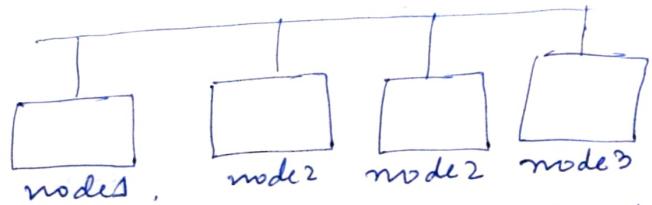
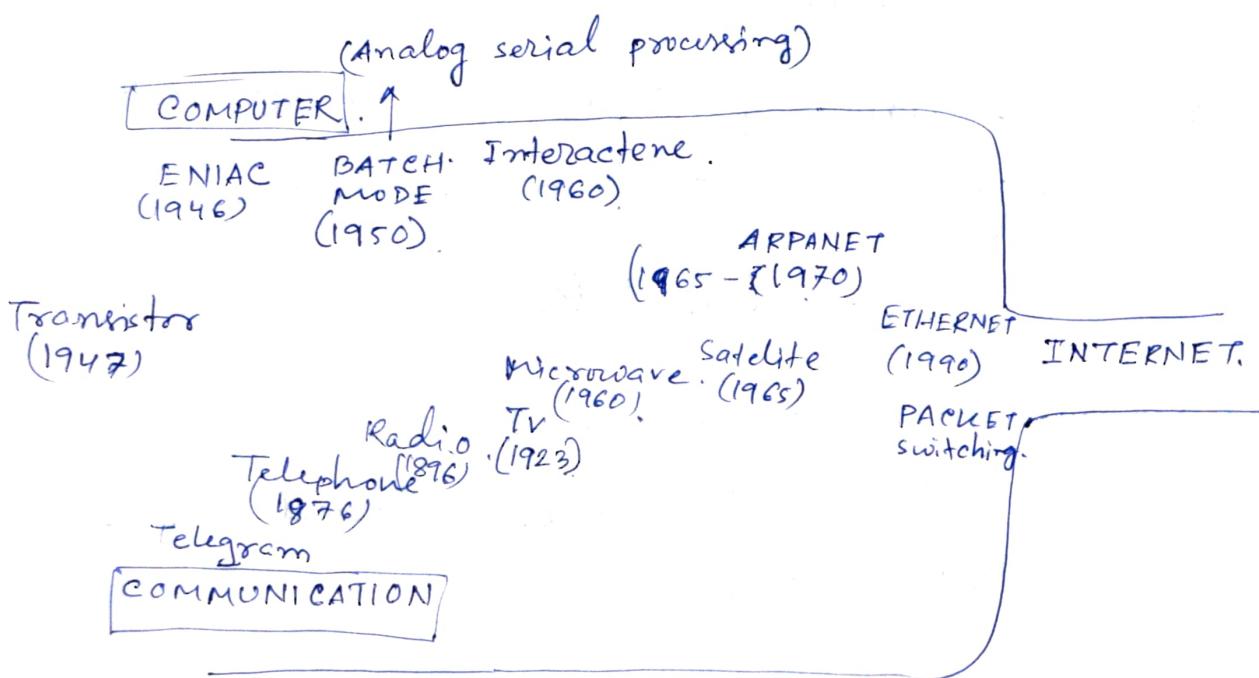


30-07-13  
Tuesday.

- Licklider → vision of the internet → 1963 → ARPANET.  
→ Robert Taylor → 1963 - 1965 [Father of Internet]  
    ↳ First incarnation, ARPANET  
→ NSF NET → 2<sup>nd</sup> incarnation.  
(National Science Foundation) → 1986.  
    →



- NSTNET NET → 3<sup>rd</sup> incarnation → commercialization

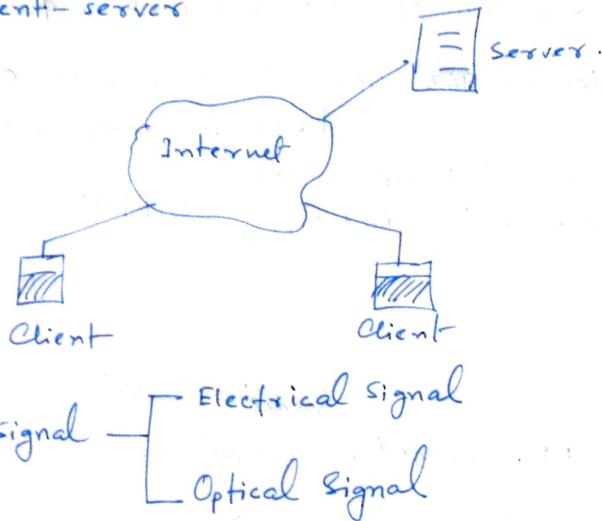


# Assignment 01:

Evolution of Internet.  
; before 6th August.

05.07.13.  
Monday

### \* client-server

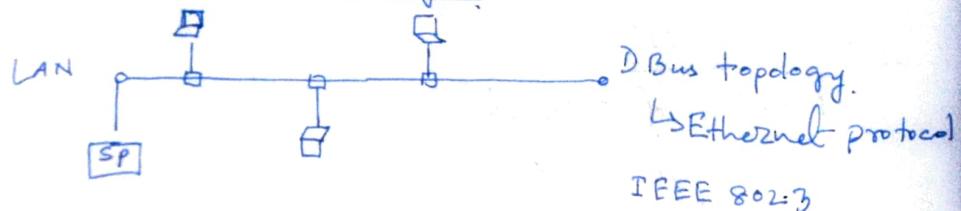


\* Digital Signal provides more bandwidth than analog signal (Proof mathematically)  
↳ Data rate = bandwidth.

### \* Communication Devices:

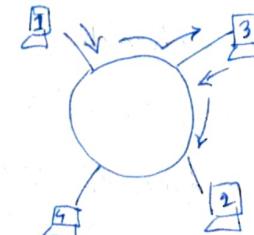
- [1] Co-axial cable [2] Twisted cable.
- [3] Microwave [4]

### # Network topology:



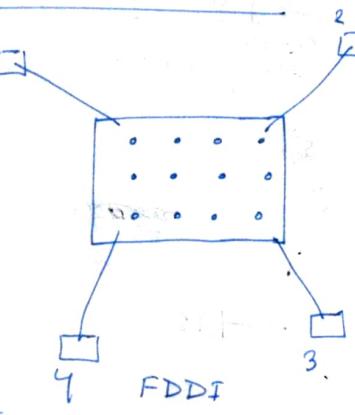
### 2. Ring topology:

IEEE 802.3



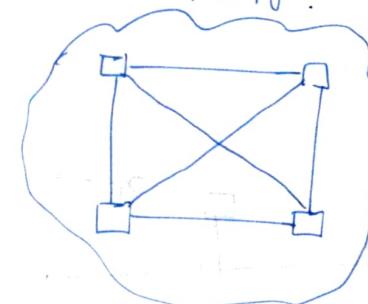
### 3. Fiber distribution:

IEEE 802.3



(Fiber Distribution)  
Data Interface

### 4. Mesh topology:



WAN Technology.

