

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

This is one of the greatest books for beginners in quantum computing, in my opinion, especially for those who might feel intimidated by the abundance of complicated online resources. The book's Analogies section, which simplified complicated topics for laypeople, was its strongest point. The majority of the book's issues are resolved below with little to no prior knowledge required, and repeated reads of the book made nearly every topic easier.

Note:- This book is solved according to my understanding. please email me for any comments or feedback.

Book link

<https://link.springer.com/book/10.1007/978-3-030-98339-0>

Chapter 1

Basic intro

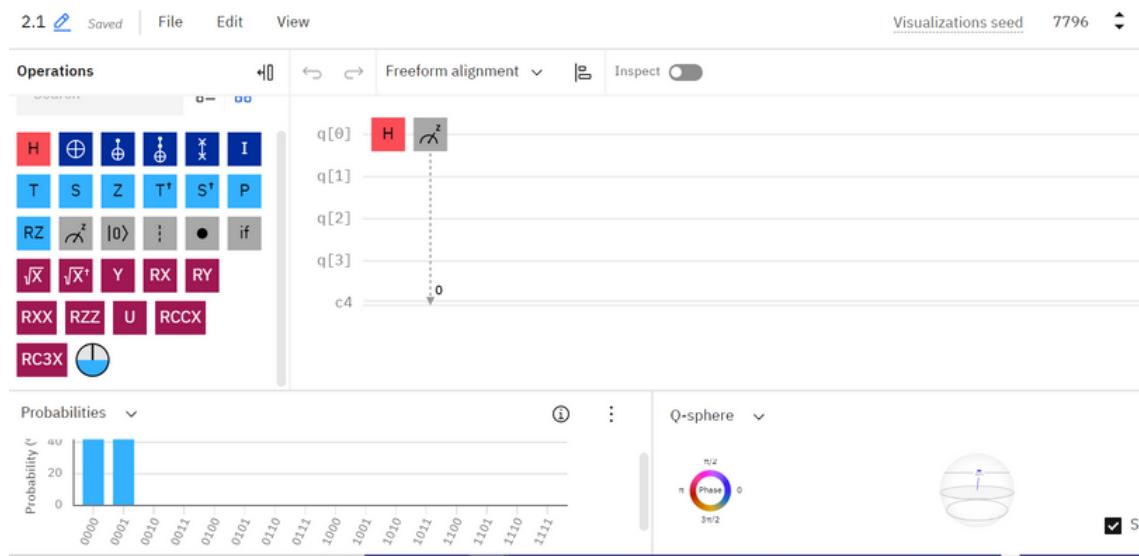
Code link

<https://github.com/nithinGovindugari/IntrotoQC>

Chapter 2

2.1) Created IBM account

2.2)



2.3)

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Details

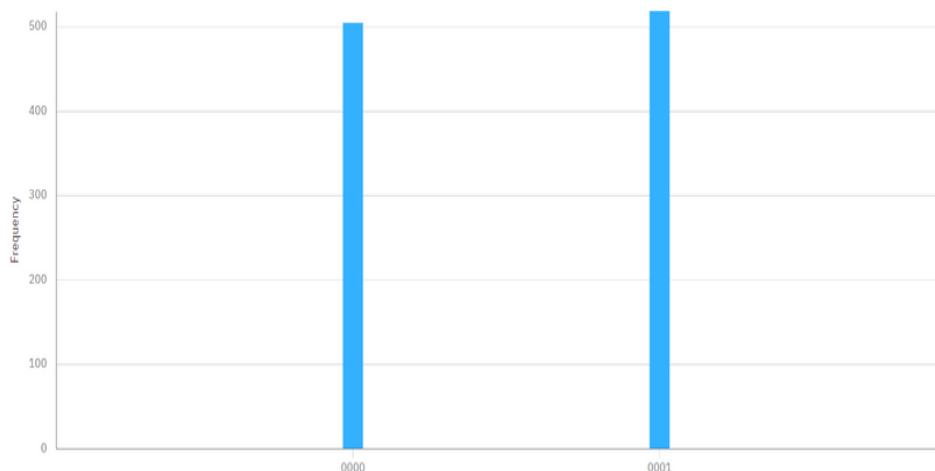
23m 57.8s

Total completion time

ibmq_manila

Backend

Histogram



Chapter 3

3.1) Had a google collab account already

3.2 & 3.3 & 3.4)

```
import numpy as np
from sklearn.preprocessing import normalize

#problem3
v1 = np.array([[np.sqrt(3)],[1]])
v2 = np.array([[1],[1]])
inner_result = np.vdot(v1,v2)
print(inner_result)
print(v1.shape)

v3 = np.array([[1j],[1]])
v4 = np.array([[ -1j],[1]])
inner_result1 = np.vdot(v3,v4)
print(inner_result1)

v5 = np.array([[1+2j],[1]])
v6 = np.array([[1-1j],[1]])
inner_result2 = np.vdot(v5,v6)
print(inner_result2)
```

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Chapter 4

4.1)

From 4.18 we know

$$|V'\rangle = \frac{|V\rangle}{\sqrt{\sum_{i=0}^{n-1} |a_i|^2}}$$

We can derive numerators of above co-efficients

$$|V_1\rangle = 2|1\rangle + |1\rangle$$

$$= \left(\frac{2}{\sqrt{2^2+1}}, \frac{1}{\sqrt{2^2+1}} \right) = \left(\frac{2}{\sqrt{5}}, \frac{1}{\sqrt{5}} \right)$$

$$= (0.8944 \quad 0.4472)$$

4.3)

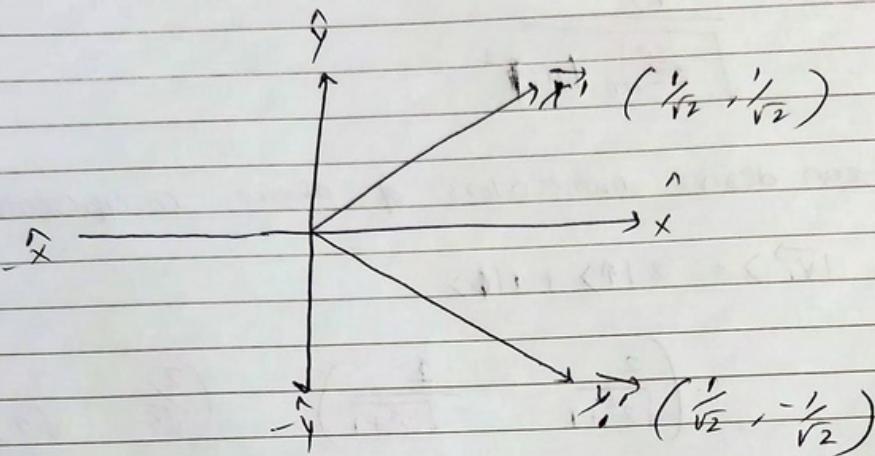
$$|V_1\rangle = 2|1\rangle + (1+i)|V\rangle$$

$$= \left(\frac{2}{\sqrt{2^2+(1+i)^2}}, \frac{1+i}{\sqrt{2^2+(1+i)^2}} \right) = \left(\frac{2}{\sqrt{4+2i}}, \frac{1+i}{\sqrt{4+2i}} \right)$$

$$= (0.816 + 0.4082i \quad 0.4082 + 0.4082i)$$

$$4.5) \quad \text{Given} \quad |\vec{x}\rangle = \frac{1}{\sqrt{2}}(|\hat{x}\rangle + |\hat{y}\rangle)$$

$$|\vec{y}\rangle = \frac{1}{\sqrt{2}}(|\hat{x}\rangle - |\hat{y}\rangle)$$



$$\langle \vec{x} | \vec{x} \rangle = \left(\frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}} \right) \left(\begin{matrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{matrix} \right) = \frac{1}{2} + \frac{1}{2} = 1 \text{ (normalized)}$$

$$\langle \vec{y} | \vec{y} \rangle = \left(\frac{1}{\sqrt{2}}, -\frac{1}{\sqrt{2}} \right) \left(\begin{matrix} \frac{1}{\sqrt{2}} \\ -\frac{1}{\sqrt{2}} \end{matrix} \right) = \frac{1}{2} + \frac{1}{2} = 1 \text{ (Normalized)}$$

$$\langle \vec{x} | \vec{y} \rangle = \left(\frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}} \right) \left(\begin{matrix} \frac{1}{\sqrt{2}} \\ -\frac{1}{\sqrt{2}} \end{matrix} \right) = \frac{1}{2} - \frac{1}{2} = 0 \text{ (orthogonal)}$$

$$\langle \vec{y} | \vec{x} \rangle = \left(\frac{1}{\sqrt{2}}, -\frac{1}{\sqrt{2}} \right) \left(\begin{matrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{matrix} \right) = \frac{1}{2} + \frac{1}{2} = 0 \text{ (orthogonal)}$$

It is satisfying both orthogonality and normalization

If satisfies orthonormality

```
#Problem 4
v1 = np.array([[2],[1]])
print(v1)
v2 = v1/np.linalg.norm(v1)
print(v2)

v3 = np.array([[2],[1+1j]])
print(v3)
v4 = v3/np.linalg.norm(v3)
print(v4)
```

```
[[2]
 [1]]
[[0.89442719]
 [0.4472136 ]]
[[2.+0.j]
 [1.+1.j]]
[[0.81649658+0.j          ]
 [0.40824829+0.40824829j]]
```

Chapter 5

5.1) we know from 36 page.

$$|0\rangle = \frac{1}{\sqrt{2}} |+\rangle + \frac{1}{\sqrt{2}} |- \rangle \quad \& \quad |1\rangle = \frac{1}{\sqrt{2}} |+\rangle - \frac{1}{\sqrt{2}} |- \rangle$$

①

$$\text{we know } |+\rangle = \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle) \quad \& \quad |- \rangle = \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

②

Given $|v_1\rangle = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$ in $|+\rangle / |- \rangle$ basis

$$|v_1\rangle = |+\rangle + 2|-\rangle$$

$$\text{we get } |0\rangle = \frac{1}{\sqrt{2}} (|+\rangle + 2|-\rangle) \quad |1\rangle = \frac{1}{\sqrt{2}} (|+\rangle - 2|-\rangle)$$

③

converting back by substituting ③ in ②

$$|+\rangle = \frac{1}{\sqrt{2}} \left(\frac{1}{\sqrt{2}} (2|+\rangle) \right) = 1 \cdot \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

$$|- \rangle = \frac{1}{\sqrt{2}} \left(\frac{1}{\sqrt{2}} (4|-\rangle) \right) = 2$$

5.2)

$$\text{Given } |\vec{v}_1\rangle = \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

Assumed it is in $|0\rangle, |1\rangle$ basis

$$|\vec{v}_1\rangle = |10\rangle + 2|11\rangle$$

Converting it to $|+\rangle, |-\rangle$ basis

$$= \left(\frac{1}{\sqrt{2}} |+\rangle + \frac{1}{\sqrt{2}} |-\rangle \right) + 2 \left(\frac{1}{\sqrt{2}} |+\rangle - \frac{1}{\sqrt{2}} |-\rangle \right)$$

$$= \frac{3}{\sqrt{2}} |+\rangle - \frac{1}{\sqrt{2}} |-\rangle$$

Normalized values are

$$= \begin{pmatrix} \frac{3}{\sqrt{2}} \\ \left(\frac{3}{\sqrt{2}}\right)^2 + \left(-\frac{1}{\sqrt{2}}\right)^2 \end{pmatrix}, \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \left(\frac{3}{\sqrt{2}}\right)^2 + \left(-\frac{1}{\sqrt{2}}\right)^2 \end{pmatrix}$$

$$= (0.948, 0.316)$$

Probability is positive as it can never be negative

5.3) Probability to find $|1+\rangle$ is $(0.948, 0.316)$ is 0.948

5.4) Probability to find $|0\rangle$ in $|0\rangle/|1\rangle$ basis

$$|\tilde{v_1}\rangle = |0\rangle + |1-\rangle$$

$$= \left(\frac{1}{\sqrt{1+4}}, \frac{2}{\sqrt{1+4}} \right) = \left(\frac{1}{\sqrt{5}}, \frac{2}{\sqrt{5}} \right)$$

$$= (0.447, 0.894)$$

Probability of finding $|0\rangle$ is 0.447

```
#Problem 5.2
v1 = np.array([[3/np.sqrt(2)],[1/np.sqrt(2)]])
print(v1)
v2 = v1/np.linalg.norm(v1)
print(v2)

# Problem 5.4
v3 = np.array([[1] , [(2 )]])
print(v3)
v4 = v3/np.linalg.norm(v3)
print(v4)
```

```
[[2.12132034]
 [0.70710678]]
[[0.9486833 ]
 [0.31622777]]
[[1]
 [2]]
[[0.4472136 ]
 [0.89442719]]
```

Chapter 6

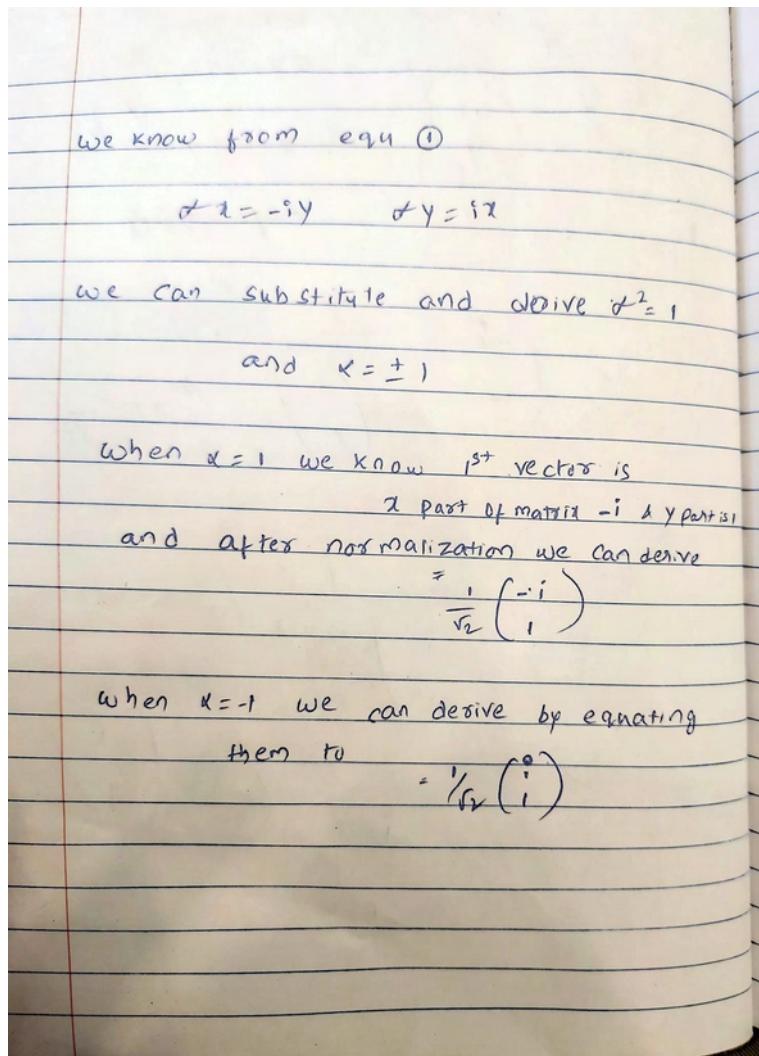
6.1
sp1Given operator σ_y matrix

$$\sigma_y \text{ |eigen vector } \rangle = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}$$
$$= \begin{pmatrix} -iy \\ ix \end{pmatrix}$$

we can assume eigen value as α

$$\begin{pmatrix} -iy \\ ix \end{pmatrix} = \alpha \begin{pmatrix} x \\ y \end{pmatrix} \quad \boxed{1}$$
$$\alpha x = -iy \quad \alpha y = ix$$

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```
#Lecture_6-Problems
import numpy as np
sigmay = np.array([[0, -1j], [1j, 0]])
print(sigmay)
eigenvalues,eigenvectors = np.linalg.eig(sigmay)
print(eigenvalues)
print(eigenvectors)
```



```
[[ 0.+0.j -0.-1.j]
 [ 0.+1.j  0.+0.j]
 [ 1.+0.j -1.+0.j]
 [[-0.         -0.70710678j  0.70710678+0.j        ]
 [ 0.70710678+0.j       0.         -0.70710678j]]]
```

Chapter 7

problems

$$7.1) \quad \sigma_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$$

$$\text{To find. } \sigma_y^2 = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$$

$$= \begin{pmatrix} 0 - i^2 & 0 + 0 \\ 0 + 0 & -i^2 + 0 \end{pmatrix} \quad (\because i^2 = -1)$$

$$= \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = I$$

$$\sigma_y^2 = I$$

$$7.2) \quad \sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

$$\sigma_z^2 = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

$$= \begin{pmatrix} 1 * 1 + 0 & 0 + 0 \\ 0 + 0 & 0 + (-1) \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = I$$

$$\sigma_z^2 = I$$

7.3) σ_y is Hermitian

we know $\sigma_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$

Firstly Conjugation

$$\sigma_y^* = \begin{pmatrix} 0 & i \\ -i & 0 \end{pmatrix}$$

Transpose

$$\sigma_y^{*T} = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$$

$$\sigma_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} = \sigma_y^{*T}$$

7.4) $\langle 0 | \sigma_y$

$$(1 \ 0) \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} = \begin{bmatrix} ix_0 + 0x_1 & ix_1 - i + 0x_0 \end{bmatrix}$$

$$\begin{pmatrix} 0 & -i \end{pmatrix}$$

It is in bra form and it is equal to $\langle 0 | \sigma_y | 0 \rangle$ only if we do transpose and complex conjugate to $(\langle 0 | \sigma_y)^* \text{ i.e., } (\langle 0 | \sigma_y)^{T*} = \langle 0 | \sigma_y | 0 \rangle$

```
#Lecture_7-Problems
#7.1
sigmay=np.array([[0,-1j],[1j,0]])
print(sigmay)
print(np.matmul(sigmay, sigmay))

#7.2
sigmaz=np.array([[1,0],[0,-1]])
print(sigmaz)
print(np.matmul(sigmaz, sigmaz))

#7.3

sigmayh = np.conjugate(np.transpose(sigmay))
print(sigmayh)

#7.4
zerospin=np.array([[1],[0]])
print(zerospin.shape)
print(np.dot(sigmay,zerospin))
```

↳ [[0.+0.j -0.-1.j]
 [0.+1.j 0.+0.j]]
 [[1.+0.j 0.+0.j]
 [0.+0.j 1.+0.j]]
 [[1 0]
 [0 -1]]
 [[1 0]
 [0 1]]
 [[0.-0.j 0.-1.j]
 [-0.+1.j 0.-0.j]]
 (2, 1)
 [[0.+0.j]
 [0.+1.j]]

Chapter 8

```
#Lecture_8-Problems

import numpy as np
sigmax=np.array([[0,1],[1,0]])
print(sigmax)
eigenvalues, eigenvectors=np.linalg.eig(sigmax)
print(eigenvalues)
print(eigenvectors)

sigmaxp=np.array([1/2,1/2])
sigmaxn=np.array([1/2,-1/2])
eigenvalues, eigenvectors=np.linalg.eig(sigmax)
print(eigenvalues)
print(eigenvectors)
print(np.dot(sigmaxp,sigmaxp))
print(np.dot(sigmaxn,sigmaxn))
```

```
↳ [[ 0.+0.j -0.-1.j]
     [ 0.+1.j  0.+0.j]]
[[1.+0.j 0.+0.j]
 [0.+0.j 1.+0.j]]
[[ 1  0]
 [ 0 -1]]
[[1 0]
 [0 1]]
[[ 0.-0.j  0.-1.j]
 [-0.+1.j  0.-0.j]]
(2, 1)
[[0.+0.j]
 [0.+1.j]]
```

8.2) Given σ_z

we know $\sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$ it is ket notation
 and even if find adjoint to get bra it is same
 so it is Hermitian matrix

we know project operator is

$$P_{\sigma_z} = |\sigma_z\rangle \langle \sigma_z|$$

$$\begin{aligned} &= \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \\ &= \begin{pmatrix} 1+0 & 0+0 \\ 0+0 & 0+1 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \end{aligned}$$

$$I = P_{\sigma_z}$$

E.3)

Given

eigen vectors of σ_x

$$|+\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} \quad |- \rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ -1 \end{pmatrix}$$

we know for probability

$$P_i = |\langle i | i \rangle|^2$$

probability of $|+\rangle$ ~~is~~ \Rightarrow

$$= \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} \cdot \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \end{pmatrix}$$

$$= \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$

so it is 50% & similarly for $|-\rangle$

is

$$= \frac{1}{2} \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$$

```
#8.4
sigmaxp=np.array([1/np.sqrt(2),1/np.sqrt(2)])
sigmaxn=np.array([1/np.sqrt(2),-1/np.sqrt(2)])
eigenvalues, eigenvectors=np.linalg.eig(sigmax)
print(eigenvalues)
print(eigenvectors)
print(np.dot(sigmaxp,sigmaxp))
print(np.dot(sigmaxn,sigmaxn))
```

.

result

```
0.9999999999999998
0.9999999999999998
```

Chapter 9

q.1 sol

we know from eq 9.15

$$\text{that } U^\dagger U = I$$

we know from page no 24 we can use
eq 4.1 $\hat{v}_i \cdot \hat{v}_j = \begin{cases} 1 & \text{if } i=j \\ 0 & \text{if } i \neq j \end{cases}$

so when we consider a row in 9.15
we can say if it is in same row like
example $b_{0,0}$ is 0th row it is 1 etc

q.2) Given q.2 used it in (i) (ii) (iii)

$$\begin{vmatrix} a_{0,0} - \lambda_i & a_{0,1} \\ a_{1,0} & a_{1,1} - \lambda_i \end{vmatrix}$$

i) $\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$

$$= \begin{vmatrix} 0 - \lambda_i & 1 \\ 1 & 0 - \lambda_i \end{vmatrix} = 0 \quad +\lambda_i^2 - 1 = 0$$

$$\boxed{\lambda_i^2 = \pm 1}$$

ii) $\sigma_y = \begin{pmatrix} 0 & i \\ -i & 0 \end{pmatrix}$

$$= \begin{vmatrix} 0 - \lambda_i & i \\ -i & 0 - \lambda_i \end{vmatrix} = 0 \Rightarrow \lambda_i^2 - (i^2) = 0$$

$$\lambda^2 = -1$$

$$\boxed{\lambda = \pm 1}$$

iii) $\sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$

$$= \begin{vmatrix} 1 - \lambda_i & 0 \\ 0 & -1 - \lambda_i \end{vmatrix} = 0 \Rightarrow -(1 - \lambda^2) = 0$$

$$\lambda^2 = 1$$

Eigen values are same for all

