Contents				3.8 Segment Tree Persistent			6.19	Theorem	19
				3.9 Sparse Table		7	Mad	d.	10
				3.10 Treap		/	Mat		19
1	Misc	2		3.11 Treap2			7.1	CRT	
	1.1 Note	2		3.12 Trie	9		7.2	Josephus Problem	
	1.2 Default Code	2					7.3	Lagrange any x	
	1.3 Run	2	4	Dynamic-Programming	9		7.4	Lagrange continuous x	
	1.4 Custom Set PQ Sort	2		4.1 Digit DP			7.5	Linear Mod Inverse	
	1.5 Dynamic Bitset	2		4.2 Knaspack On Tree			7.6	Lucas's Theorem	
	1.6 Enumerate Subset	2		4.3 SOS DP			7.7	Matrix	
	1.7 Fast Input	2		4.4 Integer Partition	10		7.8	Matrix 01	
	1.8 OEIS	2	_		4.0		7.9	Miller Rabin	
	1.9 Pragma	3	5	Geometry	10			Pollard Rho	
	1.10 Xor Basis	3		5.1 Geometry Struct				Polynomial	
	1.11 random int	3		5.2 Geometry Struct 3D				gosephus	
	1.12 Python	3		5.3 Pick's Theorem	12			5 數論分塊	
	1.13 diff	3	_		10			. 最大質因數	
	1.14 hash command	3	0	Graph	12			5 歐拉公式	
	1.15 setup	3		6.1 2-SAT				Burnside's Lemma	
	-			6.2 Augment Path				Catalan Number	
2	Convolution	3		6.3 C3C4				Matrix Tree Theorem	
	2.1 FFT any mod	3		6.4 Cut BCC				Stirling's formula	
	2.2 FFT new	4		6.5 Dinic				Theorem	
	2.3 FFT short	4		6.6 Dominator Tree				二元一次方程式	
	2.4 FWT	4		6.7 EdgeBCC				!歐拉定理	
	2.5 Min Convolution Concave Concave	5		6.8 EnumeratePlanarFace			7.23	5 錯排公式	23
	2.6 NTT mod 998244353	5		6.9 HLD					
				6.10 Kosaraju		8	Stri	8	24
3	Data-Structure	5		6.11 Kuhn Munkres			8.1	AC automation	
	3.1 BIT	5		6.12 LCA			8.2	Hash	
	3.2 Disjoint Set Persistent	5		6.13 MCMF			8.3	KMP	
	3.3 PBDS GP Hash Table	5		6.14 Tarjan			8.4	Manacher	
	3.4 PBDS Order Set	5		6.15 Tarjan Find AP			8.5	Min Rotation	
	3.5 Segment Tree Add Set	6		6.16 Tree Isomorphism			8.6	Suffix Array	
	3.6 Segment Tree Li Chao Line	6		6.17 圓方樹			8.7	Z Algorithm	25
	3.7 Segment Tree Li Chao Segment	6		6.18 最大權則合圖	10		Q Q	k_th Substring1	25

Misc

1.1 Note

開始寫題目之前,請做下面的事:

- 在「開始寫任何題目之前」、應該要先自己看過「所有」範例測資的邏輯和 演算法是否有跟範例輸出對上
- 如果你覺得別人的某段程式碼有錯誤,就應該直接講出來
- +2~+3 後就開始生測資跟對拍(根據寫 generator 跟 checker 的時間決定) 寫程式請遵照以下原則:

• 準確使用註解分段程式碼

- declare
- init
- input
- process / queries
- output
- 陣列若可以開到最大,則使用常數宣告大小

上傳之前,請依序檢查以下資訊:

- 1. 是否開啟 IO 優化
- 2. 是否有 t 筆輸入但忘了輸入
- 3. 是否有初始化容器
- 4. 題目範圍有沒有開到最大
- 5. 跑過所有範例測資,並嚴格確認是否正確

1.2 Default Code [481325]

```
i #include <bits/stdc++.h>
 using namespace std;
 #define int long long
4| #define debug(HEHE...) std::cerr << #HEHE << " = ", dout(HEHE 12| priority_queue<my_struct, vector<my_struct>, decltype(cmp)>
 void dout() { std::cerr << "\n"; }</pre>
 template <typename T, typename... U>
 void dout(T t, U... u) { std::cerr << t << ' ', dout(u...); }</pre>
 void solve(){
 signed main(){
     ios::sync_with_stdio(0), cin.tie(0);
     int t = 1;
     while (t--){
         solve();
     return 0;
```

1.3 Run

```
from os import *
f = "pA"
while 1:
    i = input("input: ")
    system("clear")
    p = listdir(".")
    if i != "":
    print(f"file = {f}")
    if system(f"g++ {f}.cpp -std=c++17 -Wall -Wextra -Wshadow
          -O2 -D LOCAL -q -fsanitize=undefined, address -o {f}
        print("CE")
        continue
    for x in sorted(p):
        if f in x and ".in" in x:
            print(x)
            if system(f"./\{f\} < \{x\}"):
                print("RE")
            print()
```

1.4 Custom Set PQ Sort [d4df55]

```
1 // 所有自訂的結構體,務必檢查相等的 case,給所有元素一個排序
  struct my_struct{
     int val;
      my struct(int val) : val( val) {}
  auto cmp = [](my_struct a, my_struct b) {
     return a.val > b.val;
  set<my_struct, decltype(cmp)> ss({1, 2, 3}, cmp);
      pq(cmp, {1, 2, 3});
map<my_struct, my_struct, decltype(cmp)> mp({{1, 4}, {2, 5},
      {3, 6}}, cmp);
```

1.5 Dynamic Bitset [c78aa8]

```
const int MAXN = 2e5 + 5;
template <int len = 1>
void solve(int n) {
    if (n > len) {
        solve<min(len*2, MAXN)>(n);
        return;
    bitset<len> a;
```

1.6 Enumerate Subset [a13e46]

```
1 // 時間複雜度 0(3^n)
2 // 枚舉每個 mask 的子集
 for (int mask=0 ; mask<(1<<n) ; mask++){</pre>
     for (int s=mask ; s>=0 ; s=(s-1)&m){
         // s 是 mask 的子集
         if (s==0) break;
```

1.7 Fast Input [6f8879]

```
1 // fast IO
2 // 6f8879
 inline char readchar(){
      static char buffer[BUFSIZ], * now = buffer + BUFSIZ, *
           end = buffer + BUFSIZ;
      if (now == end)
          if (end < buffer + BUFSIZ)</pre>
          end = (buffer + fread(buffer, 1, BUFSIZ, stdin));
          now = buffer;
      return *now++;
  inline int nextint(){
      int x = 0, c = readchar(), neg = false;
      while(('0' > c | | c > '9') && c!='-' && c!=EOF) c =
      if(c == '-') neg = true, c = readchar();
      while ('0' \le c \&\& c \le '9') x = (x << 3) + (x << 1) + (c^'0')
           , c = readchar();
      if(neg) x = -x;
      return x; // returns 0 if EOF
```

1.8 OEIS [f915c2]

```
山// 若一個線性遞迴有 k 項·給他恰好 2*k 個項可以求出線性遞迴
  // f915c2
  template <typename T>
  vector<T> BerlekampMassey(vector<T> a) {
      auto scalarProduct = [](vector<T> v, T c) {
          for (T &x: v) x *= c;
          return v;
      vector<T> s, best;
      int bestPos = 0;
      for (size_t i = 0; i < a.size(); i++) {</pre>
          \dot{T} error = a[i];
          for (size_t j = 0; j < s.size(); j++) error -= s[j] *</pre>
               a[i-1-j];
          if (error == 0) continue;
          if (s.empty()) {
              s.resize(i + 1);
17
              bestPos = i;
              best.push_back(1 / error);
```

```
continue;

continue;

vector<T> fix = scalarProduct(best, error);
    fix.insert(fix.begin(), i - bestPos - 1, 0);

if (fix.size() >= s.size()) {
    best = scalarProduct(s, - 1 / error);
    best.insert(best.begin(), 1 / error);
    bestPos = i;
    s.resize(fix.size());

}

for (size_t j = 0; j < fix.size(); j++)
    s[j] += fix[j];

return s;
}
</pre>
```

1.9 Pragma [09d13e]

```
| #pragma GCC optimize("03,unroll-loops")
| #pragma GCC target("avx,avx2,sse,sse2,sse3,sse4,popcnt")
```

1.10 Xor Basis [840136]

```
i vector<int> basis;
void add vector(int x){
     for (auto v : basis){
         x=min(x, x^v);
     if (x) basis.push back(x);
9 | // 給一數字集合 S · 求能不能 XOR 出 x
10 bool check(int x){
     for (auto v : basis){
         x=min(x, x^v);
     return 0;
17 / / 給一數字集合 S · 求能 XOR 出多少數字
18 // 答案等於 2^{basis 的大小}
20 | // 給一數字集合 S · 求 XOR 出最大的數字
21 int get_max(){
     int ans=0;
     for (auto v : basis){
         ans=max(ans, ans^v);
     return ans;
```

1.11 random int [9cc603]

```
return uniform_int_distribution<int>(1, r)(seed);
}
```

1.12 Python

```
| # Decimal
  from decimal import *
  getcontext().prec = 6
  # system setting
  sys.setrecursionlimit(100000)
  sys.set int max str digits(10000)
  from turtle import *
  N = 3000000010
  setworldcoordinates(-N, -N, N, N)
  hideturtle()
  speed (100)
  def draw line(a, b, c, d):
      teleport(a, b)
      goto(c, d)
  def write_dot(x, y, text, diff=1): # diff = 文字的偏移
      teleport(x, y)
      dot(5, "red")
      teleport(x+N/100*diff, y+N/100*diff)
      write(text, font=("Arial", 5, "bold"))
29 draw_line(*a[i], *(a[i-1]))
30 write_dot(*a[i], str(a[i]))
```

1.13 diff

1.14 hash command

```
| cat file.cpp | cpp -dD -P -fpreprocessed | tr -d "[:space:]"
| md5sum | cut -c-6
```

1.15 setup

```
se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
  :inoremap " ""<Esc>i
  :inoremap {<CR> {<CR>}<Esc>ko
  :inoremap {{ {}}<ESC>i
  function! F(...)
   execute '!./%:r < ./' . a:1
  endfunction
  command! -nargs=* R call F(<f-args>)
  map <F7> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -Wextra -
       Wshadow -02 -DLOCAL -g -fsanitize=undefined,address<CR>
13 map <F8> :!./%:r<CR>
14 map <F9> :!./%:r < ./%:r.in<CR>
  ca hash w !cpp -dD -P -fpreprocessed \| tr -d "[:space:]" \|
       md5sum \| cut -c-6
  " i+<esc>25A---+<esc>
  " o|<esc>25A |<esc>
  " "ggVGyG35pGdd
```

2 Convolution

2.1 FFT any mod [234f9e]

```
2 | 修改 const int MOD = 998244353 更改要取餘的數字
 3 \mid PolyMul(a, b) 回傳多項式乘法的結果 ( c_k = \sum_{i=1}^{n} a_i + b_j
        mod MOD )
   大約可以支援 5e5 \cdot a_i, b_i 皆在 MOD 以下的非負整數
   const int MOD = 998244353;
   typedef complex<double> cd;
   // b9c90a
   void FFT(vector<cd> &a) {
       int n = a.size(), L = 31-__builtin_clz(n);
       vector<complex<long double>> R(2, 1);
       vector<cd> rt(2, 1);
       for (int k=2; k < n; k*=2){
           R.resize(n);
           rt.resize(n);
           auto x = polar(1.0L, acos(-1.0L) / k);
           for (int i=k ; i<2*k ; i++){</pre>
20
               rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);
21
22
23
       vector<int> rev(n);
       for (int i=0 ; i<n ; i++){</pre>
           rev[i] = (rev[i/2] | (i&1) << L)/2;
27
       for (int i=0 ; i<n ; i++){</pre>
           if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
```

```
for (int k=1; k<n; k*=2){</pre>
        for (int i=0 ; i<n ; i+=2*k){</pre>
            for (int j=0; j<k; j++){</pre>
                 auto x = (double *)&rt[j+k];
                 auto y = (double *)&a[i+j+k];
                 cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]* 17
                     y[0]);
                 a[i+j+k] = a[i+j]-z;
                 a[i+j] += z;
        }
    return;
// d3c65e
vector<int> PolyMul(vector<int> a, vector<int> b){
   if (a.empty() || b.empty()) return {};
    vector<int> res(a.size()+b.size()-1);
   int B = 32- builtin clz(res.size()), n = (1<<B), cut =</pre>
         int(sqrt(MOD));
    vector<cd> L(n), R(n), outs(n), outl(n);
    for (int i=0 ; i<a.size() ; i++){</pre>
        L[i] = cd((int) a[i]/cut, (int)a[i]%cut);
    for (int i=0 ; i<b.size() ; i++){</pre>
                                                                      vector<double> PolyMul(const vector<double> a, const vector<
        R[i] = cd((int) b[i]/cut, (int)b[i]%cut);
    FFT(L);
   FFT(R);
    for (int i=0 ; i<n ; i++){</pre>
        int j = -i\&(n-1);
        \operatorname{outl}[j] = (L[i] + \operatorname{conj}(L[j])) * R[i]/(2.0*n);
        outs[j] = (L[i]-conj(L[j])) * R[i]/(2.0*n)/1i;
    FFT(outl);
   FFT(outs);
    for (int i=0 ; i<res.size() ; i++){</pre>
        int av = (int)(real(outl[i])+0.5), cv = (int)(imag(
             outs[i])+0.5);
        int bv = (int)(imag(outl[i])+0.5) + (int)(real(outs[i
        res[i] = ((av%MOD*cut+bv) % MOD*cut+cv) % MOD;
    return res;
```

2.2 FFT new [c95bb8]

```
1 typedef complex < double > cd;
 // b9c90a
4 void FFT(vector<cd> &a) {
     int n = a.size(), L = 31-__builtin_clz(n);
     vector<complex<long double>> R(2, 1);
     vector<cd> rt(2, 1);
     for (int k=2 ; k<n ; k*=2){</pre>
          R.resize(n);
         rt.resize(n);
```

2.3 FFT short [70c01a]

```
1 #define int long long
 using Cplx = complex<double>;
 const double pi = acos(-1);
 const int mod = 998244353, g = 3;
 int power(int a, int b) {
     int res = 1;
```

auto x = polar(1.0L, acos(-1.0L) / k);

rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);

for (int i=k ; i<2*k ; i++){</pre>

rev[i] = (rev[i/2] | (i&1) << L)/2;

for (int i=0; i<n; i+=2*k){

a[i+j] += z;

if (a.empty() || b.empty()) return {};

copy(a.begin(), a.end(), begin(in)); for (int i=0 ; i<b.size() ; i++){</pre>

for (int i=0 ; i<res.size() ; i++){</pre>

res[i] = imag(out[i]) / (4 * n);

vector<cd> in(n), out(n);

in[i].imag(b[i]);

for (cd& x : in) x *= x;

for (int i=0 ; i<n ; i++){</pre>

vector < double > res(a.size()+b.size()-1);

int L = 32 - __builtin_clz(res.size()), n = 1 << L;</pre>

out[i] = in[-i & (n - 1)] - conj(in[i]);

if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>

for (int j=0 ; j<k ; j++){</pre>

a[i+j+k] = a[i+j]-z;

auto x = (double *)&rt[j+k];

auto y = (double *)&a[i+j+k];

}

}

return;

double> b){

FFT(in);

FFT(out);

return res;

vector<int> rev(n); for (int i=0 ; i<n ; i++){</pre>

for (int i=0 ; i<n ; i++){</pre>

for (int k=1; k<n; k*=2){

```
while (b) {
                                                           if (b & 1) res = res * a % mod;
                                                           a = a * a % mod:
                                                           b >>= 1;
                                                       return res;
                                                  int inv(int x) { return power(x, mod - 2); }
                                                   // FFT use Cplx, NTT use ll
                                                   void FFT(vector<int> &a, int n, int op) {
                                                       // n must be 2^k
                                                       vector<int> R(n);
                                                       FOR (i, 0, n - 1)
                                                           R[i] = R[i/2]/2 + (i&1)*(n/2);
                                                       FOR (i, 0, n - 1)
                                                           if (i < R[i]) swap(a[i], a[R[i]]);</pre>
                                                       for (int m = 2; m <= n; m *= 2) {</pre>
                                                           // Cplx w1({cos(2*pi/m), sin(2*pi/m)*op});
cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]* 26
                                                           int w1 = power(g, (mod-1)/m * op + mod-1);
                                                           for (int i = 0; i < n; i += m) {</pre>
                                                               // Cplx wk({1, 0});
                                                               int wk = 1;
                                                               FOR (k, 0, m / 2 - 1) {
                                                                   auto x = a[i+k], y = a[i+k+m/2] * wk % mod;
                                                                   a[i+k] = (x+y) \% mod;
                                                                   a[i+k+m/2] = (x-y+mod) \% mod;
                                                                   w\bar{k} = wk * w1 \% mod;
                                                       if (op == -1)
                                                           FOR (i, 0, n - 1) {
                                                               // a[i] = a[i] / n;
                                                41
                                                               a[i] = a[i] * inv(n) % mod;
                                                42
                                                43
```

