Q.1 Given an Array of size N. Find the K^{th} minimum element in the Array. $K < \log_2 N$

A: { 1, 5, -1, 2, 10, 3 }

K = 3 ⇒ 2

K = 5 ⇒ 5

K = 1 ⇒ -1

Brute force

1. Sort (Arr)
2. return Arr[K-1]

A: { 1, 5, -1, 2, 10, 3 }

↓ sort

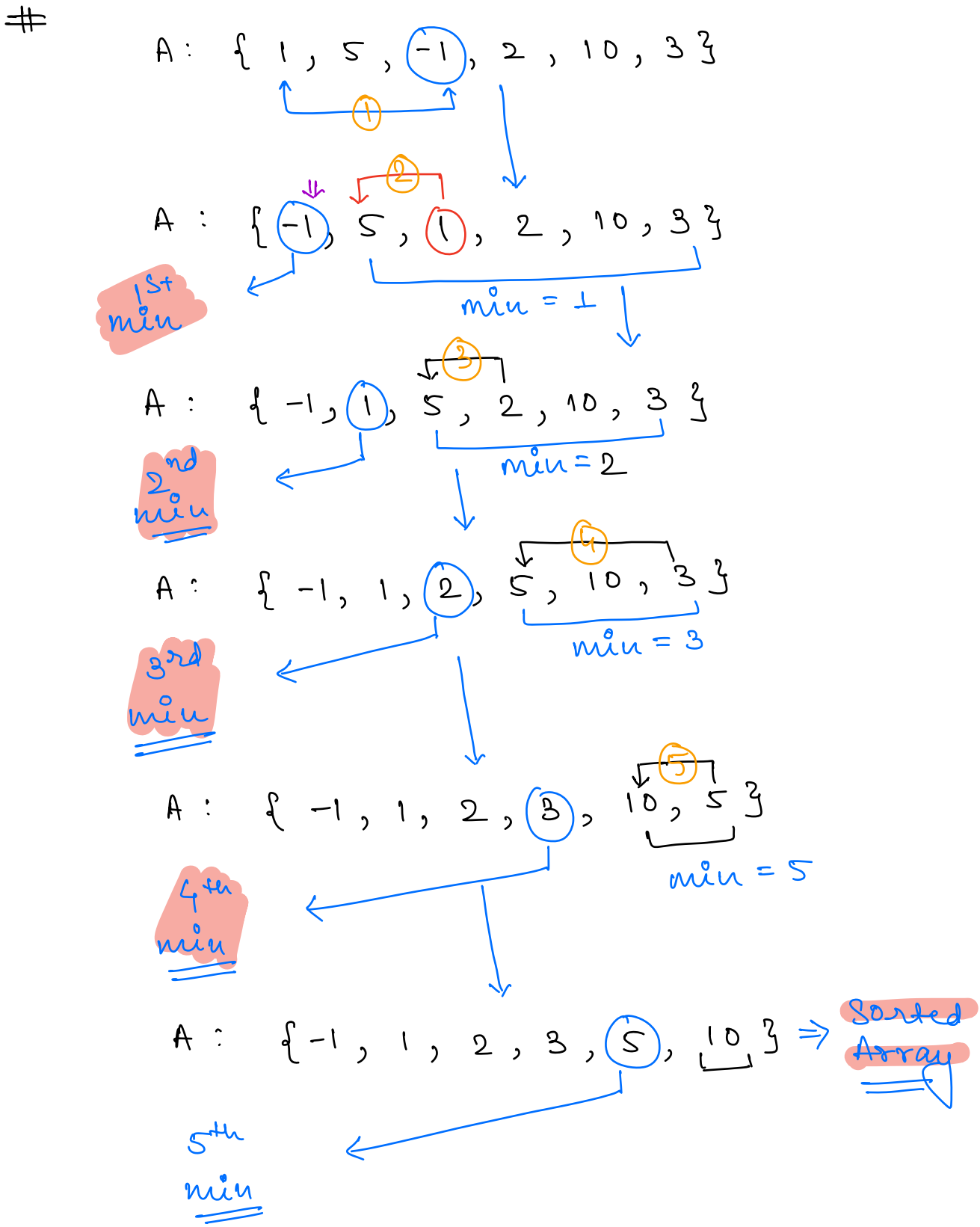
0	1	2	3	4	5
-1	1	2	3	5	10

K = 3 ⇒ A[2]

TC: O(N log N)

SC: Depends on sorting algo.