```
▶ #Lecture_9-Problems  
#9.3  
# M=np.array([[1,0],[0,-1]])  
sigmax=np.array([[0,1],[1,0]])  
sigmay=np.array([[0,-1j],[1j,0]])  
sigmaz=np.array([[1,0],[0,-1]])  
  
print(np.linalg.det(sigmax))  
print(np.linalg.det(sigmay))  
print(np.linalg.det(sigmaz))
```

```
⇨ -1.0  
(-1+0j)  
-1.0
```

9.4)

for example we can take

$$A \text{ as } \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \text{ i.e., } \sigma_x$$

We know pauli matrices are unitary
we can prove it now

$$|V_1\rangle = \frac{1}{\sqrt{2}} |A\rangle + \frac{1}{\sqrt{2}} |B\rangle$$

$$|V_1\rangle = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ -\frac{1}{\sqrt{2}} \end{pmatrix}$$

$$A|V_1\rangle = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} \frac{1}{\sqrt{2}} \\ -\frac{1}{\sqrt{2}} \end{pmatrix} = \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{pmatrix}$$

We can prove orthonormality by

Adding squares of coefficients of transformed
 $A|V_1\rangle$ and inner product with other matrices will be zero

$$\left(-\frac{1}{\sqrt{2}}\right)^2 + \left(\frac{1}{\sqrt{2}}\right)^2 = 1$$

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Chapter 10

(10.1)

Transformation matrix from basis σ_y to σ_x

We know from problem 6.1

that eigen vectors of σ_y are $\begin{pmatrix} i \\ -i \end{pmatrix}, \begin{pmatrix} i \\ i \end{pmatrix}$

$$|+\rangle_{\sigma_y} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ i \end{pmatrix} \quad |-\rangle_{\sigma_y} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ -i \end{pmatrix}$$

$$u = (x_{\text{new}}, y_{\text{old}}, *) = \begin{bmatrix} \frac{1}{2}(-i, i)^T & \frac{1}{2}(i, i)^T \\ \frac{1}{2}(i, i)^T & \frac{1}{2}(i, -i)^T \end{bmatrix}$$

$$= \left[\frac{1}{2} \begin{bmatrix} i+1 & i-1 \\ -i+1 & -i-1 \end{bmatrix} \right]$$

10.2) Given in page 87

complete basis is defined by outer product of each vectors sums to Identity matrix
i.e. outer product is nothing but ket \times bra

$$\sum_{i=0}^{n-1} |i\rangle \langle i| = |-\rangle \langle -| + |+\rangle \langle +| \\ \text{By deriving, } \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} + \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = I$$

10.3) For σ_z we know

$$|-\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ -1 \end{pmatrix} \quad |+\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

We know from 10.19

$$\sigma_z = \sum_{i=0}^{2-1} \lambda_i |i\rangle \langle i| \\ = 1 \left(\frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} \right) \left(\frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} \right)^* + -1 \left(\frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ -1 \end{pmatrix} \right) \left(\frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ -1 \end{pmatrix} \right)^*$$

$$= \left(\frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \right) - \left(\frac{1}{2} \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix} \right) \\ \sigma_z = \frac{1}{2} \begin{pmatrix} 0 & 2 \\ 2 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

#Lecture_10-Problems

#10.4

```
import numpy as np
sigmaxn=np.array([1/np.sqrt(2),-1])
sigmaxp=np.array([1/np.sqrt(2),1])
sigmayp=np.array([1/np.sqrt(2),+1j])
sigmayn=np.array([1/np.sqrt(2),-1j])

print(np.matmul(sigmaxn, sigmayn))
print(np.matmul(sigmaxn, sigmayp))
print(np.matmul(sigmaxp, sigmayn))
print(np.matmul(sigmaxp, sigmayp))
```

```
(0.4999999999999999+1j)
(0.4999999999999999-1j)
(0.4999999999999999-1j)
(0.4999999999999999+1j)
```

Chapter 11

$$11.1) \text{ Given } |b\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \quad |g\rangle_2 = \begin{pmatrix} i/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix}$$

From 11.11 & 11.14

we know

$$\begin{pmatrix} \alpha_3 \\ \beta_3 \\ \gamma_3 \\ \delta_3 \end{pmatrix} = \begin{pmatrix} \alpha_1 \alpha_2 \\ \alpha_1 \beta_2 \\ \beta_1 \alpha_2 \\ \beta_1 \beta_2 \end{pmatrix}$$

i) $|F\rangle_1 \otimes |g\rangle_2$

$$\begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} i/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ i/\sqrt{2} \\ i/\sqrt{2} \end{pmatrix}$$

ii) $|g\rangle_1 \otimes |f\rangle_2$

$$\begin{pmatrix} i/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ i/\sqrt{2} \\ 0 \\ i/\sqrt{2} \end{pmatrix}$$

Both are not same

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11.2)

$$\begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$= \begin{pmatrix} 0 \\ 0 \\ \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$\left(\begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ \frac{1}{\sqrt{2}} \\ 0 \\ \frac{i}{\sqrt{2}} \end{array} \right)$$

11.3) we need to prove distribution rule

$$11.20 \quad (|t\rangle + |e\rangle) \otimes |g\rangle = |t\rangle \otimes |g\rangle + |e\rangle \otimes |g\rangle$$

$$\left(\begin{pmatrix} 0 \\ 1 \end{pmatrix} + \begin{pmatrix} 1 \\ 0 \end{pmatrix} \right) \otimes \begin{pmatrix} 1/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \begin{pmatrix} 1/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix} + \begin{pmatrix} 1 \\ 0 \end{pmatrix} \begin{pmatrix} 1/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix}$$

$$\begin{pmatrix} 1/\sqrt{2} \\ 1/\sqrt{2} \\ 1/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 1/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix} + \begin{pmatrix} 1/\sqrt{2} \\ 1/\sqrt{2} \\ 0 \\ 0 \end{pmatrix}$$

$$= \begin{pmatrix} 1/\sqrt{2} \\ 1/\sqrt{2} \\ 1/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix} = \begin{pmatrix} 1/\sqrt{2} \\ 1/\sqrt{2} \\ 1/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix} \quad L.H.S = R.H.S$$

11.4)

```
#Lecture_11-Problems

#11.4
#proving 11.1
import numpy as np
g=np.array([[0],[1]])
h=np.array([[1/np.sqrt(2)],[1j/np.sqrt(2)]])
print(np.kron(g,h))
print(np.kron(h,g))

# proving 11.2
x = np.kron(h,g)
print(np.kron(g,x))

#proving 11.3
y=np.array([[1],[0]])
z = (g + y)
print(np.kron(z,h))
a = np.kron(z,h)
```

```
l = np.kron(g,h)
m = np.kron(y,h)
b = l+m
print(l+m)

if a.all() == b.all() :
    print(" L.H.S == R.H.S")
```

```
[[0.        +0.j        ]
 [0.        +0.j        ]
 [0.70710678+0.j      ]
 [0.        +0.70710678j]]
 [[0.        +0.j        ]
 [0.70710678+0.j      ]
 [0.        +0.j        ]
 [0.        +0.70710678j]]
 [[0.        +0.j        ]
 [0.        +0.j        ]
 [0.        +0.j        ]
 [0.        +0.j        ]
 [0.70710678+0.j      ]
 [0.        +0.j        ]
 [0.        +0.70710678j]]
 [[0.        +0.j        ]
 [0.        +0.j        ]
 [0.        +0.j        ]
 [0.        +0.j        ]
 [0.        +0.70710678j]
 [0.        +0.70710678j]]
 [[0.70710678+0.j      ]
 [0.        +0.70710678j]
 [0.70710678+0.j      ]
 [0.        +0.70710678j]]
 [[0.70710678+0.j      ]
 [0.        +0.70710678j]
 [0.70710678+0.j      ]
 [0.        +0.70710678j]]
 L.H.S == R.H.S
```

ii.5)

$$\text{Given } n_1 = n_2 = 1 \text{ & } |f\rangle = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$$|e\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}, |g\rangle = \begin{pmatrix} i \\ 1 \end{pmatrix} \text{ and } |k\rangle = \begin{pmatrix} i \\ 1 \end{pmatrix}$$

eqn ii.21

$$\langle h_1 | h_2 \rangle = \left\langle \sum_{j=1}^{n_1} |f_j\rangle \otimes |g_j\rangle \middle| \sum_{l=1}^{n_2} |e_l\rangle \otimes |k_l\rangle \right\rangle$$

$$= \sum_{j=1}^{n_1} \sum_{l=1}^{n_2} \langle f_j | e_l \rangle \langle g_j | k_l \rangle$$

$$\stackrel{\because n_1 = n_2 = 1}{=} \left\langle \begin{pmatrix} 1 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} i \\ 1 \end{pmatrix} \middle| \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} i \\ 1 \end{pmatrix} \right\rangle$$

$$= (1|i; i|1) \begin{pmatrix} 0 \\ 0 \\ i \\ 1 \end{pmatrix} \stackrel{(\because \text{It is inner product})}{=} i - 1$$

R.H.S

$$= (1|1) \begin{pmatrix} 0 \\ 1 \end{pmatrix} \cdot (1|i) \begin{pmatrix} i \\ 1 \end{pmatrix}$$

$$= (0+1) (i-1)$$

$$= i - 1$$

Chapter 12

12.1

So,

$$\text{Given } |f\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \text{ and } |g\rangle = \begin{pmatrix} i/\sqrt{2} \\ i/\sqrt{2} \end{pmatrix}$$

Initially

$$\sigma_z |t_0\rangle,$$

$$= \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$\sigma_z |g\rangle_2 = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} i/\sqrt{2} \\ -i/\sqrt{2} \end{pmatrix} = \begin{pmatrix} i/\sqrt{2} \\ i/\sqrt{2} \end{pmatrix}$$

$$\sigma_z |f\rangle_1 \otimes \sigma_z |g\rangle_2$$

$$= \begin{pmatrix} 1 \\ 0 \end{pmatrix} \begin{pmatrix} i/\sqrt{2} \\ i/\sqrt{2} \end{pmatrix} = \begin{pmatrix} i/\sqrt{2} \\ i/\sqrt{2} \\ 0 \\ 0 \end{pmatrix}$$

Now to find $|t\rangle_1 \otimes |g\rangle_2$

$$= \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} i/\sqrt{2} \\ i/\sqrt{2} \end{pmatrix} \Rightarrow \begin{pmatrix} 0 \\ 0 \\ i/\sqrt{2} \\ i/\sqrt{2} \end{pmatrix}$$

$$\sigma_x \otimes \sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$= \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

$$\sigma_z \otimes \sigma_z (|tt\rangle \otimes |gg\rangle)$$

$$= \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 1/\sqrt{2} \\ 1/\sqrt{2} \end{pmatrix}$$

$$= \begin{pmatrix} 1/\sqrt{2} \\ 1/\sqrt{2} \\ 0 \\ 0 \end{pmatrix}$$

So it is equal to $\sigma_z |tt\rangle \otimes \sigma_z |gg\rangle$

$$\sigma_z |tt\rangle \otimes \sigma_z |gg\rangle = \sigma_z \otimes \sigma_z (|tt\rangle \otimes |gg\rangle)$$

Chapter 13

13.1

As quantum computing is known for exponential speed and for quantum parallelism world can store up to 21021 quantum states due to superposition.

13.2)

Here $|\phi^+\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 1 \end{pmatrix}$ is proved already

$$|\phi^-\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 0 \\ 0 \\ -1 \end{pmatrix}$$

$$= \frac{1}{\sqrt{2}} (|100\rangle + |010\rangle + |001\rangle - |111\rangle)$$

we can assume $|\phi^+\rangle = |m\rangle \otimes |n\rangle$

$$|m\rangle = a_0|0\rangle + a_1|1\rangle \quad \& \quad |n\rangle = b_0|0\rangle + b_1|1\rangle$$

$$\text{Given } 13.11 \rightarrow |\phi^+\rangle = a_0b_0|100\rangle + a_0b_1|101\rangle + a_1b_0|010\rangle + a_1b_1|111\rangle$$

$$\text{Similar to } 13.12 \quad a_0b_0=1, \quad a_0b_1=0, \quad a_1b_0=0 \\ a_1b_1=-1$$

Similarly $a_0b_0=0$ (or) $a_1b_1=0$ should be here

So $|\phi^+\rangle$ is entangled as our assumption is false

$$|\Psi^+\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}$$

$$= \frac{1}{\sqrt{2}} (|100\rangle + |101\rangle + |110\rangle + |111\rangle)$$

$$|\Psi^-\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 \\ 1 \\ -1 \\ 0 \end{pmatrix}$$

$$= \frac{1}{\sqrt{2}} (|0100\rangle + |101\rangle - |110\rangle + |011\rangle)$$

From

.. 13.2

$$a_0 b_0 = 0, a_0 b_1 = 1$$

$$a_1 b_0 = 0, a_1 b_1 = 0$$

From 13.2

$$a_0 b_0 = 0, a_0 b_1 = 1, a_1 b_0 = 0$$

$$a_1 b_1 = 0$$

we can assume

$a_0 b_0 = 0$ either of

them should be zero

it denies our above

we can assume $a_0 b_0 = 0$

and one of them should be

zero it denies calculation

calculation

13.3

Yes, I believe entanglement is concept ~~isn't~~

sol of distance we can expect the result.

It will be profitting for him

13.4)

Here as mentioned basis states are orthonormal and if we have the same state in both two spaces then the inner product will be 1

Chapter 14

14.1

14.2

$$\text{Given } \frac{1}{2} (|00\rangle + |01\rangle + |10\rangle + |11\rangle)$$

$$\text{From 14.12} \Rightarrow \frac{1}{2} |0\rangle\langle 0| + \frac{1}{2} |0\rangle\langle 1|$$

$$+ \frac{1}{2} |1\rangle\langle 0| + \frac{1}{2} |1\rangle\langle 1|$$

$$\rho = \frac{1}{2} \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} (|0\rangle\langle 0|) +$$

$$+ \frac{1}{2} \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} (|0\rangle\langle 1|) + \frac{1}{2} \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} (|1\rangle\langle 0|)$$

$$+ \frac{1}{2} \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} (|1\rangle\langle 1|)$$

$$= \frac{1}{2} \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} + \frac{1}{2} \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} + \frac{1}{2} \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix}$$

$$+ \frac{1}{2} \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$

$$\rho = \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$

As we know entanglement entropy is zero

from previous chapter that we know our given

state is not product of two subspaces

and it should be zero

14.2)

Given $\rho = \frac{1}{2} \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & -1 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$ $\rightarrow 14.15$

* We know $\rho_B = \text{Tr}_A(\rho)$

Figure 14.2 for tracing second qubit

$$= \begin{pmatrix} (0) 0 & (0) 0 \\ 0 (1/2) & -1/2 (0) \\ (0) -1/2 & (1/2) 0 \\ 0 (0) & 0 (0) \end{pmatrix}$$

Above same shape elements can be added

From above tracing B space

$$= \begin{bmatrix} 0 & x & 0 & x \\ x & 1/2 & x & 0 \\ 0 & x & 1/2 & x \\ x & 0 & x & 0 \end{bmatrix}$$

Chapter 15

15.1)

we know

from 15.6

$$U_{NOT} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

To prove it as unitary we know

$$U^\dagger \cdot U = I$$

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \text{ should be } I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$\begin{pmatrix} 0 \times 0 + 1 \times 1 & 0 \times 0 + 1 \times 0 \\ 1 \times 0 + 0 \times 1 & 1 \times 1 + 0 \times 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

Hence proved

15.2)

we know from 15.15

$$\text{CNot is } U = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

from 9.18 we know

$$U^\dagger U = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$= \begin{pmatrix} 1 \times 1 + 0 \times 0 + 0 \times 0 + 0 \times 0 & 0 & 0 & 0 \\ 0 & 0 \times 0 + 1 \times 1 + 0 \times 0 + 0 \times 0 & 0 & 0 \\ 0 & 0 & 0 \times 0 + 0 \times 0 + 1 \times 0 + 0 \times 0 & 0 \\ 0 & 0 & 0 & 0 \times 0 + 0 \times 0 + 1 \times 1 + 0 \times 0 \end{pmatrix}$$

Hence proved

$$= I = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

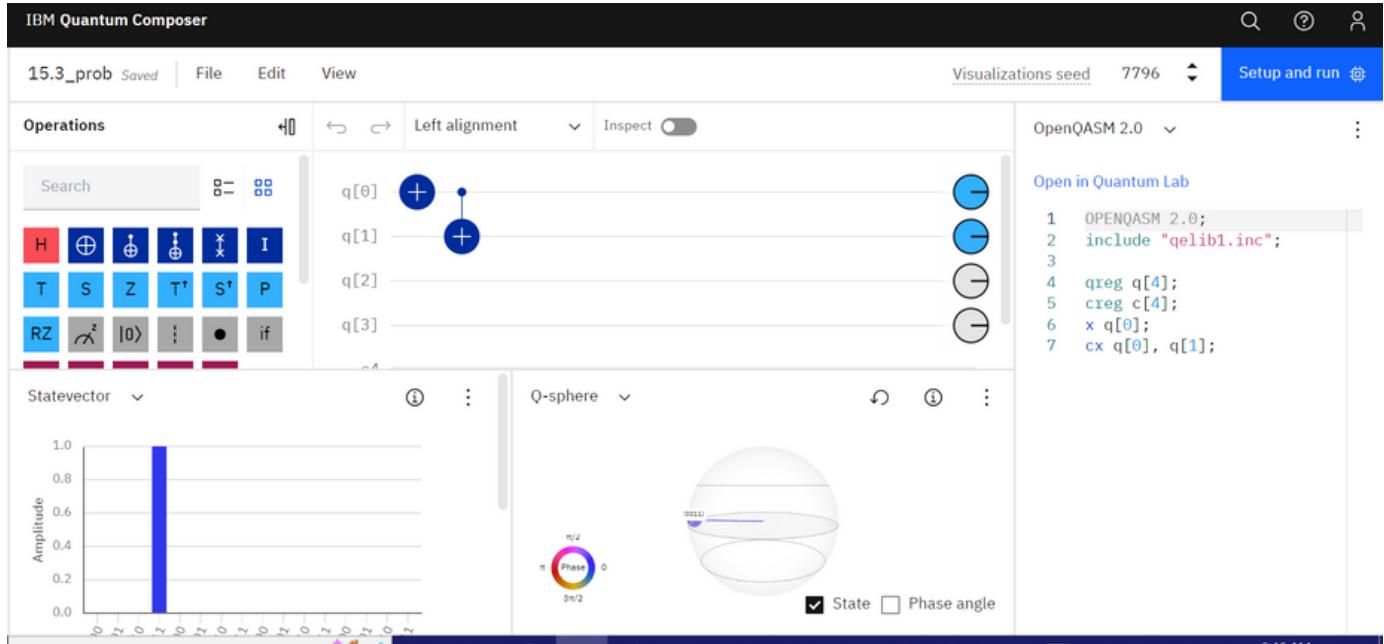
```
#Chapter 15

#15.1
import numpy as np
sigmay=np.array([[0,1],[1,0]])    #Not gate and transpose is same for this matrix
print(sigmay)
print(np.matmul(sigmay, sigmay))

#15.2
sigmac=np.array([[1,0,0,0],[0,1,0,0],[0,0,0,1],[0,0,1,0]])   #CNot gate and transpose is same for this matrix
print(sigmac)
print(np.matmul(sigmac, sigmac))
```

