ICPC Cookbook perocoders

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```
* Modular Exponentiation (a^k)%mod
11 mpow(11 a, 11 k){
                                                   * nCr w/ mod Time- O(n*k) Space- O(k)
    11 \text{ ans} = 1;
                                                   (only particular)
                                                   ll nck(ll n, ll k)
    while(k){
       if(k&1)ans = ans*a%mod;
       a = a*a\%mod;
                                                      11 C[k+1];
       k >>= 1:
                                                      memset(C, 0, sizeof(C));
    }
                                                       C[0] = 1 \% mod;
                                                      for (ll i = 1; i <= n; i++)
    return ans;
                                                      {
}
                                                          for (11 j = min(i, k); j > 0; j--)
* nCr w/o mod - O(n)(only particular)
                                                               C[j] = (C[j] + C[j-1]) \% mod;
11 ncr(11 n, 11 r) {
                                                      }
    11 \text{ res} = 1;
                                                      return C[k];
    r = min(n-r,r);
                                                   }
    for (ll i = 0; i < r; i++) {
                                                   * nCr w/ mod Time- O(n*k) Space- O(n*k)
        res = res * (n-i);
                                                   (all intermediates are calculated)
        res = res / (i+1);
                                                   11 nck(11 n, 11 k)
    }
                                                   {
    return res;
                                                      ll C[n+1][k+1];
}
                                                      ll i, j;
                                                      for (i = 0; i <= n; i++)
* nCr w/ mod Time- O(n), Space- O(n)
                                                      {
(only particular)
                                                           for (j = 0; j \le min(i, k); j++)
// make sure "mod" is prime
                                                           {
ll nCr(ll n, ll r)
                                                               if (j == 0 || j == i)
{
                                                                   C[i][j] = 1 \% mod;
   vector<ll> f(n + 1,1);
                                                               else
   for (11 i=2; i<=n;i++)
                                                                   C[i][j] = (C[i-1][j-1] +
       f[i] = (f[i-1]*i) \% mod;
                                                   C[i-1][j]) % mod;
    11 a = mpow(f[r], mod-2);
                                                           }
    11 b = mpow(f[n-r], mod-2);
                                                      }
    11 c = f[n];
                                                      return C[n][k];
   return (c*((a * b) % mod))) % mod;
                                                   * To check power of 2
* Modular Invers
                                                   bool isPower2(int v){
// Useful for large n in nCr w/ mod
                                                       return (v & (v - 1)) == 0;
11 modular_inverse(11 n, 11 mod){
                                                   }
   return mpow(n, mod-2);
```

}

```
* Sieve for Primes
                                                          if (prime[p] == true) {
void SieveOfEratosthenes(int n)
                                                               for (int i=p*p; i<=n; i += p)</pre>
{
                                                                   prime[i] = false;
    bool prime[n+1];
                                                          }
    memset(prime, true, sizeof(prime));
                                                       }
    for (int p=2; p*p<=n; p++) {</pre>
* Sieve in O(n)
                                                                SPF[i] = i;
const long long MAX_SIZE = 1000001;
                                                           }
vector<long long >isprime(MAX SIZE ,
                                                           for (long long int j=0;
true);
                                                                j < (int)prime.size() &&</pre>
vector<long long >prime;
                                                                i*prime[j] < N && prime[j] <=
vector<long long >SPF(MAX_SIZE);
                                                   SPF[i];j++) {
                                                                isprime[i*prime[j]]=false;
void manipulated seive(int N) {
    isprime[0] = isprime[1] = false ;
                                                                SPF[i*prime[j]] = prime[j] ;
    for (long long int i=2; i<N; i++) {</pre>
                                                           }
        if (isprime[i]) {
                                                       }
            prime.push_back(i);
                                                   }
* Dijkstra's Algorithm SSSP O(V + E*logV)
                                                           vis[x] = true;
// Initialize dist[] with INT MAX
                                                            auto itr = P[x].begin();
// P[x] is list of pair of vertices
                                                           while(itr!=P[x].end()){
Adjacent to x with weight. {e,w}
                                                                11 e = (*itr).first;
void Dijkstra(ll src,ll dist[],bool
                                                                11 w = (*itr).second;
vis[]){
                                                                if(dist[x] + w < dist[e]){</pre>
    dist[src] = 0;
                                                                   dist[e] = dist[x] + w;
    multiset<pair<ll,ll> > s;
    s.insert(make_pair(0,src));
                                                   s.insert(make_pair(dist[e],e));
    while(!s.empty()){
                                                                }
                                                                itr++;
        pair<11,11> p = *s.begin();
        s.erase(s.begin());
                                                           }
        11 x = p.second; 11 wei = p.first;
                                                       }
        if(vis[x]) continue;
                                                   }
```