#


A : { 1 , 5 , -1 , 2 , 10 , 3 }
1
A : { -1 , 5 , 1 , 2 , 10 , 3 }
2
1st min
min = 1
A : { -1 , 1 , 5 , 2 , 10 , 3 }
3
2nd min
min = 2
A : { -1 , 1 , 2 , 5 , 10 , 3 }
4
3rd min
min = 3
A : { -1 , 1 , 2 , 3 , 10 , 5 }
5
4th min
min = 5
A : { -1 , 1 , 2 , 3 , 5 , 10 } ⇒ Sorted Array
5th min

TC of finding the Min(Arr) = O(N)

TC: O(KN) < O(NlogN)

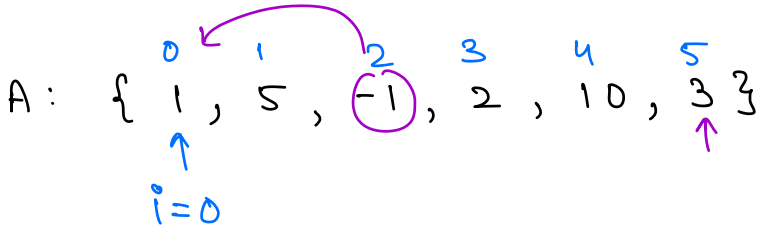
SC: O(1)

Selection Sort

Selecting the minimum element & placing this element at its right position is Selection Sort.

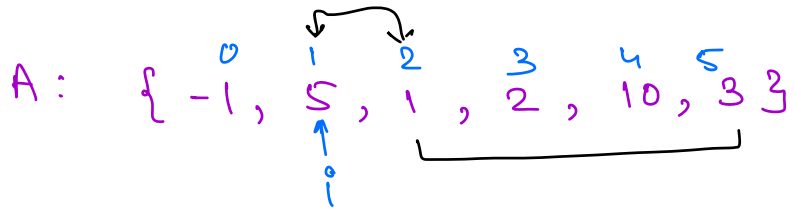
TC: $O(N^2)$

```
for (i = 0; i < N; i++) {
    min = A[i]
    min_index = i;
    for (j = i+1; j < N; j++) {
        if (A[j] < min) {
            min = A[j];
            min_index = j;
        }
    }
    swap(A[i], A[min_index]);
}
```

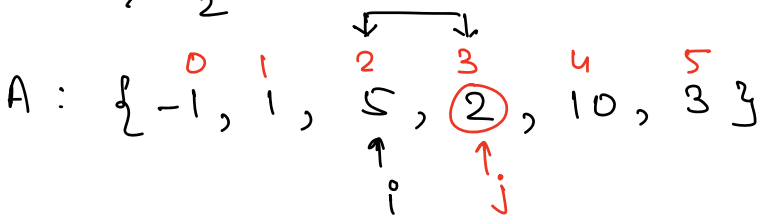


min = ~~1~~ -1
m_i = ~~0~~ 2

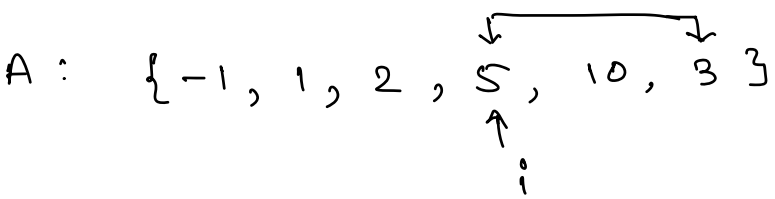
$\Rightarrow$ swap(A[0], A[2])



