



Extended PCC Vivace: Online-Learning Congestion Control

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Extended PCC Vivace: Online-Learning Congestion Control





Extension 1 - Traffic-Aware Utility Selection

Baseline Utility Function

$$U = t \cdot S(d) - \lambda \cdot t \cdot l$$

Where:

t : **Throughput** (bits/sec)

d : **Delay** (or latency)

l : **Loss rate** (fraction of packets lost)

$S(d)$: a *delay sensitivity function* reduces the score if delay is high

λ : weight controlling how strongly to penalize loss



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Solution

Traffic classifier + specialized utility functions



Extended PCC Vivace: Online-Learning Congestion Control

Extension 1 - Traffic-Aware Utility Selection

STEP 1: AUTOMATIC TRAFFIC CLASSIFICATION - *Traffic Classifier*

FEATURE EXTRACTION

- Packet Size Distribution
(avg, std, median, max, min)
- Inter-Arrival Times (IAT)
(mean, variance, coefficient of variation)
- Entropy of Packet Sizes
(measures size variability)
- Burst Detection
(identifies bursty patterns)
- Periodicity Score
(consistency in timing)



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TRAFFIC CLASSIFICATION

- BULK (Throughput Optimization)
 - Avg packet size ≥ 1200 bytes
 - Low size variance, high entropy
- STREAMING (Stability Focus)
 - Medium packets (700-1200 bytes)
 - Consistent timing, moderate bursts
 - High periodicity score
- REALTIME (Latency Optimization)
 - Small packets < 300 bytes
 - Low variance, frequent bursts
 - Periodic inter-arrivals
- DEFAULT (Fallback)
 - Low confidence threshold
 - Unknown/mixed traffic pattern



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CLASSIFICATION PROCESS FLOW

Observe Packets (min. 20) → Extract Features → Score Each Type (bulk/streaming/realtime) → Select Best Match
→ Validate Against Confidence Threshold (default: 0.5) → Output: (Traffic Type, Confidence Score)
→ MetaController Uses Classification to Select Specialized Utility Function → Optimize Network Parameters



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Extension 1 - Traffic-Aware Utility Selection

STEP 2: SPECIALIZED UTILITY FUNCTIONS - *Utility Function Bank*



FILE TRANSFER (FTP)



STREAMING



REAL TIME



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FILE TRANSFER (FTP)



VIDEO STREAMING



VOIP (Video Call)

1. Bulk Transfer Utility (U_{bulk})

$$U_{bulk} = \alpha_1 T - \alpha_2 l$$

Meaning:

T : throughput (bits/sec)

l : loss rate (fraction of packets lost)

α_1, α_2 : weights that balance importance of throughput vs. loss



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2. Real-time Utility (U_{realtime})

$$U_{\text{realtime}} = \beta_1 T \cdot S(L_{\text{max}})$$

Meaning:

T : throughput (enough to carry audio/video data)

L_{max} : tail latency (worst-case delay)

$S(L_{\text{max}})$: decreasing function - gives low score when tail latency is large

β_1 : scaling weight



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VOIP (Video Call)

3. Streaming Utility ($U_{\text{streaming}}$)

$$U_{\text{streaming}} = \gamma_1 T \cdot I(T > T_{\text{min}}) - \gamma_2 \text{Var}(T)$$

Meaning:

T : throughput

$I(T > T_{\text{min}})$: indicator = 1 if throughput above threshold, else 0

$\text{Var}(T)$: variance of throughput (instability)