```
* Bellman Ford's Algorithm SSSP O(V*E)
// Initialize dist[] with INT_MAX
void BellmanFord(int src,int dist[]){
    dist[src] = 0;
    for(int i=0;i<n;i++){</pre>
        for(int j=0;j<n;j++){</pre>
            auto itr = P[j].begin();
            while(itr!=P[j].end()){
                 int u = j;
                 int v = (*itr).first;
                 int weight = (*itr).second;
                 if (dist[u] != INF &&
                     dist[u] + weight < dist[v])</pre>
                     dist[v] = dist[u] + weight;
                 itr++;
            }
        }
    }
   for(int j=0;j<n;j++){</pre>
        auto itr = P[j].begin();
        while(itr!=P[j].end()){
            int u = j;
            int v = (*itr).first;
            int weight = (*itr).second;
            if (dist[u] != INF &&
                 dist[u] + weight < dist[v])</pre>
                 printf("negative weight cycle");
                 itr++;
        }
    }
}
* Floyd-Warshall's Algorithm APSP O(V^3)
// Don't forget to fill dist[][] before calling this procedure
void FloydWarshall(){
    for(int k = 0; k < n; k++)
        for(int i = 0; i < n; i++)</pre>
            for(int j = 0; j < n; j++)
                 dist[i][j] = min( dist[i][j], dist[i][k] + dist[k][j] );
}
```

```
* Kruskal's Algorithm MST O(E*logV)
                                                           if(root(x) != root(y)){
// Sort edges in preprocessing
                                                                minimumCost += cost;
int id[MAX], nodes, edges;
                                                                union1(x, y);
pair <long long, pair<int, int>> p[MAX];
                                                           }
void initialize(){
                                                       }
    for(int i = 0; i < MAX; ++i)
                                                       return minimumCost;
        id[i] = i;
                                                   }
}
int root(int x){
                                                   * Prim's Algorithm MST O(V + E*logV)
    while(id[x] != x){
                                                   int prim(int s)
        id[x] = id[id[x]]; x = id[x];
                                                   {
    }
                                                      multiset<pair<int,int>> pq;
    return x;
                                                      pq.insert(make pair(0,s));
}
                                                      int ans=0;
void union1(int x, int y){
                                                      while(!pq.empty())
    int p = root(x);
                                                      {
    int q = root(y);
                                                       pair <int,int> p = *pq.begin();
    id[p] = id[q];
                                                       pq.erase(pq.begin());
}
                                                       if(vis[p.second]==true)
bool find(int A,int B) {
                                                         continue;
    if( root(A)==root(B) )
                                                       vis[p.second]=true;
       return true;
                                                         auto itr = P[x].begin();
    else
                                                         while(itr!=P[x].end()) {
         return false;
                                                         int e = (*itr).first;
}
                                                         int w = (*itr).second;
long long kruskal(pair<long long,</pre>
                                                         pq.insert(make_pair(w,e));
pair<int, int> > p[]){
                                                         itr++;
    int x, y;
                                                         }
    long long cost, minimumCost = 0;
                                                       ans += p.first;
    for(int i = 0;i < edges;++i){</pre>
                                                      }
        x = p[i].second.first;
                                                      return ans;
        y = p[i].second.second;
                                                   }
        cost = p[i].first;
```

