

Unit I

Introduction



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UNIT-1:

Title and Content Layout with List

- **The Role of programming Languages:**

- Towards Higher-level Languages
- Programming Paradigms
- Criteria for good
- language design and Language implementation.

- **Language Description :**

- Expression notation
- Abstract syntax tree
- Context free Grammars

What is a Programming Language?

- a tool for instructing machines
- a means of communicating between programmers
- a vehicle for expressing high-level designs
- a notation for algorithms
- a way of expressing relationships between concepts
- a tool for experimentation
- a means for controlling computerized devices

Language Designers

- Balance
- ... making computing **convenient** for *people with*
- and making **efficient** use of computing machines

Levels

- Gross distinction between programming language
- based on readability
- based on independence
- based on purpose (specific ... general)

Generations of Programming Languages [Levels]

1GL: machine codes

2GL: symbolic assemblers

3GL: (machine-independent) imperative languages
(FORTRAN, Pascal, C ...)

4GL: domain specific application generators

5GL: AI languages ...

Each generation is at a higher level of abstraction

Machine Level [is unintelligible]

- 00000010101111001010
- 00000010101111001000
- 00000011001110101000
- Can you tell what this code fragment does?
- Can it be executed on any machine?

Assembly Language

- LD R1,"0"
- LD R2, M
- ST R2, R1
- ... real assembly used mnemonics
- Add A(M), Had to do your own indexing
- What does this program do?

Basic Concepts of a RAM machine

- Memory: addresses, contents
- Program: instructions
- input/output:(files)

2000 200A
A 20 c 10

$$A = 3 + c$$

Lvalue → address

Rvalue → contents

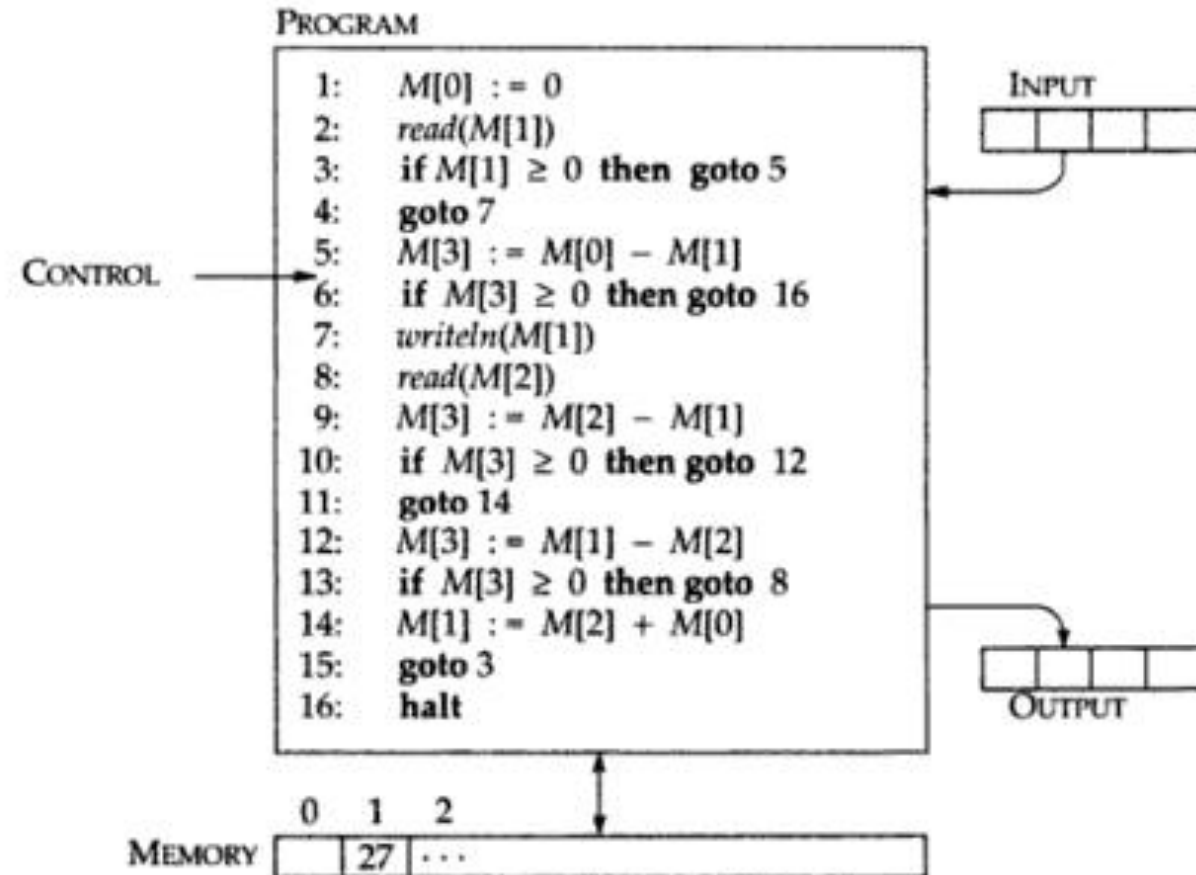


Figure 1.1 A random-access machine (RAM).

High Level [Benefits]

- Readable familiar notations
- machine independence
- availability of program libraries
- consistency check (check data types)

Problems of Scale

- Changes are easy to make
- isolated program fragments can be understood
- BUT... one small bug can lead to disaster
- read the NOT story about Mariner rockets
- Notice how the chairman does not understand that a “small” problem can lead to devastating result and why it was not caught