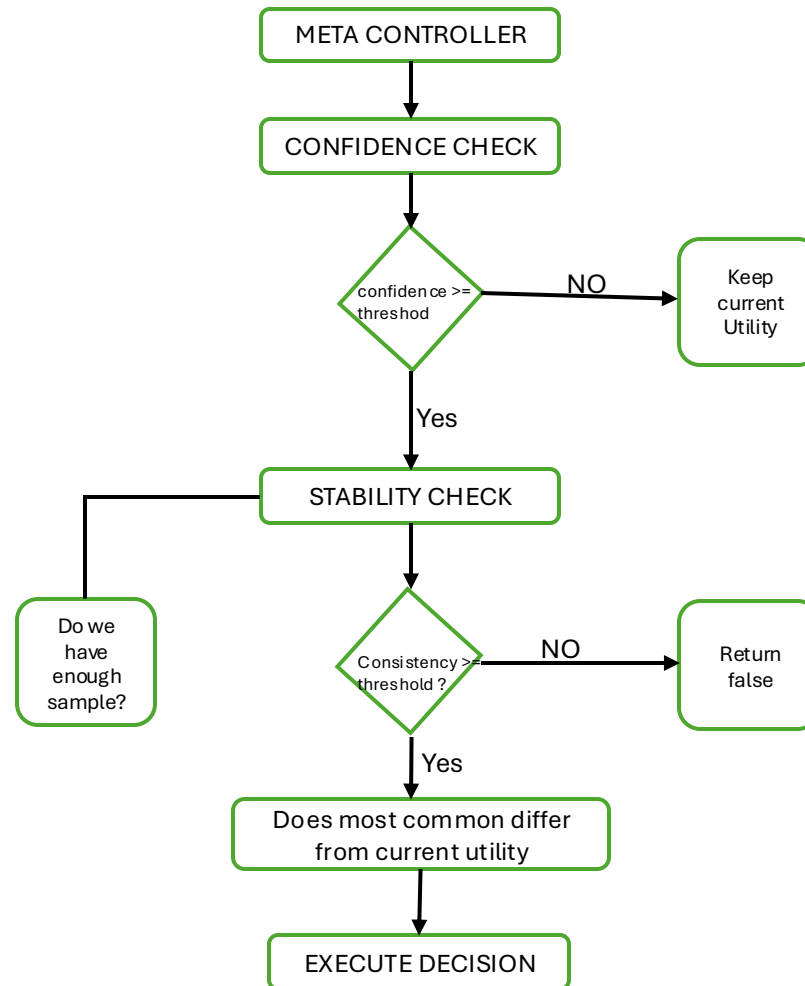
γ_1, γ_2 : weighting constants



Extension 1 - Traffic-Aware Utility Selection



STEP 3: INTELLIGENT SWITCHING - *MetaController*

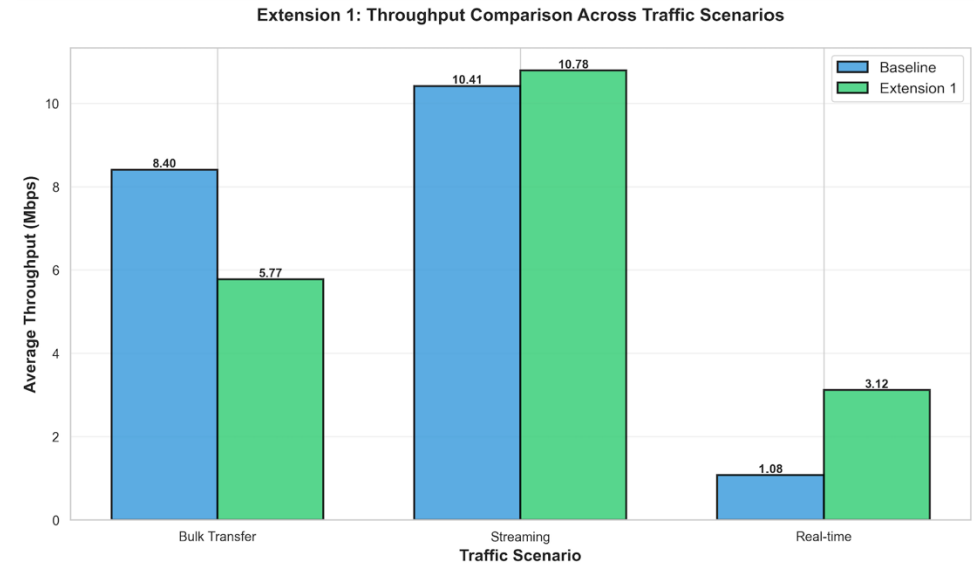
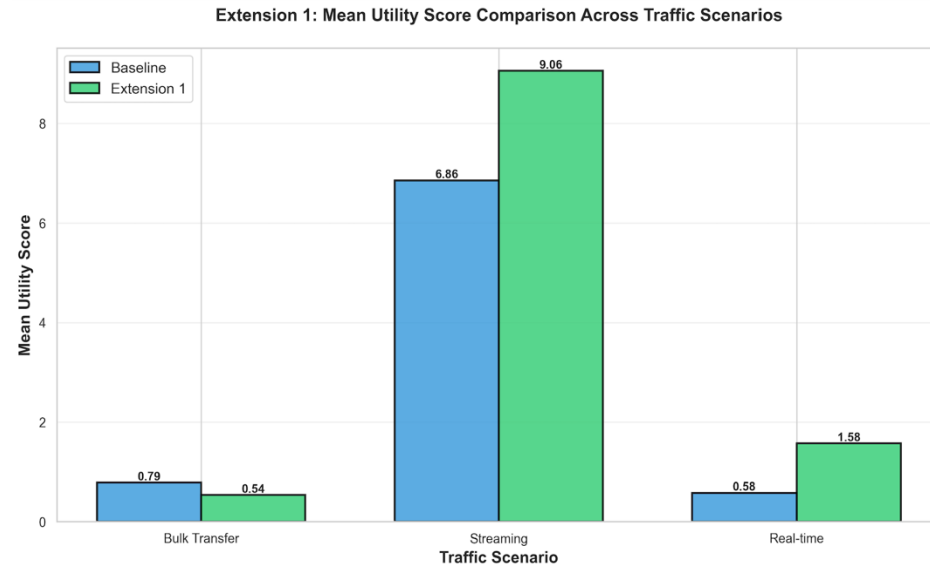




Extended PCC Vivace: Online-Learning Congestion Control

Extension 1 - Traffic-Aware Utility Selection

Results:

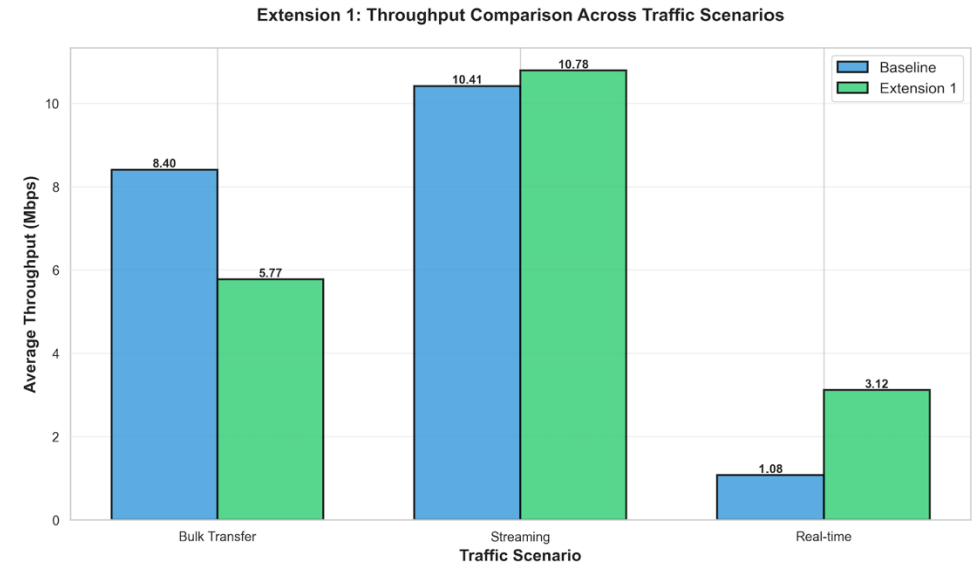
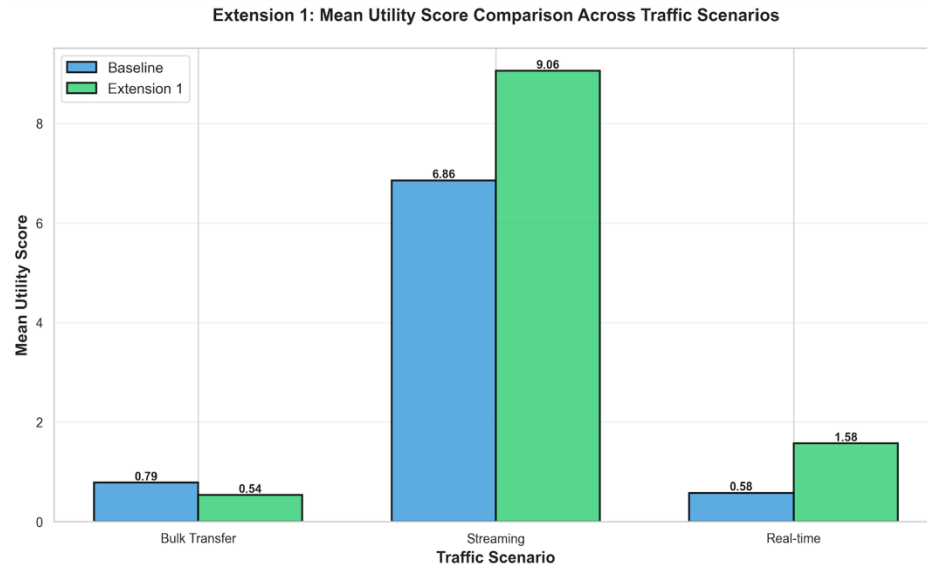




Extended PCC Vivace: Online-Learning Congestion Control

Extension 1 - Traffic-Aware Utility Selection

Results:



Summary (In One Line)

Extension 1 transforms the generic PCC baseline into an intelligent, context-aware system that automatically classifies traffic and dynamically applies the most suitable utility function-ensuring each application achieves its ideal network performance.