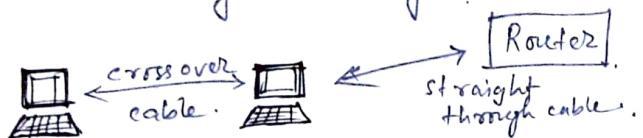
## # Types of Line Cables :-

\* LAN: organization of networks.

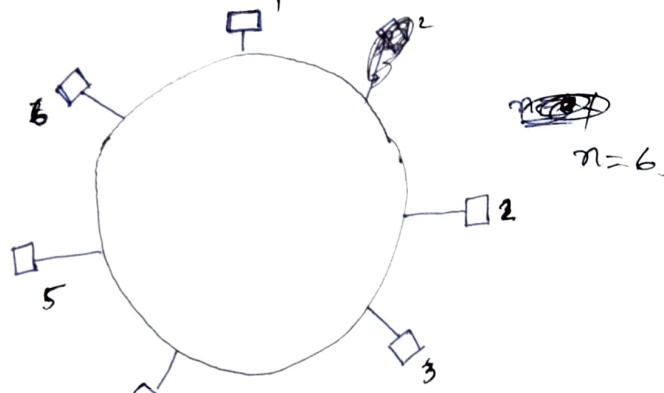
\* UTP: Type of LAN cable (unshielded twisted pair)

\* CATS: line cable  $\leftarrow$  type of UTP

- ↳ Crossover (for same device)
- ↳ straight through (for different devices)

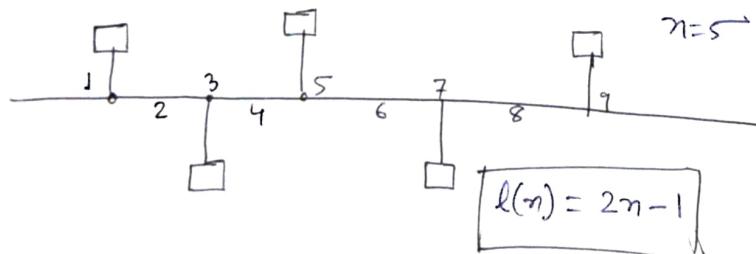


[Q]



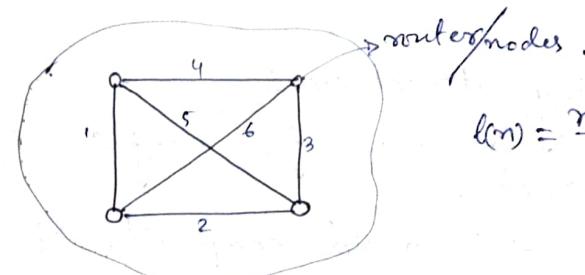
Calculate the no. of links in terms of 'n' (no. of users)

$$r(n) = 2n$$



$$l(n) = 2n - 1$$

[Q]

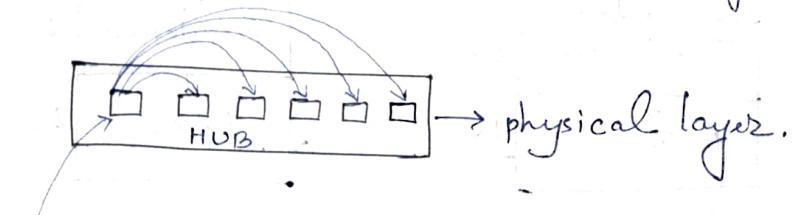


$$l(n) = \frac{n(n-1)}{2}$$

## # Component of internet:

\* HUB: It's towards the signal.

↳ repeaters : enhances the signal.



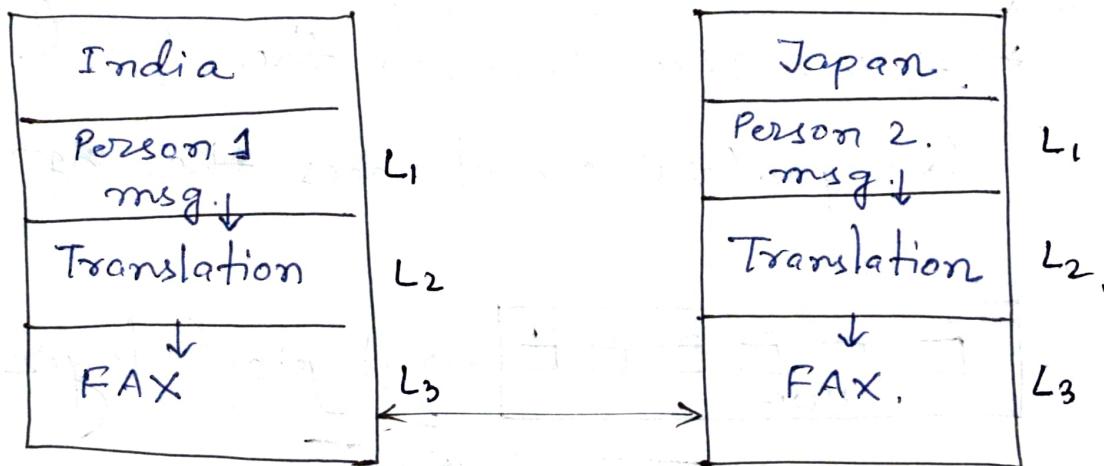
\* Switch: Data link layer  $\rightarrow$  forwarding to particular Rx (LAN)

\* Bridges: Similar f(x). as switch. but for WAN. works in data link layer.

\* Router: L3; Layer 2-3,  $\rightarrow$  Network layer.

