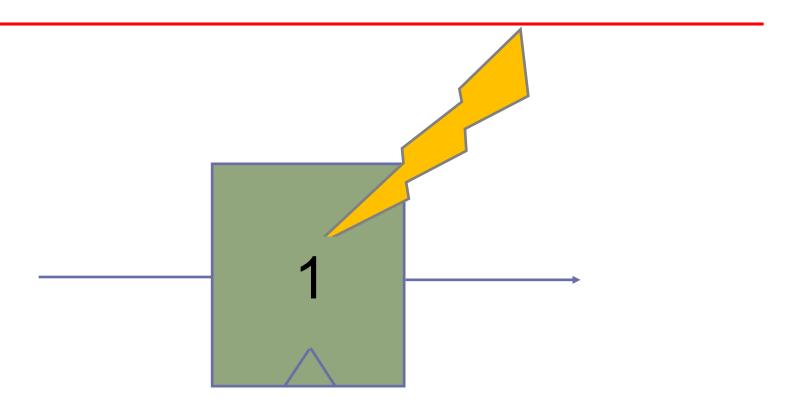
### Reliable Architectures

Joel Emer
Computer Science & Artificial Intelligence Lab
M.I.T.

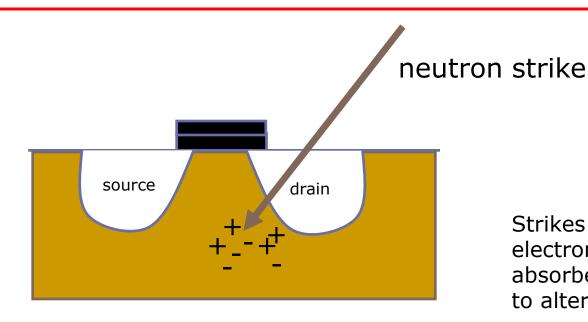
Many of the slides in this presentation are from public presentations made by Joel Emer for the AVF work

### Event Changes State of a Single Bit



- Hard Error Changes that are permanent
- Soft Error Changes that are not permanent

### Impact of Neutron Strike on a Si Device

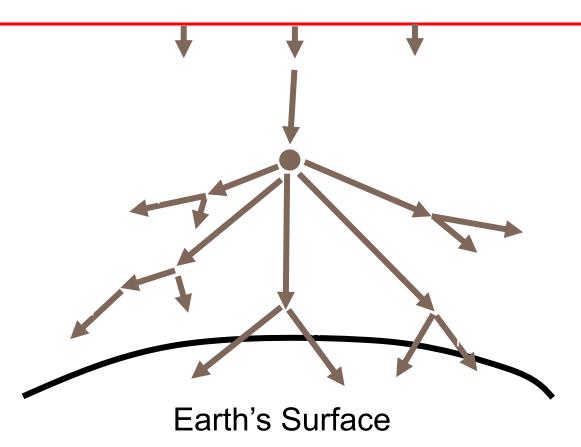


Transistor Device

Strikes release electron & hole pairs that can be absorbed by source & drain to alter the state of the device

Secondary source of upsets: Alpha particles from packaging

### Cosmic Rays Come From Deep Space



- Neutron flux is higher at higher altitudes
  - 3-5x increase in Denver at 5,000 feet
  - 100x increase in airplanes at 30,000+ feet

### Basics of Charge Generation

### Cosmic rays of >1GeV result in neutrons of >1MeV

Energy (eV)	Electron-Hole Pairs	Charge (Femtocoulombs)
3.6eV	1	3.2x10 <sup>-4</sup>
1MeV	~2.8x10 <sup>5</sup>	~44
1GeV	~2.8x10 <sup>8</sup>	~44x10 <sup>3</sup>

#### In 2010:

- Critical charge on a DRAM: ~25 fCoulomb
- Critical charge on an SRAM: <4 fCoulomb</li>

# Cosmic Ray Strikes: Evidence & Reaction

#### Publicly disclosed incidences

- Error logs in large servers, E. Normand, "Single Event Upset at Ground Level," IEEE Trans. on Nucl Sci, Vol. 43, No. 6, Dec 1996.
- Sun Microsystems found cosmic ray strikes on L2 cache with defective error protection caused Sun's flagship servers to crash, R. Baumann, IRPS Tutorial on SER, 2000.
- Cypress Semiconductor reported in 2004 a single soft error brought a billion-dollar automotive factory to a halt once a month, Zielger & Puchner, "SER – History, Trends, and Challenges," Cypress, 2004.
- In 2003, a "single-event upset" was blamed for an electronic voting error in Schaerbeekm, Belgium. A bit flip in the electronic voting machine added 4,096 extra votes to one candidate.