Extended PCC Vivace: Online-Learning Congestion Control



Extension 2 - Loss-Type Awareness





Extension 2 - Loss-Type Awareness



Problem

Not all losses mean congestion.

In reality, packet loss can come from two completely different causes:

Type of Loss	Cause	Correct Reaction	Why PCC Vivace Fails
Congestion Loss	Queue overflow at router when link is saturated	Reduce rate	✓ This is correct behavior
Wireless Loss	Random link-layer corruption in Wi-Fi or cellular (fading, interference)	Keep rate steady	✗ PCC wrongly reduces rate

Extended PCC Vivace: Online-Learning Congestion Control

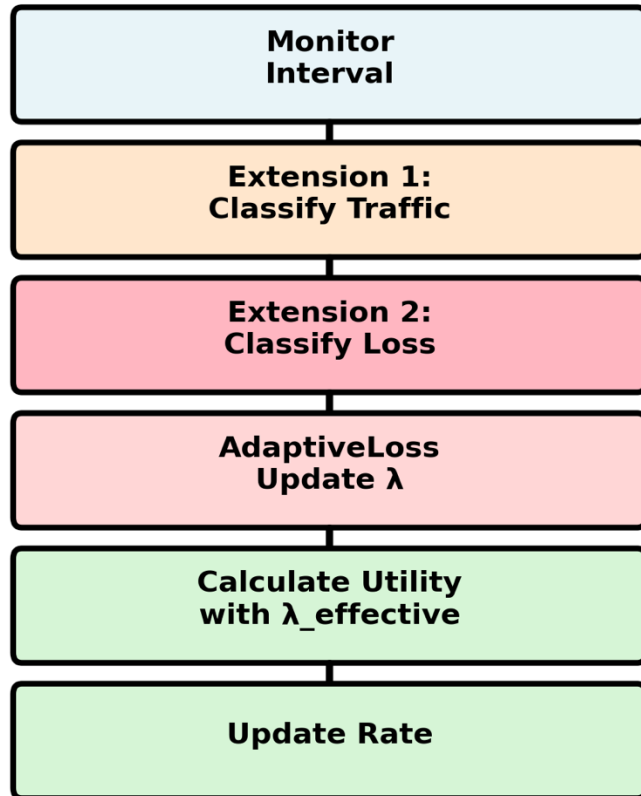


Extension 2 - Loss-Type Awareness

Approach: Correlation-based classification (Pearson coefficient)



Extension 2 Integration Flow



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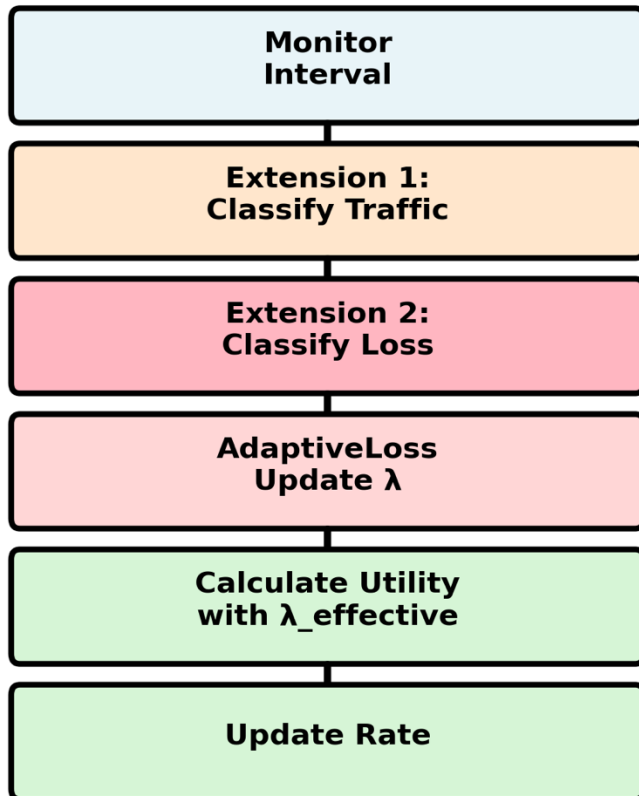


Extension 2 - Loss-Type Awareness



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Extension 2 Integration Flow



In each Monitor Interval (every ~100ms):

1. Measure loss and RTT
→ Send to Loss Classifier
2. Loss Classifier analyzes the data
→ Calculates correlation
→ Returns: $p_{\text{wireless}} = 0.8$, confidence = 0.9
3. Adaptive Loss Coefficient receives this information
→ Calculates $\lambda = 2.0 + (\text{correlation-based adjustment})$
4. When calculating utility:
→ Instead of using fixed $\lambda = 10.0$
→ Use $\lambda = 2.0$ (for wireless) or 10.0 (for congestion)
5. PCC adjusts rate based on this utility
→ If wireless loss: Maintains rate (high utility)
→ If congestion loss: Reduces rate (low utility)

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Extension 2 - Loss-Type Awareness



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Component 1: Loss Classifier — Detect What Type of Loss

Core Idea

The classifier uses **correlation between packet loss and RTT (Round-Trip Time)** to identify the nature of packet loss.

Type	Observation	Correlation
Congestion Loss	RTT increases with loss (queue builds up)	High correlation
Wireless Loss	RTT stable despite loss	Low correlation

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Extension 2 - Loss-Type Awareness

Approach: Correlation-based classification (Pearson coefficient)



Step 1: Collect Raw Data

Every time a monitor interval (MI) ends:

- Log packet losses (loss_rate)
- Log RTT samples (rtt)

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Step 2: Compute RTT Inflation

- $\text{RTT inflation} = \text{current RTT} - \text{baseline RTT}$
- If RTT is much larger than baseline, the buffer is filling up \rightarrow likely congestion.

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Step 3: Align Loss & RTT Events by Time

For each RTT event, find the nearest loss event.

