

2 pointers

→ C++ X

→ indices iterating over the Array.

Q.1  
Amazon

Given a sorted Array of size  $N$  of distinct elements. Check if there exists a pair  $(i, j)$

s.t.  $A[i] + A[j] = K$ ,  $i \neq j$ .

⇒ 2 Sum.

$A: \{ 3, 7, 8, 12, 14 \}$   $K = 15$

→ True

Approach # 1 :-

TC :  $O(N^2)$

SC :  $O(1)$

Approach # 2 :-

TC :  $O(N)$

SC :  $O(N)$

Approach # 3 :-

$A: \{ 3, 7, 8, 12, 14 \}$   $K = 15$   
↑

Search  $(K - A[i])$

in Array.

⇒

TC :  $O(N \log N)$

SC :  $O(1)$

Approach # 4 :-

A: {<sup>0</sup>-3, <sup>1</sup>0, <sup>2</sup>1, <sup>3</sup>3, <sup>4</sup>6, <sup>5</sup>8, <sup>6</sup>11, <sup>7</sup>14, <sup>8</sup>18, <sup>9</sup>25} K = 17

*(Note: In the original image, 'i' is written below -3 and 'j' is written below 25, both with red arrows pointing to their respective indices in the array.)*

i	j	A[i]	A[j]	Sum
0	9	-3	25	22 > 17
0	8	-3	18	15 < 17
1	8	0	18	18 > 17
1	7	0	14	14 < 17
2	7	1	14	15 < 17
3	7	3	14	<u>17 == 17</u>

return true.

```
bool Check ( A[], N, K ) {  
    i = 0, j = N-1;  
    while ( i < j ) {  
        if ( A[i] + A[j] == K )  
            return true;  
        else if ( A[i] + A[j] < K )  
            i++;  
        else  
            j--;  
    }  
    return false;  
}
```

3

$$TC : O(N)$$

$$SC : O(1)$$

Q.2  
Amazon/  
Visa

Given a sorted Array of size  $N$  of distinct elements. Check if there exists a pair  $(i, j)$  s.t  $A[i] - A[j] = (k)$ ,  $k > 0$

$A[] : \{-3, 0, 1, 3, 6, 8, 11, 14, 18, 25\}$

$\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ \uparrow & & \uparrow & & \uparrow & & & & & \\ P_1 & & P_2 & & & & & & & \end{matrix}$

$k = 5$   $\Rightarrow$  True

$$25 - 0 = (25)$$

	$P_1$	$P_2$	$A[P_2] - A[P_1]$
1)	0	$N-1$	$\times$ $28 > 5$ (decrease the diff) $P_1++$ $P_2--$
2)	0	$N/2$	$\times$ $11 > 5$
3)	$\frac{N}{2}$	$\frac{N}{2} + 1$	$3 < 5$
<u>4)</u>	0	1	$3 < 5$
	0	2	$4 < 5$
	0	3	$6 > 5$
	1	3	$3 < 5$
	1	4	$6 > 5$
	2	4	$5 = 5$

return  
True.

$$A[P] : \left\{ \begin{matrix} 0 \\ 1 \\ 2 \end{matrix} \right\} \quad \begin{matrix} 1 \\ 4 \\ 6 \end{matrix} \quad \begin{matrix} 2 \\ 1 \\ 2 \end{matrix} \quad K=2$$

$\underline{\underline{P1}}$        $\underline{\underline{P2}}$        $A[P2] - A[P1]$   
 0      1      3 7 2  $\Rightarrow P1++$   
  
 $\underline{1}$        $\underline{1}$       if ( $P1 == P2$ )  $\Rightarrow P2++$   
  
 1      2      2 == 2  $\Rightarrow$  return true.

```
bool check ( A[], N, K ) {
    p1 = 0
    p2 = 1
    while ( p2 < N ) {
        if ( A[p2] - A[p1] == K )
            return true;
        else if ( A[p2] - A[p1] > K ) {
            p1++
            if ( p1 == p2 ) p2++
        }
        else {
            p2++
        }
    }
    return false;
}
```

$$T_C: O(N)$$
$$Sc : O(1)$$

A:  $\{ \underset{\substack{\uparrow \\ p_1}}{3}, 7, 8, 12, \underset{\substack{\uparrow \\ p_2}}{14} \}$   $K = 16$ .

P1	P2	diff (A[P2] - A[P1])
0	1	4 < 16 P2++
0	2	5 < 16 P2++
0	3	9 < 16 P2++
0	4	16 == 16 ⇒ <u>return</u> <u>tone.</u>

## What if  $K < 0$ .

$$\{ \underset{\uparrow p_2}{3}, \underset{\uparrow p_1}{8} \}$$

abs(K)

\* Important points about 2 pointers approach

1) Where to initialize the pointers.

2) How to update your pointers.

3) When to stop?