### Physical solutions are hard

#### Shielding?

- No practical absorbent (e.g., approximately > 10 ft of concrete)
- This is unlike Alpha particles which are easily blocked

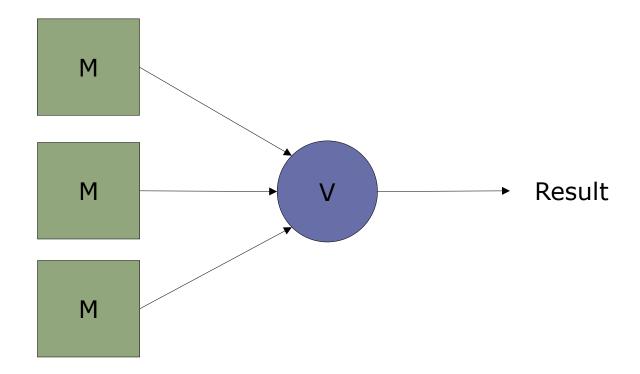
#### Technology solution?

- Partially-depleted SOI of some help, effect on logic unclear
- Fully-depleted SOI may help, but is challenging to manufacture
- FinFETs are showing significantly lower vulnerability

#### Circuit-level solution?

- Radiation-hardened circuits can provide 10x improvement with significant penalty in performance, area, cost
- 2-4x improvement may be possible with less penalty

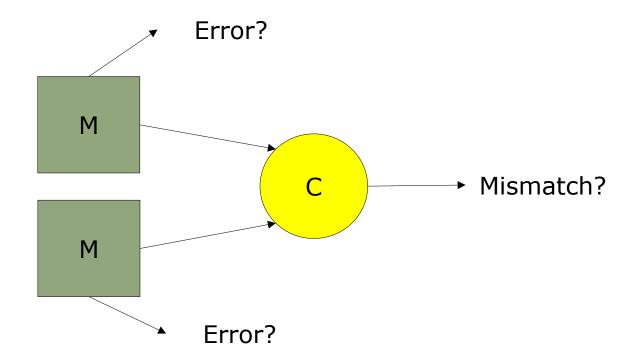
## Triple Modular Redundancy (Von Neumann, 1956)



V does a majority vote on the results

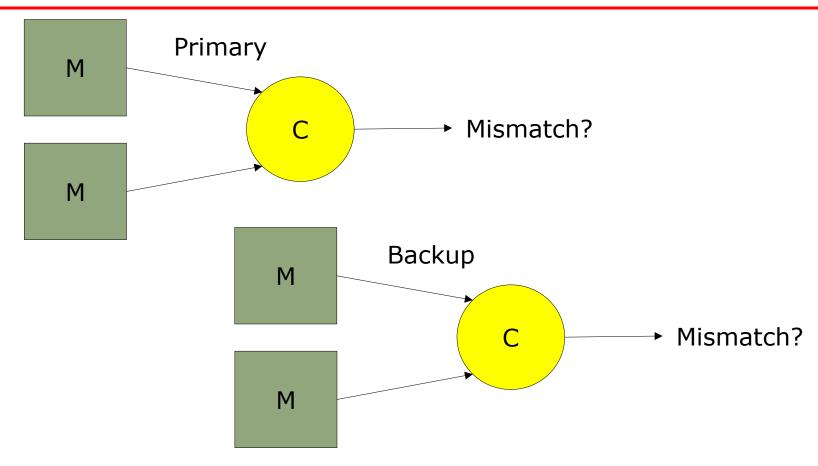
### Dual Modular Redundancy

(e.g., BINAC 1949, Stratus 1982)



- Processing stops on mismatch
- Error signal used to decide which processor be used to restore state to other

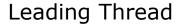
## Pair and Spare Lockstep (e.g., Tandem, 1975)