min = ~~5~~ 1
m_i = ~~0~~ 2



min = ~~5~~ 2
m_i = ~~2~~ 3



min = ~~5~~ 3
m_i = ~~2~~ 5

A: { -1 , 1 , 2 , 3 , 10 , 5 }
i j

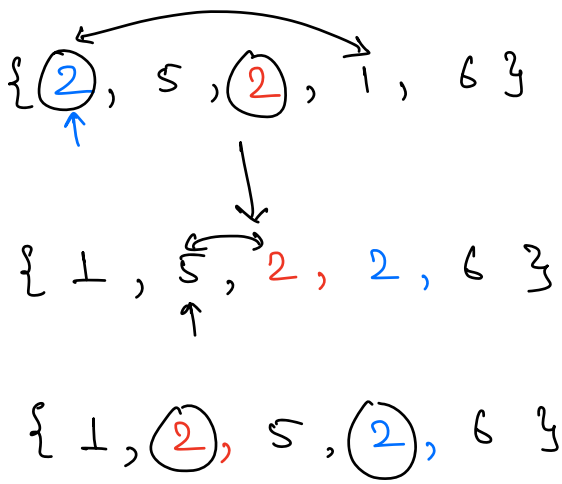
min = ~~10~~ 5
m_i = ~~4~~ 5

A : {-1, 1, 2, 3, 5, 10}

TC : $O(N^2)$

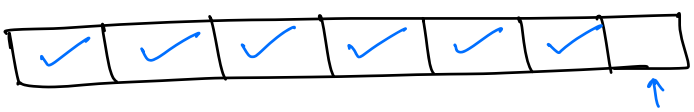
Quiz Inplace? ⇒ YES : SC = O(1)