2.4 FWT [832aa5]

```
│// 已經把 mint 刪掉‧需要增加註解
  vector<int> xor_convolution(vector<int> a, vector<int> b, int
      if (k == 0) {
          return vector<int>{a[0] * b[0]};
      vector<int> aa(1 << (k - 1)), bb(1 << (k - 1));
      FOR (i, 0, (1 << (k - 1)) - 1) {
          aa[i] = a[i] + a[i + (1 << (k - 1))];
          bb[i] = b[i] + b[i + (1 << (k - 1))];
      vector<int> X = xor convolution(aa, bb, k - 1);
      FOR (i, 0, (1 << (k - 1)) - 1) {
          aa[i] = a[i] - a[i + (1 << (k - 1))];
          bb[i] = b[i] - b[i + (1 << (k - 1))];
      vector<int> Y = xor convolution(aa, bb, k - 1);
      vector<int> c(1 << k);</pre>
      FOR (i, 0, (1 << (k - 1)) - 1) {
                              ] = (X[i] + Y[i]) / 2;
          c[i + (1 << (k - 1))] = (X[i] - Y[i]) / 2;
20
21
      return c;
```

2.5 Min Convolution Concave Concave [ffb28d]

```
1 // 需要增加註解
2 // min convolution
3 vector<int> mkk(vector<int> a, vector<int> b) {
4 vector<int> slope;
5 FOR (i, 1, ssize(a) - 1) slope.pb(a[i] - a[i - 1]);
6 FOR (i, 1, ssize(b) - 1) slope.pb(b[i] - b[i - 1]);
7 sort(all(slope));
8 slope.insert(begin(slope), a[0] + b[0]);
9 partial_sum(all(slope), begin(slope));
1 return slope;
1 }
```

 $| \text{const int MOD} = (119 \iff 23) + 1, ROOT = 62; // = 998244353$

2.6 NTT mod 998244353 [5c6335]

```
2 // For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 <<
  // and 483 << 21 (same root). The last two are > 10^9.
5 // 9cd58a
  void NTT(vector<int> &a) {
      int n = a.size();
      int L = 31- builtin clz(n);
      vector<int> rt(2, 1);
      for (int k=2, s=2; k< n; k*=2, s++){
          rt.resize(n);
          int z[] = {1, qp(ROOT, MOD>>s)};
          for (int i=k ; i<2*k ; i++){</pre>
              rt[i] = rt[i/2]*z[i&1]%MOD;
      vector<int> rev(n);
      for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2]|(i&1)<<L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]){</pre>
              swap(a[i], a[rev[i]]);
      }
      for (int k=1; k<n; k*=2){</pre>
          for (int i=0 ; i<n ; i+=2*k){</pre>
              for (int j=0 ; j<k ; j++){</pre>
                   int z = rt[j+k]*a[i+j+k]%MOD, &ai = a[i+j];
                   a[i+j+k] = ai-z+(z>ai ? MOD : 0);
                   ai += (ai+z)=MOD ? z-MOD : z);
  // 0b0e99
40 vector<int> polyMul(vector<int> &a, vector<int> &b){
      if (a.empty() || b.empty()) return {};
      int s = a.size()+b.size()-1, B = 32-__builtin_clz(s), n =
      int inv = qp(n, MOD-2);
```

3 Data-Structure

3.1 BIT [7ef3a9]

```
1 vector<int> BIT(MAX_SIZE);
2 3 // const int MAX_N = (1<<20)
4 int k_th(int k){ // 回傳 BIT 中第 k 小的元素 (based-1)
5 int res = 0;
6 for (int i=MAX_N>>1; i>=1; i>>=1)
7 if (BIT[res+i]<k)
8 k -= BIT[res+=i];
9 return res+1;
```

3.2 Disjoint Set Persistent [447002]

```
struct Persistent_Disjoint_Set{
    Persistent_Segment_Tree arr, sz;
    void init(int n){
        arr.init(n);
        vector<int> v1;
        for (int i=0 ; i<n ; i++){</pre>
            v1.push back(i);
        arr.build(v1, 0);
        sz.init(n);
        vector<int> v2;
        for (int i=0 ; i<n ; i++){</pre>
            v2.push back(1);
        sz.build(v2, 0);
    int find(int a){
        int res = arr.query_version(a, a+1, arr.version.size
             ()-1).val;
        if (res==a) return a;
        return find(res);
    bool unite(int a, int b){
        a = find(a);
        b = find(b);
```

```
if (a!=b){
              int sz1 = sz.query_version(a, a+1, arr.version.
32
                   size()-1).val;
              int sz2 = sz.query version(b, b+1, arr.version.
                   size()-1).val;
              if (sz1<sz2){
                  arr.update_version(a, b, arr.version.size()
                  sz.update_version(b, sz1+sz2, arr.version.
                       size()-1);
                  arr.update_version(b, a, arr.version.size()
                  sz.update_version(a, sz1+sz2, arr.version.
                       size()-1);
              return true;
43
          return false;
44
45
46
```

3.3 PBDS GP Hash Table [866cf6]

```
i #include <ext/pb ds/assoc container.hpp>
  using namespace __gnu_pbds;
  typedef tree<int, null type, less<int>, rb tree tag,
       tree order statistics node update> order set;
  struct custom hash {
      static uint64 t splitmix64(uint64 t x) {
          // http://xorshift.di.unimi.it/splitmix64.c
          x += 0x9e3779b97f4a7c15;
          x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
          x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
          return x ^ (x >> 31);
      size t operator()(uint64 t x) const {
          static const uint64 t FIXED RANDOM = chrono::
               steady_clock::now().time_since_epoch().count();
          return splitmix64(x + FIXED RANDOM);
15
16
17
  };
19 gp_hash_table<int, int, custom_hash> ss;
```

3.4 PBDS Order Set [231774]

3.5 Segment Tree Add Set [bb1898]

```
1 // [ll, rr), based-0
2 // 使用前記得 init(陣列大小), build(陣列名稱)
3 // add(LL, rr): 區間修改
4 // set(ll, rr): 區間賦值
5 // query(ll, rr): 區間求和 / 求最大值
6 struct SegmentTree{
     struct node{
         int add tag = 0;
         int set tag = 0;
         int sum = 0;
         int ma = 0;
     };
     vector<node> arr;
     SegmentTree(int n){
         arr.resize(n<<2);</pre>
     node pull(node A, node B){
         C.sum = A.sum + B.sum:
         C.ma = max(A.ma, B.ma);
         return C;
      // cce0c8
     void push(int idx, int ll, int rr){
         if (arr[idx].set tag!=0){
             arr[idx].sum = (rr-11)*arr[idx].set tag;
             arr[idx].ma = arr[idx].set tag;
             if (rr-ll>1){
                 arr[idx*2+1].add tag = 0;
                 arr[idx*2+1].set_tag = arr[idx].set_tag;
                 arr[idx*2+2].add tag = 0;
                 arr[idx*2+2].set tag = arr[idx].set tag;
             arr[idx].set tag = 0;
         if (arr[idx].add tag!=0){
             arr[idx].sum += (rr-ll)*arr[idx].add tag;
             arr[idx].ma += arr[idx].add tag;
                 arr[idx*2+1].add tag += arr[idx].add tag;
                 arr[idx*2+2].add_tag += arr[idx].add_tag;
             arr[idx].add_tag = 0;
     void build(vector<int> &v, int idx = 0, int ll = 0, int
          rr = n){
         if (rr-ll==1){
             arr[idx].sum = v[11]:
             arr[idx].ma = v[ll];
         }else{
```

```
int mid = (11+rr)/2;
               build(v, idx*2+1, ll, mid);
               build(v, idx*2+2, mid, rr);
               arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
       void add(int ql, int qr, int val, int idx = 0, int ll =
            0, int rr =n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return;</pre>
           if (ql<=ll && rr<=qr){
               arr[idx].add_tag += val;
               push(idx, 11, rr);
               return;
           int mid = (11+rr)/2;
           add(ql, qr, val, idx*2+1, ll, mid);
           add(ql, qr, val, idx*2+2, mid, rr);
           arr[idx]=pull(arr[idx*2+1], arr[idx*2+2]);
       void set(int ql, int qr, int val, int idx=0, int ll=0,
            int rr=n){
           push(idx, 11, rr);
           if (rr<=ql || qr<=ll) return;</pre>
           if (q1<=11 && rr<=qr){</pre>
               arr[idx].add tag = 0;
               arr[idx].set_tag = val;
               push(idx, ll, rr);
               return;
           int mid = (11+rr)/2;
           set(ql, qr, val, idx*2+1, ll, mid);
           set(ql, qr, val, idx*2+2, mid, rr);
           arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
       node query(int ql, int qr, int idx = 0, int ll = 0, int
            rr = n){
           push(idx, 11, rr);
           if (rr<=ql || qr<=ll) return node();</pre>
           if (q1<=11 && rr<=qr) return arr[idx];</pre>
           int mid = (11+rr)/2;
           return pull(query(ql, qr, idx*2+1, ll, mid), query(ql
                , qr, idx*2+2, mid, rr));
100 } ST;
```

3.6 Segment Tree Li Chao Line [45b8ba]

```
int b = INF;
    int v(int x){
       return a*x+b;
};
vector<Node> arr;
LC Segment Tree(int n = 0){
   arr.resize(4*n);
void update(Node val, int idx = 0, int ll = 0, int rr =
    if (rr-ll==0) return;
   if (rr-ll==1){
        if (val.y(ll)<arr[idx].y(ll)){</pre>
           arr[idx] = val;
        return;
   int mid = (11+rr)/2;
   if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
         的線斜率要比較小
    if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
        update(val, idx*2+1, ll, mid);
    }else{ // 交點在右邊
       swap(arr[idx], val); // 在左子樹中,新線比舊線還
       update(val, idx*2+2, mid, rr);
   return;
}
int query(int x, int idx = 0, int 11 = 0, int rr = MAX_V)
   if (rr-ll==0) return INF;
   if (rr-ll==1){
        return arr[idx].y(ll);
   int mid = (11+rr)/2;
   if (x<mid){</pre>
        return min(arr[idx].y(x), query(x, idx*2+1, ll,
       return min(arr[idx].y(x), query(x, idx*2+2, mid,
            rr));
```

struct Node{ // y = ax+b
 int a = 0;

struct LC Segment Tree{

20

27

3.7 Segment Tree Li Chao Segment [2cb0a4]

```
1 | /*
2 | 全部都是 0-based
```

60

61

```
LC Segment Tree st(n);
s|update_segment({a, b}, ql, qr):在 [ql, qr) 插入一條 y=ax+b
9 | querv(x): 查詢所有直線在位置 x 的最小值
||| const int MAX V = 1e6+10; // 值域最大值
  struct LC Segment Tree{
     struct Node{ // y = ax+b
         int a = 0;
         int b = INF;
         int y(int x){
             return a*x+b;
     };
     vector<Node> arr;
     LC Segment Tree(int n = 0){
         arr.resize(4*n):
     void update(Node val, int idx = 0, int ll = 0, int rr =
         if (rr-ll==0) return;
         if (rr-ll<=1){
             if (val.y(ll)<arr[idx].y(ll)){</pre>
                 arr[idx] = val;
             return:
         }
         int mid = (11+rr)/2;
         if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
              的線斜率要比較小
         if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
             update(val, idx*2+1, ll, mid);
         }else{ // 交點在右邊
             swap(arr[idx], val); // 在左子樹中·新線比舊線還
             update(val, idx*2+2, mid, rr);
         return:
     }
     // 在 [ql, qr) 加上一條 val 的線段
     void update segment(Node val, int ql, int qr, int idx =
          0, int 11 = 0, int rr = MAX_V){
         if (rr-ll==0) return;
         if (rr<=ql || qr<=ll) return;</pre>
         if (q1<=11 && rr<=qr){
             update(val, idx, ll, rr);
             return:
         int mid = (11+rr)/2;
         update segment(val, ql, qr, idx*2+1, ll, mid);
         update segment(val, ql, qr, idx*2+2, mid, rr);
          return;
```

```
62
63
    int query(int x, int idx = 0, int ll = 0, int rr = MAX_V)
64
    if (rr-ll==0) return INF;
65
    if (rr-ll==1){
        return arr[idx].y(ll);
    }
68
    int mid = (ll+rr)/2;
    if (xxmid){
        return min(arr[idx].y(x), query(x, idx*2+1, ll, mid));
    }
72
    return min(arr[idx].y(x), query(x, idx*2+2, mid, something arr));
73
    rr));
74
    }
75
    }
76
};
```

3.8 Segment Tree Persistent [3b5aa9]

```
全部都是 0-based
  Persistent_Segment_Tree st(n+q);
  st.build(v, 0);
  函式:
  update version(pos, val, ver): 對版本 ver 的 pos 位置改成 val 73
  query version(ql, qr, ver): 對版本 ver 查詢 [ql, qr) 的區間和 74
  clone version(ver):複製版本 ver 到最新的版本
  struct Persistent_Segment_Tree{
                                                               77
      int node cnt = 0;
                                                               78
      struct Node{
          int lc = -1;
          int rc = -1;
          int val = 0;
      };
      vector<Node> arr;
      vector<int> version;
      Persistent Segment Tree(int sz){
          arr.resize(32*sz);
          version.push back(node cnt++);
      void pull(Node &c, Node a, Node b){
          c.val = a.val+b.val;
          return:
      void build(vector<int> &v, int idx, int ll = 0, int rr =
          auto &now = arr[idx];
37
          if (rr-ll==1){
                                                              100
              now.val = v[11];
                                                              101 };
              return;
```

```
int mid = (11+rr)/2;
    now.lc = node cnt++:
    now.rc = node cnt++;
    build(v, now.lc, ll, mid);
    build(v, now.rc, mid, rr);
    pull(now, arr[now.lc], arr[now.rc]);
    return:
void update(int pos, int val, int idx, int ll = 0, int rr
    auto &now = arr[idx];
    if (rr-ll==1){
        now.val = val:
        return:
    int mid = (11+rr)/2;
    if (pos<mid){</pre>
        arr[node cnt] = arr[now.lc];
        now.lc = node cnt;
        node cnt++:
        update(pos, val, now.lc, ll, mid);
        arr[node cnt] = arr[now.rc];
        now.rc = node cnt;
        node cnt++;
        update(pos, val, now.rc, mid, rr);
    pull(now, arr[now.lc], arr[now.rc]);
    return:
void update version(int pos, int val, int ver){
    update(pos, val, version[ver]);
Node query(int ql, int qr, int idx, int ll = 0, int rr =
    auto &now = arr[idx];
    if (q1<=11 && rr<=qr) return now;</pre>
    if (rr<=ql || qr<=ll) return Node();</pre>
    int mid = (11+rr)/2;
    pull(ret, query(ql, qr, now.lc, ll, mid), query(ql,
         gr, now.rc, mid, rr));
    return ret;
Node query version(int ql, int qr, int ver){
    return query(ql, qr, version[ver]);
void clone version(int ver){
    version.push back(node cnt);
    arr[node cnt] = arr[version[ver]];
    node cnt++;
```