```
}
                                                    // Returns true if key presents in trie,
* Trie
                                                   else
const int ALPHABET SIZE = 26;
                                                   // false
struct TrieNode
                                                   bool search(struct TrieNode *root, string
{
   struct TrieNode
                                                   {
*children[ALPHABET SIZE];
                                                      struct TrieNode *pCrawl = root;
   // isEndOfWord is true if the node
                                                      for (int i = 0; i < key.length(); i++)</pre>
represents
                                                      {
   // end of a word
                                                          int index = key[i] - 'a';
   bool isEndOfWord;
                                                          if (!pCrawl->children[index])
};
                                                               return false;
// Returns new trie node (initialized to
                                                          pCrawl = pCrawl->children[index];
NULLs)
                                                      }
struct TrieNode *getNode(void)
                                                      return (pCrawl != NULL &&
{
                                                   pCrawl->isEndOfWord);
   struct TrieNode *pNode = new TrieNode;
   pNode->isEndOfWord = false;
                                                   int main()
   for (int i = 0; i < ALPHABET_SIZE; i++)</pre>
       pNode->children[i] = NULL;
                                                      // Input keys (use only 'a' through 'z'
   return pNode;
                                                      // and lower case)
                                                      string keys[] = {"the", "a", "there",
}
// If not present, inserts key into trie
                                                                       "answer", "any", "by",
                                                                        "bye", "their" };
// If the key is prefix of trie node, just
// marks leaf node
                                                      int n = sizeof(keys)/sizeof(keys[0]);
void insert(struct TrieNode *root, string
                                                       struct TrieNode *root = getNode();
                                                       // Construct trie
key)
{
                                                      for (int i = 0; i < n; i++)
   struct TrieNode *pCrawl = root;
                                                          insert(root, keys[i]);
   for (int i = 0; i < key.length(); i++)</pre>
                                                      // Search for different keys
                                                      search(root, "the")? cout << "Yes\n" :</pre>
   {
       int index = key[i] - 'a';
                                                                            cout << "No\n";</pre>
       if (!pCrawl->children[index])
                                                      search(root, "these")? cout << "Yes\n"</pre>
           pCrawl->children[index] =
getNode();
                                                                              cout << "No\n";</pre>
       pCrawl = pCrawl->children[index];
                                                      return 0;
   }
                                                   }
   // mark last node as leaf
```

pCrawl->isEndOfWord = true;

* Bipartite matching algorithms

```
#define M
#define N
bool bpm(bool bpGraph[M][N], int u,
        bool seen[], int matchR[])
{
   // Try every job one by one
   for (int v = 0; v < N; v++)
   {
       // If applicant u is interested in
       // job v and v is not visited
       if (bpGraph[u][v] && !seen[v])
       {
           // Mark v as visited
           seen[v] = true;
           if (matchR[v] < 0 || bpm(bpGraph, matchR[v],</pre>
                                     seen, matchR))
           {
               matchR[v] = u;
               return true;
           }
       }
   }
   return false;
}
int maxBPM(bool bpGraph[M][N])
{
   int matchR[N];
   memset(matchR, -1, sizeof(matchR));
   int result = 0;
   for (int u = 0; u < M; u++)</pre>
   {
       // Mark all jobs as not seen
       // for next applicant.
       bool seen[N];
       memset(seen, 0, sizeof(seen));
       // Find if the applicant 'u' can get a job
       if (bpm(bpGraph, u, seen, matchR))
           result++;
```

```
}
   return result;
}
int main()
{
   bool bpGraph[M][N]'
    cout << maxBPM(bpGraph);</pre>
   return 0;
}
* Segment Tree
int tree[2 * N];
void build( int arr[])
{
   for (int i=0; i<n; i++)</pre>
       tree[n+i] = arr[i];
   for (int i = n - 1; i > 0; --i)
       tree[i] = tree[i<<1] + tree[i<<1 | 1];</pre>
}
void updateTreeNode(int p, int value)
{
   tree[p+n] = value;
   p = p+n;
   for (int i=p; i > 1; i >>= 1)
       tree[i>>1] = tree[i] + tree[i^1];
int query(int 1, int r)
{
   int res = 0;
   for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1)
   {
       if (1&1)
           res += tree[1++];
       if (r&1)
           res += tree[--r];
   }
   return res;
}
```

