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1 Misc

1.1 Custom Set PQ Sort

```
1  // priority_queue · 務必檢查相等的 case · 給所有元素一個排序的
依據
struct cmp{
    bool operator () (Data a, Data b){
        return a.x<b.x;
    }
};
priority_queue<Data, vector<Data>, cmp> pq;
8
// set · 務必檢查相等的 case · 給所有元素一個排序的依據
auto cmp = [](int a, int b) {
    return a > b;
};
set<int, decltype(cmp)> s = {1, 2, 3, 4, 5};
cout << *s.begin() << '\n';</pre>
```

1.2 Default Code New

```
1  #include <bits/stdc++.h>
2  using namespace std;
3  #define int long long
4  const int MAX_N = 5e5 + 10;
6  const int INF = 2e18;
7  void solve(){
9  }
10  }
11  signed main(){
12  ios::sync_with_stdio(0), cin.tie(0);
13  int t = 1;
14  while (t--){
15  solve();
16  }
17  return 0;
18  }
19  return 0;
21 }
```

1.3 Default Code Old

```
#include <bits/stdc++.h>
#define int long long

#define ALL(x) x.begin(), x.end()
#define SZ(x) ((int)x.size())

#define fastio ios::sync_with_stdio(0), cin.tie(0);

using namespace std;

#ifdef LOCAL
#define cout cout << "\033[0;32m"
#define cerr cerr << "\033[0;31m"
#define endl endl << "\033[0m"</pre>
```

```
#else
#pragma GCC optimize("03,unroll-loops")
#pragma GCC target("avx,avx2,sse,sse2,sse3,sse4,popcnt")
#define endl "\n"
#endif

const int MAX_N = 5e5+10;
const int INF = 2e18;

void solve1(){

    return;
}

signed main(){

fastio;

int t = 1;
while (t--){
    solve1();
}

return 0;
}
```

1.4 Enumerate Subset

```
1 // 時間複雜度 O(3^n)
2 // 枚舉每個 mask 的子集
3 for (int mask=0; mask<(1<<n); mask++){
    for (int s=mask; s>=0; s=(s-1)&m){
        // s 是 mask 的子集
        if (s==0) break;
7     }
8 }
```

1.5 Fast Input

```
1 // fast IO
2 // 6f8879
3 inline char readchar(){
      static char buffer[BUFSIZ], * now = buffer + BUFSIZ, *
           end = buffer + BUFSIZ;
      if (now == end)
          if (end < buffer + BUFSIZ)</pre>
               return EOF;
          end = (buffer + fread(buffer, 1, BUFSIZ, stdin));
          now = buffer;
      return *now++;
14 inline int nextint(){
      int x = 0, c = readchar(), neg = false;
      while(('0' > c | | c > '9') && c!='-' && c!=EOF) c =
           readchar();
      if(c == '-') neg = true, c = readchar();
      while ('0' \le c \&\& c \le '9') x = (x << 3) + (x << 1) + (c^'0')
           , c = readchar();
```

1.6 OEIS

```
ı | / / 若一個線性遞迴有 k 項·給他恰好 2*k 個項可以求出線性遞迴
  // f915c2
  template <typename T>
  vector<T> BerlekampMassey(vector<T> a) {
      auto scalarProduct = [](vector<T> v, T c) {
          for (T &x: v) x *= c;
          return v;
      };
      vector<T> s, best;
      int bestPos = 0;
      for (size_t i = 0; i < a.size(); i++) {</pre>
          T error = a[i];
          for (size_t j = 0; j < s.size(); j++) error -= s[j] *</pre>
                a[i-1-j];
          if (error == 0) continue;
          if (s.empty()) {
              s.resize(i + 1);
              bestPos = i;
              best.push_back(1 / error);
              continue:
          vector<T> fix = scalarProduct(best, error);
          fix.insert(fix.begin(), i - bestPos - 1, 0);
          if (fix.size() >= s.size()) {
              best = scalarProduct(s, - 1 / error);
              best.insert(best.begin(), 1 / error);
              bestPos = i;
              s.resize(fix.size());
          for (size_t j = 0; j < fix.size(); j++)</pre>
              s[j] += fix[j];
31
32
      return s;
33
```

1.7 Xor Basis

```
vector<int> basis;
void add_vector(int x){
    for (auto v : basis){
        x=min(x, x^v);
    }
    if (x) basis.push_back(x);
}
// 給一數字集合 S·求能不能 XOR 出 x
bool check(int x){
    for (auto v : basis){
        x=min(x, x^v);
    }
    return 0;
}
```

```
17 / / 給一數字集合 S · 求能 XOR 出多少數字
18 // 答案等於 2^{basis 的大小}
20 // 給一數字集合 S · 求 XOR 出最大的數字
21 int get_max(){
     int ans=0;
     for (auto v : basis){
         ans=max(ans, ans^v);
     return ans;
```

1.8 random int

```
i| mt19937 seed(chrono::steady_clock::now().time_since_epoch().
      count());
2 int rng(int 1, int r){
     return uniform int distribution<int>(1, r)(seed);
```

1.9 Python

```
1 # system setting
  sys.setrecursionlimit(100000)
  sys.set_int_max_str_digits(10000)
  # turtle
6 from turtle import *
8 \mid N = 3000000010
  setworldcoordinates(-N, -N, N, N)
10 hideturtle()
  speed(100)
13 def draw_line(a, b, c, d):
      teleport(a, b)
      goto(c, d)
| | def write_dot(x, y, text, diff=1): # diff = 文字的偏移
      teleport(x, y)
      dot(5, "red")
      teleport(x+N/100*diff, y+N/100*diff)
      write(text, font=("Arial", 5, "bold"))
25 draw_line(*a[i], *(a[i-1]))
26 write_dot(*a[i], str(a[i]))
```

1.10 diff

```
ı set -e
2 g++ ac.cpp -o ac
3 g++ wa.cpp -o wa
4 for ((i=0;;i++))
5 do
```

```
echo "$i"
       python3 gen.py > input
       ./ac < input > ac.out
       ./wa < input > wa.out
      diff ac.out wa.out || break
11 done
```

1.11 hash command

```
1 cat file.cpp | cpp -dD -P -fpreprocessed | tr -d "[:space:]"
       | md5sum | cut -c-6
```

1.12 run

```
import os
  f = "pA"
  while 1:
      i = input("input: ")
       p = os.listdir(".")
      if i != "":
          f = i
       print(f"file = {f}")
       if os.system(f"g++ {f}.cpp -std=c++17 -Wall -Wextra -
           Wshadow -02 -D LOCAL -q -fsanitize=undefined, address
            -o {f}"):
           print("CE")
           continue
      os.system("clear")
       for x in sorted(p):
                                                                   21
           if f in x and ".in" in x:
               print(x)
               if os.system(f"./\{f\} < \{x\}"):
                   print("RE")
               print()
21
                                                                   26
```

1.13 **setup**

```
se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
  :inoremap " ""<Esc>i
  :inoremap {<CR> {<CR>}<Esc>ko
  :inoremap {{ {}}<ESC>i
  function! F(...)
   execute '!./%:r < ./' . a:1
  endfunction
10 command! -nargs=* R call F(<f-args>)
  map <F7> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -Wextra -
       Wshadow -02 -DLOCAL -g -fsanitize=undefined,address<CR>
13 map <F8> :!./%:r<CR>
14 map <F9> :!./%:r < ./%:r.in<CR>
```

```
16 ca hash w !cpp -dD -P -fpreprocessed \| tr -d "[:space:]" \|
       md5sum \ cut -c-6
  " i+<esc>25A---+<esc>
  " o|<esc>25A |<esc>
  " "ggVGyG35pGdd
```

Convolution

2.1 FFT any mod

28

29

30 31

35

41

44

```
_{3} | PolyMul(a, b) 回傳多項式乘法的結果(c k = \sum {i+j} a i+b j
  大約可以支援 5e5 \cdot a_i, b_i 皆在 MOD 以下的非負整數
  const int MOD = 998244353;
  typedef complex<double> cd;
  // b9c90a
  void FFT(vector<cd> &a) {
      int n = a.size(), L = 31-__builtin_clz(n);
      vector<complex<long double>> R(2, 1);
      vector<cd> rt(2, 1);
      for (int k=2; k<n; k*=2){
          R.resize(n);
          rt.resize(n);
          auto x = polar(1.0L, acos(-1.0L) / k);
          for (int i=k; i<2*k; i++){
              rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);
      vector<int> rev(n);
      for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2] | (i&1) << L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
      for (int k=1; k<n; k*=2){</pre>
          for (int i=0 ; i<n ; i+=2*k){</pre>
              for (int j=0 ; j<k ; j++){</pre>
                 auto x = (double *)&rt[j+k];
                 auto y = (double *)&a[i+j+k];
                 cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]*
                      y[0]);
                 a[i+j+k] = a[i+j]-z;
                 a[i+j] += z;
      return;
43
  // d3c65e
  vector<int> PolyMul(vector<int> a, vector<int> b){
     if (a.empty() || b.empty()) return {};
```

```
vector<int> res(a.size()+b.size()-1);
int B = 32- builtin clz(res.size()), n = (1<<B), cut =</pre>
     int(sqrt(MOD));
vector<cd> L(n), R(n), outs(n), outl(n);
for (int i=0 ; i<a.size() ; i++){</pre>
    L[i] = cd((int) a[i]/cut, (int)a[i]%cut);
for (int i=0 ; i<b.size() ; i++){</pre>
    R[i] = cd((int) b[i]/cut, (int)b[i]%cut);
FFT(L);
FFT(R);
for (int i=0 ; i<n ; i++){</pre>
    int j = -i&(n-1);
    outl[j] = (L[i]+conj(L[j])) * R[i]/(2.0*n);
    outs[j] = (L[i]-conj(L[j])) * R[i]/(2.0*n)/1i;
FFT(outl);
FFT(outs);
for (int i=0 ; i<res.size() ; i++){</pre>
    int av = (int)(real(outl[i])+0.5), cv = (int)(imag(
         outs[i])+0.5);
    int bv = (int)(imag(outl[i])+0.5) + (int)(real(outs[i
         1)+0.5);
    res[i] = ((av%MOD*cut+bv) % MOD*cut+cv) % MOD;
return res;
```

2.2 FFT new

```
1 typedef complex < double > cd;
3 // b9c90a
4 void FFT(vector<cd> &a) {
     int n = a.size(), L = 31-_ builtin clz(n);
     vector<complex<long double>> R(2, 1);
     vector<cd> rt(2, 1);
     for (int k=2 ; k<n ; k*=2){</pre>
          R.resize(n);
          rt.resize(n):
          auto x = polar(1.0L, acos(-1.0L) / k);
          for (int i=k ; i<2*k ; i++){</pre>
              rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);
     vector<int> rev(n);
     for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2] | (i&1) << L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
      for (int k=1; k<n; k*=2){</pre>
          for (int i=0; i<n; i+=2*k){
              for (int j=0 ; j<k ; j++){</pre>
                  auto x = (double *)&rt[j+k];
                  auto y = (double *)&a[i+j+k];
```

2.3 FFT short

return res;

}

return:

FFT(in);

FFT(out);

```
| #define int long long
  using Cplx = complex<double>;
  const double pi = acos(-1);
  const int mod = 998244353, g = 3;
  int power(int a, int b) {
      int res = 1:
          if (b & 1) res = res * a % mod;
          a = a * a % mod:
          b >>= 1;
      return res;
int inv(int x) { return power(x, mod - 2); }
  // FFT use Cplx, NTT use ll
  void FFT(vector<int> &a, int n, int op) {
      // n must be 2^k
      vector<int> R(n);
      FOR (i, 0, n - 1)
          R[i] = R[i/2]/2 + (i&1)*(n/2);
      FOR (i, 0, n - 1)
          if (i < R[i]) swap(a[i], a[R[i]]);</pre>
      for (int m = 2; m <= n; m *= 2) {</pre>
          // Cplx w1({cos(2*pi/m), sin(2*pi/m)*op});
```

y[0]);

a[i+j] += z;

if (a.empty() || b.empty()) return {};

copy(a.begin(), a.end(), begin(in));

for (int i=0 ; i<res.size() ; i++){</pre>

res[i] = imag(out[i]) / (4 * n);

for (int i=0 ; i<b.size() ; i++){</pre>

vector<cd> in(n), out(n);

in[i].imag(b[i]);

for (cd& x : in) x *= x;

for (int i=0 ; i<n ; i++){</pre>

vector<double> res(a.size()+b.size()-1);

a[i+j+k] = a[i+j]-z;

vector<double> PolyMul(const vector<double> a, const vector<</pre>

int L = 32 - builtin clz(res.size()), n = 1 << L;</pre>

out[i] = in[-i & (n - 1)] - conj(in[i]);

```
cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]* 26
                                                          int w1 = power(g, (mod-1)/m * op + mod-1);
                                                          for (int i = 0; i < n; i += m) {</pre>
                                                              // Cplx wk({1, 0});
                                                              int wk = 1;
                                                              FOR (k, 0, m / 2 - 1) {
                                                                  auto x = a[i+k], y = a[i+k+m/2] * wk % mod;
                                                                  a[i+k] = (x+y) \% mod;
                                                                  a[i+k+m/2] = (x-y+mod) \% mod;
                                                                  wk = wk * w1 % mod:
                                                      if (op == -1)
                                                          FOR (i, 0, n - 1) {
                                               39
                                                              // a[i] = a[i] / n;
                                                              a[i] = a[i] * inv(n) % mod;
                                               42
                                               43
```

2.4 FWT

```
1// 已經把 mint 刪掉,需要增加註解
vector<int> xor convolution(vector<int> a, vector<int> b, int
       if (k == 0) {
            return vector<int>{a[0] * b[0]};
       vector\langle int \rangle aa(1 \langle \langle (k-1) \rangle, bb(1 \langle \langle (k-1) \rangle;
       FOR (i, 0, (1 << (k - 1)) - 1) {
            aa[i] = a[i] + a[i + (1 << (k - 1))];
           bb[i] = b[i] + b[i + (1 << (k - 1))];
       vector<int> X = xor convolution(aa, bb, k - 1);
       FOR (i, 0, (1 << (k-1)) - 1)
           aa[i] = a[i] - a[i + (1 << (k - 1))];
bb[i] = b[i] - b[i + (1 << (k - 1))];</pre>
       vector<int> Y = xor convolution(aa, bb, k - 1);
       vector<int> c(1 << \overline{k});
       FOR (i, 0, (1 << (k - 1)) - 1) {
                                   ] = (X[i] + Y[i]) / 2;
            c[i + (1 << (k - 1))] = (X[i] - Y[i]) / 2;
21
       return c;
```