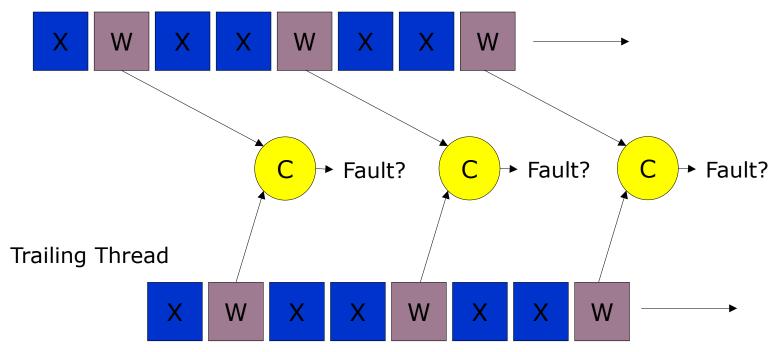


- Primary creates periodic checkpoints
- Backup restarts from checkpoint on mismatch

## Redundant Multithreading

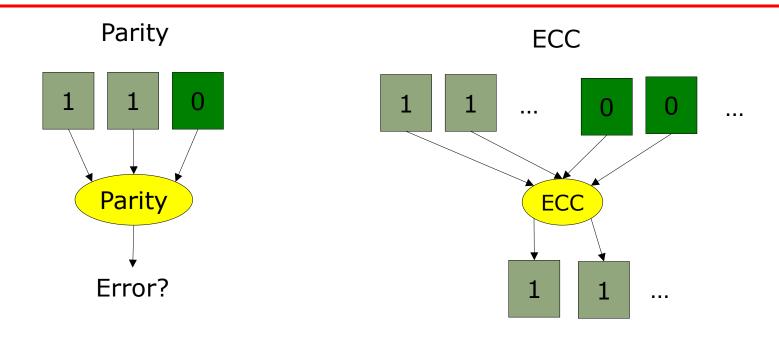
(e.g., Reinhardt, Mukherjee, 2000)





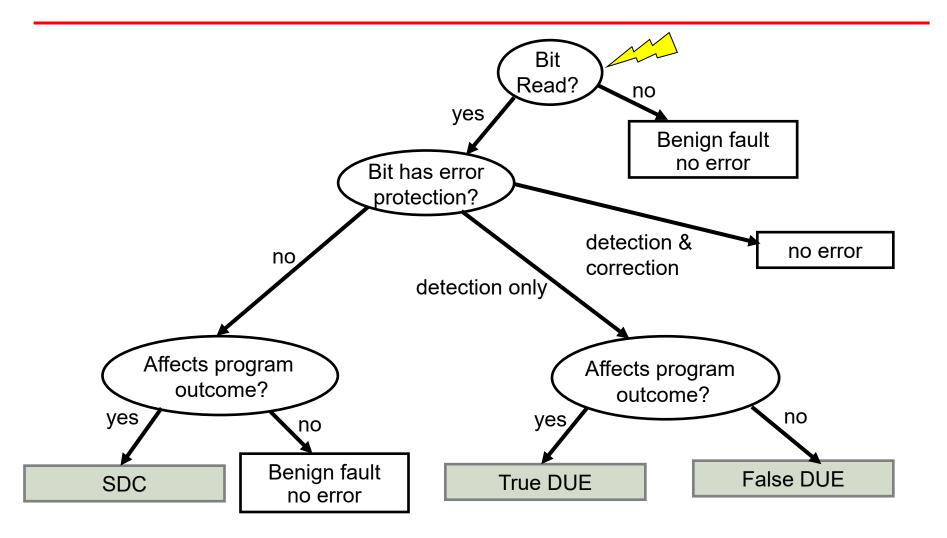
Writes are checked

### **Component Protection**



- Fujitsu SPARC in 130 nm technology (ISSCC 2003)
  - 80% of 200k latches protected with parity

## Strike on a bit (e.g., in register file)



SDC = Silent Data Corruption, DUE = Detected Unrecoverable Error

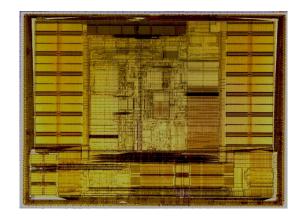
### **Metrics**

#### Interval-based

- MTTF = Mean Time to Failure
- MTTR = Mean Time to Repair
- MTBF = Mean Time Between Failures = MTTF + MTTR
- Availability = MTTF / MTBF

#### Rate-based

- FIT = Failure in Time = 1 failure in a billion hours
- $1 \text{ year MTTF} = 10^9 / (24 * 365) \text{ FIT} = 114,155 \text{ FIT}$
- SER FIT = SDC FIT + DUE FIT