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Step 4: Calculate Correlation

Use **Pearson correlation** between binary loss and RTT arrays:

$$\rho = \frac{\sum (x - \mu_x)(y - \mu_y)}{\sigma_x \sigma_y}$$

Where:

$\rho = 1$: perfectly correlated (congestion)

$\rho = 0$: uncorrelated (wireless)

$\rho = -1$: inversely correlated (rare)

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Step 5: Classify Based on Correlation

Set thresholds:

- High threshold = 0.5 → Congestion
- Low threshold = 0.2 → Wireless
- Between = Mixed

$$\begin{cases} \rho > 0.5 & \Rightarrow \text{Congestion } (p_{\text{wireless}} = 0.1) \\ \rho < 0.2 & \Rightarrow \text{Wireless } (p_{\text{wireless}} = 0.9) \\ \text{else} & \Rightarrow \text{Mixed } (p_{\text{wireless}} \approx 0.5) \end{cases}$$

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Extension 2 - Loss-Type Awareness



Component 2: Adaptive Loss Coefficient - Adjust the Penalty

In the utility function, the **loss penalty (λ)** should depend on the **loss type**.

➤ Traditional PCC Utility

$$U = T - \lambda \cdot L$$

T = Throughput

L = Loss rate

$\lambda = 10.0$ (fixed)

Problem: λ is constant, so the controller overreacts even to wireless losses.

➤ Extension 2 Utility

$$U = T - \lambda_{adaptive} \cdot L$$

Where:

$$\lambda_{adaptive} = \lambda_{base}(1 - p_{wireless}) + \lambda_{wireless}(p_{wireless})$$

Parameters:

$\lambda_{base} = 10.0$ → congestion penalty

$\lambda_{wireless} = 2.0$ → wireless penalty

$p_{wireless}$ → detected wireless fraction

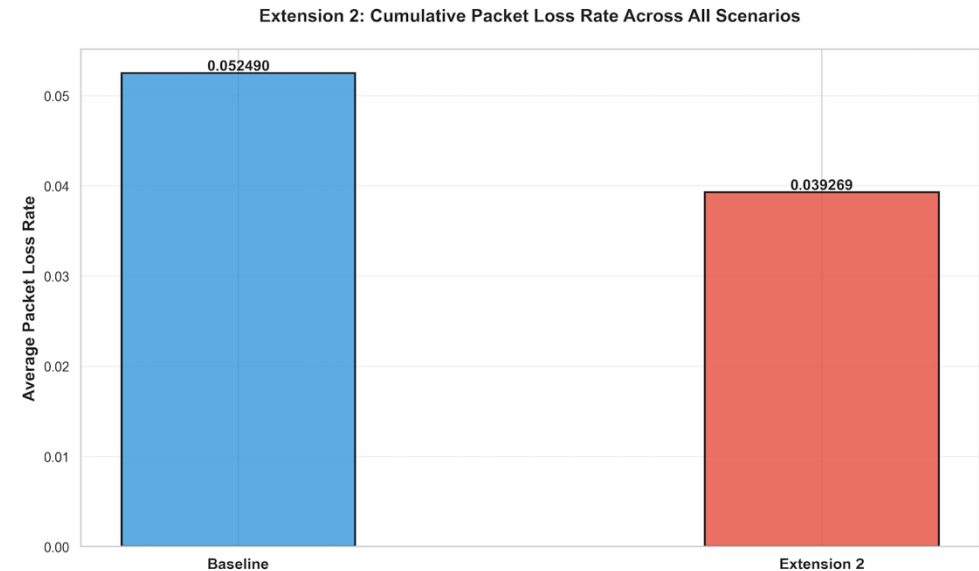
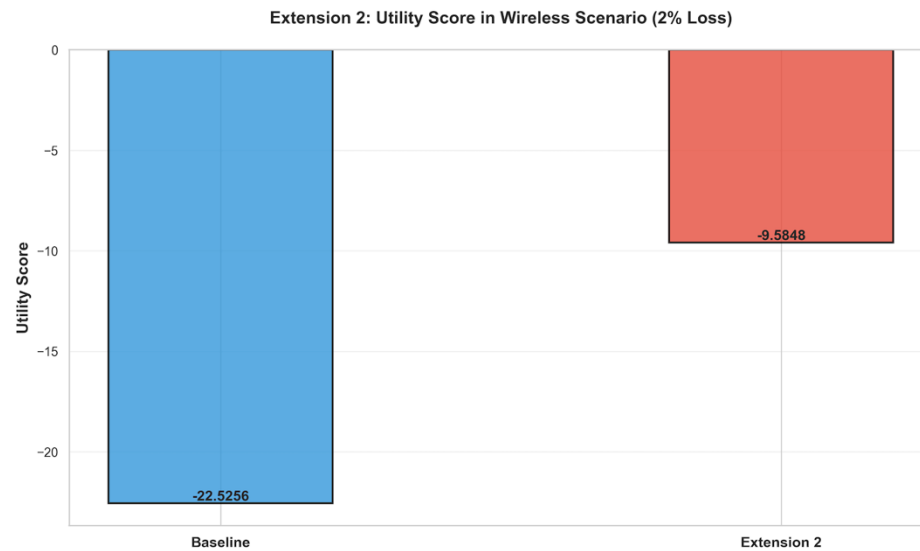
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Extension 2 - Loss-Type Awareness



Results:



Summary (In One Line)

Extension 2 makes PCC Vivace loss-type aware by detecting whether losses stem from congestion or wireless errors and adaptively adjusting the loss penalty, preventing unnecessary rate drops



Extension 3 - Distributed Fairness Control

Baseline Problem: Unfair Bandwidth Sharing

Why It Happens - Gradient Oscillation Problem

PCC Vivace adjusts rate via *gradient feedback* (utility changes)

When multiple flows explore simultaneously:

- Their actions interfere
- Gradients oscillate (+, -, +, -)
- No stable convergence

The more aggressive flow keeps increasing → monopolizes capacity Other flows keep backing off
→ starve

Result: Persistent unfairness & slow convergence.



Extension 3 - Distributed Fairness Control

Baseline Problem: Unfair Bandwidth Sharing

Goal: Achieve *Distributed Fairness*

→ Each flow detects contention, cooperates, and balances bandwidth.