2.5 NTT mod 998244353

```
| \text{const} \text{ int } \text{MOD} = (119 << 23) + 1, \text{ROOT} = 62; // = 998244353
 // For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 <<
 // and 483 << 21 (same root). The last two are > 10^9.
 // 9cd58a
 void NTT(vector<int> &a) {
      int n = a.size();
      int L = 31-__builtin_clz(n);
      vector<int> rt(2, 1);
      for (int k=2, s=2; k<n; k*=2, s++){
          rt.resize(n);
```

```
int z[] = {1, qp(ROOT, MOD>>s)};
        for (int i=k ; i<2*k ; i++){</pre>
            rt[i] = rt[i/2]*z[i&1]%MOD;
   }
    vector<int> rev(n);
    for (int i=0 ; i<n ; i++){</pre>
        rev[i] = (rev[i/2]|(i&1)<<L)/2;
    for (int i=0 ; i<n ; i++){</pre>
        if (i<rev[i]){</pre>
            swap(a[i], a[rev[i]]);
   }
    for (int k=1 ; k<n ; k*=2){</pre>
        for (int i=0 ; i<n ; i+=2*k){</pre>
            for (int j=0 ; j<k ; j++){</pre>
                 int z = rt[j+k]*a[i+j+k]%MOD, &ai = a[i+j];
                a[i+j+k] = ai-z+(z>ai ? MOD : 0);
                ai += (ai+z)=MOD ? z-MOD : z);
vector<int> polyMul(vector<int> &a, vector<int> &b){
    if (a.empty() || b.empty()) return {};
   int s = a.size()+b.size()-1, B = 32-__builtin_clz(s), n =
    int inv = qp(n, MOD-2);
    vector<int> L(a), R(b), out(n);
   L.resize(n), R.resize(n);
   NTT(L), NTT(R);
    for (int i=0 ; i<n ; i++){</pre>
        out[-i&(n-1)] = L[i]*R[i]%MOD*inv%MOD;
    NTT(out);
    out.resize(s);
    return out;
```

2.6 Min Convolution Concave Concave

3 Data-Structure

3.1 BIT

3.2 Disjoint Set Persistent

Persistent_Segment_Tree arr, sz;

struct Persistent Disjoint Set{

```
void init(int n){
   arr.init(n);
    vector<int> v1;
    for (int i=0 ; i<n ; i++){</pre>
        v1.push back(i);
   arr.build(v1, 0);
    sz.init(n);
    vector<int> v2;
   for (int i=0 ; i<n ; i++){</pre>
        v2.push_back(1);
    sz.build(v2, 0);
int find(int a){
   int res = arr.query_version(a, a+1, arr.version.size
        ()-1).val;
   if (res==a) return a;
   return find(res);
bool unite(int a, int b){
   a = find(a);
   b = find(b);
   if (a!=b){
        int sz1 = sz.query version(a, a+1, arr.version.
             size()-1).val;
        int sz2 = sz.query_version(b, b+1, arr.version.
             size()-1).val;
        if (sz1<sz2){
            arr.update version(a, b, arr.version.size()
            sz.update version(b, sz1+sz2, arr.version.
                 size()-1);
        }else{
```

3.3 PBDS GP Hash Table

```
| #include <ext/pb ds/assoc container.hpp>
  using namespace __gnu_pbds;
  typedef tree<int, null type, less<int>, rb tree tag,
       tree order statistics node update> order set;
  struct custom hash {
      static uint64 t splitmix64(uint64 t x) {
          // http://xorshift.di.unimi.it/splitmix64.c
          x += 0x9e3779b97f4a7c15;
          x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
          x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
          return x ^ (x >> 31);
      size t operator()(uint64 t x) const {
          static const uint64_t FIXED_RANDOM = chrono::
               steady clock::now().time since epoch().count();
          return splitmix64(x + FIXED RANDOM);
15
16
17
  };
19 gp hash table (int, int, custom hash) ss;
```

3.4 PBDS Order Set

```
| /*
| .find_by_order(k) 回傳第 k 小的值(based-0)
| .order_of_key(k) 回傳有多少元素比 k 小
| 不能在 #define int Long Long 後 #include 檔案
| */
| *include <ext/pb_ds/assoc_container.hpp>
| #include <ext/pb_ds/tree_policy.hpp>
| using namespace __gnu_pbds;
| typedef tree<int,null_type,less<int>,rb_tree_tag,
| tree_order_statistics_node_update> order_set;
```

3.5 Segment Tree Add Set

```
      1 | // [ll, rr), based-0

      2 | // 使用前記得 init(陣列大小), build(陣列名稱)

      3 | // add(ll, rr): 區間修改

      4 | // set(ll, rr): 區間賦值
```

5 | // query(ll, rr): 區間求和 / 求最大值

```
6 struct SegmentTree{
     struct node{
         int add tag = 0;
         int set tag = 0;
         int sum = 0:
         int ma = 0;
     };
     vector<node> arr;
     SegmentTree(int n){
         arr.resize(n<<2);</pre>
     node pull(node A, node B){
         node C:
         C.sum = A.sum+B.sum;
         C.ma = max(A.ma, B.ma):
         return C;
      // cce0c8
     void push(int idx, int ll, int rr){
         if (arr[idx].set tag!=0){
             arr[idx].sum = (rr-ll)*arr[idx].set_tag;
             arr[idx].ma = arr[idx].set tag;
             if (rr-ll>1){
                  arr[idx*2+1].add_tag = 0;
                  arr[idx*2+1].set tag = arr[idx].set tag;
                  arr[idx*2+2].add_tag = 0;
                  arr[idx*2+2].set_tag = arr[idx].set_tag;
             arr[idx].set_tag = 0;
         if (arr[idx].add_tag!=0){
             arr[idx].sum += (rr-ll)*arr[idx].add_tag;
             arr[idx].ma += arr[idx].add tag;
             if (rr-ll>1){
                  arr[idx*2+1].add_tag += arr[idx].add_tag;
                  arr[idx*2+2].add tag += arr[idx].add tag;
             arr[idx].add tag = 0;
     void build(vector<int> &v, int idx = 0, int ll = 0, int
          rr = n){
         if (rr-11==1){
             arr[idx].sum = v[l1];
             arr[idx].ma = v[ll];
         }else{
             int mid = (11+rr)/2;
             build(v, idx*2+1, l1, mid);
             build(v, idx*2+2, mid, rr);
             arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
     }
     void add(int al, int ar, int val, int idx = 0, int ll =
          0, int rr =n){
         push(idx, 11, rr);
if (rr<=ql || qr<=ll) return;</pre>
         if (ql<=11 && rr<=qr){
             arr[idx].add tag += val:
             push(idx, ll, rr);
```

```
return:
           int mid = (11+rr)/2:
           add(ql, qr, val, idx*2+1, ll, mid);
           add(ql, qr, val, idx*2+2, mid, rr);
           arr[idx]=pull(arr[idx*2+1], arr[idx*2+2]);
       void set(int ql, int qr, int val, int idx=0, int ll=0,
            int rr=n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return;</pre>
           if (q1<=11 && rr<=qr){</pre>
               arr[idx].add tag = 0;
               arr[idx].set_tag = val;
               push(idx, 11, rr);
               return:
           int mid = (11+rr)/2:
           set(ql, qr, val, idx*2+1, ll, mid);
           set(ql, qr, val, idx*2+2, mid, rr);
           arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
       node query(int ql, int qr, int idx = 0, int ll = 0, int
           rr = n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return node();</pre>
           if (q1<=11 && rr<=qr) return arr[idx];</pre>
           int mid = (11+rr)/2;
           return pull(query(ql, qr, idx*2+1, ll, mid), query(ql
                , qr, idx*2+2, mid, rr));
100 } ST;
```

26

27

28

61

3.6 Segment Tree Li Chao Line

```
arr.resize(4*n);
   void update(Node val, int idx = 0, int ll = 0, int rr =
       if (rr-ll==0) return;
       if (rr-ll==1){
           if (val.y(ll) < arr[idx].y(ll)){</pre>
               arr[idx] = val;
           return;
       int mid = (11+rr)/2;
       if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
            的線斜率要比較小
       if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
           update(val, idx*2+1, ll, mid);
       }else{ // 交點在右邊
           swap(arr[idx], val); // 在左子樹中,新線比舊線還
           update(val, idx*2+2, mid, rr);
       return:
   }
   int query(int x, int idx = 0, int ll = 0, int rr = MAX_V)
       if (rr-ll==0) return INF;
       if (rr-ll==1){
           return arr[idx].v(ll):
       int mid = (11+rr)/2:
       if (x<mid){</pre>
           return min(arr[idx].y(x), query(x, idx*2+1, ll,
           return min(arr[idx].y(x), query(x, idx*2+2, mid,
                rr));
};
```

3.7 Segment Tree Li Chao Segment

int a = 0;

```
int b = INF;
                                                                             rr));
                                                                                                                                        now.val = val:
    int y(int x){
                                                                }
                                                                                                                                        return;
        return a*x+b;
                                                          76 };
};
                                                                                                                                    int mid = (11+rr)/2;
vector<Node> arr:
                                                                                                                                    if (pos<mid){</pre>
                                                            3.8 Segment Tree Persistent
LC Segment Tree(int n = 0){
                                                                                                                                        now.lc = node cnt;
    arr.resize(4*n);
                                                                                                                                        node cnt++;
                                                            全部都是 0-based
void update(Node val, int idx = 0, int ll = 0, int rr =
                                                                                                                                        now.rc = node_cnt;
    if (rr-ll==0) return;
                                                                                                                                        node cnt++;
                                                            Persistent Segment Tree st(n+q);
    if (rr-ll<=1){
                                                            st.build(v, 0);
        if (val.y(ll)<arr[idx].y(ll)){</pre>
            arr[idx] = val:
                                                                                                                                    return;
                                                            update_version(pos, val, ver): 對版本 ver 的 pos 位置改成 val 73
        return;
                                                                                                                                }
                                                            query version(ql, qr, ver): 對版本 ver 查詢 [ql, qr) 的區間和 74
    }
                                                          11 clone version(ver): 複製版本 ver 到最新的版本
    int mid = (11+rr)/2;
    if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
                                                            struct Persistent_Segment_Tree{
                                                                int node cnt = 0;
                                                                                                                          78
         的線斜率要比較小
                                                                struct Node{
    if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
                                                                    int lc = -1;
        update(val, idx*2+1, ll, mid);
                                                                    int rc = -1;
                                                                                                                                    auto &now = arr[idx];
    }else{ // 交點在右邊
                                                                    int val = 0;
        swap(arr[idx], val); // 在左子樹中,新線比舊線還
                                                                vector<Node> arr:
        update(val, idx*2+2, mid, rr);
                                                                vector<int> version;
                                                                                                                                    int mid = (11+rr)/2;
    return;
                                                                Persistent Segment Tree(int sz){
}
                                                                    arr.resize(32*sz);
                                                                    version.push back(node cnt++);
// 在 [ql, qr) 加上一條 val 的線段
                                                                    return;
void update segment(Node val, int ql, int qr, int idx =
                                                                                                                                    return ret;
    0, int 11 = 0, int rr = MAX V){
                                                                                                                                }
    if (rr-ll==0) return;
                                                                void pull(Node &c, Node a, Node b){
    if (rr<=ql || qr<=ll) return;
                                                                    c.val = a.val+b.val;
    if (q1<=11 && rr<=qr){</pre>
                                                                                                                          93
        update(val, idx, ll, rr);
        return;
                                                                                                                                void clone version(int ver){
                                                                void build(vector<int> &v. int idx, int ll = 0, int rr =
                                                                    auto &now = arr[idx]:
    int mid = (11+rr)/2:
    update segment(val, ql, qr, idx*2+1, ll, mid);
                                                                                                                                    node cnt++;
                                                                                                                          99
    update segment(val, ql, qr, idx*2+2, mid, rr);
                                                                    if (rr-ll==1){
                                                                                                                         100
    return:
                                                                        now.val = v[11];
                                                                        return;
}
int query(int x, int idx = 0, int ll = 0, int rr = MAX_V)
                                                                                                                            3.9 Sparse Table
                                                                    int mid = (11+rr)/2;
    if (rr-ll==0) return INF:
                                                                    now.lc = node cnt++;
    if (rr-ll==1){
                                                                    now.rc = node cnt++;
        return arr[idx].y(ll);
                                                                    build(v, now.lc, ll, mid);
                                                                                                                          1 struct SparseTable{
                                                                    build(v, now.rc, mid, rr);
                                                                                                                                vector<vector<int>> st;
                                                                    pull(now, arr[now.lc], arr[now.rc]);
                                                                                                                                void build(vector<int> v){
    int mid = (11+rr)/2;
                                                                    return:
                                                                                                                                    int h = __lg(v.size());
    if (x<mid){</pre>
                                                                                                                                    st.resize(h+1):
        return min(arr[idx].y(x), query(x, idx*2+1, ll,
                                                                                                                                    st[0] = v;
                                                                void update(int pos, int val, int idx, int ll = 0, int rr
    }else{
                                                                                                                                    for (int i=1 ; i<=h ; i++){</pre>
                                                                    auto &now = arr[idx];
                                                                                                                                        int gap = (1 << (i-1));
```

return min(arr[idx].y(x), query(x, idx*2+2, mid,

```
if (rr-ll==1){
        arr[node cnt] = arr[now.lc];
        update(pos, val, now.lc, ll, mid);
        arr[node cnt] = arr[now.rc];
        update(pos, val, now.rc, mid, rr);
    pull(now, arr[now.lc], arr[now.rc]);
void update_version(int pos, int val, int ver){
    update(pos, val, version[ver]);
Node query(int ql, int qr, int idx, int ll = 0, int rr =
    if (q1<=11 && rr<=qr) return now;</pre>
    if (rr<=ql || qr<=ll) return Node();</pre>
    pull(ret, query(ql, qr, now.lc, ll, mid), query(ql,
         qr, now.rc, mid, rr));
Node query_version(int ql, int qr, int ver){
    return query(ql, qr, version[ver]);
    version.push back(node cnt);
    arr[node cnt] = arr[version[ver]];
```

```
for (int j=0; j+gap<st[i-1].size(); j++){</pre>
                  st[i].push_back(min(st[i-1][j], st[i-1][j+gap 47
                       1));
      // 回傳 [ll, rr) 的最小值
      int query(int 11, int rr){
          int h = lg(rr-ll);
          return min(st[h][11], st[h][rr-(1<<h)]);</pre>
21 };
  3.10 Treap
1 struct Treap{
      Treap *1 = nullptr, *r = nullptr;
```

```
int pri = rand(), val = 0, sz = 1;
      Treap(int _val){
           val = val;
  };
int size(Treap *t){return t ? t->sz : 0;}
void pull(Treap *t){
      t\rightarrow sz = size(t\rightarrow 1) + size(t\rightarrow r) + 1;
13 }
15 Treap* merge(Treap *a, Treap *b){
      if (!a || !b) return a ? a : b;
      if (a->pri>b->pri){
           a->r = merge(a->r, b);
           pull(a);
           return a;
      }else{
           b \rightarrow 1 = merge(a, b \rightarrow 1);
           pull(b);
           return b;
      }
27
30 | pair<Treap*, Treap*> split(Treap *&t, int k){ // 1-based <前
       k 個元素, 其他元素>
      if (!t) return {};
      if (size(t->1)>=k){
           auto pa = split(t->1, k);
           t->1 = pa.second;
           pull(t);
           return {pa.first, t};
           auto pa = split(t->r, k-size(t->l)-1);
           t->r = pa.first;
           pull(t);
           return {t, pa.second};
45 // functions
```

```
46 Treap* build(vector<int> v){
      Treap* ret = nullptr;
      for (int i=0 ; i<v.size() ; i++){</pre>
          ret = merge(ret, new Treap(v[i]));
      return ret:
  array<Treap*, 3> cut(Treap *t, int 1, int r){ // 1-based <前
       1~L-1 個元素, L~r 個元素, r+1 個元素>
      array<Treap*, 3> ret;
      tie(ret[1], ret[2]) = split(t, r);
      tie(ret[0], ret[1]) = split(ret[1], 1-1);
      return ret;
  void print(Treap *t, bool flag = true){
      if (t->1!=0) print(t->1, false);
      cout << t->val;
      if (t->r!=0) print(t->r, false);
```