```
[[0 1]
 [1 0]]
[[1 0]
 [0 1]]
[[0 1]
 [1 0]
 [0 0 0 0]
 [0 1 0 0]
 [0 0 0 1]
 [0 0 1 0]
 [1 0 0 0]
 [0 1 0 0]
 [0 0 1 0]
 [0 0 0 1]]
```

15.3)



Introduction to Quantum computing (From a layperson to programmer in 30 steps)

Chapter 16

16.1)

16.1) we know Z-gate is $\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$

and $|-\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ -1 \end{pmatrix}$ $|+\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$

when we apply Z gate to $|-\rangle$

$$Z|-\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} 1 \\ -1 \end{bmatrix} \rightarrow \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \times 1 + 0 \times -1 \\ 0 \times 1 - 1 \times -1 \end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = |+\rangle$$

16.2)

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

(6.2)

we know control phase shift gate is

$$U_{CPS, \phi} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{i\phi} \end{bmatrix}$$

In z-axis when angle is π
we can assume $e^{i\phi} = \cos\phi + i\sin\phi = \cos\pi + i\sin\pi = -1$

controlled z gate is $\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix}$

It is in $|10\rangle/|11\rangle$ basis

in order to prove it is CNOT in $|+\rangle|-\rangle$ we need to convert
it \Rightarrow we know $|+\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$ $|-\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$

we know transformation matrix from $|10\rangle/|11\rangle$ to $|+\rangle|-\rangle$ is
(Hadamard gate)

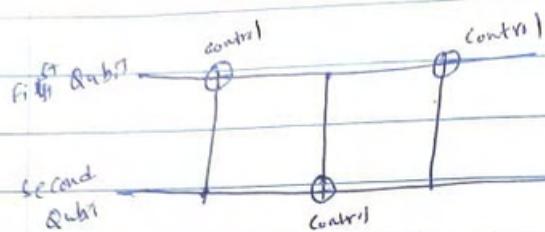
$$U = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

when we apply the above matrix to convert it to $|+\rangle|-\rangle$ is

$$(Z_{|+\rangle|-\rangle}) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

so hence C_Z in $|+\rangle|-\rangle$ is CNOT gate

16.3) Here I have analyzed 3 CNOT gates usage building swap gate



Here 1st qubit converts (First Not gate)

1st CNot gate CNot₁ ($|00\rangle = |00\rangle$, $|01\rangle = |01\rangle$, $|10\rangle = |11\rangle$)

control is with 2nd qubit

2nd CNot gate CNot₂ ($|00\rangle = |00\rangle$, $|01\rangle = |11\rangle$, $|10\rangle = |01\rangle$)

control is with 1st qubit

3rd CNot gate CNot₃ ($|00\rangle = |00\rangle$, $|11\rangle = |10\rangle$, $|01\rangle = |01\rangle$)

we can see swap is done and it is proved

For CNot 2 using state representation we can do

$$CNot_2 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} \quad \begin{array}{l} |00\rangle \\ |11\rangle \\ |01\rangle \\ |10\rangle \end{array}$$

Chapter 17

17.1)

17.1

Given

$$H|\Psi\rangle = \alpha|1+\rangle + \beta|1-\rangle \quad \text{From 17.4}$$

Normally we can directly understand that
 co-efficient of $|0+\rangle$ & $|1+\rangle$ basis is from 17.4
 $\alpha \frac{1}{\sqrt{2}}(\alpha+\beta)$ & $\beta \frac{1}{\sqrt{2}}(\alpha-\beta)$ from understanding

we can even derive from 17.4

$$= \alpha|1+\rangle + \beta|1-\rangle$$

$$\text{we know } |1+\rangle = \frac{1}{\sqrt{2}}(|0+\rangle + |1+\rangle), \quad |1-\rangle = \frac{1}{\sqrt{2}}(|0-\rangle - |1-\rangle)$$

$$= \alpha \left(\frac{1}{\sqrt{2}}(|0+\rangle + |1+\rangle) \right) + \beta \left(\frac{1}{\sqrt{2}}(|0-\rangle - |1-\rangle) \right)$$

$$= \frac{\alpha}{\sqrt{2}}|0+\rangle + \frac{\alpha}{\sqrt{2}}|1+\rangle + \frac{\beta}{\sqrt{2}}|0-\rangle - \frac{\beta}{\sqrt{2}}|1-\rangle$$

$$H|\Psi\rangle = \frac{1}{\sqrt{2}}(\alpha+\beta)|0+\rangle + \frac{1}{\sqrt{2}}(\alpha-\beta)|0-\rangle$$

17.2)

17.2

$$\text{Given } H^{\otimes n}|y\rangle = \frac{1}{2^{n/2}} \sum_{x=0}^{2^n-1} (-1)^{x \cdot y} |x\rangle$$

Here n means n qubit Hadamard gate

$$H^{\otimes 2}|z\rangle = \frac{1}{2^{2/2}} \sum_{x=0}^{2^2-1} (-1)^{x \cdot z} |x\rangle$$

$$= \frac{1}{2} \left((-1)^{0 \cdot 0}|0\rangle + (-1)^{1 \cdot 0}|1\rangle + (-1)^{2 \cdot 0}|2\rangle + (-1)^{3 \cdot 0}|3\rangle \right)$$

From 17.3 we have

$$0 \cdot 2 = (0 \text{AND} 1) \oplus (0 \text{AND} 0) = 0 \oplus 0 = 0$$

$$1 \cdot 2 = (0 \text{AND} 1) \oplus (1 \text{AND} 0) = 0 \oplus 0 = 0$$

$$2 \cdot 2 = (1 \text{AND} 1) \oplus (0 \text{AND} 0) = 1 \oplus 0 = 1$$

$$3 \cdot 2 = (1 \text{AND} 1) \oplus (1 \text{AND} 0) = 1 \oplus 0 = 1$$

$$= H^{\otimes 2}|z\rangle = \frac{1}{2} (|0\rangle + |1\rangle - |2\rangle - |3\rangle)$$

17.3)

17.3)

From 17.2 Example using 17.19

$$H^{\otimes 2} |12\rangle = H^{\otimes 2} |10\rangle$$

$$= (H|11\rangle) \otimes (H|10\rangle)$$

$$= \frac{1}{2^{2/2}} (|10\rangle + (-1)^1 |11\rangle) \otimes (|10\rangle + (-1)^0 |11\rangle)$$

$$= \frac{1}{2} (|10\rangle - |11\rangle) \otimes (|10\rangle + |11\rangle)$$

$$= \frac{1}{2} (|100\rangle + |101\rangle - |110\rangle - |111\rangle)$$

Here we proved same solution we achieved
in previous problems

17.4)

17.4

From 17.10 we know

$$H^{\otimes 2} = H \otimes H$$

$$= \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix}$$

 we need to find $H^{\otimes 2}|12\rangle$

$$\text{we know } |12\rangle = |10\rangle = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$

$$H^{\otimes 2}|10\rangle = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$

$$= \frac{1}{2} \begin{pmatrix} 1 \times 0 + 1 \times 0 + 1 \times 1 + 1 \times 0 \\ 1 \times 0 - 1 \times 0 + 1 \times 1 + 1 \times 0 \\ 1 \times 0 + 1 \times 0 - 1 \times 1 - 1 \times 0 \\ 1 \times 0 - 1 \times 0 - 1 \times 1 + 1 \times 0 \end{pmatrix}$$

$$= \frac{1}{2} \begin{pmatrix} 1 \\ 1 \\ -1 \\ -1 \end{pmatrix}$$

$$= \frac{1}{2} (|100\rangle + |101\rangle - |110\rangle - |111\rangle)$$

Hence proved

Chapter 18

18.1)

18.1)

$$|\psi\rangle = \alpha|100\rangle + \beta|101\rangle + \gamma|110\rangle + \delta|111\rangle$$

$$|\phi_1\rangle = \alpha|100\rangle + \beta|111\rangle + \gamma|110\rangle + \delta|101\rangle$$

LSB

$$|\phi_2\rangle = \alpha|100\rangle + \beta|110\rangle + \gamma|111\rangle + \delta|101\rangle$$

$$|\phi\rangle = \alpha|100\rangle + \beta|110\rangle + \gamma|101\rangle + \delta|111\rangle$$

MSB

So it does not have any effect even if we implement C NOT gates starting from LSB or MSB.

 $U_{x\oplus, \text{right}}, U_{x\oplus, \text{center}}, U_{x\oplus, \text{left}}$

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

 U_{swap}

18.2)

18.2)

Given LSB as controlling bit from 15, 13
we can derive

$$U_{XOR}|00\rangle = |0 \oplus 0, 0\rangle = |0, 0\rangle = |00\rangle$$

$$U_{XOR}|01\rangle = |0 \oplus 1, 0\rangle = |1, 0\rangle = |01\rangle$$

$$U_{XOR}|10\rangle = |1 \oplus 0, 0\rangle = |1, 0\rangle = |10\rangle$$

$$U_{XOR}|11\rangle = |1 \oplus 1, 1\rangle = |0, 1\rangle = |01\rangle$$

we can derive matrix as

$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

It matches our assumed value in previous problem

18.3)

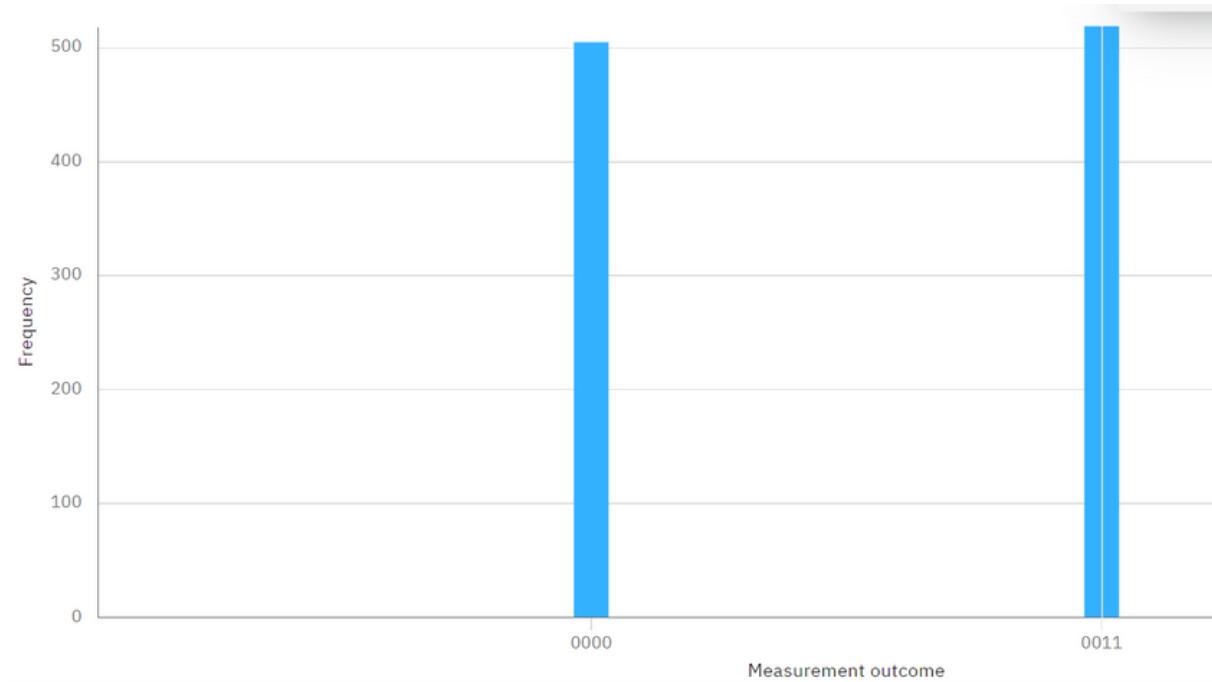
The screenshot shows the IBM Quantum Composer interface. At the top, it says "IBM Quantum Composer" and "18.3Exercise". The main area displays a quantum circuit with two qubits (q[0] and q[1]) and one classical register (c[4]). The circuit consists of an identity gate (I) on q[0], followed by a CNOT gate with control on q[0] and target on q[1]. Below the circuit, there are two probability plots. The left plot, titled "Probabilities", shows the probability of each computational basis state from 0000 to 1111. The right plot, titled "Q-sphere", shows the state vector on aBloch sphere. On the right side of the interface, the QASM code is displayed:

```

OPENQASM 2.0;
include "qelib1.inc";
qreg q[2];
creg c[4];
id q[0];
h q[1];
cx q[1], q[0];
measure q[0] -> c[0];
measure q[1] -> c[1];

```

Simulation output

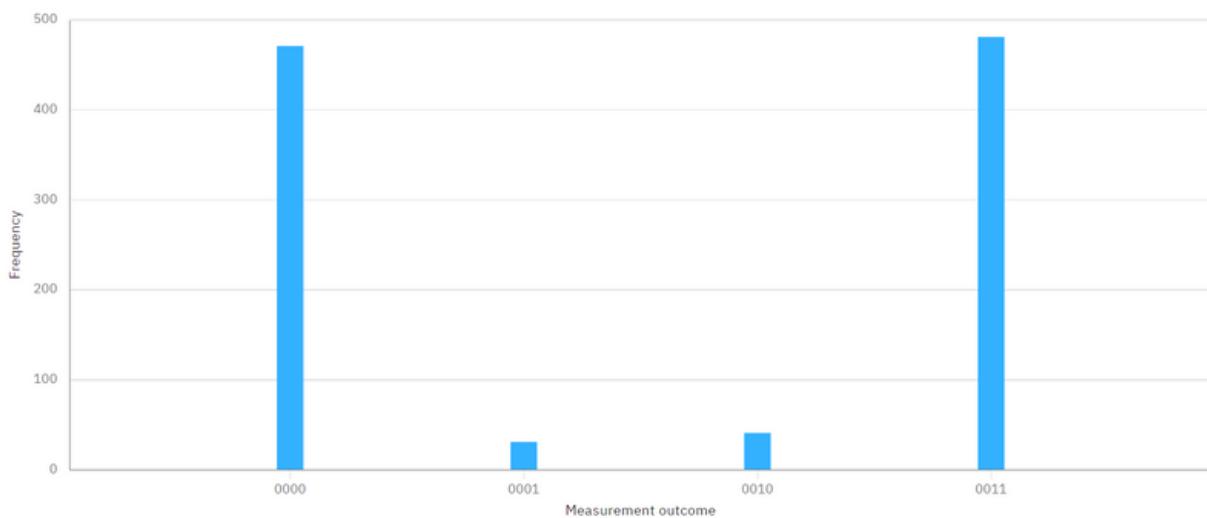


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Hardware execution result

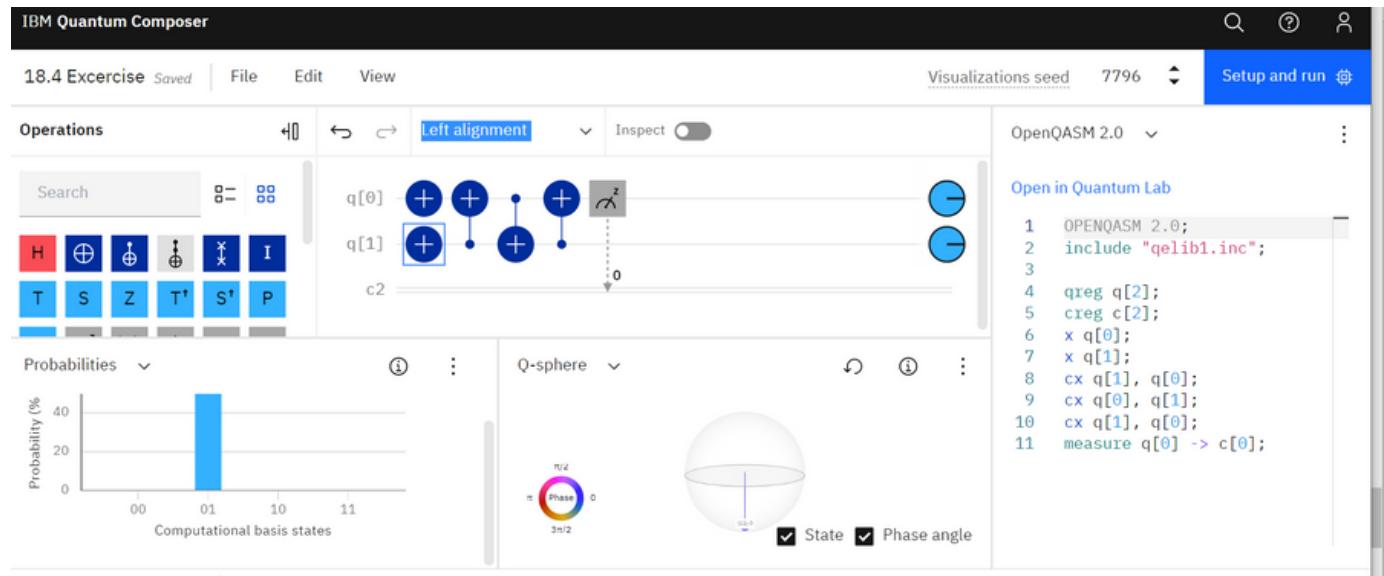
Details

36m 30.5s Total completion time	Sent from	Untitled circuit	Status Timeline
ibmq_manila Backend	Created on	Sep 07, 2022 11:50 AM	Created: Sep 07, 2022 11:50 AM
	Sent to queue	Sep 07, 2022 11:50 AM	Transpiling: 1.3s
	Provider	ibm-q/open/main	Validating: 1.1s
	Run mode	fairshare	In queue: 35m 51.6s
	# of shots	1024	Running: 4.9s
	# of circuits	1	time in system 4.9s
			Completed: Sep 07, 2022 12:26 PM



Introduction to Quantum computing (From a layperson to programmer in 30 steps)

18.4)



Simulation



Hardware execution result

Nithin Reddy Govindugari(nithinreddy1747@gmail.com)

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

Details

21h 2m 51.4s

Total completion time

Sent from

[18.4 Excercise](#)

ibm_oslo

Backend

Sent to queue

Sep 07, 2022 12:14 PM

Provider

ibm-q/open/main

Run mode

fairshare

of shots

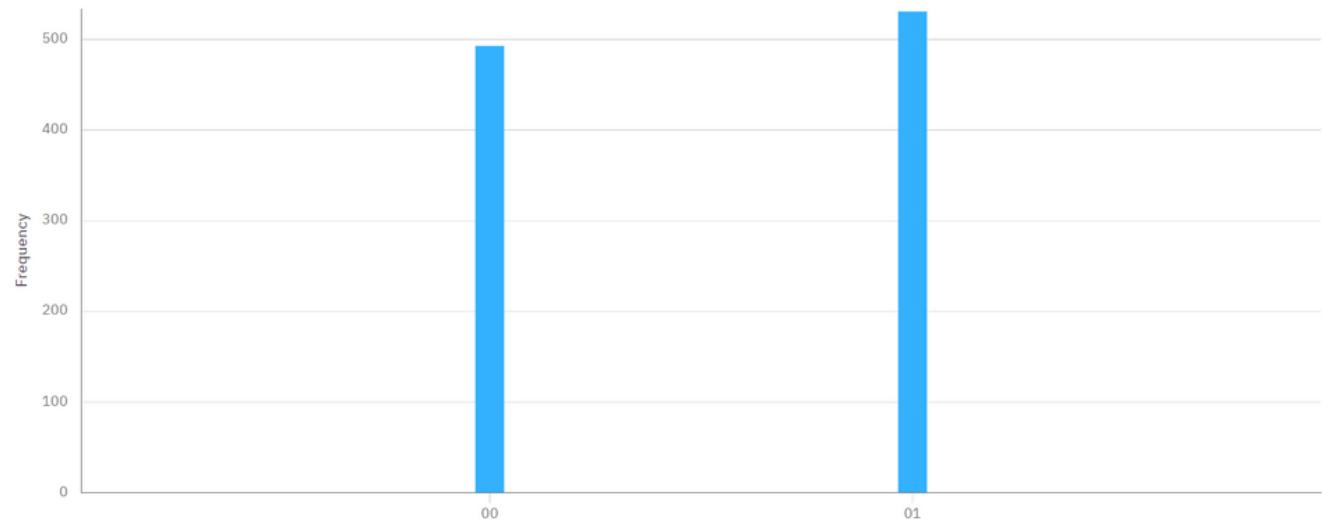
1024

of circuits

1

Status Timeline

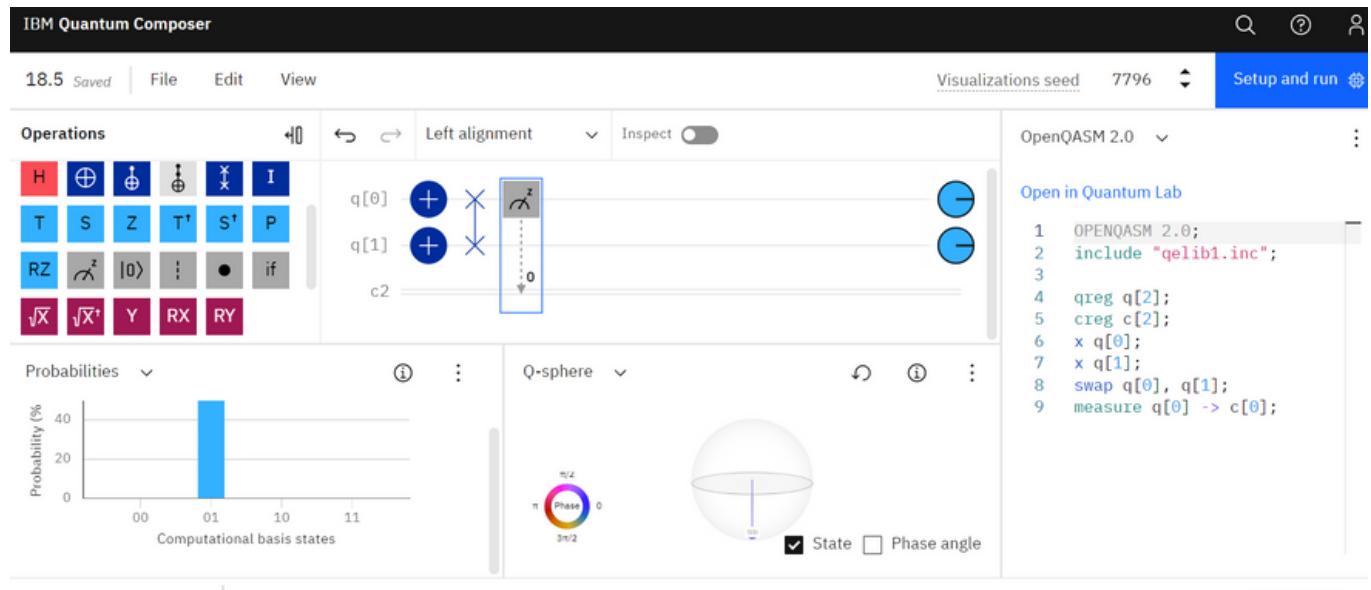
- ✓ Created: Sep 07, 2022 12:14 PM
- ✓ Transpiling: 1s
- ✓ Validating: 1.1s
- ✓ In queue: 21h 2m 41.1s
- ✓ Running: 4.4s
time in system 4.4s
- ✓ Completed: Sep 08, 2022 9:16 AM



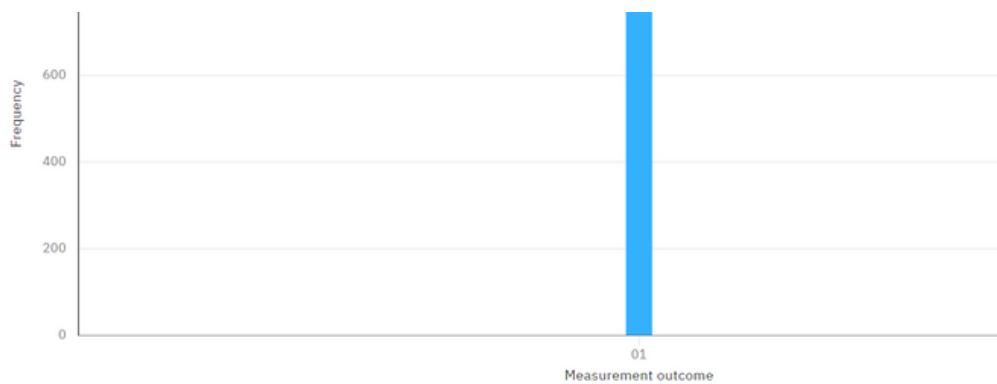
18.5)

Nithin Reddy Govindugari(nithinreddy1747@gmail.com)

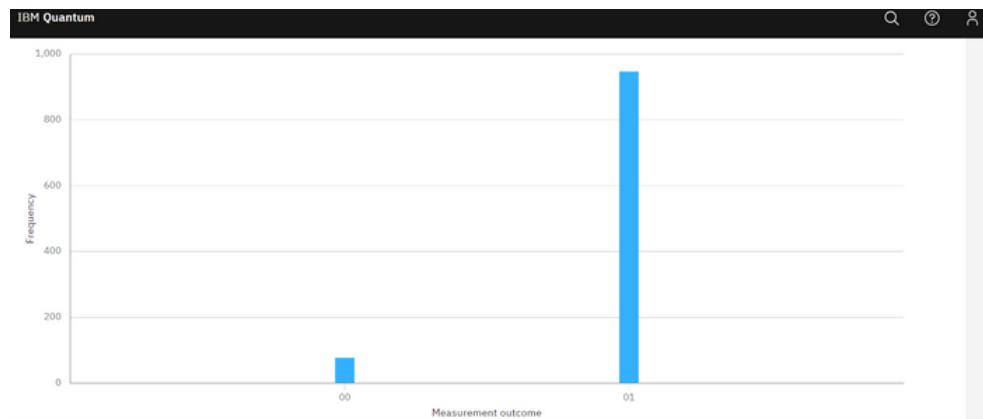
Introduction to Quantum computing (From a layperson to programmer in 30 steps)



Simulation results



Hardware execution result



Chapter 19

19.1)

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19.1

Given

$$L \left(\frac{1}{\sqrt{2}} (|a\rangle_A + |b\rangle_A) |0\rangle_B \right) = \frac{1}{\sqrt{2}} (|a\rangle_A + |b\rangle_A) \\ + \frac{1}{\sqrt{2}} (|a\rangle_B + |b\rangle_B)$$

$$L(|a\rangle_A |0\rangle_B + |b\rangle_A |0\rangle_B)$$

$$\Rightarrow |a\rangle_A |a\rangle_B + |b\rangle_A |b\rangle_B \rightarrow L.H.S - (1)$$

→ R.H.S (Tensor product)

$$\Rightarrow |a\rangle_A |a\rangle_B + |a\rangle_A |b\rangle_B + |b\rangle_A |a\rangle_B + |b\rangle_A |b\rangle_B$$

So here L.H.S \neq R.H.S because $|a\rangle_A |b\rangle_B$
 $\& |b\rangle_A |a\rangle_B$ probability is zero

Non cloning theorem survives on the concept
 of not able to copying linear combinations
 and linearity stands important

Ignored $\frac{1}{\sqrt{2}}$ for simplification

19.2)

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

(Q.2)

- a) Here we know when Hadamard gate is applied to
 i) State it will be $|+\rangle$ state and before measuring
 the entangled or Bob qubit we are applying z gate
 we know z gate is not gate in $|+\rangle/|-\rangle$ basis
 so it converts $|+\rangle$ to $|+\rangle$ before measurement

b) Here we have $H_2 \text{CNOT } H_1$
 As we know CNOT is 2 qubit we can make
 Hadamard 2 qubit by applying identity

$$\begin{aligned}
 &= (H \otimes I) \text{CNOT} (H \otimes I) \\
 &= \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{pmatrix} \\
 &= \frac{1}{2} \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & -1 \\ 1 & 0 & -1 & 0 \end{pmatrix} \\
 &= \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & -1 \\ 1 & 1 & -1 & 1 \\ 1 & -1 & 1 & 1 \\ -1 & 1 & 1 & 1 \end{pmatrix}
 \end{aligned}$$

$$\begin{aligned}
 H \otimes I &\stackrel{\text{def}}{=} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \\
 &= \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{pmatrix}
 \end{aligned}$$

- c) Here we know CNOT gate is used for entanglement
 and important part in Quantum Teleportation
 In first step Hadamard gate is applied to Alice qubit and it creates
 superposition and later CNOT gate is applied for entanglement
 and Hadamard is applied to Alice qubit to bring back to
 original state. In LSB of Bob z gate is applied to switch the state

D) Here as we have converted Hadamard gate to two qubit matrix as we are applying CNOT gate to it and entanglement avoids cancelling of two Hadamard gates

19.3)

19.3) Yes, It is possible to clone basis state as Non-cloning theorem only limits copying of linear combination of Alice to Bob but when we measure Alice state we can clone as many times as we want as when we clone measure the value remains same

You can't clone arbitrary state due to linearity and we proved it 1st question

According to 19.7 we have to change basis and we can do teleportation but Alice won't be same after teleporting

Heisenberg uncertainty principle states that you can't calculate position of a things exactly at a time in two basis. we can't have exact basis states of two different basis at a time. we need to have linear combination of one state and base state of one state

Chapter 20
 20.1)

20.1)

$$= \frac{1}{\sqrt{2}} (\alpha |1000\rangle + \beta |101\rangle + \alpha |1011\rangle + \beta |1110\rangle)$$

$$= \frac{1}{\sqrt{2}} \left(\alpha |10\rangle \otimes |10\rangle \otimes |1\rangle + \beta |1\rangle \otimes |0\rangle \otimes |11\rangle + \alpha |0\rangle \otimes |11\rangle \otimes |1\rangle + \beta |1\rangle \otimes |11\rangle \otimes |0\rangle \right)$$

$$\frac{1}{\sqrt{2}} \left(\alpha \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} \right) + \beta \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$+ \alpha \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} + \beta \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$= \frac{1}{\sqrt{2}} \left(\alpha \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} \right) + \beta \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

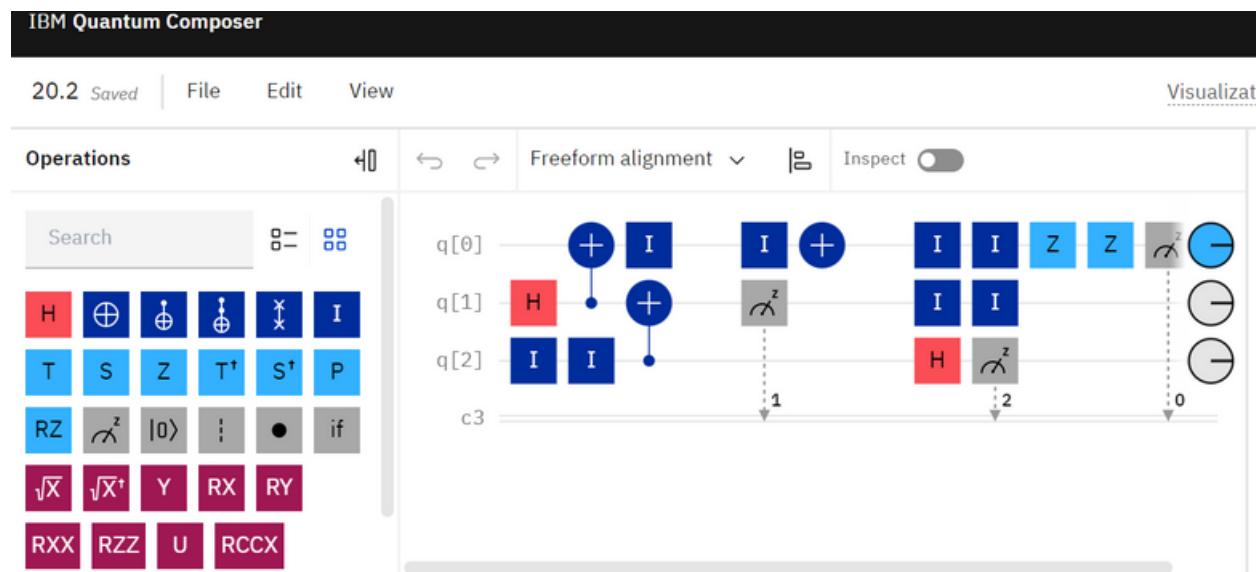
$$+\alpha \left(\left[\begin{array}{c} 0 \\ 1 \\ 0 \\ 0 \end{array} \right] \otimes \left[\begin{array}{c} 0 \\ 1 \end{array} \right] \right) + \beta \left(\left[\begin{array}{c} 0 \\ 0 \\ 0 \\ 1 \end{array} \right] \otimes \left[\begin{array}{c} 0 \\ 1 \end{array} \right] \right)$$

$$\frac{1}{\sqrt{2}} \left(\alpha \left[\begin{array}{c} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{array} \right] + \beta \left[\begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \end{array} \right] + \alpha \left[\begin{array}{c} 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \end{array} \right] + \beta \left[\begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \end{array} \right] \right)$$

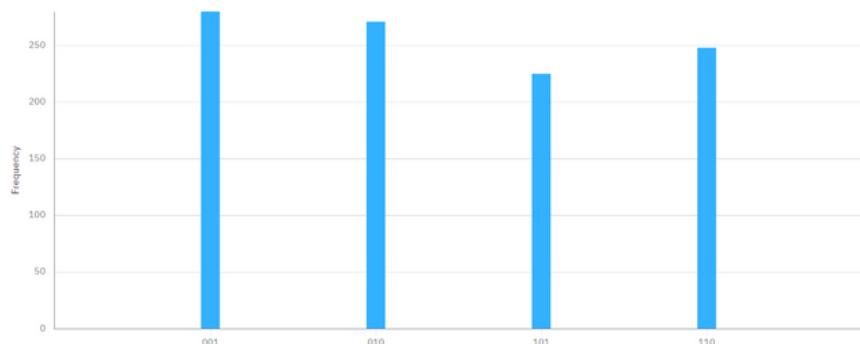
$$= \frac{1}{\sqrt{2}} \left[\begin{array}{c} \alpha \\ 0 \\ 0 \\ \alpha \\ \beta \\ \beta \\ 0 \\ 0 \end{array} \right]$$

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

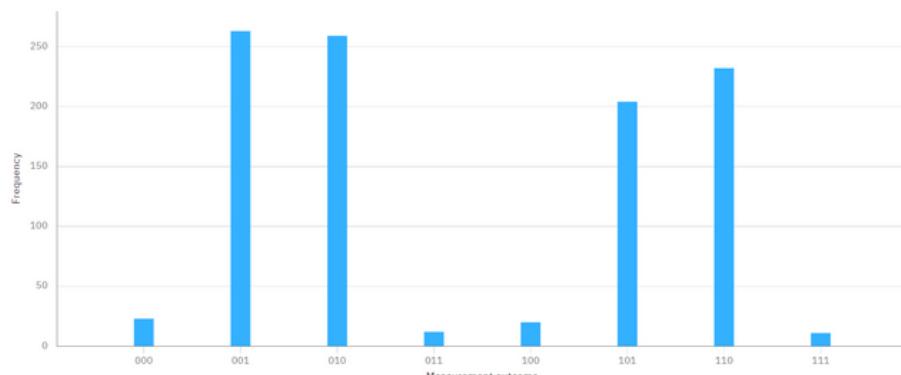
20.2)



Simulation result



IBM



20.3)

Nithin Reddy Govindugari(nithinreddy1747@gmail.com)

20.3)

1) Due to no cloning theorem it won't allow cloning of linear arbitrary states so we need to measure Alice

2) Entanglement swapping (20.5) allows us to transport / teleport

3) Default circuit represents Alice to teleport to Bob. According to 20.4. so yes Quantum teleportation happens by default

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

21.1)

$$\begin{aligned}
 & \text{i) } f_B(x) \\
 & U_f |0\rangle |0\rangle = |0\rangle |0 \oplus f(0)\rangle \\
 & U_f |0\rangle |1\rangle = |0\rangle |0 \oplus f(0)\rangle \\
 & U_f |1\rangle |0\rangle = |1\rangle |0 \oplus f(1)\rangle \\
 & U_f |1\rangle |1\rangle = |1\rangle |1 \oplus f(1)\rangle
 \end{aligned}$$

$$\begin{aligned}
 & \text{ii) } f_B(x) \Rightarrow f_B(0) = 0 \& f_B(1) = 1 \\
 & = |0\rangle |0 \oplus 0\rangle = |0\rangle |0\rangle \\
 & = |0\rangle |0 \oplus 0\rangle = |0\rangle |1\rangle \\
 & = |0\rangle |0 \oplus 1\rangle = |1\rangle |1\rangle \\
 & = |1\rangle |1 \oplus 1\rangle = |1\rangle |0\rangle
 \end{aligned}$$

$$\begin{aligned}
 & \text{iii) } f_C(x) \Rightarrow f_C(0) = 1 \& f_C(1) = 0 \\
 & = |0\rangle |0 \oplus 1\rangle = |0\rangle |1\rangle = |0\rangle \\
 & = |0\rangle |1 \oplus 1\rangle = |0\rangle |0\rangle \\
 & = |1\rangle |0 \oplus 0\rangle = |1\rangle |0\rangle \\
 & = |1\rangle |1 \oplus 0\rangle = |1\rangle |1\rangle
 \end{aligned}$$

$$\begin{aligned}
 & \text{iv) } f_D(x) \Rightarrow f_D(0) = 1 \& f_D(1) = 1 \\
 & = |0\rangle |0 \oplus 1\rangle = |0\rangle |1\rangle \\
 & = |0\rangle |1 \oplus 1\rangle = |0\rangle |0\rangle \\
 & = |1\rangle |0 \oplus 1\rangle = |0\rangle |1\rangle \\
 & = |1\rangle |1 \oplus 1\rangle = |1\rangle |0\rangle
 \end{aligned}$$

21.2)

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

21.2)

We know from 17.11 that

$$H^{\otimes 2} |100\rangle = \frac{1}{2} (|100\rangle + |101\rangle + |110\rangle + |111\rangle)$$

$$= \frac{1}{2} \begin{pmatrix} 1 \\ 1 \\ 1 \\ 1 \end{pmatrix}$$

$$L.H.S = \frac{1}{2} (H \cdot H) (I \times) |10\rangle |10\rangle$$

$$= \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

$$= \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix}$$

$$= \frac{1}{2} (H \otimes H) |10\rangle |11\rangle \rightarrow R.H.S$$

$$= \frac{1}{2} \begin{pmatrix} 1 \\ -1 \\ -1 \\ -1 \end{pmatrix}$$

$$|1+\rangle |1-\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} \otimes \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ -1 \end{pmatrix}$$

$$= \frac{1}{2} \begin{pmatrix} 1 \\ -1 \\ 1 \\ -1 \end{pmatrix}$$

$$\text{Hence } (H \otimes H) (I \otimes I) |10\rangle |10\rangle = (H \otimes H) |10\rangle |11\rangle = |1+\rangle |1-\rangle$$

21.3)

21.3) Given state $|101\rangle$

from 17.21

$$H^{\otimes n} |y\rangle = \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} (-1)^{x \cdot y} |x\rangle$$

we know $n=2$ $|101\rangle_2 = |10\rangle$.

$$H^{\otimes 2} |1\rangle = \frac{1}{\sqrt{2^2}} \left((-1)^{0 \cdot 1} |0\rangle + (-1)^{1 \cdot 1} |1\rangle \right)$$

$$+ (-1)^{2 \cdot 1} |2\rangle + (-1)^{3 \cdot 1} |3\rangle \right)$$

$$0 \cdot 1 = (0 \text{ AND } 0) \oplus (0 \text{ AND } 1) = 0 \oplus 0 = 0$$

$$1 \cdot 1 = (0 \text{ AND } 0) \oplus (1 \text{ AND } 1) = 0 \oplus 1 = 1$$

$$2 \cdot 1 = (1 \text{ AND } 0) \oplus (0 \text{ AND } 1) = 0 \oplus 0 = 0$$

$$3 \cdot 1 = (1 \text{ AND } 0) \oplus (1 \text{ AND } 1) = 0 \oplus 1 = 1$$

$$H^{\otimes 2} |1\rangle = H^{\otimes 2} |10\rangle = \frac{1}{2} (|10\rangle - |11\rangle + |12\rangle - |13\rangle)$$

$$= \frac{1}{2} (|100\rangle - |101\rangle + |110\rangle - |111\rangle) - i$$

from 21.6

$$|+\rangle_{1-2} = \frac{1}{2} (|100\rangle - |101\rangle + |110\rangle - |111\rangle) - i$$

Hence proved $\textcircled{1} = \textcircled{2}$

21.4)

From 21.5 let us try for balanced function

$f(x)$ and gives equal number of 0 & 1

$$\text{ex: } f_B(0) = 0, f_B(1) = 1$$

Q:

$$U_f |0\rangle |0\rangle = |0\rangle |0\rangle \oplus f(0)\rangle = |0\rangle |0\oplus 0\rangle = |0\rangle |0\rangle$$

$$U_f |0\rangle |1\rangle = |0\rangle |1\rangle \oplus f(0)\rangle = |0\rangle |1\oplus 0\rangle = |0\rangle |1\rangle$$

$$U_f |1\rangle |0\rangle = |1\rangle |0\rangle \oplus f(1)\rangle = |1\rangle |0\oplus 1\rangle = |1\rangle |1\rangle$$

$$U_f |1\rangle |1\rangle = |1\rangle |1\rangle \oplus f(1)\rangle = |1\rangle |1\oplus 1\rangle = |1\rangle |0\rangle$$

Writings by Nithin Reddy

Chapter 22
22.1) & 22.2)

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22.1 Given $x=5, f(5)=3, n=3, m=3$

y assumed as $|10\rangle$

we know from 22.1

$$U_f |x\rangle_n |y\rangle_m = |x\rangle_n |y \oplus f(x)\rangle_m$$

$$U_f |15\rangle_3 |10\rangle_3 = |15\rangle_3 |0 \oplus f(5)\rangle_3$$

~~|100101>~~

$$= |101\rangle_3 |0 \oplus \text{P(1)}\rangle = |101\rangle_3 |000\rangle$$

$$= |101\rangle |001\rangle$$

$$= |1101011\rangle$$

22.2 If we are using same method as phase quantum oracle

Two basis vectors $|a\rangle \times |b\rangle$ and inner product is $\langle a|b\rangle$

i) $|a\rangle$ transformed by XOF Quantum oracle
we can assume γ as 0

$$|a\rangle |0 \oplus f(a)\rangle$$

$$|a\rangle |f(a)\rangle$$

$$ii) |b\rangle \rightarrow |b\rangle; |f(b)\rangle$$

Inner product is (No matrix involved)

$$\langle a|b\rangle = |a\rangle |b\rangle$$

$$\langle f(a)|f(b)\rangle \langle a|b\rangle$$

For example if $a = b$ then we know

Inner product of same vectors is

$$\langle f(a)|f(b)\rangle \langle a|b\rangle$$

$$\langle f(a)|f(a)\rangle \langle a|a\rangle$$

$$= 1 \cdot 1 = 1$$

& If $a \neq b$, $\langle a|b\rangle = 0$ as we already know

Basic vector inner product is NOT changed

It is still unitary

22.3)

$$\begin{aligned}
 & \text{22.3} \quad U_{\text{swap}} = \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} \langle j | U | i \rangle | ij \rangle \langle i | \\
 & \qquad \qquad \qquad \boxed{N=4} \\
 & \cdot \langle 001 | U_{\text{swap}} | 100 \rangle | 100 \rangle \langle 001 | + \langle 01 | U_{\text{swap}} | 100 \rangle \\
 & \qquad \qquad \qquad \langle 001 | \\
 & \qquad \qquad \qquad + \langle 101 | U_{\text{swap}} | 100 \rangle | 110 \rangle \langle 001 | \\
 & \qquad \qquad \qquad + \langle 111 | U_{\text{swap}} | 100 \rangle | 110 \rangle \langle 001 | \\
 & \qquad \qquad \qquad + \dots + \dots \\
 & \qquad \qquad \qquad \boxed{\cancel{\langle 001 | U_{\text{swap}} | 110 \rangle | 100 \rangle \langle 101 |} + \cancel{\langle 101 | U_{\text{swap}} | 110 \rangle | 100 \rangle \langle 101 |}} \\
 & \cdot \langle 001 | U_{\text{swap}} | 101 \rangle | 100 \rangle \langle 011 | + \langle 01 | U_{\text{swap}} | 101 \rangle | 100 \rangle \\
 & \qquad \qquad \qquad \langle 011 | \\
 & \qquad \qquad \qquad + \langle 101 | U_{\text{swap}} | 101 \rangle | 110 \rangle \langle 011 | + \\
 & \qquad \qquad \qquad \langle 111 | U_{\text{swap}} | 101 \rangle | 110 \rangle \langle 011 | \\
 & \qquad \qquad \qquad + \\
 & \cdot \langle 001 | U_{\text{swap}} | 110 \rangle | 100 \rangle \langle 101 | + \langle 01 | U_{\text{swap}} | 110 \rangle | 100 \rangle \\
 & \qquad \qquad \qquad \langle 101 | \\
 & \qquad \qquad \qquad + \langle 101 | U_{\text{swap}} | 110 \rangle | 110 \rangle \langle 101 | \\
 & \qquad \qquad \qquad + \\
 & \qquad \qquad \qquad + \langle 111 | U_{\text{swap}} | 110 \rangle | 110 \rangle \langle 101 |
 \end{aligned}$$

$$+ \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 0 & 0 & i \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$
$$v_{swap} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

from 16.3 It is proved

for f_G $f_G(0) = 1$, $f_G(1) = 0$

$$22.4) \langle 01 \langle 11 | (U_f | 0 \rangle | 10 \rangle) = \langle 01 \langle 11 | 101 \rangle = 1$$

$$\langle 01 \langle 01 | (U_f | 0 \rangle | 11 \rangle) = \langle 01 \langle 01 | 100 \rangle = 1$$

$$\langle 11 \langle 01 | (U_f | 11 \rangle | 10 \rangle) = \langle 11 \langle 01 | 110 \rangle = 1$$

$$\langle 11 \langle 11 | (U_f | 11 \rangle | 11 \rangle) = \langle 11 \langle 11 | 111 \rangle = 1$$

$$= | 01 \rangle | 100 \rangle + | 100 \rangle | 01 \rangle + | 110 \rangle | 110 \rangle + | 111 \rangle | 111 \rangle$$

$$= \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} [1 \ 0 \ 0 \ 0] + \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} [0 \ 1 \ 0 \ 0] + \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix} [0 \ 0 \ 1 \ 0]$$

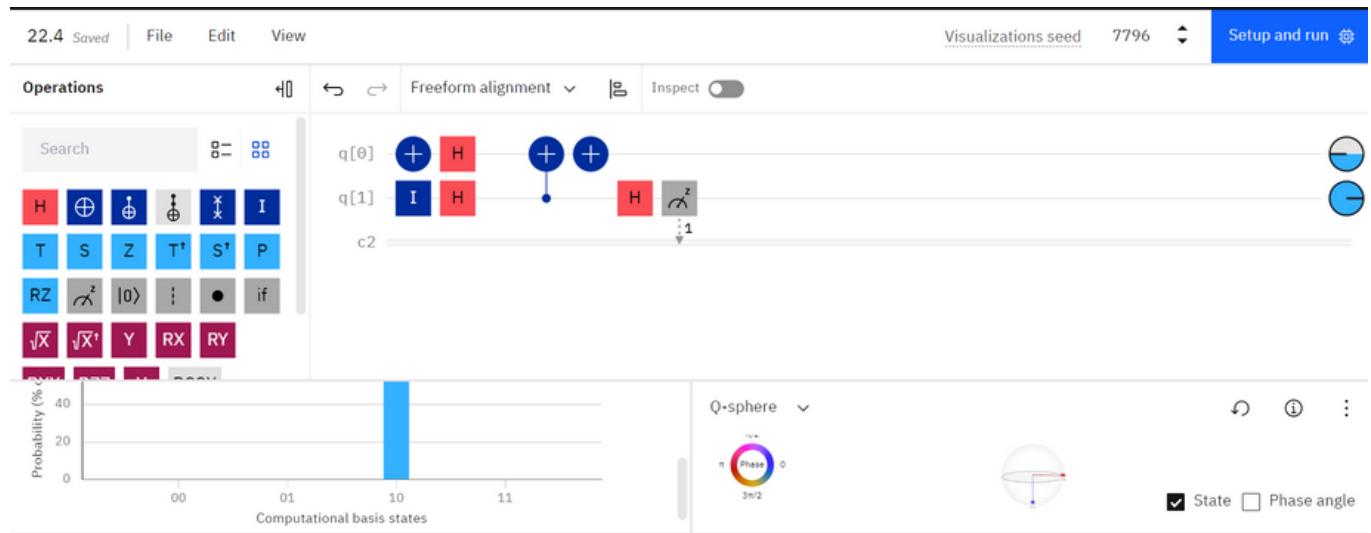
$$= \begin{pmatrix} 0 \\ 0 \\ 1 \\ 1 \end{pmatrix} [0 \ 0 \ 0 \ 1]$$

$$\begin{bmatrix} 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

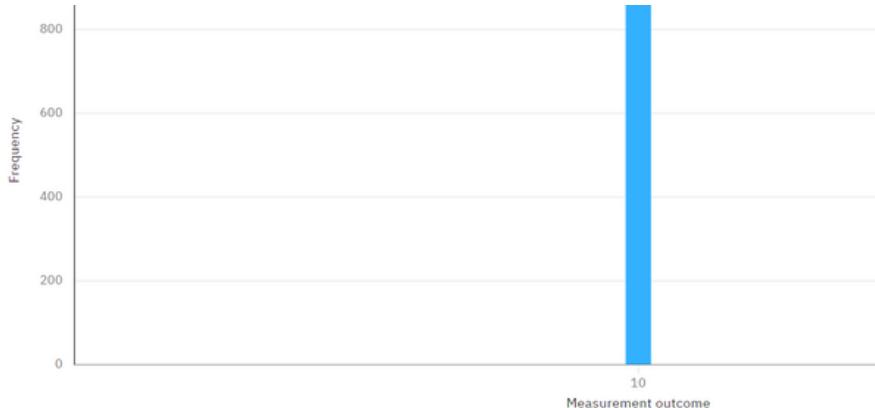
$$\begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

Introduction to Quantum computing (From a layperson to programmer in 30 steps)



Simulator outcome

It is 10 in IBMQ instead of 01 because for balanced function MSB is always 1 (page no 212)



$$f_0, f_1 |0\rangle = 1 \quad f_d(1) = 1$$

$$\langle 01| \langle 11| (U_f |0\rangle |0\rangle) = \langle 01| \langle 10| |01\rangle = 1$$

$$\langle 01| \langle 01| (U_f |0\rangle |1\rangle) = \langle 00| \langle 10| |00\rangle = 1$$

$$\langle 11| \langle 11| (U_f |1\rangle |0\rangle) = \langle 11| \langle 11| |11\rangle = 1$$

$$\langle 11| \langle 01| (U_f |1\rangle |1\rangle) = \langle 10| \langle 11| |10\rangle = 1$$

$$= |01\rangle |00\rangle + |00\rangle |01\rangle + |11\rangle |10\rangle$$

$$+ |10\rangle |11\rangle$$

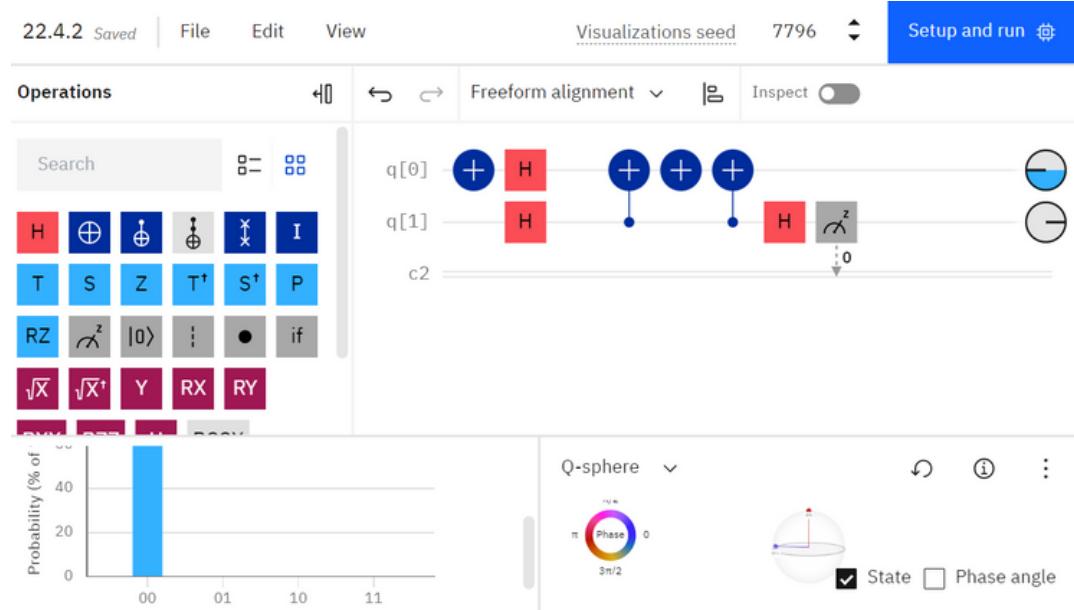
$$+ \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$+ \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

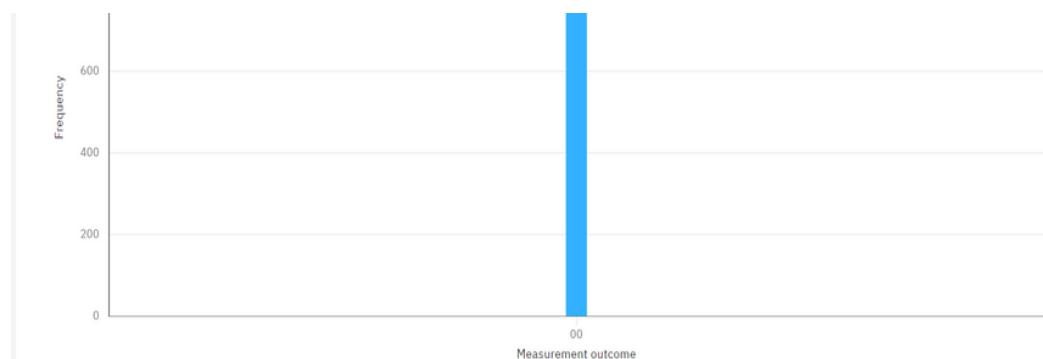
$$\begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

It is similar to above statement and as explained in page 200 it is always 0 for constant function here we can also remove both cnot gates as it wont make much difference. I have added it to describe the functionality



Simulator results



Chapter 23

(3.1) Fig 23.2

Let us assume angle between $|a\rangle$ & $|b\rangle$ as α

$\sqrt{1/\theta}$	$\alpha + 2\theta$	$-\theta$
$w\sqrt{1/\theta}$	$\alpha - 2\theta$	$+3\theta$
$\sqrt{w}\sqrt{1/\theta}$	α	$+ \theta$
$w\sqrt{w}\sqrt{1/\theta}$	$\alpha - 4\theta$	$+5\theta$

angle between $|b\rangle$ & $|a\rangle$, between $|b\rangle$ & $|a\rangle$,

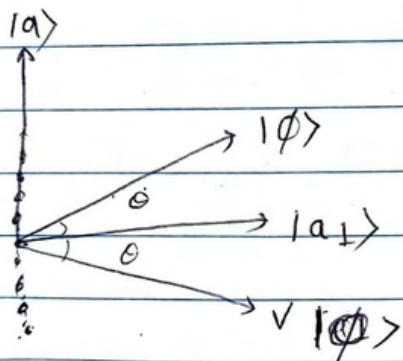
23.2)

23.2)

$$\text{Given } V = I - 2|a\rangle\langle a|$$

As we already know V is a quantum gate and it is linear combination of all basis vectors except our target entry $|a\rangle$

we know phase shift is added to target entry and it is multiplied by -1



$$\text{From } V|\phi\rangle = (I - 2|a\rangle\langle a|)|\phi\rangle$$

$$= \sum_{i=0, i \neq a}^2 |i\rangle$$

Here when operator V is applied to $|\phi\rangle$ then it reflects $|\phi\rangle$ towards negation (phase shift)

23.3)

Given $n=2, a=3$

$$|0\rangle = \frac{1}{\sqrt{2}} \sum_{x=0}^{2^n-1} |x\rangle$$

we know $n=2, a=3$

$$= \frac{1}{\sqrt{2}} \sum_{x=0}^{2^2-1} |x\rangle$$

$$= \frac{1}{\sqrt{2}} \sum_{x=0}^3 |x\rangle$$

$$= \frac{1}{\sqrt{2}} (|00\rangle + |01\rangle + |10\rangle + |11\rangle)$$

$$|a\rangle = |11\rangle$$

$$|a_+\rangle = \frac{1}{\sqrt{2^n-1}} \sum_{x=0 (x \neq a)}^{2^n-1} |x\rangle$$

$$= \frac{1}{\sqrt{3}} \sum_{x=0 (x \neq a)}^3 |x\rangle$$

$$= \frac{1}{\sqrt{3}} (|00\rangle + |01\rangle + |10\rangle)$$

23.4 & 23.5)

23.4 Sol

we know

$$|\psi\rangle = \sqrt{\frac{1}{2}} (|00\rangle + |01\rangle + |10\rangle + |11\rangle)$$

$$= \frac{1}{2} (|00\rangle + |01\rangle + |10\rangle - |11\rangle)$$

$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix}$$

23.5)

Algorithm B has linear speed up
over Algorithm A

Chapter 24

24.1)

24.1)

$$U = \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} \langle j | U | i \rangle | j \rangle \langle i | \rightarrow 22.6$$

we know from 24.4 that

$$U | 1 \rangle_n = (-1)^{f(1)} | 1 \rangle_n \rightarrow 23.9$$

$U_f | 01 \rangle = - | 01 \rangle$ and for every other base it will be same

$$\langle 01 | \langle 01 | (U_f | 10 \rangle | 10 \rangle) = \langle 001 | | 00 \rangle = 1$$

$$\langle 01 | \langle 1 | (U_f | 10 \rangle | 11 \rangle) - \langle 001 | \langle - | 01 \rangle = -1 \quad [4.10B.3]$$

$$\langle 1 | \langle 1 | (U_f | 11 \rangle | 11 \rangle) = \langle 111 | | 11 \rangle = 1$$

$$\langle 1 | \langle 0 | (U_f | 11 \rangle | 11 \rangle) = \langle 1 | \langle 01 | | 11 \rangle = \langle 10 | | 10 \rangle = 1$$

$$U_f = \langle 01 | \langle 01 | (U_f | 10 \rangle | 10 \rangle) | 100 \rangle \langle 100 | + \langle 01 | \langle 1 | | 1 | \\ (U_f | 10 \rangle | 11 \rangle) | 101 \rangle \langle 101 | + \langle 1 | \langle 0 | (U_f | 11 \rangle | 10 \rangle) | 110 \rangle \langle 110 | \\ + \langle 1 | \langle 01 | (U_f | 11 \rangle | 11 \rangle) | 111 \rangle \langle 111 |$$

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$$= |00\rangle\langle 00| - |01\rangle\langle 01| + |11\rangle\langle 11| + |10\rangle\langle 10| \\ + |10\rangle\langle 10| + \dots$$

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} - \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} + \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} = U_1 \text{ i.e., } 24.4$$

24.4)

we need to apply $H^{\otimes 2}(2|0\rangle_2\langle 0|_2 - I)H^{\otimes 2}(|1\rangle_2)$

$$\text{Firstly } 2|0\rangle_2\langle 0|_2 - I = 2 \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \end{pmatrix} \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

$$+ 2 \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} - \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$2|0\rangle_2\langle 0|_2 - I = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

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$$= \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{pmatrix} -1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix} \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix}$$

$$\frac{1}{2} \begin{pmatrix} 1 \\ -1 \\ 1 \\ -1 \end{pmatrix}$$

$$= \frac{1}{4} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{pmatrix} -1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ -1 \\ 1 \\ -1 \end{pmatrix}$$

$$\frac{1}{4} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{pmatrix} -1 \\ 1 \\ 1 \\ -1 \end{pmatrix}$$

$$= \begin{pmatrix} 1+1+1-1 \\ -1-1+1+1 \\ 1+1-1+1 \\ -1-1-1-1 \end{pmatrix} = \frac{1}{4} \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$

$|111\rangle$ over 10 tabs

24.3)

24.3)

we know from 23.9 for $a=4$

$$U_f = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix}$$

$$\sqrt{1/\rho} = U_f \left(\frac{1}{2} (|00\rangle + |01\rangle + |10\rangle + |11\rangle) \right)$$

$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix} \frac{1}{2} \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

$$\frac{1}{2} \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

for $a=3$

$$= H^{\otimes 2} (2|0\rangle_2 \langle 01_2 - I) H^{\otimes 2} (\sqrt{1/\rho})$$

$$= \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix} \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{bmatrix} \frac{1}{2} \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

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$$\begin{array}{c}
 \text{Quantum State} \\
 = \frac{1}{\sqrt{4}} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix} \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \\
 = \frac{1}{\sqrt{4}} \begin{pmatrix} 1+1+1-1 \\ 1-1+1+1 \\ 1+1-1+1 \\ 1+1-1-1 \end{pmatrix} \\
 = \frac{1}{2} \begin{pmatrix} 2 \\ 2 \\ 2 \\ -2 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 1 \\ -1 \end{pmatrix} \\
 = \frac{1}{\sqrt{4}} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \\ 1 \\ -1 \end{pmatrix} \\
 = \frac{1}{\sqrt{4}} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ -1 \\ -1 \\ 1 \end{pmatrix}
 \end{array}$$

$$\frac{1}{4} \begin{bmatrix} 1 - 1 + 1 \\ 1 + 1 - 1 \oplus 1 \\ 1 \oplus 1 + 1 - 1 \\ 1 + 1 + 1 + 1 \end{bmatrix} = \frac{1}{4} \begin{bmatrix} 0 \\ 0 \\ 0 \\ 4 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix} = |111\rangle$$

24.4)

from 24.15

$$C_Z = (I \otimes H) \cup_{\text{xor}} (I \otimes H)$$

$$= \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \otimes \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \otimes \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

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$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} & 0 & \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \\ 0 & \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} & 1 & \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} & 0 & \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \\ 1 & \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} & 1 & \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \end{pmatrix}$$

$$= \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & -1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & -1 \end{pmatrix}$$

$$= \frac{1}{2} \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & -1 \end{bmatrix} \begin{bmatrix} 1 & 1 & 0 & 0 \\ 0 & -1 & 0 & 0 \\ 0 & 0 & 1 & -1 \\ 0 & 0 & 1 & 1 \end{bmatrix}$$

$$= \frac{1}{2} \begin{bmatrix} 2 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 \\ 0 & 0 & 2 & 0 \\ 0 & 0 & 0 & -2 \end{bmatrix}$$

$$= \cancel{\frac{1}{2}} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix} = CZ$$

24.5)

