

Projected Gradient Descent Efficiently* Solves the Trust Region Subproblem

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Abstract

We show that projected gradient descent asymptotically converges to a global minimizer of the trust region subproblem. We remark on next steps in this project at the end.

1 Introduction

Trust region methods are sequential programming procedures in which heuristics are used to approximately solve a general optimization problem through multiple constrained quadratic programs. As a subroutine, these methods formulate and solve many instances of the following *trust region subproblem*

$$\begin{aligned} &\text{minimize} && (1/2)x^T A x + b^T x \\ &\text{subject to} && \|x\| \leq R \end{aligned} \tag{1}$$

with variable $x \in \mathbf{R}^n$. The problem data are a symmetric matrix $A \in \mathbf{R}^{n \times n}$, a vector $b \in \mathbf{R}^n$, and a radius parameter $R > 0$. Crucially, the matrix A is possibly indefinite.

1.1 Previous works

The trust region subproblem is well-studied, and thus there many previous works worth mentioning. In earlier papers, the problem was solved either via subspace methods such as Steihaug-Toint (where no global convergence guarantees have been proven, to our knowledge), or using fast eigenvector and eigenvalue computation procedures like the Lanczos method [CGT00, EG09, GLRT99, GRT10]. More recently, however, some authors have provided convergence guarantees for this problem. For example, by reducing the trust region subproblem to a sequence of approximate eigenvector computations, Hazan and Koren [HK16] demonstrate that $\tilde{O}(1/\sqrt{\varepsilon})^1$ matrix-vector multiplies are enough to guarantee an ε -suboptimal point. In [HK17], Nguyen and Kiling-Karzan reduce the trust region problem to a convex QCQP using eigenvector calculations, where first-order methods apply.

However, perhaps the most obvious algorithm to solve (1), is the *projected gradient method*, which we study in this paper. To our knowledge, the only previous work that analyzes the convergence properties of this procedure on (1) is [TA98], where Tao and An augment this procedure by a restarting scheme, requiring possibly $O(d)$ restarts, which could scale poorly for large-scale

*Technically, a conjecture.

¹We use the $\tilde{O}(\cdot)$ notation to hide logarithmic factors.

problems. We also mention a recent work by Carmon and Duchi [CD16], studying the closely related problem

$$\text{minimize } (1/2)x^T Ax + b^T x + (\rho/3)\|x\|_2^3, \quad (2)$$

in variable $x \in \mathbf{R}^n$, again with A symmetric, possibly indefinite, and parameter $\rho > 0$. The authors analyze gradient descent, proving that $\tilde{O}(1/\varepsilon)$ gradient steps are enough to output an ε -suboptimal point.

In this paper we demonstrate that the projected gradient method on (1) asymptotically converges to a global minimizer on the trust region subproblem.

1.2 Notation and classical results

In the sequel, we refer to the objective function as $f : \mathbf{R}^n \rightarrow \mathbf{R}$, given by $f(x) = (1/2)x^T Ax + 2b^T x$. Additionally, the constraint set is the closed ball $\mathcal{B}(R) \triangleq \{x \in \mathbf{R}^n \mid \|x\| \leq R\}$, where $\|\cdot\|$ denotes the Euclidean norm. We use the notation x^* to denote the global minimum of f when it is unique, so that $x^* = \operatorname{argmin}_{x \in \mathcal{B}(R)} f(x)$. We use f^* to denote the optimal value of f , so that $f^* = \inf_{x \in \mathcal{B}(R)} f(x)$. Hence, when x^* exists, $f^* = f(x^*)$.

We fix the eigendecomposition of $A = UDU^T$, where $D = \mathbf{diag}(\lambda_1, \dots, \lambda_n)$, and U has orthonormal columns u_i . We impose without loss that $\lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_n$. By $\|\cdot\|_{\text{op}}$, we denote the ℓ_2 -operator norm $\|M\|_{\text{op}} = \sup_{\|x\|=1} \|Mx\|$, for any $M \in \mathbf{R}^{n \times n}$. A useful identity is that $\|M\|_{\text{op}} = \max_i |\lambda_i(M)|$ when M is a symmetric $n \times n$ matrix. We will put $\beta \triangleq \|A\|_{\text{op}}$.