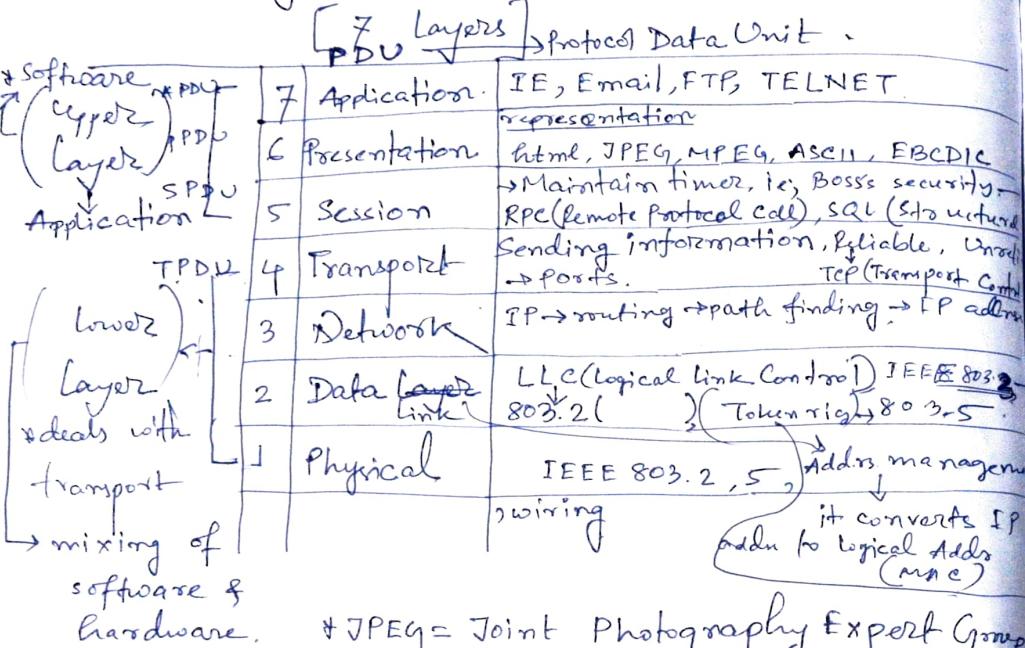
## Layering Structures.

- ↳ needed for easy management of internet.
- ↳ to avoid complexity of the network.



12.08.13  
Mon

# OSI Layer  $\Rightarrow$  Open system interconnection.  
by ISO body.



\* JPEG = Joint Photography Expert Group

\* MPEG = Moving

Data  
↓

01. PPDU  $\rightarrow$  Bits.  $\rightarrow$  raw data.  $[H_1 \dots H_n]$
  02. DPDU  $\rightarrow$  Frames.
  03. NPDU  $\rightarrow$  Packets.
  04. TPDU  $\rightarrow$  Segmentation
  05. SPDU
  06. PPDU  $\rightarrow$
  07. APDU  $\rightarrow$
- $[H_1 | H_2]$   
 $[H_1 | H_2 | H_3]$   
 $[H_1 | H_2 | H_3 | H_4]$   
 $[H_1 | H_2 | H_3 | H_4 | H_5]$   
 $[H_1 | H_2 | H_3 | H_4 | H_5 | H_6]$   
 $[H_1 | H_2 | H_3 | H_4 | H_5 | H_6 | H_7]$

\* RPC : Remote procedure call.

\* Session : Logically maintains the timing.

\* TCP : Transfer Control Protocol.

\* Transport : Ports management.

\* IP : Internet Protocol.

\* LLC : Logical Link Control.

\* Data Link : Address Management

↓  
convert to logical address.

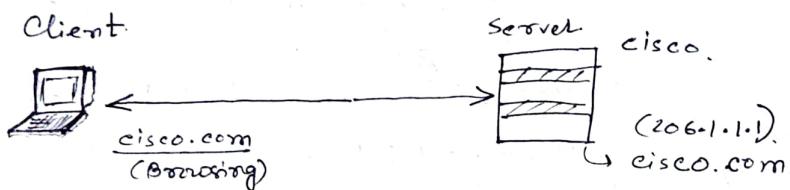
\* Physical : Wiring, physical connection

\* PDU : Protocol Data Unit.

\* APDU : Application Protocol Data Unit.

19.08.13  
Mon

## Mapping of an Application with OSI Layer.



7. # Application: Browsing → IE.  $H_1 H_2 \dots H_7 M$

6. # Presentation: html → http: //cisco.com  $H_1 H_2 \dots H_7 M$

5. \* Session : start & End of session  $H_1 H_2 \dots H_7 M$

4. # Transport Layer: Selection of port based on application.  
e.g; IE 80, Email: 25, FTP: 21.

S: 109 → will be changing router to router.  
D: 80

3. # Network: Assign IP Address/ logical address.  $H_1 H_2 \dots H_7 M$

S: 10.1.1.1 → Again it'll be changing.  
D: 206.1.1.1

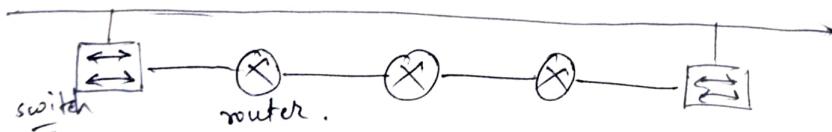
2. # Data link: Logical Address to MAC Address.  $H_1 H_2 \dots H_7 M$

10.1.1.1 →  $\text{MAC}_1 \text{MAC}_2 \dots \text{MAC}_7$

206.1.1.1 →  $\text{MAC}_1 \text{MAC}_2 \dots \text{MAC}_7$

1. # Physical

$H_1 H_2 \dots H_7 M$  (Header)



## # cmd commands:

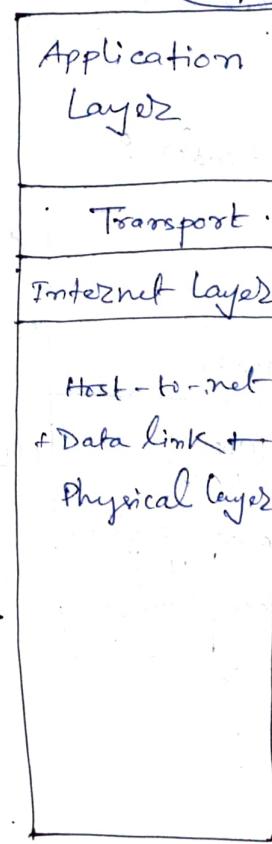
1 ipconfig /all

2 c:>> netstart

Shows  
APP: Cisco.com IP: S/D MAC: S/D  
u: u u u

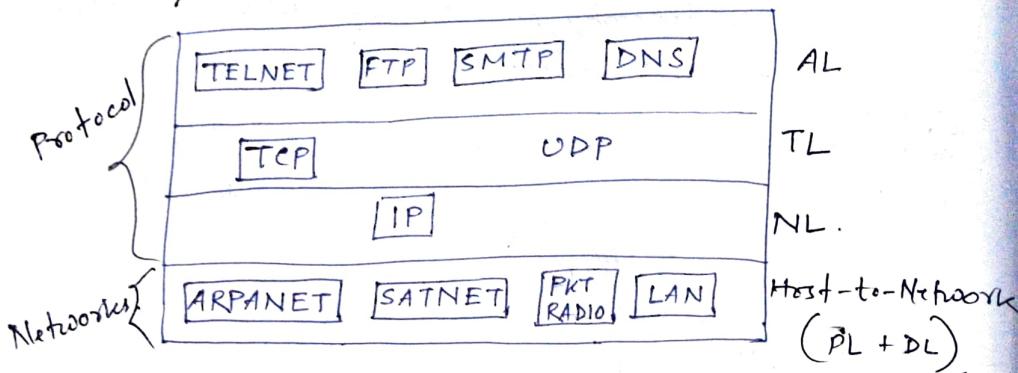
status:  
Closed.  
Started.