```
#24.5
#Keywords
#superposition is sup_pos
#amplitude is amp

#Grovers Algorithm is implemented using standard procedure.
#Steps followed in books are implemented and tried to create classical grovers algorithm in different

plots = []

# Putting all qubits in superposition using hadamard gate
def sup_pos(qubits):
    states = []
    total_states = int(math.pow(2,qubits))
    amp = 1/math.sqrt(total_states)
    for _ in range(0,total_states):
        states.append(amp)
    return states

# phase flip inverts the states amplitude
def grover_diffusion(states):
    average = sum(states)/len(states)
    for i in range(0,len(states)):
        states[i] = (average-states[i]) + average           #inversion about mean
    return states

#oracle to implement phase inversion
def oracle(states,datalist,key):
    for i in range(0,int(len(datalist))):
        if datalist[i] == key:
            states[i] *= -1                                #phase inversion
    return states
```

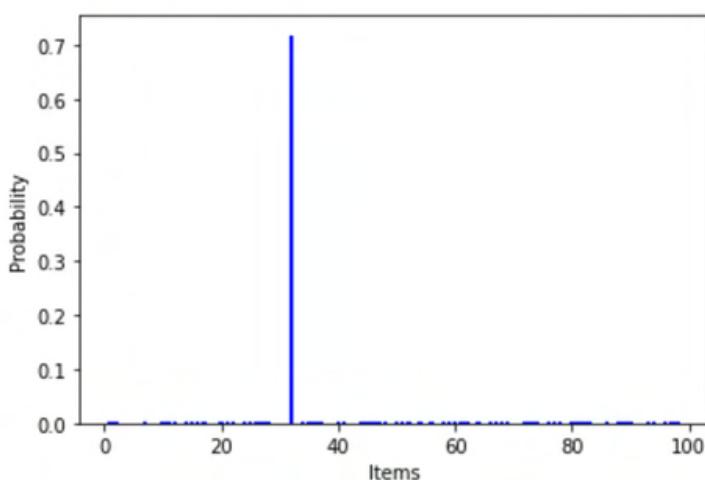
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```
#classical main grover search function
def grover_search(qubits,datalist,key):
    states = sup_pos(qubits)
    num_iterations = math.ceil(math.sqrt(math.pow(2,qubits)))           #iterations
    probability_states = []
    for _ in range(0,num_iterations):
        states = oracle(states,datalist,key)
        states = grover_diffusion(states)
        probability_states = [states[i]*states[i] for i in range(0,len(states))]
    plots.append(probability_states)
    return probability_states

def grover(datalist,key):
    size_datalist = len(datalist)
    qubits_needed = math.ceil(math.log(size_datalist,2))
    paddings_required = int(math.pow(2,qubits_needed) - size_datalist)
    #required if the number of data items is not a power of 2.
    for _ in range(0,paddings_required):
        datalist.append(0)
    grover_search(qubits_needed,datalist,key)
```

```
datalist = random.sample(range(1, 100), 64)
grover(datalist,32)
print(datalist)
print("plots showing the change in probabilities of items after every grover's iteration :")
result ={}
iteration = 1
for plot in plots:
    for i in range(len(plot)):
        result[datalist[i]] = plot[i]
    print("-----")
    print("Iteration ",iteration," :")
    plt.bar(result.keys(), result.values(),color='b')
    plt.ylabel("Probability")
    plt.xlabel("Items")
    plt.show()
    iteration +=1
    print("-----")
```

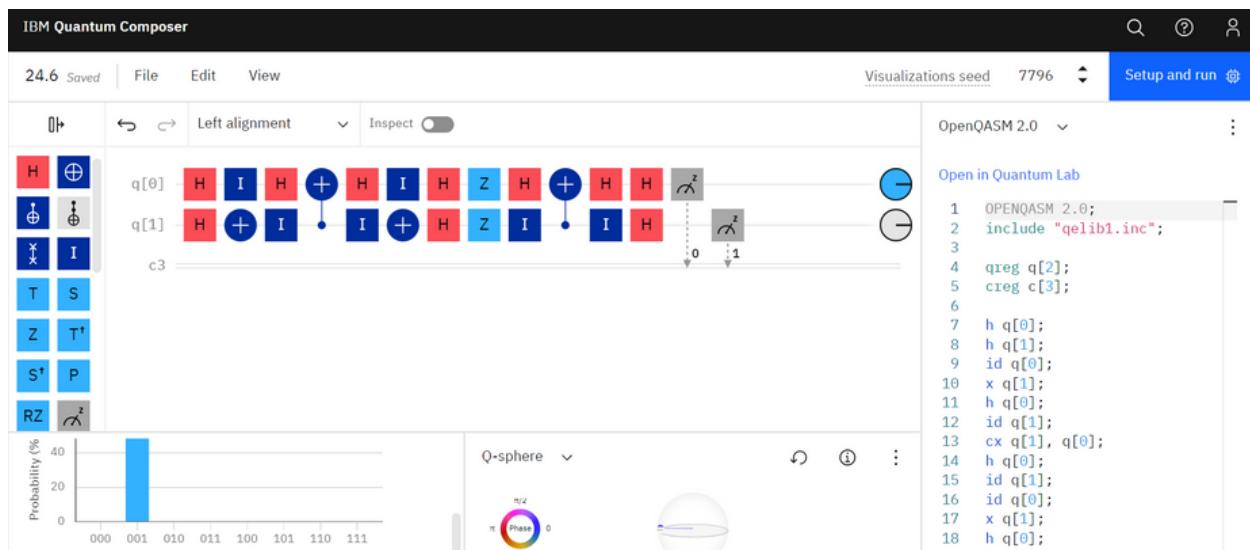
Iteration 8 :



