as in the original code.

Improved Java code:

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**Copy code** 

int n = ... // some value of n

```
int sum = n * (n + 1) / 2:
```

12: Find the value of T(2) for the recurrence relation T(n) = 3T(n-1) + 12n, given that T(0) = 5.

Ans: To find the value of T(2) for the given recurrence relation T(n) = 3T(n-1) + 12n with T(0) = 5, we can recursively apply the relation:

```
T(2) = 3T(1) + 12(2)

T(1) = 3T(0) + 12(1) = 3(5) + 12 = 27

T(2) = 3(27) + 12(2) = 81 + 24 = 105
```

So, the value of T(2) is 105.

13: Given a recurrence relation, solve it using a substitution method. Relation: T(n) = T(n-1) + c

Ans: To solve the given recurrence relation T(n) = T(n-1) + c using the substitution method, we can repeatedly substitute the expression for T(n-1) into T(n), until we reach the base case T(0) = 5.

```
T(n) = T(n-1) + c
= (T(n-2) + c) + c
= T(n-2) + 2c
= (T(n-3) + c) + 2c
= T(n-3) + 3c
...
= T(0) + nc
= 5 + nc
```

Therefore, the solution to the recurrence relation is T(n) = 5 + nc.

14: Given a recurrence relation:  $T(n) = 16T(n/4) + n_2\log n$  Find the time complexity of this relation using the master theorem.

Ans: To find the time complexity of the given recurrence relation using the master theorem, we need to first identify the values of a, b, and f(n).

```
T(n) = 16T(n/4) + n^2 * log(n)
```

Here, a = 16 (the number of recursive calls).

b = 4 (the factor by which n is divided in each recursion), and

 $f(n) = n^2 * log(n).$ 

Now, let's calculate the value of log\_b(a):

```
log_4(16) = 2, because 4^2 = 16.
```

Comparing f(n) with  $n \log_b(a)$ , we have:

```
n^2 * log(n) = n^2 * log_4(16) = n^2.
```

Since  $f(n) = n^2$  and  $n^{\log}b(a) = n^2$ , we are in case 2 of the master theorem.

The time complexity of the given recurrence relation is  $O(n^2 * log(n))$ .

15: Solve the following recurrence relation using recursion tree method T(n) = 27(n/2) + n

Ans: To solve the recurrence relation T(n) = 27(n/2) + n using the recursion tree method, we'll construct a recursion tree and sum up the costs at each level.

**Recursion tree:** 

```
SCSS
```

## Copy code

```
T(n)

|
T(n/2) + n

/ \
T(n/4) + n/2 T(n/4) + n/2

... ...
/
T(1) + 1 T(1) + 1
```

The tree has a depth of log(n) (base 2), and at each level, the cost is n. Therefore, the total cost is the sum of costs at each level, which is:

```
Total cost = n + n/2 + n/4 + ... + 1
```

This is a geometric series with a common ratio of 1/2, and the sum of a geometric series is given by the formula:

```
Sum = (first term) * (1 - (common ratio)^(number of terms)) / (1 - common ratio)
```

Plugging in the values, we get:

```
Total cost = n * (1 - (1/2)^{\log(n)}) / (1 - 1/2)
= n * (1 - 1/n) / (1/2)
= 2n * (1 - 1/n)
```

Therefore, the time complexity of the recurrence relation T(n) = 27(n/2) + n using the recursion tree method is O(n).

16. T(n) = 27(n/2) + K, Solve using Recurrence tree method.

Ans: To solve the recurrence relation T(n) = 27(n/2) + K using the recursion tree method, we'll construct a recursion tree and sum up the costs at each level.

**Recursion tree:** 

## SCSS

## **Copy code**

```
T(n)
 T
 T(n/2) + K
T(n/4) + K T(n/4) + K
 ...
 T(1) + K T(1) + K
```

The recursion tree has a depth of log(n) (base 2), and at each level, the cost is K. Therefore, the total cost is the sum of costs at each level, which is:

Total cost = K + K + K + ... (log(n) times)

This can be expressed as:

Total cost = K \* log(n)

Therefore, the time complexity of the recurrence relation T(n) = 27(n/2) + K using the recursion tree method is O(log(n)).