

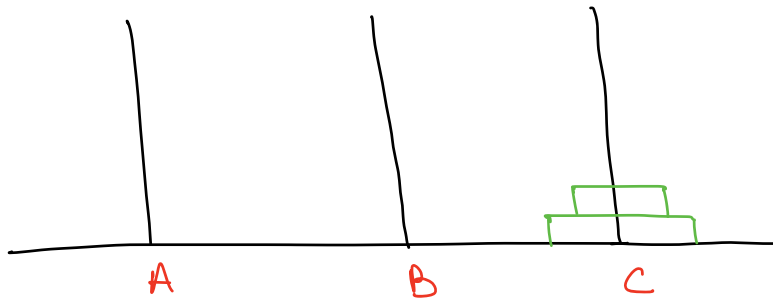
Tower of Hanoi

→ 3 towers A, B, C

→ There are **N** disks placed on A initially.

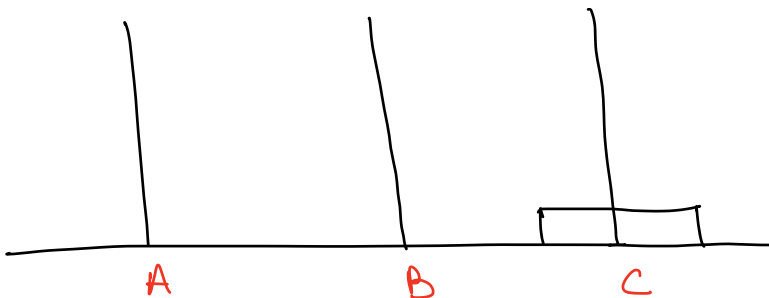
→ Move all the disks from A to C using B, making sure that at **no point** a smaller disk is below a larger disk.

Ex **N = 2**



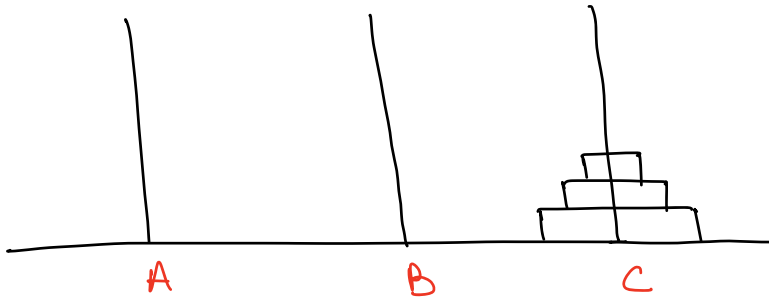
- 1) $A \rightarrow B$
- 2) $A \rightarrow C$
- 3) $B \rightarrow C$