TCP/IP

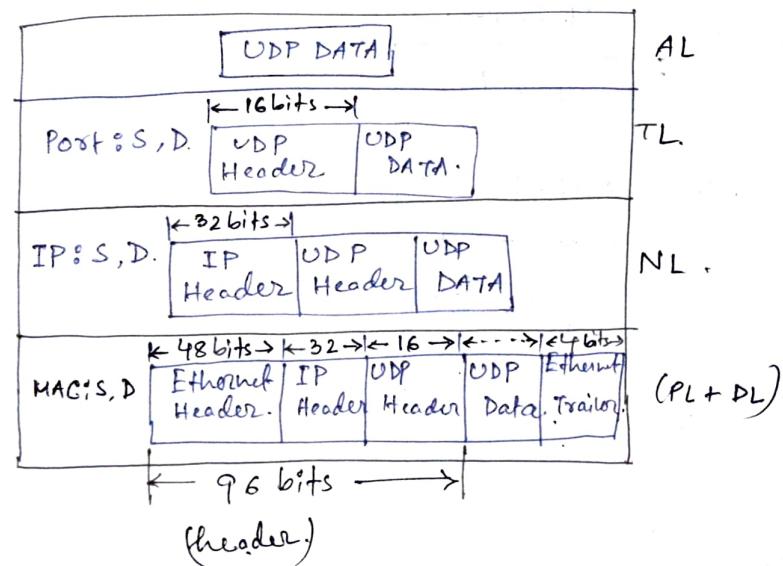


21.08.13  
wed

## ⇒ TCP/IP Protocol :



## ⇒ Encapsulation of UDP Header :



## # Physical Layer :

1. Data Communication : Mix of Analog + Digital.
2.
  - i Communication Media.
  - ii Modulation
  - iii Multiplexing & Switching.
  - iv Terminal Device.
  - v Communication Protocols & Standards.

i ex: STP (Single Twisted Pair) ] → wiring.  
UTP (Unshielded Twisted Pair)

Copper wire = 10baseT — LAN cable.



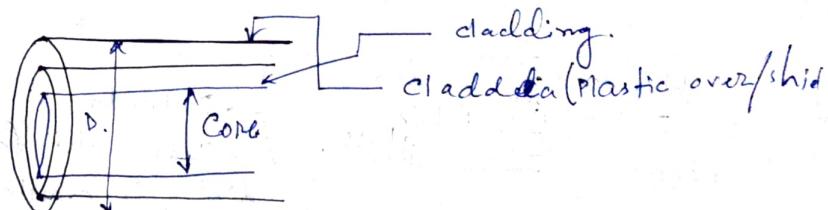
(STP).



co-axial cable.

## # Optical fiber :

Total internal reflection.



$D = 2 - 125 \mu\text{m}$ .

### \* Optical fiber

#### Single Mode

1. One light signal.

2. core  $8-9 \mu\text{m}$ .

3. wavelength 1300 to 1550 nm

$$f = \frac{c}{\lambda}$$

4. for long distance comm.

5. Less Chromatic Dispersion

#### Multimode

Multiple light signals.

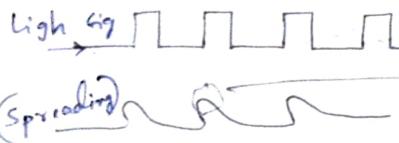
core  $62.5/50\mu\text{m}/100\mu\text{m}$ .

frequency

160 - 500 MHz

For short distance.

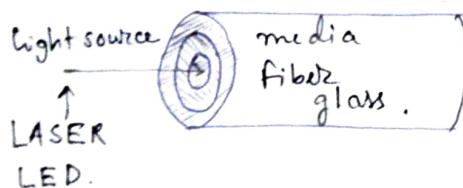
5. More Chromatic Dispersion



Solitons:

(Pulses without spreading)

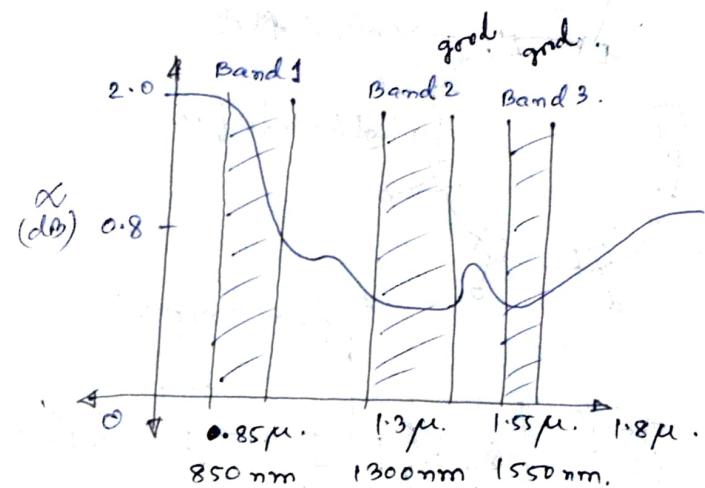
22.08.13  
Thursday



\* Here light fall on photo diod & it emits electricity.

$$\alpha = \text{Attenuation} = 10 \log_{10} \frac{P_r}{P_t}$$

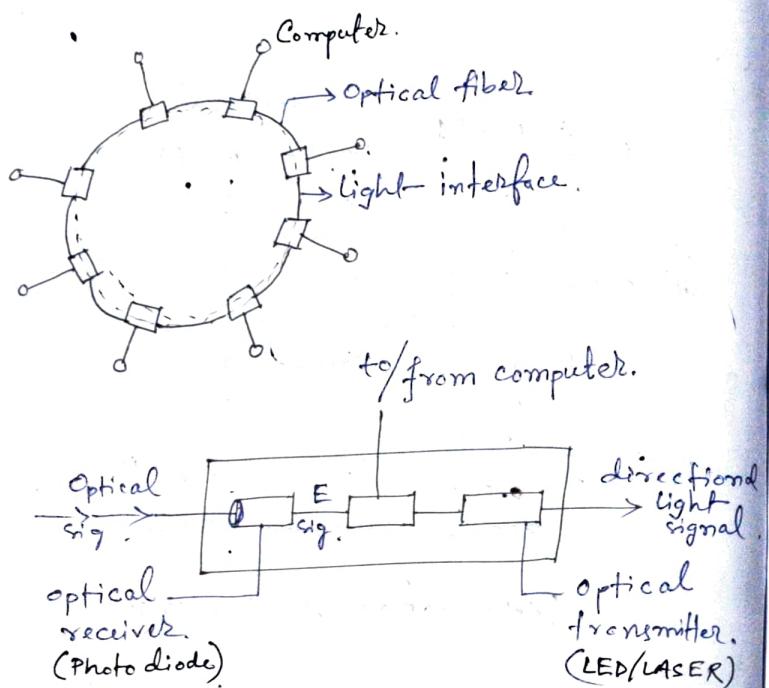
→ transmitting power  
 o A Photodiode  
 , → Receiving power



[Attenuation vs. wavelength graph].