#### Hypothetical Example

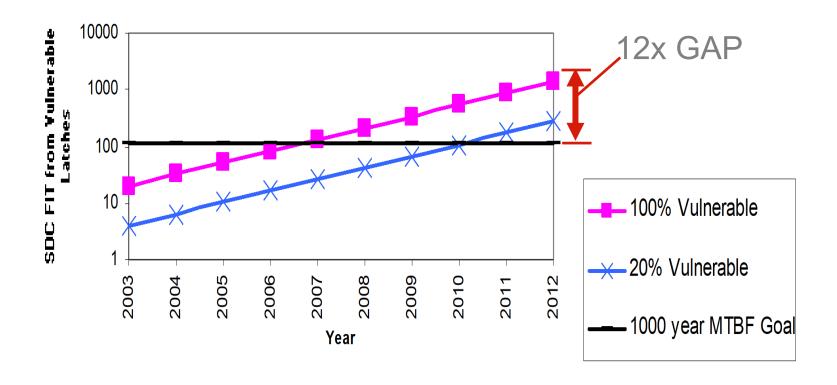
Cache: 0 FIT

+ IQ: 100K FIT

+ FU: 58K FIT

Total of 158K FIT

# Number of Vulnerable Bits Growing with Moore's Law



Typical SDC goal: 1000 year MTBF Typical DUE goal: 10-25 year MTBF

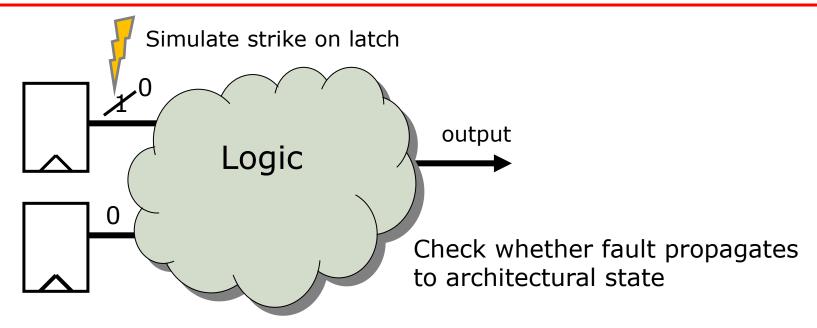
# Architectural Vulnerability Factor (AVF)

$$AVF_{bit} = Probability Bit Matters$$

$$= \frac{\# \text{ of Visible Errors}}{\# \text{ of Bit Flips from Particle Strikes}}$$

$$FIT_{bit} = intrinsic FIT_{bit} * AVF_{bit}$$

# Statistical Fault Injection (SFI) with RTL



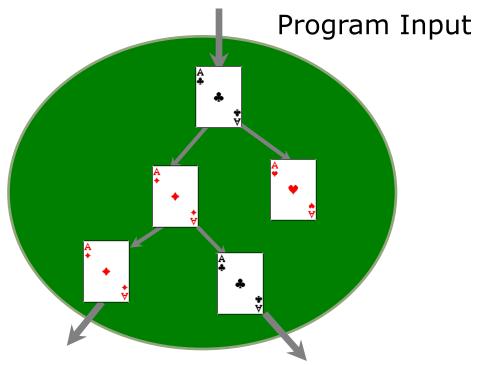
- + Naturally characterizes all logical structures
- RTL not available until late in the design cycle
- Numerous experiments to flip all bits
- Generally done at the chip level
  - Limited structural insight

# Architectural Vulnerability Factor Does a bit matter?

Branch Predictor

• Program Counter

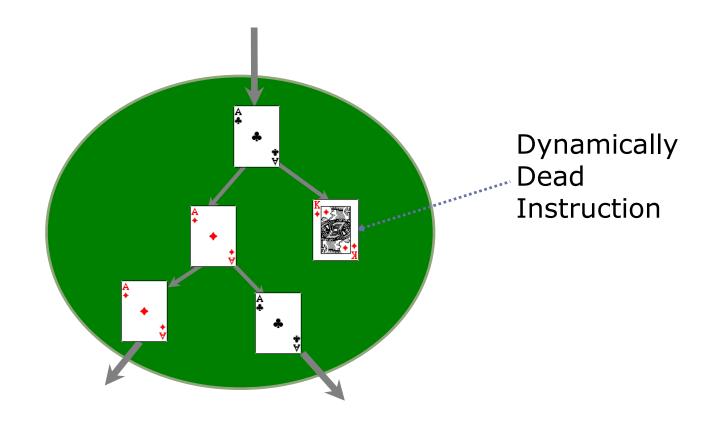
# Architecturally Correct Execution (ACE)



