# 1 Misc

#### 1.1 Xor-Basis

```
vector<int> basis:
void add vector(int x){
      for (auto v : basis){
         x=min(x, x^v);
     if (x) basis.push_back(x);
9 | // 給一數字集合 S · 求能不能 XOR 出 x
10 bool check(int x){
     for (auto v : basis){
         x=min(x, x^v);
     return x;
17 / / 給一數字集合 S · 求能 XOR 出多少數字
18 // 答案等於 2^{basis 的大小}
20 | // 給一數字集合 S · 求 XOR 出最大的數字
21 int get_max(){
     int ans=0;
     for (auto v : basis){
         ans=max(ans, ans^v);
     return ans;
```

#### 1.2 Default-Code

```
1 #include <bits/stdc++.h>
 #define int long long
 #define ALL(x) x.begin(), x.end()
 #define SZ(x) ((int)x.size())
  #define fastio ios::sync with stdio(0), cin.tie(0);
 using namespace std;
 #ifdef LOCAL
 #define cout cout << "\033[0;32m"
10 #define cerr cerr << "\033[0;31m"
#define endl "\n" << "\033[0m"
#pragma GCC optimize("03,unroll-loops")
14 #pragma GCC target("avx,avx2,sse,sse2,sse3,sse4,popcnt")
15 #define endl "\n"
16 #endif
18 const int MAX_N = 5e5+10;
19 const int INF = 2e18;
 void solve1(){
      return;
26 signed main(){
```

```
27

28 fastio;

29 int t = 1;

30 while (t--){

32 solve1();

33 }

34 return 0;

36 }
```

#### 1.3 Radix-Sort

```
1 // 值域限制:0~1073741823(2^30-1)
 inline void radix_sort(vector<int> &a, int n){
     static int cnt[32768] = {0};
      vector<int> tmpa(n);
      for(int i = 0; i < n; ++i)</pre>
         ++cnt[a[i] & 32767];
      for(int i = 1; i < 32768; ++i)
         cnt[i] += cnt[i-1];
      static int temp;
      for(int i = n-1; i >= 0; --i){
         temp = a[i] & 32767;
          --cnt[temp]:
          tmpa[cnt[temp]] = a[i];
      static int cnt2[32768] = {0};
      for(int i = 0; i < n; ++i)
         ++cnt2[(tmpa[i]>>15)];
     for(int i = 1; i < 32768; ++i)</pre>
         cnt2[i] += cnt2[i-1];
      for(int i = n-1; i >= 0; --i){
         temp = (tmpa[i]>>15);
          --cnt2[temp];
         a[cnt2[temp]] = tmpa[i];
      return;
```

# 1.4 Set-Pq-Sort

```
1  // priority_queue
2  struct cmp{
      bool operator () (Data a, Data b){
          return a.x<b.x;
      }
6  };
7  priority_queue<Data, vector<Data>, cmp> pq;
8  
9  // set
struct Data{
    int x;
12  
13     bool operator < (const Data &b){
          return x<b.x;
15     }
16  };</pre>
```

#### 1.5 2-SAT

```
| #include <bits/stdc++.h>
  using namespace std;
  struct TWO_SAT {
      int n, N;
      vector<vector<int>> G, rev_G;
      deque<bool> used;
      vector<int> order, comp;
      deque<bool> assignment;
      void init(int n) {
          n = _n;
          N = n * 2;
          G.resize(N + 5);
          rev_G.resize(N + 5);
      void dfs1(int v) {
          used[v] = true;
          for (int u : G[v]) {
              if (!used[u])
                  dfs1(u);
          order.push_back(v);
23
24
      void dfs2(int v, int cl) {
          comp[v] = c1;
          for (int u : rev G[v]) {
              if (comp[u] == -1)
                  dfs2(u, c1);
      bool solve() {
          order.clear();
          used.assign(N, false);
          for (int i = 0; i < N; ++i) {
               if (!used[i])
                  dfs1(i);
          comp.assign(N, -1);
          for (int i = 0, j = 0; i < N; ++i) {
               int v = order[N - i - 1];
               if (comp[v] == -1)
                  dfs2(v, j++);
          assignment.assign(n, false);
          for (int i = 0; i < N; i += 2) {
               if (comp[i] == comp[i + 1])
                  return false;
               assignment[i / 2] = (comp[i] > comp[i + 1]);
50
          return true;
51
52
      void add_disjunction(int a, bool na, int b, bool nb) { //
          // na means whether a is negative or not
          // nb means whether b is negative or not
          a = 2 * a ^ na;
55
          b = 2 * b ^ nb;
          int neg_a = a ^ 1;
          int neg_b = b ^ 1;
58
          G[neg_a].push_back(b);
59
          G[neg_b].push_back(a);
60
          rev G[b].push back(neg a);
          rev_G[a].push_back(neg_b);
```

```
return;
       void get result(vector<int>& res) {
           res.clear();
           for (int i = 0; i < n; i++)</pre>
               res.push_back(assignment[i]);
70 };
71 /* CSES Giant Pizza
72 3 5
|73| + 1 + 2
|74| - 1 + 3
   - + + + -
77 */
78 int main() {
       int n, m;
       cin >> n >> m;
       TWO SAT E;
       E.init(m);
       char c1, c2;
       int inp1, inp2;
       for (int i = 0; i < n; i++) {</pre>
           cin >> c1 >> inp1;
           cin >> c2 >> inp2;
           E.add disjunction(inp1 - 1, c1 == '-', inp2 - 1, c2
                == '-');
       bool able = E.solve();
       if (able) {
           vector <int> ans;
           E.get_result(ans);
           for (int i : ans)
               cout << (i == true ? '+' : '-') << ' ';
           cout << ' \ n';
       } else {
           cout << "IMPOSSIBLE\n";</pre>
101
102
       return 0;
```

#### 1.6 Enumerate-Subset

# 1.7 Fast-Input

```
1 // fast IO
2 // 6f8879
```

```
3 inline char readchar(){
      static char buffer[BUFSIZ], * now = buffer + BUFSIZ, *
           end = buffer + BUFSIZ:
      if (now == end)
          if (end < buffer + BUFSIZ)</pre>
              return EOF;
          end = (buffer + fread(buffer, 1, BUFSIZ, stdin));
          now = buffer;
      return *now++;
 inline int nextint(){
      int x = 0, c = readchar(), neg = false;
      while (('0' > c | c > '9') \&\& c! = '-' \&\& c! = EOF) c =
           readchar();
      if(c == '-') neg = true, c = readchar();
      while ('0' <= c && c <= '9') x = (x << 3) + (x << 1) + (c^{0}) 24
           , c = readchar();
      if(neg) x = -x;
      return x; // returns 0 if EOF
```

# 1.8 setup

```
se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
:inoremap " ""<Esc>i
:inoremap {<CR> {<CR>}<Esc>ko
:inoremap {{ {}}<ESC>i
function! F(...)
 execute '!./%:r < ./' . a:1
endfunction
command! -nargs=* R call F(<f-args>)
map <F8> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -Wextra -
     Wshadow -02 -DLOCAL -g -fsanitize=undefined,address<CR>
map <F9> :!./%:r<CR>
map <F10> :!./%:r < ./%:r.in<CR>
ca hash w !cpp -dD -P -fpreprocessed \| tr -d "[:space:]" \|
    md5sum \| cut -c-6
" i+<esc>25A---+<esc>
" o|<esc>25A |<esc>
" "ggVGyG35pGdd
```

# 2 Convlution

# 2.1 FFT

```
typedef complex < double > cd;
const double PI = acos(-1);

void FFT(vector < cd > &a , bool inv){
   int n = a.size();
}
```

```
for (int i=1, j=0; i<n; i++){</pre>
           int bit = (n>>1);
           for ( ; j&bit ; bit>>=1){
               j ^= bit;
           i ^= bit;
           if (i<j){
               swap(a[i], a[j]);
17
18
      for (int len=2 ; len<=n ; len<<=1){</pre>
           cd wlen = polar(1.0, (inv ? 2 : -2)*PI/len);
          for (int i=0 ; i<n ; i+=len){</pre>
               cd w(1);
               for (int j=0 ; j<len/2 ; j++){</pre>
                   cd u = a[i+j];
                   cd v = a[i+j+len/2]*w;
                   a[i+j] = u+v;
                   a[i+j+len/2] = u-v;
                   w *= wlen;
      }
      if (inv){
           for (auto &x : a){
               x /= n;
      return;
  vector<cd> polyMul(vector<cd> a, vector<cd> b){
      int sa = a.size(), sb = b.size(), n = 1;
      while (n<sa+sb-1) n *= 2;</pre>
      a.resize(n);
      b.resize(n);
      vector<cd> c(n);
      FFT(a, 0);
      FFT(b, 0);
      for (int i=0; i<n; i++) c[i] = a[i]*b[i];</pre>
      FFT(c, 1);
      c.resize(sa+sb-1);
      return c;
```

#### 2.2 FFT-2

```
typedef complex<double> cd;

void FFT(vector<cd> &a) {
   int n = a.size(), L = 31-__builtin_clz(n);
   vector<complex<long double>> R(2, 1);
   vector<cd> rt(2, 1);
   for (int k=2; k<n; k*=2){</pre>
```

```
R.resize(n);
       rt.resize(n);
        auto x = polar(1.0L, acos(-1.0L) / k):
        for (int i=k ; i<2*k ; i++){</pre>
            rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);
   }
   vector<int> rev(n);
   for (int i=0 ; i<n ; i++){</pre>
        rev[i] = (rev[i/2] | (i&1) << L)/2;
   for (int i=0 ; i<n ; i++){</pre>
       if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
   for (int k=1; k<n; k*=2){
        for (int i=0 ; i<n ; i+=2*k){</pre>
            for (int j=0; j<k; j++){
                auto x = (double *)&rt[j+k];
                auto y = (double *)&a[i+j+k];
                cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]*
                     y[0]);
                a[i+j+k] = a[i+j]-z;
                a[i+j] += z;
   return;
vector<double> PolyMul(const vector<double> a, const vector<</pre>
    double> b){
   if (a.empty() || b.empty()) return {};
   vector<double> res(a.size()+b.size()-1);
   int L = 32 - builtin clz(res.size()), n = 1 << L;</pre>
   vector<cd> in(n), out(n);
   copy(a.begin(), a.end(), begin(in));
   for (int i=0 ; i<b.size() ; i++){</pre>
        in[i].imag(b[i]);
   FFT(in);
   for (cd& x : in) x *= x;
   for (int i=0 ; i<n ; i++){</pre>
       out[i] = in[-i & (n - 1)] - conj(in[i]);
   FFT(out);
   for (int i=0 ; i<res.size() ; i++){</pre>
       res[i] = imag(out[i]) / (4 * n);
   return res;
```