3.9 Sparse Table [31f22a]

3.10 Treap [5851f5]

```
1 struct Treap{
      Treap *1 = nullptr, *r = nullptr;
      int pri = rand(), val = 0, sz = 1;
      Treap(int _val){
          val = _val;
  };
int size(Treap *t){return t ? t->sz : 0;}
void pull(Treap *t){
      t->sz = size(t->1)+size(t->r)+1;
15 Treap* merge(Treap *a, Treap *b){
      if (!a || !b) return a ? a : b;
      if (a->pri>b->pri){
          a \rightarrow r = merge(a \rightarrow r, b);
          pull(a);
          return a;
          b->1 = merge(a, b->1);
          pull(b):
          return b;
30 | pair<Treap*, Treap*> split(Treap *&t, int k){ // 1-based <前
      k 個元素, 其他元素>
      if (!t) return {};
      if (size(t->1)>=k){
          auto pa = split(t->1, k);
```

```
t->1 = pa.second;
        pull(t);
        return {pa.first, t};
        auto pa = split(t->r, k-size(t->l)-1);
        t->r = pa.first;
        pull(t);
        return {t, pa.second};
// functions
Treap* build(vector<int> v){
    Treap* ret = nullptr;
    for (int i=0 ; i<v.size() ; i++){</pre>
        ret = merge(ret, new Treap(v[i]));
    return ret;
array<Treap*, 3> cut(Treap *t, int 1, int r){ // 1-based <前
    1~l-1 個元素, l~r 個元素, r+1 個元素>
    array<Treap*, 3> ret;
    tie(ret[1], ret[2]) = split(t, r);
    tie(ret[0], ret[1]) = split(ret[1], 1-1);
    return ret;
void print(Treap *t, bool flag = true){
    if (t->1!=0) print(t->1, false);
    cout << t->val;
    if (t->r!=0) print(t->r, false);
    if (flag) cout << endl;</pre>
```

3.11 Treap2 [1bf328]

```
ı | // 1-based · 請注意 MAX_N 是否足夠大
2 int root = 0;
  int lc[MAX_N], rc[MAX_N];
  int pri[MAX_N], val[MAX_N];
  int sz[MAX_N], tag[MAX_N], fa[MAX_N];
  int new_node(int v){
      static int nodeCnt = 0:
      nodeCnt++;
      val[nodeCnt] = v;
      sz[nodeCnt] = 1;
      pri[nodeCnt] = rand();
      return nodeCnt;
  void push(int x){
      if (tag[x]){
          if (lc[x]) tag[lc[x]] ^= 1;
          if (rc[x]) tag[rc[x]] ^= 1;
      tag[x] = 0;
22 int pull(int x){
      if (x){
          fa[x] = 0;
          sz[x] = 1+sz[lc[x]]+sz[rc[x]];
```

```
if (lc[x]) fa[lc[x]] = x;
          if (rc[x]) fa[rc[x]] = x;
29
      return x;
30 }
  int merge(int a, int b){
      if (!a or !b) return a|b;
      push(a), push(b);
      if (pri[a]>pri[b]){
          rc[a] = merge(rc[a], b);
          return pull(a);
          lc[b] = merge(a, lc[b]);
          return pull(b);
43
  // [1, k] [k+1, n]
  void split(int x, int k, int &a, int &b) {
      if (!x) return a = b = 0, void();
      if (sz[lc[x]] >= k) {
          split(lc[x], k, a, lc[x]);
          b = x;
          pull(a); pull(b);
          split(rc[x], k - sz[lc[x]] - 1, rc[x], b);
          pull(a); pull(b);
  // functions
  // 回傳 x 在 Treap 中的位置
  int get pos(int x){
      vector<int> sta;
      while (fa[x]){
          sta.push back(x);
          x = fa[x];
      while (sta.size()){
          push(x);
          x = sta.back();
          sta.pop_back();
      push(x);
      int res = sz[x] - sz[rc[x]];
      while (fa[x]){
          if (rc[fa[x]]==x){
              res += sz[fa[x]]-sz[x];
          x = fa[x];
81
82
      return res;
85 // 1-based <前 [1, L-1] 個元素, [L, r] 個元素, [r+1, n] 個元
86 array<int, 3> cut(int x, int 1, int r){
      array<int, 3> ret;
      split(x, r, ret[1], ret[2]);
      split(ret[1], 1-1, ret[0], ret[1]);
```

3.12 Trie [b6475c]

```
1 struct Trie{
      struct Data{
          int nxt[2]={0, 0};
      };
      int sz=0;
      vector<Data> arr;
      void init(int n){
          arr.resize(n);
      void insert(int n){
          int now=0;
          for (int i=N; i>=0; i--){
              int v=(n>>i)&1;
              if (!arr[now].nxt[v]){
                  arr[now].nxt[v]=++sz;
              now=arr[now].nxt[v];
      }
      int query(int n){
          int now=0, ret=0;
          for (int i=N ; i>=0 ; i--){
              int v=(n>>i)&1;
              if (arr[now].nxt[1-v]){
                  ret+=(1<<i);
                  now=arr[now].nxt[1-v];
              }else if (arr[now].nxt[v]){
                  now=arr[now].nxt[v];
              }else{
                  return ret;
          return ret;
40 } tr;
```

4 Dynamic-Programming

4.1 Digit DP [133f00]

```
#include <bits/stdc++.h>
using namespace std;
long long 1, r;
long long dp[20][10][2][2]; // dp[pos][pre][limit] = 後 pos
    位 pos 前一位是 pre (是/否)有上界 (是/否)有前綴零
long long memorize_search(string &s, int pos, int pre, bool
    limit, bool lead){
    // 已經被找過了,直接回傳值
   if (dp[pos][pre][limit][lead]!=-1) return dp[pos][pre][
        limit][lead];
    // 已經搜尋完畢,紀錄答案並回傳
   if (pos==(int)s.size()){
       return dp[pos][pre][limit][lead] = 1;
   // 枚舉目前的位數數字是多少
    long long ans = 0;
    for (int now=0 ; now<=(limit ? s[pos]-'0' : 9) ; now++){</pre>
       if (now==pre){
           // 1~9 絕對不能連續出現
           if (pre!=0) continue;
           // 如果已經不在前綴零的範圍內·Ø 不能連續出現
           if (lead==false) continue;
       ans += memorize_search(s, pos+1, now, limit&(now==(s[
           pos]-'0')), lead&(now==0));
   // 已經搜尋完畢,紀錄答案並回傳
    return dp[pos][pre][limit][lead] = ans;
// 回傳 [0, n] 有多少數字符合條件
long long find_answer(long long n){
   memset(dp, -1, sizeof(dp));
   string tmp = to string(n);
    return memorize_search(tmp, 0, 0, true, true);
int main(){
    // input
    cin >> 1 >> r;
   // output - 計算 [L, r] 有多少數字任意兩個位數都不相同
    cout << find_answer(r)-find_answer(l-1) << "\n";</pre>
    return 0;
```

4.2 Knaspack On Tree [df69b1]

```
11// 需要重構、需要增加註解
  #include <bits/stdc++.h>
  #define F first
  #define S second
  #define all(x) begin(x), end(x)
  using namespace std;
  #define chmax(a, b) (a) = (a) < (b) ? (b) : (a)
  #define chmin(a, b) (a) = (a) < (b) ? (a) : (b)
  #define 11 long long
  #define FOR(i, a, b) for (int i = a; i <= b; i++)
  int N, W, cur;
  vector<int> w, v, sz;
  vector<vector<int>> adj, dp;
  void dfs(int x) {
      sz[x] = 1;
      for (int i : adj[x]) dfs(i), sz[x] += sz[i];
      cur++;
      // choose x
      for (int i=w[x]; i<=W; i++){</pre>
          dp[cur][i] = dp[cur - 1][i - w[x]] + v[x];
26
      // not choose x
      for (int i=0 ; i<=W ; i++){</pre>
           chmax(dp[cur][i], dp[cur - sz[x]][i]);
30
31
  signed main() {
      cin >> N >> W;
      adj.resize(N + 1);
      w.assign(N + 1, 0);
      v.assign(N + 1, 0);
      sz.assign(N + 1, 0);
      dp.assign(N + 2, vector < int > (W + 1, 0));
      for (int i=1; i<=N; i++){</pre>
           int p; cin >> p;
           adj[p].push_back(i);
      for (int i=1; i<=N; i++) cin >> w[i];
      for (int i=1; i<=N; i++) cin >> v[i];
      cout \langle\langle dp[N + 1][W] \langle\langle ' \rangle n';
```

4.3 SOS DP [8dfa8b]

```
1 // 總時間複雜度為 O(n 2^n)
2 // 計算 dp[i] = i 所有 bit mask 子集的和
3 for (int i=0; i<n; i++){
4 for (int mask=0; mask<(1<<n); mask++){
5 if ((mask>:i)&1){
6 dp[mask] += dp[mask^(1<<ii)];
7 }
8 }
9 }</pre>
```

Integer Partition

```
dp[i][x] = 要將整數 x 拆成 i 堆的「組合數」
dp[i+1][x+1]+=dp[i][x] ( 創造新的一堆 )
dp[i][x+i]+=dp[i][x] (把每一堆都增加 1)
```

Geometry

5.1 Geometry Struct [d9966f]

```
using ld = double;
3 // 判斷數值正負: {1:正數,0:零,-1:負數}
4 int sign(long long x) {return (x >= 0) ? ((bool)x) : -1; }
s = \frac{1}{1}  int sign(ld x) {return (abs(x) < 1e-9) ? 0 : (x>0 ? 1 : -1);}
 template<typename T>
8 struct point {
    T x, y;
     point() {}
     point(const T &x, const T &y) : x(x), y(y) {}
     explicit operator point<ld>() {return point<ld>(x, y); }
     // A [6357c4]
     point operator+(point b) {return {x+b.x, y+b.y}; }
     point operator-(point b) {return {x-b.x, y-b.y}; }
     point operator*(T b) {return {x*b, y*b}; }
     point operator/(T b) {return {x/b, y/b}; }
     bool operator==(point b) {return x==b.x && y==b.y; }
     T operator*(point b) {return x * b.x + y * b.y; }
     T operator^(point b) {return x * b.y - y * b.x; }
     // B [c415da]
     // 逆時針極角排序
     bool side() const{return (y == 0) ? (x > 0) : (y < 0); }
     bool operator<(const point &b) const {</pre>
         return side() == b.side() ?
             (x*b.y > b.x*y) : side() < b.side();
     friend ostream& operator<<(ostream& os, point p) {</pre>
         return os << "(" << p.x << ", " << p.y << ")";
     // 判斷 ab 到 ac 的方向: {1:逆時鐘,0:重疊,-1:順時鐘}
     friend int ori(point a, point b, point c) {
         return sign((b-a)^(c-a));
     friend bool btw(point a, point b, point c) {
         return ori(a, b, c) == 0 && sign((a-c)*(b-c)) <= 0;
     // 判斷線段 ab, cd 是否相交
     friend bool banana(point a, point b, point c, point d) {
         if (btw(a, b, c) || btw(a, b, d)
             || btw(c, d, a) || btw(c, d, b)) return true;
         int u = ori(a, b, c) * ori(a, b, d);
         int v = ori(c, d, a) * ori(c, d, b);
         return u < 0 && v < 0;
     } // C [09fd7c]
     // 判斷 "射線 ab" 與 "線段 cd" 是否相交
     friend bool rayHitSeg(point a, point b, point c, point d) {
         if (a == b) return btw(c, d, a); // Special case
```

```
if (((a - b) ^ (c - d)) == 0) {
           return btw(a, c, b) || btw(a, d, b) || banana(a, 115
                b, c, d):
       point u = b - a, v = d - c, s = c - a;
       return sign(s ^{\circ} v) * sign(u ^{\circ} v) >= 0 && sign(s ^{\circ} u) 119
            * sign(u ^ v) >= 0 \&\& abs(s ^ u) <= abs(u ^ v); 120
   } // D [db541a]
    // 旋轉 Arg(b) 的角度(小心溢位)
    point rotate(point b){return \{x*b.x-y*b.y, x*b.y+y*b.x\};} 123
    // 回傳極座標角度·值域:[-\pi, +\pi]
    friend ld Arg(point b) {
       return (b.x != 0 || b.y != 0) ? atan2(b.y, b.x) : 0; 126 // 可以在有 n 個點的簡單多邊形內 用 O(n) 判斷一個點:
    friend T abs2(point b) {return b * b; }
template<typename T>
struct line {
   /*----*/
    point<T> p1, p2;
    // ax + by + c = 0
                                                            135
   T a, b, c; //|a|, |b| \le 2C, |c| \le 8C^2
    line() {}
    line(const point<T> &x, const point<T> &y) : p1(x), p2(y){ 138
    void build() {
       a = p1.y - p2.y;
       b = p2.x - p1.x;
       c = (-a*p1.x)-b*p1.y;
   } // E [683239]
    // 判斷點和有向直線的關係: {1:左邊,0:在線上,-1:右邊}
    int ori(point<T> &p) {
       return sign((p2-p1) ^ (p-p1));
    // 判斷直線斜率是否相同
    bool parallel(line &1) {
       return ((p1-p2) ^ (l.p1-l.p2)) == 0;
    // 兩直線交點
                                                            153
    point<ld> line_intersection(line &1) {
                                                            154
       using P = point<ld>;
                                                            155
       point < T > u = p2-p1, v = 1.p2-1.p1, s = 1.p1-p1;
       return P(p1) + P(u) * ((ld(s^v)) / (u^v));
};
template<typename T>
struct polygon {
    vector<point<T>> v;
    polygon(const vector<point<T>> &u) : v(u) {}
    // simple 為 true 的時候會回傳任意三點不共線的凸包
    void make convex hull(int simple) {
       auto \overline{cmp} = [\overline{\&}](point<T> \&p, point<T> \&q) {
           return (p.x == q.x)? (p.y < q.y): (p.x < q.x);
       simple = (bool)simple;
       sort(v.begin(), v.end(), cmp);
       v.resize(unique(v.begin(), v.end()) - v.begin());
       if (v.size() <= 1) return;</pre>
       vector<point<T>> hull;
                                                            175
       for (int t = 0; t < 2; ++t){
                                                            176
           int sz = hull.size():
```

```
for (auto &i:v) {
                 while (hull.size() >= sz+2 && ori(hull[hull.
                       size()-2], hull.back(), i) < simple) {</pre>
                     hull.pop back();
                 hull.push back(i);
             hull.pop back();
             reverse(v.begin(), v.end());
         swap(hull, v);
      } // F [2bb3ef]
127 // {1:在多邊形內,0:在多邊形上,-1:在多邊形外}
      int in polygon(point<T> a){
          const T MAX POS = 1e9 + 5; // [記得修改] 座標的最大值
         point<T> pre = v.back(), b(MAX_POS, a.y + 1);
          int cnt = 0:
          for (auto &i:v) {
             if (btw(pre, i, a)) return 0;
             if (banana(a, b, pre, i)) cnt++;
             pre = i;
         return cnt%2 ? 1 : -1;
      } // G [f11340]
  /// 警告:以下所有凸包專用的函式都只接受逆時針排序且任三點不
       共線的凸包 ///
|142| // 可以在有 n 個點的凸包內 \cdot 用 O(\log n) 判斷一個點:
143 // {1:在凸包內, 0:在凸包邊上, -1:在凸包外}
      int in_convex(point<T> p) {
          int n = v.size();
          int a = ori(v[0], v[1], p), b = ori(v[0], v[n-1], p);
          if (a < 0 || b > 0) return -1;
         if (btw(v[0], v[1], p)) return 0;
          if (btw(v[0], v[n - 1], p)) return 0;
         int l = 1, r = n - 1, mid;
          while (1 + 1 < r) {
             mid = (1 + r) >> 1;
             if (ori(v[0], v[mid], p) >= 0) 1 = mid;
             else r = mid;
          int k = ori(v[1], v[r], p);
         if (k <= 0) return k;</pre>
         return 1;
      } // H [e64f1e]
  // 凸包專用的環狀二分搜,回傳 0-based index
      int cycle search(auto &f) {
         int n = v.size(), l = 0, r = n;
         if (n == 1) return 0;
         bool rv = f(1, 0);
          while (r - 1 > 1) {
             int m = (1 + r) / 2;
             if (f(0, m) ? rv: f(m, (m + 1) % n)) r = m;
             else 1 = m;
         return f(1, r % n) ? 1 : r % n;
      } // I [fe2f51]
172 // 可以在有 n 個點的凸包內,用 O(Log n)判斷一條直線:
173 // {1: 穿過凸包, 0: 剛好切過凸包, -1: 沒碰到凸包}
      int line cut convex(line<T> L) {
         L.build();
         point<T> p(L.a, L.b);
```

```
auto gt = [&](int neg) {
178
              auto f = [\&](int x, int y) {
                                                              238
179
                  return sign((v[x] - v[y]) * p) == neg;
180
              };
                                                               240
181
              return -(v[cycle_search(f)] * p);
                                                              241
182
183
          T x = gt(1), y = gt(-1);
          if (L.c < x \mid | y < L.c) return -1;
          return not (L.c == x || L.c == y);
      } // J [b6a4c8]
                                                               246
      可以在有 n 個點的凸包內,用 O(Log n)判斷一個線段:
                                                               247
                                                              248
      {1: 存在一個凸包上的邊可以把這個線段切成兩半,
                                                              249
       6: 有碰到凸包但沒有任何凸包上的邊可以把它切成兩半,
190 // -1: 沒碰到凸包}
191 /// 除非線段兩端點都不在凸包邊上,否則此函數回傳 Ø 的時候不一
       定表示線段沒有通過凸包內部 ///
      int segment across convex(line<T> L) {
                                                              253
193
          L.build();
194
          point<T> p(L.a, L.b);
          auto gt = [&](int neg) {
195
                                                              255
196
              auto f = [\&](int x, int y) {
                  return sign((v[x] - v[y]) * p) == neg;
198
              };
                                                              257
              return cycle search(f);
199
                                                              258
200
                                                              259
201
          int i = gt(1), j = gt(-1), n = v.size();
                                                               260
202
          T x = -(v[i] * p), y = -(v[j] * p);
                                                              261
203
          if (L.c < x \mid | y < L.c) return -1;
                                                               262
          if (L.c == x || L.c == y) return 0;
204
205
          if (i > j) swap(i, j);
206
207
          auto g = [&](int x, int lim) {
                                                               265
208
              int now = 0, nxt;
209
              for (int i = 1 << __lg(lim); i > 0; i /= 2) {
210
                  if (now + i > lim) continue;
                                                               268
                  nxt = (x + i) % n;
211
                  if (L.ori(v[x]) * L.ori(v[nxt]) >= 0) {
212
213
                      x = nxt:
                                                               271
                      now += i;
214
                                                               272
215
              } // ↓ BE CAREFUL
216
              return -(ori(v[x], v[(x + 1) % n], L.p1) * ori(v[ 275
217
                   x], v[(x + 1) % n], L.p2));
          return max(g(i, j - i), g(j, n - (j - i)));
      } // K [b4f073]
     可以在有 n 個點的凸包內 \cdot 用 O(Log n) 判斷一個線段:
      {1: 線段上存在某一點位於凸包內部(邊上不算),
       0: 線段上存在某一點碰到凸包的邊但線段上任一點均不在凸包 282
       內部.
224 // -1: 線段完全在凸包外面 }
                                                              284
      int segment_pass_convex_interior(line<T> L) {
                                                               285
226
          if (in_convex(L.p1) == 1 || in_convex(L.p2) == 1)
               return 1;
          L.build();
          point<T> p(L.a, L.b);
228
229
          auto gt = [&](int neg) {
                                                               288
                                                              289
230
              auto f = [\&](int x, int y) {
                  return sign((v[x] - v[y]) * p) == neg;
231
                                                               290
                                                              291
232
              };
233
              return cycle_search(f);
                                                              292
234
                                                               293
                                                               294
235
          int i = gt(1), j = gt(-1), n = v.size();
          T x = -(v[i] * p), y = -(v[j] * p);
                                                               295
```

```
if (L.c < x || y < L.c) return -1;
                                                                    296
           if (L.c == x || L.c == y) return 0;
                                                                    297
                                                                    298
           if (i > j) swap(i, j);
                                                                    299
           auto g = [&](int x, int lim) {
                                                                    300
                int now = 0, nxt;
                                                                    301
                for (int i = 1 << __lg(lim); i > 0; i /= 2) {
                                                                    302
                    if (now + i > \overline{lim}) continue:
                                                                    303
                    nxt = (x + i) % n;
                                                                    304
                    if (L.ori(v[x]) * L.ori(v[nxt]) > 0) {
                                                                    305
                        x = nxt:
                                                                    306
                        now += i;
                } // ↓ BE CAREFUL
                return -(ori(v[x], v[(x + 1) % n], L.p1) * ori(v[ 310
                    x], v[(x + 1) % n], L.p2));
                                                                    312
           int ret = \max(g(i, j - i), g(j, n - (j - i)));
                                                                    313
           return (ret == 0) ? (in_convex(L.p1) == 0 &&
                                                                    314
                in_convex(L.p2) == 0) : ret;
       } // L [5f45ca]
256 // 回傳點過凸包的兩條切線的切點的 0-based index (不保證兩條
                                                                   315
                                                                    316
        切線的順逆時針關係)
                                                                    317
       pair<int,int> convex_tangent_point(point<T> p) {
                                                                    318
           int n = v.size(), z = -1, edg = -1;
                                                                    319
           auto gt = [&](int neg) {
                                                                    320
                auto check = [&](int x) {
                                                                    321
                    if (v[x] == p) z = x;
                    if (btw(v[x], v[(x + 1) % n], p)) edg = x;
                                                                    322
                    if (btw(v[(x + n - 1) % n], v[x], p)) edg = (
                         x + n - 1) % n;
                };
                                                                    324
                auto f = [\&](int x, int y) {
                                                                    325
                    check(x); check(y);
                                                                    326
                    return ori(p, v[x], v[y]) == neg;
                                                                    327
                };
                                                                    328
                return cycle_search(f);
                                                                    329
                                                                    330
           int x = gt(1), y = gt(-1);
                                                                    331
           if (z != -1) {
                                                                    332
                return \{(z + n - 1) \% n, (z + 1) \% n\};
                                                                    333
                                                                    334
           else if (edg != -1) {
                                                                    335
                return {edg, (edg + 1) % n};
                                                                    336
                                                                    337
           else {
                return {x, y};
                                                                    338
                                                                    339
       } // M [a6f66b]
                                                                    340
       friend int halfplane intersection(vector<line<T>> &s,
                                                                    341
            polygon<T> &P) {
                                                                    342
           auto angle cmp = [&](line<T> &A, line<T> &B) {
                point < T > a = A.p2-A.p1, b = B.p2-B.p1;
                return (a < b);</pre>
           sort(s.begin(), s.end(), angle cmp); // 線段左側為該
                 線段半平面
           int L, R, n = s.size();
           vector<point<T>> px(n);
           vector<line<T>> q(n);
           q[L = R = 0] = s[0];
           for(int i = 1; i < n; ++i) {</pre>
               while (L < R \&\& s[i].ori(px[R-1]) <= 0) --R;
                while(L < R && s[i].ori(px[L]) <= 0) ++L;
```

