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# 1 Misc

#### 1.1 Xor-Basis

```
vector<int> basis:
void add vector(int x){
      for (auto v : basis){
         x=min(x, x^v);
     if (x) basis.push_back(x);
9 | // 給一數字集合 S · 求能不能 XOR 出 x
10 bool check(int x){
     for (auto v : basis){
         x=min(x, x^v);
     return x;
17 / / 給一數字集合 S · 求能 XOR 出多少數字
18 // 答案等於 2^{basis 的大小}
20 | // 給一數字集合 S · 求 XOR 出最大的數字
21 int get_max(){
     int ans=0;
     for (auto v : basis){
         ans=max(ans, ans^v);
     return ans;
```

#### 1.2 Default-Code

```
1 #include <bits/stdc++.h>
 #define int long long
  #define ALL(x) x.begin(), x.end()
 #define SZ(x) ((int)x.size())
  #define fastio ios::sync with stdio(0), cin.tie(0);
 using namespace std;
 #ifdef LOCAL
 #define cout cout << "\033[0;32m"
10 #define cerr cerr << "\033[0;31m"
#define endl endl << "\033[0m"</pre>
#pragma GCC optimize("03,unroll-loops")
14 #pragma GCC target("avx,avx2,sse,sse2,sse3,sse4,popcnt")
15 #define endl "\n"
16 #endif
18 const int MAX_N = 5e5+10;
19 const int INF = 2e18;
 void solve1(){
      return;
26 signed main(){
```

```
27

28 fastio;

29

30 int t = 1;

31 while (t--){

32 solve1();

33 }

34 

35 return 0;

36 }
```

#### 1.3 Radix-Sort

```
1 // 值域限制:0~1073741823(2^30-1)
 inline void radix_sort(vector<int> &a, int n){
     static int cnt[32768] = {0};
      vector<int> tmpa(n);
      for(int i = 0; i < n; ++i)</pre>
          ++cnt[a[i] & 32767];
      for(int i = 1; i < 32768; ++i)</pre>
          cnt[i] += cnt[i-1];
      static int temp;
      for(int i = n-1; i >= 0; --i){
          temp = a[i] & 32767;
          --cnt[temp]:
          tmpa[cnt[temp]] = a[i];
      static int cnt2[32768] = {0};
      for(int i = 0; i < n; ++i)
          ++cnt2[(tmpa[i]>>15)];
     for(int i = 1; i < 32768; ++i)</pre>
          cnt2[i] += cnt2[i-1];
      for(int i = n-1; i >= 0; --i){
          temp = (tmpa[i]>>15);
          --cnt2[temp];
         a[cnt2[temp]] = tmpa[i];
      return;
```

# 1.4 Set-Pq-Sort

```
1  // priority_queue
2  struct cmp{
      bool operator () (Data a, Data b){
          return a.x<b.x;
      }
6  };
7  priority_queue<Data, vector<Data>, cmp> pq;
8
9  // set
10  struct Data{
      int x;
12  bool operator < (const Data &b){
          return x<b.x;
15  }
16  };</pre>
```

#### 1.5 2-SAT

```
| #include <bits/stdc++.h>
  using namespace std;
  struct TWO_SAT {
      int n, N;
      vector<vector<int>> G, rev_G;
      deque<bool> used;
      vector<int> order, comp;
      deque<bool> assignment;
      void init(int n) {
          n = _n;
          N = n * 2;
          G.resize(N + 5);
          rev_G.resize(N + 5);
      void dfs1(int v) {
          used[v] = true;
          for (int u : G[v]) {
              if (!used[u])
                  dfs1(u);
          order.push_back(v);
23
24
      void dfs2(int v, int cl) {
          comp[v] = c1;
          for (int u : rev G[v]) {
              if (comp[u] == -1)
                  dfs2(u, c1);
      bool solve() {
          order.clear();
          used.assign(N, false);
          for (int i = 0; i < N; ++i) {
               if (!used[i])
                  dfs1(i);
          comp.assign(N, -1);
          for (int i = 0, j = 0; i < N; ++i) {
               int v = order[N - i - 1];
               if (comp[v] == -1)
                  dfs2(v, j++);
          assignment.assign(n, false);
          for (int i = 0; i < N; i += 2) {
               if (comp[i] == comp[i + 1])
                  return false:
               assignment[i / 2] = (comp[i] > comp[i + 1]);
50
          return true;
51
52
      void add_disjunction(int a, bool na, int b, bool nb) { //
          // na means whether a is negative or not
          // nb means whether b is negative or not
          a = 2 * a ^ na;
          b = 2 * b ^ nb;
          int neg_a = a ^ 1;
          int neg_b = b ^ 1;
58
          G[neg_a].push_back(b);
59
          G[neg_b].push_back(a);
          rev G[b].push back(neg a);
          rev_G[a].push_back(neg_b);
```

```
return;
       void get result(vector<int>& res) {
           res.clear();
           for (int i = 0; i < n; i++)</pre>
               res.push_back(assignment[i]);
71 /* CSES Giant Pizza
72 3 5
|73| + 1 + 2
|74| - 1 + 3
   - + + + -
77 */
78 int main() {
       int n, m;
       cin >> n >> m;
       TWO SAT E;
       E.init(m);
       char c1, c2;
       int inp1, inp2;
       for (int i = 0; i < n; i++) {</pre>
           cin >> c1 >> inp1;
           cin >> c2 >> inp2;
           E.add disjunction(inp1 - 1, c1 == '-', inp2 - 1, c2
       bool able = E.solve();
       if (able) {
           vector <int> ans;
           E.get_result(ans);
           for (int i : ans)
               cout << (i == true ? '+' : '-') << ' ';
           cout << ' \ n';
       } else {
           cout << "IMPOSSIBLE\n";</pre>
102
       return 0;
```

#### 1.6 Enumerate-Subset

```
1 // 時間複雜度 O(3^n)
2 // 枚舉每個 mask 的子集
3 for (int mask=0; mask<(1<<n); mask++){
4 for (int s=mask; s>=0; s=(s-1)&m){
5 // s 是 mask 的子集
6 if (s==0) break;
7 }
```

### 1.7 Fast-Input

```
1 // fast IO
2 // 6f8879
```

```
3 inline char readchar(){
      static char buffer[BUFSIZ], * now = buffer + BUFSIZ, *
           end = buffer + BUFSIZ:
      if (now == end)
          if (end < buffer + BUFSIZ)</pre>
              return EOF;
          end = (buffer + fread(buffer, 1, BUFSIZ, stdin));
          now = buffer;
      return *now++;
 inline int nextint(){
      int x = 0, c = readchar(), neg = false;
      while (('0' > c | c > '9') \&\& c! = '-' \&\& c! = EOF) c =
           readchar():
      if(c == '-') neg = true, c = readchar();
      while ('0' <= c && c <= '9') x = (x << 3) + (x << 1) + (c^{0}) 24
           , c = readchar();
      if(neg) x = -x;
      return x; // returns 0 if EOF
```

### 1.8 setup

```
se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
:inoremap " ""<Esc>i
:inoremap {<CR> {<CR>}<Esc>ko
:inoremap {{ {}}<ESC>i
function! F(...)
 execute '!./%:r < ./' . a:1
command! -nargs=* R call F(<f-args>)
map <F7> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -Wextra -
     Wshadow -02 -DLOCAL -g -fsanitize=undefined,address<CR>
map <F8> :!./%:r<CR>
map <F9> :!./%:r < ./%:r.in<CR>
ca hash w !cpp -dD -P -fpreprocessed \| tr -d "[:space:]" \|
    md5sum \| cut -c-6
" i+<esc>25A---+<esc>
" o|<esc>25A |<esc>
" "ggVGyG35pGdd
```

### 2 Convolution

### 2.1 FFT

```
typedef complex <double > cd;
const double PI = acos(-1);

void FFT(vector < cd > &a , bool inv){
   int n = a.size();
}
```

```
for (int i=1, j=0; i<n; i++){</pre>
        int bit = (n>>1);
        for ( ; j&bit ; bit>>=1){
            j ^= bit;
        i ^= bit;
        if (i<j){
            swap(a[i], a[j]);
   for (int len=2 ; len<=n ; len<<=1){</pre>
        cd wlen = polar(1.0, (inv ? 2 : -2)*PI/len);
        for (int i=0 ; i<n ; i+=len){</pre>
            cd w(1);
            for (int j=0 ; j<len/2 ; j++){</pre>
                cd u = a[i+j];
                cd v = a[i+j+len/2]*w;
                a[i+j] = u+v;
                a[i+j+len/2] = u-v;
                w *= wlen;
   }
   if (inv){
        for (auto &x : a){
            x /= n;
    return;
vector<cd> polyMul(vector<cd> a, vector<cd> b){
   int sa = a.size(), sb = b.size(), n = 1;
    while (n<sa+sb-1) n *= 2;</pre>
   a.resize(n);
   b.resize(n);
   vector<cd> c(n);
   FFT(a, 0);
   FFT(b, 0);
    for (int i=0; i<n; i++) c[i] = a[i]*b[i];</pre>
   FFT(c, 1);
   c.resize(sa+sb-1);
    return c;
```