3.11 Trie

if (flag) cout << endl;</pre>

```
1 struct Trie{
      struct Data{
         int nxt[2]={0, 0};
     int sz=0:
     vector<Data> arr;
      void init(int n){
         arr.resize(n);
      void insert(int n){
         int now=0;
          for (int i=N ; i>=0 ; i--){
              int v=(n>>i)&1;
              if (!arr[now].nxt[v]){
                  arr[now].nxt[v]=++sz;
              now=arr[now].nxt[v];
     int query(int n){
          int now=0, ret=0;
         for (int i=N ; i>=0 ; i--){
              int v=(n>>i)&1;
              if (arr[now].nxt[1-v]){
                  ret+=(1<<i);
                  now=arr[now].nxt[1-v];
              }else if (arr[now].nxt[v]){
                  now=arr[now].nxt[v];
              }else{
                  return ret;
         return ret;
```

```
40 } tr;
```

17

31

4 Dynamic-Programming

4.1 Digit DP

```
| #include <bits/stdc++.h>
  using namespace std;
 long long l, r;
s | long long dp[20][10][2][2]; // dp[pos][pre][limit] = 後 pos
      位·pos 前一位是 pre·(是/否)有上界·(是/否)有前綴零
  long long memorize_search(string &s, int pos, int pre, bool
      limit, bool lead){
     // 已經被找過了,直接回傳值
     if (dp[pos][pre][limit][lead]!=-1) return dp[pos][pre][
         limit][lead];
     // 已經搜尋完畢,紀錄答案並回傳
     if (pos==(int)s.size()){
         return dp[pos][pre][limit][lead] = 1;
     // 枚舉目前的位數數字是多少
     long long ans = 0;
     for (int now=0 ; now<=(limit ? s[pos]-'0' : 9) ; now++){</pre>
         if (now==pre){
             // 1~9 絕對不能連續出現
             if (pre!=0) continue;
             // 如果已經不在前綴零的範圍內·@ 不能連續出現
             if (lead==false) continue;
         ans += memorize search(s, pos+1, now, limit&(now==(s[
             pos]-'0')), lead&(now==0));
     }
     // 已經搜尋完畢,紀錄答案並回傳
33
     return dp[pos][pre][limit][lead] = ans;
34
  // 回傳 [0, n] 有多少數字符合條件
  long long find_answer(long long n){
     memset(dp, -1, sizeof(dp));
     string tmp = to string(n);
     return memorize_search(tmp, 0, 0, true, true);
42
  int main(){
     // input
```

```
cin >> 1 >> r;

// output - 計算 [l, r] 有多少數字任意兩個位數都不相同
cout << find_answer(r)-find_answer(l-1) << "\n";
return 0;
3 }
```

4.2 Knaspack On Tree

1 // 需要重構、需要增加註解

```
2 #include <bits/stdc++.h>
3 #define F first
  #define S second
  #define all(x) begin(x), end(x)
  using namespace std;
  #define chmax(a, b) (a) = (a) < (b) ? (b) : (a)
  #define chmin(a, b) (a) = (a) < (b) ? (a) : (b)
  #define ll long long
  #define FOR(i, a, b) for (int i = a; i <= b; i++)</pre>
15 int N, W, cur;
  vector<int> w, v, sz;
  vector<vector<int>> adj, dp;
  void dfs(int x) {
      sz[x] = 1;
      for (int i : adj[x]) dfs(i), sz[x] += sz[i];
      cur++;
      // choose x
      for (int i=w[x]; i<=W; i++){</pre>
          dp[cur][i] = dp[cur - 1][i - w[x]] + v[x];
      // not choose x
      for (int i=0 ; i<=W ; i++){</pre>
          chmax(dp[cur][i], dp[cur - sz[x]][i]);
33 signed main() {
      cin >> N >> W:
      adj.resize(N + 1);
      w.assign(N + 1, 0);
      v.assign(N + 1, 0);
      sz.assign(N + 1, 0);
      dp.assign(N + 2, vector<int>(W + 1, 0));
      for (int i=1 ; i<=N ; i++){</pre>
          int p; cin >> p;
          adj[p].push back(i);
      for (int i=1; i<=N; i++) cin >> w[i];
      for (int i=1; i<=N; i++) cin >> v[i];
      cout << dp[N + 1][W] << ' \ '';
```

4.3 SOS DP

```
1 // 總時間複雜度為 O(n 2^n)
2 // 計算 dp[i] = i 所有 bit mask 子集的和
3 for (int i=0; i<n; i++){
4 for (int mask=0; mask<(1<<n); mask++){
5 if ((mask>>i)&1){
6 dp[mask] += dp[mask^(1<<ii)];
7 }
8 }
9 }
```

4.4 Integer Partition

```
dp[i][x] = 要將整數 x 拆成 i 堆的「組合數」 dp[i+1][x+1] + = dp[i][x] ( 創造新的一堆 ) \\ dp[i][x+i] + = dp[i][x] ( 把每一堆都增加 1 )
```

5 Geometry

5.1 Geometry Struct

```
1 // 判斷數值正負: {1:正數,0:零,-1:負數}
 int sign(long long x) {return (x \ge 0) ? ((bool)x) : -1; }
 int sign(double x) {
     return (abs(x) < 1e-9) ? 0 : (x > 0 ? 1 : -1);
 template<tvpename T>
 struct point {
     T x, y;
     point() {}
     point(const T &x, const T &y) : x(x), y(y) {}
     point operator+(point b) {return {x+b.x, y+b.y}; }
     point operator-(point b) {return {x-b.x, y-b.y}; }
     point operator*(T b) {return {x*b, y*b}; }
     point operator/(T b) {return {x/b, y/b}; }
     bool operator==(point b) {return x==b.x && y==b.y; }
     // 逆時針極角排序
     bool operator<(point &b) {return (x*b.y > b.x*y); }
     friend ostream& operator<<(ostream& os, point p) {</pre>
         os << "(" << p.x << ", " << p.y << ")";
     // 判斷 ab 到 ac 的方向: {1:逆時鐘,0:重疊,-1:順時鐘}
     friend int ori(point a, point b, point c) {
         return sign((b-a)^(c-a));
     friend bool btw(point a, point b, point c) {
         return ori(a, b, c) == 0 && sign((a-c)*(b-c)) <= 0;
     // 判斷線段 ab. cd 是否相交
     friend bool banana(point a, point b, point c, point d) {
         int s1 = ori(a, b, c);
         int s2 = ori(a, b, d);
         int s3 = ori(c, d, a);
```

```
int s4 = ori(c, d, b);
          if (btw(a, b, c) || btw(a, b, d) || btw(c, d, a) ||
              btw(c, d, b)) return 1;
          return (s1 * s2 < 0) && (s3 * s4 < 0);
      T operator*(point b) {return x * b.x + y * b.y; }
      T operator^(point b) {return x * b.y - y * b.x; }
      T abs2() {return (*this) * (*this); }
      // 旋轉 Arg(b) 的角度(小心溢位)
      point rotate(point b) {return {x*b.x - y*b.y, x*b.y + y*b
47
  template<typename T>
  struct line {
      point<T> p1, p2;
      // ax + by + c = 0
      T a, b, c; //|a|, |b| \le 2C, |c| \le 8C^2
      line(const point<T> &x,const point<T> &y) : p1(x), p2(y){
          build();
      void build() {
      a = p1.y - p2.y;
      b = p2.x - p1.x:
      c = (-a*p1.x)-b*p1.y;
      // 判斷點和有向直線的關係: {1:左邊,0:在線上,-1:右邊}
    int ori(point<T> &p) {
      return sign((p2-p1) ^ (p-p1));
    // 判斷直線斜率是否相同
    bool parallel(line &1) {
      return ((p1-p2) ^ (l.p1-l.p2)) == 0;
    // 兩直線交點
      point<long double> line intersection(line &1) {
          using P = point<long double>;
      point < T > a = p2-p1, b = 1.p2-1.p1, s = 1.p1-p1;
      return P(p1.x,p1.y) + P(a.x,a.y) * (((long double)(s^b))
77
  };
  template<typename T>
80 struct polygon {
      vector<point<T>> v;
      polygon() {}
      polygon(const vector<point<T>> &u) : v(u) {}
      // simple 為 true 的時候會回傳任意三點不共線的凸包
      void make convex hull(int simple) {
          auto cmp = [&](point<T> &p, point<T> &q) {
              return (p.x == q.x)? (p.y < q.y): (p.x < q.x);
          simple = (bool)simple;
          sort(v.begin(), v.end(), cmp);
          v.resize(unique(v.begin(), v.end()) - v.begin());
          vector<point<T>> hull:
          for (int t = 0; t < 2; ++t){
              int sz = hull.size();
              for (auto &i:v) {
                  while (hull.size() >= sz+2 && ori(hull[hull.
                      size()-2], hull.back(), i) < simple) {</pre>
```

```
hull.pop back();
                  hull.push back(i):
101
              hull.pop_back();
102
              reverse(v.begin(), v.end());
103
104
          swap(hull, v);
105
      可以在有 n 個點的簡單多邊形內,用 O(n)判斷一個點:
      {1: 在多邊形內,0:在多邊形上,-1:在多邊形外}
      int in polygon(point<T> a){
108
           const T MAX POS = 1e9 + 5; // [記得修改] 座標的最大值
          point<T> pre = v.back(), b(MAX POS, a.y + 1);
110
                                                               172
111
          int cnt = 0;
                                                              173
112
                                                               174
          for (auto &i:v) {
113
                                                              175
114
              if (btw(pre, i, a)) return 0;
                                                               176
115
              if (banana(a, b, pre, i)) cnt++;
116
              pre = i;
                                                               178
117
118
          return cnt%2 ? 1 : -1;
119
121 | /// 警告:以下所有凸包專用的函式都只接受逆時針排序且任三點不
       共線的凸包 ///
                                                               184
      可以在有 n 個點的凸包內,用 O(\log n) 判斷一個點:
                                                               185
123 // {1:在凸包內, 0:在凸包邊上, -1:在凸包外}
                                                              186
      int in_convex(point<T> p) {
                                                               187
125
          int n = v.size();
          int a = ori(v[0], v[1], p), b = ori(v[0], v[n-1], p); 189
126
          if (a < 0 || b > 0) return -1;
127
128
          if (btw(v[0], v[1], p)) return 0;
                                                               191
          if (btw(v[0], v[n - 1], p)) return 0;
129
130
          int l = 1, r = n - 1, mid;
          while (1 + 1 < r) {
131
132
              mid = (1 + r) >> 1;
              if (ori(v[0], v[mid], p) >= 0) 1 = mid;
133
134
              else r = mid:
135
           int k = ori(v[1], v[r], p);
136
137
          if (k <= 0) return k;</pre>
          return 1;
138
139
   // 凸包專用的環狀二分搜·回傳 0-based index
      int cycle search(auto &f) {
141
142
          int n = v.size(), l = 0, r = n;
          bool rv = f(1, 0);
143
144
          while (r - 1 > 1) {
                                                              204
              int m = (1 + r) / 2;
145
                                                               205
146
              if (f(0, m) ? rv: f(m, (m + 1) % n)) r = m;
                                                              206
147
              else 1 = m:
                                                              207
148
                                                               208
149
          return f(1, r % n) ? 1 : r % n;
                                                               209
                                                              210
151 // 可以在有 n 個點的凸包內 · 用 O(Log n) 判斷一條直線:
                                                              211
      {1: 穿過凸包,0: 剛好切過凸包,-1: 沒碰到凸包}
                                                              212
      int line cut convex(line<T> L) {
                                                              213
          point<T> p(L.a, L.b); // 記得 L 要 build
154
                                                              214
          auto gt = [&](int neg) {
155
                                                              215
156
              auto f = [&](int x, int y) {
                                                              216
157
                  return sign((v[x] - v[y]) * p) == neg;
                                                              217
                                                              218
              return -(v[cycle search(f)] * p);
```

```
220
           T x = gt(1), y = gt(-1);
                                                                221
           if (L.c < x || y < L.c) return -1;
           return not (L.c == x || L.c == y);
                                                                223
                                                                224
                                                                225
   // 可以在有 n 個點的凸包內 · 用 O(Log n) 判斷一個線段:
                                                                226
166 // {1: 存在一個凸包上的邊可以把這個線段切成兩半,
                                                                227
       0: 有碰到凸包但沒有任何凸包上的邊可以把它切成兩半,
168 // -1: 沒碰到凸包}
169 1// 除非線段兩端點都不在凸包邊上,否則此函數回傳 8 的時候不一
                                                                229
       定表示線段沒有诵過凸包內部 ///
                                                                230
       int segment_across_convex(line<T> L) {
           point<T> p(L.a, L.b); // 記得 L 要 build
                                                                231
           auto gt = [&](int neg) {
                                                                232
               auto f = [&](int x, int y) {
                  return sign((v[x] - v[y]) * p) == neg;
                                                                233
                                                                234
               return cycle search(f);
                                                                235
                                                                236
           int i = gt(1), j = gt(-1), n = v.size();
                                                                237
          T x = -(v[i] * p), y = -(v[j] * p);
                                                                238
          if (L.c < x || y < L.c) return -1;
if (L.c == x || L.c == y) return 0;
                                                                239
                                                                240
           if (i > j) swap(i, j);
                                                                241
           auto g = [&](int x, int lim) {
                                                                242
               int now = 0, nxt;
                                                                243
               for (int i = 1 \leftrightarrow lg(lim); i > 0; i /= 2) {
                                                                244
                   if (now + i > lim) continue;
                                                                245
                   nxt = (x + i) % n;
                                                                246
                   if (L.ori(v[x]) * L.ori(v[nxt]) >= 0) {
                                                                247
                      x = nxt;
                                                                248
                      now += i;
                                                                249
                                                                250
               } // ↓ BE CAREFUL
               return -(ori(v[x], v[(x + 1) % n], L.p1) * ori(v[ 252
                   x], v[(x + 1) % n], L.p2));
                                                                254
           return max(g(i, j - i), g(j, n - (j - i)));
                                                                255
                                                                256
      可以在有 n 個點的凸包內,用 O(Log n)判斷一個線段:
                                                                257
       {1: 線段上存在某一點位於凸包內部(邊上不算).
                                                                258
       0: 線段上存在某一點碰到凸包的邊但線段上任一點均不在凸包
       內部.
                                                                260
201 // -1: 線段完全在凸包外面 }
                                                                261
       int segment pass convex interior(line<T> L) {
                                                                262
           if (in convex(L.p1) == 1 || in convex(L.p2) == 1)
                                                                263
           point<T> p(L.a, L.b); // 記得 L 要 build
                                                                264
           auto gt = [&](int neg) {
                                                                265
               auto f = [\&](int x, int y) {
                   return sign((v[x] - v[y]) * p) == neg;
                                                                266
                                                                267
               return cycle search(f);
                                                                268
                                                                260
           int i = gt(1), j = gt(-1), n = v.size();
                                                                270
          T x = -(v[i] * p), y = -(v[j] * p);
if (L.c < x || y < L.c) return -1;
                                                                271
                                                                272
           if (L.c == x || L.c == y) return 0;
                                                                273
                                                                274
           if (i > j) swap(i, j);
                                                                275
           auto g = [&](int x, int lim) {
                                                                276
               int now = 0. nxt;
               for (int i = 1 << __lg(lim); i > 0; i /= 2) {
```

```
if (now + i > lim) continue;
               nxt = (x + i) % n;
               if (L.ori(v[x]) * L.ori(v[nxt]) > 0) {
                    x = nxt;
                    now += i;
           } // ↓ BE CAREFUL
           return -(ori(v[x], v[(x + 1) % n], L.p1) * ori(v[
                x], v[(x + 1) % n], L.p2));
       int ret = max(g(i, j - i), g(j, n - (j - i)));
       return (ret == 0) ? (in_convex(L.p1) == 0 &&
            in_convex(L.p2) == 0) : ret;
// 回傳點過凸包的兩條切線的切點的 0-based index (不保證兩條
    切線的順逆時針關係)
   pair<int,int> convex tangent point(point<T> p) {
       int n = v.size(), z = -1, edg = -1;
       auto gt = [&](int neg) {
           auto check = [&](int x) {
               if (v[x] == p) z = x;
               if (btw(v[x], v[(x + 1) % n], p)) edg = x;
               if (btw(v[(x + n - 1) \% n], v[x], p)) edg = (
                    x + n - 1) % n;
           };
            auto f = [\&](int x, int y) {
               check(x); check(y);
               return ori(p, v[x], v[y]) == neg;
            return cycle_search(f);
       int x = gt(1), y = gt(-1);
       if (z != -1) {
            return \{(z + n - 1) \% n, (z + 1) \% n\};
       else if (edg != -1) {
            return {edg, (edg + 1) % n};
       else {
            return {x, y};
   friend int halfplane intersection(vector<line<T>> &s,
        polygon<T> &P) {
       #define neg(p) ((p.y == 0 ? p.x : p.y) < 0)
       auto angle_cmp = [&](line<T> &A, line<T> &B) {
            point < T > a = A.p2-A.p1, b = B.p2-B.p1;
            return neg(a) < neg(b) \mid \mid (neg(a) == neg(b) \&\& (a)
                ^b) > 0);
       };
       sort(s.begin(), s.end(), angle cmp); // 線段左側為該
            線段半平面
       int L, R, n = s.size();
       vector<point<T>> px(n);
       vector<line<T>> q(n);
       q[L = R = 0] = s[0];
       for(int i = 1; i < n; ++i) {</pre>
           while(L < R && s[i].ori(px[R-1]) <= 0) --R;</pre>
           while(L < R && s[i].ori(px[L]) <= 0) ++L;</pre>
           q[++R] = s[i];
           if(q[R].parallel(q[R-1])) {
               if(q[R].ori(s[i].p1) > 0) q[R] = s[i];
```

```
if(L < R) px[R-1] = q[R-1].line_intersection(q[R</pre>
279
            while(L < R && q[L].ori(px[R-1]) <= 0) --R;
280
281
            P.v.clear();
            if(R - L <= 1) return 0;
282
            px[R] = q[R].line_intersection(q[L]);
283
            for(int i = L; i <= R; ++i) P.v.push back(px[i]);</pre>
284
285
            return R - L + 1;
286
287 };
```