Bugs

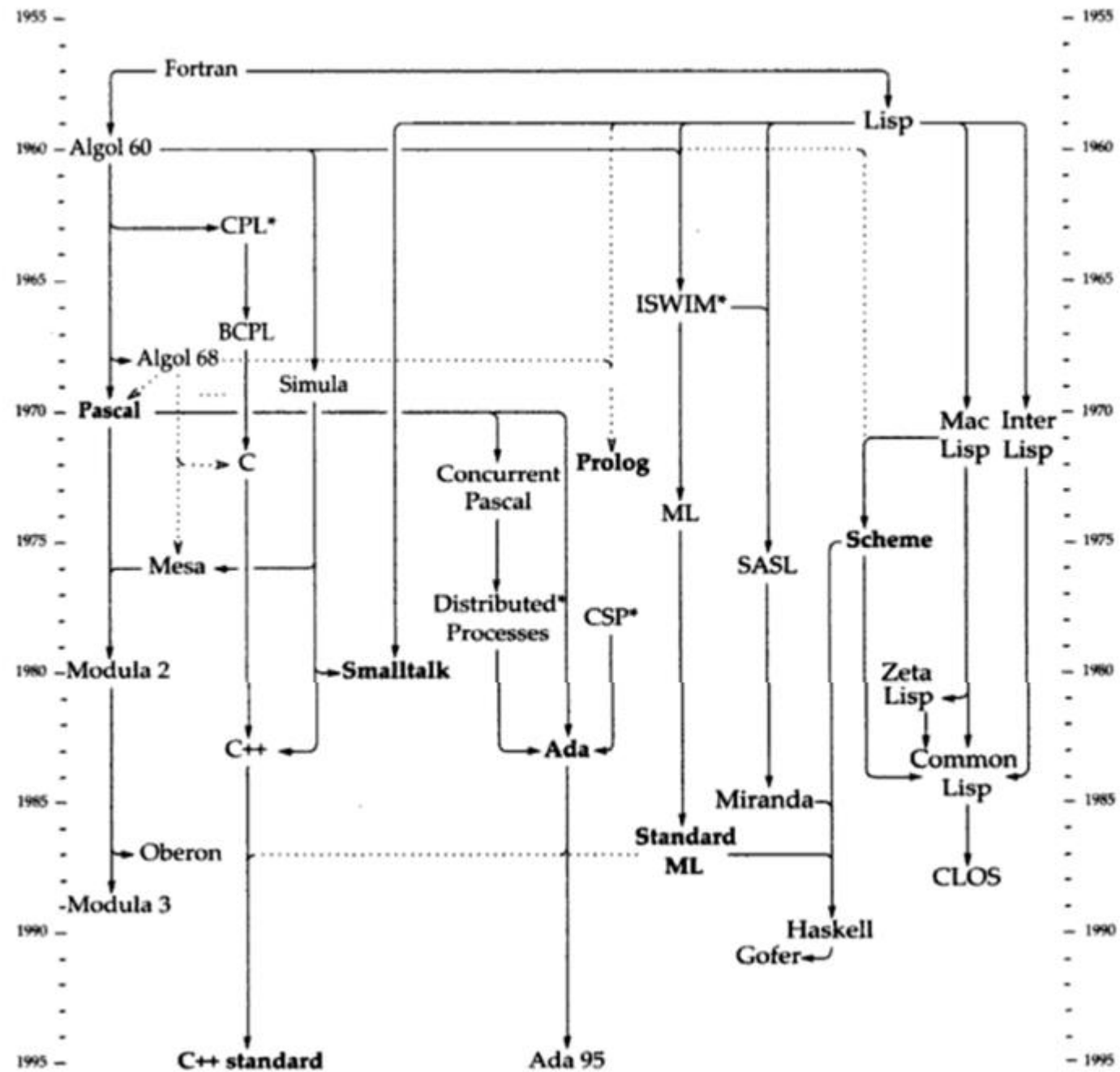
- *Programming testing can be used to show the presence of bugs, but never their absence!*
- Dijkstra
- Programming Languages can help
- readable and understandable
- organize such that parts can be understood

Role of Programming Languages

- Art (science) of programming is organizing complexity
- Must organize in such a way that our limited powers are sufficient to guarantee that the computation will establish the desired effect
- (Dijkstra - structured programming, sometimes referred to as goto-less programming)

Programming Paradigms

- **Imperative** - action oriented, sequence of actions
 - examples: C, Pascal, Basic, Fortran
- **Functional** - LISP, symbolic data processing,
- **Object-Oriented**-
 - examples: C++, Java, Smalltalk
- **Logic** - Prolog, logic reasoning



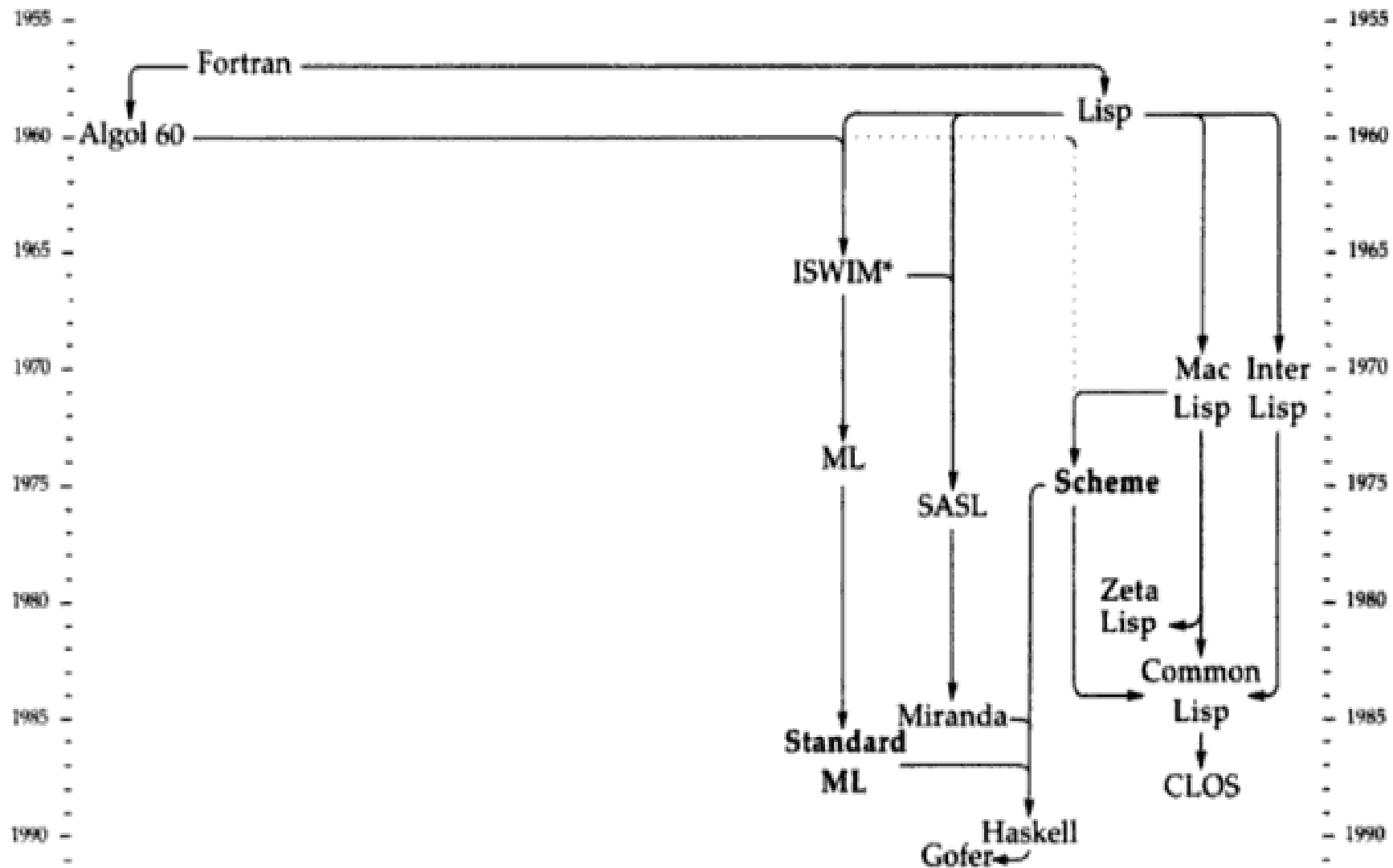
Imperative[Procedural] programming



Imperative[Procedural] programming

- Fortran 1955-1960 for scientific programming
 - familiar notations
 - efficiency
- Toward common language
- Algol, Algol 60 1960s. Imperative languages is viewed as Algol family.
- Pascal was designed as a teaching language
- C 1972, created by Dennis Ritchie. C provides a rich set of operators, a terse syntax, and efficient access to the machine. C implements UNIX system.