### 2.2 FFT-2

```
typedef complex<double> cd;

void FFT(vector<cd> &a) {
   int n = a.size(), L = 31-__builtin_clz(n);
   vector<complex<long double>> R(2, 1);
   vector<cd> rt(2, 1);
   for (int k=2; k<n; k*=2){</pre>
```

```
R.resize(n);
       rt.resize(n);
        auto x = polar(1.0L, acos(-1.0L) / k):
        for (int i=k ; i<2*k ; i++){</pre>
            rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);
   }
   vector<int> rev(n);
   for (int i=0 ; i<n ; i++){</pre>
        rev[i] = (rev[i/2] | (i&1) << L)/2;
   for (int i=0 ; i<n ; i++){</pre>
       if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
   for (int k=1; k<n; k*=2){
        for (int i=0 ; i<n ; i+=2*k){</pre>
            for (int j=0; j<k; j++){
                auto x = (double *)&rt[j+k];
                auto y = (double *)&a[i+j+k];
                cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]*
                     y[0]);
                a[i+j+k] = a[i+j]-z;
                a[i+j] += z;
   return;
vector<double> PolyMul(const vector<double> a, const vector<</pre>
    double> b){
   if (a.empty() || b.empty()) return {};
   vector<double> res(a.size()+b.size()-1);
   int L = 32 - builtin clz(res.size()), n = 1 << L;</pre>
   vector<cd> in(n), out(n);
   copy(a.begin(), a.end(), begin(in));
   for (int i=0 ; i<b.size() ; i++){</pre>
        in[i].imag(b[i]);
   FFT(in);
   for (cd& x : in) x *= x;
   for (int i=0 ; i<n ; i++){</pre>
       out[i] = in[-i & (n - 1)] - conj(in[i]);
   FFT(out);
   for (int i=0 ; i<res.size() ; i++){</pre>
       res[i] = imag(out[i]) / (4 * n);
   return res;
```

### 2.3 NTT-998244353

```
const int MOD = (119 << 23) + 1, ROOT = 62; // = 998244353

// For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 << 21

3 // and 483 << 21 (same root). The last two are > 10^9.
```

```
6 | void NTT(vector<int> &a) {
      int n = a.size();
      int L = 31- builtin clz(n);
      vector<int> rt(2, 1);
      for (int k=2, s=2; k<n; k*=2, s++){
          rt.resize(n):
          int z[] = {1, qp(ROOT, MOD>>s)};
          for (int i=k ; i<2*k ; i++){</pre>
              rt[i] = rt[i/2]*z[i&1]%MOD;
     }
      vector<int> rev(n);
      for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2] | (i&1) << L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]){</pre>
              swap(a[i], a[rev[i]]);
     }
      for (int k=1; k<n; k*=2){</pre>
          for (int i=0 ; i<n ; i+=2*k){</pre>
              for (int j=0 ; j<k ; j++){</pre>
                  int z = rt[j+k]*a[i+j+k]%MOD, &ai = a[i+j];
                  a[i+j+k] = ai-z+(z>ai ? MOD : 0);
                  ai += (ai+z)=MOD ? z-MOD : z);
 vector<int> polyMul(vector<int> &a, vector<int> &b){
      if (a.empty() || b.empty()) return {};
      int s = a.size()+b.size()-1, B = 32-_builtin_clz(s), n =
            1<<B:
      int inv = qp(n, MOD-2);
      vector<int> L(a), R(b), out(n);
      L.resize(n), R.resize(n);
      NTT(L), NTT(R);
      for (int i=0 ; i<n ; i++){</pre>
          out[-i&(n-1)] = L[i]*R[i]%MOD*inv%MOD;
      NTT(out);
      out.resize(s);
      return out;
```

### 2.4 FFT-mod

```
1 | /*
2 | 修改 const int MOD = 998244353 更改要取餘的數字
3 | PolyMul(a, b) 回傳多項式乘法的結果 ( c_k = \sum_{i+j} a_i+b_j mod MOD )
4 |
5 | 大約可以支援 5e5 · a_i, b_i 皆在 MOD 以下的非負整數
6 | */
7 | const int MOD = 998244353;
```

```
8 typedef complex < double > cd;
  // b9c90a
  void FFT(vector<cd> &a) {
      int n = a.size(), L = 31- builtin clz(n);
      vector<complex<long double>> R(2, 1);
      vector<cd> rt(2, 1);
      for (int k=2; k<n; k*=2){</pre>
          R.resize(n);
          rt.resize(n);
          auto x = polar(1.0L, acos(-1.0L) / k);
          for (int i=k ; i<2*k ; i++){</pre>
              rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);
      }
      vector<int> rev(n):
      for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2] | (i&1) << L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
      for (int k=1; k<n; k*=2){
          for (int i=0; i<n; i+=2*k){
              for (int j=0 ; j<k ; j++){</pre>
                  auto x = (double *)&rt[j+k];
                  auto y = (double *)&a[i+j+k];
                  cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]*
                       y[0]);
                  a[i+j+k] = a[i+j]-z;
                  a[i+j] += z;
      return;
  vector<int> PolyMul(vector<int> a, vector<int> b){
      if (a.empty() || b.empty()) return {};
      vector<int> res(a.size()+b.size()-1);
      int B = 32- builtin_clz(res.size()), n = (1<<B), cut =</pre>
           int(sqrt(MOD));
      vector<cd> L(n), R(n), outs(n), outl(n);
      for (int i=0 ; i<a.size() ; i++){</pre>
          L[i] = cd((int) a[i]/cut, (int)a[i]%cut);
      for (int i=0 ; i<b.size() ; i++){</pre>
          R[i] = cd((int) b[i]/cut, (int)b[i]%cut);
      FFT(L);
      FFT(R);
      for (int i=0 ; i<n ; i++){</pre>
          int j = -i&(n-1);
          outl[j] = (L[i]+conj(L[j])) * R[i]/(2.0*n);
          outs[j] = (L[i]-conj(L[j])) * R[i]/(2.0*n)/1i;
      FFT(outl);
      FFT(outs);
      for (int i=0 ; i<res.size() ; i++){</pre>
          int av = (int)(real(outl[i])+0.5), cv = (int)(imag(
               outs[i])+0.5):
```

### 3 Data-Structure

#### 3.1 GP-Hash-Table

```
| #include <ext/pb ds/assoc container.hpp>
using namespace gnu pbds;
typedef tree<int, null type, less<int>, rb tree tag,
      tree_order_statistics_node_update> order_set;
      static uint64 t splitmix64(uint64 t x) {
          // http://xorshift.di.unimi.it/splitmix64.c
          x += 0x9e3779b97f4a7c15;
          x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
          x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
          return x ^ (x >> 31);
      size t operator()(uint64 t x) const {
          static const uint64 t FIXED RANDOM = chrono::
               steady clock::now().time since epoch().count();
          return splitmix64(x + FIXED RANDOM);
16
17 };
19 gp_hash_table<int, int, custom_hash> ss;
```

# 3.2 Sparse-Table

#### 3.3 Order-Set

28

33

50

#### 3.4 BIT

```
vector<int> BIT(MAX SIZE);
  void update(int pos, int val){
      for (int i=pos ; i<MAX SIZE ; i+=i&-i){</pre>
          BIT[i]+=val;
  int query(int pos){
      int ret=0:
      for (int i=pos ; i>0 ; i-=i&-i){
          ret+=BIT[i]:
       return ret;
  // const int MAX_N = (1 << 20)
  int k th(int k){ // 回傳 BIT 中第 k 小的元素(based-1)
       int res = 0;
       for (int i=MAX N>>1 ; i>=1 ; i>>=1)
          if (bit[res+i]<k)</pre>
               k -= bit[res+=i];
       return res+1;
23 }
```