```
* Longest Common Subsequence
#include<bits/stdc++.h>
                                                    else
using namespace std;
int main()
int lcs( char *X, char *Y, int m, int n )
                                                        int mid = (start+end)/2;
{
                                                        if(a[mid]==key)
  int L[m+1][n+1];
                                                        { return; }
  int i, j;
                                                        else
    for (i=0; i<=m; i++)
                                                          if(a[mid]<key)</pre>
  {
                                                                search(a,mid,end,key);
    for (j=0; j<=n; j++)
                                                          else
                                                                search(a,start,mid,key);
      if (i == 0 || j == 0)
                                                       }
        L[i][j] = 0;
                                                    }
                                                    int main()
      else if (X[i-1] == Y[j-1])
                                                    {
        L[i][j] = L[i-1][j-1] + 1;
                                                       ios_base::sync_with_stdio(false);
                                                       cin.tie(NULL);
      else
                                                       int n;
        L[i][j] = max(L[i-1][j],
                                                       cin>>n;
L[i][j-1]);
                                                       int arr[n];
    }
                                                       for(int i=0;i<n;i++)</pre>
  }
                                                       cin>>arr[i];
  return L[m][n];
                                                       int a[n], s=0;
}
                                                       a[0] = arr[0], s=1;
                                                       for(int i=1;i<n;i++)</pre>
* Longest increasing subsequence
                                                       {
#include<bits/stdc++.h>
                                                        if(arr[i] < a[0])</pre>
using namespace std;
                                                          a[0] = arr[i];
                                                        else if(arr[i] > a[s-1])
void search(int a[],int start,int end,int
                                                          a[s++] = arr[i];
key)
{
                                                        else
   if(start+1==end)
                                                        {
                                                             search(a,0,s-1,arr[i]); }
                                                       }
```

cout<<s;

return 0;

}

if(a[start]==key || a[end]==key)

return;

}

a[end] = key;

```
* Polygon Related Algorithms
const double PI = 4*atan(1);
struct point{
    double x; double y;
};
double dist(point P0, point P1){
    return sqrt((P0.x-P1.x)*(P0.x-P1.x) + (P0.y-P1.y)*(P0.y-P1.y));
}
double angle(point pointA, point pointB, point pointC){
    double a = pointB.x - pointA.x; double b = pointB.y - pointA.y;
    double c = pointB.x - pointC.x; double d = pointB.y - pointC.y;
    double atanA = atan2(a, b); double atanB = atan2(c, d);
    return atanB - atanA;
}
int ccw (point P0, point P1, point P2) {
    double dx1 = P1.x - P0.x; double dx2 = P2.x - P0.x;
   double dy1 = P1.y - P0.y; double dy2 = P1.y - P0.y;
    if (dy1 * dx2 > dy2 * dx1) return -1;
    if (dx1 * dy2 > dy1 * dx2) return 1;
    if ((dx1 * dx2 < 0) || (dy1 * dy2 < 0)) return 1;
    if ((dx1 * dx1 + dy1 * dy1) < (dx2 * dx2 + dy2 * dy2)) return -1;
    return 0;
}
double perimeter(const vector<point> &P){
    double result = 0.0;
    for(int i = 0; i < (int)P.size()-1; i++) // remember that P[0] = P[n-1]
        result += dist(P[i], P[i+1]);
    return result;
}
double area(const vector<point> &P){
    double result = 0.0, x1, y1, x2, y2;
    for(int i = 0; i < (int)P.size()-1; i++){</pre>
        x1 = P[i].x; x2 = P[i+1].x; y1 = P[i].y; y2 = P[i+1].y;
        result += (x1 * y2 - x2 * y1);
    return fabs(result) / 2.0;
}
bool isConvex(const vector<point> &P){ // returns true if all three
    int sz = (int)P.size(); // consecutive vertices of P form the same turns
    if (sz <= 3) return false;</pre>
```

```
bool isLeft = ccw(P[0], P[1], P[2]); // remember one result
    for (int i = 1; i < sz-1; i++) // then compare with the others
        if (ccw(P[i], P[i+1], P[(i+2) == sz ? 1 : i+2]) != isLeft)
            return false; // different sign -> this polygon is concave
    return true;
}
// returns true if point p is in either convex/concave polygon P
bool inPolygon(point pt, const vector<point> &P){
    if((int)P.size() == 0) return false;
    double sum = 0; // assume the first vertex is equal to the last vertex
    for (int i = 0; i < (int)P.size()-1; i++) {</pre>
        if (ccw(pt, P[i], P[i+1])) sum += angle(P[i], pt, P[i+1]); // left turn/ccw
        else sum -= angle(P[i], pt, P[i+1]); } // right turn/cw
    double EPS = 0.1;
    return fabs(fabs(sum) - 2*PI) < EPS;</pre>
}
* Fast Fourier Transform (FFT) O(n*logn)
void init_fft(long long n){
    FFT_N = n; omega.resize(n); double angle = 2 * PI / n;
    for(int i = 0; i < n; i++)
        omega[i] = base( cos(i * angle), sin(i * angle));
}
void fft (vector<base> & a){
    long long n = (long long) a.size();
    if (n == 1) return;
    long long half = n >> 1;
   vector<base> even (half), odd (half);
   for (int i=0, j=0; i<n; i+=2, ++j){
        even[j] = a[i]; odd[j] = a[i+1];
    }
  fft (even), fft (odd);
  for (int i=0, fact = FFT_N/n; i < half; ++i){</pre>
        base twiddle = odd[i] * omega[i * fact];
        a[i] = even[i] + twiddle; a[i+half] = even[i] - twiddle;
   }
}
void multiply (const vector<long long> & a, const vector<long long> & b, vector<long long>
& res){
    vector<base> fa (a.begin(), a.end()), fb (b.begin(), b.end());
    long long n = 1;
```

