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1 Misc

1.1 2-SAT

```
| #include <bits/stdc++.h>
using namespace std;
 struct TWO SAT {
     int n, N;
     vector<vector<int>> G, rev_G;
     deque<bool> used;
     vector<int> order, comp;
     deque<bool> assignment;
     void init(int n) {
         n = _n;
         N = n * 2;
         G.resize(N + 5);
         rev_G.resize(N + 5);
     void dfs1(int v) {
         used[v] = true;
         for (int u : G[v]) {
             if (!used[u])
                 dfs1(u);
         order.push_back(v);
     void dfs2(int v, int cl) {
         comp[v] = c1;
         for (int u : rev_G[v]) {
             if (comp[u] == -1)
                 dfs2(u, c1);
     bool solve() {
         order.clear();
         used.assign(N, false);
         for (int i = 0; i < N; ++i) {
             if (!used[i])
                 dfs1(i);
         comp.assign(N, -1);
         for (int i = 0, j = 0; i < N; ++i) {
             int v = order[N - i - 1];
             if (comp[v] == -1)
                 dfs2(v, j++);
         assignment.assign(n, false);
         for (int i = 0; i < N; i += 2) {
             if (comp[i] == comp[i + 1])
                 return false;
             assignment[i / 2] = (comp[i] > comp[i + 1]);
         return true;
     void add_disjunction(int a, bool na, int b, bool nb) { //
          A or B 都是 0-based
         // na means whether a is negative or not
         // nb means whether b is negative or not
         a = 2 * a ^ na;
         b = 2 * b ^ nb;
         int neg a = a ^ 1;
         int neg_b = b ^ 1;
```

```
G[neg_a].push_back(b);
           G[neg b].push back(a);
           rev_G[b].push_back(neg_a);
           rev_G[a].push_back(neg_b);
           return:
      void get result(vector<int>& res) {
           res.clear();
           for (int i = 0; i < n; i++)</pre>
               res.push_back(assignment[i]);
  /* CSES Giant Pizza
72 3 5
73 + 1 + 2
  - 1 + 3
   - + + + -
   */
  int main() {
      int n, m;
      cin >> n >> m;
      TWO_SAT E;
      E.init(m);
       char c1, c2;
      int inp1, inp2;
      for (int i = 0; i < n; i++) {</pre>
           cin >> c1 >> inp1;
           cin >> c2 >> inp2;
           E.add_disjunction(inp1 - 1, c1 == '-', inp2 - 1, c2
       bool able = E.solve();
      if (able) {
           vector <int> ans;
           E.get_result(ans);
           for (int i : ans)
               cout << (i == true ? '+' : '-') << ' ';
           cout << '\n';
      } else {
           cout << "IMPOSSIBLE\n";</pre>
```

1.2 Custom Set PO Sort

return 0;

101

```
// priority_queue · 務必檢查相等的 case · 給所有元素一個排序的
依據
struct cmp{
bool operator () (Data a, Data b){
return a.x<b.x;
};
priority_queue<Data, vector<Data>, cmp> pq;
// set · 務必檢查相等的 case · 給所有元素一個排序的依據
struct Data{
int x;
```

```
bool operator < (const Data &b) const {
    return x<b.x;
}

}
</pre>
```

1.3 Default Code New

```
#include <bits/stdc++.h>
using namespace std;
#define int long long

const int MAX_N = 5e5 + 10;
const int INF = 2e18;

void solve(){

signed main(){
   ios::sync_with_stdio(0), cin.tie(0);
   int t = 1;
   while (t--){
        solve();
   }

return 0;
}
```

1.4 Default Code Old

```
| #include <bits/stdc++.h>
 #define int long long
 #define ALL(x) x.begin(), x.end()
 #define SZ(x) ((int)x.size())
 #define fastio ios::sync with stdio(0), cin.tie(0);
 using namespace std;
 #ifdef LOCAL
 #define cout cout << "\033[0;32m"
 #define cerr cerr << "\033[0;31m"
 #define endl endl << "\033[0m"</pre>
 #pragma GCC optimize("03,unroll-loops")
 #pragma GCC target("avx,avx2,sse,sse2,sse3,sse4,popcnt")
 #define end1 "\n"
 const int MAX N = 5e5+10;
 const int INF = 2e18;
 void solve1(){
     return:
 signed main(){
     fastio;
```

```
int t = 1;
while (t--){
solve1();
}
return 0;
```

1.5 Enumerate Subset

1.6 Fast Input

```
1 // fast IO
2 // 6f8879
3 inline char readchar(){
      static char buffer[BUFSIZ], * now = buffer + BUFSIZ, *
           end = buffer + BUFSIZ;
      if (now == end)
          if (end < buffer + BUFSIZ)</pre>
               return EOF;
           end = (buffer + fread(buffer, 1, BUFSIZ, stdin));
          now = buffer:
12
      return *now++;
14 inline int nextint(){
      int x = 0, c = readchar(), neg = false;
      while (('0' > c | c > '9') \&\& c! = '-' \&\& c! = EOF) c =
      if(c == '-') neg = true, c = readchar();
      while ('0' \le c \&\& c \le '9') x = (x << 3) + (x << 1) + (c^{'0'})
           , c = readchar();
      if(neg) x = -x;
      return x; // returns 0 if EOF
```

1.7 Radix Sort

```
cnt[i] += cnt[i-1];
static int temp;
for(int i = n-1; i >= 0; --i){
   temp = a[i] & 32767;
    --cnt[temp];
   tmpa[cnt[temp]] = a[i];
static int cnt2[32768] = {0};
for(int i = 0; i < n; ++i)</pre>
   ++cnt2[(tmpa[i]>>15)];
for(int i = 1; i < 32768; ++i)</pre>
   cnt2[i] += cnt2[i-1];
for(int i = n-1; i >= 0; --i){
   temp = (tmpa[i]>>15);
    --cnt2[temp];
   a[cnt2[temp]] = tmpa[i];
return;
```

1.8 Xor Basis

```
vector<int> basis;
  void add vector(int x){
      for (auto v : basis){
         x=min(x, x^v);
      if (x) basis.push_back(x);
  // 給一數字集合 S,求能不能 XOR 出 x
  bool check(int x){
      for (auto v : basis){
         x=min(x, x^v);
      return x;
17 // 給一數字集合 S, 求能 XOR 出多少數字
  // 答案等於 2^{basis 的大小}
20 // 給一數字集合 S,求 XOR 出最大的數字
21 int get_max(){
      int ans=0:
      for (auto v : basis){
         ans=max(ans, ans^v);
      return ans;
```

1.9 random int

1.10 hash command

1.11 run

1.12 run2

```
1 import os
  f = "pA"
  while 1:
      i = input("input: ")
      p = os.listdir(".")
      if i != "":
          f = i
      print(f"file = {f}")
      if os.system(f"g++ {f}.cpp -std=c++17 -Wall -Wextra -
           Wshadow -02 -D LOCAL -g -fsanitize=undefined,address
            -o \{f\}'')!=0:
          print("CE")
          continue
      os.system("clear")
      for x in sorted(p):
          if f in x and ".in" in x:
              print(x)
              if os.system(f"./\{f\} < \{x\}")!=0:
20
                   print("RE")
21
               print()
```

1.13 setup

```
ı| se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
  :inoremap " ""<Esc>i
  :inoremap {<CR> {<CR>}<Esc>ko
  :inoremap {{ {} < ESC >i
  function! F(...)
   execute '!./%:r < ./' . a:1
  endfunction
  command! -nargs=* R call F(<f-args>)
12 map <F7> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -Wextra -
       Wshadow -02 -DLOCAL -g -fsanitize=undefined,address<CR>
13 map <F8> :!./%:r<CR>
14 map <F9> :!./%:r < ./%:r.in<CR>
16 ca hash w !cpp -dD -P -fpreprocessed \| tr -d "[:space:]" \|
      md5sum \| cut -c-6
  " i+<esc>25A---+<esc>
   o|<esc>25A |<esc>
  " "ggVGyG35pGdd
```

2 Convolution

2.1 FFT any mod

```
2| 修改 const int MOD = 998244353 更改要取餘的數字
3 PolyMul(a, b) 回傳多項式乘法的結果 (c_k = \sum_{i+j} a_i+b_j)
      mod MOD )
s| 大約可以支援 5e5 · a_i, b_i 皆在 MOD 以下的非負整數
  const int MOD = 998244353;
  typedef complex<double> cd;
void FFT(vector<cd> &a) {
     int n = a.size(), L = 31-__builtin_clz(n);
      vector<complex<long double>> R(2, 1);
      vector<cd> rt(2, 1);
      for (int k=2; k < n; k*=2){
          R.resize(n);
          rt.resize(n);
          auto x = polar(1.0L, acos(-1.0L) / k);
          for (int i=k ; i<2*k ; i++){</pre>
              rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);
      vector<int> rev(n);
      for (int i=0 ; i<n ; i++){</pre>
          rev[i] = (rev[i/2] | (i&1) << L)/2;
      for (int i=0 ; i<n ; i++){</pre>
          if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
```

```
for (int k=1 ; k<n ; k*=2){</pre>
        for (int i=0 ; i<n ; i+=2*k){</pre>
             for (int j=0 ; j<k ; j++){</pre>
                 auto x = (double *)&rt[j+k];
                 auto y = (double *)&a[i+j+k];
                 cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]* 17
                      y[0]);
                 a[i+j+k] = a[i+j]-z;
                 a[i+j] += z;
                                                                    21
        }
    return;
vector<int> PolyMul(vector<int> a, vector<int> b){
    if (a.empty() || b.empty()) return {};
    vector<int> res(a.size()+b.size()-1);
    int B = 32- builtin clz(res.size()), n = (1<<B), cut =</pre>
         int(sqrt(MOD));
    vector<cd> L(n), R(n), outs(n), outl(n);
    for (int i=0 ; i<a.size() ; i++){</pre>
                                                                    35
        L[i] = cd((int) a[i]/cut, (int)a[i]%cut);
    for (int i=0 ; i<b.size() ; i++){</pre>
        R[i] = cd((int) b[i]/cut, (int)b[i]%cut);
    FFT(L);
    FFT(R);
    for (int i=0 ; i<n ; i++){</pre>
        int j = -i&(n-1);
        \operatorname{outl}[j] = (L[i] + \operatorname{conj}(L[j])) * R[i]/(2.0*n);
        outs[j] = (L[i]-conj(L[j])) * R[i]/(2.0*n)/1i;
    FFT(outl);
    FFT(outs);
    for (int i=0 ; i<res.size() ; i++){</pre>
        int av = (int)(real(outl[i])+0.5), cv = (int)(imag(
             outs[i])+0.5);
        int bv = (int)(imag(outl[i])+0.5) + (int)(real(outs[i
        res[i] = ((av%MOD*cut+bv) % MOD*cut+cv) % MOD;
                                                                    56
    return res;
                                                                    57
```

2.2 FFT new

```
typedef complex<double> cd;

void FFT(vector<cd> &a) {
    int n = a.size(), L = 31-_builtin_clz(n);
    vector<complex<long double>> R(2, 1);
    vector<cd> rt(2, 1);
    for (int k=2 ; k<n ; k*=2){
        R.resize(n);
        rt.resize(n);
    auto x = polar(1.0L, acos(-1.0L) / k);
        rt.resize(n), acos(-1.0L) / k);</pre>
```

2.3 FFT old

return res;

return;

FFT(in);

FFT(out);

double> b){

```
typedef complex<double > cd;
const double PI = acos(-1);

void FFT(vector<cd> &a, bool inv){
    int n = a.size();
    for (int i=1, j=0; i<n; i++){
        int bit = (n>>1);
}
```

for (int i=k ; i<2*k ; i++){</pre>

rev[i] = (rev[i/2] | (i&1) << L)/2;

for (int i=0 ; i<n ; i+=2*k){</pre>

if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>

for (int j=0 ; j<k ; j++){</pre>

y[0]);

a[i+j] += z;

if (a.empty() || b.empty()) return {};
vector<double> res(a.size()+b.size()-1);

copy(a.begin(), a.end(), begin(in));

for (int i=0 ; i<res.size() ; i++){
 res[i] = imag(out[i]) / (4 * n);</pre>

for (int i=0 ; i<b.size() ; i++){</pre>

vector<cd> in(n), out(n);

in[i].imag(b[i]);

for (cd& x : in) x *= x;

for (int i=0 ; i<n ; i++){</pre>

a[i+j+k] = a[i+j]-z;

auto x = (double *)&rt[j+k];

auto y = (double *)&a[i+j+k];

vector<double> PolyMul(const vector<double> a, const vector<</pre>

int L = 32 - __builtin_clz(res.size()), n = 1 << L;</pre>

out[i] = in[-i & (n - 1)] - conj(in[i]);

cd z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]*

vector<int> rev(n);

for (int i=0 ; i<n ; i++){</pre>

for (int i=0 ; i<n ; i++){</pre>

for (int k=1; k<n; k*=2){</pre>

rt[i] = R[i] = (i&1 ? R[i/2]*x : R[i/2]);

```
for ( ; j&bit ; bit>>=1){
            j ^= bit;
        i ^= bit;
        if (i<j){</pre>
            swap(a[i], a[j]);
   for (int len=2 ; len<=n ; len<<=1){</pre>
        cd wlen = polar(1.0, (inv ? 2 : -2)*PI/len);
        for (int i=0 ; i<n ; i+=len){</pre>
            cd w(1);
            for (int j=0 ; j<len/2 ; j++){</pre>
                 \dot{c}du = a[\dot{i}+\dot{j}];
                 cd v = a[i+j+len/2]*w;
                 a[i+j] = u+v;
                 a[i+j+len/2] = u-v;
                 w *= wlen;
   }
   if (inv){
        for (auto &x : a){
            x /= n;
   }
    return;
vector<cd> polyMul(vector<cd> a, vector<cd> b){
   int sa = a.size(), sb = b.size(), n = 1;
   while (n<sa+sb-1) n *= 2;</pre>
   a.resize(n);
   b.resize(n);
   vector<cd> c(n);
   FFT(a, 0);
   FFT(b, 0);
   for (int i=0 ; i<n ; i++) c[i] = a[i]*b[i];</pre>
   FFT(c, 1);
   c.resize(sa+sb-1);
   return c;
```

2.4 NTT mod 998244353

```
| \text{const int MOD} = (119 << 23) + 1, ROOT = 62; // = 998244353
2 // For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 <<
_3 // and 483 << 21 (same root). The last two are > 10^9.
5 // 9cd58a
6 void NTT(vector<int> &a) {
     int n = a.size();
     int L = 31-__builtin_clz(n);
     vector<int> rt(2, 1);
```

```
for (int k=2, s=2; k<n; k*=2, s++){
        rt.resize(n);
                                                                 11
        int z[] = \{1, qp(ROOT, MOD>>s)\};
                                                                 12
        for (int i=k ; i<2*k ; i++){</pre>
                                                                 13
            rt[i] = rt[i/2]*z[i&1]%MOD;
   }
    vector<int> rev(n);
    for (int i=0 ; i<n ; i++){</pre>
        rev[i] = (rev[i/2]|(i&1)<<L)/2;
    for (int i=0 ; i<n ; i++){</pre>
        if (i<rev[i]){</pre>
            swap(a[i], a[rev[i]]);
   }
    for (int k=1 ; k<n ; k*=2){</pre>
        for (int i=0; i<n; i+=2*k){
            for (int j=0 ; j<k ; j++){</pre>
                int z = rt[j+k]*a[i+j+k]%MOD, &ai = a[i+j];
                                                                  1 struct Persistent_Disjoint_Set{
                a[i+j+k] = ai-z+(z>ai ? MOD : 0);
                ai += (ai+z)=MOD ? z-MOD : z);
       }
vector<int> polyMul(vector<int> &a, vector<int> &b){
    if (a.empty() || b.empty()) return {};
    int s = a.size()+b.size()-1, B = 32-__builtin_clz(s), n = 12
    int inv = qp(n, MOD-2);
    vector<int> L(a), R(b), out(n);
    L.resize(n), R.resize(n);
    NTT(L), NTT(R);
    for (int i=0 ; i<n ; i++){</pre>
        out[-i&(n-1)] = L[i]*R[i]%MOD*inv%MOD;
    NTT(out);
    out.resize(s);
    return out;
                                                                 24
    Data-Structure
```

3.1 BIT

```
vector<int> BIT(MAX_SIZE);
void update(int pos, int val){
    for (int i=pos ; i<MAX_SIZE ; i+=i&-i){</pre>
        BIT[i]+=val;
int query(int pos){
    int ret=0;
```

```
return ret;
14 }
16 // const int MAX_N = (1 << 20)
int res = 0;
     for (int i=MAX N>>1 ; i>=1 ; i>>=1)
       if (bit[res+i]<k)</pre>
           k -= bit[res+=i];
     return res+1;
```

for (int i=pos ; i>0 ; i-=i&-i){

ret+=BIT[i];