#### 2.4 FFT-mod

#### 2.3 NTT-998244353

```
const int MOD = (119 << 23) + 1, ROOT = 62; // = 998244353

// For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 <<

21

// and 483 << 21 (same root). The Last two are > 10^9.

// 9cd58a
```

```
6 | void NTT(vector<int> &a) {
      int n = a.size();
      int L = 31- builtin clz(n);
      vector<int> rt(2, 1);
      for (int k=2, s=2; k<n; k*=2, s++){
          rt.resize(n):
          int z[] = {1, qp(ROOT, MOD>>s)};
          for (int i=k ; i<2*k ; i++){</pre>
              rt[i] = rt[i/2]*z[i&1]%MOD;
     }
      vector<int> rev(n);
      for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2]|(i&1)<<L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]){</pre>
              swap(a[i], a[rev[i]]);
         }
     }
      for (int k=1; k<n; k*=2){</pre>
          for (int i=0 ; i<n ; i+=2*k){</pre>
              for (int j=0 ; j<k ; j++){</pre>
                  int z = rt[j+k]*a[i+j+k]%MOD, &ai = a[i+j];
                  a[i+j+k] = ai-z+(z>ai ? MOD : 0);
                  ai += (ai+z)=MOD ? z-MOD : z);
 vector<int> polyMul(vector<int> &a, vector<int> &b){
      if (a.empty() || b.empty()) return {};
      int s = a.size()+b.size()-1, B = 32-__builtin_clz(s), n =
           1<<B:
      int inv = qp(n, MOD-2);
      vector<int> L(a), R(b), out(n);
      L.resize(n), R.resize(n);
      NTT(L), NTT(R);
      for (int i=0 ; i<n ; i++){</pre>
          out[-i&(n-1)] = L[i]*R[i]%MOD*inv%MOD;
      NTT(out);
      out.resize(s);
      return out;
```

```
8 typedef complex < double > cd;
  // b9c90a
  void FFT(vector<cd> &a) {
      int n = a.size(), L = 31-__builtin_clz(n);
      vector<complex<long double>> R(2, 1);
      vector<cd> rt(2, 1);
      for (int k=2; k<n; k*=2){</pre>
          R.resize(n);
          rt.resize(n);
          auto x = polar(1.0L, acos(-1.0L) / k);
          for (int i=k ; i<2*k ; i++){</pre>
              rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);
      }
      vector<int> rev(n);
      for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2] | (i&1) << L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
      for (int k=1; k<n; k*=2){
          for (int i=0 ; i<n ; i+=2*k){</pre>
              for (int j=0 ; j<k ; j++){</pre>
                  auto x = (double *)&rt[j+k];
                  auto y = (double *)&a[i+j+k];
                  cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]*
                       y[0]);
                  a[i+j+k] = a[i+j]-z;
                  a[i+j] += z;
      return;
  vector<int> PolyMul(vector<int> a, vector<int> b){
      if (a.empty() || b.empty()) return {};
      vector<int> res(a.size()+b.size()-1);
      int B = 32- builtin_clz(res.size()), n = (1<<B), cut =</pre>
           int(sqrt(MOD));
      vector<cd> L(n), R(n), outs(n), outl(n);
      for (int i=0 ; i<a.size() ; i++){</pre>
          L[i] = cd((int) a[i]/cut, (int)a[i]%cut);
      for (int i=0 ; i<b.size() ; i++){</pre>
          R[i] = cd((int) b[i]/cut, (int)b[i]%cut);
      FFT(L);
      FFT(R);
      for (int i=0 ; i<n ; i++){</pre>
          int j = -i&(n-1);
          outl[j] = (L[i]+conj(L[j])) * R[i]/(2.0*n);
          outs[j] = (L[i]-conj(L[j])) * R[i]/(2.0*n)/1i;
      FFT(outl);
      FFT(outs);
      for (int i=0 ; i<res.size() ; i++){</pre>
          int av = (int)(real(outl[i])+0.5), cv = (int)(imag(
               outs[i])+0.5):
```

28

33

#### 3 Data-Structure

#### 3.1 GP-Hash-Table

```
| #include <ext/pb ds/assoc container.hpp>
using namespace gnu pbds;
typedef tree<int, null type, less<int>, rb tree tag,
      tree_order_statistics_node_update> order_set;
      static uint64 t splitmix64(uint64 t x) {
          // http://xorshift.di.unimi.it/splitmix64.c
          x += 0x9e3779b97f4a7c15;
          x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
          x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
          return x ^ (x >> 31);
      size t operator()(uint64 t x) const {
          static const uint64 t FIXED RANDOM = chrono::
               steady clock::now().time since epoch().count();
          return splitmix64(x + FIXED RANDOM);
16
17 };
19 gp_hash_table<int, int, custom_hash> ss;
```

# 3.2 Sparse-Table

```
| vector<vector<int>> st;
void build(vector<int>> v){
    int h = __lg(v.size());
    st.resize(h+1);
    st[0] = v;

    for (int i=1; i<=h; i++){
        int gap = (1<<(i-1));
        for (int j=0; j+gap<st[i-1].size(); j++){
            st[i].push_back(min(st[i-1][j], st[i-1][j+gap]));
        }
    }
}

// 回傳 [ll, rr) 的最小值
int RMQ(int ll, int rr){
    int RMQ(int ll, int rr){
    int h = __lg(rr-ll);
    return min(st[h][ll], st[h][rr-(1<<h)]);
}
```

#### 3.3 Order-Set

#### 3.4 BIT

```
vector<int> BIT(MAX SIZE);
  void update(int pos, int val){
      for (int i=pos ; i<MAX SIZE ; i+=i&-i){</pre>
          BIT[i]+=val;
  int query(int pos){
      int ret=0:
      for (int i=pos ; i>0 ; i-=i&-i){
          ret+=BIT[i]:
       return ret;
  // const int MAX_N = (1 << 20)
  int k th(int k){ // 回傳 BIT 中第 k 小的元素(based-1)
       int res = 0;
       for (int i=MAX N>>1 ; i>=1 ; i>>=1)
          if (bit[res+i]<k)</pre>
               k -= bit[res+=i];
       return res+1;
23 }
```

# 3.5 Persistent-Segment-Tree

```
int lc = -1;
    int rc = -1;
    int val = 0:
};
vector<Node> arr;
vector<int> version;
Persistent Segment Tree(int sz){
    arr.resize(32*sz);
    version.push_back(node_cnt++);
    return;
}
void pull(Node &c, Node a, Node b){
    c.val = a.val+b.val;
    return;
void build(vector<int> &v, int idx, int ll = 0, int rr =
    auto &now = arr[idx];
    if (rr-ll==1){
        now.val = v[11];
        return;
    int mid = (11+rr)/2;
    now.lc = node_cnt++;
    now.rc = node cnt++;
    build(v, now.lc, ll, mid);
    build(v, now.rc, mid, rr);
    pull(now, arr[now.lc], arr[now.rc]);
    return;
void update(int pos, int val, int idx, int ll = 0, int rr
    auto &now = arr[idx];
    if (rr-ll==1){
        now.val = val;
        return;
    int mid = (11+rr)/2;
    if (pos<mid){</pre>
        arr[node cnt] = arr[now.lc];
        now.lc = node cnt;
        node_cnt++;
        update(pos, val, now.lc, ll, mid);
        arr[node_cnt] = arr[now.rc];
        now.rc = node cnt;
        node cnt++;
        update(pos, val, now.rc, mid, rr);
    pull(now, arr[now.lc], arr[now.rc]);
    return:
void update version(int pos, int val, int ver){
    update(pos, val, version[ver]);
```

#### **3.6** Trie

```
| struct Trie{
     struct Data{
         int nxt[2]={0, 0};
     int sz=0:
     vector<Data> arr;
     void init(int n){
         arr.resize(n);
     void insert(int n){
         int now=0;
         for (int i=N ; i>=0 ; i--){
             int v=(n>>i)&1;
             if (!arr[now].nxt[v]){
                 arr[now].nxt[v]=++sz;
             now=arr[now].nxt[v];
     }
     int query(int n){
         int now=0, ret=0;
         for (int i=N ; i>=0 ; i--){
             int v=(n>>i)&1;
             if (arr[now].nxt[1-v]){
                 ret+=(1<<i);
                 now=arr[now].nxt[1-v];
             }else if (arr[now].nxt[v]){
                 now=arr[now].nxt[v];
             }else{
                 return ret;
```

# 3.7 LC-Segment-Tree

return ret;

}

40 } tr;

```
2| 全部都是 0-based
  LC Segment Tree st(n);
  update(val): 將一個 pair <a, b> 代表插入一條 y=ax+b 的直線
  query(x): 查詢所有直線在位置 x 的最小值
  const int MAX V = 1e6+10; // 值域最大值
  struct LC_Segment_Tree{
      struct Node{ // y = ax+b
          int a = 0:
          int b = INF;
         int y(int x){
             return a*x+b;
      };
      vector<Node> arr;
      LC_Segment_Tree(int n = 0){
          arr.resize(4*n);
27
      void update(Node val. int idx = 0, int ll = 0, int rr =
          MAX V){
          if (rr-ll==1){
             if (val.y(l1) < arr[idx].y(l1)){</pre>
                 arr[idx] = val;
             return;
         int mid = (11+rr)/2;
          if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
              的線斜率要比較小
          if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
             update(val, idx*2+1, ll, mid);
         }else{ // 交點在右邊
             swap(arr[idx], val); // 在左子樹中, 新線比舊線還
             update(val, idx*2+2, mid, rr);
          return;
      int query(int x, int idx = 0, int ll = 0, int rr = MAX_V)
          if (rr-ll==1){
             return arr[idx].y(ll);
```