	LED	Semiconductor LASER:
Data rate	Low	High
Fiber type	Multimode	Single or Multimode
Distance	short	long
Life time	Long life	Short life
Temp & Sensitivity	Minor	Substantial
Cost	Less	More

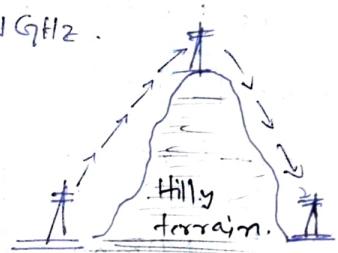
## # Use of Fiber Network:



## Radio & Microwave:

100 kHz — 1 GHz.

Here the problem in transmission is LOS.  
(Light of sight)



## Satelite Communication.

↳ 20—40 GHz

↳ Propagation delay. (250—300 μs)

## # Solitons:

Light signal



(Spreading)



Solitons are the signal without spreading  
& so its called as 'original' signal.

## Signal & Channel Properties:

Bandwidth, Data rate, Band rate, bit rate,  
S/N, BER, Constellation, signal level.

Signal & Channel properties: [27.08.13. Tue]

Medium.

\* Signals flows through the channel.

Bandwidth: Range of frequencies that can pass through the medium.

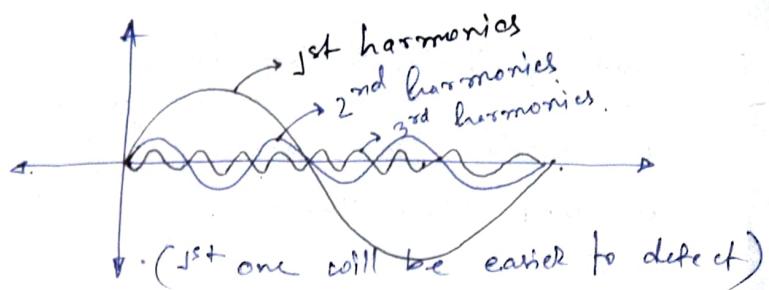
Data rate: signals per unit time is called data rate.

i) Data rates  $\propto$  Bandwidth.

ii) Data rate depends on signal types

- Analog or Digital
- Analog  $\rightarrow$  Data rate  $\propto \log_2(\text{no of constellations point})$
- Digital  $\rightarrow$  Data rate  $\propto \log_2(\text{no of levels})$

iii) Data rate depends receiver sensitivity.



N) Temperature  $\propto$  resistance.

v) Bandwidth of a signal  $\propto$  BW of a data rate. medium. (limited)

Let's say 8 bits/sec |  $\downarrow$  |  $\downarrow$  |  $\downarrow$  |  $\downarrow$  | ; it will be hard  
16 bits/sec |  $\downarrow$  |  $\downarrow$  |  $\downarrow$  |  $\downarrow$  | ; it will be hard

\* data rate  $\uparrow$  signal will be distorted.  
bandwidth is limited.

03.09.13

Tue

ii) MDR (Maximum Data Rate) for binary sig (chnl)

↳ Based on Nyquist equation.

$$\boxed{\text{MDR}_{\text{(channel)}} = 2 \times \text{BW}} ; \text{ for binary sig level.}$$

(Baud rate)

↳ Multiple Signal Level:

$$\text{MDR} = 2 \times \text{BW} \times \log_2(\text{signal level})$$

\* For real system  $\leftarrow$  System with noise (noisy chnl)

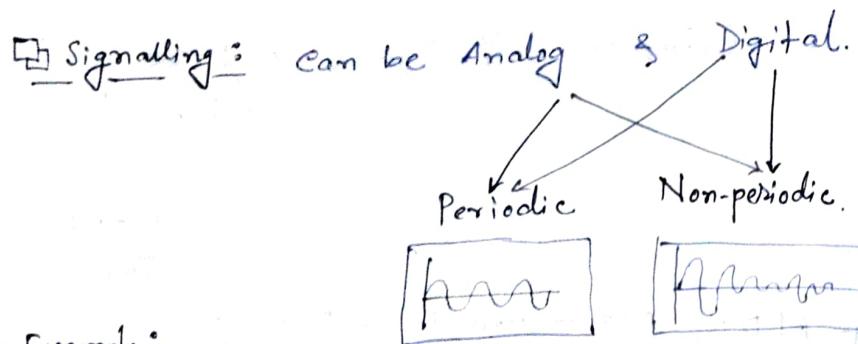
"shannon's theorem"

$$\rightarrow \boxed{\text{MDR} = \text{BW} \times \log_2(1 + \frac{S}{N})} \quad \text{--- (1)}$$

$$\boxed{\frac{S}{N} (\text{dB}) = 10 \log \left[ \frac{P_{\text{sig}}}{P_{\text{noise}}} \right]}$$

if  $\Rightarrow P_s = P_n$ ; then, from eqn (1);

$$\rightarrow \boxed{\text{MDR} = \text{BW} \times \log_2(1 + \frac{10^{\frac{S}{N}/10}}{1})} = 0$$



Example:

Voice Sample

$\hookrightarrow$  Audio CD  $\rightarrow$  Analog.

$\hookrightarrow$  MP3  $\rightarrow$  Digital.

\* Analog signal always refers to low BW.  
 \* Digital  $\square \quad \square \quad \square \quad \square \quad \square$  high  $\square$

\* Dig. sigs bw depends upon the no. of signal levels.