Additionally, say a differentiable function $g : \mathbf{R}^n \rightarrow \mathbf{R}$ is L -smooth on convex set $C \subset \mathbf{R}^n$, provided that

$$\|\nabla g(x) - \nabla g(y)\| \leq L\|x - y\| \quad \text{for any } x, y \in C.$$

It is well known that this implies

$$g(x) - g(y) \leq \nabla g(y)^T (x - y) + \frac{L}{2}\|x - y\|^2 \quad \text{for any } x, y \in C. \quad (3)$$

Equivalently, $\|g(x)\|_{\text{op}} \leq L$, for Lebesgue almost every $x \in C$. For nonempty, closed, convex sets $C \subset \mathbf{R}^n$, associate the projection operator $\Pi_C : \mathbf{R}^n \rightarrow C$ given by

$$\Pi_C(x) = \operatorname{argmin}_{y \in C} \left(\frac{1}{2}\|x - y\|^2 \right),$$

for any $x \in \mathbf{R}^n$. In the sequel we denote by $I : \mathbf{R}^n \rightarrow \mathbf{R}^n$ the identity operator on \mathbf{R}^n .

2 Asymptotic convergence to a global minimizer

2.1 Projected gradient descent

Projected gradient descent (PGD) begins at an initialization $x^{(0)} \in \mathbf{R}^n$ and generates iterates

$$y^{(k+1)} = x^{(k)} - \eta \nabla f(x^{(k)}) \quad (4)$$

$$x^{(k+1)} = \Pi_{\mathcal{B}(R)}(y^{(k+1)}), \quad (5)$$

for nonnegative integer k and step size η . We make the following assumptions about this procedure.

Assumption 2.1. In (4), the step size η satisfies $0 < \eta < \frac{1}{\beta}$.

Assumption 2.2. The initial point satisfies $x^{(0)} = 0$.

2.2 Asymptotic convergence to a global minimizer

We begin by providing a few results, which characterize the iterates of projected gradient descent.

Lemma 2.3. *Let Assumptions 2.1 and 2.2 hold. Then the iterates of gradient descent satisfy $(u_i^T x^{(k)})(u_i^T b) \leq 0$ for all $i = 1, \dots, n$ and every $k \geq 0$. 0*

Proof. Evidently, the claim holds due to Assumption 2.2 when $k = 0$. Thus, inductively assume that for some k

$$(u_i^T x^{(k)})(u_i^T b) \leq 0 \quad \text{for all } i = 1, \dots, n. \quad (6)$$

By definition, $x^{(k+1)} = cy^{(k+1)}$ for some $c \in (0, 1]$, so it suffices to ensure $(u_i^T y^{(k+1)})(u_i^T b) \leq 0$. Using (6) along with Assumption 2.1,

$$(u_i^T y^{(k+1)})(u_i^T b) = (1 - \eta\lambda_i)(u_i^T x^{(k)})(u_i^T b) - \eta(u_i^T b)^2 \leq 0,$$

since $\eta < \beta^{-1} \leq \lambda_i^{-1}$, for all $i = 1, \dots, n$. This proves the result. \square

The following result shows projected gradient descent is a descent method for (1).

Lemma 2.4. *Let Assumption 2.1 hold. Then for any $k > 0$,*

$$f(x^{(k+1)}) - f(x^{(k)}) \leq \left(\frac{\beta}{2} - \frac{1}{2\eta} \right) \|x^{(k+1)} - x^{(k)}\|^2.$$

Proof. Basic manipulations imply

$$\nabla f(x^{(k)})^T (x - x^{(k)}) + \frac{1}{2\eta} \|x - x^{(k)}\|^2 = \frac{1}{2\eta} \|x - (x^{(k)} - \eta \nabla f(x^{(k)}))\|^2 - \frac{\eta}{2} \|\nabla f(x^{(k)})\|^2.$$

Thus, as $\eta > 0$ it follows that

$$\operatorname{argmin}_{x \in \mathcal{B}(R)} \left(\nabla f(x^{(k)})^T (x - x^{(k)}) + \frac{1}{2\eta} \|x - x^{(k)}\|^2 \right) = \operatorname{argmin}_{x \in \mathcal{B}(R)} \left(\frac{1}{2} \|x - (x^{(k)} - \eta \nabla f(x^{(k)}))\|^2 \right).$$

