

Lecture 9 Linked List

◦ Linear Data structure

◦ Linear order maintained using pointer (link)

* Advantages of Array

◦ elements can be accessed randomly by using the index

◦ using Array, other data structure like linked list, stacks, queues, tree graphs etc can be implemented

◦ Two 2d array are used to represent matrices

* Disadvantages of Array

- The number of elements to be stored in an array should be known in advance
- Allocating more memory than the requirement leads to wastage of memory space and less allocation of memory also leads to a problem
- An array is a static structure. Once declared the size of the array cannot be modified
- Insertion and deletion are quite difficult

→ List Contains node

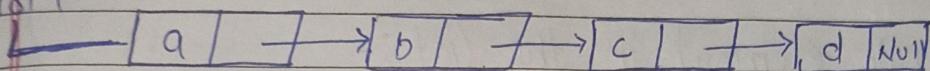
Node contain 2 fields

- Data (key or value)
- Link (next or form)

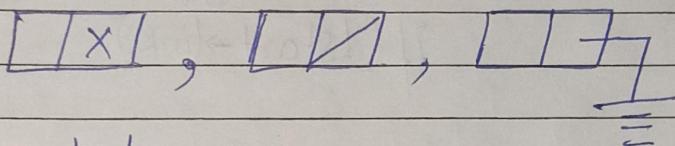
Node
[data | link]

- There is list pointer which points to the first node of list

list pointer



- Representation of NULL



* Accessing node

list → data

list → link

`printf("%c", list->data);` = a
`printf("%c", list->link->data);` = b

* Node structure

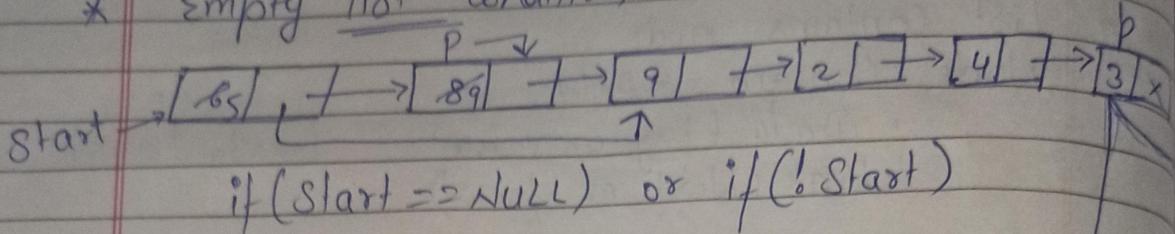
struct node {

char data;

struct node *link;

};

* Empty list Condition



- o not empty list

$\text{if}(\text{start} \neq \text{NULL}) \text{ or } \text{if}(\text{start})$

- o Single element in the list :-

~~Not possible~~ $\text{if}(\text{start} \rightarrow \text{link} == \text{NULL})$

or
 $\text{if}(!\text{start} \rightarrow \text{link})$

* NULL pointer Dereferencing

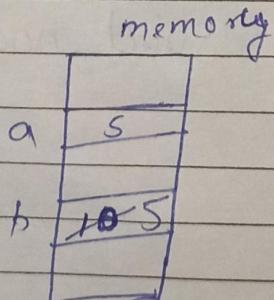
• Same list

`printf("%d", p → link → data);` ⇒ null pointer dereferencing

* Link Assignment

`int a = 5, b = 10;`

$b = a$
 Location → value



$\text{Start} \rightarrow \text{data} = 5$

Same list

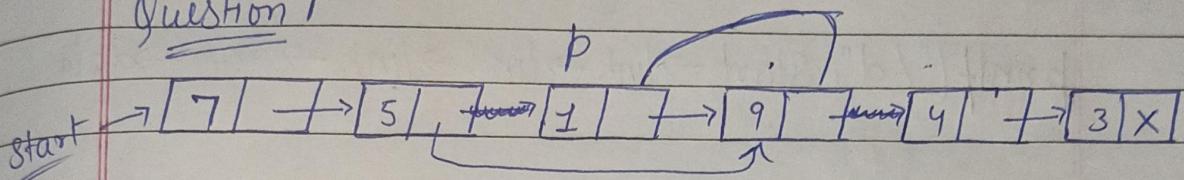
$p = \text{Start} \rightarrow \text{link}$

$p \rightarrow \text{data} = p \xrightarrow{q} \text{link} \rightarrow \text{data}$

$p \xrightarrow{q} l$

$\text{Start} \rightarrow \text{link} = p \rightarrow \text{link}$

Question 1



struct node *p;