# 3.5 Persistent-Segment-Tree

```
int lc = -1;
    int rc = -1;
    int val = 0:
};
vector<Node> arr;
vector<int> version;
Persistent Segment Tree(int sz){
    arr.resize(32*sz);
    version.push_back(node_cnt++);
    return;
}
void pull(Node &c, Node a, Node b){
    c.val = a.val+b.val;
    return;
void build(vector<int> &v, int idx, int ll = 0, int rr =
    auto &now = arr[idx];
    if (rr-ll==1){
        now.val = v[11];
        return;
    int mid = (11+rr)/2;
    now.lc = node_cnt++;
    now.rc = node cnt++;
    build(v, now.lc, ll, mid);
    build(v, now.rc, mid, rr);
    pull(now, arr[now.lc], arr[now.rc]);
    return;
void update(int pos, int val, int idx, int ll = 0, int rr
    auto &now = arr[idx];
    if (rr-ll==1){
        now.val = val;
        return;
    int mid = (11+rr)/2;
    if (pos<mid){</pre>
        arr[node cnt] = arr[now.lc];
        now.lc = node cnt;
        node_cnt++;
        update(pos, val, now.lc, ll, mid);
        arr[node_cnt] = arr[now.rc];
        now.rc = node cnt;
        update(pos, val, now.rc, mid, rr);
    pull(now, arr[now.lc], arr[now.rc]);
    return:
void update_version(int pos, int val, int ver){
    update(pos, val, version[ver]);
```

#### **3.6** Trie

```
| struct Trie{
     struct Data{
         int nxt[2]={0, 0};
     int sz=0:
     vector<Data> arr;
     void init(int n){
         arr.resize(n);
     void insert(int n){
         int now=0;
         for (int i=N ; i>=0 ; i--){
             int v=(n>>i)&1;
             if (!arr[now].nxt[v]){
                 arr[now].nxt[v]=++sz;
             now=arr[now].nxt[v];
     }
     int query(int n){
         int now=0, ret=0;
         for (int i=N ; i>=0 ; i--){
             int v=(n>>i)&1;
             if (arr[now].nxt[1-v]){
                 ret+=(1<<i);
                 now=arr[now].nxt[1-v];
             }else if (arr[now].nxt[v]){
                 now=arr[now].nxt[v];
             }else{
                 return ret;
```

# 3.7 LC-Segment-Tree

```
全部都是 0-based
  LC Segment Tree st(n);
  update(val): 將一個 pair <a, b> 代表插入一條 y=ax+b 的直線
  query(x): 查詢所有直線在位置 x 的最小值
  const int MAX V = 1e6+10; // 值域最大值
  struct LC_Segment_Tree{
      struct Node{ // y = ax+b
          int a = 0:
          int b = INF;
         int y(int x){
             return a*x+b;
      };
      vector<Node> arr;
      LC_Segment_Tree(int n = 0){
          arr.resize(4*n);
27
      void update(Node val. int idx = 0, int ll = 0, int rr =
          MAX V){
          if (rr-ll==1){
             if (val.y(l1) < arr[idx].y(l1)){</pre>
                 arr[idx] = val;
             return;
         int mid = (11+rr)/2;
          if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
              的線斜率要比較小
          if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
             update(val, idx*2+1, ll, mid);
         }else{ // 交點在右邊
             swap(arr[idx], val); // 在左子樹中·新線比舊線還
             update(val, idx*2+2, mid, rr);
          return;
      int query(int x, int idx = 0, int ll = 0, int rr = MAX_V)
          if (rr-ll==1){
             return arr[idx].y(ll);
```

# 3.8 Persistent-Disjoint-Set

```
struct Persistent Disjoint Set{
      Persistent_Segment_Tree arr, sz;
      void init(int n){
          arr.init(n);
          vector<int> v1;
          for (int i=0 ; i<n ; i++){</pre>
               v1.push_back(i);
          arr.build(v1, 0);
          sz.init(n);
          vector<int> v2;
          for (int i=0 ; i<n ; i++){</pre>
              v2.push_back(1);
          sz.build(v2, 0);
      int find(int a){
          int res = arr.query_version(a, a+1, arr.version.size
          if (res==a) return a;
          return find(res);
25
      bool unite(int a, int b){
          a = find(a):
          b = find(b);
          if (a!=b){
               int sz1 = sz.query version(a, a+1, arr.version.
                   size()-1).val;
               int sz2 = sz.query version(b, b+1, arr.version.
                   size()-1).val;
              if (sz1<sz2){</pre>
                  arr.update version(a, b, arr.version.size()
                  sz.update version(b, sz1+sz2, arr.version.
                        size()-1);
                  arr.update version(b, a, arr.version.size()
                  sz.update version(a, sz1+sz2, arr.version.
                        size()-1);
```

```
42 return true;
43 }
44 return false;
45 }
46 };
```

# 3.9 Add-Set-Segment-Tree

```
1 // [ll, rr), based-0
2 // 使用前記得 init(陣列大小), build(陣列名稱)
3 // add(LL, rr): 區間修改
4 // set(ll, rr): 區間賦值
5 // query(ll, rr): 區間求和 / 求最大值
6 struct SegmentTree{
     struct node{
         int add_tag = 0;
         int set_tag = 0;
         int sum = 0;
         int ma = 0;
     };
     vector<node> arr;
     SegmentTree(int n){
         arr.resize(n<<2);</pre>
     node pull(node A, node B){
         node C;
         C.sum = A.sum + B.sum;
         C.ma = max(A.ma, B.ma);
         return C;
     void push(int idx, int ll, int rr){
         if (arr[idx].set tag!=0){
             arr[idx].sum = (rr-ll)*arr[idx].set_tag;
             arr[idx].ma = arr[idx].set tag;
             if (rr-ll>1){
                 arr[idx*2+1].add tag = 0;
                 arr[idx*2+1].set tag = arr[idx].set tag;
                 arr[idx*2+2].add_tag = 0;
                 arr[idx*2+2].set tag = arr[idx].set tag;
             arr[idx].set_tag = 0;
         if (arr[idx].add tag!=0){
             arr[idx].sum += (rr-ll)*arr[idx].add tag;
             arr[idx].ma += arr[idx].add tag;
                 arr[idx*2+1].add tag += arr[idx].add tag;
                 arr[idx*2+2].add_tag += arr[idx].add_tag;
             arr[idx].add tag = 0;
     void build(vector<int> &v, int idx = 0, int ll = 0, int
          rr = n){}
         if (rr-ll==1){
             arr[idx].sum = v[11];
```

```
arr[idx].ma = v[ll];
           }else{
               int mid = (11+rr)/2:
               build(v, idx*2+1, ll, mid);
               build(v, idx*2+2, mid, rr);
               arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
       }
       void add(int ql, int qr, int val, int idx = 0, int ll =
            0, int rr =n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return;</pre>
           if (q1<=11 && rr<=qr){
               arr[idx].add_tag += val;
               push(idx, 11, rr);
               return;
           int mid = (11+rr)/2;
           add(ql, qr, val, idx*2+1, ll, mid);
           add(ql, qr, val, idx*2+2, mid, rr);
           arr[idx]=pull(arr[idx*2+1], arr[idx*2+2]);
       void set(int ql, int qr, int val, int idx=0, int ll=0,
            int rr=n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return;</pre>
           if (q1<=11 && rr<=qr){</pre>
               arr[idx].add_tag = 0;
               arr[idx].set_tag = val;
               push(idx, ll, rr);
               return;
           int mid = (11+rr)/2;
           set(ql, qr, val, idx*2+1, ll, mid);
           set(q1, qr, val, idx*2+2, mid, rr);
           arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
       node query(int ql, int qr, int idx = 0, int ll = 0, int
            rr = n){
           push(idx, 11, rr);
           if (rr<=ql || qr<=ll) return node();</pre>
           if (q1<=11 && rr<=qr) return arr[idx];</pre>
           int mid = (11+rr)/2;
           return pull(query(ql, qr, idx*2+1, ll, mid), query(ql
                , qr, idx*2+2, mid, rr));
100 } ST;
   3.10 Treap
  struct Treap{
       Treap *1 = nullptr, *r = nullptr;
       int pri = rand(), val = 0, sz = 1;
```

Treap(int \_val){

8 };

val = \_val;

```
int size(Treap *t){return t ? t->sz : 0;}
  void pull(Treap *t){
      t\rightarrow sz = size(t\rightarrow l) + size(t\rightarrow r) + 1:
  Treap* merge(Treap *a, Treap *b){
       if (!a | | !b) return a ? a : b;
      if (a->pri>b->pri){
           a \rightarrow r = merge(a \rightarrow r, b);
           pull(a);
21
           return a;
      }else{
           b\rightarrow 1 = merge(a, b\rightarrow 1);
           pull(b);
           return b;
27
28
  pair<Treap*, Treap*> split(Treap *&t, int k){ // 1-based <前
       k 個元素, 其他元素>
      if (!t) return {};
      if (size(t->1)>=k){
           auto pa = split(t->1, k);
           t->l = pa.second;
           return {pa.first, t};
           auto pa = split(t->r, k-size(t->l)-1);
           t->r = pa.first:
           pull(t);
           return {t, pa.second};
  // functions
  Treap* build(vector<int> v){
       Treap* ret;
       for (int i=0 ; i<SZ(v) ; i++){</pre>
           ret = merge(ret, new Treap(v[i]));
51
       return ret;
52
54 array<Treap*, 3> cut(Treap *t, int 1, int r){ // 1-based <前
       1~L-1 個元素, L~r 個元素, r+1 個元素>
       array<Treap*, 3> ret;
       tie(ret[1], ret[2]) = split(t, r);
       tie(ret[0], ret[1]) = split(ret[1], 1-1);
       return ret;
  void print(Treap *t, bool flag = true){
      if (t->1!=0) print(t->1, false);
       cout << t->val;
      if (t->r!=0) print(t->r, false);
      if (flag) cout << endl;</pre>
66 }
```