Comparing the display above to (4), (5), and the definition of $\Pi_{\mathcal{B}(R)}$,

$$x^{(k+1)} = \operatorname{argmin}_{x \in \mathcal{B}(R)} \left(\nabla f(x^{(k)})^T (x - x^{(k)}) + \frac{1}{2\eta} \|x - x^{(k)}\|^2 \right). \quad (7)$$

Appealing to the β -smoothness of f and evaluating (7) at $x^{(k)} \in \mathcal{B}(R)$,

$$f(x^{(k+1)}) - f(x^{(k)}) \leq \nabla f(x^{(k)})^T (x^{(k+1)} - x^{(k)}) + \frac{\beta}{2} \|x^{(k+1)} - x^{(k)}\|^2 \leq \left(\frac{\beta}{2} - \frac{1}{2\eta} \right) \|x^{(k+1)} - x^{(k)}\|^2.$$

\square

The following result provides a useful optimality criterion for the trust region subproblem (1).

Theorem 2.5 ([CGT00], Corollary 7.2.2.). *A point $x \in \mathcal{B}(R)$ is a global minimizer of f subject to $\|x\| \leq R$ if and only if for some $z \geq 0$,*

$$(A + zI)x = -b \quad A + zI \succeq 0 \quad z(\|x\| - R) = 0.$$

Furthermore, x is unique if and only if $A + zI \succ 0$. In this case, we write $x = x^$.*

An important special case from Theorem 2.5 is that when $\|x^*\| < R$, then $\nabla f(x^*) = 0$. Furthermore, with a simplifying assumption, we can provide a set of simpler optimality criterion.

Corollary 2.6. *Suppose that $b^T u_1 \neq 0$. Then if for some $\tilde{x} \in \mathcal{B}(R)$ and $z \geq 0$, it holds that*

$$(A + zI)\tilde{x} = -b \quad z(\|\tilde{x}\| - R) = 0 \quad (u_1^T \tilde{x})(u_1^T b) \leq 0 \quad (8)$$

then \tilde{x} is the unique global minimizer to f over $\mathcal{B}(R)$, i.e., $\tilde{x} = x^$.*

Proof. Focusing on the first condition, $b^T u_1 = -(z + \lambda_1)(u_1^T \tilde{x})$. Thus, $b^T u_1 \neq 0$ implies that $(u_1^T \tilde{x}) \neq 0$ and $z + \lambda_1 \neq 0$, strengthening the third condition to $(u_1^T \tilde{x})(u_1^T b) < 0$. But this implies that $z + \lambda_1 = -(u_1^T b)(u_1^T \tilde{x}) / (u_1^T \tilde{x})^2 > 0$, which implies that $z > \lambda_i$ for all i , whence $A + zI \succ 0$, establishing the result. \square

The assumptions along with Corollary 2.6 and Lemmas 2.3 and 2.4 give us our desired asymptotic convergence guarantee.

Proposition 2.7 (Asymptotic convergence). *Let Assumptions 2.1 and 2.2 hold, and suppose $b^T u_1 \neq 0$. Then as $k \rightarrow \infty$, the iterates of projected gradient descent satisfy $x^{(k)} \rightarrow x^*$ and $f(x^{(k)}) \downarrow f(x^*)$.*

Proof. It suffices to demonstrate that $x^{(k)} \rightarrow x^*$, because then the conclusion follows via continuity of f and to that end, Lemma 2.4. Lemma 2.4 and Assumption 2.1 yield the following bound for any integer $T \geq 1$,

$$\left(\frac{1}{2\eta} - \frac{\beta}{2}\right) \sum_{k=0}^{T-1} \|x^{(k+1)} - x^{(k)}\|^2 \leq f(x^{(0)}) - f(x^{(T)}) \leq f(x^{(0)}) - f^*. \quad (9)$$

Now, define $\phi : \mathcal{B}(R) \rightarrow \mathbf{R}^n$ by $\phi(x) = \Pi_{\mathcal{B}(R)}(x - \eta \nabla f(x)) - x$, for points $x \in \mathcal{B}(R)$. The bound in (9) implies that the displayed series is convergent as $T \rightarrow \infty$ and thus $\phi(x^{(k)}) \rightarrow 0$. Note also that the map ϕ is evidently continuous, as ∇f is β -Lipschitz and $\Pi_{\mathcal{B}(R)}$ is non-expansive, thus 1-Lipschitz.

