Problem M1.1: Self-Modifying Code (Spring 2015 Quiz 1, Part A)

Problem M1.1.A

Memory:

	…
A	A[0]
	A[1]
	…
	A[n-1]
	…
ONE	1
N	n
I	0
TMP	0

Program:

```
loop:  LD      I
       SUB     N
       BGE     done

I1:    LD      A
       BGE     cont
       ST      TMP
       CLEAR
       SUB     TMP
I2:    ST      A

cont:  LD      I1
       ADD     ONE
       ST      I1
       LD      I2
       ADD     ONE
       ST      I2







       LD      I
       ADD     ONE
       ST      I
       BGE     loop
done:  END
```

Problem M1.1.B

Memory:

	…
A	A[0]
	A[1]
	…
	A[n-1]
	…
ONE	1
N	n
I	0
IDX	A

Program:

```
loop:  LD       I
       SUB      N
       BGE      done

       LDind    IDX
       BGE      cont
       CLEAR
       SUBind   IDX
       STind    IDX

cont:  LD       IDX
       ADD      ONE
       ST       IDX









       LD       I
       ADD      ONE
       ST       I
       BGE      loop
done:  END
```

Problem M1.2: Self-modifying Code (Spring 2017 Quiz 1, Part A)

Problem M1.2.A

```
.macro LISTPUSH
      STORE _TMP      ;; store accumulator (address of the new node)
      ADD _ONE        ;; accum <- address of the new node’s next field
      STOREADR _STN   ;; address field of location _STN has the address
                      ;; of the new node’s next field
      CLEAR
      ADD _HEAD       ;; accum <- M[_HEAD], current head pointer
_STN: STORE 0         ;; 0 will be replaced with the node’s next field
                      ;; address. M[_TMP + 1] <- accum
      CLEAR
      ADD _TMP        ;; retrieve address of new node in accumulator
      STORE _HEAD     ;; M[_HEAD] <- accum; Update the head pointer
                      ;; to the new node








.end
```

Problem M1.2.B

Write a macro for **LISTPOP**, which removes the node at the head of the list and stores its address in the accumulator, or stores _INVALID (-1) in the accumulator if the list is empty. Implement the macro using the EDSACjr instruction set and macros provided above.

```
.macro LISTPOP
     CLEAR              ;; accumulator is not an input
     LD _HEAD           ;; accum <- M[_HEAD]
     STOREADDR _LH      ;;
_LH: LD 0               ;; accum <- M[M[_HEAD]]
     BLT _DONE          ;; if the head points to _INVALID, we’re done

     LD _HEAD           ;; accum contains a pointer to the old head
     STORE _TMP         ;; save old value of head
     ADD _ONE           ;; accum <- address of head node’s next field
     STOREADR _ADDN     ;; replace address field of _ADDN
                        ;; with address of head node’s next field
_ADDN: LD 0             ;; 0 will be replaced with the address of head
                        ;; node’s next field. accum <- addr of 2nd node
     STORE _HEAD        ;; update head with the list’s second node
     LD _TMP            ;; accum <- former head node pointer
_DONE:
```

```
.end
```

Problem M1.2.C

Assume there exists a macro called **FREE** that takes an address as input in the accumulator and deallocates it (just like `free(void* ptr)` in C). Write a macro for **LISTCLEAR**, which uses the **FREE** macro and your **LISTPOP** macro to remove and deallocate all nodes in the list. Assume all valid node addresses are positive, or else a pointer is `_INVALID` (-1). Implement the macro using the EDSACjr instruction set and macros provided above.

```
.macro LISTCLEAR
_LOOP:     LISTPOP         ; accum <- address of removed node
           BLT _DONE       ; exit if an _INVALID node pointer is found
           FREE            ; de-allocates the removed node
           CLEAR
           BGE _LOOP
_DONE:








.end
```

