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#### 1 Misc

#### **1.1** Note

開始寫題目之前,請做下面的事:

- 在「開始寫任何題目之前」.應該要先自己看過「所有」範例測資的邏輯和 演算法是否有跟範例輸出對上
- 如果你覺得別人的某段程式碼有錯誤,就應該直接講出來
- +2~+3後就開始生測資跟對拍(根據寫 generator 跟 checker 的時間決定) 寫程式請遵照以下原則:

• 準確使用註解分段程式碼

- declare
- init
- input
- process / queries
- output
- 陣列若可以開到最大,則使用常數宣告大小

上傳之前,請依序檢查以下資訊:

- 1. 是否開啟 IO 優化
- 2. 是否有 t 筆輸入但忘了輸入
- 3. 是否有初始化容器
- 4. 題目範圍有沒有開到最大
- 5. 跑過所有範例測資,並嚴格確認是否正確

#### 1.2 Default Code [481325]

```
#include <bits/stdc++.h>
using namespace std;
#define int long long
#define debug(HEHE...) std::cerr << #HEHE << " = ", dout(HEHE
)

void dout() { std::cerr << "\n"; }
template <typename T, typename... U>
void dout(T t, U... u) { std::cerr << t << ' ', dout(u...); }

void solve(){

signed main(){
ios::sync_with_stdio(0), cin.tie(0);

int t = 1;
while (t--){
    solve();
}

return 0;
}</pre>
```

#### 1.3 Run

```
from os import *
f = "pA"
while 1:
    i = input("input: ")
    system("clear")
    p = listdir(".")
    if i != "":
        f = i
    print(f"file = {f}")
    if system(f"g++ {f}.cpp -std=c++17 -Wall -Wextra -Wshadow
          -O2 -D LOCAL -q -fsanitize=undefined,address -o {f}
        print("CE")
        continue
    for x in sorted(p):
        if f in x and ".in" in x:
            print(x)
            if system(f"./\{f\} < \{x\}"):
                print("RE")
            print()
```

### 1.4 Custom Set PQ Sort [2892de]

```
// priority_queue · 務必檢查相等的 case · 給所有元素一個排序的
依據
struct cmp{
  bool operator () (Data a, Data b){
    return a.x<b.x;
  }
};
priority_queue<Data, vector<Data>, cmp> pq;
// set · 務必檢查相等的 case · 給所有元素一個排序的依據
uauto cmp = [](int a, int b) {
  return a > b;
};
set<int, decltype(cmp)> s = {1, 2, 3, 4, 5};
cout << *s.begin() << '\n';</pre>
```

# 1.5 Dynamic Bitset [c78aa8]

```
const int MAXN = 2e5 + 5;
template <int len = 1>
void solve(int n) {
   if (n > len) {
      solve<min(len*2, MAXN)>(n);
      return;
   }
   bitset<len> a;
   }
}
```

#### 1.6 Enumerate Subset [a13e46]

```
1 // 時間複雜度 O(3^n)
2 // 枚舉每個 mask 的子集
3 for (int mask=0; mask<(1<<n); mask++){
    for (int s=mask; s>=0; s=(s-1)&m){
        // s 是 mask 的子集
        if (s==0) break;
    }
}
```

#### 1.7 Fast Input [6f8879]

```
1 // fast IO
  // 6f8879
  inline char readchar(){
      static char buffer[BUFSIZ], * now = buffer + BUFSIZ, *
           end = buffer + BUFSIZ;
      if (now == end)
           if (end < buffer + BUFSIZ)</pre>
               return EOF;
          end = (buffer + fread(buffer, 1, BUFSIZ, stdin));
          now = buffer;
      return *now++;
  inline int nextint(){
      int x = 0, c = readchar(), neg = false;
      while(('0' > c | | c > '9') && c!='-' && c!=EOF) c =
           readchar();
      if(c == '-') neg = true, c = readchar();
      while ('0' \le c \&\& c \le '9') x = (x << 3) + (x << 1) + (c^'0')
            , c = readchar();
      if(neg) x = -x;
      return x; // returns 0 if EOF
20
21 }
```

### 1.8 OEIS [f915c2]

```
1// 若一個線性遞迴有 k 項,給他恰好 2*k 個項可以求出線性遞迴
  // f915c2
  template <typename T>
  vector<T> BerlekampMassey(vector<T> a) {
      auto scalarProduct = [](vector<T> v, T c) {
          for (T &x: v) x *= c;
          return v;
      };
      vector<T> s, best;
      int bestPos = 0;
      for (size t i = 0; i < a.size(); i++) {</pre>
          T error = a[i];
          for (size_t j = 0; j < s.size(); j++) error -= s[j] *</pre>
               a[i-1-j];
          if (error == 0) continue;
          if (s.empty()) {
              s.resize(i + 1);
              bestPos = i;
17
```

```
best.push_back(1 / error);
continue;

vector<T> fix = scalarProduct(best, error);
fix.insert(fix.begin(), i - bestPos - 1, 0);
if (fix.size() >= s.size()) {
    best = scalarProduct(s, - 1 / error);
    best.insert(best.begin(), 1 / error);
    bestPos = i;
    s.resize(fix.size());
}

for (size_t j = 0; j < fix.size(); j++)
    s[j] += fix[j];
}
return s;
}
</pre>
```

# 1.9 Pragma [09d13e]

```
1 #pragma GCC optimize("03,unroll-loops")
2 #pragma GCC target("avx,avx2,sse,sse2,sse3,sse4,popcnt")
```

# 1.10 Xor Basis [840136]

```
vector<int> basis;
  void add vector(int x){
     for (auto v : basis){
         x=min(x, x^v);
     if (x) basis.push_back(x);
9 // 給一數字集合 S · 求能不能 XOR 出 x
10 bool check(int x){
     for (auto v : basis){
         x=min(x, x^v);
     return 0;
17 // 給一數字集合 S, 求能 XOR 出多少數字
18 // 答案等於 2^{basis 的大小}
20 / / 給一數字集合 S,求 XOR 出最大的數字
21 int get max(){
     for (auto v : basis){
         ans=max(ans, ans^v);
     return ans;
```

# 1.11 random int [9cc603]

#### 1.12 Python

```
1 # system setting
  sys.setrecursionlimit(100000)
  sys.set_int_max_str_digits(10000)
  from turtle import *
  N = 3000000010
  setworldcoordinates(-N, -N, N, N)
  hideturtle()
  speed(100)
  def draw_line(a, b, c, d):
      teleport(a, b)
      goto(c, d)
  def write dot(x, y, text, diff=1): # diff = 文字的偏移
      teleport(x, y)
      dot(5, "red")
      teleport(x+N/100*diff, y+N/100*diff)
      write(text, font=("Arial", 5, "bold"))
25 draw_line(*a[i], *(a[i-1]))
26 write_dot(*a[i], str(a[i]))
```

#### 1.13 diff

#### 1.14 hash command

#### 1.15 setup

```
| se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
  :inoremap " ""<Esc>i
  :inoremap {<CR> {<CR>}<Esc>ko
  :inoremap {{ {}}<ESC>i
  function! F(...)
   execute '!./%:r < ./' . a:1
  endfunction
  command! -nargs=* R call F(<f-args>)
  map <F7> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -Wextra -
       Wshadow -02 -DLOCAL -g -fsanitize=undefined,address<CR>
13 map <F8> :!./%:r<CR>
14 map <F9> :!./%:r < ./%:r.in<CR>
  ca hash w !cpp -dD -P -fpreprocessed \| tr -d "[:space:]" \|
       md5sum \| cut -c-6
  " i+<esc>25A---+<esc>
  " o|<esc>25A |<esc>
  " "ggVGyG35pGdd
```

#### 2 Convolution

# 2.1 FFT any mod [234f9e]

```
2 | 修改 const int MOD = 998244353 更改要取餘的數字
3 \mid PolyMul(a, b) 回傳多項式乘法的結果 ( c_k = \sum_{i=1}^{n} a_i + b_j
       mod MOD )
  大約可以支援 5e5 \cdot a_i, b_i 皆在 MOD 以下的非負整數
  const int MOD = 998244353;
  typedef complex<double> cd;
  // b9c90a
  void FFT(vector<cd> &a) {
      int n = a.size(), L = 31-__builtin_clz(n);
      vector<complex<long double>> R(2, 1);
      vector<cd> rt(2, 1);
      for (int k=2; k < n; k*=2){
          R.resize(n);
          rt.resize(n);
          auto x = polar(1.0L, acos(-1.0L) / k);
          for (int i=k ; i<2*k ; i++){</pre>
20
               rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);
21
22
23
      vector<int> rev(n);
      for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2] | (i&1) << L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
```

```
for (int k=1; k<n; k*=2){</pre>
        for (int i=0 ; i<n ; i+=2*k){</pre>
            for (int j=0; j<k; j++){</pre>
                 auto x = (double *)&rt[j+k];
                 auto y = (double *)&a[i+j+k];
                 cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]* 17
                     y[0]);
                 a[i+j+k] = a[i+j]-z;
                 a[i+j] += z;
        }
    return;
// d3c65e
vector<int> PolyMul(vector<int> a, vector<int> b){
   if (a.empty() || b.empty()) return {};
    vector<int> res(a.size()+b.size()-1);
   int B = 32- builtin clz(res.size()), n = (1<<B), cut =</pre>
         int(sqrt(MOD));
    vector<cd> L(n), R(n), outs(n), outl(n);
    for (int i=0 ; i<a.size() ; i++){</pre>
        L[i] = cd((int) a[i]/cut, (int)a[i]%cut);
    for (int i=0 ; i<b.size() ; i++){</pre>
        R[i] = cd((int) b[i]/cut, (int)b[i]%cut);
    FFT(L);
   FFT(R);
    for (int i=0 ; i<n ; i++){</pre>
        int j = -i\&(n-1);
        \operatorname{outl}[j] = (L[i] + \operatorname{conj}(L[j])) * R[i]/(2.0*n);
        outs[j] = (L[i]-conj(L[j])) * R[i]/(2.0*n)/1i;
    FFT(outl);
   FFT(outs);
    for (int i=0 ; i<res.size() ; i++){</pre>
        int av = (int)(real(outl[i])+0.5), cv = (int)(imag(
             outs[i])+0.5);
        int bv = (int)(imag(outl[i])+0.5) + (int)(real(outs[i
        res[i] = ((av%MOD*cut+bv) % MOD*cut+cv) % MOD;
    return res;
```

### 2.2 FFT new [c95bb8]

```
1 typedef complex < double > cd;
 // b9c90a
4 void FFT(vector<cd> &a) {
     int n = a.size(), L = 31-__builtin_clz(n);
     vector<complex<long double>> R(2, 1);
     vector<cd> rt(2, 1);
     for (int k=2 ; k<n ; k*=2){</pre>
          R.resize(n);
         rt.resize(n);
```

# 2.3 FFT short [70c01a]

auto x = polar(1.0L, acos(-1.0L) / k);

rt[i] = R[i] = (i&1 ? R[i/2]\*x : R[i/2]);

for (int i=k ; i<2\*k ; i++){</pre>

rev[i] = (rev[i/2] | (i&1) << L)/2;

for (int i=0; i<n; i+=2\*k){

a[i+j] += z;

if (a.empty() || b.empty()) return {};

copy(a.begin(), a.end(), begin(in)); for (int i=0 ; i<b.size() ; i++){</pre>

for (int i=0 ; i<res.size() ; i++){</pre>

res[i] = imag(out[i]) / (4 \* n);

vector<cd> in(n), out(n);

in[i].imag(b[i]);

for (cd& x : in) x \*= x;

for (int i=0 ; i<n ; i++){</pre>

vector < double > res(a.size()+b.size()-1);

if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>

for (int j=0 ; j<k ; j++){</pre>

a[i+j+k] = a[i+j]-z;

auto x = (double \*)&rt[j+k];

auto y = (double \*)&a[i+j+k];

vector<double> PolyMul(const vector<double> a, const vector<</pre>

int L = 32 - \_\_builtin\_clz(res.size()), n = 1 << L;</pre>

out[i] = in[-i & (n - 1)] - conj(in[i]);

}

}

return;

double> b){

FFT(in);

FFT(out);

return res;

vector<int> rev(n); for (int i=0 ; i<n ; i++){</pre>

for (int i=0 ; i<n ; i++){</pre>

for (int k=1; k<n; k\*=2){

```
1 #define int long long
 using Cplx = complex<double>;
 const double pi = acos(-1);
 const int mod = 998244353, g = 3;
 int power(int a, int b) {
     int res = 1;
```

```
while (b) {
                                                           if (b & 1) res = res * a % mod;
                                                           a = a * a % mod:
                                                           b >>= 1;
                                                       return res;
                                                  int inv(int x) { return power(x, mod - 2); }
                                                   // FFT use Cplx, NTT use ll
                                                   void FFT(vector<int> &a, int n, int op) {
                                                       // n must be 2^k
                                                       vector<int> R(n);
                                                       FOR (i, 0, n - 1)
                                                           R[i] = R[i/2]/2 + (i&1)*(n/2);
                                                       FOR (i, 0, n - 1)
                                                           if (i < R[i]) swap(a[i], a[R[i]]);</pre>
                                                       for (int m = 2; m <= n; m *= 2) {</pre>
                                                           // Cplx w1({cos(2*pi/m), sin(2*pi/m)*op});
cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]* 26
                                                           int w1 = power(g, (mod-1)/m * op + mod-1);
                                                           for (int i = 0; i < n; i += m) {</pre>
                                                               // Cplx wk({1, 0});
                                                               int wk = 1;
                                                               FOR (k, 0, m / 2 - 1) {
                                                                   auto x = a[i+k], y = a[i+k+m/2] * wk % mod;
                                                                   a[i+k] = (x+y) \% mod;
                                                                   a[i+k+m/2] = (x-y+mod) \% mod;
                                                                   w\bar{k} = wk * w1 \% mod;
                                                       if (op == -1)
                                                           FOR (i, 0, n - 1) {
                                                               // a[i] = a[i] / n;
                                                41
                                                               a[i] = a[i] * inv(n) % mod;
                                                42
                                                43
```

### 2.4 FWT [832aa5]

```
│// 已經把 mint 刪掉‧需要增加註解
  vector<int> xor_convolution(vector<int> a, vector<int> b, int
      if (k == 0) {
          return vector<int>{a[0] * b[0]};
      vector<int> aa(1 << (k - 1)), bb(1 << (k - 1));
      FOR (i, 0, (1 << (k - 1)) - 1) {
          aa[i] = a[i] + a[i + (1 << (k - 1))];
          bb[i] = b[i] + b[i + (1 << (k - 1))];
      vector<int> X = xor convolution(aa, bb, k - 1);
      FOR (i, 0, (1 << (k - 1)) - 1) {
          aa[i] = a[i] - a[i + (1 << (k - 1))];
          bb[i] = b[i] - b[i + (1 << (k - 1))];
      vector<int> Y = xor convolution(aa, bb, k - 1);
      vector<int> c(1 << k);</pre>
      FOR (i, 0, (1 << (k - 1)) - 1) {
                              ] = (X[i] + Y[i]) / 2;
          c[i + (1 << (k - 1))] = (X[i] - Y[i]) / 2;
20
21
      return c;
```

