# A tensor-based approach to solving systems of multivariate polynomials

Nithin Govindarajan

with Raphaël Widdershoven, Shiv Chandrasekaran (UCSB), and Lieven De Lathauwer

June 15th 2023



## Overview

Motivation: noisy overdetermined polynomial systems

Polynomial root solving: from an eigenvalue to a tensor decomposition problem

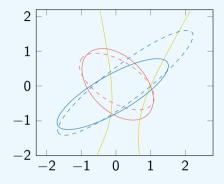
Faster Macaulay null space computations

Summar

# Noisy overdetermined systems: looking for approximative roots

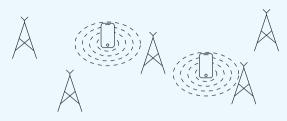
$$\begin{cases}
-3 - x - 2y + 4x^2 + 6xy + 7y^2 = 0 \\
-2 - x + y + 3x^2 - 7xy + 5y^2 = 0 \\
1 + 7x + y - 8x^2 + 3xy + y^2 = 0
\end{cases}$$

- N = 2 unknowns
- S = 3 equations  $\rightarrow$  overdetermined
- Degree d = 2



Adding noise to the red and blue equations destroys the single exact root at (0,1)

## A practical application: "blind" multi-source localization

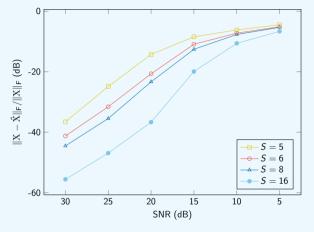


Friis transmission equation (before conversion into a polynomial expression):

$$P_i^r = \frac{A_i^r A_1^t}{\lambda^2} \frac{P_1^t}{(x_i^r - x_1^t)^2 + (y_i^r - y_1^t)^2} + \frac{A_i^r A_2^t}{\lambda^2} \frac{P_2^t}{(x_i^r - x_2^t)^2 + (y_i^r - y_2^t)^2}, \quad i = 1, \dots, S.$$
Noisy measured quantities! Unknown Given

 $\longrightarrow$  for  $S \ge 5$ , positions of transmitters can be *retrieved* up to permutation ambiguity!

# Similar to least-squares: adding more equations (i.e., antennas) yield better estimates



Median relative error of estimated transmitter positions over 200 experiments (Widdershoven, Govindarajan, et al. 2023)

## Overview

Motivation: noisy overdetermined polynomial systems

Polynomial root solving: from an eigenvalue to a tensor decomposition problem

Faster Macaulay null space computations

Summar

# Algebraic methods: "classical" vs. recent numerical (multi)-linear algebra approaches

Find all (projective) roots of the system of the multivariate polynomials:

$$\Sigma: \left\{ egin{array}{lll} p_1 &=& p_1(x_1,x_2,\ldots,x_N) \ &dots & & , & S\geq N, & \deg(p_s)=d_s. \ p_S &=& p_S(x_1,x_2,\ldots,x_N) \end{array} 
ight.$$

Auzinger, Stetter, Lazard, ...

Batselier, Dreesen....

Numerical Polynomial

Linear Algebra (NPLA)

Vanderstukken, De Lathauwer,...

Numerical Polynomial

Multi-Linear Algebra (NPMLA)

Numerical Polynomial Algebra (NPA)

Features:

Features:

- Gröbner basis construction

Features:

- Macaulay null space construction - Generalized eigenvalue problem

- Macaulay null space construction

- eigenvalue problem

- tensor decomposition problem

DISCLAIMER: The above is a very selected overview and only shows "ancestors" of our own work. It is by no means a summary of all the contributions done on this topic.