```
while (n < 2*max (a.size(), b.size())) n <<= 1;
    fa.resize (n), fb.resize (n);
    init_fft(n);
    fft (fa), fft (fb);
    for (size_t i=0; i<n; ++i) fa[i] = conj( fa[i] * fb[i]);</pre>
    fft (fa);
    res.resize (n);
    for(size_t i=0; i<n; ++i){</pre>
        res[i] = (long long) (fa[i].real() / n + 0.5); res[i]%=mod;
    }
}
* Lowest Common Ancestor (LCA) O(n)
struct Node
{
    int key;
    struct Node *left, *right;
};
// Utility function creates a new binary tree node with given key
Node * newNode(int k)
{
    Node *temp = new Node;
    temp->key = k;
    temp->left = temp->right = NULL;
    return temp;
}
// Finds the path from root node to given root of the tree, Stores the
// path in a vector path[], returns true if path exists otherwise false
bool findPath(Node *root, vector<int> &path, int k)
{
    // base case
    if (root == NULL) return false;
    // Store this node in path vector. The node will be removed if
    // not in path from root to k
    path.push_back(root->key);
    // See if the k is same as root's key
    if (root->key == k)
        return true;
```

```
// Check if k is found in left or right sub-tree
    if ( (root->left && findPath(root->left, path, k)) ||
        (root->right && findPath(root->right, path, k)) )
        return true;
    // If not present in subtree rooted with root, remove root from
    // path[] and return false
    path.pop_back();
    return false;
// Returns LCA if node n1, n2 are present in the given binary tree,
// otherwise return -1
int findLCA(Node *root, int n1, int n2)
{
    // to store paths to n1 and n2 from the root
    vector<int> path1, path2;
    // Find paths from root to n1 and root to n1. If either n1 or n2
    // is not present, return -1
    if ( !findPath(root, path1, n1) || !findPath(root, path2, n2))
        return -1;
    /* Compare the paths to get the first different value */
    int i;
    for (i = 0; i < path1.size() && i < path2.size(); i++)</pre>
        if (path1[i] != path2[i])
            break;
    return path1[i-1];
}
// Driver program to test above functions
int main()
{
    Node * root = newNode(1);
    root->left = newNode(2);
    root->right = newNode(3);
    cout << "LCA(4, 5) = " << findLCA(root, 4, 5);</pre>
    return 0;
}
```

```
* Extended Euclidean Algorithm. Calculates X,Y s.t aX+bY = gcd(a,b)
int gcdExtended(int a, int b, int *x, int *y){
   if (a == 0) // Base Case
   {
       *x = 0;
       *y = 1;
       return b;
   }
   int x1, y1; // To store results of recursive call
   int gcd = gcdExtended(b%a, a, &x1, &y1);
   // Update x and y using results of recursive
   // call
   *x = y1 - (b/a) * x1;
   *y = x1;
   return gcd;
* Modular inverse of a, modulo m
void modInverse(int a, int m)
{
   int x, y;
   int g = gcdExtended(a, m, &x, &y);
   if (g != 1)
       cout << "Inverse doesn't exist";</pre>
   else
   {
       // m is added to handle negative x
       int res = (x\%m + m) \% m;
       cout << "Modular multiplicative inverse is " << res;</pre>
       return res;
   }
}
* CHINESE REMAINDER THEOREM
// Given a1,a2 ....ak the remainders and n1,n2 ....nk
// Find min sol x which satisfies : X = a1 \pmod{n1} \dots X = ak \pmod{nk}
// N = n1*n2....nk, yi = N/ni , zi = inv(yi)mod(ni)
// X = SUMMATION for(i = 1:k) ai*yi*zi
// k is size of num[] and rem[]. Returns the smallest (a1, a2 .. ak in rem[] and n1,n2 ..
nk in num[])
// number x such that:
```