Stable ⇒ NO

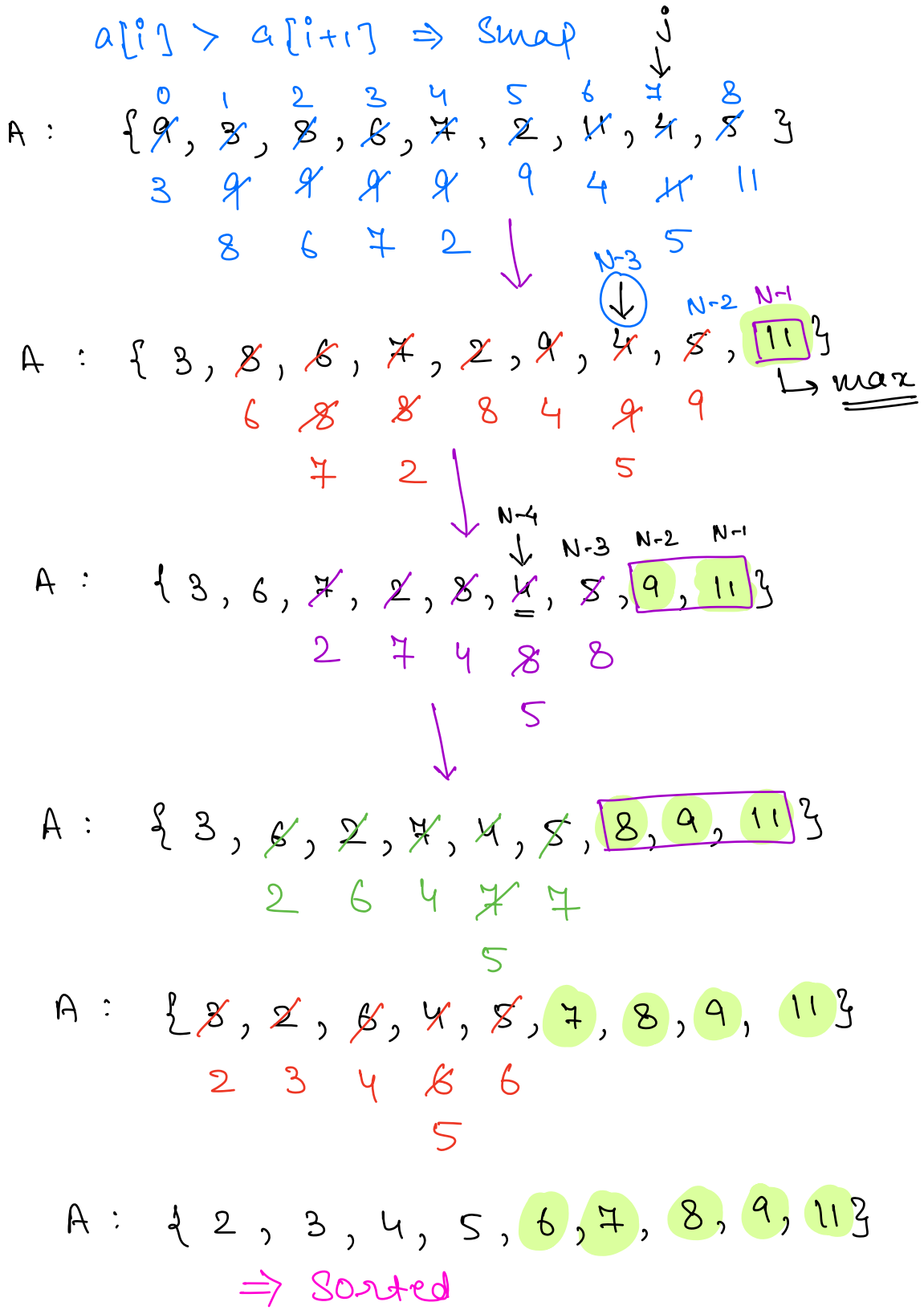


Quiz Max. no. of swaps in selection sort :-

= N-1



Constraint :- Swapping is allowed only b/w consecutive elements.



Keep on taking the Max element at the end of the Array.

⇒ Bubble Sort

Code

```
for( i=0; i<N; i++ ){
    for( j=0; j<N-1; j++ ){
        if( A[j] > A[j+1] ){
            swap(A[j], A[j+1]);
        }
    }
}
```

$j <= N-i-2$

$j < N-i-1$

		Swap
$i = 0 \Rightarrow j \in [0, N-2]$	$N-0-2$	$N-1$
$i = 1 \Rightarrow j \in [0, N-3]$	$N-1-2$	$+ N-2$
$i = 2 \Rightarrow j \in [0, N-4]$	$N-2-2$	$+ N-3$
$\vdots$		$+ \cdots$
$i \Rightarrow j \in [0, N-i-2]$		$+ 1 + 0$

$0 =< j <= N-i-2$

* If at any iteration, swap count is zero then it means Array has already become Sorted.

TC: $O(N^2)$

Quiz Max. no. of swaps in Bubble Sort.

Swaps = $1 + 2 + 3 + \text{-----} + (N-1)$

$= \frac{N(N-1)}{2}$

Inplace ? YES | SC: O(1)

Stable ? YES

A: { 2, 1, 2, 2 }
1 2 2 5

A: { 1, 2, 2, 5 }

Q: Given 2 sorted arrays of size N & M. Merge these 2 sorted arrays into 1 sorted array. Amazon

A: { 2, 5, 7, 12, 20, 24, 29 }

B: { 6, 9, 10, 14, 18, 19 }

C: { 2, 5, 6, 7, 9, 10, 12, 14, 18, 19, 20, 24, 29 }

Approach#1

⇒ int C[N+M]

⇒ Put all the elements of A & B into C.

⇒ Sort(C) ⇒ O((N+M) log(N+M)) ↳ N+M

TC: O((N+M) log(N+M))

SC: O(N+M)

Approach#2

A: {2, 5, 7, 12, 20, 24, 29} ⇒ Sorted
   ↑i

B: {6, 9, 10, 14, 18, 19} ⇒ Sorted.
   ↑j

C: {2, 5, 6, 7, 9, 10, 12, 14, 18, 19, 20, 24, 29}
   0 1 2
   ↑
(N+M)

C[0] = min(A[0], B[0])

C[k] = min(A[i], B[j])

```
merge (int A[], int M, int B[], int N){
    int C[N+M]
    i=0 // Array A
    j=0 // Array B
    K=0 // Array C
    while( i< M && j < N){
        if(A[i] < B[j]){
            C[K] = A[i];
            i++
            K++
        }
        else{
            C[K] = B[j]
            j++
            K++
        }
    }
    while( i< M){
        C[K] = A[i];
        i++
        K++
    }
    while( j< N){
        C[K] = B[j];
        j++;
        K++;
    }
    return C;
}
```

