# Assignment 1

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## Question 1

$$A_N = 1 + \frac{2}{N} \sum_{j=1}^{N} A_{j-1} for N > 0$$
 (1)

Multiplying both sides by N gives

$$NA_N = N + 2\sum_{j=1}^{N} A_{j-1} for N > 0$$
 (2)

This holds for N-1 as

$$(N-1)A_{N-1} = (N-1) + 2\sum_{j=1}^{N-1} A_{j-1} for N > 1$$
(3)

Subtracting the last two

$$NA_N = 1 + (N+1)A_{N-1} \tag{4}$$

Dividing both sides by N(N+1) gives

$$\frac{A_N}{N+1} = \frac{1}{N(N+1)} + \frac{A_{N-1}}{N} \tag{5}$$

Iterating

$$\frac{A_N}{N+1} = \sum_{k=2}^{N} \frac{1}{k(k+1)} + \frac{A_1}{2} \tag{6}$$

Assuming  $A_0 = 0$  as mentioned in the errata and solving the equation gives

$$A_N = N$$

### Question 2

Let's assume our pivot is the  $x+1^{th}$  smallest element in the array of N numbers after it is partitioned. Hence, there will be x elements to the left of the element once it gets partitioned.

Hence, we need to calculate the following -

$$\frac{1}{N}\sum_{x=0}^{N-1}E(\text{number of elements greater than the pivot in the first x elements})$$

This is so because only the number of elements greater than the pivot in the first x elements will need to be exchanged.

The quantity in the summation can be calculated using linearity of expectation. Say, there are variables  $X_1 \dots X_x$  where  $X_i = 1i$  if  $i^{th}$  element is greater than pivot else 0

Now, by linearity of expectation

E(Number of elements greater than pivot in the first x elements)

$$= E(X_1 + X_2 + \ldots + X_x)$$
(8)

But, notice that  $E(X_i) = Pr(X_i)$  and

$$P(X_i = 1) = 1 - \frac{x}{N - 1} \tag{9}$$

Hence, using the above observations and plugging it into (7), we get the required result.

Number of exchanges before partitioning  $=\frac{(N-2)}{6}$ 

### Question 3

We solve this with exactly the same method for solving the usual Quicksort recurrence. Notice that on the iteration step we can't go all the way to  $C_1$  but have to stop at  $C_M$  and this gives us the following recurrence.

$$\frac{C_N}{N+1} = 2\sum_{k=M+2} N + 1\frac{1}{k} + \frac{C_M}{M+1}$$
 (10)

Replacing the value of  $C_M$  gives us  $C_N = 2\sum_{k=M+2}^{N+1} \frac{1}{k} + \frac{M(M-1)}{4(M+1)}$ 

#### Question 4

Using the approximation to  $H_N$  gives us the following value for f(M)

$$f(M) = \frac{M(M-1)}{4(M+1)} - 2ln(M+2)$$
(11)

Throwing away  $\frac{M-1}{M+1}$  gives us the minimum to occur at M=6 Pretty cool huh!