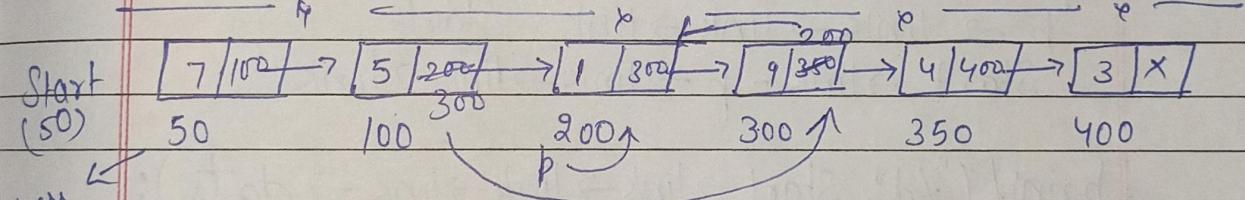
$p = \text{start} \rightarrow \text{link} \rightarrow \text{link};$

$(\text{start} \rightarrow \text{link}) \rightarrow \text{link} = p \rightarrow \text{link};$

$(p \rightarrow \text{link}) \rightarrow \text{link} = p;$

$\text{printf}(" \%d", (\text{start} \rightarrow \text{link} \rightarrow \text{link} \rightarrow \text{link} \rightarrow \text{link}) \rightarrow \text{data});$

Output = 9



Struct node *p;

$p = \text{start} \rightarrow \text{link} \rightarrow \text{link}; \quad p = 200$

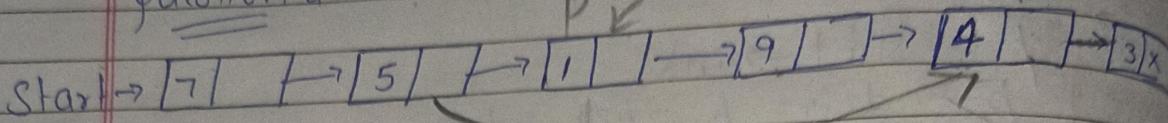
$\text{start} \rightarrow \text{link} \rightarrow \text{link} = p \xrightarrow{q} \text{link}; \quad q = 300$

$p \rightarrow \text{link} \rightarrow \text{link} = p \xrightarrow{q} 200$

$\text{printf}(" \%d", \text{start} \rightarrow \text{link} \rightarrow \text{link} \rightarrow \text{link} \rightarrow \text{link} \rightarrow \text{data});$

Output = 9

Question 2



struct node *p;

p = start → link → link;

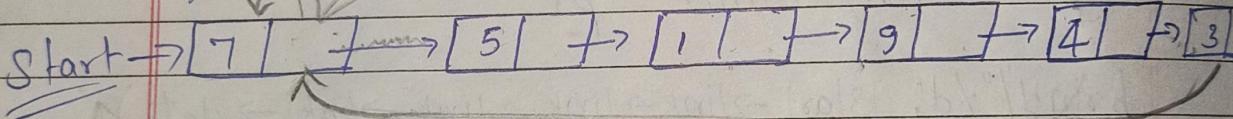
(start → link) → link = p → link → link

(p → link → link) → link = p

printf("%d", start → link → link → link → link → data)

①

Ques 3



struct node *p;

p = start → link → link → link

p → link → link → link = start

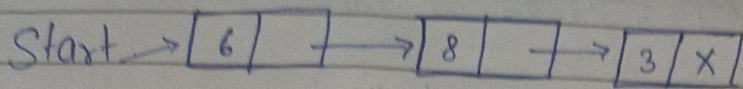
start → link = p → link → link → link

printf("%d", start → link → link → link → link → data);

③

⑦

Lecture 10 - Linked List - Traversing and Insertion



struct node *p = start | Run time complexity
 while ($p \neq \text{NULL}$) {
 process $p \rightarrow \text{data}$
 $p = p \rightarrow \text{link}$
 } } $= \Theta(n)$

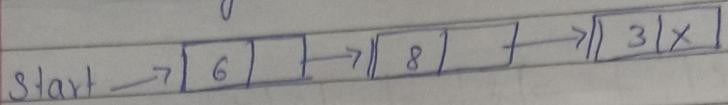
* Finding the address of last node

$p = \text{start}$
 $\text{if } (p == \text{NULL}) \text{ return;}$
 $\text{while } (p \rightarrow \text{link} \neq \text{NULL}) \text{ or while } (p \rightarrow \text{link})$
 {
 $p = p \rightarrow \text{link}$ } - Run time complexity = $\Theta(N)$
 }
 return p;

* Counting nodes in list

int count = 0;
 $p = \text{start}$
 $\text{while } (p \neq \text{NULL}) \{$ } Time Complexity = $\Theta(N)$
 $\text{Count}++;$
 $p = p \rightarrow \text{link};$ }
 }
 return count;

Searching in Linked list



- only linear search

because linked list is not stored sequence wise so we can't get the mid element to perform "Binary Search"

p = start

while (P != NULL) {

 if (p->data == item)

 return p

 } p = p->Link

} return NULL;