Ex **N = 1**



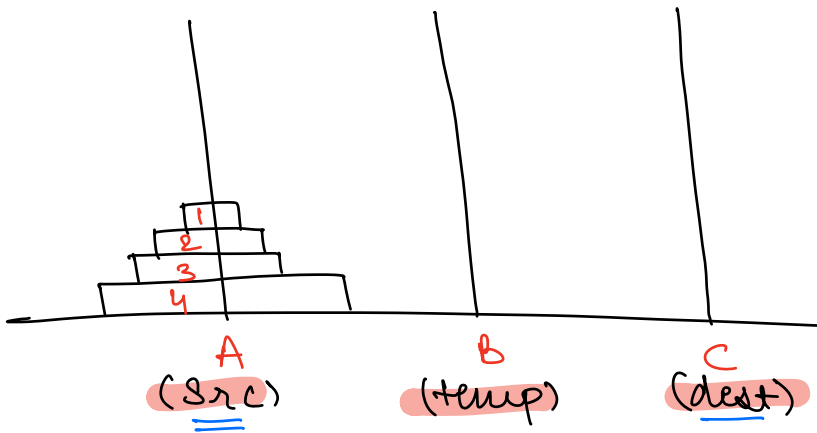
$A \rightarrow C$

Ex N=3

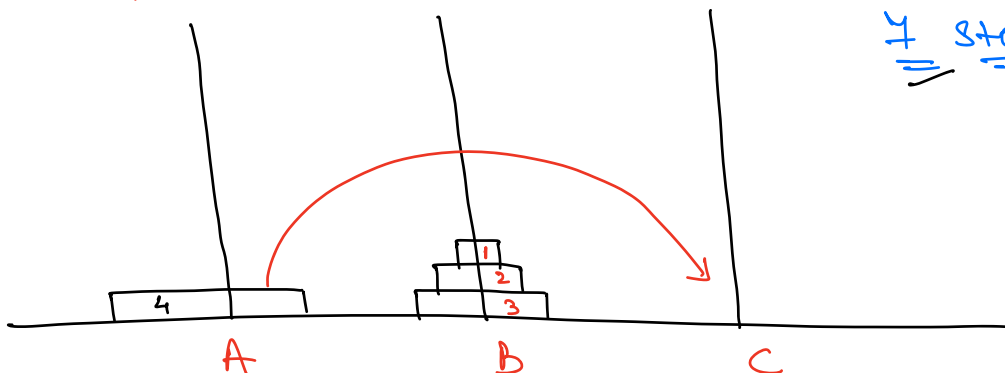


A → C
A → B
C → B
A → C
B → A
B → C
A → C

N=4 : Transfer 4 disks from tower (A) to (C) using tower B.

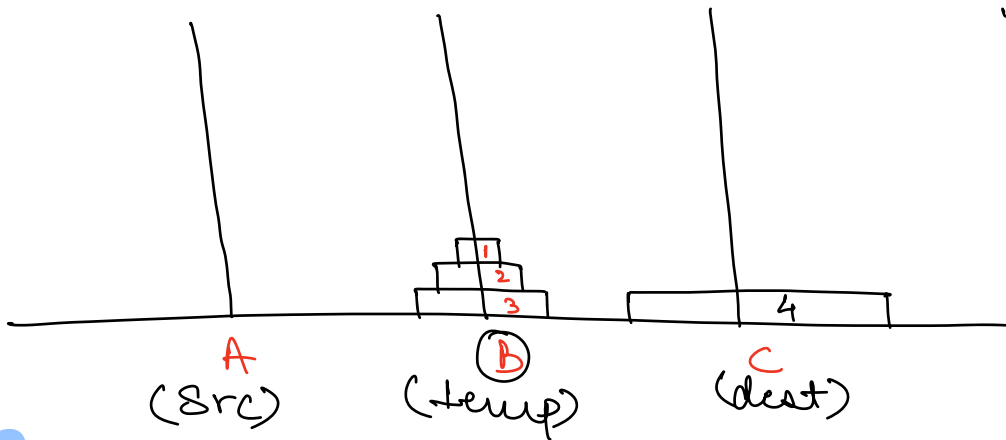


i) Transfer 1, 2, 3 disks from (A) to (B) using (C)



4 steps

ii) Transfer disk 4 from (A) to (C). ✓



iii) Transfer 1, 2, 3 disks from (B) to (C) using (A)
↳ 7 steps

⇒ Total 15 steps

// Print steps to move N disks from src to dest using temp.

```
void TOH (N, src, dest, temp) {  
    if (N == 0) {  
        return;  
    }  
    TOH (N-1, src, temp, dest); ✓  
    print ("src → dest"); ✓  
    TOH (N-1, temp, dest, src); ✓
```