# 4 Dynamic-Programming

#### 4.1 SOS-DP

```
1 // 總時間複雜度為 O(n 2^n)
2 // 計算 dp[i] = i 所有 bit mask 子集的和
3 for (int i=0; i<n; i++){
4 for (int mask=0; mask<(1<<n); mask++){
5 if ((mask>)i)&1){
6 dp[mask] += dp[mask^(1<<i)];
7 }
8 }
9 }
```

# 4.2 Digit-DP

```
i #include <bits/stdc++.h>
 using namespace std;
4 long long 1, r;
5 long long dp[20][10][2][2]; // dp[pos][pre][limit] = 後 pos
      位 pos 前一位是 pre (是/否)有上界 (是/否)有前綴零
  long long memorize_search(string &s, int pos, int pre, bool
      limit, bool lead){
     // 已經被找過了,直接回傳值
     if (dp[pos][pre][limit][lead]!=-1) return dp[pos][pre][
          limit][lead];
     // 已經搜尋完畢,紀錄答案並回傳
     if (pos==(int)s.size()){
         return dp[pos][pre][limit][lead] = 1;
     // 枚舉目前的位數數字是多少
     long long ans = 0;
     for (int now=0 ; now<=(limit ? s[pos]-'0' : 9) ; now++){</pre>
         if (now==pre){
            // 1~9 絕對不能連續出現
            if (pre!=0) continue;
            // 如果已經不在前綴零的範圍內· Ø 不能連續出現
            if (lead==false) continue;
         ans += memorize search(s, pos+1, now, limit&(now==(s[
             pos]-'0')), lead&(now==0));
     }
     // 已經搜尋完畢,紀錄答案並回傳
     return dp[pos][pre][limit][lead] = ans;
36 | // 回傳 [0, n] 有多少數字符合條件
37 long long find answer(long long n){
     memset(dp, -1, sizeof(dp));
```

```
string tmp = to_string(n);

return memorize_search(tmp, 0, 0, true, true);

int main(){

// input
cin >> 1 >> r;

// output - 計算 [l, r] 有多少數字任意兩個位數都不相同
cout << find_answer(r)-find_answer(1-1) << "\n";

return 0;

}
```

### 4.3 整數拆分

```
dp[i][x] = 要將整數 x 拆成 i 堆的「組合數」 dp[i+1][x+1] + = dp[i][x] ( 創造新的一堆 ) \\ dp[i][x+i] + = dp[i][x] ( 把每一堆都增加 1 )
```

# 5 Geometry

#### **5.1** Line-Intersection

### 5.2 Pick's-Theorem

給定頂點坐標均是整點的簡單多邊形·面積 = 內部格點數 + 邊上格點數/2 - 1

# 5.3 Point-In-Polygon

```
1 | /*
2 | 可以在有 n 個點的簡單多邊形內·用 O(n) 的時間回傳:
3 | 1: 在多邊形內, 0: 在多邊形上, -1: 在多邊形外
4 | */
```

### 5.4 Convex-Hull

### 5.5 Point-Struct

```
const int EPS = 1e-6;

struct Point{
    Point x, y;

Point(Point _x = 0, Point _y = 0){
    x = _x;
    y = _y;
}

// 純量乘、除法
Point operator * (Point a){return {a*x, a*y};};
Point operator / (Point a){return {a/x, a/y};};

// 向量加、減法
Point operator + (Point a){return {x+a.x, y*a.y};};
Point operator - (Point a){return {x-a.x, y-a.y};};

// 內積、外積
double operator * (Point a){return x*a.x+y*a.y;};
```

```
double operator ^ (Point a){return x*a.y-y*a.x;};
      // bool operator < (const Point &a) const {return (x*a.v<
           a.x*y);} // 極角排序(順時鐘)
      bool operator < (const Point &a) const {return x==a.x ? y</pre>
           \langle a.v : x \langle a.x \rangle \rangle
      bool operator == (const Point &a) const {return x==a.x &&
      double dis(Point a){return sqrtl(abs(x-a.x)*abs(x-a.x)+
           abs(y-a.y)*abs(y-a.y));}
28 };
30 // 判斷向量正負:1=正數,0=0,-1=負數
31 int sign(double a){
      if (abs(a) < EPS) return 0;</pre>
      else return (a>0 ? 1 : -1);
36 // 判斷 ab 到 ac 的方向: 1=逆時鐘, 0=重疊, -1=順時鐘
37 int ori(Point a, Point b, Point c){
      return sign((b-a)^(c-a));
```

# 6 Graph

## 6.1 Find-Bridge

```
1 vector<int> dep(MAX N), low(MAX N);
 vector<pair<int, int>> bridge;
 bitset<MAX N> vis;
 void dfs(int now, int pre){
     vis[now] = 1;
     low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
     for (auto x : G[now]){
         if (x==pre){
             continue;
         }else if (vis[x]==0){
             // 沒有走過的節點
             dfs(x, now);
             low[now] = min(low[now], low[x]);
         }else if (vis[x]==1){
             low[now] = min(low[now], dep[x]);
     }
     if (now!=1 && low[now]==dep[now]){
         bridge.push_back({now, pre});
     return:
```

### 6.2 Find-AP

```
vector<int> dep(MAX N), low(MAX N), AP;
  bitset<MAX N> vis;
  void dfs(int now, int pre){
      int cnt = 0;
      bool ap = 0;
      vis[now] = 1;
      low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
      for (auto x : G[now]){
          if (x==pre)
              continue:
          }else if (vis[x]==0){
              cnt++;
              dfs(x, now);
              low[now] = min(low[now], low[x]);
              if (low[x]>=dep[now]) ap=1;
          }else{
              low[now] = min(low[now], dep[x]);
      }
      if ((now==pre && cnt>=2) || (now!=pre && ap)){
          AP.push back(now):
26 }
```

### 6.3 HLD

```
| #include <bits/stdc++.h>
 #define int long long
 using namespace std:
 const int N = 100005;
 vector <int> G[N];
 struct HLD {
     vector<int> pa, sz, depth, mxson, topf, id;
     int n, idcnt = 0;
     HLD(int _n) : n(_n), pa(_n + 1), sz(_n + 1), depth(_n +
          1), mxson(n + 1), topf(n + 1), id(n + 1) {}
     void dfs1(int v = 1, int p = -1) {
         pa[v] = p; sz[v] = 1; mxson[v] = 0;
         depth[v] = (p == -1 ? 0 : depth[p] + 1);
         for (int u : G[v]) {
             if (u == p) continue;
             dfs1(u, v);
             sz[v] += sz[u];
             if (sz[u] > sz[mxson[v]]) mxson[v] = u;
     void dfs2(int v = 1, int top = 1) {
         id[v] = ++idcnt;
         topf[v] = top;
         if (mxson[v]) dfs2(mxson[v], top);
         for (int u : G[v]) {
             if (u == mxson[v] || u == pa[v]) continue;
             dfs2(u, u);
     // query 為區間資料結構
     int path query(int a, int b) {
         int res = 0;
```

