

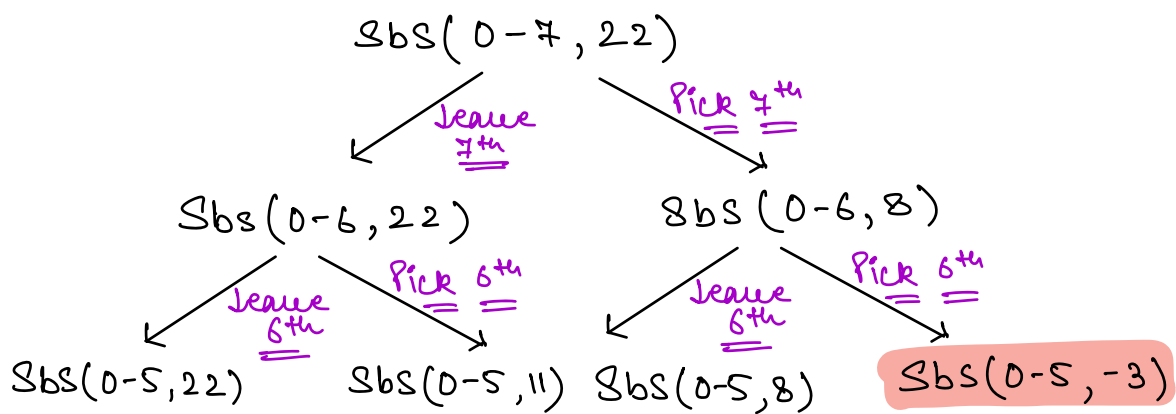
Q: Given  $arr[N]$ , find the no. of subsets with  $Sum = k$ . Array elements  $> 0$  &  $k > 0$

$arr[8] : \{7, 4, 9, 6, 10, 13, 11, 14\} \quad k = 22$

$\left. \begin{array}{l} \{7, 4, 11\} \\ \{9, 13\} \\ \{7, 9, 6\} \end{array} \right\} \underline{\underline{3}}$

Idea :- Generate all possible subsets & check if  $Sum == k$ .

Bit Masking  $\Rightarrow TC: O(N \cdot 2^N)$   
 Backtracking  $\Rightarrow TC: O(2^N)$



Ex :-

							6	10	14	2	8	
	0	1	2	3	4	5	6	7	8	9	10	$k = 45$
6 elements, $s = 25$ :							✓	*	✓	*	*	$\Rightarrow 20$
6 elements, $s = 25$ :							x	✓	x	✓	✓	$\Rightarrow 20$

→ Overlapping Subproblems. } DP.  
 → Optimal Substructure.

DP State

$dp(i, j)$  : # No. of subsets using elements from  $[0, i]$  & sum =  $j$

$$* \quad dp(i, j) = \left\{ \begin{array}{ll} \text{leave } i^{\text{th}} & \text{pick } i^{\text{th}} \\ dp(i-1, j) & + \quad dp(i-1, j - a[i]) \end{array} \right\}$$

$j - a[i] \geq 0$   
 $j \geq a[i]$

\* dp Table

No. of subsets using elements from  $[0 \text{ to } N-1]$   
 with Sum =  $K$   
 $\rightarrow$  ans:  $dp[N-1][K]$

Size :-  $dp[N][K+1]$

\* Base Case :-

$$dp(i, j) = dp(i-1, j) + dp(i-1, j - a[i])$$

$j \geq a[i]$

$\rightarrow i=0$  : Base Case

A[4]: { 4, 2, 1, 3 }  $k=4$

dp[4][8]

	0	1	2	3	4	5	6	7
0	1	0	0	0	1	0	0	0
1								
2								
3					i,j			

dp[0][0]: ①

No. of subsets using indices [0-0] with sum=0.

{4} → {3}

dp[0][1] = No. of subsets using indices [0-0] with sum=1.

dp[0][4]: 1

for (j=0; j ≤ k; j++)

dp[0][j] = 0

dp[0][0] = 1

if (A[0] ≤ k) { dp[0][A[0]] = 1 }

Base Case

\*

dp[1][0]: No. of subsets using indices [0-1] with sum=0.

