Q Find number of triplets

i, j, k { indices }

such that

$i < j < k$

&

$arr[i] < arr[j] < arr[k]$

	0	1	2	3	4
arr:	3	4	6	9	2

i	j	k
0	1	2
0	2	3
1	2	3
0	1	3

ans = 4

	0	1	2	3	4
arr:	2	6	9	4	10

0	1	2
0	1	4
0	2	4
0	3	4
1	2	4

ans = 5

Basic :- consider all possible triplets

$i < j < k$

$arr[i] < arr[j]$

```
cnt = 0;
for( i=0; i<n; i++)
{
    for( j=i+1; j<n; j++)
    {
        if( arr[i] >= arr[j]) continue;
        for(int k = j+1; k<n; k++)
        {
            if( arr[k] > arr[j])
                cnt++;
        }
    }
}
```

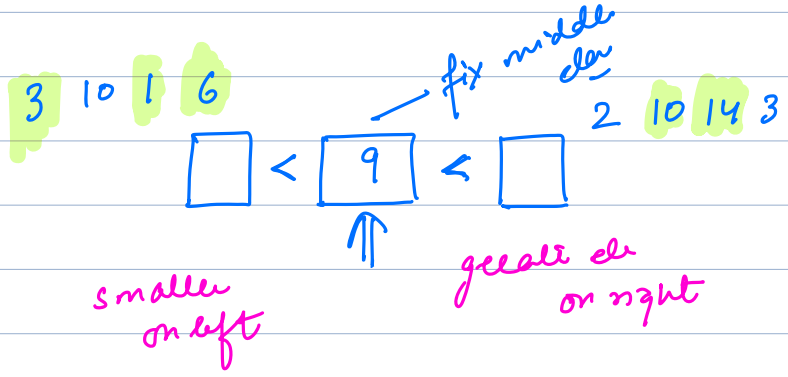
$O(N^3)$

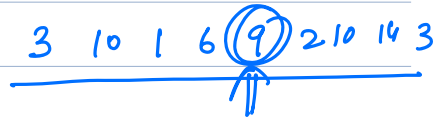
We can't sort over here.

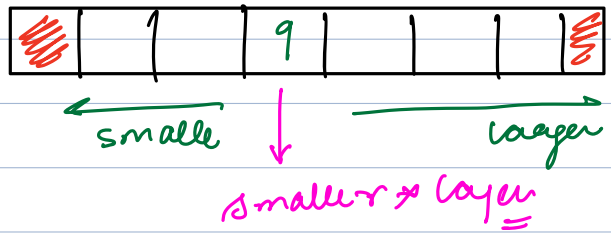
N^2 ✓

$N^2 \log_2 N$

N^2







0	1	2	3	4	5	6
3	6	9	12	16	11	1

smaller: 1 2

larger: 4 3

4

1 → n-2 — treating each element as middle

```
for( i=1; i<n-1; i++)
{
    // i+th -> middle
    ls=0;
    for(j=0; j<i; j++)
    {  if(arr[j]<arr[i])
           ls++;
    }
    rs=0;
    for( j=i+1; j<n; j++)
    {
        ____
    }

ans+= ls*rs;
}
```

T.C: O(N²)

↓

T.C: nlogn

BBST

takeaway - O(N²) ✓

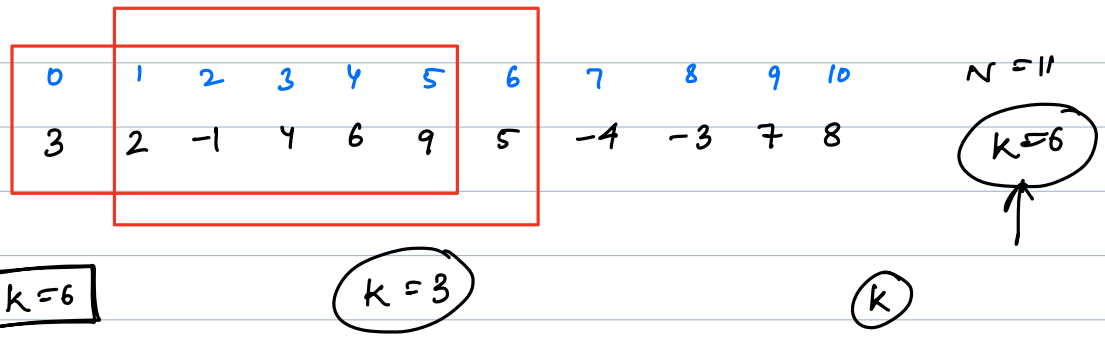
↓

O(N²)

O(N) tw

Q array of size N.