# 3.8 Persistent-Disjoint-Set

59 };

int mid = (11+rr)/2;
if (x<mid){</pre>

rr));

return min(arr[idx].y(x), query(x, idx\*2+1, ll,

return min(arr[idx].y(x), query(x, idx\*2+2, mid,

```
struct Persistent Disjoint Set{
      Persistent_Segment_Tree arr, sz;
      void init(int n){
          arr.init(n);
          vector<int> v1;
          for (int i=0 ; i<n ; i++){</pre>
              v1.push_back(i);
          arr.build(v1, 0);
          sz.init(n);
          vector<int> v2;
          for (int i=0 ; i<n ; i++){</pre>
              v2.push_back(1);
          sz.build(v2, 0);
      int find(int a){
          int res = arr.query_version(a, a+1, arr.version.size
          if (res==a) return a;
          return find(res);
25
      bool unite(int a, int b){
          a = find(a):
          b = find(b);
          if (a!=b){
              int sz1 = sz.query version(a, a+1, arr.version.
                   size()-1).val;
              int sz2 = sz.query version(b, b+1, arr.version.
                   size()-1).val;
              if (sz1<sz2){
                  arr.update version(a, b, arr.version.size()
                  sz.update version(b, sz1+sz2, arr.version.
                       size()-1);
                  arr.update version(b, a, arr.version.size()
                  sz.update version(a, sz1+sz2, arr.version.
                       size()-1);
```

```
return true;
          return false:
46 };
       Add-Set-Segment-Tree
1 // [ll, rr), based-0
2 // 使用前記得 init(陣列大小), build(陣列名稱)
3 // add(LL, rr): 區間修改
4 // set(ll, rr): 區間賦值
5 // query(ll, rr): 區間求和 / 求最大值
6 struct SegmentTree{
      struct node{
          int add_tag = 0;
          int set_tag = 0;
          int sum = 0;
          int ma = 0;
      };
      vector<node> arr;
      SegmentTree(int n){
          arr.resize(n<<2);</pre>
      node pull(node A, node B){
          node C;
          C.sum = A.sum + B.sum;
          C.ma = max(A.ma, B.ma);
          return C;
      void push(int idx, int ll, int rr){
          if (arr[idx].set tag!=0){
              arr[idx].sum = (rr-ll)*arr[idx].set_tag;
              arr[idx].ma = arr[idx].set tag;
             if (rr-ll>1){
                  arr[idx*2+1].add tag = 0;
                  arr[idx*2+1].set tag = arr[idx].set tag;
                  arr[idx*2+2].add_tag = 0;
                  arr[idx*2+2].set tag = arr[idx].set tag;
              arr[idx].set_tag = 0;
          if (arr[idx].add tag!=0){
              arr[idx].sum += (rr-ll)*arr[idx].add tag;
              arr[idx].ma += arr[idx].add tag;
                  arr[idx*2+1].add tag += arr[idx].add tag;
                  arr[idx*2+2].add_tag += arr[idx].add_tag;
```

arr[idx].add tag = 0;

arr[idx].sum = v[11];

rr = n){ if (rr-ll==1){

void build(vector<int> &v, int idx = 0, int ll = 0, int

```
arr[idx].ma = v[ll];
           }else{
               int mid = (11+rr)/2:
               build(v, idx*2+1, ll, mid);
               build(v, idx*2+2, mid, rr);
               arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
       }
       void add(int ql, int qr, int val, int idx = 0, int ll =
            0, int rr =n){
           push(idx, ll, rr);
                                                                   21
           if (rr<=ql || qr<=ll) return;</pre>
           if (q1<=11 && rr<=qr){
               arr[idx].add_tag += val;
               push(idx, 11, rr);
               return;
                                                                   27
           int mid = (11+rr)/2;
                                                                   28
           add(ql, qr, val, idx*2+1, ll, mid);
           add(ql, qr, val, idx*2+2, mid, rr);
           arr[idx]=pull(arr[idx*2+1], arr[idx*2+2]);
       void set(int ql, int qr, int val, int idx=0, int ll=0,
            int rr=n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return;</pre>
           if (q1<=11 && rr<=qr){</pre>
               arr[idx].add_tag = 0;
               arr[idx].set_tag = val;
               push(idx, ll, rr);
               return;
           int mid = (11+rr)/2;
           set(ql, qr, val, idx*2+1, ll, mid);
           set(q1, qr, val, idx*2+2, mid, rr);
           arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
       node query(int ql, int qr, int idx = 0, int ll = 0, int
            rr = n){
           push(idx, 11, rr);
                                                                   51
           if (rr<=ql || qr<=ll) return node();</pre>
                                                                   52
           if (q1<=11 && rr<=qr) return arr[idx];</pre>
           int mid = (11+rr)/2;
           return pull(query(ql, qr, idx*2+1, ll, mid), query(ql
                , qr, idx*2+2, mid, rr));
100 } ST;
   3.10 Treap
  struct Treap{
       Treap *1 = nullptr, *r = nullptr;
                                                                   66 }
       int pri = rand(), val = 0, sz = 1;
       Treap(int _val){
           val = _val;
 8 };
```

```
int size(Treap *t){return t ? t->sz : 0;}
  void pull(Treap *t){
      t\rightarrow sz = size(t\rightarrow l) + size(t\rightarrow r) + 1:
  Treap* merge(Treap *a, Treap *b){
      if (!a | | !b) return a ? a : b;
      if (a->pri>b->pri){
          a \rightarrow r = merge(a \rightarrow r, b);
          pull(a):
           return a;
      }else{
          b\rightarrow 1 = merge(a, b\rightarrow 1);
          pull(b);
           return b;
  pair<Treap*, Treap*> split(Treap *&t, int k){ // 1-based <前
       k 個元素, 其他元素>
      if (!t) return {};
      if (size(t->1)>=k){
           auto pa = split(t->1, k);
          t->l = pa.second;
           return {pa.first, t};
          auto pa = split(t->r, k-size(t->l)-1);
          t->r = pa.first:
          pull(t);
           return {t, pa.second};
  // functions
  Treap* build(vector<int> v){
      Treap* ret;
      for (int i=0 ; i<SZ(v) ; i++){</pre>
           ret = merge(ret, new Treap(v[i]));
      return ret;
54 array<Treap*, 3> cut(Treap *t, int 1, int r){ // 1-based <前
       1~L-1 個元素, L~r 個元素, r+1 個元素>
      array<Treap*, 3> ret;
      tie(ret[1], ret[2]) = split(t, r);
      tie(ret[0], ret[1]) = split(ret[1], 1-1);
      return ret;
  void print(Treap *t, bool flag = true){
      if (t->1!=0) print(t->1, false);
      cout << t->val;
      if (t->r!=0) print(t->r, false);
      if (flag) cout << endl;</pre>
```

# 4 Dynamic-Programming

#### 4.1 SOS-DP

```
1 // 總時間複雜度為 O(n 2^n)
2 // 計算 dp[i] = i 所有 bit mask 子集的和
3 for (int i=0; i<n; i++){
4 for (int mask=0; mask<(1<<n); mask++){
5 if ((mask>>i)&1){
6 dp[mask] += dp[mask^(1<<ii)];
7 }
8 }
9 }
```

# 4.2 Digit-DP

i #include <bits/stdc++.h>

```
using namespace std;
4 long long 1, r;
5 long long dp[20][10][2][2]; // dp[pos][pre][limit] = 後 pos
      位 pos 前一位是 pre (是/否)有上界 (是/否)有前綴零
  long long memorize_search(string &s, int pos, int pre, bool
      limit, bool lead){
     // 已經被找過了,直接回傳值
     if (dp[pos][pre][limit][lead]!=-1) return dp[pos][pre][
          limit][lead];
     // 已經搜尋完畢,紀錄答案並回傳
     if (pos==(int)s.size()){
         return dp[pos][pre][limit][lead] = 1;
     // 枚舉目前的位數數字是多少
     long long ans = 0;
     for (int now=0 ; now<=(limit ? s[pos]-'0' : 9) ; now++){</pre>
         if (now==pre){
            // 1~9 絕對不能連續出現
            if (pre!=0) continue;
            // 如果已經不在前綴零的範圍內· Ø 不能連續出現
            if (lead==false) continue;
         ans += memorize search(s, pos+1, now, limit&(now==(s[
             pos]-'0')), lead&(now==0));
     }
     // 已經搜尋完畢,紀錄答案並回傳
     return dp[pos][pre][limit][lead] = ans;
36 | // 回傳 [0, n] 有多少數字符合條件
37 long long find answer(long long n){
     memset(dp, -1, sizeof(dp));
```

```
string tmp = to_string(n);

return memorize_search(tmp, 0, 0, true, true);

int main(){

// input
cin >> 1 >> r;

// output - 計算 [l, r] 有多少數字任意兩個位數都不相同
cout << find_answer(r)-find_answer(l-1) << "\n";

return 0;
```

## 4.3 整數拆分

```
dp[i][x] = 要將整數 x 拆成 i 堆的「組合數」 dp[i+1][x+1] + = dp[i][x] ( 創造新的一堆 ) \\ dp[i][x+i] + = dp[i][x] ( 把每一堆都增加 1 )
```