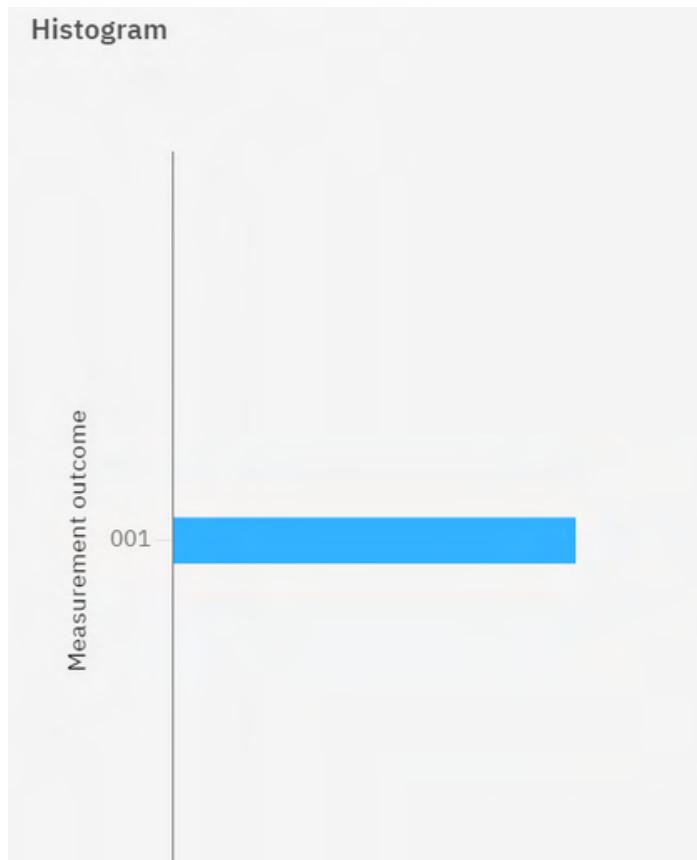
24.6)

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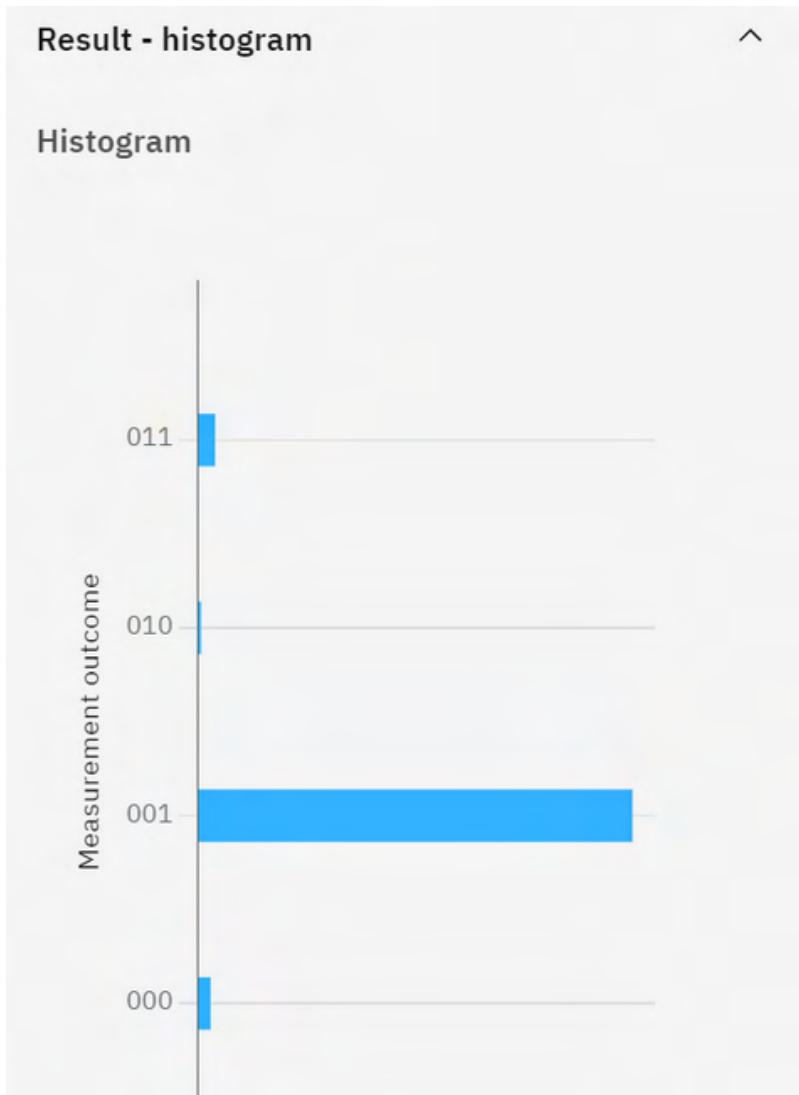


Simulation



Hardware circuit result

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Chapter 25

25.1)

Given that we have to prove

$$25.10 \Rightarrow 25.9$$

we know 25.9. is

$$y_k = \frac{1}{\sqrt{N}} \sum_{j=0}^{N-1} e^{i 2\pi k j / N} x_j$$

$$= \frac{1}{\sqrt{N}} \begin{pmatrix} e^{-i 2\pi (0,0)/N} & e^{-i 2\pi (0,1)/N} & \dots & e^{-i 2\pi (0,N-1)/N} \\ e^{-i 2\pi (1,0)/N} & e^{-i 2\pi (1,1)/N} & \dots & e^{-i 2\pi (1,N-1)/N} \\ \vdots & \vdots & \ddots & \vdots \\ e^{-i 2\pi ((N-1),0)/N} & e^{-i 2\pi ((N-1),1)/N} & \dots & e^{-i 2\pi ((N-1),(N-1))/N} \end{pmatrix}$$

$$\begin{pmatrix} x_0 \\ x_1 \\ \vdots \\ x_{N-1} \end{pmatrix} \xrightarrow{\quad} ②$$

from 25.4 (First property we know)

$$e^{i2\pi m/N} = \omega^m$$

So it equals to ① after substitution

$$\frac{1}{\sqrt{N}} \begin{pmatrix} \omega^{-0,0} & \omega^{-0,1} & \cdots & \omega^{-0,(N-1)} \\ \omega^{-1,0} & \omega^{-1,1} & \cdots & \omega^{-1,(N-1)} \\ \vdots & \vdots & \ddots & \vdots \\ \omega^{-(N-1),0} & \omega^{-(N-1),1} & \cdots & \omega^{-(N-1),(N-1)} \end{pmatrix}$$

$$\begin{pmatrix} x_0 \\ x_1 \\ \vdots \\ x_{N-1} \end{pmatrix}$$

25.2)

25.2Yes. U_{QFT} is Hermitian

From 25.20 we know

$$U_{QFT} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad \text{i.e., Hadamard gate}$$

equals to

$$U_{QFT}^T = U_{QFT} \quad (\text{Hermitian rule is satisfied})$$

(25.12)
 Here, U_{QFT} is symmetric & satisfies Hermitian requirements. I believe it should be Hermitian because according to Fourier transforms domains changes i.e., bases in QFT it should remain same

25.3)

25.3)

U_{QFT} is unitary (considering row vectors)

$$\langle x | = \frac{1}{\sqrt{N}} \begin{bmatrix} w^{-y_0}, w^{-y_1}, \dots, w^{-y_{(N-1)}} \end{bmatrix}^*$$

$$\langle y | = \frac{1}{\sqrt{N}} \begin{bmatrix} w^{-y_0}, w^{-y_1}, \dots, w^{-y_{(N-1)}} \end{bmatrix}^*$$

$$\langle x | y \rangle = \frac{1}{N} \sum_{i=0}^{N-1} (w^{-m \cdot L})^* \cdot w^{-n \cdot L}$$

$$= \frac{1}{N} \sum_{i=0}^{N-1} ((e^{-i2\pi m \cdot L/N})^*) e^{-i2\pi n \cdot L}$$

$$= \frac{1}{N} \sum_{i=0}^{N-1} w^{-m \cdot L} \cdot w^{-n \cdot L}$$

$$= \frac{1}{N} \sum_{i=0}^{N-1} (w^{-L(m-n)})$$

if $m=n$ it is 1 else it is zero

U_{QFT} is unitary

25.4)

25.3)

U_{QFT} is unitary (considering row vectors)

$$|\psi\rangle = \frac{1}{\sqrt{N}} [w^{-y_0}, w^{-y_1}, \dots, w^{-y_{(N-1)}}]^T$$

$$|\psi\rangle = \frac{1}{\sqrt{N}} [w^{-y_0}, w^{-y_1}, \dots, w^{-y_{(N-1)}}]^T$$

$$\langle x|y\rangle = \frac{1}{N} \sum_{i=0}^{N-1} (w^{-m_i})^* w^{-n_i}$$

$$= \frac{1}{N} \sum_{i=0}^{N-1} ((e^{-i2\pi m_i/N})^*) e^{-i2\pi n_i}$$

$$= \frac{1}{N} \sum_{i=0}^{N-1} w^{-m_i} \cdot w^{-n_i}$$

$$= \frac{1}{N} \sum_{i=0}^{N-1} (w^{-L(m-n)})$$

if $m=n$ it is 1 else it is zero

U_{QFT} is unitary

25.5)

25.5)

$$\text{Given } U_{\text{IQFT}} U_{\text{QFT}} = U_{\text{QFT}} U_{\text{IQFT}} = I$$

we know

$$U_{\text{QFT}} = \frac{1}{\sqrt{N}} \begin{pmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \omega^{-1} & \omega^{-2} & \dots & \omega^{-(N-1)} \\ 1 & \omega^{-2} & \omega^{-4} & \dots & \omega^{-2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{-(N-1)} & \omega^{-2(N-1)} & \dots & \omega^{-(N^2-N)} \end{pmatrix}$$

$$U_{\text{IQFT}} = \frac{1}{\sqrt{N}} \begin{pmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \omega^1 & \omega^2 & \dots & \omega^{(N-1)} \\ 1 & \omega^2 & \omega^4 & \dots & \omega^{2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{(N-1)} & \omega^{2(N-1)} & \dots & \omega^{(N-1)(N-1)} \end{pmatrix}$$

$$U_{\text{IQFT}} * U_{\text{QFT}} = \frac{1}{N} \begin{pmatrix} N(\omega^0 + \omega^1 + \omega^2 + \dots + \omega^{(N-1)}) & & & & \\ \omega^1 + \omega^2 + \dots + \omega^{(N-1)} & N & & & \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \vdots & \ddots & N \end{pmatrix}$$

 From 25.5 we know $\omega^0 + \omega^1 + \dots + \omega^{(N-1)} = 0$

$$= \frac{N}{N} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} - \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

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$$U_{\text{QFT}} = \frac{1}{\sqrt{N}} \begin{bmatrix} N & 1 + \omega^1 + \omega^2 + \dots + \omega^{(N-1)} \\ 1 + \omega^1 + \omega^2 + \dots + \omega^{(N-1)} & N \\ \vdots & \vdots \\ N & N \end{bmatrix}$$

From 25.5 & 25.6

$$= \frac{1}{\sqrt{N}} \begin{bmatrix} N & 0 & \dots & 0 \\ 0 & N & \dots & \vdots \\ \vdots & \vdots & \ddots & N \\ \vdots & \vdots & \vdots & \vdots \end{bmatrix}$$

25b)

we know

$$U_{\text{IQFT}} = \frac{1}{\sqrt{N}} \begin{bmatrix} 1 & 1 & 1 & 1 & \dots & 1 \\ 1 & \omega^1 & \omega^2 & \dots & \omega^{(N-1)} \\ 1 & \omega^2 & \omega^4 & \dots & \omega^{(N-2)(N-1)} \\ 1 & \omega^{(N-1)} & \omega^{(N-1)(N-2)} & \dots & \omega^{(N-1)(N-1)} \end{bmatrix}$$

for 2 qubit $\omega = i$

$$= \frac{1}{\sqrt{N}} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & i & i^2 & i^3 \\ 1 & i^2 & i^4 & i^6 \\ 1 & i^3 & i^6 & i^9 \end{bmatrix}$$

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$$U_{IQFT} = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix}$$

$U_{QFT} \cdot U_{I\otimes T}$

$$\frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & -i & -1 & i \\ 1 & -1 & 1 & -1 \\ 1 & i & -1 & -i \end{pmatrix} \cdot \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & i & -1 & -i \end{pmatrix}$$

$$= \frac{1}{4} \begin{pmatrix} 1x1+1x1+i1x1+i1 & 1x1+i-1+i & 1-i+1-i & 1-i+i-1 \\ 1x1-i-1+i & 1+i+i+i & 1+i-i-i & 1-i+i-i \\ 1x1+i-1 & 1+i-i-i & 1+i+i+i & 1+i-i-i \\ 1x1-i-i & 1-i+i-1 & 1-i-i+i & 1+i+i+i \end{pmatrix}$$

$$= \frac{1}{4} \begin{pmatrix} 4 & 0 & 0 & 0 \\ 0 & 4 & 0 & 0 \\ 0 & 0 & 4 & 0 \\ 0 & 0 & 0 & 4 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

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```
#25.7

from sympy.abc import w
N = int(input("Enter a value", ))
# w = 'w'

matrix=[]
row=[]
for i in range(N): #total row is 3
    row=[]
    for j in range(N): #total column is 3
        # row.append( a \N{SUPERSCRIPT} (i*j))
        row.append( (w)** (i*j) )

    matrix.append(row) #add fully defined column into the row
print( matrix )
```

☞ Enter a value4