q[++R] = s[i];

```
if(q[R].parallel(q[R-1])) {
                                               if(q[R].ori(s[i].p1) > 0) q[R] = s[i];
                                     if(L<R) px[R-1] = q[R-1].line_intersection(q[R]);</pre>
                           while (L < R \&\& q[L].ori(px[R-1]) <= 0) --R;
                           P.v.clear():
                           if(R - L <= 1) return 0;</pre>
                           px[R] = q[R].line_intersection(q[L]);
                           for(int i = L; i <= R; ++i) P.v.push_back(px[i]);</pre>
                           return R - L + 1;
                 } // N [102d48]
      };
309
      struct Cir {
                 point<ld> o; ld r;
                 friend ostream& operator<<(ostream& os, Cir c) {</pre>
                           return os \langle\langle "(x" \langle\langle "+-"[c.o.x \rangle= 0] \langle\langle abs(c.o.x) \rangle
                                       ((x'')^2 + (y'')^2 + (y'
                                       <<")^2 = " << c.r * c.r;
                 bool covers(Cir b) {
                           return sqrt((ld)abs2(o - b.o)) + b.r <= r;</pre>
                 vector<point<ld>>> Cir_intersect(Cir c) {
                           1d d2 = abs2(o - c.o), d = sqrt(d2);
                           if (d < max(r, c.r) - min(r, c.r) \mid | d > r + c.r)
                                       return {};
                           auto sqdf = [\&](1d x, 1d y) \{ return x*x - y*y; \};
                           point < ld > u = (o + c.o) / 2 + (o - c.o) * (sqdf(c.r, ))
                                       r) / (2 * d2));
                           1d A = sqrt(sqdf(r + d, c.r) * sqdf(c.r, d - r));
                           point<ld> v = (c.o - o).rotate({0,1}) * A / (2 * d2);
                           if (sign(v.x) == 0 \&\& sign(v.y) == 0) return \{u\};
                           return \{u - v, u + v\};
                 } // O [330a1c]
                 auto point_tangent(point<ld> p) {
                           vector<point<ld>> res;
                           1d d_sq = abs2(p - o);
                           if (sign(d_sq - r * r) <= 0) {</pre>
                                      res.pb(p + (p - o).rotate(\{0, 1\}));
                           } else if (d_sq > r * r) {
                                     ld s = d_sq - r * r;
                                     point < ld > v = p + (o - p) * s / d_sq;
                                     point < ld > u = (o - p).rotate({0, 1}) * sqrt(s) *
                                                 r / d sq;
                                      res.pb(v + u);
                                     res.pb(v - u);
                           return res;
                 } // P [0067e6]
343 };
```

5.2 Geometry Struct 3D [4a50c9]

```
8 struct pt3 {
      T x, y, z;
      pt3(){}
      pt3(const T &x, const T &y, const T &z):x(x),y(y),z(z){}
      explicit operator pt3<ld>() {return pt3<ld>(x, y, z); }
      pt3 operator+(pt3 b) {return {x+b.x, y+b.y, z+b.z}; }
      pt3 operator-(pt3 b) {return {x-b.x, y-b.y, z-b.z}; }
      pt3 operator*(T b) {return {x * b, y * b, z * b}; }
      pt3 operator/(T b) {return {x / b, y / b, z / b}; }
      bool operator==(pt3 b){return x==b.x&&y==b.y&&z==b.z;}
      T operator*(pt3 b) {return x*b.x+y*b.y+z*b.z; }
     pt3 operator^(pt3 b) {
          return pt3(y * b.z - z * b.y,
                    z * b.x - x * b.z,
                    x * b.y - y * b.x);
      friend T abs2(pt3 b) {return b * b; }
      friend T len (pt3 b) {return sqrt(abs2(b)); }
     friend ostream& operator<<((ostream& os, pt3 p) {</pre>
         return os << "(" << p.x << ", " << p.y << ", " << p.z << ")";
34 template<typename T>
35 struct face {
     int a, b, c; // 三角形在 vector 裡面的 index
     pt3<T> q; // 面積向量(朝外)
40 /// 警告;ν 在過程中可能被修改,回傳的 face 以修改後的為準
41 // O(n²),最多只有 2n-4 個面
42 // 當凸包退化時會回傳空的凸包,否則回傳凸包上的每個面
43 template<typename T>
44 vector<face<T>> hull3(vector<pt3<T>> &v) {
      int n = v.size();
     if (n < 3) return {};
      // don't use "==" when you use ld
      sort(all(v), [&](pt3<T> &p, pt3<T> &q) {
          return sign(p.x - q.x)? (p.x < q.x):
           (sign(p.y - q.y) ? p.y < q.y : p.z < q.z);
      v.resize(unique(v.begin(), v.end()) - v.begin());
      for (int i = 2; i <= n; ++i) {
          if (i == n) return {};
          if (sign(len(((v[1] - v[0]) ^ (v[i] - v[0]))))) {
              swap(v[2], v[i]);
              break;
      pt3<T> tmp_q = (v[1] - v[0]) ^ (v[2] - v[0]);
      for (int i = 3; i <= n; ++i) {
          if (i == n) return {};
          if (sign((v[i] - v[0]) * tmp_q)) {
              swap(v[3], v[i]);
              break:
     }
      vector<face<T>> f;
      vector<vector<int>> dead(n, vector<int>(n, true));
      auto add face = [&](int a, int b, int c) {
```

```
f.emplace_back(a, b, c, (v[b] - v[a]) ^ (v[c] - v[a]) ^ 25
    dead[a][b] = dead[b][c] = dead[c][a] = false;
add_face(0, 1, 2);
add face(0, 2, 1);
for (int i = 3; i < n; ++i) {</pre>
   vector<face<T>> f2;
    for (auto &[a, b, c, q] : f) {
        if (sign((v[i] - v[a]) * q) > 0)
            dead[a][b] = dead[b][c] = dead[c][a] = true;
        else f2.emplace_back(a, b, c, q);
   f.clear();
   for (face<T> &F : f2) {
        int arr[3] = {F.a, F.b, F.c};
        for (int j = 0; j < 3; ++j) {
            int a = arr[j], b = arr[(j + 1) % 3];
            if (dead[b][a]) add_face(b, a, i);
   f.insert(f.end(), all(f2));
return f;
```

5.3 Pick's Theorem

給定頂點坐標均是整點的簡單多邊形,面積 = 內部格點數 + 邊上格點數/2-1

Graph

6.1 2-SAT [5a6317]

```
1 struct TWO SAT {
     int n, N;
      vector<vector<int>> G, rev G;
      deque<bool> used:
      vector<int> order, comp;
      deque<bool> assignment:
      void init(int _n) {
         n = _n;
         N = n * 2;
         G.resize(N + 5);
         rev G.resize(N + 5);
      void dfs1(int v) {
          used[v] = true;
          for (int u : G[v]) {
             if (!used[u])
                 dfs1(u);
         order.push_back(v);
      void dfs2(int v, int cl) {
         comp[v] = cl;
          for (int u : rev G[v]) {
              if (comp[u] == -1)
```

```
dfs2(u, c1);
      bool solve() {
          order.clear();
          used.assign(N, false);
          for (int \bar{i} = 0; i < N; ++i) {
              if (!used[i])
                  dfs1(i):
          comp.assign(N, -1);
          for (int i = 0, j = 0; i < N; ++i) {
              int v = order[N - i - 1];
              if (comp[v] == -1)
                  dfs2(v, j++);
          assignment.assign(n, false);
          for (int i = 0; i < N; i += 2) {
              if (comp[i] == comp[i + 1])
                  return false:
              assignment[i / 2] = (comp[i] > comp[i + 1]);
          return true;
      // A or B 都是 0-based
      void add disjunction(int a, bool na, int b, bool nb) {
          // na is true => ~a, na is false => a
          // nb is true => ~b, nb is false => b
          a = 2 * a ^ na:
          b = 2 * b ^ nb;
          int neg a = a ^ 1;
          int neg b = b ^ 1;
          G[neg_a].push_back(b);
          G[neg_b].push_back(a);
          rev_G[b].push_back(neg_a);
          rev_G[a].push_back(neg_b);
          return;
      void get result(vector<int>& res) {
          res.clear();
          for (int i = 0; i < n; i++)</pre>
              res.push back(assignment[i]);
68 };
```

6.2 Augment Path [f8a5dd]

61

62

63

67

```
1 struct AugmentPath{
      int n, m;
      vector<vector<int>> G;
      vector<int> mx, my;
      vector<int> visx, visy;
      int stamp;
      AugmentPath(int _n, int _m) : n(_n), m(_m), G(n), mx(n,
           -1), my(m, -1), visx(n), visy(n){
          stamp = 0;
      void add(int x, int y){
13
          G[x].push_back(y);
```

```
// bb03e2
      bool dfs1(int now){
          visx[now] = stamp;
          for (auto x : G[now]){
              if (my[x]==-1){
                  mx[now] = x;
                  my[x] = now;
                  return true;
          for (auto x : G[now]){
              if (visx[my[x]]!=stamp && dfs1(my[x])){
                  mx[now] = x;
                  my[x] = now;
                  return true;
          return false;
      vector<pair<int, int>> find_max_matching(){
          vector<pair<int, int>> ret;
          while (true){
              stamp++;
              int tmp = 0;
              for (int i=0 ; i<n ; i++){</pre>
                  if (mx[i]==-1 && dfs1(i)) tmp++;
              if (tmp==0) break;
          for (int i=0 ; i<n ; i++){</pre>
              if (mx[i]!=-1){
                  ret.push back({i, mx[i]});
          return ret;
      // 645577
      void dfs2(int now){
          visx[now] = true;
          for (auto x : G[now]){
              if (my[x]!=-1 && visy[x]==false){
                  visv[x] = true;
                  dfs2(my[x]);
      }
      // 要先執行 find max matching 一次
      vector<pair<int, int>> find_min_vertex_cover(){
          fill(visx.begin(), visx.end(), false);
          fill(visy.begin(), visy.end(), false);
72
          vector<pair<int, int>> ret;
          for (int i=0 ; i<n ; i++){</pre>
              if (mx[i]==-1) dfs2(i);
          for (int i=0 ; i<n ; i++){</pre>
              if (visx[i]==false) ret.push_back({1, i});
```