Suppose now that $\tilde{x} \in \mathcal{B}(R)$ is a subsequential limit of $(x^{(k)})$ (indeed, one exists since this sequence is bounded), and observe by continuity $\phi(\tilde{x}) = 0$. To show that $\tilde{x} = x^*$, by Corollary 2.6, it suffices to establish the first two conditions of (8), as the third immediately holds by Lemma 2.3. Observe first that $\phi(\tilde{x}) = 0$ implies that for some $c \geq 1$,

$$\tilde{x} - \eta \nabla f(\tilde{x}) = \tilde{x} - \eta(A\tilde{x} - b) = c\tilde{x}. \quad (10)$$

Indeed, setting $z = (c - 1)\eta^{-1}$, this implies that $(A + zI)\tilde{x} = -b$. If \tilde{x} lies on the boundary of $\mathcal{B}(R)$, so that $\|\tilde{x}\| = R$, then as $z \geq 0$, this establishes (8) and hence $\tilde{x} = x^*$. On the other hand, if \tilde{x} is in the interior of $\mathcal{B}(R)$, so that $\|\tilde{x}\| < R$, then $\phi(\tilde{x}) = 0$ implies that $c = 1$ in (10), and thus $z = 0$, once again establishing (8), hence also that $\tilde{x} = x^*$. As this analysis applies to any such subsequential limit \tilde{x} of the bounded sequence $(x^{(k)})$, the claim is now proven (since the iterates lie in $\mathcal{B}(R)$, which is compact). \square

We provide some numerical evidence demonstrating the effect of Proposition 2.7 in Figure 1.

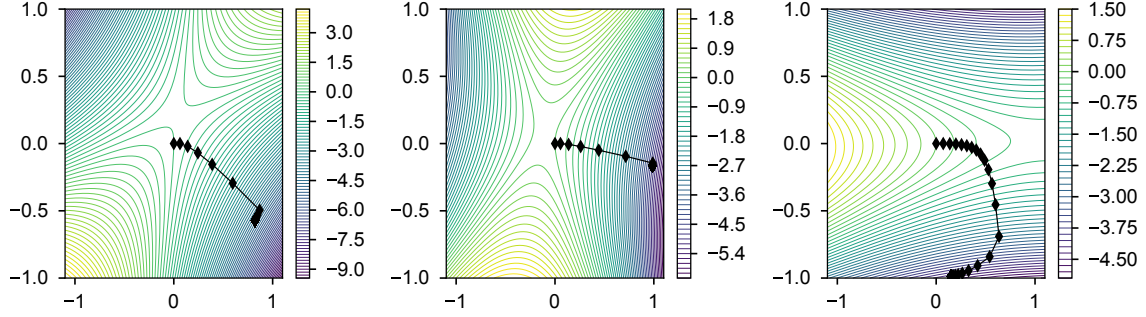


Figure 1: Three random indefinite instances of the trust region subproblem (1), with $R = 1$, $\eta = 1/(2\|A\|_{\text{op}})$ and $x^{(0)} = 0$. From left to right, the eigenvalues are $\lambda = (-8, 3)$, $\lambda = (-9, 3)$, and $\lambda = (-7, 1)$. The dots indicate iterates of projected gradient descent and the lines indicate the process $\dot{x} = -\nabla f(x)$.

3 Non-asymptotic convergence guarantees

3.1 Monotone increasing norm

We state and prove a lemma about the norms of projected gradient descent iterates, which will help us bound the time to optimum.

Proposition 3.1. *Let $x^{(0)} = 0$. Then the iterates of projected gradient descent satisfy*

$$\|x^{(t+1)}\| \geq \|x^{(t)}\| \quad (11)$$

for all $t \geq 0$.

Before proving this proposition, we prove the following lemmas about the behavior of PGD with respect to the eigenbasis of the problem. Throughout, we denote $\text{sgn}(a) = \mathbf{1}\{a \geq 0\} - \mathbf{1}\{a \leq 0\}$.

