1. Good Evening
2. Let's begin at 9:10 pm
3. Topic - Splitwise Algorithm

Agenda

1. Settlement algorithm for splitwise → 2
2. Undo algorithm for TTT → 3

Next class

1. Command Design Pattern] 1
2. Restful API] 1
3. One data flow end to end] 1
4. O(1) WS from TTT] 20

Settlement Algo

1. Pull Reqd. data ✓ → EPRepo, ESRepo
2. Transform → Make balance HM
3. Settlement Algo: Round trip, Greedy

Users → A, B, C, D, E, F ✓

Groups → Hostel, Office ✓

→ Hostel: A*, B*$, C, D ✓ →

→ Office: A, B, E*, F*$

User

id	name	phn	pwd
1	A		
2	B		
3	C		
4	D		
5	E		
6	F		

Group

id	name	cb
1	Hostel	2
2	Ofc	6

Group Participants

id	gid	uid	isAdmin
1	1	1	✓
2	1	2	✓
3	1	3	X
4	1	4	X
5	2	1	X
6	2	2	X
7	2	5	✓
8	2	6	✓

Hostel

✓ → Dinner ✓ 19/2/2023

Amount : 1000 ✓

PB : A (500) ✓, B (500) ✓

SB : A(300) ✓, B(300) ✓, C (200) ✓, D(200) ✓

✓ → New Year ✓ 01/01/2023

Amount : 2000 ✓

PB : D (2000)

SB : A (500), B (500), C (500), D(500)

PB = B(3000), E (2000), F(1000)

Q.c.

Trip 10/1/2023

Amount: 4000

PB : B(2000), E(2000)

SB : A(1000), B(1000), E(1000), F(1000)

Farewell 31/1/2023

Amount : 2000

PB : B(1000), F(1000)

SB : A(500), B(500), E(500), F(500)

Booze 2/2/2023

Amount : 3000

PB : A(1500), B(1500)

SB : A(1000), B(1000), C(1000)

Expense

id	title	amount	date
1	Dinner	1000	
2	New Year	2000	
3	Trip →	4000	
4	Farewell →	2000	
5	Booze	3000	

Group Expense

id	g_id	e_id
1	1	1
2	1	2
3	2	3
4	2	4

EP

id	uid	cid	amount
1	1	1	500
2	2	1	500
3	4	2	2000
4	2	3	2000
5	5	3	2000
6	6	4	1000
7	1	5	1500
8	2	5	1500
9	2	4	1000

ES

id	uid	eid	amount
1	1	1	300
2	2	1	200
3	3	1	200
4	4	1	200
5	1	2	500
6	2	2	500
7	3	2	500
8	4	2	500
9	1	3	1000
10	2	3	1000
11	5	3	1000
12	6	3	1000
13	1	4	500
14	2	4	500
15	5	4	500
16	6	4	500
17	1	5	1000
18	2	5	1000
19	3	5	1000

Ofc →

PB = { B(2000), E(2000), F(1000) , B(1000) }

B(3000), E(2000), F(1000)

```
Select   ep.usrid, SUM(ep.amount)
FROM     Group g
JOIN     GroupExpense ge   ON g.gid = ge.gid
JOIN     Expense e         on ge.eid = e.eid
JOIN     ExpensePaidBy ep  ON e.eid = ep.eid
GROUP BY ep.usrid
Where    g.gid = 2
```

Ofc

PB = B (3000), E (2000), F (1000)

SB = A (1500), B (1500), E (1500), F (1500)

A → +0 − 1500 = −1500

B → +3000 − 1500 = +1500

E → +2000 − 1500 = +500

F → +1000 − 1500 = −500

A = - 1500

B = + 1500

E = + 500

F = - 500

Settle a group → GC → GS → EPR, ESR

L<EP>

L<ES>

1. Expense Paid By Repository → Retrieve List of Expense Paid By using group id

2. ExpenseShouldBy Repo → Retrieve List of ExpenseShould By using group id

3. Make balance hashmap

List<EP>

B	Temp	+ 2000
E	Temp	+ 2000
F	FW	+ 1000
B	FW	+ 1000

List<ES>

A	Temp	+ 1000
B	"	+ 1000
E	"	+ 1000
F	"	+ 1000
A	FW	+ 500
B	"	+ 500
E	"	+ 500
F	"	+ 500

B → #2000 +1000 -1000 = 500

E → #2000 - 1000 - 500

F → #1000 - 1000 - 500

A → -1000 - 500

B → + 1500

E → + 500

F → - 500

A → - 1500

\+ → You paid more than your real share

\- → You paid less than your real share

1. Getting data
2. Making Bal Hashmap
3. Settle.

A → +200

B → +500

C → -300

D → -200

E → -100

F → -100

A → ~~+200~~

B → ~~+100~~

C → ~~-200~~ -100 ⊙

D → ~~-100~~

1. C 200 → A.

2. D 100 → B

+ → Paid extra

− → Paid Less

Payments should be from −ve to +ve

For a transaction add to payer

& subtract from receiver.

Algorithm 1 (Round trip)

A → +200
B → +100
C → -200
D → -100

Guarantee → For n users it will always produce n-1 transactions.