```
// x \% num[0] = rem[0],
// x \% num[1] = rem[1],
// ......
// x \% num[k-2] = rem[k-1]
// Assumption: Numbers in num[] are pairwise coprime
// (gcd for every pair is 1)
int findMinX(int num[], int rem[], int k)
{
  // Compute product of all numbers
   int prod = 1;
  for (int i = 0; i < k; i++)</pre>
       prod *= num[i];
  // Initialize result
  int result = 0;
  // Apply above formula
  for (int i = 0; i < k; i++)</pre>
  {
       int pp = prod / num[i];
       result += rem[i] * inv(pp, num[i]) * pp;
   }
   return result % prod;
}
* EULER TOTIENT SIEVE
#define SIZE 10000
int phi[SIZE];
for(int i = 0;i<SIZE;i++)</pre>
    phi[i] = i;
int EulerTotientFiller(){
  for(int i = 2;i<SIZE;i++){</pre>
    if(phi[i] == i){ // its a prime
     for(j = 2*i;j<SIZE;j+= i){</pre>
            phi[j] = phi[j] - (phi[j]/i); // i is prime factor of j
      }
    }
    else{
     // its not prime skip it.
    }
   }
}
```

```
* KMP Pattern Matching O(n+k)
                                                  }
#include <bits/stdc++.h>
                                                  void computeLPSArray(char* pat, int M,
void computeLPSArray(char* pat, int M,
                                                  int* lps)
int* lps);
                                                  { int len = 0;
void KMPSearch(char* pat, char* txt)
                                                     lps[0] = 0; // lps[0] is always 0
{
                                                     int i = 1;
    int M = strlen(pat);
                                                      while (i < M) {
    int N = strlen(txt);
                                                         if (pat[i] == pat[len]) {
    int lps[M];
                                                             len++;
    computeLPSArray(pat, M, lps);
                                                             lps[i] = len;
    int i = 0; // index for txt[]
                                                             i++;
    int j = 0; // index for pat[]
                                                         }
    while (i < N) {
                                                         else // (pat[i] != pat[len])
       if (pat[j] == txt[i]) {
                                                         {
                                                             if (len != 0) {
           j++;
           i++;
                                                                  len = lps[len - 1];
       }
                                                             }
       if (j == M) {
                                                             else // if (len == 0)
            printf("Found pattern at index
                                                             { lps[i] = 0; i++; }
%d ", i - j);
                                                  }
            j = lps[j - 1];
                                                  }
                                                  }
       else if (i < N && pat[j] != txt[i])</pre>
                                                  int main()
{
                                                  {
           if (j != 0)
                                                      char txt[] = "ABABDABACDABABCABAB";
                                                      char pat[] = "ABABCABAB";
               j = lps[j - 1];
                                                      KMPSearch(pat, txt);
           else
               i = i + 1;
                                                      return 0;
       }
                                                  }
    }
```

* Ford Fulkerson algorithm

```
#include <iostream>
#include <limits.h>
#include <string.h>
#include <queue>
using namespace std;
// Number of vertices in given graph
#define V 6
/* Returns true if there is a path from source 's' to sink 't' in
residual graph. Also fills parent[] to store the path */
bool bfs(int rGraph[V][V], int s, int t, int parent[])
{
    // Create a visited array and mark all vertices as not visited
    bool visited[V];
    memset(visited, 0, sizeof(visited));
    // Create a queue, enqueue source vertex and mark source vertex
    // as visited
    queue <int> q;
    q.push(s);
    visited[s] = true;
    parent[s] = -1;
    // Standard BFS Loop
    while (!q.empty())
    {
        int u = q.front();
        q.pop();
        for (int v=0; v<V; v++)</pre>
        {
            if (visited[v]==false && rGraph[u][v] > 0)
            {
                q.push(v);
                parent[v] = u;
                visited[v] = true;
            }
        }
    }
```