Main Components:

- ❖ **Contention Detector** -- Detect if other flows are competing
- ❖ **Cooperative Explorer** -- Coordinate exploration (take turns)
- ❖ **Virtual Queue Estimator** -- Monitor queue depth from RTT
- ❖ **Fairness Controller** -- Penalize greedy flows in utility



Extension 3 – Distributed Fairness Control

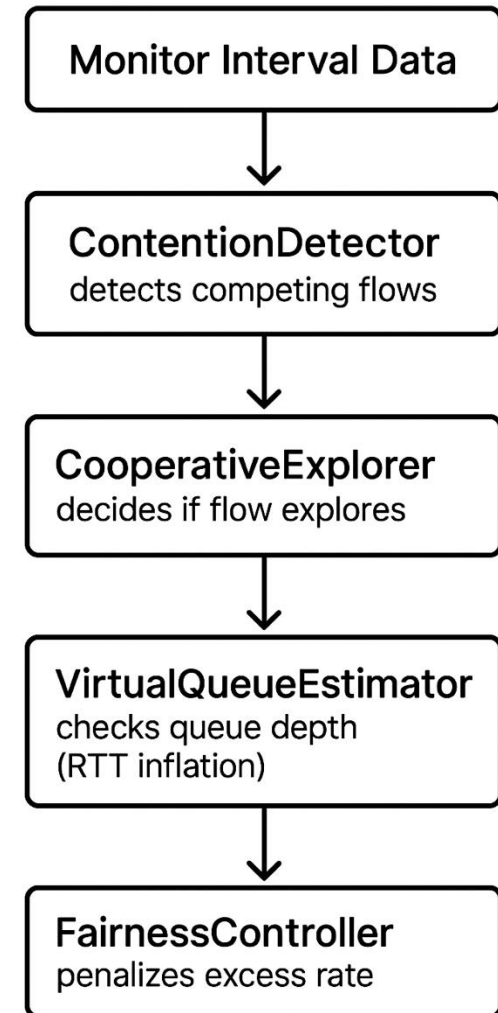
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Extension 3 – Distributed Fairness Control

❖ Contention Detector - Detect Competition

Problem:

Flows don't know when others are present.

Insight:

Gradient sign pattern reveals contention:

Solo: monotonic (+ + + +)

Competing: oscillating (+ - + -)

Algorithm:

Count gradient sign changes over last 5 intervals

Compute ratio & volatility

Classify as SOLO / LIGHT / MODERATE / HEAVY
contention

Example:

Gradients = [+0.5, -0.8, -0.2, +0.5, -0.7]

→ 4 sign changes → **HEAVY contention** (4–5 flows)



Extension 3 – Distributed Fairness Control

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❖ Cooperative Explorer — Avoid Collisions

Problem:

All flows explore together → collisions → oscillations.

Solution:

Hash-based turn-taking (no communication):

slot = current_time // cycle

if hash(flow_id, slot) % 2 == 0:

 explore = True

Effect:

Only one flow explores per slot

Others wait, then take their turn

Smooth convergence & no interference

Result:

→ Turn-based exploration → stability achieved!



Extension 3 – Distributed Fairness Control

❖ Virtual Queue Estimator — Detect Monopolization

Problem:

Dominant flows fill queues without noticing.

Idea:

Estimate queue depth from RTT inflation:

$$Queue\ Delay = RTT_{current} - RTT_{baseline}$$

Baseline RTT = 5th percentile of last samples

Queue fills → larger inflation → congestion detected

Action:

If queue depth > threshold → reduce rate.

→ Prevents buffer monopolization.



Extension 3 – Distributed Fairness Control

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❖ FairnessController — Equalize Sharing

Problem:

Even if contention detected, flows need fairness enforcement.

Solution:

Add *fairness penalty* to utility function:

r_{fair} = total bandwidth / estimated flow count

Penalize flows sending above their fair share

Leave under-utilizing flows unpenalized

Example:

Flow A: 6.5 Mbps → penalty = $(6.5 - 5)^2 = 2.25$

Flow B: 3.5 Mbps → penalty = 0

→ A's utility ↓ → reduces rate; B's utility ↑ → increases rate

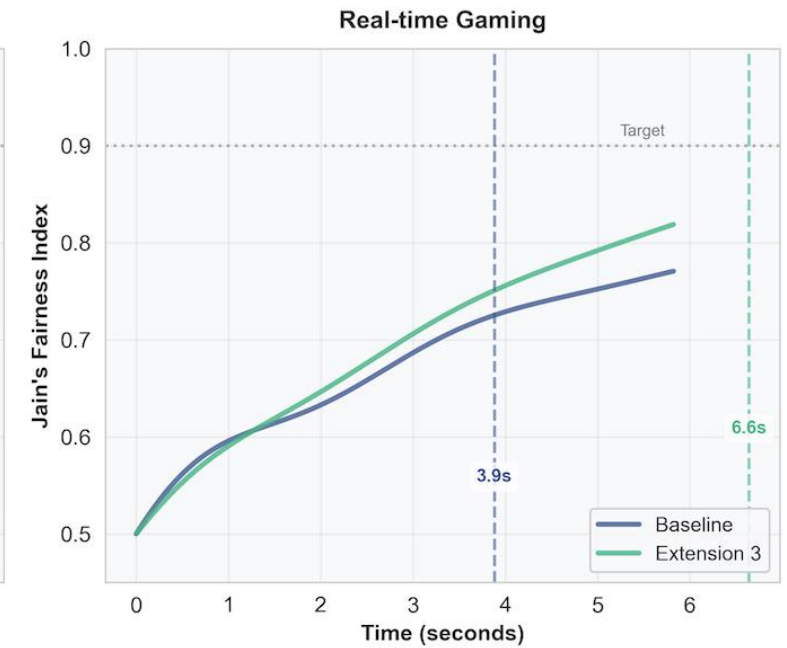
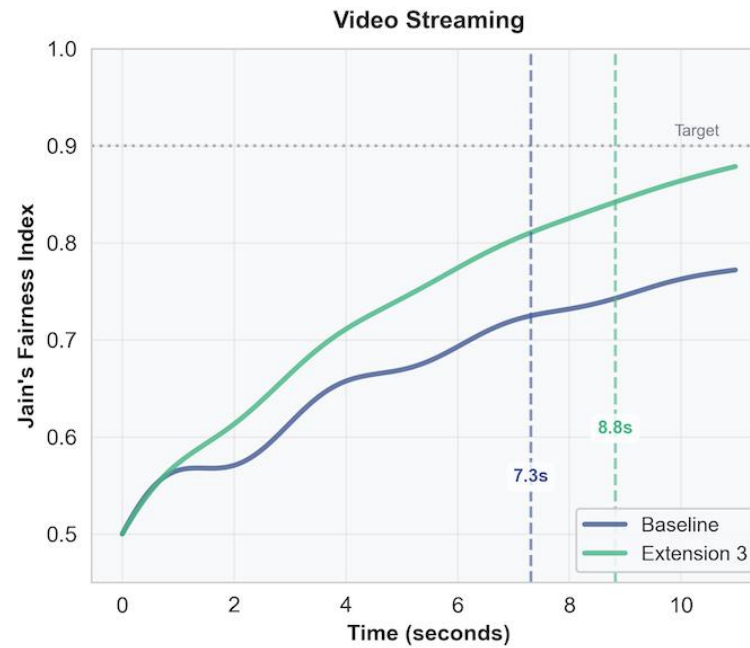
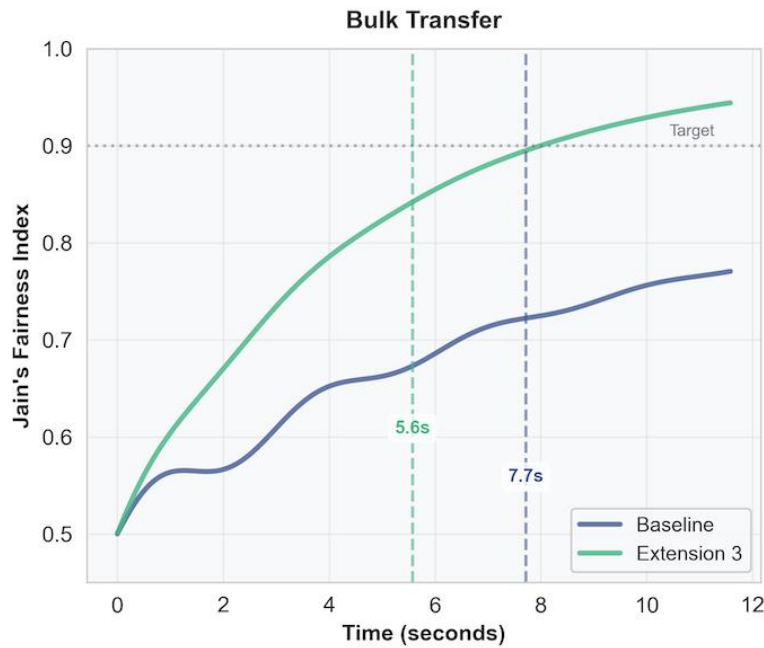
→ Fair equilibrium around 5 Mbps each.