# 5 Geometry

#### 5.1 Point-Struct

const int EPS = 1e-6;

```
typedef int pt;
struct point{
    pt x, y;
    point(pt x = 0, pt y = 0)
       x = _x;
       y = y;
    point operator * (pt a){return {a*x, a*y};};
    point operator / (pt a){return {a/x, a/y};};
   // 向量加、減法
    point operator + (point a){return {x+a.x, y*a.y};};
    point operator - (point a){return {x-a.x, y-a.y};};
   // 內積、外積
   double operator * (point a){return x*a.x+y*a.y;};
   double operator ^ (point a){return x*a.y-y*a.x;};
   // bool operator < (const point &a) const {return (x*a.y< 11)
        a.x*y);} // 極角排序(順時鐘)
    bool operator < (const point &a) const {return x==a.x ? y 13
    bool operator == (const point &a) const {return x==a.x && 15
    double dis(point a){return sqrtl(abs(x-a.x)*abs(x-a.x)+
        abs(y-a.y)*abs(y-a.y));}
```

```
29 };
30
31 // 判斷向量正負: 1=正數, 0=0, -1=負數
31 int sign(double a){
32 int sign(double a){
33 if (abs(a)<EPS) return 0;
34 else return (a>0 ? 1 : -1);
35 }
36
37 // 判斷 ab 到 ac 的方向: 1=逆時鐘, 0=重疊, -1=順時鐘
38 int ori(point a, point b, point c){
39 return sign((b-a)^(c-a));
40 }
```

#### **5.2** Line-Intersection

#### 5.3 Pick's-Theorem

給定頂點坐標均是整點的簡單多邊形,面積 = 內部格點數 + 邊上格點數/2 - 1

# 5.4 Point-In-Polygon

#### 5.5 Convex-Hull

```
int cnt = 0;
      bool ap = 0:
      vis[now] = 1;
      low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
      for (auto x : G[now]){
          if (x==pre)
              continue:
          }else if (vis[x]==0){
              cnt++:
              dfs(x, now);
              low[now] = min(low[now], low[x]);
              if (low[x]>=dep[now]) ap=1;
              low[now] = min(low[now], dep[x]);
      }
      if ((now==pre && cnt>=2) || (now!=pre && ap)){
          AP.push back(now);
25
```

void dfs(int now, int pre){

```
// 〈最大流,最小費用〉
39
      pair<int, int> flow(int s, int t) {
40
          int fl = 0, cost = 0, d;
41
          while ((d = spfa(s, t)) < LONG LONG MAX) {</pre>
42
               int cur = LONG_LONG_MAX;
43
               for (int v = t; v != s; v = par[v])
                  cur = min(cur, G[par[v]][par_eid[v]].rc);
               cost += d * cur:
              for (int v = t; v != s; v = par[v]) {
                  G[par[v]][par_eid[v]].rc -= cur;
                  G[v][G[par[v]][par eid[v]].rv].rc += cur;
52
53
          return {fl, cost};
55
```

# 6 Graph

# 6.1 Find-Bridge

```
i | vector<int> dep(MAX N), low(MAX N);
vector<pair<int, int>> bridge;
 bitset<MAX N> vis;
 void dfs(int now, int pre){
     vis[now] = 1;
     low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
     for (auto x : G[now]){
         if (x==pre){
             continue:
         }else if (vis[x]==0){
             // 沒有走過的節點
             dfs(x, now);
             low[now] = min(low[now], low[x]);
         }else if (vis[x]==1){
             low[now] = min(low[now], dep[x]);
     if (now!=1 && low[now]==dep[now]){
         bridge.push back({now, pre});
     return:
```

# 6.2 Find-AP

```
1 vector<int> dep(MAX_N), low(MAX_N), AP;
2 bitset<MAX_N> vis;
```

## **6.3** MCMF

```
11 // frostrav 會用,所以他超強
2 struct Flow {
     struct Edge {
         int u, rc, k, rv;
     };
     vector<vector<Edge>> G;
     vector<int> par, par_eid;
     Flow(int n): G(n + 1), par(n + 1), par_eid(n + 1) {}
     // v->u / 流量:c, cost: k
     void add(int v, int u, int c, int k) {
         G[v].push_back({u, c, k, (int)G[u].size()});
         G[u].push_back({v, 0, -k, (int)G[v].size() - 1});
     int spfa(int s, int t) {
         fill(par.begin(), par.end(), -1);
         vector<int> dis(par.size(), LONG LONG MAX);
         vector<bool> in_q(par.size(), false);
         aueue<int> 0:
         dis[s] = 0; in_q[s] = true;
         0.push(s);
         while (! Q.empty()) {
             int v = Q.front(); Q.pop();
             in q[v] = false;
             for (int i = 0; i < (int)G[v].size(); i++) {</pre>
                  auto [u, rc, k, rv] = G[v][i];
                 if (rc > 0 \&\& dis[v] + k < dis[u]) {
                     dis[u] = dis[v] + k;
                     par[u] = v;
                     par_eid[u] = i;
                     if (! in_q[u]) Q.push(u);
                     in_q[u] = true;
                 }
             }
         return dis[t];
```

#### 6.4 HLD

```
| #include <bits/stdc++.h>
  #define int long long
  using namespace std;
  const int N = 100005;
  vector <int> G[N];
  struct HLD {
      vector<int> pa, sz, depth, mxson, topf, id;
      int n, idcnt = 0;
      HLD(int _n) : n(_n), pa(_n + 1), sz(_n + 1), depth(_n +
           1), mxson(_n + 1), topf(_n + 1), id(_n + 1) {}
      void dfs1(int v = 1, int p = -1) {
          pa[v] = p; sz[v] = 1; mxson[v] = 0;
12
          depth[v] = (p == -1 ? 0 : depth[p] + 1);
13
          for (int u : G[v]) {
              if (u == p) continue;
              dfs1(u, v);
              sz[v] += sz[u];
              if (sz[u] > sz[mxson[v]]) mxson[v] = u;
18
19
20
      void dfs2(int v = 1, int top = 1) {
21
          id[v] = ++idcnt;
          topf[v] = top;
          if (mxson[v]) dfs2(mxson[v], top);
          for (int u : G[v]) {
              if (u == mxson[v] || u == pa[v]) continue;
              dfs2(u, u);
28
29
      // auerv 為區間資料結構
      int path_query(int a, int b) {
          int res = 0;
          while (topf[a] != topf[b]) { /// 若不在同一條鍊上
33
              if (depth[topf[a]] < depth[topf[b]]) swap(a, b);</pre>
34
              res = max(res, 011); // query : L = id[topf[a]],
                  r = id[a]
36
              a = pa[topf[a]];
37
          /// 此時已在同一條鍊上
38
39
          if (depth[a] < depth[b]) swap(a, b);</pre>
```

```
res = max(res, 011); // query : l = id[b], r = id[a] 55
                                                                                                                                  120 }
                                                                            m[v]=++id;
          return res;
                                                                                                                                  121
                                                                            return id;
                                                                                                                                     signed main(void){
43 };
                                                                                                                                         fastio;
                                                                                                                                         int t=1:
                                                                                                                                         cin >> t;
 6.5 Tree-Isomorphism
                                                                    void solve1(){
                                                                                                                                         while (t--){
                                                                                                                                  128
                                                                                                                                             solve1();
                                                                        // init
                                                                                                                                  129
i #include <bits/stdc++.h>
                                                                        id1=0:
                                                                                                                                  130
                                                                                                                                         return 0;
#pragma GCC optimize("03,unroll-loops")
                                                                        id2=0;
                                                                                                                                  131 }
 #define fastio ios::sync with stdio(0), cin.tie(0), cout.tie
                                                                        c1={0, 0};
                                                                        c2={0, 0};
4 #define dbg(x) cerr << #x << " = " << x << endl
                                                                        fill(sz1.begin(), sz1.begin()+n+1, 0);
                                                                        fill(sz2.begin(), sz2.begin()+n+1, 0);
                                                                                                                                     6.6 Bridge BCC
  #define int long long
  using namespace std;
                                                                        fill(we1.begin(), we1.begin()+n+1, 0);
                                                                        fill(we2.begin(), we2.begin()+n+1, 0);
  // declare
                                                                        for (int i=1; i<=n; i++){</pre>
                                                                                                                                    | #include <bits/stdc++.h>
                                                                            g1[i].clear();
  const int MAX_SIZE = 2e5+5;
                                                                                                                                     using namespace std;
10 const int INF = 9e18;
                                                                            g2[i].clear();
11 const int MOD = 1e9+7;
                                                                                                                                     const int N = 200005;
                                                                        m1.clear();
  const double EPS = 1e-6;
                                                                                                                                     vector <int> G[N];
                                                                        m2.clear();
  typedef vector<vector<int>> Graph;
                                                                                                                                     int low[N], depth[N];
  typedef map<vector<int>, int> Hash;
                                                                                                                                     bool vis[N];
                                                                        // input
                                                                                                                                     vector <vector <int>> bcc;
                                                                        cin >> n;
16 int n, a, b;
                                                                                                                                     stack <int> stk;
                                                                        for (int i=0 ; i<n-1 ; i++){</pre>
17 int id1, id2;
18 pair<int, int> c1, c2;
                                                                            cin >> a >> b;
                                                                                                                                     void dfs(int v, int p) {
                                                                            g1[a].push_back(b);
 vector<int> sz1(MAX_SIZE), sz2(MAX_SIZE);
                                                                                                                                         stk.push(v);
                                                                            g1[b].push_back(a);
  vector<int> we1(MAX SIZE), we2(MAX SIZE);
                                                                                                                                         vis[v] = true;
 Graph g1(MAX SIZE), g2(MAX SIZE);
                                                                                                                                         low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
                                                                        for (int i=0 ; i<n-1 ; i++){</pre>
  Hash m1, m2;
                                                                                                                                         for (int u : G[v]) {
  int testcase=0;
                                                                            cin >> a >> b;
                                                                                                                                             if (u == p) continue;
                                                                            g2[a].push_back(b);
                                                                                                                                             if (!vis[u]) {
  void centroid(Graph &g, vector<int> &s, vector<int> &w, pair
                                                                            g2[b].push back(a);
                                                                                                                                                 /// (v, u) 是樹邊
      int, int> &rec, int now, int pre){
                                                                                                                                                 dfs(u, v);
      s[now]=1;
                                                                                                                                                 low[v] = min(low[v], low[u]);
      w[now]=0;
                                                                        // get tree centroid
                                                                                                                                             } else {
      for (auto x : g[now]){
                                                                        centroid(g1, sz1, we1, c1, 1, 0);
                                                                                                                                                 /// (v, u) 是回邊
          if (x!=pre){
                                                                        centroid(g2, sz2, we2, c2, 1, 0);
                                                                                                                                                 low[v] = min(low[v], depth[u]);
              centroid(g, s, w, rec, x, now);
                                                                        // process
              s[now]+=s[x];
                                                                                                                                   25
                                                                        int res1=0, res2=0, res3=0;
              w[now]=max(w[now], s[x]);
                                                                                                                                         /// v 在不依靠父邊的情況下永遠沒辦法走到它的祖先
                                                                        if (c2.second!=0){
                                                                                                                                         if (low[v] == depth[v]) {
                                                                            res1=dfs(g1, m1, id1, c1.first, 0);
      }
                                                                                                                                             bcc.emplace back();
                                                                            m2=m1:
                                                                                                                                             while (stk.top() != v) {
                                                                            id2=id1;
      w[now]=max(w[now], n-s[now]);
                                                                                                                                                 bcc.back().push_back(stk.top());
                                                                            res2=dfs(g2, m1, id1, c2.first, 0);
      if (w[now]<=n/2){</pre>
                                                                                                                                                 stk.pop();
          if (rec.first==0) rec.first=now;
                                                                            res3=dfs(g2, m2, id2, c2.second, 0);
          else rec.second=now;
                                                                        }else if (c1.second!=0){
                                                                                                                                             bcc.back().push_back(stk.top());
                                                                            res1=dfs(g2, m1, id1, c2.first, 0);
                                                                                                                                             stk.pop();
                                                                            m2=m1;
                                                                                                                                   35
                                                                            id2=id1:
  int dfs(Graph &g, Hash &m, int &id, int now, int pre){
                                                                            res2=dfs(g1, m1, id1, c1.first, 0);
      vector<int> v;
                                                                            res3=dfs(g1, m2, id2, c1.second, 0);
      for (auto x : g[now]){
                                                                111
          if (x!=pre){
                                                                112
                                                                            res1=dfs(g1, m1, id1, c1.first, 0);
                                                                                                                                     6.7 Cut BCC
              int add=dfs(g, m, id, x, now);
                                                                            res2=dfs(g2, m1, id1, c2.first, 0);
                                                                113
              v.push_back(add);
                                                                114
                                                                115
                                                                                                                                     #include <bits/stdc++.h>
                                                                        cout << (res1==res2 | res1==res3 ? "YES" : "NO") << endl 2
      sort(v.begin(), v.end());
                                                                                                                                     using namespace std;
      if (m.find(v)!=m.end()){
                                                                                                                                     const int N = 200005;
```