A +200 ★ B +100 C -200 D -100

1. B → 200 → A

A +200-200 B +100+200 C -200 D -100

A 0 B +300 C -200 D -100

2. C → 300 → B

A 0 B +300-300 C -200+300 D -100

A 0 B 0 C +100 ✓ D -100

3. D → 100 → C

A 0 B 0 C 0+100=100 D 0-100=-100

A → +200
B → +500
C → -100
D → -300
E → -200
F → -100

5 transactions

+ $Receive
− Pay

A 200	B 500	C -100	D -300	E -200	F -100
-200	+200	—	—	—	—
0	+700	-100	-300	-200	-100
✓	-200				
0	0	600	-300	-200	-100
		0	300	-200	-100
			0	+100	-100
				0	0

1. B → 200 → A
2. C → 700 → B
3. D → 600 → C
4. E → 300 → D
5. F → 100 → E

n users → n-1 transactions

O(n)

A → +~~200~~ 0 ✓

B → +~~100~~

C → -~~100~~

D → -~~200~~ 0

D $\xrightarrow{200}$ A

C $\xrightarrow{100}$ B

A → -100 + 100 = 0

B → -200 -100 = -300 + 300 = 0

C → +100 - 300 = -200 + 200 = 0

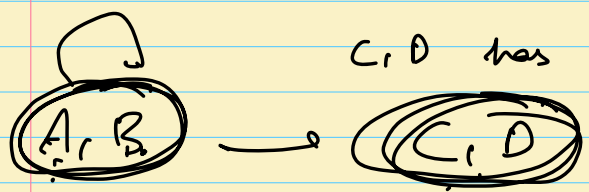
D → +200 - 200 = 0

1. A $\xrightarrow{100}$ B ✱
2. B $\xrightarrow{300}$ C
3. C $\xrightarrow{200}$ D

A -100
B -200
C +100
D +200

A, B has paid less

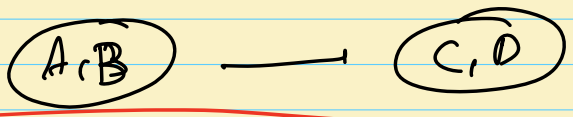
C, D has paid more



Less payers More payers

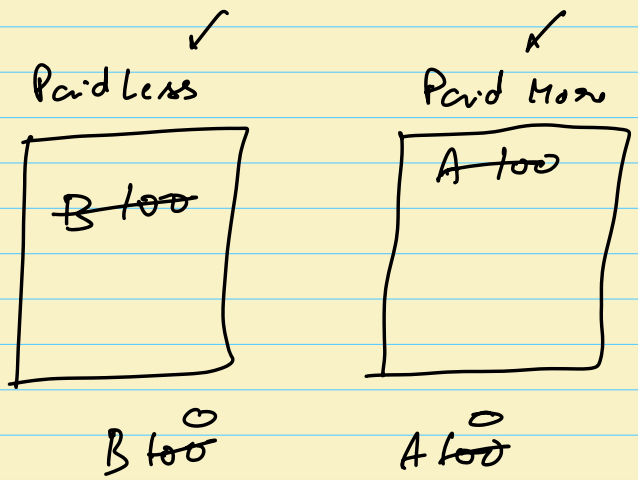
Round trip

1. Guaranteed n-1 transactions.
2. Problem = Not smart



Break = 10:24 to 10:34

- ✓ A → +500
- ✓ B → -200
- ✓ C → +600]
- ✓ D → -400
- ✓ E → -300
- ✓ F → -500
- ✓ G → +400
- ✓ H → -400
- ✓ I → +300



1. F —500→ C
2. H —400→ A
3. D —400→ G
4. E —300→ I
5. B —100→ C
6. B —100→ A

1. Add -ves to paid Less PQ [take abs of value]

2. Add +ves to paid more PQ

3. While PL PQ & PM PQ are not empty.

   a) Remove x from pl PQ

   b) " y " pm PQ

   c) Amount = Min (x.val, y.val)

   d) Trans = $x \xrightarrow{\text{amount}} y$

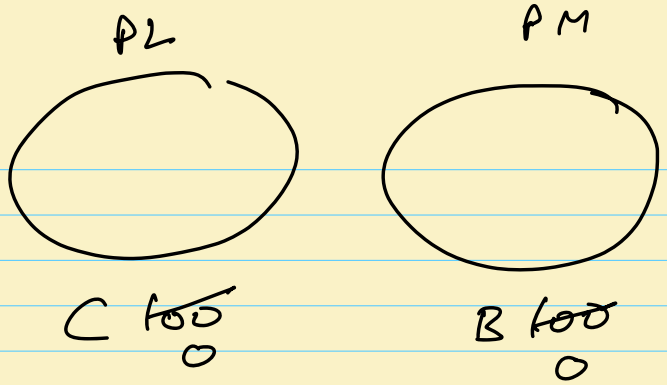
   e. Reduce x.val, by amount
      y.val

   f. If x.val != 0 add back to paid Less PQ

   g. If y.val != 0 add back to paid more PQ.

example

A −200
B +100
C −100
D +200



1. A —200→ D
2. C —100→ B

Roundtrip ✓

1. n-1 transactions
2. O(n) algo]
3. Non-smart

Greedy Algo ✓

1. ≤ n-1 transactions
2. nlogn algo.]
3. Smart { All trans are from PL to pm Sidi }

Does the greedy algo always give minimum transactions?

Not necessarily.

The algorithm that guarantees minimum transaction is NP hard [2^n time complexity]

1. ~~Round trip~~ ✓ → O(n) [n-1 transactions]
2. ~~Greedy Algo~~ ✓ → nlogn [smart enough but no guarantees]
3. ~~Backtracking Algo~~ → Guarantees minimum transaction [2^n]

NP Hard

$n = 35$ $\quad 2^{35} = 10^9 \cdot 2^6$

$= 10^9 \cdot 64$

Next class

1. Command Design Pattern
2. Restful API
3. End to End Data Flow
4. TTT Discussion.

2.5 hours.