```
[[1, 1, 1, 1], [1, w, w**2, w**3], [1, w**2, w**4, w**6], [1, w**3, w**6, w**9]]
```

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Chapter 26

26.1)

26.1)

According to method show in book
for 3 qubit matrix $N = 8$

According to 25.22

$$U_{QFT} = \frac{1}{\sqrt{8}}$$

$$\begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & w^1 & w^2 & w^3 & w^4 & w^5 & w^6 & w^7 \\ 1 & w^{-2} & w^{-4} & w^6 & w^{-8} & w^{10} & w^{-12} & w^{-14} \\ 1 & w^{-3} & w^{-6} & w^{-9} & w^{-12} & w^{15} & w^{-18} & w^{-21} \\ 1 & w^{-4} & w^{-8} & w^{12} & w^{16} & w^{20} & w^{-24} & w^{-11} \\ 1 & w^{-5} & w^{10} & w^{-15} & w^{20} & w^{25} & w^{-30} & w^{-35} \\ 1 & w^{-6} & w^{-12} & w^{-18} & w^{24} & w^{30} & w^{-36} & w^{-42} \\ 1 & w^{-7} & w^{-19} & w^{-21} & w^{25} & w^{-33} & w^{-46} & w^{-4} \end{pmatrix}$$

In page 249 It is mentioned different book has different interpretation and am following method used in this text book

$$\text{Given } w_8 = e^{i \frac{2\pi}{8}} = e^{i \frac{\pi}{4}}$$

$$\text{If we convert it } \cos \frac{\pi}{4} + i \sin \frac{\pi}{4} = \frac{1+i}{\sqrt{2}}$$

$$\text{All even } w_8 = \left(e^{i \frac{2\pi}{8}}\right)^8 = \frac{1+i}{\sqrt{2}} = \frac{i}{\sqrt{2}}$$

$$\begin{array}{cccccccccc}
 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
 & \downarrow & & & & & & & & \\
 1 & \frac{\sqrt{2}}{2} & \left(\frac{\sqrt{2}}{1+i}\right)^2 & \left(\frac{\sqrt{2}}{1+i}\right)^3 & \left(\frac{\sqrt{2}}{1+i}\right)^4 & \left(\frac{\sqrt{2}}{1+i}\right)^5 & \left(\frac{\sqrt{2}}{1+i}\right)^6 & \left(\frac{\sqrt{2}}{1+i}\right)^7 & \left(\frac{\sqrt{2}}{1+i}\right)^8 & \left(\frac{\sqrt{2}}{1+i}\right)^9 \\
 & \downarrow & & & & & & & & \\
 1 & \left(\frac{\sqrt{2}}{1+i}\right)^2 & \left(\frac{\sqrt{2}}{1+i}\right)^4 & \left(\frac{\sqrt{2}}{1+i}\right)^6 & \left(\frac{\sqrt{2}}{1+i}\right)^8 & \left(\frac{\sqrt{2}}{1+i}\right)^{10} & \left(\frac{\sqrt{2}}{1+i}\right)^{12} & \left(\frac{\sqrt{2}}{1+i}\right)^{14} & \left(\frac{\sqrt{2}}{1+i}\right)^{16} & \left(\frac{\sqrt{2}}{1+i}\right)^{18} \\
 & \downarrow & & & & & & & & \\
 1 & \left(\frac{\sqrt{2}}{1+i}\right)^3 & \left(\frac{\sqrt{2}}{1+i}\right)^6 & \left(\frac{\sqrt{2}}{1+i}\right)^9 & \left(\frac{\sqrt{2}}{1+i}\right)^{12} & \left(\frac{\sqrt{2}}{1+i}\right)^{15} & \left(\frac{\sqrt{2}}{1+i}\right)^{18} & \left(\frac{\sqrt{2}}{1+i}\right)^{21} & \left(\frac{\sqrt{2}}{1+i}\right)^{24} & \left(\frac{\sqrt{2}}{1+i}\right)^{27} \\
 & \downarrow & & & & & & & & \\
 1 & \left(\frac{\sqrt{2}}{1+i}\right)^4 & \left(\frac{\sqrt{2}}{1+i}\right)^8 & \left(\frac{\sqrt{2}}{1+i}\right)^{16} & \left(\frac{\sqrt{2}}{1+i}\right)^{32} & \left(\frac{\sqrt{2}}{1+i}\right)^{64} & \left(\frac{\sqrt{2}}{1+i}\right)^{128} & \left(\frac{\sqrt{2}}{1+i}\right)^{256} & \left(\frac{\sqrt{2}}{1+i}\right)^{512} & \left(\frac{\sqrt{2}}{1+i}\right)^{1024} \\
 & \downarrow & & & & & & & & \\
 1 & \left(\frac{\sqrt{2}}{1+i}\right)^5 & \left(\frac{\sqrt{2}}{1+i}\right)^{10} & \left(\frac{\sqrt{2}}{1+i}\right)^{15} & \left(\frac{\sqrt{2}}{1+i}\right)^{20} & \left(\frac{\sqrt{2}}{1+i}\right)^{25} & \left(\frac{\sqrt{2}}{1+i}\right)^{30} & \left(\frac{\sqrt{2}}{1+i}\right)^{35} & \left(\frac{\sqrt{2}}{1+i}\right)^{40} & \left(\frac{\sqrt{2}}{1+i}\right)^{45} \\
 & \downarrow & & & & & & & & \\
 1 & \left(\frac{\sqrt{2}}{1+i}\right)^6 & \left(\frac{\sqrt{2}}{1+i}\right)^{12} & \left(\frac{\sqrt{2}}{1+i}\right)^{18} & \left(\frac{\sqrt{2}}{1+i}\right)^{24} & \left(\frac{\sqrt{2}}{1+i}\right)^{30} & \left(\frac{\sqrt{2}}{1+i}\right)^{36} & \left(\frac{\sqrt{2}}{1+i}\right)^{42} & \left(\frac{\sqrt{2}}{1+i}\right)^{48} & \left(\frac{\sqrt{2}}{1+i}\right)^{54} \\
 & \downarrow & & & & & & & & \\
 1 & \left(\frac{\sqrt{2}}{1+i}\right)^7 & \left(\frac{\sqrt{2}}{1+i}\right)^{14} & \left(\frac{\sqrt{2}}{1+i}\right)^{21} & \left(\frac{\sqrt{2}}{1+i}\right)^{28} & \left(\frac{\sqrt{2}}{1+i}\right)^{35} & \left(\frac{\sqrt{2}}{1+i}\right)^{42} & \left(\frac{\sqrt{2}}{1+i}\right)^{49} & \left(\frac{\sqrt{2}}{1+i}\right)^{56} & \left(\frac{\sqrt{2}}{1+i}\right)^{63}
 \end{array}$$

here given input $1000 \rightarrow$

As we know when apply Hadamard gate

\rightarrow to individually it will be $(1+i + 1-i)$

($\therefore QFT = \text{Hadamard}^{\frac{N}{2}}$ for 1 qubit)

and we know their will be no phase shift

if all bits is zero then we know their
will equal probability in all states if all
the given inputs are in starting state

26.2) Given that we have to prove

$$U_{CPS,\phi} U_{CPS,-\phi} = I$$

$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{i\phi} \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{-i\phi} \end{bmatrix}$$

$$\begin{bmatrix} 1 \times 1 + 0 \times 0 + 0 \times 0 + 0 \times 0 & 0 \times [0+1+0+0] & 0 [0+0+1+0] & 0 [0+0+0+1] \\ 0 \times 1 + 0 \times 0 + 0 \times 0 + 0 \times 0 & 1 [0+1+0+0] & 0 [0+0+1+0] & 0 [0+0+0+1] \\ 0 \times 1 + 0 \times 0 + 0 \times 0 + 0 \times 0 & 0 [0+1+0+0] & 1 [0+0+1+0] & 0 [0+0+0+1] \\ 0 \times 1 + 0 \times 0 + 0 \times 0 + 0 \times 0 & 0 [0+0+0+1] & 0 [0+0+1+0] & e^{i\phi}, e^{-i\phi} \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{-i\phi} + i\phi \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

26.3)

26.3) Given that we have to prove

$$U_{\text{swap}} = U_{\text{swap}}^{-1} \Rightarrow U_{\text{swap}} \cdot U_{\text{swap}} = I$$

$$U_{\text{swap}} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$\text{1st method } U_{\text{swap}}^{-1} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

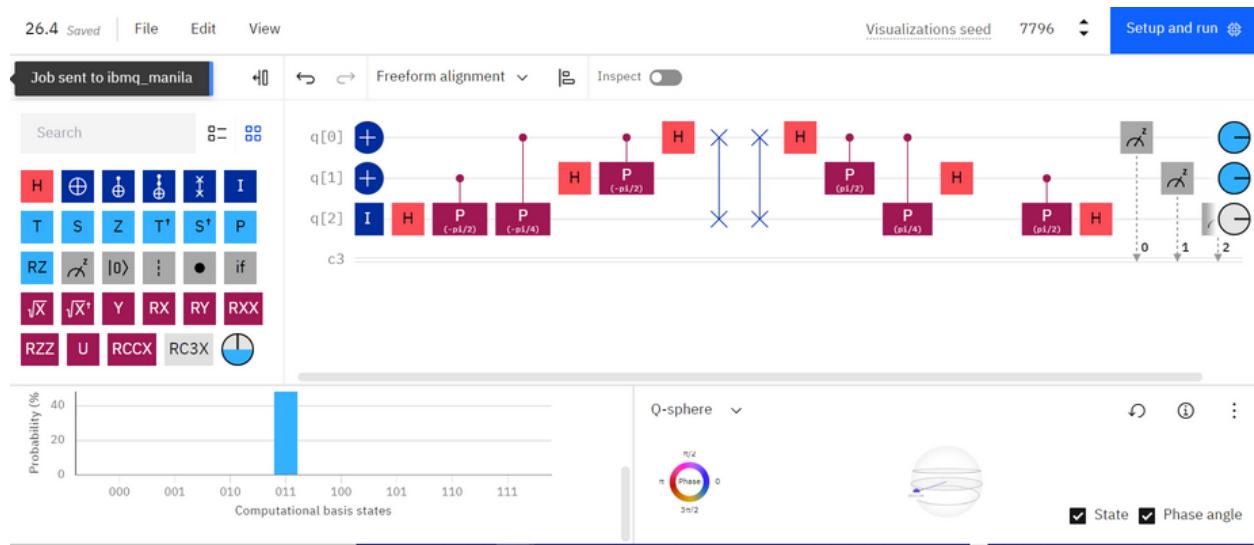
2nd method

$$U_{\text{swap}} \cdot U_{\text{swap}}^{-1} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

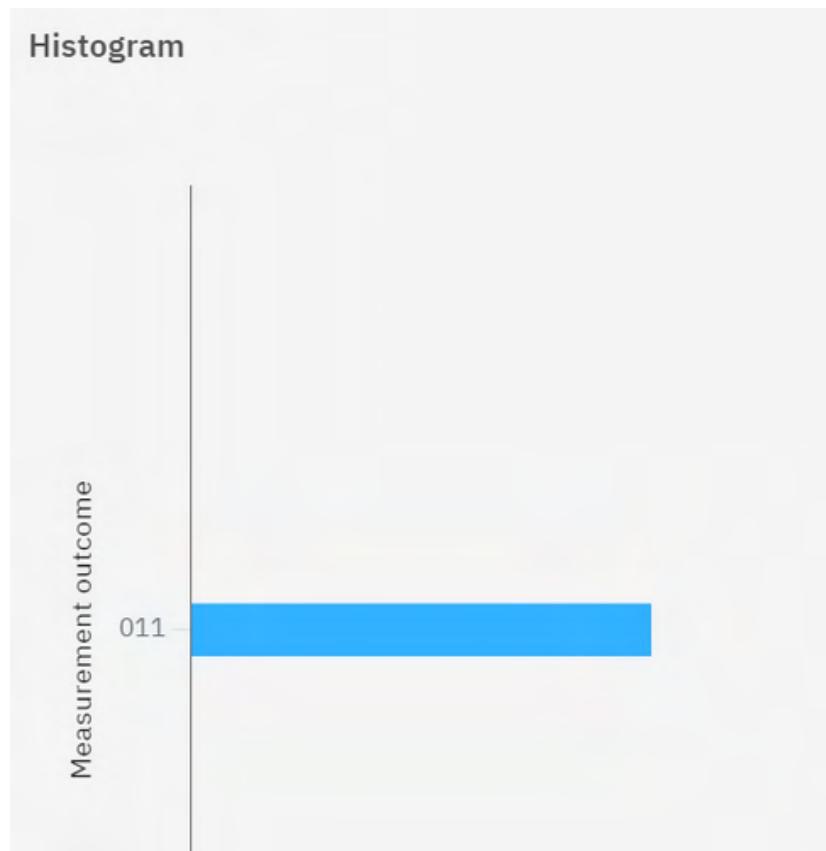
$$= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} = I$$

26.4)

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

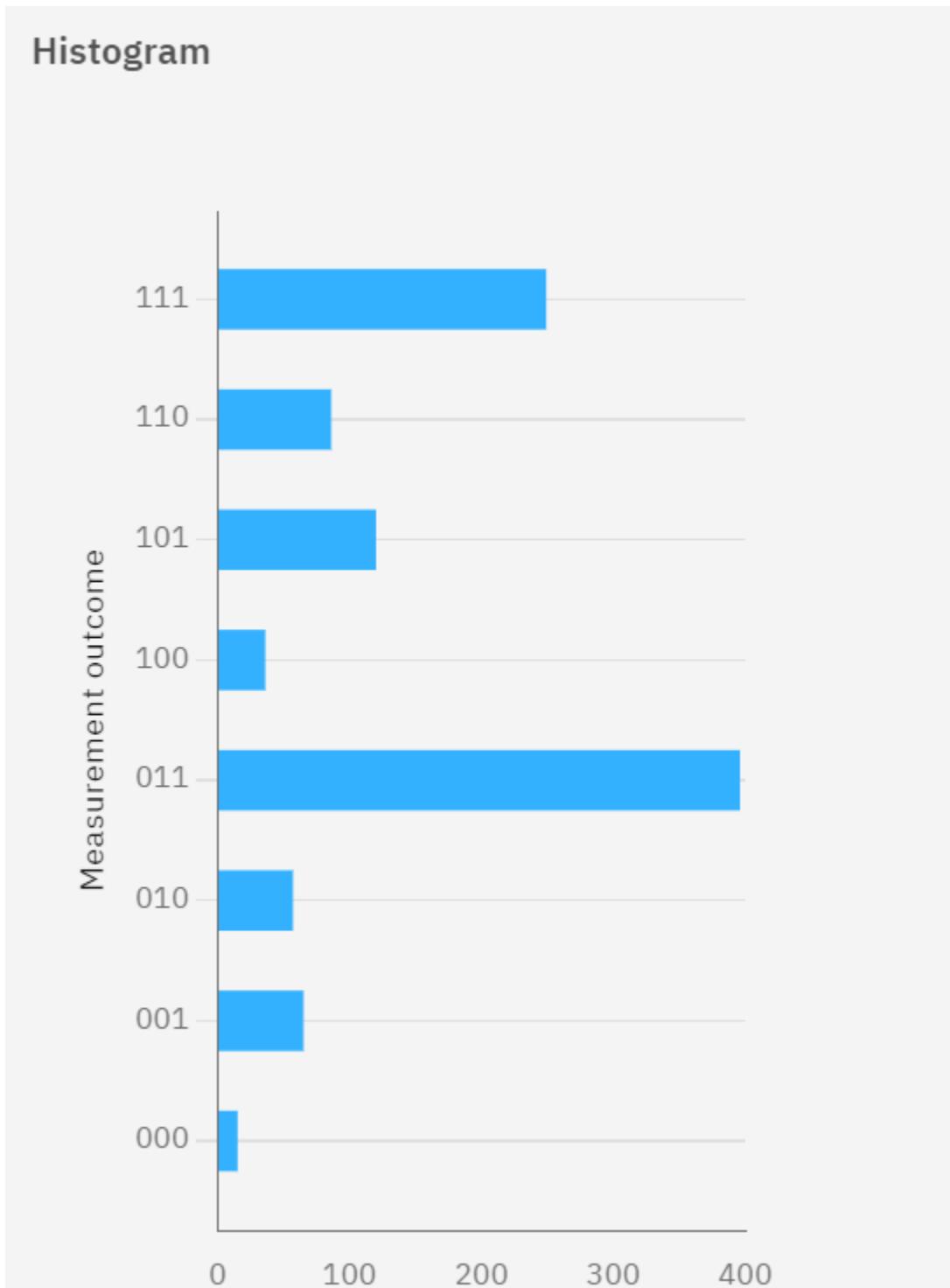


Simulation result



Hardware circuit result

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26.5) according to figure 26.7 we need

$n-1$ U_{c,p,s,p} gates with φ from $\frac{-2\pi}{2^2}$ to $\frac{-2\pi}{2^n}$

If the execution time is $y = \frac{2^n}{2}$ seconds then

I believe as we know QF has $O(n^2)$ speed

26.6) we have derived and elaborated matrix in 26.1

26.7) $|11001100\rangle \xrightarrow{\text{swap}} |0011001\rangle$

Chapter 27

27.1)

27th Chapter27.1) Given σ_z of $|1\rangle\rangle$

From example 27.2 we know

$$|1\rangle = \cos\theta/2 e^{-i\phi/2} |0\rangle + \sin\theta/2 e^{i\phi/2} |1\rangle$$

If we include phase angle as well

$$|1\rangle = e^{i\chi} (\cos\theta/2 e^{-i\phi/2} |0\rangle + \sin\theta/2 e^{i\phi/2} |1\rangle)$$

L, from [27.1]

$$\langle \psi | \sigma_z | 1\rangle = e^{-i\chi} (\cos\theta/2 e^{i\phi/2} \sin\theta/2 e^{i\phi/2})$$

$$\begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} e^{i\chi} \begin{pmatrix} \cos\theta/2 e^{-i\phi/2} \\ \sin\theta/2 e^{i\phi/2} \end{pmatrix}$$

$$= e^{-i\chi+i\chi} \cdot \cos\theta \quad \text{from } \rightarrow 27.6$$

Global phase effect has no effect proved below
 $= 1 \cdot \cos\theta = \cos\theta$

27.2)

27.2) we have three cases proved and we need to prove remaining three cases

$$i) z = -1, x = 0, y = 0 \quad i.e., \theta = \pi, \phi = 0$$

It is also called South pole

$$\begin{aligned} |\Psi_{\theta=\pi, \phi=0}\rangle &= \cos \frac{\theta}{2} e^{-i\phi/2}|0\rangle + \sin \frac{\theta}{2} e^{i(\phi)/2}|1\rangle \\ &= \cos \frac{\pi}{2} e^{-i\phi/2}|0\rangle + \sin \frac{\pi}{2} e^{i(\phi)/2}|1\rangle \\ &= 0 \cdot |0\rangle + 1 \cdot |1\rangle \end{aligned}$$

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

$\frac{\pi}{2}, \frac{\pi}{4}$
 $\frac{1}{\sqrt{2}}$

$|11\rangle$ $|3\pi/4\rangle$

i) consider $x=1, y=2=0$

$\theta = \pi/2, \phi = 0$

$$|\psi_{\theta=\pi/2, \phi=0}\rangle = \cos \frac{\theta}{2} e^{-i\phi/2} |0\rangle + \sin \frac{\theta}{2} e^{i\phi/2} |1\rangle$$

$$= \cos \frac{\pi}{4} e^{-i0} |0\rangle + \sin \frac{\pi}{4} e^{i0} |1\rangle$$

$$= \frac{1}{\sqrt{2}} (|0\rangle + \frac{i}{\sqrt{2}} |1\rangle)$$

$$= |+\rangle$$

iv) consider $y=-1, x=z=0$

$\theta = \pi/2, \phi = -\pi/2$

$$|\psi_{\theta=\pi/2, \phi=\pi/2}\rangle = \cos \frac{\pi}{4} e^{i\pi/4} |0\rangle + \sin \frac{\pi}{4} e^{-i\pi/4} |1\rangle$$

$$= \frac{1}{\sqrt{2}} e^{i\pi/4} (|0\rangle + e^{-i\pi/2} |1\rangle)$$

$$= e^{i\pi/4} \frac{1}{\sqrt{2}} (|0\rangle - i|1\rangle)$$

Here we are ignoring global phase in all scenarios

27.3)

27.3)

$$|\psi\rangle = \cos\theta/2 e^{-i\phi/2}|0\rangle + \sin\theta/2 e^{i\phi/2}|1\rangle$$

$$\begin{pmatrix} \cos\theta/2 e^{-i\phi/2} \\ \sin\theta/2 e^{i\phi/2} \end{pmatrix}$$

i) ~~o~~

$$\langle\psi| \sigma_z |\psi\rangle = (\cos\theta/2 e^{i\phi/2} \sin\theta/2 e^{-i\phi/2})$$

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} \cos\theta/2 e^{-i\phi/2} \\ \sin\theta/2 e^{i\phi/2} \end{pmatrix}$$

$$= (\cos\theta/2 e^{i\phi/2} \sin\theta/2 e^{-i\phi/2}) \begin{pmatrix} \sin\theta/2 e^{i\phi/2} \\ \cos\theta/2 e^{i\phi/2} \end{pmatrix}$$

$$= \cos\theta/2 e^{i\phi/2} \sin\theta/2 e^{-i\phi/2} + \sin\theta/2 e^{-i\phi/2} \cos\theta/2 e^{i\phi/2}$$

$$e^{i\phi} \sin \frac{\theta}{2} - \cos \frac{\theta}{2}$$

$$= \cos\theta/2 \sin\theta/2 \cdot e^{i\phi} + \sin\theta/2 \cos\theta/2 e^{i\phi}$$

$$= \cos\theta_2 \sin\theta_2 (e^{i\phi} + e^{-i\phi})$$

$$= \cos\theta_2 \sin\theta_2 (\cos\phi_2 + i\sin\phi_2 + \cos\phi_2 - i\sin\phi_2)$$

$$= \cos\theta_2 \sin\theta_2 (\cos\phi)$$

$$= \sin\theta \cos\phi$$

$$\text{ii) } |\psi\rangle = \cos\theta_2 e^{i\phi/2} \sin\theta_2 e^{-i\phi/2} \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$$

$$\begin{bmatrix} \cos\theta_2 e^{-i\phi/2} \\ \sin\theta_2 e^{i\phi/2} \end{bmatrix}$$

$$= (\cos\theta_2 e^{i\phi/2} \sin\theta_2 e^{-i\phi/2})$$

$$\begin{bmatrix} -i\sin\theta_2 e^{i\phi/2} \\ i\cos\theta_2 e^{-i\phi/2} \end{bmatrix}$$

$$= \cos\theta_2 e^{i\phi/2} - i\sin\theta_2 e^{i\phi/2} + i\cos\theta_2 e^{-i\phi/2} + \sin\theta_2 e^{-i\phi/2}$$

*

$$i \sin\theta/2 \cos\theta/2 e^{-i\phi} - i \cos\theta/2 \sin\theta/2 e^{i\phi}$$

$$= i \sin\theta/2 \cos\theta/2 (e^{-i\phi} - e^{i\phi})$$

$$= i \sin\theta/2 \cos\theta/2 (\cancel{\cos\theta} - i \sin\theta/2 \cdot \cancel{-\cos\theta} - i \sin\theta)$$

$$= i \sin\theta/2 \cos\theta/2 (-i \sin\phi)$$

$$= -i^2 \sin\theta \sin\phi$$

$$= \sin\theta \sin\phi$$

27.4

 27.4)
Sol

$$\text{Given } |\Psi\rangle = \left(\frac{3}{4} - i\frac{\sqrt{3}}{4}\right) |0\rangle + \left(\frac{\sqrt{3}}{4} + i\frac{1}{4}\right) |1\rangle$$

$$= \frac{\sqrt{3}}{2} \left(\frac{\sqrt{3}-i}{2}\right) |0\rangle + \frac{1}{2} \left(\frac{\sqrt{3}+i}{2}\right) |1\rangle$$

(1)

we know from (27.2)

$$|\Psi\rangle = \cos \theta/2 e^{-i\phi/2} |0\rangle + \sin \theta/2 e^{i\phi/2} |1\rangle \rightarrow (2)$$

we need to compare with (1)

$$\cos \theta/2 = \frac{\sqrt{3}}{2}$$

$$\theta/2 = \pi/6 \Rightarrow \theta = \pi/3$$

$$\sin \theta/2 = \frac{1}{2}$$

$$\theta/2 = \pi/6 \Rightarrow \theta = \pi/3$$

$$e^{-i\phi/2} = \cos \phi/2 - i \sin \phi/2$$

$$= \frac{\sqrt{3}-i}{2}$$

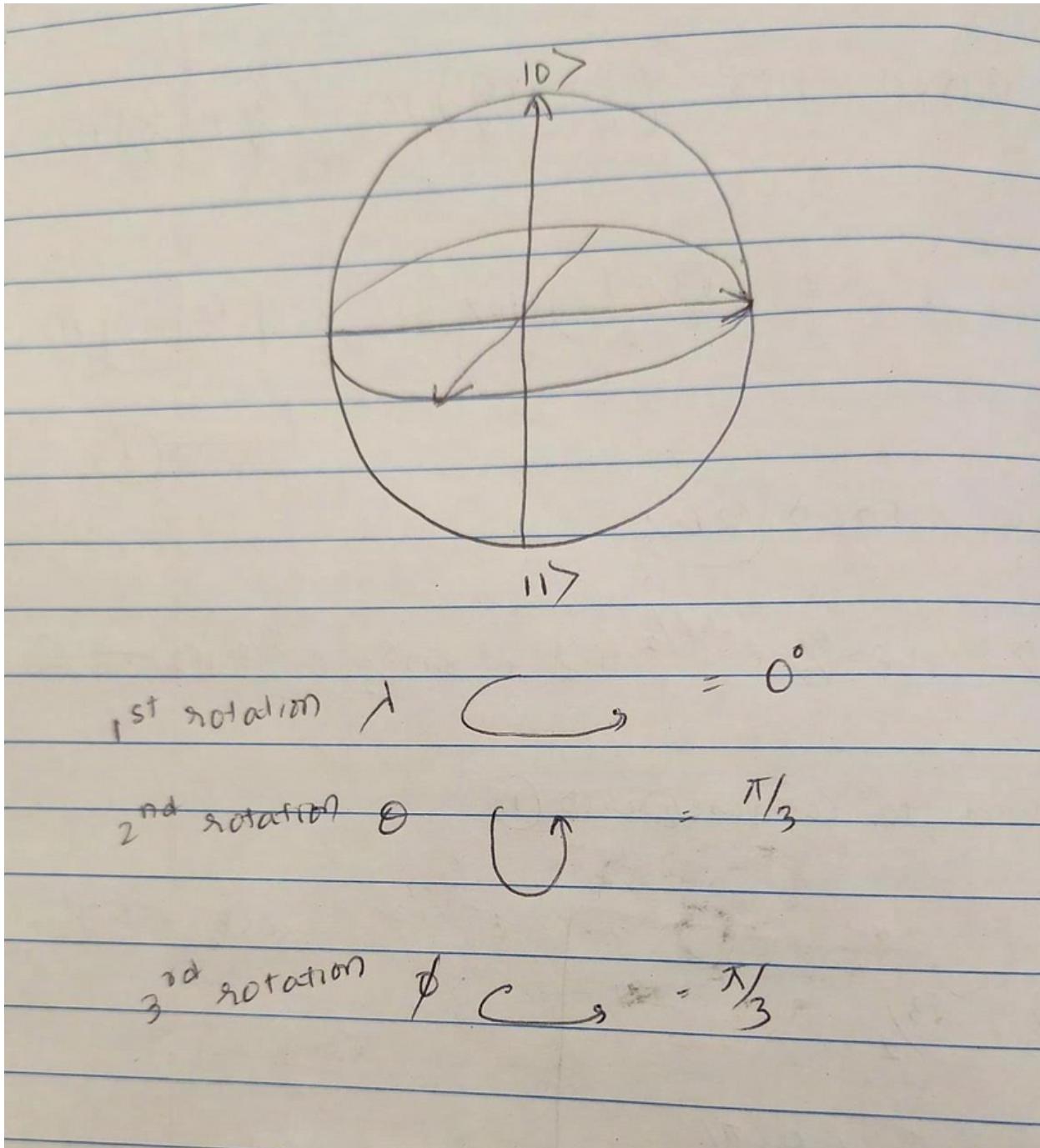
$$\frac{\sqrt{3}}{2} - \frac{i}{2}$$

$$\cos \phi/2 = \frac{\sqrt{3}}{2} \quad \sin \phi/2 = \frac{1}{2}$$

$$\phi/2 = \pi/3$$

$$\phi = \pi/3$$

$$|\Psi\rangle = \cos \pi/6 e^{-i\pi/6} |0\rangle + \sin \pi/6 e^{i\pi/6} |1\rangle$$



27.5)

27.5)

$$H(\psi) = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

$$U_{\theta, \phi, \lambda} = U_{\text{Hadamard}}$$

$$\begin{pmatrix} \cos\theta/2 & -e^{i\lambda} \sin\theta/2 \\ e^{i\phi} \sin\theta/2 & e^{i(\lambda+\phi)} \cos\theta/2 \end{pmatrix} = \begin{pmatrix} 1/\sqrt{2} & 1/\sqrt{2} \\ 1/\sqrt{2} & -1/\sqrt{2} \end{pmatrix}$$

$$\cos\theta/2 = 1/\sqrt{2}, \quad -e^{i\lambda} \sin\theta/2 = 1/\sqrt{2}$$

$$e^{i\phi} \sin\theta/2 = 1/\sqrt{2}, \quad e^{i(\lambda+\phi)} \cos\theta/2 = 1/\sqrt{2}$$

$$\theta/2 = 1/\sqrt{2} \Rightarrow \theta = \pi/2 \quad \text{we can derive } \lambda = \pi$$

$$\text{and } -e^{i\lambda} = 1 \Rightarrow \phi = 0$$

$$\Rightarrow U_{\theta, \phi, \lambda} = U_{\pi/2, 0, \pi}$$

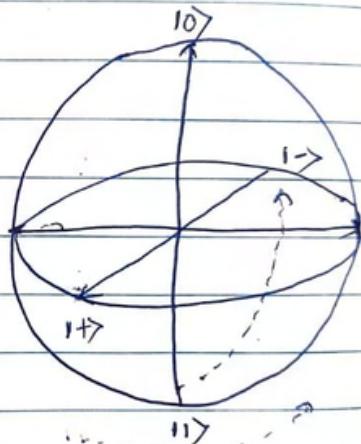
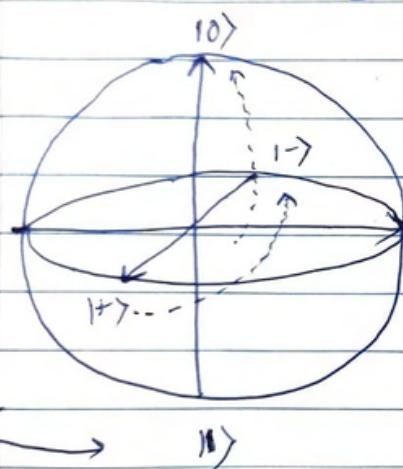
27.6)

27.6) Hadamard applied to $|+\rangle$ to $|-\rangle$

We know Hadamard's $U\theta, \phi, \lambda$ is $U_{\pi/2}, 0, \pi$

$$H \rightarrow |+\rangle \rightarrow |0\rangle$$

$$H \rightarrow |+\rangle \rightarrow |-\rangle$$



1st rotation $\lambda = \pi$

2nd rotation $\theta = \pi/2$

3rd rotation $\phi = 0$

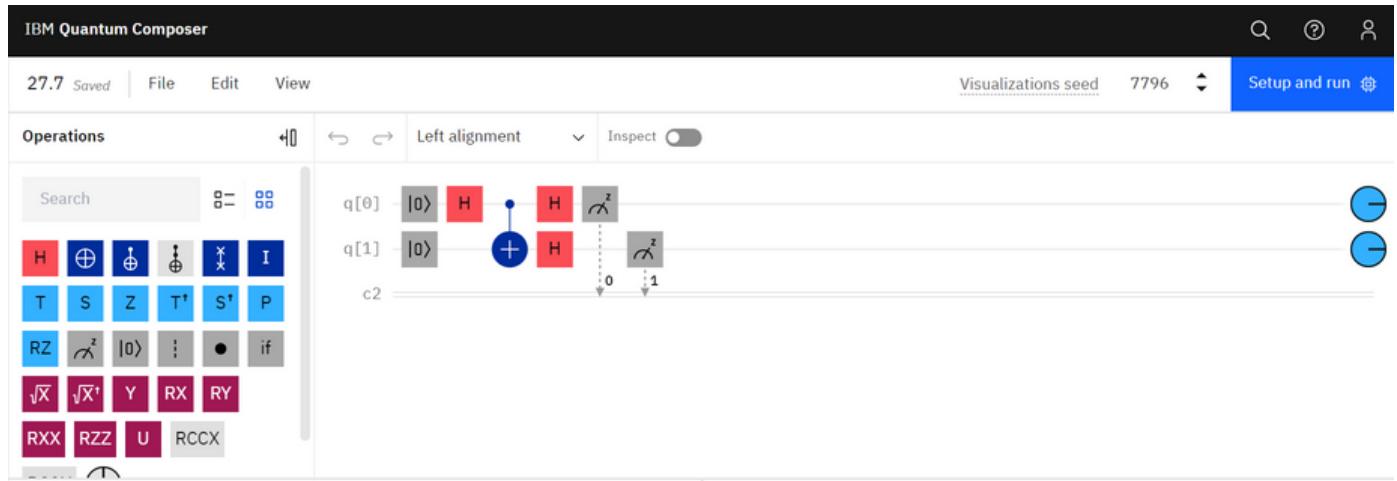
1st rotation $\lambda = \pi$ (No use)

2nd rotation $\theta = \pi/2$

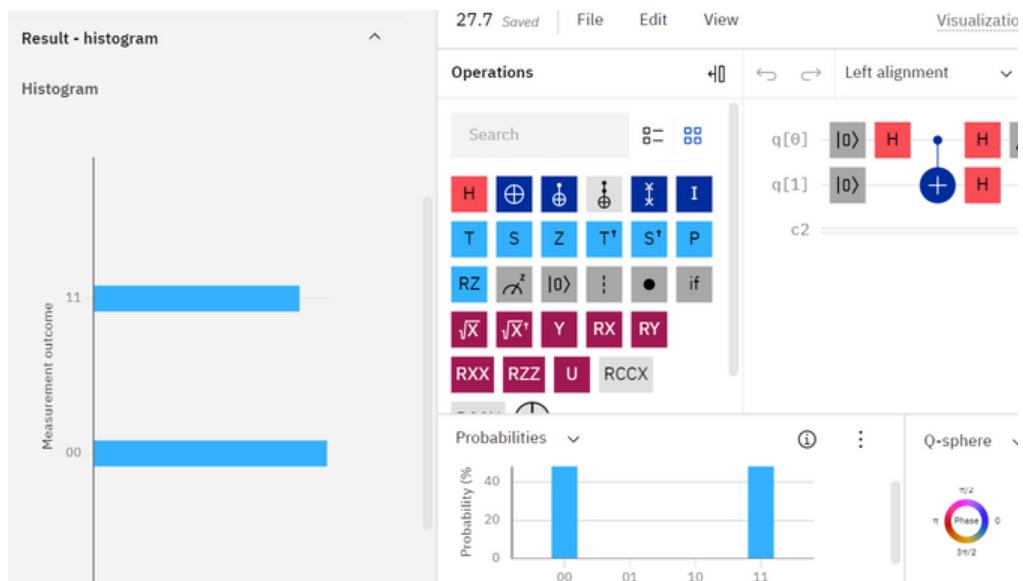
3rd rotation $\phi = 0$

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

27.7

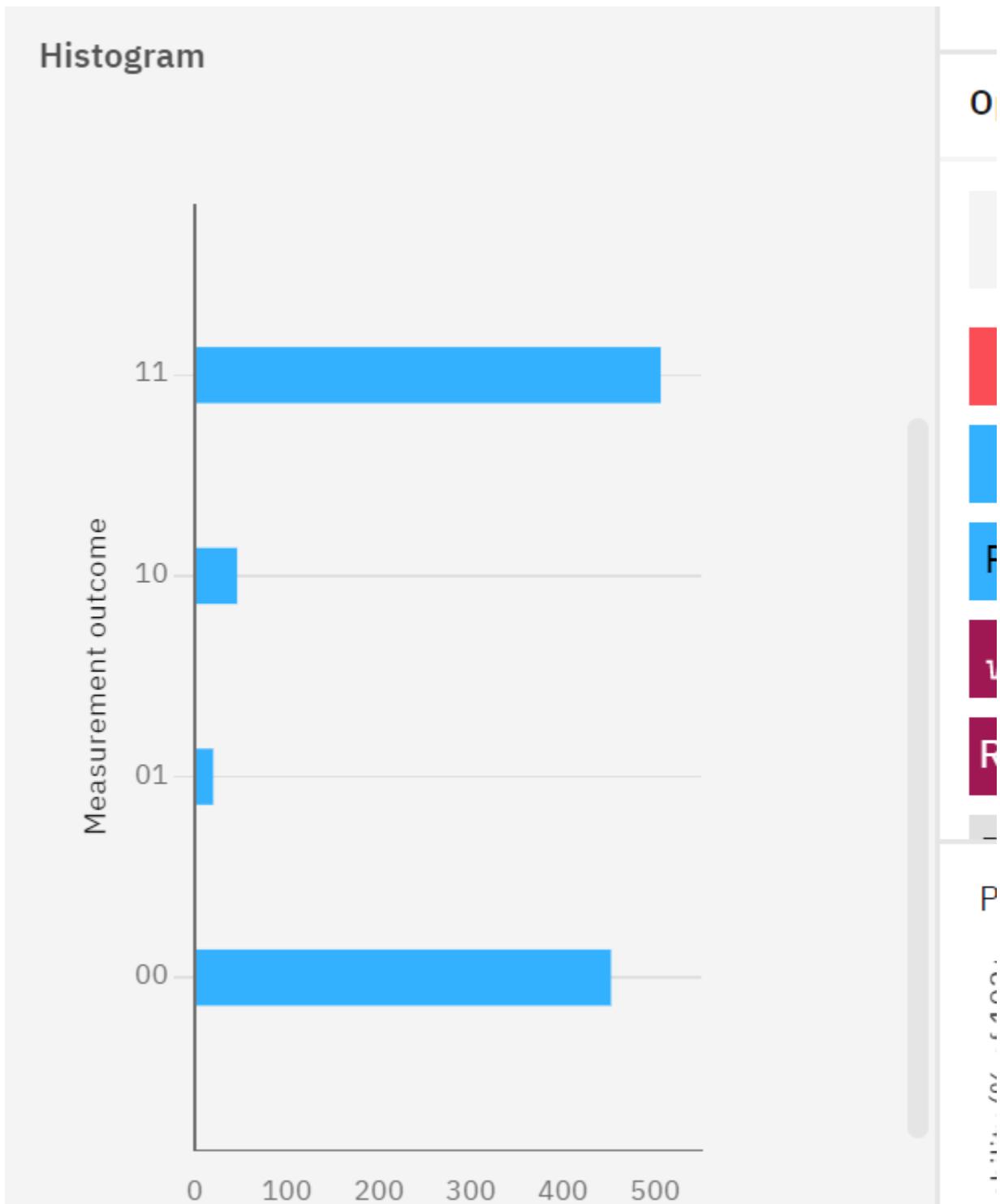


Result



Hardware execution result

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28.1)

Given $\frac{1}{\sqrt{2}} \begin{pmatrix} -i & -i \\ -i & i \end{pmatrix}$

we can extract $-i$ from the above gate it becomes hadamard gate

$$-i \cdot \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

$$-i \begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix} = e^{i\lambda} \begin{pmatrix} \cos \theta/2 & -e^{-i\lambda} \sin \theta/2 \\ e^{i\lambda} \sin \theta/2 & e^{i(\lambda+\phi)} \cos \theta/2 \end{pmatrix}$$

$$\cos \theta/2 = \frac{1}{\sqrt{2}} \Rightarrow \theta/2 = \frac{\pi}{4} \quad [\theta = \frac{\pi}{2}]$$

$$= e^{-i\lambda} \sin \theta/2 \Rightarrow \boxed{\sin \theta/2 = \frac{1}{\sqrt{2}}} \Rightarrow -e^{-i\lambda} \Rightarrow \lambda = \pi$$

$$e^{i\phi} \sin \theta/2 = e^{i\phi} \Rightarrow \phi = 0$$

$$e^{i\gamma} = -i \Rightarrow -(\cos \gamma + i \sin \gamma) = \gamma = \frac{\pi}{2}$$

$$\boxed{U_{\theta, \phi, \lambda, \gamma} = U_{\frac{\pi}{2}, 0, \pi, \frac{\pi}{2}}}$$

28.2)

Introduction to Quantum computing (From a layperson to programmer in 30 steps)

Given $\hat{U}_0, \phi, \lambda, \gamma$

$$\hat{U} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & e^{i\gamma \cos \theta / 2} & e^{-i(\lambda + \gamma) \sin \theta / 2} \\ 0 & 0 & e^{i(\phi + \gamma) \sin \theta / 2} & e^{i(\lambda + \phi + \gamma) \cos \theta / 2} \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

From 28.7 we can derive

We know :-

$$c_{2,3} = \langle 111 | 101 | \hat{U} | 111 \rangle = -e^{i(\lambda + \gamma) \sin \theta / 2}$$

which is in $|110\rangle$
basis
after inner product

$$c_{3,3} = \langle 111 | 111 | \hat{U} | 111 \rangle = e^{i(\lambda + \phi + \gamma) \cos \theta / 2}$$

$$c_{2,2} = \langle 101 | 101 | \hat{U} | 110 \rangle = e^{i\gamma \cos \theta / 2}$$

$$c_{3,2} = \langle 101 | 111 | \hat{U} | 110 \rangle = e^{i(\phi + \gamma) \sin \theta / 2}$$

28.3)

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28.3) Given

$$C - U_{\theta, \phi, \lambda, \gamma} = |0\rangle\langle 0| \otimes I + |1\rangle\langle 1| \otimes U_{\theta, \phi, \lambda, \gamma}$$

$$|0\rangle\langle 0| \Rightarrow |0\rangle\langle 0| \Rightarrow \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}$$

$$|1\rangle\langle 1| \Rightarrow |1\rangle\langle 1| \Rightarrow \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} + \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} \cdot e^{i\gamma} \begin{bmatrix} \cos \theta/2 & e^{i\phi/2} \\ e^{i(\theta+\phi)} & \sin \theta/2 \end{bmatrix} \begin{bmatrix} \cos \theta/2 & e^{i\phi/2} \\ e^{i(\theta+\phi)} & \sin \theta/2 \end{bmatrix}$$

Anything multiplied with Identity is same

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & e^{i\gamma} \cos \theta/2 & e^{i(\lambda+\gamma)} \sin \theta/2 \\ 0 & 0 & e^{i(\phi+\gamma)} \sin \theta/2 & e^{i(\lambda+\phi+\gamma)} \cos \theta/2 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & e^{i\gamma} \cos \theta/2 & e^{i(\lambda+\gamma)} \sin \theta/2 \\ 0 & 0 & e^{i(\phi+\gamma)} \sin \theta/2 & e^{i(\lambda+\phi+\gamma)} \cos \theta/2 \end{bmatrix}$$

28.4)

28.4)

We know

$$U_{ps, \pi} |e_1\rangle = +1 \begin{pmatrix} 1 \\ 0 \end{pmatrix} = +1 |e_0\rangle$$

$$Z |e_0\rangle = e^{i2\pi \cdot 0} |e_0\rangle = |e_0\rangle$$

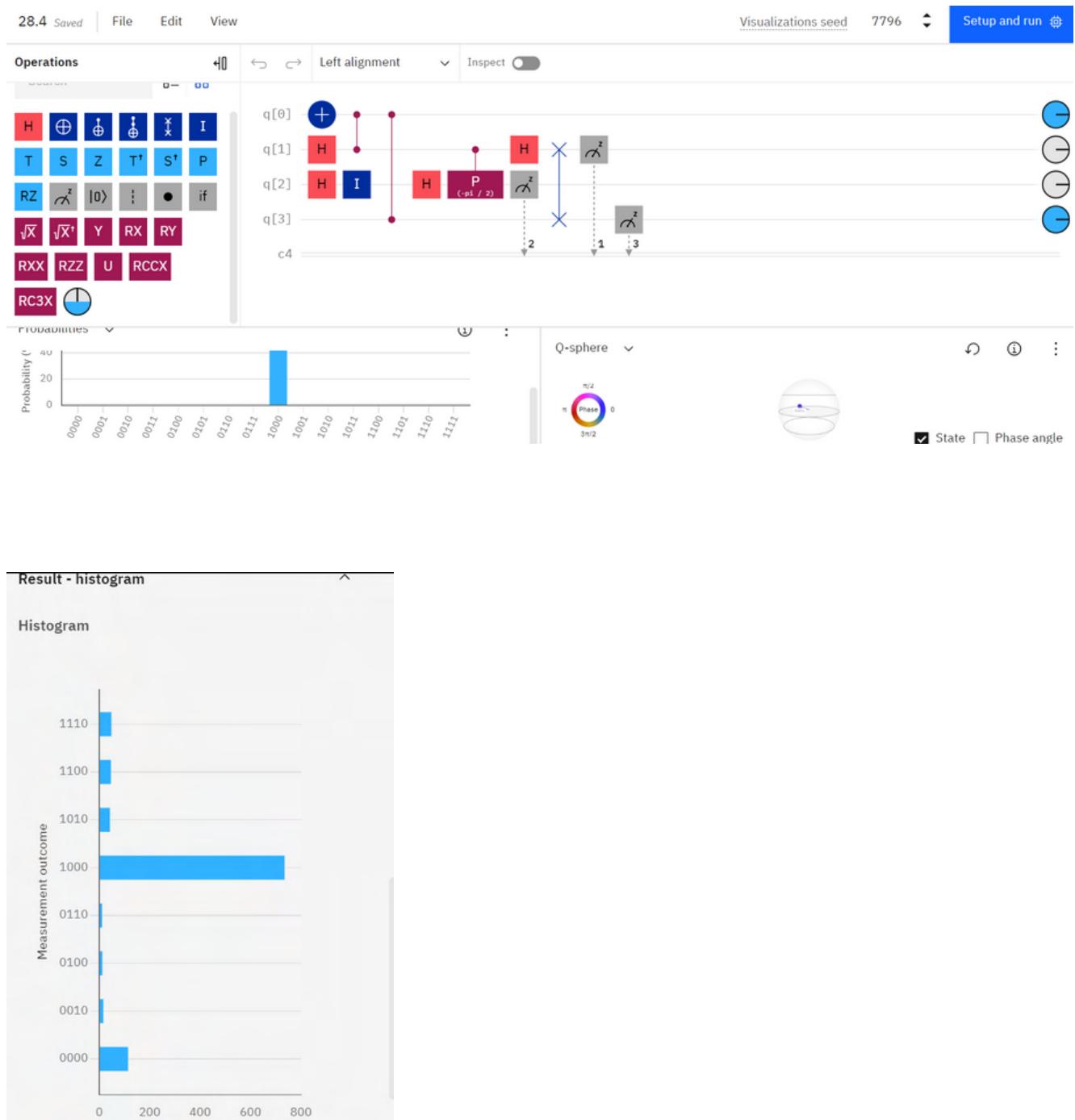
$$Z^2 |e_0\rangle = (e^{i2\pi \cdot 0})^2 |e_0\rangle = |e_0\rangle$$

$$Z = e^{i2\pi \cdot 0} \quad \text{i.e., } e^0 = \cos 0 + i \sin 0 = 1$$

If control Qubit is 0 $\otimes e^0$ it does nothing

28.5)

Introduction to Quantum computing (From a layperson to programmer in 30 steps)



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28.6)

Equation in $2\pi\psi = 2\pi k / 2^n$ is proved
in 28.3.2

$$2\pi\psi = 2\pi k / 2^n$$

$$\text{Here } |k\rangle = |10\rangle = |2\rangle$$

$$2\pi\psi = 2\pi k / 2^n = 2\pi \cdot 2 / 4 = \pi$$

which has ^{been} proved

29.1) By using method in (29.2)
 $177 \bmod 25$

$$177 = 25 \times 7 + 2$$

i.e., 2

29.2 we can find GCD by euclidean method

firstly we have 1242×242

i) Divide 1242 by 242 and we get
remainder 32

ii) Divide 242 by 32 (until last value)
 $242 \div 32 \rightarrow$ remainder is 18

iii) Divide 32 by 18
 $32 \div 18 \rightarrow$ remainder 14

iv) Divide 18 by 14

$18 \div 14 \rightarrow$ remainder is 4 $\neq 0$

v) Divide 18 by 4

$18 \div 4 \rightarrow$ remainder is 2 $\neq 0$

vi) Divides 4 $\div 2 \rightarrow$ remainder is 0



This will become GCD as remainder is zero (value \oplus GCD $\rightarrow 2$)

29.3)