5 vector <int> G[N];

return;

return m[v];

```
6| int low[N], depth[N];
7 bool vis[N];
8 vector <vector <int>> bcc:
  stack <int> stk;
  void dfs(int v, int p) {
     stk.push(v);
     vis[v] = true;
     low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
     for (int u : G[v]) {
          if (u == p) continue;
          if (!vis[u]) {
             /// (v, u) 是樹邊
             dfs(u, v);
             low[v] = min(low[v], low[u]);
             /// u 無法在不經過父邊的情況走到 v 的祖先
             if (low[u] >= depth[v]) {
                 bcc.emplace_back();
                 while (stk.top() != u) {
                     bcc.back().push_back(stk.top());
                     stk.pop();
                 bcc.back().push back(stk.top());
                 stk.pop();
                 bcc.back().push_back(v);
         } else {
             /// (v, u) 是回邊
             low[v] = min(low[v], depth[u]);
34
```

#### 6.8 圓方樹

```
i #include <bits/stdc++.h>
2 #define lp(i,a,b) for(int i=(a);i<(b);i++)</pre>
  #define pii pair<int,int>
  #define pb push back
  #define ins insert
  #define ff first
  #define ss second
  #define opa(x) cerr << #x << " = " << x << ", ";
  #define op(x) cerr << #x << " = " << x << endl:
10 #define ops(x) cerr << x;</pre>
#define etr cerr << endl;
12 #define spc cerr << ' ';
#define BAE(x) (x).begin(), (x).end()
#define STL(x) cerr << #x << " : "; for(auto &qwe:x) cerr <<</pre>
qwe << ' '; cerr << endl;
15 #define deb1 cerr << "deb1" << endl;</pre>
#define deb2 cerr << "deb2" << endl;</pre>
#define deb3 cerr << "deb3" << endl;
#define deb4 cerr << "deb4" << endl;
#define deb5 cerr << "deb5" << endl;</pre>
20 #define bye exit(0);
21 using namespace std;
23 const int mxn = (int)(2e5) + 10;
24 const int mxlg = 17:
25 int last special node = (int)(1e5) + 1;
26 vector<int> E[mxn], F[mxn];
```

```
struct edg{
    int fr. to:
    edg(int _fr, int _to){
        fr = fr;
        to = to;
ostream& operator<<(ostream& os, edg x){os << x.fr << "--" << 99
     x.to;}
vector<edg> EV;
void tarjan(int v, int par, stack<int>& S){
    static vector<int> dfn(mxn), low(mxn);
    static vector<bool> to_add(mxn);
    static int nowT = 0;
    int childs = 0;
    nowT += 1:
    dfn[v] = low[v] = nowT;
    for(auto &ne:E[v]){
        int i = EV[ne].to;
        if(i == par) continue;
        if(!dfn[i]){
            S.push(ne);
            tarjan(i, v, S);
                                                               116
            childs += 1;
            low[v] = min(low[v], low[i]);
            if(par >= 0 \&\& low[i] >= dfn[v]){
                vector<int> bcc;
                int tmp;
                do{
                    tmp = S.top(); S.pop();
                    if(!to add[EV[tmp].fr]){
                        to add[EV[tmp].fr] = true;
                        bcc.pb(EV[tmp].fr);
                    if(!to_add[EV[tmp].to]){
                        to_add[EV[tmp].to] = true;
                        bcc.pb(EV[tmp].to);
                }while(tmp != ne);
                for(auto &j:bcc){
                    to_add[j] = false;
                    F[last special node].pb(j);
                    F[j].pb(last special node);
                last special node += 1;
        else{
            low[v] = min(low[v], dfn[i]);
            if(dfn[i] < dfn[v]){ // edge i--v will be visited 144
                  twice at here, but we only need one.
                S.push(ne);
                                                               147
        }
int dep[mxn], jmp[mxn][mxlg];
void dfs lca(int v, int par, int depth){
    dep[v] = depth;
    for(auto &i:F[v]){
        if(i == par) continue;
```

```
jmp[i][0] = v;
           dfs_lca(i, v, depth + 1);
94
   inline void build lca(){
       imp[1][0] = 1;
       dfs lca(1, -1, 1);
       lp(j,1,mxlg){
           lp(i,1,mxn){
               jmp[i][j] = jmp[jmp[i][j-1]][j-1];
102
103
104
105
   inline int lca(int x, int y){
       if(dep[x] < dep[y]){ swap(x, y); }</pre>
109
       int diff = dep[x] - dep[y];
       lp(j,0,mxlg){
110
           if((diff >> j) & 1){
111
               x = jmp[x][j];
112
113
114
115
       if(x == y) return x;
       for(int j = mxlg - 1; j >= 0; j--){
117
           if(jmp[x][j] != jmp[y][j]){
118
               x = jmp[x][j];
119
               y = jmp[y][j];
121
122
123
       return jmp[x][0];
124
125
   inline bool can reach(int fr, int to){
126
       if(dep[to] > dep[fr]) return false;
       int diff = dep[fr] - dep[to];
129
130
       lp(j,0,mxlg){
           if((diff >> j) & 1){
131
132
                fr = jmp[fr][j];
133
134
135
       return fr == to;
136
137
       ios::sync with stdio(false); cin.tie(0);
      freopen("test_input.txt", "r", stdin);
       int n, m, q; cin >> n >> m >> q;
       lp(i,0,m){
           int u, v; cin >> u >> v;
           E[u].pb(EV.size());
           EV.pb(edg(u, v));
           E[v].pb(EV.size());
           EV.pb(edg(v, u));
148
       E[0].pb(EV.size());
149
150
       EV.pb(edg(0, 1));
       stack<int> S;
151
152
       tarjan(0, -1, S);
       build_lca();
153
154
       lp(aueries.0.a){
           int fr, to, relay; cin >> fr >> to >> relay;
```

```
if(fr == relay || to == relay){
158
               cout << "NO\n";</pre>
                                                                             order.push back(now);
                                                                                                                                           Flow(int n){
159
               continue:
                                                                             return:
                                                                                                                                               G.resize(n):
                                                                         }
                                                                                                                                               dis.resize(n);
160
161
           if((can_reach(fr, relay) || can_reach(to, relay)) &&
                                                                                                                                               it.resize(n);
               dep[relay] >= dep[lca(fr, to)]){
                                                                         void dfs2(vector<vector<int>> &G, int now){
                                                                             SCC id[now] = SCC.size()-1;
162
               cout << "NO \ n";
               continue:
                                                                             SCC.back().push back(now);
                                                                                                                                           int dfs(int u, int t, int f){
163
                                                                                                                                               if (u==t || f==0) return f;
                                                                             vis[now] = 1;
164
165
           cout << "YES\n";</pre>
                                                                                                                                               for (int &i=it[u]; i<G[u].size(); i++){</pre>
                                                                                                                                                    auto &[v, rc, rid] = G[u][i];
166
                                                                             for (auto x : G[now]){
167 }
                                                                                 if (vis[x]==0){
                                                                                                                                                    if (dis[v]!=dis[u]+1) continue;
                                                                                     dfs2(G, x);
                                                                                                                                                    int df = dfs(v, t, min(f, rc));
                                                                                                                                     24
                                                                                                                                                    if (df<=0) continue;</pre>
                                                                                                                                                   rc -= df:
                                                                                                                                     26
                                                                                                                                                   G[v][rid].rc += df;
  6.9 SCC 與縮點
                                                                             return;
                                                                         }
                                                                                                                                                   return df:
                                                                         void compress(){
                                                                                                                                               return 0:
                                                                             fill(vis.begin(), vis.end(), 0);
                                                                                                                                     31
                                                                                                                                           }
 2 | 給定一個有向圖·迴回傳縮點後的圖、SCC 的資訊
                                                                             for (int i=0 ; i<n ; i++){</pre>
                                                                                                                                     32
 3 所有點都以 based-0 編號
                                                                                 if (vis[i]==0){
                                                                                                                                           int flow(int s, int t){
                                                                                     dfs1(G, i);
                                                                                                                                               int ans = 0;
                                                                                                                                               while (true){
 5 函式:
                                                                                                                                                   fill(dis.begin(), dis.end(), INF);
 6| SCC compress G(n): 宣告一個有 n 個點的圖
                                                                                                                                                   queue<int> q;
 ī .add_edge(u, v): 加上一條邊 u -> v
                                                                             fill(vis.begin(), vis.end(), 0);
                                                                                                                                                   q.push(s);
 8 . compress: O(n Log n) 計算 G3、SCC、SCC id 的資訊,並把縮點後
                                                                             reverse(order.begin(), order.end());
                                                                                                                                                   dis[s] = 0;
       的結果存在 result 裡
                                                                             for (int i=0 ; i<n ; i++){</pre>
                                                                                 if (vis[order[i]]==0){
                                                                                                                                                   while (q.size()){
10 SCC[i] = 某個 SCC 中的所有點
                                                                                     SCC.push_back(vector<int>());
                                                                                                                                                       int u = q.front(); q.pop();
                                                                                                                                                       for (auto [v, rc, rid] : G[u]){
   if (rc<=0 || dis[v]<INF) continue;</pre>
                                                                                     dfs2(inv_G, order[i]);
| II | SCC id [i] = 第 i 個點在第幾個 SCC
13 // c8b146
                                                                             }
                                                                                                                                                            dis[v] = dis[u]+1;
                                                                                                                                     45
                                                                                                                                                            q.push(v);
14 struct SCC compress{
                                                                             for (int i=0 ; i<m ; i++){</pre>
      int n = 0, m = 0;
                                                                                 if (SCC_id[edges[i].first]!=SCC_id[edges[i].
      vector<vector<int>>> G, inv G, result;
      vector<pair<int, int>> edges;
                                                                                      second]){
                                                                                                                                                    if (dis[t]==INF) break;
                                                                                     result[SCC_id[edges[i].first]].push_back(
      vector<bool> vis;
                                                                                          SCC_id[edges[i].second]);
      vector<int> order;
                                                                                                                                                   fill(it.begin(), it.end(), 0);
                                                                                                                                                   while (true){
                                                                                                                                     52
                                                                                                                                                       int df = dfs(s, t, INF);
      vector<vector<int>> SCC;
                                                                                                                                     53
                                                                             for (int i=0 ; i<SCC.size() ; i++){</pre>
                                                                                                                                                       if (df<=0) break;</pre>
      vector<int> SCC id;
                                                                                 sort(result[i].begin(), result[i].end());
                                                                                                                                                       ans += df:
                                                                                 result[i].resize(unique(result[i].begin(), result 56
      SCC compress(int n){
                                                                                      [i].end())-result[i].begin());
           n = n;
           G.resize(n);
                                                                                                                                               return ans;
           inv G.resize(n):
                                                                                                                                           // the code below constructs minimum cut
           result.resize(n);
                                                                                                                                           void dfs_mincut(int now, vector<bool> &vis){
           vis.resize(n);
           SCC id.resize(n);
                                                                                                                                           vis[now] = true;
                                                                                                                                           for (auto &[v, rc, rid] : G[now]){
                                                                     6.10 Dinic
                                                                                                                                             if (vis[v]==false && rc>0){
                                                                                                                                               dfs mincut(v, vis);
      void add edge(int u, int v){
           G[u].push back(v);
                                                                                                                                     66
           inv G[v].push back(u);
                                                                   1 // 時間複雜度: O(V^2E)
                                                                                                                                     67
           edges.push_back({u, v});
                                                                     struct Flow{
           m++;
                                                                         struct Edge{
                                                                                                                                         vector<pair<int, int>> construct(int n, int s, vector<pair</pre>
                                                                             int v, rc, rid;
                                                                                                                                              int.int>> &E){
                                                                                                                                             // E is G without capacity
      void dfs1(vector<vector<int>> &G, int now){
                                                                         vector<vector<Edge>> G;
           vis[now] = 1;
                                                                                                                                             vector<bool> vis(n);
                                                                         void add(int u, int v, int c){
                                                                             G[u].push_back({v, c, G[v].size()});
           for (auto x : G[now]){
                                                                                                                                            dfs mincut(s, vis);
                                                                                                                                           vector<pair<int, int>> ret;
               if (vis[x]==0){
                                                                             G[v].push_back({u, 0, G[u].size()-1});
                                                                                                                                           for (auto &[u, v] : E){
                   dfs1(G, x);
                                                                         vector<int> dis, it;
                                                                                                                                             if (vis[u]==true && vis[v]==false){
```