# The Macaulay-based method for polynomial root solving

The rows of the Macaulay matrix M(d) span the set

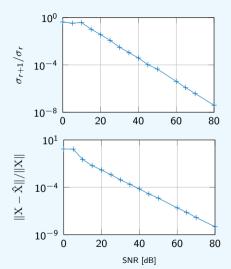
$$\mathcal{M}_d := \left\{ \sum_{s=1}^{S} h_s \cdot p_s : \quad \deg(h_s) = d - d_s 
ight\} = \left\langle p_1^h, \dots, p_s^h 
ight
angle_d.$$

For example, M(3) for the system in slide 3:

If  $d \ge d^*$  (degree of regularity), dim null M(d) = no. of projective roots of the system.

# Our system from slide 3: Macaulay method is suitable in the noisy setting

### Add noise to the nonzero coefficients:



Note: M(4) is expressed in lex ordering this time!

# The matrix view: recovering the roots from a generalized eigenvalue problem

- Let  $S_t, S_{x_1}, S_{x_2}, \dots, S_{x_N}$  denote appropriate "row selection" matrices.
- lacksquare Construct  $G_t = S_t N$  and  $G_{x_i} = S_{x_i} N$  for  $i = 1, \dots, N$  with  $N = \mathsf{null}\, M(d)$
- Solve the generalized eigenvalue decomposition (GEVD) problem:

$$(\alpha_t G_t + \alpha_{x_1} G_{x_1} + \ldots + \alpha_{x_N} G_{x_N}) \mathbf{a} = \lambda (\beta_t G_t + \beta_{x_1} G_{x_1} + \ldots + \beta_{x_N} G_{x_N}) \mathbf{a}$$

■ For i = 1, ..., R (number of roots), we have the eigenvalues:

$$\lambda_{i} = \frac{\alpha_{t} t^{(i)} + \alpha_{x_{1}} x_{1}^{(i)} + \ldots + \alpha_{x_{N}} x_{N}^{(i)}}{\beta_{t} t^{(i)} + \beta_{x_{1}} x_{1}^{(i)} + \ldots + \beta_{x_{N}} x_{N}^{(i)}}$$

■ Eigenvectors *reveal* root location, since

$$\mathbf{v}_i = (\alpha_t G_t + \alpha_{x_1} G_{x_1} + \ldots + \alpha_{x_N} G_{x_N}) \mathbf{a}_i$$

are multivariate Vandermonde vectors evaluated at the system roots!

Matrix pencils can be badly conditioned when eigenvalues coalesce...

# Theorem (See e.g., Golub and Van Loan 2012))

$$\begin{aligned} \textit{Define} \ \mathsf{sep}(T_{11}, T_{22}) &:= \mathsf{min} \, \|T_{11}X - XT_{22}\|_{\textit{F}} \, / \, \|X\|_{\textit{F}} \ \textit{and let} \ Q^*AQ = \begin{bmatrix} T_{11} & T_{12} \\ 0 & T_{22} \end{bmatrix} \\ \textit{with} \ Q &= \begin{bmatrix} Q_1 & Q_2 \end{bmatrix}. \ \textit{Then for a "sufficiently small"} \ E, \ \textit{there exists a} \ \tilde{Q}_1 \ \textit{with} \end{aligned}$$

$$\mathsf{dist}(\mathsf{col}\ \mathrm{Q}_1,\mathsf{col}\ \tilde{\mathrm{Q}}_1) \lesssim \frac{1}{\mathsf{sep}(\mathrm{T}_{11},\mathrm{T}_{22})}$$

such that  $\tilde{Q}_1$  is an invariant subspace for  $\tilde{A} = A + E$ .

A bad choice of  $\alpha$ 's and  $\beta$ 's can bring eigenvalues arbitrarily close!

# GESD principle: why limit to just one pencil, if you can exploit multiple?

**GESD** algorithm: using multiple pencils, recursively split eigenspaces corresponding to well-separated eigenvalue clusters (Evert, Vandecappelle, et al. 2022).