```
// If we reached sink in BFS starting from source, then return
    // true, else false
    return (visited[t] == true);
}
// Returns the maximum flow from s to t in the given graph
int fordFulkerson(int graph[V][V], int s, int t)
{
    int u, v;
    // Create a residual graph and fill the residual graph with
    // given capacities in the original graph as residual capacities
    // in residual graph
    int rGraph[V][V]; // Residual graph where rGraph[i][j] indicates
                    // residual capacity of edge from i to j (if there
                    // is an edge. If rGraph[i][j] is 0, then there is not)
   for (u = 0; u < V; u++)
       for (v = 0; v < V; v++)
            rGraph[u][v] = graph[u][v];
    int parent[V]; // This array is filled by BFS and to store path
    int max_flow = 0; // There is no flow initially
    // Augment the flow while tere is path from source to sink
    while (bfs(rGraph, s, t, parent))
    {
       // Find minimum residual capacity of the edges along the
       // path filled by BFS. Or we can say find the maximum flow
       // through the path found.
        int path_flow = INT_MAX;
       for (v=t; v!=s; v=parent[v])
        {
            u = parent[v];
            path_flow = min(path_flow, rGraph[u][v]);
        }
        // update residual capacities of the edges and reverse edges
       // along the path
       for (v=t; v != s; v=parent[v])
        {
            u = parent[v];
            rGraph[u][v] -= path_flow;
            rGraph[v][u] += path_flow;
        }
```

```
// Add path flow to overall flow
        max_flow += path_flow;
    }
    // Return the overall flow
    return max_flow;
// Driver program to test above functions
int main()
{
    int graph[V][V]
    cout << "The maximum possible flow is " << fordFulkerson(graph, 0, 5);</pre>
    return 0;
}
* Coordinate Compression 1D
int main() {
   int t, n, q, x1, x2, a, b, c, m;
   cin >> t;
   for (int c = 1; c <= t; c++) {
       cin >> n >> q;
       vector<int> x(n);
       scan(x, n);
       vector<int> y(n);
       scan(y, n);
       vector<int> z(q);
       scan(z, q);
       vector<int>1(n), r(n), k(q);
       for (int i = 0; i < n; i++) {
           l[i] = min(x[i], y[i]) + 1;
           r[i] = \max(x[i], y[i]) + 1;
           r[i]++;
       }
       for (int i = 0; i < q; i++) {
           k[i] = z[i] + 1;
       }
       set <int> s;
       for (int i = 0; i < n; i++) {
           s.insert(l[i]);
           s.insert(r[i]);
```

```
}
       vector<int> pos;
       set <int> :: iterator itr;
       for (itr = s.begin(); itr != s.end(); ++itr) {
           pos.push_back(*itr);
       }
       int p = pos.size();
       vector<int> m(p+1, 0);
       for (int i = 0; i < n; i++) {
           m[lower_bound(pos.begin(), pos.end(), l[i]) - pos.begin()]++;
           m[lower_bound(pos.begin(), pos.end(), r[i]) - pos.begin()]--;
       }
       for (int i = 0; i < p-1; i++) {
           m[i+1] = m[i+1] + m[i];
       }
       vector<ll> sf(p, 0);
       sf[p-1] = m[p-1];
       for (int i = p-1; i > 0; i--) {
           sf[i-1] = sf[i] + m[i-1] * 1LL * (pos[i]-pos[i-1]);
       }
       11 ans = 0;
       for (int i = 0; i < q; i++) {
           if (k[i] <= sf[0]) {</pre>
               int x1 = upper_bound(sf.begin(), sf.end(), k[i], greater<long long>()) -
sf.begin();
               if (sf[x1] == k[i]) {
                   ans += (1LL * pos[x1] * (i+1));
               }
               else {
                   X1--;
                   long long d = sf[x1] - k[i];
                   d = d / m[x1];
                   11 \text{ term} = pos[x1] + d;
                   ans += term * (i+1);
               }
           }
       }
       cout << "Case #" << c << ": " << ans << endl;</pre>
   }
}
```

* Bridge in Graph void Graph::bridgeUtil(int u, bool visited[], int disc[], int low[], int parent[]) { // A static variable is used for simplicity, we can // avoid use of static variable by passing a pointer. static int time = 0; // Mark the current node as visited visited[u] = true; // Initialize discovery time and low value disc[u] = low[u] = ++time; // Go through all vertices aadjacent to this list<int>::iterator i; for (i = adj[u].begin(); i != adj[u].end(); ++i) { int v = *i; // v is current adjacent of u // If v is not visited yet, then recur for it if (!visited[v]) { parent[v] = u; bridgeUtil(v, visited, disc, low, parent); // Check if the subtree rooted with v has a // connection to one of the ancestors of u low[u] = min(low[u], low[v]);// If the lowest vertex reachable from subtree // under v is below u in DFS tree, then u-v // is a bridge if (low[v] > disc[u]) cout << u <<" " << v << endl; } // Update low value of u for parent function calls. else if (v != parent[u]) low[u] = min(low[u], disc[v]); } }

* Articulation Point