6.3 C3C4 [d00465]

```
1 // 0-based
 void C3C4(vector<int> deg, vector<array<int, 2>> edges){
      int N = deg.size();
     int M = deges.size();
      vector<int> ord(N), rk(N);
      iota(ord.begin(), ord.end(), 0);
      sort(ord.begin(), ord.end(), [&](int x, int y) { return
          deg[x] > deg[y]; });
      for (int i=0 ; i<N ; i++) rk[ord[i]] = i;</pre>
      vector<vector<int>>> D(N), adj(N);
      for (auto [u, v] : e) {
          if (rk[u] > rk[v]) swap(u, v);
         D[u].emplace_back(v);
          adj[u].emplace_back(v);
         adj[v].emplace back(u);
      vector<int> vis(N);
     int c3 = 0, c4 = 0;
      for (int x : ord) { // c3
          for (int y : D[x]) vis[y] = 1;
         for (int y : D[x]) for (int z : D[y]){
              c3 += vis[z]; // xyz is C3
          for (int y : D[x]) vis[y] = 0;
     for (int x : ord) { // c4
          for (int y : D[x]) for (int z : adj[y])
              if (rk[z] > rk[x]) c4 += vis[z]++;
          for (int y : D[x]) for (int z : adj[y])
              if (rk[z] > rk[x]) --vis[z];
         // both are O(M*sqrt(M)), test @ 2022 CCPC guangzhou
      cout << c4 << "\n";
```

6.4 Cut BCC [2af809]

```
#include <bits/stdc++.h>
using namespace std;

const int N = 200005;
vector <int> G[N];
int low[N], depth[N];
bool vis[N];
vector <vector <int>> bcc;
stack <int>> stk;
```

```
void dfs(int v, int p) {
      stk.push(v):
      vis[v] = true;
      low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
      for (int u : G[v]) {
          if (u == p) continue;
          if (!vis[u]) {
              /// (v, u) 是樹邊
              dfs(u, v);
              low[v] = min(low[v], low[u]);
              /// u 無法在不經過父邊的情況走到 v 的祖先
21
              if (low[u] >= depth[v]) {
23
                  bcc.emplace back();
                  while (stk.top() != u) {
24
                      bcc.back().push_back(stk.top());
26
                      stk.pop();
27
                  bcc.back().push_back(stk.top());
29
                  stk.pop();
                  bcc.back().push_back(v);
31
32
          } else {
              /// (v, u) 是回邊
33
34
              low[v] = min(low[v], depth[u]);
35
36
37 }
```

6.5 Dinic [961b34]

```
1 // 一般圖: O(EV<sup>2</sup>)
2 // 二分圖: O(E√V)
3 struct Flow{
       using T = int; // 可以換成別的型別
      struct Edge{
           int v; T rc; int rid;
      };
      vector<vector<Edge>> G;
      void add(int u, int v, T c){
           G[u].push back({v, c, G[v].size()});
          G[v].push_back({u, 0, G[u].size()-1});
13
      vector<int> dis, it;
      Flow(int n){
          G.resize(n);
          dis.resize(n);
17
           it.resize(n);
19
20
      // ce56d6
      T dfs(int u, int t, T f){
           if (u == t || f == 0) return f;
           for (int &i=it[u] ; i<G[u].size() ; i++){</pre>
               auto &[v, rc, rid] = G[u][i];
               if (dis[v]!=dis[u]+1) continue;
26
27
               T df = dfs(v, t, min(f, rc));
               if (df <= 0) continue;</pre>
28
               rc -= df:
29
30
               G[v][rid].rc += df;
31
               return df;
```

```
return 0;
// e22e39
T flow(int s, int t){
    T ans = 0;
    while (true){
        fill(dis.begin(), dis.end(), INF);
        queue<int> q;
        q.push(s);
        dis[s] = 0;
        while (q.size()){
            int u = q.front(); q.pop();
            for (auto [v, rc, rid] : G[u]){
                if (rc <= 0 || dis[v] < INF) continue;</pre>
                dis[v] = dis[u] + 1;
                q.push(v);
        if (dis[t]==INF) break;
        fill(it.begin(), it.end(), 0);
        while (true){
            T df = dfs(s, t, INF);
            if (df <= 0) break;</pre>
            ans += df:
    return ans;
// the code below constructs minimum cut
void dfs mincut(int now, vector<bool> &vis){
    vis[now] = true;
    for (auto &[v, rc, rid] : G[now]){
        if (vis[v] == false && rc > 0){
            dfs_mincut(v, vis);
}
vector<pair<int, int>> construct(int n, int s, vector<</pre>
     pair<int,int>> &E){
    // E is G without capacity
    vector<bool> vis(n);
    dfs mincut(s, vis);
    vector<pair<int, int>> ret;
    for (auto &[u, v] : E){
        if (vis[u] == true && vis[v] == false){
            ret.emplace back(u, v);
    return ret;
```

6.6 Dominator Tree [52b249]

```
2| 全部都是 0-based
```

```
一開始要初始化 G(N, root),代表有 N 個節點,根是 root
6 | G[i] = i 的 idom · 也就是從 root 走到 i 時 · 一定要走到的點且離
      i 最近
 struct DominatorTree{
     int N;
     vector<vector<int>> G;
     vector<vector<int>> buckets, rg;
     // dfn[x] = the DFS otder of x
     // rev[x] = the vertex with DFS order x
     // par[x] = the parent of x
     vector<int> dfn, rev, par;
     vector<int> sdom, dom, idom;
     vector<int> fa, val;
     int stamp;
     int root;
     int operator [] (int x){
         return idom[x];
     DominatorTree(int _N, int _root) :
         N(_N),
         G(N), buckets(N), rg(N),
         dfn(N, -1), rev(N, -1), par(N, -1),
         sdom(N, -1), dom(N, -1), idom(N, -1),
         fa(N, -1), val(N, -1)
         stamp = 0;
         root = _root;
     void add_edge(int u, int v){
         G[u].push back(v);
     void dfs(int x){
         rev[dfn[x] = stamp] = x;
         fa[stamp] = sdom[stamp] = val[stamp] = stamp;
         for (int u : G[x]){
             if (dfn[u]==-1){
                 dfs(u);
                 par[dfn[u]] = dfn[x];
             rg[dfn[u]].push_back(dfn[x]);
     int eval(int x, bool first){
         if (fa[x]==x) return !first ? -1 : x;
         int p = eval(fa[x], false);
         if (p==-1) return x;
         if (sdom[val[x]]>sdom[val[fa[x]]]) val[x] = val[fa[x
              ]];
         fa[x] = p;
         return !first ? p : val[x];
     void link(int x, int y){
```

3 G 要是有向無權圖

62

63

```
fa[x] = y;
      }
      void build(){
          dfs(root);
          for (int x=stamp-1 ; x>=0 ; x--){
               for (int y : rg[x]){
                   sdom[x] = min(sdom[x], sdom[eval(y, true)]);
               if (x>0) buckets[sdom[x]].push back(x);
               for (int u : buckets[x]){
                   int p = eval(u, true);
                   if (sdom[p]==x) dom[u] = x;
                   else dom[u] = p;
              if (x>0) link(x, par[x]);
          idom[root] = root;
          for (int x=1 ; x<stamp ; x++){</pre>
               if (sdom[x]!=dom[x]) dom[x] = dom[dom[x]];
          for (int i=1; i<stamp; i++) idom[rev[i]] = rev[dom[</pre>
90
91 };
```

6.7 EdgeBCC [d09eb1]

```
1 // d09eb1
2 // 0-based · 支援重邊
  struct EdgeBCC{
      int n, m, dep, sz;
      vector<vector<pair<int, int>>> G;
      vector<vector<int>> bcc;
      vector<int> dfn, low, stk, isBridge, bccId;
      vector<pair<int, int>> edge, bridge;
      EdgeBCC(int n): n(n), m(0), sz(0), dfn(n), low(n), G(n)
           ), bcc(n), bccId(n) {}
      void add edge(int u, int v) {
          edge.push_back({u, v});
          G[u].push back({v, m});
          G[v].push_back({u, m++});
      void dfs(int now, int pre) {
          dfn[now] = low[now] = ++dep;
          stk.push back(now);
          for (auto [x, id] : G[now]){
              if (!dfn[x]){
                  dfs(x, id);
                  low[now] = min(low[now], low[x]);
26
              }else if (id!=pre){
                  low[now] = min(low[now], dfn[x]);
27
28
29
30
31
          if (low[now]==dfn[now]){
              if (pre!=-1) isBridge[pre] = true;
```

```
int u;
            u = stk.back():
            stk.pop_back();
            bcc[sz].push_back(u);
            bccId[u] = sz;
        } while (u!=now);
        SZ++:
}
void get_bcc() {
    isBridge.assign(m, 0);
    dep = 0;
    for (int i=0 ; i<n ; i++){</pre>
        if (!dfn[i]) dfs(i, -1);
    for (int i=0 ; i<m ; i++){</pre>
        if (isBridge[i]){
            bridge.push_back({edge[i].first , edge[i].
                 second});
```

6.8 EnumeratePlanarFace [e70ee1]

```
1 // 0-based
2 struct PlanarGraph{
     int n, m, id;
     vector<point<int>> v;
     vector<vector<pair<int, int>>> G;
     vector<int> conv, nxt, vis;
     PlanarGraph(int n, int m, vector<point<int>> v) :
     n(n), m(m), id(0),
     v(v), G(n),
     conv(2*m), nxt(2*m), vis(2*m) {}
     void add edge(int x, int y){
         G[x].push back({y, 2*id});
         G[y].push_back({x, 2*id+1});
         conv[2*id] = x;
         conv[2*id+1] = y;
         id++;
     vector<int> enumerate face(){
         for (int i=0 ; i<n ; i++){</pre>
             sort(G[i].begin(), G[i].end(), [&](pair<int, int>
                    a, pair<int, int> b){
                  return (v[a.first]-v[i])<(v[b.first]-v[i]);</pre>
             });
             int sz = G[i].size(), pre = sz-1;
             for (int j=0 ; j<sz ; j++){</pre>
                  nxt[G[i][pre].second] = G[i][j].second^1;
                  pre = j;
```

6.9 HLD [f57ec6]

return ret;

vector<int> ret;

for (int i=0 ; i<2*m ; i++){</pre>

vector<int> pt;

while (!vis[now]){

int area = 0, now = i;

vis[now] = true;

now = nxt[now];

pt.push_back(pt.front());

// pt = face boundary

ret.push_back(area);

// pt is outer face

if (area>0){

}else{

pt.push back(conv[now]);

for (int i=0 ; i+1<pt.size() ; i++){</pre>

area -= (v[pt[i]]^v[pt[i+1]]);

if (vis[i]==false){

```
| #include <bits/stdc++.h>
 #define int long long
 using namespace std;
 const int N = 100005;
 vector <int> G[N];
 struct HLD {
     vector<int> pa, sz, depth, mxson, topf, id;
     int n, idcnt = 0;
     HLD(int _n) : n(_n), pa(_n + 1), sz(_n + 1), depth(_n +
          1), mxson(_n + 1), topf(_n + 1), id(_n + 1) {}
     void dfs1(int v = 1, int p = -1) {
         pa[v] = p; sz[v] = 1; mxson[v] = 0;
         depth[v] = (p == -1 ? 0 : depth[p] + 1);
         for (int u : G[v]) {
             if (u == p) continue;
             dfs1(u, v);
             sz[v] += sz[u];
             if (sz[u] > sz[mxson[v]]) mxson[v] = u;
     void dfs2(int v = 1, int top = 1) {
         id[v] = ++idcnt;
         topf[v] = top;
         if (mxson[v]) dfs2(mxson[v], top);
         for (int u : G[v]) {
             if (u == mxson[v] || u == pa[v]) continue;
             dfs2(u, u);
     // query 為區間資料結構
     int path_query(int a, int b) {
```

```
int res = 0;
33
          while (topf[a] != topf[b]) { /// 若不在同一條鍊上
              if (depth[topf[a]] < depth[topf[b]]) swap(a, b);</pre>
              res = max(res, 011); // query : L = id[topf[a]],
                   r = id[a]
              a = pa[topf[a]];
37
          /// 此時已在同一條鍊上
38
39
          if (depth[a] < depth[b]) swap(a, b);</pre>
40
          res = max(res, 011); // query : l = id[b], r = id[a]
41
          return res;
42
43 };
```

6.10 Kosaraju [c7d5aa]

```
1 /* c7d5aa
  給定一個有向圖,迴回傳縮點後的圖、SCC 的資訊
  所有點都以 based-0 編號
  函式:
  SCC_compress G(n): 宣告一個有 n 個點的圖
7 | .add edge(u, v): 加上一條邊 u -> v
8 .compress: O(n log n) 計算 G3、SCC、SCC id 的資訊,並把縮點後
      的結果存在 result 裡
10 SCC[i] = 某個 SCC 中的所有點
11 SCC_id[i] = 第 i 個點在第幾個 SCC
12 */
  struct SCC compress{
     int N, M, sz;
     vector<vector<int>>> G, inv G, result;
     vector<pair<int, int>> edges;
     vector<bool> vis;
     vector<int> order;
      vector<vector<int>> SCC;
     vector<int> SCC id;
      SCC compress(int N):
     N(N), M(0), sz(0),
      G(N), inv_G(N),
      vis(N), SCC id(N)
27
28
     vector<int> operator [] (int x){
         return result[x];
30
32
33
     void add_edge(int u, int v){
         G[u].push back(v);
         inv_G[v].push_back(u);
         edges.push_back({u, v});
37
         M++;
      void dfs1(vector<vector<int>> &G, int now){
         vis[now] = 1;
         for (auto x : G[now]) if (!vis[x]) dfs1(G, x);
         order.push back(now);
```

```
void dfs2(vector<vector<int>> &G, int now){
    SCC id[now] = SCC.size()-1;
    SCC.back().push_back(now);
    vis[now] = 1;
    for (auto x : G[now]) if (!vis[x]) dfs2(G, x);
void compress(){
    fill(vis.begin(), vis.end(), 0);
    for (int i=0; i<N; i++) if (!vis[i]) dfs1(G, i);</pre>
    fill(vis.begin(), vis.end(), 0);
    reverse(order.begin(), order.end());
    for (int i=0 ; i<N ; i++){</pre>
        if (!vis[order[i]]){
            SCC.push_back(vector<int>());
            dfs2(inv_G, order[i]);
    }
    result.resize(SCC.size());
    sz = SCC.size();
    for (auto [u, v] : edges){
        if (SCC_id[u]!=SCC_id[v]) result[SCC_id[u]].
             push_back(SCC_id[v]);
    for (int i=0 ; i<SCC.size() ; i++){</pre>
        sort(result[i].begin(), result[i].end());
        result[i].resize(unique(result[i].begin(), result
             [i].end())-result[i].begin());
```

6.11 Kuhn Munkres [e66c35]

```
1 // O(n^3) 找到最大權匹配
2 struct KuhnMunkres{
     int n; // max(n, m)
     vector<vector<int>> G;
     vector<int> match, lx, ly, visx, visy;
     vector<int> slack;
     int stamp = 0;
     KuhnMunkres(int n) : n(n), G(n, vector<int>(n)), lx(n),
          ly(n), slack(n), match(n), visx(n), visy(n) {}
     void add(int x, int y, int w){
         G[x][y] = max(G[x][y], w);
     bool dfs(int i, bool aug){ // aug = true 表示要更新 match
         if (visx[i]==stamp) return false;
         visx[i] = stamp;
         for (int j=0 ; j<n ; j++){</pre>
             if (visy[j]==stamp) continue;
             int d = lx[i]+ly[j]-G[i][j];
             if (d==0){
                 visy[j] = stamp;
                 if (match[j]==-1 || dfs(match[j], aug)){
```

```
if (aug){
                     match[j] = i;
                 return true;
        }else{
            slack[j] = min(slack[j], d);
    return false:
}
bool augment(){
    for (int j=0 ; j<n ; j++){</pre>
        if (visy[j]!=stamp && slack[j]==0){
             visy[j] = stamp;
             if (match[j]==-1 || dfs(match[j], false)){
                 return true;
        }
    return false;
void relabel(){
    int delta = INF;
    for (int j=0 ; j<n ; j++){</pre>
        if (visy[j]!=stamp) delta = min(delta, slack[j]);
    for (int i=0 ; i<n ; i++){</pre>
        if (visx[i]==stamp) lx[i] -= delta;
    for (int j=0 ; j<n ; j++){</pre>
        if (visy[j]==stamp) ly[j] += delta;
        else slack[j] -= delta;
}
int solve(){
    for (int i=0 ; i<n ; i++){</pre>
        lx[i] = 0;
        for (int j=0 ; j<n ; j++){</pre>
             lx[i] = max(lx[i], G[i][j]);
    fill(ly.begin(), ly.end(), 0);
    fill(match.begin(), match.end(), -1);
    for(int i = 0; i < n; i++) {</pre>
        fill(slack.begin(), slack.end(), INF);
        if(dfs(i, true)) continue;
        while(augment()==false) relabel();
        stamp++;
        dfs(i, true);
    int ans = 0;
    for (int j=0 ; j<n ; j++){</pre>
        if (match[j]!=-1){
             ans += G[match[j]][j];
```