5.2 Pick's Theorem

給定頂點坐標均是整點的簡單多邊形,面積 = 內部格點數 + 邊上格點數/2 - 1

6 Graph

6.1 2-SAT

```
1 struct TWO_SAT {
     int n, N;
     vector<vector<int>> G, rev_G;
     deque<bool> used;
     vector<int> order, comp;
     deque<bool> assignment;
     void init(int n) {
         n = _n;
         N = n * 2;
         G.resize(N + 5);
         rev_G.resize(N + 5);
     void dfs1(int v) {
         used[v] = true;
         for (int u : G[v]) {
             if (!used[u])
                 dfs1(u);
         order.push back(v);
     void dfs2(int v, int cl) {
         comp[v] = cl;
         for (int u : rev_G[v]) {
             if (comp[u] == -1)
                 dfs2(u, c1);
     bool solve() {
         order.clear();
         used.assign(N, false);
         for (int i = 0; i < N; ++i) {
             if (!used[i])
                 dfs1(i);
         comp.assign(N, -1);
         for (int i = 0, j = 0; i < N; ++i) {
             int v = order[N - i - 1];
             if (comp[v] == -1)
                 dfs2(v, j++);
```

```
for (int i = 0; i < N; i += 2) {
               if (comp[i] == comp[i + 1])
                  return false:
                                                                  35
               assignment[i / 2] = (comp[i] > comp[i + 1]);
                                                                  36
                                                                  37
          return true:
      void add_disjunction(int a, bool na, int b, bool nb) { //
            A or B 都是 0-based
                                                                  42
          // na means whether a is negative or not
          // nb means whether b is negative or not
                                                                  43
          a = 2 * a ^ na;
          b = 2 * b ^ nb;
          int neg_a = a ^ 1;
           int neg_b = b ^ 1;
          G[neg a].push back(b);
          G[neg_b].push_back(a);
          rev_G[b].push_back(neg_a);
           rev G[a].push back(neg b);
          return:
      void get result(vector<int>& res) {
           res.clear();
           for (int i = 0; i < n; i++)
               res.push_back(assignment[i]);
67 };
```

assignment.assign(n, false);

6.2 Augment Path

```
1 struct AugmentPath{
      int n, m;
      vector<vector<int>> G;
      vector<int> mx, my;
      vector<int> visx, visy;
      AugmentPath(int _n, int _m) : n(_n), m(_m), G(n), mx(n,
           -1), my(m, -1), visx(n), visy(n){
          stamp = 0;
      }
      void add(int x, int y){
          G[x].push_back(y);
      // bb03e2
      bool dfs1(int now){
          visx[now] = stamp;
          for (auto x : G[now]){
              if (my[x]==-1){
                  mx[now] = x;
                  my[x] = now;
                  return true;
          for (auto x : G[now]){
              if (visx[my[x]]!=stamp && dfs1(my[x])){
                  mx[now] = x;
                  my[x] = now;
```

```
return false:
      }
      vector<pair<int, int>> find_max_matching(){
           vector<pair<int, int>> ret;
           while (true){
               stamp++:
               int tmp = 0;
               for (int i=0 ; i<n ; i++){</pre>
                   if (mx[i]==-1 && dfs1(i)) tmp++;
               if (tmp==0) break;
           for (int i=0 ; i<n ; i++){</pre>
               if (mx[i]!=-1){
                   ret.push_back({i, mx[i]});
           return ret;
      }
      // 645577
      void dfs2(int now){
          visx[now] = true;
           for (auto x : G[now]){
               if (my[x]!=-1 && visy[x]==false){
                   visy[x] = true;
                   dfs2(my[x]);
      // 要先執行 find_max_matching 一次
70
      vector<pair<int, int>> find min vertex cover(){
71
           fill(visx.begin(), visx.end(), false);
           fill(visy.begin(), visy.end(), false);
72
73
74
           vector<pair<int, int>> ret;
75
           for (int i=0 ; i<n ; i++){</pre>
               if (mx[i]==-1) dfs2(i);
76
79
           for (int i=0 ; i<n ; i++){</pre>
80
               if (visx[i]==false) ret.push_back({1, i});
81
           for (int i=0 ; i<m ; i++){</pre>
82
               if (visy[i]==true) ret.push back({2, i});
83
84
85
86
           return ret;
87
88 };
```

return true;

6.3 Bridge BCC

```
#include <bits/stdc++.h>
mathrmsday
using namespace std;
```

```
4 const int N = 200005;
5 vector <int> G[N]:
6 int low[N], depth[N];
 bool vis[N];
 vector <vector <int>> bcc;
 stack <int> stk;
 void dfs(int v, int p) {
     stk.push(v);
     vis[v] = true;
     low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
     for (int u : G[v]) {
         if (u == p) continue;
         if (!vis[u]) {
            /// (v, u) 是樹邊
             dfs(u, v);
             low[v] = min(low[v], low[u]);
             /// (v, u) 是回邊
             low[v] = min(low[v], depth[u]);
     /// v 在不依靠父邊的情況下永遠沒辦法走到它的祖先
     if (low[v] == depth[v]) {
         bcc.emplace_back();
         while (stk.top() != v) {
             bcc.back().push_back(stk.top());
             stk.pop();
         bcc.back().push_back(stk.top());
         stk.pop();
```

6.4 Cut BCC

```
i #include <bits/stdc++.h>
 using namespace std;
 const int N = 200005;
 vector <int> G[N];
 int low[N], depth[N];
 bool vis[N];
 vector <vector <int>> bcc;
 stack <int> stk;
 void dfs(int v, int p) {
     stk.push(v);
     vis[v] = true;
     low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
     for (int u : G[v]) {
         if (u == p) continue;
         if (!vis[u]) {
            /// (v, u) 是樹邊
             dfs(u, v);
             low[v] = min(low[v], low[u]);
            /// u 無法在不經過父邊的情況走到 v 的祖先
            if (low[u] >= depth[v]) {
                 bcc.emplace back();
                 while (stk.top() != u) {
                    bcc.back().push_back(stk.top());
```

```
stk.pop();
}
bcc.back().push_back(stk.top());
stk.pop();
bcc.back().push_back(v);
}
}
else {
    /// (v, u) 是回邊
    low[v] = min(low[v], depth[u]);
}
}
}
```

6.5 Dinic

1 // 一般圖: O(EV2)

```
2 // 二分圖: O(E√V)
3 struct Flow{
      using T = int; // 可以換成別的型別
      struct Edge{
          int v; T rc; int rid;
      vector<vector<Edge>> G;
      void add(int u, int v, T c){
          G[u].push back({v, c, G[v].size()});
          G[v].push_back({u, 0, G[u].size()-1});
      vector<int> dis, it;
      Flow(int n){
         G.resize(n);
          dis.resize(n);
          it.resize(n);
      // ce56d6
     T dfs(int u, int t, T f){
          if (u == t || f == 0) return f;
          for (int &i=it[u] ; i<G[u].size() ; i++){</pre>
              auto &[v, rc, rid] = G[u][i];
              if (dis[v]!=dis[u]+1) continue;
              T df = dfs(v, t, min(f, rc));
              if (df <= 0) continue;</pre>
              rc -= df:
              G[v][rid].rc += df;
              return df;
          return 0;
      // e22e39
     T flow(int s, int t){
          T ans = 0;
          while (true){
              fill(dis.begin(), dis.end(), INF);
              queue<int> q;
              q.push(s);
              dis[s] = 0;
              while (q.size()){
                  int u = q.front(); q.pop();
                  for (auto [v, rc, rid] : G[u]){
```

```
if (rc <= 0 || dis[v] < INF) continue;</pre>
                       dis[v] = dis[u] + 1;
                       q.push(v);
              if (dis[t]==INF) break;
              fill(it.begin(), it.end(), 0);
              while (true){
                   T df = dfs(s, t, INF);
                   if (df <= 0) break;</pre>
                   ans += df;
          return ans;
      // the code below constructs minimum cut
      void dfs mincut(int now, vector<bool> &vis){
          vis[now] = true;
          for (auto &[v, rc, rid] : G[now]){
               if (vis[v] == false && rc > 0){
                   dfs_mincut(v, vis);
      }
      vector<pair<int, int>> construct(int n, int s, vector<</pre>
           pair<int,int>> &E){
          // E is G without capacity
          vector<bool> vis(n);
          dfs_mincut(s, vis);
          vector<pair<int, int>> ret;
          for (auto &[u, v] : E){
               if (vis[u] == true && vis[v] == false){
                   ret.emplace back(u, v);
85
           return ret;
86
```

6.6 Dominator Tree

```
1 /*
2|全部都是 0-based
3 G 要是有向無權圖
  一開始要初始化 G(N, root)、代表有 N 個節點、根是 root
  用完之後要 build
 G[i] = i 的 idom · 也就是從 root 走到 i 時 · 一定要走到的點且離
      i 最近
 struct DominatorTree{
     int N;
     vector<vector<int>> G;
     vector<vector<int>> buckets, rg;
     // dfn[x] = the DFS otder of x
     // rev[x] = the vertex with DFS order x
     // par[x] = the parent of x
     vector<int> dfn, rev, par;
     vector<int> sdom, dom, idom;
     vector<int> fa, val;
```

```
int stamp;
int root;
int operator [] (int x){
    return idom[x];
DominatorTree(int N, int root) :
    G(N), buckets(N), rg(N),
    dfn(N, -1), rev(N, -1), par(N, -1),
    sdom(N, -1), dom(N, -1), idom(N, -1),
    fa(N, -1), val(N, -1)
    stamp = 0;
    root = root;
void add edge(int u, int v){
    G[u].push_back(v);
void dfs(int x){
    rev[dfn[x] = stamp] = x;
    fa[stamp] = sdom[stamp] = val[stamp] = stamp;
    stamp++;
    for (int u : G[x]){
        if (dfn[u]==-1){
            dfs(u);
            par[dfn[u]] = dfn[x];
        rg[dfn[u]].push_back(dfn[x]);
}
int eval(int x, bool first){
    if (fa[x]==x) return !first ? -1 : x;
    int p = eval(fa[x], false);
    if (p==-1) return x;
    if (sdom[val[x]]>sdom[val[fa[x]]]) val[x] = val[fa[x
    fa[x] = p;
    return !first ? p : val[x];
void link(int x, int y){
    fa[x] = y;
void build(){
    dfs(root);
    for (int x=stamp-1 ; x>=0 ; x--){
        for (int y : rg[x]){
            sdom[x] = min(sdom[x], sdom[eval(y, true)]);
        if (x>0) buckets[sdom[x]].push_back(x);
        for (int u : buckets[x]){
            int p = eval(u, true);
            if (sdom[p]==x) dom[u] = x;
            else dom[u] = p;
        if (x>0) link(x, par[x]);
```