$\hookrightarrow$  ① two different levels. ; where  $\square = +1V$   
 let's say,  $\square \square 0 0 \square \square 0 \square \square$

Orig. sig :  $(+1, +1, -1, -1, +1, +1, -1, +1)$

$\rightarrow \square \square$

$t = 0.1 \mu s = 10^{-7} \text{ sec.}$

$\therefore \text{BW} = \text{[no of bits transmitted per second]}$

=

05.09.13  
 See the Mobile comm copy  
 for this day's class note

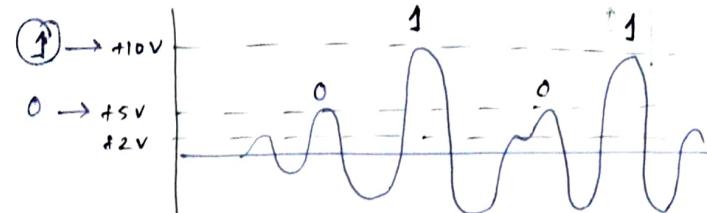
11.09.13

wed

### Coding Mechanism:

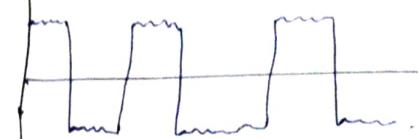
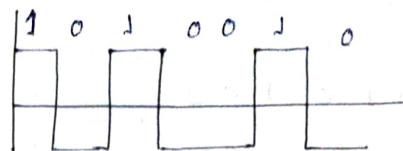
Representation of signal  $\rightarrow$  Shift keying

ASK, FSK, PSK.



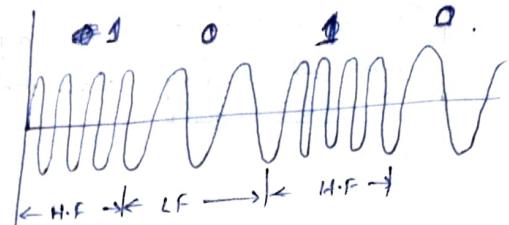
Sequence = 0 1 0 1 .

Digital



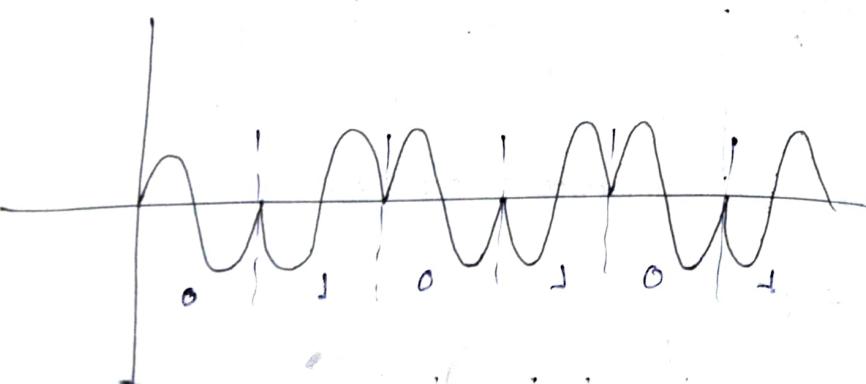
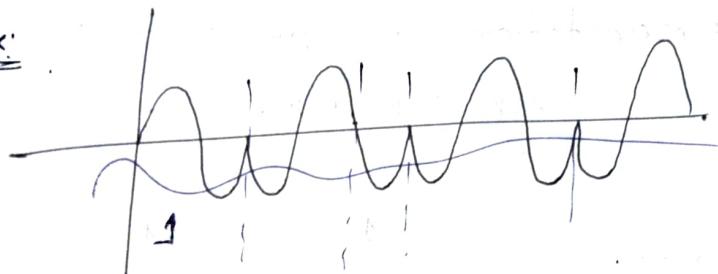
(noisy version)

# FSK:



Sequence: 01.010.

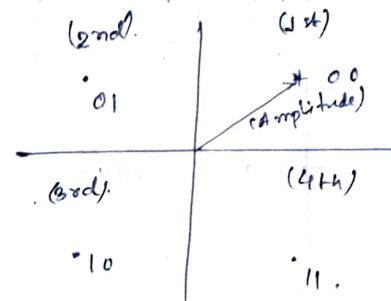
# PSK:



Sequence = 010101

QPSK. →

4 phases.  $\Rightarrow 2^x = 4 \Rightarrow x = 2$  bits  
to represent a signal.

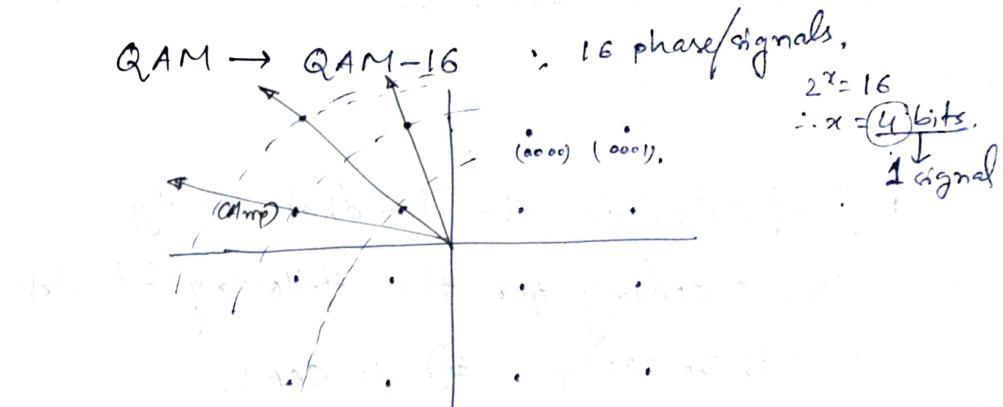


00  
01  
10  
11.

(Constellation Pattern).

[ $45^\circ, 135^\circ, 225^\circ, 315^\circ$ ]  
00 01 10 11.

QAM → QAM-16



16 phase/signals.

$$2^x = 16 \therefore x = 4 \text{ bits.}$$

1 signal

\* This channel is more noisy than QPSK.

Ex sequence  $10110110$

representation in

~~ASK~~

(ASK)

(FSK)

(PSK)

1. **Q** Calculate the baudrate & type the encoding for the following bit rates;

i) 36,000 bps,  $\Rightarrow$  4 QAM.

ii) 8000 bps, 32 QAM.

2. **Q** What is the maximum data rate for a communication channel having a 100kHz bandwidth and operating at S/N ratio of 20 dB.

$$\Rightarrow 1. \text{ Soln: } BR = \frac{R}{\log_2 N}; \text{ for 4 QAM} \\ N=4.$$

①

$$= \frac{36000}{\log_2 4} = 18,000.$$

②

$$\frac{8000}{\log_2 32} =$$

2. Soln.

$$MDR = Bw \times \log_2 (1 + S/N) \rightarrow \log_2 \frac{S}{N} = 20$$

[Assume]

$$= 100 \times 10^3 \times \log_2 \left( \frac{100}{1+100} \right)$$

$$S/N = (10)^{\frac{20}{10}} = 100$$

$$= 665.8211 \text{ kbps.}$$

12.09.13.  
Thu

## Multiplexing Scheme 8

→ Multiple channels through a single physical link.