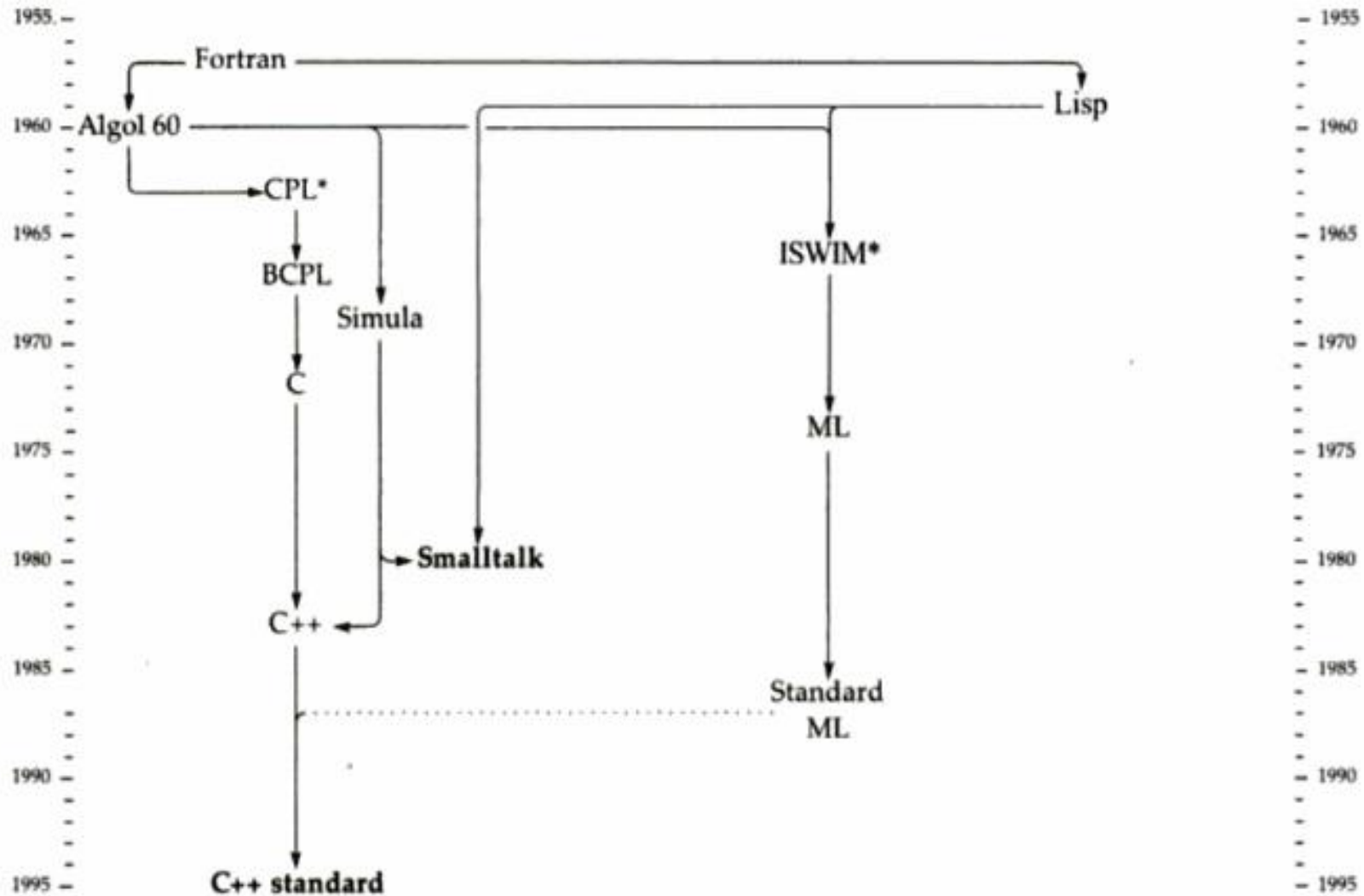
Functional programming



Functional Programming: Lisp 1958

- **Lisp** is an acronym derived from **List Processor**.
- Example:
 - (+ 2 3)
 - (Shakespeare wrote (the tempest))
 - (cdr (a b c d e))
- Program and data are both represented by lists in LISP
- Creation of common Lisp in 1984
- **CLOS** is an object-oriented extension; the full name is **Common Lisp Object System**

Object-oriented programming



Objected-Oriented Programming (1961-1967)

- From **Simula's** origins in simulation
- Key concepts from Simula is that of class of object
 - A traffic simulation can have many cars and many trucks
 - The simulation has
 - objects of the class of cars and
 - objects of the class of trucks.
 - All of which share certain properties
 - Car and trucks are all vehicles
 - The classification of objects into classes and subclasses is central to object oriented programming

Objected-Oriented Programming (1961-1967)

- Smalltalk is an interactive system with graphic user interface
- C++ was designed to bring the benefits of objects to imperative programming in C.
- Java (Network Language)
 - Embedded in your Browser
 - Platform-independent capability
 - Using virtual machine

Programming Paradigms

A programming language is a problem-solving tool.

Imperative style:	program = algorithms + data <i>good for decomposition</i>
Functional style:	program = functions ° functions <i>good for reasoning</i>
Logic programming style:	program = facts + rules <i>good for searching</i>
Object-oriented style:	program = objects + messages <i>good for modeling(!)</i>

Other styles and paradigms: blackboard, pipes and filters, constraints, lists, ...

Procedural vs Functional

- Program: a sequence of instructions for a von Neumann m/c.
 - Computation by instruction execution.
 - Iteration.
 - Modifiable or updateable variables.
- Program: a collection of function definitions (m/c independent).
 - Computation by term rewriting.
 - Recursion.
 - *Assign-only-once* variables.

Functional Style : Illustration

- Definition : Equations

$$\text{sum}(0) = 0$$

$$\text{sum}(n) = n + \text{sum}(n-1)$$

- Computation : Substitution and Replacement

$$\text{sum}(2)$$

$$= 2 + \text{sum}(2-1)$$

$$= \dots$$

$$= 3$$

Paradigm vs Language

■ Imperative Style

```
i := 0; sum := 0;  
while (i < n) do  
    i := i + 1;  
    sum := sum + i  
end;
```

- Storage efficient

■ Functional Style

```
func sum(i:int) : int;  
    if i = 0  
    then 0  
    else i + sum(i-1)  
end;
```


- No Side-effect

Role of Variables

- Imperative (read/write)

i	<input type="text"/>	0	1	2	3	...
sum	<input type="text"/>	0	1	3	6	...

- Functional (read only)

i1	<input type="text" value="3"/>		<input type="text" value="6"/>	sum1
i2	<input type="text" value="2"/>		<input type="text" value="3"/>	sum2
i3	<input type="text" value="1"/>		<input type="text" value="1"/>	sum3

...

Bridging the Gap

- Tail recursive programs can be automatically optimized for space by translating them into equivalent while-loops.

```
func sum(i : int, r : int) : int;  
    if i = 0 then r  
    else sum(i-1, n+r)  
end
```

- Scheme does not have loops.