6.12 LCA [5b6a5b]

return ans;

93

94 };

```
」 // 1-based · 可以支援森林 · 0 是超級源點 · 所有樹都要跟他建邊
  struct Tree{
      int N, M = 0, H;
      vector<int> parent, dep;
      vector<vector<int>>> G, LCA;
      Tree(int_N): N(_N+1), H(__lg(N)+1), parent(N, -1), dep(
          LCA.resize(H, vector<int>(N, 0));
      void add_edge(int u, int v){
          G[u].push_back(v);
          G[v].push back(u);
      void dfs(int now = 0, int pre = 0){
          dep[now] = dep[pre]+1;
          parent[now] = pre;
          for (auto x : G[now]){
               if (x==pre) continue;
               dfs(x, now);
      void build LCA(int root = 0){
          for (int i=0 ; i<N ; i++) LCA[0][i] = parent[i];</pre>
          for (int i=1 ; i<H ; i++){</pre>
               for (int j=0 ; j<N ; j++){</pre>
                   LCA[i][j] = LCA[i-1][LCA[i-1][j]];
      int jump(int u, int step){
          for (int i=0; i<H; i++){
               if (step&(1<<i)) u = LCA[i][u];</pre>
          return u;
      int get_LCA(int u, int v){
          if (dep[u]<dep[v]) swap(u, v);</pre>
          u = jump(u, dep[u]-dep[v]);
          if (u==v) return u;
          for (int i=H-1; i>=0; i--){
               if (LCA[i][u]!=LCA[i][v]){
                  u = LCA[i][u];
                  v = LCA[i][v];
50
52
           return parent[u];
53
55 };
```

6.13 MCMF [1b5a27]

| struct Flow {

```
struct Edge {
  int u, rc, k, rv;
vector<vector<Edge>> G;
vector<int> par, par_eid;
Flow(int n) : G(n+1), par(n+1), par_eid(n+1) {}
// v->u, capcity: c, cost: k
void add(int v, int u, int c, int k){
  G[v].push_back({u, c, k, G[u].size()});
  G[u].push_back({v, 0, -k, G[v].size()-1});
// 3701d6
int spfa(int s, int t){
  fill(par.begin(), par.end(), -1);
  vector<int> dis(par.size(), INF);
  vector<bool> in_q(par.size(), false);
  queue<int> Q;
  dis[s] = 0;
  in_q[s] = true;
  Q.push(s);
  while (!Q.empty()){
    int v = Q.front();
    0.pop();
    in q[v] = false;
    for (int i=0 ; i<G[v].size() ; i++){</pre>
      auto [u, rc, k, rv] = G[v][i];
      if (rc>0 && dis[v]+k<dis[u]){</pre>
        dis[u] = dis[v]+k;
        par[u] = v;
        par_eid[u] = i;
        if (!in q[u]) Q.push(u);
        in_q[u] = true;
  return dis[t];
// return <max flow, min cost>, 150093
pair<int, int> flow(int s, int t){
  int fl = 0, cost = 0, d;
  while ((d = spfa(s, t))<INF){</pre>
    int cur = INF:
    for (int v=t ; v!=s ; v=par[v])
      cur = min(cur, G[par[v]][par_eid[v]].rc);
    fl += cur;
    cost += d*cur;
    for (int v=t; v!=s; v=par[v]){
      G[par[v]][par_eid[v]].rc -= cur;
      G[v][G[par[v]][par_eid[v]].rv].rc += cur;
  return {fl, cost};
vector<pair<int, int>> construct(){
```

```
64
    vector<pair<int, int>> ret;
    for (int i=0; i<n; i++){
        for (auto x : G[i]){
            if (x.rc==0){
                ret.push_back({i+1, x.u-n+1});
                 break;
            }
        }
     }
    return ret;
    }
}</pre>
```

6.14 Tarjan [8b2350]

```
1 struct tarjan SCC {
     int now_T, now_SCCs;
     vector<int> dfn, low, SCC;
     stack<int> S;
     vector<vector<int>> E;
     vector<bool> vis, in stack;
      tarjan_SCC(int n) {
         init(n);
     void init(int n) {
          now_T = now_SCCs = 0;
          dfn = low = SCC = vector<int>(n);
         E = vector<vector<int>>(n);
         S = stack<int>();
          vis = in stack = vector<bool>(n);
     void add(int u, int v) {
          E[u].push back(v);
     void build() {
    for (int i = 0; i < dfn.size(); ++i) {</pre>
              if (!dfn[i]) dfs(i);
      void dfs(int v) {
          now T++;
          vis[v] = in stack[v] = true;
          dfn[v] = low[v] = now_T;
          S.push(v);
          for (auto &i:E[v]) {
              if (!vis[i]) {
                  vis[i] = true;
                  dfs(i);
                  low[v] = min(low[v], low[i]);
              else if (in_stack[i]) {
                  low[v] = min(low[v], dfn[i]);
          if (low[v] == dfn[v]) {
              int tmp;
                  tmp = S.top();
                  S.pop();
                  SCC[tmp] = now SCCs;
                  in stack[tmp] = false;
              } while (tmp != v);
```

6.15 Tarjan Find AP [1daed6]

```
1 vector<int> dep(MAX_N), low(MAX_N), AP;
  bitset<MAX N> vis;
  void dfs(int now, int pre){
      int cnt = 0;
      bool ap = 0;
      vis[now] = 1;
      low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
      for (auto x : G[now]){
          if (x==pre){
               continue;
          }else if (vis[x]==0){
              cnt++;
              dfs(x, now);
              low[now] = min(low[now], low[x]);
              if (low[x]>=dep[now]) ap=1;
          }else{
              low[now] = min(low[now], dep[x]);
      if ((now==pre && cnt>=2) || (now!=pre && ap)){
          AP.push back(now);
26
```

6.16 Tree Isomorphism [cd2bbc]

```
| #include <bits/stdc++.h>
  #pragma GCC optimize("03,unroll-loops")
  #define fastio ios::sync with stdio(0), cin.tie(0), cout.tie
  #define dbg(x) cerr << #x << " = " << x << endl
  #define int long long
  using namespace std;
  // declare
  const int MAX SIZE = 2e5+5;
  const int INF = 9e18:
  const int MOD = 1e9+7;
  const double EPS = 1e-6;
  typedef vector<vector<int>> Graph;
  typedef map<vector<int>, int> Hash;
  int n, a, b;
17 int id1, id2;
18 pair<int, int> c1, c2;
vector<int> sz1(MAX_SIZE), sz2(MAX_SIZE);
vector<int> we1(MAX_SIZE), we2(MAX_SIZE);
21 Graph g1(MAX_SIZE), g2(MAX_SIZE);
22 Hash m1, m2;
23 int testcase=0;
```

```
g2[a].push back(b);
                                                                                                                                     16 #define deb2 cerr << "deb2" << endl;
25 void centroid(Graph &g, vector<int> &s, vector<int> &w, pair< 90
                                                                                                                                     #define deb3 cerr << "deb3" << endl;
                                                                             g2[b].push back(a);
       int, int> &rec, int now, int pre){
                                                                         }
                                                                                                                                     #define deb4 cerr << "deb4" << endl:
      s[now]=1;
                                                                                                                                       #define deb5 cerr << "deb5" << endl;</pre>
      w[now]=0;
                                                                         // get tree centroid
                                                                                                                                       #define bye exit(0);
                                                                         centroid(g1, sz1, we1, c1, 1, 0);
      for (auto x : g[now]){
                                                                                                                                       using namespace std;
          if (x!=pre){
                                                                         centroid(g2, sz2, we2, c2, 1, 0);
              centroid(g, s, w, rec, x, now);
                                                                                                                                       const int mxn = (int)(2e5) + 10;
                                                                                                                                       const int mxlg = 17;
              s[now]+=s[x];
                                                                         // process
              w[now]=max(w[now], s[x]);
                                                                         int res1=0, res2=0, res3=0;
                                                                                                                                       int last_special_node = (int)(1e5) + 1;
                                                                         if (c2.second!=0){
                                                                                                                                       vector<int> E[mxn], F[mxn];
                                                                             res1=dfs(g1, m1, id1, c1.first, 0);
                                                                                                                                       struct edg{
      w[now]=max(w[now], n-s[now]);
                                                                                                                                           int fr, to;
      if (w[now] <= n/2)
                                                                             res2=dfs(g2, m1, id1, c2.first, 0);
                                                                                                                                           edg(int _fr, int _to){
          if (rec.first==0) rec.first=now;
                                                                             res3=dfs(g2, m2, id2, c2.second, 0);
                                                                                                                                               fr = fr;
          else rec.second=now;
                                                                         }else if (c1.second!=0){
                                                                                                                                               to = _to;
40
                                                                             res1=dfs(g2, m1, id1, c2.first, 0);
                                                                                                                                     33
                                                                             m2=m1:
                                                                                                                                       };
                                                                             id2=id1;
                                                                                                                                       ostream& operator<<(ostream& os, edg x){os << x.fr << "--" <<
  int dfs(Graph &g, Hash &m, int &id, int now, int pre){
                                                                             res2=dfs(g1, m1, id1, c1.first, 0);
                                                                                                                                       vector<edg> EV:
      vector<int> v;
                                                                             res3=dfs(g1, m2, id2, c1.second, 0);
                                                                 110
      for (auto x : g[now]){
                                                                 111
                                                                                                                                       void tarjan(int v, int par, stack<int>& S){
          if (x!=pre){
                                                                             res1=dfs(g1, m1, id1, c1.first, 0);
                                                                 112
                                                                                                                                           static vector<int> dfn(mxn), low(mxn);
              int add=dfs(g, m, id, x, now);
                                                                             res2=dfs(g2, m1, id1, c2.first, 0);
                                                                 113
              v.push back(add);
                                                                 114
                                                                                                                                           static vector<bool> to_add(mxn);
                                                                                                                                           static int nowT = 0;
                                                                 115
                                                                 116
                                                                         cout << (res1==res2 | res1==res3 ? "YES" : "NO") << endl 43
                                                                                                                                           int childs = 0;
      sort(v.begin(), v.end());
                                                                                                                                           nowT += 1;
      if (m.find(v)!=m.end()){
                                                                                                                                           dfn[v] = low[v] = nowT;
                                                                                                                                           for(auto &ne:E[v]){
          return m[v];
                                                                 119
                                                                         return:
      }else{
                                                                                                                                               int i = EV[ne].to;
                                                                 120
          m[v]=++id;
                                                                                                                                               if(i == par) continue;
                                                                                                                                               if(!dfn[i]){
          return id;
                                                                     signed main(void){
                                                                         fastio:
                                                                                                                                                   S.push(ne);
                                                                                                                                                   tarjan(i, v, S);
                                                                         int t=1:
                                                                                                                                                   childs += 1;
                                                                         cin >> t;
                                                                                                                                                   low[v] = min(low[v], low[i]);
  void solve1(){
                                                                         while (t--){
                                                                                                                                                   if(par >= 0 && low[i] >= dfn[v]){
                                                                 128
                                                                             solve1();
                                                                                                                                                       vector<int> bcc;
      // init
                                                                 129
      id1=0:
                                                                 130
                                                                         return 0;
                                                                                                                                                       int tmp;
      id2=0:
                                                                 131
                                                                                                                                                       do{
      c1={0, 0};
                                                                                                                                                            tmp = S.top(); S.pop();
                                                                                                                                                            if(!to add[EV[tmp].fr]){
      c2={0, 0};
      fill(sz1.begin(), sz1.begin()+n+1, 0);
                                                                                                                                                                to add[EV[tmp].fr] = true;
                                                                    6.17 圓方樹 [675aec]
      fill(sz2.begin(), sz2.begin()+n+1, 0);
                                                                                                                                                                bcc.pb(EV[tmp].fr);
      fill(we1.begin(), we1.begin()+n+1, 0);
      fill(we2.begin(), we2.begin()+n+1, 0);
                                                                                                                                                            if(!to_add[EV[tmp].to]){
      for (int i=1 ; i<=n ; i++){</pre>
                                                                   | #include <bits/stdc++.h>
                                                                                                                                                                to add[EV[tmp].to] = true;
          g1[i].clear();
                                                                                                                                                                bcc.pb(EV[tmp].to);
                                                                    #define lp(i,a,b) for(int i=(a);i<(b);i++)
          g2[i].clear();
                                                                    #define pii pair<int.int>
                                                                                                                                                       }while(tmp != ne);
                                                                    #define pb push back
                                                                                                                                                       for(auto &j:bcc){
      m1.clear();
                                                                    #define ins insert
                                                                                                                                                            to_add[j] = false;
      m2.clear();
                                                                    #define ff first
                                                                                                                                                            F[last special node].pb(j);
                                                                    #define ss second
      // input
                                                                    #define opa(x) cerr << #x << " = " << x << ", ";
                                                                                                                                                            F[i].pb(last special node);
                                                                    #define op(x) cerr << #x << " = " << x << endl;
      for (int i=0 ; i<n-1 ; i++){</pre>
                                                                  10 #define ops(x) cerr << x;</pre>
                                                                                                                                                       last special node += 1:
          cin >> a >> b;
                                                                  #define etr cerr << endl;</pre>
          g1[a].push back(b);
                                                                  12 #define spc cerr << ' ';
          g1[b].push back(a);
                                                                  #define BAE(x) (x).begin(), (x).end()
                                                                  14 #define STL(x) cerr << #x << " : "; for(auto &qwe:x) cerr <<</pre>
                                                                                                                                                    low[v] = min(low[v], dfn[i]);
      for (int i=0; i<n-1; i++){
                                                                         qwe << ''; cerr << endl;</pre>
                                                                                                                                                    if(dfn[i] < dfn[v]){ // edge i--v will be visited</pre>
          cin >> a >> b;
                                                                  #define deb1 cerr << "deb1" << endl;</pre>
                                                                                                                                                          twice at here, but we only need one.
```

```
}
83
s6 int dep[mxn], jmp[mxn][mxlg];
   void dfs lca(int v, int par, int depth){
       dep[v] = depth;
       for(auto &i:F[v]){
           if(i == par) continue;
           jmp[i][0] = v;
           dfs_lca(i, v, depth + 1);
93
   inline void build_lca(){
       jmp[1][0] = 1;
       dfs_lca(1, -1, 1);
       lp(j,1,mxlg){
           lp(i,1,mxn){
               jmp[i][j] = jmp[jmp[i][j-1]][j-1];
102
103
104
105
   inline int lca(int x, int y){
       if(dep[x] < dep[y]){ swap(x, y); }</pre>
       int diff = dep[x] - dep[y];
109
       lp(j,0,mxlg){
110
           if((diff >> j) & 1){
               x = jmp[x][j];
112
113
114
115
       if(x == y) return x;
116
       for(int j = mxlg - 1; j >= 0; j--){
117
           if(jmp[x][j] != jmp[y][j]){
118
               x = jmp[x][j];
120
               y = jmp[y][j];
121
122
123
       return jmp[x][0];
124
125
   inline bool can reach(int fr, int to){
       if(dep[to] > dep[fr]) return false;
127
128
129
       int diff = dep[fr] - dep[to];
       lp(i,0,mxlg){
130
131
           if((diff >> j) & 1){
                fr = jmp[fr][j];
132
133
134
135
       return fr == to;
136
137
       ios::sync with stdio(false); cin.tie(0);
        freopen("test_input.txt", "r", stdin);
       int n, m, q; cin >> n >> m >> q;
142
       lp(i,0,m){}
143
           int u, v; cin >> u >> v;
           E[u].pb(EV.size());
144
           EV.pb(edg(u, v));
```