## **6.4** Tree-Isomorphism

```
1 #include <bits/stdc++.h>
 #pragma GCC optimize("03,unroll-loops")
 #define fastio ios::sync with stdio(0), cin.tie(0), cout.tie
 #define dbg(x) cerr << #x << " = " << x << endl
 #define int long long
 using namespace std;
 // declare
 const int MAX SIZE = 2e5+5;
 const int INF = 9e18;
 const int MOD = 1e9+7;
 const double EPS = 1e-6;
 typedef vector<vector<int>> Graph;
 typedef map<vector<int>, int> Hash;
 int n, a, b;
 int id1, id2;
 pair<int, int> c1, c2;
 vector<int> sz1(MAX_SIZE), sz2(MAX_SIZE);
 vector<int> we1(MAX_SIZE), we2(MAX_SIZE);
 Graph g1(MAX_SIZE), g2(MAX_SIZE);
 Hash m1, m2;
 int testcase=0;
 void centroid(Graph &g, vector<int> &s, vector<int> &w, pair
      int, int> &rec, int now, int pre){
     s[now]=1:
     w[now]=0;
     for (auto x : g[now]){
         if (x!=pre){
             centroid(g, s, w, rec, x, now);
             s[now]+=s[x];
             w[now]=max(w[now], s[x]);
     w[now]=max(w[now], n-s[now]);
     if (w[now]<=n/2){</pre>
         if (rec.first==0) rec.first=now;
         else rec.second=now;
 int dfs(Graph &g, Hash &m, int &id, int now, int pre){
     vector<int> v:
     for (auto x : g[now]){
         if (x!=pre){
```

```
int add=dfs(g, m, id, x, now);
               v.push back(add);
       sort(v.begin(), v.end());
       if (m.find(v)!=m.end()){
           return m[v]:
       }else{
           m[v]=++id;
           return id;
   void solve1(){
       // init
       id1=0:
       id2=0;
       c1={0, 0};
       c2=\{0, 0\};
       fill(sz1.begin(), sz1.begin()+n+1, 0);
       fill(sz2.begin(), sz2.begin()+n+1, 0);
       fill(we1.begin(), we1.begin()+n+1, 0);
       fill(we2.begin(), we2.begin()+n+1, 0);
       for (int i=1 ; i<=n ; i++){</pre>
           g1[i].clear();
           g2[i].clear();
       m1.clear();
       m2.clear();
       // input
       cin >> n:
       for (int i=0 ; i<n-1 ; i++){</pre>
           cin >> a >> b;
           g1[a].push_back(b);
           g1[b].push_back(a);
       for (int i=0; i<n-1; i++){
           cin >> a >> b;
           g2[a].push_back(b);
           g2[b].push back(a);
       // get tree centroid
       centroid(g1, sz1, we1, c1, 1, 0);
       centroid(g2, sz2, we2, c2, 1, 0);
       // process
       int res1=0, res2=0, res3=0;
       if (c2.second!=0){
           res1=dfs(g1, m1, id1, c1.first, 0);
101
           m2=m1:
102
           id2=id1;
           res2=dfs(g2, m1, id1, c2.first, 0);
103
           res3=dfs(g2, m2, id2, c2.second, 0);
       }else if (c1.second!=0){
105
106
           res1=dfs(g2, m1, id1, c2.first, 0);
107
108
           id2=id1;
           res2=dfs(g1, m1, id1, c1.first, 0);
109
           res3=dfs(g1, m2, id2, c1.second, 0);
110
111
           res1=dfs(g1, m1, id1, c1.first, 0);
```

```
res2=dfs(g2, m1, id1, c2.first, 0);
114
       }
115
116
117
       cout << (res1==res2 || res1==res3 ? "YES" : "NO") << endl
       return:
120
121
   signed main(void){
       fastio;
       int t=1:
       cin >> t;
       while (t--){
128
           solve1();
129
130
       return 0:
131 }
   6.5 Bridge BCC
```

```
| #include <bits/stdc++.h>
 using namespace std;
 const int N = 200005;
 vector <int> G[N];
 int low[N], depth[N];
 bool vis[N];
 vector <vector <int>> bcc;
 stack <int> stk;
 void dfs(int v, int p) {
     stk.push(v);
     vis[v] = true;
     low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
     for (int u : G[v]) {
         if (u == p) continue;
         if (!vis[u]) {
             /// (v, u) 是樹邊
             dfs(u, v);
             low[v] = min(low[v], low[u]);
        } else {
             /// (v, u) 是回邊
             low[v] = min(low[v], depth[u]);
     /// v 在不依靠父邊的情況下永遠沒辦法走到它的祖先
     if (low[v] == depth[v]) {
         bcc.emplace back();
         while (stk.top() != v) {
             bcc.back().push_back(stk.top());
             stk.pop();
         bcc.back().push back(stk.top());
         stk.pop();
```

#### 6.6 Cut BCC

```
| #include <bits/stdc++.h>
  using namespace std;
  const int N = 200005;
  vector <int> G[N];
  int low[N], depth[N];
  bool vis[N];
  vector <vector <int>> bcc;
  stack <int> stk:
  void dfs(int v, int p) {
      stk.push(v);
      vis[v] = true;
      low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
      for (int u : G[v]) {
          if (u == p) continue;
          if (!vis[u]) {
              /// (v, u) 是樹邊
              dfs(u, v);
20
              low[v] = min(low[v], low[u]);
              /// u 無法在不經過父邊的情況走到 v 的祖先
              if (low[u] >= depth[v]) {
                  bcc.emplace back();
23
                  while (stk.top() != u) {
                      bcc.back().push_back(stk.top());
                      stk.pop();
                  bcc.back().push back(stk.top());
                  stk.pop();
                  bcc.back().push_back(v);
          } else {
33
              /// (v, u) 是回邊
              low[v] = min(low[v], depth[u]);
34
35
36
37 }
```

### 6.7 圓方樹

```
| #include <bits/stdc++.h>
  #define lp(i,a,b) for(int i=(a);i<(b);i++)</pre>
  #define pii pair<int,int>
  #define pb push back
  #define ins insert
  #define ff first
  #define ss second
  #define opa(x) cerr << #x << " = " << x << ", ";
  #define op(x) cerr << #x << " = " << x << endl;
10 #define ops(x) cerr << x;
#define etr cerr << endl;</pre>
#define spc cerr << ' ';</pre>
#define BAE(x) (x).begin(), (x).end()
14 #define STL(x) cerr << #x << " : "; for(auto &qwe:x) cerr <<</pre>
qwe << ' '; cerr << endl;
15 #define deb1 cerr << "deb1" << endl;</pre>
#define deb2 cerr << "deb2" << endl:
#define deb3 cerr << "deb3" << endl;
#define deb4 cerr << "deb4" << endl;</pre>
```

```
19 #define deb5 cerr << "deb5" << endl;</pre>
20 #define bye exit(0);
                                                                    84 }
21 using namespace std:
                                                                      int dep[mxn], jmp[mxn][mxlg];
                                                                       void dfs lca(int v, int par, int depth){
23 const int mxn = (int)(2e5) + 10;
24 const int mxlg = 17:
                                                                           dep[v] = depth;
25 int last special node = (int)(1e5) + 1;
                                                                           for(auto &i:F[v]){
26 vector<int> E[mxn], F[mxn];
                                                                               if(i == par) continue;
                                                                               jmp[i][0] = v;
  struct edg{
                                                                               dfs_lca(i, v, depth + 1);
      int fr. to:
                                                                          }
      edg(int _fr, int _to){
          fr = _fr;
                                                                       inline void build lca(){
          to = to;
                                                                           jmp[1][0] = 1;
                                                                           dfs lca(1, -1, 1);
                                                                          lp(\bar{j},1,mxlg){
35 ostream& operator << (ostream& os, edg x) {os << x.fr << "--" << 99
        x.to:}
                                                                               lp(i,1,mxn){
  vector<edg> EV:
                                                                                   jmp[i][j] = jmp[jmp[i][j-1]][j-1];
                                                                   102
  void tarjan(int v, int par, stack<int>& S){
                                                                          }
                                                                   103
      static vector<int> dfn(mxn), low(mxn);
                                                                   104
      static vector<bool> to add(mxn);
      static int nowT = 0:
                                                                       inline int lca(int x, int y){
                                                                           if(dep[x] < dep[y]){ swap(x, y); }</pre>
      int childs = 0;
                                                                           int diff = dep[x] - dep[y];
      nowT += 1:
      dfn[v] = low[v] = nowT;
                                                                   110
                                                                           lp(j,0,mxlg){
      for(auto &ne:E[v]){
                                                                               if((diff >> j) & 1){
                                                                   111
          int i = EV[ne].to;
                                                                                   x = jmp[x][j];
                                                                   112
          if(i == par) continue;
                                                                   113
          if(!dfn[i]){
                                                                   114
              S.push(ne);
                                                                           if(x == y) return x;
                                                                   115
               tarjan(i, v, S);
                                                                   116
              childs += 1:
                                                                   117
                                                                           for(int j = mxlg - 1; j >= 0; j--){
              low[v] = min(low[v], low[i]);
                                                                               if(jmp[x][j] != jmp[y][j]){
                                                                   118
                                                                                   x = jmp[x][j];
                                                                   119
              if(par >= 0 && low[i] >= dfn[v]){
                                                                                   y = jmp[y][j];
                                                                   120
                   vector<int> bcc;
                                                                   121
                   int tmp;
                                                                   122
                   do{
                                                                           return jmp[x][0];
                                                                   123
                       tmp = S.top(); S.pop();
                                                                   124
                       if(!to_add[EV[tmp].fr]){
                            to add[EV[tmp].fr] = true;
                                                                       inline bool can_reach(int fr, int to){
                           bcc.pb(EV[tmp].fr);
                                                                          if(dep[to] > dep[fr]) return false;
                                                                   128
                       if(!to_add[EV[tmp].to]){
                                                                           int diff = dep[fr] - dep[to];
                                                                   129
                            to add[EV[tmp].to] = true;
                                                                           lp(i,0,mxlg){
                                                                   130
                           bcc.pb(EV[tmp].to);
                                                                               if((diff >> j) & 1){
                                                                   131
                                                                                   fr = jmp[fr][j];
                                                                   132
                   }while(tmp != ne);
                                                                   133
                   for(auto &j:bcc){
                                                                   134
                       to add[i] = false;
                                                                   135
                                                                           return fr == to;
                       F[last special node].pb(j);
                                                                   136
                       F[j].pb(last special node);
                                                                   137
                   last special node += 1;
                                                                           ios::sync with stdio(false); cin.tie(0);
                                                                          freopen("test_input.txt", "r", stdin);
                                                                           int n, m, q; cin >> n >> m >> q;
          else{
                                                                           lp(i,0,m){
              low[v] = min(low[v], dfn[i]);
if(dfn[i] < dfn[v]){ // edge i--v will be visited 144
                                                                               int u, v; cin >> u >> v;
                                                                               E[u].pb(EV.size());
                     twice at here, but we only need one.
                                                                               EV.pb(edg(u, v));
                                                                   145
                   S.push(ne);
                                                                               E[v].pb(EV.size());
                                                                   146
                                                                               EV.pb(edg(v, u));
                                                                   147
                                                                   148
```

```
E[0].pb(EV.size());
150
       EV.pb(edg(0, 1));
       stack<int> S:
       tarjan(0, -1, S);
153
       build lca();
154
155
       lp(queries,0,a){
            int fr. to. relav: cin >> fr >> to >> relav:
156
            if(fr == relay || to == relay){
157
158
                cout << "NO \ n";
159
                continue:
160
            if((can_reach(fr, relay) || can_reach(to, relay)) &&
161
                 dep[relay] >= dep[lca(fr, to)]){
                cout << "NO\n";
162
                continue:
163
164
165
            cout << "YES\n";</pre>
166
167
```