3.2 Disjoint Set Persistent

```
Persistent_Segment_Tree arr, sz;
      void init(int n){
          arr.init(n);
          vector<int> v1:
          for (int i=0 ; i<n ; i++){</pre>
               v1.push back(i);
          arr.build(v1, 0);
          sz.init(n);
          vector<int> v2;
          for (int i=0 ; i<n ; i++){</pre>
              v2.push back(1);
          sz.build(v2, 0);
      int find(int a){
          int res = arr.query_version(a, a+1, arr.version.size
               ()-1).val;
          if (res==a) return a;
          return find(res);
      bool unite(int a, int b){
          a = find(a);
          b = find(b);
          if (a!=b){
              int sz1 = sz.query version(a, a+1, arr.version.
                   size()-1).val;
               int sz2 = sz.query version(b, b+1, arr.version.
                   size()-1).val;
35
              if (sz1<sz2){
                  arr.update_version(a, b, arr.version.size()
                  sz.update version(b, sz1+sz2, arr.version.
                       size()-1);
              }else{
                  arr.update version(b, a, arr.version.size()
```

3.3 PBDS GP Hash Table

```
i #include <ext/pb ds/assoc container.hpp>
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>, rb_tree_tag,
      tree order statistics node update> order set;
  struct custom_hash {
      static uint64_t splitmix64(uint64_t x) {
          // http://xorshift.di.unimi.it/splitmix64.c
          x += 0x9e3779b97f4a7c15;
          x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
          x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
          return x ^ (x >> 31);
      size t operator()(uint64 t x) const {
          static const uint64 t FIXED RANDOM = chrono::
               steady_clock::now().time_since_epoch().count();
          return splitmix64(x + FIXED RANDOM);
17 };
19 gp hash table (int, int, custom hash) ss;
```

3.4 PBDS Order Set

3.5 Segment Tree Add Set

```
1 // [ll, rr), based-0
2 // 使用前記得 init(陣列大小), build(陣列名稱)
3 // add(ll, rr): 區間修改
4 // set(ll, rr): 區間賦值
5 // query(ll, rr): 區間求和 / 求最大值
6 struct SegmentTree{
```

```
vector<node> arr:
SegmentTree(int n){
    arr.resize(n<<2);</pre>
node pull(node A, node B){
    node C:
    C.sum = A.sum+B.sum;
    C.ma = max(A.ma, B.ma);
    return C;
 // cce0c8
void push(int idx, int ll, int rr){
    if (arr[idx].set tag!=0){
        arr[idx].sum = (rr-ll)*arr[idx].set_tag;
        arr[idx].ma = arr[idx].set tag;
        if (rr-ll>1){
            arr[idx*2+1].add tag = 0;
            arr[idx*2+1].set tag = arr[idx].set tag;
            arr[idx*2+2].add_tag = 0;
            arr[idx*2+2].set tag = arr[idx].set tag;
        arr[idx].set_tag = 0;
    if (arr[idx].add_tag!=0){
        arr[idx].sum += (rr-ll)*arr[idx].add tag;
        arr[idx].ma += arr[idx].add tag;
        if (rr-ll>1){
             arr[idx*2+1].add_tag += arr[idx].add_tag;
            arr[idx*2+2].add_tag += arr[idx].add_tag;
        arr[idx].add tag = 0;
}
void build(vector<int> &v, int idx = 0, int ll = 0, int
     rr = n){
    if (rr-ll==1){
        arr[idx].sum = v[11];
        arr[idx].ma = v[ll];
    }else{
        int mid = (11+rr)/2;
        build(v, idx*2+1, ll, mid);
        build(v, idx*2+2, mid, rr);
        arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
}
void add(int gl, int gr, int val, int idx = 0, int ll =
     0, int rr =n){
    push(idx, 11, rr);
if (rr<=q1 || qr<=11) return;</pre>
    if (ql<=11 && rr<=qr){
        arr[idx].add tag += val;
        push(idx, ll, rr);
        return;
```

struct node{

};

int add tag = 0;

int set tag = 0:

int sum = 0;

int ma = 0;

```
int mid = (11+rr)/2;
           add(ql, qr, val, idx*2+1, ll, mid);
           add(ql, qr, val, idx*2+2, mid, rr);
           arr[idx]=pull(arr[idx*2+1], arr[idx*2+2]);
       void set(int ql, int qr, int val, int idx=0, int ll=0,
            int rr=n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return;</pre>
           if (q1<=11 && rr<=qr){
               arr[idx].add_tag = 0;
               arr[idx].set_tag = val;
               push(idx, ll, rr);
               return;
           int mid = (11+rr)/2;
           set(ql, qr, val, idx*2+1, ll, mid);
           set(ql, qr, val, idx*2+2, mid, rr);
           arr[idx] = pull(arr[idx*2+1], arr[idx*2+2]);
       node query(int ql, int qr, int idx = 0, int ll = 0, int
            rr = n){
           push(idx, ll, rr);
           if (rr<=ql || qr<=ll) return node();</pre>
           if (q1<=11 && rr<=qr) return arr[idx];</pre>
           int mid = (ll+rr)/2;
           return pull(query(ql, qr, idx*2+1, ll, mid), query(ql
                , qr, idx*2+2, mid, rr));
100 } ST;
```

3.6 Segment Tree Li Chao Line

```
全部都是 0-based
  LC_Segment_Tree st(n);
  update({a, b}): 插入一條 y=ax+b 的全域直線
  query(x): 查詢所有直線在位置 x 的最小值
  const int MAX V = 1e6+10; // 值域最大值
  struct LC Segment Tree{
     struct Node{ // y = ax+b
         int a = 0;
         int b = INF;
         int y(int x){
             return a*x+b;
20
      vector<Node> arr;
     LC Segment Tree(int n = 0){
25
         arr.resize(4*n);
```

```
void update(Node val, int idx = 0, int ll = 0, int rr =
    if (rr-ll==0) return;
    if (rr-ll==1){
        if (val.y(ll)<arr[idx].y(ll)){
           arr[idx] = val;
        return;
    }
    int mid = (11+rr)/2;
    if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
         的線斜率要比較小
    if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
        update(val, idx*2+1, ll, mid);
    }else{ // 交點在右邊
        swap(arr[idx], val); // 在左子樹中,新線比舊線還
        update(val, idx*2+2, mid, rr);
    return;
}
int query(int x, int idx = 0, int ll = 0, int rr = MAX_V)
    if (rr-ll==0) return INF;
    if (rr-ll==1){
        return arr[idx].y(ll);
    int mid = (11+rr)/2;
    if (x<mid){</pre>
        return min(arr[idx].y(x), query(x, idx*2+1, ll,
        return min(arr[idx].y(x), query(x, idx*2+2, mid,
            rr));
```

3.7 Segment Tree Li Chao Segment

```
75
76 };
    int y(int x){
        return a*x+b:
};
vector<Node> arr;
LC Segment Tree(int n = 0){
    arr.resize(4*n);
}
void update(Node val, int idx = 0, int ll = 0, int rr =
    if (rr-11==0) return;
    if (rr-ll<=1){</pre>
        if (val.y(ll) < arr[idx].y(ll)) {</pre>
            arr[idx] = val:
        return;
    int mid = (11+rr)/2;
    if (arr[idx].a > val.a) swap(arr[idx], val); // 原本
         的線斜率要比較小
    if (arr[idx].y(mid) < val.y(mid)){ // 交點在左邊
        update(val, idx*2+1, ll, mid);
    }else{ // 交點在右邊
        swap(arr[idx], val); // 在左子樹中,新線比舊線還
        update(val, idx*2+2, mid, rr);
    return;
}
// 在 [al, ar) 加上一條 val 的線段
void update segment(Node val, int ql, int qr, int idx =
     0, int 11 = 0, int rr = MAX_V){
    if (rr-ll==0) return;
    if (rr<=ql || qr<=ll) return;</pre>
    if (q1<=11 && rr<=qr){
        update(val, idx, ll, rr);
        return;
    int mid = (11+rr)/2;
    update_segment(val, ql, qr, idx*2+1, ll, mid);
    update_segment(val, ql, qr, idx*2+2, mid, rr);
}
int query(int x, int idx = 0, int ll = 0, int rr = MAX V)
    if (rr-ll==0) return INF;
    if (rr-ll==1){
        return arr[idx].y(ll);
    int mid = (11+rr)/2;
    if (x<mid){</pre>
        return min(arr[idx].y(x), query(x, idx*2+1, ll,
        return min(arr[idx].y(x), query(x, idx*2+2, mid,
             rr));
```

3.8 Segment Tree Persistent

```
全部都是 0-based
  Persistent Segment Tree st(n+q);
  st.build(v, 0);
  承式:
  update_version(pos, val, ver): 對版本 ver 的 pos 位置改成 val
10 query version(ql, qr, ver): 對版本 ver 查詢 [ql, qr) 的區間和
11 clone version(ver): 複製版本 ver 到最新的版本
  struct Persistent_Segment_Tree{
      int node_cnt = 0;
      struct Node{
          int lc = -1;
          int rc = -1;
          int val = 0:
      };
      vector<Node> arr;
      vector<int> version;
      Persistent Segment Tree(int sz){
          arr.resize(32*sz);
          version.push back(node cnt++);
          return;
      void pull(Node &c, Node a, Node b){
          c.val = a.val+b.val;
          return;
33
      void build(vector<int> &v, int idx, int ll = 0, int rr =
          auto &now = arr[idx];
          if (rr-ll==1){
37
              now.val = v[11]:
              return:
          int mid = (11+rr)/2;
          now.lc = node cnt++:
          now.rc = node cnt++:
          build(v, now.lc, ll, mid);
          build(v, now.rc, mid, rr);
          pull(now, arr[now.lc], arr[now.rc]);
      void update(int pos, int val, int idx, int ll = 0, int rr
          auto &now = arr[idx];
54
          if (rr-ll==1){
              now.val = val;
```

```
return;
           int mid = (11+rr)/2;
           if (pos<mid){</pre>
               arr[node cnt] = arr[now.lc];
               now.lc = node cnt;
               node cnt++;
               update(pos, val, now.lc, ll, mid);
               arr[node cnt] = arr[now.rc];
               now.rc = node_cnt;
               node_cnt++;
               update(pos, val, now.rc, mid, rr);
           pull(now, arr[now.lc], arr[now.rc]);
           return:
      }
      void update_version(int pos, int val, int ver){
           update(pos, val, version[ver]);
      Node query(int ql, int qr, int idx, int ll = 0, int rr =
           auto &now = arr[idx];
           if (q1<=11 && rr<=qr) return now;</pre>
           if (rr<=ql || qr<=ll) return Node();</pre>
           int mid = (11+rr)/2;
           Node ret:
           pull(ret, query(ql, qr, now.lc, ll, mid), query(ql,
                qr, now.rc, mid, rr));
           return ret;
      Node query_version(int ql, int qr, int ver){
           return query(ql, qr, version[ver]);
      void clone_version(int ver){
           version.push back(node cnt);
           arr[node_cnt] = arr[version[ver]];
           node cnt++;
101 };
  3.9 Sparse Table
```

3.10 Treap

struct Treap{

```
Treap *1 = nullptr, *r = nullptr;
       int pri = rand(), val = 0, sz = 1;
       Treap(int _val){
           val = _val;
  int size(Treap *t){return t ? t->sz : 0;}
  void pull(Treap *t){
       t\rightarrow sz = size(t\rightarrow l) + size(t\rightarrow r) + 1;
  Treap* merge(Treap *a, Treap *b){
       if (!a || !b) return a ? a : b;
       if (a->pri>b->pri){
           a \rightarrow r = merge(a \rightarrow r, b);
           pull(a);
           return a;
       }else{
           b - > 1 = merge(a, b - > 1);
           pull(b);
           return b;
30 pair<Treap*, Treap*> split(Treap *&t, int k){ // 1-based <前
       k 個元素, 其他元素>
       if (!t) return {};
       if (size(t->1)>=k){
           auto pa = split(t->1, k);
           t->1 = pa.second;
           pull(t);
           return {pa.first, t};
           auto pa = split(t->r, k-size(t->l)-1);
           t->r = pa.first;
           pull(t);
           return {t, pa.second};
  // functions
46 Treap* build(vector<int> v){
       Treap* ret;
       for (int i=0 ; i<SZ(v) ; i++){</pre>
```

```
ret = merge(ret, new Treap(v[i]));
51
      return ret:
52 }
53
  array<Treap*, 3> cut(Treap *t, int l, int r){ // 1-based <前
       1~l-1 個元素, l~r 個元素, r+1 個元素>
      array<Treap*, 3> ret;
      tie(ret[1], ret[2]) = split(t, r);
      tie(ret[0], ret[1]) = split(ret[1], 1-1);
      return ret:
59
  void print(Treap *t, bool flag = true){
      if (t->1!=0) print(t->1, false);
      cout << t->val;
      if (t->r!=0) print(t->r, false);
      if (flag) cout << endl;</pre>
```

3.11 Trie

```
struct Trie{
      struct Data{
          int nxt[2]={0, 0};
      vector<Data> arr;
      void init(int n){
          arr.resize(n);
      void insert(int n){
          int now=0;
          for (int i=N ; i>=0 ; i--){
              int v=(n>>i)&1;
              if (!arr[now].nxt[v]){
                  arr[now].nxt[v]=++sz;
              now=arr[now].nxt[v];
      int query(int n){
          int now=0, ret=0;
          for (int i=N ; i>=0 ; i--){
              int v=(n>>i)&1;
              if (arr[now].nxt[1-v]){
                  ret+=(1<<i);
                  now=arr[now].nxt[1-v];
              }else if (arr[now].nxt[v]){
                  now=arr[now].nxt[v];
33
              }else{
34
                  return ret;
35
36
37
          return ret;
39
40 } tr;
```

4 Dynamic-Programming

4.1 Digit DP

i #include <bits/stdc++.h>

```
using namespace std;
4 long long 1, r;
5 | long long dp[20][10][2][2]; // dp[pos][pre][limit] = 後 pos
      位·pos 前一位是 pre·(是/否)有上界·(是/否)有前綴零
7 long long memorize_search(string &s, int pos, int pre, bool
      limit, bool lead){
     // 已經被找過了,直接回傳值
     if (dp[pos][pre][limit][lead]!=-1) return dp[pos][pre][
          limit][lead];
     // 已經搜尋完畢,紀錄答案並回傳
     if (pos==(int)s.size()){
         return dp[pos][pre][limit][lead] = 1;
     // 枚舉目前的位數數字是多少
     long long ans = 0;
     for (int now=0 ; now<=(limit ? s[pos]-'0' : 9) ; now++){</pre>
         if (now==pre){
            // 1~9 絕對不能連續出現
            if (pre!=0) continue;
            // 如果已經不在前綴零的範圍內, 0 不能連續出現
            if (lead==false) continue:
         ans += memorize_search(s, pos+1, now, limit&(now==(s[
             pos]-'0')), lead&(now==0));
     // 已經搜尋完畢,紀錄答案並回傳
     return dp[pos][pre][limit][lead] = ans;
36 | // 回傳 [0, n] 有多少數字符合條件
37 long long find_answer(long long n){
     memset(dp, -1, sizeof(dp));
     string tmp = to string(n);
     return memorize search(tmp, 0, 0, true, true);
44 int main(){
     // input
     cin >> 1 >> r;
     // output - 計算 [L, r] 有多少數字任意兩個位數都不相同
     cout << find_answer(r)-find_answer(l-1) << "\n";</pre>
     return 0;
```

4.2 SOS DP

```
1 // 總時間複雜度為 O(n 2^n)
2 // 計算 dp[i] = i 所有 bit mask 子集的和
3 for (int i=0; i<n; i++){
4 for (int mask=0; mask<(1<<n); mask++){
5 if ((mask>>i)&1){
6 dp[mask] += dp[mask^(1<<i)];
7 }
8 }
9 }</pre>
```

4.3 Integer Partition

```
dp[i][x] = 要將整數 x 拆成 i 堆的「組合數」 dp[i+1][x+1]+=dp[i][x] ( 創造新的一堆 ) dp[i][x+i]+=dp[i][x] ( 把每一堆都增加 1 )
```