Analogy: Styles vs Formalisms

- Iteration
- Tail-Recursion
- General Recursion
- Regular Expression
- Regular Grammar
- Context-free Grammar

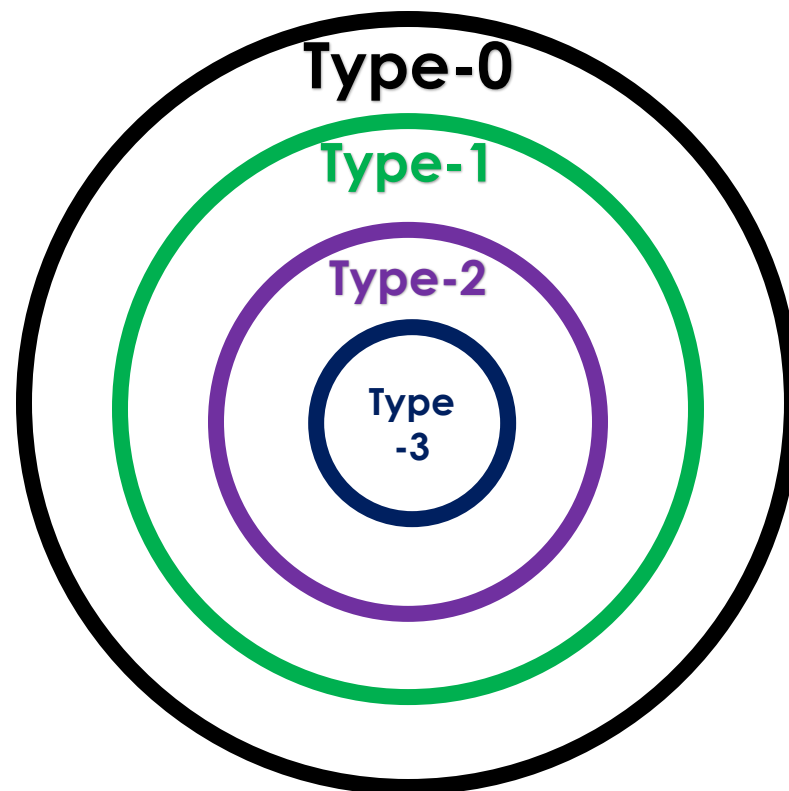
CHOMSKY HIERARCHY



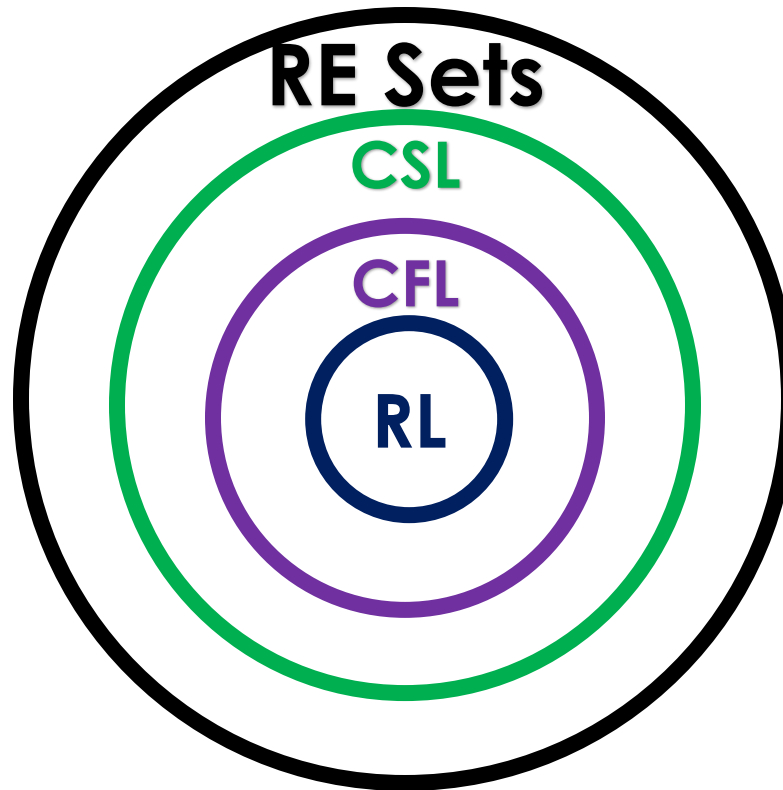
CHOMSKY HIERARCHY

- Four Language

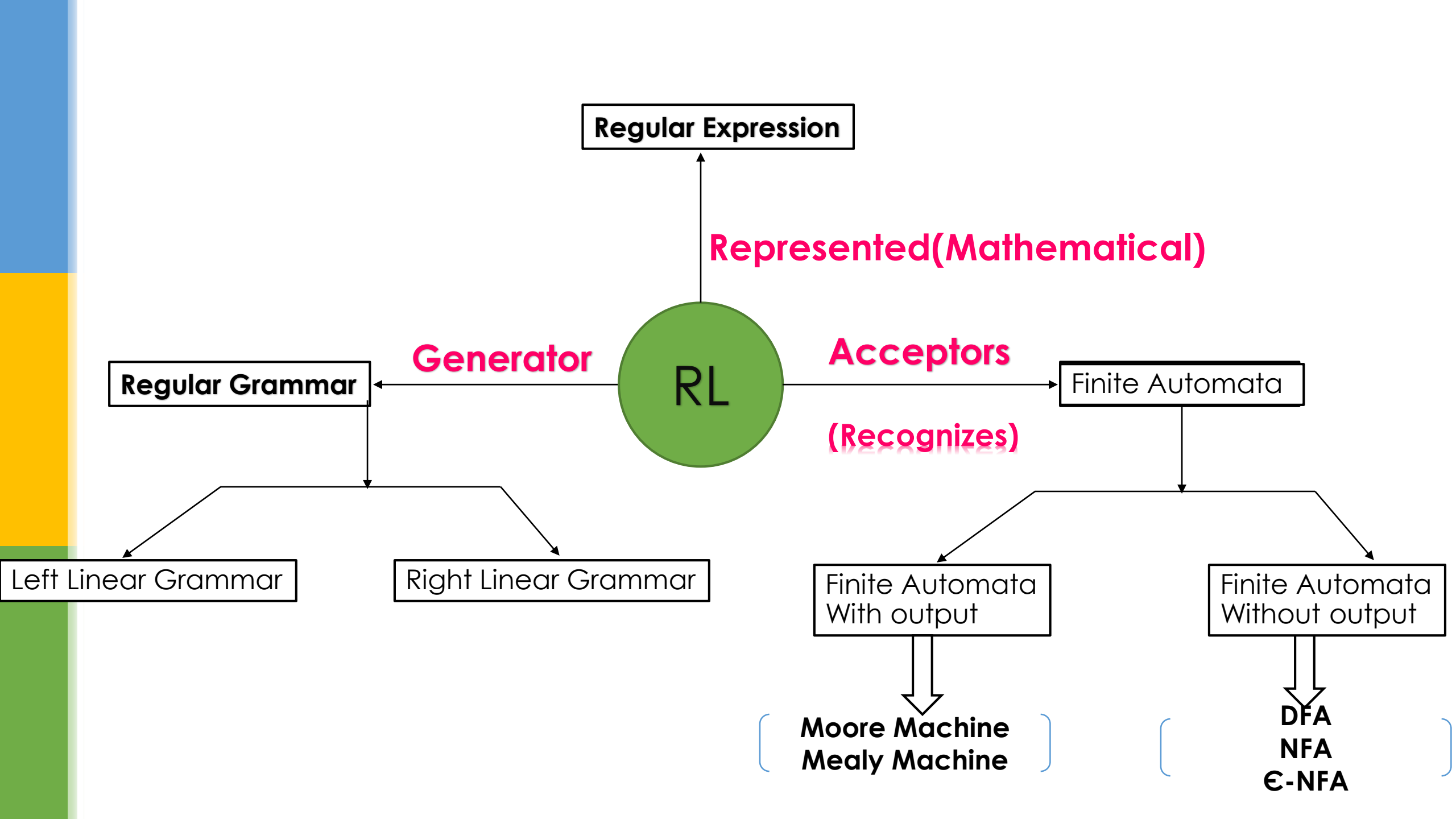
1. Type 0 → Unrestricted Lang[or Recursive Enumerable Lang]
2. Type 1 → Context sensitive Lang
3. Type 2 → Context free Lang
4. Type 3 → Regular Lang



The types of languages form a strict hierarchy:



regular languages \subset context-free languages \subset context-sensitive languages \subset recursive languages \subset recursively enumerable languages.



