## Inorder:

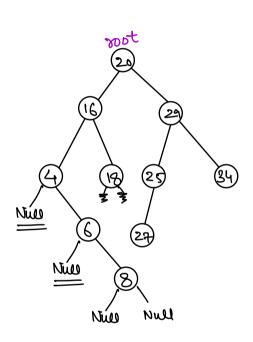
Left Root Right

void imorder (root) {

if (root == Nul)

veture;

- 2) Im Order (root. lyt);
- 3) Print (root data)
- 3 4) Imorder (root right),



4,6,8,16,18,20,27,25,29,34

34: X,Z,X,X 27: X,Z,X,X 25: X,Z,X,X 29: X,Z,X,X 18: X,Z,X,X 6: X,Z,X,X 4: X,Z,X,X 16: X,Z,X,X 20: X,Z,X,X

Call Stack

## Idea:

1) Till you get a Null on the left side, keep fushing into the Stack. 2) If root == Null, get the top element from Stack & go to right.

```
Void imOrder (xoot) {

Stack (Node) St;

Node aux = xoot;

while (aux!=Nnll||St·size(7>0) {

if (aux!=Nnll){

St·push (aux);

aux = cux·lyt;

clse {

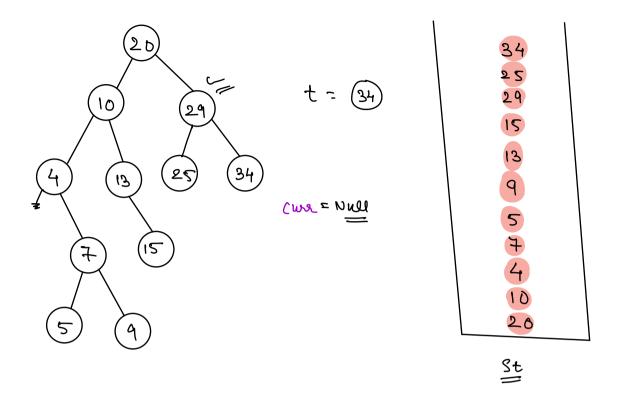
Node temp = St·top(7);

St·pop(5);

print (temp·data);

aux = temp·right;

}
```



4,5,7,9,10,18,15,20,25,29,34

TC: O(N)

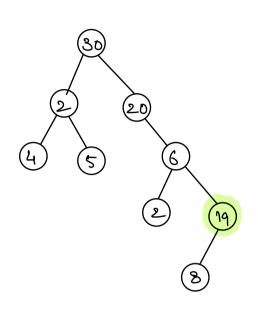
SC: 0(4)

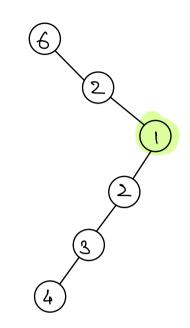
1- wc: 0(n)

HW

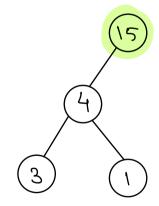
- 1) Preorder Iterative
- 2) Rostorder Iterative.

## Q: Given a tree, find the last in order node that gets printed.

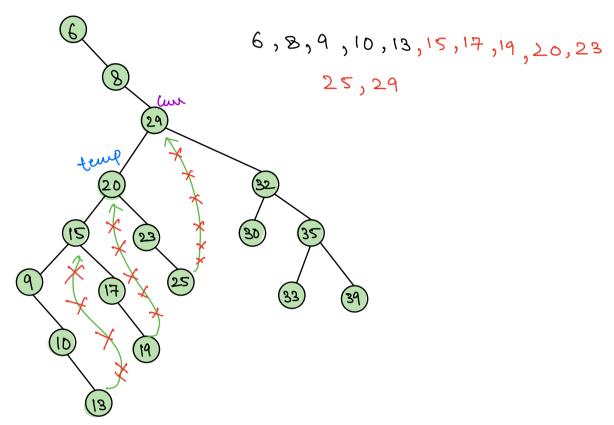




- => Keep gring ou light side until me get Null.
- → light most node.



De can me do Inorder transcral in O(1) Sc.



→ If left child is NULL, frient the over mode?

fo to right side.

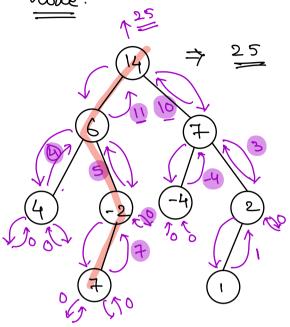
```
Void
        inorder ( root) 1
       Node cur = root;
       while ( cur != Null) {
            if ( cur. left = = NULL) {
                     frint (curr.data);
                     Curr = curr right;
             else (
                   Node temp = cur left;
                   mhile (temp. right != Null 28
                                 temp. right 1= curr) {
                         temp = temp.right;
                  if (temp. right == Null) (

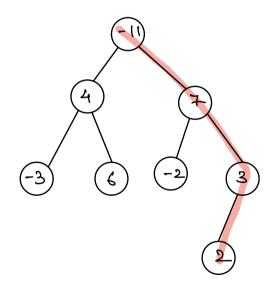
"Visiting our node for 1st time.

temp. right = curr;

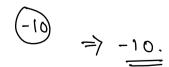
curr = curr. left;
                    else { // VIsiting our Node for 2nd time
temp. right = Null
                          Print ( wer data);
                           Cur = cur. right;
                     z
→ Morris Inorder Traversal.
                  SC: 0(1)
```

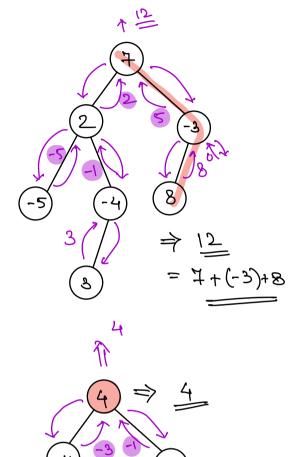
D' find the max sum fath starting from root node.





$$\Rightarrow$$
  $\bot$ 





Single node can also be an ans.

int max Path Sum (root) i

if (root == Null)

Leturn 0;

int l = max Path Sum (root · left);

int x = max Path Sum (root · right);

return root · data + max (l, x);

max (l, x, 0)

max (max (l, x), 0)

TC:O(N) SC:O(N)

# Max Path Sum Containing root ude.