Print (start, end) indexes of all subarrays of size K.



k=6 | k=3 | k

0-5
1-6
⋮
5-10

6 subs

0-2
⋮
8-10

9 subarr

✓0 → k-1
1 → k
2 → k+1
3 → k+2
⋮
:k
✓N-k → N-1

n-k+1 ⇒ total count of subarr of size k

b - a + 1 = k
(N-1) - a + 1 = k
a = N-k

```
int s = 0;
int e = k-1;
while (e < N)
{
    print(s,e);
    s++; e++;
}
```

start & end pt of first subarr of size k

⇒ consider all possible subarr of size k

Q Find the max subarray sum of size k.

0	1	2	3	4	5	6	7	8	9
-3	4	-2	5	3	-2	8	2	-1	4

k = 5

s	e	sum
0	4	7
1	5	8
2	6	12
3	7	16
4	8	10
5	9	11

```
s = 0;
e = k-1;
while (e < N)
{   // s,e   sum = 0;
    for(i = s; i <= e; i++)
        sum += a[i];
    ans = max(ans, sum);
}
```

prefix sum

T.C: O(N)
S.C: O(N)

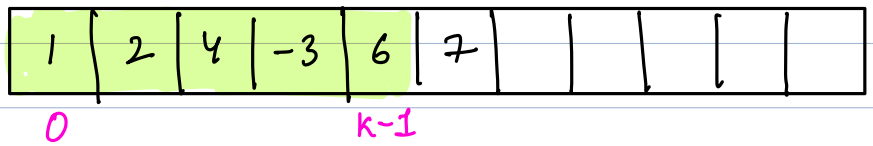
T.C: $(n-k+1)*k$

k = 1 → $(n-1+1)*1$ → O(n)

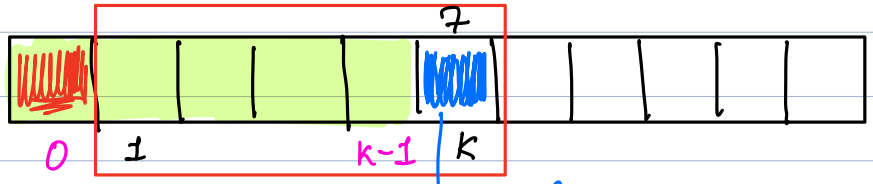
k = N → $(n-n+1)*n$ → O(n)

$k = n/2$ → $(n-n/2+1)*n/2$ ≈ $\frac{n^2}{4}$ → $O(n^2)$

k size subar



sum = 10

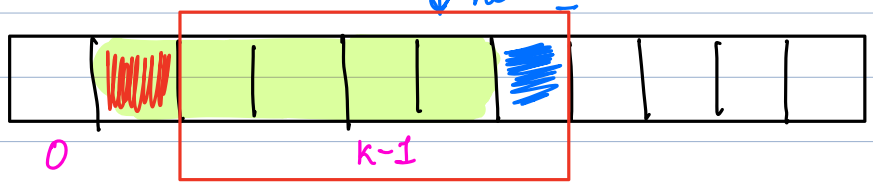


sum + 7 - 1

10 + 7 - 1

= 16

new sum



0 ⟶ k-1 iterate & calculate = sum

1 ⟶ k sum + arr[k] − arr[0] = sum

2 ⟶ k+1 sum + arr[k+1] − arr[1]

⋮

N-k ⟶ N-1 sum + arr[N-1] − arr[N-k-1]

sum + arr[e] − arr[s-1]

```
sum = 0;   ans = -∞;
// calculate for first subarr
for ( i=0; i<k; i++)
{
        sum += arr[i];
}
ans = max(ans, sum);   // ans = sum;  ✓
```

(k)

```
{ s = 1;
  e = k;
  while ( e<n )
  {
        sum = sum + arr[e] - arr[s-1];
        ans = max(ans, sum);
        s++;
        e++;
  }
```

(N-K)

~~K~~ + N - ~~K~~

T.C: O(N)