Grammar

Language

Machine

UG

CSG

CFG

RG

RE Sets

CSL

CFL

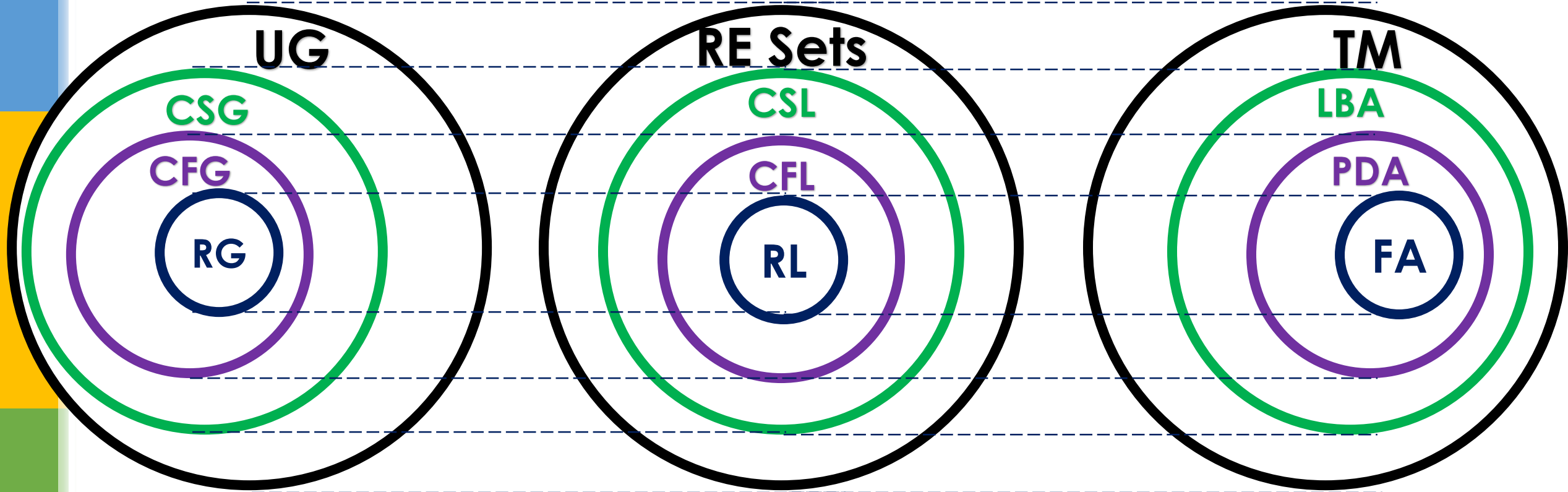
RL

TM

LBA

PDA

FA



Object-Oriented Style

- Programming with *Abstract Data Types*
 - ADTs specify/describe behaviors.
- Basic Program Unit: *Class*
 - Implementation of an ADT.
 - Abstraction enforced by encapsulation.
- Basic Run-time Unit: *Object*
 - Instance of a class.
 - Has an associated *state*.

Procedural vs Object-Oriented

- Emphasis on procedural abstraction.
- Top-down design;
Step-wise refinement.
- Suited for programming in the small.

- Emphasis on data abstraction.
- Bottom-up design;
Reusable libraries.
- Suited for programming in the large.

Integrating Heterogeneous Data

- In C, Pascal, etc., use

Union Type / Switch Statement

Variant Record Type / Case Statement

- In C++, Java, Eiffel, etc., use

Abstract Classes / Virtual Functions

Interfaces and Classes / Dynamic Binding

Comparison : *Figures* example

■ Data

■ Square

- side

■ Circle

- radius

■ Operation (area)

■ Square

- $\text{side} * \text{side}$

■ Circle

- $\text{PI} * \text{radius} * \text{radius}$

■ Classes

■ Square

- side

- area

(= $\text{side} * \text{side}$)

■ Circle

- radius

- area

(= $\text{PI} * \text{radius} * \text{radius}$)