**Program Outputs** 

- ACE path requires only a subset of values to flow correctly through the program's data flow graph (and the machine)
- Anything else (un-ACE path) can be derated away

# Example of un-ACE instruction: Dynamically Dead Instruction

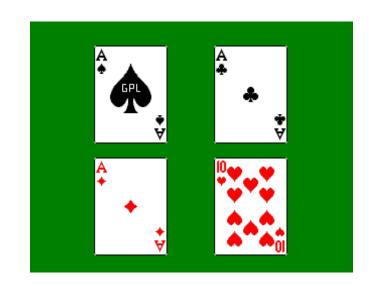


Most bits of an un-ACE instruction do not affect program output

### Vulnerability of a structure

AVF = fraction of cycles a bit contains ACE state

$$T = 3$$



ACE% = 24/4

### Vulnerability of a structure

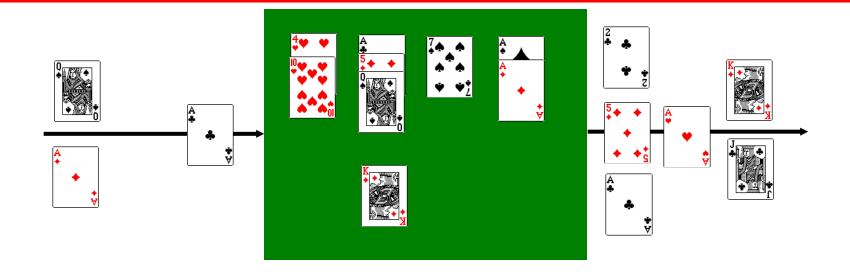
AVF = fraction of cycles a bit contains ACE state

$$=\frac{(2+1+0+3)/4}{4}$$

= Average number of ACE bits in a cycle

Total number of bits in the structure

### Little's Law for ACEs



$$\overline{N}_{ace} = \overline{T}_{ace} \times \overline{L}_{ace}$$

$$AVF = \frac{\overline{N}_{ace}}{N_{total}}$$

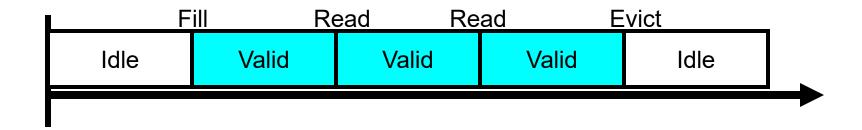
## Computing AVF

- Approach is conservative
  - Assume every bit is ACE unless proven otherwise
- Data Analysis using a Performance Model
  - Prove that data held in a structure is un-ACE
- Timing Analysis using a Performance Model
  - Tracks the time this data spent in the structure

## ACE Lifetime Analysis (1)

(e.g., write-through data cache)

#### Idle is unACE

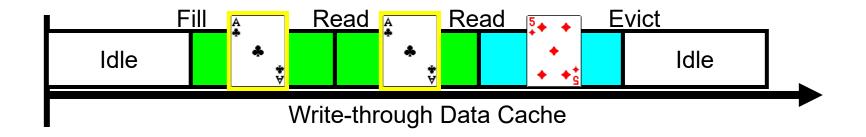


- Assuming all time intervals are equal
- For 3/5 of the lifetime the bit is valid
- Gives a measure of the structure's utilization
  - Number of useful bits
  - Amount of time useful bits are resident in structure
  - Valid for a particular trace

## ACE Lifetime Analysis (2)

(e.g., write-through data cache)

### Valid is not necessarily ACE

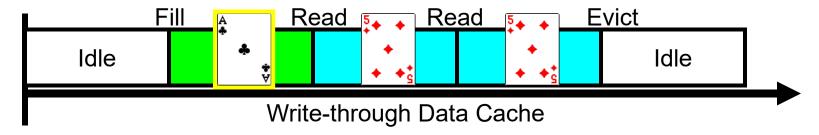


- ACE % = AVF = 2/5 = 40%
- Example Lifetime Components
  - ACE: fill-to-read, read-to-read
  - unACE: idle, read-to-evict, write-to-evict