Extension 3 – Distributed Fairness Control

Results

Extension 3: Fairness Convergence Over Time

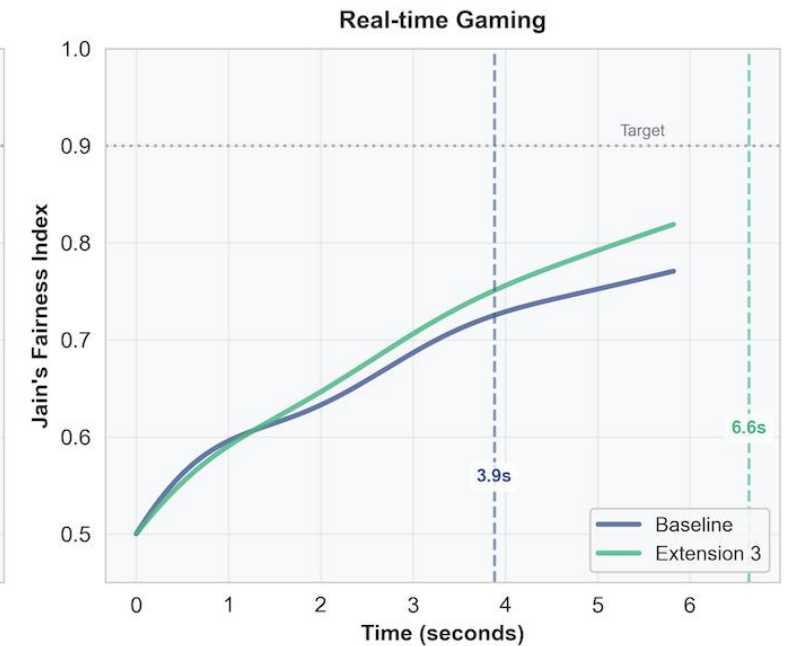
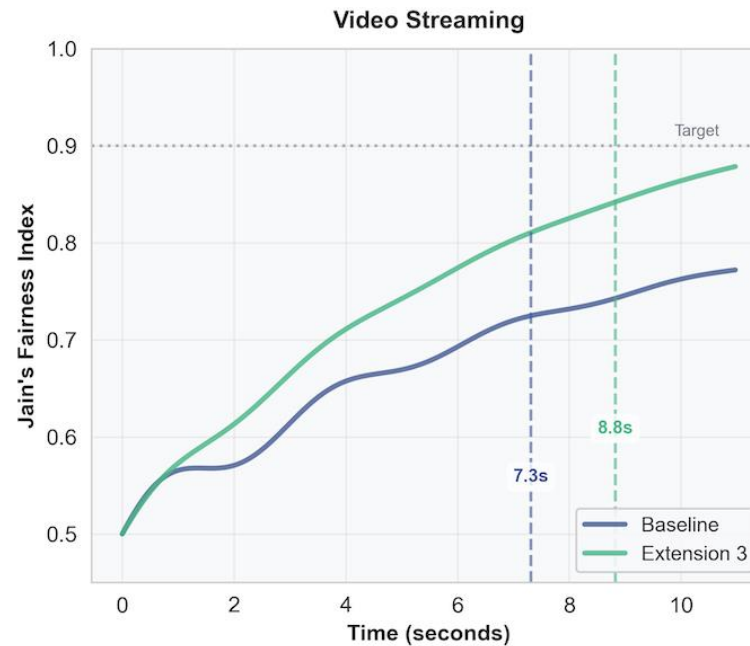
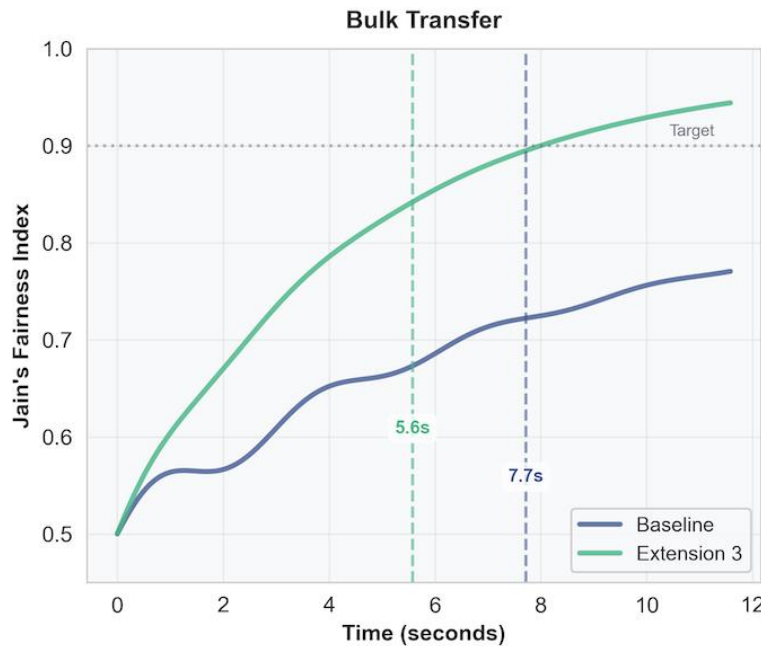




Extension 3 – Distributed Fairness Control

Results

Extension 3: Fairness Convergence Over Time



Summary (In One Line)

Extension 3 ensures distributed fairness by detecting competing flows, coordinating exploration, and enforcing fair rate allocation without any centralized coordination.