#### 2.5 Min Convolution Concave Concave [ffb28d]

```
1 // 需要增加註解
2 // min convolution
3 vector<int> mkk(vector<int> a, vector<int> b) {
4 vector<int> slope;
5 FOR (i, 1, ssize(a) - 1) slope.pb(a[i] - a[i - 1]);
6 FOR (i, 1, ssize(b) - 1) slope.pb(b[i] - b[i - 1]);
7 sort(all(slope));
8 slope.insert(begin(slope), a[0] + b[0]);
9 partial_sum(all(slope), begin(slope));
1 return slope;
1 }
```

 $| \text{const int MOD} = (119 \iff 23) + 1, ROOT = 62; // = 998244353$ 

# 2.6 NTT mod 998244353 [5c6335]

```
2 // For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 <<
  // and 483 << 21 (same root). The last two are > 10^9.
5 // 9cd58a
 void NTT(vector<int> &a) {
      int n = a.size();
      int L = 31- builtin clz(n);
      vector<int> rt(2, 1);
      for (int k=2, s=2; k< n; k*=2, s++){
          rt.resize(n);
          int z[] = {1, qp(ROOT, MOD>>s)};
          for (int i=k ; i<2*k ; i++){</pre>
              rt[i] = rt[i/2]*z[i&1]%MOD;
      vector<int> rev(n);
      for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2]|(i&1)<<L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]){</pre>
              swap(a[i], a[rev[i]]);
      }
      for (int k=1 ; k<n ; k*=2){</pre>
          for (int i=0 ; i<n ; i+=2*k){</pre>
              for (int j=0 ; j<k ; j++){</pre>
                   int z = rt[j+k]*a[i+j+k]%MOD, &ai = a[i+j];
                   a[i+j+k] = ai-z+(z>ai ? MOD : 0);
                   ai += (ai+z)=MOD ? z-MOD : z);
  // 0b0e99
40 vector<int> polyMul(vector<int> &a, vector<int> &b){
      if (a.empty() || b.empty()) return {};
      int s = a.size()+b.size()-1, B = 32-__builtin_clz(s), n =
      int inv = qp(n, MOD-2);
```

```
vector<int> L(a), R(b), out(n);
L.resize(n), R.resize(n);
NTT(L), NTT(R);
for (int i=0; i<n; i++){
    out[-i&(n-1)] = L[i]*R[i]%MOD*inv%MOD;
}
NTT(out);
out.resize(s);
return out;
}</pre>
```

#### 3 Data-Structure

# 3.1 BIT [7ef3a9]

# 3.2 Disjoint Set Persistent [447002]

```
struct Persistent_Disjoint_Set{
    Persistent_Segment_Tree arr, sz;
    void init(int n){
        arr.init(n);
        vector<int> v1;
        for (int i=0 ; i<n ; i++){</pre>
            v1.push back(i);
        arr.build(v1, 0);
        sz.init(n);
        vector<int> v2;
        for (int i=0 ; i<n ; i++){</pre>
            v2.push back(1);
        sz.build(v2, 0);
    int find(int a){
        int res = arr.query_version(a, a+1, arr.version.size
             ()-1).val;
        if (res==a) return a;
        return find(res);
    bool unite(int a, int b){
        a = find(a);
        b = find(b);
```

```
if (a!=b){
               int sz1 = sz.query_version(a, a+1, arr.version.
32
                   size()-1).val;
               int sz2 = sz.query version(b, b+1, arr.version.
                   size()-1).val;
              if (sz1<sz2){
                  arr.update_version(a, b, arr.version.size()
                  sz.update_version(b, sz1+sz2, arr.version.
                       size()-1);
                  arr.update_version(b, a, arr.version.size()
                  sz.update_version(a, sz1+sz2, arr.version.
                       size()-1);
42
              return true;
43
          return false;
44
45
46
```

# 3.3 PBDS GP Hash Table [866cf6]

```
i #include <ext/pb ds/assoc container.hpp>
  using namespace __gnu_pbds;
  typedef tree<int, null type, less<int>, rb tree tag,
       tree order statistics node update> order set;
  struct custom hash {
      static uint64 t splitmix64(uint64 t x) {
          // http://xorshift.di.unimi.it/splitmix64.c
          x += 0x9e3779b97f4a7c15;
          x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
          x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
          return x ^ (x >> 31);
      size t operator()(uint64 t x) const {
          static const uint64 t FIXED RANDOM = chrono::
               steady_clock::now().time_since_epoch().count();
          return splitmix64(x + FIXED RANDOM);
15
16
17
  };
19 gp_hash_table<int, int, custom_hash> ss;
```

### 3.4 PBDS Order Set [231774]

```
9 using namespace __gnu_pbds;
10 typedef tree<int, null type, less<int>, rb tree tag,
       tree order statistics node update> order set;
```

#### 3.5 Segment Tree Add Set [bb1898]

```
1 // [ll, rr), based-0
2 // 使用前記得 init(陣列大小), build(陣列名稱)
3 // add(LL, rr): 區間修改
4 // set(ll, rr): 區間賦值
5 // query(ll, rr): 區間求和 / 求最大值
6 struct SegmentTree{
     struct node{
         int add tag = 0;
         int set tag = 0;
         int sum = 0;
         int ma = 0;
     };
     vector<node> arr;
     SegmentTree(int n){
         arr.resize(n<<2);</pre>
     node pull(node A, node B){
         C.sum = A.sum + B.sum:
         C.ma = max(A.ma, B.ma);
         return C;
      // cce0c8
     void push(int idx, int ll, int rr){
         if (arr[idx].set tag!=0){
             arr[idx].sum = (rr-11)*arr[idx].set tag;
             arr[idx].ma = arr[idx].set tag;
             if (rr-ll>1){
                 arr[idx*2+1].add tag = 0;
                 arr[idx*2+1].set_tag = arr[idx].set_tag;
                 arr[idx*2+2].add tag = 0;
                 arr[idx*2+2].set tag = arr[idx].set tag;
             arr[idx].set tag = 0;
         if (arr[idx].add tag!=0){
             arr[idx].sum += (rr-ll)*arr[idx].add tag;
             arr[idx].ma += arr[idx].add tag;
                 arr[idx*2+1].add tag += arr[idx].add tag;
                 arr[idx*2+2].add_tag += arr[idx].add_tag;
             arr[idx].add_tag = 0;
     void build(vector<int> &v, int idx = 0, int ll = 0, int
          rr = n){
         if (rr-ll==1){
             arr[idx].sum = v[11]:
             arr[idx].ma = v[ll];
         }else{
```

```
int mid = (11+rr)/2;
               build(v, idx*2+1, ll, mid);
               build(v, idx*2+2, mid, rr);
               arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
      }
       void add(int ql, int qr, int val, int idx = 0, int ll =
            0, int rr =n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return;</pre>
           if (ql<=ll && rr<=qr){
               arr[idx].add_tag += val;
               push(idx, 11, rr);
               return;
           int mid = (11+rr)/2;
           add(ql, qr, val, idx*2+1, ll, mid);
           add(ql, qr, val, idx*2+2, mid, rr);
           arr[idx]=pull(arr[idx*2+1], arr[idx*2+2]);
       void set(int ql, int qr, int val, int idx=0, int ll=0,
            int rr=n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return;</pre>
           if (q1<=11 && rr<=qr){
               arr[idx].add tag = 0;
               arr[idx].set_tag = val;
               push(idx, ll, rr);
               return;
           int mid = (11+rr)/2;
           set(ql, qr, val, idx*2+1, ll, mid);
           set(ql, qr, val, idx*2+2, mid, rr);
           arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
       node query(int ql, int qr, int idx = 0, int ll = 0, int
            rr = n){
           push(idx, 11, rr);
           if (rr<=ql || qr<=ll) return node();</pre>
           if (q1<=11 && rr<=qr) return arr[idx];</pre>
           int mid = (11+rr)/2;
           return pull(query(ql, qr, idx*2+1, ll, mid), query(ql
                , qr, idx*2+2, mid, rr));
100 } ST;
```

### 3.6 Segment Tree Li Chao Line [45b8ba]

```
全部都是 0-based
LC_Segment_Tree st(n);
update(\{a, b\}):插入一條 y=ax+b 的全域直線
query(x): 查詢所有直線在位置 x 的最小值
```

```
vector<Node> arr;
      LC Segment Tree(int n = 0){
          arr.resize(4*n);
      void update(Node val, int idx = 0, int ll = 0, int rr =
          if (rr-ll==0) return;
          if (rr-ll==1){
              if (val.y(ll)<arr[idx].y(ll)){</pre>
                  arr[idx] = val;
              return;
          int mid = (11+rr)/2;
          if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
               的線斜率要比較小
          if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
              update(val, idx*2+1, ll, mid);
          }else{ // 交點在右邊
              swap(arr[idx], val); // 在左子樹中,新線比舊線還
              update(val, idx*2+2, mid, rr);
          return;
      }
      int query(int x, int idx = 0, int 11 = 0, int rr = MAX_V)
          if (rr-ll==0) return INF;
          if (rr-ll==1){
              return arr[idx].y(ll);
          int mid = (11+rr)/2;
          if (x<mid){</pre>
              return min(arr[idx].y(x), query(x, idx*2+1, ll,
              return min(arr[idx].y(x), query(x, idx*2+2, mid,
                  rr));
61
```

return a\*x+b;

struct Node{ // y = ax+bint a = 0;

int b = INF;

int v(int x){

struct LC Segment Tree{

20

27

};

#### 3.7 Segment Tree Li Chao Segment [2cb0a4]

```
2|全部都是 0-based
```

60

```
LC Segment Tree st(n);
s|update_segment({a, b}, ql, qr):在 [ql, qr) 插入一條 y=ax+b
9 | querv(x): 查詢所有直線在位置 x 的最小值
||| const int MAX V = 1e6+10; // 值域最大值
13 struct LC Segment Tree{
     struct Node{ // y = ax+b
         int a = 0;
         int b = INF;
         int y(int x){
             return a*x+b;
     };
     vector<Node> arr;
     LC Segment Tree(int n = 0){
         arr.resize(4*n):
     void update(Node val, int idx = 0, int ll = 0, int rr =
         if (rr-ll==0) return;
         if (rr-ll<=1){
             if (val.y(ll)<arr[idx].y(ll)){</pre>
                 arr[idx] = val;
             return:
         }
         int mid = (11+rr)/2;
         if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
              的線斜率要比較小
         if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
             update(val, idx*2+1, ll, mid);
         }else{ // 交點在右邊
             swap(arr[idx], val); // 在左子樹中·新線比舊線還
             update(val, idx*2+2, mid, rr);
         return:
     }
     // 在 [ql, qr) 加上一條 val 的線段
     void update segment(Node val, int ql, int qr, int idx =
          0, int 11 = 0, int rr = MAX_V){
         if (rr-ll==0) return;
         if (rr<=ql || qr<=ll) return;</pre>
         if (q1<=11 && rr<=qr){
             update(val, idx, ll, rr);
             return:
         int mid = (11+rr)/2;
         update segment(val, ql, qr, idx*2+1, ll, mid);
         update segment(val, ql, qr, idx*2+2, mid, rr);
          return;
```

```
62
63
    int query(int x, int idx = 0, int ll = 0, int rr = MAX_V)
64
    if (rr-ll==0) return INF;
65
    if (rr-ll==1){
        return arr[idx].y(ll);
    }
68
    int mid = (ll+rr)/2;
69
    int mid = (ll+rr)/2;
69
    if (xxmid){
        return min(arr[idx].y(x), query(x, idx*2+1, ll, sl mid));
    }
71
    return min(arr[idx].y(x), query(x, idx*2+2, mid, sl mr));
72
    return min(arr[idx].y(x), query(x, idx*2+2, mid, sl mr));
73
    return min(arr[idx].y(x), query(x, idx*2+2, mid, sl mr));
74
    }
75
    }
76
}
```

# 3.8 Segment Tree Persistent [3b5aa9]

```
全部都是 0-based
  Persistent_Segment_Tree st(n+q);
  st.build(v, 0);
  函式:
  update version(pos, val, ver): 對版本 ver 的 pos 位置改成 val 73
  query version(ql, qr, ver): 對版本 ver 查詢 [ql, qr) 的區間和 74
  clone version(ver):複製版本 ver 到最新的版本
  struct Persistent_Segment_Tree{
                                                               77
      int node cnt = 0;
                                                               78
      struct Node{
          int lc = -1;
          int rc = -1;
          int val = 0;
      };
      vector<Node> arr;
      vector<int> version;
      Persistent Segment Tree(int sz){
          arr.resize(32*sz);
          version.push back(node cnt++);
      void pull(Node &c, Node a, Node b){
          c.val = a.val+b.val;
          return:
      void build(vector<int> &v, int idx, int ll = 0, int rr =
          auto &now = arr[idx];
37
          if (rr-ll==1){
                                                              100
              now.val = v[11];
                                                              101 };
              return;
```

```
int mid = (11+rr)/2;
    now.lc = node cnt++:
    now.rc = node cnt++;
    build(v, now.lc, ll, mid);
    build(v, now.rc, mid, rr);
    pull(now, arr[now.lc], arr[now.rc]);
    return:
void update(int pos, int val, int idx, int ll = 0, int rr
    auto &now = arr[idx];
    if (rr-ll==1){
        now.val = val:
        return:
    int mid = (11+rr)/2;
    if (pos<mid){</pre>
        arr[node cnt] = arr[now.lc];
        now.lc = node cnt;
        node cnt++:
        update(pos, val, now.lc, ll, mid);
        arr[node cnt] = arr[now.rc];
        now.rc = node cnt;
        node cnt++;
        update(pos, val, now.rc, mid, rr);
    pull(now, arr[now.lc], arr[now.rc]);
    return:
void update version(int pos, int val, int ver){
    update(pos, val, version[ver]);
Node query(int ql, int qr, int idx, int ll = 0, int rr =
    auto &now = arr[idx];
    if (q1<=11 && rr<=qr) return now;</pre>
    if (rr<=ql || qr<=ll) return Node();</pre>
    int mid = (11+rr)/2;
    pull(ret, query(ql, qr, now.lc, ll, mid), query(ql,
         gr, now.rc, mid, rr));
    return ret;
Node query version(int ql, int qr, int ver){
    return query(ql, qr, version[ver]);
void clone version(int ver){
    version.push back(node cnt);
    arr[node cnt] = arr[version[ver]];
    node cnt++;
```

#### **3.9** Sparse Table [31f22a]

# 3.10 Treap [5851f5]

1 struct Treap{

```
Treap *1 = nullptr, *r = nullptr;
      int pri = rand(), val = 0, sz = 1;
      Treap(int _val){
          val = _val;
  };
int size(Treap *t){return t ? t->sz : 0;}
void pull(Treap *t){
      t->sz = size(t->1)+size(t->r)+1;
15 Treap* merge(Treap *a, Treap *b){
      if (!a || !b) return a ? a : b;
      if (a->pri>b->pri){
          a \rightarrow r = merge(a \rightarrow r, b);
          pull(a);
          return a;
          b->1 = merge(a, b->1);
          pull(b):
          return b;
30 | pair<Treap*, Treap*> split(Treap *&t, int k){ // 1-based <前
      k 個元素, 其他元素>
      if (!t) return {};
      if (size(t->1)>=k){
          auto pa = split(t->1, k);
```

```
t->1 = pa.second;
        pull(t);
        return {pa.first, t};
        auto pa = split(t->r, k-size(t->l)-1);
        t->r = pa.first;
        pull(t);
        return {t, pa.second};
// functions
Treap* build(vector<int> v){
    Treap* ret = nullptr;
    for (int i=0 ; i<v.size() ; i++){</pre>
        ret = merge(ret, new Treap(v[i]));
    return ret;
array<Treap*, 3> cut(Treap *t, int 1, int r){ // 1-based <前
    1~l-1 個元素, l~r 個元素, r+1 個元素>
    array<Treap*, 3> ret;
    tie(ret[1], ret[2]) = split(t, r);
    tie(ret[0], ret[1]) = split(ret[1], 1-1);
    return ret;
void print(Treap *t, bool flag = true){
    if (t->1!=0) print(t->1, false);
    cout << t->val;
    if (t->r!=0) print(t->r, false);
    if (flag) cout << endl;</pre>
```