## ACE Lifetime Analysis (3)

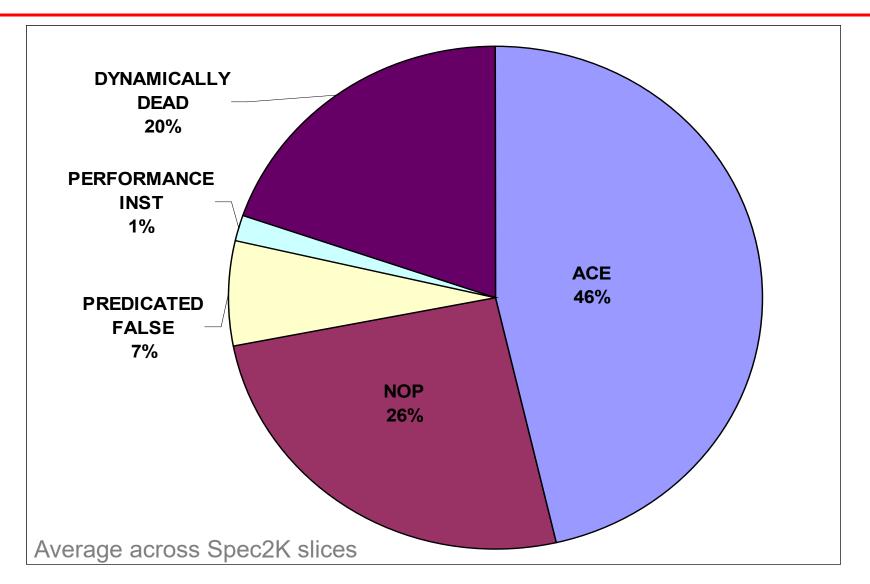
(e.g., write-through data cache)