Logic Programming Paradigm

- Integrates Data and Control Structures

`edge(a,b).`

`edge(a,c).`

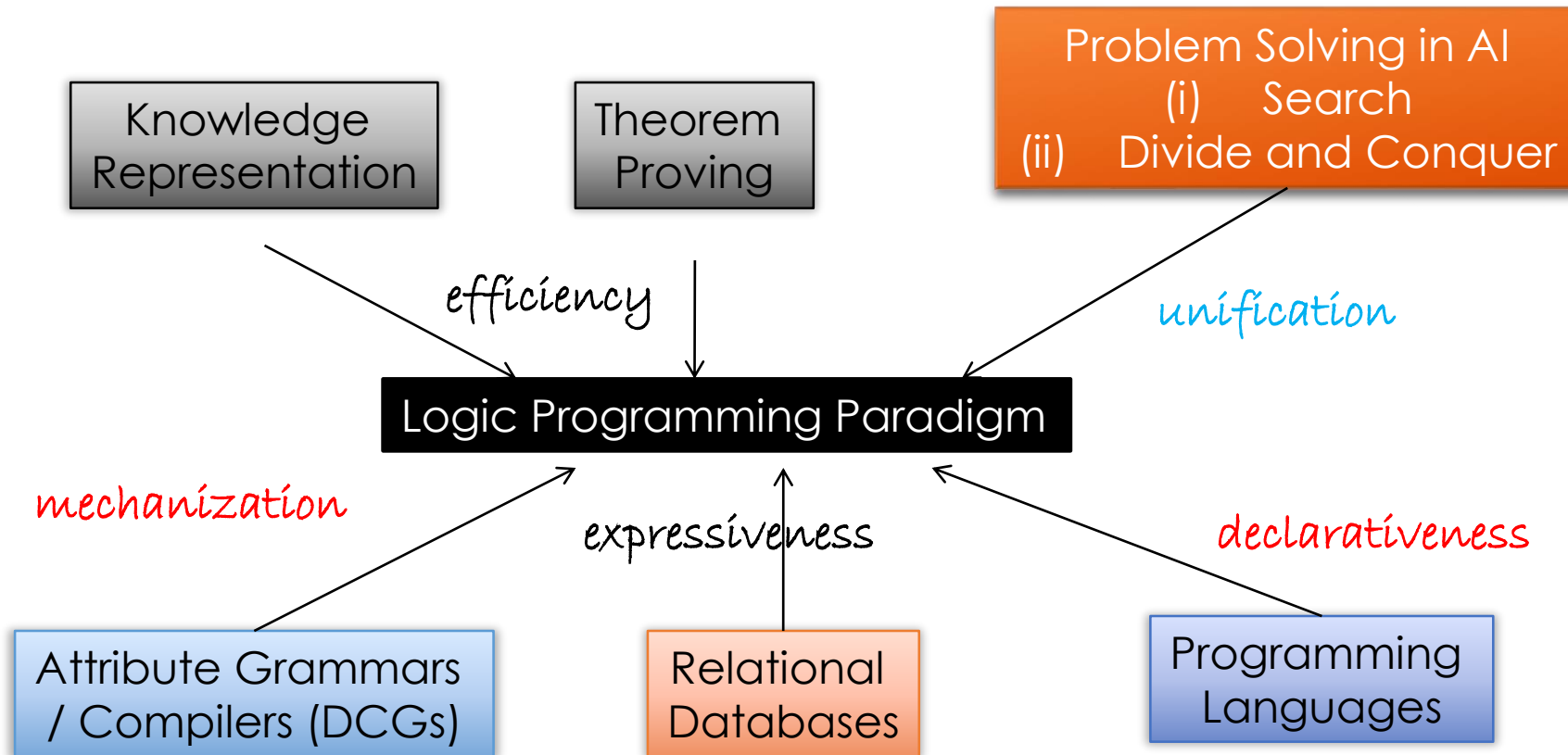
`edge(c,a).`

`path(X,X).`

`path(X,Y) :- edge(X,Y).`

`path(X,Y) :- edge(X,Z), path(Z,Y).`

Trading expressiveness for efficiency : Executable specification



Language Implementation: Bridging the Gap

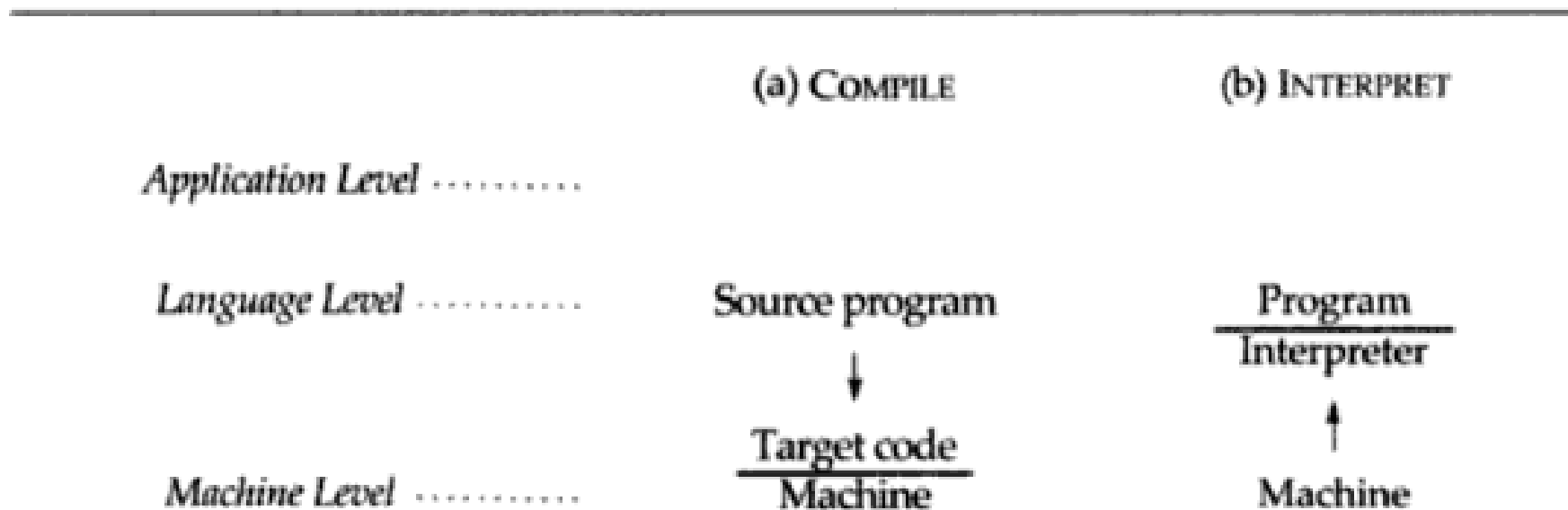
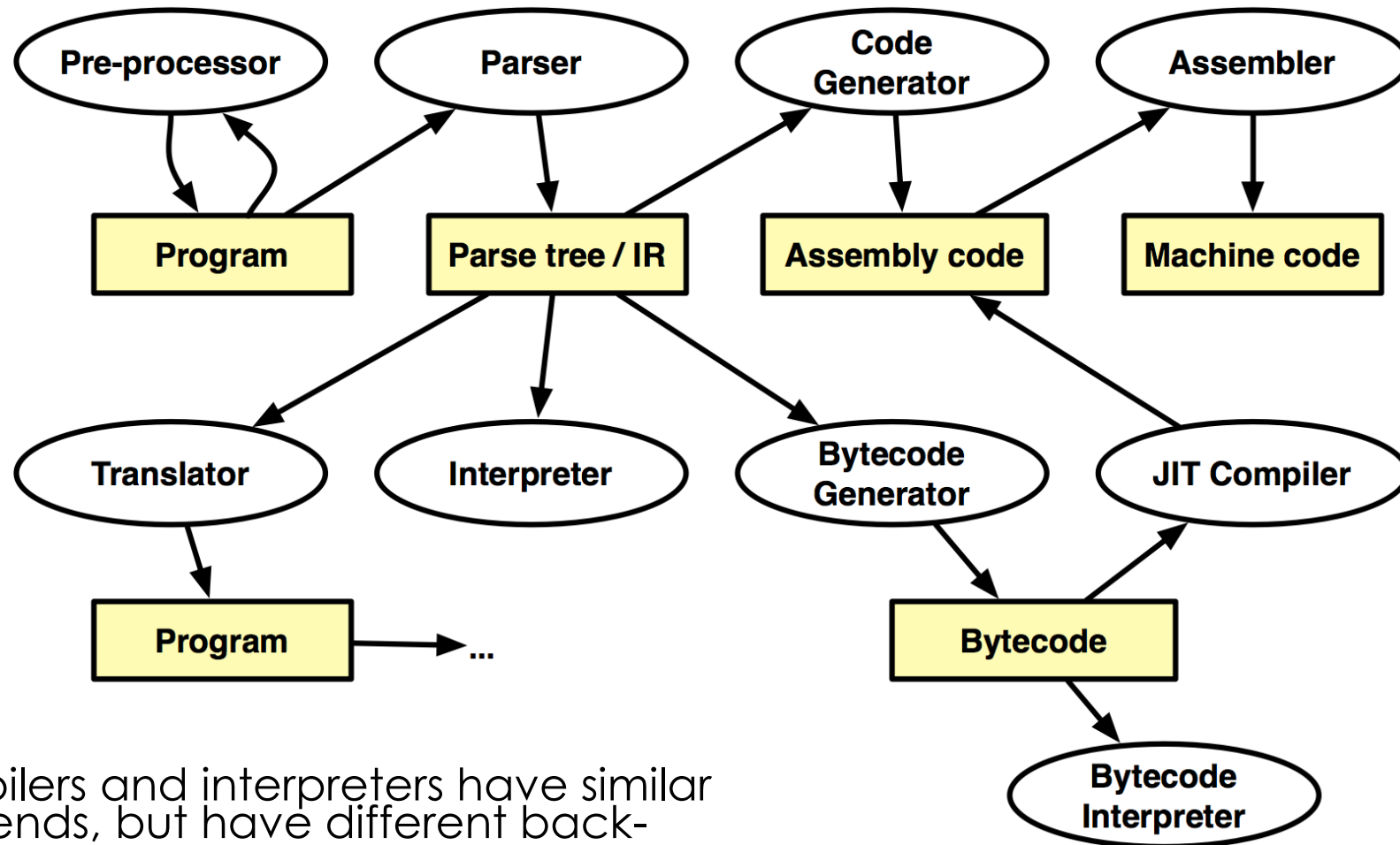


Figure 1.7 A language can be implemented by (a) translating it down to the level of the machine, or (b) by building a higher-level machine called an interpreter that can run the language directly.