## Simultaneous diagonalization:

There exists an invertible  $A \in \mathbb{C}^{R \times R}$  that simultaneously diagonalizes

$$G_t A = V \operatorname{diag}(t^{(1)}, \dots, t^{(R)}),$$
 $G_{x_1} A = V \operatorname{diag}(x_1^{(1)}, \dots, x_1^{(R)}),$ 
 $\vdots$ 
 $G_{x_N} A = V \operatorname{diag}(x_N^{(1)}, \dots, x_N^{(R)}),$ 

with V being multivariate Vandermonde matrix evaluated at the roots.

Reformulation of the root recovery as a tensor decomposition problem

## Theorem (Vanderstukken and De Lathauwer 2021)

Let  $\mathcal G$  have frontal slices  $\mathrm G_t, \mathrm G_{\mathsf x_1}, \dots, \mathrm G_{\mathsf x_2}$ , and assume  $\Sigma$  has only simple roots. If  $\mathcal G$  is constructed from  $\mathsf{null}\, \mathrm M(d)$  with  $d \geq d^* + 1$ , then  $\mathcal G$  has the essentially unique CPD

$$\mathcal{G} = [\![ \mathbf{V}, \mathbf{A}^{-1}, \mathbf{X} ]\!], \quad X = \begin{bmatrix} t^{(1)} & t^{(2)} & \cdots & t^{(R)} \\ x_1^{(1)} & x_1^{(2)} & \cdots & x_1^{(R)} \\ \vdots & \vdots & & \vdots \\ x_N^{(1)} & x_N^{(2)} & \cdots & x_N^{(R)} \end{bmatrix}.$$

$$\mathcal{G}$$
 =  $+ \cdots +$ 

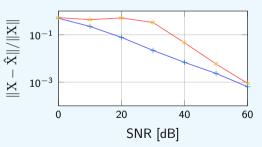
If the polynomial system has roots of multiplicity greater than one, the theorem can be generalized with the introduction of block-term decompositions (Vanderstukken, Kürschner, et al. 2021)

# Observed numerical benefits of the tensor approach for noisy overdetermined systems

Take N = 10 noisy copies of the square system:

$$\Sigma: \left\{ \begin{array}{l} f_1(x_1, x_2) = x_1^3 + x_2^3 - 9x_1^2x_2 + 20x_1x_2 - 3x_1 - 20 = 0 \\ f_2(x_1, x_2) = x_1^2 + 4x_2^2 - x_1x_2 - 80 = 0 \end{array} \right.$$

median forward error over 200 trials



The tensor-based method that relies on simultaneous diagonalization is better capable of recovering roots in noisy conditions than a pure matrix-based method which relies solely on GEVD (Vanderstukken and De Lathauwer 2021).

## Overview

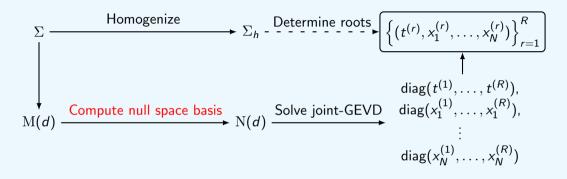
Motivation: noisy overdetermined polynomial systems

Polynomial root solving: from an eigenvalue to a tensor decomposition problem

Faster Macaulay null space computations

Summar

# Null space computation is the major computational bottleneck in many algorithms!



Exploit the *Toeplitz* structures in Macaulay matrices?

# Bivariate systems: Macaulay matrix is almost Toeplitz block-(block-)Toeplitz

		1	x	$x^2$	<sub>x</sub> 3	$x^4$	у	xy	$x^2y$	$x^3y$	y <sup>2</sup>	$xy^2$	$x^2y^2$	у3	$xy^3$	y <sup>4</sup>
	$\rho_1$	1	6	4			2	5			3					
	$p_2$	9	1	3			8	7			2					
	$xp_1$		1	6	4			2	5			3				
	xp <sub>2</sub>		9	1	3			8	7			2				
	$x^2p_1$			1	6	4			2	5			3			
M(4) =	$x^{2}p_{2}$			9	1	3			8	7			2			
W(4) =	ур1						1	6	4		2	5		3		
	yp <sub>2</sub>						9	1	3		8	7		2		
	xyp <sub>1</sub>							1	6	4		2	5		3	
	xyp <sub>2</sub>							9	1	3		8	7		2	
	$y^{2}p_{1}$										1	6	4	2	5	3
	$y^2p_2$										9	1	3	8	7	2