Multipath Rate Allocation

■ Problem (Compared to Extension 3)

Extension 3: great fairness & fast convergence.

- But still **single-path routing** → all traffic on one route.
- If that path:
 - becomes congested → throughput drops.
 - fails → complete interruption.
- No multi-path utilization → poor flexibility, under-used capacity.
- System lacks resilience to link failures or congestion.



Multipath Rate Allocation

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Core Idea - Path-Based Load Balancing

Split and adapt traffic across multiple network paths dynamically.

Key Concepts:

- **Path Distribution:** use multiple paths instead of one.
- **Dynamic Rebalancing:** shift load as conditions change.
- **Path Recovery:** reroute instantly if a path fails.
- **Better Utilization:** use every available route efficiently.



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■ Goal

Enable **multi-path routing** to increase throughput, use spare capacity, and survive path failures.

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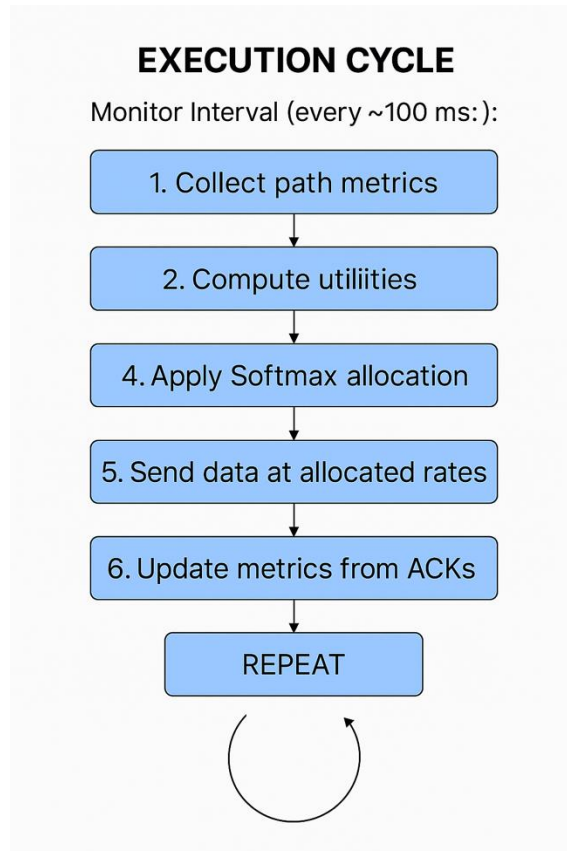
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IMPLEMENTATION FLOW

☐ Path Discovery

Enumerate available interfaces (e.g., eth0, wlan0, lte0).

Create Path objects with baseline RTT, bandwidth, loss.

Maintain ACTIVE / IDLE / FAILED states.



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❑ Path Monitoring

Continuously measure throughput, RTT, loss.
Apply Exponential Moving Average ($\alpha = 0.3$).
Detect degradation → trigger recovery or switch.



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❑ Utility Calculation

For each path:

$$U_p = 0.4T_p + 0.3(1 - L_p) + 0.2(1 - loss_p) + 0.1stability_p$$

Produces a normalized score (0–1).

High utility = good path quality.



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For each path:

$$U_p = 0.4T_p + 0.3(1 - L_p) + 0.2(1 - loss_p) + 0.1stability_p$$

Produces a normalized score (0–1).

High utility = good path quality.

❑ Correlation Detection

Compute RTT correlation ρ between paths.

If $\rho > 0.8$ → same bottleneck → de-prioritize one.



Multipath Rate Allocation

IMPLEMENTATION FLOW

❑ Path Discovery

Enumerate available interfaces (e.g., eth0, wlan0, lte0).
Create Path objects with baseline RTT, bandwidth, loss.
Maintain ACTIVE / IDLE / FAILED states.

❑ Path Monitoring

Continuously measure throughput, RTT, loss.
Apply Exponential Moving Average ($\alpha = 0.3$).
Detect degradation → trigger recovery or switch.

❑ Utility Calculation

For each path:

$$U_p = 0.4T_p + 0.3(1 - L_p) + 0.2(1 - loss_p) + 0.1stability_p$$

Produces a normalized score (0–1).

High utility = good path quality.

❑ Correlation Detection

Compute RTT correlation ρ between paths.
If $\rho > 0.8$ → same bottleneck → de-prioritize one.

❑ SOFTMAX RATE ALLOCATION (CORE LOGIC)

$$r_p = R_{total} \times \frac{e^{U_p/\tau}}{\sum e^{U_{p'}/\tau}}$$

Where:

R_{total} : total sending rate

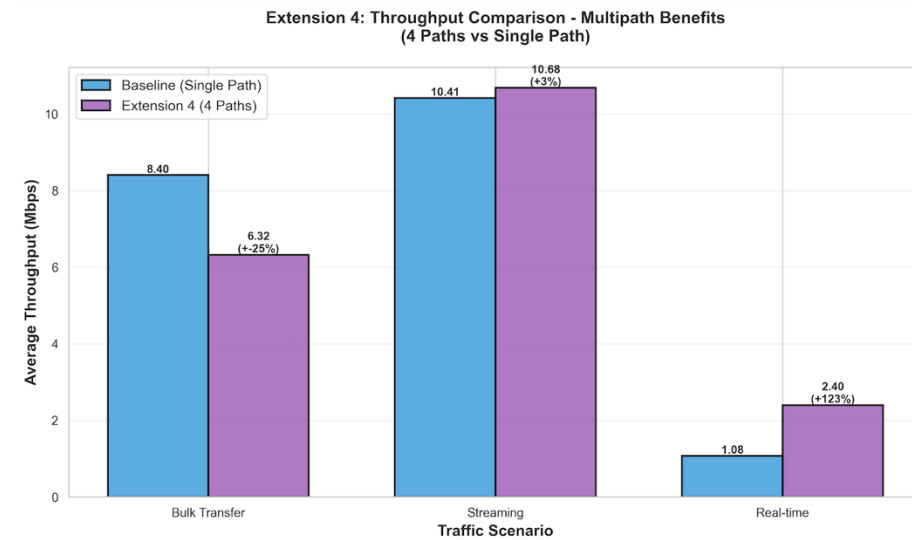
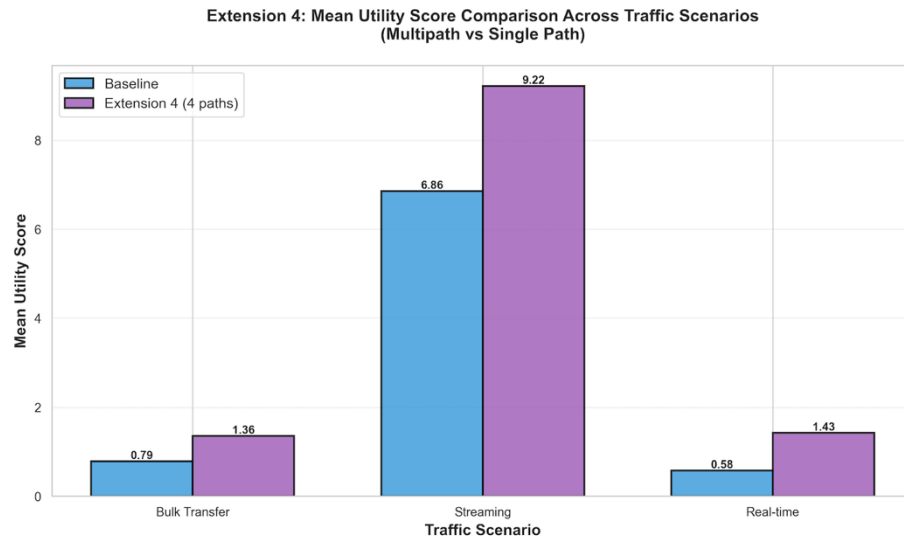
U_p : path utility

τ : temperature (exploration–exploitation balance)



Multipath Rate Allocation

Results:



Summary (In One Line)

Extension 4 enables intelligent multipath rate allocation — dynamically distributing traffic across multiple network paths using utility-based Softmax scheduling to maximize throughput, resilience, and adaptability.

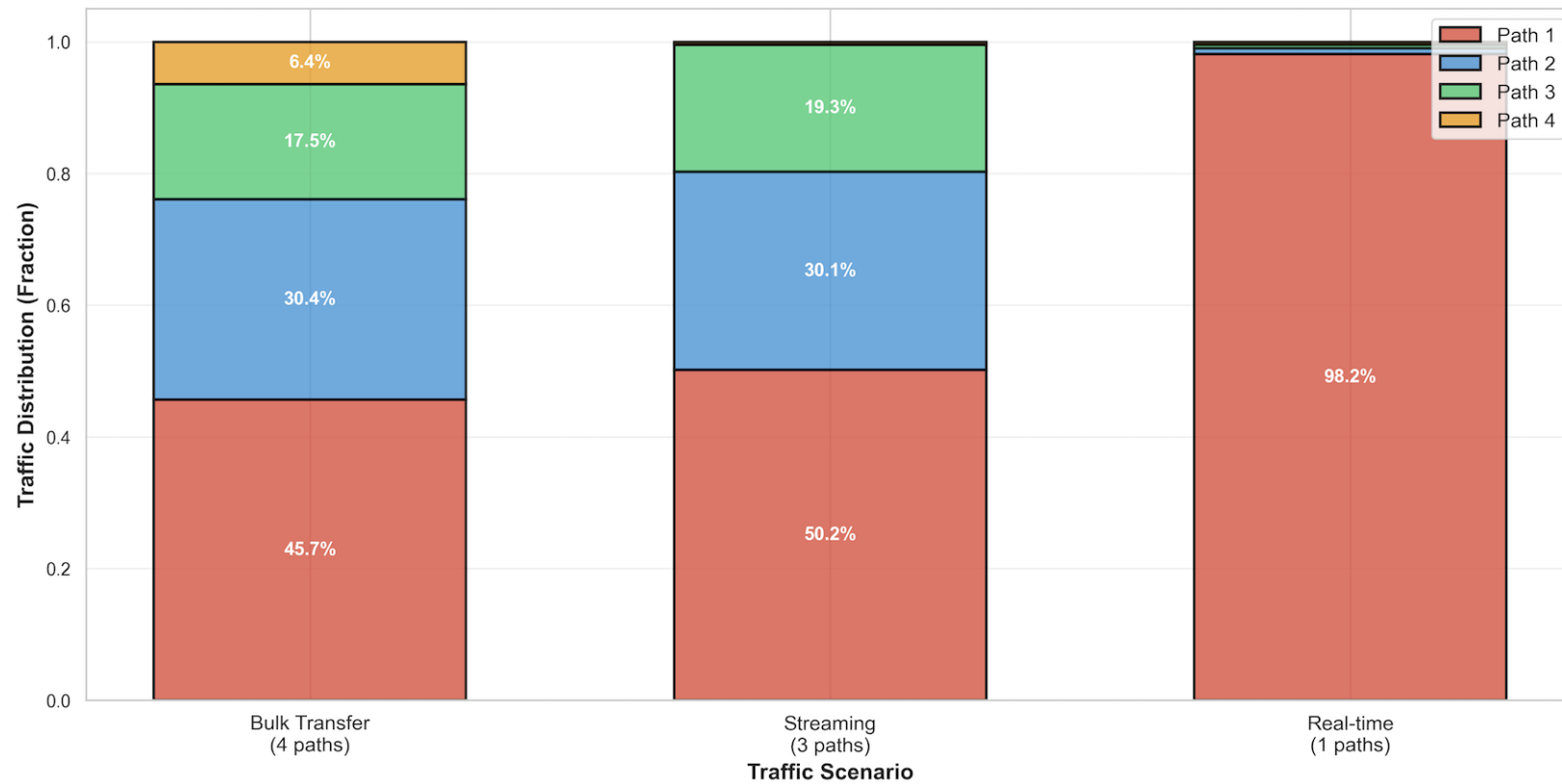


Extended PCC Vivace: Online-Learning Congestion Control

Results:

Extension 4: Real Traffic Distribution Across Paths

Extension 4 Adaptive Path Allocation
(Real Data from Softmax Scheduler)





Future Works: PCC Vivace with Multipath Extensions

- ☐ Application-Aware QoE Optimization (Extension of Extension 1)
- ☐ Parameter optimization for Extensions 1-2 (addressing throughput decreases)
- ☐ Security and Privacy for Multipath Transmission
- ☐ Integration testing
- ☐ Cross-platform validation

THANK YOU