#### 6.8 SCC 與縮點

```
2 | 給定一個有向圖, 迴回傳縮點後的圖、SCC 的資訊
  所有點都以 based-0 編號
5 函式:
 6 | SCC compress G(n): 宣告─個有 n 個點的圖
7|.add_edge(u, v): 加上一條邊 u -> v
8 .compress: O(n Log n) 計算 G3、SCC、SCC id 的資訊,並把縮點後
       的結果存在 result 裡
10 | SCC[i] = 某個 SCC 中的所有點
11 | SCC id[i] = 第 i 個點在第幾個 SCC
12 */
13 // c8b146
14 struct SCC compress{
      int n = 0, m = 0;
      vector<vector<int>>> G, inv G, result;
      vector<pair<int, int>> edges:
      vector<bool> vis;
      vector<int> order:
21
      vector<vector<int>> SCC;
      vector<int> SCC id:
      SCC compress(int n){
         n = n;
         G.resize(n);
          inv G.resize(n);
          result.resize(n);
          vis.resize(n);
         SCC id.resize(n);
30
31
33
      void add edge(int u, int v){
         G[u].push_back(v);
34
          inv G[v].push back(u);
35
          edges.push back({u, v});
          m++;
```

```
}
     void dfs1(vector<vector<int>> &G, int now){
         vis[now] = 1;
         for (auto x : G[now]){
             if (vis[x]==0){
                 dfs1(G, x);
         order.push_back(now);
         return;
     }
     void dfs2(vector<vector<int>> &G, int now){
         SCC_id[now] = SCC.size()-1;
         SCC.back().push back(now);
         vis[now] = 1;
         for (auto x : G[now]){
             if (vis[x]==0){
                 dfs2(G, x);
         return;
     void compress(){
         fill(vis.begin(), vis.end(), 0);
         for (int i=0 ; i<n ; i++){</pre>
             if (vis[i]==0){
                 dfs1(G, i);
         fill(vis.begin(), vis.end(), 0);
         reverse(order.begin(), order.end());
         for (int i=0 ; i<n ; i++){</pre>
             if (vis[order[i]]==0){
                 SCC.push_back(vector<int>());
                 dfs2(inv_G, order[i]);
         }
         for (int i=0 ; i<m ; i++){</pre>
             if (SCC_id[edges[i].first]!=SCC_id[edges[i].
                  result[SCC_id[edges[i].first]].push_back(
                      SCC_id[edges[i].second]);
         for (int i=0 ; i<SCC.size() ; i++){</pre>
             sort(result[i].begin(), result[i].end());
             result[i].resize(unique(result[i].begin(), result
                   [i].end())-result[i].begin());
       Dinic
1 // 時間複雜度: O(V^2E)
2 struct Flow{
```

struct Edge{

```
int v, rc, rid;
};
vector<vector<Edge>> G;
void add(int u, int v, int c){
    G[u].push_back({v, c, G[v].size()});
    G[v].push back({u, 0, G[u].size()-1});
vector<int> dis, it;
Flow(int n){
    G.resize(n):
    dis.resize(n);
    it.resize(n);
int dfs(int u, int t, int f){
    if (u==t || f==0) return f;
    for (int &i=it[u]; i<G[u].size(); i++){</pre>
        auto &[v, rc, rid] = G[u][i];
        if (dis[v]!=dis[u]+1) continue;
        int df = dfs(v, t, min(f, rc));
        if (df<=0) continue;</pre>
        rc -= df;
        G[v][rid].rc += df;
        return df;
    return 0;
int flow(int s, int t){
    int ans = 0;
    while (true){
        fill(dis.begin(), dis.end(), INF);
        queue<int> q;
        q.push(s);
        dis[s] = 0;
        while (q.size()){
            int u = q.front(); q.pop();
            for (auto [v, rc, rid] : G[u]){
   if (rc<=0 || dis[v]<INF) continue;</pre>
                 dis[v] = dis[u]+1;
                 q.push(v);
        if (dis[t]==INF) break;
        fill(it.begin(), it.end(), 0);
        while (true){
            int df = dfs(s, t, INF);
            if (df<=0) break;</pre>
            ans += df;
    return ans;
// the code below constructs minimum cut
void dfs mincut(int now, vector<bool> &vis){
vis[now] = true;
for (auto &[v, rc, rid] : G[now]){
 if (vis[v]==false && rc>0){
    dfs mincut(v, vis);
```

### 6.10 Dijkstra

```
1 // 可以在 O(E Log E) 的時間複雜度解決在無負權有向圖單點源最短
  const int INF = 2e18; // 要確保 INF 開的足夠大
  vector<vector<pair<int, int>>> G(n); // G[i] = <節點, 權重>
  vector<int> dis(n, INF);
  priority_queue<pair<int, int>, vector<pair<int, int>>,
      greater<pair<int, int>>> pq;
  dis[s] = 0;
  pq.push({0, s});
  while (pq.size()){
      int now_dis = pq.top().first;
      int now_node = pq.top().second;
      pq.pop();
      if (now dis>dis[now node]) continue;
      for (auto x : G[now node]){
         if (now dis+x.second<dis[x.first]){</pre>
             dis[x.first] = now_dis+x.second;
              pq.push({dis[x.first], x.first});
20
21
22
23
```

### 6.11 定理

- 最小點覆蓋 = 最大匹配 = n 最大點獨立集
  - 最小點覆蓋:選最少點讓所有的邊都有碰到一個點
  - 最大點獨立集:選最多不共邊的點
- 只有邊帶權的二分圖的定理(可能不重要)
  - w-vertex-cover ( 帶權點覆蓋 ):每條邊的兩個連接點被選中的次數總和至少要是 $w_e$ 。
  - w-weight matching ( 帶權匹配 )
  - minimum vertex count of w-vertex-cover = maximum weight count of w-weight matching (一個點可以被選很多次·但邊不行)
- 點、邊都帶權的二分圖的定理(可能不重要)

- b-matching:假設 v 的點權是  $b_v$  · 那所有 v 的匹配邊 e 的權重都要 57 滿足  $\sum w_e \leq b_v$  ° 58
- The maximum w-weight of a b-matching equals the minimum b-weight of vertices in a w-vertex-cover.