6.7 Enumerate Triangle

```
1 // O(m Log m) 枚舉無向圖所有三角形 · 0-based
void Enumerate_Triangle(vector<pair<int, int>> &edge, vector<</pre>
      int> &deg){
      int n = deg.size();
      int m = edge.size();
      vector<vector<int>> G(n):
      for (int i=0 ; i<m ; i++){</pre>
          if (deg[edge[i].first] > deg[edge[i].second]) swap(
               edge[i].first, edge[i].second);
          if (deg[edge[i].first] == deg[edge[i].second] && edge
               [i].first > edge[i].second) swap(edge[i].first,
               edge[i].second);
          G[edge[i].first].push_back(edge[i].second);
      vector<int> vis(n, false);
      for (int i=0 ; i<n ; i++){</pre>
          for (auto j : G[i]) vis[j] = true;
          for (auto j : G[i]){
              for (auto k : G[j]){
                  if (vis[k]){
                      // i, j, k is a triangle
          for (auto j : G[i]) vis[j] = false;
```

6.8 Find Bridge

```
| vector<int> dep(MAX_N), low(MAX_N);
| vector<pair<int, int>> bridge;
| bitset<MAX_N> vis;
| void dfs(int now, int pre){
| vis[now] = 1;
| low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
| for (auto x : G[now]){
| if (x==pre){
| continue;
| low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
| void dfs(int now, int pre){
| vis[now] = 1;
| low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
| for (auto x : G[now]){
| if (x==pre){
| continue;
| low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
| vector<int>| dep(nax_N), low(MAX_N);
| vector<int>| dep(nax_N), low(MAX_N);
| vector<int>| vector<int>| dep(nax_N), low(MAX_N);
| vector<int>| vector<int>| dep(nax_N), low(MAX_N);
| vector<int>| dep(nax_N), low(MAX_N);
| vector<int>| dep(nax_N), low(MAX_N);
| vector<int>| dep(nax_N), low(MAX_N);
| low(nax_N);
| vector<int>| vector<int>| dep(nax_N), low(MAX_N);
| low(nax_N);
| vector<int>| vector<int>| dep(nax_N), low(MAX_N);
| low(nax_N);
| vector<int>| vector<int>| dep(nax_N);
| low(nax_N);
| low(
```

6.9 HLD

return;

}

```
| #include <bits/stdc++.h>
  #define int long long
  using namespace std;
  const int N = 100005:
  vector <int> G[N];
  struct HLD {
      vector<int> pa, sz, depth, mxson, topf, id;
      int n, idcnt = 0;
      HLD(int _n) : n(_n), pa(_n + 1), sz(_n + 1), depth(_n +
           1), mxson((n + 1), topf((n + 1), id((n + 1)))
      void dfs1(int v = 1, int p = -1) {
          pa[v] = p; sz[v] = 1; mxson[v] = 0;
          depth[v] = (p == -1 ? 0 : depth[p] + 1);
for (int u : G[v]) {
               if (u == p) continue;
              dfs1(u, v);
               sz[v] += sz[u];
               if (sz[u] > sz[mxson[v]]) mxson[v] = u;
19
20
21
      void dfs2(int v = 1, int top = 1) {
          id[v] = ++idcnt;
          topf[v] = top;
          if (mxson[v]) dfs2(mxson[v], top);
          for (int u : G[v]) {
               if (u == mxson[v] || u == pa[v]) continue;
              dfs2(u, u);
      // query 為區間資料結構
      int path_query(int a, int b) {
          int res = 0:
          while (topf[a] != topf[b]) { /// 若不在同一條鍊上
               if (depth[topf[a]] < depth[topf[b]]) swap(a, b);</pre>
               res = max(res, 011); // query : L = id[topf[a]],
                   r = id[a]
               a = pa[topf[a]];
37
          /// 此時已在同一條鍊上
          if (depth[a] < depth[b]) swap(a, b);</pre>
39
          res = \max(\text{res}, 011); // query : l = id[b], r = id[a]
          return res;
41
42
43 };
```

low[now] = min(low[now], low[x]);

low[now] = min(low[now], dep[x]);

}else if (vis[x]==1){

if (now!=1 && low[now]==dep[now]){

bridge.push back({now, pre});

6.10 Kosaraju

return true;

```
}
                                                                           }
2 給定一個有向圖, 迴回傳縮點後的圖、SCC 的資訊
                                                                           result.resize(SCC.size());
                                                                                                                                             return false;
3 所有點都以 based-0 編號
                                                                           sz = SCC.size();
                                                                           for (auto [u, v] : edges){
                                                                               if (SCC_id[u]!=SCC_id[v]) result[SCC_id[u]].
                                                                                                                                         void relabel(){
5| 函式:
                                                                                    push back(SCC id[v]);
                                                                                                                                             int delta = INF:
6 | SCC_compress G(n): 宣告─個有 n 個點的圖
                                                                                                                                             for (int j=0 ; j<n ; j++){</pre>
ī . add edge(u, v): 加上一條邊 u -> v
                                                                           for (int i=0 ; i<SCC.size() ; i++){</pre>
                                                                                                                                                 if (visy[j]!=stamp) delta = min(delta, slack[j]);
8|.compress: O(n Log n) 計算 G3、SCC、SCC_id 的資訊,並把縮點後
                                                                               sort(result[i].begin(), result[i].end());
       的結果存在 result 裡
                                                                               result[i].resize(unique(result[i].begin(), result 55
                                                                                                                                             for (int i=0 ; i<n ; i++){</pre>
                                                                                    [i].end())-result[i].begin());
                                                                                                                                                 if (visx[i]==stamp) lx[i] -= delta;
10| SCC[i] = 某個 SCC 中的所有點
                                                                                                                                             for (int j=0 ; j<n ; j++){</pre>
| II | SCC id [i] = 第 i 個點在第幾個 SCC
                                                                                                                                                 if (visy[j]==stamp) ly[j] += delta;
                                                                   };
                                                                                                                                                 else slack[j] -= delta;
13 struct SCC compress{
     int N, M, sz;
                                                                                                                                         }
     vector<vector<int>>> G, inv_G, result;
                                                                   6.11 Kuhn Munkres
     vector<pair<int, int>> edges;
                                                                                                                                         int solve(){
     vector<bool> vis;
     vector<int> order;
                                                                 1 // O(n^2) 找到最大權匹配
                                                                                                                                             for (int i=0 ; i<n ; i++){</pre>
                                                                   struct KuhnMunkres{
      vector<vector<int>> SCC;
                                                                                                                                                 lx[i] = 0;
                                                                       int n; // max(n, m)
     vector<int> SCC id;
                                                                                                                                                 for (int j=0 ; j<n ; j++){</pre>
                                                                       vector<vector<int>> G:
                                                                                                                                                     lx[i] = max(lx[i], G[i][j]);
                                                                       vector<int> match, lx, ly, visx, visy;
      SCC_compress(int _N) :
                                                                       vector<int> slack;
     N(N), M(0), sz(0),
                                                                       int stamp = 0;
     G(N), inv_G(N),
      vis(N), SCC_id(N)
                                                                                                                                             fill(ly.begin(), ly.end(), 0);
                                                                                                                                             fill(match.begin(), match.end(), -1);
                                                                       KuhnMunkres(int n) : n(n), G(n, vector<int>(n)), lx(n),
      {}
                                                                            ly(n), slack(n), match(n), visx(n), visy(n) {}
      vector<int> operator [] (int x){
                                                                                                                                             for(int i = 0; i < n; i++) {</pre>
          return result[x];
                                                                       void add(int x, int y, int w){
                                                                                                                                                 fill(slack.begin(), slack.end(), INF);
                                                                           G[x][y] = max(G[x][y], w);
                                                                                                                                                 stamp++:
                                                                                                                                                 if(dfs(i, true)) continue;
      void add_edge(int u, int v){
                                                                       bool dfs(int i, bool aug){ // aug = true 表示要更新 match
                                                                                                                                                 while(augment()==false) relabel();
          G[u].push back(v);
                                                                           if (visx[i]==stamp) return false;
                                                                                                                                                 stamp++;
          inv G[v].push back(u);
                                                                                                                                                 dfs(i, true);
          edges.push_back({u, v});
                                                                           visx[i] = stamp;
                                                                           for (int j=0 ; j<n ; j++){</pre>
                                                                                                                                             int ans = 0:
                                                                               if (visy[j]==stamp) continue;
                                                                                                                                             for (int j=0 ; j<n ; j++){</pre>
      void dfs1(vector<vector<int>> &G, int now){
                                                                               int d = lx[i]+ly[j]-G[i][j];
                                                                                                                                                 if (match[j]!=-1){
          vis[now] = 1:
                                                                                                                                  88
                                                                                                                                                     ans += G[match[j]][j];
          for (auto x : G[now]) if (!vis[x]) dfs1(G, x);
                                                                               if (d==0){
          order.push back(now);
                                                                                   visy[j] = stamp;
                                                                                   if (match[j]==-1 || dfs(match[j], aug)){
                                                                                       if (aug){
                                                                                                                                             return ans;
      void dfs2(vector<vector<int>> &G, int now){
                                                                                           match[j] = i;
                                                                                                                                  93
          SCC id[now] = SCC.size()-1;
                                                                                                                                  94 };
          SCC.back().push back(now);
                                                                                       return true:
          vis[now] = 1;
          for (auto x : G[now]) if (!vis[x]) dfs2(G, x);
                                                                               }else{
                                                                                                                                    6.12 LCA
                                                                                   slack[j] = min(slack[j], d);
      void compress(){
          fill(vis.begin(), vis.end(), 0);
                                                                           return false;
                                                                                                                                   1 struct Tree{
          for (int i=0; i<N; i++) if (!vis[i]) dfs1(G, i);</pre>
                                                                                                                                         int N, M = 0, H;
                                                                                                                                         vector<vector<int>> G;
          fill(vis.begin(), vis.end(), 0);
                                                                       bool augment(){
                                                                                                                                         vector<vector<int>> LCA:
          reverse(order.begin(), order.end());
                                                                           for (int j=0 ; j<n ; j++){</pre>
                                                                                                                                         vector<int> parent;
          for (int i=0 ; i<N ; i++){</pre>
                                                                               if (visy[j]!=stamp && slack[j]==0){
                                                                                                                                         vector<int> dep:
              if (!vis[order[i]]){
                                                                                   visy[j] = stamp;
                  SCC.push_back(vector<int>());
                                                                                   if (match[j]==-1 || dfs(match[j], false)){
                                                                                                                                         Tree(int _N) : N(_N), H(__lg(_N)+1){
```

dfs2(inv_G, order[i]);

```
G.resize(N);
    parent.resize(N, -1);
    dep.resize(N, 0);
    LCA.resize(H, vector<int>(N, 0));
void add_edge(int u, int v){
    G[u].push_back(v);
    G[v].push_back(u);
void dfs(int now, int pre){ // root 的 pre 是自己
    dep[now] = dep[pre]+1;
    parent[now] = pre;
    for (auto x : G[now]){
        if (x==pre) continue;
        dfs(x, now);
}
void build_LCA(int root = 0){
    dfs(root, root);
    for (int i=0; i<N; i++) LCA[0][i] = parent[i];</pre>
    for (int i=1; i<H; i++){
   for (int j=0; j<N; j++){</pre>
             LCA[i][j] = LCA[i-1][LCA[i-1][j]];
}
int jump(int u, int step){
    for (int i=0 ; i<H ; i++){</pre>
        if (step&(1<<i)) u = LCA[i][u];</pre>
    return u;
int get_LCA(int u, int v){
    if (dep[u]<dep[v]) swap(u, v);</pre>
    u = jump(u, dep[u]-dep[v]);
    if (u==v) return u;
    for (int i=H-1; i>=0; i--){
        if (LCA[i][u]!=LCA[i][v]){
             u = LCA[i][u];
             v = LCA[i][v];
    return parent[u];
```

6.13 MCMF

```
struct Flow {
    struct Edge {
        int u, rc, k, rv;
    };

vector<vector<Edge>> G;
vector<iint> par, par_eid;
Flow(int n) : G(n+1), par(n+1), par_eid(n+1) {}
```

```
// v->u, capcity: c, cost: k
void add(int v, int u, int c, int k){
  G[v].push_back({u, c, k, SZ(G[u])});
  G[u].push_back({v, 0, -k, SZ(G[v])-1});
// 3701d6
int spfa(int s, int t){
  fill(ALL(par), -1);
  vector<int> dis(SZ(par), INF);
  vector<bool> in q(SZ(par), false);
  queue<int> 0;
  dis[s] = 0;
  in_q[s] = true;
  Q.push(s);
  while (!Q.empty()){
    int v = Q.front();
    Q.pop();
    in_q[v] = false;
    for (int i=0 ; i<SZ(G[v]) ; i++){</pre>
      auto [u, rc, k, rv] = G[v][i];
      if (rc>0 && dis[v]+k<dis[u]){</pre>
        dis[u] = dis[v]+k;
        par[u] = v;
        par eid[u] = i;
        if (!in_q[u]) Q.push(u);
        in_q[u] = true;
  return dis[t];
// return <max flow, min cost>, 150093
pair<int, int> flow(int s, int t){
  int fl = 0, cost = 0, d;
  while ((d = spfa(s, t))<INF){</pre>
    int cur = INF;
    for (int v=t; v!=s; v=par[v])
      cur = min(cur, G[par[v]][par_eid[v]].rc);
    fl += cur:
    cost += d*cur;
    for (int v=t ; v!=s ; v=par[v]){
      G[par[v]][par eid[v]].rc -= cur;
      G[v][G[par[v]][par_eid[v]].rv].rc += cur;
  return {fl, cost};
vector<pair<int, int>> construct(){
  vector<pair<int, int>> ret;
  for (int i=0 ; i<n ; i++){</pre>
    for (auto x : G[i]){
      if (x.rc==0){
        ret.push_back({i+1, x.u-n+1});
        break:
  return ret;
```