# FDM → Analog telephone system.

# TDM

# CDMA

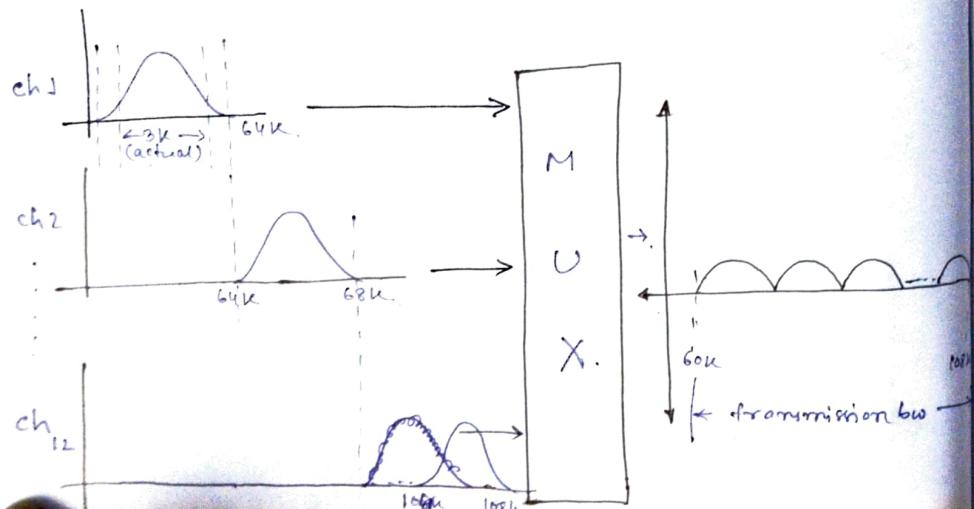
# WDM  
# DWDM ] → optional.

# FDM → Analog telephone system

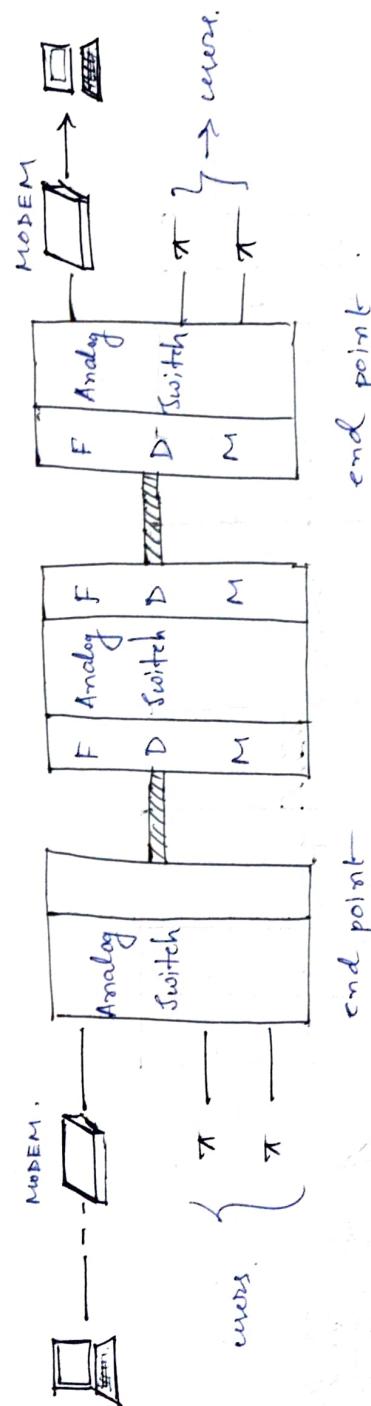
↳ 12 ch. each having a diff. frequency.

→ 4000 Hz. each.

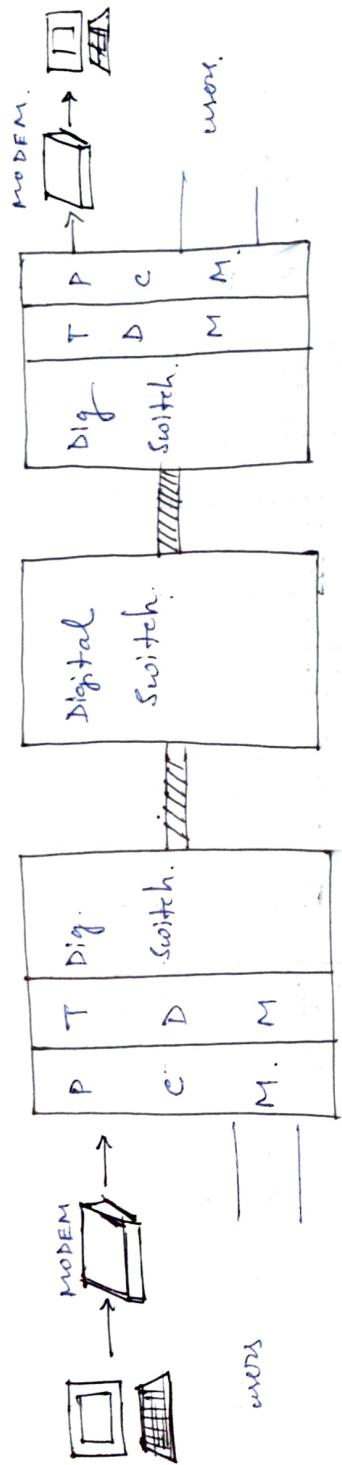
→ Transmission bandwidth  
60 kHz to 108 kHz.



BLOCK DIAGRAM OF FDM SYSTEM



## BLOCK DIAGRAM OF TDM SYSTEM.

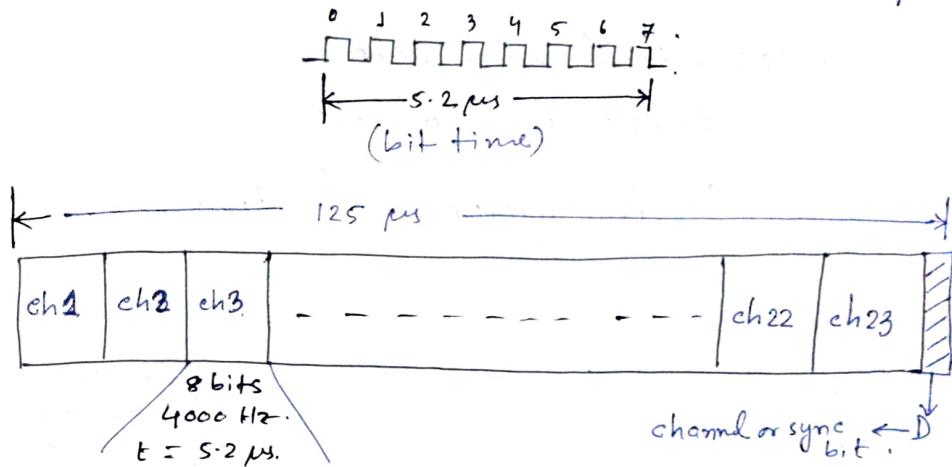


(T1 system)  
# TDM → Digital telephone system.