# 3.11 Treap2 [1bf328]

```
ı | // 1-based · 請注意 MAX_N 是否足夠大
2 int root = 0;
  int lc[MAX_N], rc[MAX_N];
  int pri[MAX_N], val[MAX_N];
  int sz[MAX_N], tag[MAX_N], fa[MAX_N];
  int new_node(int v){
      static int nodeCnt = 0:
      nodeCnt++;
      val[nodeCnt] = v;
      sz[nodeCnt] = 1;
      pri[nodeCnt] = rand();
      return nodeCnt;
  void push(int x){
      if (tag[x]){
          if (lc[x]) tag[lc[x]] ^= 1;
          if (rc[x]) tag[rc[x]] ^= 1;
      tag[x] = 0;
22 int pull(int x){
      if (x){
          fa[x] = 0;
          sz[x] = 1+sz[lc[x]]+sz[rc[x]];
```

```
if (lc[x]) fa[lc[x]] = x;
          if (rc[x]) fa[rc[x]] = x;
29
      return x;
30 }
  int merge(int a, int b){
      if (!a or !b) return a|b;
      push(a), push(b);
      if (pri[a]>pri[b]){
          rc[a] = merge(rc[a], b);
          return pull(a);
          lc[b] = merge(a, lc[b]);
          return pull(b);
43
|45| // [1, k] [k+1, n]
  void split(int x, int k, int &a, int &b) {
      if (!x) return a = b = 0, void();
      if (sz[lc[x]] >= k) {
          split(lc[x], k, a, lc[x]);
          b = x;
          pull(a); pull(b);
          split(rc[x], k - sz[lc[x]] - 1, rc[x], b);
          pull(a); pull(b);
58
  // functions
  // 回傳 x 在 Treap 中的位置
  int get pos(int x){
      vector<int> sta;
      while (fa[x]){
          sta.push back(x);
          x = fa[x];
      while (sta.size()){
          push(x);
          x = sta.back();
          sta.pop_back();
      push(x);
      int res = sz[x] - sz[rc[x]];
      while (fa[x]){
          if (rc[fa[x]]==x){
              res += sz[fa[x]]-sz[x];
          x = fa[x];
81
82
      return res;
85 // 1-based <前 [1, L-1] 個元素, [L, r] 個元素, [r+1, n] 個元
86 array<int, 3> cut(int x, int 1, int r){
      array<int, 3> ret;
      split(x, r, ret[1], ret[2]);
      split(ret[1], 1-1, ret[0], ret[1]);
```

#### 3.12 Trie [b6475c]

```
1 struct Trie{
      struct Data{
          int nxt[2]={0, 0};
      };
      int sz=0;
      vector<Data> arr;
      void init(int n){
          arr.resize(n);
      void insert(int n){
          int now=0;
          for (int i=N; i>=0; i--){
              int v=(n>>i)&1;
              if (!arr[now].nxt[v]){
                  arr[now].nxt[v]=++sz;
              now=arr[now].nxt[v];
      }
      int query(int n){
          int now=0, ret=0;
          for (int i=N ; i>=0 ; i--){
              int v=(n>>i)&1;
              if (arr[now].nxt[1-v]){
                  ret+=(1<<i);
                  now=arr[now].nxt[1-v];
              }else if (arr[now].nxt[v]){
                  now=arr[now].nxt[v];
              }else{
                  return ret;
          return ret;
40 } tr;
```

# 4 Dynamic-Programming

# 4.1 Digit DP [133f00]

```
1 #include <bits/stdc++.h>
 using namespace std;
 long long 1, r;
 long long dp[20][10][2][2]; // dp[pos][pre][limit] = 後 pos
      位 pos 前一位是 pre (是/否)有上界 (是/否)有前綴零
 long long memorize_search(string &s, int pos, int pre, bool
     limit, bool lead){
     // 已經被找過了,直接回傳值
     if (dp[pos][pre][limit][lead]!=-1) return dp[pos][pre][
         limit][lead];
     // 已經搜尋完畢,紀錄答案並回傳
     if (pos==(int)s.size()){
         return dp[pos][pre][limit][lead] = 1;
     // 枚舉目前的位數數字是多少
     long long ans = 0;
     for (int now=0 ; now<=(limit ? s[pos]-'0' : 9) ; now++){</pre>
        if (now==pre){
            // 1~9 絕對不能連續出現
            if (pre!=0) continue;
            // 如果已經不在前綴零的範圍內·Ø 不能連續出現
            if (lead==false) continue;
        ans += memorize_search(s, pos+1, now, limit&(now==(s[
             pos]-'0')), lead&(now==0));
     // 已經搜尋完畢,紀錄答案並回傳
     return dp[pos][pre][limit][lead] = ans;
 // 回傳 [0, n] 有多少數字符合條件
 long long find_answer(long long n){
     memset(dp, -1, sizeof(dp));
     string tmp = to string(n);
     return memorize_search(tmp, 0, 0, true, true);
 int main(){
     // input
     cin >> 1 >> r;
     // output - 計算 [L, r] 有多少數字任意兩個位數都不相同
     cout << find_answer(r)-find_answer(l-1) << "\n";</pre>
     return 0;
```

### 4.2 Knaspack On Tree [df69b1]

```
11// 需要重構、需要增加註解
  #include <bits/stdc++.h>
  #define F first
  #define S second
  #define all(x) begin(x), end(x)
  using namespace std;
  #define chmax(a, b) (a) = (a) < (b) ? (b) : (a)
  #define chmin(a, b) (a) = (a) < (b) ? (a) : (b)
  #define 11 long long
  #define FOR(i, a, b) for (int i = a; i <= b; i++)
  int N, W, cur;
  vector<int> w, v, sz;
  vector<vector<int>> adj, dp;
  void dfs(int x) {
      sz[x] = 1;
      for (int i : adj[x]) dfs(i), sz[x] += sz[i];
      cur++;
      // choose x
      for (int i=w[x]; i<=W; i++){</pre>
          dp[cur][i] = dp[cur - 1][i - w[x]] + v[x];
26
      // not choose x
      for (int i=0 ; i<=W ; i++){</pre>
           chmax(dp[cur][i], dp[cur - sz[x]][i]);
30
31
  signed main() {
      cin >> N >> W;
      adj.resize(N + 1);
      w.assign(N + 1, 0);
      v.assign(N + 1, 0);
      sz.assign(N + 1, 0);
      dp.assign(N + 2, vector < int > (W + 1, 0));
      for (int i=1; i<=N; i++){</pre>
           int p; cin >> p;
           adj[p].push_back(i);
      for (int i=1; i<=N; i++) cin >> w[i];
      for (int i=1; i<=N; i++) cin >> v[i];
      cout \langle\langle dp[N + 1][W] \langle\langle ' \rangle n';
```

#### 4.3 SOS DP [8dfa8b]

```
1 // 總時間複雜度為 O(n 2^n)
2 // 計算 dp[i] = i 所有 bit mask 子集的和
3 for (int i=0; i<n; i++){
4 for (int mask=0; mask<(1<<n); mask++){
5 if ((mask>:i)&1){
6 dp[mask] += dp[mask^(1<<ii)];
7 }
8 }
9 }</pre>
```

# **Integer Partition**

```
dp[i][x] = 要將整數 x 拆成 i 堆的「組合數」
dp[i+1][x+1]+=dp[i][x] ( 創造新的一堆 )
dp[i][x+i]+=dp[i][x] (把每一堆都增加 1)
```

# Geometry

#### **5.1** Geometry Struct [7e8017]

```
using ld = double;
3 // 判斷數值正負: {1:正數,0:零,-1:負數}
4 int sign(long long x) {return (x \ge 0) ? ((bool)x) : -1; }
s int sign(ld x) {return (abs(x) < 1e-9) ? 0 : (x>0 ? 1 : -1);}
 template<typename T>
 struct point {
    T x, y;
     point() {}
     point(const T &x, const T &y) : x(x), y(y) {}
     explicit operator point<ld>() {return point<ld>(x, y); }
     point operator+(point b) {return {x+b.x, y+b.y}; }
     point operator-(point b) {return {x-b.x, y-b.y}; }
     point operator*(T b) {return {x*b, y*b}; }
     point operator/(T b) {return {x/b, y/b}; }
     bool operator==(point b) {return x==b.x && y==b.y; }
     T operator*(point b) {return x * b.x + y * b.y; }
     T operator^(point b) {return x * b.y - y * b.x; }
     // 逆時針極角排序
     bool side() { return (y == 0) ? (x > 0) : (y < 0); }
     bool operator<(point &b) {</pre>
         return side() == b.side() ?
            (x*b.y > b.x*y) : side() < b.side();
     friend ostream& operator<<(ostream& os, point p) {</pre>
         return os << "(" << p.x << ", " << p.y << ")";
     // 判斷 ab 到 ac 的方向: {1:逆時鐘.0:重疊.-1:順時鐘}
     friend int ori(point a, point b, point c) {
         return sign((b-a)^(c-a));
     friend bool btw(point a, point b, point c) {
         return ori(a, b, c) == 0 && sign((a-c)*(b-c)) <= 0;
     // 判斷線段 ab. cd 是否相交
     friend bool banana(point a, point b, point c, point d) {
         if (btw(a, b, c) || btw(a, b, d)
             || btw(c, d, a) || btw(c, d, b)) return true;
         int u = ori(a, b, c) * ori(a, b, d);
                                                             108
         int v = ori(c, d, a) * ori(c, d, b);
         return u < 0 && v < 0;
                                                             111
     // 旋轉 Arg(b) 的角度(小心溢位)
     point rotate(point b){return {x*b.x-y*b.y, x*b.y+y*b.x};} 113
     // 回傳極座標角度,值域:[-π, +π]
```

```
friend ld Arg(point b) {
      friend T abs2(point b) {return b * b; }
54 };
                                                                 119
                                                                 120
  template<typename T>
                                                                 121
  struct line {
                                                                 122
      point<T> p1, p2;
                                                                 123
      // ax + by + c = 0
                                                                 124
      T a, b, c; //|a|, |b| \le 2C, |c| \le 8C^2
                                                                 125
      line() {}
      line(const point<T> &x, const point<T> &y) : p1(x), p2(y){ _{127}
          build();
      void build() {
          a = p1.y - p2.y;
          b = p2.x - p1.x;
          c = (-a*p1.x)-b*p1.y;
      // 判斷點和有向直線的關係: {1:左邊,0:在線上,-1:右邊}
      int ori(point<T> &p) {
          return sign((p2-p1) ^ (p-p1));
73
      // 判斷直線斜率是否相同
      bool parallel(line &1) {
          return ((p1-p2) ^ (l.p1-l.p2)) == 0;
                                                                 141
      // 兩直線交點
                                                                 142
      point<ld> line intersection(line &1) {
                                                                 143
          using P = point<ld>;
                                                                 144
          point < T > u = p2-p1, v = 1.p2-1.p1, s = 1.p1-p1;
                                                                 145
          return P(p1) + P(u) * ((ld(s^v)) / (u^v));
                                                                 147
  };
                                                                 148
                                                                 149
  template<typename T>
  struct polygon {
      vector<point<T>> v;
      polygon() {}
      polygon(const vector<point<T>> &u) : v(u) {}
                                                                 154
      // simple 為 true 的時候會回傳任意三點不共線的凸包
                                                                 155
      void make convex hull(int simple) {
                                                                 156
          auto cmp = [\bar{\&}](point<T> &p, point<T> &q) {
              return (p.x == q.x)? (p.y < q.y): (p.x < q.x);
                                                                 159
          simple = (bool)simple;
                                                                 160
          sort(v.begin(), v.end(), cmp);
          v.resize(unique(v.begin(), v.end()) - v.begin());
          /// 警告: v.size() 是 1 的時候會回傳空的凸包 ///
                                                                 163
          vector<point<T>> hull;
          for (int t = 0; t < 2; ++t){
                                                                 165
              int sz = hull.size();
              for (auto &i:v) {
                  while (hull.size() >= sz+2 && ori(hull[hull.
                        size()-2], hull.back(), i) < simple) {</pre>
                      hull.pop back();
                                                                 170
                                                                 171
                  hull.push back(i);
                                                                 172
                                                                 173
              hull.pop_back();
                                                                 174
              reverse(v.begin(), v.end());
                                                                 175
          swap(hull, v);
```

```
115 // 可以在有 n 個點的簡單多邊形內 · 用 O(n) 判斷一個點:
return (b.x != 0 || b.y != 0) ? atan2(b.y, b.x) : 0; 116 // {1: 在多邊形內, 0: 在多邊形上, -1: 在多邊形外}
                                                      int in_polygon(point<T> a){
                                                          const T MAX POS = 1e9 + 5; // [記得修改] 座標的最大值
                                                          point<T> pre = v.back(), b(MAX_POS, a.y + 1);
                                                          int cnt = 0:
                                                          for (auto &i:v) {
                                                              if (btw(pre, i, a)) return 0;
                                                              if (banana(a, b, pre, i)) cnt++;
                                                              pre = i;
                                                          return cnt%2 ? 1 : -1;
                                                130 1 / / / 警告:以下所有凸包專用的函式都只接受逆時針排序且任三點不
                                                       共線的凸包 ///
                                                |131| // 可以在有 n 個點的凸包內 n · 用 O(\log n) 判斷一個點 :
                                                132 // {1: 在凸包内, 0: 在凸包邊上, -1: 在凸包外}
                                                       int in convex(point<T> p) {
                                                          int n = v.size();
                                                          int a = ori(v[0], v[1], p), b = ori(v[0], v[n-1], p);
                                                          if (a < 0 \mid | b > 0) return -1;
                                                          if (btw(v[0], v[1], p)) return 0;
                                                          if (btw(v[0], v[n - 1], p)) return 0;
                                                          int l = 1, r = n - 1, mid:
                                                          while (1 + 1 < r) {
                                                              mid = (1 + r) >> 1;
                                                              if (ori(v[0], v[mid], p) >= 0) l = mid;
                                                              else r = mid:
                                                          int k = ori(v[1], v[r], p);
                                                          if (k <= 0) return k;</pre>
                                                          return 1;
                                                   // 凸包專用的環狀二分搜,回傳 0-based index
                                                      int cycle search(auto &f) {
                                                          int n = v.size(), l = 0, r = n;
                                                          if (n == 1) return 0;
                                                          bool rv = f(1, 0);
                                                          while (r - 1 > 1) {
                                                              int m = (1 + r) / 2;
                                                              if (f(0, m) ? rv: f(m, (m + 1) % n)) r = m;
                                                              else 1 = m;
                                                          return f(1, r % n) ? 1 : r % n;
                                                161 // 可以在有 n 個點的凸包內,用 O(Log n)判斷一條直線:
                                                162 | // {1: 穿過凸包, 0: 剛好切過凸包, -1: 沒碰到凸包}
                                                       int line cut convex(line<T> L) {
                                                          L.build();
                                                          point<T> p(L.a, L.b);
                                                          auto gt = [&](int neg) {
                                                              auto f = [\&](int x, int y) {
                                                                 return sign((v[x] - v[y]) * p) == neg;
                                                              return -(v[cycle_search(f)] * p);
                                                          T x = gt(1), y = gt(-1);
                                                          if (L.c < x || y < L.c) return -1;
                                                          return not (L.c == x || L.c == y);
                                                176 // 可以在有 n 個點的凸包內 · 用 O(Log n) 判斷一個線段:
                                                177 // {1: 存在一個凸包上的邊可以把這個線段切成兩半,
```

```
0:有碰到凸包但沒有任何凸包上的邊可以把它切成兩半。
179 // -1: 沒碰到凸包}
                                                                 239
180 /// 除非線段兩端點都不在凸包邊上,否則此函數回傳 0 的時候不一 240
        定表示線段沒有通過凸包內部 ///
       int segment across convex(line<T> L) {
                                                                242
182
           L.build():
                                                                243
           point<T> p(L.a, L.b);
183
           auto gt = [&](int neg) {
                                                                244
               auto f = [\&](int x, int y) {
186
                   return sign((v[x] - v[y]) * p) == neg;
              };
              return cycle_search(f);
                                                                246
                                                                247
                                                                248
190
           int i = gt(1), j = gt(-1), n = v.size();
                                                                240
191
          T x = -(v[i] * p), y = -(v[j] * p);
           if (L.c < x || y < L.c) return -1;
                                                                250
192
                                                                251
          if (L.c == x || L.c == y) return 0;
193
                                                                252
194
           if (i > j) swap(i, j);
195
196
           auto g = [&](int x, int lim) {
                                                                253
                                                                254
197
               int now = 0, nxt;
               for (int i = 1 << __lg(lim); i > 0; i /= 2) {
                                                                 255
198
                   if (now + i > lim) continue;
                                                                256
199
                   nxt = (x + i) % n;
                                                                 257
200
                                                                258
201
                   if (L.ori(v[x]) * L.ori(v[nxt]) >= 0) {
                                                                259
202
                       x = nxt;
203
                       now += i:
                                                                 260
                                                                 261
204
              } // ↓ BE CAREFUL
205
              return -(ori(v[x], v[(x + 1) % n], L.p1) * ori(v[ 263
                   x], v[(x + 1) % n], L.p2));
                                                                 266
           return max(g(i, j - i), g(j, n - (j - i)));
208
                                                                 267
210 | // 可以在有 n 個點的凸包內 · 用 O(Log n) 判斷一個線段:
                                                                 269
211 // {1: 線段上存在某一點位於凸包內部(邊上不算),
      0: 線段上存在某一點碰到凸包的邊但線段上任一點均不在凸包
213 // -1: 線段完全在凸包外面 }
                                                                272
       int segment_pass_convex_interior(line<T> L) {
                                                                273
           if (in convex(L.p1) == 1 || in convex(L.p2) == 1)
                return 1;
          L.build();
216
                                                                 276
217
           point<T> p(L.a, L.b);
218
           auto gt = [&](int neg) {
                                                                277
               auto f = [\&](int x, int y) {
219
                                                                278
                   return sign((v[x] - v[y]) * p) == neg;
220
                                                                279
221
              };
                                                                280
222
               return cycle search(f);
                                                                281
223
                                                                282
224
           int i = gt(1), j = gt(-1), n = v.size();
                                                                283
225
          T x = -(v[i] * p), y = -(v[j] * p);
                                                                284
          if (L.c < x || y < L.c) return -1;
226
                                                                285
227
          if (L.c == x || L.c == y) return 0;
                                                                 286
228
                                                                287
229
          if (i > j) swap(i, j);
                                                                 288
           auto g = [&](int x, int lim) {
230
                                                                289
231
               int now = 0, nxt;
                                                                 290
               for (int i = 1 << __lg(lim); i > 0; i /= 2) {
232
                                                                 291
233
                   if (now + i > lim) continue;
                                                                 292
234
                   nxt = (x + i) % n;
                                                                293
235
                   if (L.ori(v[x]) * L.ori(v[nxt]) > 0) {
                                                                 294
236
                       x = nxt:
                                                                295
                       now += i;
                                                                296
```

```
} // ↓ BE CAREFUL
                                                              298 };
            return -(ori(v[x], v[(x + 1) % n], L.p1) * ori(v[
                 x], v[(x + 1) % n], L.p2));
        int ret = \max(g(i, j - i), g(j, n - (j - i)));
        return (ret == 0) ? (in convex(L.p1) == 0 &&
             in convex(L.p2) == 0) : ret:
// 回傳點過凸包的兩條切線的切點的 0-based index (不保證兩條
    pair<int,int> convex tangent point(point<T> p) {
        int n = v.size(), z = -1, edg = -1;
        auto gt = [&](int neg) {
            auto check = [&](int x) {
                if (v[x] == p) z = x;
                if (btw(v[x], v[(x + 1) % n], p)) edg = x;
                if (btw(v[(x + n - 1) % n], v[x], p)) edg = (
                     x + n - 1) % n;
            };
            auto f = [\&](int x, int y) {
                check(x); check(y);
                return ori(p, v[x], v[y]) == neg;
            return cycle search(f);
        int x = gt(1), y = gt(-1);
        if (z != -1) {
            return \{(z + n - 1) \% n, (z + 1) \% n\};
        else if (edg != -1) {
            return {edg, (edg + 1) % n};
        else {
            return {x, y};
    friend int halfplane_intersection(vector<line<T>> &s,
         polygon<T> &P) {
        auto angle cmp = [&](line<T> &A, line<T> &B) {
            point < T > a = A.p2-A.p1, b = B.p2-B.p1;
            return (a < b);</pre>
        sort(s.begin(), s.end(), angle cmp); // 線段左側為該
             線段半平面
        int L, R, n = s.size();
        vector<point<T>> px(n);
        vector<line<T>> q(n);
        q[L = R = 0] = s[0];
        for(int i = 1; i < n; ++i) {</pre>
            while(L < R && s[i].ori(px[R-1]) <= 0) --R;
            while(L < R && s[i].ori(px[L]) <= 0) ++L;</pre>
            q[++R] = s[i];
            if(q[R].parallel(q[R-1])) {
                if(q[R].ori(s[i].p1) > 0) q[R] = s[i];
            if(L < R) px[R-1] = q[R-1].line intersection(q[R]);
        while(L < R && q[L].ori(px[R-1]) <= 0) --R;
        P.v.clear();
        if(R - L <= 1) return 0;
        px[R] = q[R].line_intersection(q[L]);
        for(int i = L; i <= R; ++i) P.v.push back(px[i]);</pre>
        return R - L + 1;
```

