- 1. Good Evening
- 2. We will begin at 9:07 pm
- 3. Topic Synchronization

Agenda

- 1. Problem of Synchronization
 - 1. Adder-Subtractor problem
 - 2. What is problem of synchronization?
- > 3. Conditions that lead to synchronization problems.
- 2. Solution to synchronization problems
 - 1. A theoretical solution; Features.
 - 2. Practical Solutions
 - 1. Muter Locks V
 - 2. Synchronized keyword V
 - 3. Atomic variables V
 - 4. Semaphores: Producer-Consumer Problem

Adder - Subtractor Problem

loop

Count:0

Loop

Count = i

Subtractor

moin Count c Thread to = res thread (a) Thread t2 = new Thread(s) ti. startc). t. join() t(. finish() 14 / t2. join() () t2. finish() cl fl C2 _ +2

Payister Adder [count +1]

Payister Subtractor [count -]

1.
$$X^{0} - count$$

2. $X^{0} = X + 1$

3. $X^{0} = count$

4. $X^{-1} = X - 1$

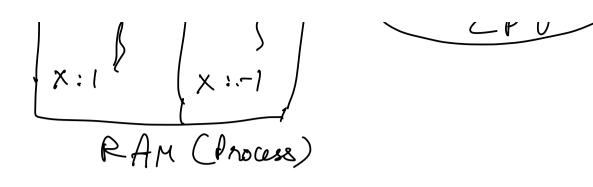
5. $X^{1} \rightarrow count$

6. $X \rightarrow count$

PC ti | PC ti

3. $X^{0} = count$

6. $X \rightarrow count$



Problem of synchronization: If multiple threads will work on shared data, at the same time, we might get inconsistent/ wrong results because of CPU Scheduling.

When will synchronization problems happen?

Or the conditions that lead to synchronization

problems?

O. Multiple threads & shored?

date.

1. Critical Section: Part of code working on shared piece of data.

Adder

1. paint hi

2. Count += 1

3. paint bye

Count

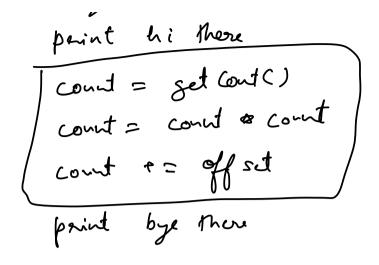
Subfractor

1 print hello

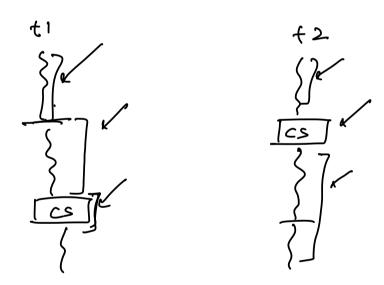
2. Count -= 1

3. print see you.