5 Geometry

5.1 Geometry Struct

```
1 // 判斷數值正負: {1:正數,0:零,-1:負數}
 int sign(long long x) {return (x \ge 0) ? ((bool)x) : -1; }
 int sign(double x) {
     return (abs(x) < 1e-9) ? 0 : (x > 0 ? 1 : -1);
 template<tvpename T>
 struct point {
     T x, y;
     point() {}
     point(const T &x, const T &y) : x(x), y(y) {}
     point operator+(point b) {return {x+b.x, y+b.y}; }
     point operator-(point b) {return {x-b.x, y-b.y}; }
     point operator*(T b) {return {x*b, y*b}; }
     point operator/(T b) {return {x/b, y/b}; }
     bool operator==(point b) {return x==b.x && y==b.y; }
     // 逆時針極角排序
     bool operator<(point &b) {return (x*b.y > b.x*y); }
     friend ostream& operator<<(ostream& os, point p) {</pre>
         os << "(" << p.x << ", " << p.y << ")";
     // 判斷 ab 到 ac 的方向: {1:逆時鐘,0:重疊,-1:順時鐘}
     friend int ori(point a, point b, point c) {
         return sign((b-a)^(c-a));
     friend bool btw(point a, point b, point c) {
         return ori(a, b, c) == 0 \&\& sign((a-c)*(b-c)) <= 0;
     // 判斷線段 ab. cd 是否相交
     friend bool banana(point a, point b, point c, point d) {
         int s1 = ori(a, b, c);
         int s2 = ori(a, b, d);
         int s3 = ori(c, d, a);
```

```
int s4 = ori(c, d, b);
          if (btw(a, b, c) || btw(a, b, d) || btw(c, d, a) ||
              btw(c, d, b)) return 1;
          return (s1 * s2 < 0) && (s3 * s4 < 0);
      T operator*(point b) {return x * b.x + y * b.y; }
      T operator^(point b) {return x * b.y - y * b.x; }
      T abs2() {return (*this) * (*this); }
      // 旋轉 Arg(b) 的角度(小心溢位)
      point rotate(point b) {return {x*b.x - y*b.y, x*b.y + y*b
47
  template<typename T>
  struct line {
      point<T> p1, p2;
      // ax + by + c = 0
      T a, b, c; //|a|, |b| \le 2C, |c| \le 8C^2
      line(const point<T> &x, const point<T> &y) : p1(x), p2(y){
          build();
      void build() {
      a = p1.y - p2.y;
      b = p2.x - p1.x:
      c = (-a*p1.x)-b*p1.y;
      // 判斷點和有向直線的關係: {1:左邊,0:在線上,-1:右邊}
    int ori(point<T> &p) {
      return sign((p2-p1) ^ (p-p1));
    // 判斷直線斜率是否相同
    bool parallel(line &1) {
      return ((p1-p2) ^ (l.p1-l.p2)) == 0;
    // 兩直線交點
      point<long double> line intersection(line &1) {
          using P = point<long double>:
      point < T > a = p2-p1, b = 1.p2-1.p1, s = 1.p1-p1;
      return P(p1.x,p1.y) + P(a.x,a.y) * (((long double)(s^b))
77
  };
  template<tvpename T>
80 struct polygon {
      vector<point<T>> v;
      polygon() {}
      polygon(const vector<point<T>> &u) : v(u) {}
      // simple 為 true 的時候會回傳任意三點不共線的凸包
      void make convex hull(int simple) {
          auto cmp = [&](point<T> &p, point<T> &q) {
              return (p.x == q.x)? (p.y < q.y): (p.x < q.x);
          simple = (bool)simple;
          sort(v.begin(), v.end(), cmp);
          v.resize(unique(v.begin(), v.end()) - v.begin());
          vector<point<T>> hull:
          for (int t = 0; t < 2; ++t){
              int sz = hull.size();
              for (auto &i:v) {
                 while (hull.size() >= sz+2 && ori(hull[hull.
                      size()-2], hull.back(), i) < simple) {</pre>
```

```
hull.pop back();
                  hull.push back(i):
101
              hull.pop_back();
102
              reverse(v.begin(), v.end());
103
104
          swap(hull, v);
105
      可以在有 n 個點的簡單多邊形內,用 O(n)判斷一個點:
      {1: 在多邊形內,0:在多邊形上,-1:在多邊形外}
      int in polygon(point<T> a){
108
          const T MAX POS = 1e9 + 5; // [記得修改] 座標的最大值
          point<T> pre = v.back(), b(MAX POS, a.y + 1);
110
                                                               172
111
          int cnt = 0;
                                                              173
112
                                                               174
          for (auto &i:v) {
113
                                                              175
114
              if (btw(pre, i, a)) return 0;
                                                               176
115
              if (banana(a, b, pre, i)) cnt++;
116
              pre = i;
                                                               178
117
118
          return cnt%2 ? 1 : -1;
119
121 | /// 警告:以下所有凸包專用的函式都只接受逆時針排序且任三點不
       共線的凸包 ///
                                                               184
      可以在有 n 個點的凸包內,用 O(\log n) 判斷一個點:
                                                               185
123 // {1:在凸包內, 0:在凸包邊上, -1:在凸包外}
                                                              186
      int in_convex(point<T> p) {
                                                               187
125
          int n = v.size();
          int a = ori(v[0], v[1], p), b = ori(v[0], v[n-1], p); 189
126
          if (a < 0 || b > 0) return -1;
127
128
          if (btw(v[0], v[1], p)) return 0;
                                                               191
          if (btw(v[0], v[n - 1], p)) return 0;
129
130
          int l = 1, r = n - 1, mid;
          while (1 + 1 < r) {
131
132
              mid = (1 + r) >> 1;
              if (ori(v[0], v[mid], p) >= 0) 1 = mid;
133
134
              else r = mid:
135
          int k = ori(v[1], v[r], p);
136
137
          if (k <= 0) return k;</pre>
          return 1;
138
139
   // 凸包專用的環狀二分搜·回傳 0-based index
      int cycle search(auto &f) {
141
142
          int n = v.size(), l = 0, r = n;
          bool rv = f(1, 0);
143
144
          while (r - 1 > 1) {
                                                              204
              int m = (1 + r) / 2;
145
                                                               205
146
              if (f(0, m) ? rv: f(m, (m + 1) % n)) r = m;
                                                              206
147
              else 1 = m:
                                                              207
148
                                                               208
149
          return f(1, r % n) ? 1 : r % n;
                                                               209
                                                              210
151 // 可以在有 n 個點的凸包內 · 用 O(Log n) 判斷一條直線:
                                                              211
      {1: 穿過凸包,0: 剛好切過凸包,-1: 沒碰到凸包}
                                                              212
      int line cut convex(line<T> L) {
                                                              213
          point<T> p(L.a, L.b); // 記得 L 要 build
154
                                                              214
          auto gt = [&](int neg) {
155
                                                              215
156
              auto f = [&](int x, int y) {
                                                              216
157
                  return sign((v[x] - v[y]) * p) == neg;
                                                              217
                                                              218
              return -(v[cycle search(f)] * p);
```

```
220
           T x = gt(1), y = gt(-1);
                                                                221
           if (L.c < x || y < L.c) return -1;
           return not (L.c == x || L.c == y);
                                                                223
                                                                224
                                                                225
   // 可以在有 n 個點的凸包內 · 用 O(Log n) 判斷一個線段:
                                                                226
166 // {1: 存在一個凸包上的邊可以把這個線段切成兩半,
                                                                227
       0: 有碰到凸包但沒有任何凸包上的邊可以把它切成兩半,
168 // -1: 沒碰到凸包}
169 1// 除非線段兩端點都不在凸包邊上,否則此函數回傳 8 的時候不一
                                                                229
       定表示線段沒有通過凸包內部 ///
                                                                230
       int segment_across_convex(line<T> L) {
           point<T> p(L.a, L.b); // 記得 L 要 build
                                                                231
           auto gt = [&](int neg) {
                                                                232
               auto f = [&](int x, int y) {
                  return sign((v[x] - v[y]) * p) == neg;
                                                                233
                                                                234
               return cycle search(f);
                                                                235
                                                                236
           int i = gt(1), j = gt(-1), n = v.size();
                                                                237
          T x = -(v[i] * p), y = -(v[j] * p);
                                                                238
          if (L.c < x || y < L.c) return -1;
if (L.c == x || L.c == y) return 0;
                                                                239
                                                                240
           if (i > j) swap(i, j);
                                                                241
           auto g = [&](int x, int lim) {
                                                                242
               int now = 0, nxt;
                                                                243
               for (int i = 1 \leftrightarrow lg(lim); i > 0; i /= 2) {
                                                                244
                   if (now + i > lim) continue;
                                                                245
                   nxt = (x + i) % n;
                                                                246
                   if (L.ori(v[x]) * L.ori(v[nxt]) >= 0) {
                                                                247
                      x = nxt;
                                                                248
                      now += i;
                                                                249
                                                                250
               } // ↓ BE CAREFUL
               return -(ori(v[x], v[(x + 1) % n], L.p1) * ori(v[ 252
                   x], v[(x + 1) % n], L.p2));
                                                                254
           return max(g(i, j - i), g(j, n - (j - i)));
                                                                255
                                                                256
      可以在有 n 個點的凸包內,用 O(Log n)判斷一個線段:
                                                                257
       {1: 線段上存在某一點位於凸包內部(邊上不算).
                                                                258
       Ø: 線段上存在某一點碰到凸包的邊但線段上任一點均不在凸包
       內部.
                                                                260
201 // -1: 線段完全在凸包外面 }
                                                                261
       int segment pass convex interior(line<T> L) {
                                                                262
           if (in convex(L.p1) == 1 || in convex(L.p2) == 1)
                                                                263
           point<T> p(L.a, L.b); // 記得 L 要 build
                                                                264
           auto gt = [&](int neg) {
                                                                265
               auto f = [\&](int x, int y) {
                   return sign((v[x] - v[y]) * p) == neg;
                                                                266
                                                                267
               return cycle search(f);
                                                                268
                                                                260
           int i = gt(1), j = gt(-1), n = v.size();
                                                                270
          T x = -(v[i] * p), y = -(v[j] * p);
if (L.c < x || y < L.c) return -1;
                                                                271
                                                                272
           if (L.c == x || L.c == y) return 0;
                                                                273
                                                                274
           if (i > j) swap(i, j);
                                                                275
           auto g = [&](int x, int lim) {
                                                                276
               int now = 0. nxt;
               for (int i = 1 << __lg(lim); i > 0; i /= 2) {
```

```
if (now + i > lim) continue;
               nxt = (x + i) % n;
               if (L.ori(v[x]) * L.ori(v[nxt]) > 0) {
                    x = nxt;
                    now += i;
           } // ↓ BE CAREFUL
           return -(ori(v[x], v[(x + 1) % n], L.p1) * ori(v[
                x], v[(x + 1) % n], L.p2));
       int ret = max(g(i, j - i), g(j, n - (j - i)));
       return (ret == 0) ? (in_convex(L.p1) == 0 &&
            in_convex(L.p2) == 0) : ret;
// 回傳點過凸包的兩條切線的切點的 0-based index (不保證兩條
    切線的順逆時針關係)
   pair<int,int> convex tangent point(point<T> p) {
       int n = v.size(), z = -1, edg = -1;
       auto gt = [&](int neg) {
           auto check = [&](int x) {
               if (v[x] == p) z = x;
               if (btw(v[x], v[(x + 1) % n], p)) edg = x;
               if (btw(v[(x + n - 1) \% n], v[x], p)) edg = (
                    x + n - 1) % n;
           };
            auto f = [\&](int x, int y) {
               check(x); check(y);
               return ori(p, v[x], v[y]) == neg;
           };
            return cycle_search(f);
       int x = gt(1), y = gt(-1);
       if (z != -1) {
            return \{(z + n - 1) \% n, (z + 1) \% n\};
       else if (edg != -1) {
            return {edg, (edg + 1) % n};
       else {
            return {x, y};
   friend int halfplane intersection(vector<line<T>> &s,
        polygon<T> &P) {
       #define neg(p) ((p.y == 0 ? p.x : p.y) < 0)
       auto angle_cmp = [&](line<T> &A, line<T> &B) {
            point < T > a = A.p2-A.p1, b = B.p2-B.p1;
            return neg(a) < neg(b) \mid \mid (neg(a) == neg(b) \&\& (a)
                ^b) > 0);
       };
       sort(s.begin(), s.end(), angle cmp); // 線段左側為該
            線段半平面
       int L, R, n = s.size();
       vector<point<T>> px(n);
       vector<line<T>> q(n);
       q[L = R = 0] = s[0];
       for(int i = 1; i < n; ++i) {</pre>
           while(L < R && s[i].ori(px[R-1]) <= 0) --R;</pre>
           while(L < R && s[i].ori(px[L]) <= 0) ++L;</pre>
           q[++R] = s[i];
           if(q[R].parallel(q[R-1])) {
               if(q[R].ori(s[i].p1) > 0) q[R] = s[i];
```

```
if(L < R) px[R-1] = q[R-1].line_intersection(q[R 47])
                     1);
279
            while(L < R && q[L].ori(px[R-1]) <= 0) --R;
280
281
           P.v.clear();
           if(R - L <= 1) return 0:
282
283
           px[R] = q[R].line intersection(q[L]);
           for(int i = L; i <= R; ++i) P.v.push back(px[i]);</pre>
285
           return R - L + 1;
286
287 };
```

