

# ALL PAIRS SHORTEST PATHS

# Presentation Outline

## LIST OF TOPICS

- What is All Pairs Shortest Path Problem?
- Floyd-Warshall Algorithm
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# WHAT IS ALL PAIRS SHORTEST PATH PROBLEM?

- *The all-pairs shortest path problem is the determination of the shortest graph distances between every pair of vertices in a given graph.*
- *We have to calculate the minimum cost to find the shortest path.*

**THE ALL PAIR  
SHORTEST PATH  
ALGORITHM IS  
ALSO KNOWN AS  
FLOYD - WARSHALL  
ALGORITHM**

Floyd-Warshall Algorithm is an algorithm for finding the shortest path between all the pairs of vertices in a weighted graph.

- This algorithm follows the dynamic programming approach to find the shortest paths.
- This algorithm works for both the directed and undirected weighted graphs. But, it does not work for the graphs with negative cycles (where the sum of the edges in a cycle is negative).

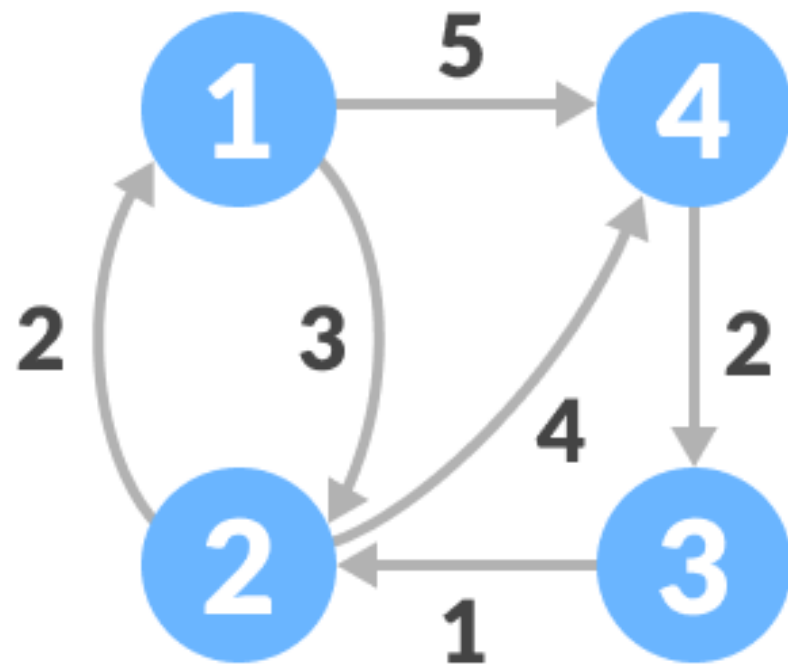
**FLOYD-WARSHALL ALGORITHM**

# PSEUDOCODE

## FLOYD - WARSHALL (W)

1.  $n \leftarrow \text{rows } [W]$ .
2.  $D0 \leftarrow W$
3. for  $k \leftarrow 1$  to  $n$
4. do for  $i \leftarrow 1$  to  $n$
5. do for  $j \leftarrow 1$  to  $n$
6. do  $\text{dij}(k) \leftarrow \min (\text{dij}(k-1), \text{dik}(k-1) + \text{dkj}(k-1) )$
7. return  $D(n)$

Let the given graph be:



Follow the steps below to find the shortest path between all the pairs of vertices.

*EXAMPLE*

CREATE A  
MATRIX  $A^0$  OF DIMENSION  $N \times N$ ,  
WHERE 'N' IS THE NO. OF  
VERTICES.

Each cell  $A[i][j]$  is filled with the  
distance from the  $i$ th vertex to  
the  $j$ th vertex. If there is no path  
from  $i$ th vertex to  $j$ th vertex, the  
cell is left as infinity.

$$A^0 = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 3 & \infty & 5 \\ 2 & 0 & \infty & 4 \\ \infty & 1 & 0 & \infty \\ \infty & \infty & 2 & 0 \end{bmatrix} \end{matrix}$$



Now, create a matrix A1 using matrix A0.

$$\mathbf{A}^1 = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 3 & \infty & 5 \\ 2 & 0 & & \\ \infty & & 0 & \\ \infty & & & 0 \end{bmatrix} \end{matrix} \rightarrow \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 3 & \infty & 5 \\ 2 & 0 & 9 & 4 \\ \infty & 1 & 0 & 8 \\ \infty & \infty & 2 & 0 \end{bmatrix} \end{matrix}$$

The elements in the first column and row are left as they are.

$$A^2 = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 3 & & \\ 2 & 0 & 9 & 4 \\ & 1 & 0 & \\ & \infty & & 0 \end{bmatrix} \end{matrix} \rightarrow \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 3 & 9 & 5 \\ 2 & 0 & 9 & 4 \\ 3 & 3 & 1 & 0 & 5 \\ 4 & \infty & \infty & 2 & 0 \end{bmatrix} \end{matrix}$$

*In a similar way,  $A_2$  is created using  $A_3$ .  
The elements in the second column and the second row are left as they are.*

SIMILARLY, A3 AND A4 IS ALSO CREATED.

$$\mathbf{A}^3 = \begin{array}{c} \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & & \infty & \\ & 0 & 9 & \\ \infty & 1 & 0 & 8 \\ & & 2 & 0 \end{bmatrix} \end{array} \rightarrow \begin{array}{c} \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 3 & 9 & 5 \\ 2 & 0 & 9 & 4 \\ 3 & 1 & 0 & 5 \\ 5 & 3 & 2 & 0 \end{bmatrix} \end{array}$$

$$A^4 = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & & & 5 \\ & 0 & & 4 \\ & & 0 & 5 \\ 5 & 3 & 2 & 0 \end{bmatrix} \end{matrix}$$



$$\begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 0 & 3 & 7 & 5 \\ 2 & 0 & 6 & 4 \\ 3 & 1 & 0 & 5 \\ 5 & 3 & 2 & 0 \end{bmatrix} \end{matrix}$$

**4 GIVES THE  
SHORTEST PATH  
BETWEEN EACH PAIR  
OF VERTICES.**

# Time Complexity

*There are three loops. Each loop has constant complexities. So, the time complexity of the Floyd-Warshall algorithm is  $O(n^3)$ .*

# Space Complexity

*The space complexity of the Floyd-Warshall algorithm is  $O(n^2)$ .*

# Floyd Warshall Algorithm Applications

- *To find the shortest path in a directed graph*
- *To find the transitive closure of directed graphs*
- *To find the Inversion of real matrices*
- *For testing whether an undirected graph is bipartite*