# 6.11 Dijkstra

```
1 // 可以在 O(E Log E) 的時間複雜度解決在無負權有向圖單點源最短
2 const int INF = 2e18; // 要確保 INF 開的足夠大
4 vector<vector<pair<int, int>>> G(n); // G[i] = <節點, 權重>
 vector<int> dis(n, INF);
 priority_queue<pair<int, int>, vector<pair<int, int>>,
      greater<pair<int, int>>> pq;
 dis[s] = 0;
 pq.push({0, s});
 while (pq.size()){
     int now_dis = pq.top().first;
     int now node = pq.top().second;
     pq.pop();
     if (now_dis>dis[now_node]) continue;
     for (auto x : G[now node]){
         if (now_dis+x.second<dis[x.first]){</pre>
            dis[x.first] = now dis+x.second;
            pq.push({dis[x.first], x.first});
```

# 6.12 定理

- 最小點覆蓋 = 最大匹配 = n 最大點獨立集
  - 最小點覆蓋:選最少點讓所有的邊都有碰到一個點
  - 最大點獨立集:選最多不共邊的點
- 只有邊帶權的二分圖的定理(可能不重要)
  - w-vertex-cover (帶權點覆蓋): 每條邊的兩個連接點被選中的次數總和至少要是 $w_e$ 。
  - w-weight matching ( 帶權匹配 )
  - minimum vertex count of w-vertex-cover = maximum weight count of w-weight matching ( 一個點可以被選很多次・但邊不行 )
- 點、邊都帶權的二分圖的定理(可能不重要)
  - b-matching : 假設 v 的點權是  $b_v$  · 那所有 v 的匹配邊 e 的權重都要滿足  $\sum w_e \leq b_v$  。
  - The maximum w-weight of a b-matching equals the minimum b-weight of vertices in a w-vertex-cover.

#### 7 Math

#### 7.1 Burnside's-Lemma

```
\sum_{k=1}^{n} \frac{c(k)}{n}
```

- n:有多少種置換方式(例如:旋轉方式)
- c(k): 所有可能中·經過 k 次旋轉後·仍不會和別人相同的方式的數量

#### 7.2 線性篩

# 7.3 Lucas's-Theorem

```
1 // 對於很大的 C^n_{m} 對質數 p 取模·只要 p 不大就可以用。
2 int Lucas(int n, int m, int p){
3 if (m==0) return 1;
4 return (C(n%p, m%p, p)*Lucas(n/p, m/p, p)%p);
5 }
```

# 7.4 Matrix

```
struct Matrix{
   int n, m;
   vector<vector<int>> arr;

Matrix(int _n, int _m){
        n = _n;
        m = _m;
        arr.resize(n, vector<int>(m));
}

Matrix operator * (const Matrix B){
```

```
Matrix ret(n, B.m);

for (int i=0; i<n; i++){
    for (int j=0; j<B.m; j++){
        for (int k=0; k<m; k++){
            ret.arr[i][j] += arr[i][k]*B.arr[k][j];
            ret.arr[i][j] %= MOD;
    }

    }

return ret;
}
</pre>
```

## 7.5 最大質因數

```
void max_fac(int n, int &ret){
   if (n<=ret || n<2) return;
   if (isprime(n)){
      ret = max(ret, n);
      return;
   }
   int p = Pollard_Rho(n);
   max_fac(p, ret), max_fac(n/p, ret);
}</pre>
```

# 7.6 中國剩餘定理(m 不互質)

```
int extgcd(int a, int b, int &x, int &y){
     if (b==0){
         x=1, y=0;
          return a;
     int ret=extgcd(b, a%b, y, x);
     y -= a/b*x;
      return ret;
12 // 對於方程組的式子兩兩求解
13 // {是否有解, {a, m}}
14 pair < bool, pair < int, int >> CRT(int a1, int m1, int a2, int m2
      int g=__gcd(m1, m2);
     if ((a2-a1)%g!=0) return {0, {-1, -1}};
      int x, y;
      extgcd(m1, m2, x, y);
     x=(a2-a1)*x/g; // 兩者不能相反
     a1=x*m1+a1;
     m1=m1*m2/g;
     a1=(a1%m1+m1)%m1;
     return {1, {a1, m1}};
```

if(b == 0) return a;