51

5.2 Geometry 卦長

```
const double PI=atan2(0.0,-1.0);
 2 template<typename T>
  struct point{
   T x,y;
    point(){}
    point(const T&x,const T&y):x(x),y(y){}
    point operator+(const point &b)const{
      return point(x+b.x,y+b.y); }
    point operator-(const point &b)const{
      return point(x-b.x,y-b.y); }
    point operator*(const T &b)const{
      return point(x*b,y*b); }
    point operator/(const T &b)const{
      return point(x/b,y/b); }
    bool operator == (const point &b)const{
      return x==b.x&&y==b.y; }
    T dot(const point &b)const{
      return x*b.x+y*b.y; }
    T cross(const point &b)const{
      return x*b.y-y*b.x; }
    point normal()const{//求法向量
      return point(-y,x); }
    T abs2()const{//向量長度的平方
      return dot(*this); }
   T rad(const point &b)const{//兩向量的弧度
  return fabs(atan2(fabs(cross(b)),dot(b))); }
   T getA()const{//對x軸的弧度
      T A=atan2(y,x);//超過180度會變負的
      if(A<=-PI/2)A+=PI*2;</pre>
      return A:
32 };
  template<typename T>
  struct line{
   line(){}
    point<T> p1,p2;
    T a,b,c;//ax+by+c=0
    line(const point<T>&x,const point<T>&y):p1(x),p2(y){}
    void pton(){//轉成一般式
      a=p1.v-p2.v:
      b=p2.x-p1.x;
      c=-a*p1.x-b*p1.y;
   T ori(const point<T> &p)const{//點和有向直線的關係, >0左
         邊、=0在線上<0右邊
      return (p2-p1).cross(p-p1);
```

```
T btw(const point<T> &p)const{//點投影落在線段上<=0
       return (p1-p).dot(p2-p);
                                                                105
    bool point on segment(const point<T>&p)const{//點是否在線段
      return ori(p) == 0&&btw(p) <= 0;</pre>
                                                                110
52
    T dis2(const point<T> &p,bool is_segment=0)const{//點跟直線 111
         /線段的距離平方
                                                                113
       point<T> v=p2-p1,v1=p-p1;
                                                                114
       if(is segment){
                                                                115
        point<T> v2=p-p2;
        if(v.dot(v1)<=0)return v1.abs2();</pre>
        if(v.dot(v2)>=0)return v2.abs2();
                                                                118
       T tmp=v.cross(v1);
       return tmp*tmp/v.abs2();
                                                                121
    T seg dis2(const line<T> &1)const{//兩線段距離平方
      return min({dis2(l.p1,1),dis2(l.p2,1),l.dis2(p1,1),l.dis2 123
           (p2,1)});
                                                                125
65
    point<T> projection(const point<T> &p)const{//點對直線的投
       point<T> n=(p2-p1).normal();
                                                                129
       return p-n*(p-p1).dot(n)/n.abs2();
                                                                131
    point<T> mirror(const point<T> &p)const{
                                                                132
      //點對直線的鏡射,要先呼叫pton轉成一般式
                                                                133
      point<T> R:
                                                                134
      T d=a*a+b*b;
                                                                135
      R.x=(b*b*p.x-a*a*p.x-2*a*b*p.y-2*a*c)/d;
      R.y=(a*a*p.y-b*b*p.y-2*a*b*p.x-2*b*c)/d;
      return R;
                                                                137
77 J
                                                                138
    bool equal(const line &1)const{//直線相等
                                                                139
      return ori(1.p1)==0&&ori(1.p2)==0:
                                                                140
                                                                141
    bool parallel(const line &1)const{
      return (p1-p2).cross(l.p1-l.p2)==0;
                                                                142
                                                                143
    bool cross seg(const line &1)const{
      return (p2-p1).cross(l.p1-p1)*(p2-p1).cross(l.p2-p1)<=0;</pre>
           //直線是否交線段
    int line intersect(const line &1)const{//直線相交情況,-1無
         限多點、1交於一點、0不相交
       return parallel(1)?(ori(1.p1)==0?-1:0):1;
                                                                149
                                                                150
    int seg intersect(const line &1)const{
                                                                151
      T c1=ori(l.p1), c2=ori(l.p2);
                                                                152
      T c3=1.ori(p1), c4=1.ori(p2);
                                                                153
      if(c1==0&&c2==0){//共線
                                                                154
        bool b1=btw(1.p1)>=0,b2=btw(1.p2)>=0;
                                                                155
        T a3=1.btw(p1),a4=1.btw(p2);
                                                                156
        if(b1&&b2&&a3==0&&a4>=0) return 2;
                                                                157
        if(b1&&b2&&a3>=0&&a4==0) return 3;
        if(b1&&b2&&a3>=0&&a4>=0) return 0;
                                                                159
        return -1;//無限交點
       }else if(c1*c2<=0&&c3*c4<=0)return 1;</pre>
       return 0;//不相交
    point<T> line_intersection(const line &1)const{/*直線交點*/ 164
```

```
point<T> a=p2-p1,b=l.p2-l.p1,s=l.p1-p1;
   //if(a.cross(b)==0)return INF;
   return p1+a*(s.cross(b)/a.cross(b));
 point<T> seg intersection(const line &1)const{//線段交點
   int res=seg intersect(1);
   if(res<=0) assert(0);</pre>
   if(res==2) return p1;
   if(res==3) return p2;
   return line intersection(1);
template<typename T>
struct polygon{
 polygon(){}
 vector<point<T> > p;//逆時針順序
 T area()const{//面積
   T ans=0:
   for(int i=p.size()-1,j=0;j<(int)p.size();i=j++)</pre>
     ans+=p[i].cross(p[i]);
   return ans/2:
 point<T> center of mass()const{//重心
   T cx=0, cy=0, w=0;
   for(int i=p.size()-1,j=0;j<(int)p.size();i=j++){</pre>
     T a=p[i].cross(p[j]);
     cx+=(p[i].x+p[i].x)*a;
     cy += (p[i].y + p[j].y)*a;
   return point<T>(cx/3/w,cy/3/w);
 char ahas(const point<T>& t)const{//點是否在簡單多邊形內
      是的話回傳1、在邊上回傳-1、否則回傳0
   bool c=0:
   for(int i=0,j=p.size()-1;i<p.size();j=i++)</pre>
     if(line<T>(p[i],p[j]).point_on_segment(t))return -1;
     else if((p[i].y>t.y)!=(p[j].y>t.y)&&
     t.x<(p[j].x-p[i].x)*(t.y-p[i].y)/(p[j].y-p[i].y)+p[i].x
       c=!c;
   return c;
 char point_in_convex(const point<T>&x)const{
   int l=1,r=(int)p.size()-2;
   while(1 < = r){//點是否在凸多邊形內,是的話回傳1 < r在邊上回傳
        -1、否則回傳@
     int mid=(1+r)/2;
     T a1=(p[mid]-p[0]).cross(x-p[0]);
     T a2=(p[mid+1]-p[0]).cross(x-p[0]);
     if(a1>=0&&a2<=0){
       T res=(p[mid+1]-p[mid]).cross(x-p[mid]);
       return res>0?1:(res>=0?-1:0);
     }else if(a1<0)r=mid-1;</pre>
     else l=mid+1;
   return 0;
 vector<T> getA()const{//凸包邊對x軸的夾角
   vector<T>res://一定是遞增的
   for(size t i=0;i<p.size();++i)</pre>
     res.push_back((p[(i+1)%p.size()]-p[i]).getA());
   return res:
```

```
bool line intersect(const vector<T>&A,const line<T> &1)
                                                                              while (now.dot(p[1+1]-p[i]) < =now.dot(p[1]-p[i])) = (1+1)% 279
          const{//O(LogN)
                                                                                                                                             point<T> barycenter()const{//重心
166
       int f1=upper bound(A.begin(), A.end(), (1.p1-1.p2).getA()) - 220
                                                                             T d=now.abs2():
                                                                                                                                               return (a+b+c)/3;
                                                                              T tmp=now.cross(p[t]-p[i])*(now.dot(p[r]-p[i])-now.dot(_{282}
       int f2=upper_bound(A.begin(),A.end(),(1.p2-1.p1).getA())-
                                                                                  p[1]-p[i]))/d;
167
                                                                                                                                        283
                                                                                                                                             point<T> circumcenter()const{//外心
            A.begin():
                                                                              ans=min(ans,tmp);
                                                                                                                                        284
                                                                                                                                               static line<T> u.v:
       return 1.cross_seg(line<T>(p[f1],p[f2]));
168
                                                                                                                                        285
                                                                                                                                               u.p1=(a+b)/2;
169
                                                                           return p.pop back(),ans;
                                                                                                                                               u.p2=point<T>(u.p1.x-a.y+b.y,u.p1.y+a.x-b.x);
                                                                                                                                        286
     polygon cut(const line<T> &1)const{//凸包對直線切割,得到直 225
                                                                                                                                        287
                                                                                                                                               v.p1=(a+c)/2;
          線し左側的凸包
                                                                         T dis2(polygon &pl){//凸包最近距離平方
                                                                                                                                               v.p2=point<T>(v.p1.x-a.y+c.y,v.p1.y+a.x-c.x);
                                                                                                                                        288
                                                                           vector<point<T> > &P=p,&Q=pl.p;
                                                                                                                                        289
                                                                                                                                               return u.line intersection(v);
       polygon ans;
171
                                                                           int n=P.size(),m=0.size(),l=0,r=0;
       for(int n=p.size(),i=n-1,j=0;j<n;i=j++){</pre>
172
                                                                                                                                        290
                                                                         for(int i=0;i<n;++i)if(P[i].y<P[1].y)l=i;</pre>
173
         if(1.ori(p[i])>=0){
                                                                                                                                        291
                                                                                                                                              point<T> incenter()const{//內心
                                                                         for(int i=0;i<m;++i)if(0[i].y<0[r].y)r=i;</pre>
174
           ans.p.push_back(p[i]);
                                                                                                                                               T A=sqrt((b-c).abs2()),B=sqrt((a-c).abs2()),C=sqrt((a-b).
                                                                                                                                        292
           if(1.ori(p[j])<0)</pre>
                                                                           P.push back(P[0]), Q.push back(Q[0]);
175
                                                                                                                                                     abs2());
             ans.p.push_back(1.line_intersection(line<T>(p[i],p[ 232
                                                                           T ans=1e99:
176
                                                                                                                                               return point<T>(A*a.x+B*b.x+C*c.x,A*a.y+B*b.y+C*c.y)/(A+B
                                                                                                                                        293
                                                                           for(int i=0;i<n;++i){</pre>
                  il)));
                                                                                                                                                    +C);
                                                                              while((P[1]-P[1+1]).cross(Q[r+1]-Q[r])<0)r=(r+1)%m;
         }else if(l.ori(p[j])>0)
                                                                                                                                        294
                                                                              ans=min(ans,line<T>(P[1],P[1+1]).seg_dis2(line<T>(Q[r],
           ans.p.push_back(1.line_intersection(line<T>(p[i],p[j 235
178
                                                                                                                                        295
                                                                                                                                              point<T> perpencenter()const{//垂心
                                                                                  0[r+1])));
                                                                                                                                               return barycenter()*3-circumcenter()*2;
                                                                             l=(1+1)%n;
179
                                                                                                                                        297
                                                                    237
180
       return ans;
                                                                                                                                        298
                                                                                                                                           };
                                                                    238
                                                                           return P.pop back(),Q.pop back(),ans;
181
                                                                                                                                           template<typename T>
                                                                    239
     static bool monotone_chain_cmp(const point<T>& a,const
                                                                                                                                           struct point3D{
                                                                         static char sign(const point<T>&t){
                                                                    240
          point<T>& b){//凸包排序函數
                                                                                                                                             T x,y,z;
                                                                                                                                        301
                                                                    241
                                                                           return (t.y==0?t.x:t.y)<0;</pre>
183
       return (a.x<b.x)||(a.x==b.x&&a.y<b.y);
                                                                                                                                             point3D(){}
                                                                    242
184
                                                                                                                                              point3D(const T&x,const T&y,const T&z):x(x),y(y),z(z){}
                                                                    243
                                                                         static bool angle cmp(const line<T>& A,const line<T>& B){
185
     void monotone chain(vector<point<T> > &s){//凸包
                                                                                                                                              point3D operator+(const point3D &b)const{
                                                                           point<T> a=A.p2-A.p1.b=B.p2-B.p1:
                                                                    244
186
       sort(s.begin(),s.end(),monotone chain cmp);
                                                                                                                                               return point3D(x+b.x,y+b.y,z+b.z);}
                                                                           return sign(a)<sign(b)||(sign(a)==sign(b)&&a.cross(b)>0);
                                                                    245
       p.resize(s.size()+1);
                                                                                                                                              point3D operator-(const point3D &b)const{
187
                                                                    246
                                                                                                                                               return point3D(x-b.x,y-b.y,z-b.z);}
188
                                                                         int halfplane_intersection(vector<line<T> > &s){//半平面交
189
       for(size t i=0:i<s.size():++i){</pre>
                                                                                                                                              point3D operator*(const T &b)const{
                                                                           sort(s.begin(),s.end(),angle cmp);//線段左側為該線段半平
         while(m \ge 2\&\&(p[m-1]-p[m-2]).cross(s[i]-p[m-2]) <= 0)--m;
                                                                                                                                               return point3D(x*b,y*b,z*b);}
190
191
         p[m++]=s[i];
                                                                                                                                              point3D operator/(const T &b)const{
                                                                           int L,R,n=s.size();
                                                                                                                                               return point3D(x/b,y/b,z/b);}
192
                                                                           vector<point<T> > px(n);
                                                                    250
193
       for(int i=s.size()-2,t=m+1;i>=0;--i){
                                                                                                                                              bool operator==(const point3D &b)const{
                                                                           vector<line<T> > q(n);
         while(m \ge t \& (p[m-1]-p[m-2]).cross(s[i]-p[m-2])<=0)--m;
                                                                                                                                               return x==b.x&&v==b.v&&z==b.z:}
194
                                                                           q[L=R=0]=s[0];
                                                                                                                                              T dot(const point3D &b)const{
195
         p[m++]=s[i];
                                                                           for(int i=1;i<n;++i){</pre>
196
                                                                                                                                               return x*b.x+v*b.v+z*b.z:}
                                                                    254
                                                                             while(L<R&&s[i].ori(px[R-1])<=0)--R;</pre>
197
       if(s.size()>1)--m;
                                                                                                                                        316
                                                                                                                                              point3D cross(const point3D &b)const{
                                                                              while(L<R&&s[i].ori(px[L])<=0)++L;</pre>
198
       p.resize(m);
                                                                    255
                                                                                                                                        317
                                                                                                                                               return point3D(y*b.z-z*b.y,z*b.x-x*b.z,x*b.y-y*b.x);}
                                                                    256
                                                                              q[++R]=s[i];
199
                                                                                                                                             T abs2()const{//向量長度的平方
                                                                    257
                                                                              if(q[R].parallel(q[R-1])){
                                                                                                                                               return dot(*this);}
     T diam(){//直徑
                                                                                                                                        319
                                                                    258
201
       int n=p.size(),t=1;
                                                                                                                                        320
                                                                                                                                             T area2(const point3D &b)const{//和b、原點圍成面積的平方
                                                                               if(q[R].ori(s[i].p1)>0)q[R]=s[i];
       T ans=0;p.push_back(p[0]);
                                                                                                                                               return cross(b).abs2()/4;}
                                                                                                                                        321
203
       for(int i=0;i<n;i++){</pre>
                                                                                                                                        322
                                                                                                                                           };
                                                                             if(L<R)px[R-1]=q[R-1].line_intersection(q[R]);</pre>
         point<T> now=p[i+1]-p[i];
                                                                                                                                        323
                                                                                                                                           template<typename T>
         while(now.cross(p[t+1]-p[i])>now.cross(p[t]-p[i]))t=(t
                                                                                                                                           struct line3D{
                                                                           while(L<R&&q[L].ori(px[R-1])<=0)--R;</pre>
                                                                                                                                        325
                                                                                                                                             point3D<T> p1,p2;
                                                                           p.clear();
206
         ans=max(ans,(p[i]-p[t]).abs2());
                                                                                                                                             line3D(){}
                                                                                                                                        326
                                                                           if(R-L<=1)return 0;</pre>
                                                                    265
207
                                                                                                                                        327
                                                                                                                                             line3D(const point3D<T> &p1,const point3D<T> &p2):p1(p1),p2
                                                                           px[R]=q[R].line intersection(q[L]);
                                                                    266
208
       return p.pop back(),ans;
                                                                           for(int i=L;i<=R;++i)p.push_back(px[i]);</pre>
209
                                                                                                                                             T dis2(const point3D<T> &p, bool is_segment=0)const{//點跟直
                                                                                                                                        328
                                                                    268
                                                                           return R-L+1;
     T min_cover_rectangle(){//最小覆蓋矩形
210
                                                                                                                                                   線/線段的距離平方
                                                                    269
211
       int n=p.size(),t=1,r=1,l;
                                                                                                                                                point3D < T > v = p2 - p1, v1 = p - p1;
                                                                                                                                        329
212
       if(n<3)return 0;//也可以做最小周長矩形
                                                                                                                                        330
                                                                                                                                               if(is segment){
                                                                       template<typename T>
       T ans=1e99;p.push_back(p[0]);
213
                                                                                                                                        331
                                                                                                                                                 point3D<T> v2=p-p2;
                                                                       struct triangle{
214
       for(int i=0;i<n;i++){</pre>
                                                                                                                                                 if(v.dot(v1)<=0)return v1.abs2();</pre>
                                                                         point<T> a,b,c;
         point<T> now=p[i+1]-p[i];
215
                                                                                                                                        333
                                                                                                                                                 if(v.dot(v2)>=0)return v2.abs2();
                                                                         triangle(){}
         while(now.cross(p[t+1]-p[i])>now.cross(p[t]-p[i]))t=(t
216
                                                                         triangle(const point<T> &a,const point<T> &b,const point<T> 334
                                                                                                                                        335
                                                                                                                                               point3D<T> tmp=v.cross(v1);
                                                                               &c):a(a),b(b),c(c){}
         while(now.dot(p[r+1]-p[i])>now.dot(p[r]-p[i]))r=(r+1)%n _{276}
217
                                                                                                                                        336
                                                                                                                                               return tmp.abs2()/v.abs2();
                                                                         T area()const{
                                                                                                                                        337
                                                                           T t=(b-a).cross(c-a)/2;
         if(!i)l=r;
                                                                           return t>0?t:-t;
```

```
pair<point3D<T>,point3D<T> > closest pair(const line3D<T> & 393
                                                                         return (a+b+c+d)/4;
         1)const{
339
      point3D<T> v1=(p1-p2), v2=(1.p1-1.p2);
                                                                  395
                                                                       bool point in(const point3D<T> &p)const{
                                                                         return triangle3D<T>(a,b,c).point_in(p)&&triangle3D<T>(c,
340
      point3D<T> N=v1.cross(v2),ab(p1-l.p1);
                                                                              d,a).point_in(p);
      //if(N.abs2()==0)return NULL;平行或重合
342
      T tmp=N.dot(ab),ans=tmp*tmp/N.abs2();//最近點對距離
      point3D<T> d1=p2-p1,d2=l.p2-l.p1,D=d1.cross(d2),G=l.p1-p1 398
343
                                                                     template<typename T>
                                                                     struct convexhull3D{
      T t1=(G.cross(d2)).dot(D)/D.abs2();
344
                                                                       static const int MAXN=1005;
      T t2=(G.cross(d1)).dot(D)/D.abs2();
345
                                                                       struct face{
      return make_pair(p1+d1*t1,l.p1+d2*t2);
346
                                                                         int a,b,c;
347
                                                                         face(int a,int b,int c):a(a),b(b),c(c){}
348
    bool same_side(const point3D<T> &a,const point3D<T> &b)
                                                                       };
         const{
                                                                       vector<point3D<T>> pt;
       return (p2-p1).cross(a-p1).dot((p2-p1).cross(b-p1))>0;
349
                                                                       vector<face> ans;
350
                                                                       int fid[MAXN][MAXN];
351 };
                                                                       void build(){
352 template<typename T>
                                                                         int n=pt.size();
353 struct plane{
                                                                         ans.clear();
    point3D<T> p0,n;//平面上的點和法向量
                                                                         memset(fid,0,sizeof(fid));
    plane(){}
                                                                         ans.emplace back(0,1,2);//注意不能共線
    plane(const point3D<T> &p0,const point3D<T> &n):p0(p0),n(n)
                                                                         ans.emplace back(2,1,0);
                                                                  415
                                                                         int ftop = 0:
    T dis2(const point3D<T> &p)const{//點到平面距離的平方
                                                                         for(int i=3, ftop=1; i<n; ++i,++ftop){</pre>
                                                                  416
      T tmp=(p-p0).dot(n);
358
                                                                  417
                                                                           vector<face> next;
359
      return tmp*tmp/n.abs2();
                                                                  418
                                                                           for(auto &f:ans){
360
                                                                  419
                                                                             T d=(pt[i]-pt[f.a]).dot((pt[f.b]-pt[f.a]).cross(pt[f.
    point3D<T> projection(const point3D<T> &p)const{
361
                                                                                  c]-pt[f.a]));
362
      return p-n*(p-p0).dot(n)/n.abs2();
                                                                             if(d<=0) next.push_back(f);</pre>
                                                                  420
363
                                                                  421
                                                                             int ff=0;
    point3D<T> line intersection(const line3D<T> &1)const{
                                                                  422
                                                                             if(d>0) ff=ftop;
      T tmp=n.dot(1.p2-1.p1);//等於 Ø表示平行或重合該平面
                                                                  423
                                                                             else if(d<0) ff=-ftop;</pre>
      return 1.p1+(1.p2-1.p1)*(n.dot(p0-1.p1)/tmp);
366
                                                                             fid[f.a][f.b]=fid[f.b][f.c]=fid[f.c][f.a]=ff;
                                                                  424
367
                                                                  425
368
    line3D<T> plane_intersection(const plane &pl)const{
                                                                           for(auto &f:ans){
                                                                  426
369
      point3D<T> e=n.cross(pl.n),v=n.cross(e);
                                                                             if(fid[f.a][f.b]>0 && fid[f.a][f.b]!=fid[f.b][f.a])
                                                                  427
370
      T tmp=pl.n.dot(v);//等於 Ø表示平行或重合該平面
                                                                               next.emplace_back(f.a,f.b,i);
                                                                  428
      point3D<T> q=p0+(v*(pl.n.dot(pl.p0-p0))/tmp);
371
                                                                  429
                                                                             if(fid[f.b][f.c]>0 && fid[f.b][f.c]!=fid[f.c][f.b])
      return line3D<T>(q,q+e);
372
                                                                  430
                                                                               next.emplace back(f.b,f.c,i);
373
                                                                             if(fid[f.c][f.a]>0 && fid[f.c][f.a]!=fid[f.a][f.c])
                                                                  431
374
                                                                  432
                                                                               next.emplace back(f.c,f.a,i);
   template<typename T>
                                                                  433
   struct triangle3D{
                                                                  434
                                                                           ans=next;
    point3D<T> a,b,c;
                                                                  435
    triangle3D(){}
    triangle3D(const point3D<T> &a,const point3D<T> &b,const
                                                                       point3D<T> centroid()const{
          point3D<T> &c):a(a),b(b),c(c){}
                                                                         point3D<T> res(0,0,0);
    bool point in(const point3D<T> &p)const{//點在該平面上的投
                                                                         T vol=0;
380
                                                                  439
          影在三角形中
                                                                         for(auto &f:ans){
                                                                           T tmp=pt[f.a].dot(pt[f.b].cross(pt[f.c]));
      return line3D<T>(b,c).same_side(p,a)&&line3D<T>(a,c).
                                                                  441
                                                                           res=res+(pt[f.a]+pt[f.b]+pt[f.c])*tmp;
            same_side(p,b)&&line3D<T>(a,b).same_side(p,c);
                                                                  442
                                                                  443
                                                                           vol+=tmp;
382
383 };
                                                                  444
384 template<typename T>
                                                                  445
                                                                         return res/(vol*4);
                                                                  446
385 struct tetrahedron{//四面體
    point3D<T> a,b,c,d;
387
    tetrahedron(){}
    tetrahedron(const point3D<T> &a,const point3D<T> &b,const
          point3D<T> &c,const point3D<T> &d):a(a),b(b),c(c),d(d)
    T volume6()const{//體積的六倍
      return (d-a).dot((b-a).cross(c-a));
390
    point3D<T> centroid()const{
```