6.14 Tarjan

```
1 struct tarjan_SCC {
      int now_T, now_SCCs;
      vector<int> dfn, low, SCC;
      stack<int> S;
      vector<vector<int>> E;
      vector<bool> vis, in stack;
      tarjan_SCC(int n) {
          init(n);
      void init(int n) {
          now T = now SCCs = 0;
          dfn = low = SCC = vector<int>(n);
          E = vector<vector<int>>(n);
          S = stack<int>();
          vis = in stack = vector<bool>(n);
      void add(int u, int v) {
          E[u].push_back(v);
          for (int i = 0; i < dfn.size(); ++i) {</pre>
               if (!dfn[i]) dfs(i);
      void dfs(int v) {
          now T++;
          vis[v] = in stack[v] = true;
          dfn[v] = low[v] = now_T;
          S.push(v);
          for (auto &i:E[v]) {
               if (!vis[i]) {
                  vis[i] = true;
                  dfs(i);
                  low[v] = min(low[v], low[i]);
               else if (in_stack[i]) {
                  low[v] = min(low[v], dfn[i]);
38
39
          if (low[v] == dfn[v]) {
              int tmp;
43
                  tmp = S.top();
45
                  S.pop();
46
                  SCC[tmp] = now_SCCs;
47
                  in stack[tmp] = false;
              } while (tmp != v);
               now_SCCs += 1;
50
51
52 };
```

6.15 Tarjan Find AP

```
1 vector<int> dep(MAX_N), low(MAX_N), AP;
2 bitset<MAX_N> vis;
3 4 void dfs(int now, int pre){
```

```
int cnt = 0;
bool ap = 0;
vis[now] = 1;
low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);

for (auto x : G[now]){
    if (x==pre){
        continue;
    }else if (vis[x]==0){
        cnt++;
        dfs(x, now);
        low[now] = min(low[now], low[x]);
        if (low[x]>=dep[now]) ap=1;
    }else{
        low[now] = min(low[now], dep[x]);
}

if ((now==pre && cnt>=2) || (now!=pre && ap)){
        AP.push_back(now);
}
```

6.16 Tree Isomorphism

```
i #include <bits/stdc++.h>
2 #pragma GCC optimize("03,unroll-loops")
  #define fastio ios::sync with stdio(0), cin.tie(0), cout.tie
  #define dbg(x) cerr << #x << " = " << x << endl
  #define int long long
  using namespace std;
8 // declare
  const int MAX_SIZE = 2e5+5;
10 const int INF = 9e18;
11 const int MOD = 1e9+7;
12 const double EPS = 1e-6;
13 typedef vector<vector<int>> Graph;
  typedef map<vector<int>, int> Hash;
16 int n, a, b;
17 int id1, id2:
18 pair<int, int> c1, c2;
19 vector<int> sz1(MAX SIZE), sz2(MAX SIZE);
20 vector<int> we1(MAX_SIZE), we2(MAX_SIZE);
21 Graph g1(MAX_SIZE), g2(MAX_SIZE);
22 Hash m1, m2;
23 int testcase=0;
void centroid(Graph &g, vector<int> &s, vector<int> &w, pair<
      int, int> &rec, int now, int pre){
      s[now]=1;
      w[now]=0;
      for (auto x : g[now]){
          if (x!=pre){
              centroid(g, s, w, rec, x, now);
              s[now]+=s[x];
              w[now]=max(w[now], s[x]);
      w[now]=max(w[now], n-s[now]);
```

```
int dfs(Graph &g, Hash &m, int &id, int now, int pre){
    vector<int> v:
    for (auto x : g[now]){
        if (x!=pre){
            int add=dfs(g, m, id, x, now);
            v.push back(add);
    sort(v.begin(), v.end());
    if (m.find(v)!=m.end()){
        return m[v];
    }else{
        m[v]=++id;
        return id;
void solve1(){
    // init
    id1=0;
    id2=0;
    c1={0, 0};
    c2={0, 0};
    fill(sz1.begin(), sz1.begin()+n+1, 0);
    fill(sz2.begin(), sz2.begin()+n+1, 0);
    fill(we1.begin(), we1.begin()+n+1, 0);
    fill(we2.begin(), we2.begin()+n+1, 0);
    for (int i=1; i<=n; i++){
        g1[i].clear();
        g2[i].clear();
   m1.clear();
    m2.clear();
    // input
    for (int i=0 ; i<n-1 ; i++){</pre>
        cin >> a >> b;
        g1[a].push_back(b);
        g1[b].push back(a);
    for (int i=0 ; i<n-1 ; i++){</pre>
        cin >> a >> b;
        g2[a].push back(b);
        g2[b].push_back(a);
    // get tree centroid
    centroid(g1, sz1, we1, c1, 1, 0);
    centroid(g2, sz2, we2, c2, 1, 0);
    // process
    int res1=0, res2=0, res3=0;
    if (c2.second!=0){
        res1=dfs(g1, m1, id1, c1.first, 0);
        m2=m1:
        id2=id1;
```

if (w[now]<=n/2){</pre>

}

else rec.second=now;

if (rec.first==0) rec.first=now;

```
res2=dfs(g2, m1, id1, c2.first, 0);
104
           res3=dfs(g2, m2, id2, c2.second, 0);
       }else if (c1.second!=0){
           res1=dfs(g2, m1, id1, c2.first, 0);
106
107
108
           res2=dfs(g1, m1, id1, c1.first, 0);
109
           res3=dfs(g1, m2, id2, c1.second, 0);
110
111
112
           res1=dfs(g1, m1, id1, c1.first, 0);
113
           res2=dfs(g2, m1, id1, c2.first, 0);
114
       }
115
116
       cout << (res1==res2 || res1==res3 ? "YES" : "NO") << end1</pre>
117
118
119
       return;
120 }
121
   signed main(void){
122
       fastio:
       int t=1:
       cin >> t;
       while (t--){
128
           solve1();
129
130
       return 0;
131
```

6.17 圓方樹

```
1 #include <bits/stdc++.h>
  #define lp(i,a,b) for(int i=(a);i<(b);i++)</pre>
  #define pii pair<int,int>
  #define pb push back
  #define ins insert
  #define ff first
  #define ss second
  #define opa(x) cerr << #x << " = " << x << ", ";</pre>
  #define op(x) cerr << #x << " = " << x << endl;
  #define ops(x) cerr << x;</pre>
  #define etr cerr << endl;</pre>
12 #define spc cerr << ' ';
#define BAE(x) (x).begin(), (x).end()
  #define STL(x) cerr << #x << " : "; for(auto &qwe:x) cerr <<</pre>
  qwe << ' '; cerr << endl;
#define deb1 cerr << "deb1" << endl;</pre>
  #define deb2 cerr << "deb2" << endl:
  #define deb3 cerr << "deb3" << endl;</pre>
#define deb4 cerr << "deb4" << endl;</pre>
#define deb5 cerr << "deb5" << endl;</pre>
20 #define bye exit(0);
  using namespace std;
  const int mxn = (int)(2e5) + 10;
  const int mxlg = 17;
  int last_special_node = (int)(1e5) + 1;
  vector<int> E[mxn], F[mxn];
  struct edg{
      int fr, to;
```

```
edg(int _fr, int _to){
          fr = fr;
          to = to;
                                                                      inline void build lca(){
33
                                                                          jmp[1][0] = 1;
                                                                          dfs lca(1, -1, 1);
34 };
35 ostream& operator << (ostream& os, edg x) {os << x.fr << "--" <<
                                                                          lp(j,1,mxlg){
                                                                              lp(i,1,mxn){
  vector<edg> EV:
                                                                                  jmp[i][j] = jmp[jmp[i][j-1]][j-1];
  void tarjan(int v, int par, stack<int>& S){
                                                                  103
                                                                          }
      static vector<int> dfn(mxn), low(mxn);
                                                                  104
      static vector<bool> to add(mxn);
      static int nowT = 0;
                                                                      inline int lca(int x, int y){
                                                                          if(dep[x] < dep[y]){ swap(x, y); }</pre>
      int childs = 0;
      nowT += 1:
                                                                          int diff = dep[x] - dep[y];
      dfn[v] = low[v] = nowT;
                                                                          lp(j,0,mxlg){
                                                                  110
      for(auto &ne:E[v]){
                                                                              if((diff >> j) & 1){
                                                                  111
          int i = EV[ne].to;
                                                                  112
                                                                                  x = jmp[x][j];
          if(i == par) continue;
                                                                  113
          if(!dfn[i]){
                                                                  114
              S.push(ne);
                                                                          if(x == y) return x;
                                                                  115
              tarjan(i, v, S);
                                                                  116
              childs += 1;
                                                                          for(int j = mxlg - 1; j >= 0; j--){
                                                                  117
              low[v] = min(low[v], low[i]);
                                                                              if(jmp[x][j] != jmp[y][j]){
                                                                  118
                                                                                  x = jmp[x][j];
                                                                  119
              if(par >= 0 && low[i] >= dfn[v]){
                                                                  120
                                                                                  y = jmp[y][j];
                   vector<int> bcc;
                                                                  121
                   int tmp;
                                                                  122
                                                                          return jmp[x][0];
                   do{
                                                                  123
                       tmp = S.top(); S.pop();
                                                                  124
                       if(!to_add[EV[tmp].fr]){
                           to_add[EV[tmp].fr] = true;
                                                                      inline bool can reach(int fr, int to){
                           bcc.pb(EV[tmp].fr);
                                                                          if(dep[to] > dep[fr]) return false;
                       if(!to add[EV[tmp].to]){
                                                                          int diff = dep[fr] - dep[to];
                                                                  129
                           to_add[EV[tmp].to] = true;
                                                                          lp(j,0,mxlg){
                                                                  130
                           bcc.pb(EV[tmp].to);
                                                                              if((diff >> j) & 1){
                                                                  131
                                                                                  fr = jmp[fr][j];
                                                                  132
                   }while(tmp != ne);
                                                                  133
                   for(auto &j:bcc){
                                                                  134
                       to_add[j] = false;
                                                                  135
                                                                          return fr == to;
                       F[last special node].pb(j);
                                                                  136
                       F[j].pb(last special node);
                                                                  137
                                                                          ios::sync with stdio(false); cin.tie(0);
                   last special node += 1;
                                                                          freopen("test input.txt", "r", stdin);
                                                                          int n, m, q; cin >> n >> m >> q;
          else{
                                                                          lp(i,0,m){
               low[v] = min(low[v], dfn[i]);
                                                                              int u, v; cin >> u >> v;
               if(dfn[i] < dfn[v]){ // edge i--v will be visited 144</pre>
                                                                              E[u].pb(EV.size());
                     twice at here, but we only need one.
                                                                              EV.pb(edg(u, v));
                   S.push(ne);
                                                                              E[v].pb(EV.size());
                                                                              EV.pb(edg(v, u));
                                                                  147
                                                                  148
                                                                          E[0].pb(EV.size());
      }
                                                                  149
                                                                          EV.pb(edg(0, 1));
                                                                  150
                                                                  151
                                                                          stack<int> S;
  int dep[mxn], jmp[mxn][mxlg];
                                                                          tarjan(0, -1, S);
                                                                  152
  void dfs lca(int v, int par, int depth){
                                                                  153
                                                                          build lca():
      dep[v] = depth;
                                                                  154
      for(auto &i:F[v]){
                                                                  155
                                                                          lp(queries,0,q){
          if(i == par) continue;
                                                                              int fr, to, relay; cin >> fr >> to >> relay;
                                                                  156
          imp[i][0] = v;
                                                                              if(fr == relay || to == relay){
                                                                  157
          dfs lca(i, v, depth + 1);
                                                                                  cout << "NO\n":
                                                                  158
                                                                                  continue:
```

6.18 最大權閉合圖

```
1 /*
2 Problem:
      Given w = [w_0, w_1, ..., w_{n-1}] (which can be
      either positive or negative or 0), you can choose
      to take w i (0 < i < n) or not, but if edge u \rightarrow v
      exists, you must take w_v if you want to take w_u
      (in other words, you can't take w u without taking
       w v), this function returns the maximum value(> 0)
       you can get. If you need a construction, you can
       output the minimum cut of the S(source) side.
      MaxFlow(n, m) (Non-Biparte:O(n²m) / Bipartite:O(m√n))
13
  int maximum_closure(vector<int> w, vector<pair<int,int>> EV)
      int n = w.size(), S = n + 1, T = n + 2;
      Flow G(T + 5); // Graph/Dinic.cpp
      int sum = 0:
      for (int i = 0; i < n; ++i) {</pre>
          if (w[i] > 0) {
               G.add(S, i, w[i]);
               sum += w[i];
           else if (w[i] < 0) {
24
               G.add(i, T, abs(w[i]));
25
26
      for (auto &[u, v] : EV) { // You should make sure that
27
           INF > \Sigma / W i
          G.add(u, v, INF);
28
29
      int cut = G.flow(S, T):
30
31
      return sum - cut;
32
```

6.19 Theorem

- 仟意圖
 - 不能有孤點、最大匹配 + 最小邊覆蓋 = n 點覆蓋的補集是獨立集。 最小點覆蓋 + 最大獨立集 = n
- 二分圖
 - 最小點覆蓋 = 最大匹配 = n 最大獨立集
- 只有邊帶權的一分圖

- w-vertex-cover (帶權點覆蓋): 每條邊的兩個連接點被選中的次數總和至少要是 w_e 。
- w-weight matching (帶權匹配)
- minimum vertex count of w-vertex-cover = maximum weight count of 12 // 對於方程組的式子兩兩求解 w-weight matching (一個點可以被選很多次・但邊不行) 13 // {是否有解。{a. m}}
- 點、邊都帶權的二分圖的定理
 - b-matching:假設 v 的點權是 b_v · 那所有 v 的匹配邊 e 的權重都要滿足 $\sum w_e < b_v$ 。
 - The maximum w-weight of a b-matching equals the minimum b-weight
 of vertices in a w-vertex-cover.