```

# Shors simplified algorithm in python but cant handle big prime numbers
#29.3
# gcf is calculated
def get_gcf(a, b):
    while a != b:
        if a > b:
            a -= b
        elif b > a:
            b -= a
    return a

# I have used below integers as a in our text book we can add any number of numbers

# 12 is not suitable a so i have used incrementation to get better a

Randomint = [8, 10, 12]

for g in Randomint:
    N = 63
    done = False
    p = 0
    p_values = []
    while not done:
        p += 1
        r = g**p % N
        if g**p % N == 1:
            print(g, p)
            p_values.append(p)
        if len(p_values) >= 1:
            p = p_values[0]
            if p_values[0] % 2 != 0:
                g += 1
                p = 0
                p_values = []
            else:
                x = g**(p/2)+1
                y = g**((p/2)-1)
                a = int(get_gcf(N, x))
                b = int(get_gcf(N, y))
                if a == 1 or b == 1:
                    g += 1
                    p = 0
                    p_values = []
                else:
                    print(g, p, a, b)
                    done = True

```

```
if p > 500:  
    g += 1  
    p = 0  
    p_values = []  
  
print(a, "x", b, "=", N)
```

) 8 2
8 2 9 7
9 x 7 = 63
10 6
10 6 7 9
7 x 9 = 63
13 6
13 6 7 9
7 x 9 = 63

29.4)

29.4)

Given $n=1$

$|\varepsilon_0\rangle$ is initial setup and we have $n=1$ i.e., 2 for MSB & 1 for LSB

Basis state is

$$|\varepsilon_0\rangle = |0\rangle^{\otimes 2} \otimes |0\rangle = |00\rangle \otimes |0\rangle = |000\rangle$$

In first step

we apply two Hadamard gates in MSB

$$|\varepsilon_1\rangle = (H^{\otimes 2} \otimes I) (|0\rangle^{\otimes 2} |0\rangle)$$

$$\text{we know } = \frac{1}{2^n} \sum_{x=0}^{2^n-1} |x\rangle_{2n} |0\rangle_n$$

Here n is L so when $n=1$

$$\begin{aligned} &= \frac{1}{2} \sum_{x=0}^3 |x\rangle_2 |0\rangle = \frac{1}{2} (|0\rangle + |1\rangle + |2\rangle + |3\rangle) \otimes |0\rangle \\ &= \frac{1}{2} (|00\rangle + |01\rangle + |10\rangle + |11\rangle) \otimes |0\rangle \end{aligned}$$

Second step

$|\varepsilon_2\rangle$ is basically applying oracle to the above state

$$|\varepsilon_2\rangle = U_f |\varepsilon_1\rangle$$

Here we are implementing quantum oracle as XOR oracle

Here we know $f(|00\rangle)$, $f(|01\rangle)$ and $f(|11\rangle)$ is zero (0) and $f(|10\rangle)$ is 1

$$\begin{aligned} &= \frac{1}{2} \sum_{x=0}^{2^n-1} |x\rangle_{2n} |f(x)\rangle_n \\ &= \frac{1}{2} (|00\rangle \otimes f(|00\rangle) + |01\rangle \otimes f(|01\rangle) + |10\rangle \otimes f(|10\rangle) \\ &\quad + |11\rangle \otimes f(|11\rangle)) \\ &= \frac{1}{2} (|00\rangle \otimes |0\rangle + |01\rangle \otimes |0\rangle + |10\rangle \otimes |1\rangle + |11\rangle \otimes |0\rangle) \end{aligned}$$

$$|\varepsilon_2\rangle = \frac{1}{2} (|000\rangle + |010\rangle + |101\rangle + |110\rangle)$$

Third step

QFT is applied to above state and it is the last step before measurement

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We know $U_{QFT} |x\rangle = \frac{1}{\sqrt{2^n}} \sum_{k=0}^{2^n-1} \omega^{-kx} |k\rangle$

$$|\varepsilon_3\rangle = (U_{QFT} \otimes I^{\otimes n}) |\varepsilon_2\rangle$$

$$= \frac{1}{2^n} \sum_{x=0}^{2^n-1} \sum_{k=0}^{2^n-1} \omega^{-kx} |k\rangle_{2^n} |t(x)\rangle_n \quad (\because \omega = e^{i 2\pi / 2^n})$$

$$= \frac{1}{2^2} \sum_{x=0}^{2-1} \sum_{k=0}^{2-1} e^{-i 2\pi kx / 4} |k\rangle |t(x)\rangle$$

$$= \frac{1}{4} \sum_{x=0}^3 \sum_{k=0}^3 e^{-i 2\pi kx / 4} |k\rangle |t(x)\rangle$$

$$= \frac{1}{4} \left[(|000\rangle + |010\rangle + |110\rangle + |110\rangle) \right.$$

$$+ [|000\rangle + i|010\rangle + |100\rangle + i|110\rangle) + [|001\rangle - |011\rangle + |110\rangle + |111\rangle] \right]$$

$$\left. + [|000\rangle + i|010\rangle - |100\rangle - i|110\rangle) \right]$$

If is not possible to have (in my opinion)
 as we need to have atleast 5 qubits to execute
 Shors Algorithm

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