#### 5.2 Pick's Theorem

給定頂點坐標均是整點的簡單多邊形,面積 = 內部格點數 + 邊上格點數/2 - 1

# 6 Graph

#### 6.1 2-SAT [5a6317]

```
1 struct TWO_SAT {
      int n, N;
      vector<vector<int>> G, rev G;
      deque<bool> used;
      vector<int> order, comp;
      deque<bool> assignment;
      void init(int _n) {
          n = n;
          N = n * 2;
          G.resize(N + 5);
          rev G.resize(N + 5);
      void dfs1(int v) {
          used[v] = true;
          for (int u : G[v]) {
              if (!used[u])
                  dfs1(u);
          order.push back(v);
      void dfs2(int v, int cl) {
21
          comp[v] = c1;
          for (int u : rev G[v]) {
              if (comp[u] == -1)
                   dfs2(u, c1);
      bool solve() {
28
          order.clear();
29
          used.assign(N, false);
31
          for (int i = 0; i < N; ++i) {</pre>
               if (!used[i])
                  dfs1(i):
33
34
35
          comp.assign(N, -1);
          for (int i = 0, j = 0; i < N; ++i) {
36
               int v = order[N - i - 1];
37
               if (comp[v] == -1)
                  dfs2(v, j++);
          assignment.assign(n, false);
          for (int i = 0; i < N; i += 2) {
43
               if (comp[i] == comp[i + 1])
                   return false:
               assignment[i / 2] = (comp[i] > comp[i + 1]);
          return true;
```

```
// A or B 都是 0-based
      void add disjunction(int a, bool na, int b, bool nb) {
          // na is true => ~a, na is false => a
          // nb is true => ~b, nb is false => b
          a = 2 * a ^ na;
          b = 2 * b ^ nb;
          int neg a = a ^ 1;
          int neg b = b ^ 1;
          G[neg a].push back(b);
          G[neg_b].push_back(a);
          rev_G[b].push_back(neg_a);
          rev_G[a].push_back(neg_b);
          return;
      void get_result(vector<int>& res) {
          res.clear();
          for (int i = 0; i < n; i++)</pre>
              res.push_back(assignment[i]);
68 };
```

#### 6.2 Augment Path [f8a5dd]

```
i struct AugmentPath{
     int n, m;
     vector<vector<int>> G;
     vector<int> mx, my;
     vector<int> visx, visy;
     int stamp;
     AugmentPath(int _n, int _m) : n(_n), m(_m), G(n), mx(n,
          -1), my(m, -1), visx(n), visy(n){
         stamp = 0;
     void add(int x, int y){
         G[x].push back(y);
     bool dfs1(int now){
         visx[now] = stamp;
         for (auto x : G[now]){
             if (my[x]==-1){
                 mx[now] = x;
                 my[x] = now;
                 return true;
         for (auto x : G[now]){
             if (visx[my[x]]!=stamp && dfs1(my[x])){
                 mx[now] = x;
                 my[x] = now;
                 return true;
         return false;
     vector<pair<int, int>> find max matching(){
         vector<pair<int, int>> ret;
```

```
stamp++;
        int tmp = 0:
        for (int i=0 ; i<n ; i++){</pre>
            if (mx[i]==-1 && dfs1(i)) tmp++;
        if (tmp==0) break;
    for (int i=0 ; i<n ; i++){</pre>
        if (mx[i]!=-1){
            ret.push_back({i, mx[i]});
    return ret;
}
// 645577
void dfs2(int now){
    visx[now] = true;
    for (auto x : G[now]){
        if (my[x]!=-1 && visy[x]==false){
            visy[x] = true;
            dfs2(my[x]);
}
// 要先執行 find_max_matching 一次
vector<pair<int, int>> find min vertex cover(){
    fill(visx.begin(), visx.end(), false);
    fill(visy.begin(), visy.end(), false);
    vector<pair<int, int>> ret;
    for (int i=0 ; i<n ; i++){</pre>
        if (mx[i]==-1) dfs2(i);
    for (int i=0 ; i<n ; i++){</pre>
        if (visx[i]==false) ret.push_back({1, i});
    for (int i=0 ; i<m ; i++){</pre>
```

while (true){

# 6.3 C3C4 [d00465]

return ret;

87

88 };

if (visy[i]==true) ret.push\_back({2, i});

```
vector<vector<int>>> D(N), adj(N);
      for (auto [u, v] : e) {
12
          if (rk[u] > rk[v]) swap(u, v);
          D[u].emplace_back(v);
          adj[u].emplace_back(v);
          adj[v].emplace back(u);
      vector<int> vis(N);
      int c3 = 0, c4 = 0;
      for (int x : ord) { // c3
          for (int y : D[x]) vis[y] = 1;
          for (int y : D[x]) for (int z : D[y]){
              c3 += vis[z]; // xyz is C3
          for (int y : D[x]) vis[y] = 0;
28
      for (int x : ord) { // c4
          for (int y : D[x]) for (int z : adj[y])
              if (rk[z] > rk[x]) c4 += vis[z]++;
          for (int y : D[x]) for (int z : adj[y])
              if (rk[z] > rk[x]) --vis[z];
      } // both are O(M*sqrt(M)), test @ 2022 CCPC guangzhou
      cout << c4 << "\n";
```

#### 6.4 Cut BCC [2af809]

```
#include <bits/stdc++.h>
  using namespace std;
  const int N = 200005;
  vector <int> G[N]:
  int low[N], depth[N];
  bool vis[N];
  vector <vector <int>> bcc;
  stack <int> stk;
  void dfs(int v, int p) {
      stk.push(v);
      vis[v] = true;
      low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
      for (int u : G[v]) {
          if (u == p) continue;
          if (!vis[u]) {
              /// (v, u) 是樹邊
              dfs(u, v);
              low[v] = min(low[v], low[u]);
21
              /// u 無法在不經過父邊的情況走到 v 的祖先
              if (low[u] >= depth[v]) {
                  bcc.emplace_back();
23
                  while (stk.top() != u) {
24
                      bcc.back().push_back(stk.top());
                      stk.pop();
27
                  bcc.back().push_back(stk.top());
28
                  stk.pop();
                  bcc.back().push back(v);
32
          } else {
              /// (v, u) 是回邊
```

```
34 | low[v] = min(low[v], depth[u]);
35 | }
36 | }
37 | }
```

# 6.5 Dinic [961b34]

```
1 | // 一般圖: O(EV2)
2 // 二分圖: O(E√V)
3 struct Flow{
     using T = int; // 可以換成別的型別
     struct Edge{
          int v; T rc; int rid;
     vector<vector<Edge>> G;
     void add(int u, int v, T c){
          G[u].push_back({v, c, G[v].size()});
          G[v].push_back({u, 0, G[u].size()-1});
     vector<int> dis, it;
     Flow(int n){
         G.resize(n):
          dis.resize(n);
         it.resize(n);
     // ce56d6
     T dfs(int u, int t, T f){
          if (u == t || f == 0) return f;
          for (int &i=it[u]; i<G[u].size(); i++){</pre>
              auto &[v, rc, rid] = G[u][i];
              if (dis[v]!=dis[u]+1) continue;
              T df = dfs(v, t, min(f, rc));
              if (df <= 0) continue;</pre>
              rc -= df;
              G[v][rid].rc += df;
              return df;
          return 0;
     // e22e39
     T flow(int s, int t){
          T ans = 0;
          while (true){
              fill(dis.begin(), dis.end(), INF);
              queue<int> q;
              q.push(s);
              dis[s] = 0;
              while (q.size()){
                  int u = q.front(); q.pop();
                  for (auto [v, rc, rid] : G[u]){
   if (rc <= 0 || dis[v] < INF) continue;</pre>
                      dis[v] = dis[u] + 1;
                      q.push(v);
              if (dis[t]==INF) break:
              fill(it.begin(), it.end(), 0);
```

```
while (true){
            T df = dfs(s, t, INF);
            if (df <= 0) break;</pre>
            ans += df;
    return ans;
}
// the code below constructs minimum cut
void dfs mincut(int now, vector<bool> &vis){
    vis[now] = true;
    for (auto &[v, rc, rid] : G[now]){
        if (vis[v] == false && rc > 0){
            dfs_mincut(v, vis);
   }
}
vector<pair<int, int>> construct(int n, int s, vector
     pair<int,int>> &E){
    // E is G without capacity
    vector<bool> vis(n);
    dfs_mincut(s, vis);
    vector<pair<int, int>> ret;
    for (auto &[u, v] : E){
        if (vis[u] == true && vis[v] == false){
            ret.emplace back(u, v);
    return ret;
```