7 Math

7.1 CRT m Coprime

```
1 vector<int> a, m;
  int extgcd(int a, int b, int &x, int &y){
      if (b==0){
          x=1, y=0;
          return a;
      int ret=extgcd(b, a%b, y, x);
      y-=a/b*x;
      return ret;
14 // n = 有幾個式子,求解 x \equiv a_i \bmod m_i
int CRT(int n, vector<int> &a, vector<int> &m){
      int p=1, ans=0;
      vector<int> M(n), inv_M(n);
      for (int i=0 ; i<n ; i++) p*=m[i];</pre>
      for (int i=0 ; i<n ; i++){</pre>
          M[i]=p/m[i];
      int tmp;
          extgcd(M[i], m[i], inv_M[i], tmp);
          ans+=a[i]*inv_M[i]*M[i];
          ans%=p;
      return (ans%p+p)%p;
```

7.2 CRT m Not Coprime

```
int extgcd(int a, int b, int &x, int &y){
   if (b==0){
      x=1, y=0;
      return a;
   }
int ret=extgcd(b, a%b, y, x);
```

7.3 Josephus Problem

```
1  // 有 n 個人·第偶數個報數的人被刪掉·問第 k 個被踢掉的是誰
2  int solve(int n, int k){
    if (n==1) return 1;
    if (k<=(n+1)/2){
        if (2*k>n) return 2*k%n;
        else return 2*k;
    }else{
        int res=solve(n/2, k-(n+1)/2);
        if (n&1) return 2*res+1;
        else return 2*res-1;
    }
}
```

7.4 Lagrange any x

```
1 / / init: (x1, y1), (x2, y2) in a vector
 struct Lagrange{
      int n:
      vector<pair<int, int>> v;
      Lagrange(vector<pair<int, int>> & v){
          n = _v.size();
          v = v;
      // O(n^2 \log MAX A)
      int solve(int x){
          int ret = 0;
          for (int i=0 ; i<n ; i++){</pre>
              int now = v[i].second;
              for (int j=0 ; j<n ; j++){</pre>
                  if (i==j) continue;
                  now *= ((x-v[j].first)+MOD)%MOD;
                  now %= MOD:
                  now *= (qp((v[i].first-v[j].first+MOD)%MOD,
                       MOD - 2) + MOD) % MOD;
```

```
7.5 Lagrange continuous x
```

return ret;

now %= MOD;

ret = (ret+now)%MOD;

```
1 #include <bits/stdc++.h>
  using namespace std;
  const int MAX N = 5e5 + 10;
  const int mod = 1e9 + 7;
  long long inv_fac[MAX_N];
  inline int fp(long long x, int y) {
      int ret = 1;
      for (; y; y >>= 1) {
          ret = (y & 1) ? (ret * x % mod) : ret;
          x = x * x % mod:
15
      return ret;
16
  // TO USE THIS TEMPLATE, YOU MUST MAKE SURE THAT THE MOD
       NUMBER IS A PRIME.
19 struct Lagrange {
  /*
20
      Initialize a polynomial with f(x_0), f(x_0 + 1), ..., f(
      This determines a polynomial f(x) whose degree is at most
      Then you can call sample(x) and you get the value of f(x)
      Complexity of init() and sample() are both O(n).
24
25
      int m, shift; // m = n + 1
26
      vector<int> v, mul;
  // You can use this function if you don't have inv_fac array
      void construct inv fac() {
          long long fac = 1;
31
          for (int i = 2; i < MAX_N; ++i) {</pre>
              fac = fac * i % mod:
          inv fac[MAX N - 1] = fp(fac, mod - 2);
          for (int i = MAX N - 1; i >= 1; --i) {
              inv_fac[i - 1] = inv_fac[i] * i % mod;
36
37
  // You call init() many times without having a second
       instance of this struct.
      void init(int X_0, vector<int> &u) {
          shift = ((1 - X 0) \% mod + mod) \% mod;
          if (v.size() == 1) v.push_back(v[0]);
          m = v.size();
          mul.resize(m);
```

```
47 // You can use sample(x) instead of sample(x % mod).
      int sample(int x) {
                                                                  13
          x = ((long long)x + shift) % mod;
          x = (x < 0) ? (x + mod) : x;
          long long now = 1;
          for (int i = m; i >= 1; --i) {
              mul[i - 1] = now;
              now = now * (x - i) % mod;
          int ret = 0;
          bool neg = (m - 1) & 1;
          now = 1;
          for (int i = 1; i <= m; ++i) {</pre>
              int up = now * mul[i - 1] % mod;
              int down = inv_fac[m - i] * inv_fac[i - 1] % mod; 25
              int tmp = ((long long)v[i - 1] * up % mod) * down 26
              ret += (neg && tmp) ? (mod - tmp) : (tmp);
              ret = (ret >= mod) ? (ret - mod) : ret;
              now = now * (x - i) % mod;
              neg ^= 1;
          return ret;
  };
  int main() {
      int n; cin >> n;
      vector<int> v(n);
      for (int i = 0; i < n; ++i) {</pre>
          cin >> v[i];
      Lagrange L;
      L.construct_inv_fac();
      L.init(0, v);
      int x; cin >> x;
      cout << L.sample(x);</pre>
```

7.6 Lucas's Theorem

```
1 // 對於很大的 C^n_{m} 對質數 p 取模·只要 p 不大就可以用。
2 int Lucas(int n, int m, int p){
3     if (m==0) return 1;
4     return (C(n%p, m%p, p)*Lucas(n/p, m/p, p)%p);
5 }
```

7.7 Matrix

```
struct Matrix{
    int n, m;
    vector<vector<int>> arr;

Matrix(int _n, int _m){
        n = _n;
        m = _m;
        arr.assign(n, vector<int>(m));
}

vector<int> & operator [] (int i){
```

```
Matrix operator * (Matrix b){
    Matrix ret(n, b.m);
    for (int i=0 ; i<n ; i++){</pre>
        for (int j=0 ; j<b.m ; j++){</pre>
            for (int k=0 ; k<m ; k++){</pre>
                 ret.arr[i][j] += arr[i][k]*b.arr[k][j]%
                 ret.arr[i][j] %= MOD;
            }
        }
    return ret;
Matrix pow(int p){
    Matrix ret(n, n), mul = *this;
    for (int i=0 ; i<n ; i++){</pre>
        ret.arr[i][i] = 1;
    for ( ; p ; p>>=1){
        if (p&1) ret = ret*mul;
        mul = mul*mul;
    return ret;
}
int det(){
    vector<vector<int>> arr = this->arr;
    bool flag = false;
    for (int i=0 ; i<n ; i++){</pre>
        int target = -1;
        for (int j=i ; j<n ; j++){</pre>
            if (arr[j][i]){
                 target = j;
                 break;
        if (target==-1) return 0;
        if (i!=target){
            swap(arr[i], arr[target]);
             flag = !flag;
        for (int j=i+1 ; j<n ; j++){</pre>
             if (!arr[j][i]) continue;
             int freq = arr[j][i]*qp(arr[i][i], MOD-2)%MOD
             for (int k=i ; k<n ; k++){</pre>
                 arr[j][k] -= freq*arr[i][k];
                 arr[j][k] = (arr[j][k]%MOD+MOD)%MOD;
        }
    int ret = !flag ? 1 : MOD-1;
```

for (int i=0 ; i<n ; i++){</pre>

ret *= arr[i][i];

ret %= MOD;

return ret;

return arr[i];

}

7.8 Matrix 01

76 77 };

```
1 \mid const int MAX N = (1LL << 12);
  struct Matrix{
       int n, m;
       vector<bitset<MAX N>> arr;
       Matrix(int _n, int _m){
           n = _n;
           m = _m;
           arr.resize(n);
       Matrix operator * (Matrix b){
           Matrix b_t(b.m, b.n);
           for (int i=0 ; i<b.n ; i++){</pre>
                for (int j=0 ; j<b.m ; j++){</pre>
                    b_t.arr[j][i] = b.arr[i][j];
           Matrix ret(n, b.m);
           for (int i=0 ; i<n ; i++){</pre>
                for (int j=0 ; j<b.m ; j++){</pre>
                    ret.arr[i][j] = ((arr[i]&b t.arr[j]).count()
           return ret;
27
28 };
```

7.9 Miller Rabin

```
1 // O(\log n)
  typedef Uint unsigned long long
  Uint modmul(Uint a, Uint b, Uint m) {
      int ret = a*b - m*(Uint)((long double)a*b/m);
      return ret + m*(ret < 0) - m*(ret>=(int)m);
  int qp(int b, int p, int m){
      int ret = 1;
      for (; p; p>>=1){
          if (p&1){
              ret = modmul(ret, b, m);
          b = modmul(b, b, m);
16
      return ret;
  // ed23aa
  vector<int> llsprp = {2, 325, 9375, 28178, 450775, 9780504,
       1795265022};
21 bool isprime(int n, vector(int) sprp = llsprp){
      if (n==2) return 1;
```

```
if (n<2 || n%2==0) return 0;
int t = 0:
int u = n-1;
for (; u%2==0; t++) u>>=1;
for (int i=0 ; i<sprp.size() ; i++){</pre>
    int a = sprp[i]%n;
    if (a==0 || a==1 || a==n-1) continue;
    int x = qp(a, u, n);
    if (x==1 || x==n-1) continue;
    for (int j=0 ; j<t ; j++){</pre>
        x = modmul(x, x, n);
        if (x==1) return 0;
        if (x==n-1) break;
    if (x==n-1) continue;
    return 0:
}
return 1;
```

7.10 Pollard Rho

```
i | mt19937 seed(chrono::steady_clock::now().time_since_epoch().
      count());
int rnd(int 1, int r){
     return uniform int distribution<int>(1, r)(seed);
6 // O(n^{1/4}) 回傳 1 或自己的因數、記得先判斷 n 是不是質數
      (用 Miller-Rabin)
7 // c1670c
8 int Pollard Rho(int n){
     int s = 0, t = 0;
     int c = rnd(1, n-1);
     int step = 0, goal = 1;
     int val = 1:
     for (goal=1 ; ; goal<<=1, s=t, val=1){</pre>
         for (step=1 ; step<=goal ; step++){</pre>
             t = ((int128)t*t+c)%n;
             val = (int128)val*abs(t-s)%n;
             if ((step % 127) == 0){
                 int d = gcd(val, n);
                 if (d>1) return d;
         int d = __gcd(val, n);
         if (d>1) return d;
```

```
7.11 Polynomial
```

```
1 struct Poly {
      int len, deg;
      int *a:
      // len = 2^k >= the original length
      Poly(): len(0), deg(0), a(nullptr) {}
      Poly(int n) {
          len = 1;
          deg = n - 1;
          while (len < n) len <<= 1;</pre>
          a = (ll*) calloc(len, sizeof(ll));
      Poly(int 1, int d, int *b) {
          len = 1;
          deg = d;
          a = b;
      void resize(int _n) {
          int len1 = 1;
          while (len1 < n) len1 <<= 1;</pre>
          int *res = (ll*) calloc(len1, sizeof(ll));
          for (int i = 0; i < min(len, _n); i++) {</pre>
              res[i] = a[i];
          len = len1:
          deg = n - 1;
          free(a);
          a = res:
      Poly& operator=(const Poly rhs) {
          this->len = rhs.len;
          this->deg = rhs.deg;
          this->a = (11*)realloc(this->a, sizeof(11) * len);
          copy(rhs.a, rhs.a + len, this->a);
          return *this:
                                                                    100
                                                                    101
      Poly operator*(Poly rhs) {
                                                                    102
          int 11 = this->len, 12 = rhs.len;
                                                                    103
          int d1 = this->deg, d2 = rhs.deg;
          while (11 > 0 and this->a[11 - 1] == 0) 11--;
                                                                    105
          while (12 > 0 \text{ and } rhs.a[12 - 1] == 0) 12--;
                                                                    106
                                                                    107
          while (1 < max(11 + 12 - 1, d1 + d2 + 1)) 1 <<= 1;
          int *x, *y, *res;
                                                                    109
          x = (11*) calloc(1, sizeof(11));
                                                                    110
          y = (11*) calloc(1, sizeof(11));
                                                                    111
          res = (11*) calloc(1, sizeof(11));
                                                                    112
          copy(this->a, this->a + 11, x);
                                                                    113
          copy(rhs.a, rhs.a + 12, y);
                                                                    114
          ntt.tran(1, x); ntt.tran(1, y);
                                                                    115
          FOR (i, 0, 1 - 1)
          res[i] = x[i] * y[i] % mod;
ntt.tran(l, res, true);
          free(x); free(y);
                                                                    119
          return Poly(1, d1 + d2, res);
                                                                    120
                                                                    121
      Poly operator+(Poly rhs) {
          int 11 = this->len, 12 = rhs.len;
                                                                    123
          int 1 = \max(11, 12);
                                                                    124
          Poly res;
                                                                    125
          res.len = 1:
                                                                    126
          res.deg = max(this->deg, rhs.deg);
                                                                    127
          res.a = (ll*) calloc(l, sizeof(ll));
                                                                    128
          FOR (i, 0, 11 - 1) {
                                                                    129
              res.a[i] += this->a[i];
                                                                    130
              if (res.a[i] >= mod) res.a[i] -= mod;
                                                                    131
```

```
FOR (i, 0, 12 - 1) {
        res.a[i] += rhs.a[i];
        if (res.a[i] >= mod) res.a[i] -= mod;
    return res;
Poly operator-(Poly rhs) {
    int 11 = this->len, 12 = rhs.len;
    int 1 = \max(11, 12);
    Poly res;
    res.len = 1;
    res.deg = max(this->deg, rhs.deg);
    res.a = (ll*) calloc(l, sizeof(ll));
    FOR (i, 0, 11 - 1) {
        res.a[i] += this->a[i];
        if (res.a[i] >= mod) res.a[i] -= mod;
    FOR (i, 0, 12 - 1) {
        res.a[i] -= rhs.a[i];
        if (res.a[i] < 0) res.a[i] += mod;</pre>
    return res;
Poly operator*(const int rhs) {
    Poly res;
    res = *this;
    FOR (i, 0, res.len - 1) {
        res.a[i] = res.a[i] * rhs % mod;
        if (res.a[i] < 0) res.a[i] += mod;</pre>
    return res;
Poly(vector<int> f) {
    int _n = f.size();
    len = 1:
    deg = n - 1;
    while (len < _n) len <<= 1;</pre>
    a = (11*) calloc(len, sizeof(11));
    FOR (i, 0, deg) a[i] = f[i];
Poly derivative() {
    Poly g(this->deg);
    FOR (i, 1, this->deg) {
        g.a[i - 1] = this->a[i] * i % mod;
    return g;
Poly integral() {
    Poly g(this->deg + 2);
    FOR (i, 0, this->deg) {
        g.a[i + 1] = this -> a[i] * ::inv(i + 1) % mod;
    return g;
Poly inv(int len1 = -1) {
    if (len1 == -1) len1 = this->len;
    Poly g(1); g.a[0] = ::inv(a[0]);
    for (int l = 1; l < len1; l <<= 1) {</pre>
        Poly t; t = *this;
        t.resize(1 << 1);
        t = g * g * t;
        t.resize(1 << 1);
        Poly g1 = g * 2 - t;
        swap(g, g1);
    return g;
```

```
Poly ln(int len1 = -1) {
134
135
            if (len1 == -1) len1 = this->len:
            auto g = *this;
136
137
            auto x = g.derivative() * g.inv(len1);
            x.resize(len1);
138
139
            x = x.integral();
            x.resize(len1);
140
            return x;
141
142
143
       Poly exp() {
144
            Poly g(1);
            g.a[0] = 1;
145
            for (int 1 = 1; 1 < len; 1 <<= 1) {
146
147
                Poly t, g1; t = *this;
                t.resize(1 << 1); t.a[0]++;
148
149
                g1 = (t - g.ln(1 << 1)) * g;
                g1.resize(1 << 1);
150
151
                swap(g, g1);
152
153
            return g;
154
       Poly pow(ll n) {
155
156
            Poly &a = *this;
            int i = 0;
157
158
            while (i <= a.deg and a.a[i] == 0) i++;</pre>
159
            if (i and (n > a.deg or n * i > a.deg)) return Poly(a
                 .deg + 1);
            if (i == a.deg + 1) {
                Poly res(a.deg + 1);
161
                res.a[0] = 1;
162
                return res;
163
164
165
            Poly b(a.deg - i + 1);
            int inv1 = ::inv(a.a[i]);
166
            FOR (j, 0, b.deg)
                b.a[j] = a.a[j + i] * inv1 % mod;
            Poly res1 = (b.ln() * (n % mod)).exp() * (::power(a.a
            Poly res2(a.deg + 1);
FOR (j, 0, min((ll)(res1.deg), (ll)(a.deg - n * i)))
171
                res2.a[j + n * i] = res1.a[j];
172
173
            return res2;
174
175 };
```