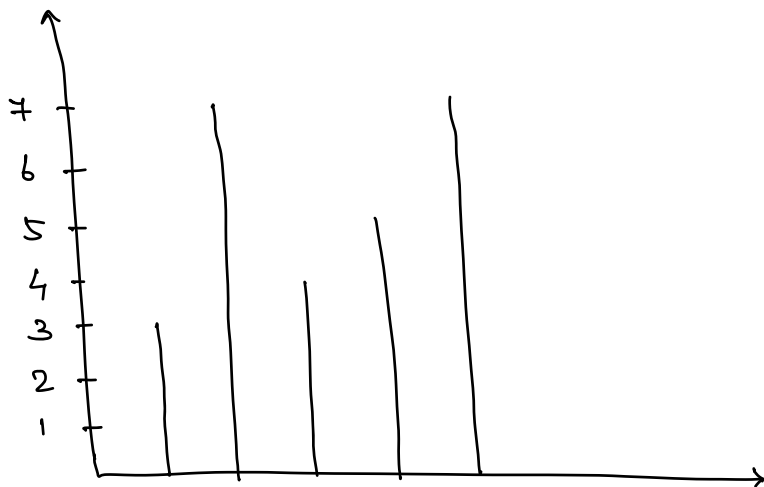
Q.3

Arceium  
Swiggy  
Intuit

## Rain Water Trapping

Given an Array of size  $N$ , where  $A[i]$  represents the height of  $i^{\text{th}}$  wall.  
Pick any 2 walls s.t max water can be stored b/w them.

Ex:  $\{ \overset{0}{3}, \overset{1}{7}, \overset{2}{4}, \overset{3}{5}, \overset{4}{7} \}$



$i, j$

water = Height \* width

$$\text{water} = \min(A[i], A[j]) * (j - i)$$

Brute force

```
for ( i = 0 ; i < N ; i++ ) {  
    for ( j = i+1 ; j < N ; j++ ) {  
        H = min ( A[i] , A[j] )  
        W = j - i  
        water = H * W  
        ans = max ( ans , water ) ;  
    }  
}
```

TC :  $O(N^2)$

SC :  $O(1)$

A : { <sup>0</sup>3, <sup>1</sup>5, <sup>2</sup>4, <sup>3</sup>7, <sup>4</sup>3, <sup>5</sup>6, <sup>6</sup>5, <sup>7</sup>4, <sup>8</sup>1, <sup>9</sup>2 }

P1

P2

$$\text{water} = \overbrace{\min(A[i], A[j])}^H * \underbrace{(j-i)}_{W.}$$

25

P1	P2	$H = \min(A[P1], A[P2])$	$W = P2 - P1$	water
0	9	2	9	18.
0	8	1	8	8
0	7	3	7	21
1	7	4	6	24
1	6	5	5	25
2	6	4	4	16
3	6	5	3	15
3	5	6	2	12
3	4	3	1	3
3	3	7	0	0

- \* Update pointer which is having lesser height wall.
- \* If the heights of the wall at P1 & P2 are equal then we can update any of them.

Code :

Todo.



Q.4 Given 3 sorted Arrays  $A[]$ ,  $B[]$  &  $C[]$  of sizes  $N$ . Find  $i, j$  &  $k$  s.t  
MS:  $\max(A[i], B[j], C[k]) - \min(A[i], B[j], C[k])$  is minimized.

$A: \{ 3, 14, 16, 20, 29, 40 \}$

$B: \{ -6, 23, 24, 30, 35, 50 \}$

$C: \{ -15, 15, 26, 31, 39, 42 \}$

i	j	k	A[i]	B[j]	C[k]	max	min	ans
0	0	0	3	-6	-15	3	-15	18
0	1	1	3	23	15	23	3	20
4	3	3	29	30	31	31	29	2
0	5	5	3	50	42	50	3	47
5	5	5	40	50	42	50	40	10
5	4	4	40	35	39	40	35	5
5	5	4	40	50	39	50	39	11

Brute force

TC:  $O(N^3)$

SC:  $O(1)$

A: { <sup>0</sup>3, <sup>1</sup>14, <sup>2</sup>16, <sup>3</sup>20, <sup>4</sup>29, <sup>5</sup>40 }

B: { <sup>-6</sup>, 23, 24, 30, 35, 50 }

C: { -15, 15, 26, 31, 39, 42 }

$$A[P_1] = 3$$

$$B[P_2] = -6$$

$$C[P_3] = -15$$

$$\max = 3$$

$$\min = -15$$

$$\text{ans} = \underline{\underline{18}}$$

$$\begin{array}{c} X - Y \\ \swarrow \quad \searrow \\ \max(A_i, B_j, C_k) \quad \min(A_i, B_j, C_k) \end{array}$$

\* To minimize the value

$X - Y$ , we are trying to increase the value of Y.

A: { <sup>0</sup>3, <sup>1</sup>14, <sup>2</sup>16, <sup>3</sup>20, <sup>4</sup>29, <sup>5</sup>40 }

B: { -6, 23, 24, 30, 35, 50 }

C: { -15, 15, 26, 31, 39, 42 }

P1	P2	P3	A[P1]	B[P2]	C[P3]	max	min	ans.
0	0	0	3	-6	-15	3	-15	18
0	0	1	3	-6	15	15	-6	21
0	1	1	3	23	15	23	3	20
1	1	1	14	23	15	23	14	9
2	1	1	16	23	15	23	15	8
2	1	2	16	23	26	26	16	10
3	1	2	20	23	26	26	20	6
4	1	2	29	23	26	29	23	6
4	2	2	29	24	26	29	24	5
4	3	2	29	30	26	30	26	4
4	3	3	29	30	31	31	29	2
5	3	3	40	30	31	40	30	10
5	4	3	40	35	31	40	31	9
5	4	4	40	35	39	40	35	5
5	5	4	40	50	39	50	39	11
5	5	5	40	50	42	50	40	10
5	5	5	40	50	42	50	40	10
6	5	5						

$P1 == N$   
break.

ans = 2

Code

(HW)

$P1 = P2 = P3 = 0$

while( $P1 < N \&\& P2 < N \&\& P3 < N$ ) {

}

Tc:  $O(N+N+N) \Rightarrow O(N)$

Sc:  $O(1)$

← \*