Graph

33

35

36

37

56

6.1 Augment Path

```
1 struct AugmentPath{
     int n, m;
     vector<vector<int>> G;
     vector<int> mx, my;
     vector<int> visx, visy;
     int stamp;
     AugmentPath(int _n, int _m) : n(_n), m(_m), G(n), mx(n,
           -1), my(m, -1), visx(n), visy(n){
          stamp = 0;
     void add(int x, int y){
         G[x].push_back(y);
     // bb03e2
     bool dfs1(int now){
          visx[now] = stamp;
          for (auto x : G[now]){
              if (my[x] == -1){
                  mx[now] = x;
                  my[x] = now;
                  return true;
          for (auto x : G[now]){
              if (visx[my[x]]!=stamp && dfs1(my[x])){
                  mx[now] = x;
                  my[x] = now;
                  return true;
          return false;
     vector<pair<int, int>> find max matching(){
          vector<pair<int, int>> ret;
          while (true){
              stamp++;
              int tmp = 0;
              for (int i=0 ; i<n ; i++){</pre>
                  if (mx[i]==-1 && dfs1(i)) tmp++;
              if (tmp==0) break;
          for (int i=0 ; i<n ; i++){</pre>
              if (mx[i]!=-1){
                  ret.push_back({i, mx[i]});
          return ret;
     // 645577
     void dfs2(int now){
```

5.3 Pick's Theorem

給定頂點坐標均是整點的簡單多邊形,面積 = 內部格點數 + 邊上格點數/2-1

```
visx[now] = true;
          for (auto x : G[now]){
              if (my[x]!=-1 && visy[x]==false){
                  visy[x] = true;
                  dfs2(my[x]);
              }
      }
      // 要先執行 find_max_matching 一次
      vector<pair<int, int>> find_min_vertex_cover(){
          fill(visx.begin(), visx.end(), false);
          fill(visy.begin(), visy.end(), false);
          vector<pair<int, int>> ret;
          for (int i=0 ; i<n ; i++){</pre>
              if (mx[i]==-1) dfs2(i);
          for (int i=0 ; i<n ; i++){</pre>
              if (visx[i]==false) ret.push_back({1, i});
          for (int i=0 ; i<m ; i++){</pre>
              if (visy[i]==true) ret.push_back({2, i});
          return ret;
88 };
```

6.2 Bridge BCC

```
| #include <bits/stdc++.h>
 using namespace std;
 const int N = 200005;
 vector <int> G[N];
 int low[N], depth[N];
 bool vis[N];
 vector <vector <int>> bcc;
 stack <int> stk;
 void dfs(int v, int p) {
     stk.push(v);
     vis[v] = true;
     low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
     for (int u : G[v]) {
        if (u == p) continue;
        if (!vis[u]) {
            /// (v, u) 是樹邊
            dfs(u, v);
            low[v] = min(low[v], low[u]);
            /// (v, u) 是回邊
            low[v] = min(low[v], depth[u]);
     /// v 在不依靠父邊的情況下永遠沒辦法走到它的祖先
     if (low[v] == depth[v]) {
         bcc.emplace back();
         while (stk.top() != v) {
```

6.3 Cut BCC

```
#include <bits/stdc++.h>
using namespace std;
const int N = 200005;
vector <int> G[N];
int low[N], depth[N];
bool vis[N];
vector <vector <int>> bcc;
stack <int> stk;
void dfs(int v, int p) {
    stk.push(v);
    vis[v] = true;
    low[v] = depth[v] = (p == -1 ? 1 : depth[p] + 1);
    for (int u : G[v]) {
       if (u == p) continue;
       if (!vis[u]) {
           /// (v, u) 是樹邊
            dfs(u, v);
            low[v] = min(low[v], low[u]);
            /// u 無法在不經過父邊的情況走到 v 的祖先
            if (low[u] >= depth[v]) {
               bcc.emplace_back();
               while (stk.top() != u) {
                   bcc.back().push_back(stk.top());
                   stk.pop();
               bcc.back().push back(stk.top());
               stk.pop();
               bcc.back().push_back(v);
       } else {
            /// (v, u) 是回邊
            low[v] = min(low[v], depth[u]);
```

6.4 Dijkstra

```
    1 // 可以在 O(E Log E) 的時間複雜度解決在無負權有向圖單點源最短路

    2 const int INF = 2e18; // 要確保 INF 開的足夠大

    3 vector<vector<pair<int, int>>> G(n); // G[i] = <節點, 權重> vector<int> dis(n, INF); priority_queue<pair<int, int>, vector<pair<int, int>>, greater<pair<int, int>>> pq; dis[s] = 0;
```

6.5 Dinic

```
1 // 一般圖: O(EV2)
2 // 二分圖: O(E√V)
  struct Flow{
       struct Edge{
           int v, rc, rid;
       vector<vector<Edge>> G;
       void add(int u, int v, int c){
           G[u].push_back({v, c, G[v].size()});
           G[v].push_back({u, 0, G[u].size()-1});
       vector<int> dis, it;
       Flow(int n){
           G.resize(n);
           dis.resize(n);
           it.resize(n);
19
       int dfs(int u, int t, int f){
    if (u==t || f==0) return f;
           for (int &i=it[u] ; i<G[u].size() ; i++){</pre>
               auto &[v, rc, rid] = G[u][i];
               if (dis[v]!=dis[u]+1) continue;
               int df = dfs(v, t, min(f, rc));
25
26
               if (df<=0) continue;</pre>
27
               rc -= df;
28
               G[v][rid].rc += df;
               return df;
30
31
           return 0;
32
33
       int flow(int s, int t){
           int ans = 0;
           while (true){
               fill(dis.begin(), dis.end(), INF);
37
               queue<int> q;
38
               q.push(s);
               dis[s] = 0;
40
42
               while (q.size()){
                    int u = q.front(); q.pop();
43
                    for (auto [v, rc, rid] : G[u]){
44
```

```
if (rc<=0 || dis[v]<INF) continue;</pre>
                      dis[v] = dis[u]+1;
                      q.push(v);
              if (dis[t]==INF) break;
              fill(it.begin(), it.end(), 0);
              while (true){
                  int df = dfs(s, t, INF);
                  if (df<=0) break;</pre>
                  ans += df;
          return ans;
      // the code below constructs minimum cut
      void dfs_mincut(int now, vector<bool> &vis){
      vis[now] = true;
      for (auto &[v, rc, rid] : G[now]){
        if (vis[v]==false && rc>0){
          dfs mincut(v, vis);
      }
    }
    vector<pair<int, int>> construct(int n, int s, vector<pair< 47</pre>
         int.int>> &E){
        // E is G without capacity
        vector<bool> vis(n);
      dfs_mincut(s, vis);
      vector<pair<int, int>> ret;
      for (auto &[u, v] : E){
        if (vis[u]==true && vis[v]==false){
                  ret.emplace back(u, v);
      return ret;
83 };
        Dinic with double
```

```
const double double_INF = 1e18;
const int INF = (int)(1e9 + 10);

struct Flow{
    const double eps = 1e-9;
    struct Edge{
        int v; double rc; int rid;
    };
    vector<vector<Edge>> G;
    void add(int u, int v, double c){
        G[u].push_back({v, c, G[v].size()});
        G[v].push_back({u, 0, G[u].size()-1});
    }
    vector<int> dis, it;

Flow(int n){
        G.resize(n);
        dis.resize(n);
        it.resize(n);
}
```

```
double dfs(int u, int t, double f){
    if (u == t || abs(f) < eps) return f;</pre>
    for (int &i=it[u]; i<G[u].size(); i++){</pre>
         auto &[v, rc, rid] = G[u][i];
         if (dis[v]!=dis[u]+1) continue;
         double df = dfs(v, t, min(f, rc));
        if (abs(df) <= eps) continue;</pre>
         G[v][rid].rc += df;
         return df:
    return 0;
double flow(int s, int t){
    double ans = 0:
    while (true){
         fill(dis.begin(), dis.end(), INF);
         queue<int> q;
         q.push(s);
         dis[s] = 0;
         while (q.size()){
             int u = q.front(); q.pop();
             for (auto [v, rc, rid] : G[u]){
   if (abs(rc) <= eps || dis[v] < INF)</pre>
                       continue;
                 dis[v] = dis[u] + 1;
                 q.push(v);
         if (dis[t]==INF) break;
         fill(it.begin(), it.end(), 0);
         while (true){
             double df = dfs(s, t, double_INF);
             if (abs(df) <= eps) break;</pre>
             ans += df;
    return ans;
// the code below constructs minimum cut
void dfs mincut(int now, vector<bool> &vis){
    vis[now] = true;
    for (auto &[v, rc, rid] : G[now]){
   if (vis[v] == false && rc > eps){
             dfs_mincut(v, vis);
vector<pair<int, int>> construct(int n, int s, vector
     pair<int,int>> &E){
    // E is G without capacity
    vector<bool> vis(n);
    dfs mincut(s, vis);
    vector<pair<int, int>> ret:
    for (auto &[u, v] : E){
   if (vis[u] == true && vis[v] == false){
             ret.emplace back(u, v);
    return ret;
```

```
85 | 3
86 };
```

6.7 Dominator Tree

```
1 /*
2|全部都是 0-based
3 一開始要初始化 G(N. root),代表有 N 個節點,根是 root
4 用完之後要 build
S[G[i] = i 的 idom · 也就是從 root 走到 i 時 · 一定要走到的點且離
       最新
  struct DominatorTree{
      int N;
      vector<vector<int>> G;
      vector<vector<int>> buckets, rg;
      // dfn[x] = the DFS otder of x
      // rev[x] = the vertex with DFS order x
      // par[x] = the parent of x
      vector<int> dfn, rev, par;
      vector<int> sdom, dom, idom;
      vector<int> fa, val;
      int stamp:
      int root;
      int operator [] (int x){
          return idom[x];
      DominatorTree(int N, int root) :
          G(N), buckets(N), rg(N),
         dfn(N, -1), rev(N, -1), par(N, -1),
          sdom(N, -1), dom(N, -1), idom(N, -1),
          fa(N, -1), val(N, -1)
         stamp = 0;
32
          root = _root;
      void add edge(int u, int v){
         G[u].push back(v);
      void dfs(int x){
          rev[dfn[x] = stamp] = x;
          fa[stamp] = sdom[stamp] = val[stamp] = stamp;
          for (int u : G[x]){
              if (dfn[u]==-1){
                 dfs(u);
                 par[dfn[u]] = dfn[x];
              rg[dfn[u]].push back(dfn[x]);
52
      int eval(int x, bool first){
53
          if (fa[x]==x) return !first ? -1 : x;
          int p = eval(fa[x], false);
```

```
if (p==-1) return x;
          if (sdom[val[x]]>sdom[val[fa[x]]]) val[x] = val[fa[x
          fa[x] = p;
          return !first ? p : val[x];
      void link(int x, int y){
          fa[x] = y;
      void build(){
          dfs(root);
          for (int x=stamp-1 ; x>=0 ; x--){
              for (int y : rg[x]){
                  sdom[x] = min(sdom[x], sdom[eval(y, true)]);
              if (x>0) buckets[sdom[x]].push_back(x);
              for (int u : buckets[x]){
                  int p = eval(u, true);
                  if (sdom[p]==x) dom[u] = x;
                  else dom[u] = p;
              if (x>0) link(x, par[x]);
          idom[root] = root;
          for (int x=1 ; x<stamp ; x++){</pre>
              if (sdom[x]!=dom[x]) dom[x] = dom[dom[x]];
          for (int i=1; i<stamp; i++) idom[rev[i]] = rev[dom[</pre>
               i]];
90 };
```

6.8 Find Bridge

```
| vector<int> dep(MAX_N), low(MAX_N);
 vector<pair<int, int>> bridge;
 bitset<MAX_N> vis;
 void dfs(int now, int pre){
     vis[now] = 1;
     low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);
     for (auto x : G[now]){
         if (x==pre){
             continue;
         }else if (vis[x]==0){
             // 沒有走過的節點
             dfs(x, now);
             low[now] = min(low[now], low[x]);
         }else if (vis[x]==1){
             low[now] = min(low[now], dep[x]);
     if (now!=1 && low[now]==dep[now]){
         bridge.push_back({now, pre});
     return;
```

6.9 HLD

1 #include <bits/stdc++.h>

#define int long long

```
using namespace std;
const int N = 100005;
vector <int> G[N];
struct HLD {
    vector<int> pa, sz, depth, mxson, topf, id;
    int n, idcnt = 0;
    HLD(int _n) : n(_n), pa(_n + 1), sz(_n + 1), depth(_n +
        1), mxson(n + 1), topf(n + 1), id(n + 1) {}
    void dfs1(int v = 1, int p = -1) {
        pa[v] = p; sz[v] = 1; mxson[v] = 0;
        depth[v] = (p == -1 ? 0 : depth[p] + 1);
        for (int u : G[v]) {
            if (u == p) continue;
            dfs1(u, v);
            sz[v] += sz[u];
            if (sz[u] > sz[mxson[v]]) mxson[v] = u;
    void dfs2(int v = 1, int top = 1) {
        id[v] = ++idcnt;
        topf[v] = top;
        if (mxson[v]) dfs2(mxson[v], top);
        for (int u : G[v]) {
            if (u == mxson[v] || u == pa[v]) continue;
            dfs2(u, u);
    // query 為區間資料結構
    int path_query(int a, int b) {
        int res = 0;
        while (topf[a] != topf[b]) { /// 若不在同一條鍊上
            if (depth[topf[a]] < depth[topf[b]]) swap(a, b);</pre>
            res = max(res, 011); // query : L = id[topf[a]],
                r = id[a]
            a = pa[topf[a]];
        /// 此時已在同一條鍊上
        if (depth[a] < depth[b]) swap(a, b);</pre>
        res = \max(\text{res}, 011); // query : l = id[b], r = id[a]
        return res;
```