Compilers and Interpreters



Compilers and interpreters have similar front-ends, but have different back-ends.

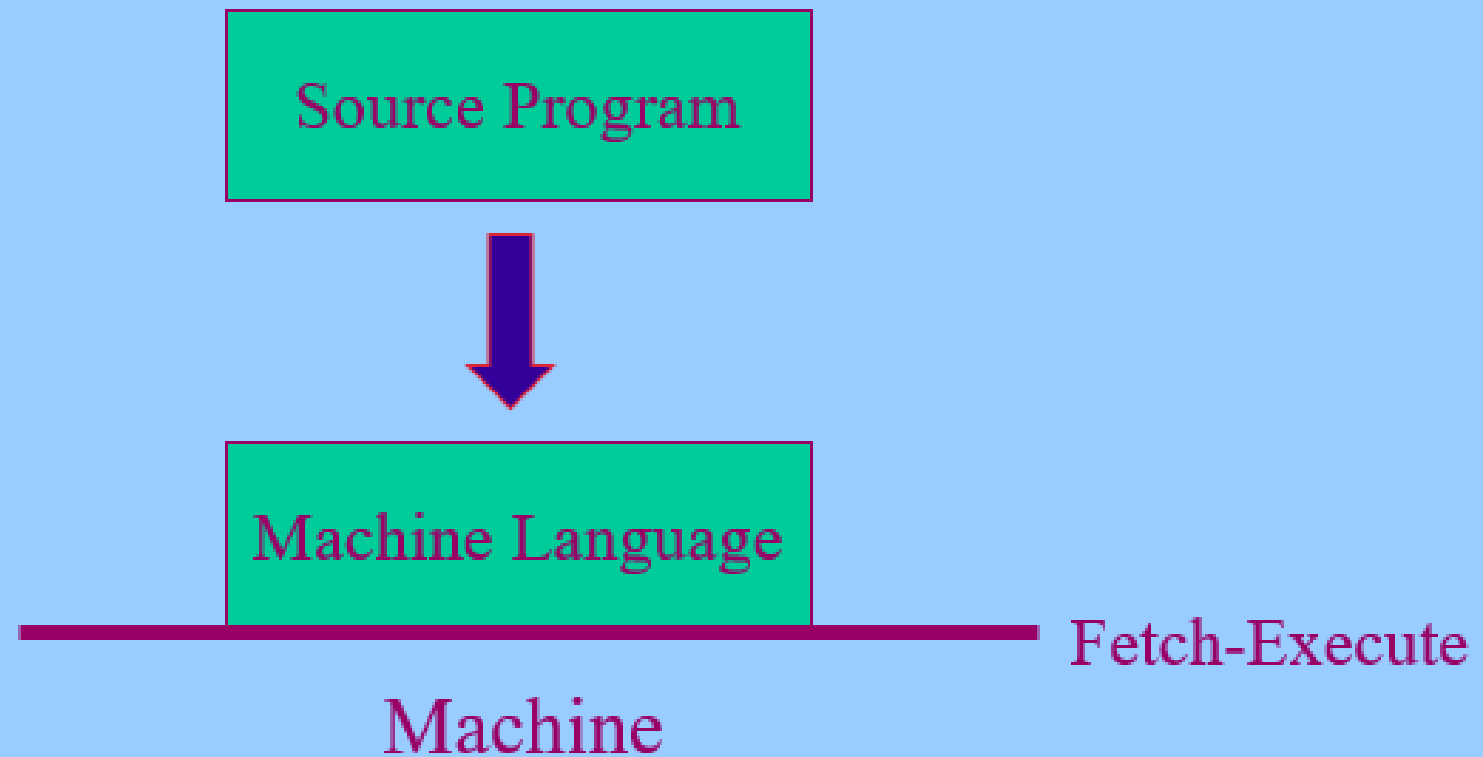
Language Implementation

- **Compiler** - source code is translated into machine code (all at once)
- **Interpreter** - machine is brought up to the language (one statement at a time)

Compiler

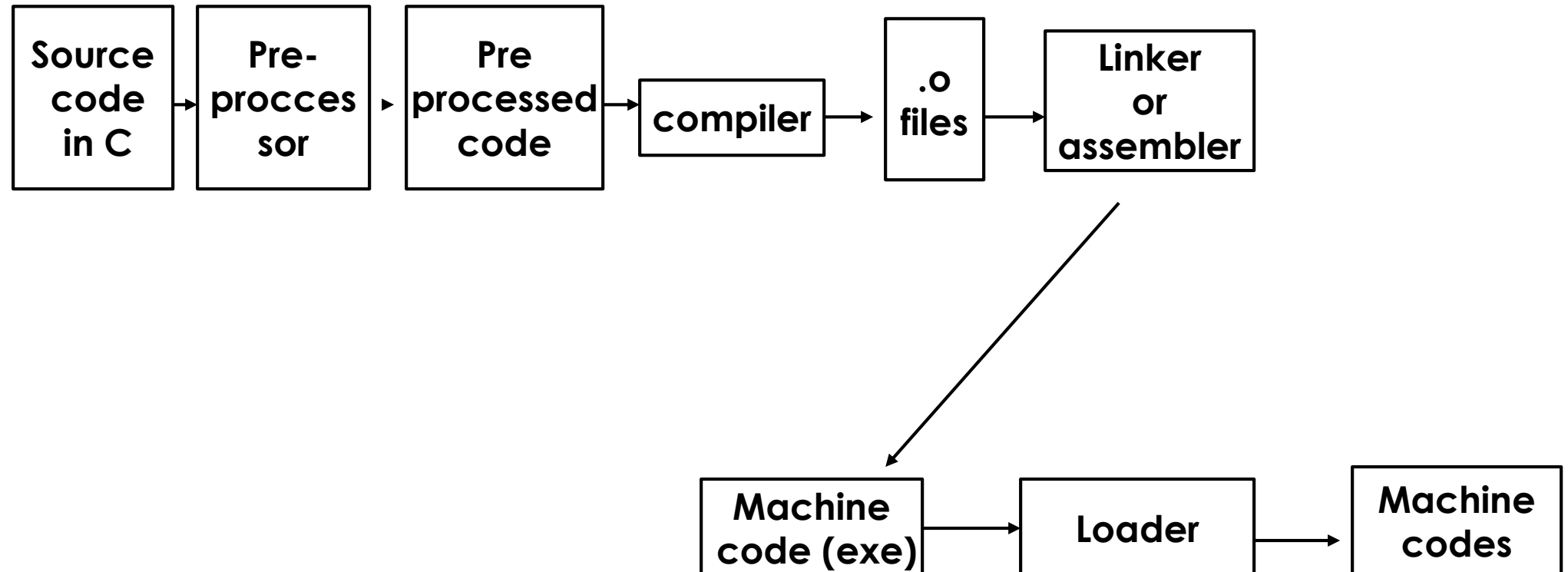
- Compiler translates source program into an equivalent program in the target language.
- Once the translation is complete, the target code is run at a later time, called run time.
- Compilation can be more efficient than interpreter