 Data ACEness is a function of instruction **ACEness** 

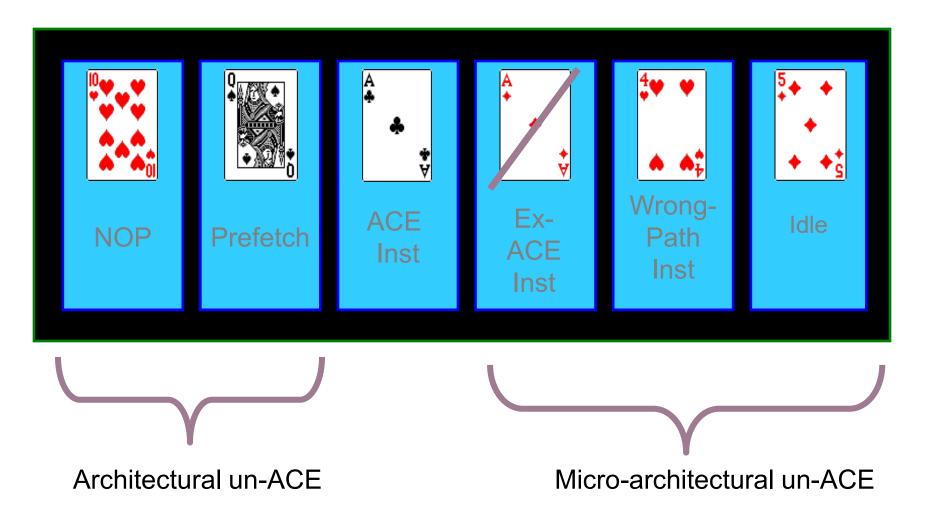


- Second Read is by an unACE instruction
- AVF = 1/5 = 20%

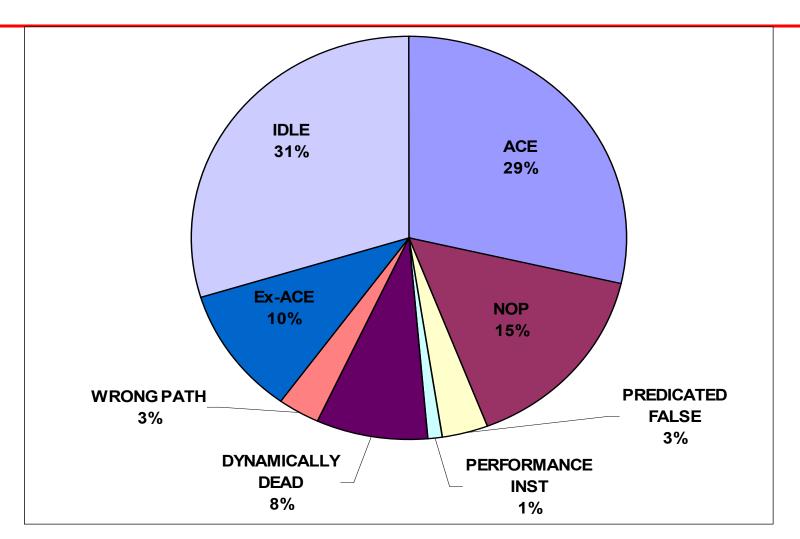
### Dynamic Instruction Breakdown



# Mapping ACE & un-ACE Instructions to the Instruction Queue

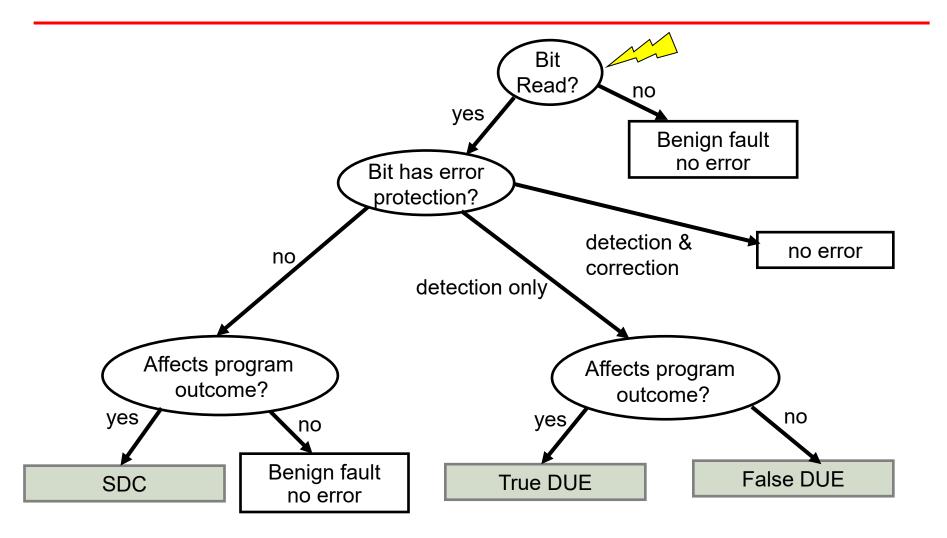


### Instruction Queue



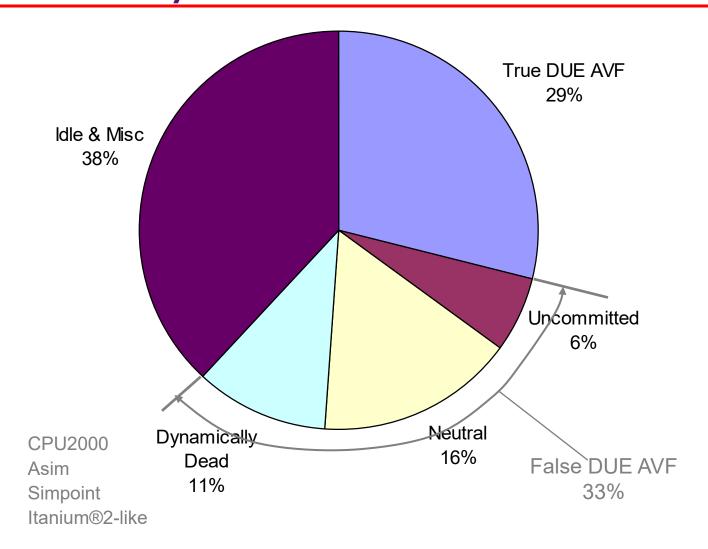
ACE percentage = AVF = 29%

## Strike on a bit (e.g., in register file)

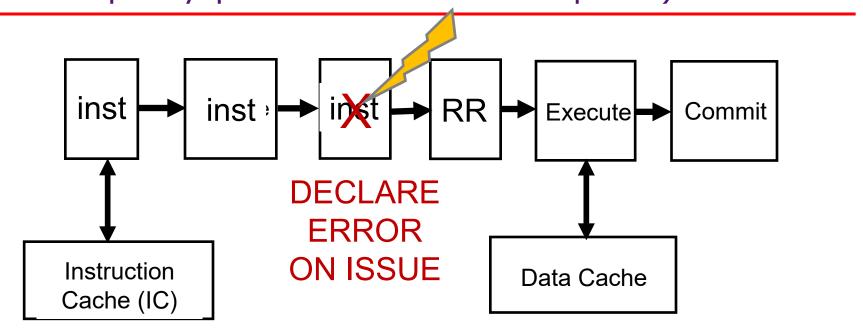


SDC = Silent Data Corruption, DUE = Detected Unrecoverable Error

# DUE AVF of Instruction Queue with Parity



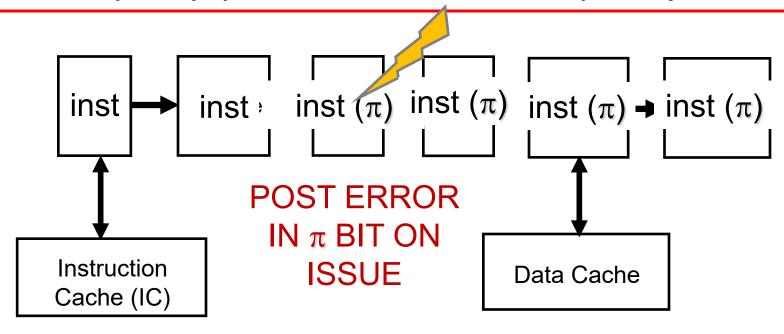
## Coping with Wrong-Path Instructions (assume parity-protected instruction queue)



Problem: not enough information at issue

## The $\pi$ (Possibly Incorrect) Bit

(assume parity-protected instruction queue)



At commit point, declare error only if not wrong-path instruction and  $\pi$  bit is set

# Sources of False DUE in an Instruction Queue

- Instructions with uncommitted results
  - e.g., wrong-path, predicated-false