#### 7.7 歐拉公式

```
1 | // phi(n) = 小於 n 並與 n 互質的正整數數量。
2 // O(sqrt(n)) · 回傳 phi(n)
3 int phi(int n){
      int ret = n;
      for (int i=2 ; i*i<=n ; i++){</pre>
          if (n%i==0){
              while (n%i==0) n /= i;
              ret = ret*(i-1)/i;
      if (n>1) ret = ret*(n-1)/n;
      return ret;
17 // O(n Log n) · 回傳 1~n 的 phi 值
vector<int> phi_1_to_n(int n){
      vector<int> phi(n+1);
      phi[0]=0;
      phi[1]=1;
21
      for (int i=2 ; i<=n ; i++){</pre>
          phi[i]=i-1;
      for (int i=2 ; i<=n ; i++){</pre>
          for (int j=2*i; j<=n; j+=i){ // 枚舉所有倍數
              phi[j]-=phi[i];
      }
      return phi;
```

## **7.8** 歐拉定理

```
若 a,m 互質 \cdot 則: a^n \bmod m = a^{n \bmod \varphi(m)} \bmod m 若 a,m 可能是任何數 \cdot 則: a^{\varphi(m)+[n \bmod \varphi(m)]} \bmod m
```

# 7.9 Fraction

```
#include <bits/stdc++.h>
using namespace std;

/// Fraction template starts ///
#define fraction_template_bonus_check
const long long ll_overflow_warning_value = (long long)(3e9);

long long gcd(long long a, long long b){
   if(a == 0) return 0;
```

```
if(a < b) return gcd(b,a);</pre>
       return gcd(b, a%b);
14 struct frac{
       long long a, b;
       frac(long long _a = 0, long long _b = 1){
           a = a; b = b;
           if(b == 0){
               cerr << "Error: division by zero\n";</pre>
               cerr << "Called : Constructor(" << a << ", " << 83</pre>
                    _b << ")\n";
               return;
           if(a == 0){b = 1; return;}
           if(b < 0){a = -a; b = -b;}
           long long gcd_ab = gcd(std::abs(a), b);
           if(gcd_ab != 1){a /= gcd_ab; b /= gcd_ab;}
           #ifdef fraction_template_bonus_check
           if(std::abs(a) > ll_overflow_warning_value || b >
                11_overflow_warning_value){
               cerr << "Overflow warning : " << a << "/" << b << 95
           #endif // fraction template bonus check
       frac operator+(frac const &B){
           return frac(a*(B.b)+(B.a)*b, b*(B.b));}
       frac operator-(frac const &B){
           return frac(a*(B.b)-(B.a)*b, b*(B.b));}
       frac operator*(frac const &B){
           return frac(a*(B.a), b*(B.b));}
       frac operator/(frac const &B){
           return frac(a*(B.b), b*(B.a));}
       frac operator+=(frac const &B){
           *this = frac(a*(B.b)+(B.a)*b, b*(B.b));}
       frac operator -= (frac const &B){
           *this = frac(a*(B.b)-(B.a)*b, b*(B.b));}
       frac operator*=(frac const &B){
           *this = frac(a*(B.a), b*(B.b));}
       frac operator/=(frac const &B){
           *this = frac(a*(B.b), b*(B.a));}
       frac abs(){
           a = std::abs(a);
           return *this;
       bool operator<(frac const &B){</pre>
           return a*B.b < B.a*b;}</pre>
       bool operator<=(frac const &B){</pre>
           return a*B.b <= B.a*b;}</pre>
       bool operator>(frac const &B){
           return a*B.b > B.a*b;}
       bool operator>=(frac const &B){
           return a*B.b >= B.a*b;}
       bool operator==(frac const &B){
           return a * B.b == B.a * b;}
       bool operator!=(frac const &B){
           return a * B.b != B.a * b;}
69 ostream& operator << (ostream &os, const frac& A){
       os << A.a << "/" << A.b;
       return os:
```

```
73 /// Fraction template ends ///
    void test(frac A, frac B){
          cout << "A =" << A << endl;
          cout << "B = " << B << endl;
          cout << endl:
          cout \langle\langle "A + B = " \langle\langle A + B \langle\langle endl;
          cout \langle \langle "A - B = " \langle \langle A - B \langle \langle endl:
          cout << ^{\prime\prime}A * B = ^{\prime\prime} << A * B << endl;
          cout << "A / B = " << A / B << endl;
          cout << endl:
          cout \langle\langle "(A \langle B) = " \langle\langle (A \langle B) \langle\langle endl;
          cout \langle\langle "(A \langle = B) = " \langle\langle (A \langle = B) \rangle\langle\langle endl;
          cout \langle\langle "(A > B) = " \langle\langle (A > B) \rangle\langle\langle endl;
          cout \langle\langle "(A \rangle = B) = " \langle\langle (A \rangle = B) \langle\langle endl;
          cout << "(A == B) = " << (A == B) << end1;
          cout \langle\langle "(A != B) = " \langle\langle (A != B) \langle\langle endl;
          cout << "----\n":
          return:
   int main(){
          frac tmp1(-7, 2);
          frac tmp2(5, 3);
          test(tmp1, tmp2);
          frac tmp3(-7);
          frac tmp4(0);
          test(tmp3, tmp4);
          return 0;
102
```

#### 7.10 錯排公式

錯排公式:  $(n \, \text{個人中} \cdot \text{每個人皆不再原來位置的組合數})$ 

$$dp_i = \begin{cases} 1 & i = 0 \\ 0 & i = 1 \\ (i-1)(dp_{i-1} + dp_{i-2}) & \text{otherwise} \end{cases}$$

# 7.11 Quick-Pow

```
int qp(int b, int p, int m = MOD){
   int ret = 1;
   for (; p; p>>=1){
       if (p&1) ret = ret*b%m;
       b = b*b%m;
   }
   return ret;
}
```

# 7.12 二元一次方程式

```
\begin{cases} ax + by = c \\ dx + ey = f \end{cases} = \begin{cases} x = \frac{ed - bf}{ad - bc} \\ y = \frac{af - ec}{ad - bc} \end{cases} 若 x = \frac{0}{0} \perp y = \frac{0}{0} \cdot y 則代表無限多組解。若 x = \frac{*}{0} \perp y = \frac{*}{0} \cdot y 則代表無解
```

# 7.13 Josephus

#### 7.14 數論分塊

#### 7.15 Miller-Rabin

```
1 // 0(1)
2 typedef Uint unsigned long long
3 Uint modmul(Uint a, Uint b, Uint m) {
      int ret = a*b - m*(Uint)(1.L/m*a*b);
      return ret + m*(ret < 0) - m*(ret>=(int)m);
  int qp(int b, int p, int m){
      int ret = 1;
      for (; p; p>>=1){
          if (p&1){
              ret = modmul(ret, b, m);
          b = modmul(b, b, m);
      return ret;
20 vector<int> llsprp = {2, 325, 9375, 28178, 450775, 9780504,
      1795265022};
21 bool isprime(int n, vector<int> sprp = llsprp){
      if (n==2) return 1;
      if (n<2 || n%2==0) return 0;
      int t = 0;
      int u = n-1:
      for (; u%2==0; t++) u>>=1;
```

```
for (int i=0; i<sprp.size(); i++){
    int a = sprp[i]%n;
    if (a==0 || a==1 || a==n-1) continue;
    int x = qp(a, u, n);
    if (x==1 || x==n-1) continue;
    for (int j=0; j<t; j++){
        x = modmul(x, x, n);
        if (x==1) return 0;
        if (x==n-1) break;
    }

if (x==n-1) continue;
    return 0;
}

return 1;
}</pre>
```

#### 7.16 Pollard-Rho

```
i | mt19937 seed(chrono::steady_clock::now().time_since_epoch().
       count());
  int rnd(int 1, int r){
      return uniform_int_distribution<int>(1, r)(seed);
  // O(n^{1/4}) 回傳 1 或自己的因數、記得先判斷 n 是不是質數
       (用 Miller-Rabin)
  // c1670c
  int Pollard_Rho(int n){
      int s = 0, t = 0;
      int c = rnd(1, n-1);
      int step = 0, goal = 1;
      int val = 1;
      for (goal=1;; goal<<=1, s=t, val=1){</pre>
          for (step=1 ; step<=goal ; step++){</pre>
              t = ((__int128)t*t+c)%n;
              val = (int128)val*abs(t-s)%n;
              if ((step % 127) == 0){
                  int d = __gcd(val, n);
                  if (d>1) return d;
          }
          int d = __gcd(val, n);
27
          if (d>1) return d;
28
30 }
```

# 7.17 中國剩餘定理 (m 互質)

```
vector<int> a, m;

int extgcd(int a, int b, int &x, int &y){
   if (b==0){
```

```
x=1, y=0;
          return a;
      int ret=extgcd(b, a%b, y, x);
      y-=a/b*x;
      return ret;
12
14 // n = 有幾個式子·求解 x \equiv a i \bmod m i
  int CRT(int n, vector<int> &a, vector<int> &m){
      int p=1, ans=0;
      vector<int> M(n), inv M(n);
      for (int i=0 ; i<n ; i++) p*=m[i];</pre>
      for (int i=0 ; i<n ; i++){</pre>
          M[i]=p/m[i];
      int tmp;
          extgcd(M[i], m[i], inv_M[i], tmp);
          ans+=a[i]*inv_M[i]*M[i];
          ans%=p;
26
28
29
      return (ans%p+p)%p;
30
```

#### 7.18 Catalan

任意括號序列:  $C_n = \frac{1}{n+1} {2n \choose n}$ 

## 7.19 數論定理

- 1.  $1 \sim x$  質數的數量  $\approx \frac{x}{\ln x}$
- 2.  $1 \sim x$  的因數的數量  $\approx x^{\frac{1}{3}}$
- 3. x 的質因數的數量  $\approx \log \log x$
- 4. p is a prime number  $\Leftrightarrow (p-1)! \equiv -1 \pmod{p}$
- 5. 每個正整數都可以表示成四個整數的平方和
- 6. 任何大於 2 的整數都可以表示成兩個質數的和