Implementation Methods 1. Compilation

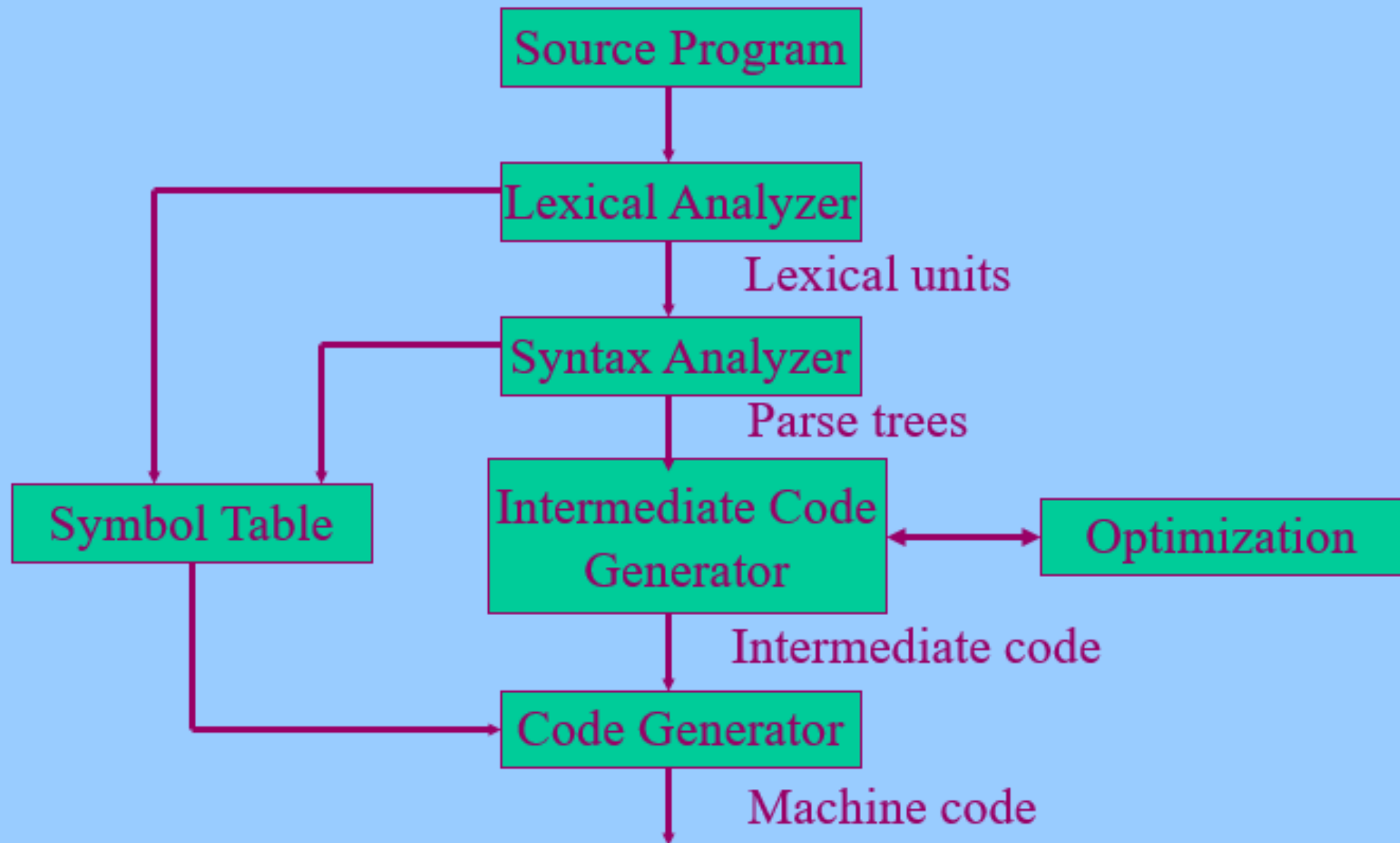


- Source programs are translated to machine language

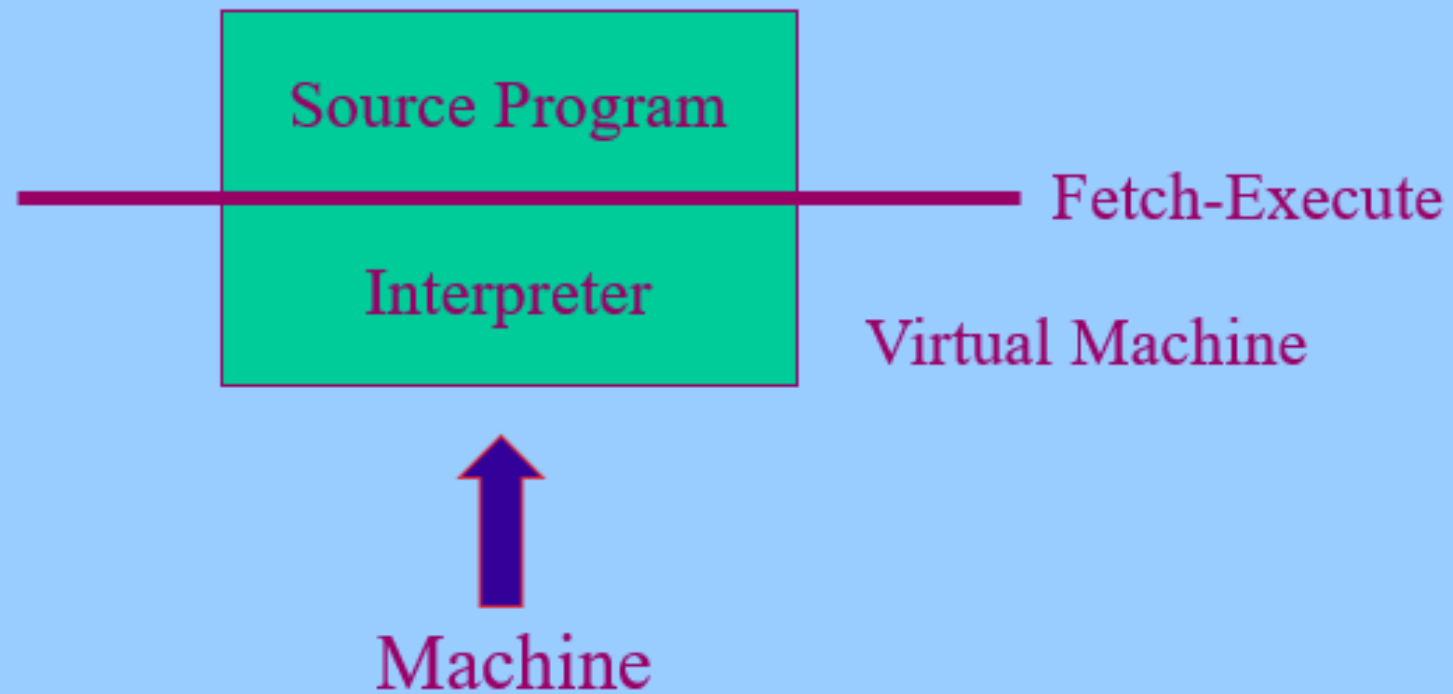
Compiled C



Implementation Methods 1. Compilation



Implementation Methods 2. Pure Interpretation



- Source programs are directly interpreted by another program without any translation.

Interpreted Code

- It takes a program and its input at the same time
- It scans the program, implementing operations as it encounters them, and doing input/output as needed.
- Interpreter can be more flexible than compilation.
- Each instruction is interpreted by machine interpreter
- does not produce object code

Comparisons

- Compilation more efficient
- Interpreted more flexible

Part 2