#### **6.12** MCMF

struct Edge {

| struct Flow {

```
int u, rc, k, rv;
};
vector<vector<Edge>> G;
vector<int> par, par_eid;
Flow(int n): G(n+1), par(n+1), par eid(n+1) {}
// v->u, capcity: c, cost: k
void add(int v, int u, int c, int k){
  G[v].push_back({u, c, k, SZ(G[u])});
  G[u].push_back({v, 0, -k, SZ(G[v])-1});
// 3701d6
int spfa(int s, int t){
  fill(ALL(par), -1);
  vector<int> dis(SZ(par), INF);
  vector<bool> in_q(SZ(par), false);
  queue<int> 0;
  dis[s] = 0;
  in_q[s] = true;
  Q.push(s);
  while (!O.empty()){
    int v = Q.front();
    Q.pop();
    in_q[v] = false;
    for (int i=0 ; i<SZ(G[v]) ; i++){</pre>
      auto [u, rc, k, rv] = G[v][i];
      if (rc>0 && dis[v]+k<dis[u]){</pre>
        dis[u] = dis[v]+k;
        par[u] = v;
        par eid[u] = i;
        if (!in_q[u]) Q.push(u);
        in_q[u] = true;
  return dis[t];
// return <max flow, min cost>, 150093
pair<int, int> flow(int s, int t){
  int fl = 0, cost = 0, d;
  while ((d = spfa(s, t))<INF){</pre>
    int cur = INF;
    for (int v=t ; v!=s ; v=par[v])
      cur = min(cur, G[par[v]][par_eid[v]].rc);
    fl += cur;
    cost += d*cur:
    for (int v=t ; v!=s ; v=par[v]){
      G[par[v]][par_eid[v]].rc -= cur;
```

### 7 Math

#### 7.1 Burnside's-Lemma

$$\sum_{k=1}^{n} \frac{c(k)}{n}$$

- n:有多少種置換方式(例如:旋轉方式)
- c(k): 所有可能中,經過 k 次旋轉後,仍不會和別人相同的方式的數量

### 7.2 線性篩

```
| const int MAX_N = 5e5;
| const int MAX_N =
```

### 7.3 Lucas's-Theorem

```
1 // 對於很大的 C^n_{m} 對質數 p 取模·只要 p 不大就可以用。
2 int Lucas(int n, int m, int p){
3 if (m==0) return 1;
4 return (C(n%p, m%p, p)*Lucas(n/p, m/p, p)%p);
5 }
```

#### 7.4 Matrix

```
| struct Matrix{
      int n, m;
      vector<vector<int>> arr;
      Matrix(int _n, int _m){
          n = _n;
          m = _m;
          arr.resize(n, vector<int>(m));
      Matrix operator * (const Matrix B){
          Matrix ret(n, B.m);
          for (int i=0 ; i<n ; i++){</pre>
               for (int j=0 ; j<B.m ; j++){</pre>
                   for (int k=0; k<m; k++){
                       ret.arr[i][j] += arr[i][k]*B.arr[k][j];
                       ret.arr[i][j] %= MOD;
21
23
          return ret;
24
```

### 7.5 最大質因數

```
void max_fac(int n, int &ret){
    if (n<=ret || n<2) return;
    if (isprime(n)){
        ret = max(ret, n);
        return;
}

int p = Pollard_Rho(n);
max_fac(p, ret), max_fac(n/p, ret);
}</pre>
```

# 7.6 中國剩餘定理(m 不互質)

```
int extgcd(int a, int b, int &x, int &y){
   if (b==0){
      x=1, y=0;
      return a;
}
```

### 7.7 歐拉公式

```
1 | // phi(n) = 小於 n 並與 n 互質的正整數數量。
2 // O(sqrt(n)) · 回傳 phi(n)
3 int phi(int n){
     int ret = n;
      for (int i=2 ; i*i<=n ; i++){</pre>
          if (n%i==0){
              while (n%i==0) n /= i;
              ret = ret*(i-1)/i:
     if (n>1) ret = ret*(n-1)/n;
      return ret:
17 // O(n Log n) · 回傳 1~n 的 phi 值
vector<int> phi_1_to_n(int n){
      vector<int> phi(n+1);
      phi[0]=0;
      phi[1]=1;
      for (int i=2 ; i<=n ; i++){</pre>
          phi[i]=i-1;
      for (int i=2 ; i<=n ; i++){</pre>
          for (int j=2*i; j<=n; j+=i){ // 枚舉所有倍數
              phi[j]-=phi[i];
     }
      return phi;
```

```
7.8 歐拉定理
```

```
若 a,m 互質 \cdot 則: a^n \bmod m = a^{n \bmod \varphi(m)} \bmod m 若 a,m 可能是任何數 \cdot 則: a^{\varphi(m)+[n \bmod \varphi(m)]} \bmod m
```

#### 7.9 Fraction

```
#include <bits/stdc++.h>
using namespace std;
/// Fraction template starts ///
#define fraction template bonus check
const long long long loverflow_warning_value = (long long)(3e9);
long long gcd(long long a, long long b){
    if(a == 0) return 0;
    if(b == 0) return a;
    if(a < b) return gcd(b,a);</pre>
    return gcd(b, a%b);
struct frac{
    long long a, b;
    frac(long long _a = 0, long long _b = 1){
        a = _a; b = _b;
        if(b == 0){
            cerr << "Error: division by zero\n";</pre>
            cerr << "Called : Constructor(" << _a << ", " <<</pre>
                 _b << ")\n";
            return;
        if(a == 0){b = 1; return;}
        if(b < 0){a = -a; b = -b;}
        long long gcd_ab = gcd(std::abs(a), b);
        if(gcd ab != 1){a /= gcd ab; b /= gcd ab;}
        #ifdef fraction template bonus check
        if(std::abs(a) > 11 overflow warning value || b >
             11 overflow_warning_value){
            cerr << "Overflow warning : " << a << "/" << b << 95
        #endif // fraction_template_bonus_check
    frac operator+(frac const &B){
        return frac(a*(B.b)+(B.a)*b, b*(B.b));}
    frac operator-(frac const &B){
        return frac(a*(B.b)-(B.a)*b, b*(B.b));}
    frac operator*(frac const &B){
        return frac(a*(B.a), b*(B.b));}
    frac operator/(frac const &B){
        return frac(a*(B.b), b*(B.a));}
    frac operator+=(frac const &B){
        *this = frac(a*(B.b)+(B.a)*b, b*(B.b));}
    frac operator -= (frac const &B){
        *this = frac(a*(B.b)-(B.a)*b, b*(B.b));}
    frac operator*=(frac const &B){
```

\*this = frac(a\*(B.a), b\*(B.b));}

```
frac operator/=(frac const &B){
          *this = frac(a*(B.b), b*(B.a));}
     frac abs(){
         a = std::abs(a);
          return *this:
     bool operator<(frac const &B){</pre>
          return a*B.b < B.a*b;}</pre>
     bool operator <= (frac const &B){</pre>
          return a*B.b <= B.a*b;}
     bool operator>(frac const &B){
          return a*B.b > B.a*b;}
     bool operator>=(frac const &B){
          return a*B.b >= B.a*b;}
     bool operator == (frac const &B){
          return a * B.b == B.a * b;}
     bool operator!=(frac const &B){
          return a * B.b != B.a * b;}
ostream& operator<<(ostream &os, const frac& A){
    os << A.a << "/" << A.b;
    return os:
/// Fraction template ends ///
void test(frac A, frac B){
    cout << "A = " << A << endl;
cout << "B = " << B << endl;
    cout << endl;</pre>
    cout << "A + B = " << A + B << endl;
    cout << "A - B = " << A - B << endl;
    cout << "A * B = " << A * B << endl;
    cout << "A / B = " << A / B << endl;
    cout << endl;</pre>
    cout \langle\langle "(A \langle B) = " \langle\langle (A \langle B) \langle\langle endl;
    cout \langle\langle "(A \langle = B) = " \langle\langle (A \langle = B) \langle\langle endl;
    cout \langle\langle "(A > B) = " \langle\langle (A > B) \langle\langle endl;
    cout << "(A >= B) = " << (A >= B) << endl;
cout << "(A == B) = " << (A == B) << endl;
cout << "(A != B) = " << (A != B) << endl;
    cout << "----\n":
    return:
int main(){
     frac tmp1(-7, 2);
     frac tmp2(5, 3);
    test(tmp1, tmp2);
     frac tmp3(-7);
     frac tmp4(0);
    test(tmp3, tmp4);
     return 0;
```

### 7.10 錯排公式

錯排公式:  $(n \oplus 1)$  個人中,每個人皆不再原來位置的組合數)

```
dp_i = \begin{cases} 1 & i = 0\\ 0 & i = 1\\ (i-1)(dp_{i-1} + dp_{i-2}) & \text{otherwise} \end{cases}
```

### 7.11 Quick-Pow

```
int qp(int b, int p, int m = MOD){
   int ret = 1;
   for (; p; p>>=1){
      if (p&1) ret = ret*b%m;
      b = b*b%m;
   }
   return ret;
}
```

# 7.12 二元一次方程式

```
\begin{cases} ax+by=c\\ dx+ey=f\\ \exists \ x=\frac{0}{0}\ \exists\ y=\frac{0}{0}^{\bullet}\cdot \mathbb{R}, \ \mathbb{R} \end{cases} = \begin{cases} x=\frac{ed-bf}{ad-bc}\\ y=\frac{af-cc}{ad-bc}\\ \exists \ x=\frac{0}{0}\ \exists\ y=\frac{*}{0}\cdot \mathbb{R}, \ \mathbb{R} \end{cases}
```