S.push(ne);

}

```
E[v].pb(EV.size());
147
            EV.pb(edg(v, u));
148
        E[0].pb(EV.size());
149
150
        EV.pb(edg(0, 1));
        stack<int> S;
151
152
        tarjan(0, -1, S);
        build lca();
154
155
        lp(queries,0,q){
156
            int fr, to, relay; cin >> fr >> to >> relay;
157
            if(fr == relay || to == relay){
                cout << "NO\n";
158
                continue;
159
160
            if((can_reach(fr, relay) || can_reach(to, relay)) &&
                 dep[relay] >= dep[lca(fr, to)]){
162
                cout << "NO\n";
163
                continue:
164
            cout << "YES\n";</pre>
165
166
167 }
```

6.18 最大權閉合圖 [6ca663]

```
邊 u → v 表示選 u 就要選 v (0-based)
      保證回傳值非負
      構造:從 S 開始 dfs,不走最小割的邊
           所有經過的點就是要選的那些點。
      一般圖: O(n²m) / 二分圖: O(m√n)
  template<typename U>
  U maximum_closure(vector<U> w, vector<pair<int,int>> EV) {
      int n = w.size(), S = n + 1, T = n + 2;
      Flow G(T + 5); // Graph/Dinic.cpp
     U sum = 0;
      for (int i = 0; i < n; ++i) {</pre>
         if (w[i] > 0) {
             G.add(S, i, w[i]);
             sum += w[i];
          else if (w[i] < 0) {</pre>
             G.add(i, T, abs(w[i]));
      for (auto &[u, v] : EV) { // 請務必確保 INF > Σ/w i/
23
         G.add(u, v, INF);
      U cut = G.flow(S, T);
      return sum - cut:
```

6.19 Theorem

- 任意圖
 - 最大匹配 + 最小邊覆蓋 = n (不能有孤點)

- 點覆蓋的補集是獨立集。最小點覆蓋 + 最大獨立集 = n
- $-w(最小權點覆蓋)+w(最大權獨立集)=\sum w_v$
- (帶點權的二分圖可以用最小割解,構造請參考 Augment Path.cpp)
- 二分圖
 - 最小點覆蓋 = 最大匹配 = n 最大獨立集
- 只有邊帶權的二分圖
 - w-vertex-cover (帶權點覆蓋): 每條邊的兩個連接點被選中的次數總和至少要是 w_e 。
 - w-weight matching (帶權匹配)
 - minimum vertex count of w-vertex-cover = maximum weight count of w-weight matching (一個點可以被選很多次・但邊不行)
- 點、邊都帶權的二分圖的定理
 - b-matching:假設 v 的點權是 b_v · 那所有 v 的匹配邊 e 的權重都要滿足 $\sum w_e < b_v$ 。
 - The maximum w-weight of a b-matching equals the minimum b-weight of vertices in a w-vertex-cover.

7 Math

1 // ax + by = c

7.1 CRT [2b18f3]

```
int extgcd abc(int a, int b, int c, int &x, int &y) {
    if (b == 0) {
        if (c % a) return INF;
        x = c / a, y = 0;
        return abs(a);
    int x1, y1;
    int g = extgcd_abc(b, a % b, c, x1, y1);
    x = y1;
    y = x1 - a / b * y1;
    return g;
// Sorry that I don't know how to merge two functions
int extgcd(int a, int b, int &x, int &y) {
    if (b == 0) {
        x = 1, y = 0;
        return a;
    int ret = extgcd(b, a % b, y, x);
    y -= a / b * x;
    return ret;
// 有 n 個式子, 求解 x ≡ a i (mod m i)
int CRT_m_coprime(int n, vector<int> &a, vector<int> &m) {
    int p = 1, ans = 0;
    vector<int> M(n), inv_M(n);
    for (int i = 0; i < n; i++) p *= m[i];</pre>
    for (int i = 0; i < n; i++) {</pre>
        M[i] = p / m[i];
```

```
int tmp;
          extgcd(M[i], m[i], inv_M[i], tmp);
          ans += a[i] * inv_M[i] * M[i];
          ans %= p;
      return (ans % p + p) % p;
41 // 對於方程組的式子兩兩求解
42 // 回傳: {是否有解, {a, m}}
43 pair<bool, pair<int, int>> CRT_m_NOT_coprime(int a1, int m1,
      int a2, int m2) {
      int g = __gcd(m1, m2);
     if ((a2 - a1) % g != 0) return {0, {-1, -1}};
     int x, y; extgcd(m1, m2, x, y);
     x = (a2 - a1) * x / g; // 兩者不能相反
      a1 = x * m1 + a1:
      m1 = m1 * m2 / g;
      a1 = (a1 \% m1 + m1) \% m1;
      return {1, {a1, m1}};
56 // ans = a / b (mod m)
57 // ans = ret.F + k * ret.S, k is integer
58 pair<int, int> div(int a, int b, int m) {
     int flag = 1;
     if (a < 0) { a = -a; flag *= -1; }</pre>
     if (b < 0) { b = -b; flag *= -1; }</pre>
     int t = -1, k = -1;
     int res = extgcd_abc(b, m, a, t, k);
     if (res == INF) return {INF, INF};
     m = abs(m / res);
     t = t * flag;
     t = (t \% m + m) \% m;
```

7.2 Josephus Problem [e0ed50]

return {t, m};

7.3 Lagrange any x [1f2c26]

```
| // init: (x1, y1), (x2, y2) in a vector
| struct Lagrange{
| int n;
```

```
vector<pair<int, int>> v;
       Lagrange(vector<pair<int, int>> & v){
          n = _v.size();
          v = v;
      // O(n^2 \log MAX A)
      int solve(int x){
           int ret = 0;
           for (int i=0 ; i<n ; i++){</pre>
               int now = v[i].second;
               for (int j=0 ; j<n ; j++){</pre>
                   if (i==j) continue;
                   now *= ((x-v[j].first)+MOD)%MOD;
                   now %= MOD;
                   now *= (qp((v[i].first-v[j].first+MOD)%MOD,
                        MOD-2)+MOD)%MOD;
                   now %= MOD;
               ret = (ret+now)%MOD;
           return ret;
28 };
```

7.4 Lagrange continuous x [57536a]

```
#include <bits/stdc++.h>
  using namespace std;
  const int MAX N = 5e5 + 10;
  const int mod = 1e9 + 7;
  long long inv_fac[MAX_N];
  inline int fp(long long x, int y) {
      int ret = 1;
      for (; y; y >>= 1) {
          ret = (y & 1) ? (ret * x % mod) : ret;
          x = x * x % mod;
15
      return ret;
  // TO USE THIS TEMPLATE, YOU MUST MAKE SURE THAT THE MOD
       NUMBER IS A PRIME.
  struct Lagrange {
      Initialize a polynomial with f(x_0), f(x_0 + 1), ..., f(
      This determines a polynomial f(x) whose degree is at most 83 }
      Then you can call sample(x) and you get the value of f(x)
      Complexity of init() and sample() are both O(n).
25 */
      int m, shift; // m = n + 1
      vector<int> v, mul;
28 // You can use this function if you don't have inv fac array
```

void construct_inv_fac() {

```
fac = fac * i % mod;
        inv_fac[MAX_N - 1] = fp(fac, mod - 2);
        for (int i = MAX N - 1; i >= 1; --i) {
            inv fac[i - 1] = inv fac[i] * i % mod;
// You call init() many times without having a second
     instance of this struct.
    void init(int X_0, vector<int> &u) {
        shift = ((1 - X 0) \% mod + mod) \% mod;
        if (v.size() == 1) v.push_back(v[0]);
        m = v.size();
        mul.resize(m);
// You can use sample(x) instead of sample(x \% mod).
    int sample(int x) {
        x = ((long long)x + shift) % mod;
        x = (x < 0) ? (x + mod) : x;
        long long now = 1;
        for (int i = m; i >= 1; --i) {
            mul[i - 1] = now;
            now = now * (x - i) % mod;
        int ret = 0;
        bool neg = (m - 1) & 1;
        now = 1;
        for (int i = 1; i <= m; ++i) {</pre>
            int up = now * mul[i - 1] % mod;
            int down = inv fac[m - i] * inv fac[i - 1] % mod;
            int tmp = ((long long)v[i - 1] * up % mod) * down
                  % mod;
            ret += (neg && tmp) ? (mod - tmp) : (tmp);
            ret = (ret >= mod) ? (ret - mod) : ret;
            now = now * (x - i) % mod;
            neg ^= 1;
        return ret;
int main() {
    int n; cin >> n;
    vector<int> v(n);
    for (int i = 0; i < n; ++i) {
        cin >> v[i];
    Lagrange L;
    L.construct inv fac();
    L.init(0, v);
    int x; cin >> x;
    cout << L.sample(x);</pre>
7.5 Linear Mod Inverse [ecf71e]
```

long long fac = 1;

for (int i = 2; i < MAX N; ++i) {</pre>

```
1 // 線性求 1-based a[i] 對 p 的乘法反元素
2 vector<int> s(n+1, 1), invS(n+1), invA(n+1);
3 for (int i=1; i<=n; i++) s[i] = s[i-1]*a[i]%p;
4 invS[n] = qp(s[n], p-2, p);</pre>
```

```
5| for (int i=n; i>=1; i--) invS[i-1] = invS[i]*a[i]%p;
6 for (int i=1; i<=n; i++) invA[i] = invS[i]*s[i-1]%p;</pre>
 7.6 Lucas's Theorem [b37dcf]
1 \mid //  對於很大的 C^n_{m} 對質數 p 取模·只要 p 不大就可以用。
int Lucas(int n, int m, int p){
     if (m==0) return 1;
     return (C(n%p, m%p, p)*Lucas(n/p, m/p, p)%p);
       Matrix [8d1a23]
| struct Matrix{
     int n, m;
     vector<vector<int>> arr;
     Matrix(int n, int m){
         n = _n;
         m = _m;
         arr.assign(n, vector<int>(m));
     vector<int> & operator [] (int i){
         return arr[i];
     Matrix operator * (Matrix b){
         Matrix ret(n, b.m);
         for (int i=0 ; i<n ; i++){</pre>
             for (int j=0 ; j<b.m ; j++){</pre>
                 for (int k=0 ; k<m ; k++){</pre>
                     ret.arr[i][j] += arr[i][k]*b.arr[k][j]%
                     ret.arr[i][j] %= MOD;
         return ret;
     Matrix pow(int p){
         Matrix ret(n, n), mul = *this;
         for (int i=0 ; i<n ; i++){</pre>
             ret.arr[i][i] = 1;
         for ( ; p ; p>>=1){
             if (p&1) ret = ret*mul;
             mul = mul*mul;
         return ret;
```

vector<vector<int>> arr = this->arr;

int det(){

bool flag = false;

for (int i=0 ; i<n ; i++){</pre>

```
int target = -1;
    for (int j=i ; j<n ; j++){</pre>
        if (arr[j][i]){
            target = j;
            break:
    if (target==-1) return 0;
    if (i!=target){
        swap(arr[i], arr[target]);
        flag = !flag;
    for (int j=i+1 ; j<n ; j++){</pre>
        if (!arr[j][i]) continue;
        int freq = arr[j][i]*qp(arr[i][i], MOD-2)%MOD
        for (int k=i ; k<n ; k++){</pre>
            arr[j][k] -= freq*arr[i][k];
            arr[j][k] = (arr[j][k]%MOD+MOD)%MOD;
   }
int ret = !flag ? 1 : MOD-1;
for (int i=0 ; i<n ; i++){</pre>
    ret *= arr[i][i];
    ret %= MOD;
return ret;
```

7.8 Matrix 01 [8d542a]

```
const int MAX_N = (1LL<<12);</pre>
struct Matrix{
    vector<bitset<MAX_N>> arr;
    Matrix(int _n, int _m){
         n = _n;
        m = _m;
         arr.resize(n);
    Matrix operator * (Matrix b){
         Matrix b_t(b.m, b.n);
         for (int i=0 ; i<b.n ; i++){</pre>
             for (int j=0 ; j<b.m ; j++){</pre>
                 b_t.arr[j][i] = b.arr[i][j];
         Matrix ret(n, b.m);
         for (int i=0 ; i<n ; i++){</pre>
             for (int j=0 ; j<b.m ; j++){</pre>
                 ret.arr[i][j] = ((arr[i]&b_t.arr[j]).count()
         return ret;
```

7.9 Miller Rabin [24bd0d]

28 };

```
1 // O(k \log^3 n), k = llsprp.size()
  typedef Uint unsigned long long;
  Uint modmul(Uint a, Uint b, Uint m) {
      int ret = a*b - m*(Uint)((long double)a*b/m);
      return ret + m*(ret < 0) - m*(ret>=(int)m);
  int qp(int b, int p, int m){
      int ret = 1;
      for ( ; p ; p>>=1){
          if (p&1) ret = modmul(ret, b, m);
          b = modmul(b, b, m);
14
      return ret;
15
16
  vector<int> llsprp = {2, 325, 9375, 28178, 450775, 9780504,
       1795265022};
  bool is_prime(int n, vector<int> sprp = llsprp){
      if (n==2) return 1:
      if (n<2 || n%2==0) return 0;
      int t = 0;
      int u = n-1;
      for (; u%2==0; t++) u>>=1;
      for (int i=0 ; i<sprp.size() ; i++){</pre>
          int a = sprp[i]%n;
          if (a==0 || a==1 || a==n-1) continue;
          int x = qp(a, u, n);
          if (x==1 || x==n-1) continue;
          for (int j=0 ; j<t ; j++){</pre>
              x = modmul(x, x, n);
              if (x==1) return 0;
               if (x==n-1) break;
          if (x==n-1) continue;
38
          return false;
39
40
42
      return true;
43
```