# The Macaulay matrix for the general bivariate case

Let  $\Delta d := d - d_{\Sigma}$ . Then,

$$\mathrm{M}(d) := \begin{bmatrix} \mathrm{M}_{0,0} & \mathrm{M}_{1,0} & \cdots & \mathrm{M}_{d_{\Sigma},0} \\ & \mathrm{M}_{0,1} & \mathrm{M}_{1,1} & \cdots & \mathrm{M}_{d_{\Sigma},1} \\ & & \ddots & \ddots & & \ddots \\ & & & \mathrm{M}_{0,\Delta d} & \mathrm{M}_{1,\Delta d} & \cdots & \mathrm{M}_{d_{\Sigma},\Delta d} \end{bmatrix} \in \mathbb{C}^{\frac{S}{2}(\Delta d + 1)(\Delta d + 2) \times \frac{1}{2}(d + 1)(d + 2)},$$

with

$$\mathrm{M}_{i,j} := egin{bmatrix} oldsymbol{c}_{0i} & oldsymbol{c}_{1i} & \cdots & oldsymbol{c}_{(d_{\Sigma}-i)i} & & & & & & \\ & oldsymbol{c}_{0i} & oldsymbol{c}_{1i} & \cdots & oldsymbol{c}_{(d_{\Sigma}-i)i} & & & & & \\ & & \ddots & \ddots & & & & \ddots & & \\ & & oldsymbol{c}_{0i} & oldsymbol{c}_{1i} & \cdots & oldsymbol{c}_{(d_{\Sigma}-i)i} \end{bmatrix} \in \mathbb{C}^{S(\Delta d+1-j) imes (d+1-i-j)}.$$

## The key observation that shall allow for a faster algorithm

Consider the displacement operator

$$\mathscr{D}\left\{\mathrm{M}(d)
ight\} = egin{bmatrix} \mathrm{Z}_{d+1,1} \otimes \mathrm{I}_{\mathcal{S}} & & & & \\ & & \ddots & & & \\ & & & \mathrm{Z}_{1,1} \otimes \mathrm{I}_{\mathcal{S}} \end{bmatrix} \mathrm{M}(d) - \mathrm{M}(d) egin{bmatrix} \mathrm{Z}_{d+1,arphi_{d+1}} & & & & \\ & & & \ddots & & \\ & & & & \mathrm{Z}_{1,arphi_{1}} \end{bmatrix}.$$

## M(d) has relative "low" displacement rank

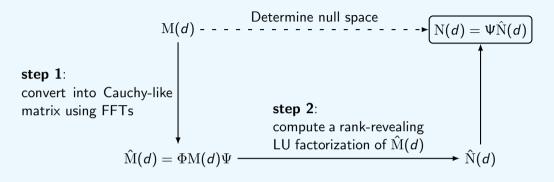
Dimensions of  $M(d) \in \mathbb{C}^{\frac{5}{2}(\Delta d+1)(\Delta d+2) \times \frac{1}{2}(d+1)(d+2)}$  grow quadratically w.r.t. d, but

$$\operatorname{\mathsf{rank}} \mathscr{D}\left\{\operatorname{M}(d)\right\} \leq S(\Delta d + 1) = S\left(d + 1 - d_{\Sigma}\right).$$

grows only *linearly* with d.