Specifications:

i) 24 channels ( $23B + 1D$ ),  $125\mu s$ .  
↳ Barrier.  
↳ [fixed for signalling]

ii) 1 channel bw :  $4000\text{ Hz}$ , 8 bits,  $t = 5.2\mu s$ .



Total Data rate:  $2 \times \text{BW} \times \text{No. of bits}$ .

$$= 2 \times 4000 \times (23 \times 8 + 1) = 1.544 \text{ Mbps}$$

channel Data rate:  $2 \times 4000 \times 8 = 64000 \text{ bps}$ .

Transmission BW =  $4000 \times 24 = 96000 \text{ Hz}$ .

### Assignment

E1 system:

specification:

8 bit each

(1 bit)

i) 32 ch. [30B + 1D + 1 (Time & Alarm)]  
(8 bit)

ii) 1 ch. has 8 bits.,  $t = 5.2 \mu s.$

iii) 1 channel bandwidth: 4000 Hz

calc: i) total data rate.

ii) channel data rate.

iii) Transmission bandwidth.

16-09-13  
Monday

CDMA: 40 years old technology.

\* Applied to cellular telecomm.

\* Data rate: 9600 kbps.

\* Data bit + <sup>digital</sup> original code =  $s$  (transmitted sig)  
original sig.  
'1' or '0'

Bipolar coding: checks the orthogonality

inner product of signals = 1 or 0.

$$\text{Orthogonality condition: } c_i \cdot c_j = \frac{1}{n} \sum_{k=1}^n c_{i,k} c_{j,k} = \begin{cases} 1, & \text{if } i=j \\ 0, & \text{if } i \neq j \end{cases}$$

"vector code"

$$c = \begin{bmatrix} c_1 & c_2 & \dots & c_n \end{bmatrix}$$

$$\begin{matrix} \uparrow \\ c_1 \\ \downarrow \\ \vdots \\ \downarrow \\ c_n \end{matrix}$$

Digital code.

(PN) code.

Sequence of code for  $n$ -channel.

$$\text{Transmitted Signal: } s = c_1 + c_2 + c_3 + \dots + c_n$$

# Receiving the signal:

getting  $c_i$  from  $s$ .

$$\Rightarrow c_i \cdot s = c_i \cdot (c_1 + c_2 + \dots + c_n)$$

$$\text{or } \bar{c}_i = \begin{cases} +1 & \\ -1 & \end{cases}$$

$$\bar{c}_i = c_i$$

$$\bar{c}_i = -c_i$$

when channel is idle/silent

Wednesday  
18-09-18

	$c_{ij}$	$\bar{c}_{ij}$	
i	channel	Vector code to transmit "1"	Vector code to transmit "0"
$c_1$ (vector codes)	ch 1	(+1, +1, +1, +1)	(-1, -1, -1, -1)
	ch 2	(-1, -1, +1, +1)	(+1, +1, -1, +1)
	ch 3	(-1, +1, -1, +1)	(+1, -1, +1, +1)
	ch 4	(-1, +1, +1, -1)	(+1, -1, -1, +1)

1 Check the orthogonality.

Trx → 2 Compute the encoded ~~row~~ signals to be transmitted.

Rx → 3 Retrieve the original signal.

Sol:

1 Orthogonality check:

For  $i \neq j$

$$c_1 c_3 = \frac{1}{4} \sum_{k=1}^4 c_{i,k} c_{j,k}$$

$$= \frac{1}{4} (+1, +1, +1, +1)(-1, +1, -1, +1)$$

$$= \frac{1}{4} (-1 + 1 - 1 + 1) = 0$$

Similarly,

$$c_2 c_4 = \frac{1}{4} (-1, -1, +1, +1)(-1, +1, +1, -1)$$

$$= \frac{1}{4} (+1 - 1 + 1 - 1) = 0$$

Now for  $i=j$ ,

$$c_1 c_1 = \frac{1}{4} (+1, +1, +1, +1)(+1, +1, +1, +1)$$

$$= \frac{1}{4} \times 4 = 1.$$

$$c_2 c_2 = \frac{1}{4} (-1, -1, +1, +1)^T = 1.$$

$$c_3 c_3 = c_4 c_4 = 1.$$

∴ proves orthogonality.

ii Encode/transmitted signal:  $c_1, c_2, c_3, c_4$ .

$$c_1 = '0' \text{ (see data bit in table)} \rightarrow (-1, -1, -1, -1)$$

$$c_2 \rightarrow 1 \rightarrow (-1, -1, +1, +1)$$

$$c_3 \rightarrow \text{silent} \rightarrow (0, 0, 0, 0)$$

$$c_4 \rightarrow 0 \rightarrow (+1, -1, -1, +1)$$

Now row signal:  $s = c_1 + c_2 + c_3 + c_4$  (row wise addition)

$c_1 \rightarrow (-1, -1, -1, -1)$
$c_2 \rightarrow (-1, -1, +1, +1)$
$c_3 \rightarrow (0, 0, 0, 0)$
$c_4 \rightarrow (+1, -1, -1, +1)$

$$= (-1, -3, -1, +1)$$

#### 14 Retrieving data:

$$\text{for } \text{ch}1 = c_1 \cdot s = \frac{1}{4} (+1, +1, +1, +1) (-1, -3, -1, +1)$$

$$= -1 ; \text{ which refers to '0'}$$

↓  
data bit  
(Actual Data)

$$\text{for } \text{ch}2 = \frac{1}{4} (-1, -1, +1, +1) (-1, -3, -1, +1)$$

$$= \frac{1}{4} \times (1 + 3 - 1 + 1) = \frac{1}{4} \times 4$$

$$= 1 ; \text{ which refers to "1" (data bit)}$$

$$\text{for } \text{ch}3 = \frac{1}{4} c_3 \cdot s$$

$$= 0 ; \text{ which refers "silent" (data bit)}$$

$$\text{for } \text{ch}4 = \frac{1}{4} c_4 \cdot s$$

$$= -1 , \text{ which refers '0' (data bit)}$$