7.10 Pollard Rho [a5daef]

7.11 Polynomial [51ca3b]

```
| struct Poly {
     int len, deg;
     // len = 2^k >= the original length
     Poly(): len(0), deg(0), a(nullptr) {}
     Poly(int n) {
         len = 1;
         deg = _n - 1;
         while (len < n) len <<= 1;
         a = (11*) calloc(len, sizeof(11));
     Poly(int 1, int d, int *b) {
         len = 1;
         deg = d;
         a = b;
     void resize(int _n) {
         int len1 = 1:
         while (len1 < _n) len1 <<= 1;</pre>
         int *res = (ll*) calloc(len1, sizeof(ll));
         for (int i = 0; i < min(len, _n); i++) {</pre>
             res[i] = a[i];
         len = len1:
         deg = _n - 1;
         free(a);
         a = res;
     Poly& operator=(const Poly rhs) {
         this->len = rhs.len;
         this->deg = rhs.deg;
         this->a = (ll*)realloc(this->a, sizeof(ll) * len);
         copy(rhs.a, rhs.a + len, this->a);
         return *this:
     Poly operator*(Poly rhs) {
```

```
int 11 = this->len, 12 = rhs.len;
    int d1 = this->deg, d2 = rhs.deg;
                                                              104
    while (11 > 0 \text{ and this} -> a[11 - 1] == 0) 11--:
                                                              105
    while (12 > 0 \text{ and } rhs.a[12 - 1] == 0) 12--;
                                                              106
                                                              107
    while (1 < max(11 + 12 - 1, d1 + d2 + 1)) 1 <<= 1;
                                                              108
    int *x, *v, *res;
                                                              109
    x = (11*) calloc(1, sizeof(11));
                                                              110
   y = (11*) calloc(1, sizeof(11));
                                                              111
    res = (ll*) calloc(l, sizeof(ll));
                                                              112
    copy(this->a, this->a + l1, x);
                                                              113
    copy(rhs.a, rhs.a + 12, y);
                                                              114
    ntt.tran(1, x); ntt.tran(1, y);
                                                              115
    FOR (i, 0, 1 - 1)
                                                              116
        res[i] = x[i] * y[i] % mod;
                                                              117
    ntt.tran(1, res, true);
                                                              118
    free(x); free(y);
                                                              119
    return Poly(1, d1 + d2, res);
                                                              120
                                                              121
Poly operator+(Poly rhs) {
                                                              122
    int 11 = this->len, 12 = rhs.len;
                                                              123
    int 1 = max(11, 12);
                                                              124
    Poly res;
                                                              125
    res.len = 1:
                                                              126
    res.deg = max(this->deg, rhs.deg);
                                                              127
    res.a = (ll*) calloc(l, sizeof(ll));
                                                              128
    FOR (i, 0, 11 - 1) {
                                                              129
        res.a[i] += this->a[i];
                                                              130
        if (res.a[i] >= mod) res.a[i] -= mod;
                                                              131
                                                              132
    FOR (i, 0, 12 - 1) {
                                                              133
        res.a[i] += rhs.a[i];
                                                              134
        if (res.a[i] >= mod) res.a[i] -= mod;
                                                              135
                                                              136
    return res;
                                                              137
                                                              138
Poly operator-(Poly rhs) {
                                                              139
    int l1 = this->len, l2 = rhs.len;
                                                              140
    int 1 = \max(11, 12);
                                                              141
    Poly res;
                                                              142
    res.len = 1;
                                                              143
    res.deg = max(this->deg, rhs.deg);
                                                              144
    res.a = (11*) calloc(1, sizeof(11));
                                                              145
    FOR (i, 0, 11 - 1) {
        res.a[i] += this->a[i];
                                                              147
        if (res.a[i] >= mod) res.a[i] -= mod;
                                                              148
                                                              149
    FOR (i, 0, 12 - 1) {
                                                              150
        res.a[i] -= rhs.a[i];
                                                              151
        if (res.a[i] < 0) res.a[i] += mod;</pre>
                                                              152
                                                              153
    return res;
                                                              154
                                                              155
Poly operator*(const int rhs) {
                                                              156
    Poly res;
                                                              157
    res = *this;
    FOR (i, 0, res.len - 1) {
                                                              159
        res.a[i] = res.a[i] * rhs % mod;
        if (res.a[i] < 0) res.a[i] += mod;</pre>
                                                              160
                                                              161
    return res;
                                                              162
                                                              163
Poly(vector<int> f) {
                                                              164
    int n = f.size();
                                                              165
    len = 1:
                                                              166
    deg = _n - 1;
```

```
while (len < n) len <<= 1;
    a = (11*) calloc(len, sizeof(11));
    FOR (i, 0, deg) a[i] = f[i];
Poly derivative() {
    Poly g(this->deg);
    FOR (i, 1, this->deg) {
        g.a[i - 1] = this->a[i] * i % mod;
    return g;
Poly integral() {
    Poly g(this->deg + 2);
    FOR (i, 0, this->deg) {
        g.a[i + 1] = this -> a[i] * ::inv(i + 1) % mod;
    return g;
Poly inv(int len1 = -1) {
    if (len1 == -1) len1 = this->len;
    Poly g(1); g.a[0] = ::inv(a[0]);
    for (int 1 = 1; 1 < len1; 1 <<= 1) {
        Poly t; t = *this;
        t.resize(1 << 1);
        t = g * g * t;
        t.resize(1 << 1);
        Poly g1 = g * 2 - t;
        swap(g, g1);
    return g;
Poly ln(int len1 = -1) {
    if (len1 == -1) len1 = this->len;
    auto g = *this;
    auto x = g.derivative() * g.inv(len1);
    x.resize(len1);
    x = x.integral();
    x.resize(len1);
    return x;
Poly exp() {
    Poly g(1);
    g.a[0] = 1;
    for (int 1 = 1; 1 < len; 1 <<= 1) {</pre>
        Poly t, g1; t = *this;
        t.resize(1 << 1); t.a[0]++;
        g1 = (t - g.ln(1 << 1)) * g;
        g1.resize(1 << 1);
        swap(g, g1);
    return g;
Poly pow(ll n) {
    Poly &a = *this;
    while (i <= a.deg and a.a[i] == 0) i++;</pre>
    if (i and (n > a.deg or n * i > a.deg)) return Poly(a
         .deg + 1);
    if (i == a.deg + 1) {
        Polv res(a.deg + 1):
        res.a[0] = 1;
        return res;
    Poly b(a.deg - i + 1);
    int inv1 = ::inv(a.a[i]);
    FOR (j, 0, b.deg)
```

7.12 josephus [0be067]

7.13 數論分塊 [8ccab5]

7.14 最大質因數 [ca5e52]

```
void max_fac(int n, int &ret){
    if (n<=ret || n<2) return;
    if (isprime(n)){
        ret = max(ret, n);
        return;
    }
    int p = Pollard_Rho(n);
    max_fac(p, ret), max_fac(n/p, ret);
}</pre>
```

7.15 歐拉公式 [85f3b1]

```
1 \mid // phi(n) = 小於 n 並與 n 互質的正整數數量。
2 // O(sqrt(n)), 回傳 phi(n)
  int phi(int n){
      int ret = n;
      for (int i=2; i*i<=n; i++){</pre>
          if (n%i==0){
              while (n%i==0) n /= i;
              ret = ret*(i-1)/i;
      if (n>1) ret = ret*(n-1)/n;
      return ret;
  // O(n log n) · 回傳 1~n 的 phi 值
  vector<int> phi_1_to_n(int n){
      vector<int> phi(n+1);
      phi[0]=0;
      phi[1]=1;
      for (int i=2 ; i<=n ; i++){</pre>
          phi[i]=i-1;
25
      for (int i=2 ; i<=n ; i++){</pre>
          for (int j=2*i; j<=n; j+=i){ // 枚舉所有倍數
              phi[j]-=phi[i];
      }
      return phi;
```

7.16 Burnside's Lemma

$$\sum_{k=1}^{n} \frac{c(k)}{n}$$

- n:有多少種置換方式(例如:旋轉方式)
- c(k): 所有可能中,經過 k 次旋轉後,仍不會和別人相同的方式的數量

7.17 Catalan Number

任意括號序列: $C_n = \frac{1}{n+1} \binom{2n}{n}$

7.18 Matrix Tree Theorem

目標:給定一張無向圖·問他的生成樹數量。 方法:先把所有自環刪掉·定義 Q 為以下矩陣

接著刪掉 Q 的第一個 row 跟 column · 它的 determinant 就是答案。目標:給定一張有向圖 · 問他的以 r 為根 · 可以走到所有點生成樹數量。

方法:先把所有自環刪掉,定義 Q 為以下矩陣

$$Q_{i,j} = \begin{cases} \deg_{in}(v_i) & \text{if } i = j \\ -(\frac{1}{8}v_i v_j \text{ bbm}) & \text{otherwise} \end{cases}$$

接著刪掉 Q 的第 r 個 row 跟 column \cdot 它的 determinant 就是答案。

7.19 Stirling's formula

 $n! \approx \sqrt{2\pi n} (\frac{n}{\epsilon})^n$

7.20 Theorem

- 1. $1 \sim x$ 質數的數量 $\approx \frac{x}{\ln x}$
- 2. x 的因數的數量 $\approx x^{\frac{1}{3}}$
- 3. x 的質因數的數量 $\approx \log \log x$
- 4. p is a prime number $\Leftrightarrow (p-1)! \equiv -1 \pmod{p}$
- 5. 每個正整數都可以表示成四個整數的平方和
- 6. 任何大於 2 的整數都可以表示成兩個質數的和
- 7. $n^{k-2} \cdot \prod_{i=1}^{k} s_i n$ 個點、k 的連通塊・加上 k-1 條邊使得變成一個連通圖的方法數、其中每個連通塊有 s_i 個點

7.21 二元一次方程式

 $\begin{cases} ax + by = e \\ cx + dy = f \end{cases} = \begin{cases} x = \frac{ed - bf}{ad - bc} \\ y = \frac{af - ec}{ad - bc} \end{cases}$

若 $x=\frac{0}{0}$ 且 $y=\frac{0}{0}$ · 則代表無限多組解。若 $x=\frac{*}{0}$ 且 $y=\frac{*}{0}$ · 則代表無解。

7.22 歐拉定理

若 a, m 互質 · 則:

$$a^n \equiv a^{n \bmod \varphi(m)} \pmod{m}$$

若 a, m 不互質 \cdot 則:

$$a^n \equiv a^{\varphi(m) + [n \mod \varphi(m)]} \pmod{m}$$

7.23 錯排公式

錯排公式: (n) 個人中,每個人皆不再原來位置的組合數)

$$dp_i = \begin{cases} 1 & i = 0\\ 0 & i = 1\\ (i-1)(dp_{i-1} + dp_{i-2}) & \text{otherwise} \end{cases}$$

8 String

8.1 AC automation [6ece7f]

```
1 \mid const int MAXN = 5e5 + 5;
  struct ac automation {
      int go[MAXN][26], fail[MAXN], is end[MAXN];
      int sz;
      void add(string s) {
          int now = 0;
          for (char c : s) {
              if (!go[now][c - 'a'])
                  go[now][c - 'a'] = ++sz;
              now = go[now][c - 'a'];
          is end[now]++;
      vector<int> que;
      void build() {
          que.pb(0);
          for (int i = 0; i < ssize(que); i++) {</pre>
              auto u = que[i];
              FOR (c, 0, 25) {
                   if (go[u][c]) {
                       int v = go[u][c];
                       fail[v] = !u ? 0 : go[fail[u]][c];
                       is end[v] += is end[fail[v]];
                       que.pb(v);
                   else {
                       go[u][c] = go[fail[u]][c];
35 } AC;
```

8.3 KMP [e5b7ce]

8.2 Hash [942f42]

8.4 Manacher [9a4b4d]

```
string Manacher(string str) {
    string tmp = "$#";
    for(char i : str) {
        tmp += i;
        tmp += '#';
    }

vector<int> p(tmp.size(), 0);
    int mx = 0, id = 0, len = 0, center = 0;
    for(int i=1 ; i<(int)tmp.size() ; i++) {
        p[i] = mx > i ? min(p[id*2-i], mx-i) : 1;

while(tmp[i+p[i]] == tmp[i-p[i]]) p[i]++;
    if(mx<i+p[i]) mx = i+p[i], id = i;
    if(len<p[i]) len = p[i], center = i;
    }

return str.substr((center-len)/2, len-1);
}</pre>
```

8.5 Min Rotation [9d296f]

8.6 Suffix Array [6352b3]

```
ı|// 注意,當 /s/=1 時,Lcp 不會有值,務必測試 /s/=1 的 case
  struct SuffixArray {
      string s;
      vector<int> sa, lcp;
      SuffixArray(string s, int lim = 256) {
          int n = s.size()+1, k = 0, a, b;
          vector<int> x(s.begin(), s.end()), y(n), ws(max(n,
               lim)), rank(n);
          x.push back(0);
          sa = 1cp = v;
          iota(sa.begin(), sa.end(), 0);
          for (int j=0, p=0; p<n; j=max(1LL, j*2), lim=p) {</pre>
              iota(y.begin(), y.end(), n-j);
              for (int i=0; i<n; i++) if (sa[i] >= j) y[p++]
                   = sa[i] - j;
              fill(ws.begin(), ws.end(), 0);
              for (int i=0; i<n; i++) ws[x[i]]++;</pre>
              for (int i=1; i<lim; i++) ws[i] += ws[i - 1];</pre>
              for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
              swap(x, y), p = 1, x[sa[0]] = 0;
              for (int i=1 ; i<n ; i++){</pre>
                  a = sa[i - 1];
                  b = sa[i];
                  x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
                        ? p - 1 : p++;
          for (int i=1; i<n; i++) rank[sa[i]] = i;</pre>
          for (int i=0, j ; i<n-1 ; lcp[rank[i++]]=k)</pre>
31
              for (k && k--, j=sa[rank[i]-1]; i+k<s.size() &&</pre>
                   j+k<s.size() && s[i+k]==s[j+k] ; k++);
          sa.erase(sa.begin());
          lcp.erase(lcp.begin(), lcp.begin()+2);
      // f49583
      vector<int> pos; // pos[i] = i 這個值在 pos 的哪個地方
      SparseTable st:
      void init lcp(){
          pos.resize(sa.size());
```

if (nowRank+add>=k){

prePrefix = sa.lcp[i];

nowRank += add;

return s.substr(sa[i], prePrefix+k-nowRank);

```
for (int i=0 ; i<sa.size() ; i++){</pre>
                                                                          }else{
        pos[sa[i]] = i;
                                                        100
                                                                               if (i<lcp.size()) now = lcp[i];</pre>
                                                                                                                        12
                                                                                                                        13
    if (lcp.size()){
                                                        102
                                                                       }
                                                                                                                        14
       st.build(lcp);
                                                                                                                        15
                                                        103
                                                                   { // build suf
                                                        104
                                                                                                                        16
}
                                                        105
                                                                       int now = 0;
                                                                                                                        17
                                                                       for (int i=s.size()-1; i>=0; i--){
                                                                                                                        18
                                                                           if (sa[i]<=p){</pre>
// 用之前記得 init
                                                        107
                                                                                                                        19
                                                        108
                                                                               suf[sa[i]] = now;
// 回傳 [l1, r1] 跟 [l2, r2] 的 Lcp·0-based
                                                        109
                                                                               if (i-1>=0) now = min(now, lcp[i-1]);
int get_lcp(int l1, int r1, int l2, int r2){
                                                        110
    int pos_1 = pos[l1], len_1 = r1-l1+1;
                                                                               if (i-1>=0) now = lcp[i-1];
                                                        111
    int pos_2 = pos[12], len_2 = r2-12+1;
                                                        112
    if (pos 1>pos 2){
                                                        113
        swap(pos_1, pos_2);
                                                        114
        swap(len_1, len_2);
                                                        115
                                                        116
                                                                   return {pre, suf};
                                                        117
    if (11==12){
                                                        118 };
        return min(len_1, len_2);
        return min({st.query(pos_1, pos_2), len_1, len_2
                                                           8.7 Z Algorithm [bcfbd6]
}
// 檢查 [l1, r1] 跟 [l2, r2] 的大小關係·0-based
                                                         1 \mid // 定義一個長度為 n 的文本為 T · 則陣列 Z 的 Z[i] 代表 T[0:n]
// 如果前者小於後者,就回傳 <0,相等就回傳 =0,否則回傳
                                                                 和 T[i:n] 最長共同前綴
                                                           // bcfbd6
// 5b8db0
                                                           vector<int> z_function(string s){
int substring cmp(int l1, int r1, int l2, int r2){
                                                               vector<int> ret(s.size());
    int len 1 = r1-l1+1;
                                                               int 11 = 0, rr = 0;
    int len_2 = r2-l2+1;
    int res = get lcp(l1, r1, l2, r2);
                                                               for (int i=1; i<s.size(); i++){</pre>
                                                                   int j = 0;
    if (res<len_1 && res<len_2){</pre>
       return s[11+res]-s[12+res];
                                                                   if (i<rr) j = min(ret[i-ll], rr-i);</pre>
    }else if (len_1==res && len_2==res){
                                                                   while (s[j]==s[i+j]) j++;
       // 如果不需要以 index 作為次要排序參數, 這裡要回
                                                                   ret[i] = j;
                                                                   if (i+j>rr){
        return 11-12;
                                                                       11 = i:
    }else{
                                                                       rr = i+j;
        return len 1==res ? -1 : 1;
}
                                                               ret[0] = s.size();
// 對於位置在 <=p 的後綴·找離他左邊/右邊最接近位置 >p 的
                                                               return ret:
     後綴的 Lcp · 0-based
// pre[i] = s[i] 離他左邊最接近位置 >p 的後綴的 Lcp · 0-
// suf[i] = s[i] 離他右邊最接近位置 >p 的後綴的 Lcp · 0-
     based
                                                           8.8 k-th Substring1 [61f66b]
// da12fa
pair<vector<int>, vector<int>> get_left_and_right_lcp(int
                                                          1// 回傳 s 所有子字串(完全不同)中,第 k 大的
    vector<int> pre(p+1);
                                                           string k_th_substring(string &s, int k){
    vector<int> suf(p+1);
                                                               int n = s.size();
                                                               SuffixArray sa(s);
    { // build pre
                                                               sa.init_lcp();
       int now = 0:
       for (int i=0 ; i<s.size() ; i++){</pre>
                                                               int prePrefix = 0, nowRank = 0;
           if (sa[i]<=p){
                                                               for (int i=0 ; i<n ; i++){</pre>
               pre[sa[i]] = now;
                                                                   int len = n-sa[i];
               if (i<lcp.size()) now = min(now, lcp[i]);</pre>
                                                                   int add = len-prePrefix;
```