```
void Graph::APUtil(int u, bool visited[], int disc[],
                                   int low[], int parent[], bool ap[])
{
   // A static variable is used for simplicity, we can avoid use of static
   // variable by passing a pointer.
   static int time = 0;
   // Count of children in DFS Tree
   int children = 0;
   // Mark the current node as visited
   visited[u] = true;
   // Initialize discovery time and low value
   disc[u] = low[u] = ++time;
   // Go through all vertices aadjacent to this
   list<int>::iterator i;
   for (i = adj[u].begin(); i != adj[u].end(); ++i) {
       int v = *i; // v is current adjacent of u
       // If v is not visited yet, then make it a child of u
       // in DFS tree and recur for it
       if (!visited[v]) {
           children++;
           parent[v] = u;
           APUtil(v, visited, disc, low, parent, ap);
           // Check if the subtree rooted with v has a connection to
           // one of the ancestors of u
           low[u] = min(low[u], low[v]);
           // u is an articulation point in following cases
           // (1) u is root of DFS tree and has two or more chilren.
           if (parent[u] == NIL && children > 1)
           ap[u] = true;
           // (2) If u is not root and low value of one of its child is more
           // than discovery value of u.
           if (parent[u] != NIL && low[v] >= disc[u])
           ap[u] = true;
       // Update low value of u for parent function calls.
       else if (v != parent[u])
           low[u] = min(low[u], disc[v]);
   }
}
```

```
* Segment Tree with Lazy Propogation build: O(n*logn), query: O(logn + k)
int tree[MAX] = {0};int lazy[MAX] = {0};
void updateRangeUtil(int si,int ss,int se,int us,int ue,int diff){
   if (lazy[si] != 0){
      tree[si] += (se-ss+1)*lazy[si];
      if (ss != se){
          lazy[si*2 + 1] += lazy[si]; lazy[si*2 + 2] += lazy[si];
      }
      lazy[si] = 0;
  }
  if (ss>se || ss>ue || se<us) return; // out of range</pre>
  if (ss>=us && se<=ue){</pre>
      tree[si] += (se-ss+1)*diff;
      if (ss != se){
          lazy[si*2 + 1] += diff; lazy[si*2 + 2] += diff;
      }
      return;
  }
  int mid = (ss+se)/2;
 updateRangeUtil(si*2+1, ss, mid, us, ue, diff);
  updateRangeUtil(si*2+2, mid+1, se, us, ue, diff);
 tree[si] = tree[si*2+1] + tree[si*2+2];
}
void updateRange(int n, int us, int ue, int diff){updateRangeUtil(0, 0, n-1, us, ue,
diff);}
int getSumUtil(int ss, int se, int qs, int qe, int si){
  if (lazy[si] != 0){
      tree[si] += (se-ss+1)*lazy[si];
       if (ss != se){lazy[si*2+1] += lazy[si]; lazy[si*2+2] += lazy[si];}
      lazy[si] = 0;
  }
 if (ss>se || ss>qe || se<qs) return 0;</pre>
 if (ss>=qs && se<=qe) return tree[si];</pre>
  int mid = (ss + se)/2;
  return getSumUtil(ss,mid,qs,qe,2*si+1)+getSumUtil(mid+1,se,qs,qe,2*si+2);
int getSum(int n, int qs, int qe){
  if (qs < 0 || qe > n-1 || qs > qe) return -1;
 return getSumUtil(0, n-1, qs, qe, 0);
}
```

```
void constructSTUtil(int arr[], int ss, int se, int si){
  if (ss > se) return;
  if (ss == se){ tree[si] = arr[ss]; return;}
  int mid = (ss + se)/2; constructSTUtil(arr, ss, mid, si*2+1);
  constructSTUtil(arr, mid+1, se, si*2+2);
  tree[si] = tree[si*2 + 1] + tree[si*2 + 2];
}
void constructST(int arr[], int n){constructSTUtil(arr, 0, n-1, 0);}
int main(){
  int arr[] = {1, 3, 5, 7, 9, 11}; int n = sizeof(arr)/sizeof(arr[0]);
  constructST(arr, n);
  printf("Sum of values in given range = %d\n",getSum(n, 1, 3));
  updateRange(n, 1, 5, 10); // Add 10 to all nodes at indexes from 1 to 5.
  printf("Updated sum of values in given range = %d\n",getSum( n, 1, 3));
  return 0;
}
```