32

85

87

# 6.6 Dominator Tree [52b249]

```
1 /*
2|全部都是 0-based
3 G 要是有向無權圖
4 一開始要初始化 G(N, root),代表有 N 個節點,根是 root
5|用完之後要 build
G[i] = i 的 idom · 也就是從 root 走到 i 時 · 一定要走到的點且離
       i 最近
  struct DominatorTree{
      int N;
      vector<vector<int>> G;
      vector<vector<int>> buckets, rg;
     // dfn[x] = the DFS otder of x
     // rev[x] = the vertex with DFS order x
      // par[x] = the parent of x
      vector<int> dfn, rev, par;
      vector<int> sdom, dom, idom;
      vector<int> fa, val;
      int stamp;
      int root;
      int operator [] (int x){
         return idom[x];
22
23
      DominatorTree(int _N, int _root) :
```

```
N(N)
    G(N), buckets(N), rg(N),
    dfn(N, -1), rev(N, -1), par(N, -1),
    sdom(N, -1), dom(N, -1), idom(N, -1),
    fa(N, -1), val(N, -1)
    stamp = 0;
    root = root;
void add edge(int u, int v){
    G[u].push back(v);
void dfs(int x){
    rev[dfn[x] = stamp] = x;
    fa[stamp] = sdom[stamp] = val[stamp] = stamp;
    stamp++;
    for (int u : G[x]){
        if (dfn[u]==-1){
            dfs(u);
            par[dfn[u]] = dfn[x];
        rg[dfn[u]].push_back(dfn[x]);
int eval(int x, bool first){
    if (fa[x]==x) return !first ? -1 : x;
    int p = eval(fa[x], false);
    if (p==-1) return x;
    if (sdom[val[x]]>sdom[val[fa[x]]]) val[x] = val[fa[x
    fa[x] = p;
    return !first ? p : val[x];
void link(int x, int y){
    fa[x] = y;
void build(){
    dfs(root);
    for (int x=stamp-1; x>=0; x--){
        for (int y : rg[x]){
            sdom[x] = min(sdom[x], sdom[eval(y, true)]);
        if (x>0) buckets[sdom[x]].push back(x);
        for (int u : buckets[x]){
            int p = eval(u, true);
            if (sdom[p]==x) dom[u] = x;
            else dom[u] = p;
        if (x>0) link(x, par[x]);
    idom[root] = root;
    for (int x=1 ; x<stamp ; x++){</pre>
        if (sdom[x]!=dom[x]) dom[x] = dom[dom[x]];
    for (int i=1; i<stamp; i++) idom[rev[i]] = rev[dom[</pre>
         i]];
```

```
6.7 EdgeBCC [d09eb1]
```

91 };

1 // d09eb1

#### 2| // 0-based · 支援重邊 3 struct EdgeBCC{ int n, m, dep, sz; vector<vector<pair<int, int>>> G; vector<vector<int>> bcc; vector<int> dfn, low, stk, isBridge, bccId; vector<pair<int, int>> edge, bridge; EdgeBCC( $int_n$ ): n(n), m(0), sz(0), dfn(n), low(n), G(n)), bcc(n), bccId(n) {} void add\_edge(int u, int v) { edge.push\_back({u, v}); G[u].push\_back({v, m}); $G[v].push back({u, m++});$ void dfs(int now, int pre) { dfn[now] = low[now] = ++dep;stk.push back(now); for (auto [x, id] : G[now]){ **if** (!dfn[x]){ dfs(x, id); low[now] = min(low[now], low[x]); }else if (id!=pre){ low[now] = min(low[now], dfn[x]); if (low[now]==dfn[now]){ if (pre!=-1) isBridge[pre] = true; int u; do{ u = stk.back(); stk.pop back(); bcc[sz].push back(u); bccId[u] = sz; } while (u!=now); SZ++; void get bcc() { isBridge.assign(m, 0); for (int i=0 ; i<n ; i++){</pre> if (!dfn[i]) dfs(i, -1); for (int i=0 ; i<m ; i++){</pre> if (isBridge[i]){ bridge.push\_back({edge[i].first , edge[i]. second});

# 6.8 EnumeratePlanarFace [e70ee1]

57 };

```
1 // 0-based
 struct PlanarGraph{
     int n, m, id;
      vector<point<int>> v;
     vector<vector<pair<int, int>>> G;
     vector<int> conv, nxt, vis;
      PlanarGraph(int n, int m, vector<point<int>> v) :
     n(n), m(m), id(0),
     v(v), G(n),
      conv(2*m), nxt(2*m), vis(2*m) {}
      void add_edge(int x, int y){
          G[x].push_back({y, 2*id});
          G[y].push_back({x, 2*id+1});
          conv[2*id] = x;
          conv[2*id+1] = y;
          id++;
     }
      vector<int> enumerate_face(){
          for (int i=0; i<n; i++){
              sort(G[i].begin(), G[i].end(), [&](pair<int, int>
                    a, pair<int, int> b){
                  return (v[a.first]-v[i])<(v[b.first]-v[i]);</pre>
              int sz = G[i].size(), pre = sz-1;
              for (int j=0 ; j<sz ; j++){</pre>
                  nxt[G[i][pre].second] = G[i][j].second^1;
                  pre = j;
          vector<int> ret;
          for (int i=0 ; i<2*m ; i++){</pre>
              if (vis[i]==false){
                  int area = 0, now = i;
                  vector<int> pt;
                  while (!vis[now]){
                      vis[now] = true;
                      pt.push back(conv[now]);
                      now = nxt[now];
                  pt.push back(pt.front());
                  for (int i=0 ; i+1<pt.size() ; i++){</pre>
                      area -= (v[pt[i]]^v[pt[i+1]]);
                  // pt = face boundary
                  if (area>0){
                      ret.push_back(area);
                  }else{
                      // pt is outer face
```

```
59 return ret;
60 }
61 };
```

#### 6.9 HLD [f57ec6]

```
| #include <bits/stdc++.h>
  #define int long long
  using namespace std;
  const int N = 100005;
  vector <int> G[N];
  struct HLD {
      vector<int> pa, sz, depth, mxson, topf, id;
      int n, idcnt = 0;
      HLD(int _n) : n(_n), pa(_n + 1), sz(_n + 1), depth(_n +
           1), mxson((n + 1), topf((n + 1), id((n + 1)))
      void dfs1(int v = 1, int p = -1) {
          pa[v] = p; sz[v] = 1; mxson[v] = 0;
          depth[v] = (p == -1 ? 0 : depth[p] + 1);
          for (int u : G[v]) {
              if (u == p) continue;
              dfs1(u, v);
              sz[v] += sz[u];
              if (sz[u] > sz[mxson[v]]) mxson[v] = u;
      void dfs2(int v = 1, int top = 1) {
          id[v] = ++idcnt;
          topf[v] = top;
          if (mxson[v]) dfs2(mxson[v], top);
          for (int u : G[v]) {
              if (u == mxson[v] || u == pa[v]) continue;
              dfs2(u, u);
30
      // query 為區間資料結構
      int path_query(int a, int b) {
          int res = 0:
          while (topf[a] != topf[b]) { /// 若不在同一條鍊上
              if (depth[topf[a]] < depth[topf[b]]) swap(a, b);</pre>
              res = max(res, 011); // query : L = id[topf[a]],
                   r = id[a]
              a = pa[topf[a]];
          /// 此時已在同一條鍊上
          if (depth[a] < depth[b]) swap(a, b);</pre>
          res = \max(\text{res}, 011); // query : l = id[b], r = id[a]
41
42
43 };
```

#### 6.10 Kosaraju [c7d5aa]

```
6| SCC compress G(n): 宣告一個有 n 個點的圖
                                                                               if (SCC_id[u]!=SCC_id[v]) result[SCC_id[u]].
7 | . add edge(u, v): 加上一條邊 u -> v
                                                                                    push back(SCC id[v]);
s.compress: O(n Log n) 計算 G3、SCC、SCC_id 的資訊,並把縮點後
                                                                           for (int i=0 ; i<SCC.size() ; i++){</pre>
       的結果存在 result 裡
                                                                               sort(result[i].begin(), result[i].end());
10 SCC[i] = 某個 SCC 中的所有點
|| SCC id[i] = 第 i 個點在第幾個 SCC
13 struct SCC compress{
                                                                76 };
      int N, M, sz;
      vector<vector<int>>> G, inv G, result;
      vector<pair<int, int>> edges;
                                                                  6.11 Kuhn Munkres [e66c35]
      vector<bool> vis;
      vector<int> order:
      vector<vector<int>> SCC;
                                                                 1 // O(n^3) 找到最大權匹配
      vector<int> SCC id;
                                                                   struct KuhnMunkres{
                                                                       int n; // max(n, m)
      SCC compress(int N):
                                                                       vector<vector<int>> G:
      N(N), M(0), sz(0),
                                                                       vector<int> match, lx, ly, visx, visy;
      G(N), inv_G(N),
                                                                       vector<int> slack;
      vis(N), SCC id(N)
                                                                       int stamp = 0;
27
      {}
28
      vector<int> operator [] (int x){
          return result[x];
                                                                       void add(int x, int y, int w){
32
                                                                           G[x][y] = max(G[x][y], w);
      void add_edge(int u, int v){
          G[u].push back(v);
          inv G[v].push back(u);
          edges.push_back({u, v});
                                                                           if (visx[i]==stamp) return false;
          M++;
                                                                           visx[i] = stamp;
                                                                           for (int j=0 ; j<n ; j++){</pre>
      void dfs1(vector<vector<int>> &G, int now){
                                                                               if (visy[j]==stamp) continue;
          vis[now] = 1;
                                                                               int d = lx[i]+ly[j]-G[i][j];
          for (auto x : G[now]) if (!vis[x]) dfs1(G, x);
          order.push back(now):
                                                                               if (d==0){
                                                                                   visy[i] = stamp;
                                                                                   if (match[j]==-1 || dfs(match[j], aug)){
      void dfs2(vector<vector<int>> &G, int now){
                                                                                       if (aug){
          SCC id[now] = SCC.size()-1;
                                                                                           match[j] = i;
          SCC.back().push back(now);
          vis[now] = 1;
                                                                                       return true;
          for (auto x : G[now]) if (!vis[x]) dfs2(G, x);
                                                                               }else{
                                                                                  slack[j] = min(slack[j], d);
      void compress(){
          fill(vis.begin(), vis.end(), 0);
          for (int i=0; i<N; i++) if (!vis[i]) dfs1(G, i);</pre>
                                                                           return false:
                                                                       }
          fill(vis.begin(), vis.end(), 0);
          reverse(order.begin(), order.end());
                                                                       bool augment(){
          for (int i=0 : i<N : i++){</pre>
                                                                           for (int j=0 ; j<n ; j++){</pre>
              if (!vis[order[i]]){
                                                                               if (visy[j]!=stamp && slack[j]==0){
                  SCC.push back(vector<int>());
                                                                                   visy[j] = stamp;
                  dfs2(inv G, order[i]);
                                                                                   if (match[j]==-1 || dfs(match[j], false)){
              }
                                                                                       return true;
          }
          result.resize(SCC.size());
          sz = SCC.size();
                                                                           return false:
          for (auto [u, v] : edges){
```

```
result[i].resize(unique(result[i].begin(), result 55
                                                                        for (int i=0 ; i<n ; i++){</pre>
             [i].end())-result[i].begin());
                                                                            if (visx[i]==stamp) lx[i] -= delta;
                                                                        for (int j=0 ; j<n ; j++){</pre>
                                                                            if (visy[j]==stamp) ly[j] += delta;
                                                                            else slack[j] -= delta;
                                                                   }
                                                                   int solve(){
                                                                        for (int i=0 ; i<n ; i++){</pre>
                                                                            lx[i] = 0;
                                                                            for (int j=0 ; j<n ; j++){</pre>
                                                                                lx[i] = max(lx[i], G[i][j]);
                                                                        fill(ly.begin(), ly.end(), 0);
KuhnMunkres(int n) : n(n), G(n, vector<int>(n)), lx(n),
                                                                        fill(match.begin(), match.end(), -1);
    ly(n), slack(n), match(n), visx(n), visy(n) {}
                                                             76
                                                                        for(int i = 0; i < n; i++) {</pre>
                                                             77
                                                                            fill(slack.begin(), slack.end(), INF);
                                                                            stamp++:
                                                                            if(dfs(i, true)) continue;
bool dfs(int i, bool aug){ // aug = true 表示要更新 match 81
                                                                            while(augment()==false) relabel();
                                                                            stamp++:
                                                                            dfs(i, true);
                                                                        int ans = 0;
                                                                        for (int j=0 ; j<n ; j++){</pre>
                                                             87
                                                                            if (match[j]!=-1){
                                                             88
                                                                                ans += G[match[j]][j];
```

92

93

94 };

51

52

53

# 6.12 LCA [4e91da]

return ans;

void relabel(){

int delta = INF;

for (int j=0 ; j<n ; j++){</pre>

if (visy[j]!=stamp) delta = min(delta, slack[j]);

```
1 struct Tree{
      int N. M = 0. H:
      vector<vector<int>> G:
      vector<vector<int>> LCA;
      vector<int> parent;
      vector<int> dep;
      Tree(int _N) : N(_N), H(__lg(_N)+1){
          G.resize(N);
          parent.resize(N, -1);
          dep.resize(N, 0);
11
          LCA.resize(H, vector<int>(N, 0));
12
13
15
      void add_edge(int u, int v){
```

```
G[u].push back(v);
    G[v].push back(u);
void dfs(int now, int pre){ // root 的 pre 是自己
    dep[now] = dep[pre]+1;
    parent[now] = pre;
    for (auto x : G[now]){
        if (x==pre) continue;
        dfs(x, now);
void build LCA(int root = 0){
    dfs(root, root);
    for (int i=0; i<N; i++) LCA[0][i] = parent[i];</pre>
    for (int i=1; i<H; i++){
        for (int j=0 ; j<N ; j++){</pre>
            LCA[i][j] = LCA[i-1][LCA[i-1][j]];
int jump(int u, int step){
    for (int i=0 ; i<H ; i++){</pre>
        if (step&(1<<i)) u = LCA[i][u];</pre>
    return u;
}
int get_LCA(int u, int v){
    if (dep[u]<dep[v]) swap(u, v);</pre>
    u = jump(u, dep[u]-dep[v]);
    if (u==v) return u;
    for (int i=H-1; i>=0; i--){
        if (LCA[i][u]!=LCA[i][v]){
            u = LCA[i][u];
            v = LCA[i][v];
    return parent[u];
```