# 8 String

# 8.1 Hash

```
ntl9937 rnd(chrono::steady_clock::now().time_since_epoch().
count());
int A = rnd(), B = 1000000007;

vector<int> myPow, myPre;
void hash_init(string s){
    myPow.resize(s.size());
    myPre.resize(s.size());
```

```
for (int i=0; i<s.size(); i++){
    if (i==0){
        myPow[i] = 1;
        myPre[i] = s[i];
    }else{
        myPow[i] = myPow[i-1]*A%B;
        myPre[i] = (myPre[i-1]*A+s[i])%B;
    }
}

return;

int hash_value(int l, int r){ // 取得 s[l..r] 的數值
    if (l==0) return myPre[r];
    return ((myPre[r]-myPre[l-1]*myPow[r-l+1])%B+B)%B;
}
```

#### 8.2 Manacher

```
string Manacher(string str) {
    string tmp = "$#";
    for(char i : str) {
        tmp += i;
        tmp += '#';
    }

    vector<int> p(tmp.size(), 0);
    int mx = 0, id = 0, len = 0, center = 0;
    for(int i=1; i<(int)tmp.size(); i++) {
        p[i] = mx > i ? min(p[id*2-i], mx-i) : 1;

        while(tmp[i+p[i]] == tmp[i-p[i]]) p[i]++;
        if(mx<i+p[i]) mx = i+p[i], id = i;
        if(len<p[i]) len = p[i], center = i;
    }

    return str.substr((center-len)/2, len-1);
}</pre>
```

# 8.3 **Z-Function**

#### 8.4 KMP

# 8.5 Suffix-Array

```
ı| // 注意·當 /s/=1 時·Lcp 不會有值·務必測試 /s/=1 的 case
  struct SuffixArray {
       vector<int> sa, lcp;
       SuffixArray(string& s, int lim = 256) {
           // 49c4d2
          int n = SZ(s)+1, k = 0, a, b;
          vector<int> x(ALL(s)), y(n), ws(max(n, lim)), rank(n)
          x.push_back(0);
           sa = 1cp = y;
           iota(ALL(sa), 0);
           for (int j=0, p=0 ; p<n ; j=max(1LL, j*2), lim=p) {</pre>
               p = j;
               iota(ALL(y), n-j);
               for (int i=0 ; i<n ; i++) if (sa[i] >= j) y[p++]
                    = sa[i] - j;
               fill(ALL(ws), 0);
               for (int i=0 ; i<n ; i++) ws[x[i]]++;</pre>
               for (int i=1; i<lim; i++) ws[i] += ws[i - 1];</pre>
               for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
               swap(x, y), p = 1, x[sa[0]] = 0;
               for (int i=1 ; i<n ; i++){</pre>
                   a = sa[i - 1];
                  b = sa[i];
                   x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
                         ? p - 1 : p++;
              }
          }
25
26
          for (int i=1; i<n; i++) rank[sa[i]] = i;</pre>
```

#### 8.6 Min-Rotation

$\mathbf{C}$	Conte	ents			3.9	Add-Set-Segment-Tree	6	7	Math		12
					3.10	Treap	6		7.1 <b>B</b>	durnside's-Lemma	. 12
									7.2 約	象性篩	. 12
1	Mis	c	1	4	Dyn	namic-Programming	7		7.3 L	ucas's-Theorem	. 12
	1.1	Xor-Basis	1		4.1	SOS-DP	7		7.4 M	Matrix	. 12
	1.2	Default-Code	1		4.2	Digit-DP			7.5 最	· 表大質因數	. 12
	1.3	Radix-Sort	1		4.3	整數拆分	7		7.6 🕈	□國剩餘定理(m 不互質)	. 12
	1.4	Set-Pq-Sort	1						7.7 🖫	次拉公式	. 13
	1.5	2-SAT	1	5		ometry	7		7.8 🗒	次拉定理	. 13
	1.6	Enumerate-Subset	2			Point-Struct			7.9 F	raction	. 13
	1.7	Fast-Input			5.2	Line-Intersection			7.10 銉	昔排公式	. 13
	1.8	setup	2		5.3	Pick's-Theorem				Quick-Pow	
		•			5.4	Point-In-Polygon					
2	Con	vlution	2		5.5	Convex-Hull	8			osephus	
									1.15 5		
	2.1	FFT	2	_	~					牧論 <sup>ˆ</sup> 分塊	
	2.1 2.2	FFT	2 2	6	Gra	•	8		7.14 數		. 14
	2.1 2.2 2.3		2 2 3	6	6.1	Find-Bridge			7.14 數 7.15 M	效論分塊	. 14 . 14
		FFT-2	2 3	6	-	Find-Bridge	8		7.14 數 7.15 M 7.16 P	<b>対論</b> 分塊	. 14 . 14 . 14
	2.4	FFT-2	2 3	6	6.1	Find-Bridge	8 8		7.14 數 7.15 M 7.16 Pe 7.17 中	效論分塊 Miller-Rabin	. 14 . 14 . 14 . 14
3	2.4	FFT-2	2 3	6	6.1	Find-Bridge	8 8 8		7.14 數 7.15 M 7.16 P 7.17 中 7.18 C	效論分塊 Miller-Rabin	. 14 . 14 . 14 . 14
3	2.4	FFT-2	2 3 3 4 4	6	6.1 6.2 6.3	Find-Bridge Find-AP  MCMF  HLD  Tree-Isomorphism	8 8 8 9		7.14 數 7.15 M 7.16 P 7.17 中 7.18 C	效論分塊 Miller-Rabin	. 14 . 14 . 14 . 14
3	2.4	FFT-2	2 3 3 4 4	6	6.1 6.2 6.3	Find-Bridge Find-AP  MCMF  HLD  Tree-Isomorphism  Bridge BCC	8 8 8 9 9	8	7.14 數 7.15 M 7.16 P 7.17 中 7.18 C	效論分塊 Miller-Rabin	. 14 . 14 . 14 . 14
3	2.4	FFT-2	2 3 3 4 4	6	6.1 6.2 6.3 6.4 6.5	Find-Bridge Find-AP  MCMF  HLD  Tree-Isomorphism  Bridge BCC  Cut BCC	8 8 8 9 9	8	7.14 婁 7.15 M 7.16 P 7.17 中 7.18 C 7.19 婁 String	效論分塊 Miller-Rabin	. 14 . 14 . 14 . 14 . 14
3	2.4 Data 3.1 3.2	FFT-2	2 3 3 4 4 4 4	6	6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8	Find-Bridge Find-AP  MCMF  HLD  Tree-Isomorphism  Bridge BCC  Cut BCC  圓方樹	8 8 8 9 9 9	8	7.14 數 7.15 M 7.16 P 7.17 中 7.18 C 7.19 數 <b>String</b> 8.1 H	対論分塊	. 14 . 14 . 14 . 14 . 14 . 14
3	2.4 Data 3.1 3.2 3.3	FFT-2 NTT-998244353 FFT-mod  a-Structure GP-Hash-Table Sparse-Table Order-Set	2 3 3 4 4 4 4	6	6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	Find-Bridge Find-AP  MCMF HLD  Tree-Isomorphism Bridge BCC Cut BCC 圓方樹 SCC 與縮點	8 8 8 9 9 9 10	8	7.14 數 7.15 M 7.16 P 7.17 中 7.18 C 7.19 數 <b>String</b> 8.1 H 8.2 M	y論分塊  Miller-Rabin  Collard-Rho  □國剩餘定理(m 互質)  Catalan  対論定理  Mash	. 14 . 14 . 14 . 14 . 14 . 14 . 14
3	2.4  Date 3.1 3.2 3.3 3.4	FFT-2 NTT-998244353 FFT-mod  a-Structure GP-Hash-Table Sparse-Table Order-Set BIT	2 3 3 4 4 4 4	6	6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10	Find-Bridge Find-AP  MCMF  HLD  Tree-Isomorphism  Bridge BCC  Cut BCC  圓方樹  SCC 與縮點	8 8 9 9 9 10 11	8	7.14 數 7.15 M 7.16 P 7.17 中 7.18 C 7.19 數 <b>String</b> 8.1 H 8.2 M 8.3 Z	y論分塊  Miller-Rabin  Collard-Rho  回國剩餘定理(m 互質)  Catalan  対論定理  Mash  Manacher	. 14 . 14 . 14 . 14 . 14 . 14 . 15 . 15
3	2.4 Data 3.1 3.2 3.3 3.4 3.5 3.6 3.7	FFT-2 NTT-998244353 FFT-mod  a-Structure GP-Hash-Table Sparse-Table Order-Set BIT Persistent-Segment-Tree Trie LC-Segment-Tree	2 3 3 4 4 4 4 4 5 5	6	6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11	Find-Bridge Find-AP MCMF HLD Tree-Isomorphism Bridge BCC Cut BCC 圓方樹 SCC 與縮點 Dinic Dijkstra	8 8 8 9 9 9 10 11 11	8	7.14 數 7.15 M 7.16 P 7.17 中 7.18 C 7.19 數 <b>String</b> 8.1 H 8.2 M 8.3 Z 8.4 K 8.5 St	対論分塊 Miller-Rabin Collard-Rho の間なけられる のでは、 Miller - Rabin のでは、 M	. 14 . 14 . 14 . 14 . 14 . 14 . 15 . 15 . 15
3	2.4 Data 3.1 3.2 3.3 3.4 3.5 3.6 3.7	FFT-2 NTT-998244353 FFT-mod  a-Structure GP-Hash-Table Sparse-Table Order-Set BIT Persistent-Segment-Tree Trie	2 3 3 4 4 4 4 4 5 5	6	6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9 6.10 6.11	Find-Bridge Find-AP  MCMF  HLD  Tree-Isomorphism  Bridge BCC  Cut BCC  圓方樹  SCC 與縮點	8 8 8 9 9 9 10 11 11	8	7.14 數 7.15 M 7.16 P 7.17 中 7.18 C 7.19 數 <b>String</b> 8.1 H 8.2 M 8.3 Z 8.4 K 8.5 St	対論分塊 Miller-Rabin Collard-Rho 回 國刺餘定理(m 互質) Catalan 対論定理 Manacher C-Function MMP	. 14 . 14 . 14 . 14 . 14 . 14 . 15 . 15 . 15