Lemma 3.2. *Let $x^{(0)} = 0$. Then, for $t \geq 1$ all iterates $x^{(t)}$ satisfy*

$$\text{sgn}(u_i^T x^{(t)}) = -\text{sgn}(u_i^T b). \quad (12)$$

Furthermore, if $\lambda_i \neq 0$, then for $t \geq 1$,

$$\text{sgn}\left(u_i^T x^{(t)} + \frac{u_i^T b}{\lambda_i}\right) = \text{sgn}\left(\frac{b_i}{\lambda_i}\right). \quad (13)$$

Proof. Since $y^{(1)} = -\eta b$, the first claim is clearly true for $t = 1$. Assume it is true up to $x^{(t-1)}$. Then (noting that the projection operator does not alter sign)

$$\begin{aligned} \text{sgn}(u_i^T x^{(t)}) &= \text{sgn}\left(u_i^T \left[x^{(t-1)} - \eta \nabla f(x^{(t-1)})\right]\right) \\ &= \text{sgn}\left((1 - \eta \lambda_i) u_i^T x^{(t-1)} - \eta u_i^T b\right) \end{aligned}$$

By assumption 2.1, $1 - \eta \lambda_i > 0$, so $\text{sgn}\left((1 - \eta \lambda_i) u_i^T x^{(t-1)}\right) = \text{sgn}(u_i^T x^{(t-1)})$. Therefore,

$$\text{sgn}\left((1 - \eta \lambda_i) u_i^T x^{(t-1)} - \eta u_i^T b\right) = -\text{sgn}(u_i^T b),$$

establishing the first claim.

The second claim is also clearly true for $t = 0$. Assume it is true up to $x^{(t-1)}$. Then, for some positive constant $c \leq 1$,

$$\begin{aligned}\mathbf{sgn}\left(u_i^T x^{(t)} + \frac{u_i^T b}{\lambda_i}\right) &= \mathbf{sgn}\left(cu_i^T \left[x^{(t-1)} - \eta \nabla f(x^{(t-1)})\right] + \frac{u_i^T b}{\lambda_i}\right) \\ &= \mathbf{sgn}\left(c(1 - \eta\lambda_i)u_i^T x^{(t-1)} + (1 - c\eta\lambda_i)\frac{u_i^T b}{\lambda_i}\right)\end{aligned}$$

By assumption 2.1, $(1 - c\eta\lambda_i) > 0$ and $c(1 - \eta\lambda_i) > 0$. Therefore,

$$\begin{aligned}\mathbf{sgn}\left(c(1 - \eta\lambda_i)u_i^T x^{(t-1)} + (1 - c\eta\lambda_i)\frac{u_i^T b}{\lambda_i}\right) &= \mathbf{sgn}\left(u_i^T x^{(t-1)} + \frac{u_i^T b}{\lambda_i}\right) \\ &= \mathbf{sgn}\left(\frac{b_i}{\lambda_i}\right)\end{aligned}$$

□

We are now ready to prove 3.1,

Proof. If $\lambda_i = 0$, then $\nabla f(x^{(t)}) = \eta b$, and so $\mathbf{sgn}(u_i^T \nabla f(x^{(t)})) = \mathbf{sgn}(u_i^T b)$. If $\lambda_i \neq 0$, then

$$\begin{aligned}\mathbf{sgn}\left(u_i^T \nabla f(x^{(t)})\right) &= \mathbf{sgn}\left(\lambda_i u_i^T x^{(t)} + u_i^T b_i\right) \\ &= \mathbf{sgn}(\lambda_i) \mathbf{sgn}\left(u_i^T x^{(t)} + \frac{u_i^T b}{\lambda_i}\right) \\ &= \mathbf{sgn}(\lambda_i) \mathbf{sgn}\left(\frac{u_i^T b}{\lambda_i}\right) \\ &= \mathbf{sgn}(u_i^T b)\end{aligned}$$

From this, it follows that

$$(x^{(t)})^T \nabla f(x^{(t)}) = \sum_{i=1}^n (u_i^T x^{(t)})(u_i^T \nabla f(x^{(t)})) \leq 0$$

and therefore,

$$\begin{aligned}\|x^{(t)} - \eta \nabla f(x^{(t)})\|^2 &= \|x^{(t)}\|^2 + \eta^2 \|\nabla f(x^{(t)})\|^2 - \eta(x^{(t)})^T \nabla f(x^{(t)}) \\ &\geq \|x^{(t)}\|^2\end{aligned}$$

□

3.2 Projected gradient descent converges to opt linearly on the boundary

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