# 6.13 MCMF [1b5a27]

```
struct Flow {
    struct Edge {
        int u, rc, k, rv;
    };

vector<vector<Edge>> G;
vector<iint> par, par_eid;
Flow(int n) : G(n+1), par(n+1), par_eid(n+1) {}

// v->u, capcity: c, cost: k
void add(int v, int u, int c, int k){
    G[v].push_back({u, c, k, G[u].size()});
    G[u].push_back({v, 0, -k, G[v].size()-1});
}

// 3701d6
```

```
int spfa(int s, int t){
  fill(par.begin(), par.end(), -1);
  vector<int> dis(par.size(), INF);
  vector<bool> in_q(par.size(), false);
  queue<int> 0;
  dis[s] = 0;
  in_q[s] = true;
  Q.push(s);
  while (!Q.empty()){
    int v = Q.front();
    Q.pop();
    in_q[v] = false;
    for (int i=0 ; i<G[v].size() ; i++){</pre>
      auto [u, rc, k, rv] = G[v][i];
      if (rc>0 && dis[v]+k<dis[u]){</pre>
        dis[u] = dis[v]+k;
        par[u] = v;
        par_eid[u] = i;
        if (!in_q[u]) Q.push(u);
        in_q[u] = true;
  }
  return dis[t];
// return <max flow, min cost>, 150093
pair<int, int> flow(int s, int t){
  int fl = 0, cost = 0, d;
  while ((d = spfa(s, t))<INF){</pre>
    int cur = INF;
    for (int v=t; v!=s; v=par[v])
      cur = min(cur, G[par[v]][par eid[v]].rc);
    cost += d*cur;
    for (int v=t ; v!=s ; v=par[v]){
      G[par[v]][par_eid[v]].rc -= cur;
      G[v][G[par[v]][par eid[v]].rv].rc += cur;
  return {fl, cost};
vector<pair<int, int>> construct(){
  vector<pair<int, int>> ret;
  for (int i=0 ; i<n ; i++){</pre>
    for (auto x : G[i]){
      if (x.rc==0){
        ret.push back({i+1, x.u-n+1});
        break:
  return ret;
```

# 6.14 Tarjan [8b2350]

```
1 struct tarjan_SCC {
```

```
int now_T, now_SCCs;
      vector<int> dfn, low, SCC;
      stack<int> S;
      vector<vector<int>> E;
      vector<bool> vis, in_stack;
      tarjan_SCC(int n) {
          init(n);
      void init(int n) {
          now_T = now_SCCs = 0;
          dfn = low = SCC = vector<int>(n);
          E = vector<vector<int>>(n);
          S = stack<int>();
          vis = in_stack = vector<bool>(n);
      void add(int u, int v) {
19
          E[u].push_back(v);
      void build() {
          for (int i = 0; i < dfn.size(); ++i) {</pre>
              if (!dfn[i]) dfs(i);
24
25
      void dfs(int v) {
          now T++;
          vis[v] = in stack[v] = true;
          dfn[v] = low[v] = now_T;
          S.push(v);
          for (auto &i:E[v]) {
              if (!vis[i]) {
                   vis[i] = true;
                   dfs(i);
                   low[v] = min(low[v], low[i]);
              else if (in stack[i]) {
                   low[v] = min(low[v], dfn[i]);
          if (low[v] == dfn[v]) {
              int tmp;
                   tmp = S.top();
                   S.pop();
                   SCC[tmp] = now_SCCs;
                   in stack[tmp] = false;
47
              } while (tmp != v);
               now SCCs += 1;
50
51
52 };
```

#### 6.15 Tarjan Find AP [1daed6]

```
vector<int> dep(MAX_N), low(MAX_N), AP;
bitset<MAX_N> vis;

void dfs(int now, int pre){
   int cnt = 0;
   bool ap = 0;
   vis[now] = 1;
   low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
```

```
for (auto x : G[now]){
    if (x==pre){
        continue;
    }else if (vis[x]==0){
        cnt++;
        dfs(x, now);
        low[now] = min(low[now], low[x]);
        if (low[x]>=dep[now]) ap=1;
}else{
        low[now] = min(low[now], dep[x]);
}

if ((now==pre && cnt>=2) || (now!=pre && ap)){
        AP.push_back(now);
}
```

#### 6.16 Tree Isomorphism [cd2bbc]

```
1 #include <bits/stdc++.h>
2 #pragma GCC optimize("03,unroll-loops")
  #define fastio ios::sync with stdio(0), cin.tie(0), cout.tie
  #define dbg(x) cerr << #x << " = " << x << endl
  #define int long long
  using namespace std;
  // declare
  const int MAX_SIZE = 2e5+5;
10 const int INF = 9e18;
11 const int MOD = 1e9+7;
  const double EPS = 1e-6;
  typedef vector<vector<int>> Graph;
  typedef map<vector<int>, int> Hash;
16 int n, a, b;
17 int id1, id2;
18 pair<int, int> c1, c2;
19 vector<int> sz1(MAX SIZE), sz2(MAX SIZE);
20 vector<int> we1(MAX_SIZE), we2(MAX_SIZE);
21 Graph g1(MAX_SIZE), g2(MAX_SIZE);
22 Hash m1, m2;
23 int testcase=0;
  void centroid(Graph &g, vector<int> &s, vector<int> &w, pair
       int, int> &rec, int now, int pre){
      s[now]=1;
      w[now]=0;
      for (auto x : g[now]){
          if (x!=pre){
              centroid(g, s, w, rec, x, now);
              s[now]+=s[x];
              w[now]=max(w[now], s[x]);
      w[now]=max(w[now], n-s[now]);
      if (w[now] <= n/2)
          if (rec.first==0) rec.first=now;
          else rec.second=now;
```

```
if (x!=pre){
            int add=dfs(g, m, id, x, now);
            v.push back(add);
    sort(v.begin(), v.end());
    if (m.find(v)!=m.end()){
        return m[v];
    }else{
        m[v]=++id;
        return id;
void solve1(){
    // init
    id1=0:
    id2=0;
    c1={0, 0};
    c2={0, 0};
    fill(sz1.begin(), sz1.begin()+n+1, 0);
    fill(sz2.begin(), sz2.begin()+n+1, 0);
    fill(we1.begin(), we1.begin()+n+1, 0);
    fill(we2.begin(), we2.begin()+n+1, 0);
    for (int i=1; i<=n; i++){</pre>
        g1[i].clear();
        g2[i].clear();
    m1.clear();
    m2.clear();
    // input
    for (int i=0; i<n-1; i++){
        cin >> a >> b;
        g1[a].push_back(b);
        g1[b].push back(a);
    for (int i=0; i<n-1; i++){
        cin >> a >> b;
        g2[a].push_back(b);
        g2[b].push back(a);
    // get tree centroid
    centroid(g1, sz1, we1, c1, 1, 0);
    centroid(g2, sz2, we2, c2, 1, 0);
    // process
    int res1=0, res2=0, res3=0;
    if (c2.second!=0){
        res1=dfs(g1, m1, id1, c1.first, 0);
        m2=m1:
        res2=dfs(g2, m1, id1, c2.first, 0);
        res3=dfs(g2, m2, id2, c2.second, 0);
    }else if (c1.second!=0){
        res1=dfs(g2, m1, id1, c2.first, 0);
```

int dfs(Graph &g, Hash &m, int &id, int now, int pre){

vector<int> v:

for (auto x : g[now]){

```
id2=id1;
109
           res2=dfs(g1, m1, id1, c1.first, 0);
110
           res3=dfs(g1, m2, id2, c1.second, 0);
111
112
           res1=dfs(g1, m1, id1, c1.first, 0);
           res2=dfs(g2, m1, id1, c2.first, 0);
113
114
115
116
       // output
117
       cout << (res1==res2 | res1==res3 ? "YES" : "NO") << endl
118
119
       return;
120
121
   signed main(void){
       fastio:
       int t=1;
       cin >> t;
       while (t--){
128
           solve1();
129
130
       return 0;
131 }
```

# 6.17 圓方樹 [675aec]

```
| #include <bits/stdc++.h>
  #define lp(i,a,b) for(int i=(a);i<(b);i++)</pre>
  #define pii pair<int,int>
  #define pb push back
  #define ins insert
  #define ff first
  #define ss second
  #define opa(x) cerr << #x << " = " << x << ", ";
  #define op(x) cerr << #x << " = " << x << endl;
  #define ops(x) cerr << x;</pre>
  #define etr cerr << endl;</pre>
  #define spc cerr << ' ';</pre>
  #define BAE(x) (x).begin(), (x).end()
14 #define STL(x) cerr << #x << " : "; for(auto &qwe:x) cerr <<
  qwe << ' '; cerr << endl;
#define deb1 cerr << "deb1" << endl;</pre>
  #define deb2 cerr << "deb2" << endl:
  #define deb3 cerr << "deb3" << endl;</pre>
  #define deb4 cerr << "deb4" << endl;</pre>
  #define deb5 cerr << "deb5" << endl;</pre>
  #define bye exit(0);
  using namespace std;
  const int mxn = (int)(2e5) + 10;
  const int mxlg = 17;
  int last_special_node = (int)(1e5) + 1;
  vector<int> E[mxn], F[mxn];
  struct edg{
      int fr, to;
       edg(int _fr, int _to){
           fr = _fr;
31
32
           to = _to;
33
34 };
```

```
35 ostream& operator << (ostream& os, edg x) {os << x.fr << "--" << 99
                                                                          lp(j,1,mxlg){
        x.to;}
                                                                              lp(i,1,mxn){
  vector<edg> EV:
                                                                                  jmp[i][j] = jmp[jmp[i][j-1]][j-1];
                                                                  102
  void tarjan(int v, int par, stack<int>& S){
                                                                  103
                                                                         }
      static vector<int> dfn(mxn), low(mxn);
                                                                  104
      static vector<bool> to add(mxn);
      static int nowT = 0:
                                                                      inline int lca(int x, int y){
                                                                          if(dep[x] < dep[y]){ swap(x, y); }</pre>
      int childs = 0;
                                                                         int diff = dep[x] - dep[y];
      nowT += 1:
      dfn[v] = low[v] = nowT;
                                                                  110
                                                                         lp(j,0,mxlg){
      for(auto &ne:E[v]){
                                                                              if((diff >> j) & 1){
                                                                  111
          int i = EV[ne].to;
                                                                                  x = jmp[x][j];
                                                                  112
          if(i == par) continue;
                                                                  113
          if(!dfn[i]){
                                                                  114
              S.push(ne):
                                                                          if(x == y) return x;
                                                                  115
              tarjan(i, v, S);
                                                                  116
              childs += 1:
                                                                  117
                                                                          for(int j = mxlg - 1; j >= 0; j--){
              low[v] = min(low[v], low[i]);
                                                                  118
                                                                              if(jmp[x][j] != jmp[y][j]){
                                                                                  x = jmp[x][j];
                                                                  119
              if(par >= 0 && low[i] >= dfn[v]){
                                                                                  y = jmp[y][j];
                                                                  120
                  vector<int> bcc;
                                                                  121
                  int tmp:
                                                                  122
                  do{
                                                                  123
                                                                          return jmp[x][0];
                       tmp = S.top(); S.pop();
                                                                  124
                       if(!to add[EV[tmp].fr]){
                                                                  125
                           to add[EV[tmp].fr] = true;
                                                                      inline bool can reach(int fr, int to){
                           bcc.pb(EV[tmp].fr);
                                                                         if(dep[to] > dep[fr]) return false;
                       if(!to_add[EV[tmp].to]){
                                                                          int diff = dep[fr] - dep[to];
                                                                  129
                           to add[EV[tmp].to] = true;
                                                                  130
                                                                          lp(j,0,mxlg){
                           bcc.pb(EV[tmp].to);
                                                                              if((diff >> j) & 1){
                                                                  131
                                                                                  fr = jmp[fr][j];
                                                                  132
                  }while(tmp != ne);
                                                                  133
                  for(auto &j:bcc){
                                                                  134
                       to add[i] = false;
                                                                  135
                                                                          return fr == to;
                       F[last_special_node].pb(j);
                                                                  136
                       F[j].pb(last_special_node);
                                                                  137
                                                                      int main(){
                                                                  138
                  last special node += 1;
                                                                          ios::sync with stdio(false); cin.tie(0);
                                                                         freopen("test_input.txt", "r", stdin);
              }
                                                                          int n, m, q; cin >> n >> m >> q;
          else{
                                                                          lp(i,0,m){
              low[v] = min(low[v], dfn[i]);
                                                                              int u, v; cin >> u >> v;
              if(dfn[i] < dfn[v]){ // edge i--v will be visited 144
                                                                              E[u].pb(EV.size());
                    twice at here, but we only need one.
                                                                              EV.pb(edg(u, v));
                  S.push(ne);
                                                                              E[v].pb(EV.size());
                                                                              EV.pb(edg(v, u));
                                                                  147
                                                                  148
                                                                  149
                                                                          E[0].pb(EV.size());
                                                                          EV.pb(edg(0, 1));
                                                                  150
                                                                          stack<int> S;
                                                                  151
  int dep[mxn], jmp[mxn][mxlg];
                                                                          tarjan(0, -1, S);
                                                                  152
  void dfs lca(int v, int par, int depth){
                                                                          build lca();
                                                                  153
      dep[v] = depth;
                                                                  154
      for(auto &i:F[v]){
                                                                          lp(queries,0,q){
                                                                  155
          if(i == par) continue;
                                                                  156
                                                                              int fr, to, relay; cin >> fr >> to >> relay;
          jmp[i][0] = v;
                                                                              if(fr == relay || to == relay){
                                                                  157
                                                                                  cout << "NO\n":
          dfs_lca(i, v, depth + 1);
                                                                  158
                                                                                  continue:
                                                                  159
                                                                  160
                                                                              if((can reach(fr, relay) || can reach(to, relay)) &&
  inline void build lca(){
                                                                                   dep[relay] >= dep[lca(fr, to)]){
      jmp[1][0] = 1;
                                                                                  cout << "NO\n":
                                                                  162
      dfs_lca(1, -1, 1);
                                                                                  continue:
```

```
164 } cout << "YES\n";
166 }
167 }
```

#### 6.18 最大權閉合圖 [6ca663]

```
邊 u → v 表示選 u 就要選 v (0-based)
      構造:從 S 開始 dfs,不走最小割的邊,
            所有經過的點就是要選的那些點。
      一般圖: O(n²m) / 二分圖: O(m√n)
  template<typename U>
  U maximum closure(vector<U> w, vector<pair<int,int>> EV) {
      int n = w.size(), S = n + 1, T = n + 2;
      Flow G(T + 5); // Graph/Dinic.cpp
      for (int i = 0; i < n; ++i) {</pre>
          if (w[i] > 0) {
             G.add(S, i, w[i]);
             sum += w[i];
         else if (w[i] < 0) {</pre>
             G.add(i, T, abs(w[i]));
21
22
      for (auto &[u, v] : EV) { // 請務必確保 INF > Σ/w i/
23
         G.add(u, v, INF);
24
25
      U cut = G.flow(S, T);
      return sum - cut;
26
```

#### 6.19 Theorem

- 任意圖
  - 最大匹配 + 最小邊覆蓋 = n (不能有孤點)
  - 點覆蓋的補集是獨立集。最小點覆蓋 + 最大獨立集 = n
  - -w(最小權點覆蓋)+w(最大權獨立集 $)=\sum w_v$
  - (帶點權的二分圖可以用最小割解,構造請參考 Augment Path.cpp)
- 二分圖
  - 最小點覆蓋 = 最大匹配 = n 最大獨立集
- 只有邊帶權的二分圖
  - w-vertex-cover ( 帶權點覆蓋 ): 每條邊的兩個連接點被選中的次數總和至少要是  $w_{s}$ 。
  - w-weight matching ( 帶權匹配 )
  - minimum vertex count of w-vertex-cover = maximum weight count of w-weight matching (一個點可以被選很多次·但邊不行)
- 點、邊都帶權的二分圖的定理

```
- b-matching:假設 v 的點權是 b_v · 那所有 v 的匹配邊 e 的權重都要 s_1 滿足 \sum w_e \leq b_v 。
```

 The maximum w-weight of a b-matching equals the minimum b-weight of vertices in a w-vertex-cover.