Here 
$$Z_{p,\varphi}:=\begin{bmatrix} 1 & & \varphi \\ & \ddots & \\ & & 1 \end{bmatrix}\in\mathbb{C}^{p\times p}.$$

# Overview of the fast algorithm



Both steps can be done *fast*! (Govindarajan, Widdershoven, et al. 2023)

# Rank-revealing LU factorization of $\hat{\mathrm{M}}(d)$ (Miranian and Gu 2003)

Let  $r(d) := \operatorname{rank} M(d)$ . Compute a rank-revealing LU (RRLU) factorization

$$\begin{split} \Pi_{1}\hat{\mathbf{M}}(d)\Pi_{2} &= \begin{bmatrix} \mathbf{I}_{r(d)} \\ \hat{\mathbf{M}}_{21}\hat{\mathbf{M}}_{11}^{-1} & \mathbf{I}_{d_{\Sigma}^{2}} \end{bmatrix} \begin{bmatrix} \hat{\mathbf{M}}_{11} \\ & \hat{\mathbf{M}}_{22} - \hat{\mathbf{M}}_{21}\hat{\mathbf{M}}_{11}^{-1}\hat{\mathbf{M}}_{12} \end{bmatrix} \begin{bmatrix} \mathbf{I}_{r(d)} & \hat{\mathbf{M}}_{11}^{-1}\hat{\mathbf{M}}_{12} \\ & \mathbf{I}_{d_{\Sigma}^{2}} \end{bmatrix} \\ &\approx \begin{bmatrix} \hat{\mathbf{M}}_{11} \\ \hat{\mathbf{M}}_{21} \end{bmatrix} \begin{bmatrix} \mathbf{I}_{r(d)} & \hat{\mathbf{M}}_{11}^{-1}\hat{\mathbf{M}}_{12} \end{bmatrix} \end{split}$$

# Expression for the null space N(d)

$$\mathrm{N}(\mathit{d}) = \Psi \Pi_2 \begin{bmatrix} \tilde{\mathrm{N}} \\ \mathrm{I}_{\mathit{d}_{\Sigma}^2} \end{bmatrix}, \qquad \tilde{\mathrm{N}} := -\hat{\mathrm{M}}_{11}^{-1} \hat{\mathrm{M}}_{12}.$$

# The classical Schur algorithm: making things work for us

- Run Schur algorithm on Cauchy-like matrices, which satisfy the displacement relation  $\operatorname{diag}(\mu)\mathrm{C}-\mathrm{Cdiag}(\nu)=\mathrm{RS}^*$  (Heinig 1995)
- Replace total pivoting with approximate total pivoting
- Let  $j_{max}$  denote the column with largest 2-norm in RS\*. Then,

$$\max_{1 \leq i \leq n} \left| c_{ij_{\max}} \right| \geq \frac{1}{K\sqrt{n}} \max_{1 \leq i,j \leq n} \left| c_{ij} \right|, \quad K := \max_{1 \leq i,j,i,j \leq n} \left| \mu_i - \nu_j \right| / \left| \mu_i - \nu_j \right|.$$

- $\blacksquare$  Keeping R orthogonal during the Gaussian elimination process allows for fast pivot selection.
- Use clever updating strategies to keep cost low.

RESULT: from 
$$\mathcal{O}(d^6)$$
 to  $\mathcal{O}(d^5)$ 

## Algorithm stability: error grows linearly with problem size

Median error  $\epsilon := \|\mathrm{M}(d)\mathrm{Q}\|_2 / \|\mathrm{M}(d)\|_2 \ge \sigma_{r(d)+1}/\sigma_1$ , with  $\mathrm{Q} \in \mathbb{C}^{n(d) \times d_{\Sigma}^2}$  an orthonormal basis for col  $\mathrm{N}(d)$ , over 100 runs for randomly generated *square* systems.