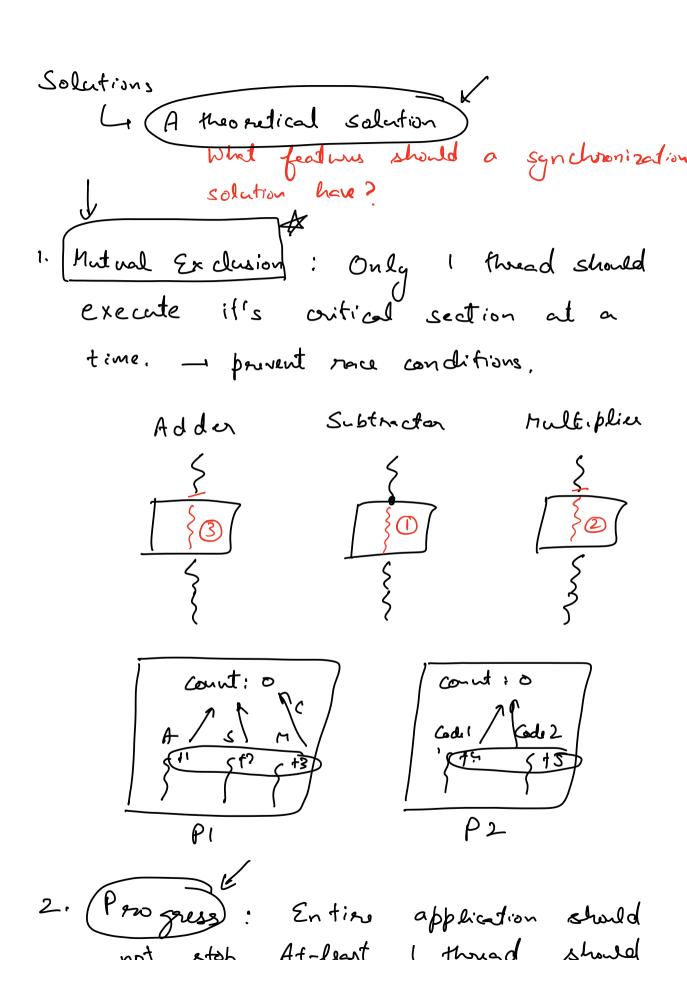
Squarer



2. Race Condition: More than I thread trying to enter critical section at the same time.



3. Pre-emption! Of a thread executing it's critical section sets "content-switched" out we may have synchronization problems.



be making progress. 3. Bounded Waiting - No thread should vait indéfinitely to enter 11's critical soction. Threads should be given access to critical section in FCFS manney of request. to should not access critical section before t1, +2 > +5 4. No busy waiting: Busy waiting code should not be there.

while (! allowedentry to critical section)? Notification mechanism instead of busy vaiting. Practical Salutions Practical un.

Muter Locks Synchronized keyword

Practical Solutions

1. Mutex Lock:

Adder

Print hi

Lock. Lock()

Conut +
Lock, unlock()

Print bye

Count & Key

Port

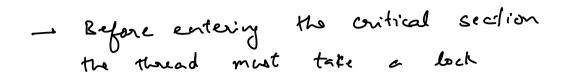
Subtractor

Print hello

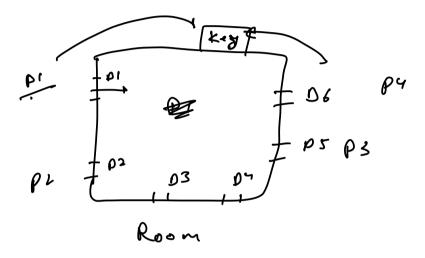
lock. lock()

Count -
lock. unlock()

Print see you



while leaving the oritical Section though should release the lock



Lock on entire for Loop

(+1 +2 +3 +4 -... +10020 /-1 -2 -3 -..-10020)

7 T T T

2. Synchronized keyword

A Mater is an OS concept

Java specific kegward

In Java locks are available on all objects.

Adder

Prin hi

Synchronized (courd) ?

Count * 1

Paint Byo

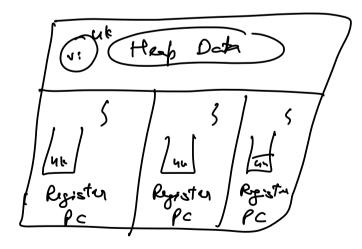
Subtractor

Paint. hello c

Synchronized (cound) ?>

Count -
Paint See you.

A Threads can show data on the heap.



Volatile

Synchronized Methods
L 2 Perspectives
Duse an API developed You have to develop the API used by Other people.
You want to make some that your API/DS will work perfectly in multi-threaded applications
methods synchronized.
Count Value X increment Value() Set Value ()

Synchronized Method

- of a thread enter a synchronized methon on an object, no other thread

can enter (any syriumonized
memods on the same object.
Connt SincValue() of Ci, C2 c dic Value() of Set value() of (1) 42 Will they would for
Cl. inc Value () Cl. inc Value () Cl. de Value () Cl. get Value () Cl. get Value () Cl. inc Value () Cl. jet Value ()
V [Hashmap] - Take locks owself
Hashfable - Take locks owself Hashfable - To we have synchronized of methods
String Ruffer Vs (String Ruffer)

J Synchronized not synchronized Size-Threaded code - Use Hashmep & String Builder Mull: -- Use Hashfable & 11 String Ruffor. Agenda Producer - Consumer problem 2. Semaphores Atomic Vaniables 4. Dead Locks. Memory Management main t1. stad()

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