TC : $O(N+M)$

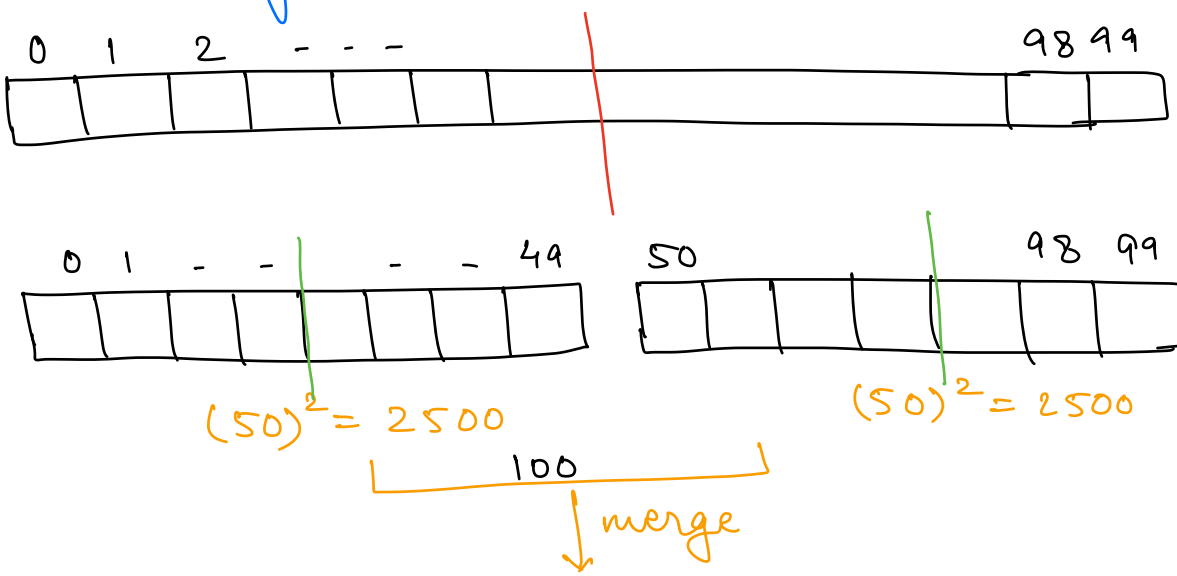
SC : $O(N+M)$ (Output Array)

Old machine

Task :- Sort an Array of size $N = 100$.