	$d_{\Sigma}$									
	2	4	8	16	32					
SVD on $\mathrm{M}(d)$	2.23e-16	3.75e-16	5.70e-16	7.94e-16	9.51e-16					
GECP on $M(d)$	1.40e-16	3.11e-16	8.33e-16	1.02e-14	1.40e-13					
GECP on $\mathscr C$	4.35e-16	1.51e-15	1.35e-14	1.72e-13	2.81e-12					
GEAP on &	4.21e-16	3.63e-15	3.88e-14	3.19e-13	4.48e-12					

#### Sources of error:

- switching to LU instead of an SVD
- lacktriangle working with the compact Cauchy representation  $\mathscr C$
- switching to approximate pivoting ← Surprisingly not so bad!

## Extending the method

- $lue{}$  Generalizations to Chebyshev systems possible; Toeplitz ightarrow Toeplitz-plus-Hankel  $\odot$
- Displacement rank theory does not generalize nicely to higher dimensions with diminishing returns  $\mathcal{O}(d^{3N})$  to  $\mathcal{O}(d^{3N-1})$   $\odot$
- Assume additional (block) sparsity of Macaulay matrix to make breakthrough with n-dimensional systems?

$$p_i(x,y) = \sum_{k=0}^{d_{\Sigma}-r} a_{ikr} x^k y^r + \sum_{r \in \mathscr{S}} \sum_{k=0}^{d_{\Sigma}-r} a_{ikr} x^k y^r, \quad \mathscr{S} \subset \{1,\ldots,d_{\Sigma}\}, \quad |\mathscr{S}| = \rho \ll d_{\Sigma}$$

## Overview

Motivation: noisy overdetermined polynomial system:

Polynomial root solving: from an eigenvalue to a tensor decomposition problem

Faster Macaulay null space computations

Summary

#### What we have discussed in this talk

- Solving polynomial systems in the *noisy* overdetermined setting.
- Benefits of taking on a "tensor" view towards polynomial root solving.
- Progress and challenges towards (asymptotically) *faster* Macaulay null space algorithms.

### References I

- Evert, Eric, Michiel Vandecappelle, and Lieven De Lathauwer (2022). "A Recursive Eigenspace Computation for the Canonical Polyadic Decomposition". In: SIAM J. Matrix Anal. Appl. 43.1, pp. 274–300.
- Golub, G.H. and C.F. Van Loan (2012). *Matrix Computations*. 4th ed. Johns Hopkins University Press.
- Govindarajan, Nithin et al. (2023). "A fast algorithm for computing Macaulay nullspaces of bivariate polynomial systems". In: *Technical Report 23-16, ESAT-STADIUS, KU Leuven (Leuven, Belgium)*.
- Heinig, Georg (1995). "Inversion of Generalized Cauchy Matrices and other Classes of Structured Matrices". In: *Linear Algebra for Signal Processing*. Ed. by A. Bojanczyk and G. Cybenko. Vol. 69. The IMA Volumes in Mathematics and its Applications. New York, NY, USA: Springer.

### References II

- Miranian, L and Ming Gu (2003). "Strong rank revealing LU factorizations". In: Linear Algebra Appl. 367, pp. 1–16.
- Vanderstukken, Jeroen and Lieven De Lathauwer (2021). "Systems of Polynomial Equations, Higher-Order Rensor Decompositions and Multidimensional Harmonic Retrieval: A Unifying Framework. Part I: The Canonical Polyadic Decomposition". eng. In: SIAM J. Matrix Anal. Appl. 42.2, pp. 883–912.
- Vanderstukken, Jeroen et al. (2021). "Systems of Polynomial Equations, Higher-Order Tensor Decompositions, and Multidimensional Harmonic Retrieval: A Unifying Framework. Part II: The Block Term Decomposition". In: SIAM J. Matrix Anal. Appl. 42.2, pp. 913–953.

#### References III



Widdershoven, Raphaël, Nithin Govindarajan, and Lieven De Lathauwer (2023). "Overdetermined systems of polynomial equations: tensor-based solution and application". In: Technical Report 23-37, ESAT-STADIUS, KU Leuven (Leuven, Belgium). Accepted for publication in European Signal Processing Conference (EUSIPCO) 2023.