## 7 Math

1 // ax + by = c

# 7.1 CRT [8d7c58]

```
int extgcd(int a, int b, int c, int &x, int &y) {
                        if (b == 0) {
                                       if (c % a) return INF;
                                       x = c / a, y = 0;
                                       return abs(a);
                        int x1, y1;
                        int g = extgcd(b, a % b, c, x1, y1);
                        x = y1;
                        y = x1 - a / b * y1;
                        return g;
 15 // 有 n 個式子·求解 x ≡ a i (mod m i)
 16 int CRT m coprime(int n, vector<int> &a, vector<int> &m) {
                        int p = 1, ans = 0;
                        vector<int> M(n), inv M(n);
                        for (int i = 0; i < n; i++) p *= m[i];</pre>
                        for (int i = 0; i < n; i++) {</pre>
                                       M[i] = p / m[i];
                                        extgcd(M[i], m[i], inv_M[i], tmp);
                                        ans += a[i] * inv M[i] * M[i];
                                        ans %= p:
                        return (ans % p + p) % p;
31 / / 對於方程組的式子兩兩求解
32 // 回傳: {是否有解, {a, m}}
 pair<bool, pair<int, int>> CRT_m_NOT_coprime(int a1, int m1,
                          int a2, int m2) {
                        int g = __gcd(m1, m2);
                        if ((a2 - a1) % g != 0) return {0, {-1, -1}};
                        int x, y; extgcd(m1, m2, x, y);
                        x = (a2 - a1) * x / g; // 兩者不能相反
                        a1 = x * m1 + a1:
                        m1 = m1 * m2 / g;
                        a1 = (a1 \% m1 + m1) \% m1;
                        return {1, {a1, m1}};
 \frac{1}{46} / \frac{1}{2}  ans = \frac{1}{2} 
47 // ans = ret.F + k * ret.S, k is integer
 48 pair<int, int> div(int a, int b, int m) {
                        int flag = 1;
                        if (a < 0) { a = -a; flag *= -1; }</pre>
```

# 7.2 Josephus Problem [e0ed50]

**if**  $(b < 0) \{ b = -b; flag *= -1; \}$ 

int res = extgcd(b, m, a, t, k);

if (res == INF) return {INF, INF};

int t = -1, k = -1;

m = abs(m / res);

t = t \* flag; t = (t % m + m) % m;

return {t, m}:

# **7.3** Lagrange any x [1f2c26]

```
1 / / init: (x1, y1), (x2, y2) in a vector
 struct Lagrange{
      int n;
      vector<pair<int, int>> v;
      Lagrange(vector<pair<int, int>> & v){
          n = v.size();
          v = v;
      // O(n^2 \log MAX A)
      int solve(int x){
          int ret = 0;
          for (int i=0 ; i<n ; i++){</pre>
              int now = v[i].second;
              for (int j=0 ; j<n ; j++){</pre>
                  if (i==j) continue;
                  now *= ((x-v[j].first)+MOD)%MOD;
                  now %= MOD;
                  now *= (qp((v[i].first-v[j].first+MOD)%MOD,
                       MOD - 2) + MOD) % MOD;
                  now %= MOD;
              ret = (ret+now)%MOD;
          return ret;
```

# 7.4 Lagrange continuous x [57536a]

```
| #include <bits/stdc++.h>
  using namespace std;
  const int MAX N = 5e5 + 10:
  const int mod = 1e9 + 7:
  long long inv fac[MAX N];
  inline int fp(long long x, int y) {
      int ret = 1;
      for (; y; y >>= 1) {
          ret = (y & 1) ? (ret * x % mod) : ret;
          x = x * x % mod:
      return ret;
16
  // TO USE THIS TEMPLATE, YOU MUST MAKE SURE THAT THE MOD
       NUMBER IS A PRIME.
19 struct Lagrange {
20
      Initialize a polynomial with f(x_0), f(x_0 + 1), ..., f(
           x \ 0 + n).
      This determines a polynomial f(x) whose degree is at most
      Then you can call sample(x) and you get the value of f(x)
      Complexity of init() and sample() are both O(n).
25
      int m, shift; // m = n + 1
      vector<int> v, mul;
  // You can use this function if you don't have inv fac array
      void construct inv fac() {
          long long fac = 1;
          for (int i = 2; i < MAX_N; ++i) {</pre>
              fac = fac * i % mod;
          inv_fac[MAX_N - 1] = fp(fac, mod - 2);
          for (int i = MAX N - 1; i >= 1; --i) {
              inv_fac[i - 1] = inv_fac[i] * i % mod;
  // You call init() many times without having a second
       instance of this struct.
      void init(int X_0, vector<int> &u) {
          v = u;
          shift = ((1 - X 0) \% mod + mod) \% mod;
          if (v.size() == 1) v.push back(v[0]);
          m = v.size():
          mul.resize(m);
  // You can use sample(x) instead of sample(x \% mod).
      int sample(int x) {
          x = ((long long)x + shift) % mod;
          x = (x < 0) ? (x + mod) : x;
          long long now = 1;
          for (int i = m; i >= 1; --i) {
53
              mul[i - 1] = now;
              now = now * (x - i) % mod;
54
55
          int ret = 0:
          bool neg = (m - 1) & 1;
```

```
for (int i = 1; i <= m; ++i) {</pre>
               int up = now * mul[i - 1] % mod;
              int down = inv_fac[m - i] * inv_fac[i - 1] % mod; 25
              int tmp = ((long long)v[i - 1] * up % mod) * down 26
              ret += (neg && tmp) ? (mod - tmp) : (tmp);
              ret = (ret >= mod) ? (ret - mod) : ret;
              now = now * (x - i) % mod;
              neg ^= 1;
          return ret;
  };
72 int main() {
      int n; cin >> n;
      vector<int> v(n);
      for (int i = 0; i < n; ++i) {</pre>
          cin >> v[i];
      Lagrange L;
      L.construct_inv_fac();
      L.init(0, v);
      int x; cin >> x;
      cout << L.sample(x);</pre>
```

# 7.5 Lucas's Theorem [b37dcf]

```
1 // 對於很大的 C^n_{m} 對質數 p 取模·只要 p 不大就可以用。
2 int Lucas(int n, int m, int p){
3    if (m==0) return 1;
    return (C(n%p, m%p, p)*Lucas(n/p, m/p, p)%p);
5 }
```

# 7.6 Matrix [8d1a23]

```
i struct Matrix{
     int n. m:
     vector<vector<int>> arr;
     Matrix(int _n, int _m){
         n = _n;
         m = m;
         arr.assign(n, vector<int>(m));
     vector<int> & operator [] (int i){
          return arr[i];
     Matrix operator * (Matrix b){
          Matrix ret(n, b.m);
          for (int i=0 ; i<n ; i++){</pre>
              for (int j=0 ; j<b.m ; j++){</pre>
                  for (int k=0 ; k<m ; k++){</pre>
                      ret.arr[i][j] += arr[i][k]*b.arr[k][j]%
                      ret.arr[i][j] %= MOD;
```

# break; if (target==-1) return 0; if (i!=target){ swap(arr[i], arr[target]); flag = !flag; for (int j=i+1; j<n; j++){</pre> if (!arr[j][i]) continue; int freq = arr[j][i]\*qp(arr[i][i], MOD-2)%MOD for (int k=i ; k<n ; k++){</pre> arr[j][k] -= freq\*arr[i][k]; arr[j][k] = (arr[j][k]%MOD+MOD)%MOD;} int ret = !flag ? 1 : MOD-1; for (int i=0 ; i<n ; i++){</pre> ret \*= arr[i][i]; ret %= MOD; return ret;

# 7.7 Matrix 01 [8d542a]

return ret;

Matrix pow(int p){

return ret;

}

int det(){

Matrix ret(n, n), mul = \*this;
for (int i=0; i<n; i++){</pre>

if (p&1) ret = ret\*mul;

vector<vector<int>> arr = this->arr;

for (int j=i ; j<n ; j++){
 if (arr[j][i]){
 target = j;</pre>

ret.arr[i][i] = 1;

for (; p; p>>=1){

bool flag = false;

for (int i=0 ; i<n ; i++){</pre>

int target = -1;

mul = mul\*mul;

```
const int MAX_N = (1LL<<12);
struct Matrix{
  int n, m;</pre>
```

```
vector<bitset<MAX_N>> arr;
       Matrix(int _n, int _m){
           n = _n;
           m = _m;
           arr.resize(n);
       Matrix operator * (Matrix b){
           Matrix b_t(b.m, b.n);
           for (int i=0 ; i<b.n ; i++){</pre>
               for (int j=0 ; j<b.m ; j++){</pre>
                   b_t.arr[j][i] = b.arr[i][j];
           Matrix ret(n, b.m);
           for (int i=0 ; i<n ; i++){</pre>
               for (int j=0 ; j<b.m ; j++){</pre>
                   ret.arr[i][j] = ((arr[i]&b_t.arr[j]).count()
           return ret;
27
28 };
```

# **7.8** Miller Rabin [2748d2]

```
1 // O(\log n)
  typedef Uint unsigned long long
  Uint modmul(Uint a, Uint b, Uint m) {
      int ret = a*b - m*(Uint)((long double)a*b/m);
      return ret + m*(ret < 0) - m*(ret>=(int)m);
  int qp(int b, int p, int m){
      int ret = 1;
      for (; p; p>>=1){
          if (p&1){
              ret = modmul(ret, b, m);
          b = modmul(b, b, m);
16
      return ret;
  vector<int> llsprp = {2, 325, 9375, 28178, 450775, 9780504,
       1795265022}:
  bool isprime(int n, vector(int) sprp = llsprp){
      if (n==2) return 1;
      if (n<2 || n%2==0) return 0;
      int t = 0;
      int u = n-1;
      for ( ; u%2==0 ; t++) u>>=1;
      for (int i=0 ; i<sprp.size() ; i++){</pre>
          int a = sprp[i]%n;
          if (a==0 || a==1 || a==n-1) continue;
31
          int x = qp(a, u, n);
          if (x==1 || x==n-1) continue;
```

### 7.9 Pollard Rho [a5daef]

```
i | mt19937 seed(chrono::steady clock::now().time since epoch().
      count());
2 int rnd(int 1, int r){
     return uniform int distribution<int>(1, r)(seed);
6 // O(n^{1/4}) 回傳 1 或自己的因數、記得先判斷 n 是不是質數
       (用 Miller-Rabin)
7 // c1670c
8 int Pollard Rho(int n){
     int s = 0, t = 0;
     int c = rnd(1, n-1);
     int step = 0, goal = 1;
     int val = 1;
     for (goal=1 ; ; goal <<=1, s=t, val=1){</pre>
         for (step=1 ; step<=goal ; step++){</pre>
             t = ((\__int128)t*t+c)%n;
             val = ( int128)val*abs(t-s)%n;
             if ((step % 127) == 0){
                 int d = __gcd(val, n);
                 if (d>1) return d;
         }
         int d = __gcd(val, n);
         if (d>1) return d;
```

### 7.10 Polynomial [51ca3b]

```
a = (ll*) calloc(len, sizeof(ll));
Polv(int 1, int d, int *b) {
    deg = d;
    a = b:
void resize(int n) {
    int len1 = 1;
    while (len1 < _n) len1 <<= 1;</pre>
    int *res = (11*) calloc(len1, sizeof(11));
    for (int i = 0; i < min(len, _n); i++) {</pre>
        res[i] = a[i];
    len = len1;
    deg = _n - 1;
    free(a);
    a = res;
Poly& operator=(const Poly rhs) {
    this->len = rhs.len;
    this->deg = rhs.deg;
    this->a = (ll*)realloc(this->a, sizeof(ll) * len);
    copy(rhs.a, rhs.a + len, this->a);
    return *this;
                                                             100
                                                             101
Poly operator*(Poly rhs) {
                                                             102
    int l1 = this->len, l2 = rhs.len;
                                                             103
    int d1 = this->deg, d2 = rhs.deg;
                                                             104
    while (11 > 0 \text{ and this} -> a[11 - 1] == 0) 11--;
                                                             105
    while (12 > 0 and rhs.a[12 - 1] == 0) 12--;
                                                             106
    while (1 < max(11 + 12 - 1, d1 + d2 + 1)) 1 <<= 1;
    int *x, *y, *res;
    x = (11*) calloc(1, sizeof(11));
                                                             110
    y = (11*) calloc(1, sizeof(11));
                                                             111
    res = (ll*) calloc(l, sizeof(ll));
                                                             112
    copy(this->a, this->a + 11, x);
                                                             113
    copy(rhs.a, rhs.a + 12, y);
                                                             114
    ntt.tran(1, x); ntt.tran(1, y);
                                                             115
    FOR (i, 0, 1 - 1)
                                                             116
        res[i] = x[i] * y[i] % mod;
                                                             117
    ntt.tran(1, res, true);
                                                             118
    free(x); free(y);
                                                             119
    return Poly(1, d1 + d2, res);
                                                             120
                                                             121
Poly operator+(Poly rhs) {
                                                             122
    int 11 = this->len, 12 = rhs.len;
    int 1 = \max(11, 12);
                                                             124
    Poly res;
                                                             125
    res.len = 1;
                                                             126
    res.deg = max(this->deg, rhs.deg);
                                                             127
    res.a = (11*) calloc(1, sizeof(11));
                                                             128
    FOR (i, 0, 11 - 1) {
                                                             129
        res.a[i] += this->a[i];
                                                             130
        if (res.a[i] >= mod) res.a[i] -= mod;
                                                             131
                                                             132
    FOR (i, 0, 12 - 1) {
                                                             133
        res.a[i] += rhs.a[i];
                                                             134
        if (res.a[i] >= mod) res.a[i] -= mod;
                                                             135
                                                             136
    return res;
                                                             137
                                                             138
Poly operator - (Poly rhs) {
                                                             139
    int l1 = this->len, l2 = rhs.len;
    int 1 = \max(11, 12);
```

```
Poly res;
    res.len = 1;
    res.deg = max(this->deg, rhs.deg);
    res.a = (ll*) calloc(l, sizeof(ll));
    FOR (i, 0, 11 - 1) {
        res.a[i] += this->a[i];
        if (res.a[i] >= mod) res.a[i] -= mod;
    FOR (i, 0, 12 - 1) {
        res.a[i] -= rhs.a[i];
        if (res.a[i] < 0) res.a[i] += mod;</pre>
    return res;
Poly operator*(const int rhs) {
    Poly res;
    res = *this:
    FOR (i, 0, res.len - 1) {
        res.a[i] = res.a[i] * rhs % mod;
        if (res.a[i] < 0) res.a[i] += mod;</pre>
    return res;
Poly(vector<int> f) {
    int n = f.size();
    len = 1;
    deg = n - 1;
    while (len < _n) len <<= 1;</pre>
    a = (ll*) calloc(len, sizeof(ll));
    FOR (i, 0, deg) a[i] = f[i];
Poly derivative() {
    Poly g(this->deg);
    FOR (i, 1, this->deg) {
        g.a[i - 1] = this - a[i] * i % mod;
    return g;
Poly integral() {
    Poly g(this->deg + 2);
    FOR (i, 0, this->deg) {
        g.a[i + 1] = this -> a[i] * :: inv(i + 1) % mod;
    return g;
Poly inv(int len1 = -1) {
    if (len1 == -1) len1 = this->len;
    Poly g(1); g.a[0] = ::inv(a[0]);
for (int 1 = 1; 1 < len1; 1 <<= 1) {
        Poly t; t = *this;
        t.resize(1 << 1);
        t = g * g * t;
        t.resize(1 << 1);
        Poly g1 = g * 2 - t;
        swap(g, g1);
    return g;
Poly ln(int len1 = -1) {
    if (len1 == -1) len1 = this->len:
    auto g = *this;
    auto x = g.derivative() * g.inv(len1);
    x.resize(len1);
    x = x.integral();
    x.resize(len1):
    return x;
```

```
143
       Poly exp() {
144
           Poly g(1);
145
           g.a[0] = 1;
146
           for (int l = 1; l < len; l <<= 1) {</pre>
               Poly t, g1; t = *this;
147
148
               t.resize(1 << 1); t.a[0]++;
               g1 = (t - g.ln(1 << 1)) * g;
149
               g1.resize(1 << 1);</pre>
150
151
               swap(g, g1);
152
153
           return g;
154
155
       Poly pow(ll n) {
           Poly &a = *this;
156
           int i = 0;
157
           while (i <= a.deg and a.a[i] == 0) i++;</pre>
158
           if (i and (n > a.deg or n * i > a.deg)) return Poly(a
159
                .deg + 1);
           if (i == a.deg + 1) {
               Poly res(a.deg + 1);
161
               res.a[0] = 1;
162
               return res;
163
           Poly b(a.deg - i + 1);
165
166
           int inv1 = ::inv(a.a[i]);
           FOR (j, 0, b.deg)
               b.a[j] = a.a[j + i] * inv1 % mod;
           Poly res1 = (b.ln() * (n % mod)).exp() * (::power(a.a 7.14 歐拉公式 [85f3b1]
                [i], n));
           Poly res2(a.deg + 1);
           FOR (j, 0, min((ll)(res1.deg), (ll)(a.deg - n * i)))
               res2.a[j + n * i] = res1.a[j];
172
           return res2;
174
175 };
```

