DATE : / /
PAGE :

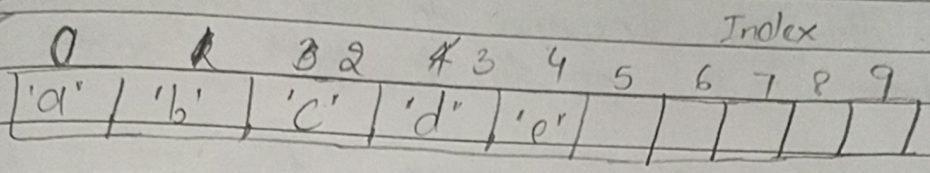
Lecture 5 - Array Insetion & deletion

Insetion

In last

Based on Index

In last

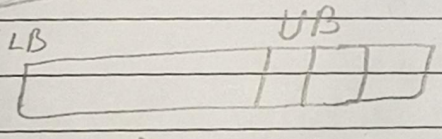


LB = 0
UB = 4
n = 5

```
Insert (arr, item) {
  if (UB = maxlimit) {
    overflow & return
  }
  A[UB+] = item
  UB++
  n++
}
```

Time complexity = $O(1)$ ⇒ Constant

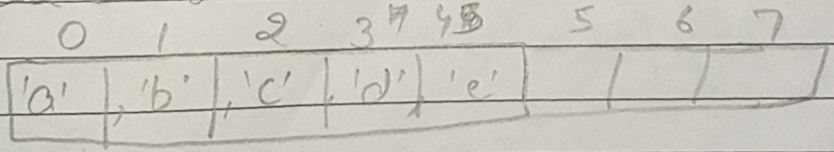
Based on Index



i = index at where new element to be inseted

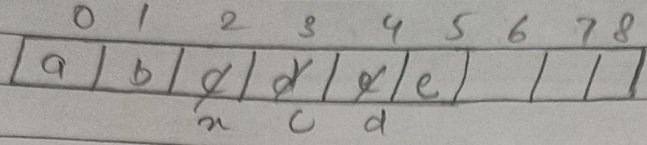
insetion(A[], item, i) { $LB \le i \le UB+1$

Case I $i = 5$



DATE : __/__/__
PAGE :

Case 2 : $i = 2$, item = 'x'



In this case we need to shift every element from $i = 2$ to the right side.

only & only if $i = UB+1$, then no shifting will be required

Algo:

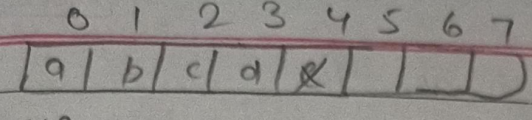
```
insert(A[], item, i) {
  if (UB == maxlimit)
      overflow & return;
  for(K = UB ; K >= i ; k--) {
      A[K+1] = A[K]
  }
  A[i] = item
  UB++;
  n++;
}
```

- `if (UB == maxlimit) overflow & return;` — Constant
- `for(K = UB ; K >= i ; k--) { A[K+1] = A[K] }` — dependent on the shifting, maximum Shifting = n
- `A[i] = item; UB++; n++;` → Constant

Time Complexity : $O(n)$ - Big Oh of n

Deletion

- Based of index
- from last

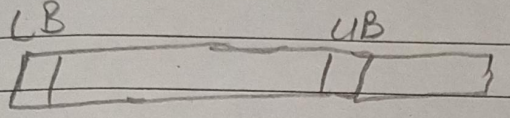


UB--

n--

```
if(n==0){
    underflow
}
```

* Based on Index



i = index from, where the element to be deleted

$LB \leq i \leq UB$

Run time Complexity

```
delete(A[], i){
  if(n==0){
      underflow & return
  }
  for(k=i; k<UB; k++){
      A[k] = A[k+1]
  }
  UB--;
  n--;
}
```

Time Complexity = O(n) - Big oh of n

DATE: / /
PAGE:

Quiz Number of shifting of the element to insert a new element at starting index in an array of 100 elements is?

Ans $n = 100$

(2) No. of shifting required to delete an element from start

A $N-1$

(3) Consider an array A of n elements ($n \geq 1$). In this array A, n insetion and n deletion are performed in arbitrary order. What should be the best and worst case complexity of all operation on array?

Ans
n insetion
n deletion
2n operation — $O(n)$ (best case)

worst case
each operation $O(n)$
for 2n operations $O(n^2)$

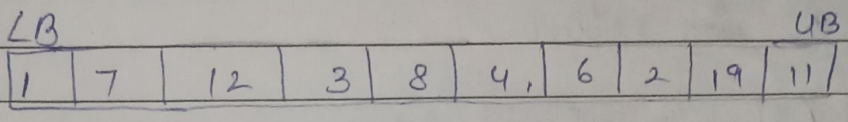
DATE : __/__/__
PAGE :

Lecture 6 - Searching

- Linear Search
- Binary Search

Prerequisite for BS : array should be sorted

* Linear Search

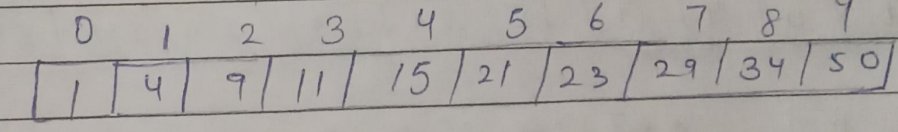



```
for(i=LB ; i<=UB; i++) {
    if(A[i] == item) {
        return i;
    }
return LB-1
```

Time Complexity = O(n) - Big oh of n

* Binary Search

* array should be sorted to perform this



Search = 15

$$mid = \frac{LB + UB}{2} = 4$$

$$\frac{0+9}{2} = 4$$

```
if (A[mid] = item) {
    return mid;
}
```

- Search 23

  mid = 4

  if (A[mid] = 23) { (A[mid])

  return

  if (item > A[mid]) {
      LB = mid
  } else if (item < A[mid]) {
      UB = mid - 1
  }

  when (LB > UB)
      Stop the loop

- Case 3: Search 10

	0	1	2	3	4	5	6	7	8	9
	L									H
	1	4	9	11	15	21	23	29	24	50

mid = $\frac{0+9}{2}$ = 4 ✗ 3

L = 0 ✗ 3
H = 9 ✗ 2
⇒ L > H ⇒ unsuccessful search

DATE: / /
PAGE:

Time complexity = $O(\log_2 n)$

Quiz Consider a sorted array of size n with duplicate elements. you have been given an element k, what is the time complexity to find the element k is appeared atleast n/2 times in the array or not?

a) O(1)
b) O(log n) ✓ Ans
c) $O(n^2)$
d) None

n = 8
k = 2

0	1	2	3	4	5	6	7
1,	2,	2,	2,	2,	2,	3,	9

log n
- apply modified binary search & find index i at which k appears first time in array from start

O(1)
- if (A[i+n/2] == k) {
  return true;
  }

DATE : / /
PAGE :

Q Consider a sorted array of size n with duplicate elements. what is the time complexity to find that any element is appeared at least n/2 times in the array or not?

a) O(1)
b) O(log n) ✓
c) O(n)
d) None

mid
↓
2, 5, 5, 5, 5, 5, 6, 10

what if

1, 2, 3, 4, 5, 5, 5, 5