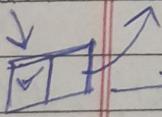
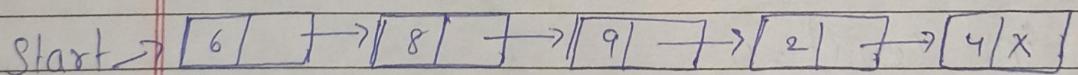
} Time complexity =

O(n)

* Insertion in linked list

- creating a new node

struct node* = (struct node*) malloc (sizeof(struct node))



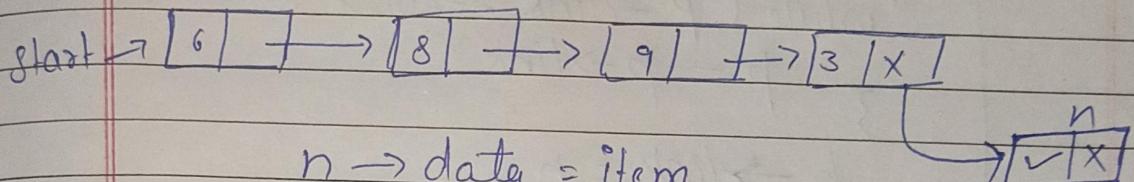
→ update links of node

n → — a — " existing node
— " list pointer

insertion at beginning (start , item) {

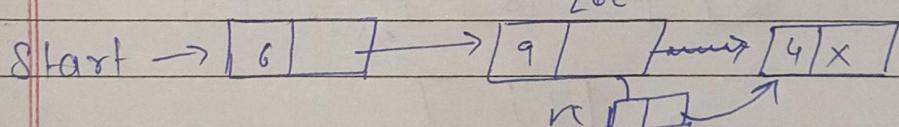
$n \rightarrow \text{data} = \text{item}$
 $n \rightarrow \text{link} = \text{Start}$
 $\text{Start} = n$
} Time Complexity = O(1)

* Insertion at the end

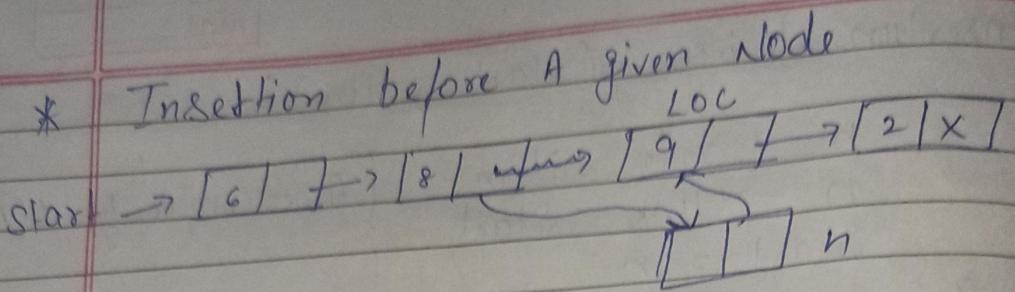


 $n \rightarrow \text{data} = \text{item}$
 $n \rightarrow \text{link} = \text{NULL}$
 $\text{if } (\text{start} == \text{NULL}) \text{ start} = n;$
 $\text{p} = \text{start}$ return
 $\text{while } (\text{p} \neq \text{NULL}) \{$
 $\quad \text{p} = \text{p} \rightarrow \text{link};$
 $\}$
 $\text{p} \rightarrow \text{link} = n$
} Time complexity = O(n)

* Insertion After given node



Time complexity $O(n)$ $n \rightarrow \text{data} = \text{item}$
 $O(1)$ $n \rightarrow \text{link} = \text{Loc} \rightarrow \text{link}$
 $\text{Loc} \rightarrow \text{link} = n$

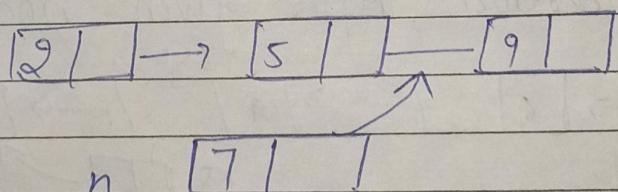


Time Complexity } $n \rightarrow \text{data} = \text{item}$
 $n \rightarrow \text{link} = \text{loc}$
 $\text{if } (\text{loc} == \text{start}) \text{ start} = n; \text{return}$

$O(n)$ $p = \text{start}$
 $\text{while } (p \rightarrow \text{link} != \text{loc}) \{$
 $\quad \quad \quad \quad p = p \rightarrow \text{link}$
 $\}$ $p \rightarrow \text{link} = n$

Quiz ① What is the complexity to insert a new node in a stored linked list

- A $O(1)$
X B $O(\log n)$
C $O(n)$ ✓
P D $O(n \log n)$



We need traverse the linked list and compare the n with existing node and according the we need to insert the element, Hence it will take $O(N)$ time

② what is the complexity to find k^{th} node from starting in a singly linked list ?

A

$O(n)$ ✓