### 7.11 **josephus [0be067]**

```
ı | // n 個人,每 k 個人就刪除的約瑟夫遊戲
int josephus(int n, int k) {
    if (n == 1)
         return 0;
     if (k == 1)
         return n-1;
     if(k > n)
         return (josephus(n-1, k) + k) % n;
     int cnt = n / k;
    int res = josephus(n - cnt, k);
    res -= n % k;
    if (res < 0)
         res += n;
     else
         res += res / (k - 1);
     return res;
```

# 7.12 數論分塊 [8ccab5]

```
2 時間複雜度為 O(sqrt(n))
 區間為 [l, r]
 for(int i=1 ; i<=n ; i++){</pre>
     int l = i, r = n/(n/i);
     i = r;
     ans.push_back(r);
```

# 7.13 最大質因數 [ca5e52]

```
void max_fac(int n, int &ret){
     if (n<=ret || n<2) return;</pre>
     if (isprime(n)){
          ret = max(ret, n);
          return;
     int p = Pollard_Rho(n);
     max_fac(p, ret), max_fac(n/p, ret);
```

```
1 | // phi(n) = 小於 n 並與 n 互質的正整數數量。
2 // O(sqrt(n)) · 回傳 phi(n)
3 int phi(int n){
      int ret = n;
      for (int i=2; i*i<=n; i++){
          if (n%i==0){
              while (n%i==0) n /= i;
              ret = ret*(i-1)/i;
      if (n>1) ret = ret*(n-1)/n;
      return ret;
  // O(n Log n) · 回傳 1~n 的 phi 值
  vector<int> phi_1_to_n(int n){
      vector<int> phi(n+1);
      phi[0]=0;
      phi[1]=1;
      for (int i=2 ; i<=n ; i++){</pre>
          phi[i]=i-1;
      for (int i=2 ; i<=n ; i++){</pre>
          for (int j=2*i; j<=n; j+=i){ // 枚舉所有倍數
              phi[j]-=phi[i];
      }
      return phi;
```

#### 7.15 Burnside's Lemma

$$\sum_{k=1}^{n} \frac{c(k)}{n}$$

- n:有多少種置換方式(例如:旋轉方式)
- c(k): 所有可能中,經過 k 次旋轉後,仍不會和別人相同的方式的數量

#### 7.16 Catalan Number

任意括號序列:  $C_n = \frac{1}{n+1} \binom{2n}{n}$ 

#### 7.17 Matrix Tree Theorem

目標:給定一張無向圖,問他的生成樹數量。 方法: 先把所有自環刪掉, 定義 Q 為以下矩陣

$$Q_{i,j} = \begin{cases} \deg(v_i) & \text{if } i = j \\ -( \underbrace{\$v_i v_j} \text{ 的數量}) & \text{otherwise} \end{cases}$$

接著刪掉 Q 的第一個 row 跟 column · 它的 determinant 就是答案 。 目標:給定一張有向圖,問他的以 r 為根,可以走到所有點生成樹數量

方法: 先把所有自環刪掉, 定義 Q 為以下矩陣

接著刪掉 Q 的第 r 個 row 跟 column · 它的 determinant 就是答案。

#### 7.18 Stirling's formula

 $n! \approx \sqrt{2\pi n} (\frac{n}{2})^n$ 

#### 7.19 Theorem

- 1.  $1 \sim x$  質數的數量  $\approx \frac{x}{\ln x}$
- 2. x 的因數的數量  $\approx x^{\frac{1}{3}}$
- 3. x 的質因數的數量  $\approx \log \log x$
- 4. p is a prime number  $\Leftrightarrow (p-1)! \equiv -1 \pmod{p}$
- 5. 每個正整數都可以表示成四個整數的平方和
- 6. 任何大於 2 的整數都可以表示成兩個質數的和
- 7.  $n^{k-2} \cdot \prod_{i=1}^k s_i n$  個點、k 的連通塊、加上 k-1 條邊使得變成一個連通 圖的方法數,其中每個連通塊有 s; 個點

#### 7.20 二元一次方程式

```
\begin{cases} ax+by=e\\ cx+dy=f \end{cases} = \begin{cases} x=\frac{ed-bf}{ad-bc}\\ y=\frac{af-ec}{ad-bc} \end{cases} 若 x=\frac{0}{0} 且 y=\frac{0}{0} · 則代表無限多組解。若 x=\frac{*}{0} 且 y=\frac{*}{0} · 則代表無解。
```

#### **7.21** 歐拉定理

```
若 a,m 互質 \cdot 則: a^n\equiv a^{n \bmod \varphi(m)} \pmod m 若 a,m 不互質 \cdot 則: a^n\equiv a^{\varphi(m)+[n \bmod \varphi(m)]} \pmod m
```

#### 7.22 錯排公式

錯排公式:  $(n \oplus 1)$  個人中·每個人皆不再原來位置的組合數)

$$dp_i = \begin{cases} 1 & i = 0 \\ 0 & i = 1 \\ (i-1)(dp_{i-1} + dp_{i-2}) & \text{otherwise} \end{cases}$$

# 8 String

# 8.1 AC automation [6ece7f]

```
1 const int MAXN = 5e5 + 5;
 struct ac automation {
     int go[MAXN][26], fail[MAXN], is end[MAXN];
     int sz:
     void add(string s) {
         int now = 0;
         for (char c : s) {
             if (!go[now][c - 'a'])
                 go[now][c - 'a'] = ++sz;
             now = go[now][c - 'a'];
         is_end[now]++;
     vector<int> que;
     void build() {
         que.pb(0);
         for (int i = 0; i < ssize(que); i++) {</pre>
             auto u = que[i];
             FOR (c, 0, 25) {
                 if (go[u][c]) {
                     int v = go[u][c];
                     fail[v] = !u ? 0 : go[fail[u]][c];
                     is_end[v] += is_end[fail[v]];
```

### 8.2 Hash [942f42]

```
i| mt19937 seed(chrono::steady_clock::now().time_since_epoch().
       count());
  int rng(int 1, int r){
      return uniform_int_distribution<int>(1, r)(seed);
  int A = rng(1e5, 8e8);
  const int B = 1e9+7;
  // 2f6192
  struct RollingHash{
      vector<int> Pow, Pre;
      RollingHash(string s = ""){
          Pow.resize(s.size());
          Pre.resize(s.size());
          for (int i=0 ; i<s.size() ; i++){</pre>
              if (i==0){
                  Pow[i] = 1;
                  Pre[i] = s[i];
              }else{
                  Pow[i] = Pow[i-1]*A%B;
                  Pre[i] = (Pre[i-1]*A+s[i])%B;
          return;
      int get(int 1, int r){ // 取得 [l, r] 的數值
          if (l==0) return Pre[r];
          int res = (Pre[r]-Pre[l-1]*Pow[r-l+1])%B;
          if (res<0) res += B;
32
          return res;
34 };
```

# 8.3 KMP [e5b7ce]

#### **8.4** Manacher [9a4b4d]

```
string Manacher(string str) {
    string tmp = "$#";
    for(char i : str) {
        tmp += i;
        tmp += '#';
    }

vector<int> p(tmp.size(), 0);
    int mx = 0, id = 0, len = 0, center = 0;
    for(int i=1; i<(int)tmp.size(); i++) {
        p[i] = mx > i ? min(p[id*2-i], mx-i) : 1;
        while(tmp[i+p[i]] == tmp[i-p[i]]) p[i]++;
        if(mx(i+p[i]) mx = i+p[i], id = i;
        if(len<p[i]) len = p[i], center = i;
    }
    return str.substr((center-len)/2, len-1);
}</pre>
```

#### 8.5 Min Rotation [9d296f]

# 8.6 Suffix Array [6352b3]

```
1  // 注意·當 /s/=1 時·Lcp 不會有值·務必測試 /s/=1 的 case
2  struct SuffixArray {
3     string s;
4     vector<int> sa, lcp;
6     // 69ced9
7     SuffixArray(string _s, int lim = 256) {
8         s = _s;
9     int n = s.size()+1, k = 0, a, b;
9     vector<int> x(s.begin(), s.end()), y(n), ws(max(n, lim)), rank(n);
```

```
x.push back(0);
12
          sa = lcp = y;
          iota(sa.begin(), sa.end(), 0);
          for (int j=0, p=0 ; p<n ; j=max(1LL, j*2), lim=p) {</pre>
              iota(y.begin(), y.end(), n-j);
              for (int i=0 ; i<n ; i++) if (sa[i] >= j) y[p++]
                   = sa[i] - j;
              fill(ws.begin(), ws.end(), 0);
              for (int i=0; i<n; i++) ws[x[i]]++;</pre>
              for (int i=1; i<lim; i++) ws[i] += ws[i - 1];</pre>
              for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
              swap(x, y), p = 1, x[sa[0]] = 0;
              for (int i=1 ; i<n ; i++){</pre>
                  a = sa[i - 1];
                  b = sa[i];
                  x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
                        ? p - 1 : p++;
          }
          for (int i=1; i<n; i++) rank[sa[i]] = i;</pre>
          for (int i=0, j ; i<n-1 ; lcp[rank[i++]]=k)</pre>
              for (k && k--, j=sa[rank[i]-1]; i+k<s.size() &&</pre>
                   j+k<s.size() && s[i+k]==s[j+k]; k++);
          sa.erase(sa.begin());
          lcp.erase(lcp.begin(), lcp.begin()+2);
      }
      // f49583
      vector<int> pos; // pos[i] = i 這個值在 pos 的哪個地方
      SparseTable st;
      void init lcp(){
          pos.resize(sa.size());
          for (int i=0 ; i<sa.size() ; i++){</pre>
              pos[sa[i]] = i;
                                                                 101
                                                                 102
          if (lcp.size()){
                                                                 103
              st.build(lcp);
                                                                 104
                                                                 105
                                                                 106
                                                                 107
      // 用之前記得 init
                                                                 108
      // 回傳 [l1, r1] 跟 [l2, r2] 的 Lcp · 0-based
                                                                 109
      int get_lcp(int l1, int r1, int l2, int r2){
                                                                 110
          int pos_1 = pos[l1], len_1 = r1-l1+1;
                                                                 111
          int pos_2 = pos[12], len_2 = r2-12+1;
                                                                 112
          if (pos_1>pos_2){
                                                                 113
              swap(pos 1, pos 2);
                                                                 114
              swap(len_1, len_2);
                                                                 115
                                                                 116
                                                                 117
          if (11==12){
                                                                 118 };
              return min(len 1, len 2);
              return min({st.query(pos_1, pos_2), len_1, len_2
                   });
      // 檢查 [L1, r1] 跟 [L2, r2] 的大小關係 · 0-based
      // 如果前者小於後者,就回傳 <0,相等就回傳 =0,否則回傳
      int substring cmp(int 11, int r1, int 12, int r2){
```

```
int len 1 = r1-l1+1;
    int len 2 = r2-12+1;
    int res = get lcp(l1, r1, l2, r2);
    if (res<len_1 && res<len_2){</pre>
        return s[l1+res]-s[l2+res];
    }else if (len_1==res && len_2==res){
        // 如果不需要以 index 作為次要排序參數,這裡要回
        return 11-12;
    }else{
        return len 1==res ? -1 : 1;
}
// 對於位置在 <=p 的後綴,找離他左邊/右邊最接近位置 >p 的
     後綴的 Lcp · 0-based
// pre[i] = s[i] 離他左邊最接近位置 >p 的後綴的 Lcp · 0-
// suf[i] = s[i] 離他右邊最接近位置 >p 的後綴的 Lcp · 0-
     hased
// da12fa
pair<vector<int>, vector<int>> get_left_and_right_lcp(int
    vector<int> pre(p+1);
    vector<int> suf(p+1);
    { // build pre
        int now = 0;
        for (int i=0 ; i<s.size() ; i++){</pre>
            if (sa[i]<=p){</pre>
                pre[sa[i]] = now;
                if (i<lcp.size()) now = min(now, lcp[i]);</pre>
                if (i<lcp.size()) now = lcp[i];</pre>
                                                          13
                                                          14
       }
                                                          15
                                                          16
    { // build suf
                                                          17
        int now = 0:
        for (int i=s.size()-1; i>=0; i--){
                                                          19
            if (sa[i]<=p){</pre>
                suf[sa[i]] = now;
                if (i-1>=0) now = min(now, lcp[i-1]);
            }else{
                if (i-1>=0) now = lcp[i-1];
    return {pre, suf};
```

for (int i=1; i<s.size(); i++){</pre>

while (s[j]==s[i+j]) j++;

if (i<rr) j = min(ret[i-l1], rr-i);</pre>

**int** i = 0:

ret[i] = j;

### 8.8 k-th Substring1 [61f66b]

```
1  // 回傳 s 所有子字串 (完全不同)中·第 k 大的
2  string k_th_substring(string &s, int k){
    int n = s.size();
    SuffixArray sa(s);
    sa.init_lcp();

6  int prePrefix = 0, nowRank = 0;
    for (int i=0 ; i<n; i++){
        int len = n-sa[i];
        int add = len-prePrefix;

12     if (nowRank+add>=k){
        return s.substr(sa[i], prePrefix+k-nowRank);
    }

15     prePrefix = sa.lcp[i];
    nowRank += add;
18 }
```

# 8.7 Z Algorithm [bcfbd6]

```
1 // 定義一個長度為 n 的文本為 T · 則陣列 Z 的 Z[i] 代表 T[0:n]
和 T[i:n] 最長共同前綴
// bcfbd6
vector<int> z_function(string s){
vector<int> ret(s.size());
int ll = 0, rr = 0;
```