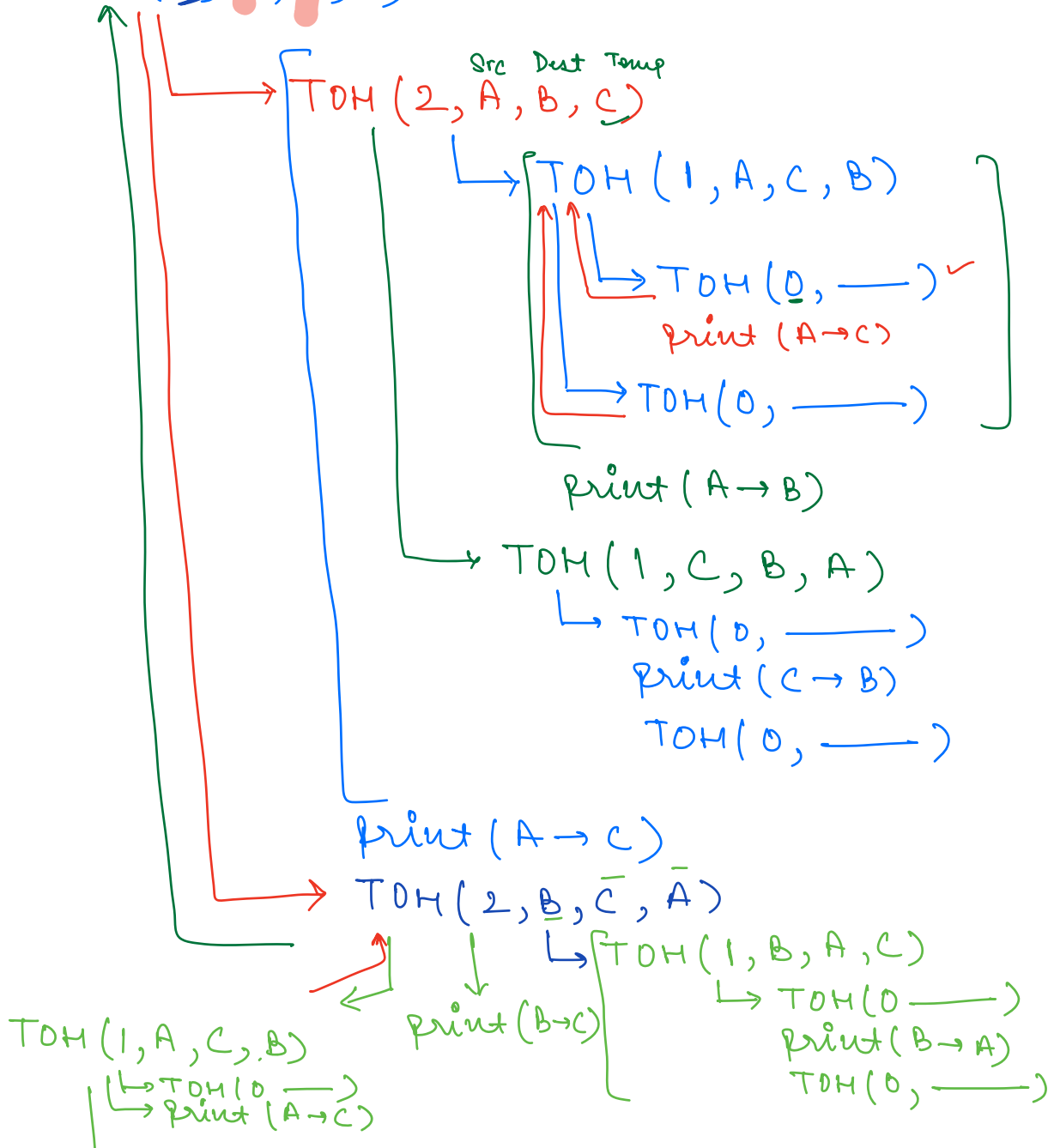
3

N=3

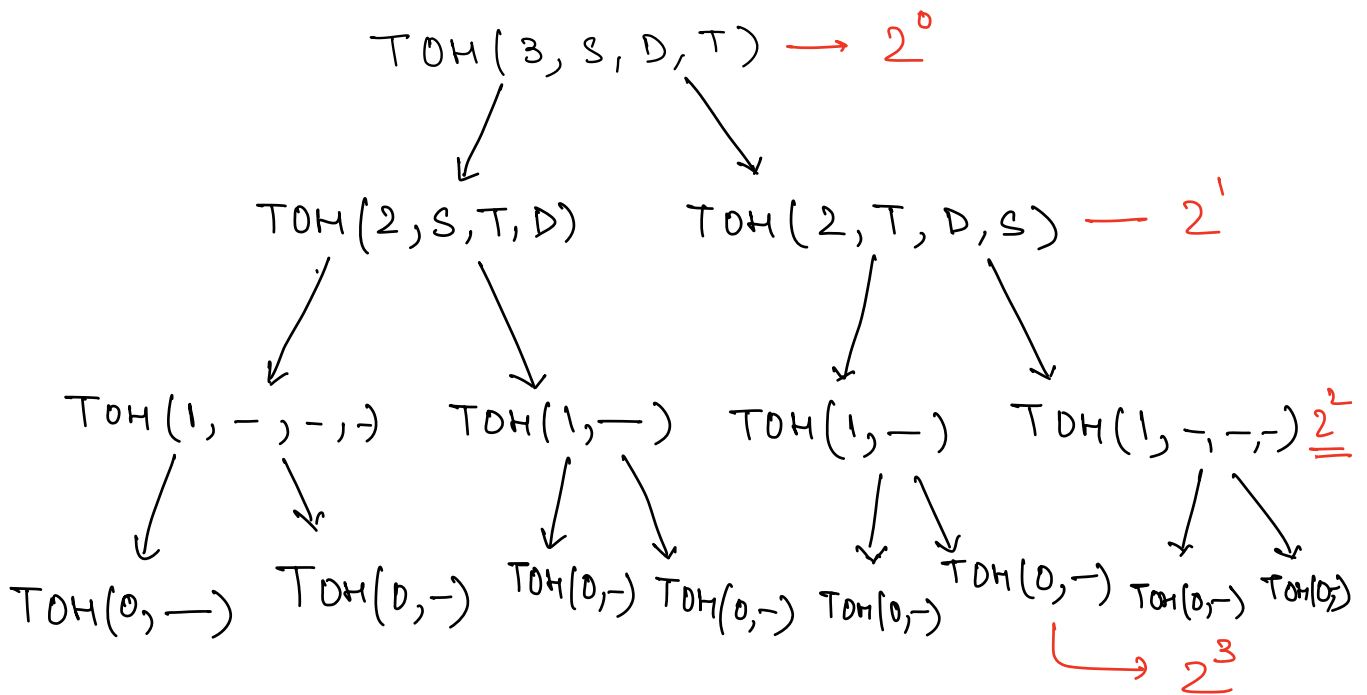


$A \rightarrow C$
 $A \rightarrow B$
 $C \rightarrow B$
 $A \rightarrow C$
 $B \rightarrow A$
 $B \rightarrow C$
 $A \rightarrow C$

TOH(3, A, C, B)



↳ $\text{TOH}(0, -)$



TC: $\left(\begin{array}{l} \text{\# of function} \\ \text{calls} \end{array} \right) * \left(\begin{array}{l} \text{Work done in} \\ \text{one fun}^{\wedge} \end{array} \right)$

No. of fun^{\wedge} calls = no. of nodes in Recursion Tree

$$= 2^0 + 2^1 + 2^2 + \dots + 2^{N-1}$$

$$a = 2^0$$

$$r = 2$$

of terms = $r = N$.

$$\text{Sum} = \frac{2^0 (1 - 2^N)}{1 - 2} = 2^N - 1.$$

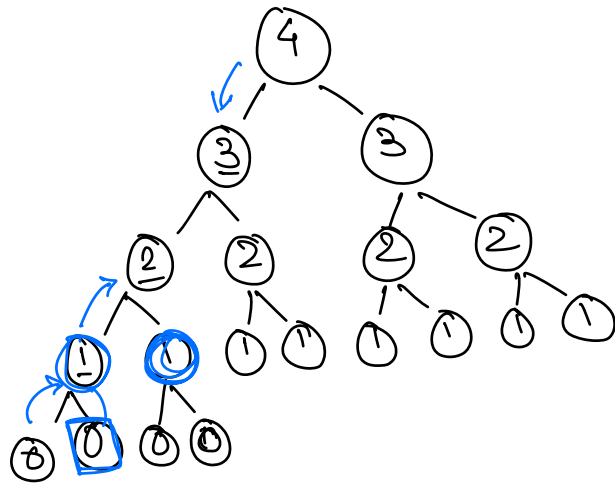
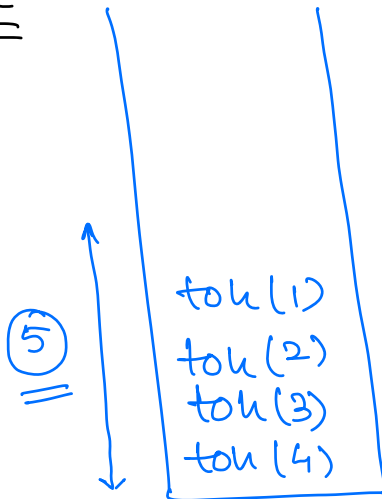
$$TC : O(2^N)$$

Recurrence Relation

$$T(N) = 2T(N-1) + 1$$

$$TC : O(2^N)$$

SC



$$SC : O(\text{Height})$$

$$SC : O(N)$$

Sorting Basics

Arranging the data in a specific order based on same parameter.

{ 1, 2, 3, 4, 5 } \Rightarrow Ascending order based on the value.

{ 5, 4, 3, 2, 1 } \Rightarrow Descending order based on the value.

{ 7, 2, 4, 9, 6 } \Rightarrow Asc. order based on the no. of factors.

↓ ↓ ↓ ↓ ↓
2 2 3 3 4

Why Sorting?

\rightarrow It makes searching faster.

Ex

12th %

Sauran	77
Siddhesh	70
Deepti	76
Shiv	98
Ayush	70
Bhavesh	57
Shamblu	78
Sandhya	85
Deepak	85

Sort \rightarrow

Bhavesh	57
Siddhesh	70
Ayush	70
Deepti	76
Sauran	77
Shamblu	78
Sandhya	85
Deepak	85
Shiv	98

Stable Sort

⇒ If two data points are equal then their relative order in sorted data should be maintained.

Inplace Algorithm

↳ SC: $O(1)$
No extra space.