Problem M1.3: Self Modifying Code on the EDSACjr

Problem M1.3.A — Writing Macros For Indirection

One way to implement `ADDind n` is as follows:

```
.macro ADDind(n)
      STORE       orig_accum  ; Save original accum
      CLEAR                   ; accum <- 0
      ADD         n           ; accum <- M[n]
      ADD         _add_op     ; accum <- ADD M[n]
      STORE       _L1         ; M[_L1] <- ADD M[n]
      CLEAR                   ; accum <- 0
_L1:  CLEAR                   ; This will be replaced by
                              ; ADD M[n] and will have
                              ; the effect: accum <- M[M[n]]
      ADD         _orig_accum ; accum <- M[M[n]] + original accum
.end macro
```

The first thing we do is save the original accumulator value. This is necessary since the instructions we are going to use within the macro are going to destroy the value in the accumulator. Next, we load the contents of `M[n]` into the accumulator. We assume that M[n] is a legal address and fits in 11 bits.

After getting the value of `M[n]` into the accumulator, we add it to the `ADD` template at `_add_op`. Since the template has 0 for its operand, the resulting number will have the `ADD` opcode with the value of `M[n]` in the operand field, and thus will be equivalently an `ADD M[n].` By storing the contents of the accumulator into the address `_L1`, we replace the `CLEAR` with what is equivalently an `ADD M[n]` instruction. Then we clear the accumulator so that when the instruction at `_L1` is executed, `accum` will get `M[M[n]]`. Finally, we add the original accumulator value to get the desired result, `M[M[n]]` plus the original content of the accumulator.

`STOREind n` can be implemented in a very similar manner.

```
.macro STOREind(n)
      STORE       _orig_accum ; Save original accum
      CLEAR                   ; accum <- 0
      ADD         n           ; accum <- M[n]
      ADD         _store_op   ; accum <- STORE M[n]
      STORE       _L1         ; M[_L1] <- STORE M[n]
      CLEAR                   ; accum <- 0
      ADD         _orig_accum ; accum <- original accum
_L1:  CLEAR                   ; This will be replaced by
                              ; STORE M[n], and will have the
                              ; effect: M[M[n]]<- orig. accum
.end macro
```

After getting the value of `M[n]` into the accumulator, we add it to the `STORE` template at `_store_op`. Since the template has 0 for its operand, the resulting number will have the `STORE` opcode with the value of `M[n]` in the operand field, and thus will be equivalently a `STORE M[n]` instruction. As before, we store this into `_L1` and then restore the accumulator value to its original value. When the PC reaches `_L1`, it then stores the original value of the accumulator into `M[M[n]]`.

`BGEind` and `BLTind` are very similar to `STOREind`. `BGEind` is shown below. `BLTind` is the same except that we use `_blt_op` instead of `_bge_op`.

```
.macro BGEind(n)
      STORE       _orig_accum ; Save original accum
      CLEAR                   ; accum <- 0
      ADD         n           ; accum <- M[n]
      ADD         _bge_op     ; acuum <- BGE M[n]
      STORE       _L1         ; M[_L1] <- BGE M[n]
      CLEAR                   ; accum <- 0
      ADD         _orig_accum ; accum <- original accum
_L1:  CLEAR                   ; This is replaced by BGE M[n]
.end macro
```

Problem M1.3.B Subroutine Calling Conventions

We implement the following contract between the caller and the callee:

1. The caller places the argument in the address slot between the function-calling jump instruction and the return address. Just before jumping to the subroutine, the caller loads the return address into the accumulator.
2. In the beginning of a subroutine, the callee receives the return address in the accumulator. The argument can be accessed by reading the memory location preceding the return address. The code below shows pass-by-value as we create a local copy of the argument. Since the subroutine receives the address of the argument, it's easy to eliminate the dereferencing and deal only with the address in a pass-by-reference manner.
3. When the computation is done, the callee puts the return value in the accumulator and then jumps to the return address.