S.C: O(1)

Takeaway → whenever subarray size is fixed

↓

sliding window

10:43

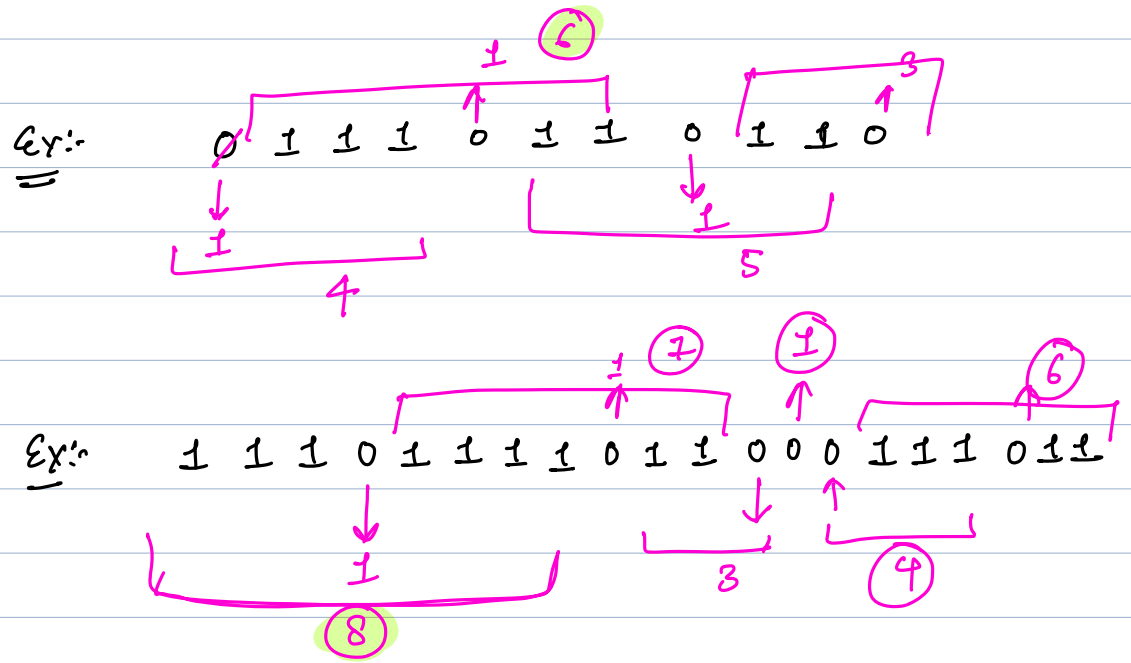
Q. Binary array of size N

↓

0 & 1's

You are allowed to replace atmost one 0 with a (1)

length of longest consecutive 1's ?



Go to every zero → consecutive 1's ?

length of consecutive one's on left + length of consecutive on right

ls + rs + 1

```
ans = 0;

for(int i=0; i<n; i++)
{
    if(arr[i]==0)
    {
        j = i-1; ls = 0;
        while(j>=0 && arr[j]==1)
        {
            ls++;
            j--;
        }
        // calculate rs
        ans = max(ans, ls+rs+1);
    }
}
```

i-1 i

0

1 1 1 1 1 1

```
if(ans==0)
    return ans-left(c);
```

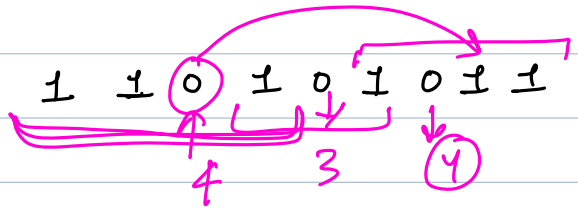
0 1 1 1 1 0 1 1 1 0 1 1 0 1

1

(3N) T.C: O(N)

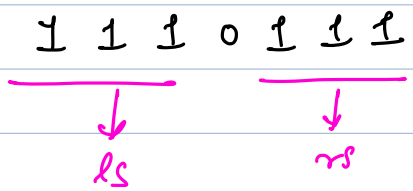
replace 0 with 1

swap atmost one '0' with '1'



lstrs != total

ans = max(ans, lstrs);



lstrs == total no of one's

↓

no extra one

ans = max(ans, (strs));

extra 1