  - solution:  $\pi$  (possibly incorrect) bit till commit
- Instruction types neutral to errors
  - e.g., no-ops, prefetches, branch predict hints
  - solution: anti- $\pi$  bit
- Dynamically dead instructions
  - instructions whose results will not be used in future
  - solution:  $\pi$  bit beyond commit

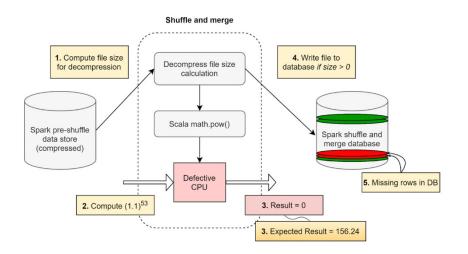
### Reliability Problems in 2020s

### Silent Data Corruption (SDC)

- Cloud companies noticed SDC is a widespread problem for largescale infrastructure systems.
- "Cores that don't count" by Google, HotOS, 2021
- "Silent data corruption at Scale" by Facebook, Arxiv, 2021

#### Problems

- Long error detection latencies: taking days to weeks
- Scalability



#### **Example errors:**

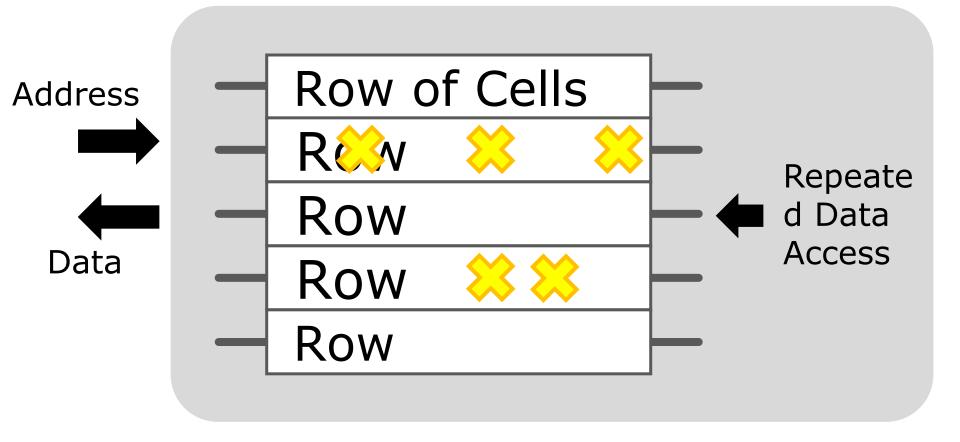
$$Int[(1.1)^3] = 0$$
,  $expected = 1$ 

$$Int[(1.1)^{107}] = 32809$$
,  $expected = 26854$ 

$$Int[(1.1)^{-3}] = 1$$
,  $expected = 0$ 

### Reliability Problems in 2020s

 Rowhammer: Repeatedly accessing a row enough times can cause disturbance errors in nearby rows



## Thank you!

Next Lecture: GPUs