# 7.13 Josephus

```
1 // 有 n 個人 · 第偶數個報數的人被刪掉 · 問第 k 個被踢掉的是誰
2 int solve(int n, int k){
3    if (n==1) return 1;
4    if (k<=(n+1)/2){
5        if (2*k>n) return 2*k%n;
6        else return 2*k;
7    }else{
6        int res=solve(n/2, k-(n+1)/2);
7        if (n&1) return 2*res+1;
8        else return 2*res-1;
9    }
1    }
2 }
```

## 7.14 數論分塊

```
| /*
| 時間複雜度為 O(sqrt(n))
| 區間為 [L, r]
| */
| for(int i=1; i<=n; i++){
| int l = i, r = n/(n/i);
| i = r;
| ans.push_back(r);
| 9 }
```

### 7.15 Pollard-Rho

```
i | mt19937 seed(chrono::steady clock::now().time since epoch().
 int rnd(int 1, int r){
     return uniform int distribution<int>(1, r)(seed);
6 | // O(n^{1/4}) 回傳 1 或自己的因數、記得先判斷 n 是不是質數
       (用 Miller-Rabin)
 // c1670c
8 int Pollard_Rho(int n){
     int s = 0, t = 0;
     int c = rnd(1, n-1);
     int step = 0, goal = 1;
     int val = 1;
      for (goal=1 ; ; goal<<=1, s=t, val=1){</pre>
          for (step=1 ; step<=goal ; step++){</pre>
              t = ((int128)t*t+c)%n;
              val = (__int128)val*abs(t-s)%n;
             if ((step % 127) == 0){
                 int d = __gcd(val, n);
                 if (d>1) return d;
         int d = __gcd(val, n);
         if (d>1) return d;
```

# 7.16 中國剩餘定理 (m 互質)

```
vector<int> a, m;
  int extgcd(int a, int b, int &x, int &y){
      if (b==0){
          x=1, y=0;
          return a;
      int ret=extgcd(b, a%b, y, x);
      y-=a/b*x;
      return ret;
14 // n = 有幾個式子, 求解 x \equiv a i \bmod m i
  int CRT(int n, vector<int> &a, vector<int> &m){
      int p=1, ans=0;
      vector<int> M(n), inv_M(n);
       for (int i=0 ; i<n ; i++) p*=m[i];</pre>
       for (int i=0 ; i<n ; i++){</pre>
22
          M[i]=p/m[i];
          extgcd(M[i], m[i], inv_M[i], tmp);
```

#### 7.17 Catalan

任意括號序列:  $C_n = \frac{1}{n+1} \binom{2n}{n}$ 

### 7.18 數論定理

- 1.  $1 \sim x$  質數的數量  $\approx \frac{x}{\ln x}$
- 2.  $1 \sim x$  的因數的數量  $\approx x^{\frac{1}{3}}$
- 3. x 的質因數的數量  $\approx \log \log x$
- 4. p is a prime number  $\Leftrightarrow (p-1)! \equiv -1 \pmod{p}$
- 5. 每個正整數都可以表示成四個整數的平方和
- 6. 任何大於 2 的整數都可以表示成兩個質數的和

#### 7.19 Miller-Rabin

```
1 // O(\log n)
 typedef Uint unsigned long long
 Uint modmul(Uint a, Uint b, Uint m) {
     int ret = a*b - m*(Uint)(1.L/m*a*b);
     return ret + m*(ret < 0) - m*(ret>=(int)m);
  int qp(int b, int p, int m){
     int ret = 1;
     for ( ; p ; p>>=1){
         if (p&1){
              ret = modmul(ret, b, m);
         b = modmul(b, b, m);
     return ret;
 vector<int> 11sprp = {2, 325, 9375, 28178, 450775, 9780504,
      1795265022};
 bool isprime(int n, vector(int) sprp = llsprp){
     if (n==2) return 1;
     if (n<2 || n%2==0) return 0;</pre>
     int t = 0;
     int u = n-1;
     for ( ; u%2==0 ; t++) u>>=1;
     for (int i=0 ; i<sprp.size() ; i++){</pre>
         int a = sprp[i]%n;
         if (a==0 || a==1 || a==n-1) continue;
```

```
int x = qp(a, u, n);
if (x==1 || x==n-1) continue;
for (int j=0; j<t; j++){
    x = modmul(x, x, n);
    if (x==1) return 0;
    if (x==n-1) break;
}

if (x==n-1) continue;
    return 0;
}

return 1;
}</pre>
```

### 7.20 Stirling's formula

 $n! \approx \sqrt{2\pi n} (\frac{n}{e})^n$ 

# 8 String

#### **8.1** Hash

```
i | mt19937 rnd(chrono::steady_clock::now().time_since_epoch().
       count());
  int A = rnd(), B = 1000000007;
  vector<int> myPow, myPre;
  void hash_init(string s){
      myPow.resize(s.size());
      myPre.resize(s.size());
      for (int i=0 ; i<s.size() ; i++){</pre>
          if (i==0){
              myPow[i] = 1;
              myPre[i] = s[i];
              myPow[i] = myPow[i-1]*A%B;
              myPre[i] = (myPre[i-1]*A+s[i])%B;
      }
      return;
22 int hash value(int 1, int r){ // 取得 s[l..r] 的數值
      if (l==0) return myPre[r];
      return ((myPre[r]-myPre[1-1]*myPow[r-1+1])%B+B)%B;
```

# 8.2 Manacher

```
string Manacher(string str) {
    string tmp = "$#";
    for(char i : str) {
        tmp += i;
}
```

```
tmp += '#';
}

vector<int> p(tmp.size(), 0);
int mx = 0, id = 0, len = 0, center = 0;
for(int i=1; i<(int)tmp.size(); i++) {
    p[i] = mx > i ? min(p[id*2-i], mx-i) : 1;

while(tmp[i+p[i]] == tmp[i-p[i]]) p[i]++;
    if(mx<i+p[i]) mx = i+p[i], id = i;
    if(len<p[i]) len = p[i], center = i;
}
return str.substr((center-len)/2, len-1);
}</pre>
```

### 8.3 Z-Function

```
1// 定義一個長度為 n 的文本為 T ,則陣列 Z 的 Z[i] 代表 T[0:n]
       和 T[i:n] 最長共同前綴
2 // bcfbd6
 vector<int> z_function(string s){
     vector<int> ret(s.size());
     int 11 = 0, rr = 0;
     for (int i=1; i<s.size(); i++){</pre>
         int j = 0;
         if (i<rr) j = min(ret[i-ll], rr-i);</pre>
         while (s[j]==s[i+j]) j++;
         ret[i] = j;
         if (i+j>rr){
             11 = i:
             rr = i+j;
     ret[0] = s.size();
     return ret;
```

### 8.4 KMP

### 8.5 Suffix-Array

```
ı|// 注意,當 /s/=1 時,Lcp 不會有值,務必測試 /s/=1 的 case
  struct SuffixArray {
      vector<int> sa, lcp;
      SuffixArray(string& s, int lim = 256) {
          // 49c4d2
          int n = SZ(s)+1, k = 0, a, b;
          vector<int> x(ALL(s)), y(n), ws(max(n, lim)), rank(n)
          x.push back(0);
          sa = 1cp = v;
          iota(ALL(sa), 0);
          for (int j=0, p=0 ; p<n ; j=max(1LL, j*2), lim=p) {</pre>
               p = j;
               iota(ALL(y), n-j);
               for (int i=0 ; i<n ; i++) if (sa[i] >= j) y[p++]
                   = sa[i] - j;
               fill(ALL(ws), 0);
               for (int i=0 ; i<n ; i++) ws[x[i]]++;</pre>
               for (int i=1; i<lim; i++) ws[i] += ws[i - 1];</pre>
               for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
               swap(x, y), p = 1, x[sa[0]] = 0;
               for (int i=1 ; i<n ; i++){</pre>
21
                  a = sa[i - 1];
                   b = sa[i];
23
                   x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
                         ? p - 1 : p++;
          // 7181dd
          for (int i=1; i<n; i++) rank[sa[i]] = i;</pre>
          for (int i=0, j ; i<n-1 ; lcp[rank[i++]]=k)</pre>
               for (k && k--, j=sa[rank[i]-1]; i+k<SZ(s) && j+k</pre>
                    SZ(s) \& s[i+k]==s[j+k]; k++);
          sa.erase(sa.begin());
          lcp.erase(lcp.begin(), lcp.begin()+2);
32
33
34 };
```

#### 8.6 Min-Rotation