6.10 Kosaraju

```
7|.add edge(u, v): 加上一條邊 u -> v
8|.compress: O(n log n) 計算 G3 \ SCC \ SCC id 的資訊,並把縮點後
       的結果存在 result 裡
10 | SCC[i] = 某個 SCC 中的所有點
11 | SCC id[i] = 第 i 個點在第幾個 SCC
13 struct SCC_compress{
      int N, M, sz;
      vector<vector<int>>> G, inv G, result;
      vector<pair<int, int>> edges;
      vector<bool> vis;
      vector<int> order;
      vector<vector<int>> SCC;
      vector<int> SCC id;
      SCC compress(int N):
      N(N), M(0), sz(0),
      G(N), inv G(N),
      vis(N), SCC id(N)
      vector<int> operator [] (int x){
          return result[x];
      void add edge(int u, int v){
          G[u].push back(v);
          inv_G[v].push_back(u);
          edges.push back({u, v});
38
39
      void dfs1(vector<vector<int>> &G, int now){
          vis[now] = 1;
          for (auto x : G[now]) if (!vis[x]) dfs1(G, x);
          order.push_back(now);
43
44
      void dfs2(vector<vector<int>> &G, int now){
          SCC id[now] = SCC.size()-1;
          SCC.back().push back(now);
          vis[now] = 1;
          for (auto x : G[now]) if (!vis[x]) dfs2(G, x);
50
51
      void compress(){
          fill(vis.begin(), vis.end(), 0);
          for (int i=0; i<N; i++) if (!vis[i]) dfs1(G, i);</pre>
          fill(vis.begin(), vis.end(), 0);
          reverse(order.begin(), order.end());
          for (int i=0 ; i<N ; i++){</pre>
              if (!vis[order[i]]){
                  SCC.push back(vector<int>());
                  dfs2(inv_G, order[i]);
          result.resize(SCC.size());
          sz = SCC.size();
          for (auto [u, v] : edges){
              if (SCC_id[u]!=SCC_id[v]) result[SCC_id[u]].
                   push_back(SCC_id[v]);
```

```
if (visy[j]!=stamp) delta = min(delta, slack[j]); 19
                                                                                                                                            vector<int> dis(SZ(par), INF);
          for (int i=0 ; i<SCC.size() ; i++){</pre>
                                                                                                                                            vector<bool> in_q(SZ(par), false);
72
              sort(result[i].begin(), result[i].end());
                                                                             for (int i=0 ; i<n ; i++){</pre>
                                                                                                                                            aueue<int> 0:
              result[i].resize(unique(result[i].begin(), result 55
                                                                                 if (visx[i]==stamp) lx[i] -= delta;
                                                                                                                                            dis[s] = 0;
                                                                                                                                            in_q[s] = true;
                   [i].end())-result[i].begin());
                                                                             for (int j=0 ; j<n ; j++){</pre>
                                                                                                                                            Q.push(s);
                                                                                 if (visy[j]==stamp) ly[j] += delta;
                                                                                 else slack[j] -= delta;
                                                                                                                                            while (!Q.empty()){
76 };
                                                                                                                                             int v = Q.front();
                                                                         }
                                                                                                                                              Q.pop();
                                                                                                                                              in q[v] = false;
  6.11 Kuhn Munkres
                                                                         int solve(){
                                                                                                                                              for (int i=0 ; i<SZ(G[v]) ; i++){</pre>
                                                                             for (int i=0 ; i<n ; i++){</pre>
                                                                                                                                                auto [u, rc, k, rv] = G[v][i];
i struct KuhnMunkres{
                                                                                                                                                if (rc>0 && dis[v]+k<dis[u]){</pre>
                                                                                 lx[i] = 0;
      int n; // max(n, m)
                                                                                 for (int j=0 ; j<n ; j++){</pre>
                                                                                                                                                  dis[u] = dis[v]+k;
      vector<vector<int>> G;
                                                                                     lx[i] = max(lx[i], G[i][j]);
                                                                                                                                                  par[u] = v;
      vector<int> match, lx, ly, visx, visy;
                                                                                                                                                  par_eid[u] = i;
      vector<int> slack;
                                                                                                                                                  if (!in_q[u]) Q.push(u);
      int stamp = 0;
                                                                                                                                                  in_q[u] = true;
                                                                             fill(ly.begin(), ly.end(), 0);
      KuhnMunkres(int n): n(n), G(n, vector(int)(n)), lx(n),
                                                                             fill(match.begin(), match.end(), -1);
           ly(n), slack(n), match(n), visx(n), visy(n) {}
                                                                                                                                            }
                                                                             for(int i = 0; i < n; i++) {</pre>
      void add(int x, int y, int w){
                                                                                 fill(slack.begin(), slack.end(), INF);
                                                                                                                                            return dis[t];
          G[x][y] = max(G[x][y], w);
                                                                                 stamp++;
                                                                                 if(dfs(i, true)) continue;
                                                                                                                                          // return <max flow, min cost>, 150093
                                                                                                                                          pair<int, int> flow(int s, int t){
      bool dfs(int i, bool aug){ // aug = true 表示要更新 match
                                                                                 while(augment()==false) relabel();
                                                                                                                                            int fl = 0, cost = 0, d;
          if (visx[i]==stamp) return false;
                                                                                 stamp++;
                                                                                 dfs(i, true);
                                                                                                                                            while ((d = spfa(s, t))<INF){</pre>
          visx[i] = stamp;
                                                                                                                                              int cur = INF;
                                                                                                                                              for (int v=t ; v!=s ; v=par[v])
          for (int j=0 ; j<n ; j++){</pre>
                                                                                                                                               cur = min(cur, G[par[v]][par_eid[v]].rc);
              if (visy[j]==stamp) continue;
                                                                             int ans = 0;
                                                                             for (int j=0 ; j<n ; j++){</pre>
                                                                                                                                              fl += cur;
              int d = lx[i]+ly[j]-G[i][j];
                                                                                 if (match[j]!=-1){
                                                                                                                                              cost += d*cur;
                                                                                                                                              for (int v=t ; v!=s ; v=par[v]){
                                                                                     ans += G[match[j]][j];
              if (d==0){
                                                                                                                                                G[par[v]][par_eid[v]].rc -= cur;
                  visy[j] = stamp;
                                                                                                                                                G[v][G[par[v]][par_eid[v]].rv].rc += cur;
                  if (match[j]==-1 || dfs(match[j], aug)){
                      if (aug){
                                                                             return ans;
                           match[j] = i;
                                                                  93 };
                                                                                                                                            return {fl, cost};
                      return true;
                                                                                                                                          vector<pair<int, int>> construct(){
              }else{
                                                                     6.12 MCMF
                  slack[j] = min(slack[j], d);
                                                                                                                                            vector<pair<int, int>> ret;
                                                                                                                                            for (int i=0 ; i<n ; i++){</pre>
                                                                                                                                              for (auto x : G[i]){
                                                                                                                                                if (x.rc==0){
          return false;
                                                                   1 struct Flow {
                                                                                                                                                  ret.push_back({i+1, x.u-n+1});
                                                                       struct Edge {
                                                                         int u, rc, k, rv;
                                                                                                                                                  break:
      bool augment(){
          for (int j=0 ; j<n ; j++){</pre>
              if (visy[j]!=stamp && slack[j]==0){
                                                                       vector<vector<Edge>> G;
                  visy[j] = stamp;
                                                                       vector<int> par, par_eid;
                                                                                                                                     73
                                                                                                                                            return ret;
                  if (match[j]==-1 || dfs(match[j], false)){
                                                                       Flow(int n): G(n+1), par(n+1), par eid(n+1) {}
                                                                                                                                     74
                      return true;
                                                                       // v->u, capcity: c, cost: k
                                                                       void add(int v, int u, int c, int k){
                                                                         G[v].push_back({u, c, k, SZ(G[u])});
                                                                                                                                       6.13 Tarjan
                                                                         G[u].push_back({v, 0, -k, SZ(G[v])-1});
          return false;
                                                                       // 3701d6
      void relabel(){
                                                                                                                                     i struct tarjan_SCC {
          int delta = INF;
                                                                      int spfa(int s, int t){
                                                                                                                                            int now T, now SCCs;
          for (int j=0 ; j<n ; j++){</pre>
                                                                         fill(ALL(par), -1);
                                                                                                                                            vector<int> dfn, low, SCC;
```

```
stack<int> S;
vector<vector<int>> E;
vector<bool> vis, in stack;
tarjan_SCC(int n) {
    init(n);
void init(int n) {
    now T = now SCCs = 0;
    dfn = low = SCC = vector<int>(n);
    E = vector<vector<int>>(n);
    S = stack<int>();
    vis = in_stack = vector<bool>(n);
void add(int u, int v) {
    E[u].push back(v);
void build() {
    for (int i = 0; i < dfn.size(); ++i) {</pre>
        if (!dfn[i]) dfs(i);
void dfs(int v) {
    now T++;
    vis[v] = in stack[v] = true;
    dfn[v] = low[v] = now_T;
    S.push(v);
    for (auto &i:E[v]) {
        if (!vis[i]) {
            vis[i] = true;
            dfs(i);
            low[v] = min(low[v], low[i]);
        else if (in_stack[i]) {
            low[v] = min(low[v], dfn[i]);
    if (low[v] == dfn[v]) {
        int tmp;
        do {
            tmp = S.top();
            S.pop();
            SCC[tmp] = now_SCCs;
            in stack[tmp] = false;
        } while (tmp != v);
        now SCCs += 1;
```

6.14 Tarjan Find AP

```
vector<int> dep(MAX_N), low(MAX_N), AP;
bitset<MAX_N> vis;

void dfs(int now, int pre){
   int cnt = 0;
   bool ap = 0;
   vis[now] = 1;
   low[now] = dep[now] = (now==1 ? 0 : dep[pre]+1);

for (auto x : G[now]){
   if (x==pre){
```

6.15 Tree Isomorphism

```
| #include <bits/stdc++.h>
  #pragma GCC optimize("03,unroll-loops")
  #define fastio ios::sync with stdio(0), cin.tie(0), cout.tie
  #define dbg(x) cerr << #x << " = " << x << endl
  #define int long long
  using namespace std;
  // declare
  const int MAX SIZE = 2e5+5;
  const int INF = 9e18;
  const int MOD = 1e9+7;
  const double EPS = 1e-6;
  typedef vector<vector<int>> Graph;
  typedef map<vector<int>, int> Hash;
  int n, a, b;
  int id1, id2;
  pair<int, int> c1, c2;
  vector<int> sz1(MAX_SIZE), sz2(MAX_SIZE);
  vector<int> we1(MAX_SIZE), we2(MAX_SIZE);
  Graph g1(MAX SIZE), g2(MAX SIZE);
  Hash m1, m2;
  int testcase=0;
  void centroid(Graph &g, vector<int> &s, vector<int> &w, pair< 90</pre>
       int, int> &rec, int now, int pre){
      s[now]=1;
      w[now]=0;
      for (auto x : g[now]){
          if (x!=pre){
               centroid(g, s, w, rec, x, now);
               s[now]+=s[x];
               w[now]=max(w[now], s[x]);
      w[now]=max(w[now], n-s[now]);
      if (w[now]<=n/2){</pre>
                                                                  103
          if (rec.first==0) rec.first=now;
          else rec.second=now;
                                                                  106
                                                                  107
43 int dfs(Graph &g, Hash &m, int &id, int now, int pre){
```

```
vector<int> v;
    for (auto x : g[now]){
        if (x!=pre){
            int add=dfs(g, m, id, x, now);
            v.push_back(add);
   sort(v.begin(), v.end());
   if (m.find(v)!=m.end()){
        return m[v];
   }else{
        m[v]=++id;
        return id;
void solve1(){
    // init
   id1=0:
   id2=0;
   c1={0, 0};
   c2={0, 0};
    fill(sz1.begin(), sz1.begin()+n+1, 0);
    fill(sz2.begin(), sz2.begin()+n+1, 0);
    fill(we1.begin(), we1.begin()+n+1, 0);
    fill(we2.begin(), we2.begin()+n+1, 0);
    for (int i=1; i<=n; i++){</pre>
       g1[i].clear();
       g2[i].clear();
   m1.clear();
   m2.clear();
   // input
    cin >> n;
    for (int i=0 ; i<n-1 ; i++){</pre>
       cin >> a >> b;
        g1[a].push back(b);
       g1[b].push_back(a);
   for (int i=0 ; i<n-1 ; i++){</pre>
       cin >> a >> b;
        g2[a].push back(b);
       g2[b].push_back(a);
   // get tree centroid
    centroid(g1, sz1, we1, c1, 1, 0);
    centroid(g2, sz2, we2, c2, 1, 0);
   // process
   int res1=0, res2=0, res3=0;
    if (c2.second!=0){
        res1=dfs(g1, m1, id1, c1.first, 0);
       m2=m1;
        res2=dfs(g2, m1, id1, c2.first, 0);
        res3=dfs(g2, m2, id2, c2.second, 0);
    }else if (c1.second!=0){
       res1=dfs(g2, m1, id1, c2.first, 0);
        m2=m1;
        res2=dfs(g1, m1, id1, c1.first, 0);
```

```
res3=dfs(g1, m2, id2, c1.second, 0);
                                                                     36 vector<edg> EV;
                                                                                                                                                        jmp[i][j] = jmp[jmp[i][j-1]][j-1];
111
       }else{
                                                                                                                                         102
112
           res1=dfs(g1, m1, id1, c1.first, 0);
                                                                        void tarian(int v, int par, stack<int>& S){
                                                                                                                                         103
           res2=dfs(g2, m1, id1, c2.first, 0);
                                                                            static vector<int> dfn(mxn), low(mxn);
113
                                                                                                                                         104
114
       }
                                                                            static vector<bool> to add(mxn);
                                                                                                                                         105
                                                                            static int nowT = 0:
                                                                                                                                            inline int lca(int x, int y){
115
116
       // output
                                                                                                                                                if(dep[x] < dep[y]){ swap(x, y); }</pre>
       cout << (res1==res2 || res1==res3 ? "YES" : "NO") << endl 43
                                                                            int childs = 0:
                                                                            nowT += 1:
                                                                                                                                                int diff = dep[x] - dep[y];
                                                                                                                                         109
118
                                                                            dfn[v] = low[v] = nowT;
                                                                                                                                         110
                                                                                                                                                lp(j,0,mxlg){
119
       return;
                                                                            for(auto &ne:E[v]){
                                                                                                                                         111
                                                                                                                                                    if((diff >> j) & 1){
120 }
                                                                                int i = EV[ne].to;
                                                                                                                                         112
                                                                                                                                                        x = jmp[x][j];
                                                                                if(i == par) continue;
121
                                                                                                                                         113
                                                                                if(!dfn[i]){
   signed main(void){
                                                                                                                                         114
       fastio:
                                                                                    S.push(ne);
                                                                                                                                                if(x == y) return x;
123
                                                                                                                                         115
124
                                                                                    tarjan(i, v, S);
                                                                                                                                         116
       int t=1:
                                                                                    childs += 1;
                                                                                                                                                for(int j = mxlg - 1; j >= 0; j--){
125
                                                                                                                                         117
       cin >> t;
126
                                                                                    low[v] = min(low[v], low[i]);
                                                                                                                                         118
                                                                                                                                                    if(jmp[x][j] != jmp[y][j]){
127
       while (t--){
                                                                                                                                         119
                                                                                                                                                        x = jmp[x][j];
128
           solve1();
                                                                                    if(par >= 0 && low[i] >= dfn[v]){
                                                                                                                                                        y = jmp[y][j];
                                                                                                                                         120
129
                                                                                        vector<int> bcc;
                                                                                                                                         121
                                                                                        int tmp;
130
       return 0;
                                                                                                                                         122
                                                                                        do{
131 }
                                                                                                                                         123
                                                                                                                                                return jmp[x][0];
                                                                                             tmp = S.top(); S.pop();
                                                                                                                                         124
                                                                                            if(!to_add[EV[tmp].fr]){
                                                                                                                                         125
                                                                                                 to_add[EV[tmp].fr] = true;
                                                                                                                                         126
                                                                                                                                            inline bool can_reach(int fr, int to){
           圓方樹
   6.16
                                                                                                 bcc.pb(EV[tmp].fr);
                                                                                                                                                if(dep[to] > dep[fr]) return false;
                                                                                            if(!to_add[EV[tmp].to]){
                                                                                                                                                int diff = dep[fr] - dep[to];
                                                                                                                                         129
                                                                                                 to add[EV[tmp].to] = true;
 i #include <bits/stdc++.h>
                                                                                                                                         130
                                                                                                                                                lp(j,0,mxlg){
                                                                                                 bcc.pb(EV[tmp].to);
                                                                                                                                                    if((diff >> j) & 1){
  #define lp(i,a,b) for(int i=(a);i<(b);i++)
                                                                                                                                         131
   #define pii pair<int,int>
                                                                                                                                                         fr = jmp[fr][j];
                                                                                                                                         132
                                                                                        }while(tmp != ne);
   #define pb push_back
                                                                                                                                         133
                                                                                        for(auto &j:bcc){
  #define ins insert
                                                                                                                                         134
  #define ff first
                                                                                            to add[j] = false;
                                                                                                                                         135
                                                                                                                                                return fr == to;
                                                                                            F[last special node].pb(j);
   #define ss second
                                                                                                                                         136
   #define opa(x) cerr << #x << " = " << x << ", ";
                                                                                            F[j].pb(last_special_node);
                                                                                                                                         137
  #define op(x) cerr << #x << " = " << x << endl;
                                                                                                                                            int main(){
                                                                                                                                         138
 10 #define ops(x) cerr << x;
                                                                                        last_special_node += 1;
                                                                                                                                                ios::sync_with_stdio(false); cin.tie(0);
                                                                                                                                                freopen("test_input.txt", "r", stdin);
 #define etr cerr << endl;</pre>
                                                                                                                                                int n, m, q; cin >> n >> m >> q;
12 #define spc cerr << ' ';
#define BAE(x) (x).begin(), (x).end()
                                                                                else{
                                                                                                                                                lp(i,0,m){
                                                                                    low[v] = min(low[v], dfn[i]);
 #define STL(x) cerr << #x << " : "; for(auto &qwe:x) cerr <<</pre>
                                                                                                                                                    int u, v; cin >> u >> v;
qwe << ''; cerr << endl;
15 #define deb1 cerr << "deb1" << endl;
                                                                                    if(dfn[i] < dfn[v]){ // edge i--v will be visited 144</pre>
                                                                                                                                                    E[u].pb(EV.size());
                                                                                          twice at here, but we only need one.
                                                                                                                                                    EV.pb(edg(u, v));
 16 #define deb2 cerr << "deb2" << endl:
                                                                                        S.push(ne);
                                                                                                                                                    E[v].pb(EV.size());
#define deb3 cerr << "deb3" << endl;
                                                                                                                                         147
                                                                                                                                                    EV.pb(edg(v, u));
 #define deb4 cerr << "deb4" << endl:
                                                                                                                                         148
19 #define deb5 cerr << "deb5" << endl;
                                                                                                                                                E[0].pb(EV.size());
                                                                                                                                         149
20 #define bye exit(0);
                                                                                                                                         150
                                                                                                                                                EV.pb(edg(0, 1));
  using namespace std;
                                                                                                                                                stack<int> S;
                                                                                                                                         151
                                                                       int dep[mxn], jmp[mxn][mxlg];
                                                                                                                                         152
                                                                                                                                                tarjan(0, -1, S);
   const int mxn = (int)(2e5) + 10:
                                                                        void dfs lca(int v, int par, int depth){
                                                                                                                                         153
                                                                                                                                                build lca();
                                                                            dep[v] = depth;
   const int mxlg = 17;
                                                                                                                                         154
  int last special node = (int)(1e5) + 1;
                                                                            for(auto &i:F[v]){
                                                                                                                                         155
                                                                                                                                                lp(queries,0,q){
26 vector<int> E[mxn], F[mxn];
                                                                                if(i == par) continue;
                                                                                                                                                    int fr, to, relay; cin >> fr >> to >> relay;
                                                                                                                                         156
                                                                                jmp[i][0] = v;
                                                                                                                                         157
                                                                                                                                                    if(fr == relay || to == relay){
28 struct edg{
                                                                                dfs_lca(i, v, depth + 1);
                                                                                                                                                        cout << "NO \ n";
                                                                                                                                                         continue;
       int fr, to;
                                                                                                                                         159
       edg(int _fr, int _to){
                                                                                                                                         160
                                                                                                                                                    if((can_reach(fr, relay) || can_reach(to, relay)) &&
           fr = _fr;
                                                                                                                                         161
                                                                       inline void build lca(){
                                                                                                                                                         dep[relay] >= dep[lca(fr, to)]){
           to = _to;
32
                                                                            jmp[1][0] = 1;
                                                                                                                                                         cout << "NO\n":
                                                                                                                                         162
33
                                                                            dfs_lca(1, -1, 1);
                                                                                                                                                         continue:
34 };
                                                                                                                                         163
35 ostream& operator<<(ostream& os, edg x){os << x.fr << "--" << 99
                                                                            lp(j,1,mxlg){
                                                                                                                                         164
                                                                                lp(i,1,mxn){
                                                                                                                                         165
                                                                                                                                                    cout << "YES\n";</pre>
```

```
166 | }
```