A call looks like

```
            ......                  ; preceding code sequence
            clear
            add         _THREE      ; accum <- 3
            bge         _here       ; skip over pointer
_hereptr    .fill       _here       ; hereptr = &here
_here       add         _hereptr    ; accum <- here+3 = return addr
            bge         _sub        ; jump to subroutine
                                    ; The following address location is
                                    ; reserved for argument passing and
                                    ; should never be executed as code:
_argument   .fill 6                 ; argument slot
            ......                  ; rest of program
```

(note that without an explicit program counter, a little work is required to establish the return address).

The subroutine begins:

```
_sub        store       _return     ; save the return address
            sub         _ONE        ; accum <- &argument = return address-1
            store       _arg        ; M[_arg] <- &argument = return address-1
            clear
            ADDind      _arg        ; accum <- *(&arg0)
            store       _arg        ; M[_arg] <- arg
```

And ends (with the return value in the accumulator):

```
            BGEind      _return
```

The subroutine uses some local storage:

```
_arg        clear                     ; local copy of argument
_return     clear                     ; reserved for return address
```

We need the following global constants:

```
_ONE        or          1             ; recall that OR's opcode is 00000
_THREE      or          3             ; so positive constants are easy to form
```

The following program uses this convention to compute fib(n) as specified in the problem set. It uses the indirection macros, templates, and storage from part M1.3.A.

```
;; The Caller Code Section
;;    ......                          ; preceding code sequence
_caller     clear
            add         _THREE        ; accum <- 3
            bge         _here
_hereptr    .fill       _here
_here       add         _hereptr      ; accum <- here+3 = return addr
            bge         _fib          ; jump to subroutine

;; The following address location is reserved for
;; argument passing and should never be executed as code
arg0        .fill       4             ; arg 0 slot.  N=4 in this example

_rtpnt      end

;; The fib Subroutine Code Section

; function call prelude
_fib        store       _return       ; save the return address
            sub         _ONE
            store       _n            ; M[_n] <- &arg0 = return address-1
            clear
            ADDind      _n            ; accum <- *(&arg0)
            store       _n            ; M[_n] <- arg0

; fib body
            clear
            store       _x            ; x=0
            add         _ONE
            store       _y            ; y=1

            clear                     ; if(n<2)
            add         _n
            sub         _TWO
            blt         _retn

            clear
            store       _i            ; for (i = 0;

_forloop    clear                     ; i < n-1;
            add         _n
            sub         _ONE
            sub         _i
            sub         _ONE
            blt         _done
_compute    clear
```

```
            add         _x
            add         _y
            store       _z          ; z = x+y
            clear
            add         _y
            store       _x          ; x = y
            clear
            add         _z
            store       _y          ; y = z

_next       clear                   ; i++)
            add         _i
            add         _ONE
            store       _i
            bge         _forloop

_retn       clear
            add   _n
            BGEind      _return     ; return n

_done       clear
            add         _z
            BGEind      _return     ; return z

;; Global constants (remember that OR's opcode is 00000)

_ONE        or 1
_TWO        or 2
_THREE      or 3
_FOUR       or 4

These memory locations are private to the subroutine

_return     clear       ; return address
_n          clear       ; n
_x          clear
_y          clear
_z          clear
_i          clear       ; index
_result     clear       ; fib
```

Now we can see how powerful this indirection addressing mode is! It makes programming much simpler.

The 1 argument-1 result convention could be extended to variable number of arguments and results by

1. Leaving as many argument slots in the caller code between the subroutine call instruction and the return address. This works as long as both the caller and callee agree on how many arguments are being passed.
2. Multiple results can be returned as a pointer to a vector (or a list) of the results. This implies an indirection, and so, yet another chance for self-modifying code.

Problem M1.3.C Subroutine Calling Other Subroutines

The subroutine calling convention implemented in Problem M1.3.B stores the return address in a fixed memory location (`_return`). When `fib_recursive` is first called, the return address is stored there. However, this original return address will be overwritten when `fib_recursive` makes its first recursive call. Therefore, your program can never return to the original caller!