7.12 數論分塊

7.13 最大質因數

```
void max_fac(int n, int &ret){
    if (n<=ret || n<2) return;
    if (isprime(n)){
        ret = max(ret, n);
        return;
}

int p = Pollard_Rho(n);
max_fac(p, ret), max_fac(n/p, ret);
}</pre>
```

7.14 歐拉公式

```
1 / / phi(n) = 小於 n 並與 n 互質的正整數數量。
2 // O(sqrt(n)) · 回傳 phi(n)
3 int phi(int n){
     int ret = n;
      for (int i=2 ; i*i<=n ; i++){</pre>
         if (n%i==0){
              while (n%i==0) n /= i;
              ret = ret*(i-1)/i;
     if (n>1) ret = ret*(n-1)/n;
      return ret;
  // O(n Log n) · 回傳 1~n 的 phi 值
 vector<int> phi 1 to n(int n){
     vector<int> phi(n+1);
      phi[0]=0;
      phi[1]=1;
      for (int i=2 ; i<=n ; i++){</pre>
         phi[i]=i-1;
      for (int i=2 ; i<=n ; i++){</pre>
         for (int j=2*i; j<=n; j+=i){ // 枚舉所有倍數
             phi[j]-=phi[i];
     }
      return phi;
```

7.15 Burnside's Lemma

$$\sum_{k=1}^{n} \frac{c(k)}{n}$$

- n:有多少種置換方式(例如:旋轉方式)
- c(k): 所有可能中,經過 k 次旋轉後,仍不會和別人相同的方式的數量

7.16 Catalan Number

任意括號序列: $C_n = \frac{1}{n+1} \binom{2n}{n}$

7.17 Matrix Tree Theorem

目標:給定一張無向圖,問他的生成樹數量。 方法:先把所有自環刪掉,定義 Q 為以下矩陣

$$Q_{i,j} = \begin{cases} \deg(v_i) & \text{if } i = j \\ -(邊v_iv_j \text{ 的數量}) & \text{otherwise} \end{cases}$$

接著刪掉 Q 的第一個 row 跟 column \cdot 它的 determinant 就是答案。目標:給定一張有向圖 \cdot 問他的以 r 為根 \cdot 可以走到所有點生成樹數量 \cdot

方法:先把所有自環刪掉,定義 Q 為以下矩陣

$$Q_{i,j} = \begin{cases} \deg_{in}(v_i) & \text{if } i = j \\ -(\frac{1}{8}v_i v_j \text{ bbm}) & \text{otherwise} \end{cases}$$

接著刪掉 Q 的第 r 個 row 跟 column · 它的 determinant 就是答案。

7.18 Stirling's formula

 $n! \approx \sqrt{2\pi n} (\frac{n}{\epsilon})^n$

7.19 Theorem

- 1. $1 \sim x$ 質數的數量 $\approx \frac{x}{\ln x}$
- 2. x 的因數的數量 $\approx x^{\frac{1}{3}}$
- 3. x 的質因數的數量 $\approx \log \log x$
- 4. p is a prime number $\Leftrightarrow (p-1)! \equiv -1 \pmod{p}$
- 5. 每個正整數都可以表示成四個整數的平方和
- 6. 仟何大於 2 的整數都可以表示成兩個質數的和
- 7. $n^{k-2} \cdot \prod_{i=1}^{k} s_i n$ 個點、k 的連通塊 · 加上 k-1 條邊使得變成一個連通 圖的方法數 · 其中每個連通塊有 s_i 個點

7.20 二元一次方程式

$$\begin{cases} ax + by = c \\ dx + ey = f \end{cases} = \begin{cases} x = \frac{ed - bf}{ad - bc} \\ y = \frac{af - ec}{ad - bc} \end{cases}$$

若 $x=\frac{0}{0}$ 且 $y=\frac{0}{0}$ · 則代表無限多組解。若 $x=\frac{*}{0}$ 且 $y=\frac{*}{0}$ · 則代表無 \mathbf{w} 。

歐拉定理 7.21

```
若 a, m 互質 · 則:
                            a^n \equiv a^{n \bmod \varphi(m)} \pmod{m}
若a, m不互質,則:
                       a^n \equiv a^{\varphi(m) + [n \mod \varphi(m)]} \pmod{m}
```

錯排公式

錯排公式: (n 個人中,每個人皆不再原來位置的組合數)

$$dp_i = \begin{cases} 1 & i = 0\\ 0 & i = 1\\ (i-1)(dp_{i-1} + dp_{i-2}) & \text{otherwise} \end{cases}$$

String

8.1 Hash

```
i | mt19937 seed(chrono::steady clock::now().time since epoch().
       count());
 int rng(int 1, int r){
      return uniform_int_distribution<int>(1, r)(seed);
 int A = rng(1e5, 8e8);
  const int B = 1e9+7;
8 // 2f6192
  struct RollingHash{
      vector<int> Pow, Pre;
      RollingHash(string s = ""){
          Pow.resize(s.size());
          Pre.resize(s.size());
          for (int i=0 ; i<s.size() ; i++){</pre>
              if (i==0){
                  Pow[i] = 1;
                  Pre[i] = s[i];
              }else{
                  Pow[i] = Pow[i-1]*A%B;
                  Pre[i] = (Pre[i-1]*A+s[i])%B;
          return:
      int get(int 1, int r){ // 取得 [l, r] 的數值
          if (l==0) return Pre[r];
          int res = (Pre[r]-Pre[l-1]*Pow[r-l+1])%B;
          if (res<0) res += B;
          return res;
34 };
```

8.2 KMP

```
ı|// 給一個字串 S·定義函數 \pi(i) = k 代表 S[1 ... k] = S[i-k
      +1 ... i](最長真前後綴)
 // e5b7ce
 vector<int> KMP(string &s){
     int n = s.size();
     vector<int> ret(n);
     for (int i=1 ; i<n ; i++){</pre>
         int j = ret[i-1];
         while (j>0 && s[i]!=s[j]) j = ret[j-1];
         j += (s[i]==s[j]);
         ret[i] = j;
     return ret;
```

Manacher

```
i string Manacher(string str) {
      string tmp = "$#";
      for(char i : str) {
         tmp += i;
          tmp += '#';
      vector<int> p(tmp.size(), 0);
     int mx = 0, id = 0, len = 0, center = 0;
     for(int i=1; i<(int)tmp.size(); i++) {</pre>
         p[i] = mx > i ? min(p[id*2-i], mx-i) : 1;
          while(tmp[i+p[i]] == tmp[i-p[i]]) p[i]++;
          if(mx<i+p[i]) mx = i+p[i], id = i;
          if(len<p[i]) len = p[i], center = i;</pre>
     return str.substr((center-len)/2, len-1);
```

8.4 Min Rotation

```
1 // 9d296f
2 int minRotation(string s) {
      int a=0, N=SZ(s); s += s;
      for (int b=0 ; b<N ; b++){</pre>
           for (int k=0 ; k<N ; k++){</pre>
               if (a+k == b \mid | s[a+k] < s[b+k]) \{b += max(0LL, k 50)\}
                    -1); break;}
               if (s[a+k] > s[b+k]) { a = b; break; }
      return a;
```

8.5 Suffix Array

```
u|// 注意·當 /s/=1 時·lcp 不會有值·務必測試 /s/=1 的 case
2 struct SuffixArray {
      string s;
      vector<int> sa, lcp;
      SuffixArray(string _s, int lim = 256) {
          s = s;
          int n = s.size()+1, k = 0, a, b;
          vector<int> x(s.begin(), s.end()), y(n), ws(max(n,
               lim)), rank(n);
          x.push back(0);
          sa = 1cp = y;
          iota(sa.begin(), sa.end(), 0);
          for (int j=0, p=0 ; p<n ; j=max(1LL, j*2), lim=p) {</pre>
              p = j;
              iota(y.begin(), y.end(), n-j);
              for (int i=0; i<n; i++) if (sa[i] >= j) y[p++]
                   = sa[i] - j;
              fill(ws.begin(), ws.end(), 0);
              for (int i=0 ; i<n ; i++) ws[x[i]]++;</pre>
              for (int i=1; i<lim; i++) ws[i] += ws[i - 1];</pre>
              for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
              swap(x, y), p = 1, x[sa[0]] = 0;
              for (int i=1 ; i<n ; i++){</pre>
                  a = sa[i - 1];
                  b = sa[i];
                  x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
                        ? p - 1 : p++;
          for (int i=1; i<n; i++) rank[sa[i]] = i;</pre>
          for (int i=0, j ; i<n-1 ; lcp[rank[i++]]=k)</pre>
              for (k && k--, j=sa[rank[i]-1]; i+k<s.size() &&</pre>
                   j+k<s.size() && s[i+k]==s[j+k]; k++);
          sa.erase(sa.begin());
          lcp.erase(lcp.begin(), lcp.begin()+2);
      // f49583
      vector<int> pos; // pos[i] = i 這個值在 pos 的哪個地方
      SparseTable st:
      void init lcp(){
          pos.resize(sa.size());
          for (int i=0 ; i<sa.size() ; i++){</pre>
              pos[sa[i]] = i;
          if (lcp.size()){
              st.build(lcp);
      // 用之前記得 init
      // 回傳 [l1, r1] 跟 [l2, r2] 的 lcp·0-based
      int get lcp(int l1, int r1, int l2, int r2){
          int pos 1 = pos[11], len 1 = r1-11+1;
          int pos_2 = pos[12], len_2 = r2-12+1;
          if (pos 1>pos 2){
              swap(pos 1, pos 2);
              swap(len_1, len_2);
          if (11==12){
              return min(len 1, len 2);
```

26

60

```
8.6 Z Algorithm
          }else{
              return min({st.query(pos_1, pos_2), len_1, len_2
                   });
                                                               1// 定義一個長度為 n 的文本為 T ,則陣列 Z 的 Z[i] 代表 T[0:n]
      }
                                                                       和 T[i:n] 最長共同前綴
                                                                 // bcfbd6
      // 檢查 [L1, r1] 跟 [L2, r2] 的大小關係 · 0-based
                                                                 vector<int> z_function(string s){
                                                                     vector<int> ret(s.size());
      // 如果前者小於後者,就回傳 <0,相等就回傳 =0,否則回傳
                                                                     int 11 = 0, rr = 0;
      // 5b8db0
                                                                     for (int i=1; i<s.size(); i++){</pre>
      int substring_cmp(int l1, int r1, int l2, int r2){
                                                                         int j = 0;
          int len 1 = r1-l1+1;
          int len_2 = r2-l2+1;
                                                                         if (i<rr) j = min(ret[i-ll], rr-i);</pre>
          int res = get lcp(l1, r1, l2, r2);
                                                                         while (s[j]==s[i+j]) j++;
                                                                         ret[i] = j;
          if (res<len_1 && res<len_2){</pre>
              return s[11+res]-s[12+res];
                                                                         if (i+j>rr){
          }else if (len_1==res && len_2==res){
                                                                             11 = i;
              // 如果不需要以 index 作為次要排序參數,這裡要回
                                                                             rr = i+j;
              return 11-12;
          }else{
              return len 1==res ? -1 : 1;
                                                                     ret[0] = s.size();
                                                                     return ret;
      // 對於位置在 <=p 的後綴·找離他左邊/右邊最接近位置 >p 的
           後綴的 Lcp · 0-based
                                                                 8.7 k-th Substring1
      // pre[i] = s[i] 離他左邊最接近位置 >p 的後綴的 Lcp · 0-
      // suf[i] = s[i] 離他右邊最接近位置 >p 的後綴的 Lcp · 0-
                                                               1// 回傳 s 所有子字串(完全不同)中,第 k 大的
           based
      // da12fa
                                                               2 string k_th_substring(string &s, int k){
      pair<vector<int>, vector<int>> get_left_and_right_lcp(int
                                                                     int n = s.size();
                                                                     SuffixArray sa(s);
          vector<int> pre(p+1);
                                                                     sa.init_lcp();
          vector<int> suf(p+1);
                                                                     int prePrefix = 0, nowRank = 0;
          { // build pre
                                                                     for (int i=0 ; i<n ; i++){</pre>
              int now = 0;
                                                                         int len = n-sa[i];
              for (int i=0 ; i<s.size() ; i++){</pre>
                                                                         int add = len-prePrefix;
                  if (sa[i]<=p){</pre>
                      pre[sa[i]] = now;
                                                                         if (nowRank+add>=k){
                      if (i<lcp.size()) now = min(now, lcp[i]);</pre>
                                                                             return s.substr(sa[i], prePrefix+k-nowRank);
                  }else{
                      if (i<lcp.size()) now = lcp[i];</pre>
101
                                                                         prePrefix = sa.lcp[i];
              }
102
                                                                         nowRank += add;
103
          { // build suf
104
              int now = 0;
105
              for (int i=s.size()-1; i>=0; i--){
106
107
                  if (sa[i]<=p){</pre>
108
                      suf[sa[i]] = now;
                      if (i-1>=0) now = min(now, lcp[i-1]);
                      if (i-1>=0) now = lcp[i-1];
113
              }
115
          return {pre, suf};
118 };
```