6.17 最大權閉合圖

```
2 Problem:
      Given w = [w_0, w_1, ..., w_{n-1}] (which can be
      either positive or negative or 0), you can choose
      to take w i (0 < i < n) or not, but if edge u -> v
      exists, you must take w_v if you want to take w_u
      (in other words, you can't take w u without taking
       w_v), this function returns the maximum value(> 0)
       you can get. If you need a construction, you can
       output the minimum cut of the S(source) side.
      MaxFlow(n, m) (Non-Biparte:O(n²m) / Bipartite:O(m√n))
int maximum_closure(vector<int> w, vector<pair<int,int>> EV)
      int n = w.size(), S = n + 1, T = n + 2;
      Flow G(T + 5); // Graph/Dinic.cpp
      int sum = 0;
      for (int i = 0; i < n; ++i) {</pre>
          if (w[i] > 0) {
              G.add(S, i, w[i]);
              sum += w[i];
          else if (w[i] < 0) {</pre>
              G.add(i, T, abs(w[i]));
      for (auto &[u, v] : EV) { // You should make sure that
           INF > \Sigma / w i /
          G.add(u, v, INF);
      int cut = G.flow(S, T);
      return sum - cut;
```

6.18 Theorem

- 任意圖
 - 不能有孤點.最大匹配 + 最小邊覆蓋 = n 點覆蓋的補集是獨立集。 最小點覆蓋 + 最大獨立集 = n
- 二分圖
 - 最小點覆蓋 = 最大匹配 = n 最大獨立集
- 只有邊帶權的二分圖
 - w-vertex-cover (帶權點覆蓋): 每條邊的兩個連接點被選中的次數總和至少要是 w_o 。
 - w-weight matching (帶權匹配)
 - minimum vertex count of w-vertex-cover = maximum weight count of 12 // 對於方程組的式子兩兩求解 w-weight matching (一個點可以被選很多次·但邊不行) 13 // {是否有解,{a, m}}
- 點、邊都帶權的二分圖的定理

- b-matching:假設 v 的點權是 b_v · 那所有 v 的匹配邊 e 的權重都要 15 滿足 $\sum w_e \leq b_v$ °
- The maximum w-weight of a b-matching equals the minimum b-weight of vertices in a w-vertex-cover.

7 Math

7.1 CRT m Coprime

```
vector<int> a, m;
int extgcd(int a, int b, int &x, int &y){
    if (b==0){
        x=1, y=0;
        return a;
    int ret=extgcd(b, a%b, y, x);
    v-=a/b*x:
    return ret;
// n = 有幾個式子,求解 x \equiv a i \bmod m i
int CRT(int n, vector<int> &a, vector<int> &m){
    int p=1, ans=0;
    vector<int> M(n), inv_M(n);
    for (int i=0 ; i<n ; i++) p*=m[i];</pre>
    for (int i=0 ; i<n ; i++){
        M[i]=p/m[i];
        extgcd(M[i], m[i], inv_M[i], tmp);
        ans+=a[i]*inv M[i]*M[i];
        ans%=p;
    return (ans%p+p)%p;
```

7.2 CRT m Not Coprime

```
int extgcd(int a, int b, int &x, int &y){
    if (b==0){
        x=1, y=0;
        return a;
    }

int ret=extgcd(b, a%b, y, x);
    y-=a/b*x;
    return ret;

y-a/b*x;
return ret;

// 對於方程組的式子兩兩求解
// {是否有解, {a, m}}

pair<bool>
pair<int, int>> CRT(int a1, int m1, int a2, int m2 44 45)
```

```
int g=__gcd(m1, m2);
if ((a2-a1)%g!=0) return {0, {-1, -1}};

int x, y;
extgcd(m1, m2, x, y);

x=(a2-a1)*x/g; // 兩者不能相反
a1=x*m1+a1;
m1=m1*m2/g;
a1=(a1%m1+m1)%m1;
return {1, {a1, m1}};
```

7.3 Fraction

21

```
1 #include <bits/stdc++.h>
  using namespace std;
  /// Fraction template starts ///
  #define fraction_template_bonus_check
  const long long long loverflow_warning_value = (long long)(3e9);
  long long gcd(long long a, long long b){
      if(a == 0) return 0:
      if(b == 0) return a;
      if(a < b) return gcd(b,a);</pre>
      return gcd(b, a%b);
  struct frac{
      long long a, b;
      frac(long long a = 0, long long b = 1){
          a = a; b = b;
          if(b == 0){
              cerr << "Error: division by zero\n";</pre>
              cerr << "Called : Constructor(" << _a << ", " <<</pre>
                   b << ")n";
              return;
          if(a == 0){b = 1; return;}
23
          if(b < 0){a = -a; b = -b;}
          long long gcd_ab = gcd(std::abs(a), b);
          if(gcd_ab != 1){a /= gcd_ab; b /= gcd_ab;}
          #ifdef fraction_template_bonus_check
          if(std::abs(a) > 11 overflow warning value || b >
               11_overflow_warning_value){
              cerr << "Overflow warning: " << a << "/" << b <<
          #endif // fraction_template_bonus_check
31
      frac operator+(frac const &B){
33
          return frac(a*(B.b)+(B.a)*b, b*(B.b));}
      frac operator-(frac const &B){
          return frac(a*(B.b)-(B.a)*b, b*(B.b));}
      frac operator*(frac const &B){
          return frac(a*(B.a), b*(B.b));}
      frac operator/(frac const &B){
          return frac(a*(B.b), b*(B.a));}
      frac operator+=(frac const &B){
          *this = frac(a*(B.b)+(B.a)*b, b*(B.b));}
      frac operator -= (frac const &B){
          *this = frac(a*(B.b)-(B.a)*b, b*(B.b));}
```

```
frac operator*=(frac const &B){
            *this = frac(a*(B.a), b*(B.b));}
       frac operator/=(frac const &B){
            *this = frac(a*(B.b), b*(B.a));}
       frac abs(){
            a = std::abs(a);
            return *this:
       bool operator<(frac const &B){</pre>
            return a*B.b < B.a*b;}</pre>
       bool operator <= (frac const &B){</pre>
            return a*B.b <= B.a*b;}</pre>
       bool operator>(frac const &B){
            return a*B.b > B.a*b;}
       bool operator>=(frac const &B){
            return a*B.b >= B.a*b;}
       bool operator == (frac const &B){
            return a * B.b == B.a * b;}
       bool operator!=(frac const &B){
            return a * B.b != B.a * b;}
69 ostream& operator << (ostream &os, const frac& A){
       os << A.a << "/" << A.b;
       return os;
   /// Fraction template ends ///
75 void test(frac A, frac B){
       cout << "A = " << A << endl;
       cout << "B = " << B << endl;
       cout << endl:
       cout << "A + B = " << A + B << endl;
       cout \langle\langle "A - B = " \langle\langle A - B \langle\langle endl;
       cout << "A * B = " << A * B << endl;
       cout \langle\langle "A / B = " \langle\langle A / B \langle\langle endl;
       cout << endl;</pre>
       cout \langle\langle "(A \langle B) = " \langle\langle (A \langle B) \langle\langle endl;
       cout \langle\langle "(A \langle = B) = " \langle\langle (A \langle = B) \rangle\langle\langle endl;
       cout \langle\langle "(A > B)' = " \langle\langle (A > B) \langle\langle endl;
       cout << "(A >= B) = " << (A >= B) << endl;
cout << "(A == B) = " << (A == B) << endl;
       cout << "(A != B) = " << (A != B) << end1;
       cout << "-----\n":
       return;
   int main(){
       frac tmp1(-7, 2);
       frac tmp2(5, 3);
       test(tmp1, tmp2);
       frac tmp3(-7);
       frac tmp4(0);
       test(tmp3, tmp4);
       return 0;
```

7.4 Josephus Problem

```
1 \mid //  有 n 個人·第偶數個報數的人被刪掉·問第 k 個被踢掉的是誰 2 \mid \text{int solve(int } n, \text{ int } k) {
```

```
if (n==1) return 1;
if (k<=(n+1)/2){
    if (2*k>n) return 2*k%n;
    else return 2*k;
}else{
    int res=solve(n/2, k-(n+1)/2);
    if (n&1) return 2*res+1;
    else return 2*res-1;
}
```

7.5 Lagrange any x

```
1 / / init: (x1, y1), (x2, y2) in a vector
2 struct Lagrange{
      int n;
      vector<pair<int, int>> v;
      Lagrange(vector<pair<int, int>> &_v){
         n = v.size();
         v = v;
      // O(n^2 \log MAX_A)
      int solve(int x){
          int ret = 0;
          for (int i=0 ; i<n ; i++){</pre>
              int now = v[i].second;
              for (int j=0 ; j<n ; j++){</pre>
                  if (i==j) continue;
                  now *= ((x-v[j].first)+MOD)%MOD;
                  now %= MOD;
                  now *= (qp((v[i].first-v[j].first+MOD)%MOD,
                       MOD - 2) + MOD) MOD;
                  now %= MOD;
              ret = (ret+now)%MOD;
          return ret;
```

7.6 Lagrange continuous x

```
#include <bits/stdc++.h>
using namespace std;

const int MAX_N = 5e5 + 10;
const int mod = 1e9 + 7;

long long inv_fac[MAX_N];

inline int fp(long long x, int y) {
    int ret = 1;
    for (; y; y >>= 1) {
        ret = (y & 1) ? (ret * x % mod) : ret;
        x = x * x % mod;
    }

return ret;
```

```
18 // TO USE THIS TEMPLATE. YOU MUST MAKE SURE THAT THE MOD
       NUMBER IS A PRIME.
19 struct Lagrange {
      Initialize a polynomial with f(x_0), f(x_0 + 1), ..., f(
      This determines a polynomial f(x) whose degree is at most
      Then you can call sample(x) and you get the value of f(x)
      Complexity of init() and sample() are both O(n).
25
      int m, shift; // m = n + 1
      vector<int> v, mul;
  // You can use this function if you don't have inv_fac array
      void construct inv fac() {
          long long fac = 1;
          for (int i = 2; i < MAX_N; ++i) {</pre>
              fac = fac * i % mod;
          inv fac[MAX N - 1] = fp(fac, mod - 2);
          for (int i = MAX N - 1; i >= 1; --i) {
              inv_fac[i - 1] = inv_fac[i] * i % mod;
38
  // You call init() many times without having a second
       instance of this struct.
      void init(int X_0, vector<int> &u) {
          shift = ((1 - X_0) \% mod + mod) \% mod;
          if (v.size() == 1) v.push_back(v[0]);
          m = v.size();
          mul.resize(m);
  // You can use sample(x) instead of sample(x % mod).
47
      int sample(int x) {
          x = ((long long)x + shift) % mod;
          x = (x < 0) ? (x + mod) : x;
          long long now = 1;
          for (int i = m; i >= 1; --i) {
              mul[i - 1] = now;
              now = now^* (x - i) \% mod;
          int ret = 0;
          bool neg = (m - 1) & 1;
          now = 1;
          for (int i = 1; i <= m; ++i) {</pre>
               int up = now * mul[i - 1] % mod;
               int down = inv fac[m - i] * inv_fac[i - 1] % mod;
               int tmp = ((long long)v[i - 1] * up % mod) * down
                    % mod:
               ret += (neg && tmp) ? (mod - tmp) : (tmp);
              ret = (ret >= mod) ? (ret - mod) : ret;
              now = now * (x - i) % mod;
              neg ^= 1;
          return ret:
  };
  int main() {
      int n: cin >> n:
      vector<int> v(n);
```

return ret;