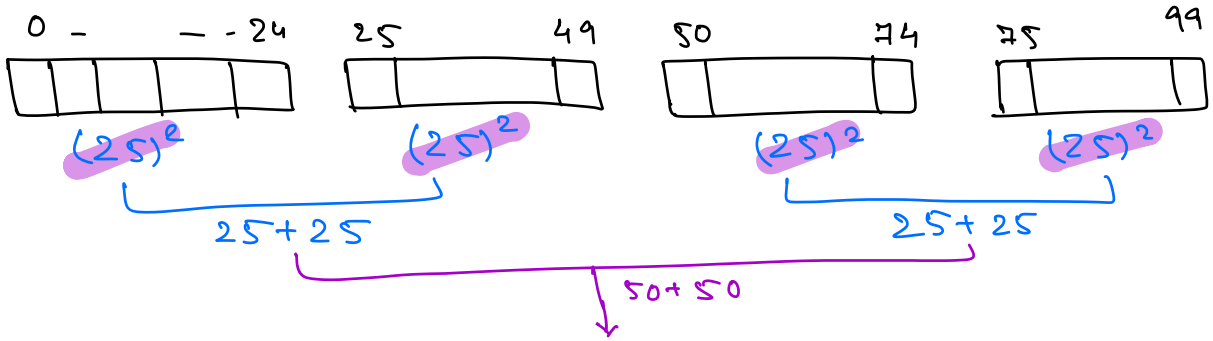
Bubble Sort, Selection Sort $> O(N^2)$

# of iterations $= (100)^2 = 10000$



# of iterations $= 2500 + 2500 + 100$

$= 5100$



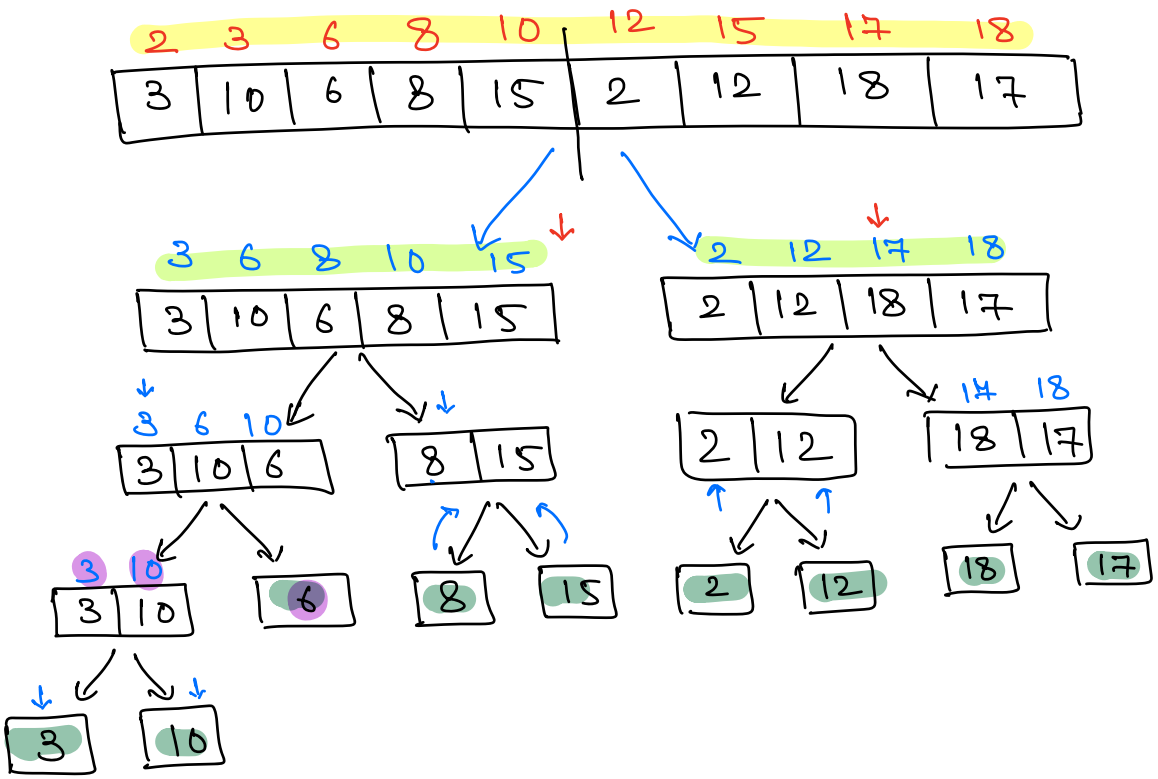
# of Iterations $= (25)^2 * 4 + 25*4 + 50 + 50$

$= 2500 + 100 + 100$

$= 2700$

$10K \longrightarrow 5K \longrightarrow 2.5K \longrightarrow \cdots$

MERGE SORT (Divide & Conquer)



Code :-

Assumption :- mergeSort (A, l, r) sorts the Array from l to r.

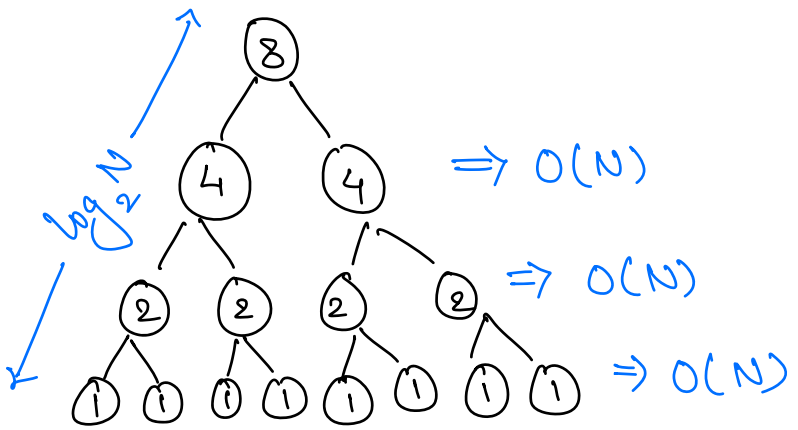
```
mergeSort ( int A[], int l, int r) {
    if ( l == r)
        return;
    int mid = (l+r) / 2 ;
    mergeSort (A, l, mid) ;
    mergeSort ( A, mid+1, r) ;
    merge ( A, l, mid, r)
}
```

Recurrence Relation :-

HW $T(N) = 2T(N/2) + N$

TC : $O(N \log N)$

Recursion Tree



TC: $O(N * \log_2 N)$

HW

SC: O(N)