⇒ {4, 2}

{3} ⇒ sum=0

\*  
Σn

{ 4 0 1 }

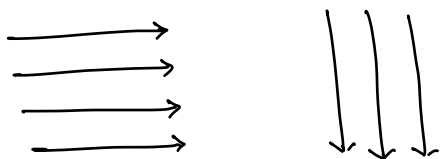
No. of subsets with sum = 0 ⇒ 2. {3}, {0}

Sum {4, 0, 1, 0}

No. of subsets with sum = 0  $\Rightarrow 2^2$

{ }  
{0, 1} {0, 0}  
{0, 0}

Code :-



```
for (i = 1; i < N; i++) {  
    for (j = 0; j <= K; j++) {  
        dp[i][j] = dp[i-1][j]  
        if (j >= a[i])  
            dp[i][j] += dp[i-1][j-a[i]]  
    }  
}
```

return dp[N-1][K]

TC:  $O(N \cdot K)$ , SC:  $O(\underline{N \cdot K})$

Example

A[4]: {4, 2, 1, 3} K=7

dp[4][8]

	0	1	2	3	4	5	6	7
0	1	0	0	0	1	0	0	0
1								
2								
3								

## \* Space Optimization

$$dp[2][K] : TC : O(N \cdot K), SC : \underline{\underline{O(K)}}$$

Q: Given  $N$  array elements, find the length of smallest subset with sum =  $K$ .

$$\begin{array}{cccccccc} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ arr[8] : \{ 7, 4, 9, 6, 10, 13, 11, 14 \} & K=22 \\ \Rightarrow \{ 13, 9 \} \Rightarrow \underline{\underline{2}} \end{array}$$

$dp[i, j]$  : length of smallest subset with  $sum = j$  using the elements from  $[0-i]$

$$dp[i, j] : \underline{\underline{\min}} \left\{ \begin{array}{l} dp[i-1, j], \quad dp[i-1, j-A[i]] + 1 \\ \text{leave } i^{th} \quad \text{pick } i^{th} \end{array} \right\}$$

$i^{th}$   
 $\uparrow$  index

$j \geq A[i]$

\* DP table.

$$dp[N][K+1]$$

\* Base Case :-  $i = 0$

Ex       $A[4] = \{4, 3, 2, 6, 3\}$        $K = 7$

$dp[4][8]$

	0	1	2	3	4	5	6	7
0	0	5	5	5	1	5	5	5
1								2
2								
3								

$dp[0][0] : 0$

length of the smallest subset with sum = 0 from  $[0-0]$ .

$\{4, 3\} \rightarrow \{3\}$  len = 0.

$dp[0][1] :$  length of the smallest subset with sum = 1 from  $[0-0]$ .

$dp[0][4] : 1$

$dp[0][5] :$

$$dp[1][7] = \min(dp[0][7], dp[0][7 - \frac{3}{4}] + 1)$$

$$dp[1][7] = \min(5, 6) = \underline{5}$$

\* { for ( $j = 0; j \leq K; j++$ )  
      $dp[0][j] = N+1;$   $\rightarrow$  No subset.  
      $dp[0][0] = 0$  // Empty subset  
     if ( $a[i] \leq K$ ) {  
          $dp[0][a[i]] = 1$   
     }  
     3

```

for (i = 1; i < N; i++) {
    for (j = 0; j <= K; j++) {
        dp[i][j] = dp[i-1][j]
        if (j >= a[i])
            dp[i][j] = min(dp[i][j], dp[i-1][j-a[i]]+1);
    }
}
return dp[N-1][K]

```

TC :  $O(N \cdot K)$

SC :  $O(N \cdot K) \longrightarrow O(K) \{ \underline{dp[2][K+1]} \}$

Q: Check if there exist a subset with sum = K.

$dp[i][j]$ : If there is a subset with sum = j using elements  $[0-i]$ .

$dp[i][j]$ :  $dp[i-1][j] \parallel dp[i-1][j - \underbrace{a[i]}_{j \geq a[i]}]$

DP Table :  $dp[N][K+1]$

Base Case  $\underline{i = 0}$ .

$A[4]: \{4, 3, 2, 6, 3\} \quad K=7$

$dp[4][8]$

	0	1	2	3	4	5	6	7
0	T	F	F	F	T	F	F	F
1	T							
2	T							
3	T							

$dp[0][0]: \{4\} \leftarrow \{3\}_{sum=0}$   
 $\underline{\underline{\{4\}}}$

$dp[0][1]: F$

$dp[0][4]:$

for ( $j = 0; j \leq K; j++$ ) {  
 $dp[0][j] = false;$

$dp[0][0] = true;$

if ( $a[i] \leq K$ ) {

$dp[0][a[i]] = true;$

}

TC:  $O(N \cdot K)$

SC:  $O(N \cdot K)$

$\rightarrow dp[2][K+1]$

SC:  $O(K)$

✖