Matrix pow(int p){

Matrix ret(n, n), mul = *this;
for (int i=0; i<n; i++){</pre>

if (p&1) ret = ret*mul;

ret.arr[i][i] = 1;

for (; p; p>>=1){

mul = mul*mul;

```
for (int i = 0; i < n; ++i) {</pre>
         cin >> v[i];
                                                                           return ret;
                                                                       }
     Lagrange L;
     L.construct_inv_fac();
     L.init(0, v);
                                                                       int det(){
     int x; cin >> x;
                                                                           vector<vector<int>> arr = this->arr;
     cout << L.sample(x);</pre>
                                                                           bool flag = false:
                                                                           for (int i=0 ; i<n ; i++){</pre>
                                                                               int target = -1;
                                                                               for (int j=i ; j<n ; j++){</pre>
                                                                                   if (arr[j][i]){
       Lucas's Theorem
                                                                                       target = j;
                                                                                       break;
1 \mid // 對於很大的 C^n_{m} 對質數 p 取模,只要 p 不大就可以用。
int Lucas(int n, int m, int p){
                                                                               if (target==-1) return 0;
     if (m==0) return 1;
                                                                               if (i!=target){
     return (C(n%p, m%p, p)*Lucas(n/p, m/p, p)%p);
                                                                                   swap(arr[i], arr[target]);
                                                                                   flag = !flag;
                                                                               for (int j=i+1 ; j<n ; j++){</pre>
 7.8 Matrix
                                                                                   if (!arr[j][i]) continue;
                                                                                   int freq = arr[j][i]*qp(arr[i][i], MOD-2)%MOD
| struct Matrix{
                                                                                   for (int k=i ; k<n ; k++){</pre>
                                                                                       arr[j][k] -= freq*arr[i][k];
     int n, m;
                                                                                       arr[j][k] = (arr[j][k]%MOD+MOD)%MOD;
     vector<vector<int>> arr;
     Matrix(int _n, int _m){
                                                                           }
         n = n;
         arr.resize(n, vector<int>(m));
                                                                           int ret = !flag ? 1 : MOD-1;
                                                                           for (int i=0; i<n; i++){</pre>
                                                                               ret *= arr[i][i];
     Matrix operator * (Matrix b){
                                                                               ret %= MOD;
         Matrix b_t(b.m, b.n);
         for (int i=0 ; i<b.n ; i++){</pre>
                                                                           return ret;
             for (int j=0 ; j<b.m ; j++){</pre>
                 b_t.arr[j][i] = b.arr[i][j];
                                                                   7.9 Matrix 01
         Matrix ret(n, b.m);
         for (int i=0 ; i<n ; i++){</pre>
             for (int j=0; j<b.m; j++){</pre>
                 for (int k=0; k<m; k++){</pre>
                                                                 ret.arr[i][j] += arr[i][k]*b_t.arr[j][k];
                                                                   struct Matrix{
                     ret.arr[i][j] %= MOD;
                                                                       int n, m;
                                                                       vector<bitset<MAX N>> arr;
             }
```

7.10 Miller Rabin

```
1 // O(\log n)
  typedef Uint unsigned long long
  Uint modmul(Uint a, Uint b, Uint m) {
      int ret = a*b - m*(Uint)((long double)a*b/m);
      return ret + m*(ret < 0) - m*(ret>=(int)m);
  int qp(int b, int p, int m){
      int ret = 1;
      for ( ; p ; p>>=1){
          if (p&1){
              ret = modmul(ret, b, m);
          b = modmul(b, b, m);
      return ret;
  vector<int> llsprp = {2, 325, 9375, 28178, 450775, 9780504,
       1795265022};
  | bool isprime(int n, vector<int> sprp = llsprp){
      if (n==2) return 1;
      if (n<2 || n%2==0) return 0;
      int t = 0;
      int u = n-1;
      for ( ; u%2==0 ; t++) u>>=1;
      for (int i=0; i<sprp.size(); i++){</pre>
          int a = sprp[i]%n;
          if (a==0 | | a==1 | | a==n-1) continue;
          int x = qp(a, u, n);
          if (x==1 || x==n-1) continue;
          for (int j=0 ; j<t ; j++){</pre>
              x = modmul(x, x, n);
               if (x==1) return 0;
               if (x==n-1) break:
          if (x==n-1) continue:
41
          return 0;
42
43
      return 1;
```

7.11 Pollard Rho

```
i | mt19937 seed(chrono::steady clock::now().time since epoch().
  int rnd(int 1, int r){
      return uniform int distribution<int>(1, r)(seed);
 6 // O(n^{1/4}) 回傳 1 或自己的因數、記得先判斷 n 是不是質數
        (用 Miller-Rabin)
 7 // c1670c
 s int Pollard_Rho(int n){
      int s = 0, t = 0;
      int c = rnd(1, n-1);
      int step = 0, goal = 1;
      int val = 1;
      for (goal=1 ; ; goal<<=1, s=t, val=1){</pre>
          for (step=1 ; step<=goal ; step++){</pre>
              t = ((__int128)t*t+c)%n;
              val = (__int128)val*abs(t-s)%n;
              if ((step % 127) == 0){
                  int d = __gcd(val, n);
                  if (d>1) return d;
          int d = __gcd(val, n);
          if (d>1) return d;
29
```

7.12 Quick Pow

```
int qp(int b, int p, int m = MOD){
   int ret = 1;
   for (; p; p>>=1){
       if (p&1) ret = ret*b%m;
       b = b*b%m;
   }
   return ret;
}
```

7.13 數論分塊

7.14 最大質因數

```
void max_fac(int n, int &ret){
    if (n<=ret || n<2) return;
    if (isprime(n)){
        ret = max(ret, n);
        return;
}

int p = Pollard_Rho(n);
max_fac(p, ret), max_fac(n/p, ret);
}</pre>
```

7.15 歐拉公式

```
1 / / phi(n) = 小於 n 並與 n 互質的正整數數量。
2 // O(sqrt(n)) · 回傳 phi(n)
3 int phi(int n){
      int ret = n;
      for (int i=2; i*i<=n; i++){
          if (n%i==0){
              while (n%i==0) n /= i;
              ret = ret*(i-1)/i;
      if (n>1) ret = ret*(n-1)/n;
      return ret;
  // O(n Log n) · 回傳 1~n 的 phi 值
  vector<int> phi 1 to n(int n){
      vector<int> phi(n+1);
      phi[0]=0;
      phi[1]=1;
      for (int i=2 ; i<=n ; i++){</pre>
          phi[i]=i-1;
      for (int i=2 ; i<=n ; i++){</pre>
          for (int j=2*i; j<=n; j+=i){ // 枚舉所有倍數
              phi[j]-=phi[i];
      }
31
      return phi;
```

7.16 線性篩

```
const int MAX_N = 5e5;

// Lpf[i] = i 的最小質因數
vector<int> prime, lpf(MAX_N);

void prime_init(){
```

```
for (int i=2; i<MAX_N; i++){
    if (lpf[i]==0){
        lpf[i] = i;
        prime.push_back(i);
}

for (int j: prime){
    if (i*j>=MAX_N) break;
    lpf[i*j] = j;
    if (i%j==0) break;
}

}
```

7.17 Burnside's Lemma

$$\sum_{k=1}^{n} \frac{c(k)}{n}$$

- n:有多少種置換方式(例如:旋轉方式)
- c(k): 所有可能中,經過 k 次旋轉後,仍不會和別人相同的方式的數量

7.18 Catalan Number

任意括號序列: $C_n = \frac{1}{n+1} {2n \choose n}$

7.19 Matrix Tree Theorem

目標:給定一張無向圖·問他的生成樹數量。 方法:先把所有自環刪掉·定義 Q 為以下矩陣

$$Q_{i,j} = \begin{cases} \deg(v_i) & \text{if } i = j \\ -(\frac{e}{2}v_iv_j) & \text{otherwise} \end{cases}$$

接著刪掉 Q 的第一個 row 跟 column · 它的 determinant 就是答案。 目標:給定一張有向圖 · 問他的以 r 為根 · 可以走到所有點生成樹數量

方法:先把所有自環刪掉,定義 Q 為以下矩陣

$$Q_{i,j} = \begin{cases} \deg_{in}(v_i) & \text{if } i = j \\ -(\frac{1}{8}v_iv_j \text{ bb}) & \text{otherwise} \end{cases}$$

接著刪掉 Q 的第 r 個 row 跟 column · 它的 determinant 就是答案。

7.20 Stirling's formula

```
n! \approx \sqrt{2\pi n} (\frac{n}{e})^n
```

7.21 Theorem

- 1. $1 \sim x$ 質數的數量 $\approx \frac{x}{\ln x}$
- 2. $1 \sim x$ 的因數的數量 $\approx x^{\frac{1}{3}}$
- 3. x 的質因數的數量 $\approx \log \log x$
- 4. p is a prime number $\Leftrightarrow (p-1)! \equiv -1 \pmod{p}$
- 5. 每個正整數都可以表示成四個整數的平方和
- 6. 任何大於 2 的整數都可以表示成兩個質數的和

7.22 二元一次方程式

```
\begin{cases} ax+by=c\\ dx+ey=f \end{cases} = \begin{cases} x=\frac{ed-bf}{ad-bc}\\ y=\frac{af-ec}{ad-bc} \end{cases} 若 x=\frac{0}{0} 且 y=\frac{0}{0} · 則代表無限多組解。若 x=\frac{*}{0} 且 y=\frac{*}{0} · 則代表無解。
```

7.23 歐拉定理

```
若 a,m 互質 \cdot 則: a^n\equiv a^{n \bmod \varphi(m)}\pmod m 若 a,m 不互質 \cdot 則: a^n\equiv a^{\varphi(m)+[n \bmod \varphi(m)]}\pmod m
```

7.24 錯排公式

錯排公式: $(n \oplus 1)$ 個人中·每個人皆不再原來位置的組合數)

$$dp_i = \begin{cases} 1 & i = 0\\ 0 & i = 1\\ (i-1)(dp_{i-1} + dp_{i-2}) & \text{otherwise} \end{cases}$$

8 String

8.1 Hash

```
int A = rng(1e5, 8e8);
const int B = 1e9+7;

struct RollingHash{
   vector<int> Pow, Pre;
   RollingHash(string s = ""){
        Pow.resize(s.size());
        Pre.resize(s.size());
}
```

```
for (int i=0 ; i<s.size() ; i++){
    if (i==0){
        Pow[i] = 1;
        Pre[i] = s[i];
    }
}

| Pow[i] = Pow[i-1]*A%B;
    Pre[i] = (Pre[i-1]*A+s[i])%B;

| Pre[i] = (Pre[i-1]*A+s[i])%B;

| Preturn;
    }

| Pow[i] = Pow[i-1]*A%B;
    | Pre[i] = (Pre[i-1]*A+s[i])%B;

| Preturn;
    | Preturn;
    | Preturn;
    | Preturn | Pret
```

8.2 KMP

8.3 Manacher

```
string Manacher(string str) {
    string tmp = "$#";
    for(char i : str) {
        tmp += i;
        tmp += '#';
    }

vector<int> p(tmp.size(), 0);
    int mx = 0, id = 0, len = 0, center = 0;
    for(int i=1; i<(int)tmp.size(); i++) {
        p[i] = mx > i ? min(p[id*2-i], mx-i) : 1;

    while(tmp[i+p[i]] == tmp[i-p[i]]) p[i]++;
    if(mx<i+p[i]) mx = i+p[i], id = i;
    if(len<p[i]) len = p[i], center = i;
    }

return str.substr((center-len)/2, len-1);
}</pre>
```

8.4 Min Rotation

8.5 Suffix Array

```
ı|// 注意,當 /s/=1 時,Lcp 不會有值,務必測試 /s/=1 的 case
2 struct SuffixArray {
     string s;
      vector<int> sa, lcp:
      SuffixArray(string _s, int lim = 256) {
          int n = s.size()+1, k = 0, a, b;
          vector<int> x(s.begin(), s.end()), y(n), ws(max(n,
               lim)), rank(n);
         x.push back(0);
          sa = 1cp = v;
          iota(sa.begin(), sa.end(), 0);
          for (int j=0, p=0 ; p<n ; j=max(1LL, j*2), lim=p) {</pre>
              p = j;
              iota(y.begin(), y.end(), n-j);
              for (int i=0; i<n; i++) if (sa[i] >= j) y[p++]
                  = sa[i] - j;
              fill(ws.begin(), ws.end(), 0);
              for (int i=0; i<n; i++) ws[x[i]]++;</pre>
              for (int i=1; i<lim; i++) ws[i] += ws[i - 1];</pre>
              for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
              swap(x, y), p = 1, x[sa[0]] = 0;
              for (int i=1; i<n; i++){</pre>
                 a = sa[i - 1];
                  b = sa[i];
                  x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
                       ? p - 1 : p++;
          for (int i=1; i<n; i++) rank[sa[i]] = i;</pre>
          for (int i=0, j ; i<n-1 ; lcp[rank[i++]]=k)</pre>
              for (k && k--, j=sa[rank[i]-1]; i+k<s.size() &&</pre>
                   j+k<s.size() && s[i+k]==s[j+k] ; k++);
          sa.erase(sa.begin());
         lcp.erase(lcp.begin(), lcp.begin()+2);
     vector<int> pos; // pos[i] = i 這個值在 pos 的哪個地方
     SparseTable st;
     void init lcp(){
         pos.resize(sa.size());
         for (int i=0 ; i<sa.size() ; i++){</pre>
              pos[sa[i]] = i;
```

```
if (lcp.size()){
                                                                        { // build suf
43
             st.build(lcp);
                                                             100
                                                                            int now = 0;
                                                                            for (int i=s.size()-1; i>=0; i--){
     }
                                                                                if (sa[i]<=p){</pre>
                                                             102
                                                             103
                                                                                    suf[sa[i]] = now;
                                                                                    if (i-1>=0) now = min(now, lcp[i-1]);
                                                             104
     // 用之前記得 init
     // 回傳 [l1, r1] 跟 [l2, r2] 的 Lcp·0-based
                                                                                    if (i-1>=0) now = lcp[i-1];
     int get_lcp(int l1, int r1, int l2, int r2){
                                                             107
          int pos_1 = pos[l1], len_1 = r1-l1+1;
                                                             108
          int pos_2 = pos[12], len_2 = r2-12+1;
         if (pos_1>pos_2){
                                                             110
              swap(pos_1, pos_2);
                                                                        return {pre, suf};
                                                             111
             swap(len_1, len_2);
                                                             112
                                                             113 };
         if (11==12){
              return min(len 1, len 2);
                                                                8.6 Z Algorithm
             return min({st.query(pos_1, pos_2), len_1, len_2
                                                               1 \mid // 定義一個長度為 n 的文本為 T · 則陣列 Z 的 Z[i] 代表 T[0:n]
     }
62
                                                                      和 T[i:n] 最長共同前綴
                                                               2 // bcfbd6
     // 檢查 [l1, r1] 跟 [l2, r2] 的大小關係·0-based
                                                               3 vector<int> z_function(string s){
      // 如果前者小於後者,就回傳 <0,相等就回傳 =0,否則回傳
                                                                    vector<int> ret(s.size());
                                                                    int 11 = 0, rr = 0;
      int substring cmp(int l1, int r1, int l2, int r2){
          int len 1 = r1-l1+1;
                                                                     for (int i=1; i<s.size(); i++){</pre>
          int len 2 = r2-12+1;
                                                                        int j = 0;
         int res = get_lcp(l1, r1, l2, r2);
                                                                        if (i<rr) j = min(ret[i-ll], rr-i);</pre>
         if (res<len_1 && res<len_2){</pre>
                                                                        while (s[j]==s[i+j]) j++;
             return s[l1+res]-s[l2+res];
                                                                        ret[i] = j;
         }else if (len 1==res && len 2==res){
             // 如果不需要以 index 作為次要排序參數,這裡要回
                                                                        if (i+j>rr){
                                                                            11 = i;
             return 11-12;
                                                                            rr = i+j;
         }else{
             return len_1==res ? -1 : 1;
                                                                    }
     }
                                                                    ret[0] = s.size();
                                                                    return ret;
     // 對於位置在 <=p 的後綴·找離他左邊/右邊最接近位置 >p 的
          後綴的 Lcp · 0-based
     // pre[i] = s[i] 離他左邊最接近位置 >p 的後綴的 Lcp · 0-
     // suf[i] = s[i] 離他右邊最接近位置 >p 的後綴的 Lcp · 0-
      pair<vector<int>, vector<int>> get left and right lcp(int
           p){
          vector<int> pre(p+1);
         vector<int> suf(p+1);
         { // build pre
             int now = 0:
             for (int i=0 ; i<s.size() ; i++){</pre>
                 if (sa[i]<=p){</pre>
                     pre[sa[i]] = now;
                     if (i<lcp.size()) now = min(now, lcp[i]);</pre>
                     if (i<lcp.size()) now = lcp[i];</pre>
```